Core limitation 1: many solutions - set-valued

Shapley Value

- Single-valued solution concept.
- Based on axioms (known as Shapley axioms) similar to Nash bargaining.

Notation: ϕ be a single-valued solution concept $\phi_i(N, v)$ is called the allocation of player $i \in N$.

Axioms:

(1) Efficiency: A solution concept ϕ satisfies efficiency if for every TU game (N, 12) $\sum \phi_{i}(N,12) = \mathcal{N}(N), \quad [m-wastage property]$ $i \in N$

2) Symmetry:

Defn: Players i and j are symmetric players if

for any every coalition $S \subseteq N(\{i,j\})$, $v(SU\{i\}) = v(SU\{j\})$

Symmetrie players give same marginal contribution to every coalition.

Defor: A solution concept of satisfier symmetry if for every coalitional game (N, ve) and every pair of symmetric players i and j in the game $(N, ve) = \phi_j(N, ve)$.

[equal treatment for equals]

3 Null player property;

Defu: A player i is called a mull player in (N, v)

'y for every SCN, v(S) = v(Sv &i3)

— clearly v(i) = 0.

Defu! A solution concept ϕ satisfies mult player property if for every coalitional game (N,10) and for every rull player i, ϕ : (N,10)=0,

Additivity! A sention concept ϕ satisfies additivity if for every pair of coalitional games (N, ν) and (N, ω) $\phi(N, \nu+\omega) = \phi(N, \nu) + \phi(N, \omega)$.

"To what extent a single game is equivalent to playing two games individually?"

This property says independence — The share/allocation from a game with added valuation is exactly the same as the playing the games independently and Collecting the Hewards.

Examples

1) $\Psi_{i}(N,\nu) = \nu(i)$ additivity - $\Psi_{i}(N_{\bullet},\nu+\omega) = (\nu+\omega)(i) = \nu(i) + \omega(i)$ = $\Psi_{i}(N,\nu) + \Psi_{i}(N,\omega)$

Symmetry - If $\forall S \subseteq N \setminus \{i,j\}$ $v(su_{ij}) = v(su_{ij})$ | Then apply $S = \phi \Rightarrow v(i) = v(j) \Rightarrow \psi_i(N,v) = \psi_i(N,v)$. mull player - for every mull player v(i) = 0, $\psi_i(N, v) = 0$. efficiency - not necessarily. $\Sigma v(i) \neq v(N)$.

(2) A player i is called a dummy player if v(sv{i}) = v(s) + v(i) ∀s ⊆ N({i})

Every mull player is a dummy player.

Let d(re) be the number of dummy players in (N, v)

Consider a solution concept. $v(i) + \sum v(i) + \frac{v(i)}{v(i)}$ v(i)

i is not a durmny player i is a dummy player.

efficiency - yer, Null-yes, mill is dummy-gets zero. symmetry - dearly it both the players are dummy Then this is time, both players are non-dummy then also time.

What about i non-dummy and 's dummy

Can they be symmetric?

- let D be The dummy set, clearly v(D) = \(\times \tag{V}(j) -'y i ∈ ND and j ∈ D and They are symmetric

Then we have $v(Su\{i\}) = v(Su\{j\}) \forall S \subseteq N(\{i,j\})$

clearly
$$v(i) = v(j)$$
 = $v(s) + v(j)$
 $s = \phi$ = $v(s) + v(i)$

= v(s) +v(j) = v(s) +v(j) +v(i)

= u(suis) + u(i)

コ v (まりをif) = v (ま) + v(i) サミCN(をif

=) i is a dummy player -> <

does not satisfy additivity. V(1) = V(2) = V(3) = V(1,2) = V(1,3) = 0 - - - - (1) ~(2,3) = ~(1,2,3) = 1 and n(1) = 12(2) = n(3) = u(1,3) = 0, u(1,2) = 12(2,3) = u(1,2,3)=1 $(0,\frac{1}{2},\frac{1}{2})$ (N, v) -> player 1 is the dummy (N, u) no dummy (3/3/3) Consider (N, K+V) (u+v)(1) = (u+v)(2) = (u+v)(3) = (u+v)(1,3) = 0(u+v)(1,2)=1, (u+v)(2,3)=2=(u+v)(1,2,3)W-dumy $\Psi(N, u+v) = \left(\frac{2}{3}, \frac{2}{3}, \frac{2}{3}\right)$ 4(N,U)+4(N,U)=(13,至,至). $\Psi_i(N, v) = \max_{\{s: i \notin s\}} \{v(sv\{i\}) - v(s)\}$ (3)symmetry, mill player efficiency, additivity x. $\Psi_{i}(N,v) = w(1,2,...,i) - w(1,2,...,i-1)$ efficiency - V additivity - V mill player - ~

not symmetry — consider the first game of example (2) eqn. (1) $\Psi(N, v) = (0, 0, 1)$ but 2,3 are symmetric.

The solution concept can be defined for any order of The players (not just The identity order) Say TI(N) denote the set of all possible orders.

De Call π ∈ TT(N) to be one ordering/permutation of The playerro.

Call the predecesson of player i in the permutation

 π $P_i(\pi) = \{j \in N : \pi(j) / \pi(i)\}$

 $P_i(\pi) = \phi' \psi \pi(i) = 1$.

 $P_i(\pi) U\{i\} = P_k(\pi) \Leftrightarrow \pi(k) = \pi(i) + 1$

Now, we can define the solution concept $\Psi_{i}^{\pi}(N,v) = v(P_{i}(\pi)U\{i\}) - v(P_{i}(\pi)).$

As we saw in cas example (4) before, this solution concept satisfier efficiency, mull-player property, additivity, not symmetry.

Shapley value

Q. Is there a solution concept that satisfies all & four properties?

A. Yes, and it is unique.

Defor (Shapley 1953) The Shapley value is The solution concept Sh defined as,

 $Sh_i(N, v) = \int \sum_{n} \left(P_i(\pi) \cup \{i\}\right)$ - v(Pi(n)) HIEN.

A simple average over all
$$\psi_i^{\pi}$$
's

Hence $Sh_{\mathbf{k}}(N, \nu) = \frac{1}{n!} \sum \psi^{\pi}(N, \nu)$

* $\pi \in \Pi(N)$

Theorem: The Shapley value is the only single-valued solution concept satisfying efficiency, additivity, null player, and symmetry.

Aco An equivalent formula for Shapley value

$$\frac{1}{m!} \sum_{\pi \in \Pi(N)} \Psi_{i}^{\pi}(N, u) = \frac{1}{n!} \sum_{S \subseteq N \setminus \{i\}} \sum_{\pi \in \pi(N)} (\mathcal{V}(P_{i}(\pi) \cup \{i\}) - \mathcal{V}(P_{i}(\pi)))$$

$$= \frac{1}{n!} \sum_{S \subseteq N \setminus \{i\}} \sum_{\pi \in \pi(N)} (\mathcal{V}(S \cup \{i\}) - \mathcal{V}(S))$$

$$= \sum_{S \subseteq N \setminus \{i\}} \frac{1}{n!} [\mathcal{V}(S \cup \{i\}) - \mathcal{V}(S)]$$

$$S \subseteq N \setminus \{i\}$$

Interpretation: "Average marginal contribution to all other mets walitions".

Proof of Shapley theorem; Part

Part 1: Shapley value satisfier the four axioms.

Each of 4th satisfies efficiency, additivity, mull player, so does their average [Exercise]

Symmetry: Let i and j be two symmetric players. Given a permutation TC, define the tollowing permutation f(TC), s.t.

 $f: T(N) \rightarrow T(N)$

 $f(\pi) \text{ just swaps the positions of } i \text{ and } j$ $(f(\pi))(k) = \begin{cases} \pi(j) & \text{if } k=i \\ \pi(i) & \text{if } k=j \end{cases}$ $\pi(k) & \text{if } k \neq i,j$

Claim: $Y_{i}^{T}(N, v) = Y_{j}^{T}(N, v)$

 $(P_{i}(\pi) \cup \{i\}) - \nu(P_{i}(\pi)) = \nu(P_{j}(f(\pi)) \cup \{j\}) - \nu(P_{j}(f(\pi)))$

Case 1: Player i appears before j'in Ti, i.e., j ∉ Pi(Ti)

clearly $P_j(f(\pi)) = P_i(\pi) = v(P_j(f(\pi))) = v(P_i(\pi))$

and since i and j are symmetric players, $v(P_i(\pi) \cup \{i\}) = v(P_j(f(\pi)) \cup \{j\})$

equ (2) in holds.

Eq. (2) holds