

Question: 1 (A)

Time Complexity:-

→ Passes:-

```
i = 0 {  
  j = 0 => sum[0] += entry[0][0];  
  ...  
  j = n+1 => sum[0] += entry[0][1];  
}
```

i = 1

```
j = 0 => sum[1] += entry[1][0];  
j = n-1 => sum[1] += entry[1][j-1];
```

i = 2

```
j = 0 => sum[2] += entry[2][0];  
j = n-1 => sum[2] += entry[2][n-1]
```

i = 3

```
j = 0 => sum[3] += entry[3][0]  
j = n-1 => sum[3] += entry[3][n-1]
```

i = 4

```
j = 0 => sum[4] += entry[4][n-1];
```

$$\Rightarrow [(n+1)(n+1)] + n^2$$

$$\Rightarrow 2n^2 + 2n + 1 = f(n) = O(n^2) \text{ Ans.}$$

Question: 2 (B)

Passes:

i = 0 → j = 0

i = 1 → j = 0 = m+0
j = 1

i = 2 → j = 0 = m+0
j = 1 = m+1

i = 3 → j = 0 = m+0
j = 1 = m+1
j = 2 = m+2

i = 4 → j = 0 = m+0
j = 1 = m+1
j = 2 = m+2
j = 3 = m+3

Time complexity:

$$f(n) = \frac{n(n+1)}{2} = \frac{n^2+n}{2}$$

$$f(n) = O(n^2) \text{ Ans}$$

C)

Passes:

$$i = 1 \rightarrow \text{tot} += 1$$

$$i = 1 \times 2$$

$$i = 2 \rightarrow \text{tot} += 2$$

$$i = 2 \times 2$$

$$i = 4 \rightarrow \text{tot} += 4$$

$$i = 4 \times 2$$

$$i = 8 \rightarrow \text{tot} += 8$$

$$i = 8 \times 2$$

$$i = 16 \rightarrow \text{tot} += 16$$

$$i = 16 \times 2$$

Time complexity:

$$f(n) = O(\log n)$$

due to i increases twice
time on every iteration.

D)

Passes

$$i = 32 \rightarrow \text{tot} += 32$$

$$i = 16 \rightarrow \text{ } \text{ } \text{ } += 16$$

$$i = 8 \rightarrow \text{ } \text{ } \text{ } += 8$$

$$i = 4 \rightarrow \text{ } \text{ } \text{ } += 4$$

$$i = 2 \rightarrow \text{ } \text{ } \text{ } += 2$$

Time complexity

$$f(n) = O(\log n)$$

E)

Passes:

Same as Part (A)

Time complexity.

$$f(n) = (n^2 + n + 1)(n + 1)$$

$$f(n) = O(n^3)$$

F)

Passes

$$i = 0 \rightarrow j = 0, n-1$$

$$i = 1 \rightarrow j = 0, n-1$$

$$i = 2 \rightarrow j = 0, n-1$$

$$i = 3 \rightarrow j = 0, n-1$$

$$i = 4 \rightarrow j = 0, n-1$$

Time complexity

$$f(n) = O(n^2)$$

9)

Passes.

- $i=0 \rightarrow k=0, n-1$
- $i=1 \rightarrow k=0, n-2$
- $i=2 \rightarrow k=0, n-3$
- $i=3 \rightarrow k=0, n-4$
- $i=4 \rightarrow k=0, n-5$

Time complexity

$$f(n) = O(n^2)$$

h)

Passes

- $i=0 \rightarrow j=0, \text{sev}(n)-1$
- $i=1 \rightarrow j=0, \text{sev}(n)-1$
- $i=2 \rightarrow j=0, \text{sev}(n)-1$
- $i=3 \rightarrow j=0, \text{sev}(n)-1$
- $i=4 \rightarrow j=0, \text{sev}(n)-1$

Time complexity

$$f(n) = O(n \times \text{sev}(n))$$

i)

Passes

- $i=0$
- $i=1$
- $i=2$
- $i=3$
- $i=4$

Time complexity

$$f(n) = O(n)$$

j)

Passes

- $i=0$
- $i=1$
- $i=2$
- $i=3$
- $i=4 \rightarrow i=0$
- $i=1$
- $i=2$
- $i=3$
- ~~$i=4$~~

Time complexity

$$f(n) = O(n^2)$$

k)

Passes

- $i=0 \rightarrow j=0$
- $i=1 \rightarrow j=0$
- $i=2 \rightarrow j=0, 1$
- $i=3 \rightarrow j=0, 1, 2$
- $i=4 \rightarrow j=0, 1, 2, 3$

Time complexity

$$f(n) = O(n^2)$$

Question: 2

A) Time complexity:

$$\begin{aligned}f(n) &= [(n+1)(n+1)] + n^2 \\&= 2n^2 + 2n + 1 \\&= O(n^2) \quad \underline{Ans}\end{aligned}$$

B) Time complexity:

$$\begin{aligned}f(n) &= [(n+1)(2n+1)] + n^2 & \therefore i = n \\&= 3n^2 + 3n + 1 \\&= O(n^2) \quad \underline{Ans}\end{aligned}$$

C) Time complexity:

$$f(n) = O(n \log n) \rightarrow \text{increases (i) twice times \& (n) times.}$$

Question: 3

A) Time complexity:

$$\begin{aligned}f(n) &= (n \log n) \times (n+1) + n \log n \\&= O(n \log n), \rightarrow \text{here also (i) is increased 2 times.}\end{aligned}$$

B) Time complexity:

$$\begin{aligned}f(n) &= [(n+1)(n+1)(n+1)] + n^3 \\&= O(n^3) \rightarrow \text{dominating term is } (n^3).\end{aligned}$$

Question: 4

A)

↳ Insertion is the best choice to it's efficiently with nearly sorted data. Selection & Bubble sort are less efficient for this scenario due to their higher numbers of unnecessary comparisons & swaps.

5)

1) In-Place Sorting:

- Insertion sort: yes
- Bubble sort: yes
- Selection sort: yes

2) online Sorting:

- Insertion sort: yes
- Bubble sort: no
- Selection sort: no

3) Stable Sorting:

- Insertion sort: yes
- Bubble sort: yes
- Selection sort: no