

Question 2

a.

Applying the Master theorem, in this case, $a = 3$, $b = 2$, $d = 0$.
 Since $\log_2 3 > 0$, we can conclude that $T(n) = O(n^{\log_2 3})$.

b.

i.

$$S \rightarrow L11|11R|L11R$$

$$L \rightarrow \varepsilon|0|L0|L10$$

$$R \rightarrow \varepsilon|0|0R|01R$$

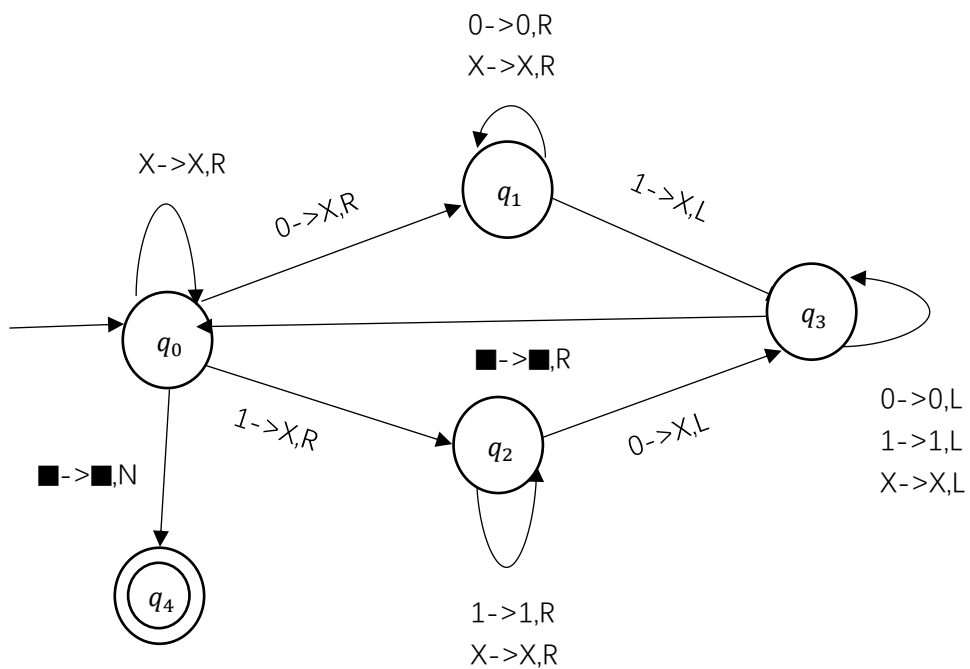
ii.

$$S \rightarrow L|R|LR|L11|11R|L11R$$

$$L \rightarrow \varepsilon|0|L0|L10$$

$$R \rightarrow \varepsilon|0|0R|01R$$

c.



d.

step 1: Divide the array into two equal halves: [5,7,12,23], [8,1,2,4]

step 2: Divide each sub-array into two equal halves: [5,7], [12,23], [8,1], [2,4]

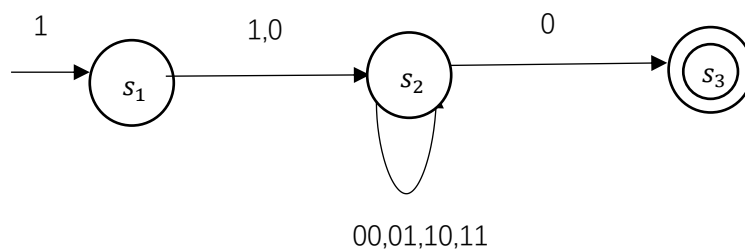
step 3: Divide each sub-array into two equal halves, and now they become arrays of unit length: [5], [7], [12], [23], [8], [1], [2], [4]

step 4: merge [5] and [7] to [5, 7] since they are already in the ascending order; merge [12] and [23] to [12, 23] since they are already in the ascending order; merge [8] and [1] to [1, 8] since $1 < 8$; merge [2] and [4] to [2, 4] since they are already in the ascending order.

Step 5: merge [5,7] and [12,23] to [5,7,12,23] since they are already in the ascending order; merge [1,8] and [2,4] to [1,2,4,8] with merge function.

Step 6: merge [5,7,12,23] and [1,2,4,8] to [1,2,4,5,7,8,12,23] with merge function.

e.



Question 3

a.

i.

after one iteration by using bubble sort, we get $A = [3, 11, 12, 2, 7, 4, 1, 17]$

after the second parsing, $A = [3, 11, 2, 7, 4, 1, 12, 17]$

ii.

$A = [7, 5, 3, 1, 2, 4, 6, 8]$

Because the pivot happens to be the smallest element in the list for each sort.

iii.

$A = [1, 2, 3, 4, 5, 6, 7, 8]$

Because the list is already in the correct order, so it will only take once for each sort.

b.

i.

$$S = A|B$$

$$A = 001|0A1|0A$$

$$B = 011|0B1|B1$$

ii.

$$L = \varepsilon|0L1|1L0|01L|10L|L01|L10$$

c.

1)for $f(n)=100n+n\log n+5$, the asymptotic function is $g(n)= n\log n$

there exist $c=101$, $k=100$, such that $f(x) \leq c \cdot g(x)$ for all $x > k$.

Proof: $c \cdot g(x) = 101x \log x$.

$c \cdot g(x) - f(x) = 101x \log x - 100x - x \log x - 5 = 100x(\log x - 1) - 5$, since $\log x > 2$ when $x > 100$, so $100x(\log x - 1) - 5 > 100 \cdot 100 \cdot (2 - 1) - 5 = 9995 > 0$

so we conclude that $g(n) = n \log n$ is the asymptotic function.

2)for $f(n) = n^{0.5} + 10 \log n + 40$, the asymptotic function is $g(n) = n^{0.5}$

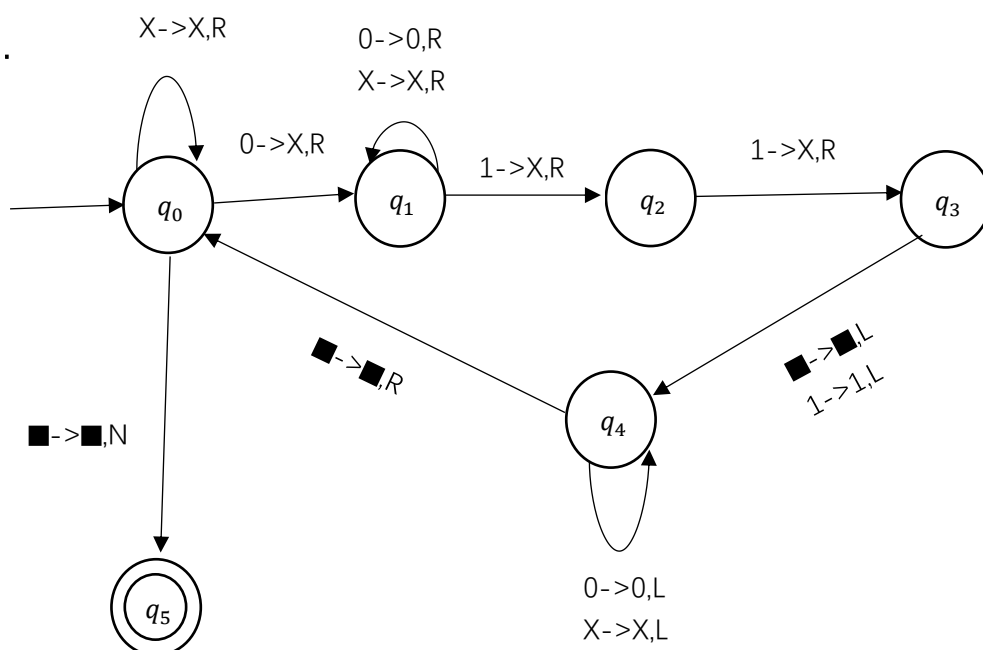
There exist $c=11$, $k=100$, such that $f(x) \leq c \cdot g(x)$ for all $x > k$.

Proof: $c \cdot g(x) = 11x^{0.5}$

$c \cdot g(x) - f(x) = 10x^{0.5} - 10 \log x + 40 = 10(x^{0.5} - \log x) + 40$, since $x^{0.5} > 10$ when $x > 100$, and $\log x > 2$ when $x > 100$, and $x^{0.5}$ is growing faster than $\log x$, so $10(x^{0.5} - \log x) + 40 > 10 \cdot (10 - 2) + 40 = 120 > 0$

so we conclude that $g(n) = n^{0.5}$ is the asymptotic function.

d.



e.

if there exists zero 0, then the expression should be 1^+

if there exists one 0, then the expression should be $1^+01^+ \cup 1^*01^+$

if there exists two 0, then the expression should be $1^+01^*01^+ \cup 1^*01^+01^+ \cup 1^*01^*01^+$

In all, the regular expression will be $1^+ \cup 1^*01^+ \cup 1^*01^+ \cup 1^+01^*01^+ \cup 1^*01^+01^+ \cup 1^*01^*01^+$