

Silverman - Lecture 3

Lower Bounds for Canonical Heights

K number field

A/K abelian variety

D ample, symmetric divisor

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Question: If $\hat{h}_D(P) \neq 0$,
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Older Analogue

$$\xrightarrow{\text{"canonical"}} h : \overline{\mathbb{Q}}^* \rightarrow [0, \infty)$$

$$h(\alpha^m) = |m| h(\alpha)$$

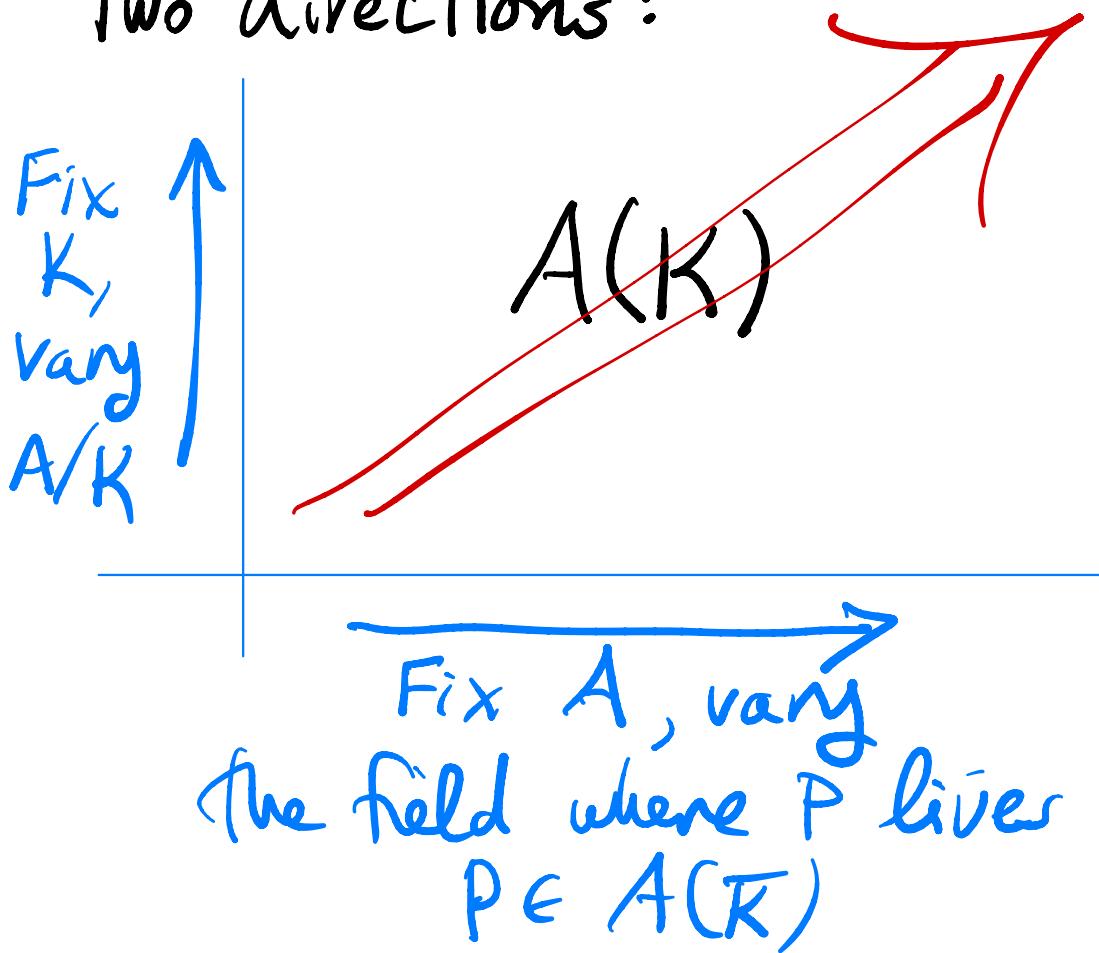
Theorem: (Kronecker)

$$h(\alpha) = 0 \iff \alpha \text{ is a root of } 1 \\ \alpha \in \mu(\overline{\mathbb{Q}})$$

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Question: How small can
 $0 \neq \hat{h}_{\mathcal{A}^D}(P)$ be?

Two directions:



Part I : Fix A/K

Vary the field of def. of P

Analogy :

If $\alpha \in \bar{\mathbb{Q}}^*$, $h(\alpha) \geq ?$
 $\alpha \notin \mu$

$$h(2'^n) = \frac{1}{n} h(2) = \frac{\log 2}{n}$$

$$(h(\alpha^n) = |n| h(\alpha)) \xrightarrow[n \rightarrow \infty]{} 0$$

$$[\mathbb{Q}(2'^n) : \mathbb{Q}] = n$$

Dicho for A/K $\widehat{h} = \widehat{h}_{A,D}$

Fix $P \in A(K)$ - tors

Let $Q \in A(\bar{K})$ and $mQ = P$.

$$\widehat{h}(Q) = \frac{1}{m^2} \underbrace{\widehat{h}(P)}_{\neq 0}$$

$$\begin{aligned}[K(Q):K] &\approx \# [m]^{-1}(P) \\ &= \# A[m] \\ &= m^{2g}, \quad g = \dim A\end{aligned}$$

$$\widehat{h}(Q) \approx \frac{1}{[K(Q):K]^g}$$

Lehmer's Conjecture

There is $C > 0$ s.t.,
for all $\alpha \in \overline{\mathbb{Q}}^* - \mu$,

$$h(\alpha) \geq \frac{C}{[\mathbb{Q}(\alpha) : \mathbb{Q}]}$$

Then (Dobrowolski 1979)

$$d = [\mathbb{Q}(\alpha) : \mathbb{Q}]$$

$$h(\alpha) \geq \frac{c}{d} \left(\frac{\log \log d}{\log d} \right)^3$$

Lehmer Conj for AVs
(Masser, 1984)

$$\exists C = C(A/K, D) > 0 \text{ s.t.}$$

$$\forall P \in A(\bar{K}) - \text{tors}$$

$$\hat{h}_D(P) \geq \frac{C}{[K(P): K]}^{1/g}$$

Conj

$$\widehat{h}_D(P) \geq \frac{C}{[K(P):K]}^{1/g}$$

Some known results: fix $\epsilon > 0$,

$$d = [K(P):K]$$

$$C = C(A/K, D, \epsilon) > 0$$

$P \in$ non-torsion pt $\in A(E)$

(a) $\widehat{h}(P) \geq C/d^{2g+1+\epsilon}$

(b) If A has CM, $\widehat{h}(P) \geq C/d^{1+\epsilon}$

(c) If $g=1$ and $\delta(A) \notin R_K$,

then $\widehat{h}(P) \geq C/d^{2+\epsilon}$

Really $d^\epsilon \rightsquigarrow$ power of $\log(d)$,

Proof methods:

- ① Transcendence Theory
- ② Fourier averaging
Harmonic analysis

$$L/K, d = [L : K] \quad B = \frac{C_1}{d^2}$$

$$A(L, B) = \{P \in A(L) : h_D(P) \leq B\}$$

Goal: Show

$$\# A(L, \frac{C_1}{d}) \leq C_2 d$$

Applies if to $\{m_P : m \leq \frac{\text{sup}}{ad}\}$
if $h^1(P)$ is small

Goal: Bd $A(L, B) = \{P \in A(L) : \hat{h}(P) \leq B\}$

$$A(L, B)^{(m)} := \{P_1 + \dots + P_m : P_i \in A(L, B)\}$$

$\# A(L, B)^{(m)} \approx \frac{\# A(L, B)^{(m)}}{m!}$ lots of pts

$$A(L, B)^{(m)} \subseteq A(L, m^2 B)$$

- ① Constructs metafine F
 s.f. F vanishes $A(L, B)^{(m)}$
 to high order

② Upper bd $|2F(Q)|$ for $Q \in$ Cauchy Reim.

- ③ Zeros estimate
 Lower bd for non-zero $A(L, B)$
- ④ ② & ③ \Rightarrow $A(L, B)$ is big

II Fix K , vary A/K .

Intuition: $\hat{h}(P) \neq 0$

$\Rightarrow \hat{h}(P) >$ complexity
of A/K

$h(A/K)$ = complexity of A/K

① $E: y^2 = x^3 + ax + b$

$h(E/K) = \min_{u \in K^\times} h([1, u^4 a, u^6 b])$

② Fix: $L_g \xrightarrow{\delta} \mathbb{P}^N$

$h(A/K) = h(f(A/K)) + \log N_{L_g/K}(\text{condenser of } A/K)$

③ Faltings

Conj: (Dem'janenko-Lang) $g = \dim A$
 $P \in A(K)$, and
 Then $\mathbb{Z} \cdot P$ are \mathbb{Z} -dense in A
 $\hat{h}_{A,D}(P) \geq C_1(K, g) h(A/K) - C_2(K, g)$

Results:

- ① True if $j(A) \in \Lambda_g(C)$
 if ε -distinct from the
 geo simple locns in $\Lambda_g(C)$
- ② True if fix A_0 , let
 $A = A_0^{x_0}$ twist
- ③ $g=1$. ABC conj \Rightarrow DL conj.

Fourier Averaging

$\{P_1, P_2, \dots, P_N\}$ small
ht.

$$\sum_{i \neq j} h(P_i - P_j)$$

$$= \sum_{v \in M_K}$$

$$\sum_{i \neq j} \lambda_v(P_i - P_j)$$

v good red ≥ 0
 v adlertee $\cancel{\geq 0}$
 v arch $\cancel{\geq 0}$

V non-arch, mult red.

$$E(K) \xrightarrow{t} E(K) \cong \frac{1}{m} \mathbb{Z}/\mathbb{Z}$$

$E_0(K)$

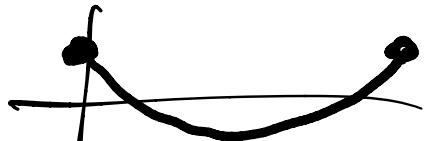
\hookrightarrow good red

$$M = \text{ord}_v(\Delta_E)$$

$$\lambda_r(P) = \frac{1}{2} B_2(t(P)) + \text{formal gp part}$$

$$B_2(T) = T^2 - T + \frac{1}{6} \geq 0$$

$T \in [0, 1]$



$$B_2(T) = \frac{1}{2\pi^2} \sum_{n \neq 0} \frac{e^{(nT)}}{n^2}$$

$e^{(ix)} = e^{2\pi ix}$

N^2 terms

$$\sum_{i \neq j} \lambda_{ij} (P_i - P_j)$$

$$\geq \sum_{i \neq j} \frac{1}{2} B_L(t(P_i - P_j))$$

$t(P_i) - t(P_j)$
 $t_i - t_j$

$$= \sum_{i \neq j} \frac{1}{2\pi^2} \frac{1}{2} \sum_{n \neq 0} \frac{e(t_i - t_j)}{n^2}$$

$$= \frac{1}{4\pi^2} \sum_{n \neq 0} \frac{1}{n^2} \sum_{i \neq j} e(t_i - t_j)$$

$$\sum_{i \neq j} e(t_i / e(-t_j)) \rightarrow N \sqrt{\sum_{i=1}^N e(t_i)^2}$$

$$L(A/K, s) \sim \mathcal{O}^{(r-1)^f} \dots$$

$$\underbrace{\text{Reg}(A/K)}_r \gg_{\min h_{A/K}(P)}^{\mathcal{R}} r$$

Conj E/\mathbb{Q} $y^2 = x^3 - ax + b$

$$\# E(\mathbb{Z}) \leq \mathcal{C}^{\text{rank } E(\mathbb{Q})}$$

Known Time

$$C \left(\frac{h(E/\mathbb{Q})}{\min h(P)} \right)$$