Understanding Recursion

A Step-by-Step Approach

Programming Fundamentals

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What is Recursion?

- A method where a function calls itself
- Solves problems by breaking them into smaller subproblems
- Two main components:
 - Base case(s)
 - Recursive case(s)

Key Concepts

Base Case

- Condition where recursion stops
- Returns value directly without recursion

• Recursive Case

- Problem broken into smaller subproblems
- Function calls itself with modified parameters

Example 1: Factorial

```
• Definition: n! = n \times (n-1)!
  • Base case: 0! = 1! = 1
function fact = factorial_recursive(n)
     if n = 0 \mid \mid n = 1
         fact = 1;
     else
         fact = n * factorial_recursive(n-1);
    end
end
```

Factorial Call Stack

5! calculation:

- 1. factorial(5) \rightarrow 5 \times factorial(4)
- 2. factorial(4) \rightarrow 4 \times factorial(3)
- 3. factorial(3) \rightarrow 3 \times factorial(2)
- 4. factorial(2) \rightarrow 2 \times factorial(1)
- 5. factorial(1) \rightarrow 1

Then unwind: 1 \rightarrow 2 \rightarrow 6 \rightarrow 24 \rightarrow 120

Example 2: Fibonacci Sequence

- Each number is sum of previous two
- Base cases: F(1) = F(2) = 1

```
function fib = fibonacci_recursive(n)
  if n <= 2
      fib = 1;
  else
      fib = fibonacci_recursive(n-1) +
            fibonacci_recursive(n-2);
  end
end</pre>
```

Recursion vs Iteration

Recursion

- Elegant and clear
- Memory intensive
- Natural for tree structures

Iteration

- More efficient
- Less memory usage
- Better for linear problems

Example 3: Binary Search

```
function index = binary_search_recursive(arr, target, le
    if left > right
        index = -1:
        return:
    end
    mid = floor((left + right)/2);
    if arr(mid) == target
        index = mid;
    elseif arr(mid) > target
        index = binary_search_recursive(arr, target, lef
    else
        index = binary_search_recursive(arr, target, mid
    end
end
```

Classic Example: Tower of Hanoi

- Problem:
 - Move n disks from source to target
 - Using auxiliary pole
 - Larger disk cannot be on smaller disk
- Recursive solution:
 - 1. Move n-1 disks to auxiliary
 - 2. Move largest disk to target
 - 3. Move n-1 disks to target

Tower of Hanoi Implementation

```
function tower_of_hanoi(n, source, auxiliary, target)
    if n == 1
        fprintf('Move_disk_1_from_%s_to_%s\n',
                source, target);
        return;
    end
    tower_of_hanoi(n-1, source, target, auxiliary);
    fprintf('Move_disk_%d_from_%s_to_%s\n',
            n, source, target);
    tower_of_hanoi(n-1, auxiliary, source, target);
end
```