



Sanjivani Rural Education Society's

Sanjivani College of Engineering, Kopargaon-423603

(An Autonomous Institute Affiliated to Savitribai Phule Pune University, Pune)

NAAC 'A' Grade Accredited, ISO 9001:2015 Certified

Department of Information Technology

(NBA Accredited)

Cryptography and Cyber Security

[IT311]



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Unit 3: Message Digest & Key Management

- Hash Algorithms: SHA-1, MD5, Key Management: Introduction, Key Management: Generations, Distribution, Updation, Digital Certificate, Digital Signature, Kerberos 5.0.



Secure Hash Algorithm (SHA)

- Secure Hash Algorithms (SHA) was developed by National Institute of Standards and Technology (NIST) along with NSA
- Published as a Federal Information Processing Standards Publications (FIPS 180 PUBS) in 1993
- A revised version was issued as FIPS PUB 180-1 in 1995 and is referred to as SHA-1
- SHA is a modified version of MD5
- Name of Standard: Secure Hash Signature Standard (SHS)
- In 2002 , NIST produced a revised version of the standard, FIPS 180-2 that defined three new versions of SHA as SHA-256, SHA-384, and SHA-512.



Secure Hash Algorithm (SHA) : Purpose

- Purpose of SHA is authentication and not the encryption
- Verify that received messages come from the alleged source and have not been altered.
- Verify the sequence and timing.
- Digital Signature is used to combat denial of receipt of a message by either the source or destination.
- Impossible to recreate a message given a message digest.



Secure Hash Algorithm (SHA) : Applications

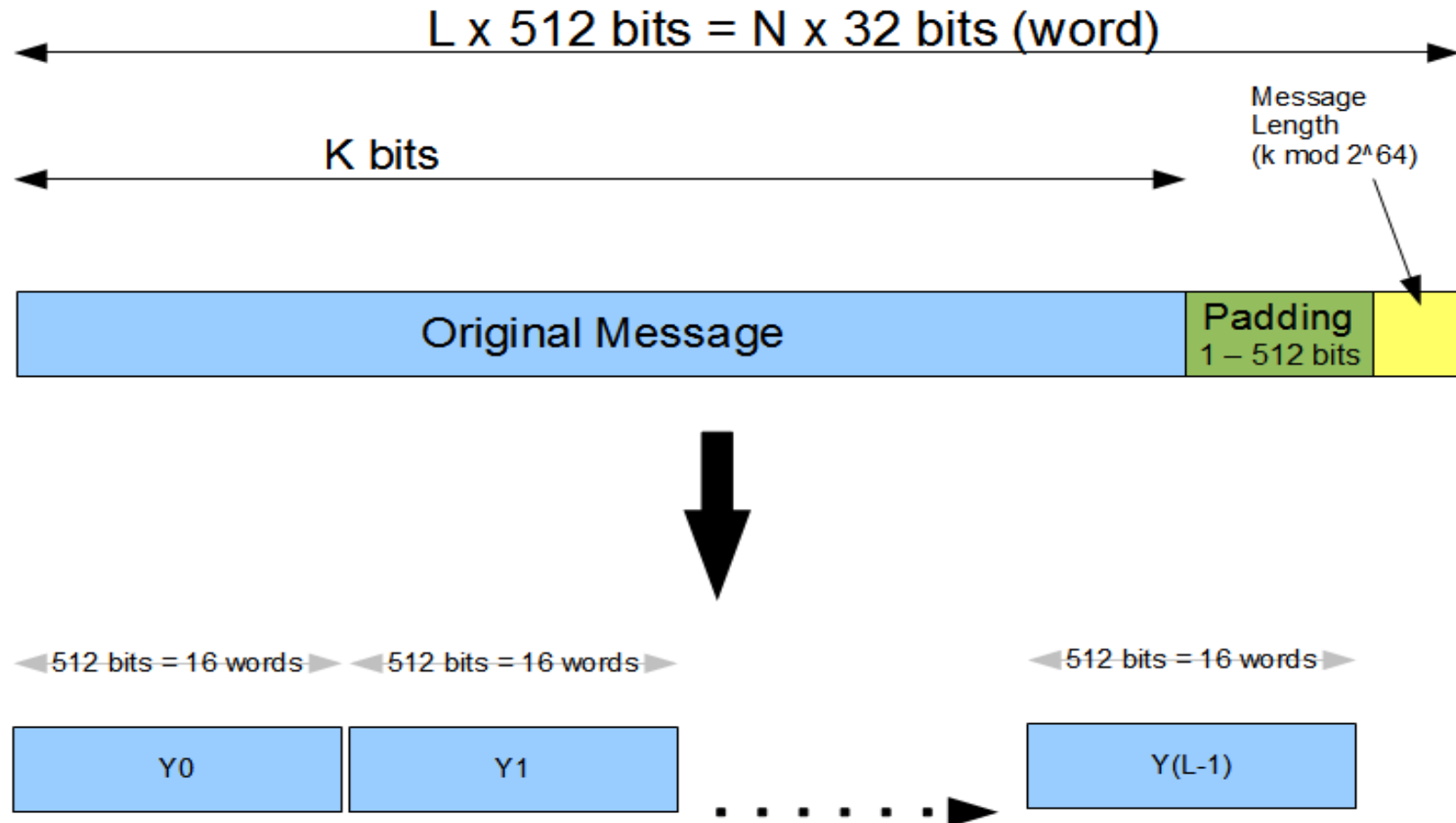
- SHA uses one way hash function. The applications are as follows:
- Public Key Algorithms
 - Password Logins
 - Encryption Key Management
 - Digital Signatures
- Integrity Checking
 - Virus and Malware Scanning
- Authentication
 - Secure Web Connections (PGP, SSL, SSH, S/MIME)



Secure Hash Algorithm (SHA) : Variants

- MD4 and MD5 by Ron Rivest (1990,1994)
- SHA-0, SHA-1 by NSA (1993, 1995)
- RIPEMD-160 (1996)
- SHA-2 (2002 – 224, 256, 385, 512)
- Whirlpool
- Tiger
- GOST-3411
- SHA-3
 - Winner selected from solicitations in 2012

Structure of SHA





Working of SHA

- SHA is closely modeled after MD5
- **Step 1: Padding**
 - To add padding to the end of the original message in such a way that the length of the message is **64 bits short of a multiple of 512**.
 - Like MD5, the padding always added, even if the message is already 64 bits short of a multiple of 512.



Working of SHA

- **Step 2: Append length**
 - The length of the message excluding the length of the padding is calculated and appended to the end of the padding as a 64-bit block.
- **Step 3: Divide the input into 512-bit blocks**
 - The input message is now divided into blocks, of length 512 bits.
 - These blocks become the input to the message digest processing logic.

Working of SHA

- **Step 4: Initialize chaining variables**
 - Five chaining variables A through E
 - In the case of SHA want to produce a message digest of length 160 bits, we need to have **five chaining variables** here ($5 \times 32 = 160$ bits).
 - In SHA, the variables A through D have the same values as they had in MD5

A	Hex	01	23	45	67
B	Hex	89	AB	CD	EF
C	Hex	FE	DC	BA	98
D	Hex	76	54	32	10

- Additionally, **E is initialized to Hex C3 D2 E1 F0.**



Working of SHA

- **Step 5: Process Blocks**
 - **Step 5.1:**
 - Copy the chaining variables A-E into variables a-e.
 - The combination of a-e, called as **abede** will be considered as a **single register** for storing the temporary intermediate as well as the final results.
 - **Step 5.2:**
 - Now, divide the current 512-bit block into 16 sub-blocks, each consisting of 32 bits.



Working of SHA

- **Step 5: Process Blocks**
 - **Step 5.3:** SHA has four rounds, each round consisting of 20 steps.
 - Each round takes three inputs
 - Current 512- bit block
 - Register abcde
 - A constant $K[t]$ (where $t=0$ to 79)
 - It then updates the contents of the register abcde using the SHA algorithm steps.
 - We have only four constants (in case of MD5-64 constants) defined for $K[t]$, one used in each of the four rounds.

Working of SHA

- **Step 5: Process Blocks**
 - **Step 5.3:** We have only four constants (in case of MD5- 64 constants) defined for $K[t]$, one used in each of the four rounds.

Round	Value of t between	$K[t]$ in hexadecimal	$K[t]$ in decimal (Only integer portion of the value shown)
1	1 and 19	5A 92 79 99	$2^{30} \times \sqrt{2}$
2	20 and 39	6E D9 EB A1	$2^{30} \times \sqrt{3}$
3	40 and 59	9F 1B BC DC	$2^{30} \times \sqrt{5}$
4	60 and 79	CA 62 C1 D6	$2^{30} \times \sqrt{10}$



Working of SHA

- **Step 5.4**

- SHA consists of four rounds, each round containing 20 iterations.
- This makes it a total of 80 iterations.
- Mathematically, an iteration consists of the following operations:

$$abcde = (e + \text{Process } P + s^{30}(a) + W[t] + K[t]), a, s^{30}(b), c, d$$

Where,

abcde = The register made up of the five variables a, b, c, d and e

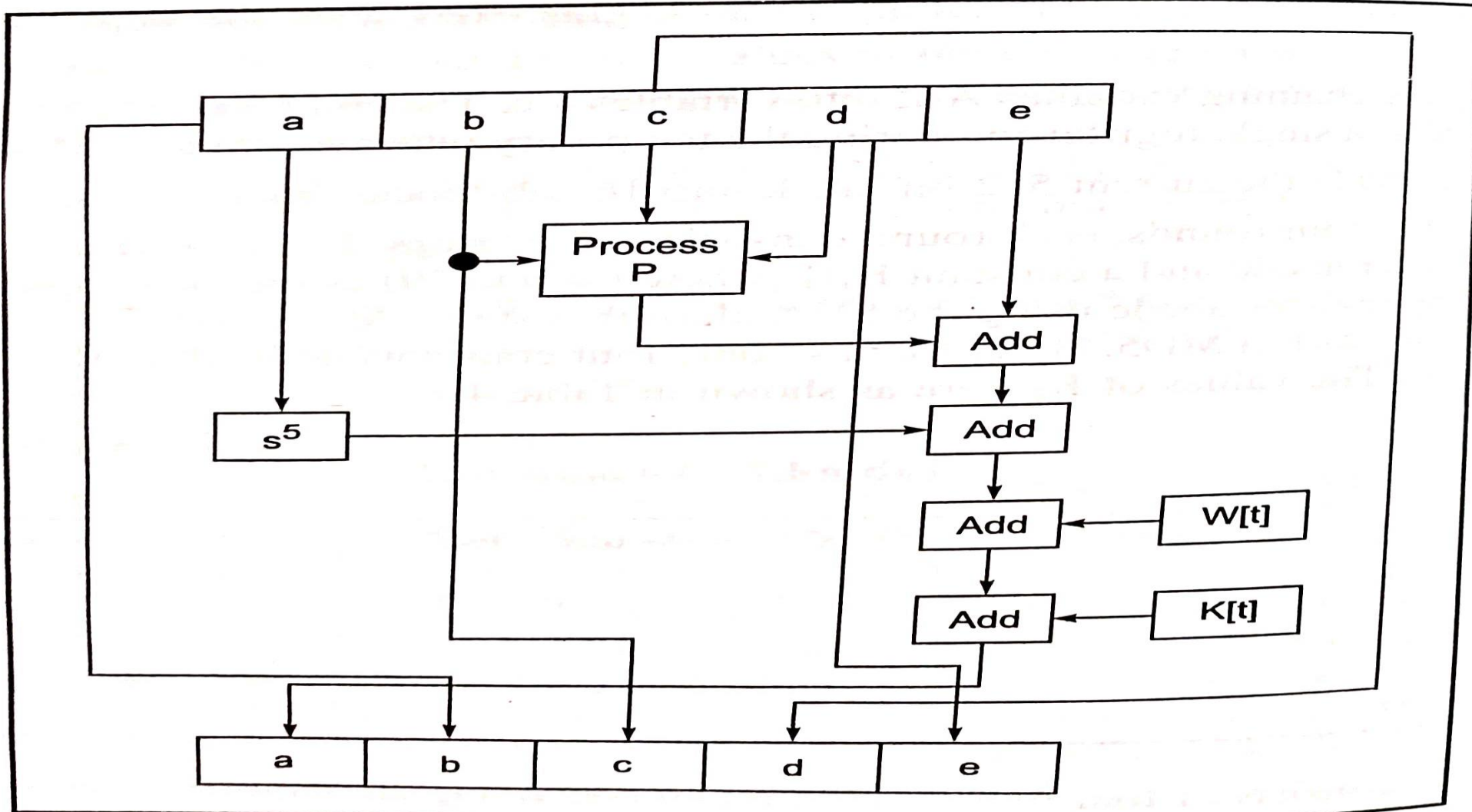
Process P = The logical operation

S^t = Circular-left shift of the 32-bit sub-block by t bits

$W[t]$ = A 32-bit derived from the current 32-bit sub block

$K[t]$ = One of the five additive constants

Working of SHA : Single SHA-1 Iteration



Working of SHA

- **Step 5.4** : Process P in each SHA-1 round

<i>Round</i>	<i>Process P</i>
1	$(b \text{ AND } c) \text{ OR } ((\text{NOT } b) \text{ AND } (d))$
2	$B \text{ XOR } c \text{ XOR } d$
3	$(b \text{ AND } c) \text{ OR } (b \text{ AND } D) \text{ OR } (c \text{ AND } d)$
4	$B \text{ XOR } c \text{ XOR } d$



Working of SHA

- **Step 5.4**

- The values of $W[t]$ can be calculated as follows:
- For the first 16 words of W (ie. $t = 0$ to 15), the contents of the input message sub-block $M[t]$ become the contents of $W[t]$ straightaway.
- That is, the first 16 blocks of the input message M copied to W .
- The remaining 64 values of W are derived using the equation:

$$W[t] = s' (W[t-16] \text{ XOR } W[t-14] \text{ XOR } W[t-8] \text{ XOR } W[t-3])$$

s' indicates a circular-left shift (i.e. rotation) by 1 bit position.



Cryptanalysis and Limitation

- Key Premises for Hash Functions:
 - Impossible to re-create a message given a fingerprint
 - Collision Free
- SHA-1 failure using brute force attack in 2^{80} operations

Comparison of SHA Parameters

	SHA-1	SHA-256	SHA-384	SHA-512
Message digest size	160	256	384	512
Message size	$< 2^{64}$	$< 2^{64}$	$< 2^{128}$	$< 2^{128}$
Block size	512	512	1024	1024
Word size	32	32	64	64
Number of steps	80	64	80	80
Security	80	128	192	256

- Notes:
1. All sizes are measured in bits.
 2. Security refers to the fact that a birthday attack on a message digest of size n produces a collision with a workfactor of approximately $2^{n/2}$.



Difference Between MD5 and SHA-1

Sr. No.	Points of Discussion	MD5	SHA-1
1	Message digest length in bits	128	160
2	Attack to try and find the original message given a message digest	Requires 2^{128} operations to break in	Requires 2^{160} operations to break in. more secure
3	Attack to try and find two messages producing the same message digest	Requires 2^{64} operations to break in	Requires 2^{80} operations to break in
4	Successful attacks so far	Attempts reported so far	No reported yet
5	Speed	Faster (64 iterations and 128-bit buffer)	Slower (80 iterations and 160-bit buffer)
6	Software implementation	Simple. Does not need any large programs or complex tables	Simple. Does not need any large programs or complex tables



References:

- Atul Kahate, "Cryptography and Network Security", second edition, Tata McGraw Hill
- William Stallings, "Cryptography and Network Security-Principles and practice"