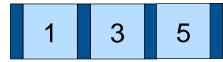
Database Systems, Even 2020-21



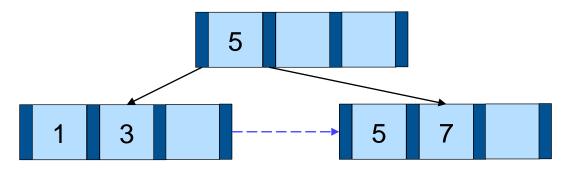
- Order: 4
- Insert: 1, 3, 5, 7, 9, 2, 4, 6, 8, 10



Insert 1, 3, 5

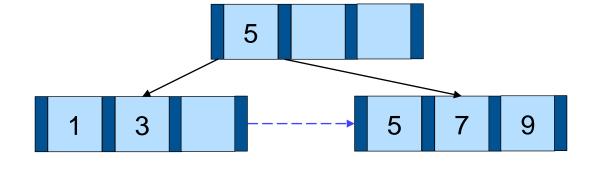


Insert 7

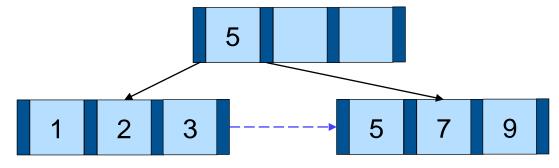


- Order: 4
- Insert: 1, 3, 5, 7, 9, 2, 4, 6, 8, 10

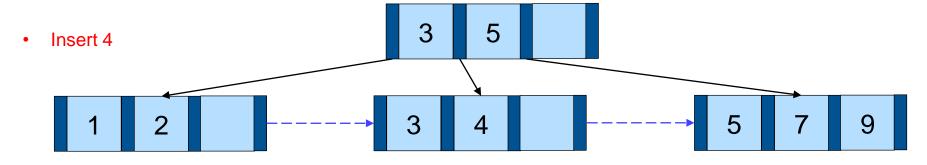




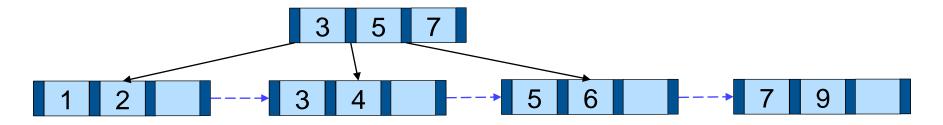
Insert 2



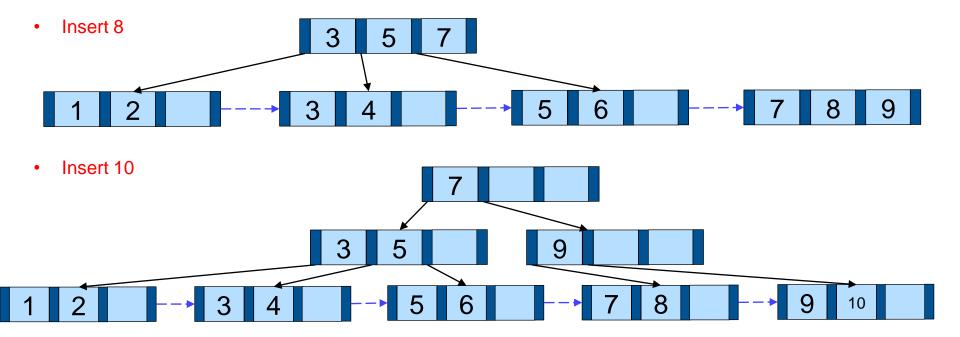
- Order: 4
- Insert: 1, 3, 5, 7, 9, 2, 4, 6, 8, 10



Insert 6



- Order: 4
- Insert: 1, 3, 5, 7, 9, 2, 4, 6, 8, 10

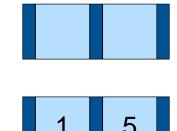


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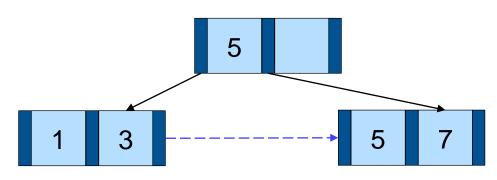


- Order: 3
- Insert: 1, 5, 3, 7, 9, 2, 4, 6, 8, 10

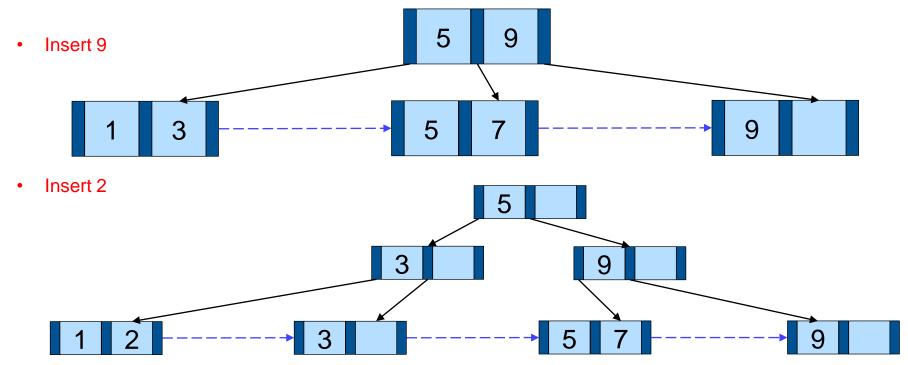
• Insert 1, 5



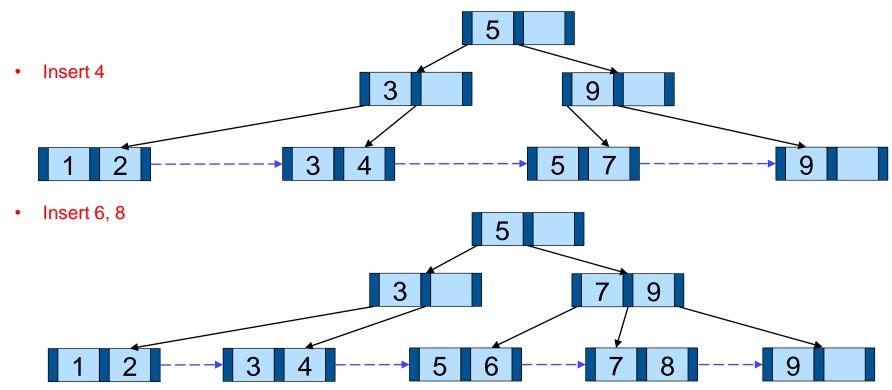
• Insert 3, 7



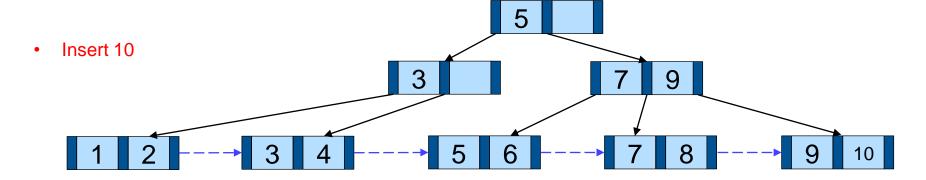
- Order: 3
- Insert: 1, 5, 3, 7, 9, 2, 4, 6, 8, 10



- Order: 3
- Insert: 1, 5, 3, 7, 9, 2, 4, 6, 8, 10

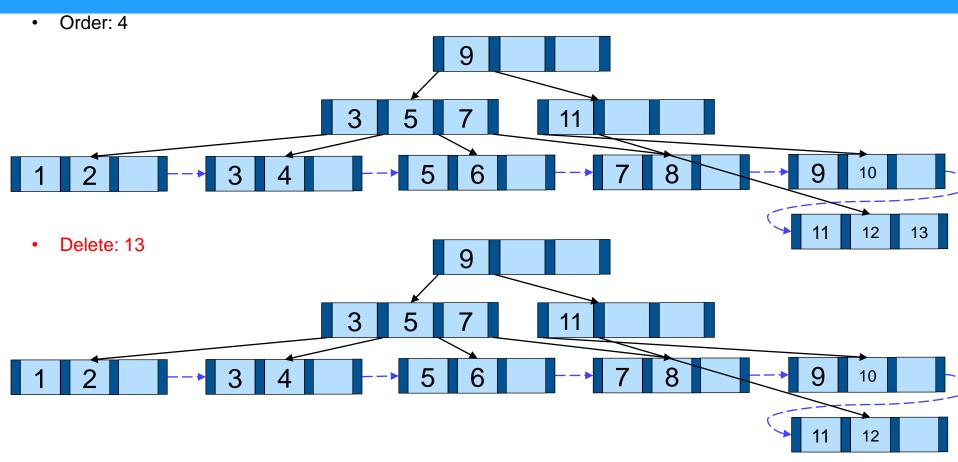


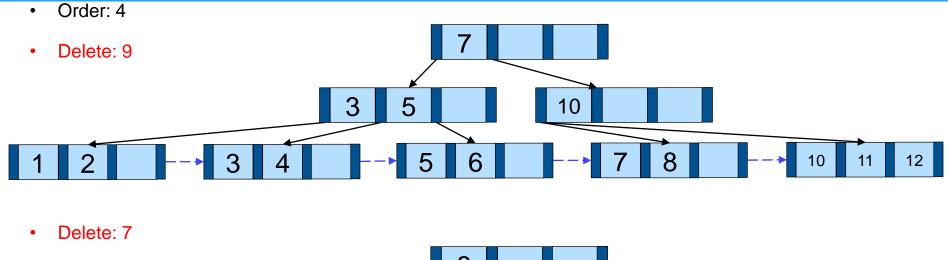
- Order: 3
- Insert: 1, 5, 3, 7, 9, 2, 4, 6, 8, 10

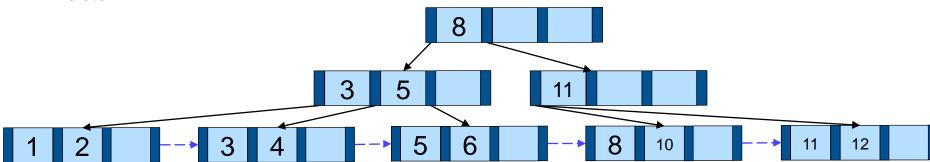


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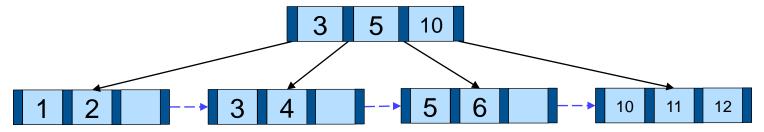








- Order: 4
- Delete: 8

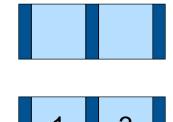


Database Systems, Even 2020-21

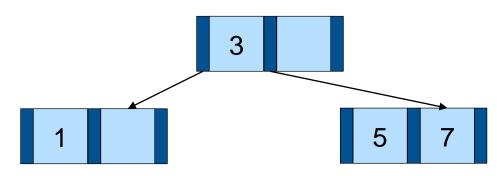


- Order: 3
- Insert: 1, 3, 5, 7, 9, 2, 4, 6, 8, 10

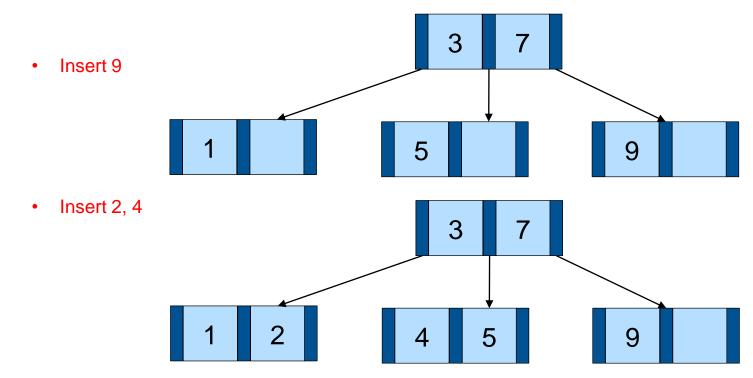
• Insert 1, 3



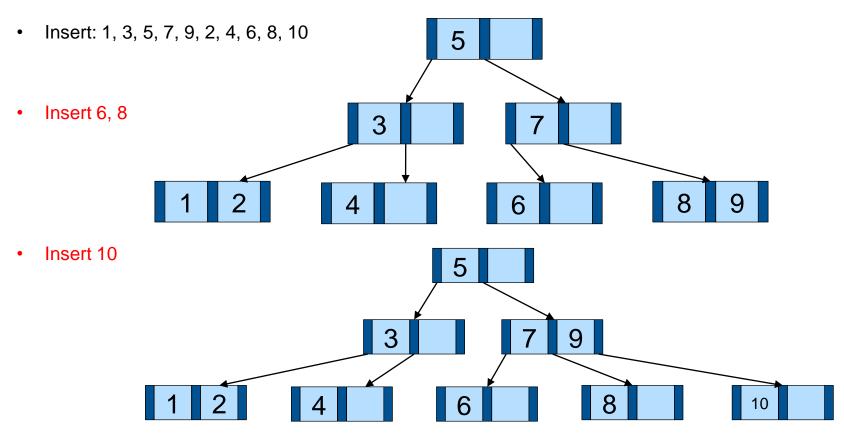
• Insert 5, 7



- Order: 3
- Insert: 1, 3, 5, 7, 9, 2, 4, 6, 8, 10



Order: 3



Database Systems, Even 2020-21



B Tree Deletion

Case-I

- If the key is already in a leaf node, and removing it doesn't cause that leaf node to have too few keys, then simply remove the key to be deleted
- key k is in node x and x is a leaf, simply delete k from x

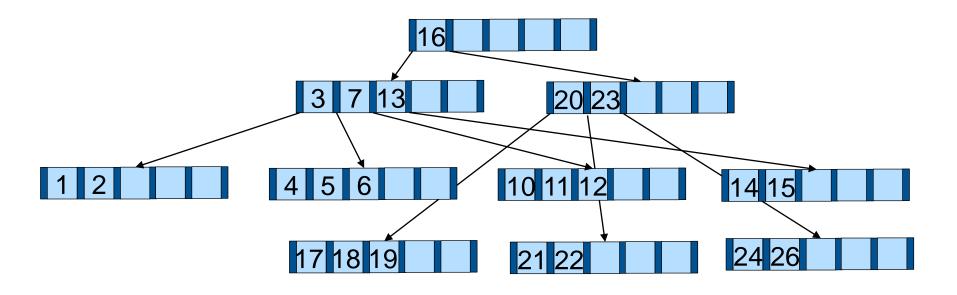
Case-II

- If key **k** is in node x and x is an internal node, there are three cases to consider:
- <u>Case-II-a:</u> If the child *y* that precedes *k* in node *x* has at least *t* keys (more than the minimum), then find the predecessor key *k'* in the subtree rooted at *y*; recursively delete *k'* and replace *k* with *k'* in *x*
- <u>Case-II-b:</u> Symmetrically, if the child *z* that follows *k* in node *x* has at least *t* keys, find the successor *k'* and delete and replace as before; note that finding *k'* and deleting it can be performed in a single downward pass
- <u>Case-II-c:</u> Otherwise, if both y and z have only t-1 (minimum number) keys, merge k and all of z into y, so that both k and the pointer to z are removed from x; y now contains 2t-1 keys, and subsequently k is deleted

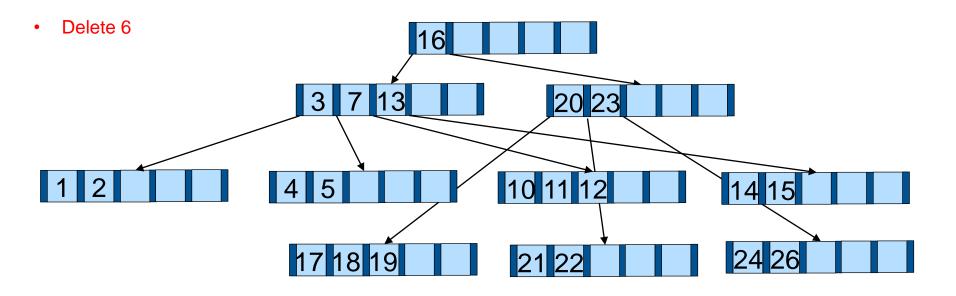
Case-III

- <u>Case-III-a:</u> If the root has only *t*–1 keys but has a sibling with *t* keys, give the root an extra key by moving a key from *x* to the root, moving a key from the roots immediate left or right sibling up into *x*, and moving the appropriate child from the sibling to *x*
- <u>Case-III-b:</u> If the root and all of its siblings have *t*−1 keys, merge the root with one sibling; this involves moving a key down from *x* into the new merged node to become the median key for that node

Order: 6

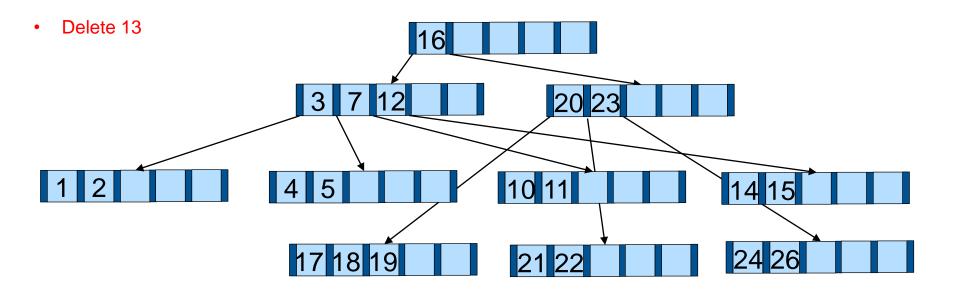


Order: 6



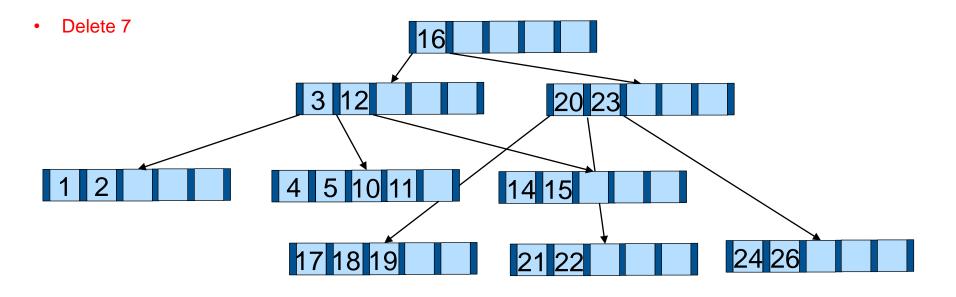
(Case-I)

Order: 6



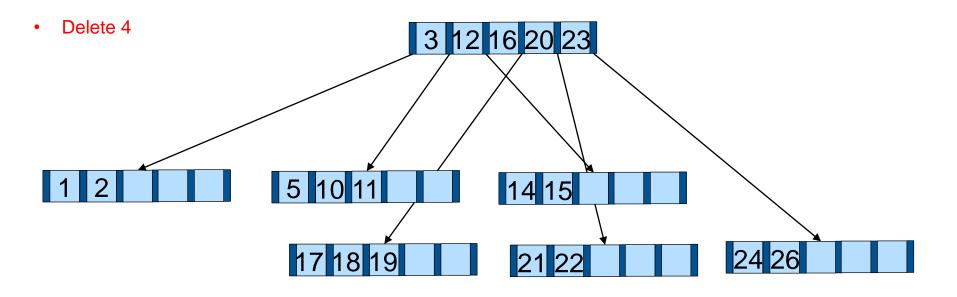
(Case-II-a)

Order: 6



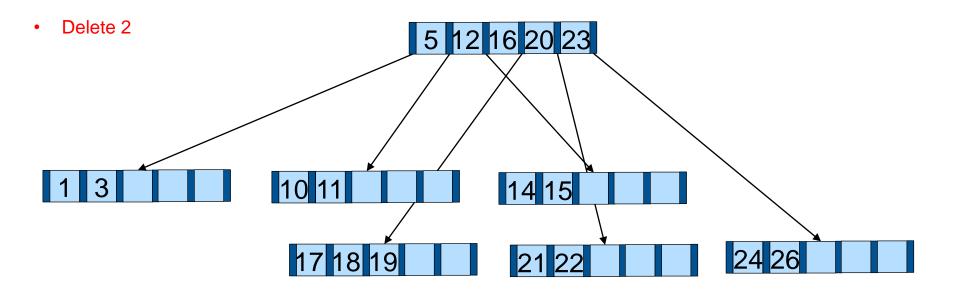
(Case-II-c)

Order: 6



(Case-III-b)

Order: 6



(Case-III-a)

Thank you for your attention...

Any question?

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