

Equilibrium Concepts

Ex-ante: before observing own type

Nash Equilibrium (σ^*, P)

$$U_i(\sigma_i^*, \underline{\sigma}_{-i}^*) \geq U_i(\sigma_i', \underline{\sigma}_{-i}^*) \quad \forall \sigma_i', \forall i$$

[Expectation taken over θ]

Ex-interim: after observing own type

Bayesian Equilibrium (σ^*, P)

$$U_i(\sigma_i^*, \underline{\sigma}_{-i}^* | \theta_i) \geq U_i(\sigma_i', \underline{\sigma}_{-i}^* | \theta_i), \quad \forall \sigma_i' \quad \forall \theta_i \in \Theta_i \quad \forall i \in N.$$

↑
this can be replaced by pure strategies $a_i \quad \forall a_i \in A_i$
both definitions are equivalent.

[Expectation taken over $P(\theta_{-i} | \theta_i)$]

Equivalence of the two equilibrium concepts

Theorem: In a finite Bayesian game, a strategy profile is a Bayesian equilibrium iff it is a Nash equilibrium.

\Rightarrow Suppose (σ^*, P) is a Bayesian equilibrium

$$U_i(\sigma_i^*, \underline{\sigma}_{-i}^* | \theta_i) \geq U_i(a_i, \underline{\sigma}_{-i}^* | \theta_i), \quad \forall a_i \in A_i, \forall \theta_i \in \Theta_i \quad \forall i \in N$$

$$U_i(\sigma_i^*, \underline{\sigma}_{-i}^*) \stackrel{\text{Eq. (3)}}{=} \sum_{\theta_i \in \Theta_i} P(\theta_i) U_i(\sigma_i^*(\theta_i), \underline{\sigma}_{-i}^* | \theta_i)$$

$$\leq \sum_{\theta_i \in \Theta_i} P(\theta_i) U_i(\sigma_i^*, \underline{\sigma}_{-i}^* | \theta_i)$$

$$= U_i(\sigma_i^*, \underline{\sigma}_{-i}^*)$$

Hence (σ^*, P) is a Nash equilibrium.

(13-2)

\Leftarrow Suppose (σ^*, p) is a Nash equilibrium. Assume for contradiction that (σ^*, p) is not a Bayesian equilibrium.

$\exists a_i \in A_i$ s.t. and some $\theta_i \in \Theta_i$, some $i \in N$

$$u_i(a_i, \underline{\sigma}_i^* | \theta_i) > u_i(\sigma_i^*, \underline{\sigma}_i^* | \theta_i) \quad \dots \textcircled{1}$$

consider the strategy $\hat{\sigma}_i$ of i

$$\hat{\sigma}_i(\theta_i') = \sigma_i^*(\theta_i') \quad \forall \theta_i' \in \Theta_i \setminus \{\theta_i\}$$

$$\hat{\sigma}_i(\theta_i, a_i) = 1 \quad \text{and} \quad \hat{\sigma}_i(\theta_i, b_i) = 0 \quad \forall b_i \in A_i \setminus \{a_i\}$$

$$u_i(\hat{\sigma}_i, \underline{\sigma}_i^*) = \sum_{\tilde{\theta}_i \in \Theta_i} P(\tilde{\theta}_i) u_i(\hat{\sigma}_i, \underline{\sigma}_i^* | \tilde{\theta}_i)$$

$$= \sum_{\tilde{\theta}_i \in \Theta_i \setminus \{\theta_i\}} P(\tilde{\theta}_i) u_i(\hat{\sigma}_i, \underline{\sigma}_i^* | \tilde{\theta}_i)$$

$$+ P(\theta_i) u_i(\hat{\sigma}_i, \underline{\sigma}_i^* | \theta_i)$$

=

$$> \sum_{\tilde{\theta}_i \in \Theta_i \setminus \{\theta_i\}} P(\tilde{\theta}_i) u_i(\sigma_i^*, \underline{\sigma}_i^* | \tilde{\theta}_i)$$

$$+ P(\theta_i) u_i(\sigma_i^*, \underline{\sigma}_i^* | \theta_i)$$

$$= u_i(\sigma_i^*, \underline{\sigma}_i^*)$$

$\rightarrow \leftarrow$ Nash equilibrium of $(\sigma_i^*, \underline{\sigma}_i^*)$

□

Existence of a Bayesian equilibrium.

Theorem: Every finite Bayesian game has a Bayesian equilibrium.

Proof: Idea: transform the Bayesian game into a complete information game treating each type a player.

$$N = \bigcup_{i \in N} \Theta_i = \{ \theta_1^1, \theta_1^2, \dots, \theta_1^{|\Theta_1|}, \theta_2^1, \dots, \theta_2^{|\Theta_2|}, \dots, \theta_n^1, \dots, \theta_n^{|\Theta_n|} \}$$

This is finite by assumption.

$$A_{\theta_i} = A_i \quad \forall \theta_i \in \Theta_i \quad \forall i \in N.$$

$$u_{\theta_i}(a_{\theta_i}, a_{-\theta_i}) = \sum_{\underline{\theta}_i \in \Theta_i} P(\underline{\theta}_i | \theta_i) u_i(a_i(\theta_i), \underline{a}_i(\underline{\theta}_i), \theta_i, \underline{\theta}_i)$$

Note: A mixed strategy of player θ_i , σ_{θ_i} is a probability distribution over A_{θ_i} , which is a mixed strategy of player i at type θ_i , $\sigma_i(\theta_i)$ in the original Bayesian game.

Hence a MSNE in the transformed game is a Bayesian equilibrium in the original game.

By Nash theorem MSNE exists in the transformed game \Rightarrow Bayesian equilibrium exists in the original game. \square