ALGORITHMS AND DATA STRUCTURES II

Lecture 1
Algorithms and their Complexity

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Course Overview

- **Schedule:**

4/09 – L1,	5/07 – MidTerm,	5/31 – L12,
4/12 – L2,	5/10 – L7,	6/04 – L13,
4/16 – L3,	5/14 – L8,	6/XX – Final
4/19 – L4,	5/21 – L9,	
4/23 – L5	5/24 – L10,	
4/26 – L6,	5/28 – L11,	

- **Exams:**
  - MidTerm – Lectures 1 to 6.
  - Final – Lectures 7 to 12.

Course Overview

- **Grading**
  - Exercises – 40%
  - MidTerm Exam– 30%
  - Final Exam – 30%

- **Exercises**
  - Text problems.
  - Programming tasks.

Course Management

- Using **Moodle** system.
- Need an account.
- Exercises downloaded from **Moodle.**
- Answers uploaded to **Moodle.**
- Scores, comments – from **Moodle.**

Today's Outline

- **Algorithms:**
  - Definition.
  - Basic concepts.
- **Function growth.**
  - Upper bound.
  - Lower bound.
  - Tight bound.
- **Algorithm complexity.**
- **Merge sort algorithm.**

Algorithms

- To solve any problem by a computer, we need an algorithm.
- Given an algorithm for the problem, we want to know the efficiency the algorithm.
- We are most interested in how much time and how much memory ***space*** the algorithm takes to solve the problem.

Algorithms

- What is an algorithm?

An <u>algorithm</u> is a well-defined computational **procedure** that transforms inputs into outputs, achieving the desired input-output relationship.

ALGORITHMS

- The **computation time** of an algorithm depends on the number of computational steps of the algorithm and the computer used.
- To evaluate the efficiency of algorithms, it is ideal to use an **unique computer** to measure their computation time.

Algorithms

- The computation time of an algorithm for a problem **depends** on the size of the problem.
- Important! How the computation time of the algorithm grows when the size of the problem increases.

Algorithms

- The size of a problem is denoted by an integer ***n***, which is a measure of the quantity of input data.
  - The size of a matrix multiplication problem might be the **largest dimension** of the matrices.
  - The size of a sorting problem might be the **number of data** to be sorted.
  - The size of a graph problem might be the **number of vertices** or edges.

Algorithms Complexity

- The computation time needed by an algorithm expressed as a function of the size of a problem is called **time complexity** of the algorithm.
- Analogous definition can be made for **space complexity.**

Algorithms Complexity

- Given an algorithm for a problem of size ***n***, it is important to find the time complexity and how the time complexity grows when ***n*** increases.

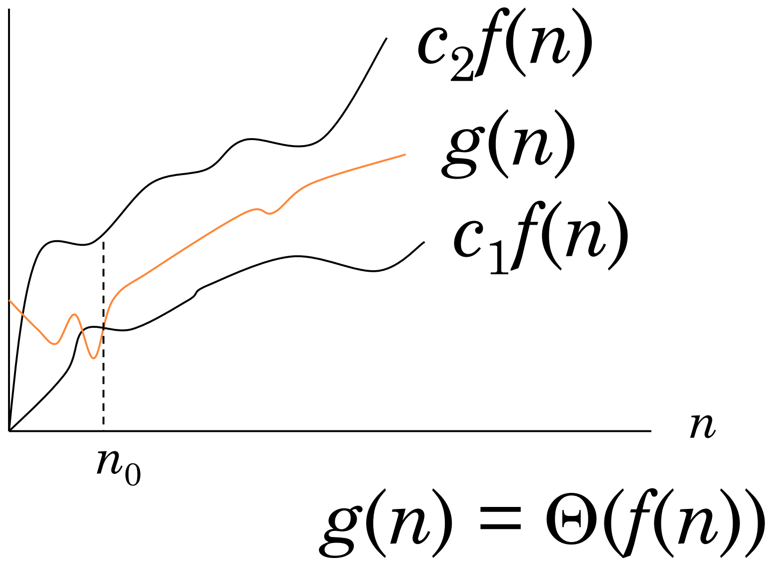
- It is the growth rate of the **time complexity** (space complexity) of an algorithm which ultimately determines the **size** of problems that can be solved by the algorithm.

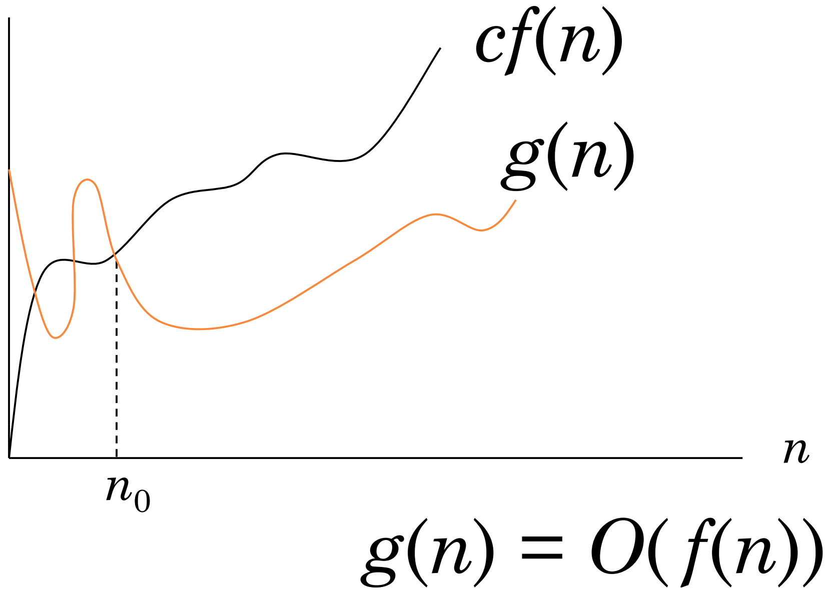
Growth of Functions

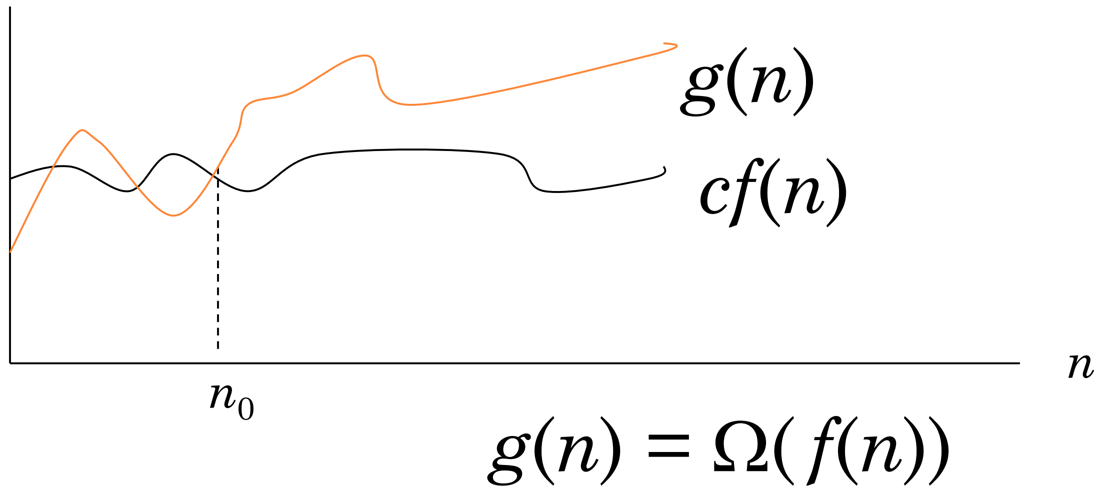
- **Upper bound.** $g(n) = O(f(n))$ if there exist constants $c > 0$ and $n_0 \geq 0$ such that for all $n \geq n_0$ we have $g(n) \leq c \cdot f(n)$.
- **Lower bound.** $g(n) = \Omega(f(n))$ if there exist constants $c > 0$ and $n_0 \geq 0$ such that for all $n \geq n_0$ we have $g(n) \geq c \cdot f(n)$.
- **Tight bound.** $g(n) = \Theta(f(n))$ if $g(n)$ is both $O(f(n))$ and $\Omega(f(n))$.

- Example: $g(n) = 32n^2 + 17n + 32$.
  - $g(n)$ is $O(n^2)$, $O(n^3)$, $\Omega(n^2)$, $\Omega(n)$, and $\Theta(n^2)$.
  - $g(n)$ is not $O(n)$, $\Omega(n^3)$, $\Theta(n)$, or $\Theta(n^3)$.

Growth of Functions







Growth of Functions

- **Upper bound** says that if constant factors are ignored $f(n)$ is at least as large as $g(n)$.
- $g(n) = O(f(n))$ means that the growth rate of $g(n)$ is smaller than or equal to the growth rate of $f(n)$.
- $O(\ldots)$ is read ``order ...'' or ``Big-Oh ....''

Growth of Functions

- **Lower bound** $g(n) = \Omega(f(n))$ (read "omega") means that the growth rate of $g(n)$ is *greater than or equal to* the growth rate of $f(n)$.
- **Tight bound** $g(n) = \Theta(f(n))$ means that for all n right of n_0, the value of $g(n)$ lies at or above $c_1 f(n)$ and at or below $c_2 f(n)$.

Growth of Functions

- Prove $g(n)=an^2+bn+c=\Theta(n^2)$
  - a, b, c are constants and $a>0$.
  - Find c_1, and c_2 (and n_0) such that
    - $c_1n^2 \leq g(n) \leq c_2n^2$ for all $n \geq n_0$.
  - It turns out: $c_1 = a/4$, $c_2 = 7a/4$ and
    - $n_0 = 2 \cdot max(|b|/a, sqrt(|c|/a))$
  - Here we also can see that lower terms and constant coefficient can be ignored.
  - How about $g(n)=an^3+bn^2+cn+d$?

Growth of Functions

- **Properties:**

If $g_1(n)$ is $O(f_1(n))$, $g_2(n)$ is $O(f_2(n))$ then

- $g_1(n)+g_2(n)$ is $O(\max(f_1(n),f_2(n)))$
- $g_1(n)g_2(n)$ is $O(f_1(n)f_2(n))$
- $ag_1(n)$ is $O(f_1(n))$ for any constant a.

Growth of Functions

- Bounds for some functions.
  - Polynomials.

$a_0 + a_1 n + \ldots + a_d n^d$ is $O(n^d)$ if $a_d > 0$.

  - Logarithms.

$O(\log_a n) = O(\log_b n) = O(n)$ for $a, b > 0$.

  - Exponentials.

For every $r > 1$ and every $d > 0$, $n^d = O(r^n)$.

  - Constant.

$O(c) = O(1)$ for any c.

Growth of Functions

- Typical growth functions.

Function	Name
c	Constant
$\log n$	Logarithmic
n	Linear
$n \log n$	Loglinear
n^2	Quadratic
n^3	Cubic
2^n	Exponential

ALGORITHMS

- Running time examples

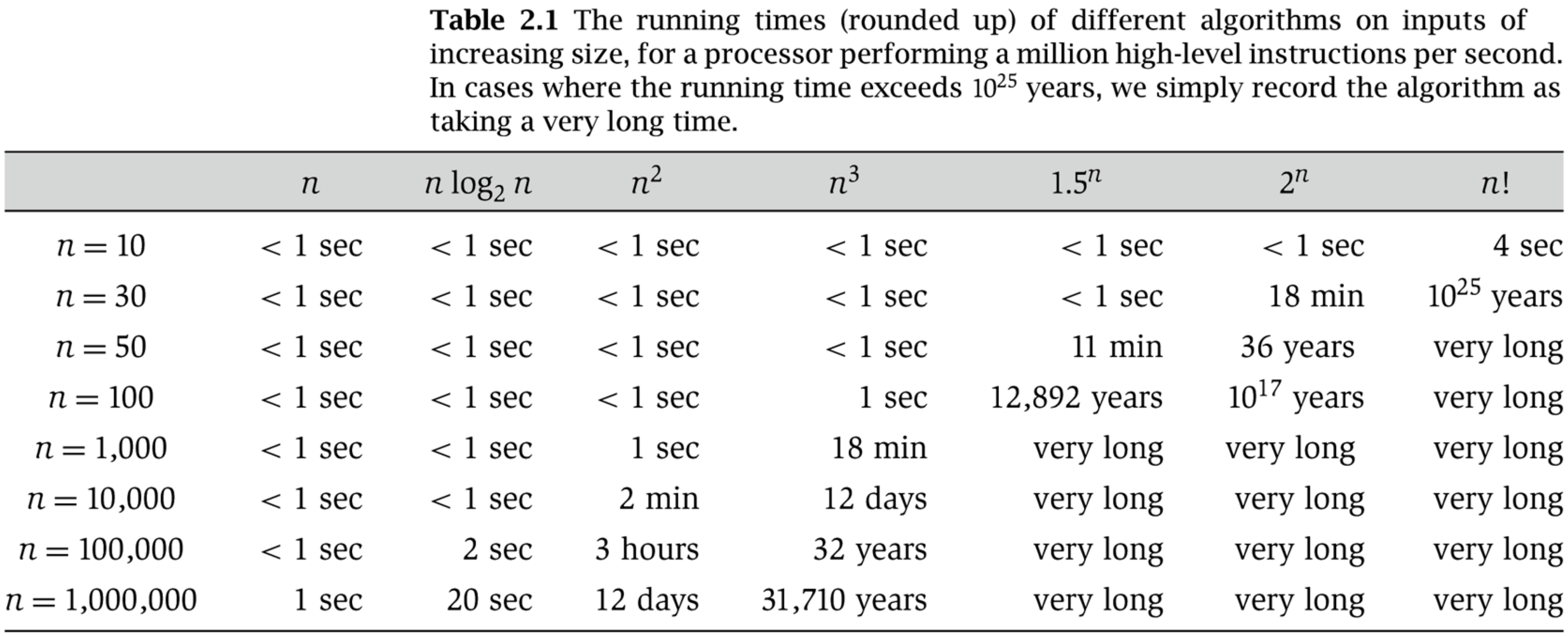


Table 2.1 The running times (rounded up) of different algorithms on inputs of increasing size, for a processor performing a million high-level instructions per second. In cases where the running time exceeds 10^{25} years, we simply record the algorithm as taking a very long time.

	n	$n \log_2 n$	n^2	n^3	1.5^n	2^n	$n!$
$n = 10$	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	4 sec
$n = 30$	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	18 min	10^{25} years
$n = 50$	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	11 min	36 years	very long
$n = 100$	$<$ 1 sec	$<$ 1 sec	$<$ 1 sec	1 sec	12,892 years	10^{17} years	very long
$n = 1{,}000$	$<$ 1 sec	$<$ 1 sec	1 sec	18 min	very long	very long	very long
$n = 10{,}000$	$<$ 1 sec	$<$ 1 sec	2 min	12 days	very long	very long	very long
$n = 100{,}000$	$<$ 1 sec	2 sec	3 hours	32 years	very long	very long	very long
$n = 1{,}000{,}000$	1 sec	20 sec	12 days	31,710 years	very long	very long	very long

Algorithms

- Statement complexity.
  - ***for/while*** loop:

```
for i = 1 to m:
        S
```

if the computation time of ***S*** is $t_i(n)$ for each i then the computation time of the for statement is $\sum_{i=1}^{m} t_i(n)$.
If $t_i(n) = t(n)$ for all i then the computation time of the loop is $mt(n)$.

ALGORITHMS

- Statement complexity.
  - ***if/else*** statement:

```
        if (condition):
                $S_1$
        else:
                $S_2$,
```

let $t_1(n)$ and $t_2(n)$ be the computation times of $\boldsymbol{S_1}$ and $\boldsymbol{S_2}$, respectively. The computation time of the **if** statement is $max\{t_1(n),t_2(n)\}$.

Algorithms

- Statement complexity.
  - **Consecutive** statements:

…
S_1
S_2
…

Let $t_1(n)$ and $t_2(n)$ be the computation times of two consecutive statements, respectively. The total computation time of the two statements is

$t_1(n)+t_2(n)$.

Algorithms

- Linear Time: O(n)
  - Running time is at most a constant factor times the size of the input.

```
max = $a_1$
for i = 2 to n {
    if ($a_i$ > max)
        max = $a_i$
}
```

  - Computing the maximum. Compute maximum of n numbers $a_1, \ldots, a_n$.

Algorithms

- Quadratic Time: $O(n^2)$
  - Closest pair of points. Given a list of n points in the plane $(x_1, y_1), \ldots, (x_n, y_n)$, find the pair that is closest.
  - $O(n^2)$ solution. Try all pairs of points.

```
min = (x1 - x2)^2 + (y1 - y2)^2
for i = 1 to n {
   for j = i+1 to n {
      d = (xi - xj)^2 + (yi - yj)^2
      if (d < min)
         min = d
   }
}
```

Algorithms

Polynomial Time: $O(n^k)$ Time

- Independent set of size k. Given a graph, are there k nodes such that no two are joined by an edge?
- $O(n^k)$ solution. Enumerate all subsets of k nodes.

```
foreach subset S of k nodes {
   check whether S in an independent set
   if (S is an independent set)
      report S is an independent set
   }
}
```

- Checking whether S is an independent set is $O(k^2)$.
- Number of k element subsets is $\binom{n}{k} = \frac{n\,(n-1)\,(n-2)\cdots(n-k+1)}{k\,(k-1)\,(k-2)\cdots(2)\,(1)} \leq \frac{n^k}{k!}$
- Total complexity is $O(n^k / k!\, O(k^2)) = O(n^k)$.

Algorithms

- Exponential Time
  - Given a graph, which is the largest independent set?
  - O(n^2 2^n) solution. Enumerate all subsets.

```
S* = ϕ
foreach subset S of nodes {
   check whether S in an independent set
   if (S is largest independent set seen so far)
      update S* = S
   }
}
```

Merge-Sort Algorithm

- **Step 1:** divide the *n*-element sequence into two sub-problems of *n*/2 elements each.
- **Step 2**: sort the two subsequences recursively using merge sort. If the length of a sequence is *1*, do nothing since it is already in order.
- **Step 3**: merge the two sorted subsequences to produce the sorted answer.

Merge-Sort Algorithm

- Pseudo-code

```
def merge_sort(A):
    middle = len(A) / 2
    left = merge_sort (A[1:middle])
    right = merge_sort (A[middle+1:end])
    return merge (left, right)
```

MERGE-SORT ALGORITHM

- Pseudo-code

```
def merge (A,B):
    result = <empty>
        while len(A) > 0 or len(B) > 0:
            if len(A) > 0 and len(B) > 0:
                if A[1] <= B[1]:
                    append result with A[1], delete A[1]
                else:
                    append result with B[1], delete B[1]
            else if len(A) > 0:
                append result with A[1], delete A[1]
            else if len(B) > 0:
                append result with B[1], delete B[1]
    return result
```

Merge-Sort Algorithm

- Animated example

MERGE-SORT ALGORITHM

- **Time complexity**

- There are two recursive calls, each of them sorts a sequence of $n/2$, and the statements after the two recursive calls take $O(n)$ time.
- Let $t(n)$ be the time complexity of the algorithm.

$$\boldsymbol{t(n) = 2t(n/2)+cn}$$

and $t(2) = O(1)$, where c is a constant. Solving the equation, $\boldsymbol{t(n) = O(n\log n)}$**.**

That's All for Today!