

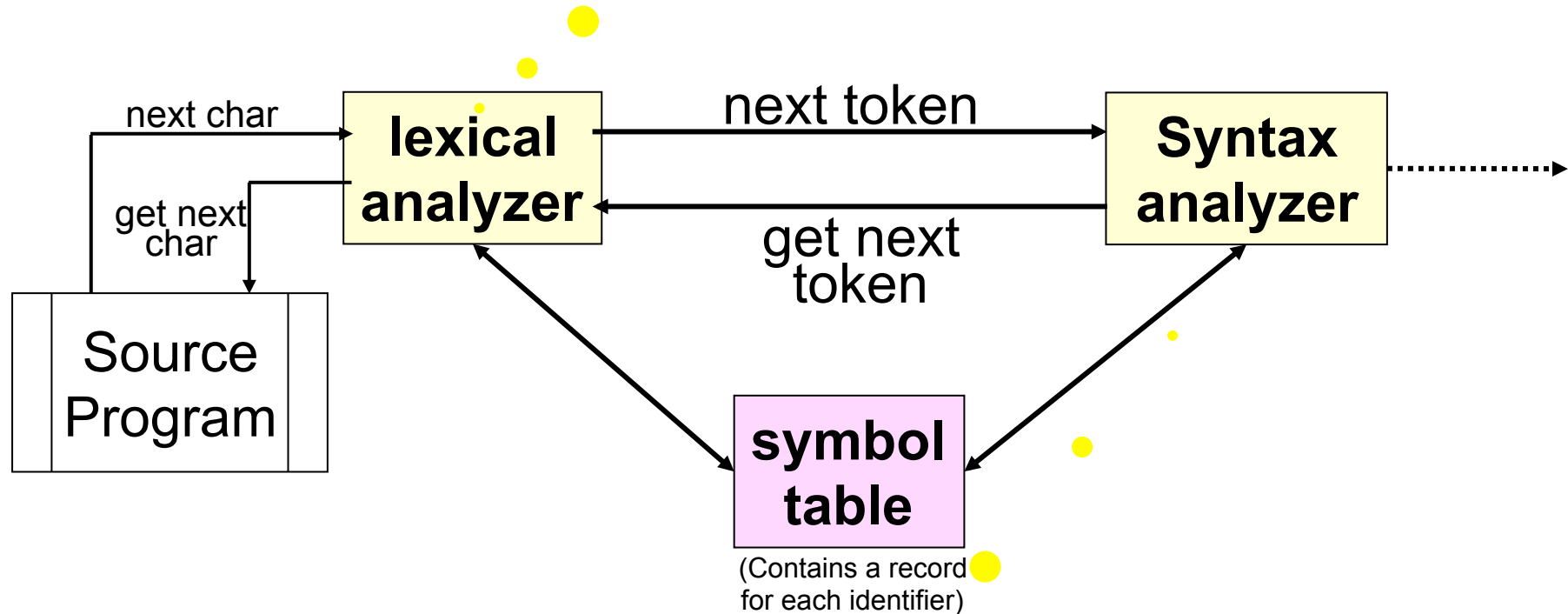
# **Language Processing Systems**

**Prof. Mohamed Hamada**

**Software Engineering Lab.  
The University of Aizu  
Japan**

# **Syntax Analysis (Parsing)**

1. Uses **Regular Expressions** to define **tokens**
2. Uses **Finite Automata** to recognize **tokens**



Uses **Top-down** parsing or **Bottom-up** parsing  
To construct a **Parse tree**

# Parsing

```
graph TD; Parsing --> TopDown[Top Down Parsing]; Parsing --> BottomUp[Bottom Up Parsing]; TopDown --> Predictive[Predictive Parsing]; TopDown --> LeftRecursion[Left Recursion]; TopDown --> LeftFactoring[Left Factoring]; Predictive --> LLk[LL(k) Parsing]; BottomUp --> ShiftReduce[Shift-reduce Parsing]; ShiftReduce --> LRk[LR(k) Parsing];
```

Top Down Parsing

Predictive Parsing

LL(k) Parsing

Left Recursion

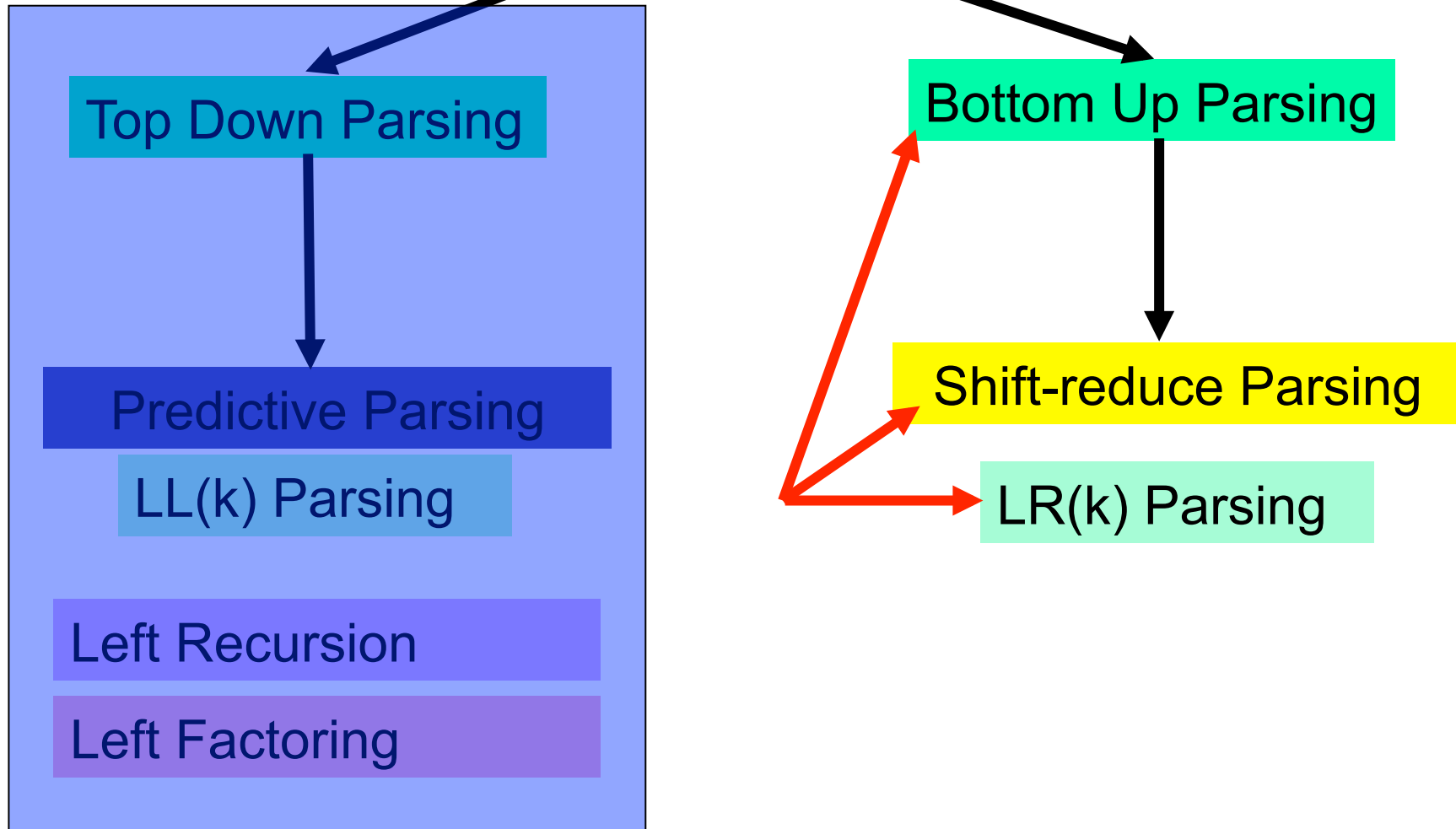
Left Factoring

Bottom Up Parsing

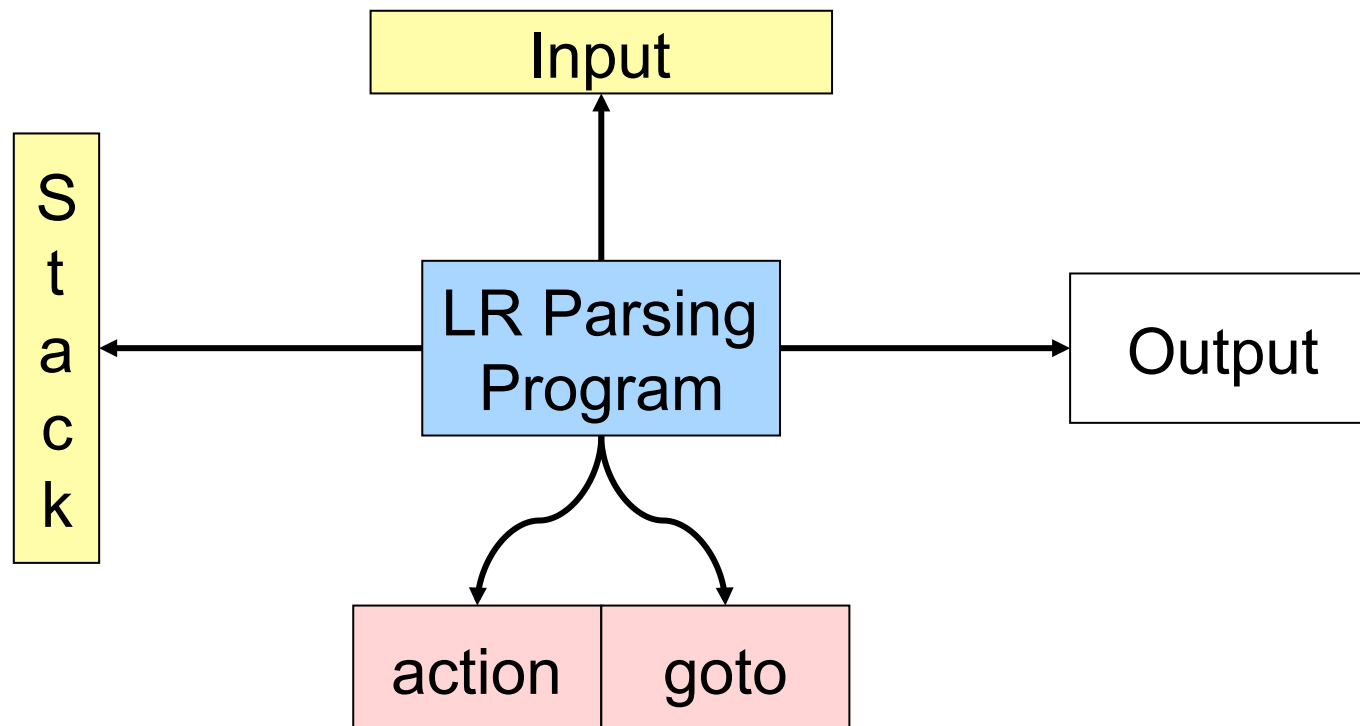
Shift-reduce Parsing

LR(k) Parsing

# Parsing



# LR Parser Example



# Shift reduce parser

1. Construct the *action-goto table* from the given grammar

2. Apply the shift-reduce parsing algorithm to construct the parse tree

# Shift reduce parser

1. Construct the *action-goto table* from the given grammar

This is what make difference between different typs of shift reduce parsing such as SLR, CLR, LALR

In this course due to short of time we will not study how to construct the *action-goto table*



# Shift reduce parser

## 2. Apply the shift-reduce parsing algorithm to construct the parse tree

The following algorithm shows how we can construct the move parsing table for an input string  $w\$$  with respect to a given grammar  $G$ .

```
set  $ip$  to point to the first symbol of the input string  $w\$$ 
repeat forever
begin
    if action[top(stack), current-input( $ip$ )] = shift(s) then begin
        push current-input( $ip$ ) then s on top of the stack
        advance  $ip$  to the next input symbol
    end
    else if action[top(stack), current-input( $ip$ )] = reduce  $A \rightarrow \beta$  then
        begin
            pop  $2*|\beta|$  symbols off the stack;
            push A then goto[top(stack), A] on top of the stack;
            output the production  $A \rightarrow \beta$ 
        end
    else if action[top(stack), current-input( $ip$ )] = accept then
        return
    else error()
end
```

# LR Parser Example

The following grammar:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

**s** represents shift

**r** represents reduce

**acc** represents accept

**empty** represents error

Can be parsed with this action and goto table

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

OUTPUT:

INPUT:

id + id \* id \$

STACK:

0

LR Parsing  
Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

STACK:

5  
id  
0

LR Parsing Program

OUTPUT:

F  
|  
id

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

OUTPUT:

STACK:

0

LR Parsing  
Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

F  
|  
id

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

STACK:

3  
F  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:

T  
|  
F  
|  
id

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

STACK:

0

LR Parsing Program



State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:

T  
|  
F  
|  
id

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

STACK:

2  
T  
0

LR Parsing  
Program

OUTPUT:

T  
|  
F  
|  
id

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

STACK:

7  
\*  
2  
T  
0

LR Parsing  
Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:

T  
|  
F  
|  
id

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

STACK:

5  
id  
7  
\*  
2  
T  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:

T | F  
F | id  
id

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

OUTPUT:

STACK:

7  
\*  
2  
T  
0

LR Parsing  
Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

T  
|  
F  
|  
id

F  
|  
id

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

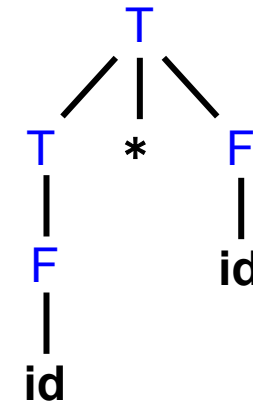
STACK:

10  
F  
7  
\*  
2  
T  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

(1)  $E \rightarrow E + T$

(2)  $E \rightarrow T$

(3)  $T \rightarrow T * F$

(4)  $T \rightarrow F$

(5)  $F \rightarrow ( E )$

(6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

STACK:

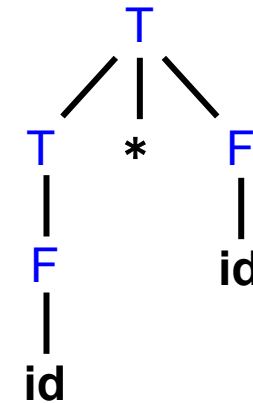
0

LR Parsing Program



State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT: 

id	*	id	+	id	\$
----	---	----	---	----	----

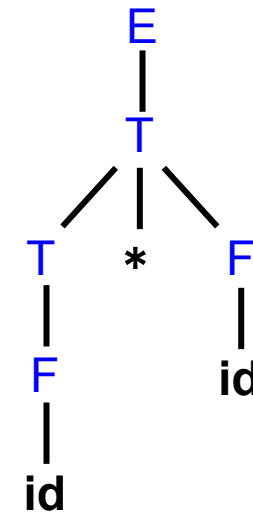
STACK: 

2
T
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

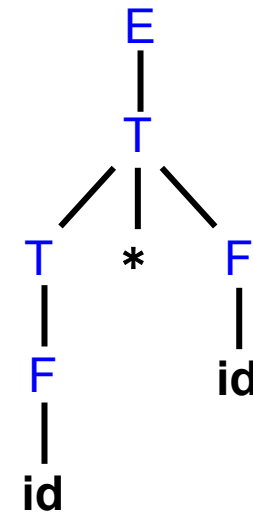
STACK:

0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT: 

id	*	id	+	id	\$
----	---	----	---	----	----

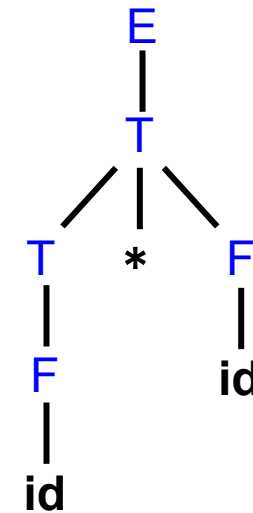
STACK: 

1
E
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:





GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

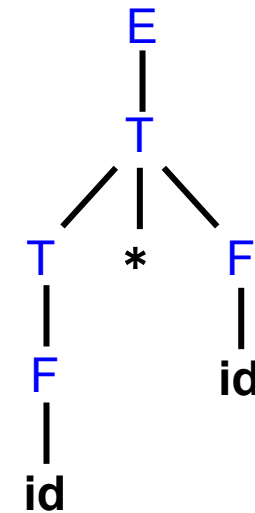
STACK:

6  
+  
1  
E  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

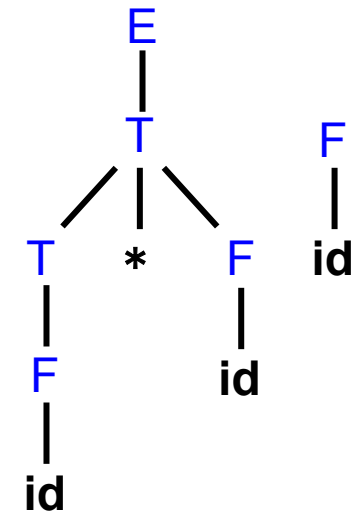
STACK:

5  
id  
6  
+  
1  
E  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

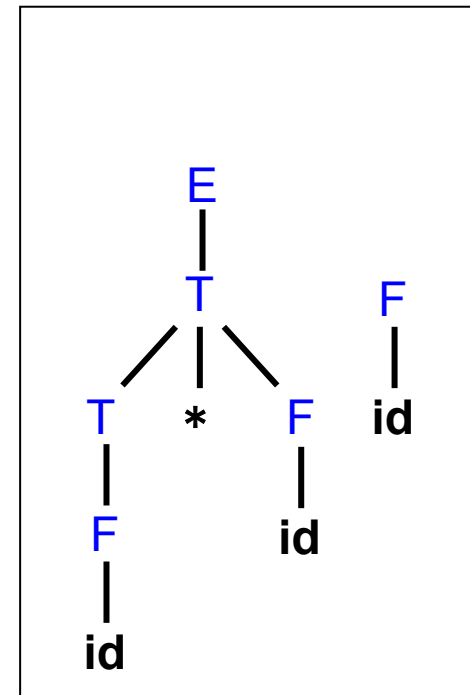
- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

OUTPUT:



STACK:

6  
+  
1  
E  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

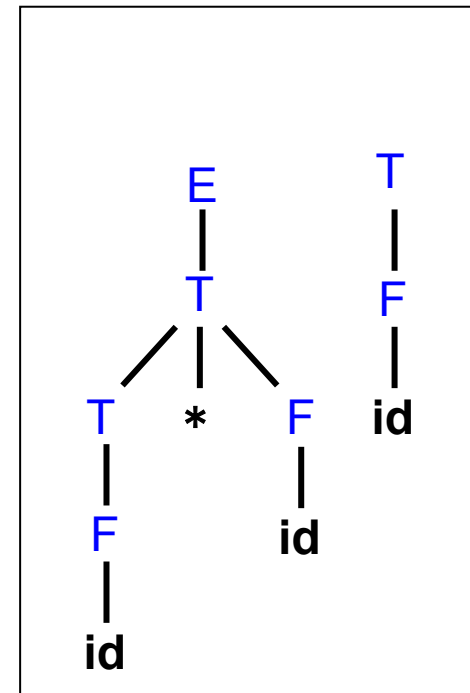
STACK:

3  
F  
6  
+  
1  
E  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

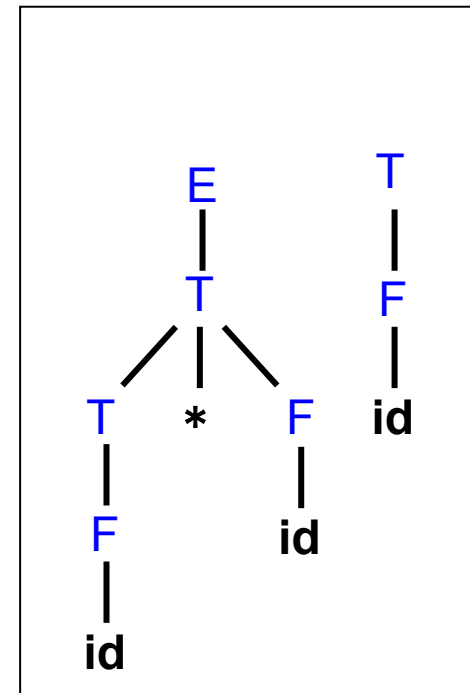
STACK:

6  
+  
1  
E  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

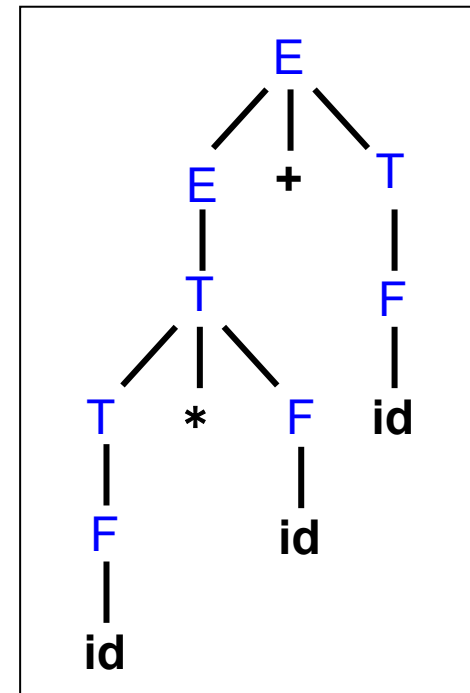
STACK:

9  
T  
6  
+  
1  
E  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

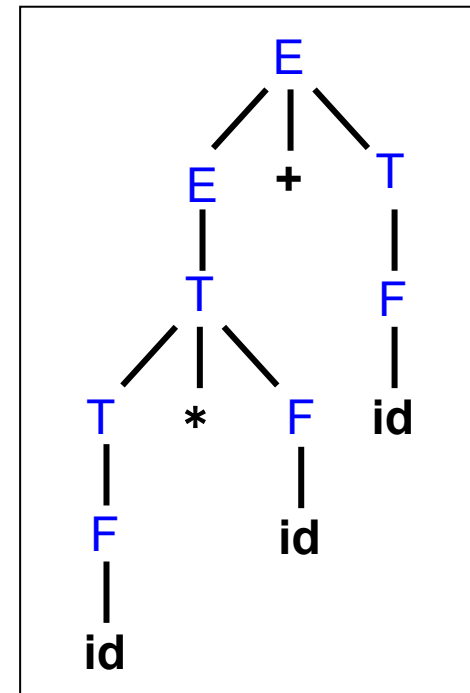
STACK:

0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:



GRAMMAR:

- (1)  $E \rightarrow E + T$
- (2)  $E' \rightarrow T$
- (3)  $T \rightarrow T * F$
- (4)  $T \rightarrow F$
- (5)  $F \rightarrow ( E )$
- (6)  $F \rightarrow id$

# Parser Example

INPUT:

id \* id + id \$

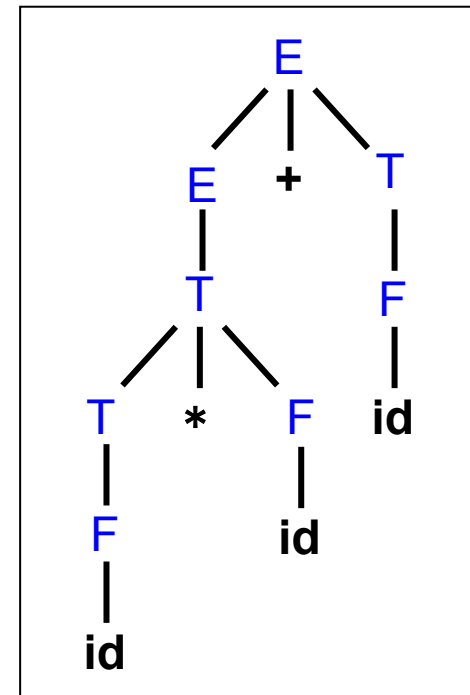
STACK:

1  
E  
0

LR Parsing Program

State	action						goto		
	id	+	*	(	)	\$	E	T	F
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

OUTPUT:





# Constructing Parsing Tables

All LR parsers use the same parsing program that we demonstrated in the previous slides.

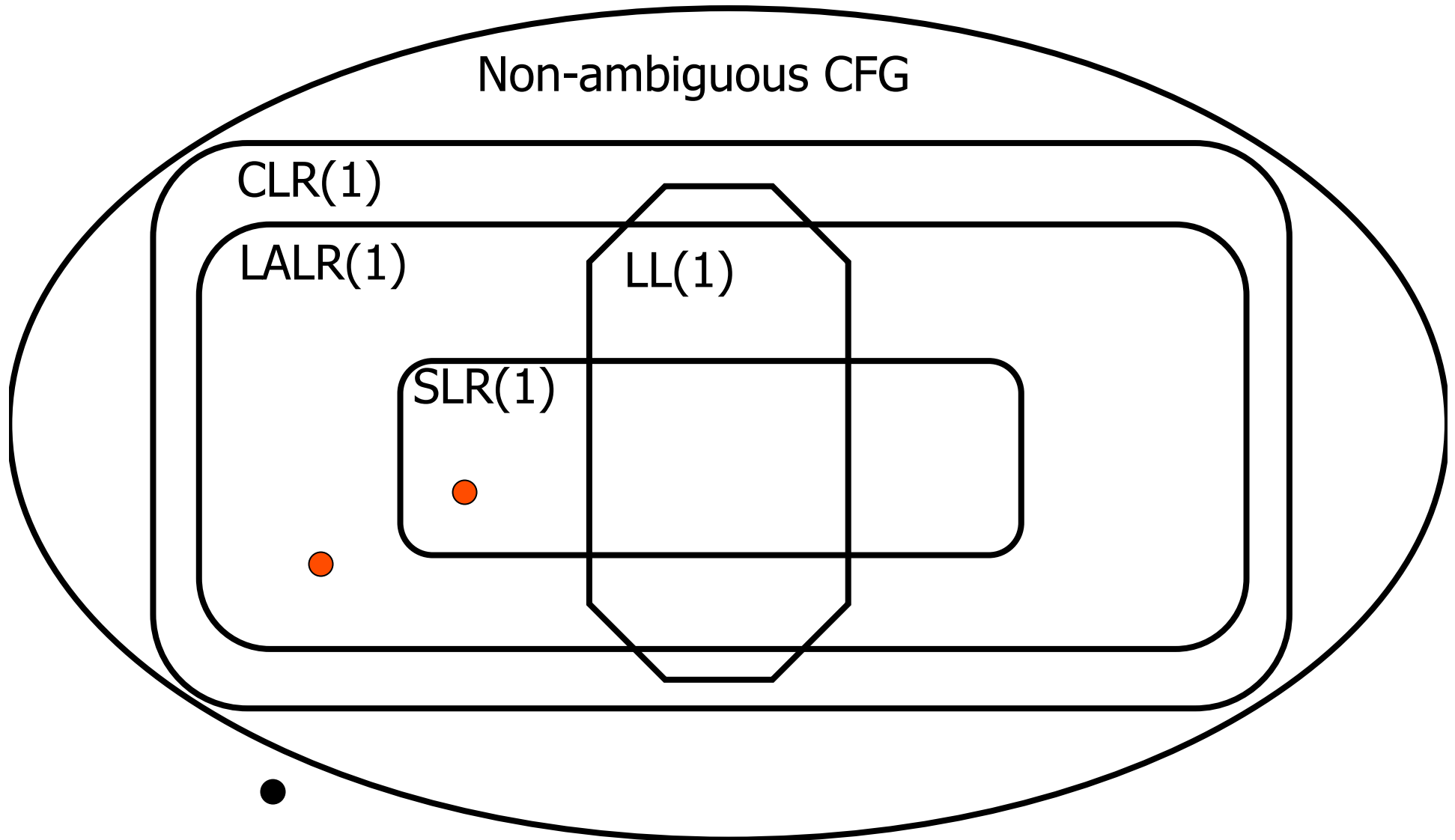
What differentiates the LR parsers are the action and the goto tables:

**Simple LR (SLR):** succeeds for the fewest grammars, but is the easiest to implement.

**Canonical LR:** succeeds for the most grammars, but is the hardest to implement. It splits states when necessary to prevent reductions that would get the parser stuck.

**Lookahead LR (LALR):** succeeds for most common syntactic constructions used in programming languages, but produces LR tables much smaller than canonical LR.

# Grammar Hierarchy



# Parsing

How parser works?

Top Down Parsing

Predictive Parsing

LL(k) Parsing

Left Recursion

Left Factoring

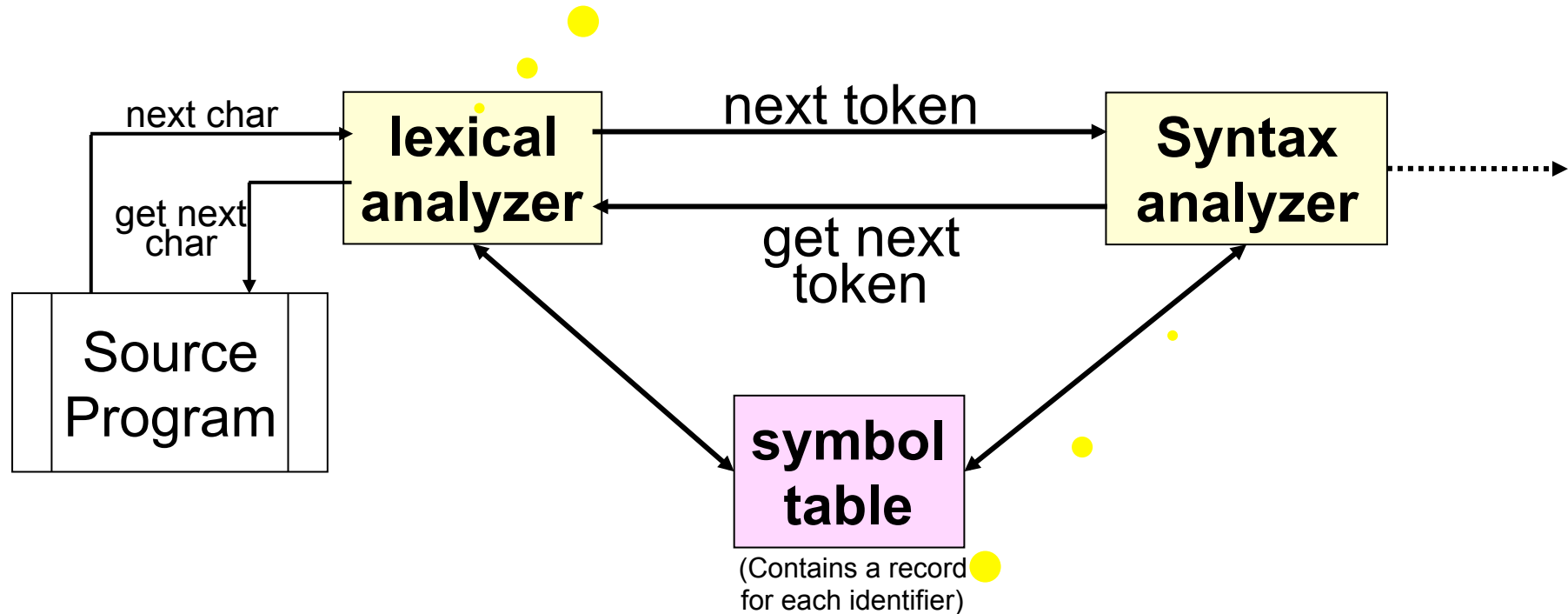
Bottom Up Parsing

Shift-reduce Parsing

LR(k) Parsing

How to write parser?

1. Uses **Regular Expressions** to define **tokens**
2. Uses **Finite Automata** to recognize **tokens**



Uses **Top-down** parsing or **Bottom-up** parsing  
To construct a **Parse tree**

# How to write a parser?

## Yacc

