

Develop the notation and definitions.

Example of a two-player game

Prisoner's dilemma

		2	
		Cooperate	Defect
1	Cooperate	-5, -5	0, -10
	Defect	-10, 0	-1, -1

Neighboring Kingdom's dilemma

		2	
		A	D
1	A	5, 5	0, 6
	D	6, 0	2, 2

Representation : Normal form / Strategic Form Game
- simultaneous move, one shot games.

Notation: Set of players $N = \{1, 2, \dots, n\}$

Set of actions A_i for player i , specific action $a_i \in A_i$

utility / payoff of agent i , $u_i : \underbrace{A_1 \times A_2 \times \dots \times A_n}_{\prod_{i \in N} A_i = A} \rightarrow \mathbb{R}$

Warning: Not every preference may have utility representation.
Special case that the preferences are expressible using functions as described.

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Example: (Preferences without utility representation)

Student - F - M - relative cooks a, b, c (any two items)

$a \succ b \succ c$

$b \succ c \succ a$

$c \succ a \succ b$

The preference of the majority of this family does not have an utility representation.

Axioms for utility representation (Concept #1)
use outcomes rather than alternatives

① Completeness: $\forall a_1, a_2$, either $a_1 \succ a_2$ or $a_2 \succ a_1$ or $a_1 \sim a_2$.

② Transitivity: If $a_1 \succ a_2$ and $a_2 \succ a_3 \Rightarrow a_1 \succ a_3$

③ Substitutability: If $a_1 \sim a_2$, then \forall sequences of outcomes a_3, \dots, a_k and sets of probabilities p, p_3, \dots, p_k s.t.

$$p + \sum_{i=3}^k p_i = 1, \quad p, p_i \geq 0$$

$$[p : a_1, p_3 : a_3, \dots, p_k : a_k] \sim [p : a_2, p_3 : a_3, \dots, p_k : a_k]$$

(2-2)

④ Decomposability: If $\forall a_i \in A$, $P_{\ell_1}(a_i) = P_{\ell_2}(a_i) \Rightarrow \ell_1 \sim \ell_2$

⑤ Monotonicity: If $a_1 \succ a_2$ and $1 > p > q > 0$

$$[p : a_1, 1-p : a_2] \succ [q : a_1, 1-q : a_2]$$

⑥ Continuity: If $a_1 \succ a_2$ and $a_2 \succ a_3$ then $\exists p \in [0, 1]$

$$\text{s.t. } a_2 \sim [p : a_1, 1-p : a_3]$$

(von-Neumann, Morgenstern 1944)

Theorem: If a preference relation \succsim satisfies axioms ①-⑥

then $\exists u : A \rightarrow [0, 1]$ s.t.

$$\textcircled{1} \quad u(a_1) \geq u(a_2) \Leftrightarrow a_1 \succsim a_2$$

$$\textcircled{2} \quad u([p_1 : a_1, p_2 : a_2, \dots, p_k : a_k]) = \sum_{i=1}^k p_i u(a_i)$$

Proof: Indifferent - Part 1 is obvious $u(a_i) = 0 \forall a_i$ (everything in simplex is same)
② follows from ~~substitutability~~ - all permutation $u(\dots)$ same
decomposability define $u(a_i) = \text{constant}$.

Case^v \exists ~~one~~ at least two alternatives that are preferred over one another
 \bar{a} most preferred, \underline{a} least preferred.

for any alternative a_i , define $u(a_i)$ as p_i s.t.
 $a_i \sim [p_i : \bar{a}, 1-p_i : \underline{a}]$ by continuity.

Part 1: $u(a_1) \geq u(a_2) \Leftrightarrow a_1 \succsim a_2$

(\Rightarrow) define $a_1 \sim \ell_1 = [u(a_1) : \bar{a}, 1-u(a_1) : \underline{a}] \rightarrow$ applies to lotteries too
 $a_2 \sim \ell_2 = [u(a_2) : \bar{a}, 1-u(a_2) : \underline{a}]$

(a) $u(a_1) > u(a_2)$ by monotonicity,
 $a_1 \sim \ell_1 \succ \ell_2 \sim a_2 \Rightarrow a_1 \succ a_2$ (transitivity)

(b) $u(a_1) = u(a_2)$ $a_1 \sim \ell_1 \equiv \ell_2 \sim a_2 \Leftrightarrow a_1 \sim a_2$

(\Leftarrow) $u(a_1) \geq u(a_2) \Leftrightarrow a_1 \succsim a_2$

$\Leftrightarrow u(a_1) < u(a_2) \Rightarrow a_1 \prec a_2$ already shown.

✓ Completeness.

part 2 : $u^* = u([p_1 : a_1, p_2 : a_2, \dots, p_k : a_k]) = \sum_{i=1}^k p_i u(a_i)$

$a_i \sim l_i = [u(a_i) : \bar{a}, 1 - u(a_i) : \underline{a}]$

✓ Substitutability:

$u^* = u([p_1 : l_1, \dots, p_k : l_k]) \rightarrow \sum_{i=1}^k p_i u(a_i)$

✓ decomposability

$= u([p_1 u(a_1) + p_2 u(a_2) + \dots + p_k u(a_k) : \bar{a}, 1 - \sum_{i=1}^k p_i u(a_i) : \underline{a}])$

$u^* = u(l^*) = u(\underbrace{[\sum_{i=1}^k p_i u(a_i) : \bar{a}, 1 - p : \underline{a}]}_{= u(l^*)}) = u^*$

Concept # 2 : Rationality and Intelligence

- Agents pick actions to maximize their utility
- Knows the rule of the game and pick action to maximize utility.

Concept # 3 : Common Knowledge behavioral model

blue/black eyed people

everybody knows the rules, everybody knows that everybody knows it, and so on.