My section will cover the first half of Chapter 3 which is the heart of the paper. NEXT SLIDE.

A sigma protocol is an interactive proof of knowledge where a prover needs to prove that it knows some piece of information. Zero knowledge means it needs to prove this without revealing the information itself. As an example, consider Alice and Bob where Alice knows the password to a locked door in a cave with 2 paths. She randomly chooses a path. Bob, who doesn't know which path she took, randomly shouts out one of the paths. If Alice is able to reliably return to the entrance using the path Bob randomly shouted out when this experiment is repeated many times, it is extremely likely that she knows the password to open the door in the cave that allows her to appear at either side. Note that she never reveals the password. NEXT SLIDE.

Now let's discuss key generation. Here you can see the formula that will be used to find a secret ideal.  $O_0$  is a maximal order meaning it is the largest possible proper subset of the quaternion algebra that also forms a ring with addition and multiplication. Gamma is a random element of  $O_0$ , 'a' is a random positive scalar less than  $D_{\text{secret}}$ , i is the quaternion element such that  $i^2 = -1$ .  $D_{\text{secret}}$  acts as the norm of  $I_{\text{secret}}$  meaning it is the gcd of the elements of  $I_{\text{secret}}$ . Then alpha which connects  $I_{\text{secret}}$  to an equivalent ideal with a power of 2 norm called  $J_{\text{secret}}$  is computed. NEXT SLIDE

Here you can see the computation of the secret isogeny from the secret ideal, the maximal order, and the basis for  $E_0[T]$ .  $E_0[T]$  is the T-torsion subgroup of  $E_0$  meaning the set of all points on  $E_0$  such that scalar multiplication with T yields infinity, the identity element. In  $SQI_{SIGN}$  we assume T is smooth meaning it has only small prime factors. For this case, an algorithm like the Pohlig-Hellman algorithm guarantees an efficient solution. The public key,  $E_A$ , is generated and  $\phi_{secret}$  is modified to map to this normalized curve. Then  $B_{A,T}$  (a basis for  $E_A$ 's T-torsion subgroup) is found by applying this secret isogeny to  $B_{0,T}$ . NEXT SLIDE

Here P is one of the basis points for the intersection of the kernel of the secret isogeny and the  $2^f$ -Torsion subgroup of  $E_0$ . The other basis point, Q, is found using the CompleteBasis algorithm. The generating point Q is mapped from where it is on  $E_0$  to a point on  $E_A$ . After a few more steps, the signing key and public key are returned and the key generation phase is complete.