



CS180 Discussion 1B

Week 3: Algorithm Analysis

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Outline

- Review of Asymptotic notation
- Survey of common running times
- Exercises

Big O notation (O) vs. Little o notation (o)

$O(g(n)) = \{f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq f(n) \leq cg(n) \text{ for all } n \geq n_0\} .$

$o(g(n)) = \{f(n) : \text{for any positive constant } c > 0, \text{ there exists a constant } n_0 > 0 \text{ such that } 0 \leq f(n) < cg(n) \text{ for all } n \geq n_0\} .$

➡ Eg: $2n = o(n^2)$, but $2n^2 \neq o(n^2)$

➡ $f(n) = o(g(n))$ implies: $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 .$

Big Omega notation (Ω) vs. Little Omega notation (ω)

$\Omega(g(n)) = \{f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq cg(n) \leq f(n) \text{ for all } n \geq n_0\} .$

$\omega(g(n)) = \{f(n) : \text{for any positive constant } c > 0, \text{ there exists a constant } n_0 > 0 \text{ such that } 0 \leq cg(n) < f(n) \text{ for all } n \geq n_0\} .$

➡ Eg: $n^2/2 = \omega(n)$, but $n^2/2 \neq \omega(n^2)$

➡ $f(n) = \omega(g(n))$ implies: $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$

Big Theta notation (Θ)

$\Theta(g(n)) = \{f(n) : \text{there exist positive constants } c_1, c_2, \text{ and } n_0 \text{ such that}$
 $0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \text{ for all } n \geq n_0\} .^1$

➡ No “little” Theta!!

Asymptotic bounds and limits

6

Proposition. If $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = c$ for some constant $0 < c < \infty$ then $f(n)$ is $\Theta(g(n))$.

Pf.

- By definition of the limit, for any $\varepsilon > 0$, there exists n_0 such that

$$c - \varepsilon \leq \frac{f(n)}{g(n)} \leq c + \varepsilon$$

for all $n \geq n_0$.

- Choose $\varepsilon = \frac{1}{2} c > 0$.
- Multiplying by $g(n)$ yields $\frac{1}{2} c \cdot g(n) \leq f(n) \leq \frac{3}{2} c \cdot g(n)$ for all $n \geq n_0$.
- Thus, $f(n)$ is $\Theta(g(n))$ by definition, with $c_1 = \frac{1}{2} c$ and $c_2 = \frac{3}{2} c$. ■

$$f(n) = o(g(n))$$

Proposition. If $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$, then $f(n)$ is $O(g(n))$ but not $\Omega(g(n))$.

$$f(n) = \omega(g(n))$$

Proposition. If $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty$, then $f(n)$ is $\Omega(g(n))$ but not $O(g(n))$.

Properties

Transitivity:

$f(n) = \Theta(g(n))$ and $g(n) = \Theta(h(n))$ imply $f(n) = \Theta(h(n))$,

$f(n) = O(g(n))$ and $g(n) = O(h(n))$ imply $f(n) = O(h(n))$,

$f(n) = \Omega(g(n))$ and $g(n) = \Omega(h(n))$ imply $f(n) = \Omega(h(n))$,

$f(n) = o(g(n))$ and $g(n) = o(h(n))$ imply $f(n) = o(h(n))$,

$f(n) = \omega(g(n))$ and $g(n) = \omega(h(n))$ imply $f(n) = \omega(h(n))$.

Properties

Reflexivity:

$$f(n) = \Theta(f(n)) ,$$

$$f(n) = O(f(n)) ,$$

$$f(n) = \Omega(f(n)) .$$

Symmetry:

$$f(n) = \Theta(g(n)) \text{ if and only if } g(n) = \Theta(f(n)) .$$

Transpose symmetry:

$$f(n) = O(g(n)) \text{ if and only if } g(n) = \Omega(f(n)) ,$$

$$f(n) = o(g(n)) \text{ if and only if } g(n) = \omega(f(n)) .$$

Analogy between asymptotic comparison and the comparison of two real numbers

$f(n) = O(g(n))$ is like $a \leq b$,

$f(n) = \Omega(g(n))$ is like $a \geq b$,

$f(n) = \Theta(g(n))$ is like $a = b$,

$f(n) = o(g(n))$ is like $a < b$,

$f(n) = \omega(g(n))$ is like $a > b$.

Analogy between asymptotic comparison and the comparison of two real numbers

- One property of real numbers, however, **DOES NOT** carry over to asymptotic notation:

Trichotomy: For any two real numbers a and b , exactly one of the following must hold: $a < b$, $a = b$, or $a > b$.

- E.g: we cannot compare the function n and $n^{1+\sin n}$ using asymptotic notation, since the value of the exponent in $n^{1+\sin n}$ oscillates between 0 and 2, taking on all values in between.

Survey of common running times

(See separate slides)



12

Exercises

Thank you!