

Longest Path for Knight on 32x32 Chess Board

Problem importance: A 32x32 chess board has 1024 cells/nodes, and a pure exploration approach requires searching up to depths of 1024. This combined that each node have a maximum exploration factor 8 child nodes make a direct search approach infeasible (8^{1024}). Tried applying some heuristics in addition to Wansdorff approach [1], but could not attain performance good enough for present available computing power.

Background: As proved by Conrada [3], for any chess board with of dimension (n), and $n \geq 6$ there will always exist a Hamilton path [2] from any start s to any end e .

Basic approach: Using Conrada's theorem, the problem for *longest path* for given 32x32 chess board can now be rephrased as *finding Hamilton path*.

As finding Hamilton path for a 32x32 chess board is in itself very expensive following divide and conquer approach is applied.

Basic concept: *The 32x32 board is divided into a set of 6x6, 6x8, 8x6 and 8x8 sub-boards, as shown figure-1. Now a Hamilton path computed in each of the sub-boards such that, these smaller path can then be stitched together (combined) to form one complete path covering almost all of 32x32 chess board cells.*

There are more details and scenarios to consider to this presumably simple solution. Some such questions are:

- How to divide 32x32 board into smaller boards such that start and end nodes fall in 6x6 sub-boards.
- How to find the start and end points of each sub-boards such that they can stitched with other paths.
- In what sequence should the path cross over these sub-boards.
- How to optimally find a Hamilton path within each sub-board.

Here are two of the major scenarios:

a) **Length of the longest path:** Observing that (i) there are equal number of black and white nodes/cells on a chessboard with even size dimensions, and (ii) knight's movement is such that at each step it jumps to a cell with a complimentary color. That is a knight in white cell will only reach cells of black color in single step. Hence, if $color(start\ cell) == color(end\ cell)$, the longest possible path will only cover $N-1$ cells. Where N is total number of cells on the board. That is:

$$\begin{aligned} \text{Length of longest path} &= N \text{ if } (color(start) \neq color(end)) \\ &= N-1 \text{ if } (color(start) == color(end)) \end{aligned}$$

Each of these situations are handled separately when finding Hamilton paths within sub-boards

b) When start and end cells fall on different sub-boards, it is easy to find individual Hamilton paths within each sub-board and join them. But when both the start and end cells fall within same sub-board it requires additional. The Hamilton path within this common sub-board needs to be divided into two chunks one which forms the head of the final path, and other which forms the tail of the final path.

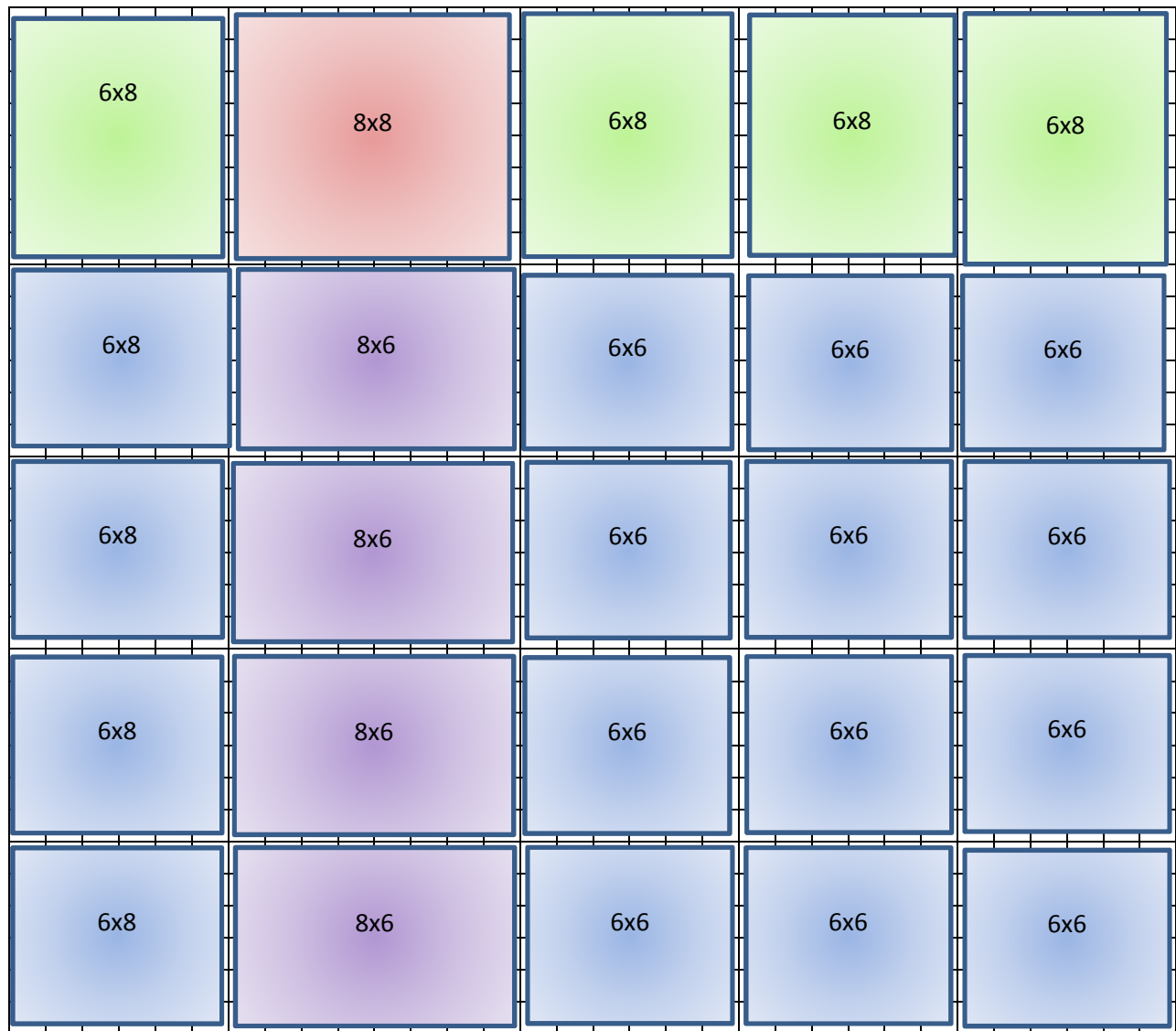


Figure-1: One of the scenario of dividing a 32x32 board into smaller sub-boards, such that each dimension of 32d is covered by (four of 6d + one of 8d). (Axel Conrada n.d.)

Reference:

- [1] Warnsdorff Approach: <http://www.slac.stanford.edu/cgi-wrap/getdoc/slac-pub-0261.pdf>
- [2] Hamilton Path: http://en.wikipedia.org/wiki/Hamiltonian_path
- [3] "Solution of the knight's Hamiltonian path problem on chessboards" Axel Conrada, Tanja Hindrichsb, Hussein Morsyc, Ingo Wegener <http://www.sciencedirect.com/science/article/pii/S0166218X9200170Q>

Appendix:

Some notes on diagonal and side way stitching sub-boards:

