

(Monotonic stack / queue)

1

Monotonic stack :- which specify some order
+ Greater Element :-

1	2	5	3	1	2	5	3	1	2	4	6
---	---	---	---	---	---	---	---	---	---	---	---

gip

-1	6	5	2	5	6	4	2	4	6	-1
----	---	---	---	---	---	---	---	---	---	----

o/p

pt:- Decreasing order



let:-

<int> find NGE (arr[])

{ nge[], stack st

for (i = n-1 → 0)

{

while (!st.empty() && st.top() ≤ arr[i])

st.pop()

if (st.empty()) nge[i] = -1

else nge[i] = st.top()

st.push(arr[i])

}

return nge

TC :- $O(2N)$

SC :- $O(N)$

4) Next Greater Element - 2 :- (Circular)

$$\text{arr}[2, 10, 12, 1, 11] \rightarrow [10, 12, 1, 11, 12]$$

→ Brute :-

arr = [2, 10, 12, 1, 11] [2, 10, 12, 1, 11] hypothetical copy

0 1 2 3 4 5 6 7 8 9

$$\Rightarrow \log(n)$$

```
for (i=0 → n-1) {
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for($j=i+1 \rightarrow i+N-1$)

$$i_{nd} = j \cdot 1 \cdot N$$

```
if (arr[ind] > arr[j]) {
```

$$\arg(z) = \arctan\left(\frac{y}{x}\right)$$

```
break }
```

3

3

TC:- $O(N^2)$

$$\text{SCl}_2 - \text{O}(\text{CN})$$

→ Optimal :-

for ($i = n-1 \rightarrow 0$) {

```
while (!st.empty() && st.top() <= arr[i+1..N]) {
    st.pop(); }
```

```

    st pop(); }

```

if ($i < N$) \rightarrow real array

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{
    if (nge[i] == empty()) - : st.top()
}

```

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St. push (arr [1..N]) $\xrightarrow{\quad}$ Push every element real or hypothetical

2

return ~~to~~ nge

Smaller Element:-

arr = [4, 5, 2, 10, 8] → [-1, 4, -1, 2, 2]

use increasing order

rate:-
ge[n]



for (i=0 → n-1) {

for (j=i-1 → 0) {

if (arr[j] < arr[i])

nge[i] = arr[j], break

}

TC:- $O(N^2)$

SC:- $O(N)$

return nge

final:-

for (i=0 → n-1) {

while (!st.empty() && st.top() >= arr[i])

{ st.pop();

}

TC:- $O(2N)$

SC:- $O(N)$

nge[i] = st.empty() ? -1 : st.top()

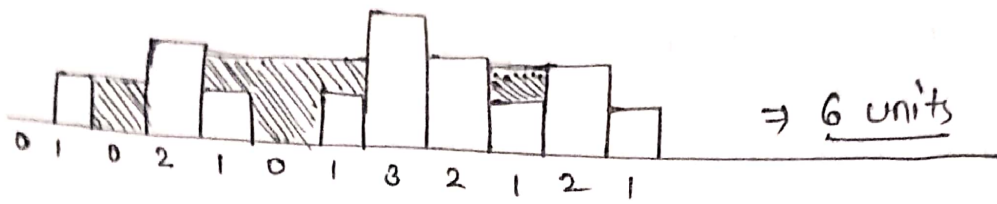
st.push(arr[i]);

}

return nge;

9) Trapping Rain water :-

arr = [0, 1, 0, 2, 1, 0, 1, 3, 2, 1, 2, 1]



⇒ By Prefix sum Arrays (left max / right max) :-

$$\sum_{i=0}^{n-1} \min(\text{leftmax}, \text{rightmax}) - \text{arr}[i]$$

⇒ total = 0

for (i = 0 → n-1) { — O(N)

leftmax = PrefixMax(i) rightmax = suffixMax(i)

if (arr[i] < leftmax && arr[i] < rightmax)

total += min(leftmax, rightmax) - arr[i];

}

return total

→ Prefix Max :-

arr = [2, 1, 0, 5, 3]

prefix = [2, 2, 2, 5, 5]

prefix[0] = arr[0]

for (i = 1 → n-1) { — O(N)

prefix[i] = max(prefix[i-1], arr[i])

}

TC :- O(3N)

SC :- O(2N)

$$arr = [1, 11, 2, 10]$$

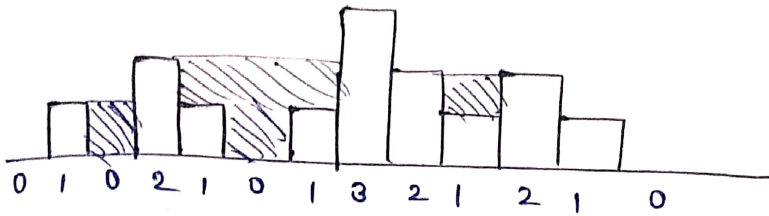
$$suffix = [11, 11, 10, 10]$$

$$arr[n-1] = arr[n-1] \quad \text{--- } O(N)$$

$$(i = n-2 \rightarrow 0) \{$$

$$suffix[i] = \max(suffix[i+1], arr[i])$$

So pointer's Approach :-



$$l = r = \text{max} = \text{total} = 0 \quad l = 0, r = n-1$$

$$\text{while } (l < r) \{ \quad \text{--- } O(N)$$

$$\text{if } (arr[l] \leq arr[r]) \{$$

$$\text{if } (lmax > arr[l])$$

$$\text{total} += lmax - arr[l];$$

else

$$lmax = arr[l]$$

$$l = l + 1$$

}

else {

$$\text{if } (rmax > arr[r])$$

$$\text{total} += rmax - arr[r]$$

else

$$rmax = arr[r]$$

}

$$r = r - 1$$

$$TC \sim O(N)$$

$$SC \sim O(1)$$