Problem 1: Calculating L2 and L3 misses for following configurations:

| Trace File | A. Inclusive | | B. Non-Inclusive Non- Exclusive | | C. Exclusive | |
|------------|--------------|-----------|------------------------------------|-----------|--------------|-----------|
| | L2 Misses | L3 Misses | L2 Misses | L3 Misses | L2 Misses | L3 Misses |
| Bzip | 5398166 | 1446388 | 5397576 | 1445846 | 5397576 | 889221 |
| GCC | 3036461 | 1373402 | 3029809 | 1366248 | 3029809 | 1242824 |
| Gromacs | 336851 | 170531 | 336724 | 170459 | 336724 | 159302 |
| h264ref | 969678 | 342146 | 965624 | 333583 | 965624 | 143681 |
| hmmer | 1743421 | 391226 | 1735322 | 376344 | 1735322 | 300046 |
| Sphinix | 8820349 | 8207362 | 8815130 | 8205144 | 8815130 | 7220776 |

- → L2 misses in Inclusive will be more than NINE and Exclusive.
 - Reason: In **inclusive** we invalidate the data from L2, if it is evicted from L3. It might be the case that evicted block is **hot block** and it results in the more misses. Whereas in **NINE** and **Exclusive** L3 miss eviction doesn't invalidate block in L2.
- **→ Exclusive policy has less L3 Misses than NINE and Inclusive** as it has more combined capacity of L2 and L3 cache due to exclusive entries in L2 and L3.
- → L2 Misses in NINE and Exclusive are same because the effect of replacement policy on L2 cache is similar in both configuration. Because in both cases for every hit or miss in L2, we do same operations on L2 cache and after every access we have same blocks in both cache.

Source Code file: Prob1.c

Problem 2: L3 Misses in Inclusive

| Trace File | Cold Misses | Cold+Capacity | Capacity Misses | Conflict Misses |
|------------|-------------|---------------|-----------------|------------------------|
| Bzip | 119753 | | | |
| GCC | 773053 | | | |
| Gromacs | 107962 | 143254 | 35292 | 27277 |
| h264ref | 63703 | 111605 | 47902 | 230541 |
| hmmer | 75884 | 153447 | 77563 | 237779 |
| Sphinix | 122069 | | | |

NOTE: Not able to run program for some files, because its slow for large files.

Source Code: Prob2.c

Cold Misses are calculated in Prob1.c file.

Cold misses doesn't depend on **associativity, capacity of cache** and will be same for all caches unless **block size** is changed.

Prob2.c file gives the **Cold+Capacity** Misses occured in fully assossiative L3 with MIN algorithm and L2 with LRU.

Calculating Conflict Misses: Given a cache, We can't calculate **conflict misses** directly. For this , We Run our address trace to **fully associative cache using MIN algo** and treating this as ideal configuration **(A)** where **conflict misses** are zero and uses this to calculate conflict misses for given **cache(B)** as follow:

Total misses in cache A=capacity+cold
Total misses in cache B=capacity+cold+conflict
Given that capacity and cold misses are same for both A and B.
Conflict misses in B=Total misses in cache **B**-Total misses in cache **A**