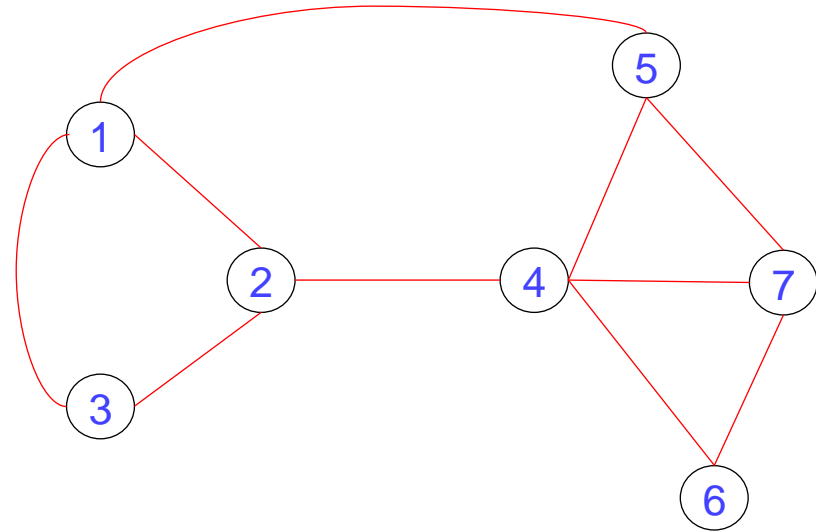
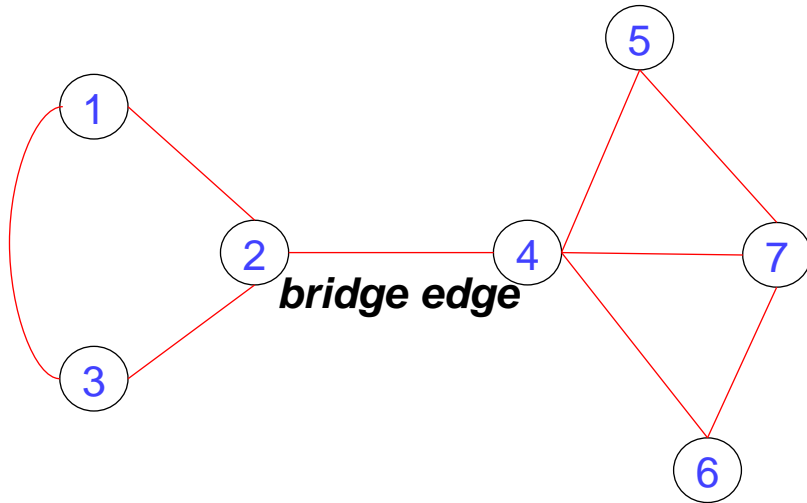




Applications of Depth-First Search

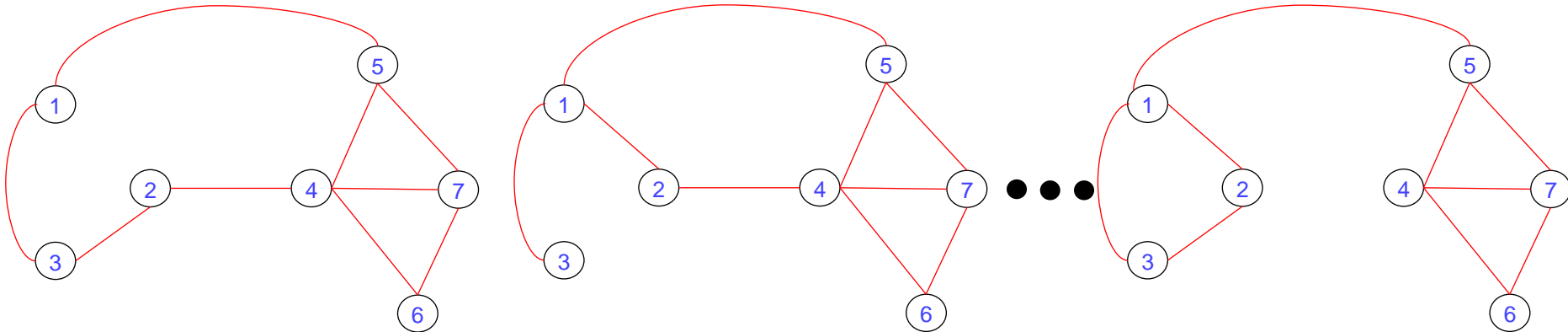
Application: 2-Edge Connected using DFS

- A graph is **2-Edge connected** iff it remains connected after the removal of any one edge
- Graph $G = (V, E)$ is **2-Edge connected** iff $G - \{e\}$ is still connected $\forall e \in E$



Application: 2-Edge Connected using DFS

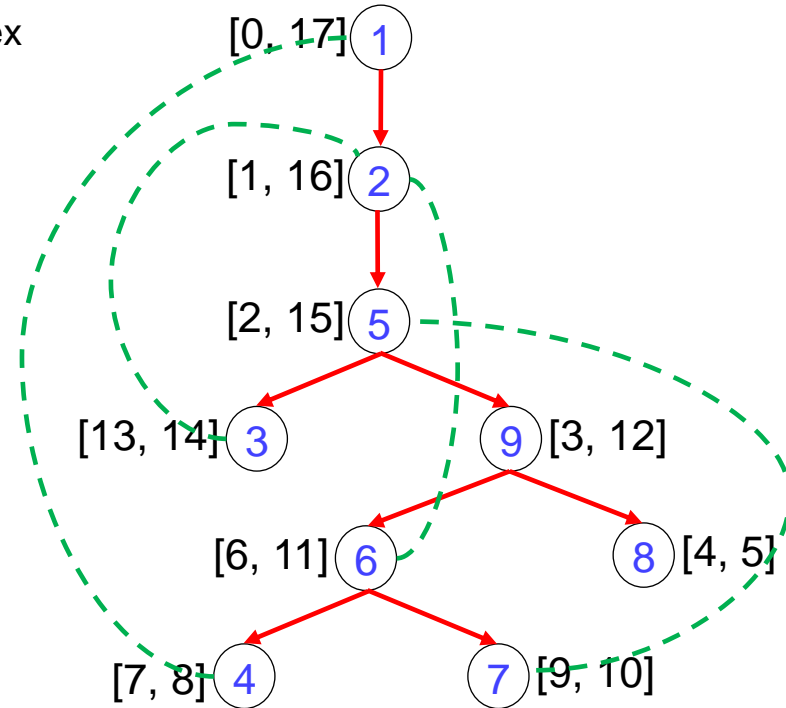
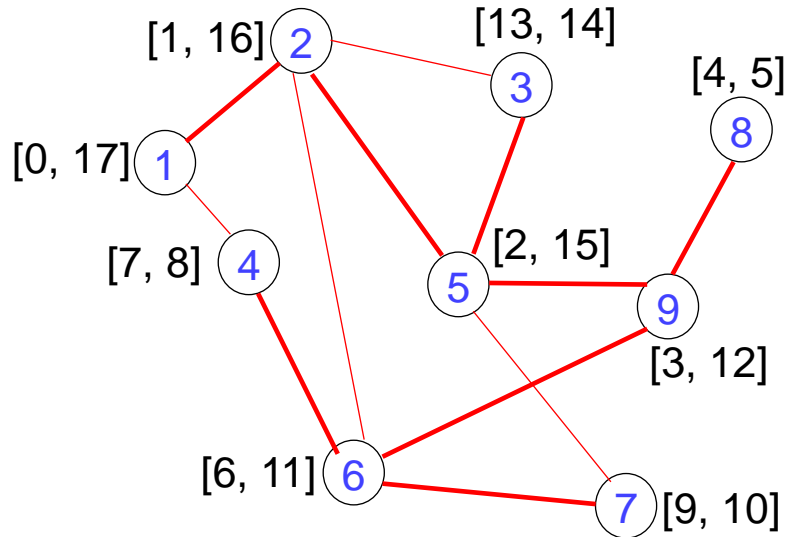
- Given a graph $G = (V, E)$, how will you check if it is a **2-Edge connected** graph?
 - Can be done in $O(E^2)$ time by removing every edge and checking if the resulting graph is connected



- How can we do it in linear time?

Application: 2-Edge Connected using DFS

- Can be done using DFS
 - To check while backtracking from a vertex v , the edge between this vertex and its immediate predecessor vertex (parent in DFS tree) is not a bridge
 - How to ensure that it is not a bridge?
 - At least there is one **back edge** from v or the descendent tree of v to any of its ancestors

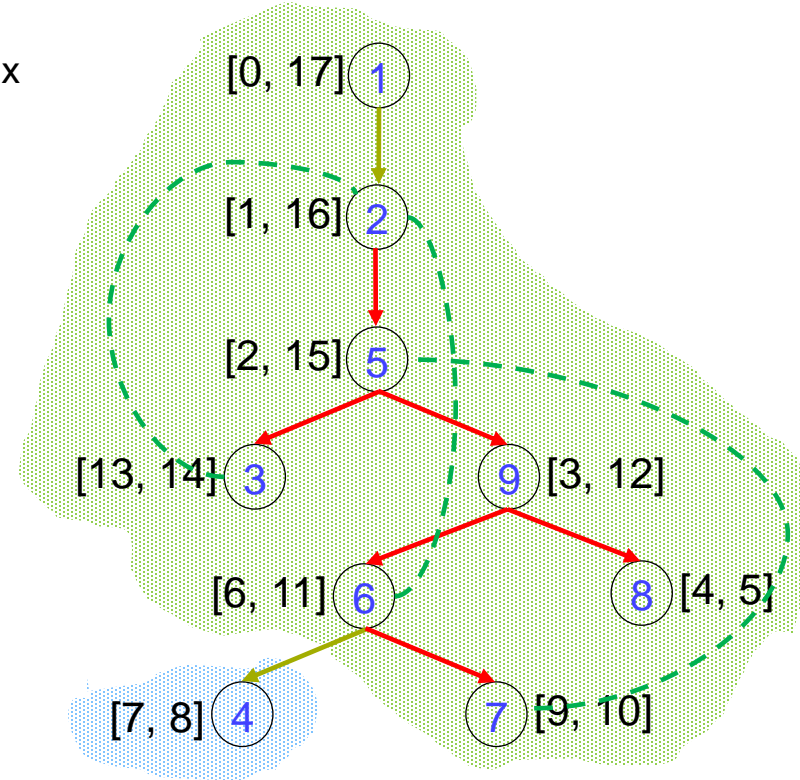
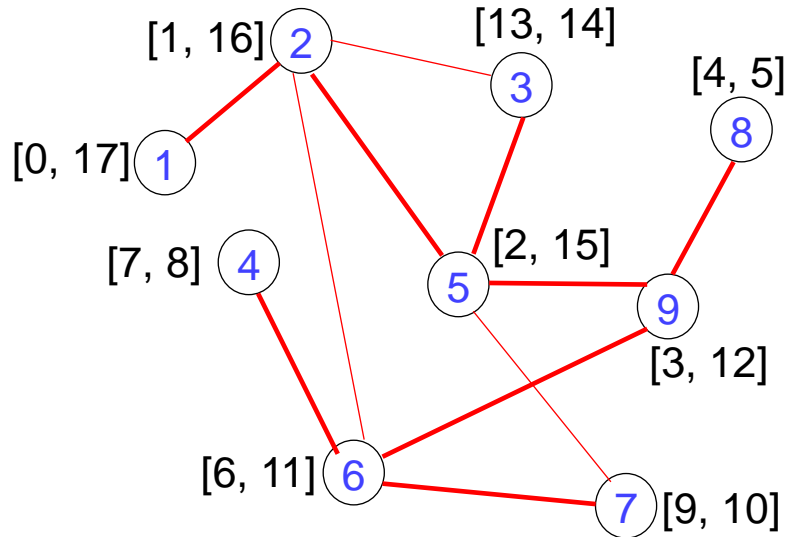


-
- A graph with 9 nodes labeled 1 through 9, each with an associated interval. The nodes are connected by red lines. The intervals are: 1: [0, 17], 2: [1, 16], 3: [13, 14], 4: [7, 8], 5: [2, 15], 6: [6, 11], 7: [9, 10], 8: [4, 5], 9: [3, 12]. The connections are: 1-2, 2-3, 2-5, 2-6, 3-5, 4-6, 5-6, 5-7, 5-9, 6-7, 6-9, 7-9, 8-9.



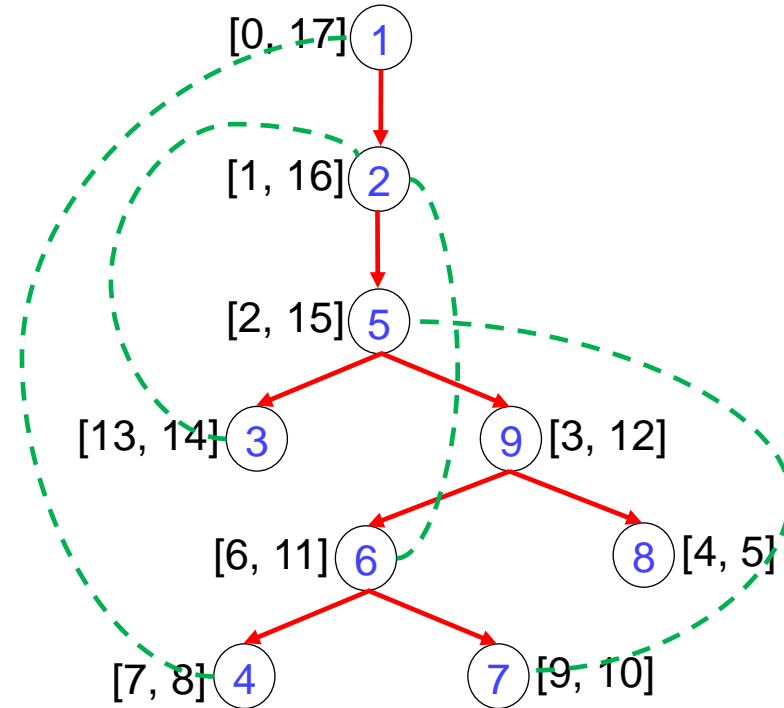
Application: 2-Edge Connected using DFS

- Can be done using DFS
 - To check while backtracking from a vertex v , the edge between this vertex and its immediate predecessor vertex (parent in DFS tree) is not a bridge
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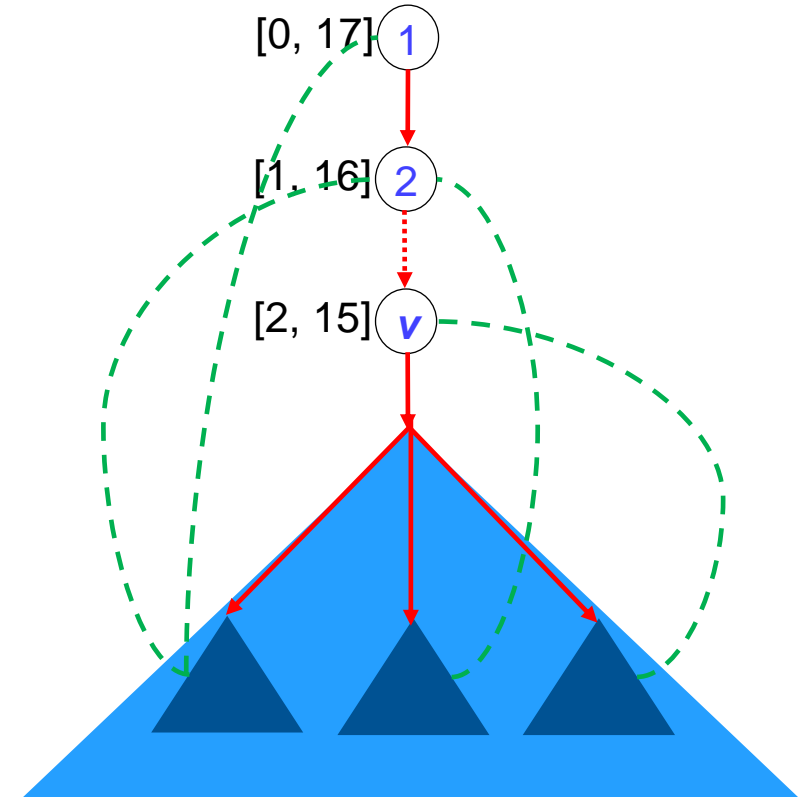
Application: 2-Edge Connected using DFS

- Do we need to keep track of all back edges?
- If we keep track of every back edge then we are going to be spending a lot of time
- The back edge which is going to a node which is closest to the root is of our interest:
deepest back edge
- How can I figure out which is the ***deepest back edge***?
 - By looking at the arrival time of the other end point, as it is closest to the root



Application: 2-Edge Connected using DFS

- **Deepest back edge**
 - All the subtree rooted at v may have **back edges** which can be explored recursively using DFS
 - We must take the one with minimum arrival time (including **back edges** from v)



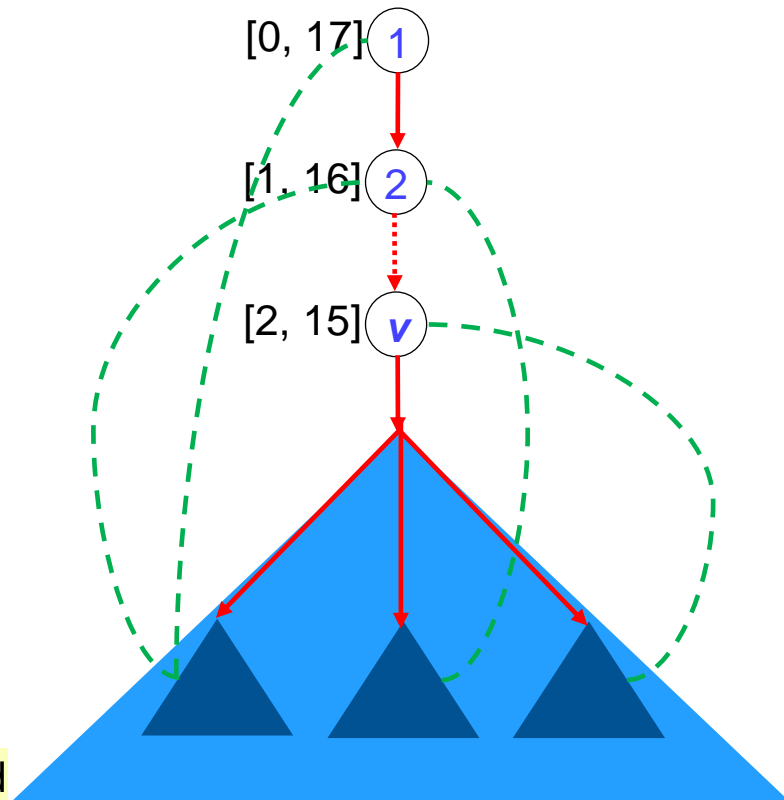
Application: 2-Edge Connected using DFS

```

time = 0;
2EC(v)
{
    visited[v] = 1;
    arr[v] = time++;
    deepest_BE = arr[v];
    for (for all vertex u adjacent from v) do
        if !visited[u] then
            deepest_BE = min(deepest_BE, 2EC(u));
        else
            deepest_BE = min(deepest_BE, arr[u]);
    if deepest_BE = arr[v] then
        return deepest_BE;
}

```

A bridge is found



Application: 2-Edge Connected using DFS

```

time = 0;
2EC(v)
{
    ....
    ....
    if !visited[u] then
        (v, u) is a tree edge;
        deepest_BE = min(deepest_BE, 2EC(u));
    else
        if !((v, u) is a tree edge)
            deepest_BE = min(deepest_BE, arr[u]);
    ....
    ....
}
```

- Two special cases:
 - If (**v**, **u**) is a tree edge
 - If **v** is the starting vertex

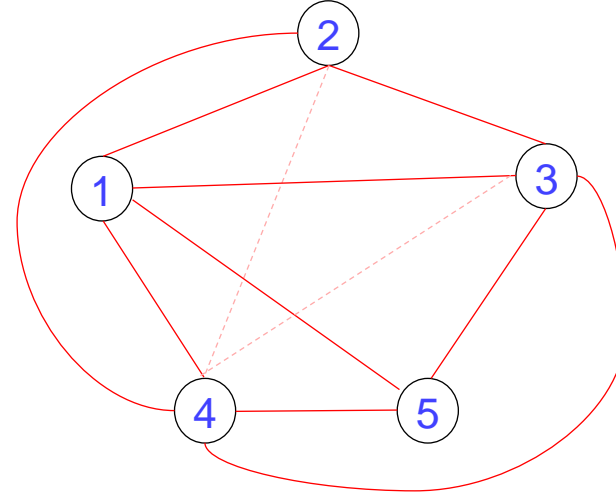
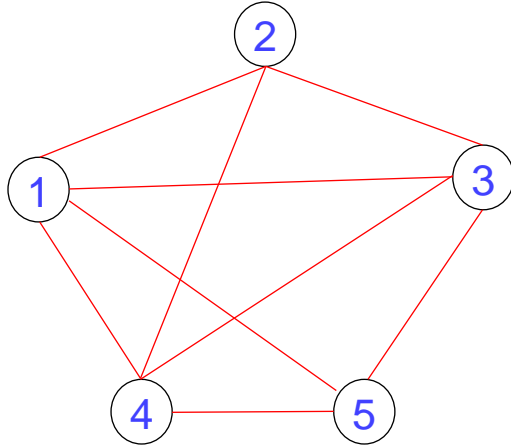
```

time = 0;
2EC(v)
{
    ....
    ....
    if deepest_BE = arr[v] && v != starting vertex then
        return deepest_BE;
}
```

- Time complexity
 - Same as DFS, $O(V + E)$
 - If the graph is connected, then $O(E)$

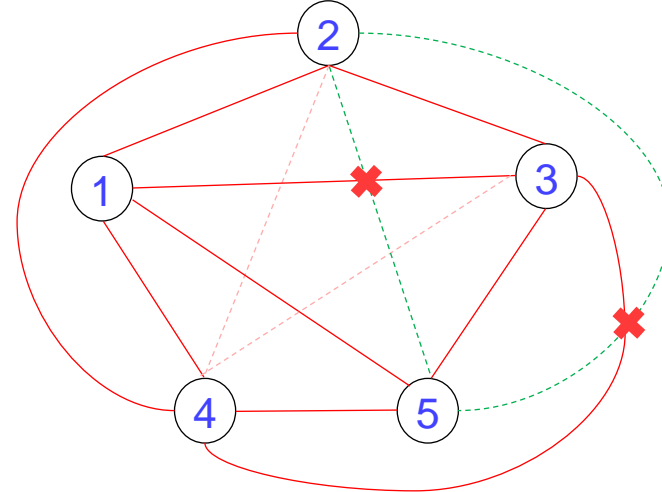
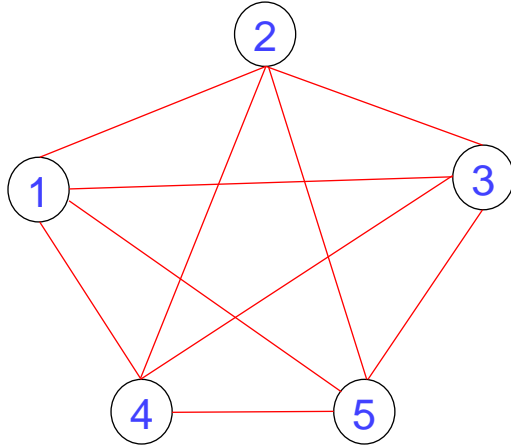
Application: Planner Graph using DFS

- Given a graph $G = (V, E)$, how will you check if it is a **planner graph**?
- A **planar graph** is a graph which can be drawn in the plane such that its edges do not intersect: planer drawing of a graph
- We can draw it whichever way I want but the edges should not intersect
- Is this graph a planer graph?



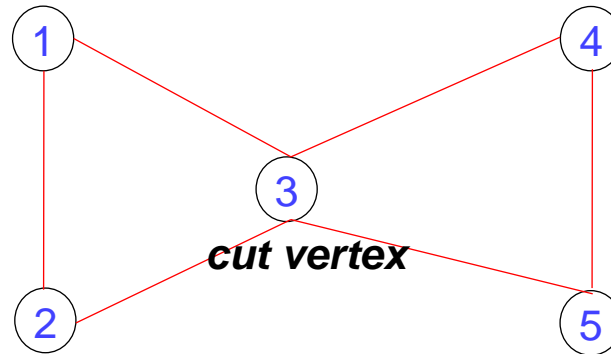
Application: Planner Graph using DFS

- Now let us consider the complete graph with 5 vertices
- Is this graph a **planer graph**?
- DFS algorithm can be used to check if a given graph is a **planer graph**



Application: 2-Vertex Connected using DFS

- A graph is **2-vertex connected**, if removing any vertex still keeps the graph connected
- In removing a vertex, we also have to remove the edges incident to that vertex
- Is this graph **2-edge connected**?
- Is this graph **2-vertex connected**?
- DFS algorithm can be used to check if a given graph is a **2-vertex connected**



Next Lecture

Depth-First Search in Directed Graphs

Thank you for your attention...

Any question?

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