

Genetic Algorithms

Theory and Practice

A Comprehensive Guide to Evolutionary Optimization

Course Materials Collection

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Chapter 1

Introduction to Optimization and Evolutionary Computation

Genetic algorithms (GAs) are population-based search methods inspired by natural selection and genetics. They maintain a population of candidate solutions encoded as strings, iteratively producing new generations by selecting the fittest individuals, recombining their information, and occasionally introducing random variation. Although stochastic in their operators, GAs are not blind random walks: they retain and exploit historical information about good solutions to generate promising new search points and thereby drive efficient exploration and exploitation of complex spaces.

This family of methods was developed from foundational work by Holland and colleagues to both model adaptive processes observed in nature and to design artificial systems that embody those mechanisms. A central aim has been robustness—the ability to balance efficiency with reliability across a wide range of problem environments—which makes GAs attractive when redesign costs are high or when problem structure violates common assumptions (e.g., continuity, differentiability, or unimodality). Because they are conceptually simple, widely applicable, and empirically effective in optimization and control, genetic algorithms have become a practical tool across engineering, science, and business domains [21].

To place genetic algorithms in context, we first give a concise definition of optimization—the class of problems GAs are commonly used to solve. We define optimization as the process of finding the best solution from a set of available alternatives. In mathematical terms, an optimization problem can be formulated as:

$$\begin{aligned} & \text{minimize (or maximize)} && f(x) \\ & \text{subject to} && g_i(x) \leq 0, \quad i = 1, 2, \dots, m \\ & && h_j(x) = 0, \quad j = 1, 2, \dots, p \\ & && x \in X \end{aligned} \tag{1.1}$$

where:

- $f(x)$ is the objective function to be optimized
- $g_i(x)$ are inequality constraints
- $h_j(x)$ are equality constraints
- X is the feasible region

Optimization problems differ by variable types (discrete, continuous, or mixed-integer) and by structural properties such as linearity, convexity, and the number of objectives. Traditional solution methods include gradient-based techniques (e.g., Newton and quasi-Newton methods) for smooth continuous problems, linear programming (Simplex and interior-point methods) for linear models, and discrete methods (branch-and-bound, dynamic programming) for combinatorial problems. However, these traditional methods

have limitations when applied to many complex real-world problems. In the following sections we highlight three key challenges where conventional approaches often struggle, and explain how genetic algorithms can help address them.

The first problem with traditional optimization methods is their tendency to get trapped in local optima. In multi-modal landscapes with many peaks and valleys, gradient-based searches can converge to a local optimum rather than the global optimum. This happens because these methods rely on local gradient information to guide the search process. When the search reaches a local optimum, the gradient becomes zero, causing the algorithm to stop progressing. This limitation is particularly problematic in high-dimensional spaces where the number of local optima can grow exponentially.

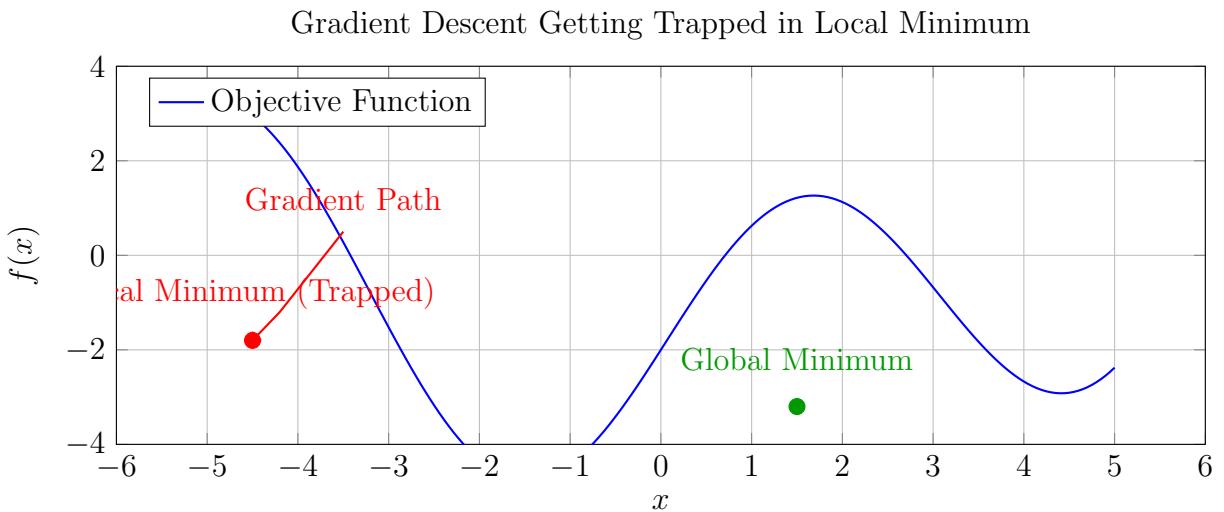


Figure 1.1: Traditional gradient-based methods follow the local gradient and become trapped in local optima, unable to escape to find the global optimum.

The second problem with gradient-based methods is that they require the objective function to be differentiable. This becomes a significant limitation when dealing with real-world problems that involve discontinuities, sharp corners, or discrete jumps. Such problems are common in engineering design, scheduling, and combinatorial optimization.

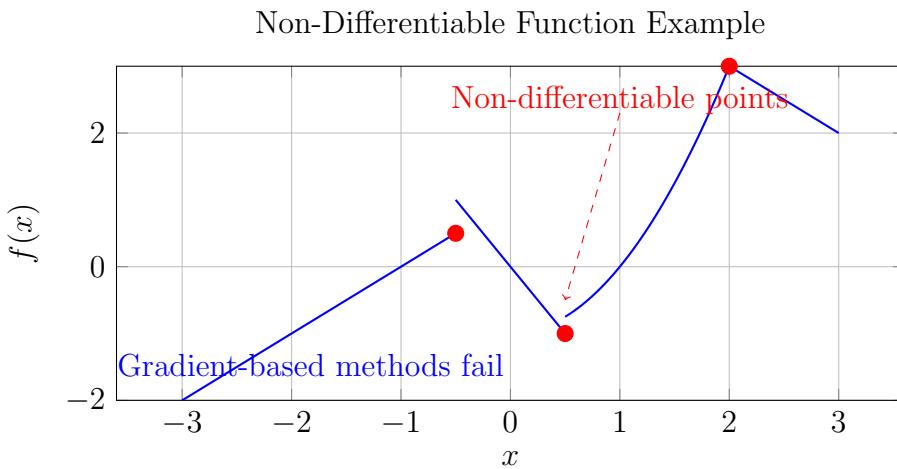


Figure 1.2: Functions with discontinuities, sharp corners, or discrete jumps cannot be optimized using gradient-based methods.

While discrete methods such as dynamic programming can handle discontinuities and combinatorial structure, both gradient-based and exact discrete algorithms suffer from the curse of dimensionality: computation typically becomes intractable as problem dimensionality grows.

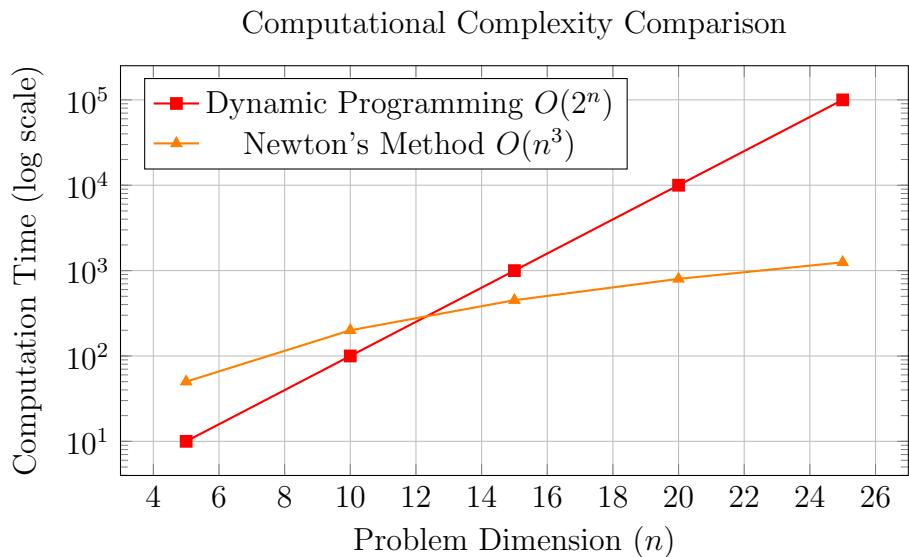


Figure 1.3: Traditional optimization methods often exhibit exponential or high polynomial growth in computation time as problem dimensionality increases, making them impractical for large-scale problems.

Therefore, we can summarize how genetic algorithms effectively address the fundamental limitations of traditional optimization methods in the following table:

Feature	Traditional	GA
Search Strategy	Local	Global
Parallelization	Difficult	Natural
Multi-objective	Complex	Built-in
Constraint Handling	Moderate	Flexible

Figure 1.4: Comprehensive comparison showing how GAs address the fundamental limitations of traditional optimization methods.

1.1 Introduction to Evolutionary Computation

Evolutionary computation is a family of algorithms inspired by biological evolution [25, 13]. These algorithms use mechanisms such as selection, which implements the principle of survival of the fittest, reproduction for creating offspring, mutation to introduce random changes, and crossover for combining genetic material from different individuals.

1.1.1 Advantages of Evolutionary Approaches

Evolutionary approaches offer several compelling advantages that make them attractive for complex optimization problems. They require no gradient information, which is particularly valuable when dealing with black-box optimization scenarios. These algorithms can effectively handle discontinuous, noisy, and multi-modal functions that would challenge traditional optimization methods. Their versatility extends to both continuous and discrete optimization problems, making them applicable across diverse domains. The population-based search strategy provides inherent robustness against local optima, and when properly configured, these algorithms can find global optima in challenging search spaces.

1.1.2 Types of Evolutionary Algorithms

The field of evolutionary computation encompasses several distinct algorithmic families, each with its own characteristics and strengths. Genetic Algorithms (GA) are inspired by natural selection and were among the first evolutionary approaches developed [25, 21]. Evolution Strategies (ES) focus specifically on real-valued optimization problems and have proven particularly effective in continuous domains [6]. Evolutionary Programming (EP) places emphasis on behavioral evolution rather than genetic representations [17, 18]. Finally, Genetic Programming (GP) extends evolutionary principles to the evolution of computer programs themselves, enabling the automatic generation of code [27].

1.1.3 Applications of Evolutionary Computation

Evolutionary algorithms have demonstrated remarkable success across a wide range of practical applications [20, 38, 13]. In engineering design optimization, they have been used to find optimal configurations for complex systems [47]. The field of machine learning has benefited from evolutionary approaches, particularly in neural network training where they can optimize both architectures and weights [32, 33]. Scheduling and timetabling problems, which are notoriously difficult combinatorial optimization challenges, have been successfully addressed using evolutionary methods [23, 16]. Beyond these areas, evolutionary algorithms have found applications in financial modeling, bioinformatics, and game playing strategy development, demonstrating their versatility across diverse problem domains.

1.1.4 Simple Examples of Genetic Algorithm Applications

Genetic Algorithms possess an extraordinary capability in finding optimal solutions for problems with vast search spaces.

1.1.5 Maximizing a Quadratic Function

A classic example often used to illustrate how Genetic Algorithms work is the problem of maximizing a simple quadratic function, such as $f(x) = x^2$ in the range $0 \leq x \leq 31$. In this case, the optimal solution is clear: $x = 31$, yielding $f(x) = 961$. However, this problem is used to demonstrate how Genetic Algorithms, despite starting with a random population (for example, four 5-bit binary chromosomes), can gradually combine features (building blocks) from existing solutions to achieve better results.

Through cycles of selection, crossover, and mutation, Genetic Algorithms show how both the maximum and average performance of the population increases from one generation to the next. Within a few generations, individuals representing $x = 31$ (binary 11111) will dominate the population.

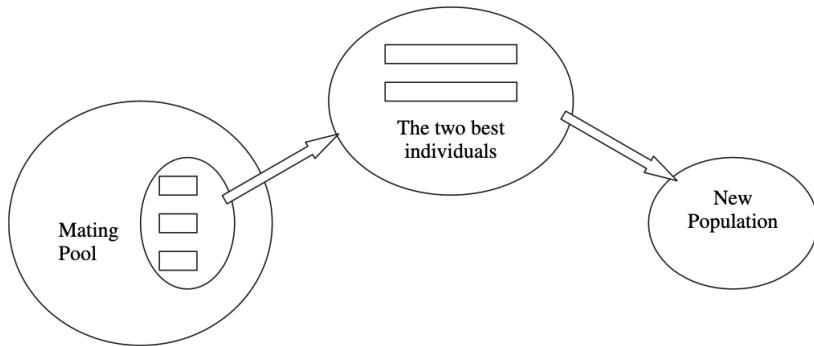


Figure 1.5: Illustration of GA cycle and variations

1.1.6 Travelling Salesman Problem

Another example of genetic algorithm application is the Travelling Salesman Problem (TSP), where Genetic Algorithms are used to find the shortest route visiting a number of cities—a highly complex combinatorial optimization problem.

1.2 Genetic Algorithm Flow

Genetic algorithms maintain a population of individuals, denoted as $P(t)$ for generation t , where each individual represents a potential solution to the problem at hand. This cycle runs iteratively through steps that mimic the evolution process.

After the initial population $P(0)$ is initialized and its fitness evaluated, the selection process begins, where fitter individuals are probabilistically selected to enter the mating pool. Selected individuals then undergo stochastic transformation through genetic operators. The crossover operator is responsible for exploiting the best information from parents, while mutation is tasked with exploring the search space by introducing new genetic material.

Newly generated individuals, called offspring $C(t)$, are then evaluated for their fitness. A new population $P(t + 1)$ is formed by selecting fitter individuals from both the parent and offspring populations through a replacement scheme. This process is repeated for several generations until the algorithm converges, pointing to the best individual, which is expected to represent an optimal or suboptimal solution to the problem.

1.3 Genetic Algorithm Variations

Although the Simple Generational Genetic Algorithm serves as the basic framework, various modifications have been developed to improve performance in dealing with problem complexity. Some of these modifications include:

1.3.1 Hybrid GA

Hybrid GA (HGA) combines Genetic Algorithms with conventional local search techniques [29, 33]. In a hybrid approach, Genetic Algorithms perform global exploration across the population, while local search is used for refinement around promising solutions. This approach often outperforms single methods because it leverages the complementary advantages of both search techniques.

1.3.2 Adaptive GA

Adaptive GA (AGA) is a Genetic Algorithm where strategic parameters (such as crossover or mutation probabilities) are dynamically adjusted during the evolution process, often using feedback from population performance to balance exploitation and exploration effectively [37, 41].

1.3.3 Parallel GA

Parallel Genetic Algorithms divide the population into different sub-populations across various processors, allowing periodic information exchange (migration) to increase diversity and convergence speed.

1.4 Chapter Summary

This chapter introduced the fundamental concepts of optimization and evolutionary computation. We explored the limitations of traditional optimization methods and highlighted the advantages of evolutionary approaches. We also examined specific examples of GA applications including function optimization, TSP, scheduling, network design, and VLSI design. The GA flow and various modifications were discussed to provide a comprehensive understanding of genetic algorithms.

1.5 Key Concepts

- Optimization problem formulation
- Objective functions and constraints
- Local vs. global optima
- Evolutionary computation principles
- Population-based search

1.6 Further Reading

- Deb, K. (2001). Multi-objective optimization using evolutionary algorithms.
- Eiben, A. E., & Smith, J. E. (2015). Introduction to evolutionary computing.

- Goldberg, D. E. (1989). Genetic algorithms in search, optimization, and machine learning.

Chapter 2

What is a Genetic Algorithm?

2.1 Introduction

Genetic Algorithms (GAs) are randomized search and optimization methods that take inspiration from natural evolution [25, 21]. Rather than improving a single candidate solution, a GA maintains a population of potential solutions and applies biologically inspired operators—selection, recombination (crossover), and mutation—to create successive generations of solutions. This population-based approach enables exploration of multiple regions of the search space in parallel and, together with stochastic variation, often helps the algorithm avoid becoming trapped in local optima.

GAs are particularly useful for problems with large, complex, or poorly understood search spaces where gradient information is unavailable or unreliable. Their flexibility in representation and operators makes them applicable to a wide variety of domains, from combinatorial optimization to continuous parameter tuning and symbolic regression.

oprule	Individual	Binary	Decimal	Fitness
1		01101	13	169
2		11000	24	576
3		01000	8	64
4		10011	19	361

Table 2.1: Initial Population Example

The table above shows a small initial population encoded in binary, together with each individual’s decoded decimal value and its fitness. This simple example illustrates how candidate solutions are represented and evaluated, which is the first step in any genetic algorithm implementation.

After initialization, the GA iteratively evaluates individuals, selects parents based on fitness, applies crossover and mutation to produce offspring, and then forms the next generation. Through these repeated cycles, the population tends to improve and the algorithm converges toward high-quality solutions, subject to the chosen encoding, fitness function, and operator settings.

2.2 Biological Inspiration

Genetic algorithms borrow their core ideas from the theory of natural selection and adaptation. In biological populations, variation arises through recombination and mutation, individuals compete for limited resources, and those with heritable traits that confer higher reproductive success tend to leave more offspring. Over many generations this process leads to populations that are better adapted to their environment; the GA framework abstracts these mechanisms to drive improvement of candidate solutions in a search process [25, 31].

Within the GA metaphor, selection favors higher-fitness individuals as parents for the next generation, crossover combines genetic material from parents to explore new regions of the search space, and mutation introduces random changes that maintain genetic diversity and allow previously unseen solutions to arise. These mechanisms—selection, recombination, and mutation—work together to balance exploration and exploitation during the search, enabling gradual adaptation of the population toward higher-quality solutions [13].

2.3 Basic Terminology

2.3.1 Genetic Algorithm Terms

Understanding common terminology helps bridge the biological metaphor and its algorithmic implementation [31, 21]. An **individual** or **chromosome** denotes a single candidate solution; it is composed of one or more **genes**, where each gene represents a component of the solution and an **allele** is the specific value held by a gene. A **population** is the collection of individuals that the algorithm maintains at any given time, and a **generation** refers to one iteration of the evolutionary cycle in which selection, recombination, and mutation produce the next population. The **fitness** of an individual quantifies its quality with respect to the optimization objective and is used to bias selection toward better solutions. The **genotype** describes the encoded representation used by the algorithm (for example, a binary string or a vector of real values), while the **phenotype** is the decoded or interpreted form of that genotype (the actual solution instance evaluated by the fitness function). Clarifying these terms is useful when designing representations and operators, because implementation choices at the genotype level determine what phenotypes can be expressed and therefore influence the search behavior and effectiveness of the GA [13].

2.4 Basic Structure of a Genetic Algorithm

A genetic algorithm (GA) is a population-based, stochastic search procedure that transforms a set of candidate solutions through repeated application of variation and selection operators. Formally, a GA can be described by the tuple (X, Φ, f, S, C, M, R) where X is the search space (phenotypes), Φ is an encoding that maps genotypes to phenotypes, $f: X \rightarrow \mathbb{R}$ is the fitness function, S is a selection operator, C a recombination (crossover) operator, M a mutation operator, and R a replacement (survivor selection) operator. At generation t the algorithm maintains a population $P_t \subseteq \Gamma$ of genotypes (where Γ denotes the set of encodings); the operators act to produce a new population P_{t+1} according to the scheme

$$P_{t+1} = R(P_t, \{C \circ M(\pi) : \pi \in \Pi(S(P_t))\}),$$

where $S(P_t)$ denotes the multiset of parent selections drawn from P_t , $C \circ M$ indicates that offspring are produced by applying mutation and crossover to selected parents, and R determines which individuals survive into the next generation. This abstract description captures the canonical loop of initialization, evaluation, selection, variation, and replacement which repeats until a termination condition (for example, a fixed computational budget, a target fitness, or lack of improvement) is satisfied [25, 31, 13].

In practice the design of each component strongly influences search behavior. The encoding Φ determines what solutions can be represented and how variation operators

explore the phenotype space; the fitness function f defines the optimization objective and provides the selection signal; the selection operator S controls the selective pressure toward higher-fitness individuals (examples include fitness-proportionate, tournament, and rank-based schemes); recombination C mixes information between parents to explore new regions of the search space; mutation M introduces random perturbations to preserve diversity and enable local exploration; and the replacement policy R balances retention of good solutions with introduction of fresh offspring. These design choices embody the exploration–exploitation trade-off discussed in Section ?? and are the primary levers for adapting a GA to a particular problem domain [21, 13].

Viewed algorithmically, the GA performs the following high-level steps each generation: evaluate f on P_t , select parents using S , produce offspring via C and M , and form P_{t+1} via R . Although many variants exist (steady-state updates, island models, hybrid schemes combining local search), this canonical structure explains both the empirical flexibility of GAs and the reasons for their computational cost: repeated fitness evaluations over a population can be expensive, but the population-based approach enables parallel exploration and robustness to multimodality and noise [31, 21].

2.5 Advantages of Genetic Algorithms

Because a GA manipulates a population of candidate solutions using selection, recombination, and mutation, it brings several practical advantages that follow directly from that population-based, variation-driven design.

First, genetic algorithms provide an effective global search mechanism: by exploring many points in the search space simultaneously and combining information from multiple parents, GAs can escape local optima and discover diverse basins of attraction in multimodal landscapes [21]. Recombination enables the mixing of useful building blocks from different individuals, while mutation injects novel variation that can lead the search into previously unexplored regions.

Second, the population-based nature of GAs makes them naturally parallelizable. Fitness evaluations for different individuals are independent and can be distributed across processors or machines, which mitigates the computational cost of evaluating large populations and enables efficient use of modern parallel hardware [13].

Third, GAs are flexible in representation and operator design. The encoding (genotype) can be chosen to suit combinatorial, continuous, or structured search spaces, and operators can be tailored to preserve problem-specific constraints or exploit domain knowledge. This representational flexibility means GAs can be applied to a wide range of problem types where more specialized optimizers would require substantial reworking [38].

Fourth, because GAs do not rely on gradient information, they work well with discontinuous, noisy, or non-differentiable objective functions. This makes them a good choice when derivative-based methods are inapplicable or unreliable [31].

Finally, GAs tend to be robust in the presence of noise and uncertainty: population diversity and stochastic variation help prevent premature convergence to spurious solutions when fitness evaluations are noisy or imprecise [24]. Taken together, these advantages explain why genetic algorithms are widely used as general-purpose optimization tools, while also highlighting that their suitability depends on the problem structure and available computational resources.

2.6 Disadvantages of Genetic Algorithms

Despite their strengths, genetic algorithms also have practical limitations that follow from the same design features highlighted in the previous section. Most notably, the reliance on populations and repeated fitness evaluations makes GAs computationally expensive for problems where a single fitness evaluation is costly. Running a large population over many generations can require substantial CPU time or wall-clock time unless evaluations are parallelized or otherwise accelerated [31].

Another important drawback is the sensitivity to parameter settings. GAs expose many tunable parameters—population size, crossover and mutation rates, selection pressure, replacement strategy, and termination criteria—and the choice of these parameters strongly influences performance. Finding a good parameter configuration often requires experimentation, automated tuning, or domain expertise; poor settings can lead to inefficient search or failure to converge to satisfactory solutions [13].

In addition, there is no formal guarantee that a GA will find the global optimum in finite time. Like most heuristic search methods, GAs are stochastic and provide probabilistic rather than deterministic assurances; they are best viewed as powerful search heuristics rather than exact optimizers. This limitation is especially relevant when optimality certificates are required by the application or when the search space has pathological features that mislead population-based exploration [21].

A closely related issue is premature convergence: the population can lose diversity and become dominated by similar individuals, which reduces the algorithm’s ability to explore new regions of the search space. Premature convergence is often caused by excessive selection pressure, overly disruptive recombination, or too-small populations, and it can be mitigated through strategies such as maintaining diversity (niching, crowding), adaptive parameter control, hybridization with local search, or using island models that preserve separate subpopulations [13, 24].

Recognizing these disadvantages clarifies the trade-offs discussed in Section 2.5: the same mechanisms that give GAs their robustness and flexibility also create costs that must be managed through careful algorithm design, parameter tuning, and computational resources. For many practical problems the benefits outweigh the costs, but evaluating that balance is an essential step when choosing whether to apply a genetic algorithm to a given task.

2.7 When to Use Genetic Algorithms

Choosing to use a genetic algorithm depends on an assessment of the problem structure, the available computational resources, and the goals of the search. GAs are most attractive when the search space is large, complex, or poorly understood: their population-based exploration and representational flexibility allow them to discover solutions where derivative information is unavailable or conventional optimizers struggle.

When little is known about the problem structure or when objective functions are discontinuous, noisy, or multimodal, GAs provide a practical alternative to gradient-based or problem-specific methods. The absence of a requirement for differentiability and the ability to operate on combinatorial and structured encodings make GAs useful in engineering design, scheduling, symbolic regression, and similar domains where classical optimizers are inapplicable [31, 38].

GAs are also a natural choice when multiple, often conflicting objectives must be explored simultaneously. Multi-objective variants produce diverse Pareto-approximate solution sets, enabling decision makers to inspect trade-offs rather than forcing a single scalarized objective [12, 13].

However, the advantages listed in Section 2.5 must be weighed against the disadvantages discussed in Section 2.6. If fitness evaluations are extremely costly and parallel resources are unavailable, the computational burden of maintaining and evolving a population may outweigh the benefits. Similarly, if strict optimality guarantees are required, a GA's heuristic, stochastic nature may be inappropriate. In such cases, hybrid approaches (combining GAs with local search or surrogate models), careful parameter tuning, or the use of specialized optimizers may offer better trade-offs.

In practice, a useful decision rule is to prefer genetic algorithms when robustness, flexibility, and the ability to handle complex or poorly behaved search spaces are more important than raw efficiency or formal optimality guarantees. Where these conditions hold, applying the variations and mitigation strategies described earlier (parallel evaluation, island models, hybrids, and adaptive control) often yields practical, high-quality solutions.

2.8 Chapter Summary

This chapter introduced genetic algorithms as optimization tools inspired by natural evolution. We covered the basic components, terminology, and a simple example. The key insight is that GAs use population-based search with selection, crossover, and mutation to evolve solutions toward optimality.

2.9 Key Concepts

- Biological inspiration and evolution metaphor
- Basic GA structure and components
- Representation, fitness, selection, crossover, mutation
- Advantages and limitations of GAs
- When to apply genetic algorithms

Chapter 3

GA Cycle and Holland Schema Theory

3.1 The Genetic Algorithm Cycle

The genetic algorithm follows a cyclic process that mimics natural evolution. Understanding this cycle is crucial for implementing and analyzing GA performance.

3.1.1 Detailed GA Cycle

The genetic algorithm follows a cyclic process that mimics natural evolution and is designed to iteratively improve a population of candidate solutions. This process begins with a population of potential solutions and repeatedly applies evaluation and variation operators to move toward higher-quality solutions. The overall cycle is intentionally modular: initialization sets up the search, evaluation measures current quality, selection favors promising solutions, crossover and mutation introduce new genetic combinations, and replacement forms the next generation. These steps together form a loop that continues until a termination condition is met.

The flow of operations is illustrated by the diagram below, which captures the typical order of actions in a classical GA: initialize, evaluate, check termination, select parents, apply crossover and mutation, perform replacement, and return to evaluation. Each step has many possible implementations and parameters that influence exploration and exploitation; for example, population size, selection pressure, crossover type, and mutation rate all affect how the search navigates the solution space. Although the diagram shows a straightforward sequence, practical algorithms often include enhancements such as elitism, adaptive rates, or steady-state updates that change details of how offspring and parents are combined.

The figure that follows summarizes this canonical cycle and highlights where control decisions occur (for instance, whether the termination condition has been reached). Keep in mind that the figure is a conceptual guide: different problem domains and representations may require tailored operators or additional bookkeeping (such as maintaining constraints or auxiliary data). Nevertheless, the high-level structure depicted is useful as a template when designing or analyzing GA behavior.

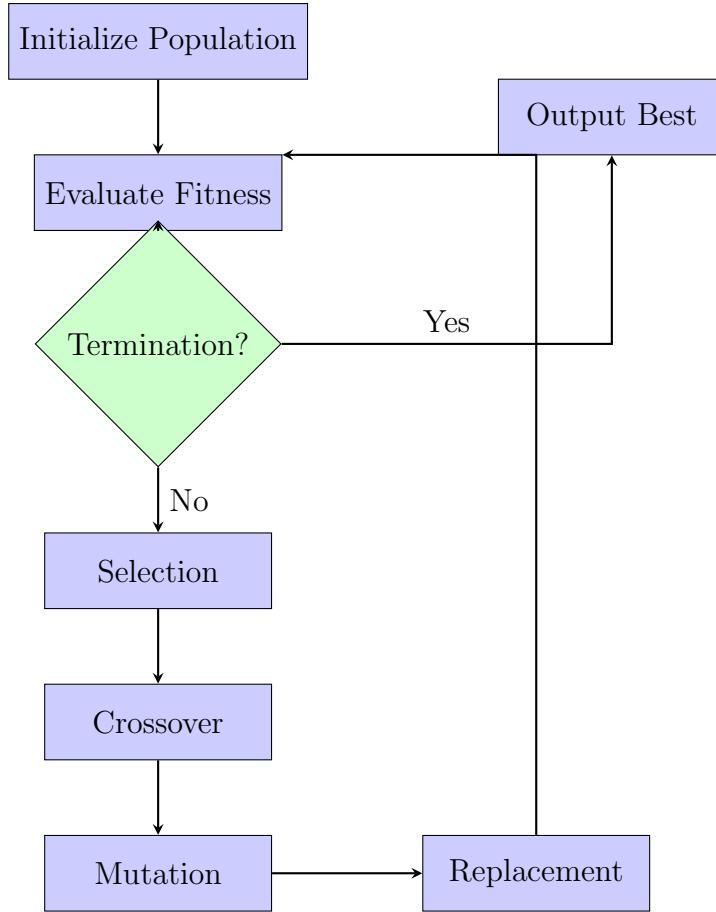


Figure 3.1: Genetic Algorithm Cycle

Initialization is the first practical step of the GA and determines the starting points for the search. An initial population of size N can be generated randomly to provide broad coverage of the search space, or it can be seeded using problem-specific heuristics to give the search a head start near promising regions. Ensuring diversity in the initial population is important because it reduces the risk of premature convergence and increases the chances that useful building blocks are present at the start. In addition to the candidate solutions themselves, initialization commonly includes setting any algorithm-level counters or parameters, such as the generation index $t = 0$, and recording any state needed for adaptive operators.

Evaluation quantifies how well each individual solves the problem at hand and converts raw solutions into fitness values used by the GA. This step typically computes an objective or fitness function for every member of the population, and may also collect summary statistics such as the population mean, variance, and the best and worst fitnesses. Those statistics are useful for monitoring progress, diagnosing issues like stagnation, and driving adaptive mechanisms (for example, adjusting mutation rates if diversity drops). Because evaluation is often the most expensive part of a GA—especially when each fitness computation involves simulation or complex calculations—practitioners pay close attention to efficient evaluation and to reusing computations where possible.

Termination is a control decision checked after evaluation to decide whether the algorithm should stop or continue. Common stopping conditions include reaching a pre-set maximum number of generations, achieving a fitness threshold that is satisfactory for the application, observing population convergence or very low diversity, detecting no improve-

ment for a fixed number of generations, or exhausting a budget of function evaluations. Choosing a termination criterion is a trade-off between computational cost and solution quality: stopping too early risks missing better solutions, while running too long wastes resources with diminishing returns. In practice it is common to combine several conditions (for example, stop when either the target fitness is reached or the generation limit is exceeded).

Selection chooses which individuals will contribute genetic material to the next generation by biasing reproduction toward fitter solutions while attempting to preserve useful variation. Selection methods vary—tournament selection, roulette-wheel (fitness-proportionate) selection, rank selection, and stochastic universal sampling are common examples—but they all aim to increase the proportion of above-average individuals over time. Well-designed selection balances exploitation (amplifying good solutions) with exploration (keeping diversity) so that the search does not lock prematurely onto suboptimal regions. Additional mechanisms such as fitness sharing or diversity preservation can be incorporated to maintain a healthy variety of candidates.

Crossover (recombination) is the operator that combines genetic material from selected parents to form new offspring, exchanging pieces of parent chromosomes to create novel solutions. The specific crossover operator (one-point, two-point, uniform, PMX, cycle crossover, etc.) and the crossover probability p_c control how often and how aggressively information is mixed. Crossover facilitates the combination of good building blocks discovered in different individuals, enabling the GA to assemble higher-quality solutions from simpler parts. However, crossover can also disrupt beneficial structures, so its design and application rate are tuned to the representation and problem characteristics.

Mutation introduces small, random changes to offspring to maintain genetic diversity and to allow the algorithm to explore regions of the search space that recombination alone might not reach. Mutation typically operates with a low per-bit probability p_m so that it makes conservative changes, preventing wholesale destruction of promising solutions while still enabling occasional novel innovations. Properly calibrated mutation helps the GA escape local optima, complements crossover by exploring orthogonal directions, and supports long-term adaptability of the population.

Replacement forms the population that will be evaluated in the next generation by deciding how parents and offspring are combined. Strategies range from generational replacement—where the entire population is replaced by the offspring—to steady-state or elitist schemes that retain some parents or the best individuals across generations. Replacement policy affects convergence speed, genetic diversity, and the risk of losing high-quality solutions; for example, elitism guarantees that the best found solution is never lost, while other policies may emphasize turnover and exploration. After replacement the generation counter is incremented ($t = t + 1$) and the cycle returns to evaluation for the next iteration.

3.1.2 What is a Schema?

A schema is a formal template defined over the alphabet $\{0, 1, *\}$. For a fixed string length l a schema

$$H \in \{0, 1, *\}^l$$

specifies required values at some loci and leaves other loci unspecified (the “don’t care” positions). Let $\Sigma = \{0, 1\}$ and denote by Σ^l the set of all length- l binary strings. We say a concrete string $s \in \Sigma^l$ matches schema H (write $s \in [H]$) exactly when every defined

position of H agrees with s :

$$s \in [H] \Leftrightarrow \forall i \in \{1, \dots, l\}, H_i \neq * \Rightarrow s_i = H_i. \quad (3.1)$$

The set $[H] = \{s \in \Sigma^l : s \text{ matches } H\}$ is the equivalence class of concrete genomes represented by H . If $k(H)$ denotes the number of don't-care symbols in H (so $k(H) = |\{i : H_i = *\}|$), then the cardinality of this class is

$$|[H]| = 2^{k(H)}. \quad (3.2)$$

Two other commonly used quantities are the order and the defining length. The order $o(H)$ equals the number of fixed positions (non-* symbols), so $o(H) = l - k(H)$. If i_{\min} and i_{\max} are the indices of the first and last fixed positions in H , the defining length is

$$\delta(H) = i_{\max} - i_{\min}. \quad (3.3)$$

Example (preserved): for $H = 1 * 0 * 1$ with $l = 5$ we have fixed positions at indices 1, 3, 5, $k(H) = 2$, $o(H) = 3$, and

$$|[H]| = 2^2 = 4, \quad \delta(H) = 5 - 1 = 4,$$

with the matching strings $\{10001, 10011, 11001, 11011\}$.

3.1.3 Schema Properties

Order of a Schema

The order $o(H)$ is the number of fixed positions (non-* symbols):

$$o(H) = \text{number of defined bits in } H \quad (3.4)$$

For $H = 1 * 0 * 1$: $o(H) = 3$

Defining Length

The defining length $\delta(H)$ of a schema measures the span between the earliest and latest fixed (non-*) positions in the pattern. Intuitively, it captures how spread out the important bits of the schema are along the chromosome and therefore how exposed the schema is to recombination events. Formally, it is computed as the index difference between the last and first fixed positions:

$$\delta(H) = \text{last fixed position} - \text{first fixed position} \quad (3.5)$$

A small defining length means the schema's fixed bits are clustered closely together. Such clustering tends to make a schema more robust under common crossover operators because fewer crossover cut points lie between the fixed positions; consequently the schema is less likely to be split apart during recombination. By contrast, a schema with a large defining length distributes its fixed bits across a longer region of the chromosome and thus has a higher probability of being disrupted by crossover, even if each fixed bit individually is unlikely to mutate.

The practical significance of defining length is closely tied to the building-block view of genetic algorithms: short, tightly-linked blocks of genes that confer above-average fitness

are easier for a GA to preserve and recombine successfully. When designing encodings or choosing crossover operators, minimizing unnecessary spread of interdependent genes can reduce the effective defining lengths of useful schemas and improve the likelihood that beneficial combinations survive and propagate.

For $H = 1 * 0 * 1$: $\delta(H) = 5 - 1 = 4$

Examples from Buku Ajar:

- $S_1 = (** * 001 * 110)$: $\delta(S_1) = 10 - 4 = 6$
- $S_2 = (** * * 00 * * 0*)$: $\delta(S_2) = 9 - 5 = 4$
- $S_3 = (11101 * * 001)$: $\delta(S_3) = 10 - 1 = 9$

3.1.4 Schema Theorem (Fundamental Theorem)

The schema theorem describes how the expected number of strings matching a schema changes from generation to generation.

Selection Effect

If $m(H, t)$ is the number of strings matching schema H at generation t , and $f(H)$ is the average fitness of strings matching H , then:

$$E[m(H, t + 1)] \geq m(H, t) \cdot \frac{f(H)}{\bar{f}} \quad (3.6)$$

where \bar{f} is the average fitness of the population.

This means schemas with above-average fitness will increase in representation.

Crossover Effect

Crossover can disrupt a schema if the crossover point falls between the defining positions. The probability of schema survival is:

$$P_s = 1 - p_c \cdot \frac{\delta(H)}{l - 1} \quad (3.7)$$

where:

- p_c is the crossover probability
- l is the string length

Mutation Effect

The probability that a schema survives mutation is:

$$P_m = (1 - p_m)^{\delta(H)} \quad (3.8)$$

where p_m is the mutation probability per bit.

Combined Schema Theorem

Combining all effects:

$$E[m(H, t + 1)] \geq m(H, t) \cdot \frac{f(H)}{\bar{f}} \cdot \left(1 - p_c \cdot \frac{\delta(H)}{l-1}\right) \cdot (1 - p_m)^{o(H)} \quad (3.9)$$

3.1.5 Building Block Hypothesis

The building-block hypothesis can be stated precisely in terms of Holland's schema formalism: a building block is a schema H with small defining length $\delta(H)$, low order $o(H)$, and above-average fitness $f(H) > \bar{f}$. The schema theorem gives a quantitative criterion for such a schema to grow in expectation from generation t to $t + 1$:

$$E[m(H, t + 1)] \geq m(H, t) \cdot \frac{f(H)}{\bar{f}} \cdot \left(1 - p_c \frac{\delta(H)}{l-1}\right) \cdot (1 - p_m)^{o(H)}. \quad (3.10)$$

Consequently, a necessary (but not sufficient) condition for expected growth of H is that the multiplicative factor on the right-hand side exceeds one. Rearranging gives the informal threshold condition

$$\frac{f(H)}{\bar{f}} > \frac{1}{\left(1 - p_c \frac{\delta(H)}{l-1}\right) (1 - p_m)^{o(H)}}. \quad (3.11)$$

This inequality makes explicit the trade-offs: higher relative fitness, smaller defining length, lower order, smaller crossover probability (or tightly clustered loci), and lower mutation rate all improve the chance that a schema will multiply.

Two additional quantitative constraints are important for the hypothesis to be operational. First, the population must initially sample the schema with sufficient multiplicity: for a random initialisation the expected initial count is $E[m(H, 0)] = n 2^{-o(H)}$, so n must be large enough that $m(H, 0)$ is not effectively zero for all useful building blocks. Second, crossover must be able to assemble complementary building blocks: if two short schemas H_1 and H_2 occur on disjoint loci and survive disruption, recombination can produce individuals that contain both, enabling a constructive “assembly” mechanism that builds higher-order structure from lower-order pieces.

Finally, the hypothesis is not an unconditional theorem; it describes a plausible mechanism rather than a universal guarantee. Finite-population sampling noise, strong epistasis (tight interactions among distant loci), and deliberately deceptive fitness functions can violate the premises above. In practice the building-block perspective remains valuable because it highlights representation and operator choices: encodings and crossover operators that keep interdependent genes close (reducing δ), maintain adequate sampling (sufficient n), and balance p_c and p_m to favour preservation-and-recombination will be more likely to exploit short, fit schemas effectively.

3.1.6 Role of Schema Order in Genetic Algorithms

Pattern Specificity Level

The order of a schema indicates the level of specificity of the pattern represented in a chromosome. The higher the order, the more specific the pattern described. For example,

a low-order schema like $1 * * * *$ still represents many possible chromosomes because only one position is fixed, while a high-order schema like 101011 is very specific and only matches one particular chromosome. Thus, order plays a role in determining how broad a schema's representation is within the population.

Survival Probability in Evolution

The order of a schema greatly influences the schema's chance of surviving through the evolution process. In genetic algorithms, two main operators that often cause changes are crossover and mutation.

Low-order schemas are relatively safer against these changes because they have few fixed positions. For example, the schema $1 * * * *$ only locks one bit at the beginning. If mutation occurs at another position or crossover cuts through the middle, the chance of schema disruption is very small. In other words, the fewer fixed bits that must be maintained, the greater the likelihood that the schema will survive and be inherited to the next generation.

Conversely, high-order schemas like 101011 have many fixed bits that must be exactly the same to remain valid. In such conditions, even one small mutation at one of the fixed positions can destroy the entire schema. Similarly, in the crossover process, the probability of being cut between fixed positions becomes larger. As a result, high-order schemas often quickly disappear from the population because they struggle to survive the combination of genetic variations that occur.

Relationship with Selection

Although they appear simple, low-order schemas actually play a very important role in genetic algorithms. The simplicity of this structure allows schemas to be more resistant to damage from crossover and mutation, so these patterns survive more often and are inherited to the next generation. If a simple schema contains patterns relevant to the optimal solution—for example, certain bit combinations that increase fitness value—then that schema will appear repeatedly on various chromosomes in the population.

This phenomenon aligns with the building block hypothesis put forward by Holland. Genetic algorithms do not directly search for complex solutions as a whole, but work by maintaining and arranging simple pattern blocks that have a positive contribution to fitness. These blocks are then combined through selection, crossover, and mutation, thus forming more complex genetic structures that approach the optimal solution.

The selection process plays a central role here. Individuals who have schemas with high contribution to fitness will be selected more often to reproduce. Thus, beneficial simple schemas can spread widely in the population. Over time, combinations of building blocks from various simple schemas produce solutions that are not only more complex, but also more efficient in solving problems.

3.1.7 Relationship Between Holland Schema and Genome

The number of genome sequences that can be represented by a schema depends on the number of don't care symbols (*). Each don't care symbol can have a value of 0 or 1, so a schema with k don't care symbols will produce 2^k possible genome sequences.

Examples:

- S_4 has 1 don't care, so there are $2^1 = 2$ possible genome sequences: $110010, 110110$

- S_5 has 2 don't care, so there are $2^2 = 4$ possible genome sequences: 1110000, 1110100, 0110000, 0110100
- S_6 has 3 don't care, so there are $2^3 = 8$ possible genome sequences: 101100111, 101100110, 101100010, 101100011, 100100111, 100100110, 100100011, 100100111

3.1.8 Schema Functions in Genetic Algorithms

Schemas perform several distinct functions in the analysis and practice of genetic algorithms. Formally, a schema H defines an equivalence class $[H] \subseteq \Sigma^l$ of concrete genomes; the population-level quantity $m(H, t)$ (the number of individuals matching H at generation t) summarizes how the GA samples that class. By operating on these counts and their expectations rather than on individual genotypes, schema-based analysis compresses a population's combinatorial structure into tractable quantities that can be used to prove bounds (for example, the schema theorem) and to reason about the aggregate effects of selection, crossover, and mutation.

Second, schemas provide an explanatory bridge between micro-level operators and macro-level behaviour. The combined schema theorem expresses how selection amplifies above-average patterns while crossover and mutation attenuate them; this lets us write expected-update formulas for $m(H, t)$ and identify the parameter regimes in which particular schemas are likely to increase. That analytic viewpoint underpins the notion of implicit parallelism: a single population of size n simultaneously samples an exponentially large set of schemas, and the GA's operators act on many of these equivalence classes in parallel through their effect on the underlying individuals.

Third, schema thinking informs algorithm design and diagnostics. Monitoring $m(H, t)$ for candidate schemas helps detect whether useful patterns are being discovered and preserved; choosing encodings and crossover operators that reduce defining lengths for interdependent loci increases the survivability of important schemas; and linkage-learning methods or estimation-of-distribution algorithms can be seen as modern generalisations that explicitly model and exploit schema-like dependencies rather than relying on blind recombination. In practice, schema-level measures also motivate choices for population size and mutation rate because they make sampling and disruption risks explicit.

Finally, it is important to recognise the limitations of schema functions as an analysis tool. Schema counts discard detailed positional interactions and so can obscure strong epistatic effects where the contribution of a locus depends nonlinearly on distant loci; they are most informative for low-order, short-defining-length patterns and for binary encodings. Consequently, while schemas are powerful for explaining certain GA behaviours and for guiding representation/operator decisions, they are not a universal modelling device—empirical evaluation and, where appropriate, richer dependency models are required to handle deception, dense epistasis, or non-binary representations.

3.2 Implicit Parallelism

GAs operate on populations of concrete genomes, but each concrete genome simultaneously instantiates an exponential family of schemata. Concretely, for a binary string of length l there are 3^l possible schemata in total (each position may be 0, 1, or “don't care”). A single length- l string matches exactly 2^l distinct schemata because each locus

may either be left fixed or replaced by a don't-care symbol. Consequently, a population of size n provides direct instances of at most $n2^l$ schemata (counting multiplicity), and—ignoring overlaps between individuals—this number can be exponentially large in l . This combinatorial fact underpins the intuitive idea that a GA evaluates many schemata in parallel.

A more refined accounting groups schemata by their order (the number of fixed loci). The number of schemata of order r equals

$$\binom{l}{r} 2^r, \quad (3.12)$$

since one chooses which r loci are fixed and then assigns a bit (0 or 1) to each fixed locus. The number of schemata whose order does not exceed k is therefore

$$S_k = \sum_{r=0}^k \binom{l}{r} 2^r, \quad (3.13)$$

which, for fixed small k , grows polynomially in l (degree k) rather than exponentially. This observation is central: the schema theorem and the building-block hypothesis emphasise short, low-order schemata (small $o(H)$ and small defining length) because those are both numerous enough to be sampled and robust enough to survive recombination and mutation with reasonable probability.

Holland's celebrated phrase “implicit parallelism” summarises the heuristic that a population of modest size can collectively evaluate and process a very large number of low-order, short-defining-length schemata in each generation. In his original exposition he offered the rule-of-thumb that a population of n strings can effectively process on the order of n^3 schemata; this statement should be read as a heuristic, not a strict combinatorial identity. The $O(n^3)$ figure arises from assumptions about the typical orders and defining lengths of schemata that materially influence fitness, together with plausible estimates of how many distinct short schemata a population samples and how selection amplifies above-average instances. Different choices of n , l , representation, and operator settings change the constant factors and the practical reach of this parallelism.

The practical implication is twofold. First, by working with a population rather than a single search trajectory, a GA can explore and propagate many candidate building blocks simultaneously, enabling constructive recombination of useful short patterns. Second, this implicit coverage is selective: the algorithm gives effective processing power to schemata that are both sufficiently sampled (appear often enough in the population) and sufficiently resilient (have small defining length and low order so that crossover and mutation do not destroy them). As a result, implicit parallelism is not a magic bullet that inspects every possible schema equally; instead, it directs computational effort toward a large, structured subset of schemata that are most relevant under the chosen encoding and operators.

Finally, it is important to recognise limitations. Finite-population sampling error, strong epistasis (where fitness depends on complex interactions among distant loci), deceptive fitness functions, and inappropriate operator settings can all undermine the effective parallelism a GA achieves in practice. Therefore, exploiting implicit parallelism requires careful encoding, a sensible population size, and operator choices that favour the preservation and recombination of short, meaningful building blocks.

3.3 Deception and Schema Theory

3.3.1 Deceptive Problems

Deception occurs when selection, operating on short-term fitness cues, systematically drives the population away from genotypes that lead to the global optimum. Informally, a fitness function is deceptive with respect to a set of low-order schemata when the schemata that appear to be “best” locally (i.e. have above-average fitness among their instances) are those whose recombination does not assemble the global solution but rather promotes genotypes that are hard to improve upon. In other words, selection rewards building blocks that guide the search toward local optima instead of toward the global optimum.

This phenomenon can be expressed in schema terms. Consider two disjoint schema classes H_1 and H_2 defined on disjoint sets of loci. If the most fit instances of H_1 and H_2 tend to produce offspring whose combined fitness is lower than alternative combinations (or if combining them is unlikely under the chosen operators), then selection acting on H_1 and H_2 may increase their frequency even though their joint presence is unfavourable for reaching the global optimum. Such a mismatch between short-term schema fitness and long-term constructive value is the essence of deception.

A canonical benchmark that illustrates deception is the concatenated deceptive-trap problem. The genome is partitioned into m disjoint blocks of size k . Each block contributes a block-wise fitness that is maximal for a specific configuration (the block’s global optimum) but otherwise assigns higher fitness to a locally-attractive configuration that is not on the path to the block optimum. When blocks are simple and independent, recombination can assemble block optima; when blocks are deceptive, selection preferentially amplifies wrong local configurations and recombination alone may not rescue the global solution.

While the term “deceptive” has a precise intent, it is not an absolute property of a fitness function alone but of the combination of function, representation, population size, and operators. A partition that appears deceptive for one recombination operator may not be deceptive for another that respects linkage; likewise, increasing population size or changing selection pressure can alter whether a particular problem instance behaves deceptively in practice.

3.3.2 Why Deception Matters for Schema Theory

Schema theory predicts that low-order, short-defining-length schemata with above-average fitness will increase in expectation under selection and survive recombination/mutation with non-negligible probability. Deception subverts this reasoning by making the locally above-average schemata lead away from genotypes that contain the globally optimal combination of building blocks. Thus, although the schema theorem’s algebraic statement remains valid as an expectation, its constructive interpretation (that selection will assemble good building blocks into better solutions) can fail when the sampled short schemata are misleading.

Two concrete consequences follow:

- Sampling risk: finite populations may not sample the rare, globally-useful schemata often enough for recombination to assemble them before deceptive schemata dominate.

- Misleading gradients: selection amplifies schemata that locally increase fitness even if these schemata decrease the probability of reaching the global optimum when combined.

3.3.3 Detecting and Measuring Deception

Practically, deception is assessed by analysing how local improvements correlate with global progress. Common diagnostics include:

- Fitness-distance correlation (FDC): the correlation between fitness and distance to a known global optimum. Strong negative correlation suggests easier search; weak or positive correlation may signal deception.
- Empirical block analysis: for decomposable problems (e.g. concatenated traps), studying the block-wise fitness landscape (unitation plots) reveals whether local optima attract search within blocks.
- Performance sensitivity: measuring success probability as a function of population size and operator settings can indicate whether modest changes remove or expose deceptive behaviour.

3.3.4 Strategies to Overcome Deception

Because deception arises from a mismatch between representation, operators, and the problem's modular structure, countermeasures generally fall into three categories: improve sampling, preserve or encourage useful diversity, and increase the algorithm's ability to respect or learn linkage.

- **Increase effective sampling:** Larger populations and conservative selection pressure reduce the chance that deceptive schemata quickly fix, giving recombination and mutation more opportunity to assemble globally-useful combinations.
- **Diversity-preserving methods:** Niching (fitness sharing, crowding), island models, and restricted tournament selection retain multiple competing alleles or subpopulations so that alternative building-block combinations are not lost prematurely.
- **Linkage-aware recombination:** Using crossover operators that respect known linkage, or designing encodings that keep interdependent loci close, reduces the probability that crossover breaks important combinations and thereby reduces apparent deception.
- **Linkage learning and estimation methods:** Algorithms that learn dependencies—messy GAs, linkage-tree GAs, hierarchical Bayesian optimization algorithm (hBOA), and dependency-structure modelling—explicitly detect and preserve interacting genes instead of relying on blind recombination.
- **Hybridisation and local search:** Combining GAs with problem-specific local search (memetic algorithms) or constructive heuristics helps repair or complete partial solutions that pure recombination cannot assemble.
- **Adaptive operators and parameter control:** Dynamically tuning mutation rates, crossover rates, and selection pressure in response to measured diversity or progress can mitigate regimes in which deception is most damaging.

3.3.5 Limitations and Practical Advice

No single remedy eliminates deception in every domain — the phenomenon is problem-dependent. The No-Free-Lunch theorems imply that algorithmic choices trade off performance across problem classes, so the right countermeasure depends on prior knowledge about structure (if available) or on adaptive techniques that infer structure during the run. In practice, good engineering steps are:

- Prefer encodings and operators that keep suspected interacting loci linked.
- Use modestly large populations and conservative selection until building blocks are reliably sampled.
- Monitor diversity and performance metrics (e.g. FDC, genotypic/phenotypic variance) and apply linkage learning or niching when signs of premature convergence appear.

Understanding deception and planning for it when designing encodings and operators is essential for applying GAs to problems with strong epistasis or deliberately deceptive structure.

3.4 Practical Implications

3.4.1 Encoding and Representation

Representation is the single most important design choice when attempting to exploit schema-like behaviour. The following principles are recommended:

- **Reduce unnecessary epistasis:** Wherever possible, choose encodings that make the contribution of a small set of loci to fitness approximately additive. Lower epistasis reduces the chance that short schemata are misleading and increases the utility of recombination.
- **Preserve linkage of interacting loci:** Place genes that interact closely in the problem domain near each other in the representation, or use positional encodings that respect natural modularity. Doing so reduces defining lengths $\delta(H)$ for important schemata and raises their survivability under common crossover operators.
- **Prefer modular encodings:** When a problem decomposes into largely independent subproblems, design encodings (or problem decompositions) that reflect those modules so that recombination can assemble global solutions from block-wise building blocks.
- **Mind the genotype–phenotype map:** Nonlinear mappings from genotype to phenotype (e.g. big-radix encodings, Gray codes, or indirect encodings) change the effective schema structure and must be evaluated empirically for how they affect building-block preservation.

3.4.2 Operator and Parameter Choices

Schema-level reasoning suggests particular trade-offs when choosing operators and their parameters:

- **Population size (n):** Larger populations reduce sampling noise and increase the probability that useful low-order schemata are present in sufficient multiplicity. Use population-sizing rules or experiments to ensure reliable sampling for the expected order of important schemata.
- **Selection pressure:** Strong selection accelerates exploitation but increases the risk of premature loss of diversity (and of useful schemata). Moderate selection or tournament sizes and techniques like fitness scaling can balance exploration and exploitation.
- **Crossover rate and type (p_c):** High crossover frequency promotes recombination of building blocks but also increases disruption of long defining-length schemata. Choose crossover operators that respect problem linkage (e.g. blockwise or problem-specific crossover) when possible.
- **Mutation rate (p_m):** Keep mutation low enough to avoid destroying short, beneficial schemata but high enough to introduce necessary variation and help escape local optima. Adaptive mutation schedules can be effective when problem structure is unknown.
- **Elitism and replacement:** Elitism preserves best-so-far solutions and thus protects discovered building blocks; however, excessive elitism can reduce diversity. Use small elitist fractions to balance stability and exploration.

3.4.3 Practical Monitoring and Diagnostics

To apply schema-informed design in practice, monitor population statistics regularly:

- Track genotypic diversity (e.g. per-locus allele frequencies) and phenotypic variance to detect premature convergence.
- Compute simple schema counts or sample candidate schemata to verify whether expected building blocks are being found and preserved.
- Measure progress metrics (best, median, and mean fitness) alongside diversity indicators; slow improvement with collapsed diversity often signals that schema recombination is failing.

3.5 Limitations of Schema Theory

Schema theory provides a valuable conceptual and analytic framework, but its assumptions and scope impose important limitations that practitioners must recognise.

3.5.1 Expectation vs finite-population dynamics

The core algebraic statements of schema theory are statements about expectations. In finite populations, stochastic sampling noise, genetic drift, and sampling error can cause realised dynamics to deviate substantially from expectation. Consequently, schema-theoretic predictions should be interpreted as tendencies rather than deterministic outcomes.

3.5.2 Operator dependence and representation sensitivity

Results derived from schema analysis depend on the choice of representation and genetic operators. The survival probabilities and constructive recombination arguments assume particular crossover and mutation models; different operators (or nonstandard genotype–phenotype maps) change these probabilities and may invalidate naive conclusions.

3.5.3 Limited handling of epistasis and complex interactions

Schema theory is most informative for low-order, short-defining-length patterns. When fitness arises from high-order interactions (strong epistasis) or from complex, distributed dependencies among loci, schema counts obscure the relevant structure and offer limited predictive power.

3.5.4 Restriction to simple alphabets and fixed-length encodings

Classical schema analyses assume binary alphabets and fixed-length strings. Extensions to richer alphabets, variable-length genomes, or indirect encodings require careful reformulation; naive application of binary-based intuition can be misleading.

3.5.5 Lack of prescriptive specificity

While schema theory explains why short, fit building blocks can be useful, it does not provide a general, prescriptive algorithm for discovering the best representation or operators for an arbitrary problem. Modern methods—linkage learning, estimation-of-distribution algorithms (EDAs), and probabilistic model-building approaches—explicitly attempt to learn and exploit problem structure that schema counts alone cannot reveal.

3.5.6 Practical takeaway

Schema theory remains a useful lens for understanding GA behaviour and for guiding design choices (representation, linkage, population sizing, and operator tuning). However, effective application requires empirical validation, diagnostic monitoring, and, when necessary, methods that explicitly learn and preserve dependencies rather than relying solely on blind recombination.

3.5.7 Deeper Limitations and Practical Consequences

Beyond the conceptual caveats above, there are several deeper limitations that affect both theoretical analysis and practical algorithm design.

Stochastic fluctuations and reliability

Because the schema theorem is an expectation, single-run trajectories can differ widely from the expected behaviour. This is not merely a small-sample issue: genetic drift, sampling variance, and the discrete nature of selection events can produce systematic divergences (for example, loss of a useful low-frequency schema by chance). Practitioners should therefore treat schema-based predictions as probabilistic statements and quantify reliability via multiple independent runs, confidence intervals on observed schema counts, or resampling (bootstrap) methods.

Difficulty of measuring schemas in practice

Explicitly tracking large numbers of schemata is computationally expensive and of limited usefulness unless the tracked schemata are chosen carefully. The number of possible schemata grows combinatorially with l , and naive enumeration is infeasible for realistic genome sizes. Practical measurements therefore focus on low-order schemata, sampled candidate patterns, or aggregate statistics (allele frequencies, linkage disequilibrium measures). Any empirical claim about schema dynamics should specify the sampling protocol, statistical uncertainty, and potential biases introduced by the choice of monitored schemata.

Finite-population bounds and worst-case behaviour

While the schema theorem gives lower bounds on expected schema counts under simplified operator models, it does not provide tight finite-population guarantees or worst-case complexity results. There exist problem instances (deceptive or highly epistatic) where the runtime to discover the global optimum grows exponentially in problem size for many GA variants. Where possible, complement schema-theoretic intuition with finite-population analyses, convergence proofs (for restricted algorithm variants), or empirically-derived scaling laws for the specific problem class.

Operator modelling limitations

Common schema-theoretic derivations assume simple, memoryless operators (single- or two-point crossover, independent bit-flip mutation, generational replacement) and ignore implementation details such as selection with replacement, stochastic universal sampling, or repair procedures for constraint handling. These implementation choices alter survival probabilities and the effective recombination patterns; hence, conclusions drawn under idealised operator models must be revalidated for production implementations.

Interpretation pitfalls

It is easy to overinterpret the schema theorem as a prescriptive justification for arbitrary GA design choices. For example, citing the schema theorem to justify an arbitrary low mutation rate or a particular crossover operator without reference to empirical evidence or

problem structure risks poor performance. Use schema reasoning as a guide for hypotheses to be tested, not as a substitute for empirical validation.

3.5.8 Connections to alternative analytic frameworks

Because of the limitations above, modern theoretical and practical work often complements schema thinking with other frameworks:

- Markov-chain and dynamical-systems analyses that characterise full-population dynamics under specific operators.
- Probabilistic model-building approaches (EDAs, hBOA) that explicitly learn and exploit dependency structure instead of relying on the passive effects of crossover.
- Linkage-detection and hierarchical decomposition methods that aim to discover interacting variable groups and preserve them explicitly during recombination.

These perspectives provide more actionable mechanisms for domains with strong epistasis or where blind recombination fails.

3.5.9 Recommended empirical protocol

When using schema-based reasoning to guide algorithm design, follow an empirical protocol that reduces the risk of incorrect conclusions:

- Run multiple independent trials and report variance, not just mean performance.
- Use controlled ablation studies to measure the effect of representation and operator choices on sampling and survival of candidate schemata.
- When feasible, visualise allele-frequency trajectories and block-wise unitation plots to diagnose where and when useful schemata are lost or preserved.
- Compare with model-building baselines (simple EDAs, linkage-aware GAs) to evaluate whether blind recombination suffices for the target problem.

These steps help translate schema-theoretic insights into reliable engineering decisions.

3.5.10 No Free Lunch Theorem

States that no algorithm is superior across all possible problems.

3.6 Chapter Summary

This chapter covered the genetic algorithm cycle and Holland's schema theory. The GA cycle provides a systematic approach to evolutionary search, while schema theory explains why GAs work by processing building blocks in parallel. Understanding these concepts is essential for effective GA design and application.

3.7 Key Concepts

- GA cycle phases and their purposes
- Schema representation and properties
- Schema theorem and its implications
- Building block hypothesis
- Implicit parallelism in GAs
- Deception and its challenges

Chapter 4

Genetic Algorithm Encoding

4.1 Introduction to Encoding

Encoding (also called representation) formalises how candidate solutions are described for a genetic algorithm (GA). Let G denote the discrete set of genotypes (the representation space) and P denote the set of phenotypes (the solution space). Encoding is the specification of a mapping

$$\phi : G \rightarrow P,$$

that assigns to each genotype a phenotype that can be evaluated by a fitness function. In practice G is usually a finite or countable combinatorial space (for example, bit-strings, vectors of integers, permutations, trees or real-valued vectors) and P is the domain of problem solutions (for example, real vectors, schedules, tours, or programs).

Two aspects of encoding must be distinguished: (i) the representational language used to construct genotypes (bits, integers, reals, nodes in a tree, etc.), and (ii) the genotype–phenotype mapping ϕ . The search performed by a GA operates in G , while fitness and problem constraints are defined on P ; therefore the properties of ϕ crucially determine how variation in genotype space translates to meaningful changes in solution quality.

Well-chosen encodings expose structure that the search procedures can exploit, reduce the incidence of infeasible solutions, and control representational redundancy and epistasis. Poor encodings can render local improvements invisible to variation, produce pathological fitness landscapes, or require expensive repair procedures. In later sections we discuss concrete encoding families (binary, gray, real-valued, permutation, tree) and the practical implications they have for representation design and algorithm performance.

4.2 Requirements for Good Encoding

When designing an encoding and its associated mapping $\phi : G \rightarrow P$, it is useful to state desiderata precisely. The following properties capture core representational requirements and trade-offs; they guide the selection or construction of encodings for a given problem.

4.2.1 Completeness

Completeness requires that the encoding be able to express every feasible phenotype of interest: formally, the image of ϕ should cover the feasible region $F \subseteq P$ of solutions, i.e. $\phi(G) \supseteq F$. If completeness fails then some valid solutions are unreachable by the GA, which introduces representational bias and can prevent the algorithm from finding optimal solutions that lie outside $\phi(G)$.

In practice completeness is balanced against representational compactness: a fully complete encoding may be large or inefficient, whereas a restricted encoding can greatly simplify search if it excludes uninteresting parts of P . Designers should explicitly state which subset of P must be reachable and ensure $\phi(G)$ contains it.

4.2.2 Soundness

Soundness (also called validity) stipulates that every genotype should map to a well-defined, constraint-satisfying phenotype: $\forall g \in G, \phi(g) \in P_{valid}$. Sound encodings avoid or minimise the production of infeasible solutions so that fitness evaluations are meaningful without costly repair. When strict soundness is impossible or impractical, designers may allow infeasible genotypes but must provide an efficient, well-defined decoding and repair strategy and ensure the search can still progress.

Soundness and completeness are orthogonal: an encoding can be sound but incomplete (every genotype valid, but not all phenotypes representable), or complete but unsound (all phenotypes representable but many genotypes invalid) depending on G and ϕ .

4.2.3 Non-redundancy

Non-redundancy means reducing (or eliminating) multiple distinct genotypes that map to the same phenotype. Formally, one prefers ϕ to be injective on the set of representationally relevant genotypes. Redundancy (many-to-one mapping) increases the effective search volume and can bias sampling: some phenotypes may be over-represented in G , making them more likely to be sampled even if they are not superior.

However, redundancy is sometimes deliberately introduced for robustness (e.g. neutral networks that allow neutral drift) or to simplify representation. When redundancy is present, it should be understood and controlled: quantify the degree of redundancy and consider its interaction with search dynamics and variation.

4.2.4 Locality

Locality formalises the intuition that small genotypic changes should produce small phenotypic changes. Let d_G and d_P be distance measures on G and P respectively (e.g. Hamming distance on bit-strings, Euclidean distance on real vectors). High locality means that $d_G(g_1, g_2)$ small $\Rightarrow d_P(\phi(g_1), \phi(g_2))$ small. Locality is important because common variation procedures make small changes in G ; if these do not correspond to small, correlated changes in P the search becomes effectively random and building-block recombination fails. Encoding choices such as Gray coding for integers or real-valued representations aim to improve locality.

Locality cannot always be achieved with other desirable properties; for example, injective, compact encodings with perfect locality may not exist for some combinatorial domains. Designers should therefore prioritise which properties matter most for the problem and for the procedures they plan to use.

4.2.5 Additional Practical Requirements

Beyond the four formal properties above, useful encodings should also satisfy several pragmatic constraints:

- **Operator Closure:** Variation procedures should, as much as possible, produce genotypes within a region of G that decodes to feasible or easily repaired phenotypes.
- **Computational Efficiency:** Decoding ϕ and any repair procedures should be computationally inexpensive relative to fitness evaluation.

- **Scalability:** The encoding should scale gracefully with problem size; representation length should not grow superlinearly without justification.
- **Low Epistasis:** The representation should aim to minimise destructive interactions between genes (epistasis) so that beneficial building blocks can be recombined reliably.
- **Interpretability and Prior Knowledge:** When available, incorporate problem-specific structure (symmetries, invariants, constraints) to simplify search and reduce unnecessary degrees of freedom.

Designing an encoding is therefore a matter of formal requirements, procedure compatibility, and empirical validation. Later sections examine common encoding families and trade-offs, and discuss practical choices that respect the desiderata above.

4.3 Binary Encoding

Binary encoding represents genotypes as fixed-length vectors over the binary alphabet: $G = \{0, 1\}^l$. A genotype $g = (b_{l-1}, \dots, b_0)$ is commonly interpreted as an unsigned integer

$$\text{bin}(g) = \sum_{i=0}^{l-1} b_i 2^i,$$

which is then mapped to a phenotype by an affine decoding when the phenotype is numeric. For a real-valued variable $x \in [x_{\min}, x_{\max}]$ the usual decoding is

$$x = x_{\min} + \frac{\text{bin}(g)}{2^l - 1} (x_{\max} - x_{\min}). \quad (4.1)$$

This formula makes the representational resolution explicit: the quantisation step is

$$\Delta = \frac{x_{\max} - x_{\min}}{2^l - 1},$$

so choosing l trades off precision against search dimensionality and behavioural properties of variation procedures.

Binary encodings are attractive because they are compact and simple to manipulate with bitwise operators. Schema theory and many early theoretical results were developed for binary representations, which aids theoretical reasoning about convergence and building-block propagation [25, 21].

However, binary encodings also introduce specific problems that must be addressed in practice:

- **Hamming cliffs and locality:** Adjacent numeric values can differ in many bits under standard binary positional encodings, breaking locality. Gray codes are a common remedy when preserving adjacency is important.
- **Precision versus length:** High precision requires long bit-strings, which increases the search space exponentially and can make positional recombination disruptive.
- **Epistasis:** Bit positions may interact non-linearly with respect to phenotype quality; correlated bits reduce the effectiveness of simple recombination.

Practical recommendations for binary encodings:

- Choose length l from the desired resolution Δ and range $[x_{\min}, x_{\max}]$ using $2^l - 1 \geq (x_{\max} - x_{\min})/\Delta$.
- If adjacency matters, consider Gray coding for integer variables and convert to binary only for procedures that work on bitstrings.
- Use procedure choices that respect gene boundaries for multi-variable concatenations (e.g. align recombination points to variable boundaries when appropriate).
- Tune per-bit modification rates as a starting heuristic; decrease when using local search or strong selective dynamics.

4.4 Overview of Encoding Types

Encodings can be organised according to the structure of G and the intended phenotype domain P . Below we summarise principal families and their canonical use-cases, with practical guidance for representation design and common pitfalls.

Taxonomy and mapping to problem classes

- **Binary (bit-strings):** $G = \{0, 1\}^l$. Good for combinatorial choices and when schema analysis is desired. Use Gray code or problem-specific bit ordering to improve locality for numeric phenotypes.
- **Integer (value) encodings:** Vectors of integers; natural for count and allocation problems. Use integer-aware variation procedures (random reset, creep) and discrete recombination methods.
- **Real-valued encodings:** Continuous vectors \mathbb{R}^n . Preferable for continuous optimisation; supports arithmetic combination and BLX- α style recombination, stochastic perturbations, and gradient-informed hybrids.
- **Permutation encodings:** Represent orderings (TSP, scheduling). Require specialised variation procedures (e.g. order-preserving transforms, mapping-based procedures, local reordering moves) that preserve permutation feasibility.
- **Tree and graph encodings:** Variable-size structures used in genetic programming, evolving expressions, or circuit topologies. Use subtree exchange and constrained growth controls to avoid bloat.
- **Indirect / developmental encodings:** Genotypes specify construction rules or grammars that generate phenotypes; useful when compact genotypes should produce structured phenotypes (e.g. neural architectures, L-systems).

Choosing an encoding

Select an encoding by matching the problem’s combinatorial structure, constraint set, and desired procedure toolkit. Key questions:

- Does the problem require an ordering, a multiset, or real-valued parameters? Choose permutation, integer/multiset, or real encodings respectively.
- Are feasibility constraints hard (must be satisfied) or soft (violations penalised)? For hard constraints prefer sound encodings or constructive decoders; for soft constraints penalisation may be acceptable.
- Is locality important for effective recombination? If so, prefer encodings (or transformations, e.g. Gray) that increase correlation between small genotypic and phenotypic changes.
- Will procedures be custom or standard? Use encodings that keep procedure implementation simple unless domain structure mandates otherwise.

Operator compatibility and empirical validation

An encoding is only useful if paired with procedures that preserve useful structure. After selecting an encoding, design or choose variation procedures that maintain feasibility, limit destructive epistasis, and respect meaningful gene boundaries. Finally, validate encoding choices empirically: compare performance across a small benchmark (different encodings, procedure sets, and variation rates) and select the combination that gives robust progress on representative instances.

The following sections give concrete representational examples and references for the main encoding families discussed here.

4.5 Real-valued Encoding

Real-valued encoding represents individuals as vectors in \mathbb{R}^n , i.e. $\mathbf{x} = (x_1, \dots, x_n)$ with each coordinate taking values on a continuous domain. This direct representation is the natural choice for continuous optimisation problems and for parameter tuning tasks where the phenotype is inherently numeric. By operating in a continuous space, real-valued encodings avoid the quantisation artefacts of fixed-length binary encodings and allow variation operators to express arbitrarily small adjustments to candidate solutions (subject to floating-point precision limits) [6, 30].

The principal practical advantage of real-valued encodings is operator compatibility: arithmetic recombination (weighted averages), BLX- α style interval recombination, simulated binary crossover (SBX) and Gaussian or Cauchy perturbation mutations all act naturally on real vectors and can be designed to respect bounds or known structure. These operators produce offspring that lie in the convex hull (or a controlled extension) of the parents, which typically yields smoother search trajectories and better exploitation of local gradients in the fitness landscape. Evolution Strategies (ES) and many modern continuous optimisers exploit these properties by combining self-adaptive step-size control with recombination to navigate rugged but differentiable landscapes efficiently [6].

There are, however, theoretical and practical trade-offs. Classical schema arguments developed for binary representations do not carry over directly to continuous encodings:

building-block notions must be reformulated in terms of regions of \mathbb{R}^n and operator-induced correlations between coordinates. Real encodings also place greater emphasis on algorithmic choices for step-size control and constraint handling — poor mutation scales or unbounded recombination can lead to slow progress or numerical instability. Consequently, practitioners must tune or adapt mutation magnitudes (fixed schedules, self-adaptation, or covariance matrix adaptation) and choose recombination parameters that match problem smoothness and scale.

From an implementation viewpoint, several pragmatic recommendations improve robustness and performance. Always normalise or scale variables to comparable ranges before applying generic operators; this prevents single coordinates from dominating recombination statistics and simplifies parameter transfer between problems. Use bounded operators or projection schemes when constraints are present, and prefer adaptive mutation strategies (e.g. log-normal step-size adaptation or CMA-style covariance updates) when the search landscape exhibits anisotropy. When local gradients are available or can be approximated cheaply, hybridising evolutionary updates with gradient-based refinement often accelerates convergence while preserving global exploration.

Finally, like any encoding choice, real-valued representations should be validated empirically against alternatives. For many smooth, low-to-moderate dimensional continuous problems they substantially outperform binary encodings in both convergence speed and final solution quality; for highly multimodal or combinatorial problems a real-valued parameterisation may be inappropriate. We therefore recommend starting with a simple real-valued operator set (arithmetic/BLX recombination and Gaussian mutation with a tuned standard deviation), run small factorial experiments to select adaptation mechanisms, and escalate to more sophisticated adaptations (self-adaptive step sizes, CMA) when warranted by problem scale or observed search behaviour.

4.6 Integer Encoding

Integer encoding represents solutions whose variables take discrete integer values. Formally an individual is a vector $\mathbf{x} = (x_1, \dots, x_n)$ with each coordinate $x_i \in \mathbb{Z}$ and, in practice, bounded to a finite domain $[a_i, b_i] \cap \mathbb{Z}$. This representation is appropriate for allocation problems, counts, and many combinatorial substructures (for example quantities in knapsack-like models, resource allocations, and discretised control parameters). The discrete nature of the variables changes the character of the search: neighbourhoods are naturally defined by integer steps, and the search landscape is inherently non-continuous and often non-convex.

Operators for integer encodings must respect integrality and any problem-specific bounds or feasibility constraints. Common mutation strategies include random-reset mutation (replace a coordinate with a uniformly sampled integer in its domain) and small-step or "creep" mutation (increment or decrement by a small integer drawn from a short-tailed distribution). Recombination can be performed directly in the integer domain (for example, discrete uniform crossover or coordinate-wise selection), or by temporarily lifting values to a continuous surrogate (arithmetic recombination followed by rounding) when an operator that benefits from averaging is desirable. When using surrogate continuous recombination, stochastic rounding or bias-corrected rounding helps reduce systematic rounding artefacts.

Integer encodings present trade-offs compared to real-valued representations. Because the domain is discrete, many analytical assumptions (e.g. smooth gradients or continuous

convexity) do not apply, and standard continuous step-size adaptation mechanisms require adaptation to discrete step scales. On the other hand, integer representations can encode feasibility directly, avoiding expensive repair procedures: e.g. representing quantities with integrality enforces natural constraints, and specialised discrete crossover/mutation operators can be designed to preserve feasibility or near-feasibility by construction.

From an algorithm design and implementation perspective several pragmatic recommendations improve robustness. First, exploit problem structure: if variables have small integer ranges prefer enumerative neighbourhood moves and small-step local search hybrids; if ranges are large, favour operators that explore broadly (random-reset, large-step proposals) combined with adaptive reduction in step magnitude. Second, enforce bounds and invariants in the decoder or by projection after variation rather than relying on implicit truncation; explicit constraint-aware operators are usually clearer and less error prone. Third, when mixing integer and continuous variables use mixed-integer operators or decoupled schedules so that each variable type receives appropriately scaled variation.

Finally, validate encoding choices empirically. Compare a direct integer encoding to alternatives (binary-encoded integers, real-valued surrogate with rounding) on small representative instances to measure convergence speed, robustness, and the cost of constraint handling. In many allocation or scheduling tasks a well-chosen integer representation plus tailored discrete operators outperforms generic continuous surrogates; however, for problems where fine-grained search behaviour is important, surrogate continuous strategies with careful rounding and step-size adaptation can be competitive. Use these empirical results to select mutation/recombination scales and to decide whether to hybridise the evolutionary loop with deterministic local search on integer neighbourhoods.

4.7 Permutation Encoding

Permutation encoding represents candidate solutions as permutations of a finite set of elements, i.e. the genotype space is the set of bijections on $\{1, \dots, n\}$. The genotype–phenotype mapping ϕ is usually the identity map: a permutation directly specifies an ordering that is interpreted by the problem-specific evaluator (for example a tour in the travelling salesman problem, a job sequence in single-machine scheduling, or an ordered list of tasks for a flow line). Because permutations inherently enforce ordering constraints, permutation encodings are sound for ordering problems and avoid many feasibility repairs required by naive encodings.

Although permutations are formally one-to-one with respect to orderings, practical representations often introduce equivalence classes and redundancies that must be recognised. A circular tour (as in symmetric TSP) admits rotational symmetry: cyclic shifts of a permutation represent the same tour and reflections may also be equivalent. Such symmetries do not change validity but affect sampling and selection probabilities; designers should either choose a canonical representative (fixing the first city) or use operators and fitness comparisons that account for the equivalence class to avoid representational bias.

Distance and locality in permutation spaces differ markedly from vector spaces. Hamming distance or simple positional metrics do not capture meaningful neighbourhood structure for order-based problems. Distances such as Kendall tau (number of pairwise disagreements), inversion distance, or edge-based metrics (number of differing adjacency relationships) better reflect the kinds of small, interpretable changes that preserve problem structure. Operator design should therefore be guided by which aspects of a permutation

constitute useful building blocks for the problem — position-based blocks, adjacency/edge blocks, or precedence relations — because different operators preserve different structures.

Variation operators for permutation encodings must preserve feasibility (i.e. produce valid permutations) and ideally respect the chosen notion of locality. Typical mutation moves are swap, insert (take-one-and-insert-at-another-position), and inversion/reversal of a subsequence; these have clear interpretations as small, local reorders. Recombination operators are designed to combine parent orderings while maintaining permutation validity: examples include partially mapped crossover (PMX), order crossover (OX), cycle crossover (CX), and edge recombination. Each emphasises different preserved structures (position, order, or adjacency) and the choice should match problem-specific building blocks (for instance, edge-based recombiners are natural for TSP where edges matter more than absolute positions).

An alternative is to use indirect encodings and constructive decoders when constraints or constructive heuristics are important. Priority or random-key encodings map real-valued keys to a permutation via a stable sorting decoder; constructive decoders build feasible schedules or tours greedily from a genotype that encodes preferences. Indirect encodings can drastically reduce design complexity by separating the genetic search from feasibility enforcement and can incorporate domain heuristics directly into the decoder, but they shift the design burden to the decoder and may obscure locality properties of the genetic operators.

Practical recommendations: initialise populations using a mixture of random permutations and problem-specific heuristics to seed useful structure; measure diversity with permutation-aware metrics (Kendall tau or edge overlap) rather than Hamming distance; prefer operators that preserve the notion of building blocks relevant to your problem; and combine global permutation-based search with local optimisation (e.g. 2-opt or 3-opt for TSP, or specialised neighbourhood search for scheduling) to exploit fine-grained improvements. When symmetries exist, use canonicalisation or equivalence-aware evaluation to avoid bias. Finally, validate choices empirically on representative instances, since operator effectiveness is strongly problem-dependent in permutation spaces.

4.8 Tree Encoding

Tree encoding represents genotypes as labelled, rooted trees whose nodes carry symbols drawn from one or more alphabets (for example function/operator symbols for internal nodes and terminal symbols for leaves). The phenotype is obtained by interpreting the tree according to problem semantics: in genetic programming the tree denotes an expression or program, in syntactic optimisation it denotes a parse tree, and in hierarchical design it denotes a composition of components. Formally the genotype space is the set of finite ordered trees over a ranked alphabet, and the decoder ϕ is the evaluation or instantiation function that maps a tree to the problem-specific object in P .

Tree encodings introduce representational choices that strongly affect operator behaviour and search dynamics. One must decide on the node alphabets (typed or untyped), arity constraints (fixed or variable arity), and linearisation for storage (pointer structures, bracketed strings, prefix/postfix notations, or explicit child lists). Typed (strongly-typed) trees enforce syntactic constraints at the representation level, preventing many invalid offspring and reducing the need for repair; untyped trees are more flexible but often require additional feasibility checks or decoders. Representation impacts locality: small subtree

replacements may induce large semantic changes when node semantics are non-linear or context-sensitive.

A central concern with tree encodings is bloat — unbounded growth in tree size without commensurate fitness improvement. Bloat arises from neutral or weakly selective regions where larger trees are not penalised, and it degrades performance by increasing evaluation cost and reducing effective population variability. Common countermeasures include static limits on depth/size, parsimony pressure (explicit size or complexity penalties in the fitness), and operator-aware controls (limiting offspring size in crossover and mutation). When using growth controls, balance is required: overly aggressive pruning can eliminate useful structural variation, while permissive settings lead to resource exhaustion.

Variation operators for trees must preserve tree well-formedness. Standard operators include subtree crossover (swap subtrees between parents), point mutation (replace a node or small subtree with a randomly generated subtree), and hoist mutation (replace a tree by one of its subtrees to reduce size). There are also context-preserving operators designed for typed languages or grammars (constrained subtree exchange, grammar-guided mutation) that maintain syntactic correctness by construction. Operator design should align with the semantics of the node alphabets: for expression trees, promoting commutative/associative-aware recombination or algebraic simplifications can increase meaningful offspring generation.

Indirect and grammar-based encodings are especially useful when the phenotype must satisfy rich syntactic or semantic constraints. In grammatical evolution and grammar-guided GP the genotype typically encodes a derivation or a sequence of production choices, and a deterministic decoder maps this to a tree that is guaranteed syntactically valid. These indirect encodings can produce compact genotypes and enable incorporation of domain knowledge, but they may obscure the locality properties of operators and require careful decoder design to avoid biased sampling of the phenotype space.

Practical recommendations: enforce feasibility early using typing or grammar constraints when the domain requires syntactic correctness; combine global tree-based search with local simplification passes (constant folding, algebraic reductions) to improve evaluation efficiency; use mixed initialisation strategies (ramped half-and-half, grow/full) to seed diverse tree sizes and shapes; and adopt explicit complexity control (parsimony pressure, depth limits, or adaptive operator bias) to manage bloat. Finally, validate operator choices empirically on representative instances and instrument tree-size, depth, and evaluation cost during experiments to detect emergent bloat or pathological behaviours.

4.9 Chapter Summary

This chapter covered various encoding schemes for genetic algorithms. The choice of encoding is crucial and should match the problem characteristics. Binary encoding provides theoretical foundation, real-valued encoding suits continuous optimization, permutation encoding handles ordering problems, and specialized encodings address domain-specific requirements.

4.10 Key Concepts

- Encoding requirements: completeness, soundness, non-redundancy

- Binary vs Gray code encoding
- Real-valued representation and procedures
- Permutation encoding for ordering problems
- Tree encoding for hierarchical structures
- Problem-specific encoding considerations

Chapter 5

Selection Methods in Genetic Algorithms

5.1 Introduction to Selection

Selection is the mechanism within a genetic algorithm that determines which individuals from a population are chosen to contribute genetic material to the next generation. At its core, selection converts fitness information into reproductive opportunities: individuals with relatively higher fitness are given greater chances to produce offspring, thereby biasing the search process toward promising regions of the solution space. This bias must be managed carefully so that the algorithm both exploits high-quality solutions and continues to explore diverse alternatives.

An essential concept associated with selection is selection pressure, which quantifies the degree to which better individuals are favored. High selection pressure accelerates convergence by amplifying the reproductive advantage of the fittest individuals, but it increases the risk of premature convergence where the population loses diversity and becomes trapped in suboptimal regions. Low selection pressure preserves diversity and encourages exploration, yet may slow the algorithm's progress toward improved solutions. Practical algorithm design therefore requires balancing these competing effects, often by tuning selection parameters or combining selection schemes with diversity-preserving mechanisms.

Selection operators come in several families, each offering different trade-offs between simplicity, selection pressure control, and sensitivity to fitness scaling. Common approaches include fitness-proportionate methods (e.g., roulette wheel and stochastic universal sampling), rank-based schemes that ensure controlled and scale-independent pressure, tournament selection which provides an adjustable and efficient means of imposing pressure, and truncation or elitist strategies that deterministically preserve top individuals. Later sections of this chapter examine these methods in detail, including their algorithms, statistical properties, and practical advantages or drawbacks.

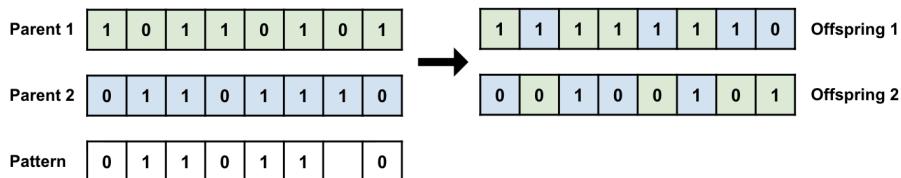


Figure 5.1: Basic selection process in Genetic Algorithms

5.2 Selection Pressure

Selection pressure quantifies how strongly a selection mechanism favors individuals with higher fitness when producing the next generation. Intuitively, it measures the expected reproductive advantage of good solutions relative to the population average. Selection pressure can be formalized in several ways; common operational measures include selection intensity (the standardized difference between parent and population means) and takeover time (the number of generations required for the best individual to dominate under repeated selection). These measures allow practitioners to compare different selection operators and parameterizations in a principled manner.

The magnitude of selection pressure has direct and predictable effects on search dynamics. Strong pressure accelerates the propagation of beneficial alleles and shortens the time to apparent convergence, which is useful when the fitness landscape is smooth and unimodal. However, excessive pressure reduces genetic diversity and increases the risk of premature convergence to local optima. Conversely, weak pressure preserves diversity and supports broader exploration of the search space but slows progress toward high-fitness regions. Therefore, the choice of selection operator and its parameters should be informed by the problem’s modality, population size, and the available number of generations.

Practical controls for selection pressure include algorithmic choices (e.g., tournament size, ranking slope, truncation fraction), fitness scaling techniques (e.g., linear or sigma scaling, Boltzmann selection), and hybrid strategies that adapt pressure during the run (e.g., start with low pressure for exploration and increase pressure for exploitation). Monitoring selection-related statistics — such as mean and variance of fitness, diversity measures (e.g., average Hamming distance in binary encodings), and takeover time estimates — provides actionable feedback for tuning. In applied settings, a common heuristic is to begin with moderate pressure (e.g., small tournament sizes, conservative ranking parameters) and adjust based on empirical performance and observed loss of diversity.

5.3 Fitness Proportionate Selection (FPS)

The Genetic Algorithm developed by Holland uses Fitness Proportionate Selection (FPS) [25, 21], where the expected value of an individual (i.e., the expected number of times that individual will be selected for reproduction) is calculated as that individual’s fitness divided by the population’s average fitness.

In this method, each individual can be selected as a parent with a probability proportional to its fitness value. Therefore, individuals with higher fitness have greater opportunities to reproduce and spread their characteristics to the next generation. Thus, this method provides selection pressure on fitter individuals in the population, thus driving evolution toward better individuals over time.

5.3.1 Roulette Wheel Selection

Also known as fitness proportionate selection, individuals are selected with probability proportional to their fitness [21, 3, 4].

The simplest selection schema is roulette-wheel selection, also called stochastic sampling with replacement. This is a stochastic algorithm and involves the following technique:

Individuals are mapped to contiguous segments on a line, where the size of each segment is equal to that individual's fitness value. A random number is generated, and the individual whose segment spans that random number is selected. This process is repeated until the desired number of individuals is reached, called the mating population. This technique is analogous to a roulette wheel, where each slice is proportional in size to the fitness value.

Number of Individual	Fitness Value	Selection Probability	Interval
1	2.0	0.18	[0.00, 0.18]
2	1.8	0.16	[0.18, 0.34]
3	1.6	0.15	[0.34, 0.49]
4	1.4	0.13	[0.49, 0.62]
5	1.2	0.11	[0.62, 0.73]
6	1.0	0.09	[0.73, 0.82]
7	0.8	0.07	[0.82, 0.89]
8	0.6	0.06	[0.89, 0.95]
9	0.4	0.03	[0.95, 0.98]
10	0.2	0.02	[0.98, 1.00]
11	0.0	0.0	—

Table 5.1: Selection probability and fitness value (from Buku Ajar)

Table 5.1 shows the selection probabilities for 11 individuals, with linear ranking with selective pressure of 2, along with their fitness values. Individual 1 is the individual with the highest fitness and occupies the largest interval, while individual 10 as the individual with the second lowest fitness has the smallest interval on the line. Individual 11, with the lowest fitness, has fitness value = 0 and gets no chance for reproduction.

To select the mating population, a number of uniformly distributed random numbers (uniformly distributed between 0.0 and 1.0) are generated independently.

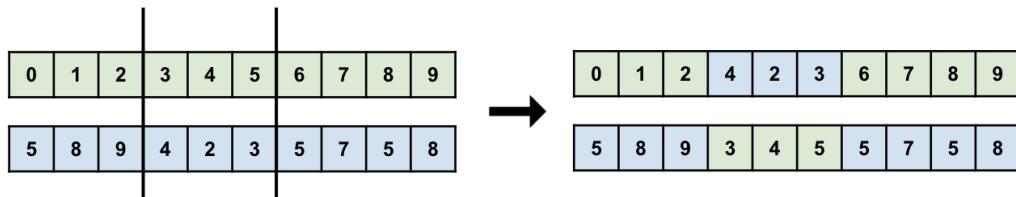


Figure 5.2: Roulette-wheel selection process with sample trials

The disadvantage of roulette-wheel selection is that although it provides zero bias, it does not guarantee minimum spread.

Algorithm

Selection Probability

The probability of selecting individual i is:

$$P_i = \frac{f_i}{\sum_{j=1}^N f_j} \quad (5.1)$$

Algorithm 1 Roulette Wheel Selection

```

Calculate total fitness:  $F = \sum_{i=1}^N f_i$ 
Generate random number:  $r \sim U[0, F]$ 
Set cumulative fitness:  $sum = 0$ 
for  $i = 1$  to  $N$  do
     $sum = sum + f_i$ 
    if  $sum \geq r$  then
        Select individual  $i$ 
        break
    end if
end for

```

Example

Individual	Fitness	Probability	Cumulative
1	10	0.25	0.25
2	20	0.50	0.75
3	5	0.125	0.875
4	5	0.125	1.0
Total	40	1.0	

Table 5.2: Roulette Wheel Selection Example

If random number $r = 0.6$, individual 2 is selected.

Advantages

Roulette-wheel selection is straightforward to implement and directly realizes fitness-proportionate sampling, so higher-fitness individuals are naturally more likely to contribute genetic material. Its probabilistic nature ensures that every individual retains a nonzero chance of selection, which helps maintain some exploration even when fitter individuals dominate.

Disadvantages

Roulette-wheel selection can suffer when fitness values are highly skewed: individuals with very large fitness may dominate and drive premature convergence, while very similar fitness values lead to weak selection pressure. Practical use therefore often requires fitness scaling or normalization, and care must be taken with negative or non-comparable fitness measures.

5.3.2 Stochastic Universal Sampling (SUS)

Improved version of roulette wheel selection that reduces variance [8].

Baker's SUS

Stochastic Universal Sampling (SUS) provides zero bias and minimum spread [8]. Individuals are mapped to contiguous segments on a line, where the size of each segment equals its fitness value, exactly as in roulette-wheel selection. In this method, equally spaced pointers are placed on the line equal to the number of individuals to be selected.

Let $N_{Pointer}$ be the number of individuals to be selected, then the distance between pointers is $1/N_{Pointer}$, and the position of the first pointer is determined by a random number generated in the range $[0, 1/N_{Pointer}]$.

For example, to select 6 individuals, the distance between pointers is $1/6 = 0.167$.

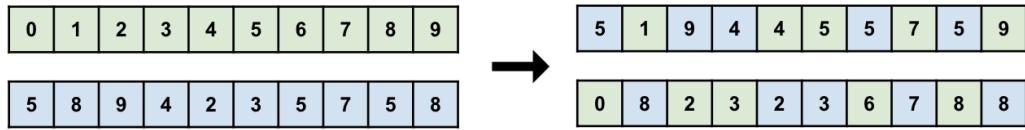


Figure 5.3: Stochastic universal sampling with equally spaced pointers

Stochastic universal sampling ensures offspring selection that is closer to the expected values compared to roulette-wheel selection.

Algorithm

Algorithm 2 Stochastic Universal Sampling

```

Calculate total fitness:  $F = \sum_{i=1}^N f_i$ 
Calculate pointer distance:  $distance = F/N$ 
Generate random start:  $start \sim U[0, distance]$ 
Create pointers:  $pointer_i = start + i \times distance$  for  $i = 0, 1, \dots, N - 1$ 
for each pointer do
    Select individual using roulette wheel logic
end for

```

Advantages over Roulette Wheel

Stochastic Universal Sampling reduces sampling variance relative to independent roulette draws by using evenly spaced pointers; this produces selection counts that are closer to their expected values and yields a more uniform coverage of the population. As a consequence, SUS better preserves the expected representation of individuals across repeated samplings, improving stability of the selection operator.

5.4 Rank-based Selection

Rank-based selection assigns selection probabilities based on fitness rank rather than raw fitness values [22, 5].

5.4.1 Overview

Ranked-Based Selection introduces a different approach to selection in Genetic Algorithms. Instead of directly using fitness values to determine selection probability, individuals in the population are first sorted (ranked) based on their fitness values, then each individual is assigned a rank. The selection probability is then calculated based on that rank, not the actual fitness value.

This rank-based approach helps reduce problems associated with direct fitness-based selection, such as premature convergence and domination by a few very fit individuals in the early stages of the optimization process. By assigning ranks and using them for selection, Ranked-Based Selection provides more balanced and controlled selection pressure, allowing better exploration of the search space and maintaining diversity in the population.

Rankings are typically assigned linearly or exponentially, where the best individual receives the highest rank and the worst individual receives the lowest rank. Selection probability is then calculated based on that ranking using a predetermined formula or mapping function. This mapping function can be adjusted to control selection pressure, where higher pressure will favor individuals with the highest ranks, while lower pressure provides a more even distribution of selection probabilities.

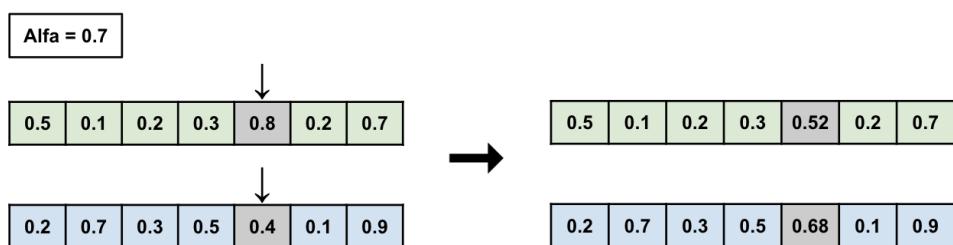


Figure 5.4: How the situation changes after converting fitness to order number (rank)

5.4.2 Linear Ranking

$$P_i = \frac{1}{N} \left[\eta^- + (\eta^+ - \eta^-) \frac{rank_i - 1}{N - 1} \right] \quad (5.2)$$

where:

- $rank_i$ is the rank of individual i ($1 = \text{worst}$, $N = \text{best}$)
- η^+ is the expected number of copies for best individual
- η^- is the expected number of copies for worst individual
- $\eta^+ + \eta^- = 2$ (to maintain population size)
- Typically: $\eta^+ = 2.0$, $\eta^- = 0.0$

5.4.3 Exponential Ranking

$$P_i = \frac{1 - e^{-rank_i}}{c} \quad (5.3)$$

where c is a normalization constant ensuring $\sum P_i = 1$.

5.4.4 Advantages of Rank Selection

By basing probabilities on rank rather than raw fitness, rank-based selection enforces a predictable and bounded selection pressure that is insensitive to the scale or distribution of fitness values. This makes it robust to outliers and negative fitness values and helps prevent a few extreme individuals from dominating the population.

5.4.5 Disadvantages

Rank-based schemes discard the magnitude information contained in fitness differences, which can slow convergence when those magnitudes are informative; additionally, they require sorting the population each generation, introducing an $O(N \log N)$ cost and some implementation complexity compared to simpler, linear-time samplers.

5.5 Tournament Selection

Tournament selection randomly selects k individuals and chooses the best among them [21, 7].

5.5.1 Overview

Tournament selection is a strong and widely used selection mechanism in Genetic Algorithms because it can maintain a balance between diversity maintenance and selective pressure [7]. Unlike roulette-wheel selection, which directly depends on an individual's fitness relative to the total population fitness, tournament selection works by holding "tournaments" among subsets of individuals, and the winner of each tournament is selected for reproduction.

The main concept of tournament selection is quite simple: instead of considering the entire population at once, a subset of individuals is randomly selected to compete with each other. The individual with the highest fitness in that "tournament" is then selected. This process is repeated until the desired number of individuals for reproduction is reached.

This method has several advantages: tournament selection maintains diversity because individuals with low fitness still have the opportunity to participate in tournaments. Additionally, this method allows selective pressure to be adjusted by setting the tournament size.

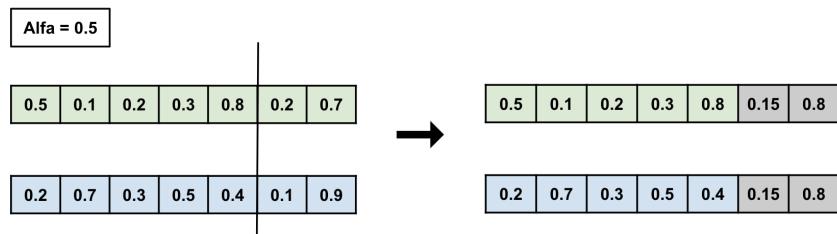


Figure 5.5: Tournament selection mechanism

5.5.2 Tournament Selection Mechanism

1. Determine tournament size (k), i.e., the number of individuals participating in each tournament.
2. Randomly select k individuals from the population.
3. Compare the fitness values of these individuals and select the individual with the highest fitness as the winner.
4. Add the winner to the mating pool.
5. Repeat steps 2–4 until the desired number of individuals is reached.

5.5.3 Binary Tournament

Most common form with $k = 2$.

Algorithm 3 Binary Tournament Selection

```

Randomly select individual  $i$ 
Randomly select individual  $j$  (where  $j \neq i$ )
if  $f_i > f_j$  then
    Select individual  $i$ 
else
    Select individual  $j$ 
end if
```

5.5.4 k-Tournament Selection

Algorithm 4 k-Tournament Selection

```

Create empty tournament set  $T$ 
for  $i = 1$  to  $k$  do
    Randomly select individual and add to  $T$ 
end for
Select best individual from  $T$ 
```

5.5.5 Tournament Size Effects

- $k = 1$: Random selection (no pressure)
- Small k : Low selection pressure
- Large k : High selection pressure
- $k = N$: Always selects best individual

5.5.6 Selection Probability

For individual with rank r out of N ($1 = \text{worst}$, $N = \text{best}$):

$$P_i = \frac{1}{N} \binom{N}{k} \sum_{j=0}^{r-1} \binom{j}{k-1} \binom{N-j-1}{0} \quad (5.4)$$

For binary tournament ($k = 2$):

$$P_i = \frac{2r-1}{N^2} \quad (5.5)$$

5.5.7 Advantages

Tournament selection is easy to implement, requires only local comparisons (no global fitness normalization), and provides a direct knob for selection pressure via the tournament size. Its independence from global statistics also makes it naturally parallelizable and tolerant of arbitrary fitness scales, including negative values.

5.5.8 Disadvantages

Tournament selection's behavior depends strongly on the chosen tournament size: large tournaments can impose very high pressure and reduce diversity, while very small tournaments approach random selection. The method can also repeatedly choose the same individual in multiple tournaments, which—without complementary diversity mechanisms—may accelerate loss of genetic variety.

5.6 Truncation Selection

Truncation selection is a deterministic policy that retains only the top fraction of the population for reproduction: given a population of size λ , the best μ individuals (by fitness) are selected and used to produce the next generation. The central parameter is the selection ratio

$$\rho = \frac{\mu}{\lambda}, \quad (5.6)$$

which directly controls selection pressure — smaller values of ρ correspond to stronger pressure because fewer individuals are permitted to reproduce. Truncation is typically implemented by sorting individuals by fitness (complexity $O(\lambda \log \lambda)$) and slicing the top μ entries; the deterministic nature makes its effect on the population easy to analyze and predict.

The principal effect of truncation selection is strong and immediate directional pressure: good solutions rapidly dominate the gene pool, which can greatly speed convergence in smooth, unimodal landscapes. This same property is its main drawback in multimodal or deceptive landscapes, where aggressive truncation reduces genetic diversity and increases the likelihood of premature convergence to suboptimal peaks. To mitigate these risks, practitioners often choose moderate values of ρ (common heuristics place ρ between 0.1 and 0.5 depending on problem scale) or combine truncation with diversity-preserving measures such as fitness sharing, crowding, or occasional stochastic replacement.

In practice, truncation is well suited for exploitation phases of an evolutionary run or for algorithms that require deterministic selection behavior (for reproducibility or theoretical analysis). When using truncation, monitor diversity statistics (e.g., genotype variance or average pairwise distance) and consider adaptive schedules that relax truncation early in the run and tighten it later as the search focuses on refinement.

5.7 Boltzmann Selection

Boltzmann selection adapts the familiar Boltzmann (Gibbs) distribution from statistical mechanics to map fitness values to selection probabilities. Under this scheme each individual i is assigned selection probability

$$P_i = \frac{e^{f_i/T}}{\sum_{j=1}^N e^{f_j/T}}, \quad (5.7)$$

where f_i is the fitness of individual i and $T > 0$ is the temperature parameter that controls the degree of randomness in selection. When T is large the distribution approaches uniform sampling (promoting exploration); as $T \rightarrow 0$ the distribution concentrates on the best individuals (promoting exploitation). This explicit temperature control makes Boltzmann selection a natural mechanism for smoothly interpolating between exploration and exploitation during an evolutionary run.

Practical use of Boltzmann selection requires a temperature schedule $T(t)$ that varies with generation t . A common choice is exponential cooling, for example $T(t) = T_0 \alpha^t$ with $0 < \alpha < 1$, but linear or problem-specific schedules may be preferable depending on landscape characteristics. The choice of initial temperature T_0 and the annealing rate determine how quickly selection pressure increases; poor choices can either leave the search unfocused for too long or force premature convergence.

Compared with simpler selection schemes, Boltzmann selection has two notable trade-offs. Its benefits are principled control of pressure and the ability to schedule a gradual transition from exploration to exploitation. Its costs are additional parameter tuning (temperature schedule) and slightly higher computational overhead due to exponentials and normalization. In practice, Boltzmann selection is most valuable when a controlled annealing strategy is desirable—e.g., when combining global exploration early with focused refinement later—or when fitness magnitudes vary widely and a temperature parameter offers robust scaling. When using Boltzmann selection, monitor fitness variance and population diversity, and be prepared to adjust the temperature schedule empirically for best results.

5.8 Elitist Selection

Elitist selection refers to selection policies that guarantee the survival of one or more top-ranked individuals from one generation to the next. The rationale is simple: stochastic variation in reproduction and replacement can accidentally discard the best solutions; preserving a small set of elites ensures that high-quality genetic material is never lost, which in turn makes measured, monotonic progress possible in many implementations.

There are two common variants. Pure elitism explicitly copies the best e individuals unchanged into the next generation; this is the most conservative approach and is typically

used with very small e (often $e = 1$). Elitist replacement is a softer variant in which after the usual selection and reproduction steps the worst individuals in the new population are replaced by the best individuals from the previous generation if those elites are strictly better. Both variants preserve improved solutions but differ in determinism and how aggressively they constrain the population.

The principal benefit of elitism is reliability: it prevents the accidental loss of the best-so-far solutions and often accelerates practical convergence by retaining proven building blocks. However, elitism also affects population diversity and exploration. If too many elites are preserved or elites are preserved for too long, the search can become overly exploitative, reducing the population's capacity to discover alternate peaks. Good practice is to keep the elite count small relative to population size (for instance, $e/\lambda \ll 0.1$) and, when necessary, combine elitism with explicit diversity-preserving techniques (e.g., occasional random immigrants, fitness sharing, or controlled mutation rates).

When using elitist strategies, monitor both the best fitness and diversity indicators (e.g., genotype variance, number of unique individuals). Consider adaptive policies that reduce elitism early to promote exploration or temporarily increase elitism late in a run for final refinement. These pragmatic controls help preserve the safety that elitism provides while mitigating its tendency to narrow the search prematurely.

5.9 Diversity-Preserving Selection

Diversity-preserving selection encompasses techniques intended to maintain useful genetic variation in the population so that evolution can continue to explore multiple promising regions of the search space. Maintaining diversity is particularly important for multimodal and deceptive problems where premature loss of variation can cause the run to converge to suboptimal peaks. The following paragraphs summarize commonly used mechanisms and practical guidance for their use.

One widespread approach is fitness sharing, which reduces the effective fitness of individuals that are close to many others in genotype or phenotype space. The shared fitness of individual i is computed as

$$f'_i = \frac{f_i}{\sum_{j=1}^N sh(d_{ij})}, \quad (5.8)$$

where d_{ij} is a distance between individuals i and j (Hamming distance for binary encodings or Euclidean distance for real-valued representations) and the sharing function is often defined as

$$sh(d) = \begin{cases} 1 - \left(\frac{d}{\sigma_{share}}\right)^\alpha & \text{if } d < \sigma_{share}, \\ 0 & \text{otherwise.} \end{cases} \quad (5.9)$$

The parameter σ_{share} defines the niche radius and α controls the decline of sharing; typical practice uses $\alpha = 1$ and selects σ_{share} by testing or domain knowledge. Fitness sharing encourages the population to occupy multiple niches and reduces the advantage of densely populated regions.

Crowding methods provide an alternative that directly controls replacement: offspring are preferentially compared and possibly replace similar individuals rather than random or worst ones. Deterministic crowding and probabilistic crowding are common variants; both aim to preserve local subpopulations by ensuring that newly created individuals

compete with genetically similar members, thereby preventing a single genotype from quickly sweeping the population.

Speciation, niching, and island models explicitly partition the population into subpopulations (species or islands) that evolve semi-independently, with occasional migration of individuals between groups. These structures preserve diversity by allowing different regions of the search space to be explored in parallel and are especially useful for problems with many well-separated optima. Practical design choices include migration rate, topology (ring, fully connected, etc.), and migration policy (best individuals, random migrants, or fitness-proportionate migrants).

When choosing and tuning diversity-preserving mechanisms, consider computational cost and measurement: fitness sharing requires $O(N^2)$ pairwise distance computations unless approximations or clustering are used; crowding typically runs in $O(N)$ per generation with careful bookkeeping; speciation and island models scale linearly per island but require configuration of migration parameters. Monitor population statistics (e.g., average pairwise distance, number of distinct genotypes, fitness variance) to detect excessive loss of diversity and adjust parameters dynamically (for example, increase mutation rate or relax selection pressure when diversity drops). Combining modest diversity-preserving methods with a well-calibrated selection pressure often yields the best practical results.

5.10 Multi-objective Selection

Multi-objective selection treats optimization problems where the quality of a solution is described by a vector of objectives rather than a single scalar. Let $\mathbf{f}(\mathbf{x}) = (f_1(\mathbf{x}), \dots, f_m(\mathbf{x}))$ denote the objective vector to be maximized; the concepts below are straightforwardly adapted to minimization by sign reversal. Central to multi-objective selection is the notion of Pareto dominance: a solution \mathbf{x} dominates \mathbf{y} (written $\mathbf{x} \prec \mathbf{y}$) if

$$\forall k \in \{1, \dots, m\} : f_k(\mathbf{x}) \geq f_k(\mathbf{y}), \quad \text{and} \quad \exists k : f_k(\mathbf{x}) > f_k(\mathbf{y}). \quad (5.10)$$

This partial order induces a front structure on the population: nondominated solutions form the Pareto front (rank 1), the nondominated set of the remainder forms rank 2, and so on. Non-dominated sorting partitions the population by repeatedly extracting the current nondominated set; the naive procedure has worst-case cost $O(mN^2)$ for N individuals and m objectives, while optimized algorithms and data structures can offer empirical improvements for many practical sizes.

Selection in multi-objective evolutionary algorithms (MOEAs) must both promote convergence toward the Pareto front and preserve diversity along it. A widely used strategy, popularized by NSGA-II, performs selection in two stages: (1) apply non-dominated sorting to assign a rank to each individual, and (2) within the same rank prefer individuals that increase population spread using a crowding measure. The crowding distance for an individual is computed by summing normalized gaps between neighboring solutions for each objective after sorting by that objective; larger crowding distance indicates a less crowded region and is therefore preferred when ranks tie. Practically, parent or survivor selection can be implemented as a binary tournament that compares first by rank (lower is better) and then by crowding distance (higher is better), which yields a simple, effective rule that balances convergence and diversity.

Several practical considerations arise when applying multi-objective selection. For many objectives (the many-objective case, $m \gtrsim 5$), dominance relations become less dis-

criminating and alternative approaches—indicator-based methods (e.g., IBEA), decomposition techniques (e.g., MOEA/D), or reference-point strategies—often perform better. Computational cost is also important: while NSGA-II is efficient and robust for moderate N and m , indicator-based selection (hypervolume-based) can be costly for large fronts. Finally, parameter choices (population size, replacement policy, mating selection) affect both the algorithm’s ability to approximate the Pareto front and the distribution of solutions; monitor convergence (e.g., IGD, hypervolume) and spread metrics, and consider hybridizing selection with archiving strategies or adaptive population sizing when sustained exploration of multiple trade-offs is required.

5.11 Selection Comparison

Method	Pressure	Diversity	Complexity	Scalability	Parameters
Roulette Wheel	Variable	Poor	$O(N)$	Poor	None
SUS	Variable	Good	$O(N)$	Poor	None
Rank Linear	Constant	Good	$O(N \log N)$	Good	η^+, η^-
Tournament	Adjustable	Good	$O(1)$	Excellent	k
Truncation	High	Poor	$O(N \log N)$	Good	μ/λ
Boltzmann	Adaptive	Excellent	$O(N)$	Good	$T(t)$

Table 5.3: Comparison of Selection Methods

The table above summarizes key practical properties of commonly used selection methods. It condenses four dimensions that guide method choice: the effective selection pressure (how strongly the operator favors high-fitness individuals), the operator’s tendency to preserve or erode population diversity, the asymptotic computational complexity for a single generation, and how well the method scales with population size or parallel implementations. The final column lists the principal tuning parameters that practitioners must set or schedule.

Interpreting the table requires combining its quantitative entries with problem-specific considerations. Methods labelled “variable” pressure (roulette wheel and SUS) depend directly on the raw fitness distribution and so are sensitive to scaling and outliers; when fitness values are skewed these methods either collapse diversity (large gaps) or provide negligible pressure (small differences). Stochastic Universal Sampling reduces the sampling variance of roulette draws and therefore yields selection counts closer to expectation, but it does not by itself remove sensitivity to fitness scaling.

Rank-based linear selection deliberately discards raw magnitude information in favor of ordinal information, producing predictable and bounded pressure that is robust to arbitrary fitness scales and outliers. The trade-off is loss of useful signal when fitness magnitudes are meaningful, plus the $O(N \log N)$ cost of sorting each generation. Tournament selection provides an efficient, local-comparison mechanism whose pressure is adjusted directly by the tournament size k ; it is simple, parallel-friendly, and insensitive to global normalization, which explains its widespread use in large-scale and distributed implementations.

Truncation selection is the most aggressive deterministic policy: keeping only the top fraction (μ/λ) applies very high pressure and rapidly concentrates the population, which can be desirable in unimodal problems or late-stage exploitation but harmful in multi-modal or deceptive landscapes absent strong diversity-preservation. Boltzmann selection

offers an explicit, continuous control knob via the temperature schedule $T(t)$: when tuned well, it interpolates smoothly between exploration and exploitation and handles widely varying fitness magnitudes, at the cost of an additional scheduling parameter and the need to monitor annealing behavior.

From a practical standpoint, selection should rarely be chosen on a single criterion. For problems where fitness scaling is unreliable or unknown, prefer rank-based or tournament methods. For applications that demand reproducible, deterministic behavior or very fast convergence on a known unimodal problem, truncation (with small μ/λ) or elitist augmentation can be appropriate. When maintaining a diverse Pareto of solutions or exploring rugged landscapes, combine selection with explicit diversity mechanisms (fitness sharing, crowding, speciation) and keep selection pressure moderate. Finally, parameter choices (tournament size, ranking slope, truncation fraction, temperature schedule) should be validated empirically and monitored with diversity statistics (e.g., average pairwise distance, takeover time) to avoid premature convergence.

The remainder of the chapter provides guidelines and hybrid strategies that implement these recommendations in practice.

5.12 Selection Guidelines

Choice of selection operator should be informed first by the problem’s modality and deception. For problems that are essentially unimodal or where a single peak is the objective, stronger selection pressure (for example, truncation or larger tournament sizes) accelerates convergence and is often appropriate. For multimodal problems, where multiple peaks may contain useful solutions, moderate pressure such as binary tournament or rank-based selection helps preserve alternative lineages and reduces the risk of premature loss of useful niches. In deceptive landscapes—where locally attractive solutions mislead the search—favor lower pressure and couple selection with explicit diversity-preserving mechanisms (fitness sharing, crowding, speciation or island models) so that the algorithm can continue exploring promising but initially low-frequency regions.

Population size interacts with selection pressure in important ways. Small populations are more susceptible to sampling error and genetic drift, so conservative pressure settings (smaller tournament sizes, gentler ranking parameters, and limited elitism) help maintain useful variation. Larger populations can sustain stronger pressure without as much risk of accidental loss of rare but valuable genotypes, and they are better suited to schemes that rely on sampling stability (e.g., rank-based selection or modest truncation). In all cases, monitor diversity statistics (average pairwise distance, number of unique genotypes, fitness variance) and adjust selection parameters if diversity falls faster than expected.

Selection intensity should also vary over the run rather than remain fixed. Early generations benefit from lower effective pressure—larger temperature in Boltzmann schedules, smaller tournament sizes, or milder ranking—so that the search emphasizes exploration and discovers multiple basins of attraction. As the run progresses and the population accumulates evidence about promising regions, gradually increase pressure (reduce temperature, enlarge tournaments, or tighten truncation) to focus effort on exploitation and refinement. Adaptive schedules, occasional re-introduction of random immigrants, or multi-level selection (different operators for parent selection and survivor selection) are practical ways to implement this temporal modulation while guarding against premature convergence.

5.13 Hybrid Selection Strategies

Hybrid selection strategies combine multiple selection ideas to obtain better practical performance than any single method in isolation. One common approach is adaptive selection, where operators or their parameters are adjusted automatically during the run in response to population statistics. Examples include annealing a Boltzmann temperature $T(t)$, increasing tournament size as diversity drops, or switching from rank-based to truncation selection during late exploitation. Adaptive rules can be simple (predefined schedules) or feedback-driven (triggered by measured diversity, fitness improvement rate, or takeover time). The central advantage of adaptation is that it permits different phases of the search—exploration and exploitation—to use different pressure regimes without manual retuning for each problem instance.

Multi-level selection splits selection responsibilities across different stages or hierarchies. For instance, parent selection (which chooses who mates) can use a low-pressure method to preserve variety of mating combinations, while survivor selection (which decides who remains in the population) can be more aggressive to retain progress. Similarly, island or hierarchical population models run semi-independent subpopulations with occasional migration; within each island different selection policies may be used to encourage complementary search behavior. Multi-level designs are particularly effective when search must balance global exploration with local refinement or when computational resources are distributed across nodes.

Combined methods explicitly mix selection operators to exploit complementary strengths. Practical examples include performing tournament selection for most parents but reserving a fraction of survivors for rank-based or fitness-shared selection to preserve niches, or applying stochastic universal sampling with elitist replacement to reduce sampling variance while guaranteeing the best-so-far individuals survive. When combining operators, carefully manage interactions — for example, ensure that deterministic components (elitism, truncation) do not negate the diversity benefits of stochastic components. Empirical validation, monitoring of diversity metrics, and conservative parameterization (small elite fractions, limited truncation windows) help achieve robust gains.

In practice, hybrid strategies are most useful when the problem exhibits multiple regimes of difficulty (early exploration, mid-run discovery of promising basins, late-stage refinement) or when robustness across problem instances is required. Implement hybrids incrementally, instrument population statistics to guide choices, and prefer simple, interpretable combinations over elaborate schemes unless justified by experimental results.

5.14 Chapter Summary

This chapter covered various selection methods in genetic algorithms. Selection balances exploration and exploitation, with different methods offering different selection pressures and characteristics. Tournament selection is often preferred for its simplicity and effectiveness, while rank-based methods provide consistent pressure. The choice depends on problem characteristics, population size, and desired convergence behavior.

5.15 Key Concepts

- Selection pressure and its effects

- Proportional vs. rank-based selection
- Tournament selection and its variants
- Elitism and diversity preservation
- Multi-objective selection methods
- Guidelines for choosing selection methods

Chapter 6

Crossover (Recombination) in Genetic Algorithms

6.1 Introduction to Crossover

Crossover, also known as recombination, is the primary genetic operator in genetic algorithms [25, 21, 31]. It combines genetic material from two or more parent solutions to create offspring, potentially inheriting beneficial traits from multiple parents. Crossover exploits existing solutions to explore new regions of the search space.

6.2 Biological Inspiration

In nature, sexual reproduction combines genetic material from two parents:

- **Crossing over:** Exchange of genetic segments between homologous chromosomes
- **Independent assortment:** Random distribution of chromosomes
- **Genetic diversity:** Offspring differ from parents
- **Building blocks:** Beneficial gene combinations are preserved and mixed

6.3 Crossover Principles

6.3.1 Exploration vs. Exploitation

- **Exploitation:** Combines good building blocks from parents
- **Exploration:** Creates new combinations not present in parents
- **Heritability:** Offspring resemble parents but with variations

6.3.2 Crossover Probability

Crossover is typically applied with probability p_c (usually 0.6-0.9):

- High p_c : More exploration, faster convergence
- Low p_c : More exploitation of current solutions
- $p_c = 1.0$: Always apply crossover
- $p_c = 0.0$: No crossover (mutation-only evolution)

6.4 Binary Crossover Operators

6.4.1 Definition and Function of Crossover Operator

Crossover is a genetic operator used to vary the arrangement of chromosomes from one generation to the next [19, 42, 1]. The crossover method used depends on the encoding method applied.

Crossover occurs in three stages:

1. The reproduction operator randomly selects a pair of individual strings for the mating process.
2. A cross site is randomly selected along the length of the string.
3. The values after the cross site are exchanged between the two strings to form new offspring.

In binary chromosome representation, each individual in the population is represented as a sequence of bits (0 and 1) that express a potential solution to a problem.

6.4.2 One-Point Crossover

Single-point crossover is the simplest form of crossover. One crossover position is randomly selected, then the part of the two parents after that position is exchanged to form two new offspring. The strings resulting from this process have Positional Bias characteristics.

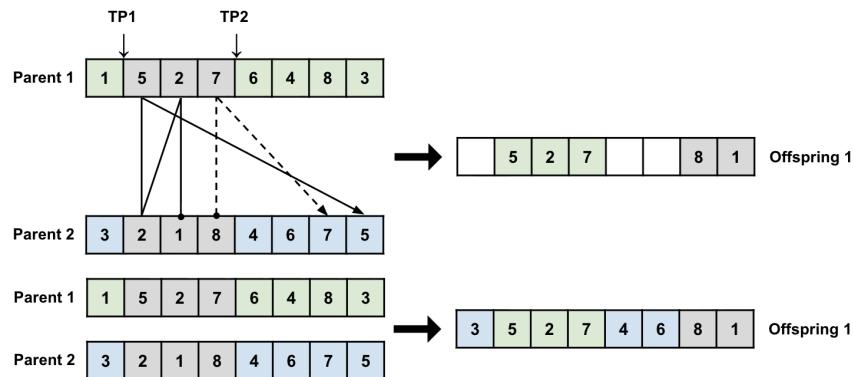


Figure 6.1: Single-Point Crossover for binary chromosomes

6.4.3 One-Point Crossover

Single crossover point divides chromosomes into two segments.

Algorithm

Algorithm 5 One-Point Crossover

Select random crossover point $k \in [1, l - 1]$

Create offspring:

$$\text{child}_1 = \text{parent}_1[1 : k] + \text{parent}_2[k + 1 : l]$$

$$\text{child}_2 = \text{parent}_2[1 : k] + \text{parent}_1[k + 1 : l]$$

Example

$$\text{Parent 1: } 1|1010011 \tag{6.1}$$

$$\text{Parent 2: } 0|0111100 \tag{6.2}$$

$$\text{Child 1: } 1|0111100 \tag{6.3}$$

$$\text{Child 2: } 0|1010011 \tag{6.4}$$

Crossover point at position 1.

Characteristics

- Simple and efficient
- Preserves long building blocks near chromosome ends
- May disrupt building blocks crossing the crossover point
- Positional bias (end positions less likely to be separated)

6.4.4 Two-Point Crossover

Two crossover points create three segments.

Multi-point crossover has a mechanism like single-point crossover; the difference is the number of randomly selected positions. In Multi-point crossover, a number N of positions are randomly selected along the length of the chromosome. These positions are exchanged to form two new offspring.

The objectives of multi-point crossover include increasing genetic diversity in the population, reducing positional bias, and increasing the chances of recombination from various schemas or different solution blocks. However, using too many cut points can become too disruptive, as it can damage already good gene combinations (coadapted alleles).

Algorithm

Algorithm 6 Two-Point Crossover

Select two random points k_1, k_2 where $1 \leq k_1 < k_2 \leq l - 1$

Create offspring:

$$\text{child}_1 = \text{parent}_1[1 : k_1] + \text{parent}_2[k_1 + 1 : k_2] + \text{parent}_1[k_2 + 1 : l]$$

$$\text{child}_2 = \text{parent}_2[1 : k_1] + \text{parent}_1[k_1 + 1 : k_2] + \text{parent}_2[k_2 + 1 : l]$$

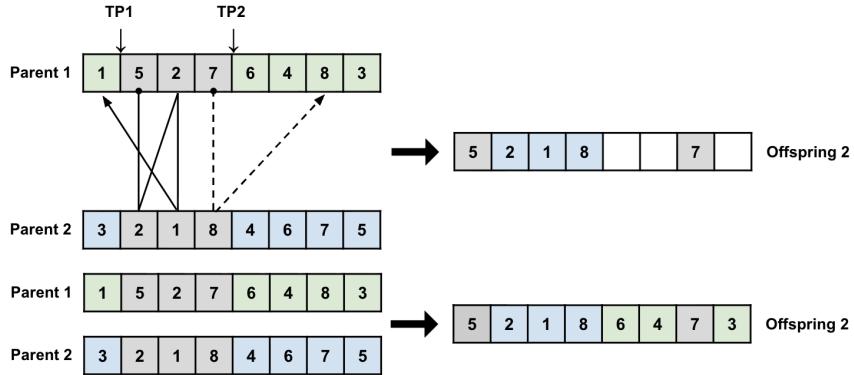


Figure 6.2: Multi-point Crossover for binary chromosomes

Example

$$\text{Parent 1: } 11|010|011 \quad (6.5)$$

$$\text{Parent 2: } 00|111|100 \quad (6.6)$$

$$\text{Child 1: } 11|111|011 \quad (6.7)$$

$$\text{Child 2: } 00|010|100 \quad (6.8)$$

Crossover points at positions 2 and 5.

Advantages

- Reduces positional bias
- Can preserve building blocks at chromosome ends
- More disruptive than one-point crossover

6.4.5 Uniform Crossover

Each gene is independently chosen from either parent [40, 15].

In uniform crossover, each gene (bit) is randomly selected from one of the corresponding genes on the parent chromosome. This selection process is done independently for each gene position. This method can be analogized to tossing a coin. If the result is "head", the gene is taken from parent 1; if the result is "tail", the gene is taken from parent 2.

This method provides equal opportunity for each gene to be inherited from one of the parents, thereby increasing genetic diversity and eliminating positional bias that typically appears in single-point crossover or multi-point crossover.

Algorithm

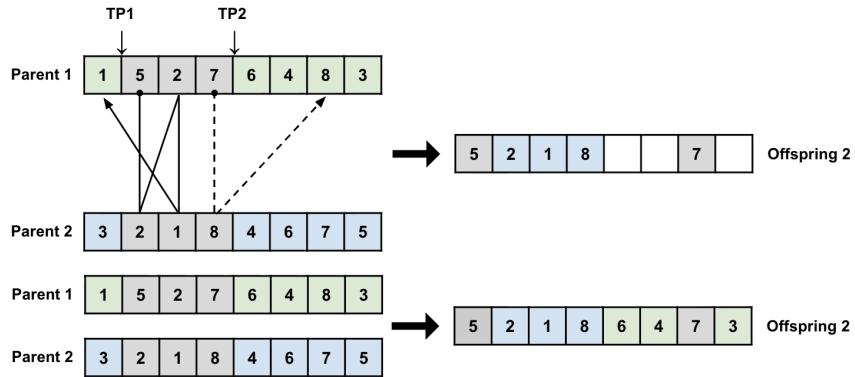


Figure 6.3: Uniform Crossover for binary chromosomes

Algorithm 7 Uniform Crossover

```

for each gene position  $i$  do
    Generate random number  $r \in [0, 1]$ 
    if  $r < 0.5$  then
         $child_1[i] = parent_1[i]$ ,  $child_2[i] = parent_2[i]$ 
    else
         $child_1[i] = parent_2[i]$ ,  $child_2[i] = parent_1[i]$ 
    end if
end for

```

Example with Mask

$$\text{Parent 1: } 11010011 \quad (6.9)$$

$$\text{Parent 2: } 00111100 \quad (6.10)$$

$$\text{Mask: } 10110100 \quad (6.11)$$

$$\text{Child 1: } 10111011 \quad (6.12)$$

$$\text{Child 2: } 01010100 \quad (6.13)$$

Mask bit 1: take from Parent 1, Mask bit 0: take from Parent 2.

Properties

- Maximum disruption potential
- No positional bias
- Good for problems where gene positions are independent
- May destroy long building blocks

6.4.6 Multi-Point Crossover

Generalization with k crossover points.

Characteristics

- $k = 0$: No crossover (copy parents)
- $k = 1$: One-point crossover
- $k = l - 1$: Uniform crossover (in expectation)
- As k increases, approaches uniform crossover

6.5 Integer Chromosome Crossover

Unlike binary chromosomes that use bits 0 and 1, integer representation is more suitable for problems involving discrete parameters or numerical values that have quantitative meaning, such as scheduling, sequencing, or combinatorial optimization.

6.5.1 Single-Point Crossover for Integer

Single-Point Crossover is the simplest form of crossover. One crossover position is randomly selected, then the part of the two parents after that position is exchanged to form two new offspring. The strings resulting from this process have Positional Bias characteristics.

6.5.2 Multi-point Crossover for Integer

Multi-point crossover has a mechanism like single-point crossover; the difference is the number of randomly selected positions. In Multi-point crossover, a number N of positions are randomly selected along the length of the chromosome. These positions are exchanged to form two new offspring.

The objectives of multi-point crossover include increasing genetic diversity in the population, reducing positional bias, and increasing the chances of recombination from various schemas or different solution blocks. However, using too many cut points can become too disruptive, as it can damage already good gene combinations (coadapted alleles).

6.5.3 Uniform Crossover for Integer

In uniform crossover, each gene is randomly selected from one of the corresponding genes on the parent chromosome. This selection process is done independently for each gene position. This method can be analogized to tossing a coin. If the result is "head", the gene is taken from parent 1; if the result is "tail", the gene is taken from parent 2.

This method provides equal opportunity for each gene to be inherited from one of the parents, thereby increasing genetic diversity and eliminating positional bias that typically appears in single-point crossover or multi-point crossover.

6.6 Real-Valued Crossover Operators

Crossover on real chromosomes is a genetic recombination process in Genetic Algorithms applied to chromosomes represented in real number form (floating-point representation). This representation is commonly used to solve continuous optimization problems, where

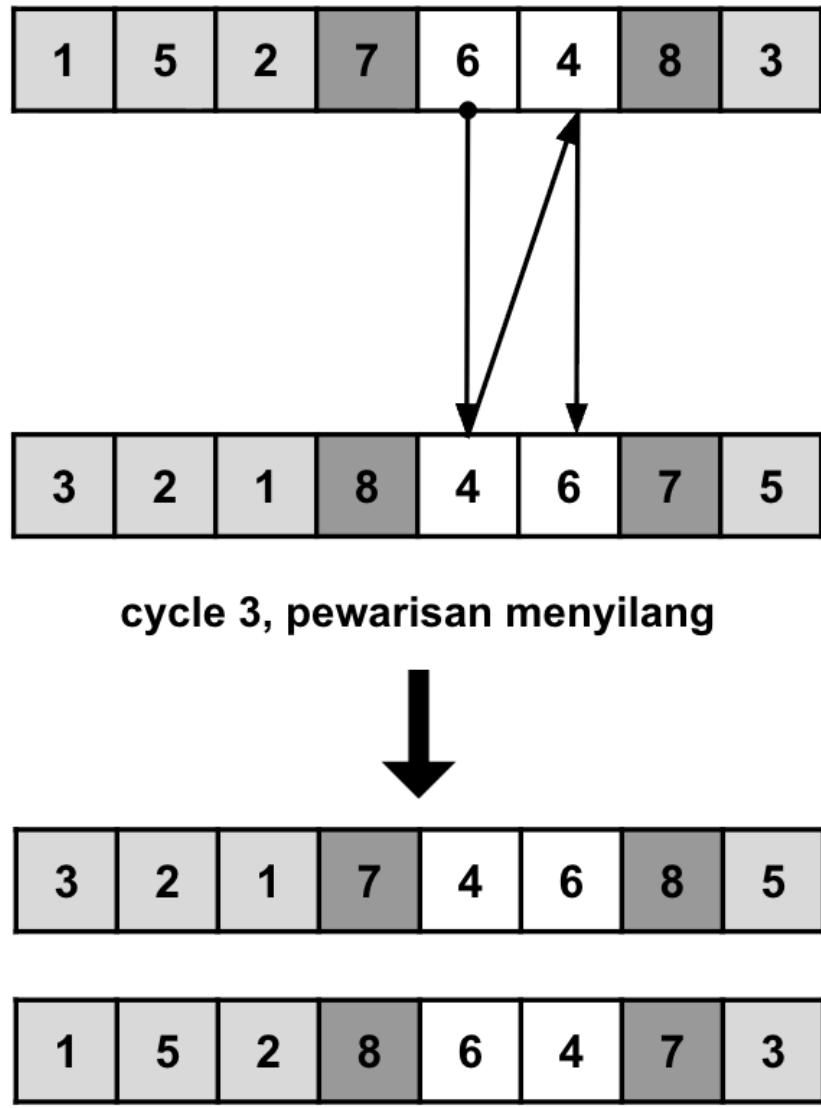


Figure 6.4: Single-Point Crossover for integer chromosomes

decision variables have values within a certain range and are not limited to integers or binary.

Unlike crossover on binary or integer chromosomes, the crossover mechanism for real chromosomes involves arithmetic operations on gene values between parents. This method allows the creation of offspring with gene values that are between or around the parent gene values, thus maintaining the continuity and stability of the evolution process.

6.6.1 Arithmetic Crossover

Linear combination of parent vectors.

Single Arithmetic Crossover

- Two parents are represented as:

- Parent 1: $\langle x_1, \dots, x_n \rangle$
- Parent 2: $\langle y_1, \dots, y_n \rangle$

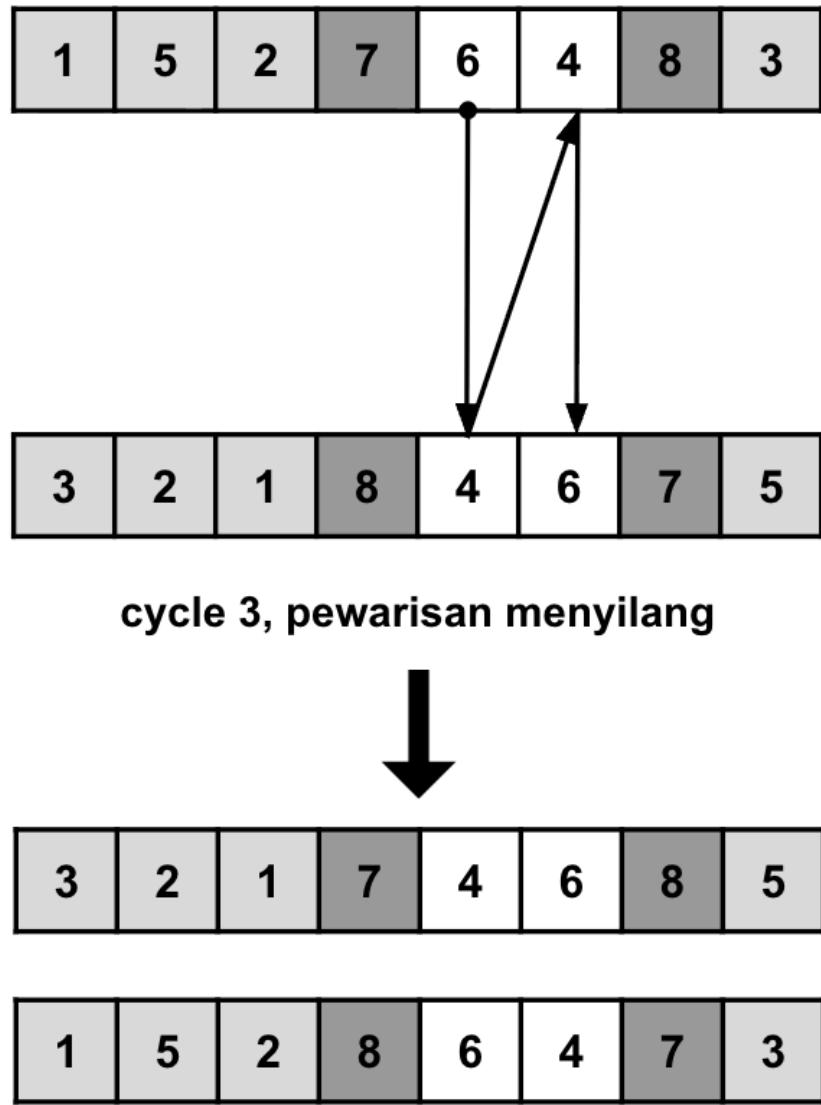


Figure 6.5: Multi-point Crossover for integer chromosomes

2. Randomly select one gene (k) to undergo crossover operation
3. The result is two offspring formed based on a linear combination of the k -th gene of those parents with control parameter α , where $0 \leq \alpha \leq 1$:
 - Offspring 1: $\langle x_1, \dots, x_{k-1}, \alpha \cdot y_k + (1 - \alpha) \cdot x_k, x_{k+1}, \dots, x_n \rangle$
 - Offspring 2: $\langle y_1, \dots, y_{k-1}, \alpha \cdot x_k + (1 - \alpha) \cdot y_k, y_{k+1}, \dots, y_n \rangle$

Simple Arithmetic Crossover

1. Two parents are represented as:
 - Parent 1: $\langle x_1, \dots, x_n \rangle$
 - Parent 2: $\langle y_1, \dots, y_n \rangle$
2. Randomly select one gene (k) to become the crossover boundary point

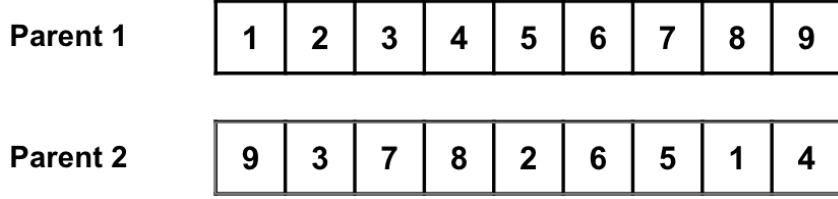


Figure 6.6: Uniform Crossover for integer chromosomes

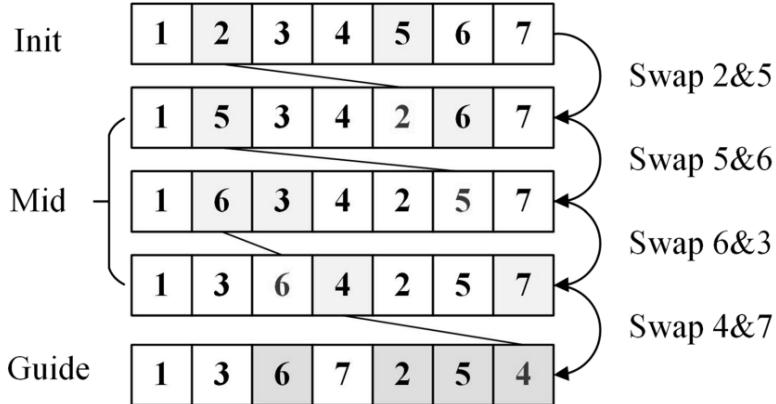


Figure 6.7: Single Arithmetic Crossover for real chromosomes

3. The result is two offspring formed based on a linear combination from gene $k+1$ to gene n with control parameter α , where $0 \leq \alpha \leq 1$:

- Offspring 1: $\langle x_1, \dots, x_k, \alpha \cdot y_{k+1} + (1 - \alpha) \cdot x_{k+1}, \dots, \alpha \cdot y_n + (1 - \alpha) \cdot x_n \rangle$
- Offspring 2: $\langle y_1, \dots, y_k, \alpha \cdot x_{k+1} + (1 - \alpha) \cdot y_{k+1}, \dots, \alpha \cdot x_n + (1 - \alpha) \cdot y_n \rangle$

Parent1	4	7	1	5	6	9	0	2	8	3
Parent2	5	1	8	9	2	7	0	4	6	3
Parent3	6	1	2	4	8	0	3	5	7	9
Parent4	5	9	6	1	2	4	8	7	3	0
Parent5	7	9	0	3	1	8	6	4	2	5
Child	4									
Parent1	4	7	1	5	6	9	0	2	8	3
Parent2	5	1	8	9	2	7	0	4	6	3
Parent3	6	1	2	4	8	0	3	5	7	9
Parent4	5	9	6	1	2	4	8	7	3	0
Parent5	7	9	0	3	1	8	6	4	2	5
Child	4	7	1	5	6	9	0	2	8	3

Figure 6.8: Simple Arithmetic Crossover for real chromosomes

Whole Arithmetic Crossover

1. Two parents are represented as:
 - Parent 1: $\langle x_1, \dots, x_n \rangle$
 - Parent 2: $\langle y_1, \dots, y_n \rangle$
2. For each gene i ($i = 1, 2, \dots, n$), offspring are formed with a linear combination of genes from both parents with control parameter α , where $0 \leq \alpha \leq 1$:

- Offspring 1: $z_i^1 = \alpha \cdot y_i + (1 - \alpha) \cdot x_i$
- Offspring 2: $z_i^2 = \alpha \cdot x_i + (1 - \alpha) \cdot y_i$

Parent1	4	7	1	5	6	9	0	2	8	3	4	7	1	5	6	9	0	2	8	3	4	7	1	5	6	9	0	2	8	3
Parent2	3	5	1	8	9	2	7	0	4	6	3	5	1	8	9	2	7	0	4	6	3	5	1	8	9	2	7	0	4	6
Parent3	6	1	2	4	8	0	3	5	7	9	6	1	2	4	8	0	3	5	7	9	6	1	2	4	8	0	3	5	7	9
Parent4	5	9	6	1	2	4	8	7	3	0	5	9	6	1	2	4	8	7	3	0	5	9	6	1	2	4	8	7	3	0
Parent5	7	9	0	3	1	8	6	4	2	5	7	9	0	3	1	8	6	4	2	5	7	9	0	3	1	8	6	4	2	5
Child	4										4	7									4	7	1	5	6					

Parent1	4	7	1	5	6	9	0	2	8	3	4	7	1	5	6	9	0	2	8	3	4	7	1	5	6	9	0	2	8	3
Parent2	3	5	1	8	9	2	7	0	4	6	3	5	1	8	9	2	7	0	4	6	3	5	1	8	9	2	7	0	4	6
Parent3	6	1	2	4	8	0	3	5	7	9	6	1	2	4	8	0	3	5	7	9	6	1	2	4	8	0	3	5	7	9
Parent4	5	9	6	1	2	4	8	7	3	0	5	9	6	1	2	4	8	7	3	0	5	9	6	1	2	4	8	7	3	0
Parent5	7	9	0	3	1	8	6	4	2	5	7	9	0	3	1	8	6	4	2	5	7	9	0	3	1	8	6	4	2	5
Child	4										4	7	1								4	7	1	5	6					

Figure 6.9: Whole Arithmetic Crossover for real chromosomes

Whole Arithmetic Crossover

$$\text{child}_1 = \alpha \text{parent}_1 + (1 - \alpha) \text{parent}_2 \quad (6.14)$$

$$\text{child}_2 = (1 - \alpha) \text{parent}_1 + \alpha \text{parent}_2 \quad (6.15)$$

where $\alpha \in [0, 1]$ is a random weight.

Simple Arithmetic Crossover

Apply arithmetic crossover to a random subset of genes.

Single Arithmetic Crossover

Apply arithmetic crossover to one randomly selected gene.

Example

$$\text{Parent 1: } (2.1, 5.7, 1.3, 8.9) \quad (6.16)$$

$$\text{Parent 2: } (4.2, 3.1, 6.8, 2.4) \quad (6.17)$$

$$\text{Child 1 } (\alpha = 0.3): \quad (3.57, 4.49, 4.98, 4.17) \quad (6.18)$$

$$\text{Child 2 } (\alpha = 0.3): \quad (2.73, 4.32, 2.98, 6.17) \quad (6.19)$$

6.6.2 BLX- α Crossover (Blend Crossover)

Creates offspring in an interval around the parents.

Algorithm

For each gene i :

1. Calculate $c_{min} = \min(\text{parent}_{1i}, \text{parent}_{2i})$
2. Calculate $c_{max} = \max(\text{parent}_{1i}, \text{parent}_{2i})$
3. Calculate interval $I = c_{max} - c_{min}$
4. Generate offspring in $[c_{min} - \alpha \cdot I, c_{max} + \alpha \cdot I]$

Parameters

- $\alpha = 0$: Offspring between parents
- $\alpha = 0.5$: Standard BLX-0.5
- Larger α : More exploration beyond parents

6.6.3 SBX (Simulated Binary Crossover)

Simulates the behavior of one-point crossover for real-valued genes.

Formula

$$child_{1i} = 0.5[(1 + \beta_i)parent_{1i} + (1 - \beta_i)parent_{2i}] \quad (6.20)$$

$$child_{2i} = 0.5[(1 - \beta_i)parent_{1i} + (1 + \beta_i)parent_{2i}] \quad (6.21)$$

where β_i is calculated from:

$$\beta_i = \begin{cases} (2u_i)^{1/(\eta_c+1)} & \text{if } u_i \leq 0.5 \\ \left(\frac{1}{2(1-u_i)}\right)^{1/(\eta_c+1)} & \text{if } u_i > 0.5 \end{cases} \quad (6.22)$$

$u_i \sim U[0, 1]$ and η_c is the distribution index.

6.7 Permutation Crossover Operators

6.7.1 Order Crossover (OX)

Preserves relative order of elements from one parent [34, 28].

Algorithm

Algorithm 8 Order Crossover

Select two random crossover points

Copy segment between points from Parent 1 to Child

Fill remaining positions with elements from Parent 2 in order they appear, skipping already placed elements

Example

$$\text{Parent 1: } (1, 2, 3, 4, 5, 6, 7, 8, 9) \quad (6.23)$$

$$\text{Parent 2: } (9, 3, 7, 8, 2, 6, 5, 1, 4) \quad (6.24)$$

$$\text{Copy segment: } (-, -, 3, 4, 5, 6, -, -, -) \quad (6.25)$$

$$\text{Fill from P2: } (7, 8, 3, 4, 5, 6, 2, 1, 9) \quad (6.26)$$

6.7.2 Partially Mapped Crossover (PMX)

Creates mapping between elements in the crossover segment [21, 28].

Algorithm

Algorithm 9 Partially Mapped Crossover

Select two crossover points

Exchange segments between parents

For conflicts outside segment, use mapping relationship to resolve

Example

$$\text{Parent 1: } (1, 2, 3, 4, 5, 6, 7, 8, 9) \quad (6.27)$$

$$\text{Parent 2: } (5, 4, 6, 9, 2, 3, 7, 1, 8) \quad (6.28)$$

$$\text{Mapping: } 3 \leftrightarrow 6, 4 \leftrightarrow 9, 5 \leftrightarrow 2, 6 \leftrightarrow 3 \quad (6.29)$$

$$\text{Child 1: } (1, 5, 6, 9, 2, 3, 7, 8, 4) \quad (6.30)$$

6.7.3 Cycle Crossover (CX)

Preserves absolute positions of elements from both parents [34, 35].

Algorithm

Algorithm 10 Cycle Crossover

Start with first element of Parent 1

Follow cycle: find element in Parent 2 at same position, locate it in Parent 1, repeat

Copy cycle elements from Parent 1

Copy non-cycle elements from Parent 2

Create second child by swapping parent roles

6.7.4 Edge Recombination Crossover

Preserves edge information from both parents (useful for TSP).

Algorithm

6.8 Crossover Analysis

6.8.1 Schema Disruption

The probability that a schema H is disrupted by crossover:

Algorithm 11 Edge Recombination

Create edge table from both parents
 Start with element having fewest edges
 Add element to offspring
 Remove element from all edge lists
 Move to element with fewest remaining edges
 If tied, choose randomly
 If no edges remain, choose unused element randomly

One-Point Crossover

$$P_{disruption} = p_c \cdot \frac{\delta(H)}{l - 1} \quad (6.31)$$

Two-Point Crossover

$$P_{disruption} = p_c \cdot \left(\frac{2\delta(H)}{l - 1} - \frac{\delta(H)(\delta(H) - 1)}{(l - 1)(l - 2)} \right) \quad (6.32)$$

Uniform Crossover

$$P_{disruption} = p_c \cdot \left(1 - \left(\frac{1}{2} \right)^{o(H)-1} \right) \quad (6.33)$$

6.8.2 Building Block Preservation

- **Short schemas:** Better preserved by all crossover types
- **Long schemas:** More disrupted, especially by uniform crossover
- **Tightly linked:** Order crossover preserves adjacency relationships

6.9 Advanced Crossover Techniques**6.9.1 Adaptive Crossover**

Adjust crossover parameters based on:

- Population diversity
- Fitness improvement rate
- Generation number
- Individual fitness levels

6.9.2 Multiple Parent Crossover

Combine genetic material from more than two parents [14, 46, 16, 9]:

- **Scanning crossover:** Scan through multiple parents
- **Voting crossover:** Majority vote among parents
- **Averaging crossover:** Average values from multiple parents

6.9.3 Problem-Specific Crossover

Design crossover operators for specific problem domains:

- **Graph problems:** Preserve graph properties
- **Scheduling:** Maintain temporal constraints
- **Neural networks:** Preserve network topology

6.10 Crossover Guidelines

6.10.1 Choosing Crossover Type

- **Binary representation:** One-point, two-point, or uniform
- **Real-valued:** Arithmetic, BLX- α , or SBX
- **Permutation:** OX, PMX, or CX depending on problem structure
- **Variable-length:** Specialized operators required

6.10.2 Parameter Setting

- **Crossover rate:** Start with $p_c = 0.8 - 0.9$
- **Population size:** Larger populations can handle higher crossover rates
- **Problem difficulty:** Harder problems may need lower rates

6.10.3 Empirical Testing

- Test multiple crossover operators
- Vary crossover parameters
- Measure diversity and convergence
- Consider problem-specific metrics

6.11 Crossover vs. Mutation

Aspect	Crossover	Mutation
Primary function	Exploitation	Exploration
Information source	Multiple parents	Random changes
Building blocks	Combines existing	Creates new
Search behavior	Convergent	Divergent
Application rate	High (0.6-0.9)	Low (0.001-0.1)
Population effect	Homogenization	Diversification

Table 6.1: Crossover vs. Mutation Comparison [40, 36]

6.12 Chapter Summary

This chapter covered crossover operators for genetic algorithms across different representation types. Crossover is the primary exploitative operator that combines beneficial traits from multiple parents. The choice of crossover operator depends on the representation and problem characteristics. Proper balance between crossover and mutation is essential for effective genetic algorithm performance.

6.13 Key Concepts

- Crossover principles and biological inspiration
- Binary crossover: one-point, two-point, uniform
- Real-valued crossover: arithmetic, BLX- α , SBX
- Permutation crossover: OX, PMX, CX
- Schema disruption analysis
- Building block preservation
- Adaptive and problem-specific crossover
- Guidelines for crossover selection and parameter setting

Chapter 7

Mutation and Generation Update

In the previous chapters, we have covered the fundamental operations of Genetic Algorithms (GA) including encoding, fitness evaluation, selection, and crossover. This chapter completes the discussion of GA operators by examining **mutation** and **generation update mechanisms** [2, 43]. These operations are crucial for maintaining genetic diversity and ensuring the algorithm's ability to explore the search space effectively.

7.1 Introduction to Mutation

After the recombination (crossover) stage has been applied to all pairs of chromosomes in the mating pool, producing N chromosomes (where N is the population size), the GA executes the mutation operator on each of these chromosomes. Mutation is a critical operator that:

- Prevents premature convergence to local optima
- Maintains genetic diversity in the population
- Introduces new genetic material that may not have been present in the initial population
- Provides a mechanism for escaping local optima

7.1.1 What is Mutation?

Mutation is the process of changing the value of one or more genes in a genome [21, 31, 6]. More specifically, it involves:

- Changing the allele of a gene at a specific locus with another allele
- Avoiding premature convergence, which is reaching a suboptimal result that is not the global maximum
- Creating offspring that are not necessarily better than their parents

Important Note: The new population resulting from mutation is not guaranteed to be better than the previous population. However, mutation provides the essential mechanism for maintaining diversity and exploring new regions of the search space.

7.1.2 Mutation in Evolutionary Algorithms vs. Biological Evolution

In biological evolution, mutation is typically considered harmful because complex organisms have highly interdependent systems. However, in Evolutionary Algorithms (EAs):

- Mutation can often lead to improvement
- Individual representations in EAs are much simpler than biological organisms
- Mutating a small portion of genes may result in better individuals
- The simplified representation makes beneficial mutations more likely

7.2 Mutation for Different Representations

Many mutation methods have been proposed in the literature [30, 6, 24]. Each method has special characteristics and may only be applicable to certain types of representations. The choice of mutation operator must be compatible with the chromosome encoding scheme.

7.2.1 Mutation for Binary Representation

Binary representation uses the simplest form of mutation: **bit-flip mutation**.

Bit-Flip Mutation

In bit-flip mutation, each bit in the chromosome has a probability P_m (mutation probability) of being flipped:

- $1 \rightarrow 0$
- $0 \rightarrow 1$

Example:

Parent: 1 0 1 1 0 1 0 0
 ^ ^

Offspring: 1 0 0 1 0 0 0 0

In this example, bits at positions 3 and 6 were selected for mutation and flipped.

Algorithm:

Algorithm 12 Bit-Flip Mutation

```

for each gene  $g_i$  in chromosome do
     $r \leftarrow$  random number in  $[0, 1]$ 
    if  $r < P_m$  then
        Flip  $g_i$ : if  $g_i = 1$  then  $g_i \leftarrow 0$ , else  $g_i \leftarrow 1$ 
    end if
end for

```

7.2.2 Mutation for Integer Representation

Integer representations require different mutation strategies. Three common approaches are:

Integer Value Flipping

Uses mathematical operations ($+$, $-$, \times , \div) to change the value of selected genes.

Example:

Parent: 8 3 7 5 2 1 9 4 6
 ^ ^
 Offspring: 8 3 2 5 2 8 9 4 6

The values at positions 3 and 6 were changed using mathematical operations.

Random Value Selection

A selected gene is replaced with a randomly chosen value from the valid range.

Example: If the valid range is [1, 9]:

Parent: 8 3 7 5 2 1 9 4 6
 ^
 Offspring: 8 3 7 9 2 1 9 4 6

Creep Mutation

Adds or subtracts a small random integer value (usually ± 1 or ± 2) to the selected gene.

Example:

Parent: 8 3 7 5 2 1 9 4 6
 ^ ^
 Offspring: 8 4 7 5 2 2 9 4 6

This method makes small, gradual changes and is particularly useful for fine-tuning solutions.

7.2.3 Mutation for Real-Valued Representation

Real-valued representations have different characteristics from binary and integer representations. Values of genes in real representations are continuous, whereas binary and integer representations are discrete. Therefore, real representations require specialized mutation operators.

Uniform Mutation

In uniform mutation, selected genes are replaced with values drawn from a uniform random distribution within the valid range $[a, b]$:

$$x'_i = a + \text{rand}(0, 1) \times (b - a) \quad (7.1)$$

where:

- x'_i is the new gene value
- a and b are the lower and upper bounds
- $\text{rand}(0, 1)$ generates a random number in $[0, 1]$

Non-Uniform Mutation with Fixed Distribution

This is the most commonly used mutation for real-valued representations. It is similar to the creep method for integer representation but uses real-valued additions instead of integer values.

The mutated value is calculated as:

$$x'_i = x_i + \mathcal{N}(0, \sigma^2) \quad (7.2)$$

where:

- x_i is the original gene value
- $\mathcal{N}(0, \sigma^2)$ is a random value from a normal (Gaussian) distribution with mean 0 and variance σ^2
- σ controls the mutation step size

Example:

Parent: 2.45 7.89 3.12 9.01 5.67
 ^

Offspring: 2.45 7.89 3.45 9.01 5.67

7.2.4 Mutation for Permutation Representation

Mutation on permutation representations must ensure that the resulting chromosome remains valid. This means that after mutation, all elements must still appear exactly once. Several specialized methods have been developed:

Swap Mutation

Two gene positions are randomly selected, and their values are exchanged.

Example:

Parent: 3 1 5 2 7 6 8 4 9
 ^ ^

Offspring: 3 1 8 2 7 6 5 4 9

Positions 3 and 7 are selected, so values 5 and 8 are swapped.

Algorithm:

Algorithm 13 Swap Mutation

```
i ← random position in chromosome
j ← random position in chromosome (different from i)
Swap values at positions i and j
```

Insert Mutation

A gene at one position is removed and inserted at another position, shifting the intermediate genes.

Example:

Parent: 3 1 5 2 7 6 8 4 9
 ^ ^
 Offspring: 3 1 5 2 7 8 6 4 9

The gene at position 7 (value 8) is removed and inserted after position 2 (value 5).

Scramble Mutation

A segment of the chromosome is selected, and the genes within that segment are randomly shuffled.

Example:

Parent: 3 1 5 2 7 6 8 4 9
 \-----/
 Offspring: 3 1 2 6 5 7 8 4 9

The segment $\{5, 2, 7, 6\}$ is selected and randomly shuffled to $\{2, 6, 5, 7\}$.

Inversion Mutation

A segment of the chromosome is selected, and the order of genes within that segment is reversed.

Example:

Parent: 3 1 5 2 7 6 8 4 9
 \-----/
 Offspring: 3 1 6 7 2 5 8 4 9

The segment $\{5, 2, 7, 6\}$ is reversed to $\{6, 7, 2, 5\}$.

7.3 Generation Update Mechanisms

After selection, crossover, and mutation operations have been applied to a population, a generation update mechanism determines which individuals survive to the next generation. This process is also called **survivor selection** or **replacement strategy**.

7.3.1 Holland's Original Model (Generational Replacement)

In Holland's original GA [25, 21]:

- All offspring replace the entire parent population
- Parents are considered "dead" and removed from the population
- The new population consists entirely of offspring

- This creates distinct, non-overlapping generations

Characteristics:

- Simple and straightforward
- Clear separation between generations
- May lose good solutions if not careful
- Often combined with elitism to preserve best solutions

7.3.2 Generational Model with Elitism

A population of size N chromosomes in one generation is replaced by N new individuals in the next generation [10, 44]. However, to preserve the best solutions:

- The best k chromosomes (elites) from the parent generation are copied directly to the next generation
- The remaining $N - k$ positions are filled with offspring
- This ensures that the best solution never gets worse across generations

Algorithm:

Algorithm 14 Generational Model with Elitism

Sort parent population by fitness

Copy top k individuals to next generation (elites)

Generate $N - k$ offspring through selection, crossover, and mutation

Add offspring to next generation

Next generation becomes current generation

Typical values: $k = 1$ or $k = 2$ (preserving 1-2 best individuals)

7.3.3 Steady-State Update

In the steady-state model [44, 39]:

- Not all chromosomes are replaced in each generation
- Only M chromosomes are replaced, where $M < N$
- Often $M = 2$ (one mating produces 2 offspring, which replace 2 individuals)

Replacement strategies for selecting which individuals to replace:

1. **Replace parents:** The two offspring replace their two parents
2. **Replace worst:** The two offspring replace the two worst individuals in the population

- 3. Replace oldest:** The two offspring replace the two oldest individuals in the population

Characteristics:

- Allows good individuals to participate in multiple matings
- More gradual evolution
- Parents and offspring coexist in the same population
- Can be more efficient computationally

7.3.4 Continuous Update

In continuous update:

- Offspring and parents can coexist in the same generation
- Individuals are selected randomly from both groups for the next generation
- Provides maximum overlap between generations
- Less commonly used than other methods

7.4 GA Parameters

The performance of a Genetic Algorithm heavily depends on proper parameter settings [22, 36, 10]. The main parameters that need to be configured are:

7.4.1 Crossover Probability (P_c)

P_c is the probability that two parents will undergo crossover.

Effects:

- $P_c = 100\%$: All offspring are produced through crossover
- $P_c = 0\%$: No crossover occurs; offspring are exact copies of parents
- Typical range: $P_c \in [0.65, 0.90]$ (65% to 90%)

Recommendations:

- Higher values (0.8-0.9) encourage exploration
- Lower values preserve good solutions but reduce diversity
- Standard setting: $P_c = 0.8$

7.4.2 Mutation Probability (P_m)

P_m is the probability that a gene in an offspring chromosome will undergo mutation.

Effects:

- $P_m = 100\%$: All genes are mutated (chaos)
- $P_m = 0\%$: No mutation occurs; no new genetic material
- Typical range: $P_m \in [0.005, 0.01]$ (0.5% to 1%)

Common formulas:

$$P_m = \frac{1}{L} \tag{7.3}$$

or

$$P_m = \frac{1}{N \times L} \tag{7.4}$$

where:

- L is the chromosome length (number of genes)
- N is the population size

Rationale: The mutation probability is often set so that, on average, one mutation occurs per chromosome.

7.4.3 Population Size (N)

The population size should be proportional to the volume of the search space.

Effects:

- Too small: Difficult to reach global optimum; may converge to local optimum
- Too large: Heavy computational burden; inconsistent with evolutionary principles
- Should not approach the size of the entire search space

Recommendations:

- Typical range: $N \in [50, 100]$
- Determined through experimentation
- Larger populations for larger, more complex problems
- Consider computational resources available

7.4.4 Number of Generations (G)

The number of generations should be proportional to population size and search space size.

Example calculation:

- If $N = 100$ and search space size $\approx 10^5$
- Then $G = 100$ might be appropriate

Stopping criteria alternatives:

1. Fixed number of generations
2. Maximum number of fitness evaluations
3. No improvement for k consecutive generations
4. Target fitness value reached
5. Combination of the above

7.4.5 General Parameter Setting Guidelines

Important Note: There are no definitive rules for setting GA parameters [45, 22]. Parameter selection relies on:

- Intuition and experience
- Experimentation (trial and error)
- Problem-specific characteristics

Common starting configuration:

- Chromosome representation: Binary/Integer/Real/Permutation (problem-dependent)
- Number of bits per variable: Based on desired precision
- Population size: $N = 50$ to 100
- Crossover probability: $P_c = 0.8$
- Mutation probability: $P_m = \frac{1}{L}$ to $\frac{1}{N \times L}$

where:

- N = Population size
- L = Chromosome length (number of genes)

7.5 Parameter Observation Study

To understand the effects of different parameters, we present a systematic observation study.

7.5.1 Test Problem

Objective: Minimize the function:

$$h(x_1, x_2) = x_1^2 + x_2^2 \quad (7.5)$$

where $x_1, x_2 \in [-10, 10]$

Fitness function:

$$\text{Fitness} = \frac{1}{x_1^2 + x_2^2 + 0.001} \quad (7.6)$$

The constant 0.001 is added to avoid division by zero at the optimal point (0, 0).

7.5.2 Experimental Setup

Parameter variations tested:

- Population size: [50, 100, 200]
- Number of bits per variable: [10, 50, 90]
- Crossover probability: [0.5, 0.7, 0.9]
- Mutation probability: $[0.5/L, 1/L, 2/L]$ where L is total chromosome length

Fairness criterion:

- Maximum number of individuals evaluated: 20,000
- Each configuration run 30 times for statistical validity

7.5.3 Sample Results

Table 7.1 shows selected results from the parameter study:

Table 7.1: GA Parameter Observation Results

Pop Size	Bits	P_c	P_m	Avg Best Fitness	Avg Evaluations
50	10	0.5	0.0250	839.55	20000
50	50	0.5	0.0050	1000.00	8301.67
50	50	0.7	0.0100	1000.00	20000
50	90	0.7	0.0056	1000.00	8780.00
100	50	0.7	0.0050	1000.00	14416.67
100	90	0.5	0.0111	1000.00	20000
200	50	0.5	0.0050	1000.00	20000
200	90	0.7	0.0056	1000.00	20000
200	90	0.9	0.0028	1000.00	19866.67

Key observations:

1. **Best configuration:** Population size = 50, Bits = 90, $P_c = 0.7$, $P_m = 0.0056$
 - Achieved optimal fitness (1000.00)

- Required only 8780 evaluations on average
- Most efficient configuration

2. Effect of bit precision:

- 10 bits: Often failed to reach optimum
- 50-90 bits: Consistently reached optimum
- Higher precision enables finer search granularity

3. Effect of population size:

- Smaller populations (50) can be very efficient
- Larger populations (200) more robust but slower
- Trade-off between speed and reliability

4. Effect of crossover probability:

- $P_c = 0.7$ performed best overall
- Moderate values balance exploration and exploitation

5. Effect of mutation probability:

- Low mutation rates ($\sim 1/L$) worked best
- Too high mutation causes chaos
- Too low mutation loses diversity

7.6 Summary and Conclusions

This chapter has covered the essential components for completing the GA cycle:

1. Mutation operators

- provide genetic diversity and prevent premature convergence:
- Binary: Bit-flip mutation
 - Integer: Flipping, random selection, creep mutation
 - Real: Uniform mutation, Gaussian mutation
 - Permutation: Swap, insert, scramble, inversion mutation

2. Generation update mechanisms

- determine how populations evolve:
- Generational replacement (with elitism)
 - Steady-state update
 - Continuous update

3. Parameter selection

- is crucial for GA performance:
- No universal rules exist
 - Requires experimentation and tuning

- Starting guidelines provide reasonable defaults

Key principles:

- Parent selection and survivor selection do not depend on chromosome representation
- Recombination and mutation operators must match the chromosome representation
- Parameter settings should be tailored to the specific problem
- Experimentation is essential for finding optimal configurations

With the completion of this chapter, we have now covered all the fundamental components of Genetic Algorithms: encoding, fitness evaluation, selection, crossover, mutation, and generation update. The next chapters will explore advanced topics and practical applications of GAs.

7.7 Exercises

- Given the two parent chromosomes for a permutation problem:
 - Parent 1: [1, 2, 7, 3, 4, 9, 8, 6, 5]
 - Parent 2: [5, 4, 3, 9, 1, 2, 6, 8, 7]
 - (a) Perform Partial-Mapped Crossover (PMX) with cut points at positions 2 and 5
 - (b) Apply inversion mutation to the offspring with mutation segment from locus 2 to 5
- For a binary-encoded GA with chromosome length $L = 50$ and population size $N = 100$:
 - (a) Calculate appropriate mutation probability using $P_m = 1/L$
 - (b) Calculate alternative mutation probability using $P_m = 1/(N \times L)$
 - (c) Discuss which might be more appropriate and why
- Design a mutation operator for a real-valued chromosome representing (x, y) coordinates where $x, y \in [-100, 100]$:
 - (a) Implement uniform mutation
 - (b) Implement Gaussian mutation with $\sigma = 5$
 - (c) Compare the expected behavior of both operators
- Implement and compare three generation update strategies:
 - (a) Generational replacement with elitism ($k = 2$)
 - (b) Steady-state with replacement of worst individuals
 - (c) Steady-state with replacement of oldest individuals

Discuss scenarios where each might be preferred.

5. For the test function $f(x_1, x_2) = x_1^2 + x_2^2$ with $x_1, x_2 \in [-10, 10]$:
 - (a) Design a complete GA including all parameters
 - (b) Run experiments with different parameter combinations
 - (c) Analyze which parameters have the most significant impact
 - (d) Propose an optimal parameter configuration based on your results

Chapter 8

Real-World Applications and Visual Examples

This chapter showcases real-world applications of genetic algorithms, directly from the course materials, demonstrating how GA concepts are applied in practice [20, 38, 13].

8.1 Game AI and Entertainment

8.1.1 Super Mario Bros Level Learning

One of the most compelling demonstrations of genetic algorithms is their application to game AI. In the course materials, we see an example of a genetic machine learning algorithm beating the first level of Super Mario Bros World at 4x speed.

Figure 8.1: Genetic Algorithm Learning to Play Super Mario Bros at 4x Speed

The GA evolves strategies by:

- **Encoding:** Button sequences as chromosomes (jump, run, duck, etc.)
- **Fitness:** Distance traveled and completion time
- **Selection:** Best-performing sequences survive
- **Crossover:** Combining successful movement patterns
- **Mutation:** Random button variations to explore new strategies [32, 27]

8.1.2 Tower Defense Game Balancing

The Towers of Reus project demonstrates how GA can balance gameplay:

- Users create maps with adjustable balancing parameters
- GA component runs until finding optimal winning solutions
- Determines if towers are too strong/weak or if levels are beatable
- Players can then test their skills against the optimized challenge

Figure 8.2: Cat Navigating Circular Maze to Reach Cheese Using Genetic Algorithm

8.2 Pathfinding and Navigation

8.2.1 Maze Navigation

This example demonstrates:

- **Problem:** Find shortest path through complex maze
- **Encoding:** Sequence of movement directions (up, down, left, right) [28]
- **Fitness:** Inverse of path length plus penalty for hitting walls
- **Crossover:** Combining successful path segments

8.2.2 Robot Navigation

Physical robot navigation showcases GA in hardware applications:

- Real-time path planning in dynamic environments
- Sensor data integration for obstacle avoidance
- Adaptive behavior evolution based on environmental feedback

8.3 Evolution Simulation

8.3.1 Simulated Evolution of Creatures

The course references simulated evolution examples from <http://www.wreck.devisland.net/ga/>:

Figure 8.3: Simulated Evolution Using Genetic Algorithm - Creatures Adapting Over Generations

Features include:

- Morphology evolution (body structure)
- Locomotion pattern optimization
- Environmental adaptation
- Multi-objective fitness (speed, stability, efficiency) [11, 12, 26]

Figure 8.4: Human Movement Evolution: From Sitting to Athletic Performance

8.4 Human Analogy Examples

8.4.1 Evolution of Movement

This analogy illustrates:

- **Population:** Different individuals with varying abilities
- **Selection:** Those who can jump higher survive
- **Inheritance:** Athletic traits passed to next generation
- **Mutation:** Random variations in technique

8.4.2 Work Journey Optimization

Figure 8.5: Optimizing Daily Commute Route Using GA Principles

Real-world application:

- Multiple route options (genes)
- Traffic conditions as environmental factors
- Time and fuel consumption as fitness criteria
- Learning from daily experiences (generations)

8.5 Academic Context

8.5.1 GA in Computational Intelligence

Figure 8.6: Position of Genetic Algorithms in Machine Learning and Soft Computing Landscape

The diagram shows GA's relationship with:

- **Machine Learning:** Kernel methods, SVM, Hidden Markov, Bayesian methods
- **Soft Computing:** Neural Networks, Fuzzy systems
- **Intersection:** Reinforcement Learning combining multiple paradigms

Figure 8.7: Comparison of Lamarck vs Darwin-Wallace Evolution Theories Using Giraffe Example

8.6 Historical Perspective

8.6.1 Natural Selection Theories

Understanding evolutionary principles:

- **Lamarck's View:** Acquired characteristics inherited
- **Darwin-Wallace View:** Natural selection favors beneficial traits
- **GA Implementation:** Follows Darwinian principles with random variation and selection

8.7 Summary

These visual examples from the course materials demonstrate the wide applicability of genetic algorithms:

1. **Entertainment:** Game AI and procedural content generation
2. **Robotics:** Path planning and adaptive behavior
3. **Simulation:** Artificial life and evolution studies
4. **Optimization:** Route planning and resource allocation
5. **Research:** Understanding natural evolutionary processes

The key insight is that GA provides a unified framework for solving complex optimization problems across diverse domains, making it one of the most versatile tools in computational intelligence.

Appendix A

Algorithm Implementations

A.1 Basic Genetic Algorithm Implementation

A.1.1 Python Implementation

Listing A.1: Basic Genetic Algorithm in Python

```
1 import numpy as np
2 import matplotlib.pyplot as plt
3 from typing import List, Tuple, Callable
4
5 class GeneticAlgorithm:
6     def __init__(self,
7                  fitness_func: Callable,
8                  chromosome_length: int,
9                  population_size: int = 100,
10                 crossover_rate: float = 0.8,
11                 mutation_rate: float = 0.01,
12                 elitism: bool = True):
13
14         self.fitness_func = fitness_func
15         self.chromosome_length = chromosome_length
16         self.population_size = population_size
17         self.crossover_rate = crossover_rate
18         self.mutation_rate = mutation_rate
19         self.elitism = elitism
20
21     # Initialize population
22     self.population = self._initialize_population()
23     self.fitness_history = []
24     self.best_individual = None
25     self.best_fitness = float('-inf')
26
27     def _initialize_population(self) -> np.ndarray:
28         """Initialize random binary population"""
29         return np.random.randint(0, 2,
30                                (self.population_size, self.
31                                         chromosome_length))
32
33     def _evaluate_fitness(self, population: np.ndarray) -> np.
34         ndarray:
35         """Evaluate fitness for all individuals"""
36         fitness_values = np.array([self.fitness_func(individual)
37                                    for individual in population])
38
39         return fitness_values
```

```

37
38     def _tournament_selection(self, population: np.ndarray,
39                               fitness_values: np.ndarray,
40                               tournament_size: int = 3) -> np.
41                               ndarray:
42         """Tournament selection"""
43         selected = []
44         for _ in range(len(population)):
45             # Select random individuals for tournament
46             tournament_indices = np.random.choice(len(population),
47                                                     ,
48                                                     tournament_size,
49                                                     replace=False)
50             tournament_fitness = fitness_values[
51                 tournament_indices]
52
53             # Select winner
54             winner_index = tournament_indices[np.argmax(
55                 tournament_fitness)]
56             selected.append(population[winner_index])
57
58         return np.array(selected)
59
60     def _one_point_crossover(self, parent1: np.ndarray,
61                             parent2: np.ndarray) -> Tuple[np.
62                                             ndarray, np.ndarray]:
63         """One-point crossover"""
64         if np.random.random() > self.crossover_rate:
65             return parent1.copy(), parent2.copy()
66
67         crossover_point = np.random.randint(1, len(parent1))
68
69         child1 = np.concatenate([parent1[:crossover_point],
70                                 parent2[crossover_point:]])
71         child2 = np.concatenate([parent2[:crossover_point],
72                                 parent1[crossover_point:]])
73
74         return child1, child2
75
76     def _bit_flip_mutation(self, individual: np.ndarray) -> np.
77                           ndarray:
78         """Bit-flip mutation"""
79         mutated = individual.copy()
80         for i in range(len(mutated)):
81             if np.random.random() < self.mutation_rate:
82                 mutated[i] = 1 - mutated[i] # Flip bit
83
84         return mutated
85
86     def _apply_elitism(self, old_population: np.ndarray,
87                       old_fitness: np.ndarray,
88                       new_population: np.ndarray) -> np.ndarray:

```

```
82     """Apply elitism by preserving best individual"""
83     if not self.elitism:
84         return new_population
85
86     best_index = np.argmax(old_fitness)
87     best_individual = old_population[best_index]
88
89     # Replace worst individual in new population with best
90     # from old
91     new_fitness = self._evaluate_fitness(new_population)
92     worst_index = np.argmin(new_fitness)
93     new_population[worst_index] = best_individual
94
95     return new_population
96
97 def evolve(self, generations: int) -> dict:
98     """Main evolutionary loop"""
99     for generation in range(generations):
100         # Evaluate fitness
101         fitness_values = self._evaluate_fitness(self.
102             population)
103
104         # Track best individual
105         max_fitness_idx = np.argmax(fitness_values)
106         if fitness_values[max_fitness_idx] > self.
107             best_fitness:
108                 self.best_fitness = fitness_values[
109                     max_fitness_idx]
110                 self.best_individual = self.population[
111                     max_fitness_idx].copy()
112
113         # Record statistics
114         self.fitness_history.append({
115             'generation': generation,
116             'best_fitness': np.max(fitness_values),
117             'avg_fitness': np.mean(fitness_values),
118             'worst_fitness': np.min(fitness_values)
119         })
120
121         # Selection
122         selected = self._tournament_selection(self.population
123             , fitness_values)
124
125         # Crossover and mutation
126         new_population = []
127         for i in range(0, len(selected), 2):
128             parent1 = selected[i]
129             parent2 = selected[(i + 1) % len(selected)]
130
131             # Crossover
```

```

126         child1, child2 = self._one_point_crossover(
127             parent1, parent2)
128
129         # Mutation
130         child1 = self._bit_flip_mutation(child1)
131         child2 = self._bit_flip_mutation(child2)
132
133         new_population.extend([child1, child2])
134
135         new_population = np.array(new_population[:self.
136             population_size])
137
138         # Apply elitism
139         self.population = self._apply_elitism(self.population
140             ,
141             fitness_values,
142             new_population)
143
144         return {
145             'best_individual': self.best_individual,
146             'best_fitness': self.best_fitness,
147             'fitness_history': self.fitness_history
148         }
149
150     def plot_fitness_history(self):
151         """Plot fitness evolution over generations"""
152         generations = [entry['generation'] for entry in self.
153             fitness_history]
154         best_fitness = [entry['best_fitness'] for entry in self.
155             fitness_history]
156         avg_fitness = [entry['avg_fitness'] for entry in self.
157             fitness_history]
158
159         plt.figure(figsize=(10, 6))
160         plt.plot(generations, best_fitness, label='BestFitness',
161             linewidth=2)
162         plt.plot(generations, avg_fitness, label='AverageFitness',
163             , linewidth=2)
164         plt.xlabel('Generation')
165         plt.ylabel('Fitness')
166         plt.title('FitnessEvolution')
167         plt.legend()
168         plt.grid(True, alpha=0.3)
169         plt.show()
170
171     # Example usage
172     def onemax_fitness(individual):
173         """OneMax problem: maximize number of 1s"""
174         return np.sum(individual)
175
176     def sphere_function_binary(individual, bounds=(-5.12, 5.12)):
```

```

169     """Sphere function with binary encoding"""
170     # Decode binary to real values
171     x = bounds[0] + (bounds[1] - bounds[0]) * np.sum(individual *
172         2**np.arange(len(individual))[:-1]) / (2**len(individual)
173         - 1)
174     return -(x**2) # Negative because we want to minimize
175
176 # Run GA on OneMax problem
177 if __name__ == "__main__":
178     ga = GeneticAlgorithm(
179         fitness_func=onemax_fitness,
180         chromosome_length=20,
181         population_size=50,
182         crossover_rate=0.8,
183         mutation_rate=0.01
184     )
185
186     result = ga.evolve(generations=100)
187
188     print(f"Best individual: {result['best_individual']}")
189     print(f"Best fitness: {result['best_fitness']}")
190
191     ga.plot_fitness_history()

```

A.2 Real-Valued Genetic Algorithm

Listing A.2: Real-Valued GA Implementation

```

1 import numpy as np
2 from typing import List, Tuple, Callable
3
4 class RealValuedGA:
5     def __init__(self,
6                  fitness_func: Callable,
7                  dimensions: int,
8                  bounds: List[Tuple[float, float]],
9                  population_size: int = 100,
10                 crossover_rate: float = 0.8,
11                 mutation_rate: float = 0.1,
12                 mutation_strength: float = 0.1):
13
14         self.fitness_func = fitness_func
15         self.dimensions = dimensions
16         self.bounds = bounds
17         self.population_size = population_size
18         self.crossover_rate = crossover_rate
19         self.mutation_rate = mutation_rate
20         self.mutation_strength = mutation_strength
21
22         self.population = self._initialize_population()

```

```

23         self.fitness_history = []
24
25     def _initialize_population(self) -> np.ndarray:
26         """Initialize random real-valued population"""
27         population = np.zeros((self.population_size, self.
28             dimensions))
29         for i in range(self.dimensions):
30             low, high = self.bounds[i]
31             population[:, i] = np.random.uniform(low, high, self.
32                 population_size)
33         return population
34
35     def _blx_alpha_crossover(self, parent1: np.ndarray,
36                             parent2: np.ndarray,
37                             alpha: float = 0.5) -> Tuple[np.
38                                 ndarray, np.ndarray]:
39         """BLX-alpha crossover"""
40         if np.random.random() > self.crossover_rate:
41             return parent1.copy(), parent2.copy()
42
43         child1 = np.zeros_like(parent1)
44         child2 = np.zeros_like(parent2)
45
46         for i in range(len(parent1)):
47             min_val = min(parent1[i], parent2[i])
48             max_val = max(parent1[i], parent2[i])
49             interval = max_val - min_val
50
51             low_bound = max(min_val - alpha * interval, self.
52                             bounds[i][0])
53             high_bound = min(max_val + alpha * interval, self.
54                             bounds[i][1])
55
56             child1[i] = np.random.uniform(low_bound, high_bound)
57             child2[i] = np.random.uniform(low_bound, high_bound)
58
59         return child1, child2
60
61     def _gaussian_mutation(self, individual: np.ndarray) -> np.
62         ndarray:
63         """Gaussian mutation"""
64         mutated = individual.copy()
65         for i in range(len(mutated)):
66             if np.random.random() < self.mutation_rate:
67                 noise = np.random.normal(0, self.
68                     mutation_strength)
69                 mutated[i] += noise
70
71                 # Ensure bounds are respected
72                 low, high = self.bounds[i]
73                 mutated[i] = np.clip(mutated[i], low, high)

```

```

67         return mutated
68
69
70     def evolve(self, generations: int) -> dict:
71         """Main evolutionary loop"""
72         for generation in range(generations):
73             # Evaluate fitness
74             fitness_values = np.array([self.fitness_func(ind)
75                                         for ind in self.population])
76
77             # Record statistics
78             self.fitness_history.append({
79                 'generation': generation,
80                 'best_fitness': np.max(fitness_values),
81                 'avg_fitness': np.mean(fitness_values),
82                 'worst_fitness': np.min(fitness_values)
83             })
84
85             # Tournament selection
86             new_population = []
87             for _ in range(self.population_size // 2):
88                 # Select parents
89                 parent1_idx = self._tournament_selection(
90                     fitness_values)
91                 parent2_idx = self._tournament_selection(
92                     fitness_values)
93
94                 parent1 = self.population[parent1_idx]
95                 parent2 = self.population[parent2_idx]
96
97                 # Crossover
98                 child1, child2 = self._blx_alpha_crossover(
99                     parent1, parent2)
100
101                 # Mutation
102                 child1 = self._gaussian_mutation(child1)
103                 child2 = self._gaussian_mutation(child2)
104
105                 new_population.extend([child1, child2])
106
107             self.population = np.array(new_population)
108
109             # Final evaluation
110             final_fitness = np.array([self.fitness_func(ind)
111                                         for ind in self.population])
112             best_idx = np.argmax(final_fitness)
113
114             return {
115                 'best_individual': self.population[best_idx],
116                 'best_fitness': final_fitness[best_idx],
117                 'fitness_history': self.fitness_history
118             }

```

```

115     }
116
117     def _tournament_selection(self, fitness_values: np.ndarray,
118                               tournament_size: int = 3) -> int:
119         """Tournament selection returning index"""
120         tournament_indices = np.random.choice(len(fitness_values),
121                                               tournament_size,
122                                               replace=False)
123         tournament_fitness = fitness_values[tournament_indices]
124         winner_idx = tournament_indices[np.argmax(
125             tournament_fitness)]
126         return winner_idx
127
128 # Example: Optimize Rastrigin function
129 def rastrigin_function(x):
130     """Rastrigin function (minimization problem)"""
131     A = 10
132     n = len(x)
133     return -(A * n + np.sum(x**2 - A * np.cos(2 * np.pi * x)))
134
135 # Usage
136 bounds = [(-5.12, 5.12)] * 2 # 2D Rastrigin
137 ga = RealValuedGA(
138     fitness_func=rastrigin_function,
139     dimensions=2,
140     bounds=bounds,
141     population_size=100,
142     mutation_strength=0.1
143 )
144
145 result = ga.evolve(generations=200)
146 print(f"Best solution: {result['best_individual']}")
147 print(f"Best fitness: {result['best_fitness']}")

```

A.3 Traveling Salesman Problem GA

Listing A.3: TSP with Genetic Algorithm

```

1 import numpy as np
2 import matplotlib.pyplot as plt
3 from typing import List, Tuple
4
5 class TSP_GA:
6     def __init__(self,
7                  cities: np.ndarray,
8                  population_size: int = 100,
9                  crossover_rate: float = 0.8,
10                 mutation_rate: float = 0.02):
11

```

```
12     self.cities = cities
13     self.num_cities = len(cities)
14     self.population_size = population_size
15     self.crossover_rate = crossover_rate
16     self.mutation_rate = mutation_rate
17
18     # Create distance matrix
19     self.distance_matrix = self._calculate_distance_matrix()
20
21     # Initialize population
22     self.population = self._initialize_population()
23
24     def _calculate_distance_matrix(self) -> np.ndarray:
25         """Calculate distance matrix between all cities"""
26         n = self.num_cities
27         distances = np.zeros((n, n))
28
29         for i in range(n):
30             for j in range(n):
31                 if i != j:
32                     distances[i][j] = np.sqrt(
33                         (self.cities[i][0] - self.cities[j][0]) **2 +
34                         (self.cities[i][1] - self.cities[j][1]) **2
35                     )
36         return distances
37
38     def _initialize_population(self) -> List[List[int]]:
39         """Initialize population with random permutations"""
40         population = []
41         for _ in range(self.population_size):
42             tour = list(range(self.num_cities))
43             np.random.shuffle(tour)
44             population.append(tour)
45         return population
46
47     def _calculate_tour_distance(self, tour: List[int]) -> float:
48         """Calculate total distance of a tour"""
49         total_distance = 0
50         for i in range(len(tour)):
51             from_city = tour[i]
52             to_city = tour[(i + 1) % len(tour)]
53             total_distance += self.distance_matrix[from_city][
54                 to_city]
55         return total_distance
56
57     def _fitness(self, tour: List[int]) -> float:
58         """Fitness function (inverse of distance)"""
59         distance = self._calculate_tour_distance(tour)
60         return 1.0 / (1.0 + distance)
```

```

60
61     def _order_crossover(self, parent1: List[int],
62                           parent2: List[int]) -> Tuple[List[int],
63                                         List[int]]:
64         """Order crossover (OX)"""
65         if np.random.random() > self.crossover_rate:
66             return parent1.copy(), parent2.copy()
67
68         size = len(parent1)
69         start, end = sorted(np.random.choice(size, 2, replace=False))
70
71         # Create children
72         child1 = [None] * size
73         child2 = [None] * size
74
75         # Copy segments
76         child1[start:end] = parent1[start:end]
77         child2[start:end] = parent2[start:end]
78
79         # Fill remaining positions
80         self._fill_remaining_ox(child1, parent2, start, end)
81         self._fill_remaining_ox(child2, parent1, start, end)
82
83         return child1, child2
84
85     def _fill_remaining_ox(self, child: List[int], parent: List[int],
86                           start: int, end: int):
87         """Helper function for order crossover"""
88         child_set = set(child[start:end])
89         parent_filtered = [city for city in parent if city not in
90                            child_set]
91
92         # Fill positions before start
93         for i in range(start):
94             child[i] = parent_filtered.pop(0)
95
96         # Fill positions after end
97         for i in range(end, len(child)):
98             child[i] = parent_filtered.pop(0)
99
100    def _swap_mutation(self, tour: List[int]) -> List[int]:
101        """Swap mutation"""
102        mutated = tour.copy()
103        if np.random.random() < self.mutation_rate:
104            i, j = np.random.choice(len(tour), 2, replace=False)
105            mutated[i], mutated[j] = mutated[j], mutated[i]
106
107        return mutated
108
109    def _tournament_selection(self, fitness_values: List[float],
110                             tournament_size: int) -> List[int]:
111        """Tournament selection"""
112        selected_tour = []
113
114        while len(selected_tour) < len(fitness_values):
115            tournament = np.random.choice(len(fitness_values),
116                                           tournament_size)
117            tournament_fitness = np.array([fitness_values[i]
118                                         for i in tournament])
119            tournament_fitness_index = np.argmax(tournament_fitness)
120            selected_tour.append(tournament[tournament_fitness_index])
121
122        return selected_tour
123
124    def _elite_selection(self, tour: List[int], elite_size: int) -
125                        > List[int]:
126        """Elite selection"""
127        elites = []
128
129        while len(elites) < elite_size:
130            elites.append(np.argmax(self.fitness_function(tour)))
131
132        return elites
133
134    def _uniform_mutation(self, tour: List[int]) -> List[int]:
135        """Uniform mutation"""
136        mutated = tour.copy()
137
138        for i in range(len(tour)):
139            if np.random.random() < self.mutation_rate:
140                mutated[i] = np.random.randint(0, len(tour))
141
142        return mutated
143
144    def _hill_climbing(self, tour: List[int]) -> List[int]:
145        """Hill climbing"""
146
147        current_tour = tour
148        current_fitness = self.fitness_function(current_tour)
149
150        while True:
151            best_tour = current_tour
152            best_fitness = current_fitness
153
154            for i in range(len(tour)):
155                for j in range(i + 1, len(tour)):
156                    mutated_tour = self._swap_mutation(best_tour, i, j)
157
158                    mutated_fitness = self.fitness_function(mutated_tour)
159
160                    if mutated_fitness > best_fitness:
161                        best_tour = mutated_tour
162                        best_fitness = mutated_fitness
163
164            if best_fitness == current_fitness:
165                break
166
167            current_tour = best_tour
168            current_fitness = best_fitness
169
170        return current_tour
171
172    def _simulated_annealing(self, tour: List[int], temperature: float,
173                            cooling_rate: float, num_iterations: int) -
174                        > List[int]:
175        """Simulated annealing"""
176
177        current_tour = tour
178        current_fitness = self.fitness_function(current_tour)
179
180        for i in range(num_iterations):
181            if temperature == 0:
182                break
183
184            mutated_tour = self._uniform_mutation(current_tour)
185
186            mutated_fitness = self.fitness_function(mutated_tour)
187
188            if mutated_fitness < current_fitness:
189                current_tour = mutated_tour
190                current_fitness = mutated_fitness
191
192            else:
193                acceptance_probability = np.exp(
194                    (current_fitness - mutated_fitness) / temperature)
195
196                if np.random.random() < acceptance_probability:
197                    current_tour = mutated_tour
198                    current_fitness = mutated_fitness
199
200            temperature *= cooling_rate
201
202        return current_tour
203
204    def _genetic_algorithm(self, population_size: int, num_generations: int,
205                          mutation_rate: float, crossover_rate: float,
206                          tournament_size: int, elite_size: int) -
207                        > List[int]:
208        """Genetic algorithm"""
209
210        population = self._initial_population(population_size)
211
212        for generation in range(num_generations):
213            population = self._selection(population, tournament_size)
214            population = self._crossover(population, crossover_rate)
215            population = self._mutation(population, mutation_rate)
216
217            elites = self._elite_selection(population, elite_size)
218            population = self._tournament_selection(population,
219                                                    tournament_size)
220
221            population.extend(elites)
222
223        return population
224
225    def _initial_population(self, population_size: int) -
226                        > List[List[int]]:
227        """Initial population"""
228
229        population = []
230
231        for i in range(population_size):
232            population.append(self._tour_generator())
233
234        return population
235
236    def _tour_generator(self) -> List[int]:
237        """Tour generator"""
238
239        tour = []
240
241        cities = list(range(len(self.cities)))
242
243        while len(tour) < len(self.cities):
244            city = np.random.choice(cities)
245            tour.append(city)
246            cities.remove(city)
247
248        return tour
249
250    def _fitness_function(self, tour: List[int]) -> float:
251        """Fitness function"""
252
253        total_distance = 0
254
255        for i in range(len(tour) - 1):
256            total_distance += self._distance_matrix[tour[i]][tour[i + 1]]
257
258        total_distance += self._distance_matrix[tour[-1]][tour[0]]
259
260        return total_distance
261
262    def _distance_matrix(self) -> List[List[int]]:
263        """Distance matrix"""
264
265        distance_matrix = []
266
267        for i in range(len(self.cities)):
268            row = []
269
270            for j in range(len(self.cities)):
271                if i == j:
272                    row.append(0)
273                else:
274                    row.append(self._distance_matrix[i][j])
275
276            distance_matrix.append(row)
277
278        return distance_matrix
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308                    row.append(0)
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311
312            distance_matrix.append(row)
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328                    row.append(self._distance_matrix[i][j])
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330            distance_matrix.append(row)
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398                    row.append(0)
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400                    row.append(self._distance_matrix[i][j])
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420            distance_matrix.append(row)
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436                    row.append(self._distance_matrix[i][j])
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761
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1005      for i in range(len(self.cities)):
1006          row = []
1007
1008          for j in range(len(self.cities)):
1009              if i == j:
1010                  row.append(0)
1011              else:
1012                  row.append(self._distance_matrix[i][j])
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1014          distance_matrix.append(row)
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1205
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1335              else:
1336                  row.append(self._distance_matrix[i][j])
1337
1338          distance_matrix.append(row)
1339
1340
```

```

107             tournament_size: int = 3) -> int:
108     """Tournament selection"""
109     tournament_indices = np.random.choice(len(fitness_values),
110                                             tournament_size,
111                                             replace=False)
112     tournament_fitness = [fitness_values[i] for i in
113                           tournament_indices]
114     winner_idx = tournament_indices[np.argmax(
115         tournament_fitness)]
116     return winner_idx
117
118
119
120
121     def evolve(self, generations: int) -> dict:
122         """Main evolutionary loop"""
123         fitness_history = []
124         best_tour = None
125         best_distance = float('inf')
126
127         for generation in range(generations):
128             # Evaluate fitness
129             fitness_values = [self._fitness(tour) for tour in
130                               self.population]
131             distances = [self._calculate_tour_distance(tour)
132                          for tour in self.population]
133
134             # Track best solution
135             min_distance_idx = np.argmin(distances)
136             if distances[min_distance_idx] < best_distance:
137                 best_distance = distances[min_distance_idx]
138                 best_tour = self.population[min_distance_idx].copy()
139
140             # Record statistics
141             fitness_history.append({
142                 'generation': generation,
143                 'best_distance': np.min(distances),
144                 'avg_distance': np.mean(distances),
145                 'worst_distance': np.max(distances)
146             })
147
148             # Create new population
149             new_population = []
150
151             # Elitism: keep best individual
152             new_population.append(best_tour.copy())
153
154             # Generate rest of population
155             while len(new_population) < self.population_size:
156                 # Selection
157                 parent1_idx = self._tournament_selection(
158                     fitness_values)

```

```

151         parent2_idx = self._tournament_selection(
152             fitness_values)
153
153         parent1 = self.population[parent1_idx]
154         parent2 = self.population[parent2_idx]
155
156         # Crossover
157         child1, child2 = self._order_crossover(parent1,
158             parent2)
159
160         # Mutation
161         child1 = self._swap_mutation(child1)
162         child2 = self._swap_mutation(child2)
163
163         new_population.extend([child1, child2])
164
165         # Trim to population size
166         self.population = new_population[:self.
167             population_size]
168
168     return {
169         'best_tour': best_tour,
170         'best_distance': best_distance,
171         'fitness_history': fitness_history
172     }
173
174     def plot_tour(self, tour: List[int], title: str = "BestTour"
175     ):
175         """Plot the tour"""
176         plt.figure(figsize=(10, 8))
177
178         # Plot cities
179         plt.scatter(self.cities[:, 0], self.cities[:, 1],
180                     c='red', s=100, zorder=2)
181
182         # Plot tour
183         tour_cities = self.cities[tour + [tour[0]]]    # Close the
184             loop
184         plt.plot(tour_cities[:, 0], tour_cities[:, 1],
185                     'b-', linewidth=2, zorder=1)
186
187         # Add city labels
188         for i, city in enumerate(self.cities):
189             plt.annotate(str(i), (city[0], city[1]),
190                         xytext=(5, 5), textcoords='offset points',
191                         )
192
192         plt.title(f"{title}\nDistance:{self.
193             _calculate_tour_distance(tour):.2f}")
193         plt.xlabel("XCoordinate")
194         plt.ylabel("YCoordinate")

```

```

195     plt.grid(True, alpha=0.3)
196     plt.show()
197
198 # Example usage
199 if __name__ == "__main__":
200     # Create random cities
201     np.random.seed(42)
202     num_cities = 20
203     cities = np.random.rand(num_cities, 2) * 100
204
205     # Initialize and run GA
206     tsp_ga = TSP_GA(cities, population_size=100, mutation_rate
207                      =0.02)
208     result = tsp_ga.evolve(generations=500)
209
210     print(f"Best distance: {result['best_distance']:.2f}")
211     print(f"Best tour: {result['best_tour']}")
212
213     # Plot best tour
214     tsp_ga.plot_tour(result['best_tour'])

```

A.4 NSGA-II for Multi-Objective Optimization

Listing A.4: NSGA-II Implementation

```

1 import numpy as np
2 from typing import List, Tuple
3
4 class NSGA2:
5     def __init__(self,
6                  objective_functions: List,
7                  num_variables: int,
8                  bounds: List[Tuple[float, float]],
9                  population_size: int = 100,
10                 crossover_rate: float = 0.9,
11                 mutation_rate: float = 0.1):
12
13         self.objective_functions = objective_functions
14         self.num_objectives = len(objective_functions)
15         self.num_variables = num_variables
16         self.bounds = bounds
17         self.population_size = population_size
18         self.crossover_rate = crossover_rate
19         self.mutation_rate = mutation_rate
20
21         # Ensure even population size
22         if self.population_size % 2 != 0:
23             self.population_size += 1
24
25     def _initialize_population(self) -> np.ndarray:

```

```

26     """Initialize random population"""
27     population = np.zeros((self.population_size, self.
28         num_variables))
29     for i in range(self.num_variables):
30         low, high = self.bounds[i]
31         population[:, i] = np.random.uniform(low, high, self.
32             population_size)
33     return population
34
35
36
37
38
39
40
41
42
43
44
45
46
47
48
49
50
51
52
53
54
55
56
57
58
59
60
61
62
63
64
65
    """
population = np.zeros((self.population_size, self.
        num_variables))
for i in range(self.num_variables):
    low, high = self.bounds[i]
population[:, i] = np.random.uniform(low, high, self.
    population_size)
return population

def _evaluate_objectives(self, population: np.ndarray) -> np.
    ndarray:
    """
Evaluate all objectives for population"""
objectives = np.zeros((len(population), self.
    num_objectives))
for i, individual in enumerate(population):
    for j, obj_func in enumerate(self.objective_functions
        ):
        objectives[i, j] = obj_func(individual)
return objectives

def _dominates(self, obj1: np.ndarray, obj2: np.ndarray) ->
    bool:
    """
Check if obj1 dominates obj2 (assuming minimization)
    """
    return np.all(obj1 <= obj2) and np.any(obj1 < obj2)

def _fast_non_dominated_sort(self, objectives: np.ndarray) ->
    Tuple[List[List[int]], np.ndarray]:
    """
Fast non-dominated sorting"""
population_size = len(objectives)
domination_count = np.zeros(population_size)
dominated_solutions = [[] for _ in range(population_size)
    ]
fronts = [[]]

# Find domination relationships
for i in range(population_size):
    for j in range(population_size):
        if i != j:
            if self._dominates(objectives[i], objectives[
                j]):
                dominated_solutions[i].append(j)
            elif self._dominates(objectives[j],
                objectives[i]):
                domination_count[i] += 1

            if domination_count[i] == 0:
                fronts[0].append(i)

# Build subsequent fronts
current_front = 0

```

```

66     while len(fronts[current_front]) > 0:
67         next_front = []
68         for i in fronts[current_front]:
69             for j in dominated_solutions[i]:
70                 domination_count[j] -= 1
71                 if domination_count[j] == 0:
72                     next_front.append(j)
73         current_front += 1
74         fronts.append(next_front)
75
76     # Remove empty last front
77     fronts.pop()
78
79     # Assign ranks
80     ranks = np.zeros(population_size)
81     for rank, front in enumerate(fronts):
82         for individual in front:
83             ranks[individual] = rank
84
85     return fronts, ranks
86
87     def _calculate_crowding_distance(self, objectives: np.ndarray
88                                     ,
89                                     front: List[int]) -> np.
90                                     ndarray:
91         """Calculate crowding distance for individuals in a front
92         """
93         if len(front) <= 2:
94             return np.full(len(front), float('inf'))
95
96         distances = np.zeros(len(front))
97
98         for obj_idx in range(self.num_objectives):
99             # Sort by objective value
100            sorted_indices = sorted(range(len(front)),
101                                key=lambda x: objectives[front[x], obj_idx])
102
103            # Set boundary points to infinity
104            distances[sorted_indices[0]] = float('inf')
105            distances[sorted_indices[-1]] = float('inf')
106
107            # Calculate distances for middle points
108            obj_range = (objectives[front[sorted_indices[-1]], obj_idx] -
109                         objectives[front[sorted_indices[0]], obj_idx])
110
111            if obj_range > 0:
112                for i in range(1, len(sorted_indices) - 1):
113                    distances[sorted_indices[i]] =
114                        (objectives[front[sorted_indices[i + 1]], obj_idx] -
115                         objectives[front[sorted_indices[i - 1]], obj_idx]) /
116                        obj_range
117
118        return distances
119
120    def _rank_population(self, fronts, ranks):
121        """Rank population based on fronts and ranks
122
123        Parameters
124        ----------
125        fronts : List[Front]
126        ranks : np.ndarray
127
128        Returns
129        -------
130        population : Population
131
132        Notes
133        -----
134        This function assumes that fronts and ranks are aligned correctly.
135        """
136
137        population = Population()
138
139        for front, rank in zip(fronts, ranks):
140            for individual in front:
141                population.add(individual, rank)
142
143        return population
144
145    def _select_parents(self, population: Population,
146                       num_parents: int) -> Population:
147        """Select parents from population
148
149        Parameters
150        ----------
151        population : Population
152        num_parents : int
153
154        Returns
155        -------
156        parents : Population
157
158        Notes
159        -----
160        This function uses a roulette wheel selection method.
161        """
162
163        parents = Population()
164
165        total_rank = population.get_total_rank()
166
167        for i in range(num_parents):
168            rank = random.uniform(0, total_rank)
169            cumulative_rank = 0
170            for individual in population:
171                cumulative_rank += individual.rank
172                if cumulative_rank > rank:
173                    parents.add(individual)
174                    break
175
176        return parents
177
178    def _create_offspring(self, parents: Population,
179                          num_offspring: int) -> Population:
180        """Create offspring from parents
181
182        Parameters
183        ----------
184        parents : Population
185        num_offspring : int
186
187        Returns
188        -------
189        offspring : Population
190
191        Notes
192        -----
193        This function uses a crossover and mutation process to create
194        new individuals.
195        """
196
197        offspring = Population()
198
199        for i in range(num_offspring):
200            parent1 = parents.select_random()
201            parent2 = parents.select_random()
202
203            offspring.add(self._crossover(parent1, parent2))
204
205            if random.random() < self.mutation_rate:
206                offspring.add(self._mutate(offspring[-1]))
207
208        return offspring
209
210    def _crossover(self, parent1: Individual,
211                  parent2: Individual) -> Individual:
212        """Perform crossover between two parents
213
214        Parameters
215        ----------
216        parent1 : Individual
217        parent2 : Individual
218
219        Returns
220        -------
221        offspring : Individual
222
223        Notes
224        -----
225        This function performs a uniform crossover.
226        """
227
228        offspring = Individual()
229
230        for i in range(self.num_objectives):
231            if random.random() < self.crossover_rate:
232                offspring[i] = parent1[i]
233            else:
234                offspring[i] = parent2[i]
235
236        return offspring
237
238    def _mutate(self, individual: Individual) -> Individual:
239        """Perform mutation on an individual
240
241        Parameters
242        ----------
243        individual : Individual
244
245        Returns
246        -------
247        mutated_individual : Individual
248
249        Notes
250        -----
251        This function performs a random mutation.
252        """
253
254        mutated_individual = Individual()
255
256        for i in range(self.num_objectives):
257            if random.random() < self.mutation_rate:
258                mutated_individual[i] = individual[i] +
259                    random.uniform(-self.mutation_std, self.mutation_std)
260            else:
261                mutated_individual[i] = individual[i]
262
263        return mutated_individual
264
265    def _evaluate(self, population: Population) -> Population:
266        """Evaluate population
267
268        Parameters
269        ----------
270        population : Population
271
272        Returns
273        -------
274        evaluated_population : Population
275
276        Notes
277        -----
278        This function evaluates each individual in the population.
279        """
280
281        evaluated_population = Population()
282
283        for individual in population:
284            evaluated_population.add(individual)
285
286            evaluated_population[i].fitness =
287                self._evaluate_fitness(individual)
288
289        return evaluated_population
290
291    def _evaluate_fitness(self, individual: Individual) -> float:
292        """Evaluate fitness of an individual
293
294        Parameters
295        ----------
296        individual : Individual
297
298        Returns
299        -------
300        fitness : float
301
302        Notes
303        -----
304        This function calculates the fitness of an individual based on
305        its objective values.
306        """
307
308        fitness = 0
309
310        for i in range(self.num_objectives):
311            fitness += individual[i]
312
313        return fitness
314
315    def _dominate(self, individual1: Individual,
316                  individual2: Individual) -> bool:
317        """Check if individual1 dominates individual2
318
319        Parameters
320        ----------
321        individual1 : Individual
322        individual2 : Individual
323
324        Returns
325        -------
326        dominates : bool
327
328        Notes
329        -----
330        This function checks if individual1 dominates individual2.
331        """
332
333        dominates = True
334
335        for i in range(self.num_objectives):
336            if individual1[i] > individual2[i]:
337                dominates = False
338                break
339
340        return dominates
341
342    def _is_dominated(self, individual: Individual,
343                      front: Front) -> bool:
344        """Check if individual is dominated by any individual in front
345
346        Parameters
347        ----------
348        individual : Individual
349        front : Front
350
351        Returns
352        -------
353        dominated : bool
354
355        Notes
356        -----
357        This function checks if individual is dominated by any individual
358        in front.
359        """
360
361        dominated = False
362
363        for other in front:
364            if self._dominate(individual, other):
365                dominated = True
366                break
367
368        return dominated
369
370    def _is_in_frontier(self, individual: Individual,
371                        front: Front) -> bool:
372        """Check if individual is in the frontier
373
374        Parameters
375        ----------
376        individual : Individual
377        front : Front
378
379        Returns
380        -------
381        in_frontier : bool
382
383        Notes
384        -----
385        This function checks if individual is in the frontier.
386        """
387
388        in_frontier = False
389
390        for other in front:
391            if self._dominate(individual, other):
392                in_frontier = True
393                break
394
395        return in_frontier
396
397    def _is_dominated_solutions(self, individual: Individual,
398                               front: Front) -> bool:
399        """Check if individual is dominated by any individual in front
400
401        Parameters
402        ----------
403        individual : Individual
404        front : Front
405
406        Returns
407        -------
408        dominated : bool
409
410        Notes
411        -----
412        This function checks if individual is dominated by any individual
413        in front.
414        """
415
416        dominated = False
417
418        for other in front:
419            if self._dominate(individual, other):
420                dominated = True
421                break
422
423        return dominated
424
425    def _is_in_dominated_solutions(self, individual: Individual,
426                                   front: Front) -> bool:
427        """Check if individual is in the dominated solutions
428
429        Parameters
430        ----------
431        individual : Individual
432        front : Front
433
434        Returns
435        -------
436        in_dominated_solutions : bool
437
438        Notes
439        -----
440        This function checks if individual is in the dominated solutions.
441        """
442
443        in_dominated_solutions = False
444
445        for other in front:
446            if self._dominate(individual, other):
447                in_dominated_solutions = True
448                break
449
450        return in_dominated_solutions
451
452    def _is_in_dominated_solutions(self, individual: Individual,
453                                   front: Front) -> bool:
454        """Check if individual is in the dominated solutions
455
456        Parameters
457        ----------
458        individual : Individual
459        front : Front
460
461        Returns
462        -------
463        in_dominated_solutions : bool
464
465        Notes
466        -----
467        This function checks if individual is in the dominated solutions.
468        """
469
470        in_dominated_solutions = False
471
472        for other in front:
473            if self._dominate(individual, other):
474                in_dominated_solutions = True
475                break
476
477        return in_dominated_solutions
478
479    def _is_in_dominated_solutions(self, individual: Individual,
480                                   front: Front) -> bool:
481        """Check if individual is in the dominated solutions
482
483        Parameters
484        ----------
485        individual : Individual
486        front : Front
487
488        Returns
489        -------
490        in_dominated_solutions : bool
491
492        Notes
493        -----
494        This function checks if individual is in the dominated solutions.
495        """
496
497        in_dominated_solutions = False
498
499        for other in front:
500            if self._dominate(individual, other):
501                in_dominated_solutions = True
502                break
503
504        return in_dominated_solutions
505
506    def _is_in_dominated_solutions(self, individual: Individual,
507                                   front: Front) -> bool:
508        """Check if individual is in the dominated solutions
509
510        Parameters
511        ----------
512        individual : Individual
513        front : Front
514
515        Returns
516        -------
517        in_dominated_solutions : bool
518
519        Notes
520        -----
521        This function checks if individual is in the dominated solutions.
522        """
523
524        in_dominated_solutions = False
525
526        for other in front:
527            if self._dominate(individual, other):
528                in_dominated_solutions = True
529                break
530
531        return in_dominated_solutions
532
533    def _is_in_dominated_solutions(self, individual: Individual,
534                                   front: Front) -> bool:
535        """Check if individual is in the dominated solutions
536
537        Parameters
538        ----------
539        individual : Individual
540        front : Front
541
542        Returns
543        -------
544        in_dominated_solutions : bool
545
546        Notes
547        -----
548        This function checks if individual is in the dominated solutions.
549        """
550
551        in_dominated_solutions = False
552
553        for other in front:
554            if self._dominate(individual, other):
555                in_dominated_solutions = True
556                break
557
558        return in_dominated_solutions
559
560    def _is_in_dominated_solutions(self, individual: Individual,
561                                   front: Front) -> bool:
562        """Check if individual is in the dominated solutions
563
564        Parameters
565        ----------
566        individual : Individual
567        front : Front
568
569        Returns
570        -------
571        in_dominated_solutions : bool
572
573        Notes
574        -----
575        This function checks if individual is in the dominated solutions.
576        """
577
578        in_dominated_solutions = False
579
580        for other in front:
581            if self._dominate(individual, other):
582                in_dominated_solutions = True
583                break
584
585        return in_dominated_solutions
586
587    def _is_in_dominated_solutions(self, individual: Individual,
588                                   front: Front) -> bool:
589        """Check if individual is in the dominated solutions
590
591        Parameters
592        ----------
593        individual : Individual
594        front : Front
595
596        Returns
597        -------
598        in_dominated_solutions : bool
599
600        Notes
601        -----
602        This function checks if individual is in the dominated solutions.
603        """
604
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606
607        for other in front:
608            if self._dominate(individual, other):
609                in_dominated_solutions = True
610                break
611
612        return in_dominated_solutions
613
614    def _is_in_dominated_solutions(self, individual: Individual,
615                                   front: Front) -> bool:
616        """Check if individual is in the dominated solutions
617
618        Parameters
619        ----------
620        individual : Individual
621        front : Front
622
623        Returns
624        -------
625        in_dominated_solutions : bool
626
627        Notes
628        -----
629        This function checks if individual is in the dominated solutions.
630        """
631
632        in_dominated_solutions = False
633
634        for other in front:
635            if self._dominate(individual, other):
636                in_dominated_solutions = True
637                break
638
639        return in_dominated_solutions
640
641    def _is_in_dominated_solutions(self, individual: Individual,
642                                   front: Front) -> bool:
643        """Check if individual is in the dominated solutions
644
645        Parameters
646        ----------
647        individual : Individual
648        front : Front
649
650        Returns
651        -------
652        in_dominated_solutions : bool
653
654        Notes
655        -----
656        This function checks if individual is in the dominated solutions.
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662            if self._dominate(individual, other):
663                in_dominated_solutions = True
664                break
665
666        return in_dominated_solutions
667
668    def _is_in_dominated_solutions(self, individual: Individual,
669                                   front: Front) -> bool:
670        """Check if individual is in the dominated solutions
671
672        Parameters
673        ----------
674        individual : Individual
675        front : Front
676
677        Returns
678        -------
679        in_dominated_solutions : bool
680
681        Notes
682        -----
683        This function checks if individual is in the dominated solutions.
684        """
685
686        in_dominated_solutions = False
687
688        for other in front:
689            if self._dominate(individual, other):
690                in_dominated_solutions = True
691                break
692
693        return in_dominated_solutions
694
695    def _is_in_dominated_solutions(self, individual: Individual,
696                                   front: Front) -> bool:
697        """Check if individual is in the dominated solutions
698
699        Parameters
700        ----------
701        individual : Individual
702        front : Front
703
704        Returns
705        -------
706        in_dominated_solutions : bool
707
708        Notes
709        -----
710        This function checks if individual is in the dominated solutions.
711        """
712
713        in_dominated_solutions = False
714
715        for other in front:
716            if self._dominate(individual, other):
717                in_dominated_solutions = True
718                break
719
720        return in_dominated_solutions
721
722    def _is_in_dominated_solutions(self, individual: Individual,
723                                   front: Front) -> bool:
724        """Check if individual is in the dominated solutions
725
726        Parameters
727        ----------
728        individual : Individual
729        front : Front
730
731        Returns
732        -------
733        in_dominated_solutions : bool
734
735        Notes
736        -----
737        This function checks if individual is in the dominated solutions.
738        """
739
740        in_dominated_solutions = False
741
742        for other in front:
743            if self._dominate(individual, other):
744                in_dominated_solutions = True
745                break
746
747        return in_dominated_solutions
748
749    def _is_in_dominated_solutions(self, individual: Individual,
750                                   front: Front) -> bool:
751        """Check if individual is in the dominated solutions
752
753        Parameters
754        ----------
755        individual : Individual
756        front : Front
757
758        Returns
759        -------
760        in_dominated_solutions : bool
761
762        Notes
763        -----
764        This function checks if individual is in the dominated solutions.
765        """
766
767        in_dominated_solutions = False
768
769        for other in front:
770            if self._dominate(individual, other):
771                in_dominated_solutions = True
772                break
773
774        return in_dominated_solutions
775
776    def _is_in_dominated_solutions(self, individual: Individual,
777                                   front: Front) -> bool:
778        """Check if individual is in the dominated solutions
779
780        Parameters
781        ----------
782        individual : Individual
783        front : Front
784
785        Returns
786        -------
787        in_dominated_solutions : bool
788
789        Notes
790        -----
791        This function checks if individual is in the dominated solutions.
792        """
793
794        in_dominated_solutions = False
795
796        for other in front:
797            if self._dominate(individual, other):
798                in_dominated_solutions = True
799                break
800
801        return in_dominated_solutions
802
803    def _is_in_dominated_solutions(self, individual: Individual,
804                                   front: Front) -> bool:
805        """Check if individual is in the dominated solutions
806
807        Parameters
808        ----------
809        individual : Individual
810        front : Front
811
812        Returns
813        -------
814        in_dominated_solutions : bool
815
816        Notes
817        -----
818        This function checks if individual is in the dominated solutions.
819        """
820
821        in_dominated_solutions = False
822
823        for other in front:
824            if self._dominate(individual, other):
825                in_dominated_solutions = True
826                break
827
828        return in_dominated_solutions
829
830    def _is_in_dominated_solutions(self, individual: Individual,
831                                   front: Front) -> bool:
832        """Check if individual is in the dominated solutions
833
834        Parameters
835        ----------
836        individual : Individual
837        front : Front
838
839        Returns
840        -------
841        in_dominated_solutions : bool
842
843        Notes
844        -----
845        This function checks if individual is in the dominated solutions.
846        """
847
848        in_dominated_solutions = False
849
850        for other in front:
851            if self._dominate(individual, other):
852                in_dominated_solutions = True
853                break
854
855        return in_dominated_solutions
856
857    def _is_in_dominated_solutions(self, individual: Individual,
858                                   front: Front) -> bool:
859        """Check if individual is in the dominated solutions
860
861        Parameters
862        ----------
863        individual : Individual
864        front : Front
865
866        Returns
867        -------
868        in_dominated_solutions : bool
869
870        Notes
871        -----
872        This function checks if individual is in the dominated solutions.
873        """
874
875        in_dominated_solutions = False
876
877        for other in front:
878            if self._dominate(individual, other):
879                in_dominated_solutions = True
880                break
881
882        return in_dominated_solutions
883
884    def _is_in_dominated_solutions(self, individual: Individual,
885                                   front: Front) -> bool:
886        """Check if individual is in the dominated solutions
887
888        Parameters
889        ----------
890        individual : Individual
891        front : Front
892
893        Returns
894        -------
895        in_dominated_solutions : bool
896
897        Notes
898        -----
899        This function checks if individual is in the dominated solutions.
900        """
901
902        in_dominated_solutions = False
903
904        for other in front:
905            if self._dominate(individual, other):
906                in_dominated_solutions = True
907                break
908
909        return in_dominated_solutions
910
911    def _is_in_dominated_solutions(self, individual: Individual,
912                                   front: Front) -> bool:
913        """Check if individual is in the dominated solutions
914
915        Parameters
916        ----------
917        individual : Individual
918        front : Front
919
920        Returns
921        -------
922        in_dominated_solutions : bool
923
924        Notes
925        -----
926        This function checks if individual is in the dominated solutions.
927        """
928
929        in_dominated_solutions = False
930
931        for other in front:
932            if self._dominate(individual, other):
933                in_dominated_solutions = True
934                break
935
936        return in_dominated_solutions
937
938    def _is_in_dominated_solutions(self, individual: Individual,
939                                   front: Front) -> bool:
940        """Check if individual is in the dominated solutions
941
942        Parameters
943        ----------
944        individual : Individual
945        front : Front
946
947        Returns
948        -------
949        in_dominated_solutions : bool
950
951        Notes
952        -----
953        This function checks if individual is in the dominated solutions.
954        """
955
956        in_dominated_solutions = False
957
958        for other in front:
959            if self._dominate(individual, other):
960                in_dominated_solutions = True
961                break
962
963        return in_dominated_solutions
964
965    def _is_in_dominated_solutions(self, individual: Individual,
966                                   front: Front) -> bool:
967        """Check if individual is in the dominated solutions
968
969        Parameters
970        ----------
971        individual : Individual
972        front : Front
973
974        Returns
975        -------
976        in_dominated_solutions : bool
977
978        Notes
979        -----
980        This function checks if individual is in the dominated solutions.
981        """
982
983        in_dominated_solutions = False
984
985        for other in front:
986            if self._dominate(individual, other):
987                in_dominated_solutions = True
988                break
989
990        return in_dominated_solutions
991
992    def _is_in_dominated_solutions(self, individual: Individual,
993                                   front: Front) -> bool:
994        """Check if individual is in the dominated solutions
995
996        Parameters
997        ----------
998        individual : Individual
999        front : Front
1000
1001       Returns
1002       -------
1003       in_dominated_solutions : bool
1004
1005       Notes
1006       -----
1007       This function checks if individual is in the dominated solutions.
1008       """
1009
1010      in_dominated_solutions = False
1011
1012      for other in front:
1013          if self._dominate(individual, other):
1014              in_dominated_solutions = True
1015              break
1016
1017      return in_dominated_solutions
1018
1019    def _is_in_dominated_solutions(self, individual: Individual,
1020                                   front: Front) -> bool:
1021        """Check if individual is in the dominated solutions
1022
1023        Parameters
1024        ----------
1025        individual : Individual
1026        front : Front
1027
1028        Returns
1029        -------
1030        in_dominated_solutions : bool
1031
1032        Notes
1033        -----
1034        This function checks if individual is in the dominated solutions.
1035        """
1036
1037        in_dominated_solutions = False
1038
1039        for other in front:
1040            if self._dominate(individual, other):
1041                in_dominated_solutions = True
1042                break
1043
1044        return in_dominated_solutions
1045
1046    def _is_in_dominated_solutions(self, individual: Individual,
1047                                   front: Front) -> bool:
1048        """Check if individual is in the dominated solutions
1049
1050        Parameters
1051        ----------
1052        individual : Individual
1053        front : Front
1054
1055        Returns
1056        -------
1057        in_dominated_solutions : bool
1058
1059        Notes
1060        -----
1061        This function checks if individual is in the dominated solutions.
1062        """
1063
1064        in_dominated_solutions = False
1065
1066        for other in front:
1067            if self._dominate(individual, other):
1068                in_dominated_solutions = True
1069                break
1070
1071        return in_dominated_solutions
1072
1073    def _is_in_dominated_solutions(self, individual: Individual,
1074                                   front: Front) -> bool:
1075        """Check if individual is in the dominated solutions
1076
1077        Parameters
1078        ----------
1079        individual : Individual
1080        front : Front
1081
1082        Returns
1083        -------
1084        in_dominated_solutions : bool
1085
1086        Notes
1087        -----
1088        This function checks if individual is in the dominated solutions.
1089        """
1090
1091        in_dominated_solutions = False
1092
1093        for other in front:
1094            if self._dominate(individual, other):
1095                in_dominated_solutions = True
1096                break
1097
1098        return in_dominated_solutions
1099
1100    def _is_in_dominated_solutions(self, individual: Individual,
1101                                   front: Front) -> bool:
1102        """Check if individual is in the dominated solutions
1103
1104        Parameters
1105        ----------
1106        individual : Individual
1107        front : Front
1108
1109        Returns
1110        -------
1111        in_dominated_solutions : bool
1112
1113        Notes
1114        -----
1115        This function checks if individual is in the dominated solutions.
1116        """
1117
1118        in_dominated_solutions = False
1119
1120        for other in front:
1121            if self._dominate(individual, other):
1122                in_dominated_solutions = True
1123                break
1124
1125        return in_dominated_solutions
1126
1127    def _is_in_dominated_solutions(self, individual: Individual,
1128                                   front: Front) -> bool:
1129        """Check if individual is in the dominated solutions
1130
1131        Parameters
1132        ----------
1133        individual : Individual
1134        front : Front
1135
1136        Returns
1137        -------
1138        in_dominated_solutions : bool
1139
1140        Notes
1141        -----
1142        This function checks if individual is in the dominated solutions.
1143        """
1144
1145        in_dominated_solutions = False
1146
1147        for other in front:
1148            if self._dominate(individual, other):
1149                in_dominated_solutions = True
1150                break
1151
1152        return in_dominated_solutions
1153
1154    def _is_in_dominated_solutions(self, individual: Individual,
1155                                   front: Front) -> bool:
1156        """Check if individual is in the dominated solutions
1157
1158        Parameters
1159        ----------
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1161        front : Front
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1163        Returns
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1166
1167        Notes
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1177                break
1178
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1182                                   front: Front) -> bool:
1183        """Check if individual is in the dominated solutions
1184
1185        Parameters
1186        ----------
1187        individual : Individual
1188        front : Front
1189
1190        Returns
1191        -------
1192        in_dominated_solutions : bool
1193
1194        Notes
1195        -----
1196        This function checks if individual is in the dominated solutions.
1197        """
1198
1199        in_dominated_solutions = False
1200
1201        for other in front:
1202            if self._dominate(individual, other):
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1204                break
1205
1206        return in_dominated_solutions
1207
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1209                                   front: Front) -> bool:
1210        """Check if individual is in the dominated solutions
1211
1212        Parameters
1213        ----------
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1215        front : Front
1216
1217        Returns
1218        -------
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1220
1221        Notes
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1224        """
1225
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1263                                   front: Front) -> bool:
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1290                                   front: Front) -> bool:
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1296        front : Front
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1298        Returns
1299        -------
1300        in_dominated_solutions : bool
1301
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1303        -----
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1306
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1312                break
1313
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1383        Notes
1384        -----
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1386        """
1387
1388        in_dominated_solutions = False
1389
13
```

```

110         distance = (objectives[front[sorted_indices[i] - 1]] - objectives[front[sorted_indices[i]]]) / obj_range
111     distances[sorted_indices[i]] += distance /
112     obj_range
113
114     return distances
115
116 def _tournament_selection(self, ranks: np.ndarray,
117                           crowding_distances: np.ndarray,
118                           population_size: int) -> List[int]:
119     """Binary tournament selection based on rank and crowding
120     distance"""
121     selected = []
122
123     for _ in range(population_size):
124         # Select two random individuals
125         candidates = np.random.choice(len(ranks), 2, replace=False)
126         i, j = candidates[0], candidates[1]
127
128         # Compare based on rank first, then crowding distance
129         if ranks[i] < ranks[j]:
130             selected.append(i)
131         elif ranks[i] > ranks[j]:
132             selected.append(j)
133         else: # Same rank, compare crowding distance
134             if crowding_distances[i] > crowding_distances[j]:
135                 selected.append(i)
136             else:
137                 selected.append(j)
138
139     return selected
140
141 def _sbx_crossover(self, parent1: np.ndarray, parent2: np.ndarray,
142                     eta: float = 20.0) -> Tuple[np.ndarray, np.ndarray]:
143     """Simulated Binary Crossover (SBX)"""
144     if np.random.random() > self.crossover_rate:
145         return parent1.copy(), parent2.copy()
146
147     child1 = np.zeros_like(parent1)
148     child2 = np.zeros_like(parent2)
149
150     for i in range(len(parent1)):
151         if np.random.random() <= 0.5:
152             if abs(parent1[i] - parent2[i]) > 1e-14:
153                 y1, y2 = min(parent1[i], parent2[i]), max(
154                     parent1[i], parent2[i])

```

```

153
154         # Calculate beta
155         rand = np.random.random()
156         if rand <= 0.5:
157             beta = (2 * rand) ** (1.0 / (eta + 1))
158         else:
159             beta = (1.0 / (2 * (1 - rand))) ** (1.0 /
160                                         (eta + 1))
161
162         child1[i] = 0.5 * ((y1 + y2) - beta * (y2 -
163                               y1))
164         child2[i] = 0.5 * ((y1 + y2) + beta * (y2 -
165                               y1))
166
167         # Ensure bounds
168         low, high = self.bounds[i]
169         child1[i] = np.clip(child1[i], low, high)
170         child2[i] = np.clip(child2[i], low, high)
171     else:
172         child1[i] = parent1[i]
173         child2[i] = parent2[i]
174
175     return child1, child2
176
177 def _polynomial_mutation(self, individual: np.ndarray,
178                         eta: float = 20.0) -> np.ndarray:
179     """Polynomial mutation"""
180     mutated = individual.copy()
181
182     for i in range(len(mutated)):
183         if np.random.random() < self.mutation_rate:
184             low, high = self.bounds[i]
185             delta1 = (mutated[i] - low) / (high - low)
186             delta2 = (high - mutated[i]) / (high - low)
187
188             rand = np.random.random()
189             mut_pow = 1.0 / (eta + 1.0)
190
191             if rand <= 0.5:
192                 xy = 1.0 - delta1
193                 val = 2.0 * rand + (1.0 - 2.0 * rand) * (xy
194                                               ** (eta + 1.0))
195                 deltaq = val ** mut_pow - 1.0
196             else:
197                 xy = 1.0 - delta2
198                 val = 2.0 * (1.0 - rand) + 2.0 * (rand - 0.5)
199                 * (xy ** (eta + 1.0))
200                 deltaq = 1.0 - val ** mut_pow

```

```

199
200         mutated[i] += deltaq * (high - low)
201         mutated[i] = np.clip(mutated[i], low, high)
202
203     return mutated
204
205 def evolve(self, generations: int) -> dict:
206     """Main NSGA-II evolution loop"""
207     # Initialize population
208     population = self._initialize_population()
209
210     for generation in range(generations):
211         # Evaluate objectives
212         objectives = self._evaluate_objectives(population)
213
214         # Non-dominated sorting
215         fronts, ranks = self._fast_non_dominated_sort(
216             objectives)
217
218         # Calculate crowding distances
219         crowding_distances = np.zeros(len(population))
220         for front in fronts:
221             if len(front) > 0:
222                 distances = self._calculate_crowding_distance(
223                     objectives, front)
224                 for i, individual_idx in enumerate(front):
225                     crowding_distances[individual_idx] =
226                         distances[i]
227
228         # Selection for mating pool
229         mating_pool_indices = self._tournament_selection(
230             ranks, crowding_distances,
231                                         self.
232                                         population_size
233                                         )
234
235         mating_pool = population[mating_pool_indices]
236
237         # Create offspring through crossover and mutation
238         offspring = []
239         for i in range(0, self.population_size, 2):
240             parent1 = mating_pool[i]
241             parent2 = mating_pool[i + 1]
242
243             child1, child2 = self._sbx_crossover(parent1,
244                                                 parent2)
245             child1 = self._polynomial_mutation(child1)
246             child2 = self._polynomial_mutation(child2)
247
248             offspring.extend([child1, child2])
249
250         offspring = np.array(offspring)

```

```

243
244     # Combine parent and offspring populations
245     combined_population = np.vstack([population,
246                                         offspring])
246     combined_objectives = self._evaluate_objectives(
247                                         combined_population)
247
248     # Environmental selection
249     combined_fronts, combined_ranks = self.
250                                         _fast_non_dominated_sort(combined_objectives)
250
251     new_population = []
252     front_idx = 0
253
254     # Add complete fronts
255     while (len(new_population) + len(combined_fronts[
256         front_idx]) <= self.population_size):
256         for individual_idx in combined_fronts[front_idx]:
257             new_population.append(individual_idx)
258         front_idx += 1
259
260         if front_idx >= len(combined_fronts):
261             break
262
263     # Add partial front if needed
264     if len(new_population) < self.population_size and
265         front_idx < len(combined_fronts):
265         last_front = combined_fronts[front_idx]
266         crowding_distances = self.
267             _calculate_crowding_distance(
268             combined_objectives, last_front)
268
269         # Sort by crowding distance (descending)
270         sorted_indices = sorted(range(len(last_front)),
271                             key=lambda x:
272                                 crowding_distances[x],
273                                 reverse=True)
273
274         remaining_slots = self.population_size - len(
275             new_population)
275         for i in range(remaining_slots):
276             new_population.append(last_front[
277                 sorted_indices[i]])
277
278         # Update population
279         population = combined_population[new_population]
280
281     # Final evaluation and return Pareto front
282     final_objectives = self._evaluate_objectives(population)
283     fronts, _ = self._fast_non_dominated_sort(
284         final_objectives)

```

```

282     pareto_front_indices = fronts[0]
283     pareto_front_solutions = population[pareto_front_indices]
284     pareto_front_objectives = final_objectives[
285         pareto_front_indices]
286
287     return {
288         'pareto_front_solutions': pareto_front_solutions,
289         'pareto_front_objectives': pareto_front_objectives,
290         'final_population': population,
291         'final_objectives': final_objectives
292     }
293
294 # Example: Minimize two objectives (ZDT1 problem)
295 def objective1(x):
296     return x[0]
297
298 def objective2(x):
299     g = 1 + 9 * np.sum(x[1:]) / (len(x) - 1)
300     h = 1 - np.sqrt(x[0] / g)
301     return g * h
302
303 # Usage
304 if __name__ == "__main__":
305     objectives = [objective1, objective2]
306     bounds = [(0, 1)] * 10 # 10-dimensional problem
307
308     nsga2 = NSGA2(objectives, 10, bounds, population_size=100)
309     result = nsga2.evolve(generations=250)
310
311     # Plot Pareto front
312     pareto_objectives = result['pareto_front_objectives']
313     plt.figure(figsize=(10, 6))
314     plt.scatter(pareto_objectives[:, 0], pareto_objectives[:, 1],
315                 c='red', alpha=0.7)
316     plt.xlabel('Objective 1')
317     plt.ylabel('Objective 2')
318     plt.title('Pareto Front')
319     plt.grid(True, alpha=0.3)
320     plt.show()

```

Appendix B

Practical Examples and Case Studies

B.1 Function Optimization Problems

B.1.1 OneMax Problem

The OneMax problem is the simplest optimization problem for binary genetic algorithms.

Problem Definition

Maximize the number of 1s in a binary string:

$$f(x) = \sum_{i=1}^n x_i \quad (\text{B.1})$$

where $x_i \in \{0, 1\}$ and n is the string length.

Expected Performance

- **Optimal solution:** All 1s string
- **Global optimum:** $f^* = n$
- **Expected convergence:** $O(n \log n)$ generations
- **Population size:** $O(\log n)$ sufficient

GA Configuration

Parameter	Value
Representation	Binary string
Population size	50 – 100
Selection	Tournament (size 3)
Crossover	One-point, $p_c = 0.8$
Mutation	Bit-flip, $p_m = 1/n$
Generations	100 – 200

Table B.1: OneMax GA Configuration

B.1.2 Sphere Function

Continuous optimization benchmark function.

Problem Definition

$$f(\mathbf{x}) = \sum_{i=1}^n x_i^2 \quad (\text{B.2})$$

where $\mathbf{x} \in [-5.12, 5.12]^n$.

Characteristics

- **Type:** Unimodal, separable
- **Global minimum:** $\mathbf{x}^* = (0, 0, \dots, 0)$
- **Global optimum:** $f^* = 0$
- **Difficulty:** Easy (convex, single optimum)

B.1.3 Rastrigin Function

Multimodal benchmark function.

Problem Definition

$$f(\mathbf{x}) = An + \sum_{i=1}^n [x_i^2 - A \cos(2\pi x_i)] \quad (\text{B.3})$$

where $A = 10$ and $\mathbf{x} \in [-5.12, 5.12]^n$.

Characteristics

- **Type:** Multimodal, separable
- **Local minima:** $A \cdot n$ local minima
- **Global minimum:** $\mathbf{x}^* = (0, 0, \dots, 0)$
- **Global optimum:** $f^* = 0$
- **Difficulty:** Medium (many local optima)

GA Challenges

- Premature convergence to local optima
- Requires high population diversity
- Benefits from diversity preservation techniques

B.1.4 Rosenbrock Function

Non-convex optimization problem.

Problem Definition

$$f(\mathbf{x}) = \sum_{i=1}^{n-1} [100(x_{i+1} - x_i^2)^2 + (1 - x_i)^2] \quad (\text{B.4})$$

Characteristics

- **Type:** Unimodal but non-convex
- **Global minimum:** $\mathbf{x}^* = (1, 1, \dots, 1)$
- **Global optimum:** $f^* = 0$
- **Difficulty:** Hard (narrow curved valley)

B.2 Combinatorial Optimization Problems

B.2.1 Traveling Salesman Problem (TSP)

Problem Description

Find the shortest route visiting all cities exactly once and returning to the starting city.

Mathematical Formulation

Minimize:

$$\sum_{i=1}^n \sum_{j=1}^n d_{ij} x_{ij} \quad (\text{B.5})$$

Subject to:

$$\sum_{j=1}^n x_{ij} = 1, \quad \forall i \quad (\text{B.6})$$

$$\sum_{i=1}^n x_{ij} = 1, \quad \forall j \quad (\text{B.7})$$

$$x_{ij} \in \{0, 1\} \quad (\text{B.8})$$

where d_{ij} is the distance between cities i and j .

GA Representation

- **Encoding:** Permutation of city indices
- **Example:** $(3, 1, 4, 2, 5)$ means visit cities in order $3 \rightarrow 1 \rightarrow 4 \rightarrow 2 \rightarrow 5 \rightarrow 3$

Specialized Operators

- **Crossover:** Order crossover (OX), Partially mapped crossover (PMX)
- **Mutation:** Swap, insert, inversion
- **Local search:** 2-opt, 3-opt improvements

Performance Tips

- Use edge recombination for better building block preservation
- Apply local search (hybrid GA)
- Consider nearest neighbor initialization
- Use elitist replacement

B.2.2 Knapsack Problem

Problem Description

Select items to maximize value while staying within weight constraint.

0/1 Knapsack Formulation

Maximize:

$$\sum_{i=1}^n v_i x_i \quad (\text{B.9})$$

Subject to:

$$\sum_{i=1}^n w_i x_i \leq W \quad (\text{B.10})$$

$$x_i \in \{0, 1\} \quad (\text{B.11})$$

where v_i is value, w_i is weight, and W is capacity.

GA Approach

- **Encoding:** Binary string (1 = include item, 0 = exclude)
- **Constraint handling:** Penalty function or repair mechanism
- **Fitness:** Value minus penalty for constraint violation

Penalty Function Example

$$fitness(x) = \sum_{i=1}^n v_i x_i - \alpha \max \left(0, \sum_{i=1}^n w_i x_i - W \right) \quad (\text{B.12})$$

where α is a penalty coefficient.

B.3 Real-World Applications

B.3.1 Neural Network Training

Problem Setup

Optimize neural network weights and biases using GA.

Representation

- **Encoding:** Real-valued vector of all weights and biases
- **Decoding:** Reshape vector into network structure

Fitness Function

$$fitness = \frac{1}{1 + MSE} \quad (B.13)$$

where MSE is mean squared error on training/validation set.

Advantages over Backpropagation

- No gradient information required
- Can optimize network topology
- Robust to local minima
- Handles discontinuous activation functions

B.3.2 Feature Selection

Problem Description

Select optimal subset of features for machine learning models.

GA Approach

- **Encoding:** Binary string (1 = include feature, 0 = exclude)
- **Fitness:** Model performance with selected features
- **Objectives:** Maximize accuracy, minimize number of features

Multi-objective Formulation

$$\text{Maximize: } accuracy(\text{selected features}) \quad (B.14)$$

$$\text{Minimize: } \text{number of selected features} \quad (B.15)$$

B.3.3 Job Shop Scheduling

Problem Description

Schedule jobs on machines to minimize makespan or total completion time.

Representation Options

1. **Priority-based:** Priority values for job-machine pairs
2. **Permutation-based:** Order of jobs for each machine
3. **Direct:** Actual schedule representation

Constraints

- Each job visits each machine exactly once
- Machines can process only one job at a time
- Jobs cannot be preempted
- Precedence constraints must be satisfied

B.4 Parameter Tuning Guidelines

B.4.1 Population Size

Problem Complexity	Population Size
Simple (OneMax)	50 – 100
Medium (TSP, 50 cities)	100 – 500
Complex (Large TSP)	500 – 2000
Multi-objective	100 – 300

Table B.2: Population Size Guidelines

B.4.2 Crossover and Mutation Rates

Problem Type	Crossover Rate	Mutation Rate
Binary optimization	0.7 – 0.9	$1/L$ to $10/L$
Real-valued	0.8 – 0.9	0.01 – 0.1
Permutation	0.8 – 0.9	0.01 – 0.05
Multi-objective	0.9	$1/L$

Table B.3: Crossover and Mutation Rate Guidelines

where L is the chromosome length.

B.4.3 Selection Pressure

- **Low pressure:** Tournament size 2-3, linear ranking
- **Medium pressure:** Tournament size 4-7
- **High pressure:** Tournament size > 7 , truncation selection

B.5 Performance Analysis

B.5.1 Convergence Metrics

- **Best fitness:** Track best solution over generations
- **Average fitness:** Monitor population quality
- **Diversity:** Measure population spread
- **Success rate:** Percentage of runs finding global optimum

B.5.2 Statistical Testing

- Run multiple independent trials (20-30)
- Report mean, standard deviation, best, worst
- Use statistical tests (t-test, Mann-Whitney U)
- Consider effect size, not just significance

B.5.3 Comparison with Other Methods

Method	Speed	Global Search	Implementation
Hill Climbing	Fast	Poor	Easy
Simulated Annealing	Medium	Good	Medium
Genetic Algorithm	Slow	Excellent	Medium
Particle Swarm	Medium	Good	Easy
Differential Evolution	Medium	Excellent	Easy

Table B.4: Algorithm Comparison

B.6 Common Pitfalls and Solutions

B.6.1 Premature Convergence

Symptoms:

- Population converges to suboptimal solution
- Low diversity after few generations
- No improvement for many generations

Solutions:

- Increase population size
- Reduce selection pressure

- Increase mutation rate
- Use diversity preservation techniques
- Apply restart strategies

B.6.2 Slow Convergence

Symptoms:

- Little improvement over many generations
- High population diversity maintained
- Random walk behavior

Solutions:

- Increase selection pressure
- Reduce mutation rate
- Apply local search (hybrid GA)
- Use better initialization
- Adjust crossover operators

B.6.3 Constraint Handling Issues

Common Problems:

- All individuals violate constraints
- Feasible region too small
- Penalty coefficients poorly set

Solutions:

- Use repair mechanisms
- Apply specialized operators
- Implement feasibility preservation
- Use multi-objective approach
- Adjust penalty weights dynamically

B.7 Advanced Techniques

B.7.1 Hybrid Genetic Algorithms

Combine GA with local search methods:

- **Memetic algorithms:** GA + local search
- **Lamarckian evolution:** Inherit improved solutions
- **Baldwinian evolution:** Use local search for fitness evaluation only

B.7.2 Adaptive Parameter Control

Automatically adjust GA parameters during evolution:

- **Deterministic:** Pre-defined schedule
- **Adaptive:** Based on population state
- **Self-adaptive:** Parameters evolve with population

B.7.3 Parallel Genetic Algorithms

Distribute computation across multiple processors:

- **Master-slave:** Parallel fitness evaluation
- **Island model:** Multiple populations with migration
- **Cellular GA:** Spatial population structure

B.8 Implementation Best Practices

B.8.1 Code Organization

- Separate representation from operators
- Use modular design for easy testing
- Implement proper random number generation
- Add logging and visualization capabilities

B.8.2 Testing and Validation

- Test on known benchmark problems
- Verify operators maintain validity
- Check random number generation quality
- Profile performance bottlenecks

B.8.3 Documentation

- Document parameter choices and reasoning
- Record experimental setup details
- Maintain version control
- Share reproducible results

B.9 Chapter Summary

This chapter provided practical examples and case studies demonstrating genetic algorithm applications across various problem domains. Key lessons include the importance of proper representation design, parameter tuning, and performance analysis. Understanding common pitfalls and their solutions is crucial for successful GA implementation.

B.10 Key Takeaways

- Problem representation is critical for GA success
- Parameter settings must match problem characteristics
- Statistical validation ensures reliable results
- Hybrid approaches often outperform pure GAs
- Domain knowledge should guide operator design
- Proper testing and documentation are essential

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