

(Monotonic Stack / queue)

Monotonic stack :- Which specify some order

+ Greater Element :-

1	2	5	3	1	2	5	3	1	2	4	6
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GIP

2	-1	6	5	2	5	6	4	2	4	6	-1
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O/P

pt:- Decreasing order



let:-

<int> find_NGE(arr[])

{ nge[n] , stack of

for (i=n-1→0)

{

 while (!st.empty () && st.top ≤ arr[i])
 st.pop()

 if (st.empty ()) nge[i] = -1

 else nge[i] = st.top()

 st.push (arr[i])

}

TC :- O(N)

SC :- O(N)

return nge

v) Next Greater Element - v :- (Circular)

arr [2, 10, 12, 1, 11] \rightarrow [10, 12, 1, 11, 12]

\rightarrow Brute :-

arr = [2, 10, 12, 1, 11] [2, 10, 12, 1, 11]
0 1 2 3 4 5 6 7 8 9
hypothesis copy

\Rightarrow nge[n]

for ($i=0 \rightarrow n-1$) {

 for ($j=i+1 \rightarrow i+N-1$) {

 ind = j % N

 if (arr[ind] > arr[i]) {

 nge[i] = arr[ind]

 break

}

TC :- $O(N^2)$

SC :- $O(N)$

\rightarrow Optimal :-

for ($i=n-1 \rightarrow 0$) {

 while (!st.empty() && st.top() \leq arr[i+1:N]) {

 st.pop();

 if ($i < N$) \rightarrow real array

{

 nge[i] = st.empty() ? -1 : st.top()

}

due to hypo
array from

TC :- $O(4N)$

SC :- $O(2N)$

st.push(arr[i+1:N]) \rightarrow Push every element real or

{

hypothetical

return nge



* Smaller Element :-

arr = [4, 5, 2, 10, 8] \rightarrow [-1, 4, -1, 2, 2]
use increasing order
note :-
 $e[n]$

for ($i=0 \rightarrow n-1$) {

 for ($j=i+1 \rightarrow n$) {

 if (arr[j] < arr[i])

 nge[i] = arr[j], break

TC:- $O(N^2)$

SC:- $O(N)$

turn nge

Final :-

for ($i=0 \rightarrow n-1$) {

 while (!st.empty() && st.top() \geq arr[i])

 st.pop();

TC:- $O(2N)$

}

SC:- $O(N)$

 nge[i] = st.empty() ? -1 : st.top();

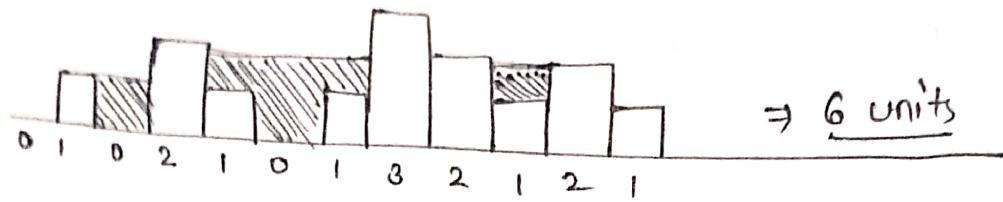
 st.push(arr[i]);

}

turn nge;

Q) Trapping Rain water :-

$$\text{arr} = [0, 1, 0, 2, 1, 0, 1, 3, 2, 1, 2, 1]$$



→ By Prefix sum Arrays (left max) right max :-

$$\sum_{i=0}^{n-1} \min(\text{leftmax} - \text{rightmax}) - \text{arr}[i]$$

$$\Rightarrow \text{total} = 0$$

for ($i=0 \rightarrow n-1$) { — $O(N)$

 leftmax = Prefix Max (i) rightmax = suffix Max (i)
 if ($\text{arr}[i] < \text{leftmax}$ & & $\text{arr}[i] < \text{rightmax}$)

 total += $\min(\text{leftmax}, \text{rightmax}) - \text{arr}[i]$;

}

→ Prefix Max :-

$$\text{arr} = [2, 1, 0, 5, 3]$$

TC :- $\leq O(3N)$

$$\text{prefix} = [2, 2, 2, 5, 5]$$

SC :- $\leq O(N)$

$$\text{prefix}[0] = \text{arr}[0]$$

for ($i=1 \rightarrow n-1$) { — $O(N)$

$$\text{prefix}[i] = \max(\text{prefix}[i-1], \text{arr}[i])$$

}

$$\text{arr} = [1, 11, 2, 10]$$

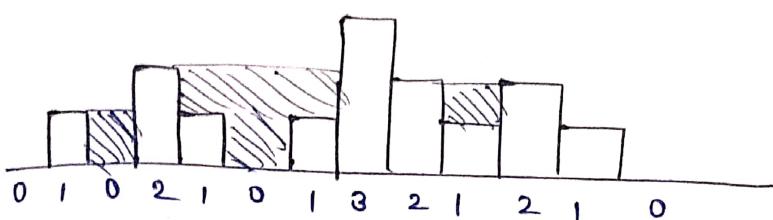
$$\text{suffix} = [11, 11, 10, 10]$$

$$\text{arr}[n-1] = \text{arr}(n-1) \rightarrow O(N)$$

$$(i=n-2 \rightarrow 0) \downarrow$$

$$\text{suffix}[i] = \max([\text{suffix}(i-1), \text{arr}(i)])$$

to pointer's Approach :-



$$\text{arr} = r\text{max} = \text{total} = 0 \quad l=0, r=n-1$$

$$\text{if } (l < r) \{ \rightarrow O(N)$$

$$\text{if } (\text{lmax} > \text{arr}(r)) \{$$

$$\text{if } (\text{lmax} > \text{arr}(l)) \}$$

$$\text{total} += \text{lmax} - \text{arr}(l);$$

else

$$TC_l \leftarrow \Sigma O(N)$$

$$\text{lmax} = \text{arr}(l)$$

$$l = l + 1$$

$$SC_l \leftarrow \Sigma O(1)$$

}

else {

$$\text{if } (r\text{max} > \text{arr}(r))$$

$$\text{total} += r\text{max} - \text{arr}(r)$$

else

$$r\text{max} = \text{arr}(r)$$

$$r = r - 1$$

}