

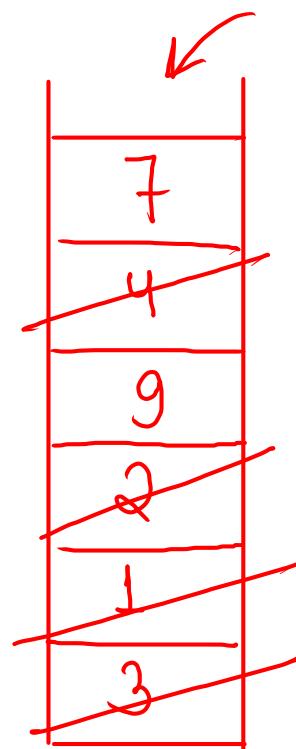
HW_Next greater element on left 1

$$\text{arr} = \begin{bmatrix} 0 & 1 & 2 & 3 & 4 & 5 \\ 3, 1, 2, 9, 4, 7 \end{bmatrix}$$

$$\text{ans} = \begin{bmatrix} -1, 3, 3, -1, 9, 9 \end{bmatrix}$$

$$A[0] = [\begin{array}{ccccccc} 0 & 1 & 2 & 3 & 4 & 5 \\ 3, & 1, & 2, & 9, & 4, & 7 \\ \uparrow & \uparrow & \uparrow & \uparrow & \uparrow & \uparrow \end{array}]$$

Ex:-
dry run



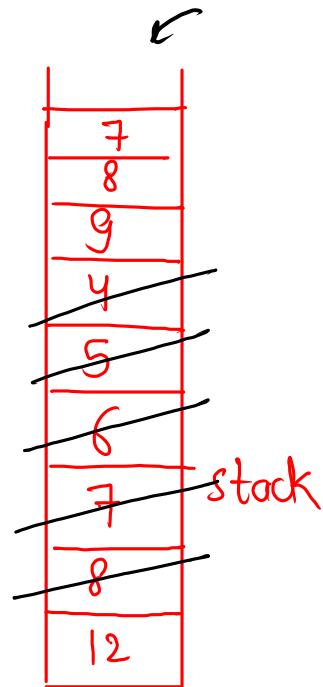
Ex:-

$$arr = [12, 8, 7, 6, 5, 4, 9, 8, 7]$$

↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓

$$ans = [-1, 12, 8, 7, 6, 5, 12, 9, 8]$$

Note :- after destroying all possible elements, whichever no. is present on top of stack that no. is your ans



pseudo code

- 1) declare stack and array ans
- 2) loop from start to end
 - 2.1) loop until peek \leq curr
st.pop()
 - 2.2) ans[i] = st.peek()
 - 2.3) st.push(curr[i]);

Code

```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();
    int[] arr = new int[n];
    for (int i = 0; i < n; i++) {
        arr[i] = scn.nextInt();
    }
    int[] ans = greaterElementOnLeft(arr, n);
    for (int i: ans) {
        System.out.print(i + " ");
    }
}
public static int[] greaterElementOnLeft(int[] arr, int n) {
    Stack<Integer> st = new Stack<>();
    int[] ans = new int[n];
    for (int i = 0; i < n; i++) {
        while ( st.size() > 0 && st.peek() <= arr[i] ) {
            st.pop();
        }
        if (st.size() > 0) {
            ans[i] = st.peek();
        } else {
            ans[i] = -1;
        }
        st.push(arr[i]);
    }
    return ans;
}
```

$$T.C = O(N)$$

$N = \text{size of arr}$

Imp

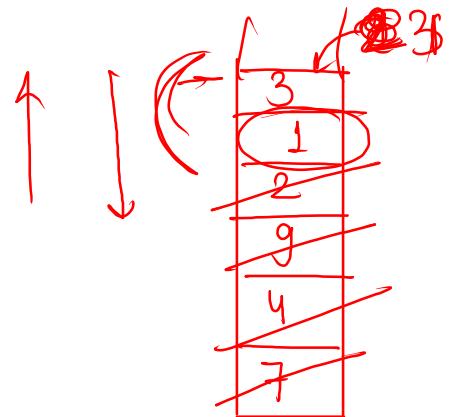
- 1) Greater element on left side
- 2) Greater element on right side :- just reverse the loop
- 3) Smaller element on left side :- change condⁿ inside inner loop.
- 4) Smaller element on right side :- use both

Next Smaller Element To The Right

```
public static int[] greaterElementOnLeft(int[] arr, int n) {  
    Stack<Integer> st = new Stack<>();  
    int[] ans = new int[n];  
    for (int i = n - 1; i >= 0; i--) {  
        while (st.size() > 0 && st.peek() >= arr[i]) {  
            st.pop();  
        }  
        if (st.size() > 0) {  
            ans[i] = st.peek();  
        } else {  
            ans[i] = -1;  
        }  
        st.push(arr[i]);  
    }  
    return ans;  
}
```

arr = [3, 1, 2, 9, 4, 7]

ans = [1, -1, -1, 4, -1, -1]



⇒ HashMap

[every operation in HM]
is having T.C constant]

↳ actually :- Amotized O(1)

Structure

hashmap

key → value

key	value
"India"	400
"England"	220
"SriLanka"	170
"Aus"	17
"Pak"	-10
"aus"	170

key (String)	Value (Integer)
Integer	Integer
Boolean	Boolean
Character	Character
String	String
Double	Double
Float	Float
Long	Long
ArrayList	ArrayList
Stack	Stack
Queue	Queue
PriorityDower	PriorityDower
HashMap	HashMap

⇒ Imp notes for hashmap

- 1) all keys are always unique
- 2) Keys are case sensitive
- 3) values can be repeated
- 4) if same key is added again in hashmap
then previous pair will be overridden
- 5) data is unorganised
(no indexing)

Syntax

HashMap< keyDataType, ValueDataType > map = new HashMap<>();

HashMap< String, Integer > map = new HashMap<>();

Inbuilt

- map.put(key , value); // to add a pair in HM
- map.get(key); // return value of given key
- map.remove(Key); // to delete entire pair

(Evergreen fn) :- [map.size() , map.isEmpty()]

↳ `map.containsKey(key);` // true if key is present
in HM, false otherwise

↳ `map.containsValue(value);` // true if value is present
in HM, false otherwise

Note:- all these tasks are completing in
 $O(1)$ time.

Ex:-

String vs Integer

"abc" → 3

"efg" → 8

"EFG" → 7

map.put("abc", 3);

~~map.put(4, "efg");~~ error

map.put("efg", 4);

map.put("EFG", 7);

map.put("efg", 8);

map.get("efg"); // 8

map.isEmpty(); // false

map.get("xyz"); // error, null

map.containsKey("abc"); // true

Map.containsValue(27); // false