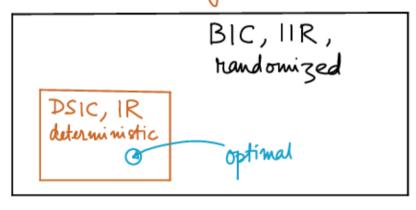
Examples of optimal mechanismo Optimal mechanism design problem:

max  $\int_{i \in \mathbb{N}} \left( \sum_{i \in \mathbb{N}} w_i(t_i) f_i(t) \right) g(t) dt$ , s.t.  $f \in \mathbb{NDE}$ 

Solution for regular Wi's:

$$f_i(t) = \begin{cases} 1 & \text{if } w_i(t_i) > w_j(t_j) \neq j \\ 0 & \text{ow} \end{cases}$$

We wanted to find an allocation that is NDE, but found an f that is non-decreasing. Also, it is deterministic.



Space of regular vintual valuations

Theorem: Suppose every agent's valuation is negular. Then, for every type profile t,

otherwise, 
$$f_i(t) = \begin{cases} 1 & \text{if } w_i(t_i) > w_j(t_j) & \text{if } t \neq 0 \end{cases}$$

ties are broken arbitrarily. Payments are given by

$$p_{i}(t) = \begin{cases} 0 & \text{if } f_{i}(t) = 0 \\ \max \{ w_{i}^{-1}(0), K_{i}^{*}(t_{i}) \} \text{ if } f_{i}(t) = 1 \end{cases}$$

then (f, t) is an optimal mechanism.

$$W_{i}^{-1}(0)$$
: The value of  $t_{i}$  where  $W_{i}(t_{i}) = 0$ .  
 $K_{i}^{*}(t_{i}) = \inf \{ t_{i} : f_{i}(t_{i}, t_{i}) = 1 \}$ 

The minimum value of ti where i begins to be the winner

Example 1: Two buyers: 
$$T_1 = [0,12]$$
,  $T_2 = [0,18]$   
Uniform, independent prior.

$$W_1(t_1) = t_1 - \frac{1 - G_1(t_1)}{g_1(t_1)} = t_1 - \frac{1 - \frac{t_1}{12}}{\frac{1}{12}} = 2t_1 - 12$$

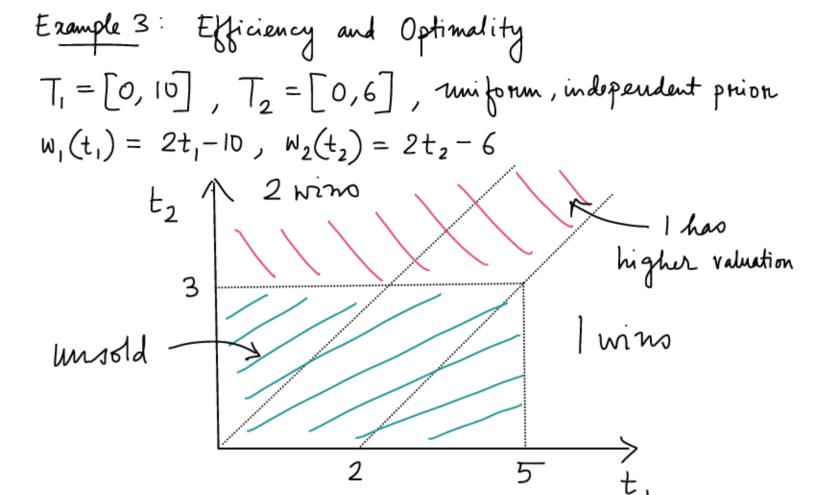
$$W_2(t_2) = 2t_2 - 18$$
 $t_1$ 
 $t_2$ 
action
 $t_1$ 
 $t_2$ 
 $t_3$ 
 $t_4$ 
 $t_5$ 
 $t_6$ 
 $t_7$ 
 $t_8$ 
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 $t_9$ 

Example 2: Symmetric bidders: The valuations are drawn from the same distribution,  $g_i = g$ ,  $T_i = T$ ,  $\forall i \in N$  virtual valuation:  $W_i = W$ .

$$W(t_i) > W(t_j) \text{ iff } t_i > t_j$$

the object goes to the highest bidder. Not sold if  $\bar{w}(0) > t_i$  $\forall i \in \mathbb{N}$ . Payment  $i = \max_i \{ \bar{w}'(0), \max_{i \neq i} t_j \}$ 

Second price auction with a reserve price, and is efficient when the object is Ald.



Unsold is inefficient, also in The negion of The plane.