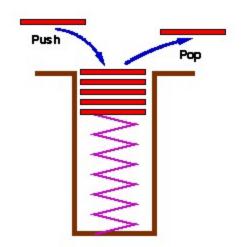
Stacks, Lists

Stacks

- Stacks are a special form of collection with LIFO semantics
- Two methods
 - int push(Stack s, void *item);
 - add item to the top of the stack
 - void *pop(Stack s);
 - remove an item from the top of the stack
- Like a plate stacker
 - Other methods

```
int IsEmpty( Stack s );
/* Return TRUE if empty */
void *Top( Stack s );
/* Return the item at the top,
    without deleting it */
```



Dr. Prakash Raghavendra – IT202- Data Structures and Algorithms-I

Stacks - Relevance

- Stacks appear in computer programs
 - Key to call / return in functions & procedures
 - Stack frame allows recursive calls
 - Call: push stack frame
 - Return: pop stack frame
- Stack frame
 - Function arguments
 - Return address
 - Local variables

Stack Frames - Functions in HLL

Program

```
function f( int x, int y) {
    int a;
                                               Stack
                                                                 parameters
                                               frame
    if ( term cond ) return ...;
                                                                 return address
                                                for f
    a = ....;
                                                                 local variables
    return g(a);
                                                                 parameters
                                               Stack
                                                                 return address
                                               frame
                                                                 local variables
                                                for g
function g( int z ) {
    int p, q;
                                                                 parameters
                                               Stack
                                               frame
                                                                 return address
    p = ....; q = ....;
                                                for f
                                                                 local variables
    return f(p,q);
                         Context
                   for execution of f
       Dr. Prakash Ragnavenura – 11202- Data Structures and Algorithms-I
```

Recursion

- Very useful technique
 - Definition of mathematical functions
 - Definition of data structures
 - Recursive structures are naturally processed by recursive functions!

Recursion

- Very useful technique
 - Definition of mathematical functions
 - Definition of data structures
 - Recursive structures are naturally processed by recursive functions!
- Recursively defined functions
 - Factorial
 - Fibonacci
 - Games
 - Towers of Hanoi
 - Chess

Recursion - Example

Fibonacci Numbers

Pseudo-code

```
fib( n ) = if ( n = 0 ) then 1
else if ( n = 1 ) then 1
else fib(n-1) + fib(n-2)
```

C

```
int fib( n ) {
   if ( n < 2 ) return 1;
   else return fib(n-1) + fib(n-2);
  }</pre>
```

Simple, elegant solution!

Recursion - Example

Fibonacci Numbers

```
C int fib(n) {
    if case, a run-time disaster!!!!

if case, a run-time disaster!!!!

But, in the Fibonacci case, a run-time disaster!!!!
```

However, many recursive functions, eg binary search, are simple, elegant and efficient!

Array Limitations

- Arrays
 - Simple,
 - Fast

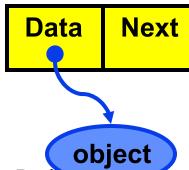
but

- Must specify size at construction time
- Murphy's law
 - Construct an array with space for n
 - *n* = twice your estimate of largest collection
 - Tomorrow you'll need n+1
- More flexible system?

- Flexible space use
 - Dynamically allocate space for each element as needed
 - Include a pointer to the next item

□ Linked list

- Each node of the list contains
 - the data item (an object pointer in our ADT)
 - a pointer to the next node



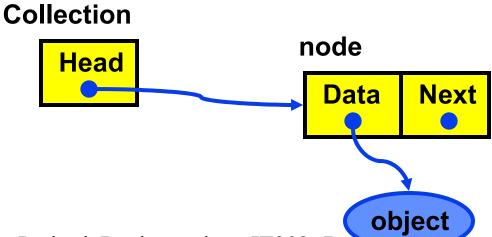
Dr. Prakası — IT202- Data Structures and Algorithms-I

- Collection structure has a pointer to the list head
 - Initially NULL

Collection



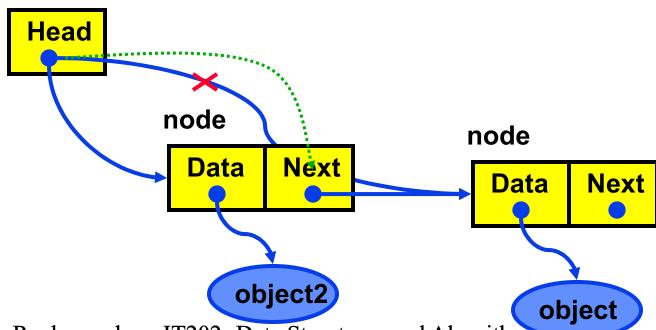
- Collection structure has a pointer to the list head
 - Initially NULL
- Add first item
 - Allocate space for node
 - Set its data pointer to object
 - Set Next to NULL
 - Set Head to point to new node



Dr. Prakash Raghavendra – IT202- Data Survey and Algorithms-I

- Add second item
 - Allocate space for node
 - Set its data pointer to object
 - Set Next to current Head
 - Set Head to point to new node

Collection



Dr. Prakash Raghavendra – IT202- Data Structures and Algorithms-

Linked Lists - Add implementation

Implementation

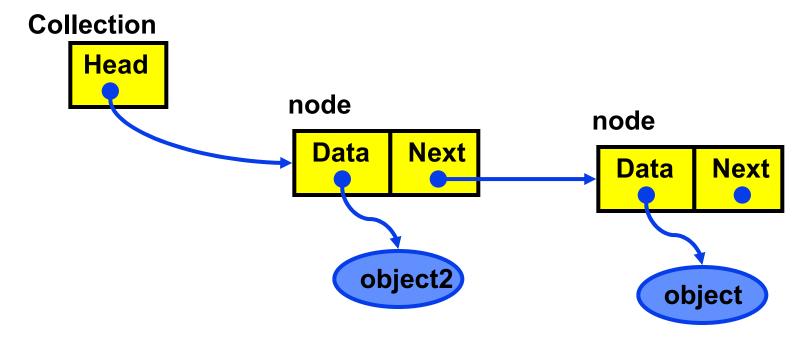
```
struct t node {
    void *item;
    struct t node *next;
    } node;
typedef struct t node *Node;
struct collection {
    Node head;
    };
int AddToCollection( Collection c, void *item ) {
    Node new = malloc( sizeof( struct t node ) );
    new->item = item:
    new->next = c->head;
    c->head = new;
    return TRUE;
 Dr. Prakash Raghavendra – IT202- Data Structures and Algorithms-I
```

Linked Lists - Add implementation

Implementation

```
struct t node {
    void *item;
                                Recursive type definition -
    struct t node *next;
                                       C allows it!
    } node;
typedef struct t node *Node;
struct collection {
    Node head;
int AddToCollection( Collection c, void *item )
    Node new = malloc( sizeof( struct t node )
    new->item = item;
    new->next = c->head;
    c->head = new;
                                    Error checking, asserts
    return TRUE;
                                      omitted for clarity!
 Dr. Prakash Raghavendra – IT202- Data Structures and Algorithms-I
```

- Add time
 - Constant independent of n
- Search time
 - Worst case n



Dr. Prakash Raghavendra – IT202- Data Structures and Algorithms-I

Linked Lists - Find implementation

Implementation

```
void *FindinCollection( Collection c, void *key ) {
   Node n = c->head;
   while ( n != NULL ) {
       if ( KeyCmp( ItemKey( n->item ), key ) == 0 ) {
            return n->item;
       n = n->next;
       }
   return NULL;
   }
```

• A recursive implementation is also possible!

Dr. Prakash Raghavendra – IT202- Data Structures and Algorithms-I

Linked Lists - Delete implementation

Implementation

```
void *DeleteFromCollection( Collection c, void *key ) {
   Node n, prev;
   n = prev = c->head;
   while ( n != NULL ) {
      if (KeyCmp(ItemKey(n->item), key) == 0) {
            prev->next = n->next;
            return n;
     prev = n;
      n = n->next;
                     head
    return NULL;
```

Dr. Prakash Raghavendra – IT202- Data Structures and Algorithms-I

Linked Lists - Delete implementation

Implementation

```
void *DeleteFromCollection( Collection c, void *key ) {
    Node n, prev;
    n = prev = c->head;
    while ( n != NULL ) {
      if (KeyCmp(ItemKey(n->item), key) == 0) {
            prev->next = n->next;
            return n;
      prev = n;
      n = n->next;
                     head
    return NULL;
                        Minor addition needed to allow
```

Dr. Prakash Raghavendr for deleting this one! An exercise!

Linked Lists - LIFO and FIFO

- Simplest implementation
 - Add to head
 - □ Last-In-First-Out (LIFO) semantics
- Modifications
 - First-In-First-Out (FIFO)

```
* Keep a tail pointer

struct t_node {
    void *item;
    struct t_node *next;
    } node;

typedef struct t_node *Node;

struct collection {
    Node head, tail;
    Dr. Prakash Raghavendra - IT202- Data Structures and Augustanas I
```

Linked Lists - Doubly linked

- Doubly linked lists
 - Can be scanned in both directions

```
struct t node {
    void *item;
    struct t node *prev,
                    *next;
    } node;
typedef struct t node *Node;
struct collection {
    Node head, tail;
                 head
    };
                             prev
                                        prev
                                                    prev
```

Dr. Prakash Raghavendra – IT202- Data Structures and Algorithms-I