

Lec 9 Part 2: Uniqueness.

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Define a game called the carrier game

A coalition is winning if it contains a distinguished set.

Intuition: say T is an influential coalition -

any coalition containing it can pass a bill/do a change.

Defn: Let $T \subseteq N$ be a non-empty coalition. The carrier game over T is the game (N, u_T) s.t. for each coalition $S \subseteq N$

$$u_T(S) = \begin{cases} 1 & \text{if } T \subseteq S \\ 0 & \text{ow.} \end{cases}$$

Theorem 1: Every game (N, v) is a linear combination of carrier games.

To define any TU game, we need to define the valuations over all non-empty subsets. Hence it has $2^n - 1$ degrees of freedom. Hence, every game (N, v) is a point in $\mathbb{R}^{2^n - 1}$. Want to show that carrier games span this space. Find carrier games that are linearly independent and forms a basis.

Suppose, carrier games are linearly dependent (for contradiction). \exists real numbers $(\alpha_T)_{\{T \subseteq N, T \neq \emptyset\}}$ not all zero, s.t.

$$\sum_{\{T \subseteq N, T \neq \emptyset\}} \alpha_T u_T(S) = 0, \quad \forall S \subseteq N.$$

Let $\mathcal{T} = \{T \subseteq N : T \neq \emptyset, \alpha_T \neq 0\}$ collections of non-empty coalitions with non-zero coefficients in the above equation. Since $\{\alpha_T\}_{\{T \subseteq N, T \neq \emptyset\}}$ are not all zero, \exists a minimal coalition in \mathcal{T} , i.e. coalition with smallest cardinality. Say $S_0 \in \mathcal{T}$ is ^{one such} the coalition. \nexists any subset of S_0 with positive coefficients. Consider

$$\begin{aligned} \sum_{\{T \subseteq N, T \neq \emptyset\}} \alpha_T u_T(S_0) &= \sum_{\{T \subseteq S_0, T \neq \emptyset\}} \alpha_T u_T(S_0) + \alpha_{S_0} u_{S_0}(S_0) \\ &\quad + \underbrace{\sum_{T \not\subseteq S_0} \alpha_T u_T(S_0)}_{=0} \\ &= \alpha_{S_0} \neq 0 \end{aligned}$$

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Theorem 2: Let T be a non-empty coalition, and $\alpha \in \mathbb{R}$. Define a game $(N, u_{T,\alpha})$ as follows

$$u_{T,\alpha}(S) = \begin{cases} \alpha & \text{if } T \subseteq S \\ 0 & \text{otherwise} \end{cases}$$

If ϕ is a solution concept that satisfies efficiency, symmetry, null player property, Then

$$\phi_i(N, u_{T,\alpha}) = \begin{cases} \frac{\alpha}{|T|} & \text{if } i \in T \\ 0 & \text{otherwise} \end{cases}$$

obs 1: If $i \notin T$ $u_{T,\alpha}(S \cup \{i\}) = u_{T,\alpha}(S) \quad \forall S \subseteq N$
 i is a null player

obs 2: If $i, j \in T$, they are symmetric.

$$u_{T,\alpha}(S \cup \{i\}) = u_{T,\alpha}(S \cup \{j\}) \quad \forall S \subseteq N \setminus \{i, j\}$$

From the fact that ϕ is efficient, symmetric, null-player compliant & solution concept, the result follows.

Finishing the proof of Part 2: Uniqueness

Shapley value satisfies the four properties.

Need to show: any ϕ satisfying these four properties is identical to Sh.

Thm 1 says that for any game (N, v) , we can write v as sum of u_{T,α_T} 's

$$\exists (\alpha_T)_{\{T \subseteq N, T \neq \emptyset\}} \text{ s.t.}$$

$$v(S) = \sum_{\{T \subseteq N, T \neq \emptyset\}} u_{T,\alpha_T}(S)$$

Thm 2 says, since both ϕ and Sh satisfy efficiency, symmetry, and null player property

$$\phi(N, u_T, \alpha_T) = Sh(N, u_T, \alpha_T), \quad \forall T \subseteq N, T \neq \emptyset.$$

Since both ϕ and Sh satisfy additivity

$$\phi(N, v) = \sum_{\{T \subseteq N, T \neq \emptyset\}} \phi(N, u_T, \alpha_T) = \sum_{\{T \subseteq N, T \neq \emptyset\}} Sh(N, u_T, \alpha_T) = Sh(N, v).$$

we started with an arbitrary game (N, v) , hence this holds for all such games. \square

Examples

① Two player bargaining: (N, v)

$$v(1) = v(2) = 0, \quad v(1, 2) = 1$$

symmetric players: 1 & 2, Shapley value is efficient

$$Sh(N, v) = \left(\frac{1}{2}, \frac{1}{2}\right)$$

② Majority game

$$v(S) = \begin{cases} 0 & \text{if } |S| \leq \frac{n}{2} \\ 1 & \text{if } |S| > \frac{n}{2} \end{cases}$$

all players are symmetric, hence Shapley values are same, together with efficiency

$$Sh(N, v) = \left(\frac{1}{n}, \dots, \frac{1}{n}\right)$$

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③ Gloves game

$$v(1) = v(2) = v(3) = v(1, 2) = 0$$

$$v(1, 3) = v(2, 3) = v(1, 2, 3) = 1$$

| Perm | Player 1 | Player 2 | Player 3 |
|---------|----------------------------|----------------------|---------------|
| 1, 2, 3 | $v(1) - v(\emptyset) = 0$ | $v(1, 2) - v(1) = 0$ | 1 |
| 1, 3, 2 | 0 | $v(1, 3) - v(1) = 0$ | 1 |
| 2, 1, 3 | $v(1, 2) - v(2) = 0$ | 0 | 1 |
| 2, 3, 1 | $v(1, 2, 3) - v(2, 3) = 0$ | 0 | 1 |
| 3, 1, 2 | $v(1, 3) - v(3) = 1$ | 0 | 0 |
| 3, 2, 1 | 0 | $v(2, 3) - v(3) = 1$ | 0 |
| | $\frac{1}{6}$ | $\frac{1}{6}$ | $\frac{4}{6}$ |

$$Sh(N, v) = \left(\frac{1}{6}, \frac{1}{6}, \frac{2}{3}\right)$$

emphasizes that ~~the~~ player 3 is the most powerful player.
but players 1 and 2 do not get zero in the allocation.

Core: $(0, 0, 1)$ is the singleton.

Hence Shapley value is not in core.

Application: Shapley - Shubik power index

Defn: (Simple, monotone games)

Simple games: The value of any coalition can either be 0 or 1.

Monotone games: If any coalition has value 1, every superset of that coalition also has value 1.

Motivation: to model legislations/decisions based on committees.

Defn: The Shapley-Shubik power index is a function associating each simple monotonic game with its Shapley value. The i^{th} co-ordinate denotes the "power" of player i in this game.

$$Sh_i(N, v) = \sum_{\substack{\{S \subseteq N \mid i \in S\} \\ S \cup \{i\} \text{ is winning} \\ S \text{ is losing}}} \frac{|S|! (n - |S| - 1)!}{n!}$$

counting all such scenarios where ~~step~~ player i is pivotal.

Case study: UN Security council

UN: Body of international political system, established in 1945 (after WWII)

till 1965: five permanent members, six non permanent members

Resolution adopted if it receives at least 7 votes but all permanent members have to be unanimous - all of them have veto powers.

Debated about unequal distribution of powers in the security council.

after 1965: five permanent members, 10 nonpermanent members
resolution needed 9 votes but veto power remains with the permanent members

This is a simple, monotonic game. Compute the Shapley-Shubik power index.

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P: permanent members, NP: non permanent members

pre-1965:

$$v(S) = \begin{cases} 1 & \text{if } S \supseteq P \text{ and } |S| \geq 7 \\ 0 & \text{otherwise} \end{cases}$$

non permanent i

$$Sh_i(N, v) = \binom{5}{1} \frac{6! 4!}{11!} = \frac{1}{462}$$

all non-permanent members are symmetric

all permanent members are symmetric

Shapley value is efficient, hence for a permanent j

$$Sh_j(N, v) = \frac{1}{5} \left(1 - \frac{6}{462} \right) = \frac{91.2}{462}$$

Power ratio of nonpermanent to permanent = 1:91.2.

post-1965:

$$v(S) = \begin{cases} 1 & \text{if } S \supseteq P \text{ and } |S| \geq 9 \\ 0 & \text{otherwise} \end{cases}$$

non-permanent i:

$$Sh_i(N, v) = \binom{9}{3} \frac{8! 6!}{15!} = \frac{4}{2145}$$

permanent j:

$$Sh_j(N, v) = \frac{1}{5} \left(1 - 10 \times \frac{4}{2145} \right) = \frac{421}{2145}$$

ratio = 1:105.25

~~Restructuring~~ Restructuring actually increased the power of the permanent members.

Convex games

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$$\forall S, T \subseteq N$$

$$v(S) + v(T) \leq v(S \cup T) + v(S \cap T)$$

Thm: If (N, v) is a convex game, the SV is in the core.

Proof: For any permutation $\pi \in \Pi(N)$, consider the imputation w^π

$$w_i^\pi = v(P_i(\pi) \cup \{i\}) - v(P_i(\pi))$$

We have shown that this imputation is in core for every $\pi \in \Pi(N)$. Since core is a convex set, any the convex combination of these points will be in core. In particular, the Shapley value

$$Sh(N, v) = \frac{1}{n!} \sum_{\pi \in \Pi(N)} w^\pi$$

will also be in the core.

Consistency of the SV.

Defn: Let ϕ be a single valued solution concept, let (N, v) be a coalitional game and $S \subseteq N$ and $S \neq \emptyset$.

The Hart-Mas-colell reduced game over S relative to ϕ is the game $(S, \tilde{v}_{S, \phi})$ s.t.

$$\tilde{v}_{S, \phi}(R) = \begin{cases} v(R \cup S^c) - \sum_{i \in S^c} \phi_i(R \cup S^c, v), & \forall R \subseteq S, R \neq \emptyset \\ 0 & \text{if } R = \emptyset \end{cases}$$

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Difference with Davis-Maschler reduced game

- solution concept for single-valued solutions
- DM selects the most beneficial coalition of S^c
- but HM considers the whole of S^c .

Defn: A solution concept ϕ is consistent w.r.t the HM reduced game if for every game (N, v) , every nonempty coalition S , and for every $i \in S$,

$$\phi_i(N, v) = \phi_i(S, \tilde{v}_{S, \phi})$$

Theorem: SV is HM reduced game consistent,