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Name : YASHRAJ SINHA Class :

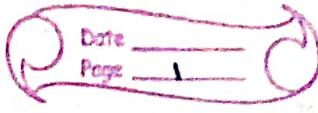
Section : Roll No. : 18M22CS335. Subject : Artificial Intelligence Lab

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LAB-1



① Implementing tic-tac-toe algorithm using python:

```
→ if isMaximizingPlayer:  
    bestValue = -infinity  
    for each child in node  
        value = minValue(child, depth+1, false)  
        bestValue = max(bestValue, value)  
    return bestValue  
  
else  
    bestValue = +infinity  
    for each child in node  
        value = maxValue(child, depth+1, true)  
        bestValue = min(bestValue, value)  
    return bestValue
```

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JAB - 2

2) Implementing vacuum cleaner using python

FUNCTION vacuum-world()

INITIALISE goal-state as ['A': '0', 'B': '0']
// 0: Clean, 1: Dirty

SET cost to 0

PROMPT for vacuum location (A/B) as location-input

PROMPT for status of location-input as status-input

SET other-location to 'B' if 'A' else 'A'

PROMPT for status of other-location as status-input-complement

IF location-input is 'A':

IF status-input is '1':

SET goal-state ['A'] to '0'

Increment cost by 2

IF status-input-complement is '1':

SET goal-state ['B'] to '0'

ELSE IF status-input-complement is '0':

Increment cost by 1

SET goal-state ['B'] to '0'

ELSE:

IF status-input is '1':

SET goal-state ['B'] to '0'

Increment cost by 2.

IF status-input-complement is '1':

SET goal-state ['A'] to '0'

ELSE

IF status-input-complement is '0':

Increment cost by 1

SET goal-state ['A'] to '0'

Print goal-state and cost

END FUNCTION

CALL Vacuum-world().

Algo

- 1) Initialize the agent's starting (x, y)
- 2) Loop until all cells are clean:
 - a. Perceive the current cell
 - b. If the cell is dirty
 - i. Clean the current cell
 - ii. Check surrounding cells (up, down, left, right) to see if any are dirty
 - iii. Move to the next dirty cell applying strategy such as BFS, DFS or other algorithm.
 - c. If no dirty cell are perceived, Stop.

3) End

Q, 18101907

3) Implementing 8-puzzle problem using python

BFS

Let fringe be a list containing the initial state.

LOOP

if fringe is empty return failure

Node ← remove-first (fringe)

if Node is goal

then return the path from initial state to Node and add generated nodes to the back of fringe

else generate all successors of Node and add generated node to the back of fringe.

END LOOP.

DFS

Let fringe be a list containing the initial state.

LOOP

if fringe is empty return failure

Node ← remove-first (fringe)

if Node is goal

then return the path from initial state to Node

else generate all successors of Node and add generated node to the front fringe.

END LOOP.

State Space tree : Using BFS

Initial

1 2 3
4 5 6
7 8 0

Final

1 2 3
4 5 6
7 8 0

Initial

Start

0 7 3

1 2 3

1 2 3
0 5 6

1 2 3
4 5 6

4 7 3

Visited

0 2 3
1 5 6
4 7 3

1 2 3
4 5 6
7 8 0

Final state.

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LAB - 3

4) Implementing A* algorithm using python.

→ A*

function A* search (problem) return a
solution or failure

node \leftarrow a node n with $n.state =$
problem.initial_state, $n.g = 0$

frontier \leftarrow a priority queue ordered
by ascending $g + h$, only
element n .

loop do

if empty? (frontier) then return
failure

$n \leftarrow \text{pop}(\text{frontier})$

if problem.goalTest ($n.state$) then
return solution (n)

for each action a in problem.
actions ($n.state$) do

$n' \leftarrow \text{ChildNode} (\text{problem}, n, a)$
insert (n' , $g(n') + h(n')$,
frontier)

(i) Using no. of misplaced tiles as heuristic
function

$$f(n) = g(n) + h(n)$$

1	2	3
4	5	6
0	7	8

Initial State

1	2	3
4	5	6
7	8	0

Goal state

Output

1	2	3
4	0	65
6	7	8

Initial state

1 2 3

1 2 3

4 5 0



4 5 8



4 5 8

6 7 8

6 7 0

6 0 7



(1 2 3)

1 2 3

1 2 3

5 0 8

0 5 8

4 5 8

4 6 7

4 6 7

0 6 7

5 6 8

5 6 8

5 6 0

4 0 7

4 0 7

4 7 8

1 2 3

1 2 3

1 2 3

5 6 8

5 6 8

5 6 0

4 0 7

4 0 7

4 7 8

1 2 3

1 2 3

1 2 3

4 5 6

0 5 6

5 0 6

0 7 8

0 7 8

4 7 8

1 2 3

1 2 3

1 2 3

4 5 6

0 5 6

5 0 6

0 7 8

0 7 8

4 7 8

1 2 3

1 2 3

1 2 3

4 5 6

0 4 5

5 0 6

0 7 8

0 7 8

4 7 8

1 2 3

1 2 3

1 2 3

4 5 6

0 4 5

5 0 6

0 7 8

0 7 8

4 7 8

1 2 3

1 2 3

1 2 3

4 5 6

0 4 5

5 0 6

0 7 8

0 7 8

4 7 8

25/10

i) Using Manhattan distance

→ Initial state

Goal state

1	2	3		1	2	3
4	0	5		4	5	6
7	8	6		7	8	0
F	E	A	O, F, S	S	E	A

Tile 1 : Manhattan distance (MD) = 0

Tile 3 : MD = 0

Tile 4 : MD = 0

Tile 2 : MD = 1

Tile 7 : MD = 1

Tile 8 : MD = 0

Tile G : MD = 1

$f(n) = \text{Total MD} = 3$

Tile 1 : Total MD = 0

Tile 2 : Total MD = 0

Tile 3 : Total MD = 0

Tile 4 : Total MD = 0

Tile 5 : Total MD = 2

Tile 6 : Total MD = 1

Tile 7 : Total MD = 0

Tile 8 : Total MD = 0

$f(n) = \text{Total MD} = 3$

$f(n) = \text{Total MD} = 0$

} 1st iteration

} 2nd iteration

→ Solution in 3rd iteration

1	2	3		1	2	3		1	2	3
4	0	5	→	4	5	0	→	4	5	6
7	8	6		7	8	6		7	3	0

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LAB - 04

Date _____
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- 5) Implementing Hill Climbing search algorithm to solve N-Queens problem.

→ function Hill-Climbing (problem) returns a state that is local maximum.

current \leftarrow Make-Node (problem, Initial-State)

loop do

neighbour \leftarrow a highest value successor of current

if neighbour.value \leq current.value
then return current.state

execute
current \leftarrow neighbour.

- 8) Show how the cost calculation of current state and neighbour nodes. And continue until you reach goal configuration of 4-Queens board.

→

state

$h(score)$

3, 1, 2, 0

2

1, 3, 2, 0

1

2, 1, 3, 0

1

0, 1, 2, 3

1

2, 3, 1, 0, 3

1

2, 0, 1, 2, 3

1

$HC_2 = 6$

6 swaps

Q.E.D.

→

State

$h(score)$

3, 1, 2, 0

2

1, 3, 2, 0

1

2, 1, 3, 0

1

0, 1, 2, 3

6

3, 2, 1, 0

6

3, 0, 2, 1

1

3, 1, 0, 2

1

state
space

0			
1			
2			
3			

$h = (2)$

0			
1			
2			

$h = (1)$

0			
1			
2			

$h = (-1)$

0			
1			
2			

$h = (6)$

0			
1			
2			

$h = (6)$

0			
1			
2			

$h = (1)$

3, 2, 1, 0

3, 0, 2, 1

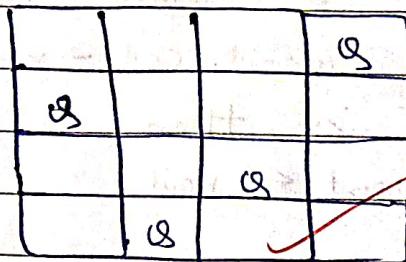


$3, 1, 0, 2$

$h = (1)$

• Next step

$1, 3, 2, 0$



$h = (1)$

state

h (score)

$1, 3, 2, 0$

1

$3, 1, 2, 0$

2

$2, 3, 1, 0$

2

$0, 3, 2, 1$

3

$1, 2, 3, 0$

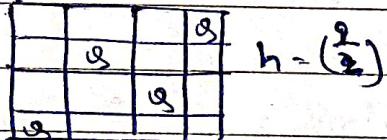
3

$1, 3, 0, 2$

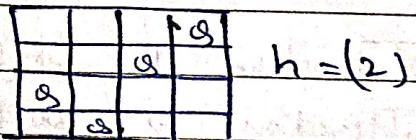
0

← return solution

state
space:



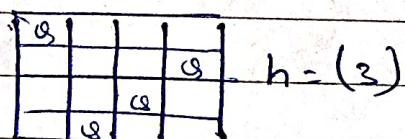
$h = \left(\frac{9}{2}\right)$



$h = (2)$

$3, 1, 2, 0$

$2, 3, 1, 0$



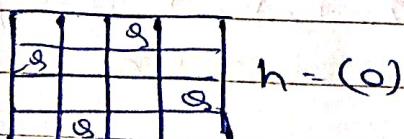
$h = (3)$

$0, 3, 2, 1$



$h = (3)$

$1, 2, 3, 0$



$h = (0)$

$1, 3, 0, 2$

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LAB - 5

- (6) Implement Simulated Annealing to solve N-Queens problem.

→

current ← initial state

T ← a large positive value

while $T > 0$ do

next ← a random neighbour of current

$\Delta E \leftarrow \text{current.cost} - \text{next.cost}$

if $\Delta E > 0$ then

current ← next

else

current ← next with probability $p = e^{\frac{\Delta E}{T}}$

end if

decrease T

end while

return current

Output

Iteration 1 : Conflicts = 24

Iteration 2 : Conflicts = 26

Iteration 4999 : Conflicts = 0

Iteration 5000 : Conflicts = 0

Final solution : [3, 6, 2, 7, 1, 4, 0, 5]

Number of conflicts in final solution : 0

⑦ Implement Unification in First Order Logic

→ i) If Ψ_1 and Ψ_2 is a variable or constant then :

a) If Ψ_1 and Ψ_2 are identical, then return NIL.

b) Else if Ψ_1 is a variable,

- then if Ψ_1 occurs in Ψ_2 then return failure

- Else return $\{(\Psi_2/\Psi_1)\}$.

c) Else if Ψ_2 is a variable,

- If Ψ_2 occurs in Ψ_1 then return failure

- Else return $\{(\Psi_1/\Psi_2)\}$.

d) Else return failure.

ii) If the initial Predicate symbol in Ψ_1 and Ψ_2 are not same, then return failure

iii) If Ψ_1 and Ψ_2 have a different number of arguments, then return failure.

iv) Set ~~substitution set (SUBST)~~ to NIL.

v) for $i=1$ to number of elements in Ψ_1

a) Call Unify function with the i th element of Ψ_1 and i th element of Ψ_2 and

put the result into S.

(b) If S = failure then return Failure.

(c) If S ≠ NIL then do.

- Apply S to the remainder of both L₁ and L₂.
- SUBST = APPEND(S, SUBST).

vi) Return SUBST.

↓ Execute

Output

Expansion 1: ("Eats", "x", "Mango", "Pizza")

Expansion 2: ("Eats", "Sumit", "y", "z")

Substitutions: { 'x': 'Sumit', 'y': 'Mango',
'z': 'Pizza' }

Ques.

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KAB-7

8) Create a knowledge base consisting of first order logic statements and prove the given query using forward reasoning.

→ Function $\text{FOL-FC-Ask}(\text{KB}, \alpha)$ returns a substitution or false

inputs: KB , the knowledge base, a set of first-order definite clauses
 α , the query, an atomic sentence.

local variables: new, the new sentences inferred on each iteration.

repeat until new is empty

$\text{new} \leftarrow \{\}$

for each rule is empty in KB do

$(p_1 \wedge \dots \wedge p_n \Rightarrow q) \leftarrow \text{STANDARDIZE-VARIABLES}(\text{rule})$

for each θ such that $\text{SUBST}(\theta, p_1 \wedge \dots \wedge p_n) = \text{SUBST}(\theta, p'_1 \wedge \dots \wedge p'_n)$

for some p'_1, \dots, p'_n in KB

$q' \leftarrow \text{SUBST}(\theta, q)$

if q' does not unify with some sentence already in KB or new then add q' to new

$\theta \leftarrow \text{unify}(q', \alpha)$

if ϕ is not fail then return

add new to KB

return false

Output

Criminal (Robert) is proven!

Factored facts:

Missile (T_1)

Weapon (T_1)

Noctile (A)

Owes (A, T_1)

Criminal (Robert)

Enemy (A, American)

American (Robert)

Sells (Robert, T_1 , A)

✓
29/11/20

Q3) Consider a vocabulary with the following symbols

(a) Occupation (Emily, Surgeon) \vee Occupation (Emily, Lawyer)

(b) Occupation (Joe, Actor) \wedge \exists_o (Occupation (Joe, o) \wedge $o \neq$ Actor)

(c) \forall_p (Occupation (p, Surgeon) \rightarrow Customer (Joe, p))

(d) \exists_p (Occupation (p, Lawyer) \wedge Customer (Joe, p))

⑥ $\exists_p (\text{Boss}(p, \text{Emily}) \wedge \text{Occupation}(p, \text{Lawyer}))$

⑦ $\exists_p (\text{Occupation}(p, \text{Lawyer}) \wedge \forall_c (\text{Customer}(c, p) \rightarrow \text{Occupation}(c, \text{Doctor})))$

⑧ $\forall_p (\text{Occupation}(p, \text{Surgeon}) \rightarrow \exists_l (\text{Occupation}(l, \text{Lawyer}) \wedge \text{Customer}(p, l)))$

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19/12/24

- 9) Create a knowledge base using Propositional logic and prove the given query using resolution.

Function: resolution(KB , query): return query if true or false

input: KB , the Knowledge base, set of propositions
query: a state to be proven

classes = convert to CNF (KB , negated query = negate(query))

new-clause = set()

apply the resolution rules:

- select two clauses contain complementary clauses
- resolve two to form a new clause
- add the new clause is empty? { contradiction is found }

if new-clause = {} print "query is true"
else print "false"

Output:

$KB : P \vee Q$

To prove: P is true

$\neg Q \vee R \perp$

$\neg P \vee \neg Q \perp$

$\neg R \vee \neg \neg P$

Resolution $\neg R \perp$

$(\neg Q \vee R), (\neg R) \rightarrow \neg Q \perp$

$(\neg P \vee \neg Q), (\neg Q) \rightarrow \neg P$

$$(\neg R \vee P), (\neg r) \rightarrow \neg R$$

contradiction / is query in true.

10) Convert a given FOL to CNF

→

Function FOL-to-CNF :

Replace $(P \leftrightarrow Q)$ with $(P \rightarrow Q) \wedge (Q \rightarrow P)$

Replace $(P \rightarrow Q)$ with $(\neg P \vee Q)$

Replace $\neg(P \vee Q)$ with $\neg P \wedge \neg Q$

Replace $\neg(P \wedge Q)$ with $\neg P \vee \neg Q$

Replace $\neg\neg P$ with P

Rename all variable to make them unique

Replace $\exists_n [p(n)]$ with $P(c)$ if n is ~~an~~ independent

Replace $\exists_n [p(n)]$ with $p(y) \vee P(j(y))$ if n depends on a universally quantified variable.

Then Distribute \vee over \wedge :

Ensure all formulas is a conjunct
of disjunction

Output

$$\forall n (\exists y (j(n, y) \rightarrow \phi(y)) \rightarrow r(n))$$

$$y \neq j(x) \quad y = j(u)$$

$$\forall n ((\neg p(n, j(n)) \vee \phi(y)) \wedge \neg r(n))$$

$$\Rightarrow (\neg p(n, j(n)) \vee (j(n) \wedge \neg r(n))$$

Q1) Implement alpha-beta pruning

→ alpha : the best option for the maximizer
 Beta : the best option for the minimizer
 pruning : Stop exploring sub-tree when it is clear that they won't affect the outcome.

Function alphabeta (node, depth, alpha, beta, is-maximizing)

if (depth == 0) or node is a terminal node :

return evaluate (node)

if is-maximizing () :

maneval = -∞

for each child in children (node) :

and = alphabeta (child, depth-1, alpha, beta, false)

maneval = max (maneval, eval)

alpha = max (alpha, eval)

if alpha >= beta :

~~break (prune)~~

~~return maneval.~~

else :

minval = +∞

for each child in children (node) :

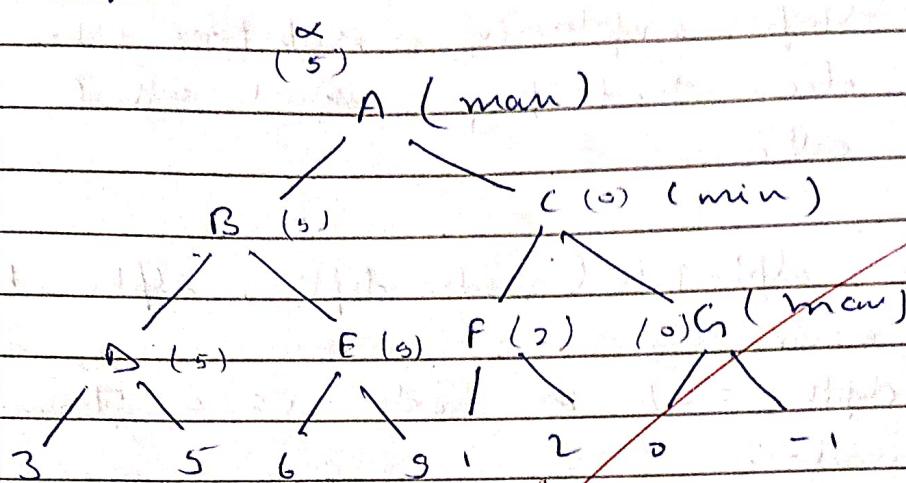
eval = alphabeta (child, depth-1, alpha, beta, eval)

beta = min (beta, eval)

if alpha >= beta :

break (prune)
return min val

Output



1) Knowledge-base using proposition logic / Entails

→ Initialize KB with proposition logic statement

If forward-chaining (KB, query) :
print "Query entailed by Knowledge base"
else

print "Query is not entailed by KB"

Initialize Forward Chaining (KB, query)

Initialize agenda with known facts from KB.

While agenda is not empty

POP fact from agenda.

If fact matches query :
return true

For each rule in KB

If fact satisfies rule's premise
add rules to agenda.

return false.

Output Q 1 2 1 2 u
2 0 1 1 2 1 2 u

For the KB = [A, B, A ∨ B => C, C => D]
query D.

Query established by Knowledge Base.