

# Distance Vector Routing Algorithm

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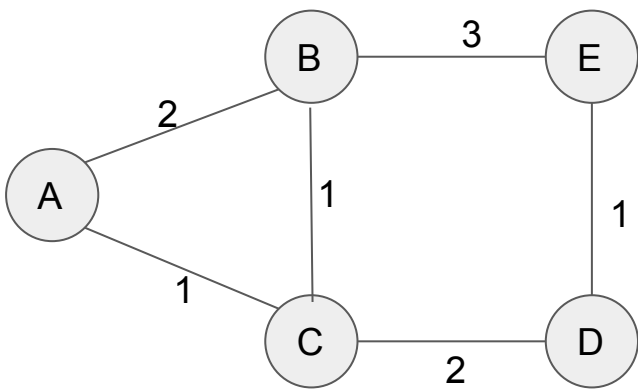
- Dynamic Routing Algorithm
- Finds the least cost path to reach the destination.
- Each router maintains a routing table
  - One entry for each router.
  - Routing Table Size - directly proportional to the No.of routers in the network.
- Also called as Belman-Ford or Ford Fulkerson algorithms.
- Introduced in 1957 - Revised in 1962.
- Used in the internet in the earlier days, but , not now.
  - Why ?

# All About DV Routing Algorithm

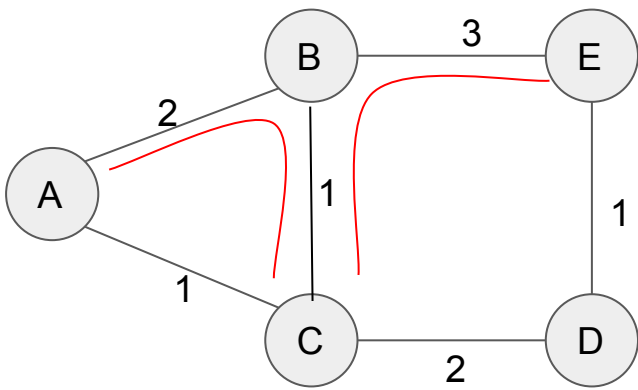
- The router contains an entry per router in the routing table
  - Distance to reach the the routers
  - Distance is maintained as a Vector, Hence Distance Vector Routing Algorithm
  - Distance - Queue length, physical distance, delay...
  - Hence, also called as Routing Information Protocol (RIP).
  - Initially, cost will be estimated, then will be converged to final cost.
- Each router exchanges its routing table to its neighbors at regular intervals.

Neighbor	Cost
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Routing Table Structure



C Rout. Tab	
A	1
B	1
D	2

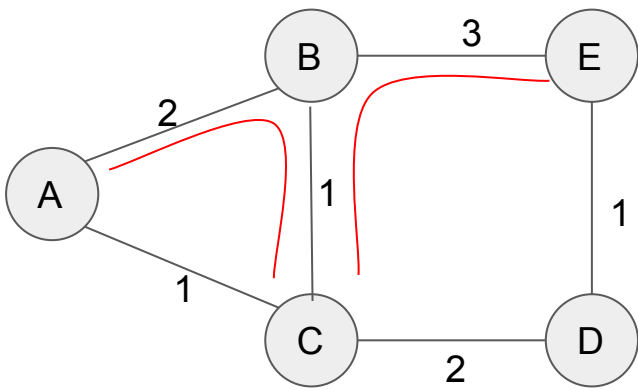


C Rout. Tab	
A	1
B	1
D	2

B Rout. Tab	
A	2
E	3
C	1

$$C \rightarrow B \rightarrow A = 1 + 2 = 3$$

$$C \rightarrow B \rightarrow E = 1 + 3 = 4$$

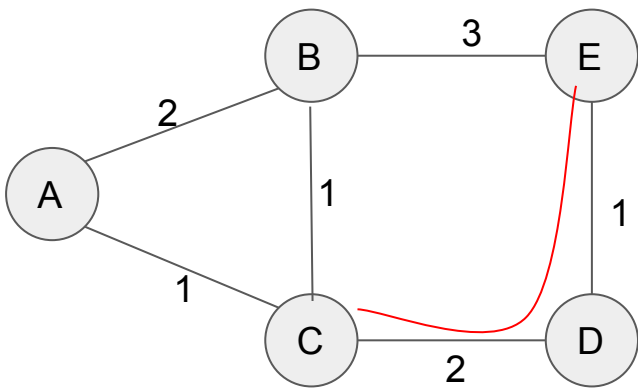


C Rout. Tab	
A	1
B	1
D	2
E	4

B Rout. Tab	
A	2
E	3
C	1

$C \rightarrow B \rightarrow A = 1 + 2 = 3$  No Changes will be done for  $C \rightarrow A$  as  $3 > 1$

$C \rightarrow B \rightarrow E = 1 + 3 = 4$  A New Entry will be made for E

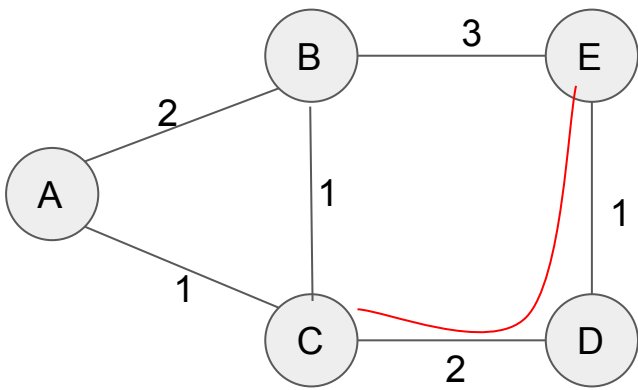


C Rout. Tab	
A	1
B	1
D	2
E	4

B Rout. Tab	
A	2
E	3
C	1

D Rout. Tab	
E	1
C	2

$$C \rightarrow D \rightarrow E = 2 + 1 = 3$$



C Rout. Tab	
A	1
B	1
D	2
E	<del>4</del> 3

B Rout. Tab	
A	2
E	3
C	1

D Rout. Tab	
E	1
C	2

$$C \rightarrow D \rightarrow E = 2 + 1 = 3$$

Entry for E is changed to 3;

$3 < 4$  , smaller cost to reach E