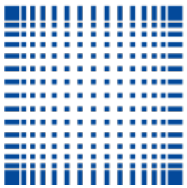

hochschule mannheim

Understanding Eventual Consistency

MSI Presentation SS2014

Horst Schneider, Patrick Beedgen
Hochschule Mannheim

June 17th, 2014

Introduction

„A prominent quote“

Anfang Hauptteil Horst

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- Two examples: Int Register **intreg**, Counter **ctr**

$$\mathrm{Op}_{\mathrm{ctr}} = \{\mathrm{rd}, \mathrm{inc}\}$$
$$\mathrm{Op}_{\mathrm{intreg}} = \{\mathrm{rd}, \mathrm{wr}(k) | k \in \mathbb{Z}\}$$

Replicated Data Types

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In a *strongly consistent system*, the semantics of a data type can be described by a function

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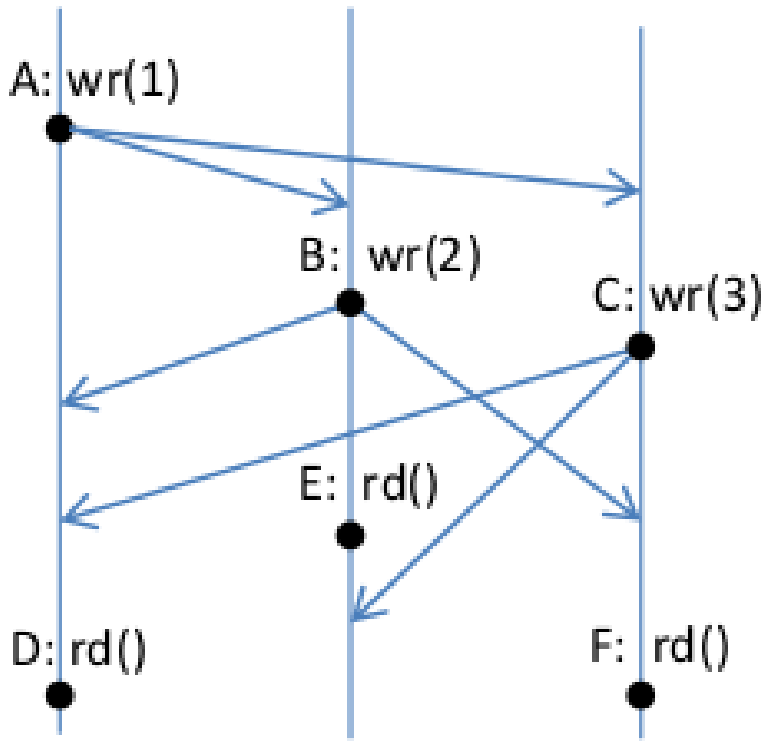
Replicated Data Types

Conflict Resolution Strategies



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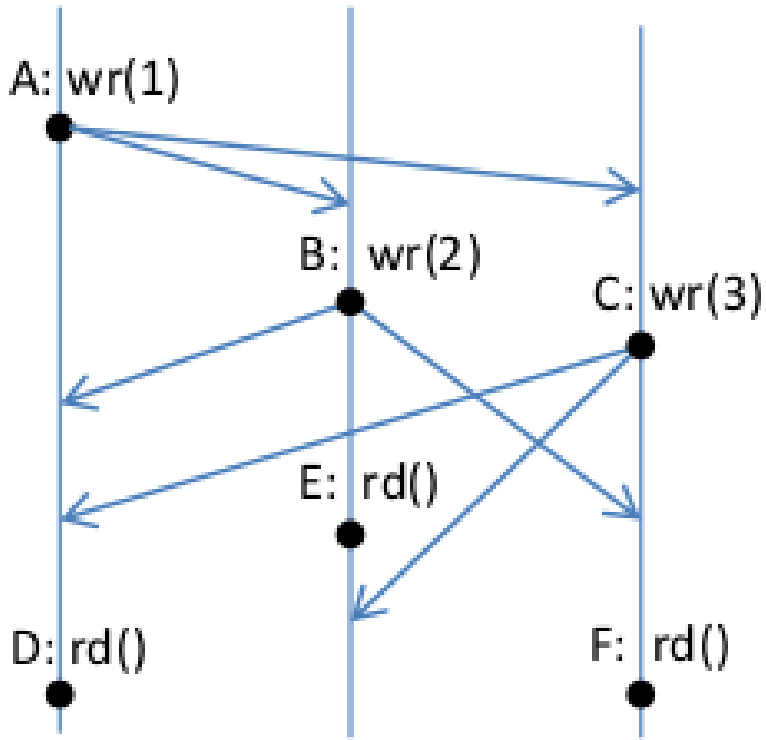

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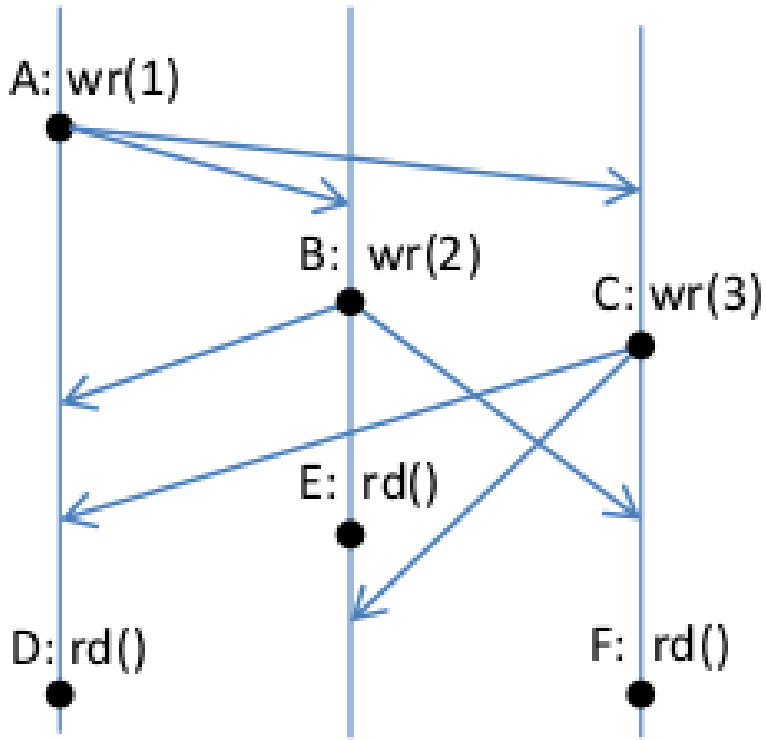
Replicated Data Types

Conflict Resolution Strategies



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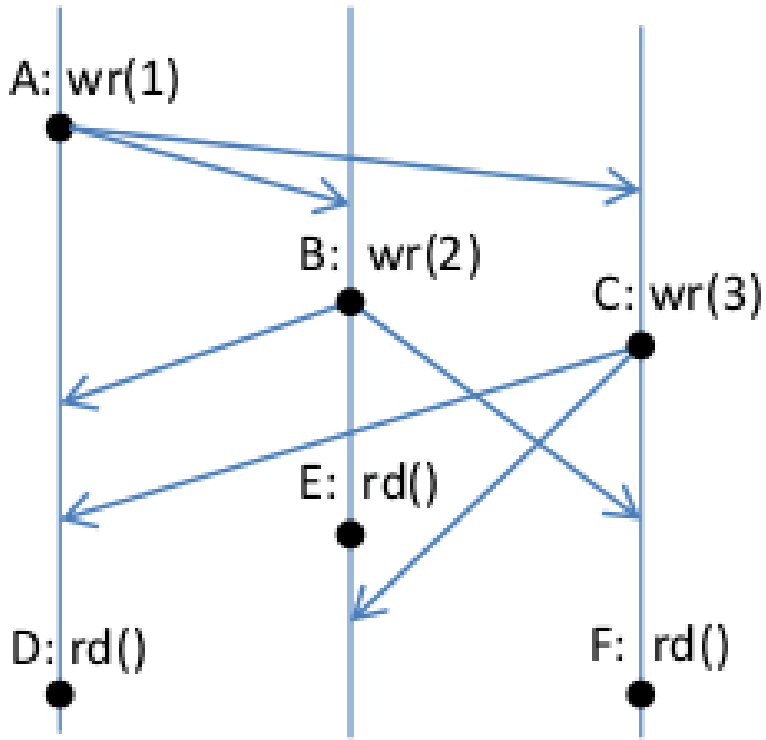

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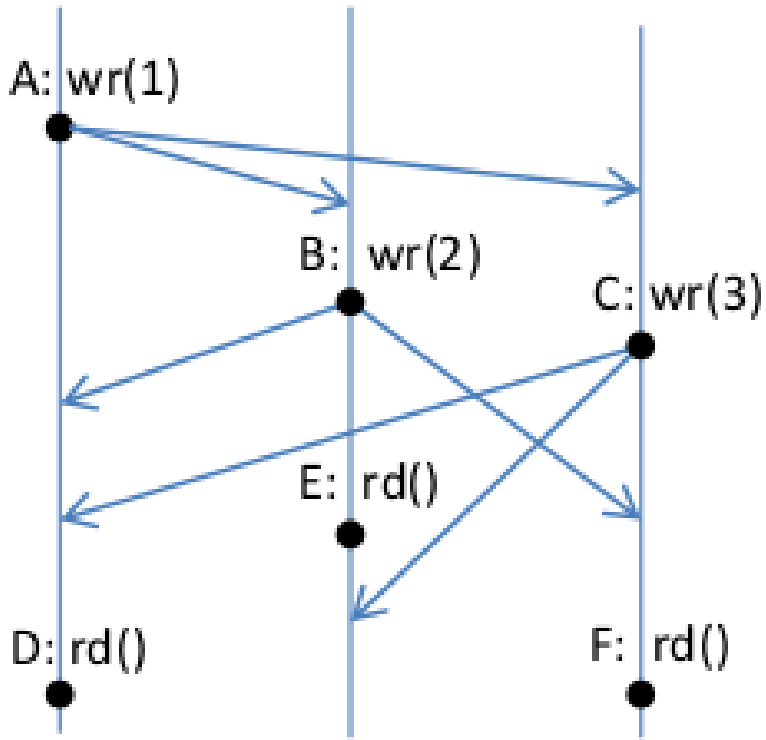

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Ende Hauptteil Horst

Anfang Hauptteil Patrick

Ende Hauptteil Patrick

Conclusion

- Which problems does the techreport solve?
- What is not solved by it?
- What do **we** think about it?