FYS3150 - project 1

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https://github.uio.no/comPhys/FYS3150/tree/project1

PROBLEM 1

We have the one-dimensional Poisson equation

$$-\frac{d^2u}{dx^2} = f(x) \tag{1}$$

where f(x) is known to be $100e^{-10x}$. We also assume $x \in [0, 1]$, that the boundary condition are u(0) = 0 = u(1) and u(x) is

$$u(x) = 1 - (1 - e^{-10})x - e^{-10x}$$
(2)

where u(x) is an exact solution to Eq. (1). We can check this analytically by differentiating u(x) twice.

$$-u''(x) = f(x)$$

$$u''(x) = -f(x)$$

$$u(x)' = 10x^{-10x} - 1 + \frac{1}{e}$$

$$u''(x) = -100e^{-10x} = -f(x)$$

PROBLEM 2

problem a)

See 'poisson_exact.cpp' and 'algorithms.cpp' in the github repository

problem b)

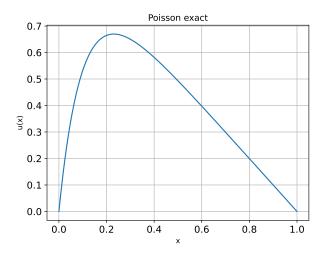


FIG. 1. A plot of the exact solution for the Poisson equation from 1 for $x \in [0,1]$

PROBLEM 3

We are discretizing the Poisson equation from 1. Discretizing x and setting up some notation:

$$x \to x_i$$

$$u(x) \to u_i$$

$$i = 0, 1, \dots, n$$

$$h = \frac{x_{max} - x_{min}}{n}$$

$$x_i = x_0 + ih$$

We're using the three-point formula to find the second derivative:

$$\frac{du^2}{dx^2} = u'' = \frac{u_{i-1} - 2u_i + u_{i+1}}{h^2} + O(h^2)$$

$$f_i = -\left(\frac{u_{i-1} - 2u_i + u_{i+1}}{h^2} + O(h^2)\right)$$

We then approximate and change the notation, $v_i \approx u_i$ and get

$$f_i = \frac{-v_{i-1} + 2v_i - v_{i+1}}{h^2} \tag{3}$$

PROBLEM 4

The equation we got in 3 isn't the most ergonomic for setting up a matrix equation, so we can rewrite it to

$$-v_{i-1} + 2v_i - v_{i+1} = h^2 f_i (4)$$

Equation 4 is a set of equations for every i

$$i = 1$$
 $-v_0$ $2v_1$ $-v_2$ $= h^2 f_i$
 $i = 2$ $-v_1$ $2v_2$ $-v_3$ $= h^2 f_2$

and so on and so forth. We can see that v_0 and v_n will end up and alone on their columns, and since we know what they are we simply move them over.

$$i = 1$$
 $2v_1 - v_2$ $= h^2 f_1 + v_0$
 $i = 2$ $-v_1$ $2v_2$ $-v_3$ $= h^2 f_2$

This can then be easily rewritten as a matrix equation $A\vec{v} = \vec{g}$

$$\begin{bmatrix} 2 & -1 & 0 & 0 & \dots & 0 & 0 \\ -1 & 2 & -1 & 0 & \dots & 0 & 0 \\ 0 & -1 & 2 & -1 & \dots & 0 & 0 \\ \vdots & \vdots & \ddots & \ddots & \ddots & \vdots & \vdots \\ 0 & 0 & \dots & -1 & 2 & -1 & 0 \\ 0 & 0 & \dots & 0 & 0 & -1 & 2 \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ \vdots \\ v_{n-3} \\ v_{n-2} \\ v_{n-1} \end{bmatrix} = \begin{bmatrix} g_1 \\ g_2 \\ g_3 \\ \vdots \\ gn - 3 \\ gn - 2 \\ g_{n-1} \end{bmatrix}$$

where g_i is $h^2 f_i$ (+ v_0 for f_1 and + v_n for f_{n-1})

PROBLEM 5

Problem a

Since we're "dropping" 2 columns, n = m - 2

problem b

we will find $\vec{v_i}$ for 1..(n-1), meaning everything but the boundary points.

${\bf problem} \ {\bf 6}$

problem a

Algorithm 1 Algorithm for solving general tridiagonal matrix

```
\begin{split} & \text{arrays a, b, c, u, f, temp of length n} \\ & \text{btemp} = b[1] \\ & u[1] = f[1]/btemp \\ & \textbf{for } i = 2,3,...,n \ \ \textbf{do} \\ & \text{temp}[i] = c[i\text{-}1] \ / \ btemp \\ & \text{btemp} = b[i] - a[i] * \text{temp}[i] \\ & u[i] = (f[i] - a[i] * u[i\text{-}1]) \ / \ btemp \\ & \textbf{for } i = n\text{-}1, n\text{-}2, ..., 1 \ \ \textbf{do} \\ & u[i] - = \text{temp}[i\text{+}1] * u[i\text{+}1] \end{split}
```

problem b

FLOPs:
$$1 + 6(n-1) + 2(n-1) = 1 + 8(n-1) = 8n - 7$$

PROBLEM 7

problem a

see 'general_tridiag.cpp' and 'algorithms.cpp' in the github repository

${\bf problem}\ {\bf b}$

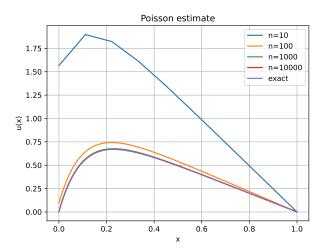


FIG. 2. Exact vs numerical comparison for the solution of equation 1

PROBLEM 8

problem a

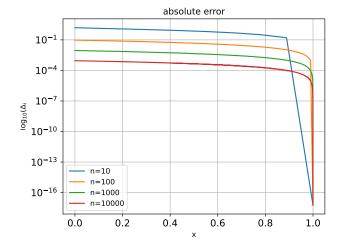


FIG. 3. The absolute error for the general tridiagonal alogrithm in \log_{10}

${\bf problem}\ {\bf b}$

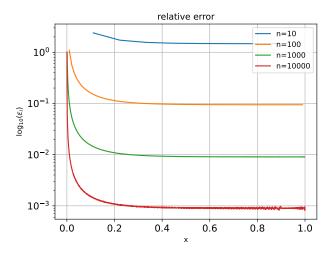


FIG. 4. The relative error of the general tridiagonal algorithm

$\mathbf{problem}\ \mathbf{c}$

n	max rel error
10	2.3903877077005538
100	1.0946543221364102
1000	1.0091303813999173
10000	1.0009115263620092
100000	1.000088892839682
1000000	1.0
10000000	1.0

 $\,$ FIG. 5. Table with the biggest relative error for each iteration

PROBLEM 9

Since the values of a, b, c never changes, we can just replace the arrays with a constants, saving us from repeatedly calculating $-(-1) \cdot something$.

problem a

Algorithm 2 Algorithm for solving special tridiagonal matrix

 $\begin{array}{l} btemp = 2 \\ u[1] = f[1]/\ 2 \\ \\ \textbf{for} \ \ i = 2,3,...,n \ \ \textbf{do} \\ temp[i] = -1 \ / \ btemp \\ btemp = b[i] + temp[i] \\ u[i] = (f[i] + u[i\text{-}1]) \ / \ btemp \\ \textbf{for} \ \ i = n\text{-}1, \, n\text{-}2, \, ..., \, 1 \ \textbf{do} \\ u[i] -= temp[i+1] \ \ ^* u[i\text{+}1] \end{array}$

arrays u, f, temp of length n

problem b

FLOPs:
$$1 + 4(n-1) + 2(n-1) = 1 + 6(n-1) = 6n - 5$$

problem c

PROBLEM 10

n	general	special
100	1.639e-06	7.334e-07
1000	1.6525e-05	7.2634e-06
10000	0.00016182	7.2397e-05
100000	0.001401	0.00073008
1000000	0.011985	0.011125
10000000	0.12067	0.11029

FIG. 6. A table showing the running time of the general and special algorithm for a given n

PROBLEM 11

LU decomposition has a complexity of $O(N^3)(+O(N^2))$ vs our algorithm that runs in O(N). for $n = 10^5$ I would expect to:

1) run out of memory, since you need to store $8*N*N = 8*10^5*10^5 = 8*10^{10} \approx 80$ GB

2) to be very slow. The alogrithm's O factor is two orders of magnitude larger, so probably around 100 times as long.