

GS Theorem allows unrestricted preferences

$$f : P^n \rightarrow A$$

One reason of such a restrictive result is that the domain of the SCF is large - a potential manipulator has many options to manipulate.

Look from the definition of strategyproofness.

$$f(P_i, P_{-i}) \succeq_i f(P_i', P_{-i}) \quad \forall P_i, P_i' \in P$$

If we now reduce the set of preferences from P to some subset of P , the

SCFs that are truthful on P continues to be truthful, but the hope is that we may find more truthful SCFs.

Domain Restrictions

In this course, we'll see three domain restrictions

- ① Single peaked preferences
- ② Divisible good allocation
- ③ Quasi-linear preferences

Each of these domains have interesting non-dictatorial SCFs that are strategyproof.

Single-peaked preferences

One common order over the alternatives

After the alternatives are ordered - agent preferences have a preferences that is single peaked w.r.t that order of the alternatives.


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Motivating examples:

- ① Facility location: School / Hospital / Post Office etc.
- ② Political ideology: Left - Center - Right
- ③ Temperature sensing

Each of these examples have a natural ordering over the alternatives. - we denote this ordering with $<$ (as in real numbers) - this ~~can~~ need not always be an ordering over the real numbers but can be any relation that is transitive and antisymmetric - for simplicity we'll discuss only alternatives on a real line.

Example: $\{a, b, c\}$ three facilities on a real line with $a < b < c$ (which means the locations of a, b, c)

	preferences (valid)				preferences (invalid)	
	a	b	b	c	a	c
	b	a	c	b	c	a
	c	c	a	a	b	b

Defn: A preference ordering P_i (strict over A) of agent i is single-peaked w.r.t. $<$ if

- for all $b, c \in A$ with $b < c \leq P_i(1)$ we have $c P_i b$
- for all $b, c \in A$ with $P_i(1) \leq b < c$ we have $b P_i c$.

Let \mathcal{S} be the set of single-peaked preferences

$\mathcal{S} \subset \mathcal{P}$ - we have seen this in the previous example

$f: \mathcal{S}^n \rightarrow A$ (social choice function)

An SCF f is manipulable if $\exists i, P_i, P_i'$ and P_{-i} s.t.
 $f(P_i', P_{-i}) P_i f(P_i, P_{-i})$

f is strategyproof if it is not manipulable.

What does it imply?

$\forall P_i, P_i', P_{-i} \forall i \in N$

either $f(P_i, P_{-i}) = f(P_i', P_{-i})$

or $f(P_i, P_{-i}) P_i f(P_i', P_{-i})$

but now P_i, P_i' belong to \mathcal{S} and not \mathcal{P}

hence there are less ^{number of} conditions to be satisfied.

How does it circumvent GS theorem?

Example: pick the leftmost peak - mechanism 1

for the leftmost agent, no reason to misreport
 for any other agent (on the right) the only way
 the outcome may change is to ~~push~~ the report
 a peak which is further left - but that is
 strictly worse for him.

Repeat this argument for any k -th peak from the left. Hence even a rightmost peak choosing SCF is strategyproof, so is the median ($k = \lfloor \frac{n}{2} \rfloor$)

Defn: (Median Voter SCF)

An SCF $f: X^n \rightarrow A$ is a median voter SCF if there exists $B = (y_1, \dots, y_{n-1})$ ~~agents~~ ^{agents} such that $f(P) = \text{median}(B, \text{peaks}(P))$ for all preference profiles $P \in X$. The ~~1~~ points in B are called the peaks of "phantom voters".

Note: B is fixed for f and does not change with P .

Why use phantom voters?

All the examples of leftmost, rightmost, median etc. can be taken into account in this single definition.

leftmost: all phantom peaks are on the leftmost point
rightmost: rightmost

Theorem: (Moulin 1980)

Every median voter SCF is strategyproof.

Proof: Only the peaks of the preferences matter.

If $P = (p_1, \dots, p_n)$ and $f(P) = a$, agent i has no reason to manipulate if $p_i(1) = a$. If the peak $p_i(1)$ is to the left of a , the only way he can change the median is by reporting his peak to be right of a - which takes the median farther away from $p_i(1)$ - hence a similar argument holds when the peak is on the right of a . Hence

strategyproof. □

Note: Mean does not have this property.