

CZ3007 Compiler Techniques

Review (Week 9)

Recap

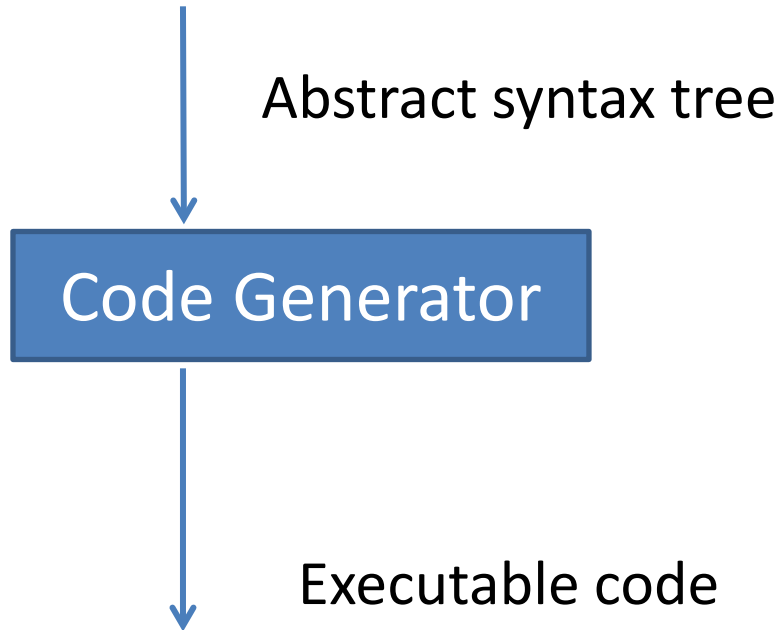
1. Introduction/Overview

- 1) What a code generator will do
- 2) Runtime environment of a program
- 3) Multi-pass process of generating and optimising code
- 4) Interpreting VM code and just-in-time compilation

2. Overview of the JVM (up to slide 36)

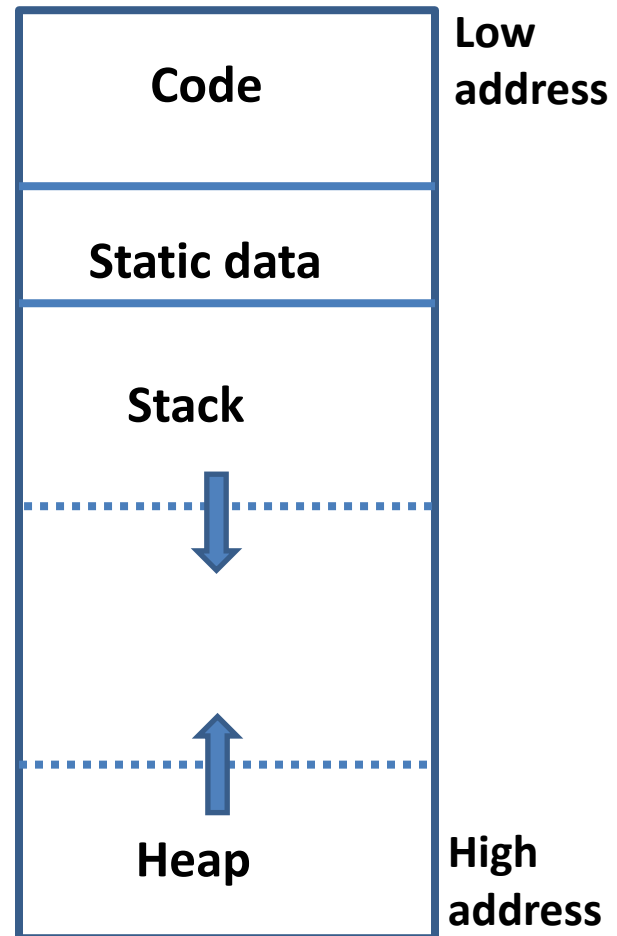
- 1) Runtime memory organisation
- 2) Bytecode
- 3) Stack-based computation

What a code generator will do



Runtime environment of a program

Memory layout



Multi-pass process of generating and optimising code



Slide 13

Interpreting VM code and just-in-time compilation



Poll 1

In the runtime memory layout of a program (in a language like C, C++)

1. The code area is of fixed size or variable size?
2. The static data area is of fixed size or variable size?
3. The stack area is of fixed size or variable size?
4. The heap area is of fixed size or variable size?

Poll 1

In the runtime memory layout of a program (in a language like C, C++)

1. The code area is of fixed size or variable size? **Fixed**
2. The static data area is of fixed size or variable size? **Fixed**
3. The stack area is of fixed size or variable size? **Variable**
4. The heap area is of fixed size or variable size? **Variable**

Poll 2

Is it possible that running code on a virtual machine like the Java Virtual Machine could be faster than running native code?

Poll 2

Is it possible that running code on a virtual machine like the Java Virtual Machine could be faster than running native code?

- **Optimization done by a compiler is based on static information. E.g.**
 - **avoid re-computation of some expressions.**
 - **$x = x^2; \Rightarrow x = x * x;$**
 - **$x = 5 * 2 * y; \Rightarrow x = 10 * y;$**
- **JVM can do more targeted optimization so it could be faster in some cases. E.g.**
 - **An example of more targeted optimization is that the JIT compilation may be able to use specific CPU model's features not generally available to achieve faster execution (like vector instructions)**
 - **Many optimizations are only feasible at runtime**

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2. Overview of the JVM (up to slide 36)

- 1) Runtime memory organisation
- 2) Bytecode
- 3) Stack-based computation

Poll 3

In the runtime memory organization of a Java program

1. Is the method area one per thread or shared?
2. Is the heap area one per thread or shared?
3. Is the stack area one per thread or shared?

Poll 3

In the runtime memory organization of a Java program

1. Is the method area one per thread or shared? **Shared.**
2. Is the heap area one per thread or shared? **Shared.**
3. Is the stack area one per thread or shared? **One per thread.**

Question

- Using JVM bytecode as an example. Why do we say that bytecode is often more compact than native code?



Instruction Set Format

31	2827	1615	87	0	Instruction type
Cond	0 0 I	Opcode	S	Rn Rd	Data processing / PSR Transfer
Cond	0 0 0 0 0 0	A S	Rd Rn	Rs 1 0 0 1 Rm	Multiply
Cond	0 0 0 0 1 U	A S	RdHi RdLo	Rs 1 0 0 1 Rm	Long Multiply (v3M / v4 only)
Cond	0 0 0 1 0 B	0 0	Rn Rd	0 0 0 0 1 0 0 1 Rm	Swap
Cond	0 1 I F U B	W L	Rn Rd	Offset	Load/Store Byte/Word
Cond	1 0 0 F U S	W L	Rn	Register List	Load/Store Multiple
Cond	0 0 0 F U 1	W L	Rn Rd	Offset1 1 S H 1 Offset2	Halfword transfer : Immediate offset (v4 only)
Cond	0 0 0 F U 0	W L	Rn Rd	0 0 0 0 1 S H 1 Rm	Halfword transfer: Register offset (v4 only)
Cond	1 0 1 L	Offset			Branch
Cond	0 0 0 1	0 0 1 0	1 1 1 1	1 1 1 1 0 0 0 1 Rn	Branch Exchange (v4T only)
Cond	1 1 0 F U	N W L	Rn CRd	CPNum Offset	Coprocessor data transfer
Cond	1 1 1 0	Op1	CRn CRd	CPNum Op2 0 CRm	Coprocessor data operation
Cond	1 1 1 0	Op1 L	CRn Rd	CPNum Op2 1 CRm	Coprocessor register transfer
Cond	1 1 1 1	SWI Number			Software interrupt

Please complete the online teaching feedback by 11.59pm, this Sunday.

Thank you!

