$tmvar$, x, y, f, m, n	variables
$covar$, c	coercion variables
$datacon$, K	
$const$, T, F, Age	
$index$, i	indices

relflag, ρ	::=			relevance flag
	|	$+$		
	|	$-$		
	|	$\mathbf{app\_rho}\,\nu$	S	
	|	(ρ)	S	
appflag, ν	::=			applicative flag
	|	R		
	|	ρ		
role, R	::=			Role
	|	$\mathbf{Nom}$		
	|	$\mathbf{Rep}$		
	|	$R_1 \cap R_2$	S	
	|	$\mathbf{param}\,R_1\ R_2$	S	
	|	$\mathbf{app\_role}\,\nu\ R$	S	
	|	(R)	S	
constraint, ϕ	::=			props
	|	$a \sim_R b : A$		
	|	(ϕ)	S	
	|	$\phi\{b/x\}$	S	
	|	$\|\phi\|$	S	
	|	$a \sim_R b$	S	
	|	$a \sim b$	S	
tm, a, b, p, v, w, A, B, C	::=			types and kinds
	|	$\star$		
	|	x		
	|	$\lambda^{\rho} x\!:\!A.b$	bind x in b	
	|	$\lambda^{\rho} x.b$	bind x in b	
	|	$a\ b^{\nu}$		
	|	$\Pi^{\rho} x\!:\!A \to B$	bind x in B	
	|	$\Lambda c\!:\!\phi.b$	bind c in b	
	|	$\Lambda c.b$	bind c in b	
	|	$a[\gamma]$		
	|	$\forall c\!:\!\phi.B$	bind c in B	
	|	$a \triangleright_R \gamma$		
	|	F		
	|	$\square$		
	|	$\mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \| _{-} \to b_2$		
	|	$\mathbf{apply}\ a\,\overline{\mu}$	M	
	|	K		
	|	$\mathbf{match}\ a\ \mathbf{with}\ brs$		
	|	$\mathbf{sub}\ R\ a$		
	|	$\mathbf{coerce}\ a$		
	|	$a\ b$		

	|	$a\{b/x\}$	S	
	|	$a\{\gamma/c\}$	S	
	|	$a\{b/x\}$	S	
	|	$a\{\gamma/c\}$	S	
	|	a	S	
	|	a	S	
	|	(a)	S	
	|	a	S	parsing precedence is hard
	|	$\|a\|_R$	S	
	|	**Int**	S	
	|	**Bool**	S	
	|	Nat	S	
	|	**Vec**	S	
	|	0	S	
	|	S	S	
	|	**True**	S	
	|	**Fix**	S	
	|	**Maybe**	S	
	|	**Just**	S	
	|	**Nothing**	S	
	|	$a \to b$	S	
	|	$\phi \Rightarrow A$	S	
	|	$\lambda x.a$	S	
	|	$\lambda x : A.a$	S	
	|	$\forall\, x : A \to B$	S	
	|	**if** ϕ **then** a **else** b	S	
brs	::=			case branches
	|	**none**		
	|	$K \Rightarrow a; brs$		
	|	$brs\{a/x\}$	S	
	|	$brs\{\gamma/c\}$	S	
	|	(brs)	S	
$co,\ \gamma$	::=			explicit coercions
	|	$\bullet$		
	|	c		
	|	**red** $a\ b$		
	|	**refl** a		
	|	$(a \mid=\mid_{\gamma}\ b)$		
	|	**sym** γ		
	|	$\gamma_1; \gamma_2$		
	|	**sub** γ		
	|	$\Pi^{R,\rho} x : \gamma_1.\gamma_2$	bind x in γ_2	
	|	$\lambda^{R,\rho} x : \gamma_1.\gamma_2$	bind x in γ_2	
	|	$\gamma_1\ \gamma_2^{R,\rho}$		

	|	$\mathbf{piFst}\,\gamma$		
	|	$\mathbf{cpiFst}\,\gamma$		
	|	$\mathbf{isoSnd}\,\gamma$		
	|	$\gamma_1 @ \gamma_2$		
	|	$\forall c{:}\gamma_1.\gamma_3$	bind c in γ_3	
	|	$\lambda c{:}\gamma_1.\gamma_3 @ \gamma_4$	bind c in γ_3	
	|	$\gamma(\gamma_1, \gamma_2)$		
	|	$\gamma @ (\gamma_1 \sim \gamma_2)$		
	|	$\gamma_1 \triangleright_R \gamma_2$		
	|	$\gamma_1 \sim_A \gamma_2$		
	|	$\mathbf{conv}\ \phi_1 \sim_\gamma \phi_2$		
	|	$\mathbf{eta}\,a$		
	|	$\mathbf{left}\,\gamma\,\gamma'$		
	|	$\mathbf{right}\,\gamma\,\gamma'$		
	|	(γ)	S	
	|	γ	S	
	|	$\gamma\{a/x\}$	S	
υ	::=			
	|	$\mathbf{Tm}\,\nu$		
	|	$\mathbf{Co}$		
$\overline{\upsilon}$	::=			
	|	$\mathbf{emptyA}$		
	|	$\overline{\upsilon}, \upsilon$	M	
	|	$\upsilon, \overline{\upsilon}$		
	|	$(\overline{\upsilon})$	M	
role_context, Ω	::=			role$_c$*ontexts*
	|	$\varnothing$		
	|	$x : R$		
	|	$\Omega, x : R$		
	|	Ω, Ω'	M	
	|	$\mathbf{var\_pat}\,p$	M	
	|	(Ω)	M	
	|	Ω	M	
roles, Rs	::=			
	|	$\cdot$		
	|	R, Rs		
	|	$\mathbf{range}\,\Omega$	S	
	|	(Rs)	S	
	|	$Rs \mathbin{++} Rs'$	S	
sig_sort	::=			signature classifier
	|	$_{Rs}\ A$		
	|	$_{Rs}\ p \sim_R a : A$		

sort	::=			binding classifier
	|	$\mathbf{Tm}\, A$		
	|	$\mathbf{Co}\, \phi$		
$context,\ \Gamma$	::=			contexts
	|	$\varnothing$		
	|	$\Gamma, x : A$		
	|	$\Gamma, c : \phi$		
	|	$\Gamma\{b/x\}$	M	
	|	$\Gamma\{\gamma/c\}$	M	
	|	Γ, Γ'	M	
	|	$\|\Gamma\|$	M	
	|	(Γ)	M	
	|	Γ	M	
$sig,\ \Sigma$	::=			signatures
	|	$\varnothing$		
	|	$\Sigma \cup \{F : sig\_sort\}$		
	|	Σ_0	M	
	|	Σ_1	M	
	|	$\|\Sigma\|$	M	
$available\_props,\ \Delta$	::=			
	|	$\varnothing$		
	|	Δ, x		
	|	Δ, c		
	|	$\mathsf{fv}\, a$	M	
	|	Δ, Δ'	M	
	|	$\widetilde{\Gamma}$	M	
	|	$\widetilde{\Omega}$	M	
	|	(Δ)	M	
$Nat,\ \mathbb{N}$	::=			
	|	0	M	
	|	$\Sigma\, \mathbb{N}$	M	
	|	$(\mathbb{N})$	M	
	|	$\|a\|$	S	
μ	::=			Pattern arguments
	|	$\mathbf{Tm}\, \nu\, a$		
	|	$\mathbf{Co}\, \gamma$		
$\overline{\mu},\ PA$	::=			
	|	**none**		
	|	μ, PA		
	|	PA, μ		

$$
\begin{array}{lcl}
\mathit{terminals} & ::= & \\
 & | & \leftrightarrow \\
 & | & \Leftrightarrow \\
 & | & \longrightarrow \\
 & | & \mathsf{min} \\
 & | & \equiv \\
 & | & \forall \\
 & | & \in \\
 & | & \notin \\
 & | & \Leftarrow \\
 & | & \Rightarrow \\
 & | & \Rightarrow^{*} \\
 & | & \rightarrow \\
 & | & \Lambda \\
 & | & \\
 & | & \square \\
 & | & \vdash \\
 & | & \dashv \\
 & | & \models \\
 & | & \models \\
 & | & \models_{\mathsf{src}} \\
 & | & \neq \\
 & | & \triangleright \\
 & | & \ \mathsf{ok} \\
 & | & \_ \\
 & | & \leadsto \\
 & | & \leadsto^{*} \\
 & | & \leadsto \\
 & | & \varnothing \\
 & | & \circ \\
 & | & \mathsf{fv} \\
 & | & \mathsf{dom} \\
 & | & \sim \\
 & | & \asymp \\
 & | & | \\
 & | & \bullet \\
 & | & \mathbf{fst} \\
 & | & \mathbf{snd} \\
 & | & \mathbf{as} \\
 & | & |\Rightarrow| \\
 & | & \vdash_{=} \\
 & | & \mathbf{refl_2} \\
 & | & ++ \\
 & | & \{ \\
 & | & \} \\
 & | & \not\rightarrow
\end{array}
$$

$formula,\ \psi$	$::=$		
	$\vert$	$judgement$	
	$\vert$	$x : A \ \in\ \Gamma$	
	$\vert$	$x : R \ \in\ \Omega$	
	$\vert$	$c : \phi \ \in\ \Gamma$	
	$\vert$	$F : sig\_sort \ \in\ \Sigma$	
	$\vert$	$x \ \in\ \Delta$	
	$\vert$	$c \ \in\ \Delta$	
	$\vert$	$c\,\textbf{not relevant} \ \in\ \gamma$	
	$\vert$	$x \ \notin\ \Delta$	
	$\vert$	$uniq\,\Gamma$	
	$\vert$	$uniq(\Omega)$	
	$\vert$	$c \ \notin\ \Delta$	
	$\vert$	$T \ \notin\ \mathsf{dom}\,\Sigma$	
	$\vert$	$F \ \notin\ \mathsf{dom}\,\Sigma$	
	$\vert$	$\mathbb{N}_1 < \mathbb{N}_2$	
	$\vert$	$\mathbb{N}_1 \leq \mathbb{N}_2$	
	$\vert$	$\nu = \rho$	
	$\vert$	$R_1 = R_2$	
	$\vert$	$a = b$	
	$\vert$	$\phi_1 = \phi_2$	
	$\vert$	$\Gamma_1 = \Gamma_2$	
	$\vert$	$\gamma_1 = \gamma_2$	
	$\vert$	$\neg\psi$	
	$\vert$	$\psi_1 \wedge \psi_2$	
	$\vert$	$\psi_1 \vee \psi_2$	
	$\vert$	$\psi_1 \Rightarrow \psi_2$	
	$\vert$	(ψ)	
	$\vert$	ψ	
	$\vert$	$c : (a : A \sim b : B) \ \in\ \Gamma$	
	$\vert$		suppress lc hypothesis generated by Ott
	$\vert$	$\Delta\#\Delta_2$	
$JSubRole$	$::=$		
	$\vert$	$R_1 \leq R_2$	Subroling judgement
$JRolePath$	$::=$		
	$\vert$	$\mathsf{RolePath}\ a = F@Rs$	Type headed by constant (partial function)
$JAppsPath$	$::=$		
	$\vert$	$\mathsf{AppsPath}\,a \nleftrightarrow F\overline{v}$	Type headed by constant (partial function)
$JSat$	$::=$		
	$\vert$	$\overline{v} \sim Rs$	
	$\vert$	$\textbf{Sat}\,F\,\overline{v}$	
$JPatCtx$	$::=$		

	|	$\Omega;\Gamma \vDash p :_F B \Rightarrow A$	Contexts generated by a p
JRename	::=		
	|	rename $p \rightarrow a$ to $p' \rightarrow a'$ excluding Δ and Δ'	rename with fresh variable
JMatchSubst	::=		
	|	match a_1 with $p \rightarrow b_1 \nrightarrow b_2$	match and substitute
JPatData	::=		
	|	**uncurry** $p = F@PA$	Pattern data (head argun
JIsPattern	::=		
	|	**pattern** p	
JSubPat	::=		
	|	**subpattern** p' p	Subpattern
JTmPatternAgree	::=		
	|	$a \leftrightarrow p$	term and pattern agree
JTmSubPatternAgree	::=		
	|	$a \sqsubseteq p$	sub-pattern agrees with te
JSubTmPatternAgree	::=		
	|	$a \sqsupseteq p$	sub-term agrees with patt
JValuePath	::=		
	|	ValuePath $a \nrightarrow F$	Type headed by constant
JCasePath	::=		
	|	$\mathsf{CasePath}_R$ $a \nrightarrow F$	Type headed by constant
JApplyArgs	::=		
	|	apply args a to $b \nrightarrow b'$	apply arguments of a (hea
JValue	::=		
	|	Value_R A	values
JValueType	::=		
	|	$\mathsf{ValueType}_R$ A	Types with head forms (e
Jconsistent	::=		
	|	$\mathsf{consistent}_R$ a b	(erased) types do not diffe
Jroleing	::=		
	|	$\Omega \vDash a : R$	Roleing judgment

JChk	::=		
	|	$(\rho = +) \vee (x \notin \mathsf{fv}\ A)$	irrelevant argument check
Jpar	::=		
	|	$\Omega \vDash a \Rightarrow_R b$	parallel reduction (implicit language)
	|	$\Omega \vDash a \Rightarrow_R^* b$	multistep parallel reduction
	|	$\Omega \vDash a \Leftrightarrow_R b$	parallel reduction to a common term
Jbeta	::=		
	|	$\vDash a >_R b$	primitive reductions on erased terms
	|	$\vDash a \rightsquigarrow_R b$	single-step head reduction for implicit languag
	|	$\vDash a \rightsquigarrow^* b/R$	multistep reduction
JBranchTyping	::=		
	|	$\Gamma \vDash \mathsf{case}\ a : A\ \mathsf{of}\ b : B \Rightarrow C \mid C'$	Branch Typing (aligning the types of case)
Jett	::=		
	|	$\Gamma \vDash \phi\ \mathsf{ok}$	Prop wellformedness
	|	$\Gamma \vDash a : A$	typing
	|	$\Gamma; \Delta \vDash \phi_1 \equiv \phi_2$	prop equality
	|	$\Gamma; \Delta \vDash a \equiv b : A/R$	definitional equality
	|	$\vDash \Gamma$	context wellformedness
Jsig	::=		
	|	$\vDash \Sigma$	signature wellformedness
Jhiding	::=		
	|	$Rs_1 \Leftarrow Rs_2$	
	|	$\Sigma_1 \Leftarrow \Sigma_2$	
JSrc	::=		
	|	$\Gamma \vDash_{\mathsf{src}} a : A$	source typing
	|	$\Gamma \vDash_{\mathsf{src}} a \rightsquigarrow a' : A$	source translation
	|	$\Gamma \vDash_{\mathsf{src}} \phi \rightsquigarrow \phi'\ \mathsf{ok}$	Prop wellformedness
Jann	::=		
	|	$\Gamma \vdash \phi\ \mathsf{ok}$	prop wellformedness
	|	$\Gamma \vdash a : A/R$	typing
	|	$\Gamma; \Delta \vdash \gamma : \phi_1 \sim \phi_2$	coercion between props
	|	$\Gamma; \Delta \vdash \gamma : A \sim_R B$	coercion between types
	|	$\vdash \Gamma$	context wellformedness
Jred	::=		
	|	$\Gamma \vdash a \rightsquigarrow b/R$	single-step, weak head reduction to values for
judgement	::=		
	|	*JSubRole*	

$$
\begin{array}{lll}
 & | & JRolePath \\
 & | & JAppsPath \\
 & | & JSat \\
 & | & JPatCtx \\
 & | & JRename \\
 & | & JMatchSubst \\
 & | & JPatData \\
 & | & JIsPattern \\
 & | & JSubPat \\
 & | & JTmPatternAgree \\
 & | & JTmSubPatternAgree \\
 & | & JSubTmPatternAgree \\
 & | & JValuePath \\
 & | & JCasePath \\
 & | & JApplyArgs \\
 & | & JValue \\
 & | & JValueType \\
 & | & Jconsistent \\
 & | & Jroleing \\
 & | & JChk \\
 & | & Jpar \\
 & | & Jbeta \\
 & | & JBranchTyping \\
 & | & Jett \\
 & | & Jsig \\
 & | & Jhiding \\
 & | & JSrc \\
 & | & Jann \\
 & | & Jred
\end{array}
$$

$$
\begin{array}{lll}
user\_syntax & ::= & \\
 & | & tmvar \\
 & | & covar \\
 & | & datacon \\
 & | & const \\
 & | & index \\
 & | & relflag \\
 & | & appflag \\
 & | & role \\
 & | & constraint \\
 & | & tm \\
 & | & brs \\
 & | & co \\
 & | & v \\
 & | & \overline{v} \\
 & | & role\_context
\end{array}
$$

| *roles*
| *sig_sort*
| *sort*
| *context*
| *sig*
| *available_props*
| *Nat*
| μ
| $\overline{\mu}$
| *terminals*
| *formula*

$\boxed{R_1 \leq R_2}$ Subroling judgement

$$\frac{}{\mathbf{Nom} \leq R} \quad \textsc{NomBot}$$

$$\frac{}{R \leq \mathbf{Rep}} \quad \textsc{RepTop}$$

$$\frac{}{R \leq R} \quad \textsc{Refl}$$

$$\frac{\begin{array}{c} R_1 \leq R_2 \\ R_2 \leq R_3 \end{array}}{R_1 \leq R_3} \quad \textsc{Trans}$$

$\boxed{\mathsf{RolePath}\ a = F@Rs}$ Type headed by constant (partial function)

$$\frac{F :_{Rs}\ A \in \Sigma_0}{\mathsf{RolePath}\ F = F@Rs} \quad \textsc{RolePath\_AbsConst}$$

$$\frac{F :_{Rs}\ p \sim_{R_1} a : A \in \Sigma_0}{\mathsf{RolePath}\ F = F@Rs} \quad \textsc{RolePath\_Const}$$

$$\frac{\mathsf{RolePath}\ a = F@R_1, Rs}{\mathsf{RolePath}\ (a\ b'^{R_1}) = F@Rs} \quad \textsc{RolePath\_TApp}$$

$$\frac{\mathsf{RolePath}\ a = F@Rs}{\mathsf{RolePath}\ (a\ b'^{+}) = F@Rs} \quad \textsc{RolePath\_App}$$

$$\frac{\mathsf{RolePath}\ a = F@Rs}{\mathsf{RolePath}\ (a\ \square^{-}) = F@Rs} \quad \textsc{RolePath\_IApp}$$

$$\frac{\mathsf{RolePath}\ a = F@Rs}{\mathsf{RolePath}\ (a[\bullet]) = F@Rs} \quad \textsc{RolePath\_CApp}$$

$\boxed{\mathsf{AppsPath} a \not\leftrightarrow F\overline{v}}$ Type headed by constant (partial function)

$$\frac{F :_{Rs}\ A \in \Sigma_0}{\mathsf{AppsPath} F \not\leftrightarrow F\mathbf{emptyA}} \quad \textsc{AppsPath\_AbsConst}$$

$$\frac{\begin{array}{l} F :_{Rs}\ p \sim_{R_1} a : A \in \Sigma_0 \\ \neg(R_1 \leq R) \end{array}}{\mathsf{AppsPath} F \not\leftrightarrow F\mathbf{emptyA}} \quad \textsc{AppsPath\_Const}$$

$$\frac{\mathsf{AppsPath} a \not\leftrightarrow F\overline{v}}{\mathsf{AppsPath}(a\ b'^{R_1}) \not\leftrightarrow F(\overline{v}, \mathbf{Tm}\, R_1)} \quad \textsc{AppsPath\_App}$$

$$\frac{\mathsf{AppsPath}\, a \nrightarrow F\overline{v}}{\mathsf{AppsPath}(a\ b^-) \nrightarrow F(\overline{v}, \mathbf{Tm}\, -)} \quad \textsc{AppsPath\_IApp}$$

$$\frac{\mathsf{AppsPath}\, a \nrightarrow F\overline{v}}{\mathsf{AppsPath}(a[\bullet]) \nrightarrow F(\overline{v}, \mathbf{Co})} \quad \textsc{AppsPath\_CApp}$$

$\boxed{\overline{v} \sim Rs}$

$$\frac{}{\mathbf{emptyA} \sim \cdot} \quad \textsc{AR\_nil}$$

$$\frac{\overline{v} \sim Rs}{\mathbf{Tm}\, R_1, \overline{v} \sim R_1, Rs} \quad \textsc{AR\_consTApp}$$

$$\frac{\overline{v} \sim Rs}{\mathbf{Tm}\, +, \overline{v} \sim Rs} \quad \textsc{AR\_consApp}$$

$$\frac{\overline{v} \sim Rs}{\mathbf{Tm}\, -, \overline{v} \sim Rs} \quad \textsc{AR\_consIApp}$$

$$\frac{\overline{v} \sim Rs}{\mathbf{Co}, \overline{v} \sim Rs} \quad \textsc{AR\_consCApp}$$

$\boxed{\mathbf{Sat}\, F\, \overline{v}}$

$$\frac{\begin{array}{l} F :_{Rs}\ A \in \Sigma_0 \\ \overline{v} \sim Rs \end{array}}{\mathbf{Sat}\, F\, \overline{v}} \quad \textsc{Sat\_Const}$$

$$\frac{\begin{array}{l} F :_{Rs}\ p \sim_{R_1} a_0 : A_1 \in \Sigma_0 \\ \neg(R_1 \leq \mathbf{Nom}) \\ \overline{v} \sim Rs \end{array}}{\mathbf{Sat}\, F\, \overline{v}} \quad \textsc{Sat\_Axiom}$$

$\boxed{\Omega; \Gamma \vDash p :_F B \Rightarrow A}$ Contexts generated by a pattern (variables bound by the pattern)

$$\frac{}{\varnothing; \varnothing \vDash F :_F A \Rightarrow A} \quad \textsc{PatCtx\_Const}$$

$$\frac{\Omega; \Gamma \vDash p :_F \Pi^+ x\!:\!A' \to A \Rightarrow B}{\Omega, x : R; \Gamma, x : A' \vDash p\ x^R :_F A \Rightarrow B} \quad \textsc{PatCtx\_PiRel}$$

$$\frac{\Omega; \Gamma \vDash p :_F \Pi^- x\!:\!A' \to A \Rightarrow B}{\Omega; \Gamma, x : A' \vDash p\ \Box^- :_F A \Rightarrow B} \quad \textsc{PatCtx\_PiIrr}$$

$$\frac{\Omega; \Gamma \vDash p :_F \forall c\!:\!\phi.A \Rightarrow B}{\Omega; \Gamma, c : \phi \vDash p[\bullet] :_F A \Rightarrow B} \quad \textsc{PatCtx\_CPi}$$

$\boxed{\mathsf{rename}\ p \to a\ \mathsf{to}\ p' \to a'\ \mathsf{excluding}\ \Delta\ \mathsf{and}\ \Delta'}$ rename with fresh variables

$$\frac{}{\mathsf{rename}\ F \to a\ \mathsf{to}\ F \to a\ \mathsf{excluding}\ \Delta\ \mathsf{and}\ \varnothing} \quad \textsc{Rename\_Base}$$

$$\frac{\begin{array}{l} \mathsf{rename}\ p_1 \to a_1\ \mathsf{to}\ p_2 \to a_2\ \mathsf{excluding}\ \Delta\ \mathsf{and}\ \Delta' \\ y \notin (\Delta, \Delta') \end{array}}{\mathsf{rename}\ (p_1\ x^R) \to a_1\ \mathsf{to}\ (p_2\ y^R) \to (a_2\{y/x\})\ \mathsf{excluding}\ \Delta\ \mathsf{and}\ (\Delta', y)} \quad \textsc{Rename\_AppRel}$$

$$\frac{\mathsf{rename}\ p_1 \to a_1\ \mathsf{to}\ p_2 \to a_2\ \mathsf{excluding}\ \Delta\ \mathsf{and}\ \Delta'}{\mathsf{rename}\ (p_1\ \Box^-) \to a_1\ \mathsf{to}\ (p_2\ \Box^-) \to a_2\ \mathsf{excluding}\ \Delta\ \mathsf{and}\ \Delta'} \quad \textsc{Rename\_AppIrrel}$$

$$\frac{\textsf{rename}\ p_1 \to a_1\ \textsf{to}\ p_2 \to a_2\ \textsf{excluding}\ \Delta\ \textsf{and}\ \Delta'}{\textsf{rename}\ (p_1[\bullet]) \to a_1\ \textsf{to}\ (p_2[\bullet]) \to a_2\ \textsf{excluding}\ \Delta\ \textsf{and}\ \Delta'} \quad \textsc{Rename\_CApp}$$

$\boxed{\textsf{match}\ a_1\ \textsf{with}\ p \to b_1 \nrightarrow b_2}$ match and substitute

$$\frac{}{\textsf{match}\ F\ \textsf{with}\ F \to b \nrightarrow b} \quad \textsc{MatchSubst\_Const}$$

$$\frac{\textsf{match}\ a_1\ \textsf{with}\ p_1 \to b_1 \nrightarrow b_2}{\textsf{match}\ (a_1\ \ a^R)\ \textsf{with}\ (p_1\ \ x^R) \to b_1 \nrightarrow (b_2\{a/x\})} \quad \textsc{MatchSubst\_AppRelR}$$

$$\frac{\textsf{match}\ a_1\ \textsf{with}\ p_1 \to b_1 \nrightarrow b_2}{\textsf{match}\ (a_1\ \ a^-)\ \textsf{with}\ (p_1\ \ \Box^-) \to b_1 \nrightarrow b_2} \quad \textsc{MatchSubst\_AppIrrel}$$

$$\frac{\textsf{match}\ a_1\ \textsf{with}\ a_2 \to b_1 \nrightarrow b_2}{\textsf{match}\ (a_1[\bullet])\ \textsf{with}\ (a_2[\bullet]) \to b_1 \nrightarrow b_2} \quad \textsc{MatchSubst\_CApp}$$

$\boxed{\mathbf{uncurry}\, p = F@PA}$ Pattern data (head arguments)

$$\frac{}{\mathbf{uncurry}\, F = F@\mathbf{none}} \quad \textsc{PatData\_Head}$$

$$\frac{\mathbf{uncurry}\, p = F@PA}{\mathbf{uncurry}\, (p\ \ a^R) = F@\mathbf{Tm}\, R\, a, PA} \quad \textsc{PatData\_Rel}$$

$\boxed{\mathbf{pattern}\, p}$

$$\frac{}{\mathbf{pattern}\, F} \quad \textsc{Pattern\_Head}$$

$$\frac{\mathbf{pattern}\, p}{\mathbf{pattern}\, (p\ \ a^R)} \quad \textsc{Pattern\_Rel}$$

$$\frac{\mathbf{pattern}\, p}{\mathbf{pattern}\, (p\ \ a^-)} \quad \textsc{Pattern\_Irrel}$$

$$\frac{\mathbf{pattern}\, p}{\mathbf{pattern}\, (p[\gamma])} \quad \textsc{Pattern\_Co}$$

$\boxed{\mathbf{subpattern}\, p'\, p}$ Subpattern

$$\frac{\mathbf{pattern}\, p}{\mathbf{subpattern}\, p\, p} \quad \textsc{SubPat\_Refl}$$

$$\frac{\mathbf{subpattern}\, p'\, p}{\mathbf{subpattern}\, p'\, (p\ \ x^R)} \quad \textsc{SubPat\_Rel}$$

$$\frac{\mathbf{subpattern}\, p'\, p}{\mathbf{subpattern}\, p'\, (p\ \ \Box^-)} \quad \textsc{SubPat\_Irrel}$$

$$\frac{\mathbf{subpattern}\, p'\, p}{\mathbf{subpattern}\, p'\, (p[\bullet])} \quad \textsc{SubPat\_Co}$$

$\boxed{a \leftrightarrow p}$ term and pattern agree

$$\frac{}{F \leftrightarrow F} \quad \textsc{tm\_pattern\_agree\_Const}$$

$$\frac{a_1 \leftrightarrow p_1}{(a_1\ \ {a_2}^R) \leftrightarrow (p_1\ \ x^R)} \quad \textsc{tm\_pattern\_agree\_AppRelR}$$

$$\frac{a_1 \leftrightarrow p_1}{(a_1\ \ a^-) \leftrightarrow (p_1\ \Box^-)} \quad \textsc{tm\_pattern\_agree\_AppIrrel}$$

$$\frac{a_1 \leftrightarrow p_1}{(a_1[\bullet]) \leftrightarrow (p_1[\bullet])} \quad \textsc{tm\_pattern\_agree\_CApp}$$

$\boxed{a \sqsubseteq p}$ sub-pattern agrees with term

$$\frac{a \leftrightarrow p}{a \sqsubseteq p} \quad \textsc{tm\_subpattern\_agree\_Base}$$

$$\frac{a \sqsubseteq p}{a \sqsubseteq (p\ \ x^R)} \quad \textsc{tm\_subpattern\_agree\_AppRelR}$$

$$\frac{a \sqsubseteq p}{a \sqsubseteq (p\ \Box^-)} \quad \textsc{tm\_subpattern\_agree\_AppIrrel}$$

$$\frac{a \sqsubseteq p}{a \sqsubseteq (p[\bullet])} \quad \textsc{tm\_subpattern\_agree\_CAppp}$$

$\boxed{a \sqsupseteq p}$ sub-term agrees with pattern

$$\frac{a \leftrightarrow p}{a \sqsupseteq p} \quad \textsc{subtm\_pattern\_agree\_Base}$$

$$\frac{a \sqsupseteq p}{a\ \ {a_2}^\nu \sqsupseteq p} \quad \textsc{subtm\_pattern\_agree\_App}$$

$$\frac{a \sqsupseteq p}{a[\bullet] \sqsupseteq p} \quad \textsc{subtm\_pattern\_agree\_CAppp}$$

$\boxed{\mathsf{ValuePath}\ \ a \nrightarrow F}$ Type headed by constant (role-sensitive partial function used in value)

$$\frac{F :_{Rs}\ \ A\ \in\ \Sigma_0}{\mathsf{ValuePath}\ \ F \nrightarrow F} \quad \textsc{ValuePath\_AbsConst}$$

$$\frac{F :_{Rs}\ \ p \sim_{R_1}\ a : A\ \in\ \Sigma_0}{\mathsf{ValuePath}\ \ F \nrightarrow F} \quad \textsc{ValuePath\_Const}$$

$$\frac{\mathsf{ValuePath}\ \ a \nrightarrow F}{\mathsf{ValuePath}\ \ (a\ \ b'^\nu) \nrightarrow F} \quad \textsc{ValuePath\_App}$$

$$\frac{\mathsf{ValuePath}\ \ a \nrightarrow F}{\mathsf{ValuePath}\ \ (a[\bullet]) \nrightarrow F} \quad \textsc{ValuePath\_CApp}$$

$\boxed{\mathsf{CasePath}_R\ \ a \nrightarrow F}$ Type headed by constant (role-sensitive partial function used in case)

$$\frac{\begin{array}{l}\mathsf{ValuePath}\ \ a \nrightarrow F\\ F :_{Rs}\ \ A\ \in\ \Sigma_0\end{array}}{\mathsf{CasePath}_R\ \ a \nrightarrow F} \quad \textsc{CasePath\_AbsConst}$$

$$\frac{\begin{array}{l}\mathsf{ValuePath}\ \ a \nrightarrow F\\ F :_{Rs}\ \ p \sim_{R_1}\ b : A\ \in\ \Sigma_0\\ \neg(R_1 \leq R)\end{array}}{\mathsf{CasePath}_R\ \ a \nrightarrow F} \quad \textsc{CasePath\_Const}$$

$$\frac{\begin{array}{l}\mathsf{ValuePath}\ \ a \nrightarrow F\\ F :_{Rs}\ \ p \sim_{R_1}\ b : A\ \in\ \Sigma_0\\ \neg(a \sqsupseteq p)\end{array}}{\mathsf{CasePath}_R\ \ a \nrightarrow F} \quad \textsc{CasePath\_UnMatch}$$

$\boxed{\mathsf{apply\ args}\ a\ \mathsf{to}\ b \not\mapsto b'}$ apply arguments of a (headed by a constant) to b

$$\frac{}{\mathsf{apply\ args}\ F\ \mathsf{to}\ b \not\mapsto b} \quad \textsc{ApplyArgs\_Const}$$

$$\frac{\mathsf{apply\ args}\ a\ \mathsf{to}\ b \not\mapsto b'}{\mathsf{apply\ args}\ (a\ a'^{R})\ \mathsf{to}\ b \not\mapsto (b'\ a'^{+})} \quad \textsc{ApplyArgs\_AppRole}$$

$$\frac{\mathsf{apply\ args}\ a\ \mathsf{to}\ b \not\mapsto b'}{\mathsf{apply\ args}\ (a\ a'^{\rho})\ \mathsf{to}\ b \not\mapsto (b'\ a'^{\rho})} \quad \textsc{ApplyArgs\_AppRho}$$

$$\frac{\mathsf{apply\ args}\ a\ \mathsf{to}\ b \not\mapsto b'}{\mathsf{apply\ args}\ a[\bullet]\ \mathsf{to}\ b \not\mapsto b'[\bullet]} \quad \textsc{ApplyArgs\_CApp}$$

$\boxed{\mathsf{Value}_R\ A}$ values

$$\frac{}{\mathsf{Value}_R\ \star} \quad \textsc{Value\_Star}$$

$$\frac{}{\mathsf{Value}_R\ \Pi^{\rho} x\!:\!A \to B} \quad \textsc{Value\_Pi}$$

$$\frac{}{\mathsf{Value}_R\ \forall c\!:\!\phi.B} \quad \textsc{Value\_CPi}$$

$$\frac{}{\mathsf{Value}_R\ \lambda^{+} x\!:\!A.a} \quad \textsc{Value\_AbsRel}$$

$$\frac{}{\mathsf{Value}_R\ \lambda^{+} x.a} \quad \textsc{Value\_UAbsRel}$$

$$\frac{\mathsf{Value}_R\ a}{\mathsf{Value}_R\ \lambda^{-} x.a} \quad \textsc{Value\_UAbsIrrel}$$

$$\frac{}{\mathsf{Value}_R\ \Lambda c\!:\!\phi.a} \quad \textsc{Value\_CAbs}$$

$$\frac{}{\mathsf{Value}_R\ \Lambda c.a} \quad \textsc{Value\_UCAbs}$$

$$\frac{\mathsf{CasePath}_R\ a \not\mapsto F}{\mathsf{Value}_R\ a} \quad \textsc{Value\_Path}$$

$\boxed{\mathsf{ValueType}_R\ A}$ Types with head forms (erased language)

$$\frac{}{\mathsf{ValueType}_R\ \star} \quad \textsc{value\_type\_Star}$$

$$\frac{}{\mathsf{ValueType}_R\ \Pi^{\rho} x\!:\!A \to B} \quad \textsc{value\_type\_Pi}$$

$$\frac{}{\mathsf{ValueType}_R\ \forall c\!:\!\phi.B} \quad \textsc{value\_type\_CPi}$$

$$\frac{\mathsf{CasePath}_R\ a \not\mapsto F}{\mathsf{ValueType}_R\ a} \quad \textsc{value\_type\_ValuePath}$$

$\boxed{\mathsf{consistent}_R\ a\ b}$ (erased) types do not differ in their heads

$$\frac{}{\mathsf{consistent}_R\ \star\ \star} \quad \textsc{consistent\_a\_Star}$$

$$\frac{}{\mathsf{consistent}_{R'}\ (\Pi^{\rho} x_1\!:\!A_1 \to B_1)\ (\Pi^{\rho} x_2\!:\!A_2 \to B_2)} \quad \textsc{consistent\_a\_Pi}$$

$$\frac{}{\mathsf{consistent}_R\ (\forall c_1\!:\!\phi_1.A_1)\ (\forall c_2\!:\!\phi_2.A_2)} \quad \textsc{consistent\_a\_CPi}$$

$$\frac{\begin{array}{c}\mathsf{CasePath}_R\ a_1 \not\mapsto F\\ \mathsf{CasePath}_R\ a_2 \not\mapsto F\end{array}}{\mathsf{consistent}_R\ a_1\ a_2}\quad \textsc{consistent\_a\_CasePath}$$

$$\frac{\neg\mathsf{ValueType}_R\ b}{\mathsf{consistent}_R\ a\ b}\quad \textsc{consistent\_a\_Step\_R}$$

$$\frac{\neg\mathsf{ValueType}_R\ a}{\mathsf{consistent}_R\ a\ b}\quad \textsc{consistent\_a\_Step\_L}$$

$\boxed{\Omega \vDash a : R}$ Roleing judgment

$$\frac{\mathit{uniq}(\Omega)}{\Omega \vDash \square : R}\quad \textsc{role\_a\_Bullet}$$

$$\frac{\mathit{uniq}(\Omega)}{\Omega \vDash \star : R}\quad \textsc{role\_a\_Star}$$

$$\frac{\begin{array}{l}\mathit{uniq}(\Omega)\\ x : R \in \Omega\\ R \leq R_1\end{array}}{\Omega \vDash x : R_1}\quad \textsc{role\_a\_Var}$$

$$\frac{\Omega, x : \mathbf{Nom} \vDash a : R}{\Omega \vDash (\lambda^\rho x.a) : R}\quad \textsc{role\_a\_Abs}$$

$$\frac{\begin{array}{l}\Omega \vDash a : R\\ \Omega \vDash b : \mathbf{Nom}\end{array}}{\Omega \vDash (a\ b^\rho) : R}\quad \textsc{role\_a\_App}$$

$$\frac{\begin{array}{l}\Omega \vDash a : R\\ \Omega \vDash b : \mathbf{param}\, R_1\, R\end{array}}{\Omega \vDash a\ b^{R_1} : R}\quad \textsc{role\_a\_TApp}$$

$$\frac{\begin{array}{l}\Omega \vDash A : R\\ \Omega, x : \mathbf{Nom} \vDash B : R\end{array}}{\Omega \vDash (\Pi^\rho x\!:\!A \to B) : R}\quad \textsc{role\_a\_Pi}$$

$$\frac{\begin{array}{l}\Omega \vDash a : R_1\\ \Omega \vDash b : R_1\\ \Omega \vDash A : \mathbf{Rep}\\ \Omega \vDash B : R\end{array}}{\Omega \vDash (\forall c\!:\!a \sim_{R_1} b : A.B) : R}\quad \textsc{role\_a\_CPi}$$

$$\frac{\Omega \vDash b : R}{\Omega \vDash (\Lambda c.b) : R}\quad \textsc{role\_a\_CAbs}$$

$$\frac{\Omega \vDash a : R}{\Omega \vDash (a[\bullet]) : R}\quad \textsc{role\_a\_CApp}$$

$$\frac{\begin{array}{l}\mathit{uniq}(\Omega)\\ F :_{Rs}\ A \in \Sigma_0\end{array}}{\Omega \vDash F : R}\quad \textsc{role\_a\_Const}$$

$$\frac{\begin{array}{l}\mathit{uniq}(\Omega)\\ F :_{Rs}\ p \sim_R a : A \in \Sigma_0\end{array}}{\Omega \vDash F : R_1}\quad \textsc{role\_a\_Fam}$$

$$\frac{\begin{array}{l}\Omega \vDash a : \mathbf{Nom} \\ \Omega \vDash b_1 : R_1 \\ \Omega \vDash b_2 : R_1\end{array}}{\Omega \vDash \mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \|_{\_} \to b_2 : R_1} \quad \textsc{role\_a\_Pattern}$$

$\boxed{(\rho = +) \vee (x \notin \mathsf{fv}\ A)}$ irrelevant argument check

$$\frac{}{(+ = +) \vee (x \notin \mathsf{fv}\ A)} \quad \textsc{Rho\_Rel}$$

$$\frac{x \notin \mathsf{fv} A}{(- = +) \vee (x \notin \mathsf{fv}\ A)} \quad \textsc{Rho\_IrrRel}$$

$\boxed{\Omega \vDash a \Rightarrow_R b}$ parallel reduction (implicit language)

$$\frac{\Omega \vDash a : R}{\Omega \vDash a \Rightarrow_R a} \quad \textsc{Par\_Refl}$$

$$\frac{\begin{array}{l}\Omega \vDash a \Rightarrow_R (\lambda^\rho x.a') \\ \Omega \vDash b \Rightarrow_{\mathbf{Nom}} b'\end{array}}{\Omega \vDash a\ b^\rho \Rightarrow_R a'\{b'/x\}} \quad \textsc{Par\_Beta}$$

$$\frac{\begin{array}{l}\Omega \vDash a \Rightarrow_R a' \\ \Omega \vDash b \Rightarrow_{(\mathbf{app\_role}\,\nu\,R)} b'\end{array}}{\Omega \vDash a\ b^\nu \Rightarrow_R a'\ b'^\nu} \quad \textsc{Par\_App}$$

$$\frac{\Omega \vDash a \Rightarrow_R (\Lambda c.a')}{\Omega \vDash a[\bullet] \Rightarrow_R a'\{\bullet/c\}} \quad \textsc{Par\_CBeta}$$

$$\frac{\Omega \vDash a \Rightarrow_R a'}{\Omega \vDash a[\bullet] \Rightarrow_R a'[\bullet]} \quad \textsc{Par\_CApp}$$

$$\frac{\Omega, x : \mathbf{Nom} \vDash a \Rightarrow_R a'}{\Omega \vDash \lambda^\rho x.a \Rightarrow_R \lambda^\rho x.a'} \quad \textsc{Par\_Abs}$$

$$\frac{\begin{array}{l}\Omega \vDash A \Rightarrow_R A' \\ \Omega, x : \mathbf{Nom} \vDash B \Rightarrow_R B'\end{array}}{\Omega \vDash \Pi^\rho x\!:\!A \to B \Rightarrow_R \Pi^\rho x\!:\!A' \to B'} \quad \textsc{Par\_Pi}$$

$$\frac{\Omega \vDash a \Rightarrow_R a'}{\Omega \vDash \Lambda c.a \Rightarrow_R \Lambda c.a'} \quad \textsc{Par\_CAbs}$$

$$\frac{\begin{array}{l}\Omega \vDash A \Rightarrow_{\mathbf{Rep}} A' \\ \Omega \vDash a \Rightarrow_{R_1} a' \\ \Omega \vDash b \Rightarrow_{R_1} b' \\ \Omega \vDash B \Rightarrow_R B'\end{array}}{\Omega \vDash \forall c\!:\!a \sim_{R_1} b : A.B \Rightarrow_R \forall c\!:\!a' \sim_{R_1} b' : A'.B'} \quad \textsc{Par\_CPi}$$

$$\frac{\begin{array}{l}F :_{Rs}\ F \sim_{R_1} b : A \in \Sigma_0 \\ R_1 \leq R \\ \mathit{uniq}(\Omega)\end{array}}{\Omega \vDash F \Rightarrow_R b} \quad \textsc{Par\_AxiomBase}$$

$$\frac{\begin{array}{l} F :_{Rs}\ p \sim_{R_1} b : A \in \Sigma_0 \\ a \sqsubseteq p \wedge \neg(a \leftrightarrow p) \\ \Omega \vDash a \Rightarrow_R a' \\ \Omega \vDash a_1 \Rightarrow_{(\mathbf{app\_role}\,\nu\,R)} a'_1 \\ \textsf{rename } p \to b \textsf{ to } p' \to b' \textsf{ excluding } (\widetilde{\Omega}, \textsf{fv}p) \textsf{ and } \Delta' \\ \textsf{match } (a'\ {a'_1}^{\nu}) \textsf{ with } p' \to b' \nrightarrow a_2 \\ R_1 \leq R \end{array}}{\Omega \vDash a\ {a_1}^{\nu} \Rightarrow_R a_2} \quad \textsc{Par\_AxiomApp}$$

$$\frac{\begin{array}{l} F :_{Rs}\ p \sim_{R_1} b : A \in \Sigma_0 \\ a \sqsubseteq p \wedge \neg(a \leftrightarrow p) \\ \Omega \vDash a \Rightarrow_R a' \\ \textsf{rename } p \to b \textsf{ to } p' \to b' \textsf{ excluding } (\widetilde{\Omega}, \textsf{fv}p) \textsf{ and } \Delta' \\ \textsf{match } (a'[\bullet]) \textsf{ with } p' \to b' \nrightarrow a_2 \\ R_1 \leq R \end{array}}{\Omega \vDash a[\bullet] \Rightarrow_R a_2} \quad \textsc{Par\_AxiomCApp}$$

$$\frac{\begin{array}{c} \Omega \vDash a \Rightarrow_{\mathbf{Nom}} a' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b'_1 \\ \Omega \vDash b_2 \Rightarrow_{R_0} b'_2 \end{array}}{\Omega \vDash (\textsf{case } a \textsf{ of } F\overline{\upsilon} \to b_1 \|_{\_} \to b_2) \Rightarrow_{R_0} (\textsf{case } a' \textsf{ of } F\overline{\upsilon} \to b'_1 \|_{\_} \to b'_2)} \quad \textsc{Par\_Pattern}$$

$$\frac{\begin{array}{l} \Omega \vDash a \Rightarrow_{\mathbf{Nom}} a' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b'_1 \\ \Omega \vDash b_2 \Rightarrow_{R_0} b'_2 \\ \textsf{AppsPath}\,a' \nrightarrow F\overline{\upsilon} \\ \textsf{apply args } a' \textsf{ to } b'_1 \nrightarrow b \\ \mathbf{Sat}\ F\,\overline{\upsilon}' \end{array}}{\Omega \vDash (\textsf{case } a \textsf{ of } F\overline{\upsilon} \to b_1 \|_{\_} \to b_2) \Rightarrow_{R_0} b[\bullet]} \quad \textsc{Par\_PatternTrue}$$

$$\frac{\begin{array}{l} \Omega \vDash a \Rightarrow_{\mathbf{Nom}} a' \\ \Omega \vDash b_1 \Rightarrow_{R_0} b'_1 \\ \Omega \vDash b_2 \Rightarrow_{R_0} b'_2 \\ \textsf{Value}_{\mathbf{Nom}}\ a' \\ \neg(\textsf{AppsPath}\,a' \nrightarrow F\overline{\upsilon}) \end{array}}{\Omega \vDash (\textsf{case } a \textsf{ of } F\overline{\upsilon} \to b_1 \|_{\_} \to b_2) \Rightarrow_{R_0} b'_2} \quad \textsc{Par\_PatternFalse}$$

$\boxed{\Omega \vDash a \Rightarrow^*_R b}$ multistep parallel reduction

$$\frac{}{\Omega \vDash a \Rightarrow^*_R a} \quad \textsc{MP\_Refl}$$

$$\frac{\begin{array}{c} \Omega \vDash a \Rightarrow_R b \\ \Omega \vDash b \Rightarrow^*_R a' \end{array}}{\Omega \vDash a \Rightarrow^*_R a'} \quad \textsc{MP\_Step}$$

$\boxed{\Omega \vDash a \Leftrightarrow_R b}$ parallel reduction to a common term

$$\frac{\begin{array}{c} \Omega \vDash a_1 \Rightarrow^*_R b \\ \Omega \vDash a_2 \Rightarrow^*_R b \end{array}}{\Omega \vDash a_1 \Leftrightarrow_R a_2} \quad \textsc{join}$$

$\boxed{\vDash a >_R b}$ primitive reductions on erased terms

$$\frac{\textsf{Value}_{R_1}\ (\lambda^{\rho} x.v)}{\vDash (\lambda^{\rho} x.v)\ b^{\rho} >_{R_1} v\{b/x\}} \quad \textsc{Beta\_AppAbs}$$

$$\frac{}{\vDash (\Lambda c.a')[\bullet] >_R a'\{\bullet/c\}} \quad \textsc{Beta\_CAppCAbs}$$

$$\frac{\begin{array}{l} F :_{Rs}\ p \sim_{R_1} b : A \in \Sigma_0 \\ \mathsf{rename}\ p \to b\ \mathsf{to}\ p_1 \to b_1\ \mathsf{excluding}\ (\mathsf{fv}a, \mathsf{fv}p)\ \mathsf{and}\ \Delta' \\ \mathsf{match}\ a\ \mathsf{with}\ p_1 \to b_1 \not\rightarrow b' \\ R_1 \leq R \end{array}}{\vDash a >_R b'} \quad \textsc{Beta\_Axiom}$$

$$\frac{\begin{array}{l} \mathsf{AppsPath}\, a \not\rightarrow F\overline{v} \\ \mathsf{apply\ args}\ a\ \mathsf{to}\ b_1 \not\rightarrow b_1' \\ \mathbf{Sat}\, F\,\overline{v}' \end{array}}{\vDash \mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \| \_ \to b_2 >_{R_0} b_1'[\bullet]} \quad \textsc{Beta\_PatternTrue}$$

$$\frac{\begin{array}{l} \mathsf{Value}_{\mathbf{Nom}}\ a \\ \neg(\mathsf{AppsPath}\, a \not\rightarrow F\overline{v}) \end{array}}{\vDash \mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \| \_ \to b_2 >_{R_0} b_2} \quad \textsc{Beta\_PatternFalse}$$

$\boxed{\vDash a \rightsquigarrow_R b}$ single-step head reduction for implicit language

$$\frac{\vDash a \rightsquigarrow_{R_1} a'}{\vDash \lambda^- x.a \rightsquigarrow_{R_1} \lambda^- x.a'} \quad \textsc{E\_AbsTerm}$$

$$\frac{\vDash a \rightsquigarrow_{R_1} a'}{\vDash a\ b^\nu \rightsquigarrow_{R_1} a'\ b^\nu} \quad \textsc{E\_AppLeft}$$

$$\frac{\vDash a \rightsquigarrow_R a'}{\vDash a[\bullet] \rightsquigarrow_R a'[\bullet]} \quad \textsc{E\_CAppLeft}$$

$$\frac{\vDash a \rightsquigarrow_{\mathbf{Nom}} a'}{\vDash \mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \| \_ \to b_2 \rightsquigarrow_{R_0} \mathsf{case}\ a'\ \mathsf{of}\ F\overline{v} \to b_1 \| \_ \to b_2} \quad \textsc{E\_Pattern}$$

$$\frac{\vDash a >_R b}{\vDash a \rightsquigarrow_R b} \quad \textsc{E\_Prim}$$

$\boxed{\vDash a \rightsquigarrow^* b/R}$ multistep reduction

$$\frac{}{\vDash a \rightsquigarrow^* a/R} \quad \textsc{Equal}$$

$$\frac{\begin{array}{l} \vDash a \rightsquigarrow_R b \\ \vDash b \rightsquigarrow^* a'/R \end{array}}{\vDash a \rightsquigarrow^* a'/R} \quad \textsc{Step}$$

$\boxed{\Gamma \vDash \mathsf{case}\, a : A\, \mathsf{of}\, b : B \Rightarrow C \mid C'}$ Branch Typing (aligning the types of case)

$$\frac{\begin{array}{l} \mathit{uniq}\,\Gamma \\ C_1\{\bullet/c\} = C_2 \end{array}}{\Gamma \vDash \mathsf{case}\, a : A\, \mathsf{of}\, b : A \Rightarrow \forall c\!:\!(a \sim_{\mathbf{Nom}} \mathbf{apply}\, b\, \overline{\mu} : A).C_1 \mid C_2} \quad \textsc{BranchTyping\_Base}$$

$$\frac{\Gamma, x : A \vDash \mathsf{case}\, a : A_1\, \mathsf{of}\, b : B \Rightarrow C \mid C'}{\Gamma \vDash \mathsf{case}\, a : A_1\, \mathsf{of}\, b : \Pi^+ x\!:\!A \to B \Rightarrow \Pi^+ x\!:\!A \to C \mid C'} \quad \textsc{BranchTyping\_PiRole}$$

$$\frac{\Gamma, x : A \vDash \mathsf{case}\, a : A_1\, \mathsf{of}\, b : B \Rightarrow C \mid C'}{\Gamma \vDash \mathsf{case}\, a : A_1\, \mathsf{of}\, b : \Pi^+ x\!:\!A \to B \Rightarrow \Pi^+ x\!:\!A \to C \mid C'} \quad \textsc{BranchTyping\_PiRel}$$

$$\frac{\Gamma, x : A \vDash \mathsf{case}\, a : A_1\, \mathsf{of}\, b : B \Rightarrow C \mid C'}{\Gamma \vDash \mathsf{case}\, a : A_1\, \mathsf{of}\, b : \Pi^- x\!:\!A \to B \Rightarrow \Pi^- x\!:\!A \to C \mid C'} \quad \textsc{BranchTyping\_PiIrrel}$$

$$\frac{\Gamma, c : \phi \vDash \mathsf{case}\ a : A\ \mathsf{of}\ b : B \Rightarrow C \mid C'}{\Gamma \vDash \mathsf{case}\ a : A\ \mathsf{of}\ b : \forall c{:}\phi.B \Rightarrow \forall c{:}\phi.C \mid C'} \quad \text{BranchTyping\_CPi}$$

$\boxed{\Gamma \vDash \phi\ \ \mathsf{ok}}$ Prop wellformedness

$$\frac{\begin{array}{l}\Gamma \vDash a : A \\ \Gamma \vDash b : A \\ \Gamma \vDash A : \star\end{array}}{\Gamma \vDash a \sim_R b : A\ \ \mathsf{ok}} \quad \text{E\_Wff}$$

$\boxed{\Gamma \vDash a : A}$ typing

$$\frac{\vDash \Gamma}{\Gamma \vDash \star : \star} \quad \text{E\_Star}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ x : A \in \Gamma\end{array}}{\Gamma \vDash x : A} \quad \text{E\_Var}$$

$$\frac{\begin{array}{l}\Gamma, x : A \vDash B : \star \\ \Gamma \vDash A : \star\end{array}}{\Gamma \vDash \Pi^\rho x{:}A \to B : \star} \quad \text{E\_Pi}$$

$$\frac{\begin{array}{l}\Gamma, x : A \vDash a : B \\ \Gamma \vDash A : \star \\ (\rho = +) \vee (x \notin \mathsf{fv}\ a)\end{array}}{\Gamma \vDash \lambda^\rho x.a : (\Pi^\rho x{:}A \to B)} \quad \text{E\_Abs}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^+ x{:}A \to B \\ \Gamma \vDash a : A\end{array}}{\Gamma \vDash b\ \ a^+ : B\{a/x\}} \quad \text{E\_App}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^+ x{:}A \to B \\ \Gamma \vDash a : A \\ \mathsf{RolePath}\ \ b = F@R, Rs\end{array}}{\Gamma \vDash b\ \ a^R : B\{a/x\}} \quad \text{E\_TApp}$$

$$\frac{\begin{array}{l}\Gamma \vDash b : \Pi^- x{:}A \to B \\ \Gamma \vDash a : A\end{array}}{\Gamma \vDash b\ \square^- : B\{a/x\}} \quad \text{E\_IApp}$$

$$\frac{\begin{array}{l}\Gamma \vDash a : A \\ \Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Rep} \\ \Gamma \vDash B : \star\end{array}}{\Gamma \vDash a : B} \quad \text{E\_Conv}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi \vDash B : \star \\ \Gamma \vDash \phi\ \ \mathsf{ok}\end{array}}{\Gamma \vDash \forall c{:}\phi.B : \star} \quad \text{E\_CPi}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi \vDash a : B \\ \Gamma \vDash \phi\ \ \mathsf{ok}\end{array}}{\Gamma \vDash \Lambda c.a : \forall c{:}\phi.B} \quad \text{E\_CAbs}$$

$$\frac{\begin{array}{l}\Gamma \vDash a_1 : \forall c{:}(a \sim_R b : A).B_1 \\ \Gamma; \widetilde{\Gamma} \vDash a \equiv b : A/R\end{array}}{\Gamma \vDash a_1[\bullet] : B_1\{\bullet/c\}} \quad \text{E\_CApp}$$

$$\frac{\vDash \Gamma \quad F :_{Rs} \ A \in \Sigma_0 \quad \varnothing \vDash A : \star}{\Gamma \vDash F : A} \ \textsc{E\_Const}$$

$$\frac{\vDash \Gamma \quad F :_{Rs} \ p \sim_{R_1} a : A \in \Sigma_0 \quad \varnothing \vDash A : \star}{\Gamma \vDash F : A} \ \textsc{E\_Fam}$$

$$\frac{\begin{array}{l}\Gamma \vDash a : A \\ \Gamma \vDash b_1 : B \\ \Gamma \vDash b_2 : C \\ \Gamma \vDash \mathsf{case}\ a : A\ \mathsf{of}\ F : A_1 \Rightarrow B \mid C \\ \Gamma \vDash F : A_1 \\ \mathbf{Sat}\ F\ \overline{v}\end{array}}{\Gamma \vDash \mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \| \_ \to b_2 : C} \ \textsc{E\_Case}$$

$\boxed{\Gamma; \Delta \vDash \phi_1 \equiv \phi_2}$ prop equality

$$\frac{\Gamma; \Delta \vDash A_1 \equiv A_2 : A/R \quad \Gamma; \Delta \vDash B_1 \equiv B_2 : A/R}{\Gamma; \Delta \vDash A_1 \sim_R B_1 : A \equiv A_2 \sim_R B_2 : A} \ \textsc{E\_PropCong}$$

$$\frac{\Gamma; \Delta \vDash A \equiv B : \star/R_0 \quad \Gamma \vDash A_1 \sim_R A_2 : A \ \ \mathsf{ok} \quad \Gamma \vDash A_1 \sim_R A_2 : B \ \ \mathsf{ok}}{\Gamma; \Delta \vDash A_1 \sim_R A_2 : A \equiv A_1 \sim_R A_2 : B} \ \textsc{E\_IsoConv}$$

$$\frac{\Gamma; \Delta \vDash \forall c\!:\!(a_1 \sim_{R_1} a_2 : A).B_1 \equiv \forall c\!:\!(b_1 \sim_{R_2} b_2 : B).B_2 : \star/R'}{\Gamma; \Delta \vDash a_1 \sim_{R_1} a_2 : A \equiv b_1 \sim_{R_2} b_2 : B} \ \textsc{E\_CPiFst}$$

$\boxed{\Gamma; \Delta \vDash a \equiv b : A/R}$ definitional equality

$$\frac{\vDash \Gamma \quad c : (a \sim_R b : A) \in \Gamma \quad c \in \Delta}{\Gamma; \Delta \vDash a \equiv b : A/R} \ \textsc{E\_Assn}$$

$$\frac{\Gamma \vDash a : A}{\Gamma; \Delta \vDash a \equiv a : A/R} \ \textsc{E\_Refl}$$

$$\frac{\Gamma; \Delta \vDash b \equiv a : A/R}{\Gamma; \Delta \vDash a \equiv b : A/R} \ \textsc{E\_Sym}$$

$$\frac{\Gamma; \Delta \vDash a \equiv a_1 : A/R \quad \Gamma; \Delta \vDash a_1 \equiv b : A/R}{\Gamma; \Delta \vDash a \equiv b : A/R} \ \textsc{E\_Trans}$$

$$\frac{\Gamma; \Delta \vDash a \equiv b : A/R_1 \quad R_1 \leq R_2}{\Gamma; \Delta \vDash a \equiv b : A/R_2} \ \textsc{E\_Sub}$$

$$\frac{\Gamma \vDash a_1 : B \quad \Gamma \vDash a_2 : B \quad \vDash a_1 >_R a_2}{\Gamma; \Delta \vDash a_1 \equiv a_2 : B/R} \ \textsc{E\_Beta}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash A_1 \equiv A_2 : \star/R' \\ \Gamma, x : A_1;\Delta \vDash B_1 \equiv B_2 : \star/R' \\ \Gamma \vDash A_1 : \star \\ \Gamma \vDash \Pi^\rho x\!:\!A_1 \to B_1 : \star \\ \Gamma \vDash \Pi^\rho x\!:\!A_2 \to B_2 : \star\end{array}}{\Gamma;\Delta \vDash (\Pi^\rho x\!:\!A_1 \to B_1) \equiv (\Pi^\rho x\!:\!A_2 \to B_2) : \star/R'} \quad \textsc{E\_PiCong}$$

$$\frac{\begin{array}{l}\Gamma, x : A_1;\Delta \vDash b_1 \equiv b_2 : B/R' \\ \Gamma \vDash A_1 : \star \\ (\rho = +) \vee (x \notin \mathsf{fv}\ b_1) \\ (\rho = +) \vee (x \notin \mathsf{fv}\ b_2)\end{array}}{\Gamma;\Delta \vDash (\lambda^\rho x.b_1) \equiv (\lambda^\rho x.b_2) : (\Pi^\rho x\!:\!A_1 \to B)/R'} \quad \textsc{E\_AbsCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \equiv b_1 : (\Pi^+ x\!:\!A \to B)/R' \\ \Gamma;\Delta \vDash a_2 \equiv b_2 : A/\mathbf{Nom}\end{array}}{\Gamma;\Delta \vDash a_1\ \ a_2{}^+ \equiv b_1\ \ b_2{}^+ : (B\{a_2/x\})/R'} \quad \textsc{E\_AppCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \equiv b_1 : (\Pi^+ x\!:\!A \to B)/R' \\ \Gamma;\Delta \vDash a_2 \equiv b_2 : A/\mathbf{param}\, R\, R' \\ \mathsf{RolePath}\ \ a_1 = F@R, Rs \\ \mathsf{RolePath}\ \ b_1 = F'@R, Rs' \\ \Gamma \vDash b_1\ \ b_2{}^R : B\{a_2/x\}\end{array}}{\Gamma;\Delta \vDash a_1\ \ a_2{}^R \equiv b_1\ \ b_2{}^R : (B\{a_2/x\})/R'} \quad \textsc{E\_TAppCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \equiv b_1 : (\Pi^- x\!:\!A \to B)/R' \\ \Gamma \vDash a : A\end{array}}{\Gamma;\Delta \vDash a_1\ \Box^- \equiv b_1\ \Box^- : (B\{a/x\})/R'} \quad \textsc{E\_IAppCong}$$

$$\frac{\Gamma;\Delta \vDash \Pi^\rho x\!:\!A_1 \to B_1 \equiv \Pi^\rho x\!:\!A_2 \to B_2 : \star/R'}{\Gamma;\Delta \vDash A_1 \equiv A_2 : \star/R'} \quad \textsc{E\_PiFst}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash \Pi^\rho x\!:\!A_1 \to B_1 \equiv \Pi^\rho x\!:\!A_2 \to B_2 : \star/R' \\ \Gamma;\Delta \vDash a_1 \equiv a_2 : A_1/\mathbf{Nom}\end{array}}{\Gamma;\Delta \vDash B_1\{a_1/x\} \equiv B_2\{a_2/x\} : \star/R'} \quad \textsc{E\_PiSnd}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \sim_R b_1 : A_1 \equiv a_2 \sim_R b_2 : A_2 \\ \Gamma, c : a_1 \sim_R b_1 : A_1;\Delta \vDash A \equiv B : \star/R' \\ \Gamma \vDash a_1 \sim_R b_1 : A_1\ \ \mathsf{ok} \\ \Gamma \vDash \forall c\!:\!a_1 \sim_R b_1 : A_1.A : \star \\ \Gamma \vDash \forall c\!:\!a_2 \sim_R b_2 : A_2.B : \star\end{array}}{\Gamma;\Delta \vDash \forall c\!:\!a_1 \sim_R b_1 : A_1.A \equiv \forall c\!:\!a_2 \sim_R b_2 : A_2.B : \star/R'} \quad \textsc{E\_CPiCong}$$

$$\frac{\begin{array}{l}\Gamma, c : \phi_1;\Delta \vDash a \equiv b : B/R \\ \Gamma \vDash \phi_1\ \ \mathsf{ok}\end{array}}{\Gamma;\Delta \vDash (\Lambda c.a) \equiv (\Lambda c.b) : \forall c\!:\!\phi_1.B/R} \quad \textsc{E\_CAbsCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash a_1 \equiv b_1 : (\forall c\!:\!(a \sim_R b : A).B)/R' \\ \Gamma;\widetilde{\Gamma} \vDash a \equiv b : A/\mathbf{param}\, R\, R'\end{array}}{\Gamma;\Delta \vDash a_1[\bullet] \equiv b_1[\bullet] : (B\{\bullet/c\})/R'} \quad \textsc{E\_CAppCong}$$

$$\frac{\begin{array}{l}\Gamma;\Delta \vDash \forall c\!:\!(a_1 \sim_R a_2 : A).B_1 \equiv \forall c\!:\!(a_1' \sim_{R'} a_2' : A').B_2 : \star/R_0 \\ \Gamma;\widetilde{\Gamma} \vDash a_1 \equiv a_2 : A/R \\ \Gamma;\widetilde{\Gamma} \vDash a_1' \equiv a_2' : A'/R'\end{array}}{\Gamma;\Delta \vDash B_1\{\bullet/c\} \equiv B_2\{\bullet/c\} : \star/R_0} \quad \textsc{E\_CPiSnd}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a \equiv b : A/R \\ \Gamma; \Delta \vDash a \sim_R b : A \equiv a' \sim_{R'} b' : A'\end{array}}{\Gamma; \Delta \vDash a' \equiv b' : A'/R'} \quad \text{E\_CAST}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a \equiv b : A/R \\ \Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Rep} \\ \Gamma \vDash B : \star\end{array}}{\Gamma; \Delta \vDash a \equiv b : B/R} \quad \text{E\_EQCONV}$$

$$\frac{\Gamma; \Delta \vDash a \sim_{R_1} b : A \equiv a' \sim_{R_1} b' : A'}{\Gamma; \Delta \vDash A \equiv A' : \star/\mathbf{Rep}} \quad \text{E\_ISOSND}$$

$$\frac{\begin{array}{l}\Gamma; \Delta \vDash a \equiv a' : A/\mathbf{Nom} \\ \Gamma; \Delta \vDash b_1 \equiv b_1' : B/R_0 \\ \Gamma; \Delta \vDash b_2 \equiv b_2' : C/R_0 \\ \Gamma \vDash \mathsf{case}\ a : A\ \mathsf{of}\ F : A_1 \Rightarrow B \mid C \\ \Gamma \vDash \mathsf{case}\ a' : A\ \mathsf{of}\ F : A_1 \Rightarrow B \mid C \\ \mathbf{Sat}\ F\ \overline{v} \\ \Gamma \vDash F : A_1\end{array}}{\Gamma; \Delta \vDash \mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \| \_ \to b_2 \equiv \mathsf{case}\ a'\ \mathsf{of}\ F\overline{v} \to b_1' \| \_ \to b_2' : C/R_0} \quad \text{E\_PATCONG}$$

$$\frac{\begin{array}{l}\mathsf{ValuePath}\ a \nrightarrow F \\ \mathsf{ValuePath}\ a' \nrightarrow F \\ \Gamma \vDash a : \Pi^+ x\!:\!A \to B \\ \Gamma \vDash b : A \\ \Gamma \vDash a' : \Pi^+ x\!:\!A \to B \\ \Gamma \vDash b' : A \\ \Gamma; \Delta \vDash a\ b^{R_1} \equiv a'\ b'^{R_1} : B\{b/x\}/R' \\ \Gamma; \widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/\mathbf{Rep}\end{array}}{\Gamma; \Delta \vDash a \equiv a' : \Pi^+ x\!:\!A \to B/R'} \quad \text{E\_LEFTREL}$$

$$\frac{\begin{array}{l}\mathsf{ValuePath}\ a \nrightarrow F \\ \mathsf{ValuePath}\ a' \nrightarrow F \\ \Gamma \vDash a : \Pi^- x\!:\!A \to B \\ \Gamma \vDash b : A \\ \Gamma \vDash a' : \Pi^- x\!:\!A \to B \\ \Gamma \vDash b' : A \\ \Gamma; \Delta \vDash a\ \Box^- \equiv a'\ \Box^- : B\{b/x\}/R' \\ \Gamma; \widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/\mathbf{Rep}\end{array}}{\Gamma; \Delta \vDash a \equiv a' : \Pi^- x\!:\!A \to B/R'} \quad \text{E\_LEFTIRREL}$$

$$\frac{\begin{array}{l}\mathsf{ValuePath}\ a \nrightarrow F \\ \mathsf{ValuePath}\ a' \nrightarrow F \\ \Gamma \vDash a : \Pi^+ x\!:\!A \to B \\ \Gamma \vDash b : A \\ \Gamma \vDash a' : \Pi^+ x\!:\!A \to B \\ \Gamma \vDash b' : A \\ \Gamma; \Delta \vDash a\ b^{R_1} \equiv a'\ b'^{R_1} : B\{b/x\}/R' \\ \Gamma; \widetilde{\Gamma} \vDash B\{b/x\} \equiv B\{b'/x\} : \star/\mathbf{Rep}\end{array}}{\Gamma; \Delta \vDash b \equiv b' : A/\mathbf{param}\ R_1\ R'} \quad \text{E\_RIGHT}$$

$$\frac{\begin{array}{l}\mathsf{ValuePath}\ a \nrightarrow F \\ \mathsf{ValuePath}\ a' \nrightarrow F \\ \Gamma \vDash a : \forall c{:}(a_1 \sim_{R_1} a_2 : A).B \\ \Gamma \vDash a' : \forall c{:}(a_1 \sim_{R_1} a_2 : A).B \\ \Gamma; \widetilde{\Gamma} \vDash a_1 \equiv a_2 : A/\mathbf{param}\, R_1\, R' \\ \Gamma; \Delta \vDash a[\bullet] \equiv a'[\bullet] : B\{\bullet/c\}/R'\end{array}}{\Gamma; \Delta \vDash a \equiv a' : \forall c{:}(a_1 \sim_{R_1} a_2 : A).B/R'} \quad \text{E\_CLeft}$$

$\boxed{\vDash \Gamma}$ context wellformedness

$$\frac{}{\vDash \varnothing} \quad \text{E\_Empty}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ \Gamma \vDash A : \star \\ x \notin \widetilde{\Gamma}\end{array}}{\vDash \Gamma, x : A} \quad \text{E\_ConsTm}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ \Gamma \vDash \phi\ \mathsf{ok} \\ c \notin \widetilde{\Gamma}\end{array}}{\vDash \Gamma, c : \phi} \quad \text{E\_ConsCo}$$

$\boxed{\vDash \Sigma}$ signature wellformedness

$$\frac{}{\vDash \varnothing} \quad \text{Sig\_Empty}$$

$$\frac{\begin{array}{l}\vDash \Sigma \\ \varnothing \vDash A : \star \\ F \notin \mathsf{dom}\,\Sigma\end{array}}{\vDash \Sigma \cup \{F :_{Rs}\ A\}} \quad \text{Sig\_ConsConst}$$

$$\frac{\begin{array}{l}\vDash \Sigma \\ F \notin \mathsf{dom}\,\Sigma \\ \varnothing \vDash A : \star \\ \Omega; \Gamma \vDash p :_F B \Rightarrow A \\ \Gamma \vDash a : B \\ \Delta \# \mathsf{fv}\, a \\ \Omega \vDash a : R\end{array}}{\vDash \Sigma \cup \{F :_{\mathbf{range}}\ \Omega\ p \sim_R a : A\}} \quad \text{Sig\_ConsAx}$$

$\boxed{Rs_1 \Leftarrow Rs_2}$

$$\frac{}{\cdot \Leftarrow \cdot} \quad \text{R\_Nil}$$

$$\frac{\begin{array}{c}R_2 \leq R_1 \\ Rs_1 \Leftarrow Rs_2\end{array}}{R_1, Rs_1 \Leftarrow R_2, Rs_2} \quad \text{R\_Cons}$$

$\boxed{\Sigma_1 \Leftarrow \Sigma_2}$

$$\frac{\Sigma_1 \Leftarrow \Sigma_2}{\Sigma_1 \Leftarrow \Sigma_2 \cup \{F : \mathit{sig\_sort}\}} \quad \text{S\_Forget}$$

$$\frac{\begin{array}{c}\Sigma_1 \Leftarrow \Sigma_2 \\ Rs_1 \Leftarrow Rs_2\end{array}}{\Sigma_1 \cup \{F :_{Rs_1}\ A\} \Leftarrow \Sigma_2 \cup \{F :_{Rs_2}\ p \sim_R a : A\}} \quad \text{S\_Hide}$$

$$\frac{\begin{array}{c}\Sigma_1 \Leftarrow \Sigma_2 \\ Rs_1 \Leftarrow Rs_2\end{array}}{\Sigma_1 \cup \{F :_{Rs_1} \ A\} \Leftarrow \Sigma_2 \cup \{F :_{Rs_2} \ A\}} \quad \textsc{S\_WeakenConst}$$

$$\frac{\begin{array}{c}\Sigma_1 \Leftarrow \Sigma_2 \\ Rs_1 \Leftarrow Rs_2\end{array}}{\Sigma_1 \cup \{F :_{Rs_1} \ p' \sim_R a : A\} \Leftarrow \Sigma_2 \cup \{F :_{Rs_2} \ p \sim_R a : A\}} \quad \textsc{S\_WeakenAxiom}$$

$$\frac{}{\varnothing \Leftarrow \varnothing} \quad \textsc{S\_Empty}$$

$$\frac{\Sigma_1 \Leftarrow \Sigma_2}{\Sigma_1 \cup \{F : sig\_sort\} \Leftarrow \Sigma_2 \cup \{F : sig\_sort\}} \quad \textsc{S\_Same}$$

$\boxed{\Gamma \vDash_{\mathsf{src}} a : A}$ source typing

$$\frac{\vDash \Gamma}{\Gamma \vDash_{\mathsf{src}} \star : \star} \quad \textsc{S\_Star}$$

$$\frac{\begin{array}{c}\vDash \Gamma \\ x : A \in \Gamma\end{array}}{\Gamma \vDash_{\mathsf{src}} x : A} \quad \textsc{S\_Var}$$

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} A : \star \\ \Gamma \vDash_{\mathsf{src}} A \rightsquigarrow A' : \star \\ \Gamma, x : A' \vDash_{\mathsf{src}} B : \star\end{array}}{\Gamma \vDash_{\mathsf{src}} (\Pi^\rho x\!:\!A \to B) : \star} \quad \textsc{S\_Pi}$$

$$\frac{\Gamma, x : A \vDash_{\mathsf{src}} a : B}{\Gamma \vDash_{\mathsf{src}} \lambda x.a : (\Pi^+ x\!:\!A \to B)} \quad \textsc{S\_Abs}$$

$$\frac{\Gamma, x : A \vDash_{\mathsf{src}} a : B}{\Gamma \vDash_{\mathsf{src}} a : (\Pi^- x\!:\!A \to B)} \quad \textsc{S\_IAbs}$$

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} a : A \\ \Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Nom}\end{array}}{\Gamma \vDash_{\mathsf{src}} a : B} \quad \textsc{S\_Conv}$$

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} a : A \\ \Gamma; \widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Rep}\end{array}}{\Gamma \vDash_{\mathsf{src}} \mathbf{coerce}\, a : B} \quad \textsc{S\_Coerce}$$

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} b : \Pi^+ x\!:\!A \to B \\ \Gamma \vDash_{\mathsf{src}} a \rightsquigarrow a' : A\end{array}}{\Gamma \vDash_{\mathsf{src}} b \ a : B\{a'/x\}} \quad \textsc{S\_App}$$

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} b : \Pi^- x\!:\!A \to B \\ \Gamma \vDash a' : A\end{array}}{\Gamma \vDash_{\mathsf{src}} b : B\{a'/x\}} \quad \textsc{S\_IApp}$$

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} \phi \rightsquigarrow \phi' \ \ \mathsf{ok} \\ \Gamma, c : \phi' \vDash_{\mathsf{src}} B : \star\end{array}}{\Gamma \vDash_{\mathsf{src}} \forall c\!:\!\phi.B : \star} \quad \textsc{S\_CPi}$$

$$\frac{\Gamma, c : \phi' \vDash_{\mathsf{src}} a : B}{\Gamma \vDash_{\mathsf{src}} a : \forall c\!:\!\phi.B} \quad \textsc{S\_CAbs}$$

$$\frac{\Gamma \vDash_{\mathsf{src}} a_1 : \forall c{:}(a \sim_R b : A).B_1 \qquad \Gamma;\widetilde{\Gamma} \vDash a \equiv b : A/R}{\Gamma \vDash_{\mathsf{src}} a_1 : B_1\{\bullet/c\}} \quad \textsc{S\_CApp}$$

$$\frac{\vDash \Gamma \qquad F :_{Rs} A \in \Sigma_0}{\Gamma \vDash_{\mathsf{src}} F : A} \quad \textsc{S\_Const}$$

$$\frac{\vDash \Gamma}{\Gamma \vDash_{\mathsf{src}} F : A} \quad \textsc{S\_Fam}$$

$$\frac{\Gamma \vDash_{\mathsf{src}} a : A \qquad \Gamma \vDash_{\mathsf{src}} b_1' : B \qquad \Gamma \vDash_{\mathsf{src}} b_2' : C}{\Gamma \vDash_{\mathsf{src}} \mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \| \_ \to b_2 : C} \quad \textsc{S\_Case}$$

$\boxed{\Gamma \vDash_{\mathsf{src}} a \rightsquigarrow a' : A}$ source translation

$$\frac{\vDash \Gamma}{\Gamma \vDash_{\mathsf{src}} \star \rightsquigarrow \star : \star} \quad \textsc{ST\_Star}$$

$$\frac{\vDash \Gamma \qquad x : A \in \Gamma}{\Gamma \vDash_{\mathsf{src}} x \rightsquigarrow x : A} \quad \textsc{ST\_Var}$$

$$\frac{\Gamma \vDash_{\mathsf{src}} A \rightsquigarrow A' : \star \qquad \Gamma, x : A' \vDash_{\mathsf{src}} B \rightsquigarrow B' : \star}{\Gamma \vDash_{\mathsf{src}} (\Pi^\rho x{:}A \to B) \rightsquigarrow (\Pi^\rho x{:}A' \to B') : \star} \quad \textsc{ST\_Pi}$$

$$\frac{\Gamma, x : A \vDash_{\mathsf{src}} a \rightsquigarrow a' : B}{\Gamma \vDash_{\mathsf{src}} \lambda x.a \rightsquigarrow \lambda^+ x.a' : (\Pi^+ x{:}A \to B)} \quad \textsc{ST\_Abs}$$

$$\frac{\Gamma, x : A \vDash_{\mathsf{src}} a \rightsquigarrow a' : B \qquad x \notin \mathsf{fv}\, a}{\Gamma \vDash_{\mathsf{src}} a \rightsquigarrow \lambda^- x.a : (\Pi^- x{:}A \to B)} \quad \textsc{ST\_IAbs}$$

$$\frac{\Gamma \vDash_{\mathsf{src}} b \rightsquigarrow b' : \Pi^+ x{:}A \to B \qquad \Gamma \vDash_{\mathsf{src}} a \rightsquigarrow a' : A}{\Gamma \vDash_{\mathsf{src}} b\ a \rightsquigarrow b'\ a'^+ : B\{a'/x\}} \quad \textsc{ST\_App}$$

$$\frac{\Gamma \vDash_{\mathsf{src}} b \rightsquigarrow b' : \Pi^+ x{:}A \to B \qquad \Gamma \vDash_{\mathsf{src}} a \rightsquigarrow a' : A \qquad \mathsf{RolePath}\ b = F@R, Rs}{\Gamma \vDash_{\mathsf{src}} b\ a \rightsquigarrow b'\ a'^R : B\{a/x\}} \quad \textsc{ST\_TApp}$$

$$\frac{\Gamma \vDash_{\mathsf{src}} b \rightsquigarrow b' : \Pi^- x{:}A \to B \qquad \Gamma \vDash a : A}{\Gamma \vDash_{\mathsf{src}} b \rightsquigarrow b'\ \Box^- : B\{a/x\}} \quad \textsc{ST\_IApp}$$

$$\frac{\Gamma \vDash_{\mathsf{src}} a \rightsquigarrow a' : A \qquad \Gamma;\widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Nom}}{\Gamma \vDash_{\mathsf{src}} a \rightsquigarrow a' : B} \quad \textsc{ST\_Conv}$$

$$\frac{\Gamma \vDash_{\mathsf{src}} a \rightsquigarrow a' : A \qquad \Gamma;\widetilde{\Gamma} \vDash A \equiv B : \star/\mathbf{Rep}}{\Gamma \vDash_{\mathsf{src}} \mathbf{coerce}\, a \rightsquigarrow a' : B} \quad \textsc{ST\_Coerce}$$

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} \phi \leadsto \phi' \ \mathsf{ok} \\ \Gamma, c : \phi' \vDash_{\mathsf{src}} B \leadsto B' : \star\end{array}}{\Gamma \vDash_{\mathsf{src}} \forall c\!:\!\phi.B \leadsto \forall c\!:\!\phi.B' : \star} \quad \text{ST\_CPI}$$

$$\frac{\Gamma, c : \phi \vDash_{\mathsf{src}} a \leadsto a' : B}{\Gamma \vDash_{\mathsf{src}} a \leadsto \Lambda c.a' : \forall c\!:\!\phi.B} \quad \text{ST\_CABS}$$

$$\frac{\begin{array}{l}\Gamma \vDash_{\mathsf{src}} a_1 \leadsto a_1' : \forall c\!:\!(a \sim_R b : A).B_1 \\ \Gamma; \widetilde{\Gamma} \vDash a \equiv b : A/R\end{array}}{\Gamma \vDash_{\mathsf{src}} a_1 \leadsto a_1'[\bullet] : B_1\{\bullet/c\}} \quad \text{ST\_CAPP}$$

$$\frac{\begin{array}{l}\vDash \Gamma \\ F :_{Rs} \ A \in \Sigma_0\end{array}}{\Gamma \vDash_{\mathsf{src}} F \leadsto F : A} \quad \text{ST\_CONST}$$

$$\frac{\vDash \Gamma}{\Gamma \vDash_{\mathsf{src}} F \leadsto F : A} \quad \text{ST\_FAM}$$

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} a \leadsto a' : A \\ \Gamma \vDash_{\mathsf{src}} b_1 \leadsto b_1' : B \\ \Gamma \vDash_{\mathsf{src}} b_2 \leadsto b_2' : C\end{array}}{\Gamma \vDash_{\mathsf{src}} \mathsf{case}\ a\ \mathsf{of}\ F\overline{v} \to b_1 \| \_ \to b_2 \leadsto \mathsf{case}\ a'\ \mathsf{of}\ F\overline{v} \to b_1' \| \_ \to b_2' : C} \quad \text{ST\_CASE}$$

$\boxed{\Gamma \vDash_{\mathsf{src}} \phi \leadsto \phi' \ \mathsf{ok}}$ Prop wellformedness

$$\frac{\begin{array}{c}\Gamma \vDash_{\mathsf{src}} a \leadsto a' : A \\ \Gamma \vDash_{\mathsf{src}} b \leadsto b' : A\end{array}}{\Gamma \vDash_{\mathsf{src}} (a \sim_{\mathbf{Nom}} b : A) \leadsto (a' \sim_{\mathbf{Nom}} b' : A) \ \mathsf{ok}} \quad \text{S\_WFF}$$

$\boxed{\Gamma \vdash \phi \ \mathsf{ok}}$ prop wellformedness

$\boxed{\Gamma \vdash a : A/R}$ typing

$\boxed{\Gamma; \Delta \vdash \gamma : \phi_1 \sim \phi_2}$ coercion between props

$\boxed{\Gamma; \Delta \vdash \gamma : A \sim_R B}$ coercion between types

$\boxed{\vdash \Gamma}$ context wellformedness

$\boxed{\Gamma \vdash a \leadsto b/R}$ single-step, weak head reduction to values for annotated language

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Definition rules:        224 good    0 bad
Definition rule clauses: 614 good    0 bad
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