

8 am

- Sort almost in "place", not require additional array
- Avg. performance $O(n \log n)$ & worst case $O(n^2)$

1 ppm

2 17, 12, 6, 19, 23, 8, 5, 10

pivot element = 12

3 increment i till find element greater than 12

4. Decrement j until find element less than 12.

5 when find swap them

6 \therefore Two parts will created
elements > 12 & elements < 12

8 Take individual part and repeat the process

Note: If we lucky, partition splits the array evenly $O(n \log n)$

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8 am



12

1 μm

26 | Sunday

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graph TD; n --> 1; n --> n_minus_1["(n-1)"]; n_minus_1 --> 1; n_minus_1 --> n_minus_2["(n-2)"]
```

How can we make sure that we are usually lucky?

- partition around 'middle' element:

- partition around random element (works well in practice)

Randomized Quicksort running time

worst $O(n^2)$

Note: Best $O(n \log_2 n)$

Avg. $0(n \log_2 n)$

Insertion is BST avg. time
is $O(n \log n)$.