

I. Databases

What is database

- Organize collection of data
- Two types of databases:
 - Operational database - collect, modify and maintain data on daily basis.
 - Analytical database - mainly for historical and time-dependant data
- Based on two mathematical theories
 - Set theory
 - First-order predicate

Relational database Model

Based on the concept of “relation”:

- **Relation/Table:** a relation (table) is a set of tuples
- **Tuple/Row:** a tuple is an individual row in a relation. Each tuple represents a single record
- **Attribute/Column:** An attribute is a named column of a relation
- **Degree:** The degree of a relation is the number of attributes it contains
- **Cardinality:** Cardinality refers to the number of tuples in a relation - колко реда има таблицата ни
- **Relational Algebra:** Procedural query language for the relational model

Domain:

- Refers to the set of allowable values that a column (attribute) can contain.
- It's the type of data that can be stored in a certain column
- Example: a domain of a `salary` column could be all real numbers, or the domain of a `gender` column could be 'male', 'female', 'prefer not to say'

Relation:

- Table in a relational database
- Set of tuples that have the same attributes
- A relation is defined by its name and the set of attributes it contains

Relational database Schema

- Describes the structure of a database in a formal way
- Provides definitions for tables, fields, relationships, views, indexes, packages, procedures, functions, queues, triggers, sequences, synonyms and other objects
- Defines constraints over the database objects
- In many cases normalization is a key requirement for a database schema

Data normalization

Нормализация на базите данни

<https://bg.myservername.com/database-normalization-tutorial>

- What is normalization
- Normal forms
 - **1NF - Used to eliminate repeating groups**

First normal form means the following conditions are met:

- Eliminate repeating groups in individual tables
- Create a separate table for each set of related repeated data
- Identify each set of related data with Primary key

Do not use multiple fields in a single table to store similar data:

1NF

employee_id	name	job_code	job	city
21313123	Spas Belev	S03	Software Engineer	Sofia
45344353	Zhorzh Raychev	S03	Software Engineer	Plovdiv
98658699	Ivan Ivanov	T01	Technical Lead	Sofia
234656643	Georgi Georgiev	P01	Product Owner	London

- **2NF - Used to eliminate redundant data**

Second normal form means the following conditions are met:

- Create separate tables for sets of values that apply to multiple records
- Relate these tables with a foreign key

Records should not depend on anything other than a table's Primary key:

employee_id	name	job_code	city
21313123	Spas Belev	S03	Sofia
45344353	Zhorzh Raychev	S03	Plovdiv
98658699	Ivan Ivanov	T01	Sofia
234656643	Georgi Georgiev	P01	London

jobs

job_code	job
S03	Software Engineer
S03	Software Engineer
T01	Technical Lead

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- **3NF - Used to eliminate columns that are not dependent on the key of the table**

Third normal form checks for **transitive dependencies**:

- Eliminate fields that do not depend on the key

Records that are not a part of the record's key do not belong in the table

3NF

employee_id	name	city
21313123	Spas Belev	Sofia
45344353	Zhorzh Raychev	Plovdiv
98658699	Ivan Ivanov	Sofia
234656643	Georgi Georgiev	London

jobs

job_code	job
S03	Software Engineer
S03	Software Engineer
T01	Technical Lead
P01	Product Owner

employee_roles

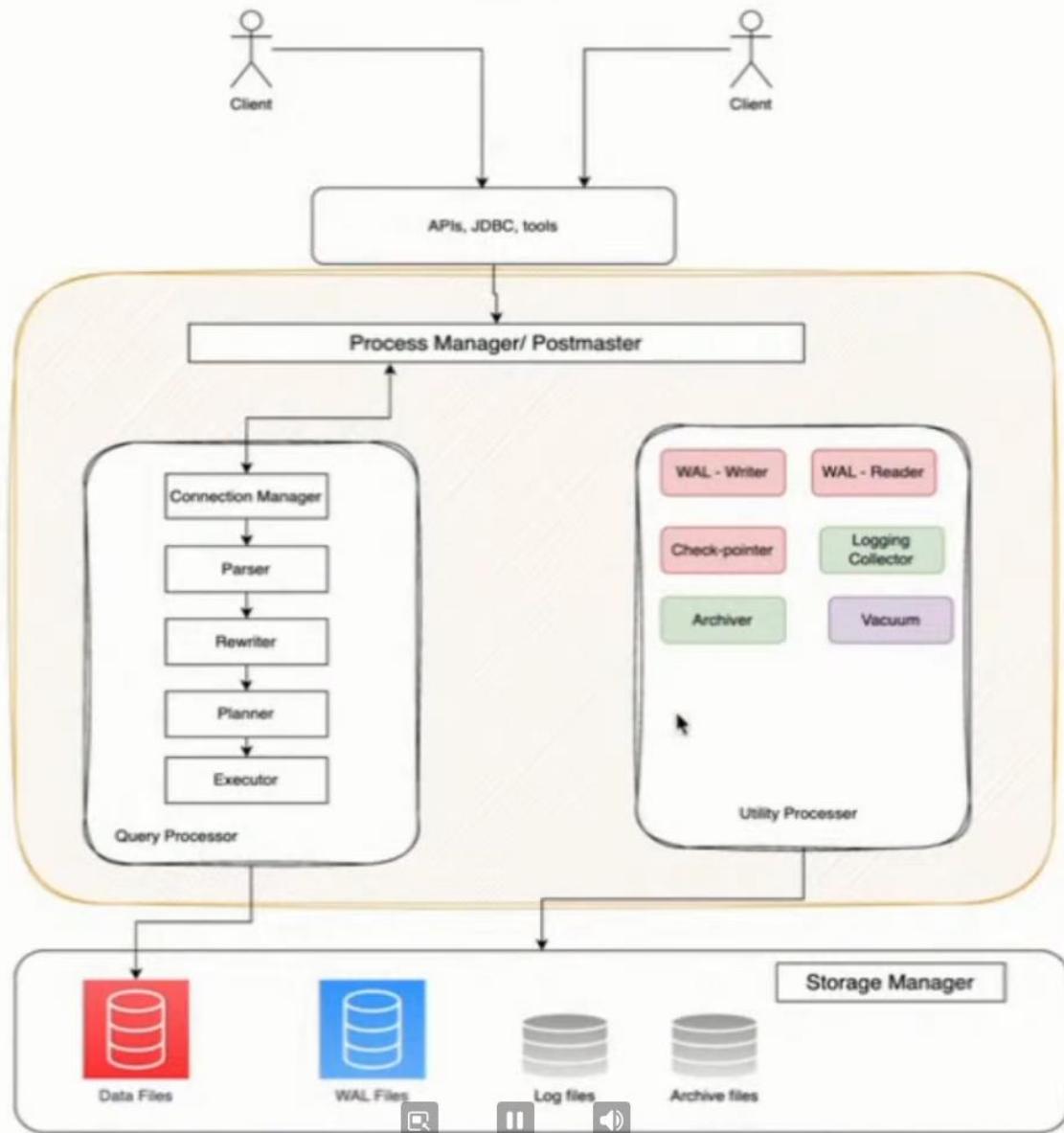
employee_id	job_code	job
21313123	S03	Software Engineer
45344353	S03	Software Engineer
98658699	T01	Technical Lead
234656643	P01	Product Owner

- 4NF - Isolate Independent Multiple Relationships
- 5NF - Isolate semantically related Multiple Relationships

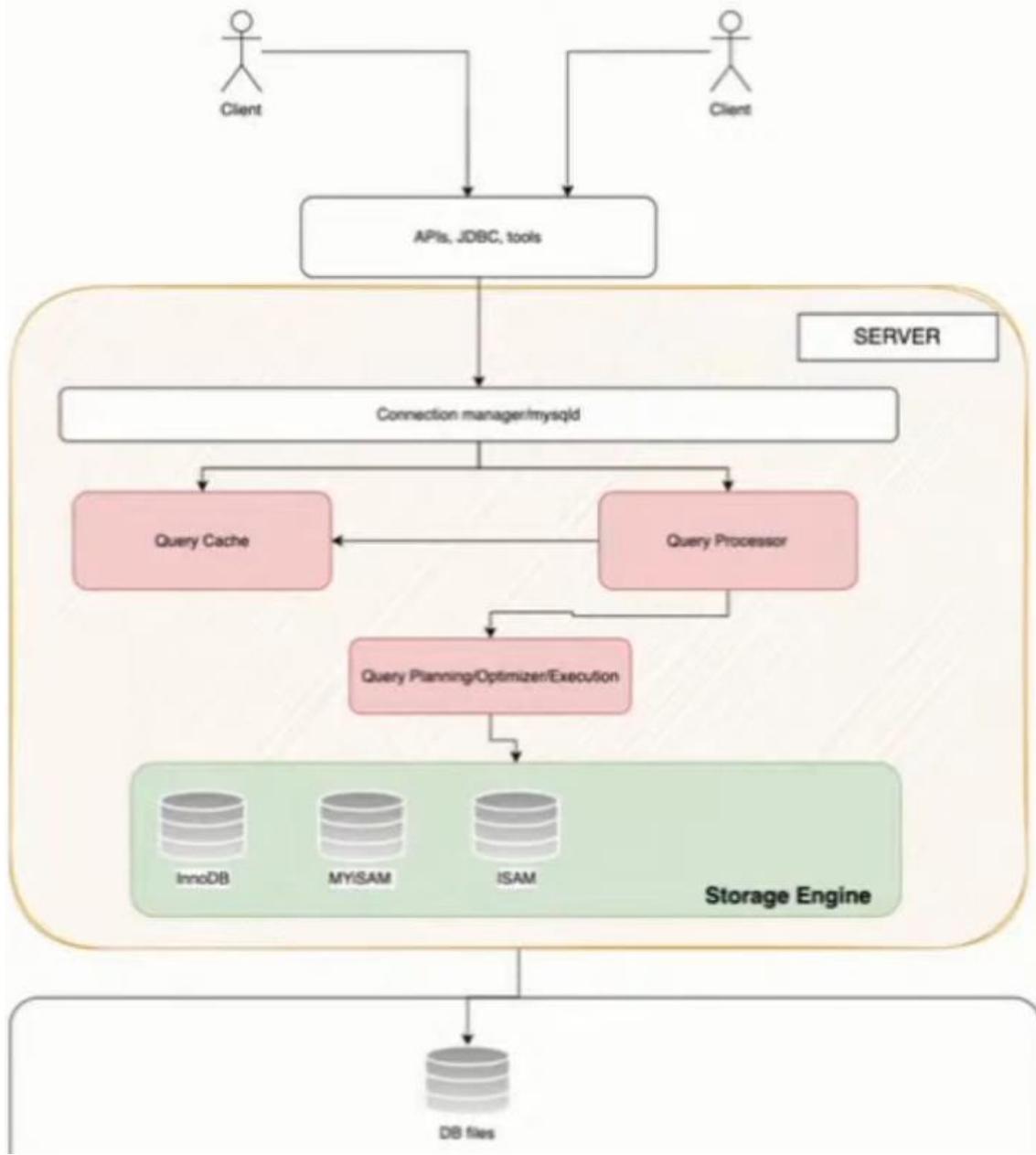
Как работи отдолу една RDBMS

Writer
Reader
Archiver
Vacuum

PostgreSQL



MySQL



Indexes

- Improve the speed of data retrieval operations on a database table at the cost of additional writes and storage space
- Indexes are used to quickly locate data without having to search every row
- Indexes can be created using one or more columns of a database table.
- The **primary key is called also clustered index**. All other indexes are called non-clustered indexes.

Key points about indexes:

- **Index Types:** There are several types of indexes, such as B-tree, bitmap, and hash indexes, each suited to particular types of queries
- **Composite index:** An index with 2 or more columns at the same time
- **Index Key:** The column(s) on which the index is created

- **Primary key:** A type of constraint used in a table to uniquely identify each record in this table. It automatically creates an index known as a clustered index.
- **Unique index:** It ensures data present in the column is unique. It happens automatically. It allows one null value for a column!

Indices и балансирано бинарно дърво

- Structures associated with a table or view that speeds retrieval of rows
 - Usually implemented as **B-trees**
- Indices can be built-in the table (**clustered**) or stored externally (**non-clustered**)
- Adding and deleting records in indexed tables is slower!
 - Indices should be used for big tables only (e.g. 50 000 rows)

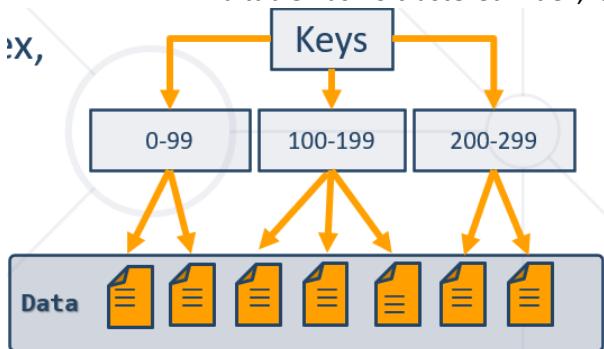
Когато имаме много на брой данни, в даден момент се налага преиндексиране на базата данни, за да се постигне балансирано бинарно дърво.

Небалансираното дърво забавя много операциите.

id PRIMARY KEY -> B-tree|

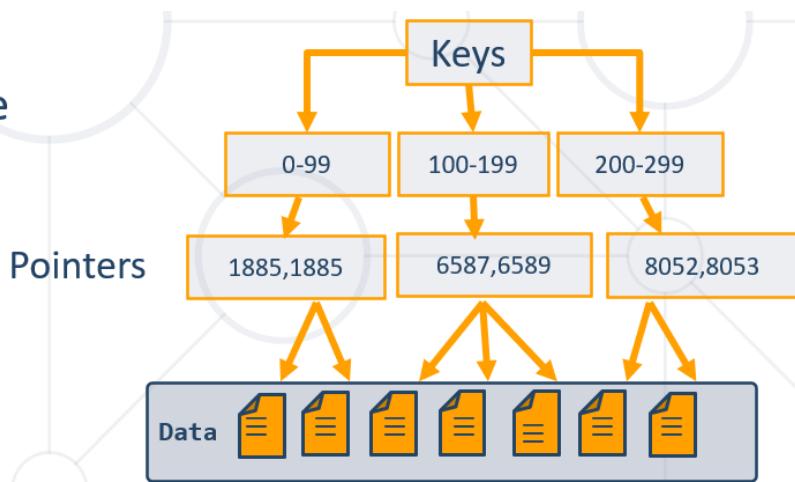
Clustered Indices

- **Clustered index(primary key) determine the order of data**
 - Very useful for fast execution of **WHERE**, **ORDER BY** and **GROUP BY** clauses
- Maximum 1 clustered index per table
 - If a table has no clustered index, its data rows are stored in an **unordered structure (heap)**



Non-Clustered Indices

- Useful for fast retrieving a **single record** or a **range** of records
 - Each **key value entry** has a pointer to the data row that contains the key value
- Maintained in a separate structure in the DB



Синтаксис на индекс

CREATE INDEX

```
ix_users_first_name_last_name
ON `users`(`first_name`, `last_name`);
```

Stored procedures

- Group of pre-compiled SQL statements that are stored in the database
- Can be invoked (or called) by:
 - User application programs
 - Triggers
 - Other stored procedures from within the database
- Encapsulate logic for data transformation, data validation, and business-specific logic

Benefits of using stored procedures:

- Performance: Because stored procedures are compiled and stored in the database, they tend to be more efficient and faster.
- Network traffic: using stored procedures can reduce network traffic between clients and servers
- Security: stored procedures provide an additional layer of security. You can grant users the permission to execute a stored procedure without giving them access to the underlying tables
- Code reuse and encapsulation: Logic created in stored procedures can be reused across multiple programs

In some cases, better in the Java code to be the logic and to have more network traffic!!!

Functions

- Precompiled SQL commands that are used to perform a specific task or operation and return a value
- Built-in(aggregate, Scalar, Date, and many more) or user-defined
- Много рядко може да се наложи да си правим наша си функция. В повечето случаи вградените функции са достатъчни

Design process

- Requirement Analysis: understand what the database will be used for, by who and what data it will store - ще пишем или ще четем повече от базата, и други подобни въпроси
- Conceptual Database Design: Identify the high-level entities and relationships among them. An E/R is often created.

- Logical Database Design: Translate the conceptual model into a logical model, which is a more detailed and precise representation of the data
- Physical Database Design
- Database Implementation

Entities

Entity - a definable thing - such as a person, object, concept or event - that can have data stored about it.

Strong entity - има си всички атрибути и е самодостатъчно/самоопределящо се. И няма нужда от foreign keys.

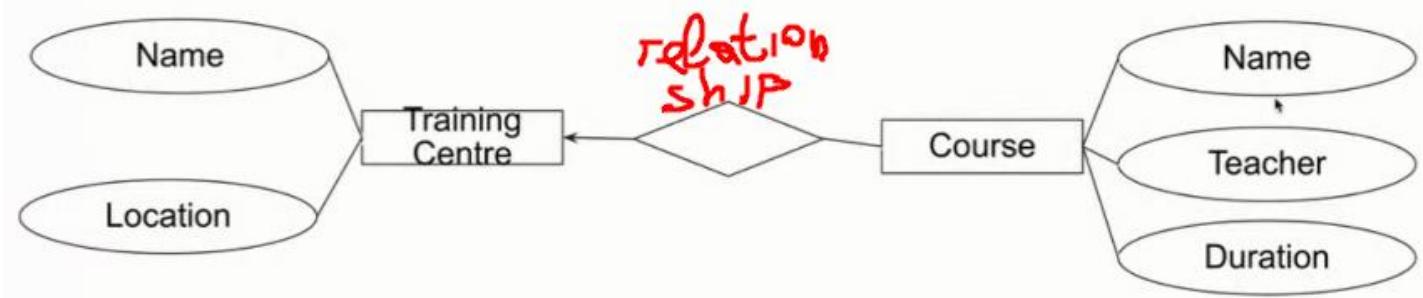
Weak entity - не може да се самоопредели от атрибутите си какво е точно

Associative entity - associates entities (or elements) within an entity set

Entity set: - same as entity type, but defined at a particular point in time, such as students enrolled in a class on the first day

Relationship

- How entities act upon each other or are associated with each other - One to Many, Many to One, Many to Many, One to One, etc.
- Relationships are typically shown as **diamonds** (on the ER diagrams) or labels directly on the **connecting lines**

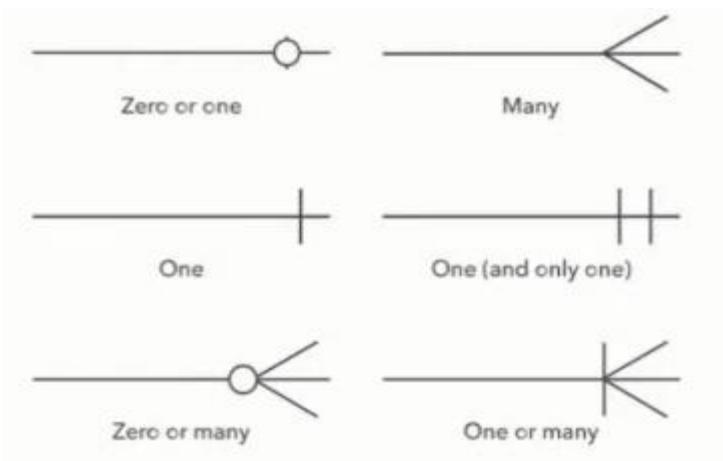


Attribute

- Attributes are categorized as simple, composite, derived, as well as single-value or multi-value:
 - Simple: attribute value is **atomic** and can't be further divided
 - Composite: Sub-attributes spring from an attribute
 - Derived: Attribute is calculated or otherwise derived from another attribute, such as age from a birthdate

Cardinality

- Defines the numerical attributes of the relationship between two entities or entity sets
- The three main cardinal relationships are one-to-one, one-to-many, and many-to-many



Tools for creating E\R or similar database diagrams (Data Modelling Tools):

- Lucidchart
- Microsoft Visio
- Draw.io
- MySQL Workbench
- ER/Studio
- <https://dbdiagram.io/home>
- IntelliJ Idea Ultimate
- pgAdmin
- <https://dbeaver.io/>

II. Installing MySQL

Installing

Please, see the installation manual from OneDrive, the first folder of MySQL course.

Stopping MySQL server/s on Windows

Когато сървър базата се инициализира в Докер/облачно за даден проект, то трябва да спрем ръчно съществуващия service (database server) на Windows.

При Windows за да пуснем контейнер, то трябва да използваме Docker Desktop (не знам дали има background приложение на Docker за Windows).

Докато при Линукс може и да не използваме Docker Desktop, а само с бекграунд приложение на Docker.

Services					
	mpssvc	4544	Windows Defender Firewall	Running	LocalService...
	MSDTC		Distributed Transaction Coordinator	Stopped	
	MSISCSI		Microsoft iSCSI Initiator Service	Stopped	netsvcs
	msiserver		Windows Installer	Stopped	
	MsKeyboardFilter		Microsoft Keyboard Filter	Stopped	netsvcs
	MySQL56	5392	MySQL56	Stopped	
	MySQL57		MySQL57	Stopped	
	MySQL80		MySQL80	Running	

III. MySQL Relational Database

Езикът за релационни база данни е един с леки диалекти – просто мениджмънт системата за управление на данни е различна – Oracle, MariaDB, MySQL, PostgreSQL, etc.

https://www.w3schools.com/sql/sql_where.asp

За Judge на СофтУни – ако не е зададено другояче, базата данни /схемата не я цитираме като подаваме решенията си в Judge.

0. Някои basic неща

MySQL е case insensitive – команди можем да пишем както с големи, така и с малки букви

Слагаме коментари с: # или /*.... */

Ctrl + / - слага в коментар по друг начин

Полета/обекти винаги ограждаме с тилда кавички `....` като по този начин escape-ваме запазени думи в SQL

Ctrl + D – добавяме един ред като горния ред

Числата въвеждаме без скоби

Текстовете ограждаме с единични обикновени скоби '....'

За пари ползваме DECIMAL вместо DOUBLE

NULL

TRUE

FALSE

= присвоява стойност – ДА и знак за сравнение - ДА

SET e_count := присвояване

!= или **<>** значат и двете различно, работи само за числови стойности. Иначе използваме IS NULL / IS NOT NULL

>= по-голямо

<= по-малко

Като кликнем колоните на таблица, то името на колоната се нанася в SQL заявката – да не си играем да пишем ръчно името на колоната

Като цяло, при базите данни, гледаме 90% предварително, и след това пишем заявките.

Ctrl + R – reverse Engineering

Ctrl + Space – Auto suggest

Реално не можем да дебъгваме в MySQL дадена функция/процедура/или какво и да е

+

-

*

/ - обикновено **дробно** делене

% или mod() модулно делене не работи по нормален начин

1. Introduction to MySQL

1.0. General info

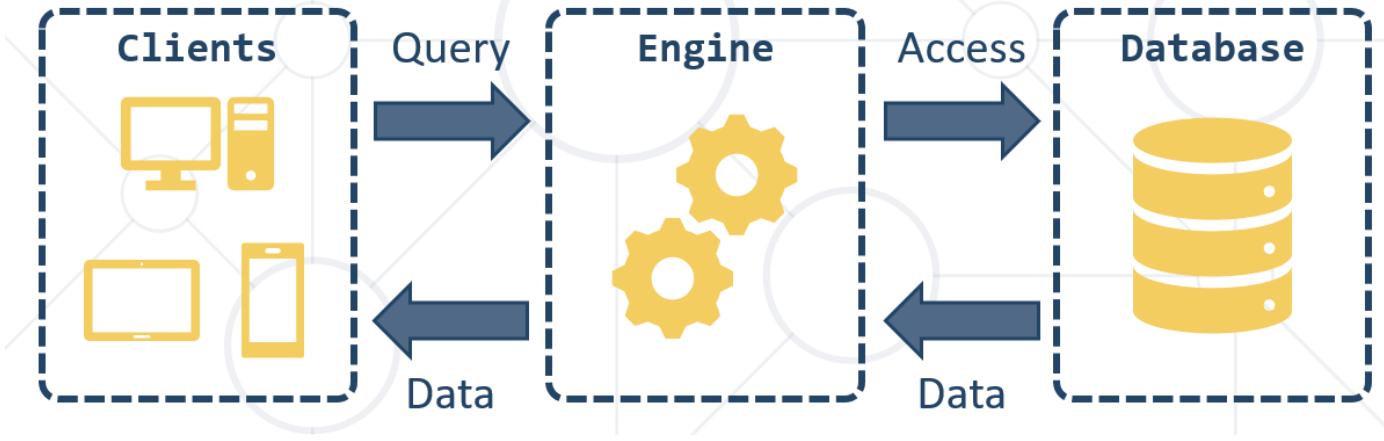
- A database is an **organized** collection of **related** information
 - It imposes **rules** on the contained data
 - Access to data is usually provided by a "**system**" (DBMS) **database management**
 - Relational storage first proposed by Edgar Codd in 1970

- **Relational Data Base Management System**
 - Database **management**
 - It **parses requests** from the user and takes the **appropriate** action
 - The user **doesn't have direct access** to the stored data

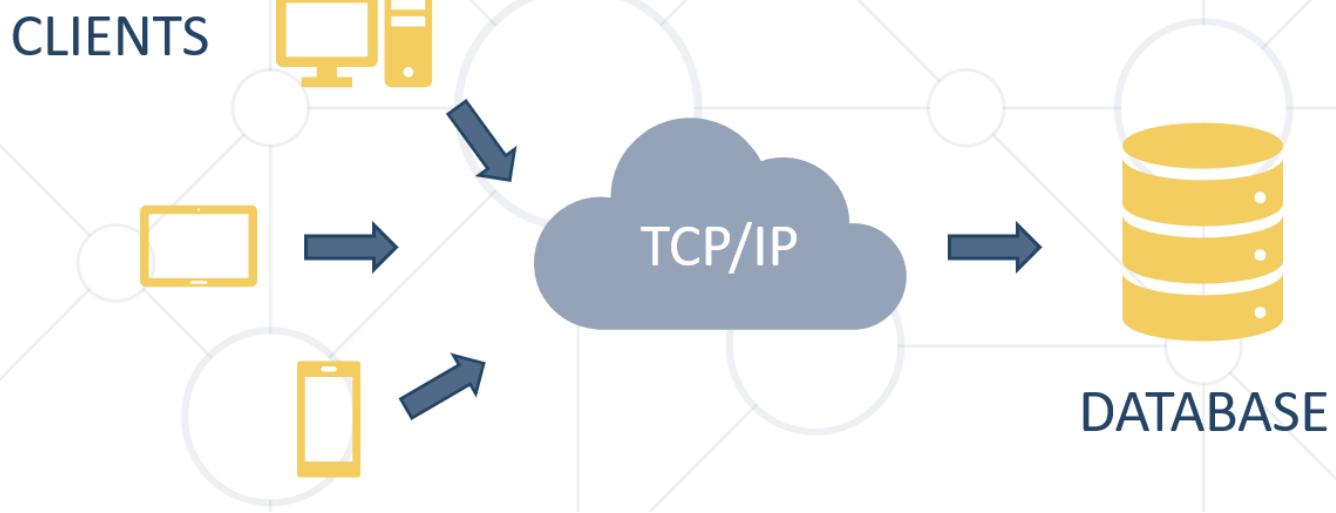
- Data is presented by **relations** – collection of tables related by **common fields**
- MS SQL Server, DB2, Oracle and MySQL

1.1. Database Engine Flow

▪ SQL Server uses the Client-Server Model



Client-Server Model



1.2. Structured Query Language = SQL

- Queries
- Clauses
- Expressions
- Predicates
- Statements



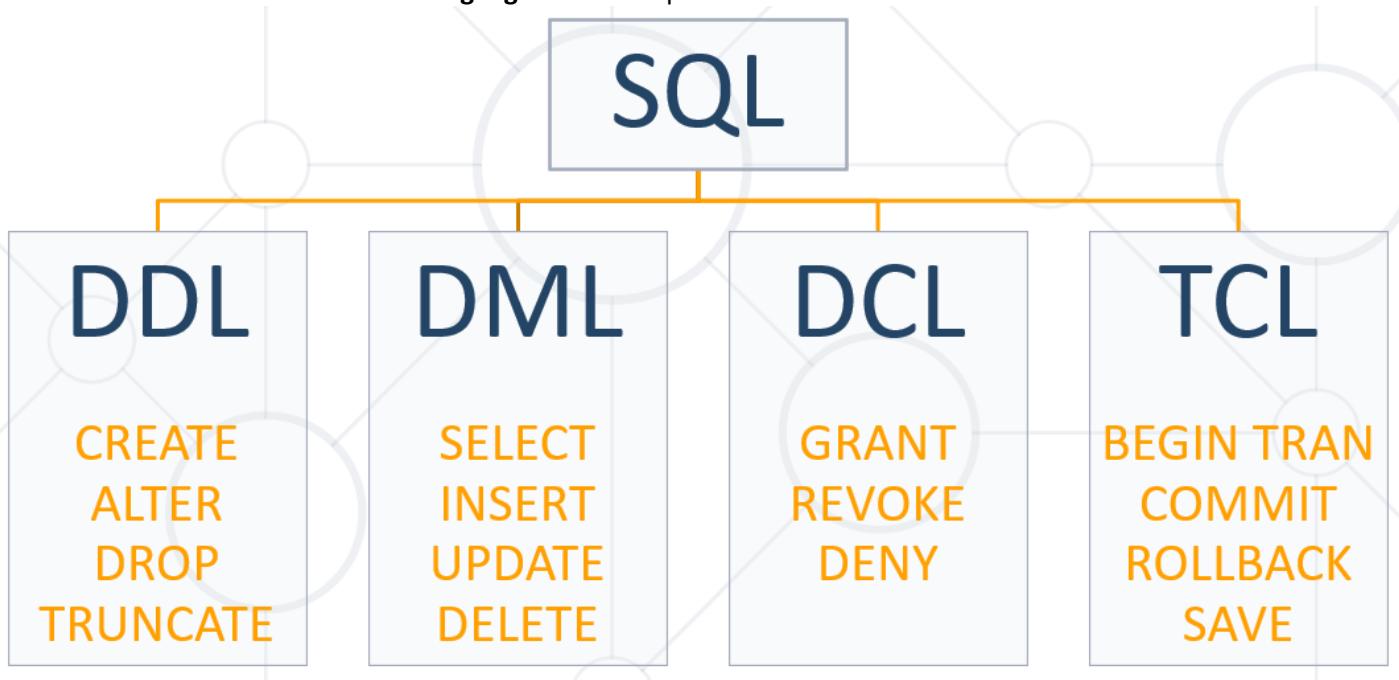
1.3. DDL, DML, DCL, TCL

Logically divided in four sections

- **Data Definition Language** – describe the structure of our data = **DDL**
- **Data Manipulation Language** – store and retrieve data = **DML**

CRUD is part of **DML** – Create, Read, Update, Delete

- **Data Control Language** – define who can access the data = **DCL**
- **Transaction Control Language** – bundle operations and allow rollback = **TCL**



1.3. DQL

Data query language (DQL) is part of the base grouping of SQL sub-languages. These sub-languages are mainly categorized into:

- a data query language (**DQL**),
- a data definition language (DDL),
- a data manipulation language (DML).
- a data control language (DCL),

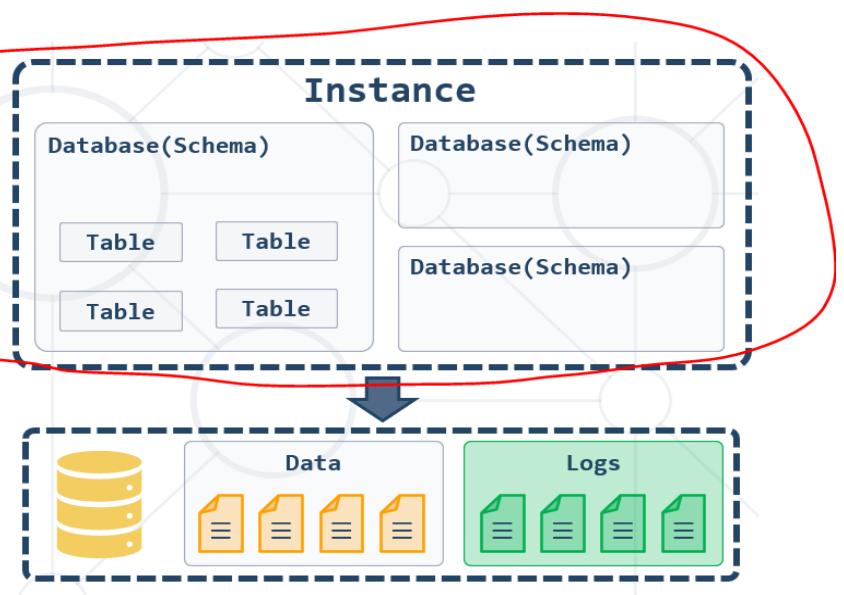
- Transaction control language (TCL)

DQL statements are used for performing queries on the data within schema objects. The purpose of DQL commands is to get the schema relation based on the query passed to it.

Although often considered part of DML, the SQL [SELECT](#) statement is strictly speaking an example of DQL. When adding FROM or WHERE data manipulators to the SELECT statement the statement is then considered part of the DML.

1.4. MySQL Server Architecture

- Logical Storage
 - Instance
 - Database/Schema
 - Table
- Physical Storage
 - Data files and Log files
 - Data pages



1.5. Database Table Elements

- The table is the main **building block** of any database
- Each **row** is called a **record** or **entity**
- Columns (**fields**) define the **type** of data they contain

1.6. Table Relationships

- We split the data and introduce **relationships** between the tables to **avoid** repeating information
- Connection via **Foreign Key** in one table pointing to the **Primary Key** in another

user_id	first	last	registered
203	David	Rivers	05/02/2016
204	Sarah	Thorne	07/17/2016
205	Michael	Walters	11/23/2015

Primary Key

Foreign Key

user_id	email
203	drivers@mail.cx
204	sarah@mail.cx
205	walters_michael@mail.cx
203	david@homedomain.cx

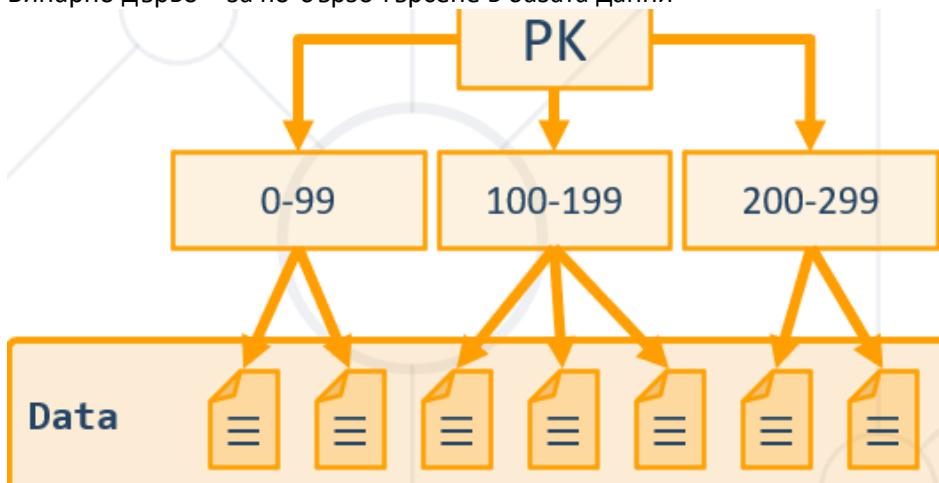
1.7. Programmability

1.7.1. Indices

- Indices make data lookup faster
 - Clustered – bound to the **primary key**, physically sorts data

- Non-Clustered – can be **any field**, references the primary index
- Structured as an **ordered tree**

Бинарно дърво – за по-бързо търсене в базата данни



1.7.2. Views

- Views are **prepared queries** for displaying **sections** of our data

```
CREATE VIEW v_employee_names AS
SELECT employee_id,
        first_name,
        last_name
FROM employees
```

```
SELECT * FROM v_employee_names
```

- Evaluated at **run time** – they do not increase performance

1.7.3. Procedures, Functions and Triggers

A database can further be customized with reusable code

- **Procedures** – carry out a predetermined **action**
 - E.g. get all employees with salary above 35000
- **Functions** – receive **parameters** and return a **result**
 - E.g. get the age of a person using their birthdate and current date
- **Triggers** – **watch** for activity in the database and **react** to it
 - E.g. when a record is deleted, write it to an archive

1.8. Data Types in MySQL Server

1.8.1. Numeric Data Types

- Numeric data types have certain range
- Their range can be changed if they are:
 - **Signed** - represent numbers both in the positive **and** negative ranges
 - **Unsigned** - represent numbers **only** in the positive range
- E.g. signed and unsigned INT:

Signed Range		Unsigned Range	
Min Value	Max Value	Min Value	Max Value
-2147483648	2147483648	0	4294967295

- **INT [(M)] [UNSIGNED]** – INT(10) – число с 10 цифри
 - TINYINT, SMALLINT, MEDIUMINT, BIGINT
- **DOUBLE [(M, D)] [UNSIGNED]**

Digits stored for value

Decimals after floating point

- E.g. DOUBLE(5, 2) – 999.99 – общо са 5 цифри, като след десетичната запетая са 2 цифри
- **DECIMAL [(M, D)] [UNSIGNED] [ZEROFILL]** – слага нули отпред ако е нужно

DECIMAL за по-голяма точност

M – общо брой цифри

D- от които след десетичната запетая

1.8.2 String Types

String column definitions include attributes that specify the **character set or collation**.

- **CHARACTER SET** (Encoding) - Determines the storage of each character (single or multiple bytes)

E.g. utf8, ucs2

- **CHARACTER COLLATION** – rules for encoding comparison - Determines the sorting order and case-sensitivity
 - E.g. latin1_general_ci, Traditional_Spanish_ci_ai etc

- **Set and collation** can be defined at the database, table or column level

Non-unicode (just English, western languages)

- **CHAR (M)** - up to 255 characters
 - fixed-length character type (example CHAR(30) или CHAR(1))
- **VARCHAR(M)** – up to 65 535 characters
 - Variable max size
- **TEXT** – up to 65 535 characters
 - TINYTEXT, MEDIUMTEXT, LONGTEXT
- **BLOB** - Binary Large Object [(M)] - 65 535 ($2^{16} - 1$) characters – когато не пазим адреса/пътя към снимката, а пазим битовата символна версия от 1000+ символа на самата снимка
 - TINYBLOB, MEDIUMBLOB, LONGBLOB

TINYBLOB : L < 2^8 = 256 Bytes

BLOB : L < 2^{16} = 65,536 Bytes

MEDIUMBLOB : L < 2^{24} = 16,777,216 Bytes

LONGBLOB : L < 2^{32} = 4,294,967,296 Bytes

fieldName Blob(size in bytes) -

The short answer is: **VARCHAR** is variable length, while **CHAR** is fixed length. **CHAR** is a fixed length string data type, so any remaining space in the field is padded with blanks. **CHAR** takes up 1 byte per character. ... **VARCHAR** is a variable length string data type, so it holds only the characters you assign to it.

Двойно повече място за Unicode (all languages worldwide) - Supports many client computers that are running different locales.

nchar/nvarchar - произлиза от national

	char	nchar	varchar	nvarchar
Character Data Type	Non-Unicode fixed-length	Unicode fixed-length can store both non-Unicode and Unicode characters (i.e. Japanese, Korean etc.)	Non-Unicode variable length	Unicode variable length can store both non-Unicode and Unicode characters (i.e. Japanese, Korean etc.)
Maximum Length	up to 8,000 characters	up to 4,000 characters	up to 8,000 characters	up to 4,000 characters
Character Size	takes up 1 byte per character	takes up 2 bytes per Unicode/Non-Unicode character	takes up 1 byte per character	takes up 2 bytes per Unicode/Non-Unicode character
Storage Size	n bytes	2 times n bytes	Actual Length (in bytes)	2 times Actual Length (in bytes)
Usage	use when data length is constant or fixed length columns	use only if you need Unicode support such as the Japanese Kanji or Korean Hangul characters due to storage overhead	used when data length is variable or variable length columns and if actual data is always way less than capacity	use only if you need Unicode support such as the Japanese Kanji or Korean Hangul characters due to storage overhead
			query that uses a varchar parameter does an index seek due to column collation sets	query that uses a nvarchar parameter does an index scan due to column collation sets

1.8.3. Date Types

- **DATE** - for values with a date part but **no time part** - 'YYYY-MM-DD' or 'YY-MM-DD'
- **TIME** - for values with time but **no date part** – 'hh: mm: ss'
- **DATETIME** - values that contain both date **and** time parts - 'YYYY-MM-DD hh: mm: ss'
- **TIMESTAMP** - both date **and** time parts

- MySQL retrieves values for a given date type in a **standard output format**
E.g. as a string in either 'YYYY-MM-DD' or 'YY-MM-DD'

ВАЖНО

Когато сравняваме дата DATE с DATETIME, изречението „hired after 1/1/1999“ го тълкуваме WHERE e.`hire_date` >= '1999-01-02'
зашпото имаме **DATETIME 1999-12-12 01:26:00.000000**

1.8.4. Boolean Types

`gender` **BOOLEAN**;

```
CREATE TABLE `people`(
`id` INT NOT NULL UNIQUE AUTO_INCREMENT PRIMARY KEY,
`name` VARCHAR(200) NOT NULL,
`picture` MEDIUMBLOB,
`height` DOUBLE(5,2),
`weight` DOUBLE(5,2),
`gender` CHAR(1) NOT NULL,
`birthdate` DATE NOT NULL,
```

```
'biography' LONGTEXT  
);
```

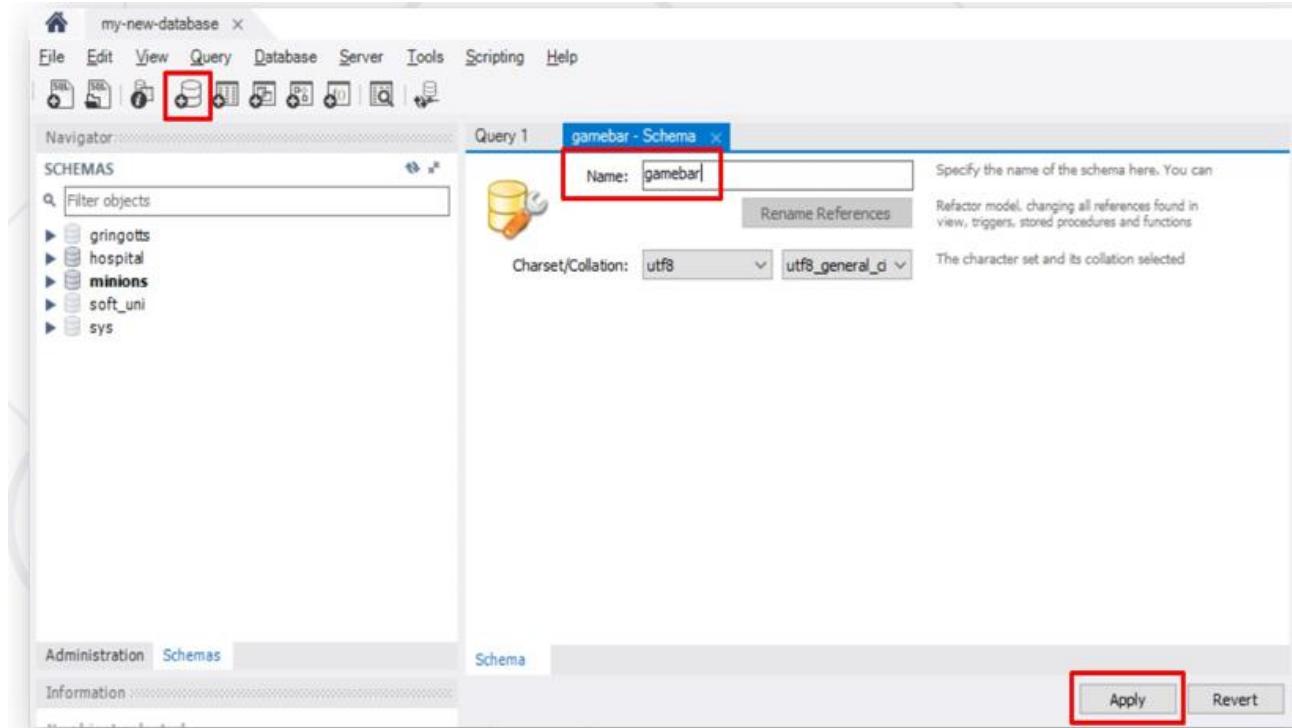
1.9. DDL – Data Definition Language - Database Modelling – using GUI or via basic SQL queries

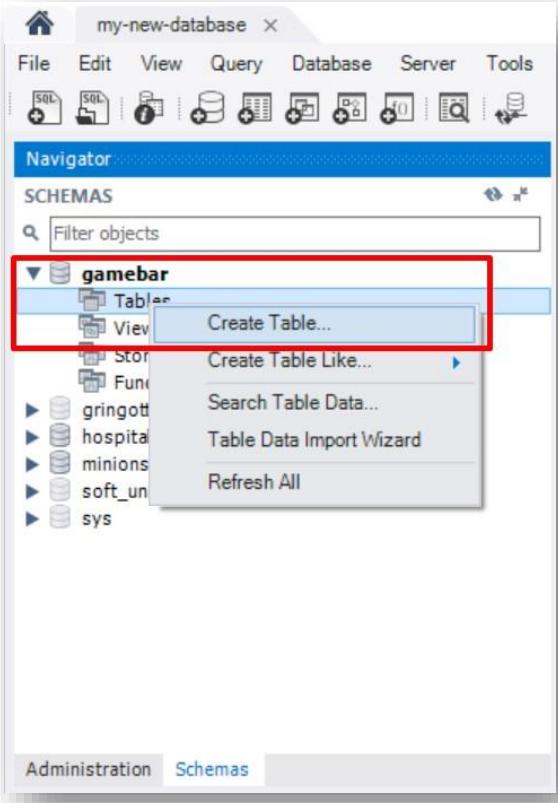
Да даваме Refresh от време на време

- **Working with IDEs – MySQL Workbench Database Management System** - we can use GUI Clients to **create** and **customize** tables
- Enables us:
 - To **create** a new database
 - To create **objects in the database** (tables, stored procedures, relationships and others)
 - To **change** the properties of objects
 - To **enter records** into the tables

1.9.0. Creating a New Database

- Select **Create new schema** from the command menu

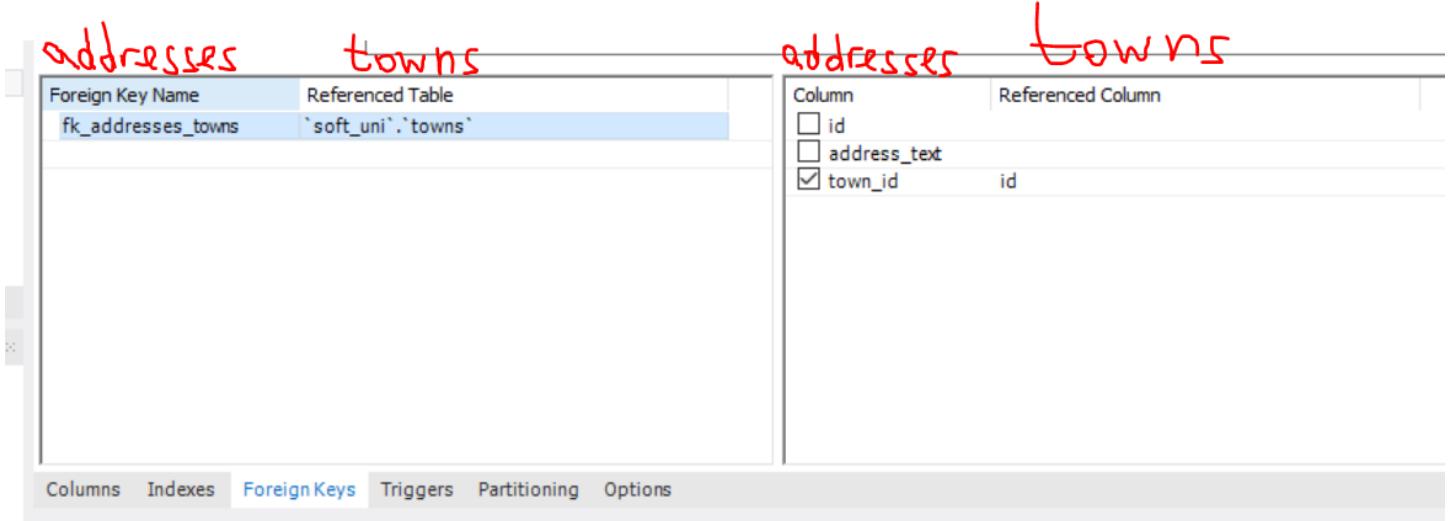




- A Primary Key is used to uniquely identify and index records

The screenshot shows a 'Query 1' window with the title 'employees - Table'. The 'Table Name:' is set to 'employees' and the 'Schema:' is 'gamebar'. In the 'Columns' tab, there is one column defined: 'id' with a 'Datatype' of 'INT(11)'. The 'PK' (Primary Key) checkbox is checked, and the 'NN' (Not Null) and 'AI' (Auto Increment) checkboxes are also checked. In the 'Indexes' tab, the 'Storage' section has several checkboxes: 'Primary Key' (checked), 'Not Null' (checked), 'Unique' (unchecked), 'Binary' (unchecked), 'Unsigned' (unchecked), 'Zero Fill' (unchecked), 'Auto Increment' (checked), and 'Generated' (unchecked). At the bottom of the window are tabs for 'Columns', 'Indexes', 'Foreign Keys', 'Triggers', 'Partitioning', and 'Options', along with 'Apply' and 'Revert' buttons.

1.9.1. Foreign keys



- Adding foreign keys

A FOREIGN KEY is a field (or collection of fields) in one table, that refers to the PRIMARY KEY in another table.

The table with the foreign key is called the child table, and the table with the primary key is called the referenced or parent table.

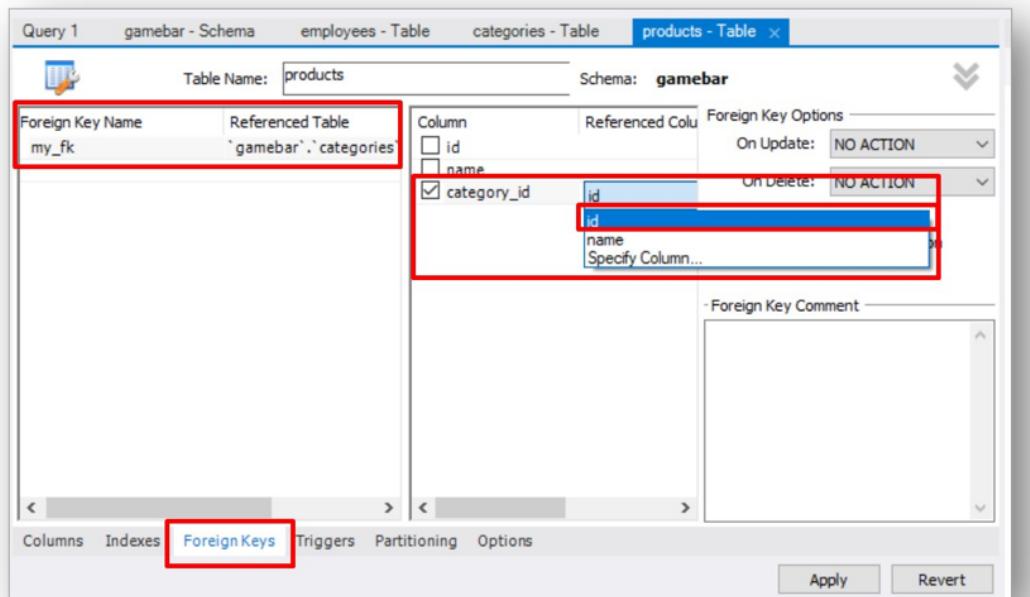
Конвенция при изписване на foreign key поле: fk_fromMinions_toTowns

В minions е чуждия ключ, а в Towns е primary ключ.

ВАЖНО: когато създаваме foreign keys, първо трябва да създадем таблицата, от която foreign key ще взема данни.

Foreign keys are created in the "Foreign keys" tab:

- **Reference table** – select the table from which you will choose a column to link your foreign key – "categories";
 - **Columns** – select the column you want to be set as foreign key – "category_id";
 - **Referenced columns** – select the column set to primary to link the foreign key – "id";



Task 11* movies - Table X

Index Name	Type
PRIMARY	PRIMARY
fk_movies_directors	INDEX
fk_movies_genres	INDEX
fk_movies_categories	INDEX

Column	#	Order
<input type="checkbox"/> id		ASC
<input type="checkbox"/> title		ASC
<input type="checkbox"/> director_id		ASC
<input type="checkbox"/> copyright_year		ASC
<input type="checkbox"/> length		ASC
<input type="checkbox"/> genre_id		ASC
<input checked="" type="checkbox"/> category_id	1	ASC
<input type="checkbox"/> rating		ASC
<input type="checkbox"/> notes		ASC

Columns Indexes Foreign Keys Triggers Partitioning Options

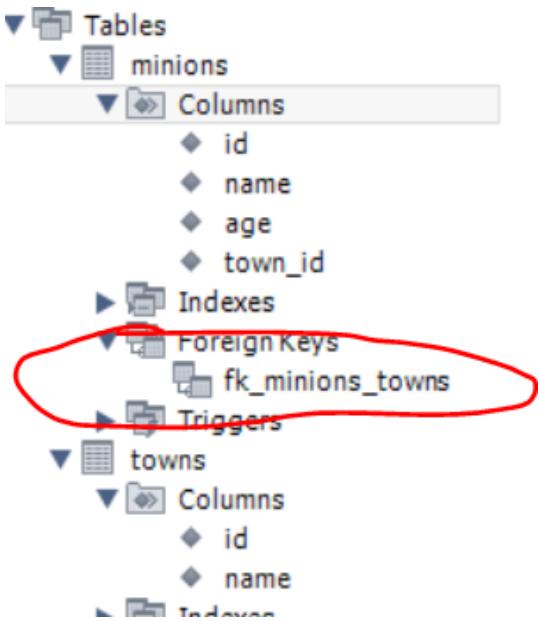
```
CREATE TABLE `Orders` (
  `OrderID` int NOT NULL,
  `OrderNumber` int NOT NULL,
  `PersonID` int,
  PRIMARY KEY (`OrderID`),

  CONSTRAINT `fk_source_target`
  FOREIGN KEY (`Orders`(`PersonID`))
  REFERENCES `Persons`(`PersonID`)
);
```

SQL FOREIGN KEY on ALTER TABLE

```
ALTER TABLE `products`
ADD CONSTRAINT `fk_products_categories`
FOREIGN KEY `products`(`category_id`)
REFERENCES `categories`(`id`);
```

```
ALTER TABLE `minions`.`minions`
ADD CONSTRAINT `fk_minions_towns`
FOREIGN KEY (`town_id`)
REFERENCES `minions`.`towns`(`id`);
```



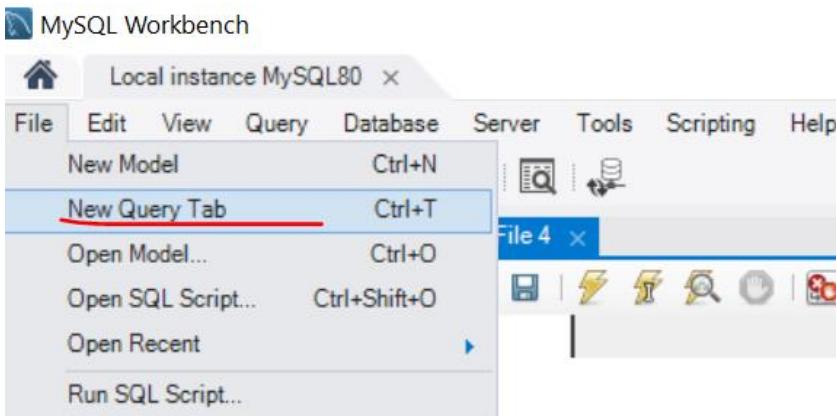
1.9.2. Where to run our SQL queries

Working with basic SQL queries

- We communicate with the database engine using SQL
- Queries provide greater **control** and **flexibility**

Queries are written in the "Query" tab.

Database creation



```
1 • CREATE schema `gamebar`;
2 • USE Execute the selected portion of the script or everything, if there is no selection
3 • create table `gamebar`.`employees`(
4     `id` INT PRIMARY KEY NOT NULL AUTO_INCREMENT,
5     `first_name` VARCHAR(45) NOT NULL,
6     `last_name` VARCHAR(45) NOT NULL);
7 • create table `gamebar`.`categories`(
8     `id` INT PRIMARY KEY NOT NULL AUTO_INCREMENT,
9     `name` VARCHAR(45) NOT NULL);
10 • USE `gamebar`;
11 • create table `gamebar`.`products`(
12     `id` INT PRIMARY KEY NOT NULL AUTO_INCREMENT,
13     `name` VARCHAR(45) NOT NULL,
14     `category_id` INT NOT NULL);
```

CREATE DATABASE `gamebar`; или **CREATE SCHEMA `gamebar`;**

Ако искаме да отворим съществуващ SQL script, то правилния начин е цъкнем **File -> Open SQL script**

1.9.3. Table Creation in SQL:

В работната част gamebar, създай таблица employees – за графично, виж скрийншота по-горе

The command **USE** – ако имаме отворени няколко база данни, да знаем с коя работим
USE `gamebar`

```
CREATE TABLE `gamebar`.`employees`(
    `id` INT NOT NULL AUTO_INCREMENT,
    `name` VARCHAR(45) NOT NULL,
    PRIMARY KEY (`id`));
```

Или

```
create table `gamebar`.`categories`(
    `id` INT PRIMARY KEY NOT NULL AUTO_INCREMENT,
    `name` VARCHAR(45) NOT NULL);
```

Sugar syntax for AUTO_INCREMENT – задаване на първоначална стойност, от която да започне да инкрементира с единица:

```
CREATE TABLE models (
    `model_id` INT AUTO_INCREMENT UNIQUE NOT NULL,
    `name` VARCHAR(20) NOT NULL,
    `manufacturer_id` INT NOT NULL
) AUTO_INCREMENT = 101; //започни от 101 като първи запис
```

1.9.4. Add records in SQL:

Option 1 – през графични interface

The screenshot shows the MySQL Workbench interface. On the left, the 'Schemas' tree view has 'gamebar' selected. Under 'Tables', 'employees' is selected. The right pane shows the SQL editor with the query: `SELECT * FROM gamebar.emp`. Below it is the 'Result Grid' showing the following data:

	id	first_name	last_name
1	1	Svilen	Velikov
2	2	Ivan	Petrov
3	3	Tsvetomir	Velikov
*	NULL	NULL	NULL

Опция 2 – с SQL заявка и hardcore-нати стойности

```
INSERT INTO `gamebar`.`employees` (`id`, `first_name`, `last_name`) VALUES ('1', 'Svilen', 'Velikov');
INSERT INTO `gamebar`.`employees` (`id`, `first_name`, `last_name`) VALUES ('2', 'Ivan', 'Petrov');
INSERT INTO `gamebar`.`employees` (`id`, `first_name`, `last_name`) VALUES ('3', 'Tsvetomir', 'Velikov');
```

Или така:

```
INSERT INTO `towns` (`id`, `name`)
VALUES
(1, 'Sofia'),
(2, 'Plovdiv'),
(3, 'Varna');
```

Или ако вкарваме всичко:

```
INSERT INTO `towns` пропускаме скобите или слагаме само празни скоби
VALUES
(1, 'Sofia'),
(2, 'Plovdiv'),
(3, 'Varna');
```

Опция 3 – с SQL заявка, без VALUES и с функция за определяне/за автоматично попълване

Пример 1

```
INSERT INTO cards(card_number, card_status, bank_account_id)
(без тези скоби в judge
SELECT REVERSE(full_name), 'Active', id
FROM clients
WHERE id>=191 AND id<=200
); без тези скоби в judge
```

Пример 2

```
INSERT INTO `coaches`(`first_name`, `last_name`, `salary`, `coach_level`)
```

```
SELECT `first_name`, `last_name`, `salary`,
CHAR_LENGTH(`first_name`)
FROM `players`
WHERE `age` >= 45;
```

Пример 3

```
INSERT INTO cards(card_status, card_number, bank_account_id)
```

(без тези скоби в judge

```
SELECT (
```

```
    CASE
```

```
        WHEN id BETWEEN 191 AND 199 THEN 'Active'
        WHEN id BETWEEN 200 AND 299 THEN 'Inactive'
        WHEN id BETWEEN 300 AND 500 THEN 'Deleted'
```

```
    END
```

```
) AS customs_status, REVERSE(full_name), id
```

```
FROM clients
```

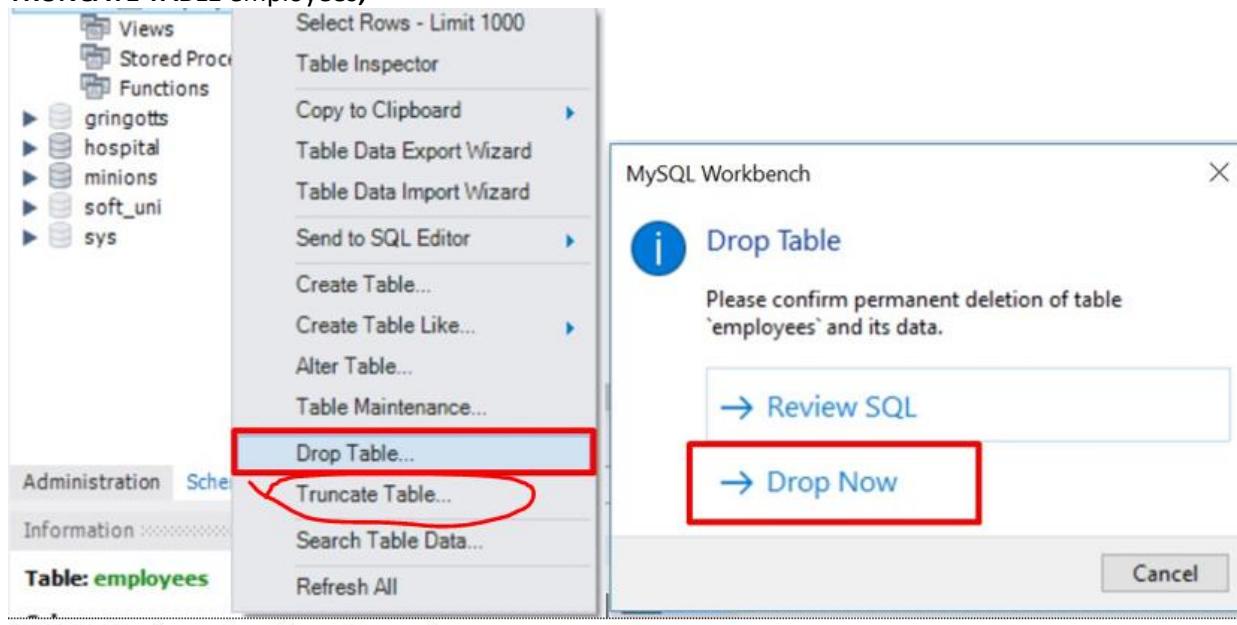
); без тези скоби в judge

1.9.5. Erasing data in SQL:

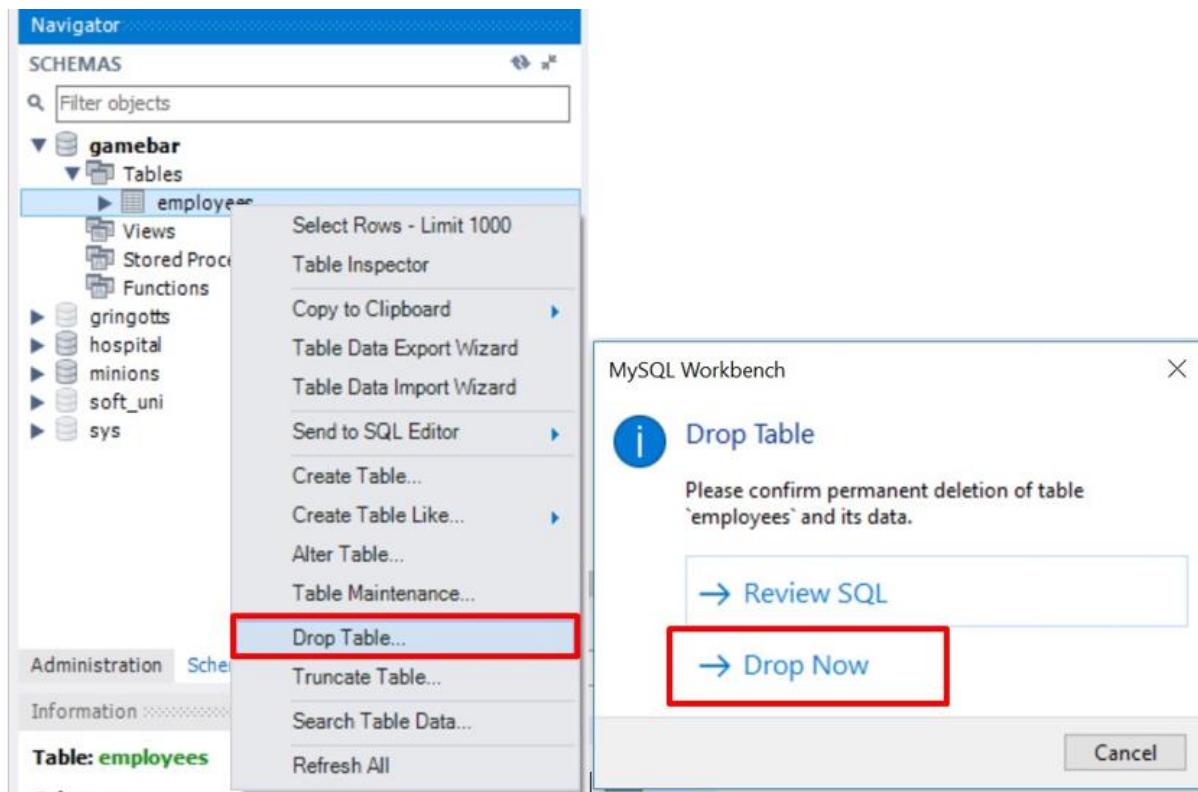
- Deleting structures is called **dropping**
 - You can drop **keys, constraints, tables** and entire **databases**
- Deleting all data in a table is called **truncating**
- Both of these actions **cannot be undone – use with caution!**

- To delete all the entries in a table, but keep the table structure

```
TRUNCATE TABLE employees;
```

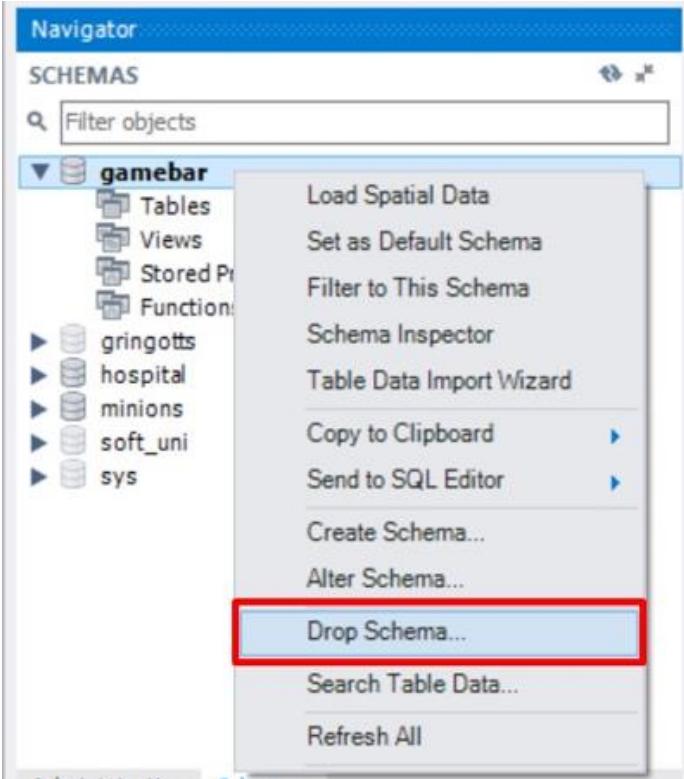


- Dropping table - delete data and structure



```
DROP Table `gamebar`.`employees`
```

- Dropping the entire Database



```
DROP DATABASE `gamebar`
```

- To remove a constraining rule from a column

- Primary keys, value constraints and unique fields

```
ALTER TABLE employess DROP CONSTRAINT pk_id;
```

- To remove **DEFAULT** value (if not specified, revert to NULL)

```
ALTER TABLE employess
```

```
ALTER COLUMN clients
```

```
DROP DEFAULT;
```

1.10. Table Customization

Primary Key

```
id INT PRIMARY KEY;
```

Not null – държим да има запис в това поле

```
id INT NOT NULL PRIMARY KEY;
```

Auto-Increment (Identity)

```
id INT AUTO_INCREMENT PRIMARY KEY;
```

Unique constraint – no repeating values in entire table

```
email VARCHAR(50) UNIQUE;
```

Default value – if not specified (otherwise set to NULL)

```
balance DECIMAL(10,2) DEFAULT 0;
```

1.11. More DDL - Altering Tables

A table can be changed using the keywords ALTER TABLE

```
ALTER TABLE employees;
```

Add new column

```
ALTER TABLE employees ADD salary DECIMAL; - добавя колона salary от тип Decimal
```

```
ALTER TABLE `gamebar`.`employees`
```

```
ADD COLUMN `middle_name` VARCHAR(45) NOT NULL AFTER `last_name`;
```

```
ALTER TABLE `users`
```

```
ADD COLUMN `pk_users` VARCHAR(45) NOT NULL AFTER `id`;
```

Changing type of a column / changing name of a column

```
ALTER TABLE `minions`.`towns`
```

```
CHANGE COLUMN `town_id` `id` INT NOT NULL AUTO_INCREMENT ;
```

Delete existing column – изтрива колона

```
ALTER TABLE people DROP COLUMN full_name;
```

```
ALTER TABLE `users` DROP COLUMN `pk_users`;
```

Modify data type of existing column

ALTER TABLE people **MODIFY COLUMN** email VARCHAR(100); - колоната email става от нов тип

Add primary key to existing column

ALTER TABLE people ADD CONSTRAINT PRIMARY KEY(); - Constraint name
PRIMARY KEY (id); - Column name (more than one for composite key)

```
ALTER TABLE `users` ADD CONSTRAINT PRIMARY KEY(`pk_users`);
```

Deleting primary key from a table

ALTER TABLE people
DROP PRIMARY KEY; - Column name (more than one for composite key)

```
ALTER TABLE `users`  
DROP PRIMARY KEY;
```

Add constraint / Add unique constraint

ALTER TABLE people **ADD CONSTRAINT** uq_email - Constraint name
UNIQUE (email); - columns names

Това не копира обединени данни в колона `pk_users`, прави следното – задай ограничение pk_users, което да е primary key от id и username

```
ALTER TABLE `users`
```

DROP PRIMARY KEY

ADD CONSTRAINT `pk_users`

PRIMARY KEY `users`(`id`, `username`);

```
ALTER TABLE `users`  
DROP PRIMARY KEY,  
ADD CONSTRAINT `pk_users`  
PRIMARY KEY `users`(`id`),  
CHANGE COLUMN `username` `username` VARCHAR(50) UNIQUE;
```

Set default value

```
ALTER TABLE people ALTER COLUMN balance SET DEFAULT 0;
```

- Set default value – вариант 2

```
ALTER TABLE `users`
```

CHANGE COLUMN `last_login_time` `last_login_time` DATETIME NULL DEFAULT CURRENT_TIMESTAMP ;

От типа данни на NOW()

ALTER TABLE `users`

```
CHANGE COLUMN `last_login_time` `last_login_time` DATETIME NULL DEFAULT NOW();
```

ALTER TABLE `users`

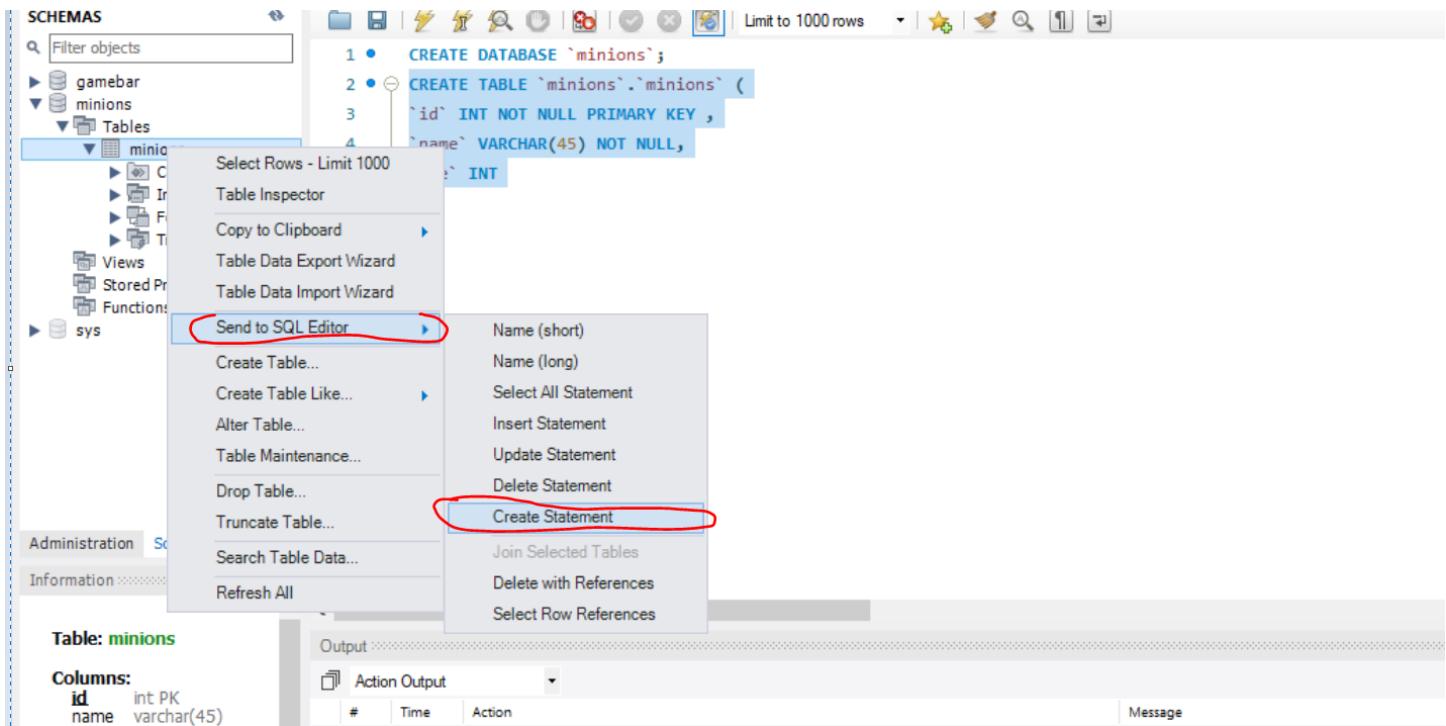
CHANGE COLUMN `username` `username` VARCHAR(30) NOT NULL DEFAULT 'Bai Peshu Starshi'

1.12. How to cheat – to see the SQL query

Create Statement – използваме го, за да създадем лесно SQL Заявка без да пишем всичко ръчно

Insert Statement – използваме го, за да видим SQL заявката, на това което сме създали

Update Statement – използваме го, за да обновим базата данни/да обновим поле/таблица



1.13. Advanced SQL queries

Сортира в alphabetic ред

SELECT `name` FROM `towns`

ORDER BY `name`; -

Сортира Double в низходящ ред

SELECT * FROM `employees`

ORDER BY `salary` DESC;

ORDER BY `salary` е същото като **ORDER BY `salary` ASC**

To show sorted only **some of the columns**

SELECT `first_name`, `last_name`, `job_title`, `salary` from `employees`

ORDER BY `salary` DESC;

Нанасяне на нова информация на даден ред за дадена колона – за всички записи

UPDATE `employees`

SET `salary` = `salary` * 1.1;

WHERE `id` > 0;

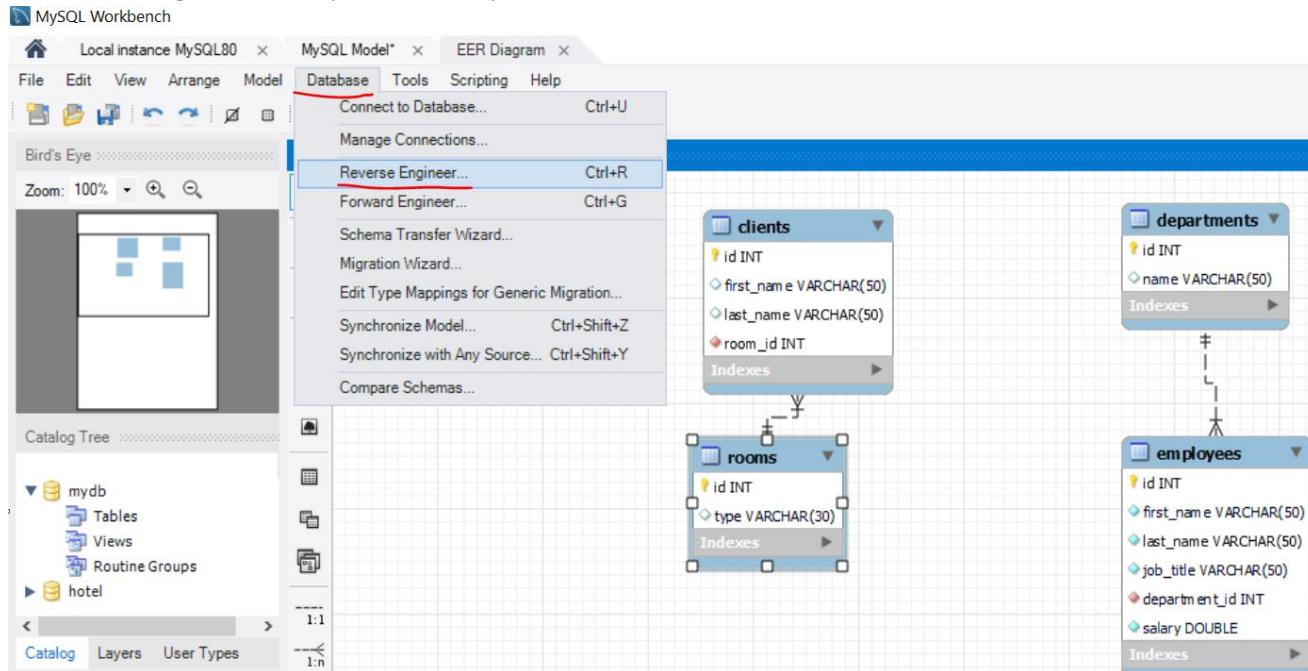
Нанасяне на нова информация на даден ред за дадена колона – за определен запис

UPDATE `employees`

SET `salary` = `salary` * 1.1;

WHERE `id` = 5;

1.14. E/R Diagram – диаграма на свързаността



2. DML - BASIC CRUD (Create, Read, Update, Delete) OPERATIONS

Which of the following is NOT a valid SELECT statement ?

- SELECT (SELECT ID FROM EMPLOYEES) EMP FROM EMP**
- SELECT * FROM Employees WHERE (SELECT SALARY FROM EMPLOYEES WHERE ID == 1) == 2000
- SELECT * FROM (SELECT SALARY FROM EMPLOYEES WHERE ID == 1)
- SELECT * FROM EMPLOYEES WHERE ID IN (1)

2.1. Query Basics

- Select first, last name and job title about employees:

SELECT `first_name`, `last_name`, `job_title` FROM `employees`;

- Select projects which start on 01-06-2003:

SELECT * FROM `projects` WHERE `start_date`='2003-06-01';

- Inserting data into table – можем да insert-нем определени колонки, но тези които изпускаме не трябва да са NOT NULL. А тези, които са AUTO_INCREMENT – сами се увеличават дори да не вкарваме данни за тях

INSERT INTO projects(`name`, `start_date`)

VALUES('Introduction to SQL Course', '2006-01-01');

Опция 3 – с SQL заявка и функция за определяне/за автоматично попълване

INSERT INTO cards(card_number, card_status, bank_account_id)

(

SELECT REVERSE(full_name), 'Active', id

FROM clients

WHERE id>=191 AND id<=200

);

- Update several cells for specific rows:

```
UPDATE `projects`
SET `end_date` = '2006-08-31', `id` = 3;
WHERE `start_date` = '2006-01-01';
```

- Update specific cells/columns for all rows/records:

```
UPDATE `employees`
SET `salary` = `salary` * 1.1;
```

- Delete specific projects – изтрива целия ред

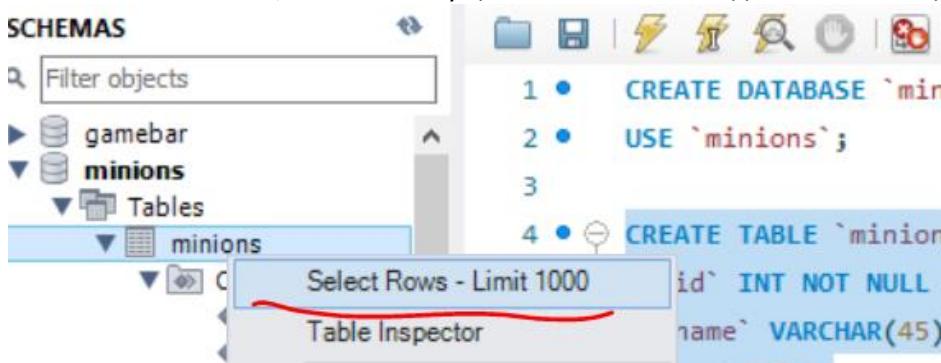
```
DELETE FROM `projects`
WHERE `start_date` = '2006-01-01';
```

2.2. Retrieving Data

Retrieve Records in SQL - using GUI:

- Get all information from a table

SELECT * FROM towns; - покажи текущите записи в базата данни от таблица towns.



The screenshot shows the 'Result Grid' window in MySQL Workbench. It displays the following data:

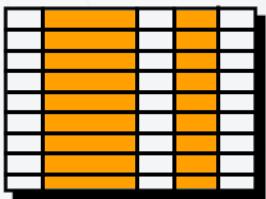
	id	name
▶	1	Sofia
▶	2	Plovdiv
▶	3	Varna
*	NULL	NULL

- You can limit the columns and number of records

SELECT first_name, last_name **FROM employees **LIMIT** 5;** - ограничи до 5 записи

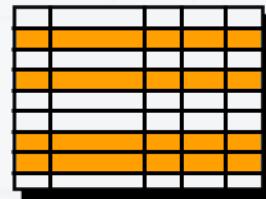
Projection

Take a subset of the columns



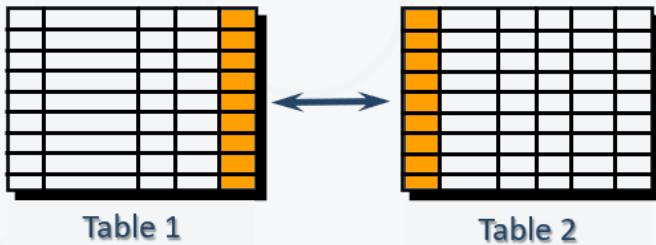
Selection

Take a subset of the rows



Join

Combine tables by some column



```
SELECT * FROM employees;
```

List of columns (* for all)

Table name

```
SELECT `id`, `first_name`, `last_name`, `job_title`  
FROM `employees`  
ORDER BY `id`;\b  
WHERE  
LIMIT 3;
```

Всички колони плюс още колони

```
SELECT *,`id`,`first_name`,`last_name`,`job_title`  
FROM `employees`  
ORDER BY `id`;\b  
WHERE  
LIMIT 3;
```

Aliases(прякор/друго име) rename a table or a column heading – използваме задължително когато работим с повече от една таблица!!!

ВАЖНО – когато искаме да работим с колоната aliases – прякото име, то след AS използваме обикновени

кавички " , но в последствие използваме специалните кавички `` - НЕЕЕ Е ТАКА, МОЖЕ ДА СИ ИЗПОЛЗВАМЕ САМО ТИЛДА КАВИЧКИ!

```

SELECT e.id AS 'No.',
e.first_name AS 'First Name',
e.last_name AS 'Last Name',
e.job_title AS 'Job Title'
FROM employees AS e
ORDER BY `Job title`;

```

Пример за Aliases когато работим едновременно с 2 таблици:

```

SELECT p.`peak_name`,
r.`river_name`,
LOWER(CONCAT(p.`peak_name`, SUBSTRING(r.`river_name`, 2))) AS `mix`
FROM `peaks` AS p, `rivers` AS r
WHERE RIGHT(LOWER(p.`peak_name`), 1) = LEFT(LOWER(r.`river_name`), 1)
ORDER BY `mix`;

```

```

SELECT
5+5 AS 'staticnumber',
`job_title` AS 'Job Title',
`id` AS 'No.'
FROM `employees`;

```

Concatenation – когато ги обединява в резултата от SELECT

concat() - returns the string that results from concatenating the arguments - предефинирана функция в MySQL

- String literals are enclosed in ['](single quotes)
- Table and column names containing special symbols use ['] (backtick)

```

SELECT concat(`first_name`, ' ', `last_name`) AS 'full_name',
`job_title` AS 'Job Title',
`id` AS 'No.'
FROM `employees`;

```

*Another function of concatenation is **concat_ws()** - stands for concatenate with separator and is a special form of CONCAT() – с първия стринг лепим останалите*

```

Separator          Arguments
SELECT concat_ws([ ], `first_name`, `last_name`, `job_title`)
AS 'full_name',
`job_title` AS 'Job Title',
`id` AS 'No.'
FROM `employees`;

```

- Skip any **NULL** values after the separator argument.

Concatenating with + : Add 2 strings together:

Не работи както трябва

Filtering the Selected Rows

- Use **DISTINCT** to eliminate duplicate results – to eliminate all the duplicate records and fetching only unique records.

```
SELECT DISTINCT `first_name`  
FROM `employees`;
```

Ако има повтарящи се имена, то покажи само веднъж повтарящото се име

- You can filter rows by specific conditions using the **WHERE** clause

```
SELECT `last_name`, `department_id`  
FROM `employees`  
WHERE `department_id` = 1;
```

- Other **logical operators** can be used for better control

```
SELECT `last_name`, `salary`  
FROM `employees`  
WHERE `salary` <= 20000;
```

- Conditions can be combined using **NOT**, **OR**, **AND** and brackets

```
SELECT `last_name` FROM `employees`  
WHERE NOT (`manager_id` = 3 OR `manager_id` = 4);
```

- Using **BETWEEN** operator to specify a range:

```
SELECT `last_name`, `salary` FROM `employees`  
WHERE `salary` BETWEEN 20000 AND 22000; - работи включително  
HAVING `max_salary` NOT BETWEEN 30000 AND 70000
```

- Using **IN / NOT IN** to specify a set of values:

```
SELECT `first_name`, `last_name`, `manager_id`  
FROM `employees`  
WHERE `manager_id` IN (109, 3, 16); - дали е измежду тези стойности
```

```
SELECT * FROM `towns`  
WHERE LOWER(SUBSTRING(`name`,1,1)) NOT IN('r', 'b', 'd')  
ORDER BY `name`;
```

Comparing with NULL

- **NULL** is a special value that means missing value
 - Not the same as **0** or a blank space
- Checking for **NULL** values

Проверка за различно по този начин **!=** и по този начин **<>** не можем да правим с **NULL!!!**

```
SELECT `last_name`, `manager_id`  
FROM `employees`  
WHERE `manager_id` = NULL; -ГРЕШНО!
```

```
SELECT `last_name`, `manager_id`  
FROM `employees`  
WHERE `manager_id` IS NULL;
```

```
SELECT `last_name`, `manager_id`
```

```
FROM `employees`  
WHERE `manager_id` IS NOT NULL;
```

Sorting with ORDER BY – може по две условия, отделяме със запетая

- Sort rows with the **ORDER BY** clause
 - **ASC**: ascending order, default
 - **DESC**: descending order

```
SELECT `last_name`, `hire_date`  
FROM `employees`  
ORDER BY `hire_date` DESC, `last_name` ASC;
```

Ако `hire_date` съвпада, то сортирай по следващ критерий `last_name`

Сортираме по много условия – отделяме със запетая

```
SELECT * FROM `employees`  
ORDER BY `salary` DESC, `first_name` ASC, `last_name` DESC, `middle_name` ASC, `employee_id`;
```

Views – все едно си запазваме предефинирана заявка

- Views are **virtual tables** made from others tables, views or joins between them
- Usage:
 - To simplify writing complex queries
 - To limit access to data for certain users

Views – Example 1

```
CREATE VIEW `v_hr_result_set` AS  
SELECT  
    CONCAT(`first_name`, ' ', `last_name`) AS 'Full Name',  
    `salary`  
FROM `employees` ORDER BY `department_id`;  
  
SELECT * FROM `v_hr_result_set`;
```

Example 2

```
CREATE VIEW `myview` AS  
SELECT `first_name`, 5 FROM `employees`  
ORDER BY `salary` DESC, `first_name` ASC, `last_name` DESC;  
  
SELECT * FROM `myview`;  
DROP VIEW `myview`; - заличи
```

2.3. Writing Data in Tables

The SQL INSERT command

INSERT INTO `towns` VALUES (33, 'Paris'); - values for all columns сме длъжни да подадем

```
INSERT INTO projects(`name`, `start_date`) - какви колко колони ще нанасяш  
VALUES ('Reflective Jacket', NOW())
```

Inserting data into table – можем да insert-нем определени колонки, но тези които изпускаме не трябва да са NOT NULL. А тези, които са AUTO_INCREMENT – сами се увеличават дори да не вкарваме данни за тях

Bulk data can be recorded in a single query, separated by comma

```
INSERT INTO `towns` (`id`, `name`)
VALUES
(1, 'Sofia'),
(2, 'Plovdiv'),
(3, 'Varna');
```

You can use existing records to create a new table – копира както структурата, така и данните

Пример 1

```
CREATE TABLE `customer_contacts` - new table name
AS SELECT `customer_id`, `first_name`, `email`, `phone`
FROM `customers`; - from existing table
```

Пример 2

```
CREATE TABLE `workers` AS
SELECT `first_name` FROM `employees`;
```

```
CREATE TABLE auto_filled AS
SELECT e.`first_name`,
d.`name` AS 'dept_name'
FROM `employees` AS e
INNER JOIN `departments` AS d
ON e.`department_id` = d.`department_id`;
```

You can write into an existing table - – копира както структура, така и данни

Пример 1

```
INSERT INTO `projects`(`name`, `start_date`)
SELECT
CONCAT(`name`,' ', 'Restructuring'),
NOW()
FROM `departments`;
```

Копира ги/добавя ги като данни за нови елементи от таблицата 😞

```
INSERT INTO `users`(`pk_users`)
SELECT
CONCAT(`id`, ",`username`) FROM `users`;
```

Пример 2

```
CREATE TABLE `workers`;
```

```
INSERT INTO `workers`
SELECT `first_name` FROM `employees` WHERE `salary` < 1000;
```

2.4. Updating Existing Records – UPDATE & DELETE

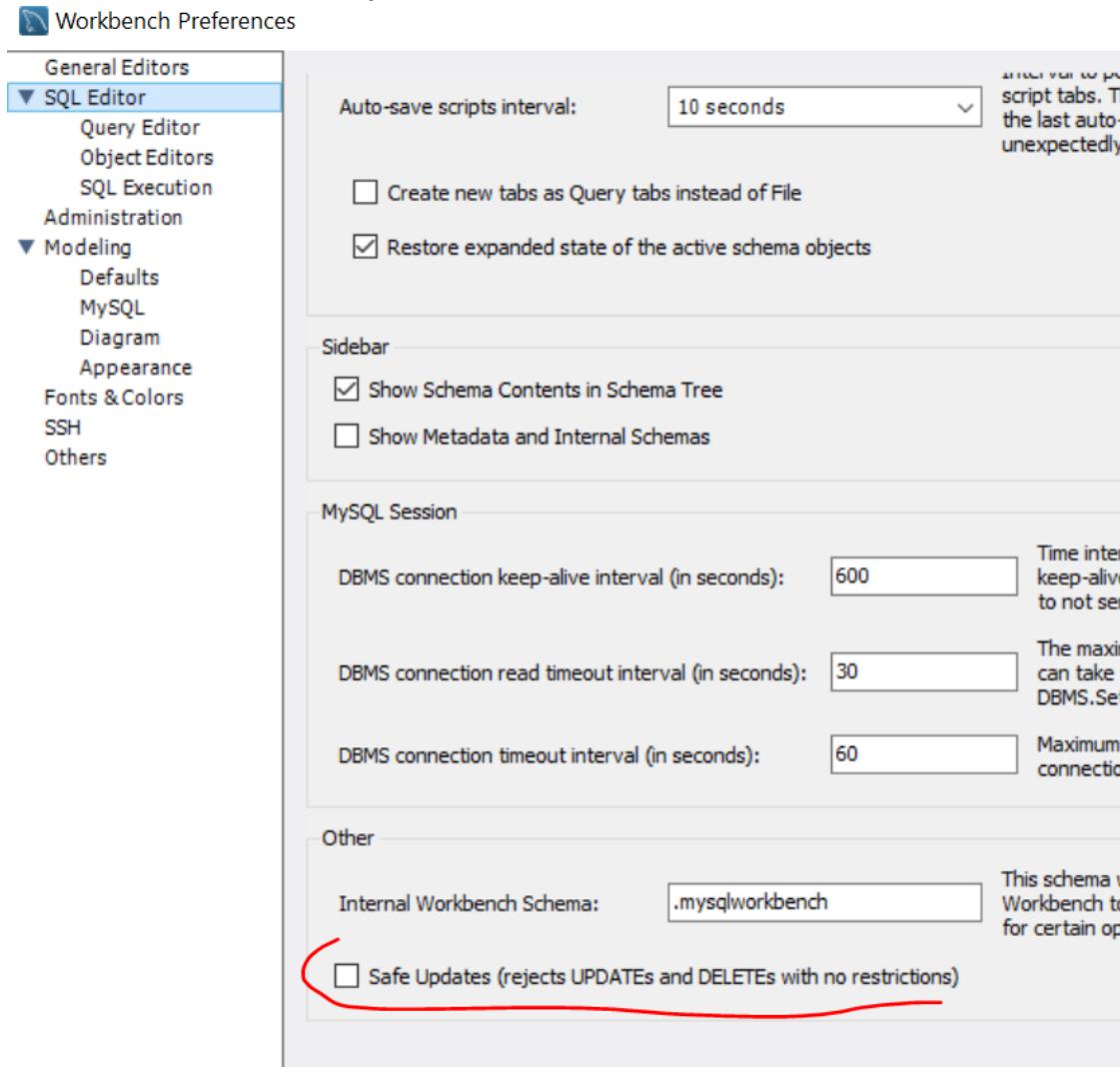
Updating data

Correcting/Updating records in SQL – using the GUI:

	id	first_name	last_name
▶	1	Svilen	Velikov
▶	2	Ivan	Ivanov
*	3	Tsvetomir	Velikov
	NULL	NULL	NULL

```
UPDATE `gamebar`.`employees` SET `last_name` = 'Ivanov' WHERE (`id` = '2');
```

Възможност в WHERE да не участва ключовото поле: Preferences -> SQL Editor



The SQL UPDATE command

```
UPDATE `employees`  
SET `last_name` = 'Brown'  
WHERE `employee_id` = 1;
```

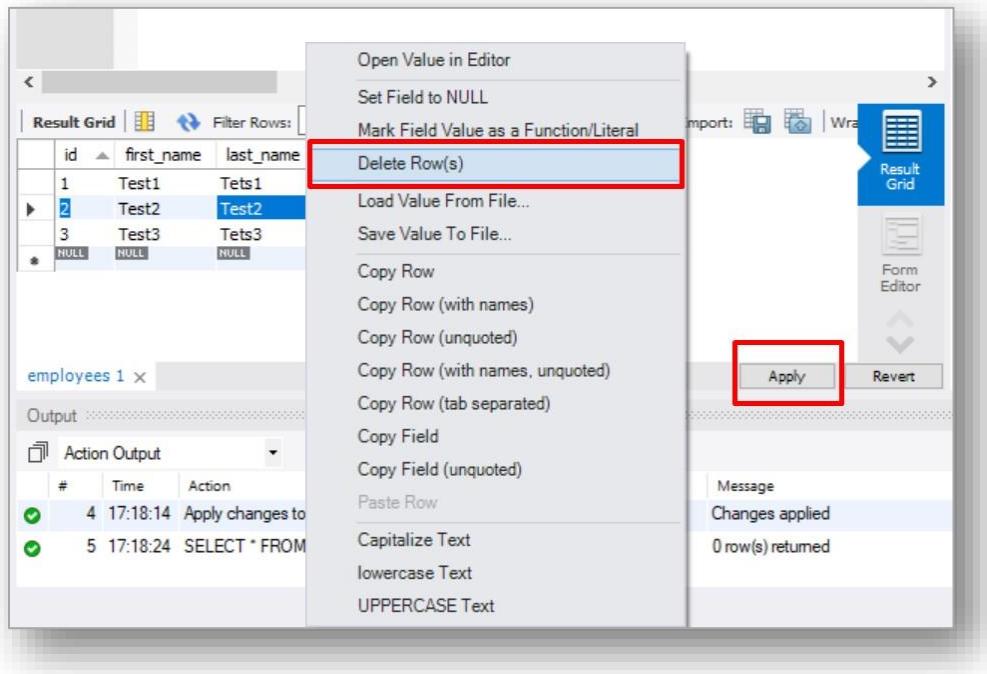
```
UPDATE `employees`  
SET `salary` = `salary` * 1.10,  
`job_title` = CONCAT('Senior', ' ', `job_title`)
```

```
WHERE `department_id` = 3;
```

- Note: Don't forget the **WHERE** clause!

Deleting Data

- Deleting row is part of DML



```
DELETE FROM `gamebar`.`employees` WHERE (`id` = '2');
```

- Deleting specific rows from a table
 - Note: Don't forget the **WHERE** clause!

```
DELETE FROM `employees`  
WHERE `employee_id` = 1;
```

- Delete all rows from a table (**TRUNCATE** works faster than **DELETE**)

```
TRUNCATE TABLE `users`;
```

3.Built-in functions

<https://dev.mysql.com/doc/refman/8.0/en/functions.html> - функции, има и .xml functions и .json functions

3.1. String functions

- **SUBSTRING()** – extracts part of a string

SUBSTRING(String, Position) – позицията/броенето започва от 1, а не от 0-левия

SUBSTRING(String, Position, Length)

SUBSTRING(String FROM Position FOR Length)

```
SELECT SUBSTRING('SoftUni', 2); - връща 'oftUni'
```

```
SELECT SUBSTRING('SoftUni', 2, 3); - връща 'oft'
```

- **REPLACE** – replaces specific string with another
 - Performs a case-sensitive match

REPLACE(String, Pattern, Replacement) Pattern - string to replace replacement – with what to replace

```
SELECT REPLACE(`title`, 'The', '***')
AS 'Title' FROM `books`
WHERE SUBSTRING(title, 1, 3) = 'The';
```

```
SELECT REPLACE(`title`, 'The', '***')
AS `title` FROM `books`
WHERE `title` LIKE 'The%'
ORDER BY `id` ASC;
```

```
SELECT `first_name`, `last_name` FROM `employees`
WHERE LOWER(`job_title`) NOT LIKE '%engineer%'
ORDER BY `employee_id`;
```

- Кастване/конвертиране от число към стринг

CAST (1 as CHAR)

- Chaining на функции – една функция в друга

- **LTRIM & RTRIM** – remove spaces from either side of string

LTRIM(String) – от началото на стринга

RTRIM(String) – от края на стринга

- **CHAR_LENGTH** – count number of characters

CHAR_LENGTH(String)

```
SELECT `name` FROM `towns`
WHERE CHAR_LENGTH(`name`) IN(5,6)
ORDER BY `name` ASC;
```

- **LENGTH** – get number of used bytes (double for Unicode)

LENGTH(String)

Кирилицата заема по 2 байта, а латиницата по един

```
SELECT LENGTH('асц'); - връща 6
SELECT LENGTH('бдт'); - връща 3
```

- **LEFT & RIGHT** – get characters from beginning or end of string

LEFT(String, Count) – от края

RIGHT(String, Count) – от началото

```
SELECT `id`, `start`,
LEFT(`name`, 3) AS 'Shorthand'
FROM `games`;
```

- **LOWER & UPPER** – change letter casing – we use it for case insensitive search

LOWER(String)

UPPER(String)

- **REVERSE** – reverse order of all characters in string

REVERSE(String)

- **REPEAT** – repeat string

REPEAT(String, Count)

- **LOCATE** – locate specific pattern (substring) in string

LOCATE(Pattern, String,[Position]) - If omitted(пропуснато), position begins at 1

SELECT LOCATE('Big', 'title') FROM `books`; - връща позицията на която се появява Big в полето `title`. **Ако не намери, връща 0**

SELECT LOCATE('@', 'chavdar.mitkov@softuni.bg'); - връща 15

SELECT LOCATE('@', 'chavdar.mitkov@softu@ni.bg', 16); - връща 22

SELECT `user_name`, SUBSTRING(`email`, LOCATE('@', `email`)+1) AS `email provider` FROM `users`

ORDER BY `email provider` ASC, `user_name` ASC;

- **INSERT** – insert substring at specific position

INSERT(StringToInsertInto, Position, Length, Substring) като Length е броят символи за унищожение

SELECT INSERT(`title`, 1, 0, 'Ordered book: ') FROM `books`; - **вмъкни на позиция 1, без да триеш нищо(0)**

'Ordedered book: ' пред всяко заглавие на книга

SELECT *,

INSERT (`title`, LOCATE('Big', `title`), 3, 'Small') AS `newtitle` - новото заглавие ако съдържа Big, то го подмени с 'Small' FROM `books`

WHERE `title` LIKE '%Big%'; - **да не започва или да не завършва с Big**

- **SUBSTRING_INDEX** – insert substring at specific position

SUBSTRING_INDEX(string, delimiter, number)

Parameter Description

string Required. The original string

delimiter Required. The delimiter to search for

number Required. The number of times to search for the delimiter. Can be both a positive or negative number. If it is a positive number, this function returns all to the left of the delimiter. If it is a negative number, this function returns all to the right of the delimiter.

SELECT

 user_name,
 SUBSTRING_INDEX(email, '@', -1) AS 'email provider'

FROM

 users;

3.2. Arithmetical Operators and Numeric Functions

Arithmetic expressions containing a null value are evaluated to null!!!

Arithmetical Operators

Name	Description
DIV	Integer division
/	Division operator
-	Minus Operator
%, MOD	Modulo operator
+	Addition operator
*	Multiplication operator
- (arg)	Change sign of argument

Numeric Functions

Used primarily for numeric **manipulation** and/or mathematical **calculations**

- **PI** – get the value of Pi (15 –digit precision)

SELECT PI() +0.000000000000000

- **ABS** – absolute value

ABS(Value)

- **SQRT** – square root

SQRT(Value)

- **POW** – raise value to desired exponent

POW(Value, Exponent)

- The **SUM()** function returns the total sum of a numeric column.

SELECT ROUND(SUM(`cost`), 2) AS `total_sum` FROM `books`;

Math Functions

- **CONV** – Converts numbers between different number bases

CONV(Value, from_base, to_base)

- **ROUND** – obtain desired precision

ROUND(Value, Precision) – Precision can be negative

SELECT ROUND(PI(), 2);

- **FLOOR & CEILING** – return the nearest integer

FLOOR(Value) - **надолу**

CEILING(Value) - **нагоре**

- **SIGN** – returns +1, -1 or 0, depending on value sign

SIGN(Value)

- **RAND** – get a random value in range [0,1)
 - If **Seed** is not specified, one is assigned at random – за хеширане, връща винаги една стойност за даден Seed

RAND() – връща random

RAND(Seed) - за хеширане, връща винаги една стойност за даден Seed

SELECT RAND('ssdfe'); винаги връща '0.15522042769493574'

3.3. Date Functions

- **EXTRACT** – extract a segment from a date as an integer

EXTRACT(Part FROM Date)

SELECT

EXTRACT(DAY FROM `born`) AS `day`,
`born` FROM `authors`;

SELECT EXTRACT(YEAR FROM '2022-05-23'); 2022

SELECT EXTRACT(month FROM '2022-05-23'); 5

SELECT EXTRACT(day FROM '2022-05-23'); 23

- **Get direct DAY, YEAR, MONTH, etc. without the function extract**

Директно си подаваме DAY, YEAR, MONTH от дадена дата

DAY(p.`date`) = 10

- **TIMESTAMPDIFF** – find difference between two dates

TIMESTAMPDIFF(Part, FirstDate, SecondDate)

- *Part* can be any part and format of date or time

year, %Y, %y

month, %M, %m

day, %w, %D

YEAR(Date)

MONTH(Date)

DAY(Date)

SELECT

timestampdiff(MONTH, `born`, `died`) AS `months_lived`,
`born` FROM `authors`;

- **DATE_FORMAT** – formats the date value according to the format

SELECT DATE_FORMAT('2017/05/31', '%Y %b %D') AS 'Date';

SELECT DATE_FORMAT('2017/05/31 23:13:00', '%Y %b %D, %h:%i:%s') AS 'Date';

2017 May 31st, 11:13:00

Хардкорнати стойности %b за месец и %i за минута

SELECT `first_name` FROM `employees`

WHERE `department_id` IN(3, 10) && DATE_FORMAT(`hire_date`, '%Y') BETWEEN 1995 AND 2005

ORDER BY `employee_id`;

Specifier	Description
%a	Abbreviated weekday name (Sun..Sat)
%b	Abbreviated month name (Jan..Dec)
%c	Month, numeric (0..12)
%D	Day of the month with English suffix (0th, 1st, 2nd, 3rd, ...)
%d	Day of the month, numeric (00..31)
%e	Day of the month, numeric (0..31)
%f	Microseconds (000000..999999)
%H	Hour (00..23)
%h	Hour (01..12)
%I	Hour (01..12)
%i	Minutes, numeric (00..59)
%j	Day of year (001..366)
%k	Hour (0..23)
%l	Hour (1..12)
%M	Month name (January..December)
%m	Month, numeric (00..12)
%p	AM or PM
%r	Time, 12-hour (<i>hh:mm:ss</i> followed by AM or PM)
%S	Seconds (00..59)
%s	Seconds (00..59)
%T	Time, 24-hour (<i>hh:mm:ss</i>)
%U	Week (00..53), where Sunday is the first day of the week; <code>WEEK()</code> mode 0
%u	Week (00..53), where Monday is the first day of the week; <code>WEEK()</code> mode 1
%V	Week (01..53), where Sunday is the first day of the week; <code>WEEK()</code> mode 2; used with %X
%v	Week (01..53), where Monday is the first day of the week; <code>WEEK()</code> mode 3; used with %x
%W	Weekday name (Sunday..Saturday)
%w	Day of the week (0=Sunday..6=Saturday)
%X	Year for the week where Sunday is the first day of the week, numeric, four digits; used with %v
%x	Year for the week, where Monday is the first day of the week, numeric, four digits; used with %v
%Y	Year, numeric, four digits
%y	Year, numeric (two digits)
%%	A literal % character
%x	x, for any “x” not listed above

- **NOW** – obtain current date and time

```
SELECT NOW();
```

- **DATE_ADD(`some_date`, INTERVAL stepValue typeStep)**

```
SELECT `product_name`, `order_date`,
DATE_ADD(`order_date`, INTERVAL 3 DAY) AS 'pay_due'
FROM `orders`;
```

MINUTE
HOUR
DAY
WEEK
MONTH
QUARTER
YEAR

- Сравнение по дата

```
SELECT `deposit_group`, `is_deposit_expired`,
AVG(`deposit_interest`) AS 'average_interest'
FROM `wizzard_deposits`
WHERE `deposit_start_date` > '1985-01-01'  когато формата на датата е този
GROUP BY `deposit_group`, `is_deposit_expired`
ORDER BY `deposit_group` DESC, `is_deposit_expired` ASC;
```

<https://dev.mysql.com/doc/refman/8.0/en/date-and-time-functions.html>

3.4. Wildcards

Системни команди / System commands – като сървър (а не като клиент)

```
USE INFORMATION_SCHEMA;
SELECT * from statistics;
SHOW tables;
```

Вземи information данни за дадена таблица

```
SELECT `COLUMN_NAME` FROM `information_schema`.`columns`
WHERE `TABLE_SCHEMA` = 'custom-orm' AND `COLUMN_NAME` != 'id' AND `TABLE_NAME` = 'users';
```

COLUMN_NAME
age
registration_date
username

Used to substitute any other character(s) in a string

- '%' - represents zero, one, or multiple characters
- '_' - represents a single character
- Can be used in combinations

- Used with **LIKE** operator in a **WHERE** clause
 - Similar to **Regular Expressions**

```
SELECT * FROM `books`
WHERE `title` LIKE '____Big%';
```

```
SELECT `user_name`, `ip_address` FROM `users`
WHERE `ip_address` LIKE '____.1%.____'  три символа.1няколко символа.няколко символа.три символа
ORDER BY `user_name` ASC;
```

- Find any values that start with "a"

```
WHERE CustomerName LIKE 'a%';
```

- Find any values that have "r" in second position

```
WHERE CustomerName LIKE '_r%';
```

- Finds any values that starts with "a" and ends with "o"

```
WHERE ContactName LIKE 'a%o';
```

- Supported characters also include – **част от regex нещата искат да опишат тук**
 - \ – specify prefix to treat special characters as normal – да го escape-нем
 - [charlist] – specifying which characters to look for
 - [!charlist] – **excluding** characters

```
SELECT * FROM `books`
WHERE `title` LIKE '____Big\%'; - обикновен процент
```

```
SELECT `first_name`, `last_name` FROM `employees`
WHERE LOWER(`job_title`) NOT LIKE '%engineer%'
ORDER BY `employee_id`;
```

3.5. Regex

```
SELECT * FROM `customers`
WHERE `city` REGEXP '[a-c]%' ; - a, b or c - връща 0 или 1 ца – дали има или няма match
```

```
SELECT * FROM `customers`
WHERE `city` NOT REGEXP '[a-c]%' ; - да не съдържа a, b or c
```

Using regular expression

- **REGEXP** - pattern matching using regular expressions

```
SELECT `employee_id`, `first_name`, `last_name`
FROM `employees`
WHERE `first_name` REGEXP '^[^K]{3}\$'; връща 0 или 1 ца – дали има или няма match
```

Пример:

```
DELIMITER $$$
```

```
CREATE FUNCTION ufn_is_word_comprised(setOfLetters VARCHAR(45), word VARCHAR(45))
RETURNS BIT //връща нула или единица, както и при BOOLEAN
DETERMINISTIC
```

```

BEGIN
    RETURN word REGEXP(concat('^[' , setOfLetters , ']+')); //поне един път трябва да има всяка
END;
$$$$

```

```

SELECT ufn_is_word_comprised('oistmiahf', 'Sofia');
SELECT ufn_is_word_comprised('oistmiahf', 'halves');

```

- **REGEXP_SUBSTR**(*expr*, *pat*[, *pos*[, *occurrence*[, *match_type*]])

Returns the substring of the string *expr* that matches the regular expression specified by the pattern *pat*, NULL if there is no match. If *expr* or *pat* is NULL, the return value is NULL.

REGEXP_SUBSTR() takes these optional arguments:

- *pos*: The position in *expr* at which to start the search. If omitted, the default is 1.
- *occurrence*: Which occurrence of a match to search for. If omitted, the default is 1.
- *match_type*: A string that specifies how to perform matching. The meaning is as described for **REGEXP_LIKE**().

```
SELECT REGEXP_SUBSTR(`title` , '[a-zA-Z]{2} ') AS `match` , `title` FROM `books`;
```

```
SELECT REGEXP_SUBSTR(`title` , '[a-zA-Z]+') AS `match` , `title` FROM `books`;
```

- **REGEXP_REPLACE**(*expr*, *pat*, *rep1*[, *pos*[, *occurrence*[, *match_type*]])

Replaces occurrences in the string *expr* that match the regular expression specified by the pattern *pat* with the replacement string *rep1*, and returns the resulting string. If *expr*, *pat*, or *rep1* is NULL, the return value is NULL.

```

SELECT
    user_name,
    REGEXP_REPLACE(email, '.*@', '') AS 'email provider'
FROM
    users;

```

3.6. Условни конструкции

*Използване на IFNULL функцията - Return the specified value IF the expression is NULL, otherwise return the expression:
IFNULL('middle_name', '')* - ако е NULL, то го замести с празен стринг, иначе върни полето

IF condition – 1 – as a function – връща резултат

```

IF(condition, value_if_true, value_if_false)
SELECT IF(1 > 0, 'true', 'false'); --returns 'true'

```

```

SELECT
    `name` AS `game`,
    /*DATE_FORMAT(start, '%k') AS `P`,*/
    IF(DATE_FORMAT(start, '%k') >= 0 && DATE_FORMAT(start, '%k') < 12, 'Morning',
        IF(DATE_FORMAT(start, '%k') >= 12 && DATE_FORMAT(start, '%k') < 18, 'Afternoon', 'Evening')) AS `Part of the
    Day`,

```

```

IF(`duration` <= 3 , 'Extra Short',
    IF(`duration` <= 6, 'Short',
        IF(`duration` <= 10, 'Long', 'Extra Long'))) AS `Duration`
FROM `games`;

```

IF condition – 2 – as a statement – не връща стойност

```

IF search_condition THEN statement_list;
[ELSEIF search_condition THEN statement_list] ...;
[ELSE statement_list];
END IF

```

DELIMITER %%

```

CREATE PROCEDURE usp_raise_salary_by_id(id int)
BEGIN
    START TRANSACTION;
    IF((SELECT count(employee_id) FROM employees WHERE employee_id like id)<>1)
    THEN ROLLBACK;
    ELSE
        UPDATE employees AS e SET salary = salary + salary*0.05
        WHERE e.employee_id = id;
    END IF;
END %%

```

DECLARE result DECIMAL;

```

IF(salary_emp < 30000) THEN SET result := 'Low'; //при DECLARE използваме := за присвояване
ELSEIF (salary_emp <= 50000) THEN SET result := 'Average'; ELSEIF слято трябва да е
ELSE SET result := 'High';
END IF;

```

CASE condition - as a function – връща резултат

Пример 1:

```

SELECT
CASE `author_id`
    WHEN 1 THEN 'Recommended for beginners' //Ако 1, върни еди какво си
    WHEN 7 THEN 'Recommended for advanced'
    ELSE 'All audiences'
END
AS `my_preference`
FROM `books`
WHERE SUBSTRING(title, 1, 3) = 'The';

```

Пример 2:

```

SELECT `name` AS 'game',
(CASE /*без променлива тук може*/
    WHEN HOUR(`start`) BETWEEN 0 AND 11 THEN 'Morning'
    WHEN HOUR(`start`) BETWEEN 12 AND 17 THEN 'Afternoon'
    ELSE 'Evening'
END) AS 'Part of the Day',
(CASE /*без променлива тук може*/

```

```

WHEN `duration` BETWEEN 0 AND 3 THEN 'Extra Short'
WHEN `duration` BETWEEN 4 AND 6 THEN 'Short'
WHEN `duration` BETWEEN 7 AND 10 THEN 'Long'
ELSE 'Extra long'
END) AS 'Duration'
FROM `games`;

```

Пример 3:

```

SELECT
CASE
    WHEN `age` BETWEEN 0 AND 10 THEN '[0-10]'
    WHEN `age` BETWEEN 11 AND 20 THEN '[11-20]'
    WHEN `age` BETWEEN 21 AND 30 THEN '[21-30]'
    WHEN `age` BETWEEN 31 AND 40 THEN '[31-40]'
    WHEN `age` BETWEEN 41 AND 50 THEN '[41-50]'
    WHEN `age` BETWEEN 51 AND 60 THEN '[51-60]'
    WHEN `age` >= 61 THEN '[61+]'
    ELSE 'OUT of RANGE'
END AS `age_group`,
COUNT(`id`) AS `wizard_count`
FROM `wizzard_deposits`
GROUP BY `age_group`
ORDER BY `age_group`;

```

3.7. Цикли - WHILE statement

Пример с функция, където използваме **WHILE statement**

```

DELIMITER %%
CREATE FUNCTION ufn_IsWordComprised(setOfLetters VARCHAR (50), word VARCHAR (50))
RETURNS INT
deterministic
BEGIN
    DECLARE index_letter INT;
    DECLARE length_word INT;
    DECLARE letter CHAR(1);

    SET index_letter := 1;
    SET length_word := CHAR_LENGTH(word);

    WHILE (index_letter <= length_word)
    DO
        SET letter := SUBSTRING(word, index_letter, 1);

        IF (LOCATE(letter, setOfLetters) > 0) THEN SET index_letter := index_letter + 1;
        ELSE
            RETURN 0;
        END IF;
    END WHILE;

    RETURN 1;
END;
%%
```

```
SELECT ufn_IsWordComprised('oistmiahf', 'Sofia');
```

3.8. Цикли - LOOP statement

LOOP

```
...  
-- terminate the loop  
IF condition THEN  
    LEAVE [label];  
END IF;  
...  
END LOOP;
```

4. Data Aggregation

4.1. Grouping - Consolidating Data Based On Criteria

- Grouping allows taking data into **separate groups** based on a **common property**
- Aggregating data allows us to group rows based on certain criteria.
- SQL provides a mechanism for aggregating data using the GROUP BY clause in a SELECT statement.
- SQL provides a standard set of functions that can be used over aggregated data

When we want to select and group by column or columns:

- Typically used with aggregate functions
- Can group by multiple columns. The result set will be grouped by the first column listed, then by the next.
- NULL values will be grouped together

Прави като представителна извадка - средната заплата на служителите в даден отдел.

```
select departmentId, avg(salary) as avgSalary  
from Employees as e  
group by departmentId;
```

```
SELECT e.job_title, count(employee_id)  
FROM `employees` AS e  
GROUP BY e.job_title;
```

Common property

Grouping column

employee	department_name	salary
Adam	Database Support	5,000
John	Database Support	15,000
Jane	Application Support	10,000
George	Application Support	15,000
Lila	Application Support	5,000
Fred	Software Support	15,000

Can be aggregated



- Групиране по 2 критерия/колони

```

SELECT
`deposit_group`,
`magic_wand_creator`,
MIN(`deposit_charge`) AS 'min_deposit_charge'
FROM
`wizzard_deposits`
GROUP BY `deposit_group`, `magic_wand_creator`
ORDER BY `magic_wand_creator` ASC , `deposit_group`;

```

- With **GROUP BY** you can get each separate group and use an "**aggregate**" function over it (like Average, Min or Max)

4.2. Aggregate Functions

- Used to operate over **one or more** groups performing **data analysis** on every one

Aggregating functions are:

COUNT(*) - count of the selected rows

SUM(column) - sum of the values in given column from the selected rows – in PostgreSQL does not allow on varchar type for example. On MySQL return 0 on varchar!

AVG(column) - average of the values in given column

MAX(column) - the maximal value in given column

MIN(column) - the minimal value in given column

- They usually **ignore NULL** values

COUNT - counts the values (not nulls) in one or more columns based on grouping criteria

- Note that when we use **COUNT** we will ignore any employee with **NULL** salary.

```

SELECT `department_id`, COUNT(`first_name`) AS 'Number of employees'
FROM `employees`
GROUP BY `department_id`
ORDER BY `department_id` ASC, `Number of employees` ASC;

```

Използваме звезда за по-мързеливо и за да ни брои всички елементи

```

SELECT `department_id`, COUNT(*) AS 'Number of employees'
FROM `employees`

```

SUM - sums the values in a column

- If any department has no salaries **NULL** will be displayed.

```
SELECT e.`department_id`,  
SUM(e.`salary`) AS 'TotalSalary'  
FROM `employees` AS e  
GROUP BY e.`department_id`;
```

MAX/MIN - takes the maximum value in a column.

```
SELECT e.`department_id`,  
MAX(e.`salary`) AS 'Max Salary'  
FROM `employees` AS e  
GROUP BY e.`department_id`;
```

AVG calculates the average value in a column.

```
SELECT e.`department_id`,  
ROUND(AVG(e.`salary`),2) AS 'Average Salary'  
FROM `employees` AS e  
GROUP BY e.`department_id`  
ORDER BY e.`department_id`;
```

Използване на AVG в секцията ORDER BY

```
SELECT w.`deposit_group`  
FROM `wizzard_deposits` AS w  
GROUP BY w.`deposit_group`  
ORDER BY AVG(w.`magic_wand_size`)  
LIMIT 1;
```

4.3. HAVING - Using Predicates While Grouping

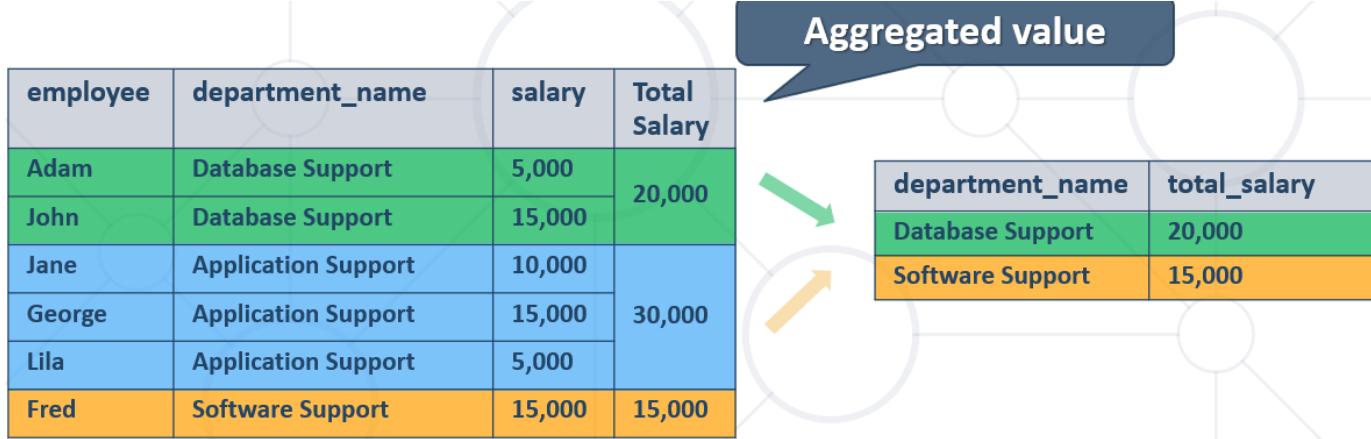
Having Clause

- The **HAVING** clause is used to filter data based on **aggregate** values.
 - We cannot use it **without** grouping **before** that
- Any Aggregate functions in the "**HAVING**" clause and in the "**SELECT**" statement are executed one time only
- **Unlike HAVING, the WHERE clause filters rows before the aggregation**

```
SELECT `deposit_group`,  
SUM(`deposit_amount`) AS `total_sum`  
FROM `wizzard_deposits`  
WHERE `magic_wand_creator` = 'Ollivander family'  
GROUP BY `deposit_group`  
HAVING `total_sum` < 150000  
ORDER BY `total_sum` DESC;
```

```
SELECT `deposit_group`,  
*  
FROM `wizzard_deposits`  
WHERE `magic_wand_creator` = 'Ollivander family'  
GROUP BY `deposit_group`  
HAVING SUM(`deposit_amount`) < 150000  
ORDER BY `total_sum` DESC;
```

Filter departments which have **total** salary **less than** 25,000.



```
SELECT `department_id`,
       SUM(`salary`) AS `TotalSalaryOfDepartment`
  FROM `employees`
 GROUP BY `department_id`
 HAVING `TotalSalaryOfDepartment` < 120000;
```

4.4. MySQL **OFFSET** and **LIMIT** is used to specify which row should be fetched first.

```
SELECT e.`department_id`,
       e.`salary` AS `third_highest_salary`
  FROM `employees` AS e
 WHERE (SELECT ine.`employee_id` FROM `employees` AS ine
       WHERE ine.`department_id` = e.`department_id`           връща employee_id когато
       GROUP BY ine.`salary`                                департамента съвпада
       ORDER BY `salary` DESC LIMIT 1 OFFSET 2 – изкарай само един резултат, започвайки да търсиш след втория
) = e.`employee_id`      когато има съвпадение по employee_id
 GROUP BY e.`department_id`
 ORDER BY e.`department_id` ASC;
```

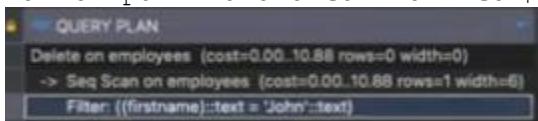
LIMIT 2, 1; - започни да търсиш след втория запис, и ограничи до 1 запис изхода

4.5. Debug – EXPLAIN ANALYZE SELECT ..

```
EXPLAIN SELECT *, SUBSTRING(`title`, 1, 4) FROM `books` LIMIT 20 OFFSET 11;
OFFSET Отмети/започни от 11тия запис нататък
```

explain

прави план на заявката без да я изпълнява – има го и в MySQL – като не сме сигурни какво прави нашата заявка и за да не стане някаква беля



explain analyze

```

QUERY PLAN
Hash Join (cost=11.57..3207.91 rows=88888 width=545) (actual time=0.126..31.709 rows=88889 loops=1)
  Hash Cond: (employees.departmentid = department.departmentid)
    -> Seq Scan on employees (cost=0.00..2948.88 rows=88888 width=33) (actual time=0.032..9.969 rows=88891)
    -> Hash (cost=10.70..10.70 rows=70 width=520) (actual time=0.036..0.037 rows=10 loops=1)
      Buckets: 1024 Batches: 1 Memory Usage: 9kB
        -> Seq Scan on department (cost=0.00..10.70 rows=70 width=520) (actual time=0.020..0.023 rows=10 loops=1)
Planning Time: 0.471 ms
Execution Time: 35.508 ms

```

4.6. Вложени агрегиращи заявки

```

SELECT e.`first_name`, e.`last_name`, e.`department_id`
FROM `employees` AS e
WHERE e.`salary` > (
  SELECT
    AVG(inn.`salary`) FROM `employees` AS inn      - намери средната заплата
    WHERE inn.`department_id` = e.`department_id`   - ако средната заплата след групиране отговаря на запла-
    GROUP BY inn.`department_id`                   - тата на всеки пореден служител от съответв. департамент
)
ORDER BY e.`department_id`, e.`employee_id`
LIMIT 10;

```

Друго решение на същата задача – използваме сега само един Alias:

```

SELECT `first_name`, `last_name`, `department_id`
FROM `employees` AS e
WHERE e.salary > (SELECT AVG(salary) FROM employees WHERE department_id = e.department_id GROUP BY department_id)
ORDER BY `department_id`, `employee_id`
LIMIT 10

```

ВАЖНО – при update и използване на нестнати операции

```

UPDATE employees_clients AS ec
SET ec.employee_id =
(
  SELECT ec.employee_id – от същата таблица ес реално
  GROUP BY ec.employee_id
  ORDER BY COUNT(ec.employee_id) ASC, ec.employee_id ASC
  LIMIT 1
)
WHERE ec.employee_id = ec.client_id;

```

4.7. Невложени заявки вършещи работа като вложени заявки

```

SELECT
  SUM(`hw`.`deposit_amount` - `gw`.`deposit_amount`) AS 'sum_difference'
FROM
  `wizzard_deposits` AS `hw`,
  `wizzard_deposits` AS `gw`
WHERE
  `gw`.`id` - `hw`.`id` = 1;

```

5. Table relations – видове връзки

5.1. Database design

Steps in database design



1. Identification of Entities

- Entity tables represent objects from the real world
 - Most often they are nouns in the specification

For example:

We need to develop a system that stores information about **students**, which are trained in various **courses**. The courses are held in different **towns**. When registering a new student the following information is entered: name, faculty number, photo and date.

Entities: **Student, Course, Town**

2. Define Table Columns

- Columns are clarifications for the entities in the text of the specification, for example:

We need to develop a system that stores information about **students**, which are trained in various **courses**. The courses are held in different **towns**. When registering a new student the following information is entered: **name**, **faculty number**, **photo** and **date**.

- Students have the following characteristics:

- Name, faculty number, photo, date of enlistment and a list of courses they visit

3. Defining primary keys

- Always define an additional column for the primary key
 - Don't use an existing column (for example SSN)
 - Can be an integer number
 - **Must be** declared as a **PRIMARY KEY**
 - Use **AUTO_INCREMENT** to implement auto-increment
 - **Put the primary key as a first column**
- Exceptions
 - Entities that have well known ID, e.g. countries (BG, DE, US) and currencies (USD, EUR, BGN)

4. Modelling relationships

- Relationships are dependencies between the entities:

We need to develop a system that stores information about **students**, which are trained in various courses. The **courses** are held in different **towns**. When registering a new student the following information is entered: name, faculty number, photo and date.

- "Students are trained in courses" – **many-to-many** relationship.
- "Courses are held in towns" – **many-to-one** (or many-to-many) relationship

5. Defining constraints

6. Filling test data

5.2. Table relations

- **Relationships** between tables are based on interconnections: **PRIMARY KEY / FOREIGN KEY**

Foreign Key

```
CREATE TABLE `Orders` (
  `OrderID` int NOT NULL,
  `OrderNumber` int NOT NULL,
  `PersonID` int,
  PRIMARY KEY (`OrderID`),

CONSTRAINT `fk_source_target`
  FOREIGN KEY (`Orders`(`PersonID`)) REFERENCES `Persons`(`PersonID`)
);
```

SQL FOREIGN KEY on ALTER TABLE
`ALTER TABLE `products``

```
ADD CONSTRAINT `fk_products_categories` FOREIGN KEY `products`(`category_id`)
REFERENCES `categories`(`id`);
```

- The **foreign key** is an **identifier** of a record located in another table (usually its primary key)
- By using relationships we avoid repeating data in the database
- Relationships have multiplicity:

One-to-many – e.g. mountains / peaks - от едната страна на foreign key полето трябва да е UNIQUE

Много модели на един производител в случая

```
CREATE TABLE `manufacturers`(
`manufacturer_id` INT AUTO_INCREMENT PRIMARY KEY NOT NULL, - реално е UNIQUE това поле
`name` VARCHAR(45),
`established_on` DATE
);
```

Foreign key е от Many страната на релацията.

```
CREATE TABLE `models`(
`model_id` INT UNIQUE NOT NULL,
`name` VARCHAR(45),
`manufacturer_id` INT NOT NULL - а това поле не е UNIQUE
);
```

```
ALTER TABLE `models`
ADD CONSTRAINT `fk_models_manufacturers`
FOREIGN KEY `models`(`manufacturer_id`)
REFERENCES `manufacturers`(`manufacturer_id`);
```

```
INSERT INTO `manufacturers`(`name`, `established_on`)
VALUES
('BMW', '1916-03-01'),
('Tesla', '2003-01-01'),
('Lada', '1966-05-01');
```

```
INSERT INTO `models`(`model_id`, `name`, `manufacturer_id`)
VALUES
(101, 'X1', 1),
(102, 'i6', 1),
(103, 'Model S', 2),
(104, 'Model X', 2),
(105, 'Model 3', 2),
(106, 'Nova', 3);
```

One-to-one – e.g. example driver / car – и от двете страни на foreign key полето трябва да е UNIQUE

```
CREATE TABLE `people`(
`person_id` INT AUTO_INCREMENT PRIMARY KEY NOT NULL,
`first_name` VARCHAR(45),
`salary` DECIMAL(10, 2) NOT NULL,
`passport_id` INT UNIQUE NOT NULL - това поле е UNIQUE
```

```
);
```

```
CREATE TABLE `passports`(  
    `passport_id` INT UNIQUE NOT NULL, - и това поле също е UNIQUE  
    `passport_number` VARCHAR(45) UNIQUE  
);
```

```
ALTER TABLE `people`  
ADD CONSTRAINT `fk_people_passports`  
FOREIGN KEY `people`(`passport_id`)  
REFERENCES `passports`(`passport_id`);
```

```
INSERT INTO `people` (`first_name`, `salary`, `passport_id`)  
VALUES  
('Roberto', 43300.00, 102),  
('Tom', 56100.00, 103),  
('Yana', 60200.00, 101);
```

```
INSERT INTO `passports` (`passport_id`, `passport_number`)  
VALUES  
(101, 'N34FG21B'),  
(102, 'K65LO4R7'),  
(103, 'ZE657QP2');
```

Всяко поле AUTO_INCREMENT и всяко поле PRIMERY KEY реално е UNIQUE

Many-to-many – e.g. student / course – изпълнява се с composite primary key и mapping table

Mapping table and composite primary key - пример за many-to-many relations

```
CREATE TABLE `students`(  
    `student_id` INT AUTO_INCREMENT PRIMARY KEY NOT NULL,  
    `name` VARCHAR(45)  
);
```

```
CREATE TABLE `exams` (  
    `exam_id` INT NOT NULL PRIMARY KEY,  
    `name` VARCHAR(30)  
);
```

Важно: за да работи foreign key, трябва да има primary key в таблиците

```
CREATE TABLE `students_exams`(  
    `student_id` INT,  
    `exam_id` INT,  
  
CONSTRAINT pk_students_exams  
PRIMARY KEY `students_exams`(`student_id`, `exam_id`),  
  
CONSTRAINT fk_students_exams_students  
FOREIGN KEY `students_exams`(`student_id`)  
REFERENCES `students`(`student_id`),
```

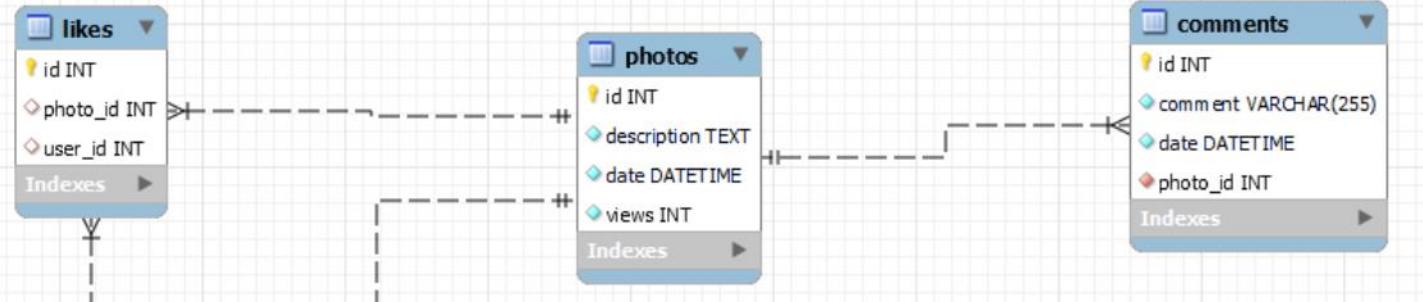
```

CONSTRAINT fk_students_exams_exams
FOREIGN KEY `students_exams`(`exam_id`)
REFERENCES `exams`(`exam_id`)
);

```

Без композитен ключ, използваме DISTINCT за да няма повторения

#08. Count Likes and Comments



В случая, като съединим 3 таблици, и за всеки различен лайк на снимка има примерно по два коментара.

Как процедирате – пускате първо заявката със звезда, и виждате какво става.

```

SELECT *
FROM `photos` AS p
LEFT JOIN `likes` AS l
ON l.`photo_id` = p.`id`
LEFT JOIN `comments` AS c
ON c.`photo_id` = p.`id`;

```

	id	description	date	views	id	photo_id	user_id	id	comment	date	photo_id
▶	1	Nullam sit amet turpis el...	2019-09-22 14:48:03	0	46	1	46	19	amet erat nulla ...	2019-06-... 1	
	1	Nullam sit amet turpis el...	2019-09-22 14:48:03	0	46	1	46	94	lacus purus aliqu...	2019-05-... 1	
	1	Nullam sit amet turpis el...	2019-09-22 14:48:03	0	52	1	10	19	amet erat nulla ...	2019-06-... 1	
	1	Nullam sit amet turpis el...	2019-09-22 14:48:03	0	52	1	10	94	lacus purus aliqu...	2019-05-... 1	
	1	Nullam sit amet turpis el...	2019-09-22 14:48:03	0	57	1	18	19	amet erat nulla ...	2019-06-... 1	
	1	Nullam sit amet turpis el...	2019-09-22 14:48:03	0	57	1	18	94	lacus purus aliqu...	2019-05-... 1	

```

SELECT p.`id`, COUNT(DISTINCT l.`id`) AS `likes_count`, COUNT(DISTINCT c.`id`) AS `comments_count`
FROM `photos` AS p
LEFT JOIN `likes` AS l
ON l.`photo_id` = p.`id`
LEFT JOIN `comments` AS c
ON c.`photo_id` = p.`id`
GROUP BY p.`id`
ORDER BY `likes_count` DESC, `comments_count` DESC, p.`id` ASC;

```

Self-referencing

```

CREATE TABLE `teachers`(
`teacher_id` INT PRIMARY KEY,
`name` VARCHAR(20),
`manager_id` INT,

```

```

CONSTRAINT fk_teachers_teachers
FOREIGN KEY `teachers`(`manager_id`)
REFERENCES `teachers`(`teacher_id`)

```

```
);
```

```
INSERT INTO `teachers` (`teacher_id`, `name`)
VALUES
(101, 'John'),
(102, 'Maya'),
(103, 'Silvia'),
(104, 'Ted'),
(105, 'Mark'),
(106, 'Greta');
```

```
UPDATE `teachers` SET `manager_id` = null WHERE (`name` = 'John');
UPDATE `teachers` SET `manager_id` = 106 WHERE (`name` = 'Maya');
UPDATE `teachers` SET `manager_id` = 106 WHERE (`name` = 'Silvia');
UPDATE `teachers` SET `manager_id` = 105 WHERE (`name` = 'Ted');
UPDATE `teachers` SET `manager_id` = 101 WHERE (`name` = 'Mark');
UPDATE `teachers` SET `manager_id` = 101 WHERE (`name` = 'Greta');
```

Друг пример за self-referencing

```
SELECT e.`employee_id`, e.`first_name`, m.`employee_id` AS 'manager_id', m.`first_name` AS 'manager_name'
FROM `employees` AS e
JOIN `employees` AS m
ON e.`manager_id` = m.`employee_id`;
```

5.3. JOIN - Retrieving Related Data

Joins

- Table relations are useful when combined with JOINS
- With JOINS we can get data from two tables **simultaneously**
 - JOINS require at least two tables and a "**join condition**"

Example:

```
SELECT * FROM table_a
JOIN table_b ON - добави таблица b към таблица a
table_b.common_column = table_a.common_column; да са равни
```

За използването на JOIN не е необходимо използването на външен ключ (foreign key)

```
SELECT v.`driver_id`, v.`vehicle_type`,
CONCAT(c.`first_name`, ' ', c.`last_name`) AS 'driver_name'
FROM `vehicles` AS v
JOIN `campers` AS c - добави таблица campers към таблица vehicles
ON v.`driver_id` = c.`id`;
```

5.4. Cascade Operations

- Cascading allows when a change is made to certain entity, this change to apply to all related entities

Ако изтриеш нещо, изтрий всичко по веригата.

Ако update-неш нещо, то го update-ни навсякъде по веригата.

CASCADE DELETE

- **CASCADE** can be either **DELETE** or **UPDATE**.
- Use **CASCADE DELETE** when:
 - The related entities are **meaningless** without the "main" one
- Do **not** use **CASCADE DELETE** when:
 - You make "**logical delete**"
 - You preserve **history**
- Keep in mind that in more complicated relations it won't work with **circular** references

Пример 1:

- Write a query to create a one-to-many relationship
- When a mountain gets removed from the database, all of his peaks are deleted too

```
CREATE TABLE `mountains`(  
    `id` INT PRIMARY KEY AUTO_INCREMENT,  
    `name` VARCHAR(20) NOT NULL  
);
```

```
CREATE TABLE `peaks`(  
    `id` INT PRIMARY KEY AUTO_INCREMENT,  
    `name` VARCHAR(20) NOT NULL,  
    `mountain_id` INT,  
    CONSTRAINT `fk_mountain_id`  
        FOREIGN KEY(`mountain_id`)  
        REFERENCES `mountains`(`id`)  
    ON DELETE CASCADE  
);
```

Пример 2:

```
CREATE TABLE drivers(  
    driver_id INT PRIMARY KEY,  
    driver_name VARCHAR(50)  
);  
  
CREATE TABLE cars(  
    car_id INT PRIMARY KEY,  
    driver_id INT,  
    CONSTRAINT fk_car_driver FOREIGN KEY(driver_id)  
    REFERENCES drivers(driver_id) ON DELETE CASCADE  
);
```

CASCADE UPDATE

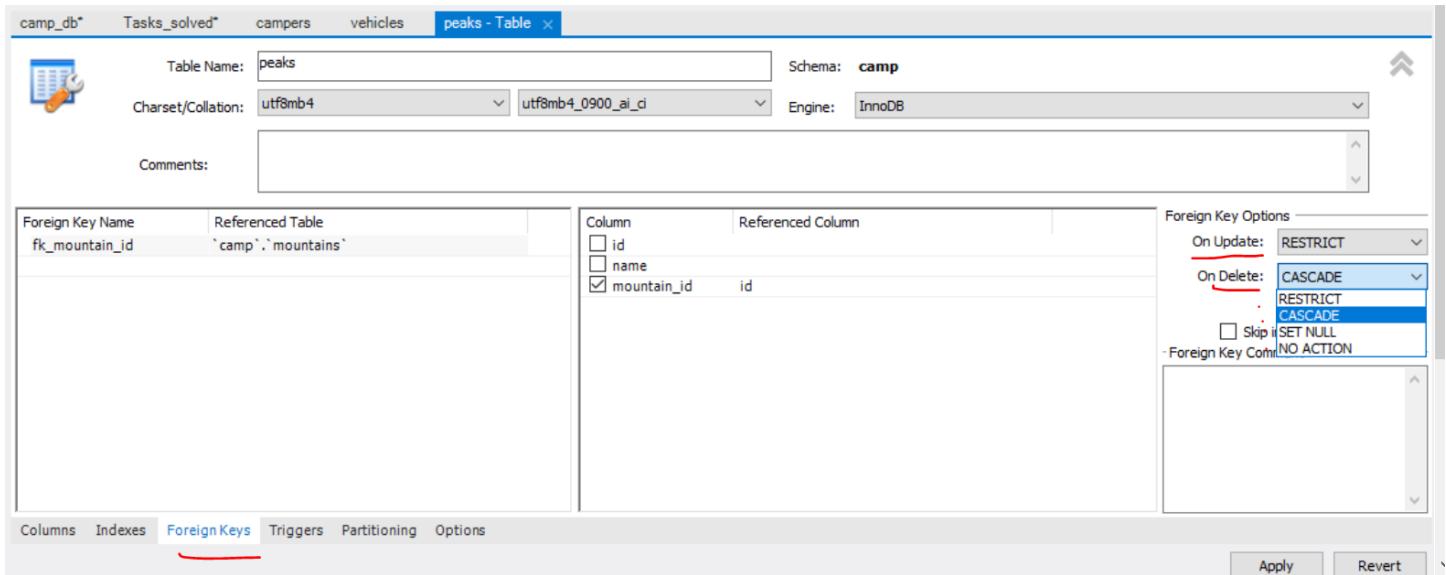
- Use **CASCADE UPDATE** when:
 - The primary key is **NOT** identity (**not auto-increment**) and therefore it **can** be changed
 - Best used with **UNIQUE** constraint
- Do **not** use **CASCADE UPDATE** when:
 - The primary is identity (**auto-increment**)
 - Cascading can be avoided using triggers or procedures

```
CREATE TABLE drivers(  
    driver_id INT PRIMARY KEY,  
    driver_name VARCHAR(50)
```

```

);
CREATE TABLE cars(
    car_id INT PRIMARY KEY,
    driver_id INT,
    CONSTRAINT fk_car_driver FOREIGN KEY(driver_id)
    REFERENCES drivers(driver_id) ON UPDATE CASCADE
);

```



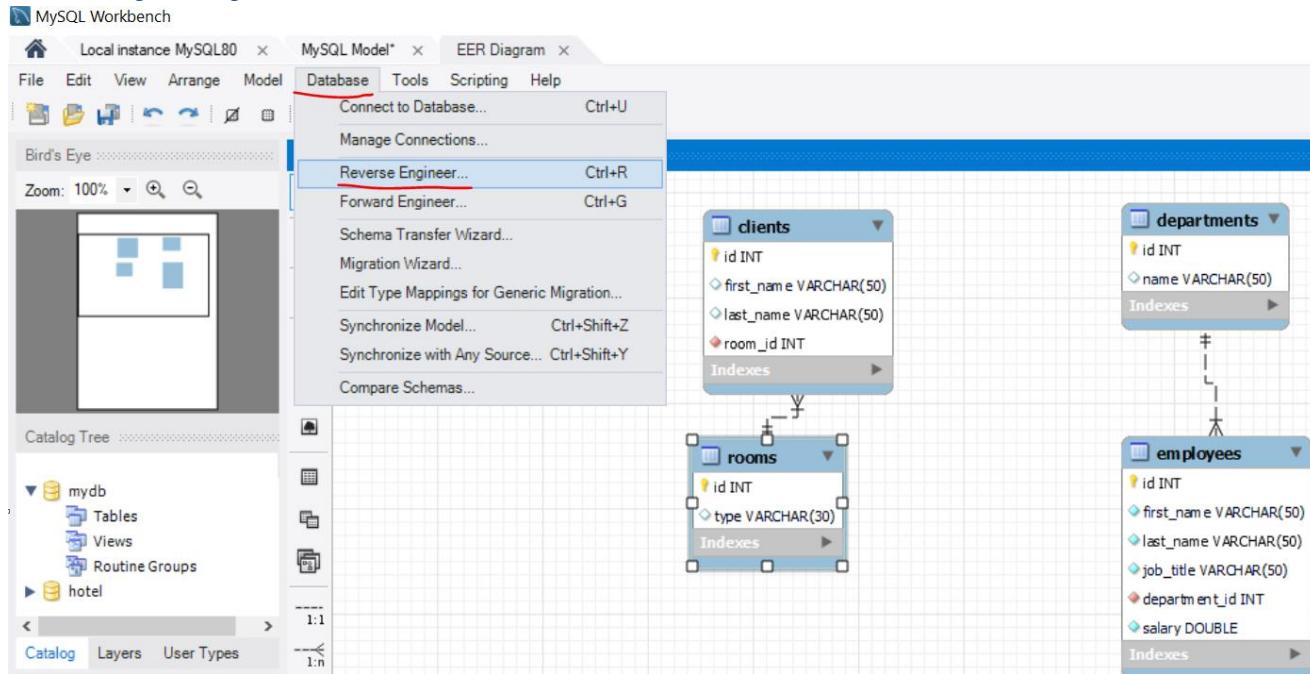
5.5. E/R Diagram

Понякога първо си чертаем Е/Р диаграмата, и след това пишем заявките

Relational Schema

- **Relational schema** of a DB is the collection of:
 - The schemas of all tables
 - Relationships between the tables
 - Any other database objects (e.g. constraints)
- The relational schema describes the **structure** of the database
 - Doesn't contain data, but **metadata**
- Relational schemas are **graphically** displayed in Entity / Relationship diagrams (**E/R Diagrams**)

Reverse engineering



Добра практика е да слагаме префикс на името на всяка таблица, която създаваме. Например:

my_Products

my_Clients

my_personnel

или

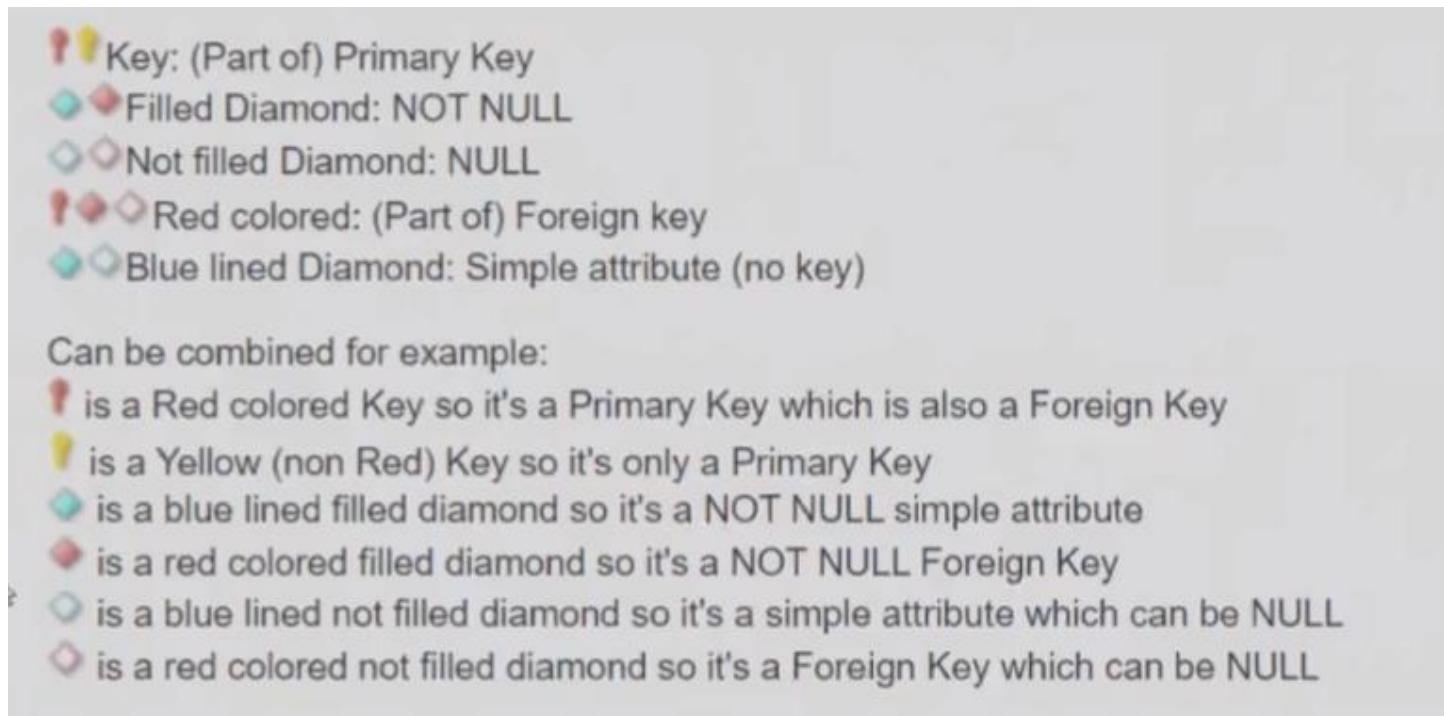
wp_posts

wp_links

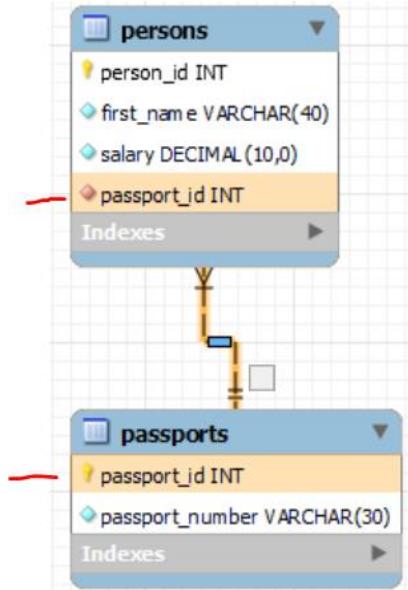
wp_commentmeta

wp_comments

wp_users



В persons таблицата пишем дефиницията за foreign key



Като посочим стрелката, и ни показва кои полета са свързани.

Как тълкуваме стрелката - в случая таблица persons (parent) има foreign key passport_id, който сочи към таблица passports (child дете) с ключ passport_id (или passport_id на persons сочи към passport_id на passports).

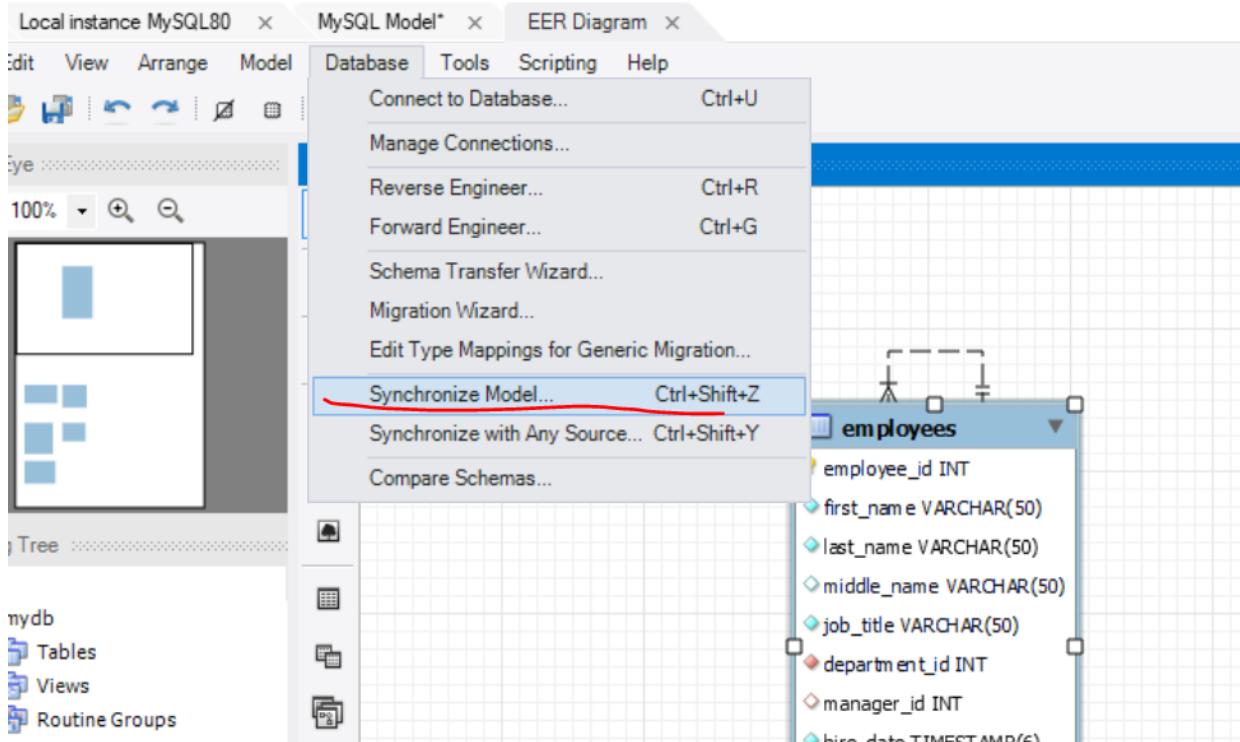
Какво означава линията:

- Пътна линия означава в двойката баща-дете, че и **двете са задължителни/и двете са primary keys**
- Прекъсната линия означава, че в двойката parent-child **не са задължителни и двете**

Synchronize model – от E/R диаграма към създаване на база данни

Каквото сме си нарисували на диаграма, го направи на истинска база

MySQL Workbench



Контрол - Януари 2022 - Веселин Бичев

Model	Update	Source
campers	→	campers
cars	↑	cars
drivers	↔	drivers
employees	→	employees
employees_projects	↑	employees_projects
mountains	↔	mountains
peaks	→	peaks
projects	↑	projects
rooms	↔	rooms
routes	↔	routes
vehicles	→	vehicles
Created from ER	→	N/A
merged_tables	↔	merged_tables

6. JOINS

6.1. Gathering Data From Multiple Tables

Можем да правим JOIN и на таблици без foreign keys

- Sometimes you need data from several tables:

Cartesian product – всяко от едната таблица с всяко от другата таблица

- Each row in the first table is paired with **all** the rows in the second table
 - When there is **no relationship** defined between the two tables

A Cartesian product is formed when:

- A join condition is omitted - когато нямаме никакъв condition
- A join condition is invalid
- All rows in the first table are joined to all rows in the second table

To avoid a Cartesian product, always include a valid join condition!

6.2. Union

- UNION combines the results from several SELECT statements
- The columns **count**, **types** and **order(порядност)** should match

```
SELECT FIRST_NAME AS NAME  
FROM EMPLOYEES
```

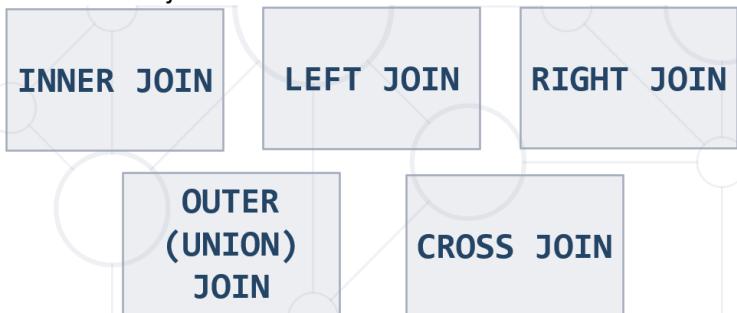
UNION

```
SELECT LAST_NAME AS NAME  
FROM EMPLOYEES
```

6.3. JOINS – used to collect data from two or more tables

Different types of joins include:

- Natural joins
- Join with USING clause
- Inner joins with ON clause
- Left, right and full outer joins
- Self joins
- Cross joins



NATURAL JOIN

The NATURAL JOIN combines the rows from two tables that have equal values in all matched by name columns.

Natural Join in SQL refers to **joining two or more tables based on common columns, which have the same name and data type**. We do not need to specify the column used for joining two tables in natural join. Natural join is used to retrieve data from more than one table in a single place.

```
SELECT * FROM ORDERS NATURAL JOIN CUSTOMERS;
```

USING clause

If several columns have the same names, we can limit the NATURAL JOIN to only one of them by the USING clause:

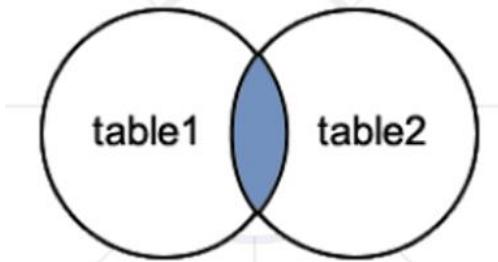
The USING clause **specifies which columns to test for equality when two tables are joined. It can be used instead of an ON clause** in the JOIN operations that have an explicit join clause.

```
SELECT E.EMPLOYEE_ID, E.LAST_NAME, D.LOCATION_ID, D.DEPARTMENT_NAME  
FROM EMPLOYEES E  
JOIN DEPARTMENTS D  
USING (DEPARTMENT_ID);
```

EMPLOYEE_ID	LAST_NAME	LOCATION_ID	DEPARTMENT_NAME
102	De Haan	1700	Executive
103	Hunold	1400	IT
104	Ernst	1400	IT
...

JOIN(== INNER JOIN) with ON clause

- To specify arbitrary conditions or specify columns to join, the ON clause is used
- Used very often
- взема сечението
- Produces a set of records which **match in both tables**



```
SELECT e.`first_name`, d.`name` AS 'dept_name'
FROM `employees` AS e           - таблица 1 = e
INNER JOIN `departments` AS d   - таблица 2 = d
ON e.`department_id` = d.`department_id`;
```

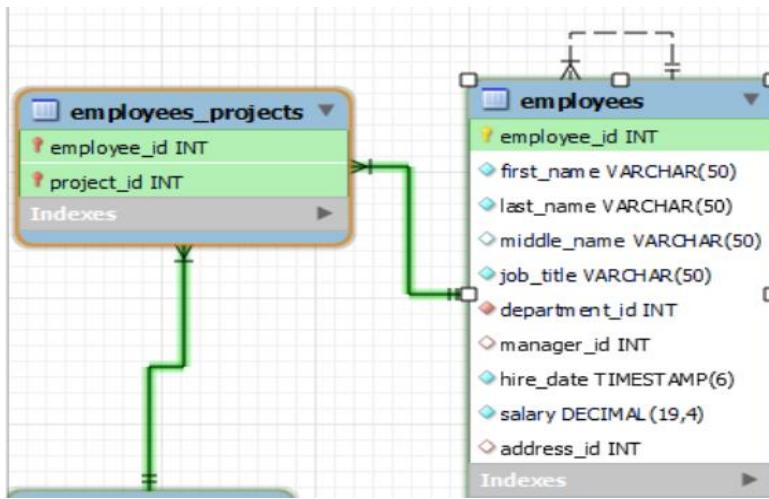
Пример 1:

//Селектирай служителите, само които имат проекти

```
SELECT e.`employee_id`, e.`first_name`, p.`name` AS 'project_name'
FROM `employees` AS e
INNER JOIN `employees_projects` AS ep
ON e.`employee_id` = ep.`employee_id`;
```

Или можем и така:

```
SELECT e.`employee_id`, e.`first_name`, p.`name` AS 'project_name'
FROM `employees` AS e
LEFT JOIN `employees_projects` AS ep //вземи всичко отляво (таблицата employees като първа таблица)
ON e.`employee_id` = ep.`employee_id`
WHERE ep.`project_id` IS NOT NULL; //НО само когато за employee_id от employees ИМА ep.`project_id`
```



Пример 2:

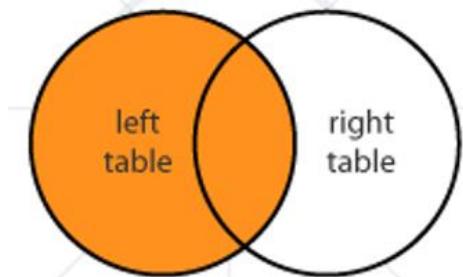
```

SELECT E.EMPLOYEE_ID, E.LAST_NAME, E.DEPARTMENT_ID, D.DEPARTMENT_ID,
D DEPARTMENT_NAME
FROM EMPLOYEES E
JOIN DEPARTMENTS D
ON (E.DEPARTMENT_ID = D.DEPARTMENT_ID)
WHERE D.DEPARTMENT_NAME != 'Department 2';
    
```

EMPLOYEE_ID	LAST_NAME	DEPARTMENT_ID	DEPARTMENT_ID	DEPARTMENT_NAME
200	Whalen	10	10	Executive
201	Hartstein	20	20	IT
202	Fay	20	20	IT

LEFT JOIN (== LEFT OUTER JOIN)

- A join that returns all the rows from the left table and the matched ones from the right
- If no match, the result is NULL on the right
- Matches every entry in **left** table regardless of match in the **right**



Example, see powerpoint presentation

Пример 1:

Тези колеги, които не са в нито един проект

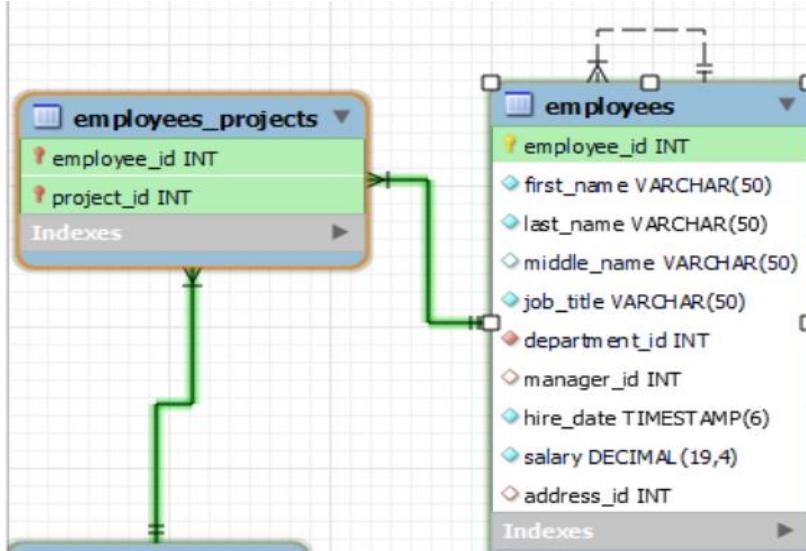
```

SELECT e.`employee_id`, e.`first_name`
FROM `employees` AS e
LEFT JOIN `employees_projects` AS ep //вземи всичко отляво (таблицата employees като първа таблица)
    
```

```

ON e.`employee_id` = ep.`employee_id`
WHERE ep.`project_id` IS NULL //когато за employee_id от employees няма ep.`project_id`
ORDER BY e.`employee_id` DESC
LIMIT 3;

```



Пример 2:

```

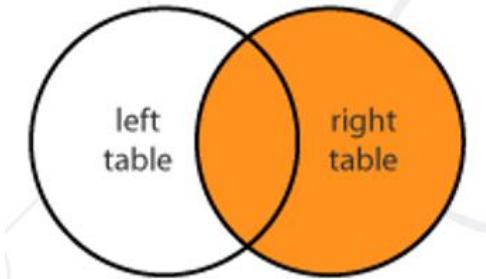
SELECT E.EMPLOYEE_ID, E.LAST_NAME, E.DEPARTMENT_ID, D.DEPARTMENT_ID,
D DEPARTMENT_NAME
FROM EMPLOYEES E
LEFT JOIN DEPARTMENTS D
ON (E.DEPARTMENT_ID = D.DEPARTMENT_ID);

```

EMPLOYEE_ID	LAST_NAME	DEPARTMENT_ID	DEPARTMENT_ID	DEPARTMENT_NAME
200	Whalen	10	10	Executive
201	Hartstein	20	20	IT
202	Fay	20	NULL ✓	NULL ✓

RIGHT JOIN (== RIGHT OUTER JOIN)

- A join that returns all the rows from the right table and the matched ones from the left
- If no match, the result is NULL on the left
- Matches every entry in **right** table regardless of match in the **left**



Example, see powerpoint presentation

```

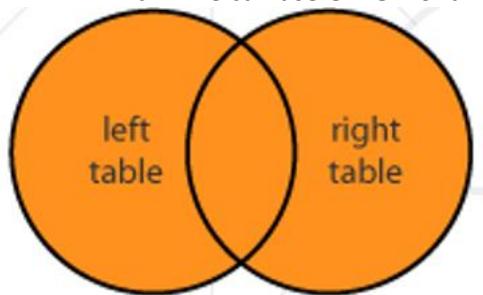
SELECT E.EMPLOYEE_NAME, D.DEPARTMENT_ID, D.DEPARTMENT_NAME
FROM EMPLOYEES E
RIGHT JOIN DEPARTMENTS D
ON (E.DEPARTMENT_ID = D.DEPARTMENT_ID);

```

EMPLOYEE_NAME	DEPARTMENT_ID	DEPARTMENT_NAME
Jennifer Whalen	10	Administration
Michael Hartstein	20	Marketing
Den Raphaely	30	Purchasing
(null) 	120	Treasury
...

FULL JOIN (==FULL OUTER JOIN)

- Not implemented in MySQL
- A join that returns all the rows where there is a match (left or right)
- If no match, the result is NULL on the respective side - left or right
- Разликата с Cartesion product е че вния пример няма дуплициране на данни
- Returns all records in both tables regardless of **any** match
 - Less useful than INNER, LEFT or RIGHT JOINS
 - We can use UNION of a LEFT and RIGHT JOIN



OUTER JOIN – взема всичко без сечението, но не е имплементирано в MySQL

```

SELECT E.EMPLOYEE_NAME, D.DEPARTMENT_ID, D.DEPARTMENT_NAME
FROM EMPLOYEES E
FULL OUTER JOIN DEPARTMENTS D
ON (E.DEPARTMENT_ID = D.DEPARTMENT_ID);

```

EMPLOYEE_NAME	DEPARTMENT_ID	DEPARTMENT_NAME
Jennifer Whalen	10	Administration
Michael Hartstein	20	Marketing
Clara Vishney	(null)	(null)
Jason Mallin	(null)	(null)
(null)	150	Shareholder Services

FULL JOIN == FULL OUTER JOIN – взема всичко **включително и сечението**, но не е имплементирано в MySQL

В MySQL нямаме нито FULL OUTER JOIN нито само FULL JOIN, и в MySQL тази операция се постига чрез UNION of LEFT and RIGHT JOIN.

UNION of LEFT and RIGHT JOIN

```
SELECT students.name, courses.name
FROM students
LEFT JOIN courses
ON students.course_id = courses.id
```

UNION

```
SELECT students.name, courses.name
FROM students
RIGHT JOIN courses
ON students.course_id = courses.id
```

Example, see powerpoint presentation

Self joining

Self join means to join a table to itself.

Always used with table aliases.

```
SELECT CONCAT (E.FIRST_NAME, ' ', E.LAST_NAME) AS EMPLOYEE_NAME,
SELECT CONCAT (M.FIRST_NAME, ' ', M.LAST_NAME) AS MANAGER_NAME
FROM EMPLOYEES E
JOIN EMPLOYEES M ON (E.MANAGER_ID = M.EMPLOYEE_ID);
```

EMPLOYEE_NAME	MANAGER_NAME
Neena Kochhar	King
Lex De Haan	King
Alexander Hunold	De Haan
Bruce Ernst	Hunold

Three way join

A three-way join is a join of three tables.

```
SELECT E.EMPLOYEE_ID, L.CITY, D.DEPARTMENT_NAME
FROM EMPLOYEES E
JOIN DEPARTMENTS D ON D.DEPARTMENT_ID = E.DEPARTMENT_ID
JOIN LOCATIONS L ON D.LOCATION_ID = L.LOCATION_ID;
```

EMPLOYEE_ID	CITY	DEPARTMENT_NAME
103	Southlake	IT
104	Southlake	IT
124	San Francisco	Administration
...

CROSS JOIN

- Produces a set of associated rows of two tables
 - Multiplication of each row in the first table with each in second
 - The result is a **Cartesian** product, when there's **no condition** in the **WHERE** clause
 - Not used often as **a lot of duplication**

```

SELECT E.LAST_NAME, D.DEPARTMENT_NAME
FROM EMPLOYEES E
CROSS JOIN DEPARTMENTS D;

```

6.4. Set operations

Set operations in SQL is a type of operations which allows **the results of multiple queries to be combined into a single result set.**

Set operators in SQL include UNION, INTERSECT, and EXCEPT, which mathematically correspond to the concepts of union, intersection and set difference.

UNION operator

sales2005

person	amount
Joe	1000
Alex	2000
Bob	5000

sales2006

person	amount
Joe	2000
Alex	2000
Zach	35000

```
SELECT * FROM sales2005
```

```
UNION
```

```
SELECT * FROM sales2006;
```

Note that there are two rows for Joe because those rows are distinct across their columns. There is only one row for Alex because those rows are not distinct for both columns.

person	amount
Joe	1000
Alex	2000
Bob	5000
Joe	2000
Zach	35000

UNION ALL gives different results, because it will not eliminate duplicates. Executing this statement:

```
SELECT * FROM sales2005
```

```
UNION ALL
```

```
SELECT * FROM sales2006;
```

person	amount
Joe	1000
Joe	2000
Alex	2000
Alex	2000
Bob	5000
Zach	35000

INTERSECT operator

The following example INTERSECT query returns all rows from the Orders table where Quantity is between **50 and 100**.

```
SELECT *
FROM Orders
WHERE Quantity BETWEEN 1 AND 100
```

INTERSECT

```
SELECT *
FROM Orders
WHERE Quantity BETWEEN 50 AND 200;
```

EXCEPT operator

The following example EXCEPT query returns all rows from the Orders table where Quantity is between 1 and 49, and those with a Quantity between 76 and 100.

Worded another way; the query returns all rows where the Quantity is between 1 and 100, apart from rows where the quantity is between 50 and 75.

```
SELECT *
FROM Orders
WHERE Quantity BETWEEN 1 AND 100
```

EXCEPT

```
SELECT *
FROM Orders
WHERE Quantity BETWEEN 50 AND 75;
```

6.5. Subqueries – една заявка в друга заявка

- Subqueries – SQL query inside a larger one
- Can be nested in **SELECT, INSERT, UPDATE, DELETE**
 - Usually added within a **WHERE** clause

```
SELECT COUNT(e.employee_id) AS `count`
```

```
FROM employees AS e
WHERE e.salary >
(
SELECT AVG(salary) AS 'average_salary' FROM employees
);
```

Конкатениране на повече от две таблици

```
SELECT e.`first_name`, e.`last_name`, t.`name` , adr.`address_text`
FROM `employees` AS e
JOIN `addresses` AS adr
ON e.`address_id` = adr.`address_id`
JOIN `towns` AS t
ON adr.`town_id` = t.`town_id`
ORDER BY e.`first_name` ASC, e.`last_name` ASC
LIMIT 5;
```

WHERE и ORDER BY ги слагаме след всички JOIN

```
SELECT c.`country_code`, m.`mountain_range` , p.`peak_name` , p.`elevation`
FROM `countries` AS c
JOIN `mountains_countries` AS mc
ON c.`country_code` = mc.`country_code`
JOIN `mountains` AS m
ON mc.`mountain_id` = m.`id`
JOIN `peaks` AS p
ON m.`id` = p.`mountain_id`
WHERE c.`country_code` = 'BG' AND p.`elevation` > 2835
ORDER BY p.`elevation` DESC;
```

Когато трябва да визуализираме променена стойност на дадена клетка

```
SELECT e.`employee_id`, e.`first_name`,
(
CASE
    WHEN YEAR(p.`start_date`) > 2004
    THEN NULL
        - променената стойност е NULL
    ELSE p.`name`
END
)
AS 'project_name'
FROM `employees` AS e
JOIN `employees_projects` AS ep
ON e.`employee_id` = ep.`employee_id`
JOIN `projects` AS p
ON p.`project_id` = ep.`project_id`
WHERE e.`employee_id` = 24
ORDER BY `project_name`;
```

#15. *Continents and Currencies

#групирана по два критерия връща за всеки континент и за всяка валута, то по колко пъти се използва дадена валута

```
SELECT contr.`continent_code`, contr.`currency_code`,
```

```

COUNT(*) AS `currency_usage`
FROM `countries` AS contr
GROUP BY contr.`continent_code`, contr.`currency_code`
HAVING `currency_usage` > 1 #всички валути, които се използват в даден континент в повече от една държава
AND
#от всички валути колко пъти са използвани, сравни само тази коя е използвана най-много пъти
`currency_usage` = (SELECT      #c.`currency_code`,
COUNT(*) AS `most_used_currency`
    FROM `countries` AS c
    WHERE c.`continent_code` = contr.`continent_code`
    GROUP BY c.`currency_code`
    ORDER BY `most_used_currency` DESC
    LIMIT 1)
ORDER BY contr.`continent_code` ASC, contr.`currency_code` ASC;

```

6.5. Subqueries example - more

Given tables Employees with columns Id, Name, Email and Salary, EmployeeDepartments with columns Id, EmployeeId, DepartmentId and Departments with columns Id, Name write an SQL query to return the names of Departments along with the number of employees in each department as Name and Count columns

Note: Use LEFT joins

```

SELECT D.Name,
(SELECT count(*) FROM EmployeeDepartments EEDD
WHERE EEDD.DepartmentId = ED.DepartmentId AND EEDD.EmployeeId IS NOT NULL) AS `COUNT`
FROM Departments D

```

```

LEFT JOIN EmployeeDepartments ED
ON ED.DepartmentId = D.Id

```

```

LEFT JOIN Employees E
ON ED.EmployeeId = E.Id

```

```

GROUP BY D.Name;
```

Групира се таблица/резултат от JOIN-овете, и там реда с Administrative си стои и е 1 на брой. Във вложената селект заявка **ED.DepartmentId** представлява тази съща резултатна таблица като в случая на Administrative записа **ED.DepartmentId** е null. Съществува и друга опция даден служител да не се води в никой отдел - **ED.DepartmentId** може да не null, а **EEDD.EmployeeId** да е null.

Name	COUNT
Administrative	0
Business Develop	1
HR	2
IT	1

6.6. More hacks

UPDATE JOIN

#Task 3 - Update - version in which we have employees with 0 clients – вложена заявка

```

UPDATE employees_clients AS ec
SET ec.employee_id =
(
    SELECT COUNT(e.id) FROM employees AS e
    LEFT JOIN (SELECT * from employees_clients) AS emcl - нова инстанция на същата таблица
    ON emcl.employee_id = e.id
    GROUP BY e.id
    ORDER BY COUNT(e.id) ASC, e.id ASC
    LIMIT 1
)
WHERE ec.employee_id = ec.client_id;

```

Или така също става

```

UPDATE `products` AS prr
JOIN `categories` AS c
ON prr.`category_id` = c.`id`
JOIN `reviews` AS r
ON r.`id` = prr.`review_id`
SET prr.`price` = prr.`price` * 0.7
WHERE c.`name` = 'Phones and tablets' AND r.`rating` < 4;

```

DELETE JOIN

Вариант 1 - Работи

```

DELETE emp FROM employees AS emp - изтрий от таблица emp
LEFT JOIN employees_clients AS ec
ON ec.employee_id = emp.id
WHERE ec.client_id IS NULL; - тези, служители, които нямат клиенти, при JOIN имат client_id да е NULL

```

Вариант 2

```
DELETE FROM employees WHERE id = - изтрий от таблица emp
```

```
(  
    SELECT emp.id FROM (SELECT * FROM employees) AS emp - нова инстанция на същата таблица  

    LEFT JOIN employees_clients AS ec  

    ON ec.employee_id = emp.id  

    WHERE ec.client_id IS NULL - тези, служители, които нямат клиенти, при JOIN имат client_id да е NULL  

)
```

GROUP BY plus JOIN

```

SELECT CONCAT(emp.first_name, ' ', emp.last_name) AS 'name',
emp.started_on,
COUNT(ec.employee_id) AS count_of_clients
FROM employees AS emp
LEFT JOIN employees_clients AS ec
ON ec.employee_id = emp.id
GROUP BY ec.employee_id
ORDER BY count_of_clients DESC, emp.id ASC
LIMIT 5;

```

7. DCL - Data Control Language

```
Grant SELECT, INSERT, UPDATE, DELETE on Employee To User1;  
Revoke INSERT On Employee To user1;  
Deny Update On Employee to user1;
```

GRANT in first case we gave privileges to user User1 to do SELECT, INSERT, UPDATE and DELETE on the table called employees.

REVOKE with this command we can take back privilege to default one, in this case, we take back command INSERT on the table employees for user User1.

DENY is a specific command. We can conclude that every user has a list of privilege which is denied or granted so command DENY is there to explicitly ban you some privileges on the database objects.:

8. Database Programmability

Да избягваме в практиката функции и процедури в MySQL

Важно: Първо разписваме праста заявка, и след това я копираме/слагаме в структурата на съответната функция/процедура/тригър

We can optimize with User-defined Functions.

Transactions improve security and consistency.

Stored Procedures encapsulate repetitive logic.

Triggers execute before certain events on tables.

8.1. User-Defined Functions

Encapsulating Custom Logic

- Extend the functionality of a MySQL Server
 - **Modular=functional** programming – write **once**, call it **any number** of times
 - Faster execution – doesn't need to be reparsed and reoptimized with each use
 - Break out complex logic into **shorter code blocks**
- Functions can be:
 - Scalar – return **single value** or **NULL**
 - Table-Valued – return a **table**

Само при използване на **DECLARE** използваме за присвояване/задаване на стойност **:=**

DRY – Don't Repeat Yourself принцип

Creating Functions

DELIMITER \$\$\$\$ - начало на разделител

CREATE FUNCTION ufn_count_employees_by_town(`town_name` VARCHAR(20))

RETURNS DOUBLE - какъв тип връща функцията, незадължителен елемент

DETERMINISTIC – при едни и същи входни данни един и същи резултат връща

BEGIN

```
DECLARE e_count DOUBLE; - деклариране на променлива  
SET e_count := (SELECT COUNT(employee_id) FROM employees AS e - за присвояване  
INNER JOIN addresses AS a ON a.address_id = e.address_id  
INNER JOIN towns AS t ON t.town_id = a.town_id  
WHERE t.name = town_name);
```

RETURN e_count; - каква стойност връща функцията, незадължителен елемент

END;

\$\$\$\$ - край на разделител

DELIMITER\$\$\$\$

```
CREATE FUNCTION ufn_count_employees_by_town(`town_name` VARCHAR(20))
```

RETURNS INT

DETERMINISTIC

BEGIN

```
DECLARE e_count INT;  
SET e_count := (SELECT COUNT(e.`employee_id`) FROM `employees` AS e  
INNER JOIN `addresses` AS a ON a.`address_id` = e.`address_id`  
INNER JOIN `towns` AS t ON t.`town_id` = a.`town_id`  
WHERE t.`name` = `town_name`);
```

RETURN e_count;

Или директно връщаме резултата

```
RETURN(  
(SELECT COUNT(e.`employee_id`) FROM `employees` AS e  
INNER JOIN `addresses` AS a ON a.`address_id` = e.`address_id`  
INNER JOIN `towns` AS t ON t.`town_id` = a.`town_id`  
WHERE t.`name` = `town_name`));
```

END;

\$\$\$\$

Deterministic vs. non-deterministic functions

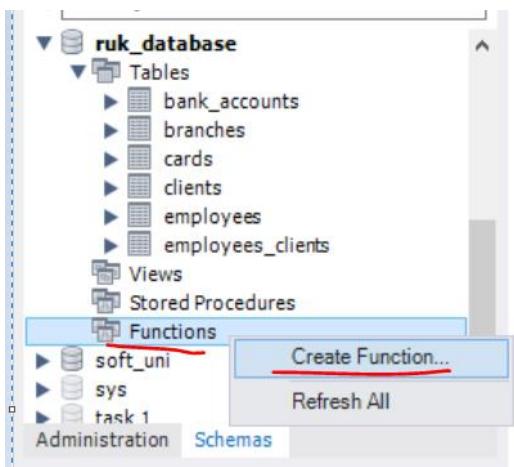
Executing and Dropping Stored Functions

И така изпълняваме функцията

```
SELECT ufn_count_employees_by_town('Sofia');
```

```
SELECT ufn_get_salary_level(51000.0);
```

```
DROP FUNCTION ufn_get_salary_level;
```



new_function - Routine

Name: new_function

DDL:

```

1 • CREATE FUNCTION `new_function` ()
2     RETURNS INTEGER
3     BEGIN
4     5     //my code here
5     RETURN 1;
6     END
7

```

8.2. Stored Procedures

Sets of Queries Stored On DB Server

- Stored procedures are logic removed from the application and placed on the database server.
 - Can greatly cut down traffic on the network
 - Improve the security of the database server
 - Separate data access routines from the business logic
- Stored procedures are accessed by programs using different platforms and API's.
- **Нямаме return при процедури**

Creating Stored Procedures

DELIMITER %%

CREATE PROCEDURE usp_select_employees_by_seniority()

BEGIN

SELECT *

FROM employees

WHERE ROUND((DATEDIFF(NOW(), hire_date) / 365.25)) < 15;

END %%

Executing and Dropping Stored Procedures

- Executing a stored procedure by **CALL**

CALL usp_select_employees_by_seniority();

- **DROP PROCEDURE**

```
DROP PROCEDURE usp_select_employees_by_seniority;
```

Defining Parameterized Procedures

- To define a parameterized procedure use the syntax:

```
CREATE PROCEDURE usp_procedure_name
```

```
(parameter_1_name parameter_type,  
parameter_2_name parameter_type,...)
```

Пример 1:

```
DELIMITER $$
```

```
CREATE PROCEDURE usp_select_employees_by_seniority(min_years_at_work INT)
```

```
BEGIN
```

```
    SELECT first_name, last_name, hire_date,  
          ROUND(DATEDIFF(NOW(), DATE(hire_date)) / 365.25, 0) AS 'years'  
     FROM employees  
    WHERE ROUND(DATEDIFF(NOW(), DATE(hire_date)) / 365.25, 0) > min_years_at_work  
    ORDER BY hire_date;  
END $$
```

```
CALL usp_select_employees_by_seniority(15);
```

Пример 2:

```
Task 2 -Employees Promotion
```

```
DELIMITER $
```

```
CREATE PROCEDURE usp_raise_salaries (dept_name VARCHAR(45))
```

```
BEGIN
```

```
/*business logic*/  
    UPDATE `employees` AS e  
   JOIN `departments` AS d  
  ON d.`department_id` = e.`department_id`  
 SET e.`salary` = e.`salary` * 1.05  
 WHERE d.`name` = dept_name;
```

```
    SELECT e.`first_name`, e.`salary` FROM `employees` AS e  
   JOIN `departments` AS d  
  ON d.`department_id` = e.`department_id`  
 WHERE d.`name` = dept_name  
 ORDER BY e.`first_name`, e.`salary`;
```

```
END $
```

```
CALL usp_raise_salaries('Finance');
```

Пример 3:

```
DELIMITER %%
```

```
CREATE PROCEDURE usp_get_towns_starting_with (starts_with VARCHAR(20))
```

```
BEGIN
```

```
    SELECT `name` FROM `towns`  
   WHERE `name` LIKE concat(starts_with, '%') - тук конкатенираме  
 ORDER BY `name`;  
END; %%
```

```
CALL usp_get_towns_starting_with('S');
```

Returning Values Using OUTPUT Parameters

```
DELIMITER $$  
CREATE PROCEDURE usp_add_numbers (first_number INT, second_number INT, OUT result INT)  
BEGIN  
    SET result = first_number + second_number;  
END $$
```

Все едно процедурата работи като функция:

```
SET @answer=0;  
CALL usp_add_numbers(5, 6,@answer);  
SELECT @answer;
```

Функции и процедури рядко се използват – те са бизнес логика, която не е редно да стои при базата данни на сървъра.

8.3. TCL - Transactions – a kind of a stored procedure

- Една процедура има една или няколко SQL заявки, които се изпълняват една след друга.
- **При транзакции, може да се създаде логика, която да изпълнява една или друга заявка или всички заявки, ИЛИ никоя от заявките да не се изпълни**
- Transaction is a **sequence of actions SQL queries** (database operations) executed as a whole
 - Either **all** of them complete successfully or **none** of them
- Example of transaction
 - A bank transfer from one account into another (withdrawal + deposit)
 - If either the withdrawal or the deposit fails **the whole operation is cancelled**
- Транзакцията се пише вътре в процедура(stored procedure)
- Ако имаме много заявки, и ако някоя се чупи, то можем да зададем ROLLBACK и на всички изпълнени до момента транзакции

ACID model is used by InnoDB engine на MySQL

Transactions Behavior

- Transactions guarantee the **consistency** and the **integrity** of the database.
 - All changes in a transaction are temporary
 - Changes are persisted when **COMMIT** is executed – всичко е временно докато не commit-нем

COMMIT е зададен като default - да

- At any time all changes can be canceled by **ROLLBACK**
- All of the operations are executed as a whole.

Само намаля в случая:

START TRANSACTION;

```
UPDATE `employees` SET `salary` = `salary` - 1000 WHERE `employee_id` = 1;
```

Отново само намаля в случая:

START TRANSACTION;

```
UPDATE `employees` SET `salary` = `salary` - 1000 WHERE `employee_id` = 1;
```

```
COMMIT;
```

Намалянето не се извършва:

```
START TRANSACTION;
```

```
    UPDATE `employees` SET `salary` = `salary` - 1000 WHERE `employee_id` = 1;
```

```
ROLLBACK;
```

#Task 3

```
DELIMITER %%
```

```
CREATE PROCEDURE usp_raise_salary_by_id(id int)
```

```
BEGIN
```

```
    DECLARE does_exist INT;
```

```
    START TRANSACTION;
```

```
        UPDATE employees SET salary = salary * 1.05 WHERE employee_id = id; - update-ни всички, но не записвай нищо още реално
```

```
    SET does_exist := (SELECT COUNT(*) FROM employees WHERE employee_id = id);
```

```
    IF (does_exist = 1)
```

```
        THEN COMMIT; - сера ги презапиши!!!
```

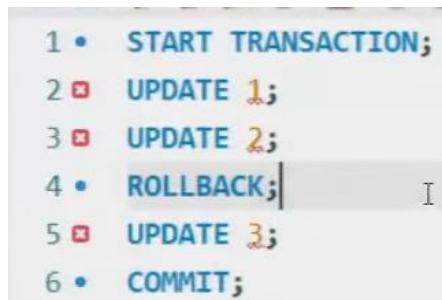
```
    ELSE
```

```
        ROLLBACK;
```

```
    END IF;
```

```
END %%
```

```
CALL usp_raise_salary_by_id(1);
```



Modern DBMS servers have built-in transaction support

- Implement "ACID" transaction
- hgfd
- ACID means:
 - Atomicity
 - Consistency
 - Isolation
 - Durability

#13. Withdraw Money

```
DELIMITER $$$
```

```
CREATE PROCEDURE usp_withdraw_money(account_id INT, money_amount DECIMAL(19, 4))
```

```
BEGIN
```

```

DECLARE bal DECIMAL(19, 4);
START TRANSACTION;
    SET bal := (SELECT a.`balance` FROM `accounts` AS a WHERE a.`id` = account_id) - money_amount;

    UPDATE `accounts`
    SET `balance` = `balance` - money_amount
    WHERE `id` = account_id;

    IF (money_amount > 0 AND bal > 0.0000) THEN COMMIT;
    ELSE ROLLBACK;
    END IF;
END;
$$$

DROP PROCEDURE usp_withdraw_money;

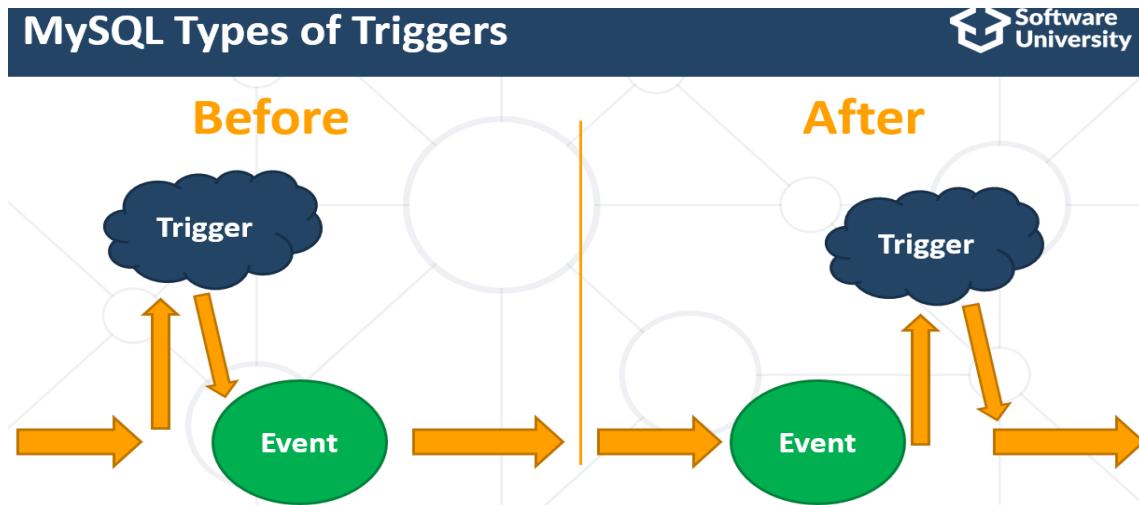
CALL usp_withdraw_money(1, 20);

```

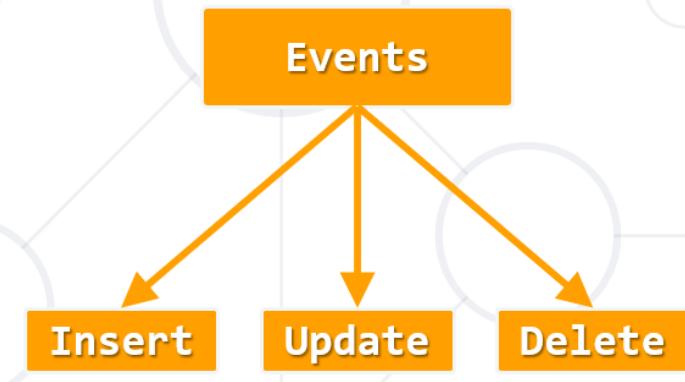
8.4. Triggers – като Event Listener

Event-listener – чакаме нещо да се случи, и тогава прави нещо

- Triggers - small programs in the database itself, activated by the database events application layer
 - UPDATE, DELETE or INSERT queries
 - Called in case of specific **event**
- We do not call triggers **explicitly**
- Triggers are **attached to a table**



There are three different events that can be applied within a trigger:



The OLD and NEW keywords allow you to access columns before/after trigger action

OLD – **before trigger action** – използва се при AFTER events

NEW – **after trigger action** – използва се при BEFORE events

After Delete

```

CREATE TABLE deleted_employees(
    employee_id INT PRIMARY KEY AUTO_INCREMENT,
    first_name VARCHAR(20),
    last_name VARCHAR(20),
    middle_name VARCHAR(20),
    job_title VARCHAR(50),
    department_id INT,
    salary DOUBLE
);

```

```

CREATE TRIGGER tr_deleted_employees
AFTER DELETE
ON employees
FOR EACH ROW
INSERT INTO deleted_employees (first_name, last_name, middle_name, job_title, department_id, salary)
VALUES(OLD.first_name, OLD.last_name, OLD.middle_name, OLD.job_title, OLD.department_id, OLD.salary);

```

Table Name: employees Schema: soft_uni
 Charset/Collation: Default Charset Default Collation Engine: InnoDB
 Comments:

```

BEFORE INSERT
AFTER INSERT
BEFORE UPDATE
AFTER UPDATE
BEFORE DELETE
AFTER DELETE
  tr_deleted_employees
  
```

Triggers

Before Insert

```

CREATE TRIGGER trigger_employee
BEFORE INSERT
ON `employees`
FOR EACH ROW
BEGIN
END;
  
```

After Update

```

CREATE DEFINER = CURRENT_USER TRIGGER `employee_AFTER_UPDATE` AFTER UPDATE
ON `employees`
FOR EACH ROW
INSERT INTO addresses_archive (old_salary, new_salary) VALUES (OLD.salary, NEW.salary)
INSERT INTO LOGS addresses_archive (old_salary, new_salary) VALUES (OLD.salary, NEW.salary)
  
```

За по-лесно, може да ползваме и този прозорец

```

BEFORE INSERT
AFTER INSERT
BEFORE UPDATE
AFTER UPDATE
  employees_AFTER_UPDATE
BEFORE DELETE
AFTER DELETE
  
```

```

1 • CREATE DEFINER = CURRENT_USER TRIGGER `soft_uni`.`employees_AFTER_UPDATE` AFTER UPDATE ON `employees` FOR EACH
2   BEGIN
3   END
4
5
  
```

Triggers

#15. Log Accounts Trigger

```

CREATE TABLE `logs`(
`log_id` INT PRIMARY KEY AUTO_INCREMENT,
  
```

```
`account_id` INT NULL,  
`old_sum` DECIMAL(19, 4) NOT NULL,  
`new_sum` DECIMAL(19, 4) NOT NULL);  
  
DELIMITER $$  
CREATE TRIGGER `tr_balance_updated`  
AFTER UPDATE ON `accounts`  
FOR EACH ROW  
BEGIN - пишем го след FOR EACH ROW  
    IF OLD.balance` != NEW.balance` THEN  
        INSERT INTO `logs`(`account_id`, `old_sum`, `new_sum`) VALUES (OLD.`id`, OLD.`balance`, NEW.`balance`);  
    END IF;  
END;$$  
  
DROP TRIGGER `tr_balance_updated`;
```

Тази проверка е излишна, освен разбира се ако има промяна не в баланса, а в други данни
#IF OLD.balance` != NEW.balance` THEN
#END IF;



```
CREATE TABLE IF NOT EXISTS `accounts` (
    `id` int(11) NOT NULL,
    `account_holder_id` int(11) NOT NULL,
    `balance` decimal(19,4) DEFAULT '0.0000',
    PRIMARY KEY (`id`),
```

```
CREATE TABLE logs(
    log_id INT PRIMARY KEY AUTO_INCREMENT,
    account_id INT,
    old_sum DECIMAL(19, 4),
    new_sum DECIMAL(19, 4)
);
```

```
CREATE DEFINER='root'@'localhost' TRIGGER `accounts_AFTER_UPDATE`
AFTER UPDATE
ON `accounts` FOR EACH ROW
INSERT INTO `logs`(`account_id`, `old_sum`, `new_sum`) - вкарай в новата таблица `logs`
VALUES (OLD.`id`, OLD.`balance`, NEW.`balance`); - вземи данни от таблица `accounts` където правим промени
```

9. Глобални настройки / стойности

```
SELECT DATABASE(); --Returns the name of the current database
SELECT USER(); --Returns the current MySQL user and host
SELECT VERSION(); --Returns the version of MySQL you are using
```

За смяна на root паролата

```
ALTER USER 'root'@'localhost' IDENTIFIED BY '';
flush privileges;
```

За проверка свързана с външни ключове

```
#Игнорирай външните ключове, за да можем да си трием спокойно :)
SET FOREIGN_KEY_CHECKS = 0;
```

#Имай в предвид външните ключове

```
SET FOREIGN_KEY_CHECKS = 1;
```

Общи

```
SET GLOBAL log_bin_trust_function_creators = 1; //да се доверявам на функциите ако аз съм ги създал
SET SQL_SAFE_UPDATES = 0; //да мањем safe updates
```

```
SELECT @@sql_mode;
```

SET sql_mode = 'ONLY_FULL_GROUP_BY'; - активен този mode

SET sql_mode = ";- премахни only full group by – да го използваме в Judge, но няма ефект 😞

Има промяна в Judge – Judge работи вече само в mode 'ONLY_FULL_GROUP_BY'

Когато групирате по PRIMARY KEY, няма проблеми обаче!!!

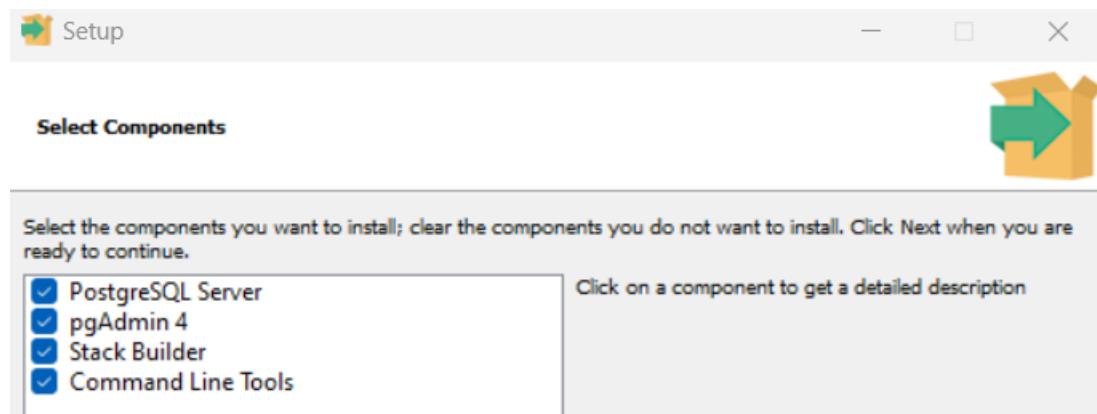
Друг вариант е да използваме вложени заявки и да не използваме GROUP BY!!!

GROUP_CONCAT

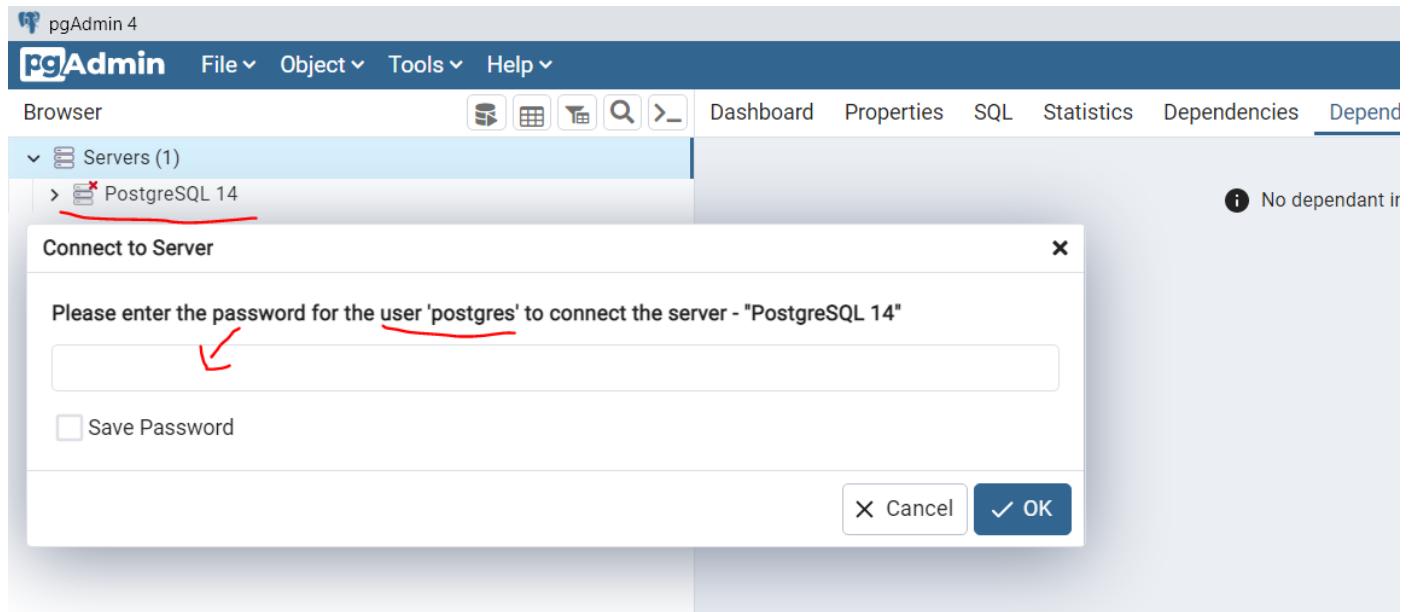
```
ANY_VALUE(muahaha) AS `bla_bla_bla`;
```

IV. Installing PostgreSQL

<https://www.postgresqltutorial.com/postgresql-getting-started/install-postgresql/>

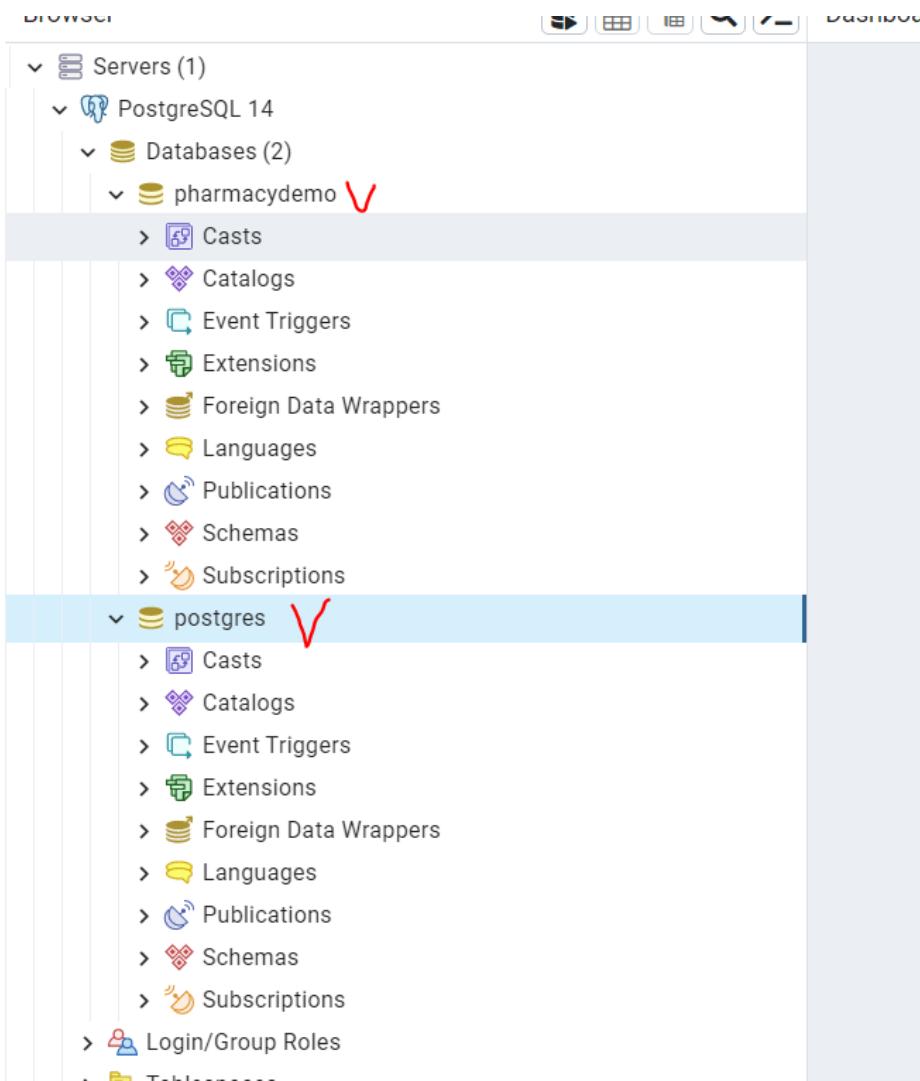


Creating new database

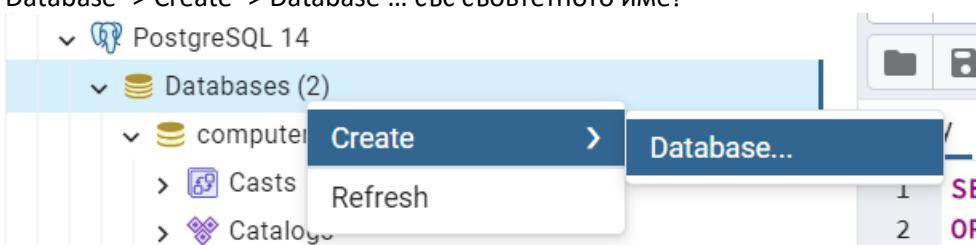


Мога да си създавам new database на server PostgreSQL 14 (user: postgres pass: 1), и да ползвам schema public винаги, и така няма проблем.

```
quarkus.datasource.db-kind=postgresql
quarkus.hibernate-orm.database.generation=update
quarkus.datasource.username = postgres
quarkus.datasource.password = 1
quarkus.datasource.jdbc.url=jdbc:postgresql://localhost:5432/pharmacydemo
```



Врътката е да си създадеш **ръчно и предварително!!!** (за разлика от WorkBench където става автоматично)
Database -> Create -> Database ... със съответното име!



И след това няма проблем

```
driverClassName: org.postgresql.Driver
url:
jdbc:postgresql://localhost:5432/computerstore?allowPublicKeyRetrieval=true&useSSL=false&createDatabaseIfNotExist=true&serverTimezone=UTC
username: postgres
password: 1111
```

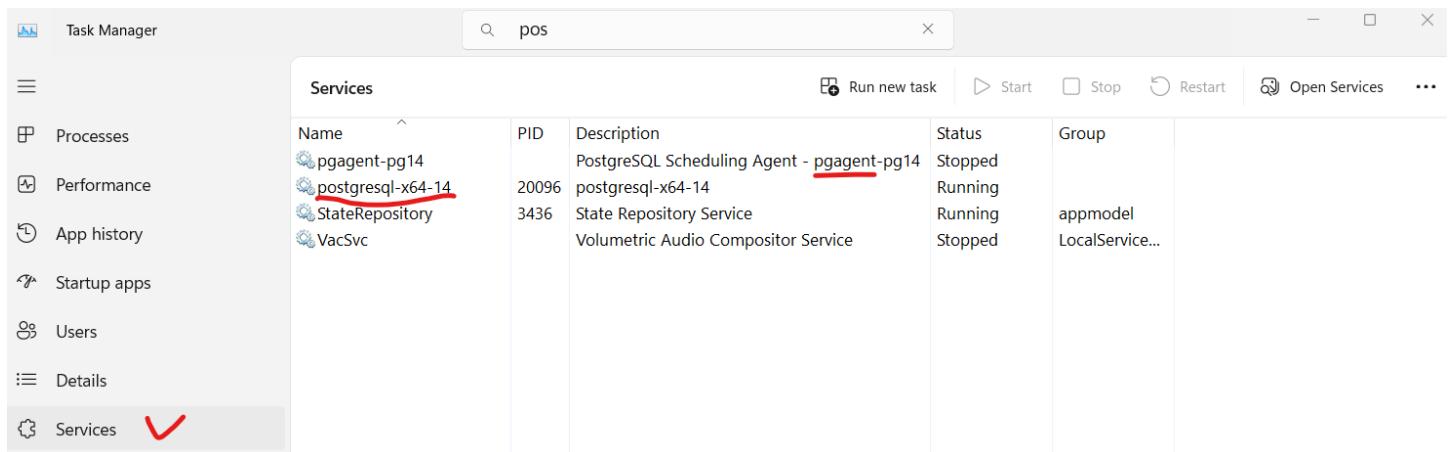
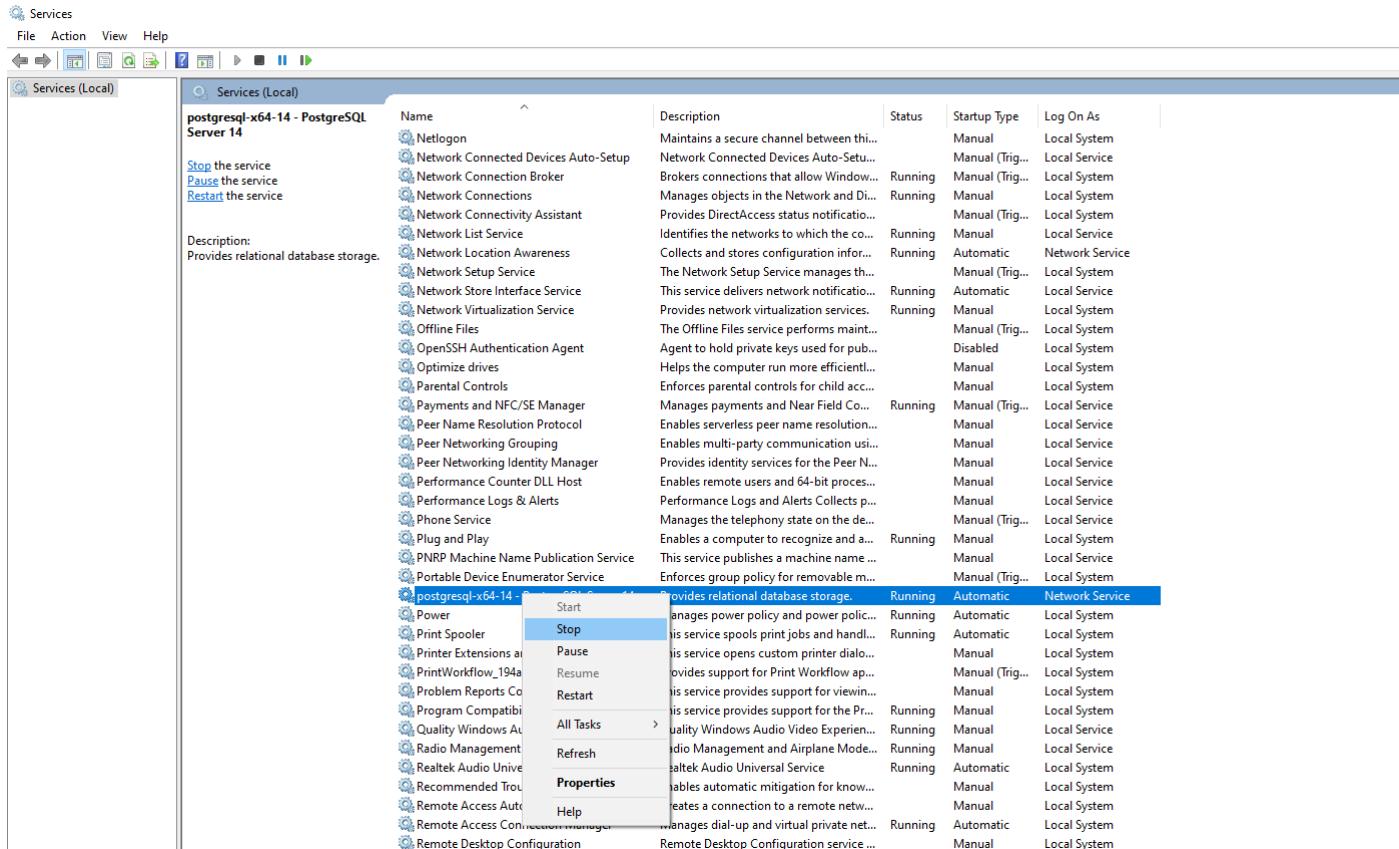
Stopping postgres server on Windows

Когато сървър базата се инициализира в Докер/облачно за даден проект, то трябва да спрем ръчно съществуващия service (database server) на Windows.

При Windows за да пуснем контейнер, то трябва да използваме Docker Desktop (не знам дали има background приложение на Docker за Windows).

Докато при Линукс може и да не използваме Docker Desktop, а само с бекграунд приложение на Docker.

1. Open Run Window by Winkey + R.
2. Type **services.msc**
3. Search Postgres service based on version installed.
4. Click stop, start or restart the service option.



V. PostgreSQL vs MySQL

В общи линии, всичко каквото има MySQL, го има и в PostgreSQL.

Но не всичко, което е в PostgreSQL го има в MySQL.

Връщане на резултата от заявката

При MySQL ще се наложи да направим втора заявка за гетване на обновения запис.

```
insert into USERS(username, email, password)
values('spas2', 'spas2@gmail.com', 'password')
returning *;
```

Nested SELECT subqueries

SELECT statements can be nested in the where clause.

- Correlated and Uncorrelated subqueries
- Note: If a subquery and a JOIN can accomplish the same task, it is often **more efficient to use a JOIN!**
 - понеже базата прави оптимизации при JOIN, а не прави оптимизации при **nested subqueries**
- **When two or more different tables, subquery possible only in PostgreSQL!**

Subquery in the SELECT clause

```
SELECT firstName, lastName,
(SELECT departmentName FROM Department WHERE Department.departmentId =
Employees.departmentId) AS departmentName
FROM Employees;
```

Subquery in the FROM clause

```
SELECT AVG(employeeCount)
FROM (
  SELECT COUNT(*) AS employeeCount
  FROM Employees
  GROUP BY departmentId
) AS subQuery;
```

Subquery in the WHERE clause

Correlated subquery

```
SELECT E.firstName, E.lastName, E.departmentId
FROM Employees E
WHERE EXISTS ( using the EXISTS function
  SELECT *
  FROM Department D
  WHERE D.departmentId = E.departmentId AND D.departmentName = 'Department 1'
);
```

The argument of EXISTS is an arbitrary SELECT statement, or subquery. The subquery is evaluated to determine whether it returns any rows. If it returns at least one row, the result of EXISTS is "true";

UPDATE JOIN

Updating data - using data from other tables

PostgreSQL can use also subqueries:

```
UPDATE table_name  
SET column1 = (SELECT column2 from other_table WHERE condition)  
WHERE condition;
```

MySQL cannot use subqueries, but join:

```
UPDATE table_name  
JOIN other_table  
ON table_name.column1 = other_table.column2  
SET table_name.column1 = new_value  
WHERE condition;
```

DELETE JOIN

Deleting data using Joins

PostgreSQL can use also subqueries:

```
DELETE FROM Employees  
WHERE id IN (SELECT id FROM Departments WHERE condition);
```

MySQL cannot use subqueries, but join:

```
DELETE Employees, Departments FROM Employees  
JOIN Departments ON Employees.id = Departments.id  
WHERE condition;
```

```
DELETE empl FROM Employees AS empl  
JOIN Departments ON empl.id = Departments.id  
WHERE condition;
```

ON CONFLICT (like a trigger)

PostgreSQL

```
explain INSERT INTO Employees(id, firstname, lastname, departmentid, dateofbirth,  
employeecardid, status)  
VALUES(34, 'Jo', 'Doe', 1111, '1990-01-01', 1234, 'Available')  
ON CONFLICT (id) DO UPDATE SET  
status='Blocked',  
timestamp=now();
```

MySQL

```
explain INSERT INTO Employees(id, firstname, lastname, departmentid, dateofbirth,  
employeecardid, status)  
VALUES(34, 'Jo', 'Doe', 1111, '1990-01-01', 1234, 'Available')  
ON DUPLICATE KEY UPDATE  
status='Blocked',  
timestamp=now();
```

A row to be inserted would cause a duplicate value in a **UNIQUE** index or **PRIMARY KEY**
/An **INSERT ... ON DUPLICATE KEY UPDATE** statement against a table having more than one
unique or primary key is also marked as unsafe. (Bug #11765650, Bug #58637)

Numbers in PostgreSQL vs numbers in Java

32 bit signed integer - positive and negative: +(2^31-1) and -(2^31) :

В java това е int
В PostgreSQL това е int4

64 bit signed integer - positive and negative: +(2⁶³-1) and -(2⁶³):
В java това е long
В PostgreSQL това е int8

BigDecimal

В java - BigDecimal
В PostgreSQL - numeric или decimal

Name	Storage Size	Description	Range
decimal	variable	user-specified precision, exact	up to 131072 digits before the decimal point; up to 16383 digits after the decimal point
numeric	variable	user-specified precision, exact	up to 131072 digits before the decimal point; up to 16383 digits after the decimal point

В postgresql няма разлика между decimal(19, 2) и numeric(19, 2), но в други SQL има разлика.

numeric(19,2)

nextval() setval()

nextval() взема следващ номер в поредицата на postgresql

```
create sequence hibernate_sequence start 1 increment 1;

create table if not exists loyalty_level (
    id int8 not null,
    level int4 not null,
    points int8 not null,
    primary key (id)
);

insert into loyalty_level(id, level, points)
values
(nextval('hibernate_sequence'), 1, 0),
(nextval('hibernate_sequence'), 2, 20),
(nextval('hibernate_sequence'), 3, 35),
(nextval('hibernate_sequence'), 4, 55),
```

Друг начин за вкарване на данни

```
create sequence if not exists Tag_SEQ increment by 50;
select setval('Tag_SEQ',  (SELECT MAX(id) FROM Tag));

insert into Tag (id, name)
select nextval('Tag_SEQ'), 'KYC_GAMEPLAY_BLOCK'
where not exists(select id from tag where name = 'KYC_GAMEPLAY_BLOCK');
```

temporary sequence

-- First you create temporary sequence. It will get dropped when the session is ended.

```
create temporary sequence s_milestone;
```

-- Next you assign the milestones with a simple update query

```
update loyalty_level  
set milestone = nextval('s_milestone')  
where level in (1,6,11,16,21)  
order by level asc;
```

Cursor (in MySQL)

Еквивалента на курсор е итератор когато говорим за колекции!

```
• CALL `hrm`.`print_employee_details`();  
  
• CALL `hrm`.`print_employee_info`('test@test.com');  
  
DELIMITER $$  
CREATE DEFINER='root'@'localhost' PROCEDURE `print_employee_info`(email varchar(200))  
BEGIN  
    DECLARE v_name varchar(100);  
    DECLARE v_email varchar(254);  
    DECLARE emp_cursor CURSOR FOR  
        select Name, Email from employees where email = email;  
    OPEN emp_cursor;  
    FETCH emp_cursor into v_name, v_email;  
    SELECT CONCAT(v_name, ' ', v_email);  
    close emp_cursor;  
END$$  
DELIMITER ;
```

Procedure

POSTGRESQL syntax

```
alter table if exists loyalty_level  
add column milestone int4 not null default 0;  
  
DO  
$$  
DECLARE  
    next_milestone int4;  
    current_level int4;  
    level_config_size int4;  
BEGIN  
    current_level := (SELECT ll.level FROM loyalty_level AS ll ORDER BY ll.level ASC  
LIMIT 1);--BO can update levels by deleting level number 1, so we take the lowest present level  
number  
    next_milestone := 1;  
    level_config_size := (SELECT COUNT(*) FROM loyalty_level);  
  
    WHILE (current_level <= level_config_size) LOOP  
        IF current_level IN (1,6,11,16,21)  
        THEN  
            --levels between milestone levels also have milestones  
            update loyalty_level  
            set milestone = next_milestone  
            where level >= current_level;
```

```

        next_milestone := next_milestone + 1;--next milestone always increase with 1
    END IF;

        current_level := current_level + 1;--levels always increase with 1
    END LOOP;
END
$$;
```

Using cast()

In PostgreSQL we can use cast

```
SELECT cast(p.id as text) FROM Profile p WHERE p.status IN (?)
```

COALESCE function syntax

PostgreSQL

The COALESCE function accepts an unlimited number of arguments. It returns the first argument that is not null. If all arguments are null, the COALESCE function will return null.

```
update wallet w
set lifetime_deposit = coalesce((select sum(d.amount)
                                    from deposit d
                                   where d.wallet_id = w.id
                                     and d.balance_change_reason in ('PLAYER_DEPOSIT',
'PLATFORM_MIGRATION')
                                     and d.status = 'ACCEPTED'), 0);
```

SELECT in WITH clause - only in PostgreSQL

Common Table Expression (CTE) subquery

The basic value of SELECT in WITH is to break down complicated queries into simpler parts. An example is:

```
WITH regional_sales AS (
    SELECT region, SUM(amount) AS total_sales
    FROM orders
    GROUP BY region
),
top_regions AS (
    SELECT region
    FROM regional_sales
    WHERE total_sales > (SELECT SUM(total_sales)/10 FROM regional_sales)
)
SELECT region,
       product,
       SUM(quantity) AS product_units,
       SUM(amount) AS product_sales
FROM orders
WHERE region IN (SELECT region FROM top_regions)
GROUP BY region, product;
```

which displays per-product sales totals in only the top sales regions. The WITH clause defines two auxiliary statements named `regional_sales` and `top_regions`, where the output of `regional_sales` is used in `top_regions` and the output of `top_regions` is used in the primary SELECT query. This example could have

been written without `WITH`, but we'd have needed two levels of nested sub-`SELECT`s. It's a bit easier to follow this way.

Common functions

String functions:

MySQL: `CONCAT()`, `SUBSTRING()`, `LENGTH()`, `UPPER()`, `LOWER()`, `TRIM()`, `REPLACE()`

PostgreSQL: `CONCAT()`, `SUBSTRING()`, `LENGTH()`, `UPPER()`, `LOWER()`, `TRIM()`, `REPLACE()`

Mathematical functions:

MySQL: `ABS()`, `CEIL()`, `FLOOR()`, `ROUND()`, `SQRT()`, `LOG()`, `EXP()`, `POW()`, `RAND()`

PostgreSQL: `ABS()`, `CEILING()`, `FLOOR()`, `ROUND()`, `SQRT()`, `LN()`, `EXP()`, `POWER()`, `RANDOM()`

Date and time functions:

MySQL: `NOW()`, `CURDATE()`, `CURTIME()`, `DATE()`, `TIME()`, `YEAR()`, `MONTH()`, `DAY()`, `HOUR()`, `MINUTE()`, `SECOND()`,
`EXTRACT(YEAR from date)`, `EXTRACT(MONTH from date)`, `EXTRACT(DAY from date)`,

PostgreSQL: `NOW()`, `CURRENT_DATE`, `CURRENT_TIME`, `DATE()`, `TIME()`,
`EXTRACT(YEAR from date)`, `EXTRACT(MONTH from date)`, `EXTRACT(DAY from date)`,
`EXTRACT(HOUR FROM time)`, `EXTRACT(MINUTE FROM time)`, `EXTRACT(SECOND FROM time)`

Conditional expressions:

MySQL: `IF()`, `CASE`

PostgreSQL: `CASE`, `COALESCE()`, `NULLIF()`

Regular expressions:

MySQL: `REGEXP`, `NOT REGEXP`

PostgreSQL:

`~`

`!~`

`~*`

`!~*`

```
SELECT firstname FROM employees WHERE firstname ~ 'a'; --Returns names containing 'a'  
SELECT firstname FROM employees WHERE firstname !~ 'a'; --Returns names not containing 'a'
```

Array to String

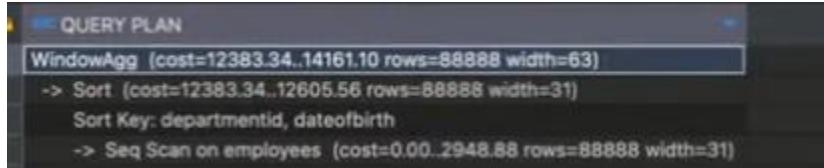
In PostgreSQL

```
SELECT ARRAY_TO_STRING(ARRAY[1, 2, 3, 4], ',');  
SELECT STRING_TO_ARRAY('1,2,3,4', ',');  
SELECT ARRAY_LENGTH(ARRAY[1, 2, 3, 4], 1);
```

Window aggregation

Both for MySQL and PostgreSQL

```
select FirstName, Salary,  
avg(Salary) over (PARTITION by DepartmentId order by DateOfBirth)  
from Employees;
```



Глобални команди

```
SELECT CURRENT_DATABASE(); --Returns the name of the current database  
SELECT CURRENT_USER; --Returns the name of the current user  
SELECT VERSION(); --Returns the version of PostgreSQL you are using
```

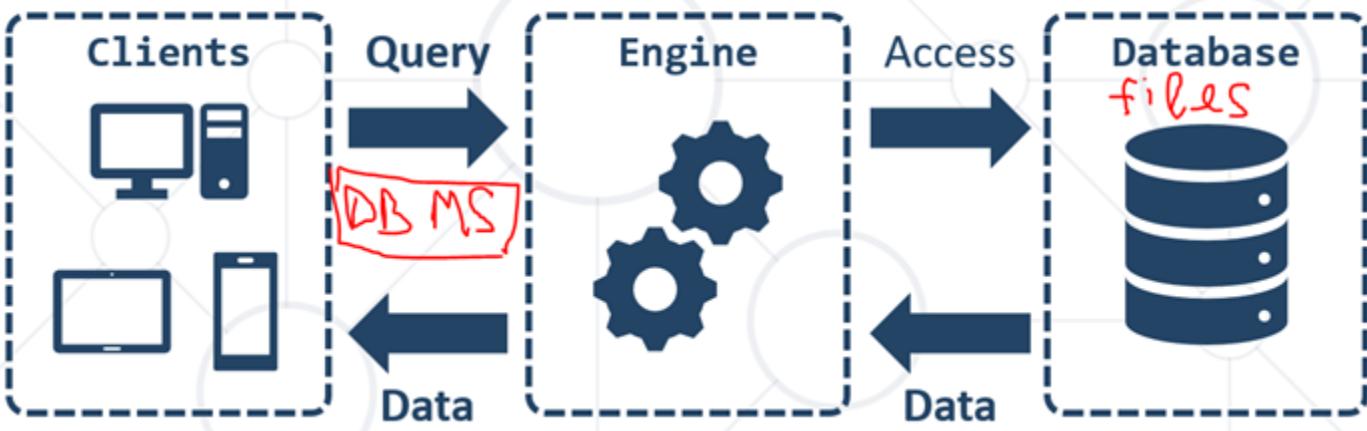
SYSTEM information functions;

VI. NoSQL

Database Engine Flow



- SQL Server uses the Client-Server Model



1. General info

- NoSQL - stands for “**Not only SQL**”, non-tabular databases
- First used in 1998
- Tackle (да се пребори с) problems of performance, volume and variety of data
- Handling of unstructured or semi-structured data
- Scale horizontally - добавяме все повече и повече load-ове, партишишъни, клъстъри и сървъри (вместо повече процесор и рам памет - вертикално скалиране)
- Flexibility, scalability and speed.

2. NoSQL vs Relational Model

Similarities:

- Both are used for storing and managing data
- Both support CRUD (Create, Read, Update, Delete) operations
- Both provide mechanisms for ensuring data consistency and integrity

Differences:

- **Schema:** Relational databases have a fixed schema while NoSQL databases are usually schema-less.
- **Scalability:** Relational databases traditionally scale vertically while NoSQL databases are designed to scale horizontally
- **Consistency:** Relational databases provide strong consistency (ACID [Atomicity-Consistency-Isolation-Durability] properties) while NoSQL databases often offer eventual consistency
- **Data model:** Relational databases use tables with rows and columns for data storage while NoSQL databases use a variety of data models: including key-value, document, columnar and graph.

NoSQL Databases (Non-Relational Databases)

A NoSQL database, has dynamic schema for unstructured data

- има задължителни полета, но има и полета, които може да не са попълнени
- Може да се променя формата на данните, които се съхраняват чрез Dynamic data – това е плюс за бизнес, който често сменя нещата в базата си от данни поради характера на бизнеса си.

<https://www.xplenty.com/blog/the-sql-vs-nosql-difference/>

SQL могат и хоризонтално, но е по-трудно

NoSQL – в 3 през нощта 300 000 души се записват и искат да се логнат. Тогава се активират хоризонтални shardins/кльстери/паралелни сървъри и те успяват да запишат всички 300 000 души, които се регистрират 2 3 през нощта.

3. Fundamental concepts in NoSQL

CAP Theorem

- We can guarantee two out of these three properties:
 - **Consistency**
 - **Availability**
 - **Partition Tolerance**
- **Consistency:** every read gets the most recent write or error
- **Availability:** every request receives a response without guaranteeing it contains most recent write (последно записаната информация) - т.е. ако даден кльстър/сървър падне, то ще вземем информацията от друг сървър и потребителят няма да разбере, че получава неактуални стари данни
- **Partition Tolerance:** system can continue operation even if there are communication breakdown between nodes

CAP Theorem - Impact

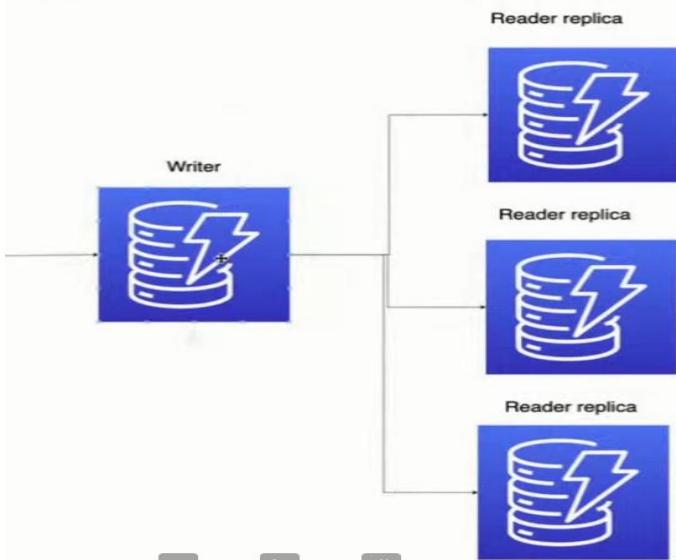
- If a system chooses consistency and availability (CA), then it can't tolerate network partitions
- If a system that opts for consistency and partition tolerance (CP), then it may not always be available
- If a system chooses availability and partition tolerance (AP), then it may occasionally serve stale or inconsistent old data
- Network partitions are almost inevitable in distributed systems (срив в комуникацията е напълно възможно да се случи) and **most databases prioritize partition tolerance** plus the expense of availability **or** consistency

Consistency

Уточнение какво имаме в предвид под сървър. Сървърно приложение, което вдига инстанция на база данни. Т.е. Когато си пуснем локално база данни е една инстанция. Но на практика могат да бъдат много инстанции на едно и също приложение. Начина по който си комуникират е въпрос на архитектура. Примерно в Кубернетис може да имаме няколко клъстъра за един и същи микросървис, по-точно говорейки за бази данни няколко клъстъра/апликации за една и съща база данни.

- **Data is the same across all nodes at any given time.**

- A write operation on one Node will be reflected by reading data from any node - като запишем една информация, то от който и да е Reader Replica Node трябва да достъпим същите новозаписани данни



- **Strong consistency**

- Write operation is not considered complete until the changes have been propagated to all nodes
 - т.е. докато не се запише записа **в/на** всяка една reader replica, то не го считаме за complete. Т.е. ако някоя reader replica е долу, може да чакаме доста или може и изобщо да не се финализира/осъществи записа
- Higher latency and less availability

- **Eventual consistency**

- Write operation is considered complete once it's written to any node - или в нашия пример когато информацията се запише в основния Writer replica, то не ни интересува дали веднага се е пропагирана към останалите Reader replica.
- Changes are propagated to other nodes eventually/в последствие
- Stale/old data read possible

Partitioning

- Breaking up large datasets into smaller ones `partitions` or also called sharding
- Each partition is stored in a separate database, node or server
- Better distribution of read/write loads and improved scalability and performance

Комуникацията между всяка реплика/партитишън най-често се осъществява чрез хеш кода за бързина.

Types of Partitioning

Horizontal (Sharding)

- Data is split **into rows** and each set of rows is stored in different partition - but **same number of columns!!**
- Each partition forms a separate table with same schema and number of columns

Vertical

- Dividing a table into smaller tables with subset of data
- Each partition contains a **different set of columns, and same number of rows possible**
- Used to split rarely accessed or wide columns

Partitioning strategies

Range partitioning:

- Created based on a certain range of the partition key (PK). E.g. Employees with IDs from 1 to 10 000 in one partition and from 10 001 to 20 000 in another

Hash partitioning:

- Applying a hash function to some attribute which results in a partition number. Ensures a **more balanced distribution of data**.

List partitioning:

- DBA explicitly maps each partition to a list of discrete values of the partition key.

Partitioning - Challenges

- Data distribution to be balanced, otherwise we get what is called '**hot partitions**', т.e. всички данни отиват в един партитишен, което не е добре - няма да бъде възможно да се възползваме от перформънса. Например ако primary key е first_name и има много Антоновци да речем.
- Data locality is crucial and operations involving multiple partitions can be slow
- Partitioning adds complexity and makes Database management and tuning harder.

Schema-less Design

- Traditional relational databases need to have the schema defined in advance and adhered (придържам / спазвам) strictly
- NoSQL provide Schema-less or schema-flexible models
- **Structure of data is not predefined** - т.e. можем да работим без схема в самото начало/ не е необходимо и да указваме тип на данните/на колоните
- **We can change the structure of the data on the fly**

Benefits:

- **Flexibility:** Schema-less models are highly flexible, enabling quick adaptation to changes.
- **Speed and Efficiency:** The lack of a rigid(твърда) schema allows for faster development and easier scalability
- **Complex and Hierarchical Data:** Schema-less models can easily store and manage complex, nested, or hierarchical data, which can be difficult to represent in a rigid schema

Challenges:

- **Inconsistency:** The flexibility of a schema-less model can lead to data inconsistency if not managed correctly.
- **Lack of Data Integrity:** In a rigid/constant schema, constraints can be applied to ensure data integrity. In a schema-less model, this responsibility often moves to the application layer. Т.e. ще се наложи в нашето Java приложение да си правим ръчни проверки за null value, за тип на колоната, и т.н. тъй като базата може да ни върне нещо което не би следвало да го обработваме. (местим един вид логика/функционалност от базата към външно приложение)
- **Query Complexity:** Without a consistent schema, querying data can become more complex. Например липсват JOIN в доста от NoSQL езиците, правим по-скъпи операции, и т.н.

NoSQL databases and implementation of schema-less design:

- **Document Databases** store data in a flexible, JSON-like format. Each document can have its own unique structure. The schema can be changed in real-time, allowing fields to be added or removed as necessary.

- **Key-Value Stores** have an entirely schema-less design. The value associated with a key can be of any type, and there's no structure enforced on the values
- **Wide Column Stores** do have a schema, but it is flexible. You can easily add or remove columns in a column family without having to modify the rest of the data
- **Graph Databases** also have flexible schemas. While they do enforce a certain level of structure (with nodes and relationships), the properties of those nodes and relationships can be very flexible

Scaling

Types of scaling:

- **Vertical Scaling**, or “**scaling up**”, involves increasing the capacity of a single machine. You can do this by adding more powerful CPUs, adding more RAM, increasing disk storage, or improving network capacity
- **Horizontal Scaling**, or “**scaling out**”, involves adding more machines to the system and distributing the load among them. This is the approach typically employed by NoSQL databases.

Normalization

Нормализацията е контра продуктивна при NoSQL. Данните ни трябва да не са нормализирани.

Можем да имам всички данни в една таблица, и вече примерно в Java или JS приложението да си ги обработваме.

Нормализацията е свързана също така и с JOIN операции, които дори и да ги има в NoSQL то са доста скъпи. Т.е. няма JOIN-ове, значи няма и нормализация.

4. Types of NoSQL databases

Key-Value

- Store data as a collection of key-value pairs, similar to a Map
 - Key is unique identifier
 - The value can be anything: number, string, **JSON**, BLOB, etc
- Example: **Redis**, **DynamoDB** - използват само/главно key-value
- Use cases:
 - Constant, fast, read/write operations, e.g. caching, session management, shopping carts, etc

Limitations:

- No complex querying
- No relationships
- No joining
- Need to handle data structure at application level

Examples:

Key-Value Example DynamoDB

```
{  
  "TableName": "Employee",  
  "KeySchema": [  
    {  
      "AttributeName": "EmployeeId",  
      "KeyType": "HASH"  
    }  
  ],  
  "AttributeDefinitions": [  
    {  
      "AttributeName": "EmployeeId",  
      "AttributeType": "N"  
    }  
  ]  
}
```

```
1  {  
2    "TableName": "Employee",  
3    "KeySchema": [  
4      {  
5        "AttributeName": "EmployeeId",  
6        "KeyType": "HASH"  
7      }  
8    ],  
9    "AttributeDefinitions": [  
10      {  
11        "AttributeName": "EmployeeId",  
12        "AttributeType": "N"  
13      }  
14    }  
15 }
```

CLI(Command Line Interface) of DynamoDB via AWS

```
sparas@Spars-MacBook-Pro docker % aws --endpoint-url=http://localhost:4566 --region=us-east-1 dynamodb put-item --table-name Employees --item "{\"EmployeeId\": {\"N\": \"1234\"}, \"Title\": {\"S\": \"Book 1234 Title\"}, \"ISBN\": {\"S\": \"111-111111111\"}, \"Authors\": {\"L\": [ {\"S\": \"John\"} ], \"Price\": {\"N\": \"20.2\"}, \"Dimensions\": {\"S\": \"8.5 x 11.0 x 0.5\"}, \"PageCount\": {\"N\": \"757\"}, \"InPublication\": {\"BOOL\": true}, \"ProductCategory\": {\"S\": \"Book\"}}}"  
usage: aws [options] <command> <subcommand> [<subcommand> ...] [parameters]  
To see help text, you can run:  
  
aws help  
aws <command> help  
aws <command> <subcommand> help  
  
Unknown options: Title}, ISBN: {S: 111-111111111}, Authors: {L: [ {S: John} ], Price: {N: 20.2 },Dimensions: {S: 8.5, x, 11.0, x, 0.5 }, PageCount: {N: 757},InPublication: {BOOL: true},P  
roductCategory: {S: Book}}, 1234
```

Document Databases

Subcategory of key-value

Store data in a semi-structured format (often **JSON**)

- Allow for nested data
- Each document has a unique key identifier and a value

Examples: **MongoDB** and CouchDB and **ElasticSearch**(<https://www.elastic.co/>)

Use cases:

- Content management systems
- Real-time analytics
- IoT

Limitations:

- Complex transactions
- Data integrity across documents (joins and complex multi-document transactions)

В MongoDB липсва използването на join-ове!

Document Example MongoDB

```
① {
②   "_id": "60df0eefdcba543afcdf7980",
③   "FirstName": "John",
④   "LastName": "Doe",
⑤   "DepartmentId":
⑥     "60df0eefdcba543afcdf7982",
⑦   "DateOfBirth":
⑧     "1990-01-01T00:00:00Z",
⑨   "EmployeeCardId":
⑩     "60df0eefdcba543afcdf7983"
⑪ }
```

```
{  
  "_id": "60df0eefdcba543afcdf7980",  
  "FirstName": "John",  
  "LastName": "Doe",  
  "DepartmentId": "60df0eefdcba543afcdf7982",  
  "DateOfBirth": "1990-01-01T00:00:00Z",  
  "EmployeeCardId": "60df0eefdcba543afcdf7983"  
}
```

CLI of MongoDB

```
Atlas atlas-pzobd1-shard-0 [primary] test> use lectures
switched to db lectures
Atlas atlas-pzobd1-shard-0 [primary] lectures> db.employees.insertOne({
...   firstName: "John",
...   lastName: "Doe",
...   dateOfBirth: new Date("1980-01-15"),
...   department: {
...     departmentId: "dep1",
...     departmentName: "IT"
...   },
...   employeeCard: {
...     employeeCardId: "card101",
...     expirationDate: new Date("2025-12-31")
...   }
... });
{
  acknowledged: true,
  insertedId: ObjectId("64c0ba4c7d47c286371954d7")
```



```
Atlas atlas-pzobd1-shard-0 [primary] lectures> db.employees.find({ _id: ObjectId("64c0ba4c7d47c286371954d7") })
[  
  {  
    _id: ObjectId("64c0ba4c7d47c286371954d7"),  
    firstName: 'John',  
    lastName: 'Doe',  
    dateOfBirth: ISODate("1980-01-15T00:00:00.000Z"),  
    department: { departmentId: 'dep1', departmentName: 'IT' },  
    employeeCard: {  
      employeeCardId: 'card101',  
      expirationDate: ISODate("2025-12-31T00:00:00.000Z")  
    }  
  }  
]
```

```
Atlas atlas-pzobd1-shard-0 [primary] lectures> db.employees.updateOne({ _id: ObjectId("64c0ba4c7d47c286371954d7") }, { $set: { firstName: "Spass" } });
{
  acknowledged: true,
  insertedId: null,
  matchedCount: 1,
  modifiedCount: 1,
  upsertedCount: 0
}
```

```
Atlas atlas-pzobd1-shard-0 [primary] lectures> db.employees.deleteOne({_id: ObjectId("64c0ba4c7d47c286371954c"),
{ acknowledged: true, deletedCount: 1 }
Atlas atlas-pzobd1-shard-0 [primary] lectures> []
```

След като сме изтрили един запис, само .find() ни намира всички записи в момента

```
Atlas atlas-pzobd1-shard-0 [primary] lectures> db.employees.find();
[
  {
    _id: ObjectId("64c0f86f7d47c286371954d8"),
    firstName: 'Jake',
    lastName: 'Doe',
    dateOfBirth: ISODate("1980-01-15T00:00:00.000Z"),
    department: { departmentId: 'dep1', departmentName: 'IT' },
    employeeCard: {
      employeeCardId: 'card101',
      expirationDate: ISODate("2025-12-31T00:00:00.000Z")
    }
  },
  {
    _id: ObjectId("64c0f9347d47c286371954d9"),
    firstName: 'Jake',
    middleName: 'Peterson',
    lastName: 'Doe',
    dateOfBirth: ISODate("1980-01-15T00:00:00.000Z"),
    department: { departmentId: 'dep1', departmentName: 'IT' },
    employeeCard: {
      employeeCardId: 'card101',
      expirationDate: ISODate("2025-12-31T00:00:00.000Z")
    }
  }
]
```

Wide Column Databases

- Store data in columns rather than stores.
 - Each column family can contain a virtually unlimited number of columns - **колона в колоната**
- Examples: Cassandra and Google's Bigtable
- Use cases:
 - Querying and analysis of large distributed data sets
- **Limitations**
 - Complex transactions
 - Data integrity across documents (joins or complex multi-document transactions)

Graph Databases

Build around the concept of Nodes and Edges

- Nodes represent entities
- Edges represent relationships between nodes

Examples: Neo4j and Amazon Neptune

Use cases:

- Social networks
- Recommendation engines

Limitations:

- System with high update rates
- Performance can degrade with data volume

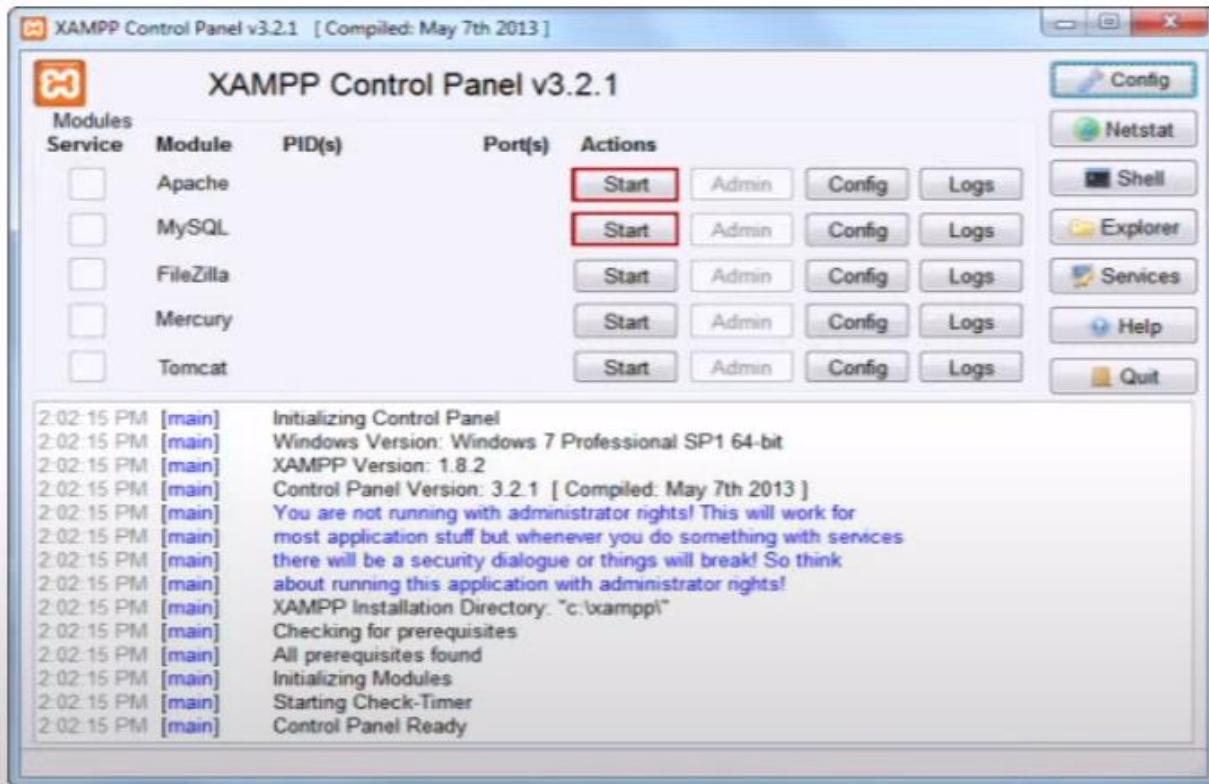
X. Other tools for using

Ако искаме, можем да ползваме и XAMPP или MariaDB или Heidi SQL вместо MySQL

Using XAMPP Control Panel

By B Lingafelter Feb 13, 2014 testserver, workflow

1. Open the XAMPP Control Panel. If you don't have a Desktop or Quick Launch icon, click **Start** > **All Programs** > **XAMPP Control Panel**.



2. Click **Start** button next to Apache. Note: Do NOT mark the Service check box.