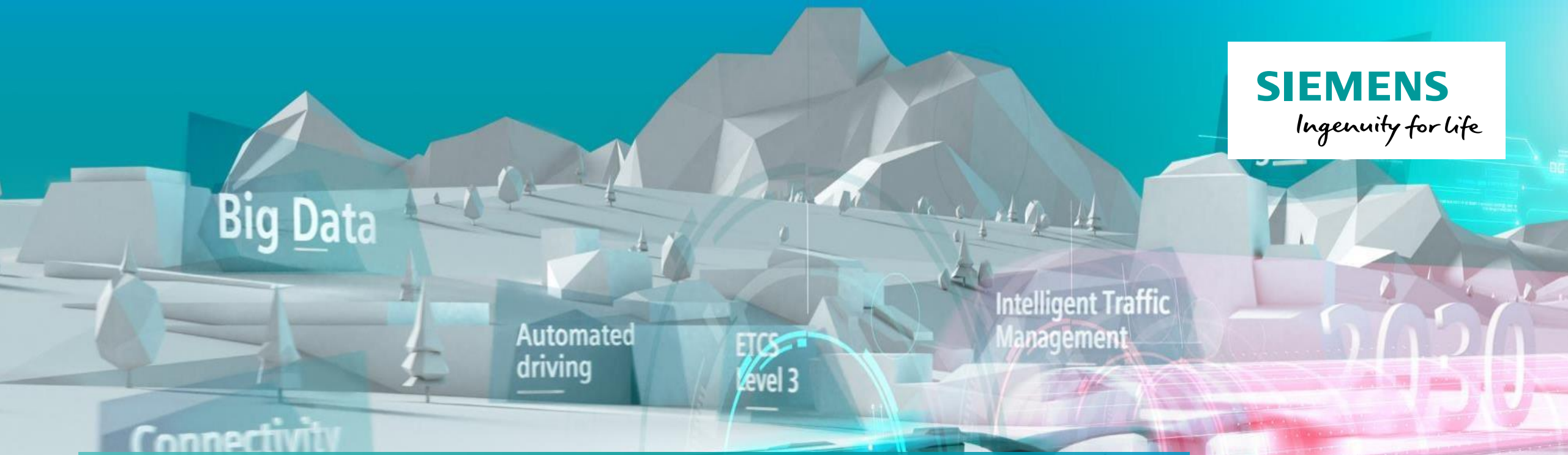

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Formal Methods From Theory towards Practice

Simon Chadwick, Tom Werner
Siemens Mobility, Chippenham, UK

Formal Verification – The Journey from Theory towards Practice
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Formal Verification – The Journey from Theory towards Practice
The Idea

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We can apply Formal Verification to Interlocking Logic

- Safety: Formal Verification of Safety Properties would improve safety
- Efficiency: Formal Verification of Safety Properties could reduce testing requirements

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Background

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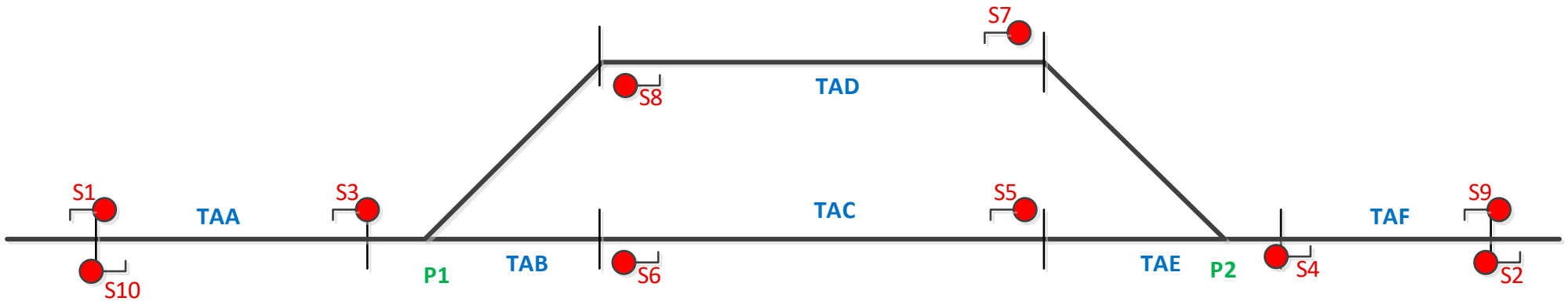
Railway signalling systems typically have a component called “Interlocking”.
Consider the passing loop below:



Some combinations of signals, points and track section occupancy are OK
Some combinations of signals, points and track section occupancy are not OK

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Background

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Different technologies have been used through the history of railway signalling.

- Mechanical interlocking
- Relay interlocking
- Electronic interlocking

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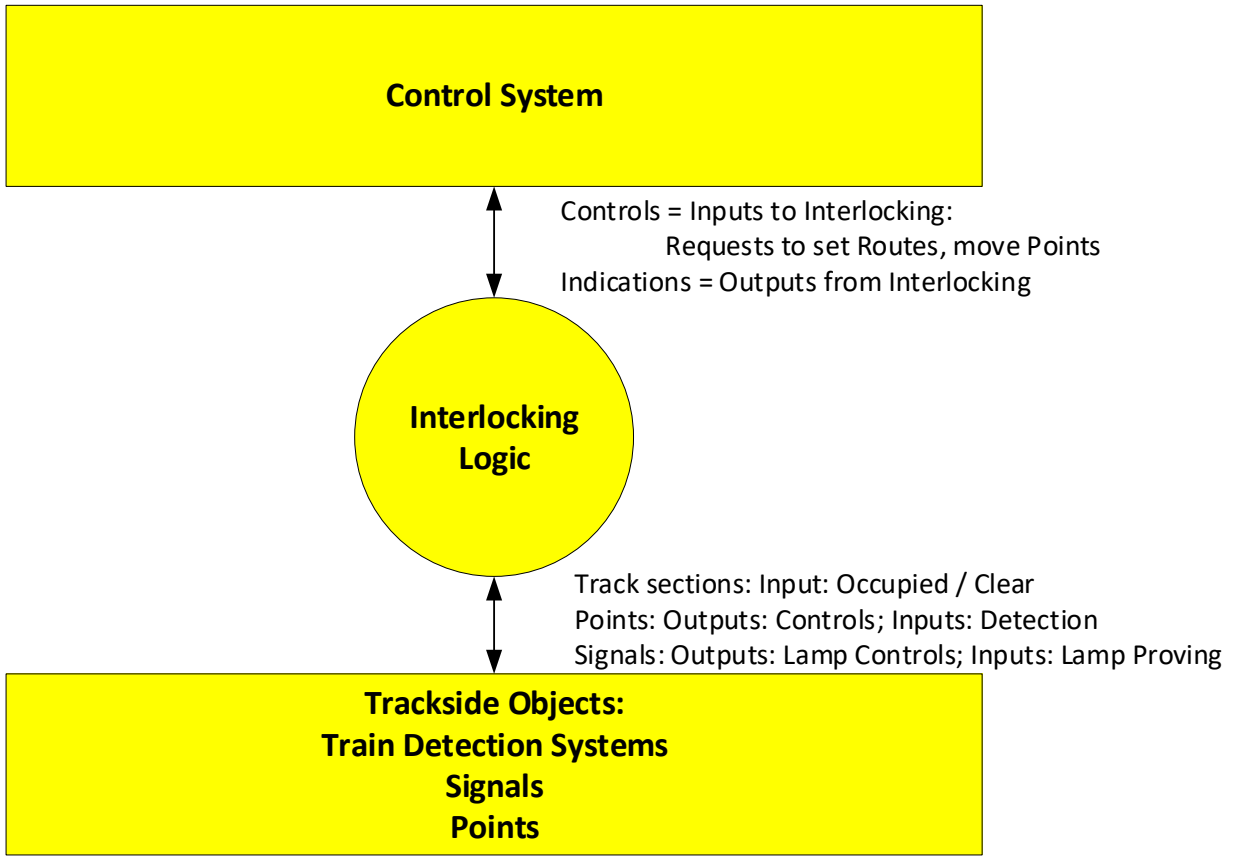
An abstraction of an electronic interlocking is:

Reminder:

- Combinations of inputs matter
- There are stored states, so sequences of combinations of inputs matter
- There are timers, so the durations of each step in a sequence of combinations of inputs matter



Some combinations of inputs and outputs are OK
Some combinations of inputs and outputs are not OK

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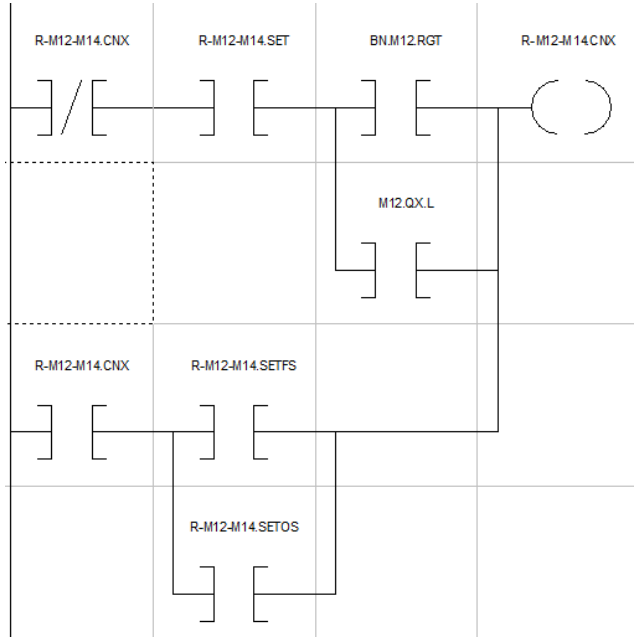
Within Siemens, we have a specific interlocking product called “WESTRACE”.

It is essentially an application-specific PLC.

It is programmed with a simplified form of Ladder Logic.



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The Idea Revisited

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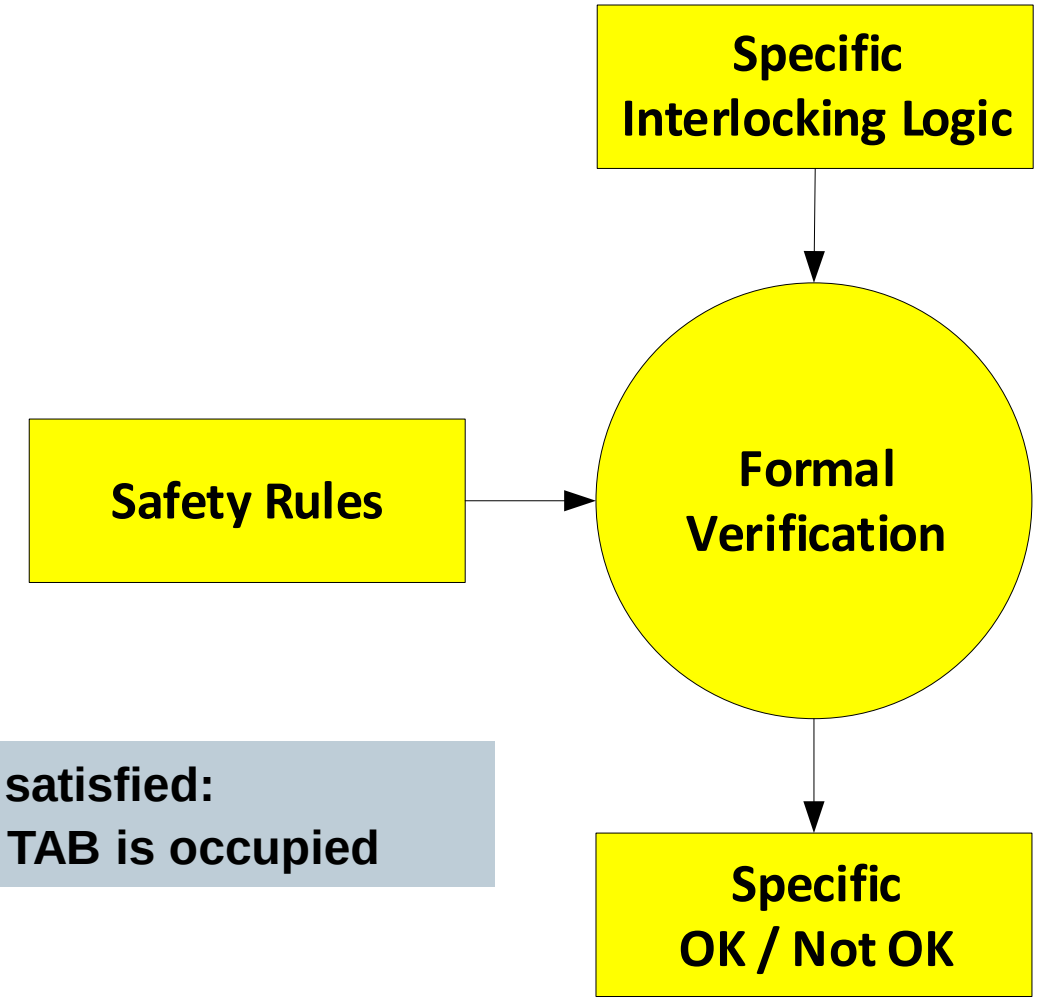
Now we can refine the idea:



Safety Rules define requirements which must be satisfied:
e.g. Point P1 must not be moved if Track Section TAB is occupied

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The Idea Revisited

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An example safety rule:
Point P1 must not be moved if Track Section TAB is occupied

Point ‘P1’ commanded normal – “P1.NL”
Point ‘P1’ commanded reverse – “P1.RL”

$$(("\mathbf{P1.NL\_0}" \wedge "\mathbf{P1.RL\_1}") \vee ("\mathbf{P1.RL\_0}" \wedge "\mathbf{P1.NL\_1}")) \Rightarrow \neg ("\mathbf{TAB.OCC(IL)\_1}")$$

Point Normal to Reverse Point Reverse to Normal

If we don’t have a ‘point is moving’ state, then in order to model this property we need to be able to consider multiple states (0 / ‘current’ and 1 / ‘next’).

This works, but it is not enough…..

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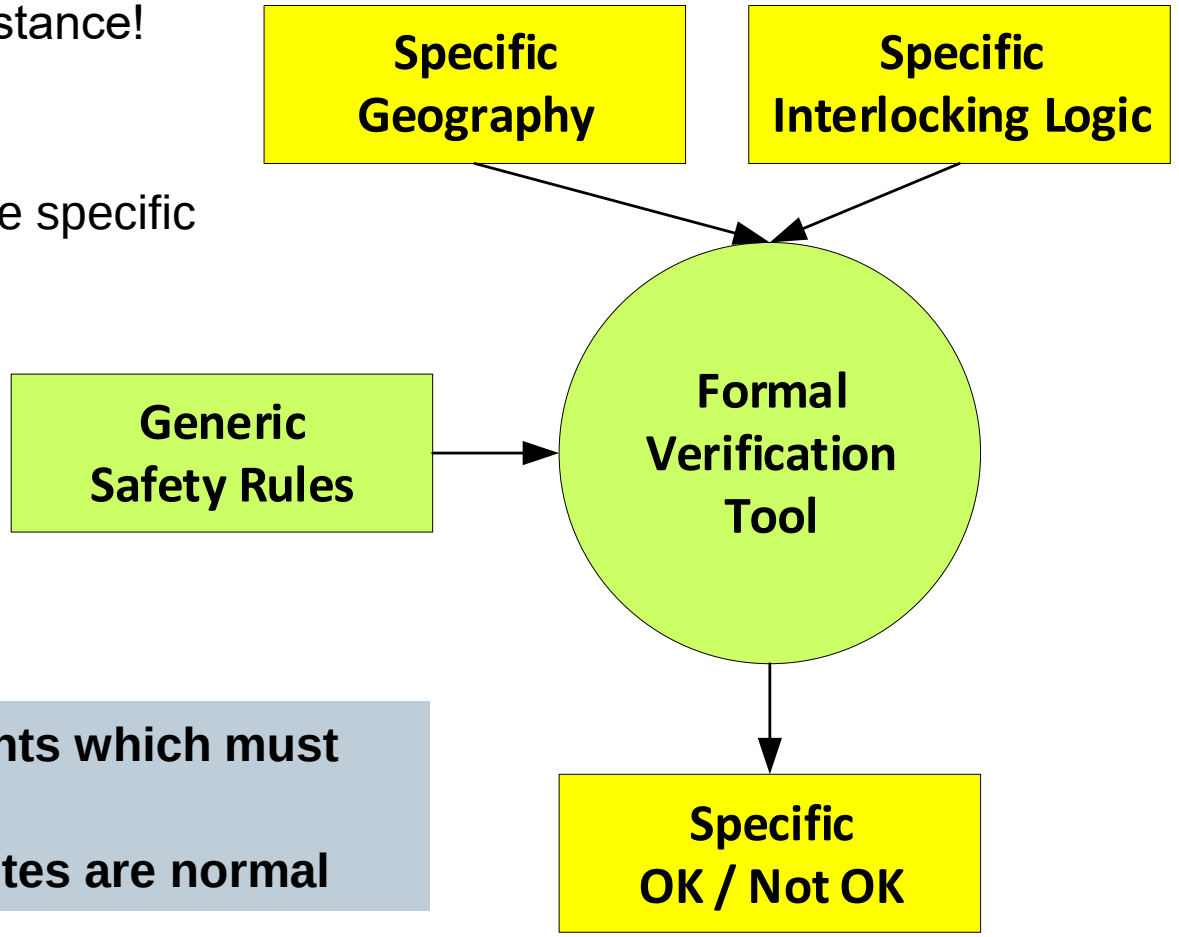
We don't want to write the Safety Rules for every instance!
We want a machine to help with that!

We need to find a machine readable way to input the specific geography



Generic Safety Rules define generic requirements which must be satisfied:
e.g. A route can only be set if all conflicting routes are normal

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Source:

2.1.3 Conflicting Routes

A route can only be set if all conflicting routes are normal.

‘Generic’ (*not specific to a track plan*)

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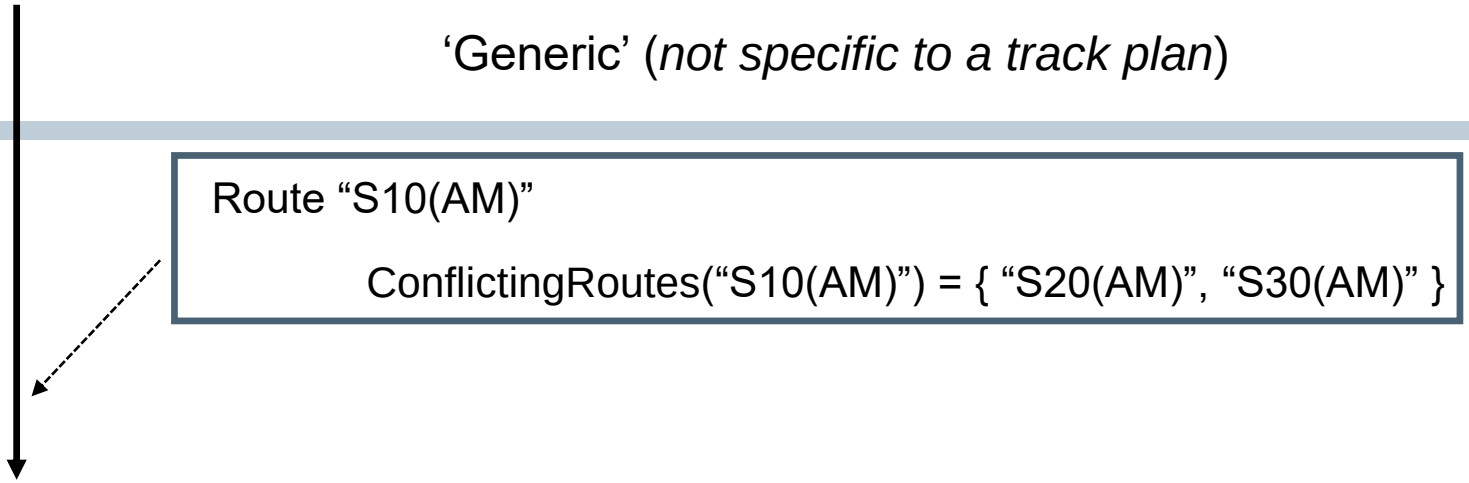
Source:

2.1.3 Conflicting Routes

A route can only be set if all conflicting routes are normal.



“S10(AM)” can only be set if “S20(AM)” and “S30(AM)” are normal.

Formal Verification – The Journey from Theory towards Practice
Safety Requirements

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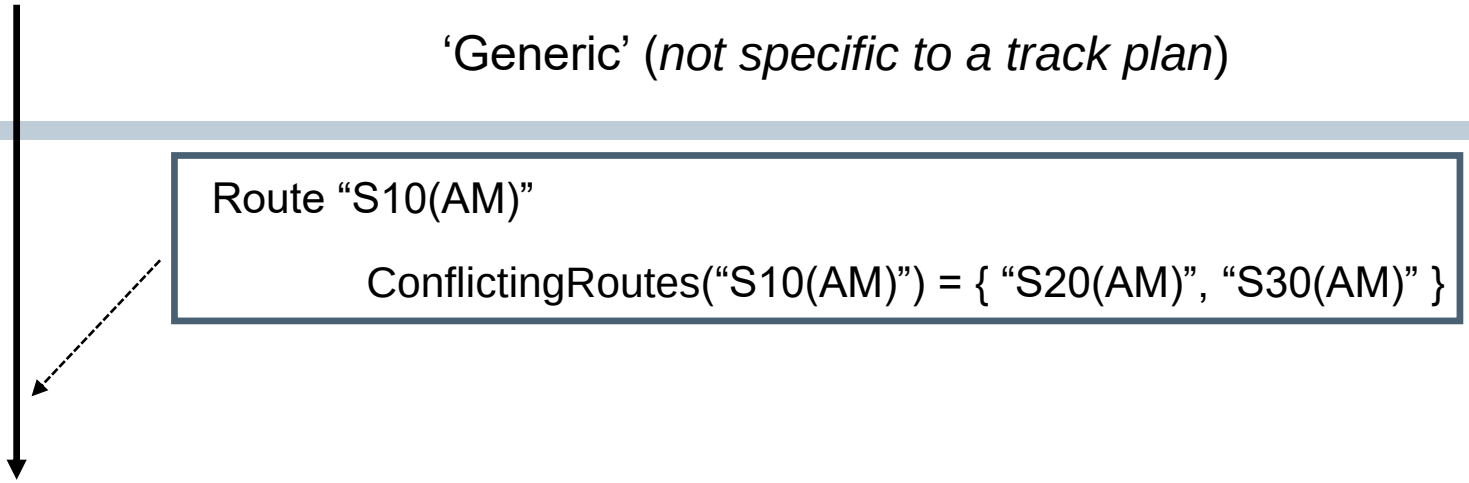
Source:

2.1.3 Conflicting Routes

A route can only be set if all conflicting routes are normal.

$$\forall\, r1 \in Route \quad routeSet\,(r1) \Rightarrow \forall\, r2 \in ConflictingRoutes\,(r1)\ routeNormal\,(r2)$$

‘Generic’ (*not specific to a track plan*)



“S10(AM)” can only be set if “S20(AM)” and “S30(AM)” are normal.

‘Concrete’ (*specific to a track plan*)

Formal Verification – The Journey from Theory towards Practice
Safety Requirements

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Source:

2.1.3 Conflicting Routes

A route can only be set if all conflicting routes are normal.

$$\forall\, r1 \in Route \quad routeSet\,(r1) \Rightarrow \forall\, r2 \in ConflictingRoutes\,(r1)\ routeNormal\,(r2)$$

‘Generic’ (*not specific to a track plan*)

Route “S10(AM)”

ConflictingRoutes(“S10(AM)”) = { “S20(AM)”, “S30(AM)” }

routeSet(“S10(AM)“) ⇒

routeNormal(“S20(AM)”) ∧ routeNormal(“S30(AM)“)

“S10(AM)” can only be set if “S20(AM)” and “S30(AM)” are normal.

“S10(AM).U” ⇒

“S20(AM).N” ∧ “S30(AM).N”

This works, but it is not enough…..

‘Concrete’ (*specific to a track plan*)

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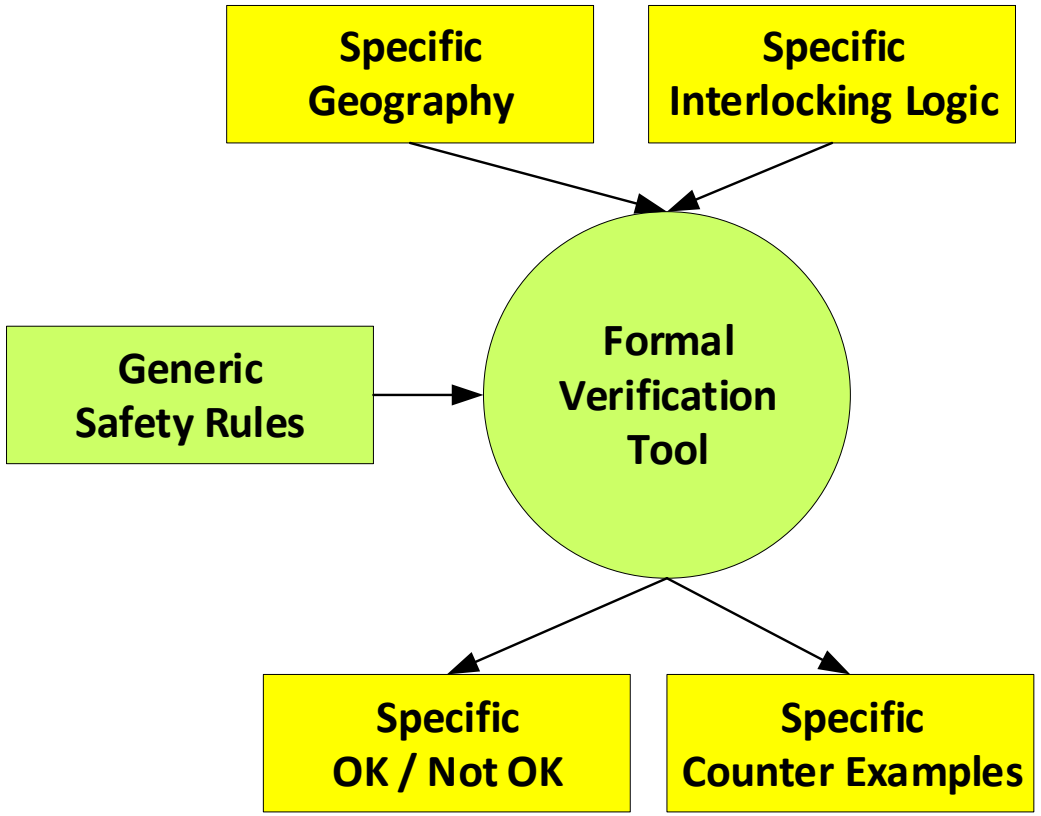
We have to be able to understand what it means if the computer says "No"!

There can be many counter examples!

Counter examples should present the steps from initialisation to the invalid state



Getting there, but it is still not enough…..

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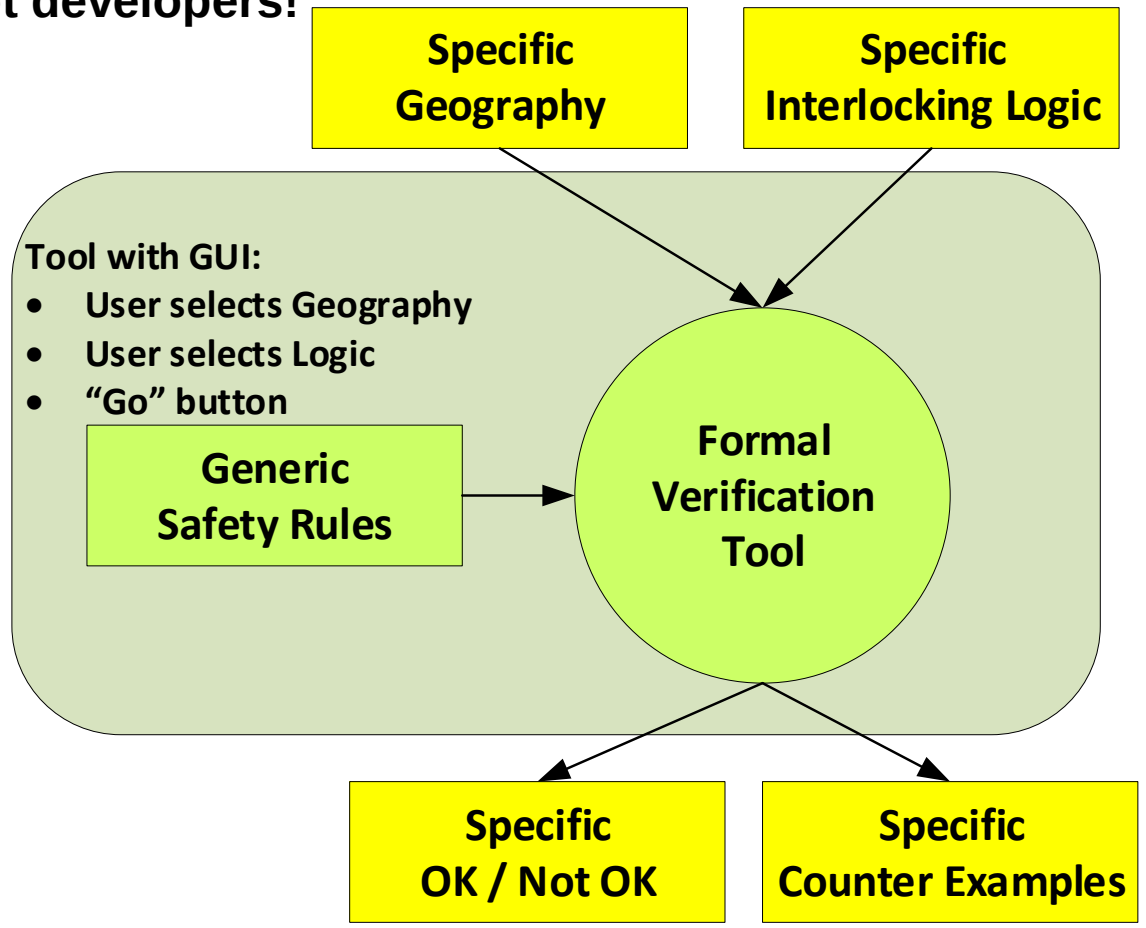
We need to provide a neat package for users, not developers!



Formal Verification – The Journey from Theory towards Practice
Stages of tool development

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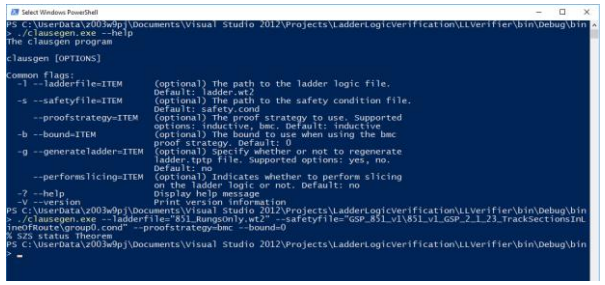
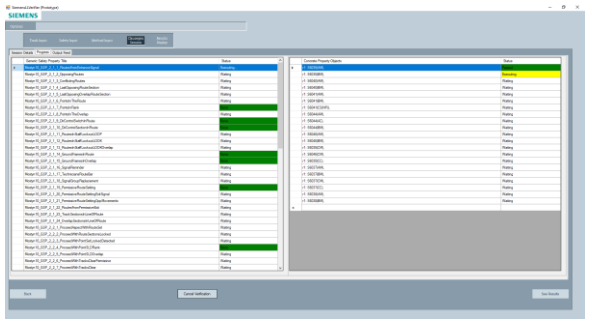
Stage 1 (Swansea - original)



- Command-line tool.
- Safety properties manually written in propositional logic for a particular track plan.

Stage 2 (Siemens - current)



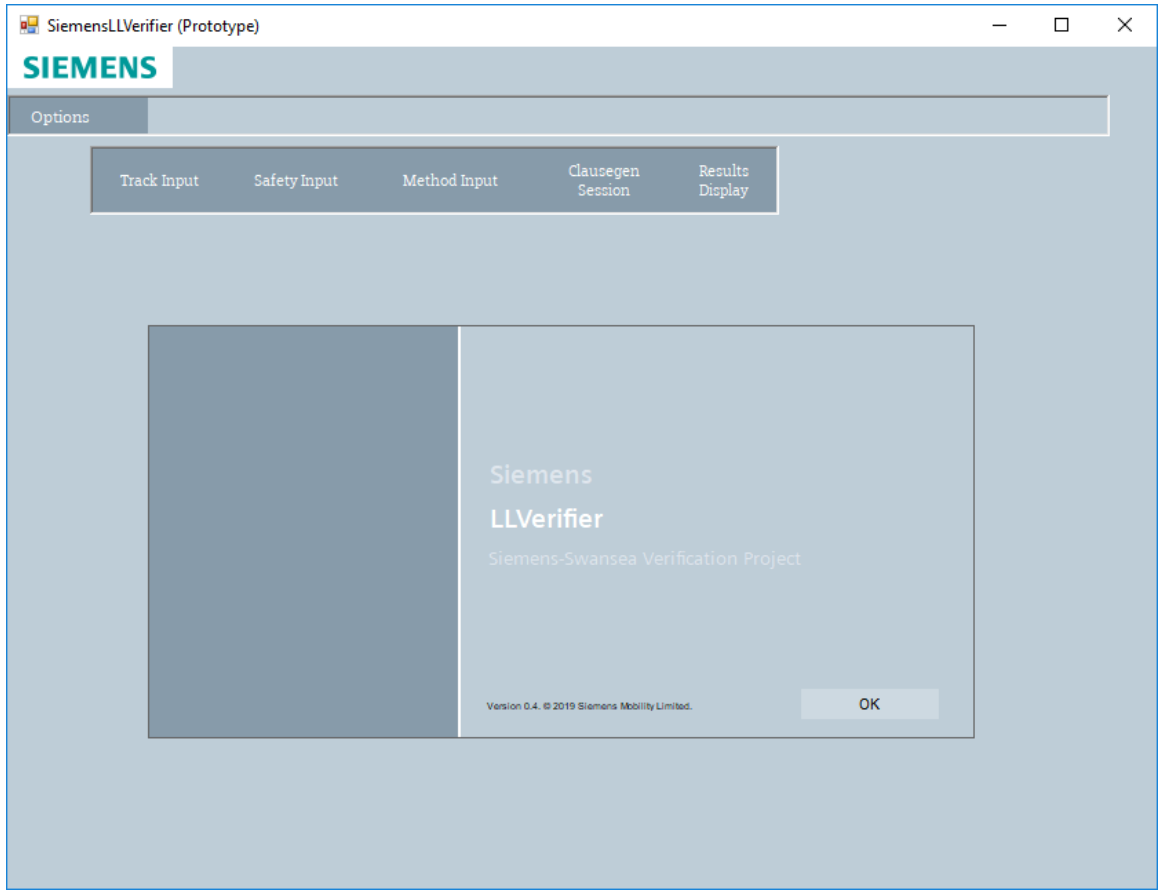
- Basic graphical front-end.
- Local windows application.
- Pre-modelled list of safety properties for particular signalling rules.
- Automatic generation of track-specific properties.

Stage 3 (Siemens - Future)

?

- Local front-end interfacing to external server for execution.
- Embedded into existing tools.
- Handling of verification outputs (reports, etc).

Formal Verification – The Journey from Theory towards Practice Demonstration



Formal Verification – The Journey from Theory towards Practice
Demonstration: Ladder logic and layout

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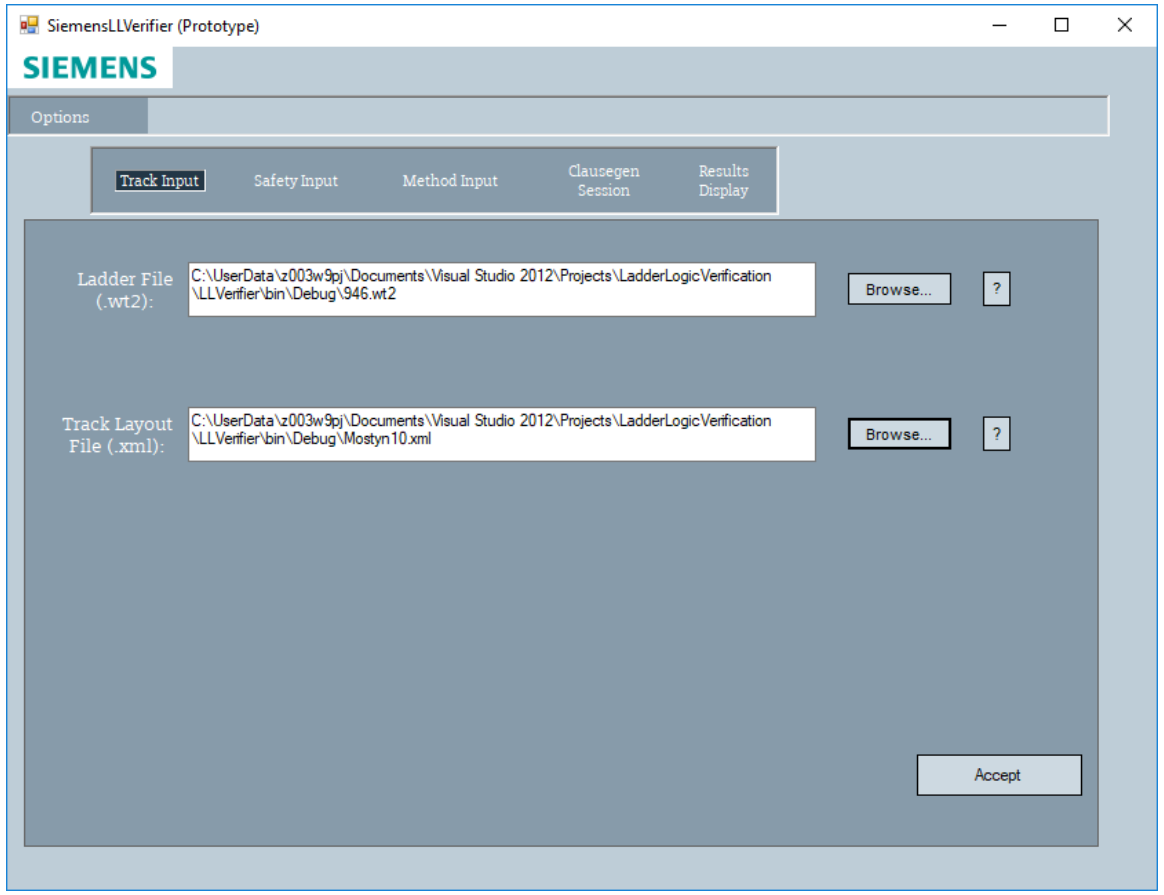

Formal Verification – The Journey from Theory towards Practice
Demonstration: Selection of safety rules

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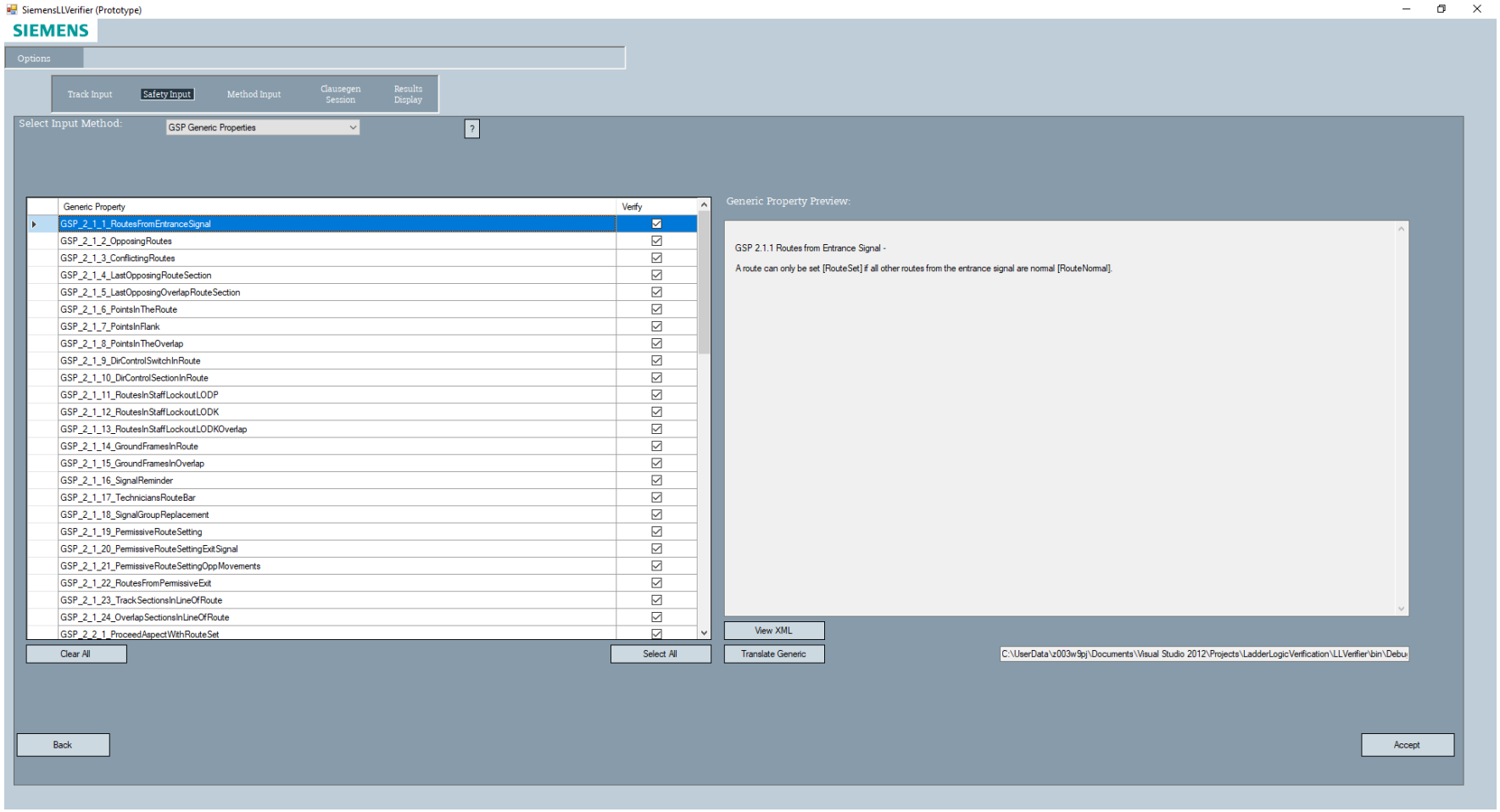

Formal Verification – The Journey from Theory towards Practice
Demonstration: Selection of verification method

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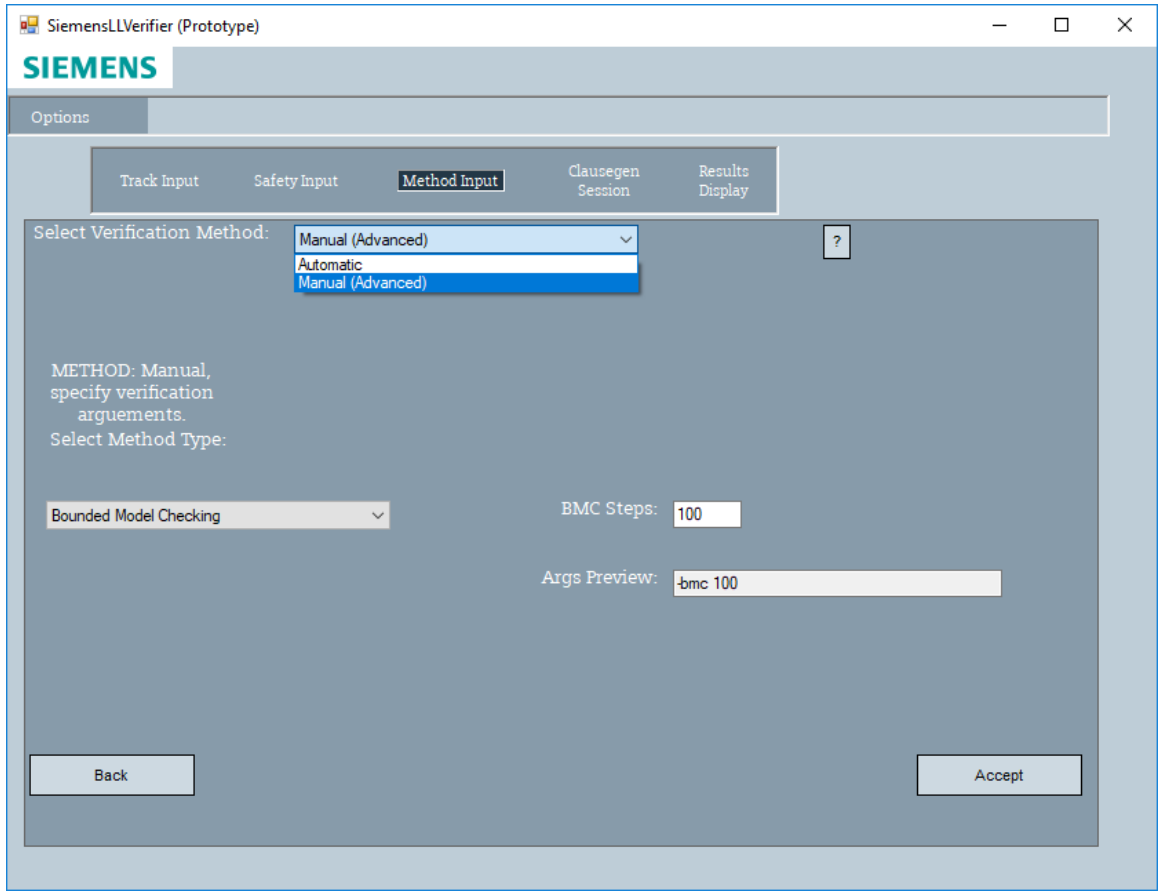
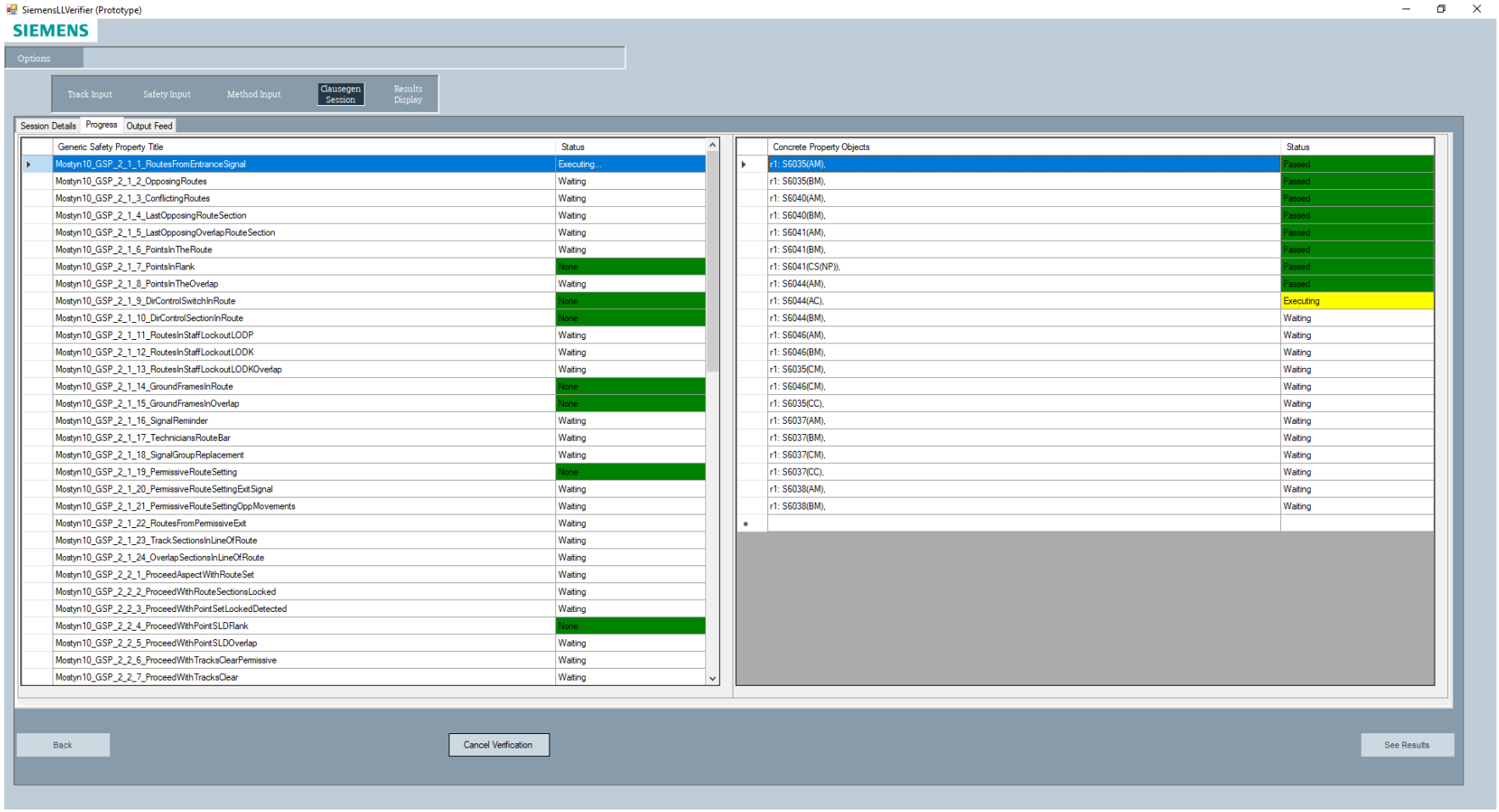


Formal Verification – The Journey from Theory towards Practice
Demonstration: Verification progress



Formal Verification – The Journey from Theory towards Practice
Demonstration: Results (1)

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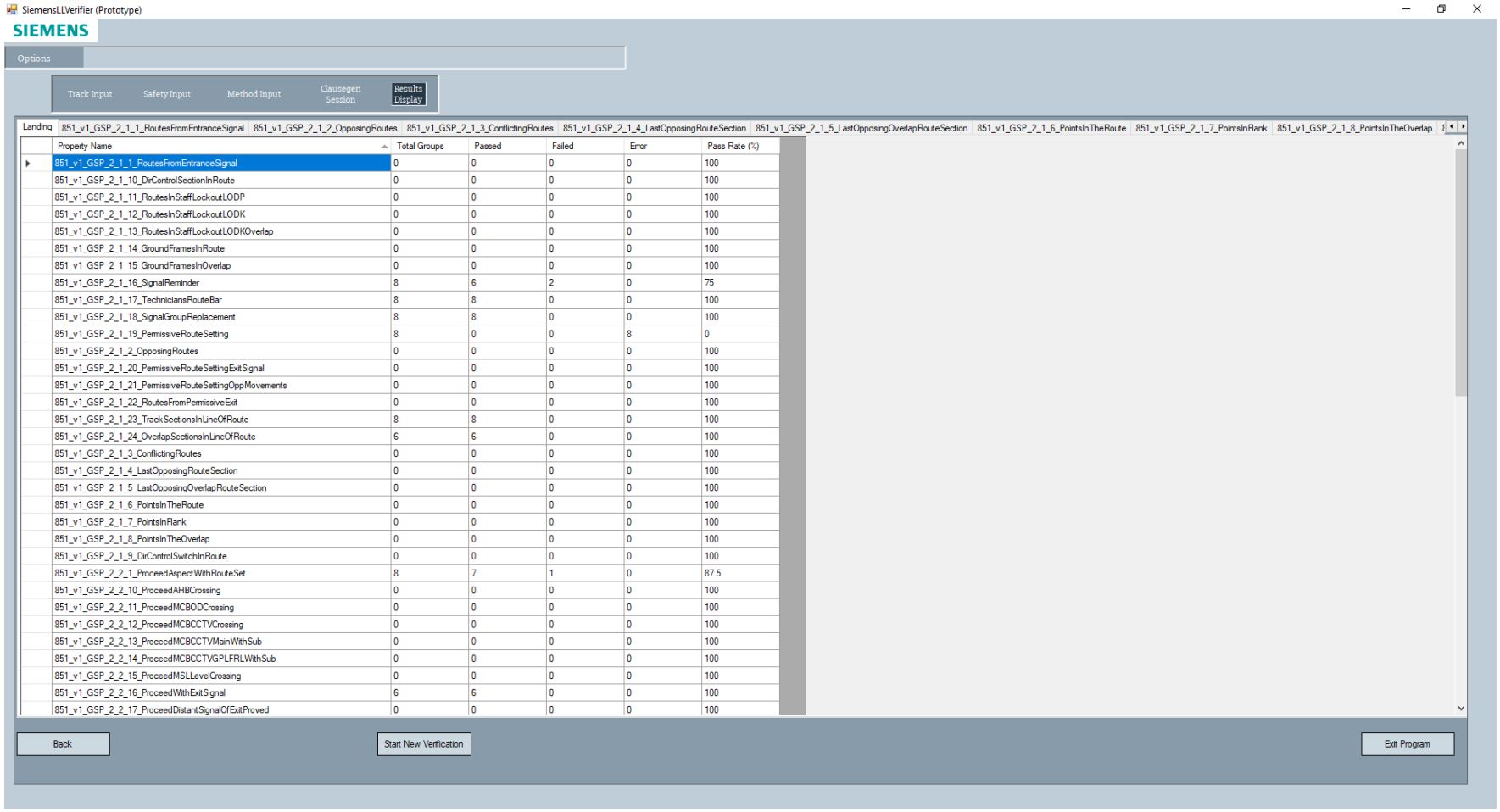

Formal Verification – The Journey from Theory towards Practice
Demonstration: Results (2)

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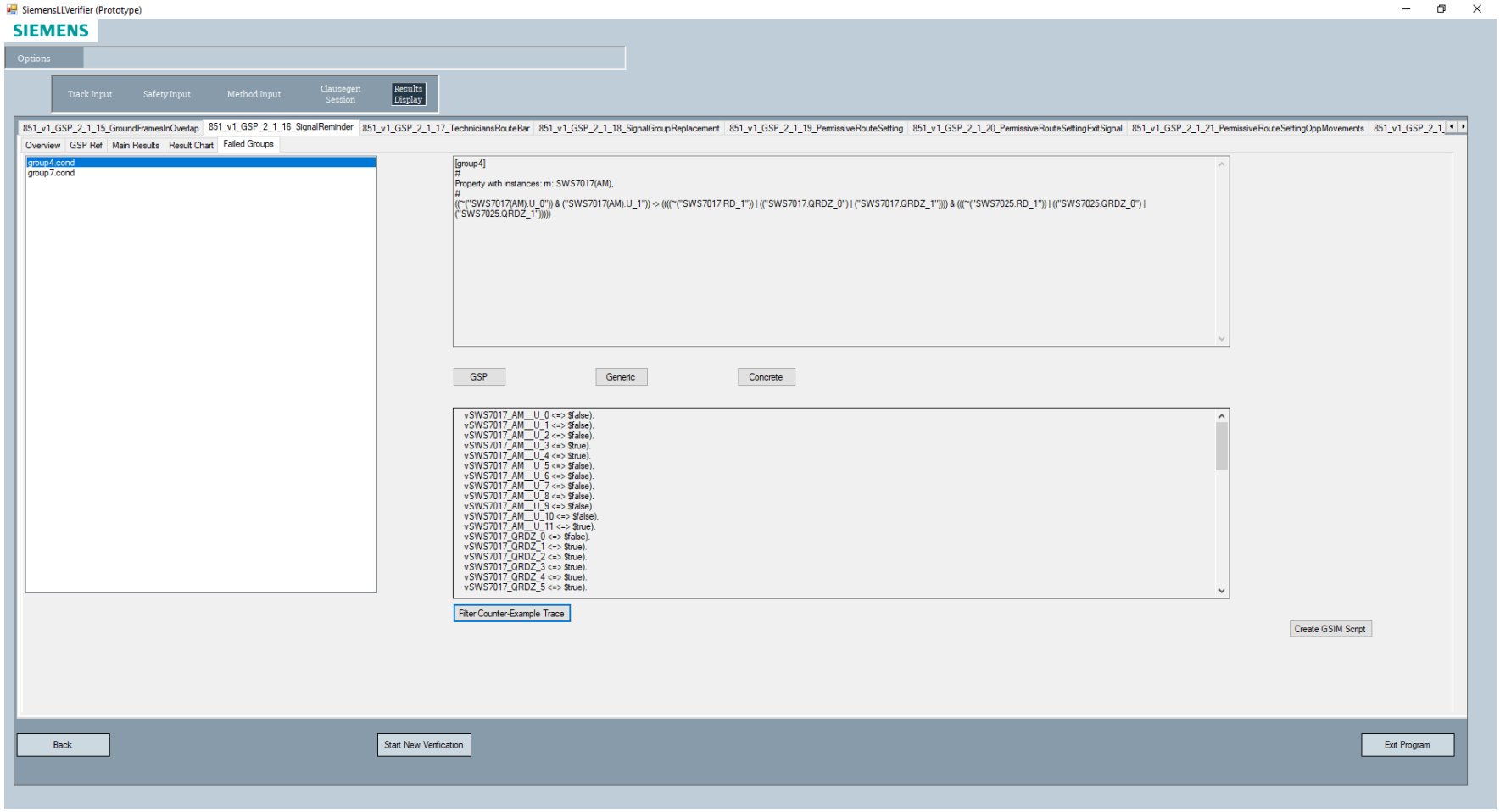

Formal Verification – The Journey from Theory towards Practice
No such thing…

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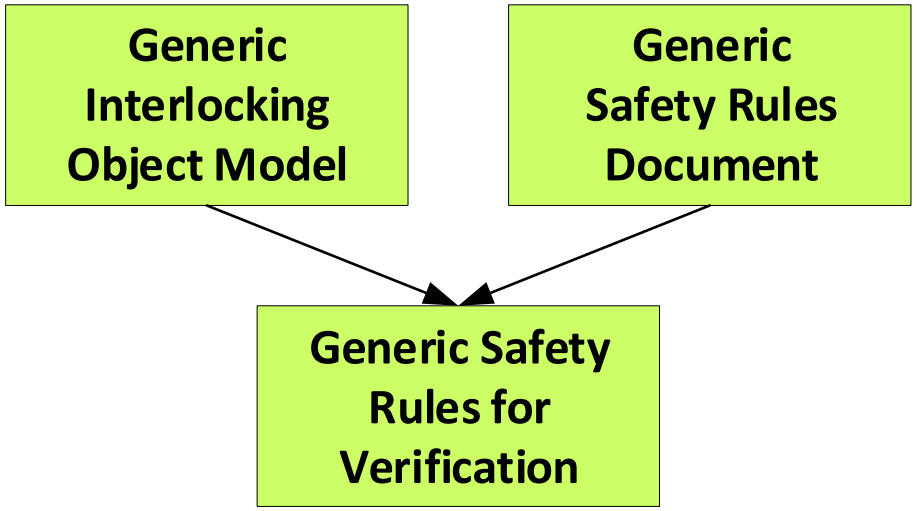
…as a free lunch!

- To set up any automation requires some effort!
- This applies to automated data generation, automated testing, formal verification
- We have done this first in document form, second in machine-readable form:



This has to be repeated for each set of signalling rules!
Some standards are available to help

Non-trivial effort!

Formal Verification – The Journey from Theory towards Practice Completeness?

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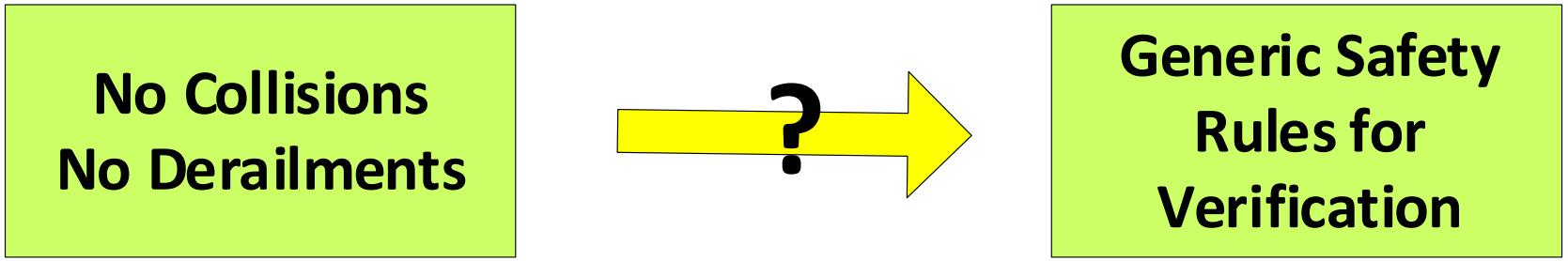
There is a completeness problem.
We have no way to demonstrate that the set of rules is complete.



Only the review by signaling experts
Review based on standards which have evolved over a long time

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Thankyou!

Simon Chadwick, Tom Werner
Siemens Mobility, Chippenham, UK