

An interpretation of the utility of the agent under VCG

$$\begin{aligned}
 u_i(\theta_i, \underline{\theta}_i | \theta_i) &= v_i(f^{AE}(\theta_i, \underline{\theta}_i), \theta_i) - p_i^{VCG}(\theta_i, \underline{\theta}_i) \\
 &= v_i(f^{AE}(\theta_i, \underline{\theta}_i), \theta_i) - \max_{a \in A} \sum_{j \neq i} v_j(a, \theta_j) \\
 &\quad + \sum_{j \neq i} v_j(f^{AE}(\theta_i, \underline{\theta}_i), \theta_j) \\
 &= \sum_{j \in N} v_j(f^{AE}(\theta_i, \underline{\theta}_i), \theta_j) - \max_{a \in A} \sum_{j \neq i} v_j(a, \theta_j) \\
 &= \underbrace{\max_{a \in A} \sum_{j \in N} v_j(a, \theta_j)}_{\text{max SW in } i\text{'s presence}} - \underbrace{\max_{a \in A} \sum_{j \neq i} v_j(a, \theta_j)}_{\text{max SW in } i\text{'s absence}}
 \end{aligned}$$

= marginal contribution of  $i$  in the SW.

### Illustration of VCG payment

- ① Single object allocation. Consider agent  $i \in N$ . Every agent has some value when the object is assigned to him — and zero otherwise. Efficiency requires that the object should go to the agent who values it the most. If  $i$ 's value/bid is the maximum, he is assigned the object,

$$\begin{aligned}
 \text{but Payment} &= \underbrace{\max_{a \in A} \sum_{j \neq i} v_j(a, \theta_j)}_{\text{second highest bid}} - \underbrace{\sum_{j \neq i} v_j(f(\theta), \theta_j)}_{=0 \text{ if } i\text{'s bid is highest}}
 \end{aligned}$$

Second price auction.

② What is Pivotal about the VCG payment? Consider a public project allocation.

First term: the social welfare of all the agents except agent  $i$  if  $i$  wasn't present

Second term: social welfare of the other agents in  $i$ 's presence.

hence another interpretation of the payment: loss in social welfare of the other agents ~~of agent  $i$~~  <sup>because</sup> because of agent  $i$ 's presence.

Example:

		Football	Library	Museum
<del>the</del>	A	0	70	50
<del>the</del>	B	100 95	10	50
<del>the</del>	C	10	50	50

VCG allocation: efficient: M

$$A = \{F, L, M\}$$

$$A \text{ pays} = 105 - 100 = 5$$

$$B \text{ pays} = 120 - 100 = 20$$

$$C \text{ pays} = 100 - 100 = 0 \leftarrow \text{non-pivotal agent}$$

The agent whose presence changes the outcome are called pivotal. VCG charges only them.

③ Combinatorial allocation: Sale of multiple objects

	$\emptyset$	$\{1\}$	$\{2\}$	$\{1,2\}$
$v_1$	0	8	6	12
$v_2$	0	9	4	14

type is the value it self

$$v_i(a, \theta_i) = \theta_i(a_i)$$

Efficient allocation  $\{1\} \rightarrow \text{Player 2}, \{2\} \rightarrow 1$

$$p_1^{\text{VCG}}(\theta_1, \theta_2) = \cancel{14} - \cancel{9} = \cancel{5}$$

$$\max_{a \in A} \sum_{j \neq i} \theta_j(a) - \sum_{j \neq i} \theta_j(f(v))$$

$$= 14 - 9 = 5, \text{ payoff} = 6 - 5 = 1$$

$$p_2^{\text{VCG}}(\theta_1, \theta_2) = 12 - 6 = 6, \text{ payoff} = 9 - 6 = 3$$

VCG mechanism in combinatorial auctions

VCG has several useful properties

- ① It is DSIC
- ② It is Efficient
- ③ Payments are non-negative (no-subsidy), hence WBB.

Combinatorial auctions:

$M = \{1, \dots, m\}$  set of objects

Set of bundles  $\Omega = \{S : S \subseteq M\}$

Type of agent  $i$  is  $\theta_i : \Omega \rightarrow \mathbb{R}$

Hence  $\theta_i(S)$ ,  $S \in \Omega$ , is the value of agent  $i$  for the bundle  $S$ . We assume  $\theta_i(S) \geq 0 \forall S \in \Omega$   
 <all bundles are "good">

An allocation in this case is a partition of the objects :  ~~$\{x_0, x_1, \dots, x_m\}$~~

$$X = \{x_0, x_1, \dots, x_m\}, x_i \in \Omega$$

$A$  is the collection of such  $X$ 's.  $x_i \cap x_j = \emptyset$  if  $i \neq j$ ,  $\bigcup_{i=0}^m x_i = M$ .

(28-4)

$X_0$  is the set of unallocated objects and  $X_i$  is the bundle allocated to agent  $i$ . Assume  $\theta_i(\emptyset) = 0$ .

Note: the model is such that the valuations ~~are~~ have no externalities (selfish valuations)

$$v_i(X, \theta_i) = \theta_i(X_i) \quad \langle \text{does not depend on the allocation of the other agents} \rangle.$$

In the allocation of goods,  $VCG$

Claim 1: ~~Any agent~~ The payment for an agent that gets no object in the efficient allocation is zero.

Fix  $\theta, M$ .

Proof: Say agent  $i$  gets no object in the efficient allocation, i.e.,  $X \in \operatorname{argmax}_{x \in A} \sum_{i \in N} v_i(x, \theta_i)$

$$X_i = \emptyset$$

VCG payment considers the allocation excluding agent  $i$ . Let that allocation is  $Y \in \operatorname{argmax}_{y \in A} \sum_{j \neq i} v_j(y, \theta_j)$

We have seen that VCG payment always yields

$$p_i^{VCG}(\theta) \geq 0 \quad [\text{no-subsidy}]$$

But we can also write

$$p_i^{VCG}(\theta) = \max_{x \in A} \sum_{j \neq i} v_j(x, \theta_j) - \sum_{j \neq i} v_j(X, \theta_j)$$

adding

$$v_i(Y, \theta_i) \geq 0$$

subtracting

$$v_i(X, \theta_i) = 0$$

$$\leq \sum_{j \in N} v_j(Y, \theta_j) - \sum_{j \in N} v_j(X, \theta_j)$$

max by definition

$$\leq 0$$

$$\boxed{p_i^{VCG}(\theta) = 0}$$

□