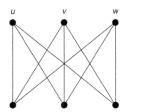
Abstraction of Graph Theory

Wael Ali

Section (2)

- A simple graph G consists of a nonempty finite set V(G) of elements called vertices (or nodes), and a finite set E(G)of distinct unordered pairs of distinct elements of V(G) called edges.
- A graph G consists of a non-empty finite set V(G) of elements called vertices, and a finite family E(G) of unordered pairs of (not necessarily distinct) elements of V(G) called edges.
- Two graphs G1 and G2 are <u>isomorphic</u> if there is a one-one correspondence between the vertices of G_x and those of G2 such that the number of edges joining any two vertices of G_x is equal to the number of edges joining the corresponding vertices of G2.



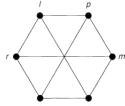


Figure 1: Two simple graphs that are isomorphic to each other

- A graph is **connected** if it cannot be expressed as the union of two graphs, and **disconnected**.
- Two vertices v and w of a graph G are adjacent if there is an edge vw joining them, and the vertices v and w are then incident with such an edge. Similarly, two distinct edges e and/are adjacent if they have a vertex in common.

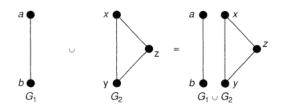
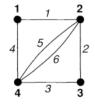
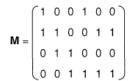


Figure 2: Each of G_1 and G_2 is a component of $G_1 \cup G_2$

- The <u>degree</u> of a vertex v of G is the number of edges incident with v, and is written deg(v). Loop-edge increase nodedegree by 2. vertex of degree 0 is an <u>isolated vertex</u> and a vertex of degree 1 is an <u>end-vertex</u>.
- Handshaking lemma; in any graph the sum of all the vertex-degrees is an even number in fact, twice the number of edges, since each edge contributes exactly 2 to the sum.
- A <u>subgraph</u> of a graph G is a graph, each of whose vertices belongs to V(G) and each of whose edges belongs to E(G).
- Matrix representation one way to represent graph is by its adjacency matrix A, and its incidence matrix M as follows;

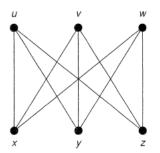


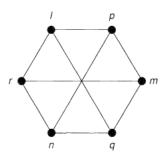
$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 2 \\ 0 & 1 & 0 & 1 \\ 1 & 2 & 1 & 0 \end{pmatrix}$$



Exercise 2

2a Write down the vertex-set and edge-set of each graph in Fig 2.5





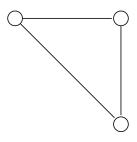
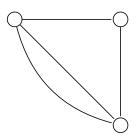


Figure 3: (i) a simple graph



The first graph G_1 is $(V(G_1), E(G_1))$

$$V(G_1) = \{x, y, z, u, v, w\}$$

$$E(G_1) = \{\{x, u\}, \{x, v\}, \{x, w\}, \{y, u\}, \{y, v\}, \{y, w\}, \{z, u\}, \{z, v\}, \{z, w\}\}$$

The second graph G_2 is $(V(G_2), E(G_2))$

$$V(G_2) = \{n, m, q, r, l, p\}$$

$$E(G_2) = \{\{n, r\}, \{n, p\}, \{n, q\}, \{m, r\}, \{m, p\}, \{m, q\}, \{l, r\}, \{l, p\}, \{l, q\}\}$$

Figure 4: (ii) a non-simple graph with no loops

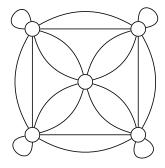


Figure 5: (iii) a 5-vertices and 8-degress each

2b Draw;

- (i) a simple graph.
- (ii) a non-simple graph with no loops.
- (iii) a non-simple graph with no multiple edges, each having 5 vertices each having 5 vertices and 8 edges.

(ii) How does the answer to part (i) changed if the degrees are 5, 5, 4, 3, 3,2?

2c Draw;

(i) Draw a graph on six vertices whose degrees are 5,5,5,5,3,3; does there exist a simple graph with these degrees?

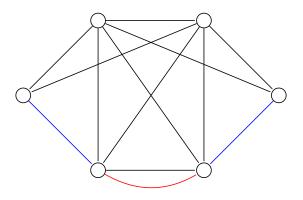


Figure 6: (i) a non-simple graph with with 6-vertices with degrees [5,5,5,5,3,3]

There isn't a simple graph with last mentioned degrees for a 6-vertices graph. But it we just remove the red-arc and one of the blue-arcs, we then get a simple graph as shown in the next figure.

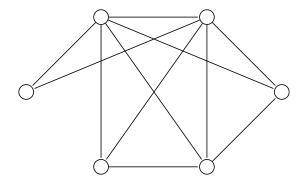
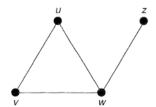


Figure 7: (i) a simple graph with with 6-vertices with degrees [5,5,4,3,3,2]

(2d) Verify that handshaking lemma is hold for figure 2.1

As we can see in the last figure, the sum



of all vertex-degrees is 2 + 2 + 3 + 1 = 8 which is an even number.

- (2f) (i) By suitably lettering the vertices, show that the two graphs in Fig. 2.20 are isomorphic.
 - (ii) Explain why two graphs in Fig. 2.21 are not isomorphic.
 - (i) As shown in the following figure ?? the two graphs are labeled with the same letters in a way to emphasize that they are both isomorphic to each other.
 - (ii) As shown in the following figure ??, we cannot find the red part of the first graph as a subset in the second graph.

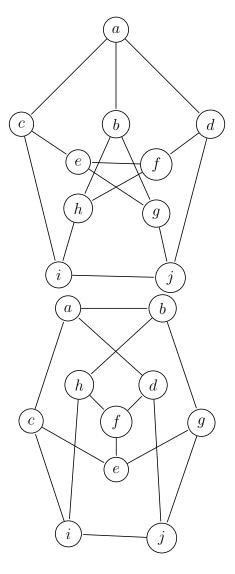


Figure 8: (i)

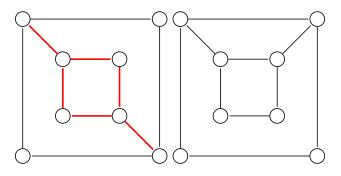
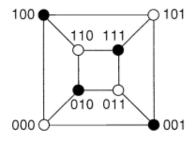


Figure 9: (ii) not isomorphic graphs

Section (3)

- Null graphs, a graph whose edge-set is empty. Complete graph, a simple graph in which each pair of distinct vertices are adjacent, in this case k-vertex must has n(n-1)/2 degree as n the total number of vertices.
- Regular graph, a graph in which each vertex has the same degree. Platonic graphs, formed from the vertices and edges of the five regular (Platonic) solids the tetrahedron, octahedron, cube.
- Bipartite graph, If the vertex set of a graph G can be split into two disjoint sets A and B so that each edge of G joins a vertex of A and a vertex of B. A complete bipartite graph is a bipartite graph in which each vertex in A is joined to each vertex in B by just one edge.
- <u>Cubes</u>, the k-cube Q_k is the graph whose vertices correspond to the sequences (a_1, a_2, \dots, a_k) , where each $a_i = 0$ or 1, and whose edges join those sequences that differ in just one place. You should check that Q_k has 2^k vertices and $k2^{k-1}$ edges, and is regular of degree k.



• Complement of a simple graph, if G is a simple graph with vertex set V(G), its complement \overline{G} is the simple graph with vertex set V(G) in which two vertices are adjacent if and only if they are not adjacent in G.

- (i) the null graph N_5 .
- (ii) the complete graph K_6 .
- (iii) the complete bipartite graph $K_{2,4}$.
- (iv) the union of $K_{1,3}$ and W_4 .
- (v) the complement of the cycle graph C_4 .

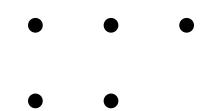


Figure 10: (i) Null graph N_5

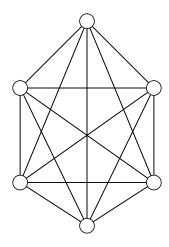


Figure 11: (ii) Complete graph K_6

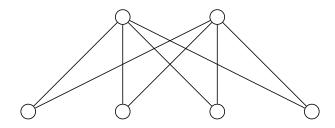


Figure 12: (iii) Complete bipartite graph $K_{2,4}$

Exercise 3

(3a) Draw the following graphs:

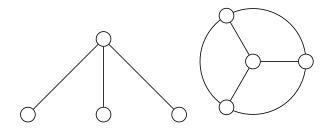


Figure 13: (iv) Union of $K_{1,3}$ and W_4

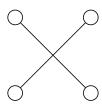


Figure 14: (iv) Complement of cycle graph C_4