

## INTRODUCTION TO NUMBER THEORY

### EUCLID'S ALGORITHM:

{ on input  $a, b$  where  
 $a, b$  are  $\mathbb{Z}$  such  
that  $a \geq b \geq 0$  }

The algorithm follows three steps:

Step 1: If  $b=0$ , then return the value of  $a$ .

Step 2: Otherwise, divide  $a$  by  $b$  and store the remainder in some variable  $r$ .

[which is nothing but modulo operation]

Step 3: Let  $b=r$  and  $a=b$  and return to step.

Step 4: Continue this process until  $b=0$ .

### FOR EXAMPLE:

Let us consider the inputs as

$$a = 25 \quad \boxed{b = 10}$$

$$r = a \% b \quad := \quad r = 5$$

$$b = 5 \quad a = 10$$

$$r = a \% b \quad := \quad r = 0$$

$$\boxed{b = 0} \quad a = 5$$

$$a \neq 0 \quad b \neq 0$$

$$r = a \% b \quad \neq$$

$$\therefore \gcd(25, 10) = a = 5$$

```
int gcd (int a, int b)
```

```
{   int r;
```

```
    if (b == 0)
```

```
        return a;
```

```
    else
```

```
        r = a % b; } a = b and b = r
```

```
        gcd(b, r); }
```

```
}
```

### Time complexity

The euclidean algorithm computes  $\text{gcd}(a, b)$ , where  $a \geq b \geq 0$ .

- \* The no. of divisions with remainder is  $O(\log(b))$
- \* The rough estimate of the total time is  $O(\log(b)^2 \log(a))$
- \* Moreover it can be proved that the Euclidean algorithm only needs  $O(\log(b) \log(a))$  time.