

MALT: A MALloc Tracker A memory profiling tool





THE QUESTION



- We have profiling tool for timing (eg. Valgrind or vtune)
- But for memory usage ?
- Memory can be an issue :
 - Failed to run (or swap) due to lack of memory resource.
 - Performance impact of memory management functions
- Three main questions :
 - How to reduce memory footprint ?
 - How to improve overhead of memory management ?
 - How to improve memory usage ?





- We want to help searching:
 - Where memory is allocated.
 - **Properties** of allocated chunks.
 - **Bad** allocation **patterns** for performance.
 - Leaks
 - **Global variables** (TLS)

```
thread Int gblVar[SIZE];
     int * func(int size)
                   child func with allocs();
 4
                   void * ptr = new char[size];
                   double* ret = new double[size*size*size];
 6
                   for (....)
 9
                                double* buffer = new double[size];
10
                                //short and quick do stuff
11
                                delete [] buffer;
12
13
14
                   return ret;
                                       Software R&D: Observing Memory - MALT,
3/30/2016
                                                    Sébastien Valat
```

Global variables and TLS

Indirect allocations

Leak

Might lead to swap for large size

Short life allocations



EXISTING TOOLS

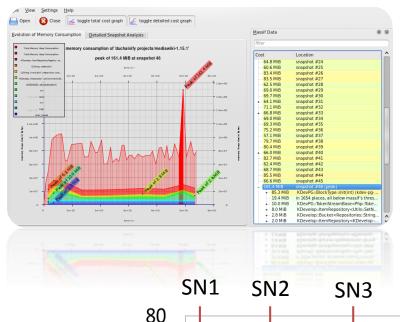


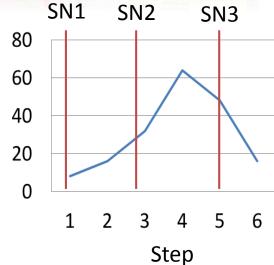
Valgrind - massif :

- Take snapshots over time.
- Link memory size to functions
- Peak might not be captured.
- Might miss short live allocations
- GUI not adequate for large code
- Slow, not parallel.

Valgrind - memcheck :

- Misuse of memory functions (malloc/new/free....)
- Leak detection
- Invalid accesses
- Slow, not parallel.

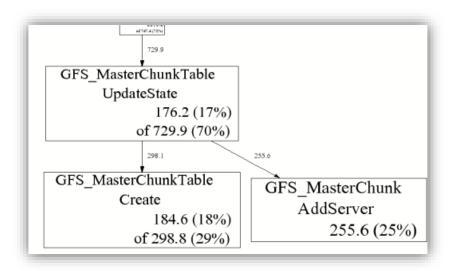






Google heap profiler (tcmalloc):

- Small overhead.
- Only provide snapshots of allocated memory per stacks.
- Peak might not be captured.
- No binary or source instrumentation.
- Output is not clear...
- Lack of a real GUI to use it.

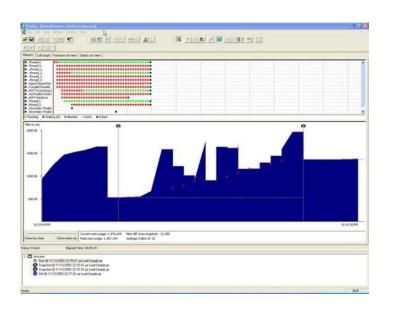


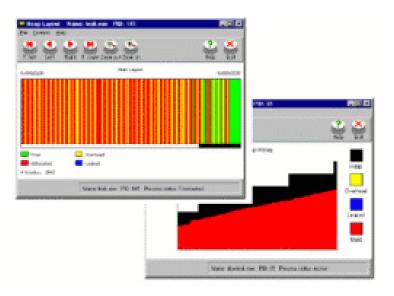
```
% pprof gfs_master profile.0100.heap
255.6 24.7% 24.7% 255.6 24.7% GFS_MasterChunk::AddServer
184.6 17.8% 42.5% 298.8 28.8% GFS_MasterChunkTable::Create
176.2 17.0% 59.5% 729.9 70.5% GFS_MasterChunkTable::UpdateState
169.8 16.4% 75.9% 169.8 16.4% PendingClone::PendingClone
76.3 7.4% 83.3% 76.3 7.4% __default_alloc_template::_S_chunk_alloc
49.5 4.8% 88.0% 49.5 4.8% hashtable::resize
```



IBM Purify++ / Parasoft Insure++

- Commercial
- Leak detection, access checking, memory debugging tools.
- Use binary or source instrumentation.
- Windows / Redhat
- Quite old GUI

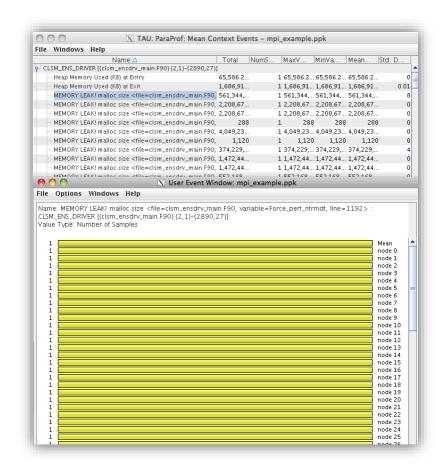






TAU memory profiler

- Provide profiles
- Follow stacks
- Track leaks
- Parallel, done for HPC/MPI
- Lack easy matching with sources





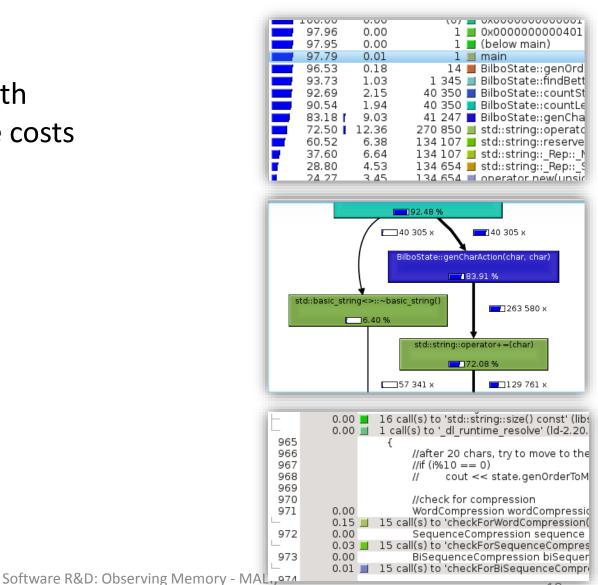
What is good in kcachgrind

Sébastien Valat

 List of functions with exclusive/inclusive costs

Nice call tree

Annotated sources





WHERE I GO



- Same approach than valgrind/kcachgind
- Mapped allocations on sources lines
- For memory resource usage :
 - Memory leaks (malloc without free)
 - Peak and total allocated memory
- For performance :
 - Allocation count
 - Allocation sizes (min/mean/max)
 - Chunk lifetime (min/mean/max)

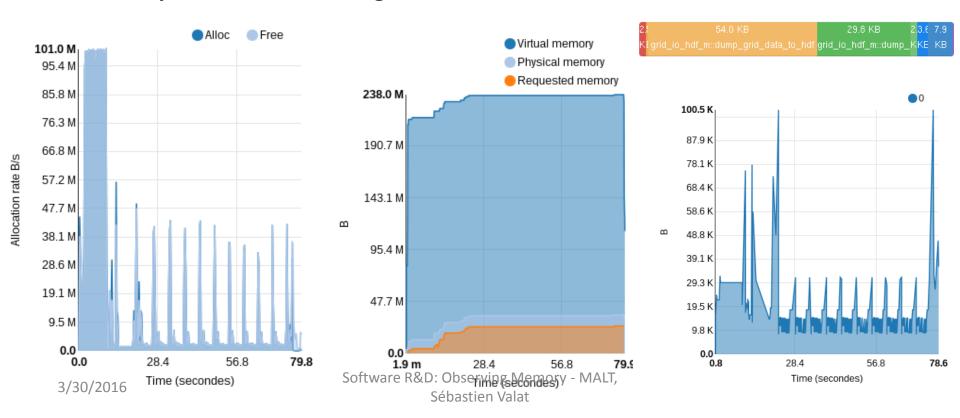




Profile over time :

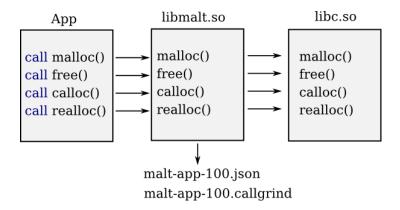
- Allocation rate
- Physical / Virtual / Requested memory
- Stack size for each thread (require function instrumentation)

Example on YALES2 with gfortran :





• Use LD_PRELOAD to intercept malloc/free/...



- Profile allocations on call stacks
- Generate JSON output file
- Build profile so size is limited by call tree



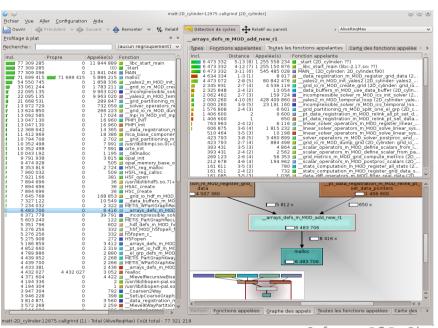
- Two approach implemented: backtrace and instrumentation
- Backtrace (default):
 - No change on binary or sources.
 - Function from glibc
 - Manage all dynamic libraries
 - No impact on compute
 - Slow on x86_64 due to use of libunwind.
 - Slow for large number of calls (~>10M)
- Instrumentation :
 - Need source recompilation (available): -finstrument-function
 - Or tools for binary instrumentation : MAQAO / Pintool (experimental)
 - Impact performance of compute only functioans
 - Faster for really large number of calls to malloc
 - Only provide stacks for the instrumented sources/binaries





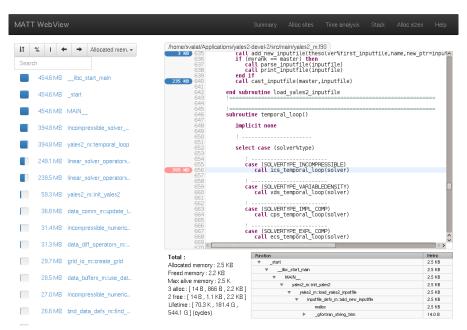
Callgrind compatibiltiy

- Can use kcachgrind
- Might be usefull for some users, cannot provide all metrics.



Own web view

- Get all metrics
- Web technology (NodeJS, D3JS, Jquery, AngularJS)
- Easier for remote usage
- Can be used for shared working





What is missing to kcachegrind

- Display human readable units
 - You prefer **15728640** of **15MB**?
 - I want to compare to what I expect.
- Cannot handle non sum cumulative metrics
 - Inclusive costs only rely on + operator
 - Some mem. metrics requires max/min (eg. local peaks, lifetime, sizes)
- No way to express time profiles
- No way to express parameter distributions (eg. sizes).



SOME VIEWS



Global summary

00:00:00.25

PHYSICAL MEMORY PEAK 2.3 MB

ALLOCATION COUNT 379

4.1 Gb

Run description

Executable:	simple-case-finstr-linked
Commande:	./simple-case-finstr-linked
Tool:	matt-0.0.0
Host:	localhost
Date:	2014-11-26 22:40
Execution time :	00:00:00.25
Ticks frequency:	1.8 GHz

Global statistics

Show all details Show help		
Physical memory peak		2.3 MB
Virtual memory peak		103.7 MB
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- Provide a small summary
- Provide some warnings

Show all details Show help				
Physical memory peak	66.7 MB			
Virtual memory peak	158.1 MB			
Requested memory peak	6.1 MB			
Cumulated memory allocations	11.5 MB			
Allocation count	172.2 K			
Recycling ratio	1.9			
Leaked memory	743.7 KB			
Largest stack	0 B			
Global variables	10.0 MB 🛕			
TLS variables	48 B			
Global variable count	421.0 K 🛕			
Peak allocation rate	37.8 MB/s			



Global summary: top 5 functions

- Summarize **top functions** for some metrics
- Points to check
- Examples on YALES2

Alloc count

Ratio	Allocs	Function
	911.9 K	data_comm_m::copy_int_comm_to_data
	896.4 K	data_comm_m::copy_data_to_int_comm
	853.2 K	data_comm_m::update_int_comm
	484.9 K	sponge_layer_m::calc_sponge_layer_mask
	296.0 K	incompressible_numerics_m:ics_diffuse_velocity_rk_4th

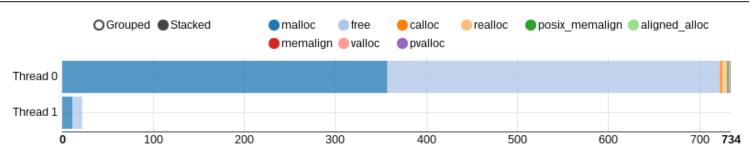
Allocated memory

Ratio	Allocs	Function
	202.4 MB	linear_solver_operators_m::solve_linear_system_deflated_pcg
	26.6 MB	bnd_data_defs_m::find_bnd_data
	21.8 MB	linear_solver_operators_m::solve_el_grp_pcg
	19.0 MB	data_comm_m::copy_int_comm_to_data
	18.1 MB	data_comm_m::update_int_comm

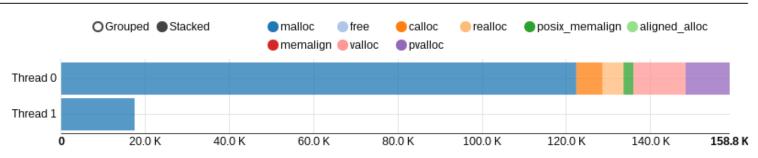
Peak memory

Per thread statistics

Call per thread

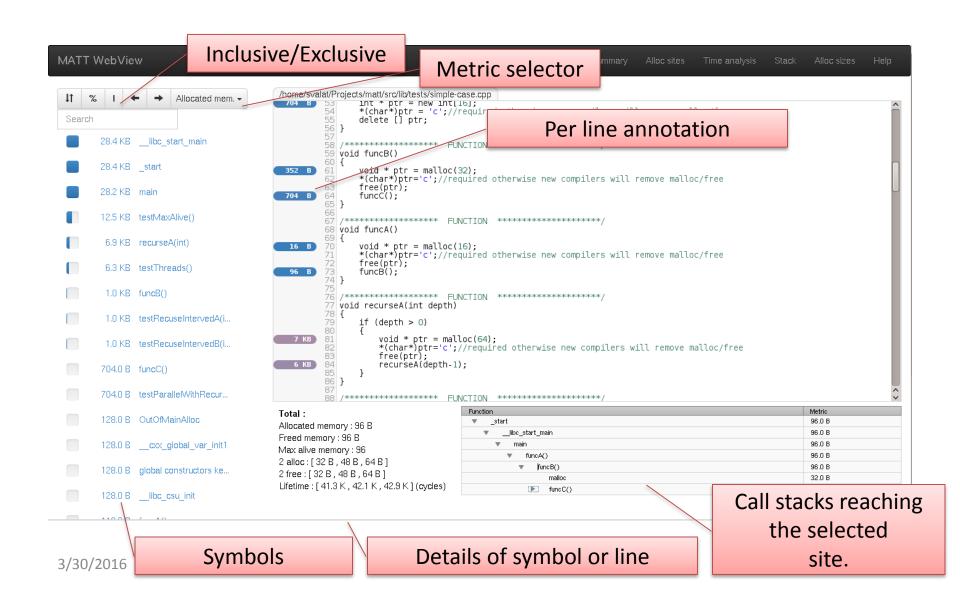


Time per thread



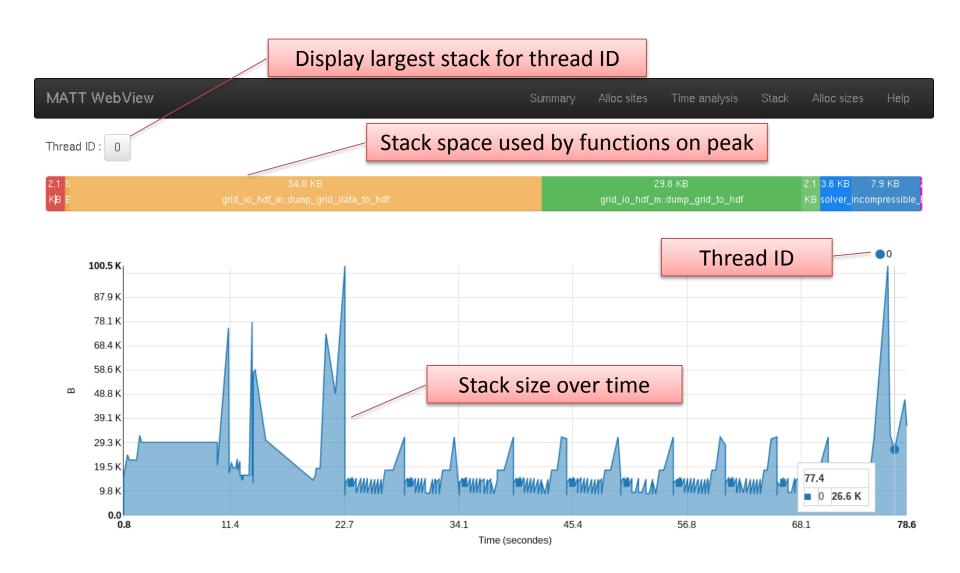


Source annotations





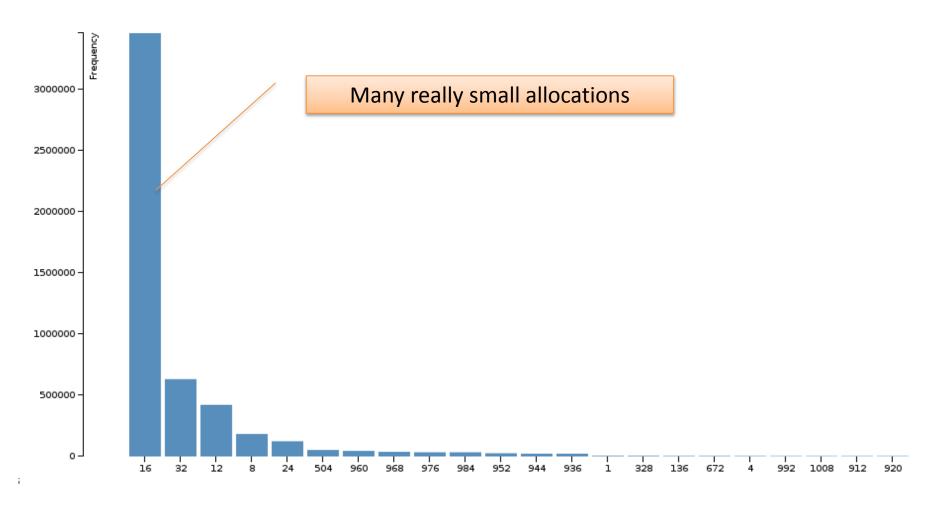
Tracking stack memory





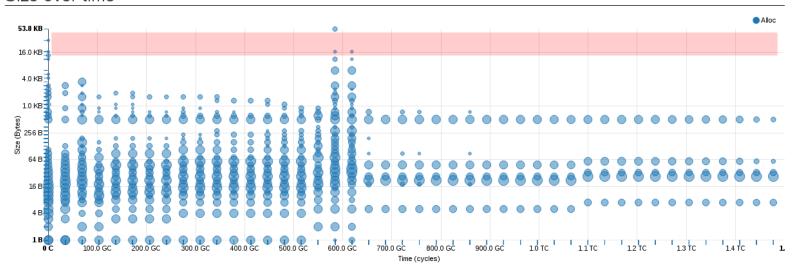


Example from YALES2 with gfortran issue

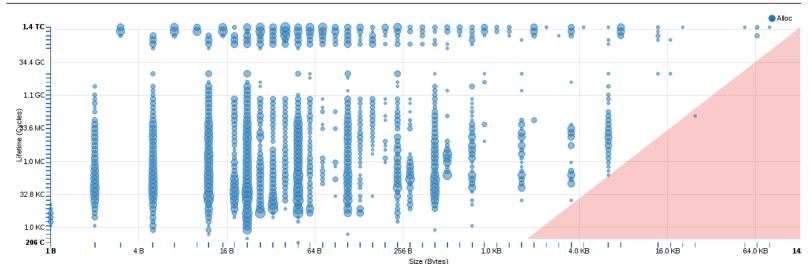




Size over time



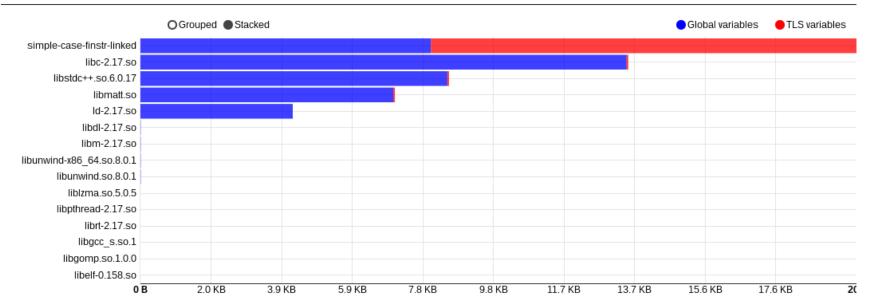
Lifetime over size



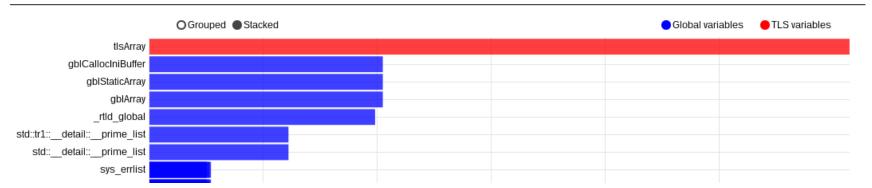


Global variables

Distribution over binaries



Distribution over variables



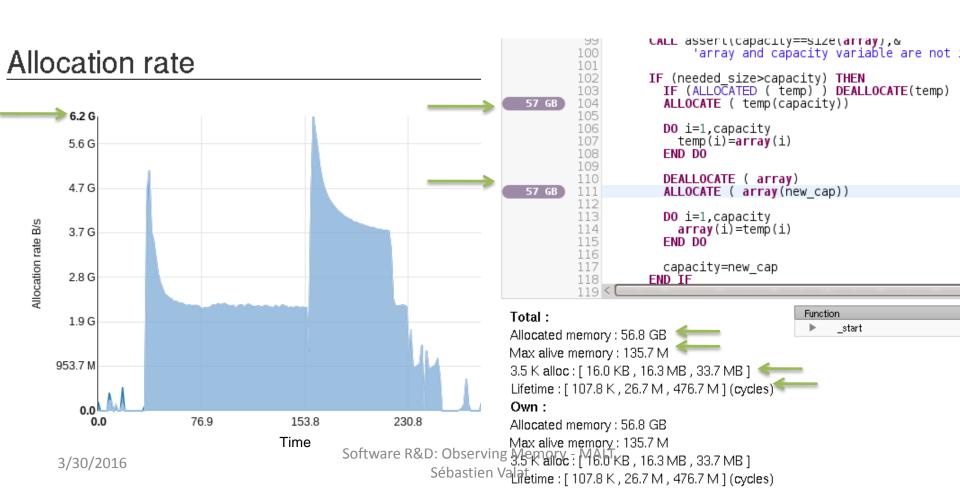


REAL CASES



Example on AVBP init phase

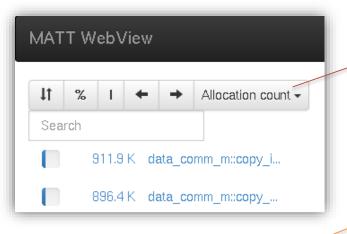
Issue with reallocation on init:





Allocatable arrays on YALES2

• Issue only occur with **gfortran**, ifort uses stack arrays.



Search intensive alloc functions

Huge number of allocation for a line programmer think it doesn't do any!

```
do i=1,nitem_el_grp

el_grp_ind = el_grp_index2int_comm_index%val(1,i)

int_comm_ind = el_grp_index2int_comm_index%val(2,i)

el_grp_r2%val(1:dim1,el_grp_ind) = int_comm_r2%val(1:dim1,int_comm_ind)

end do
```

Total:

Allocated memory: 9.5 MB Freed memory: 9.5 MB Max alive memory: 432

608.0 K allo : [16 B , 16 B , 16 B]

608.0 K free: [16 B, 16 B, 16 B]

Lifetime: [24.5 K, 39.9 K, 37.8 M] (cycles)

Own₃;/30/2016

Allocated memory: 9.5 MB

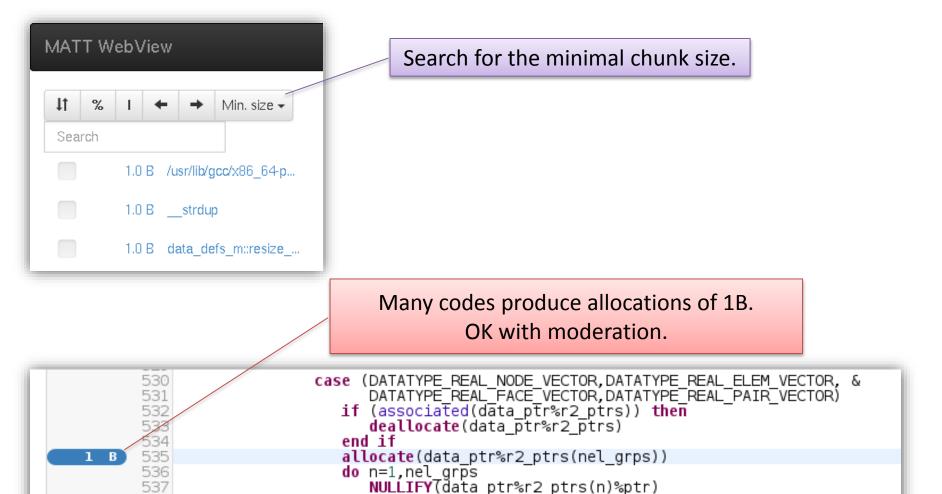
And mostly really small allocations!

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We can found allocs of 1B!

Examples on YALES 2, small allocations :



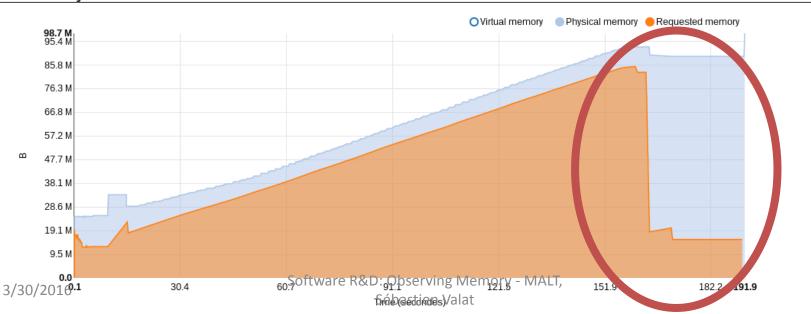
end do

538



- Example from Dassault mini-app from Loïc Thébault and Eric Petit.
- Fragmentation can prevent from returning physical pages to OS
- Solution: avoid interleaved allocation of chunks with different lifetime.
- We observed with the source annotation that most of them can be avoided.

Memory allocated over time





Missing reference on auto

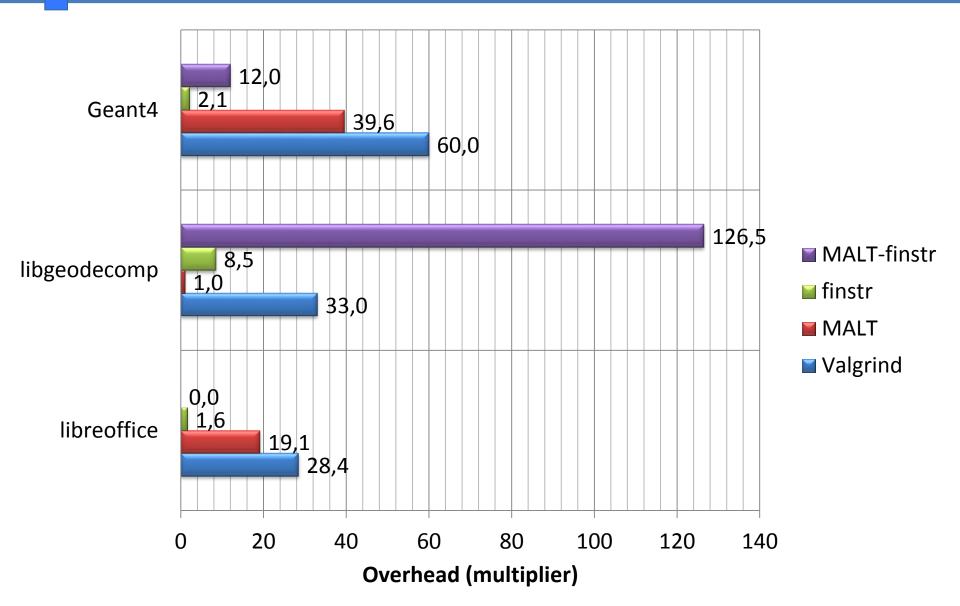
On one of my own code:

```
118
            //search none empty
            bool retry;
119
            do {
120
121
                retry = false;
                for (auto entry : channels) {
122
                    if (entry.acceptor()) {
123
                         if (entry.channel->emptyTryLock(retry) == false) {
124
                             ret = entry.channel;
125
                             break;
126
127
128
129
```

- It was in a active waiting loop....
- Divide by 300 the number of allocation

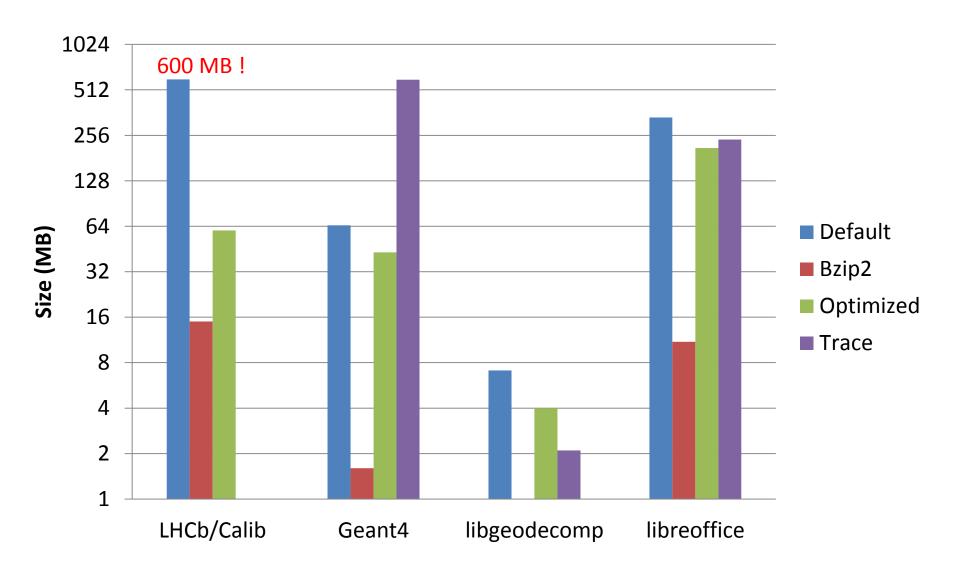


Overhead and profile size











USAGE & CONCLUSION



Backtrace mode :

```
# Optionally recompile with debug flag to get source lines : cc –g ... # Run your program ${PREFIX}/bin/malt [--config=file.ini] YOUR PRGM [OPTIONS]
```

Function tracking with -finstrument-function :

```
# Recompile with instrumentation flag :
cc -finstrument-function –g ...
# Run
${PREFIX}/bin/malt --stack=enter-exit [--config=file.ini] YOUR_PRGM [OPTIONS]
```

• Use the web view:

```
#Launch the server malt-webserver -i malt-{YOUR_PRGM}-{PID}.json # Connect with your browser on http://localhost:8080
```



- Useful tool for tracking allocation patterns
- Interesting real cases:
 - YALES2 : allocations with gfortran
 - AVBP : large allocation rate
 - PAMPA: allocation larger than expected by programmer
 - Dassault mini-app: effect and source of fragmentation
- Future work :
 - Integrate traces into the view (already get all the backend stuff)
 - Add NUMA informations (at lease statistics about mappings)
 - Hope to get Open Source release soon





Thank you.

QUESTIONS?





BACKUP



Ideas of improvement

- Add NUMA statistics
- Provide virtual/physical ratio
- Estimate page fault costs
- Exploit traces in GUI for deeper analysis
 - Alive allocations at a certain time
 - Fragmentation analysis
 - Time charts from call sites
 - Usage over threads for call sites

About global peak

- The tool maintain a call stack tree
- Profile stats on leafs
- On new global peak, need to copy each local current contribution
- Need to walk over the wall tree each time?
- Do lazy update :
 - Keep track of last local peakld on each leaf
 - On leaf update, compare the local peakid and the global one
 - If not same : remember the old local contribution.

