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Q1.

a)

P ₁	P ₂	P ₃	P ₄	P ₅	P ₆	P ₇	P ₈	P ₉	P ₁₀	P ₁ '	P ₂ '
a	c	a	b	a	c	b	c	b	a	b	a
b	b	c	c	b	a	a	a	a	c	a	b
c	a	b	a	c	b	c	b	c	b	c	c

$f(P_1, P_2) = a$ (given)

$f(P_3, P_4) = a$ (monotonicity) The position of a is getting weakly better from (P_1, P_2) to (P_3, P_4) .

$f(P_5, P_6) = a$ (monotonicity) The position of a is getting weakly better from (P_3, P_4) to (P_5, P_6)

$f(P_7, P_8)$ can only be b or c as $f(P_7, P_8)$ can be either be $P_7(1)$ or $P_8(1)$.

But if $f(P_7, P_8) = c$ then P_7 can manipulate to P_5 and get A elected.

Hence, $f(P_7, P_8) = b$.

$f(P_9, P_{10}) = b$ (monotonicity)

The position of b is getting weakly better from (P_7, P_8) to (P_9, P_{10}) .

$f(P_1', P_2') = b$ (monotonicity)

The position of b is getting weakly better from (P_9, P_{10}) to (P_1', P_2') .

b)

No. the earlier conclusion does not hold in this case as the profiles P_3, P_6 will not be allowed in single peaked preferences. There is restriction on the preferences.

Mechanism such that $f(P_1', P_2') = a$:

P _x (1)	P _y (1)	f(P _x , P _y)
a	a	a
a	b	a
a	c	a
c	a	a
c	c	c
c	b	b
b	a	a
b	c	b
b	b	b

Since preferences are single peaked, only 4 preference profiles $P_1', P_2' = P_1, P_2$ and P_4 are possible.

Showing that the above mechanism is strategyproof:

Case 1: $f(P_x, P_y)=c$ it is possible only when both $P_x(1)=P_y(1)=c$. Hence, none of them have any motivation to manipulate.

Case 2: $f(P_x, P_y)=a$ it is possible if $P_x(1)=a$ or $P_y(1)=a$ or $P_x(1)=P_y(1)=a$

If $P_x(1)=P_y(1)=a$, then none of them have any motivation to manipulate.

If $P_x(1)=a$ and $P_y(1) \neq a$, then no matter what $P_y(1)$ is $f(P_x, P_y)=a$. Hence, none of them have any motivation to manipulate.

Similarly, we can argue for $P_x(1) \neq a$ and $P_y(1)=a$.

Case 3: $f(P_x, P_y)=b$, it is possible only in the following cases:

1.) $P_x(1)=P_y(1)=b$ In this case, none of them have any motivation to manipulate.

2.) $P_x(1)=b$ and $P_y(1)=c$. In this case, $P_y(2)=b$ (because of the restriction caused by single peaked preference). Therefore, none of them have any motivation to manipulate.

3.) $P_x(1)=c$ and $P_y(1)=b$

Similar to the above argument.

Q2. Let X be the set of projects.

Let S be the set of all the non-empty subsets of X .

Now, for any agent i with preference P_i over the projects, the preference over the elements of S would always have X (i.e. the subset having all the projects) as $P_i(1)$.

Hence, the preference profiles have a restriction and hence Gibbard-Satterthwaite result will not apply here.

Q3. We need to consider only the peak preferences of all the agents. So let us denote

$P = (P_1(1), \dots, P_i(1), \dots, P_n(1))$ and let $f(P) = a \in A$ is the median of these peaks where A is set of agents.

Consider a group of agents K , such that $K \subseteq A$.

Now if agents try to manipulate, then there are three cases possible:

Case 1: If $P_i(1) = a$ for all $i \in K$, then there is no reason for the agents to manipulate.

Case 2: If $P_i(1) < a$ for all $i \in K$,

Then if the agents shift their preference to any position left of a , the median will not change.

If they shift their preference towards right of a , i.e. $(P_i, P_{-i}) \rightarrow (P_i', P_{-i})$ s.t. $a < P_i'(1)$ for all $i \in K$, since P_i is a single-peaked preference the median will move further towards right and they would be worse off.

Hence, they have no profitable manipulation.

Case 3: If $P_i(1) > a$ for all $i \in K$, then by arguments similar to case 2, there is no profitable manipulation.

Note that only these three are the only cases possible as an agent with $P_i(1) < a$ will never collude with $P_i(1) > a$ as this collusion can not be beneficial for both at the same time.

Hence, f is group strategyproof.