### **Exercise 1: System Catalog Queries**

a) One index is built on the Customers relation.

"customers\_pkey" PRIMARY KEY, btree (customerid)

```
b)
```

```
1 -- Relations ordered by disk pages occupied
2 SELECT relname, relpages
3 FROM pg_class
4 WHERE relkind = 'r'
5 ORDER BY relpages desc;
6
7 -- Indexes ordered by disk pages occupied
8 SELECT relname, relpages
9 FROM pg_class
10 WHERE relkind = 'i'
11 ORDER BY relpages desc;
```

Two relations that occupy the most number of disk pages:

customers 471 orderliness 381

Two relations that occupy the most number of disk pages:

customers\_pkey 57 orders\_pkey 35

# c)

```
13 -- Number of distinct values in each column of the Customners table
14 SELECT attname, n_distinct
15 FROM pg_stats
16 WHERE tablename = 'customers'
17 ORDER by n_distinct;|
```

Ĩ attname	n_distinct
customerid	-1
firstname	-1
lastname	-1
address1	-1
city	-1
email	-1
phone	-1
creditcard	-1
username	-1
zip	-0.475
address2	1
password	1
region	2
gender	2
creditcardtype	5
income	5
country	11
state	52
creditcardexpiration	60
age	73
(20 rows)	

In addition to the attributes that are already indexed, attributes age and creditcardexpiration are suitable for building a B-tree index. They are likely to have range-style queries and sorting.

Also, they have enough many distinct value for the index to be effective. (It is not useful to have an index for an attribute with only two values like gender)

```
d)
19 -- 1d Number of distinct values in each column of Customers table
20 -- without using the Postgres catalog
21 SELECT count(distinct customerid) as customerid,
22 count(distinct firstname) as firstname,
23 count(distinct lastname) as lastname,
24 count(distinct address1) as address1,
25 count(distinct city) as city,
26 count(distinct email) as email,
27 count(distinct phone) as phone,
28 count(distinct creditcard) as creditcard,
29 count(distinct username) as username,
30 count(distinct zip) as zip,
31 count(distinct address2) as address2,
32 count(distinct password) as password,
33 count(distinct region) as region,
34 count(distinct gender) as gender,
35 count(distinct creditcardtype) as creditcardtype,
36 count(distinct income) as income,
37 count(distinct country) as country,
38 count(distinct state) as state,
39 count(distinct creditcardexpiration) as creditcardexpiration,
40 count(distinct age) as age
41 FROM customers;
```

customerio						_			-					-
		-		-	-			-		-	20000	-	-	
address2	, ,		_											_

Table comparing expected count from the catalog compare with the actual count values

column name	n_distinct	expected count	actual count
customerid	-1	20000	20000
firstname	-1	20000	20000
lastname	-1	20000	20000
address1	-1	20000	20000
city	-1	20000	20000
email	-1	20000	20000
phone	-1	20000	20000
creditcard	-1	20000	20000
username	-1	20000	20000
zip	-0.475	9500	9500
address2	1	1	1
password	1	1	1
region	2	2	2
gender	2	2	2
creditcardtype	5	5	5
income	5	5	5

country	11	11	11
state	52	52	52
creditcardexpiration	60	60	60
age	73	73	73

The expected count from the catalog is the same with the actual count values.

## **Exercise 2: Equality Query**

```
tywei=> EXPLAIN SELECT * FROM customers WHERE country = 'Japan';
                         QUERY PLAN
Seq Scan on customers (cost=0.00..721.00 rows=995 width=156)
   Filter: ((country)::text = 'Japan'::text)
(2 rows)
The estimated cardinality of this query is 995.
tywei=> SELECT count(*) FROM customers WHERE country = 'Japan';
 count
   995
(1 row)
It is the same as the actual value.
b)
(disk pages read * seq_page_cost) + (rows scanned * cpu_tuple_cost) + (rows scanned *
cpu operator cost)
= (471 * 1.0) + (20000 * 0.01) + (20000 * 0.0025)
= 721
c)
\sigma_{country = 'Japan'}(customers)
\sigma_{country = 'Japan'}
  CUSTOMERS
```

### **Exercise 3: Equality Query with Indexes**

The index customers\_country occupies 59 pages.

The query optimizer checks the condition for each row it scans, and outputs only the ones that pass the condition. The estimated total cost of executing the best plan for this query is 721.00.

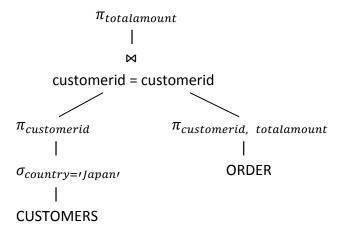
c) The estimated cost of non-clustered index is the same as using no index as in Exercise 2. The reason may be that lots of rows pass the condition. Since the index is non-clustered, number of disk pages read using the index is similar to using no index (sequential scanning). Besides, using index takes addition cost to retrieve the object.

The query optimizer fetches the rows using the index customers\_country. The estimated total cost of executing the best plan for this query is 56.66.

e) The best plan of clustered index cost much smaller than the best plan of non-clustered index for the query. Since the index is clustered, the number of pages needed to read is much smaller than non-clustered index.

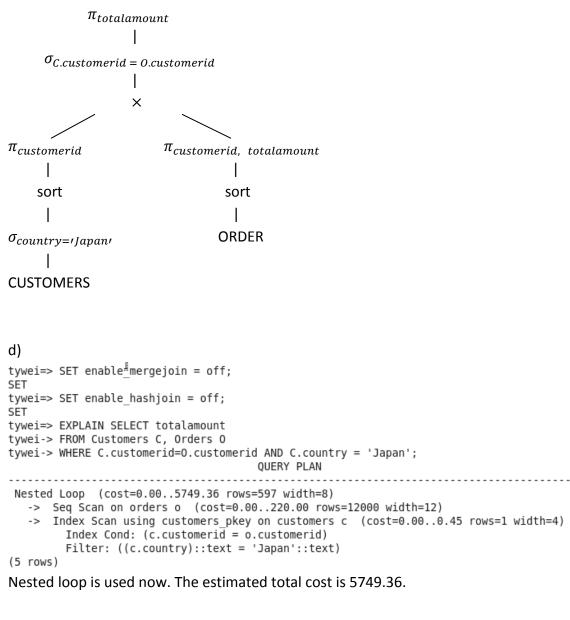
#### **Exercise 4: Join Query**

a)



b) The query optimizer uses hash join. The total estimated cost is 1004.41. The estimated result cardinality of this guery is 597.

Merge join is used by the query optimizer now. The total estimated cost is 1874.53.



#### Exercise 5

a)

```
11 -- 5a.first selected join
12 EXPLAIN SELECT AVG(totalamount) as avgOrder, country
13 FROM Customers C, Orders O
14 WHERE C.customerid = 0.customerid
15 GROUP BY country
16 ORDER BY avgOrder;
```

The query optimizer selects hash join for Query 5.1. The total expected cost is 1501.36.

```
17 -- 5a.second selected join
18 SET enable_hashjoin = off;
19 EXPLAIN SELECT AVG(totalamount) as avgOrder, country
20 FROM Customers C, Orders O
21 WHERE C.customerid = 0.customerid
22 GROUP BY country
23 ORDER BY avgOrder;
```

Disable the selected join algorithm, merge join is selected. The total expected cost is 2325.55.

b)

```
25 -- 5b.first selected join in Query 5.2
26 SET enable_hashjoin = on;
27 EXPLAIN SELECT *
28 FROM Customers C, Orders O
29 WHERE C.customerid = 0.customerid
30 ORDER BY C.customerid;
```

The query optimizer selectes merge join for Query 5.2. The total expected cost is 2265.19.

```
32 --5b.second selected jonin in Query 5.2
33 SET enable_mergejoin = off;
34 EXPLAIN SELECT *
35   FROM Customers C, Orders O
36   WHERE C.customerid = 0.customerid
37   ORDER BY C.customerid;
```

Disable the selected join algorithm, hash join is selected. The total expected cost is 3813.54.

c)

The optimizer selected different join algorithms for two queries. For Query 5.1, sorting on joining attributes is not required, so hash table is cheaper when joining large tables. For Query 5.2, sorting on joining attributes is required anyway due to the 'order by' clause, and thus merge sort is preferred.

### Exercise 6

a)

```
50 --6a.best plan for Query 6.1
51 EXPLAIN SELECT C.customerid, C.lastname
52 FROM Customers C
53 WHERE 4 < (SELECT COUNT(*)
54 FROM Orders 0
WHERE 0.customerid = C.customerid);
```

The estimated total cost of best plan is 5001021.00.

b)

```
57 -- 6b. query to create a view OrderCount
58 CREATE VIEW OrderCount(customerid, numorders)
       AS SELECT C.customerid, COUNT(*)
       FROM Customers C, Orders O
61
       WHERE O.customerid = C.customerid
62
       GROUP BY C.customerid;
c)
64 --6c.Query 6.2
65 SELECT C.customerid, C.lastname
       FROM Customers C, OrderCount OC
       WHERE OC.customerid = C.customerid
68
       AND OC.numorders > 4;
d)
70 --6d.Query 6.2 evaluation
71 EXPLAIN SELECT C.customerid, C.lastname
       FROM Customers C, OrderCount OC
73
       WHERE OC.customerid = C.customerid
       AND OC.numorders > 4;
```

The estimated total cost for Query 6.2 is 3887.44. It's much quicker than nested query in (a).

#### Exercise 7

a)

```
88 -- 7a.best plan for Query 7.1
 89 EXPLAIN SELECT customerid, lastname, numorders
 91
        SELECT C.customerid, C.lastname, count(*) as numorders
 92
            FROM Customers C, Orders O
 93
            WHERE C.customerid = O.customerid AND C.country = 'Japan'
            GROUP BY C.customerid, lastname) AS ORDERCOUNTS1
 94
 95
        WHERE 5 >= (SELECT count(*)
 96
            FROM (
 97
            SELECT C.customerid, C.lastname, count(*) as numorders
 98
            FROM Customers C, Orders O
 99
            WHERE C.customerid=0.customerid AND C.country = 'Japan'
100
            GROUP BY C.customerid, lastname) AS ORDERCOUNTS2
101
        WHERE ORDERCOUNTS1.numorders < ORDERCOUNTS2.numorders)</pre>
        ORDER BY customerid;
102
```

The estimated total cost of best plan is 614927.01.

b)

```
104 --7b.query to create a view MoreFrequentJapanCustomers
105 CREATE VIEW OrderCount(customerid, numorders)
106
        AS SELECT C.customerid, COUNT(*)
107
        FROM Customers C, Orders O
 108
        WHERE O.customerid = C.customerid
109
        AND C.country = 'Japan
110
        GROUP BY C.customerid;
111 CREATE VIEW MoreFrequentJapanCustomers(customerid, oRank)
112
       AS SELECT 01.customerid, COUNT(*)
113
        FROM OrderCount 01
114
        LEFT OUTER JOIN OrderCount 02
115
        ON 01.numorders < 02.numorders
116
        GROUP BY 01.customerid;
c)
121 -- 7c. Query 7.2
122 SELECT C.customerid, C.lastname, OC.numorders
        FROM Customers C, OrderCount OC, MoreFrequentJapanCustomers M
        WHERE OC.customerid = C.customerid
125
        AND C.customerid = M.customerid *
126
        AND M.oRank <= 5
        ORDER BY C.customerid;
127
d)
129 -- 7d. Query 7.2 evaluation
130 EXPLAIN SELECT C.customerid, C.lastname, OC.numorders
        FROM Customers C, OrderCount OC, MoreFrequentJapanCustomers M
132
        WHERE OC.customerid = C.customerid
        AND C.customerid = M.customerid
133
134
        AND M.oRank <= 5
135
        ORDER BY C.customerid;
```

The estimated total cost for Query 7.2 is 12923.09. It's much quicker than nested query in (a). In (a), the sub-query "5 >= ..." is evaluated once for each row in outer query ORDERCOUNTS1.numorders, which is inefficient.