



Parameterized Algorithms for Non-uniform All-to-all

Presenter: Jens Domke

Authors: Ke Fan, Jens Domke, Seydou Ba, and Sidharth Kumar



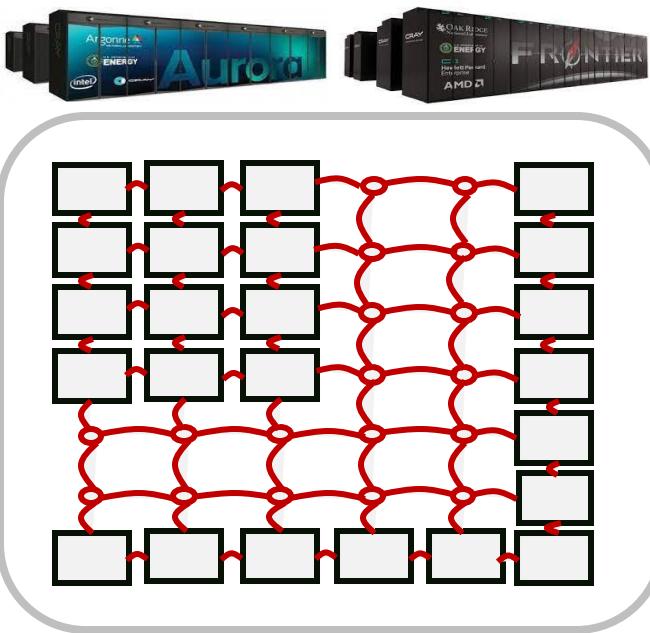
Agenda

- Introduction and Background
 - Inter-process Communication
 - All-to-all Data Exchange
 - Standard MPI All-to-all Implementations
- Challenges and Solutions
 - Tunability of Performance
 - Logarithmic Non-uniform All-to-all
- Parameterized Logarithmic Algorithm
- Parameterized Linear Algorithm
- Evaluation
- Application
- Conclusion

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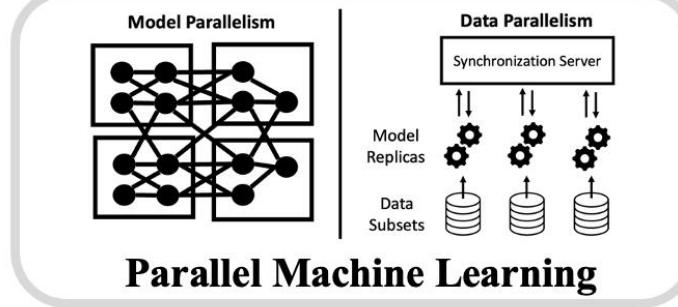
Exascale Computing Necessaries Effective Parallelism



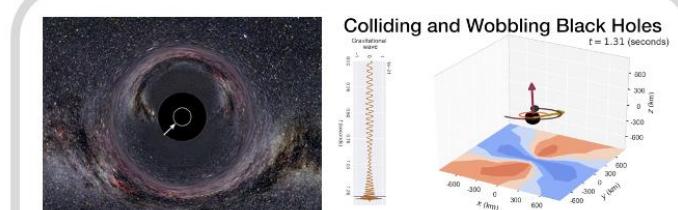
Exascale Supercomputers



Earth and Climate Modeling



Parallel Machine Learning



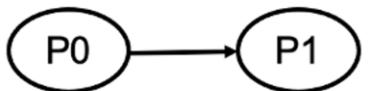
Black Holes Simulation



Inter-process Data
Movement

Three Fundamental Interfaces of Data Movement

Point-to-point



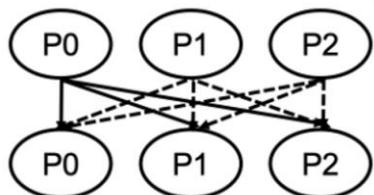
MPI_Send
MPI_Recv

One-sided

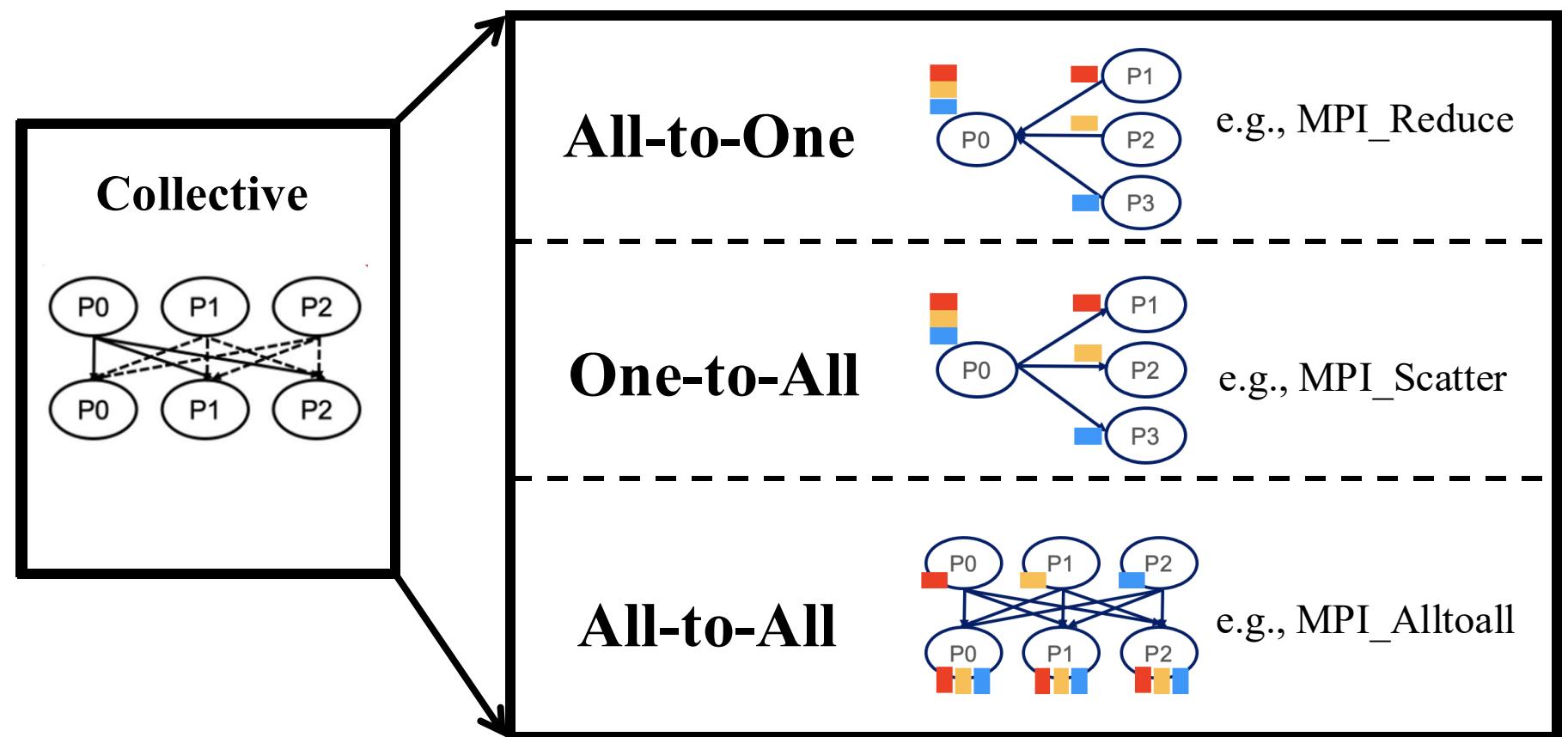
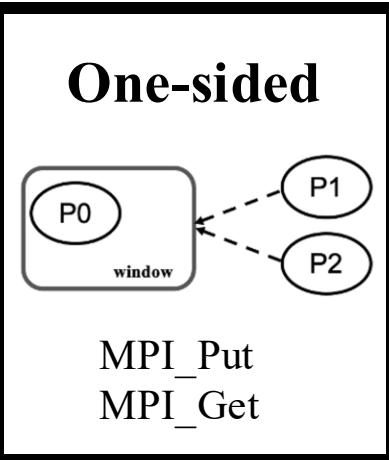
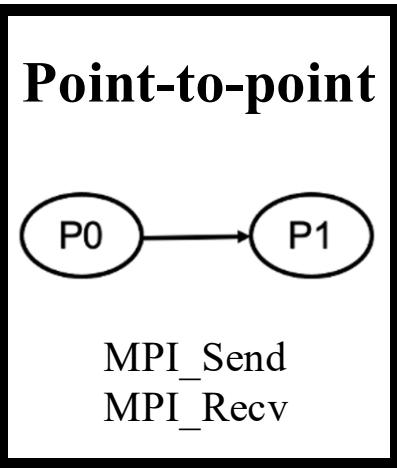


MPI_Put
MPI_Get

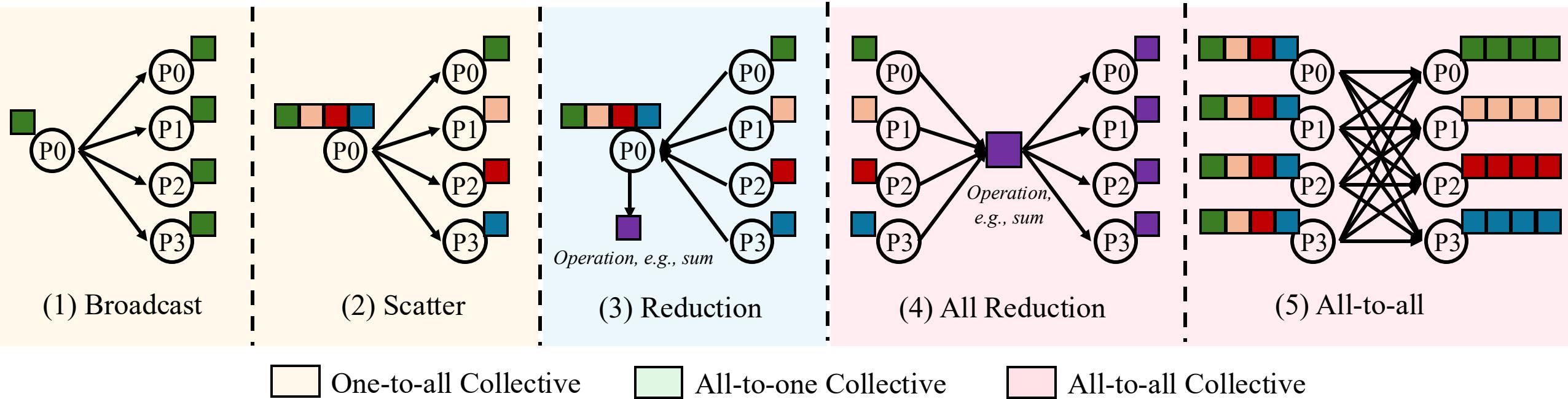
Collective



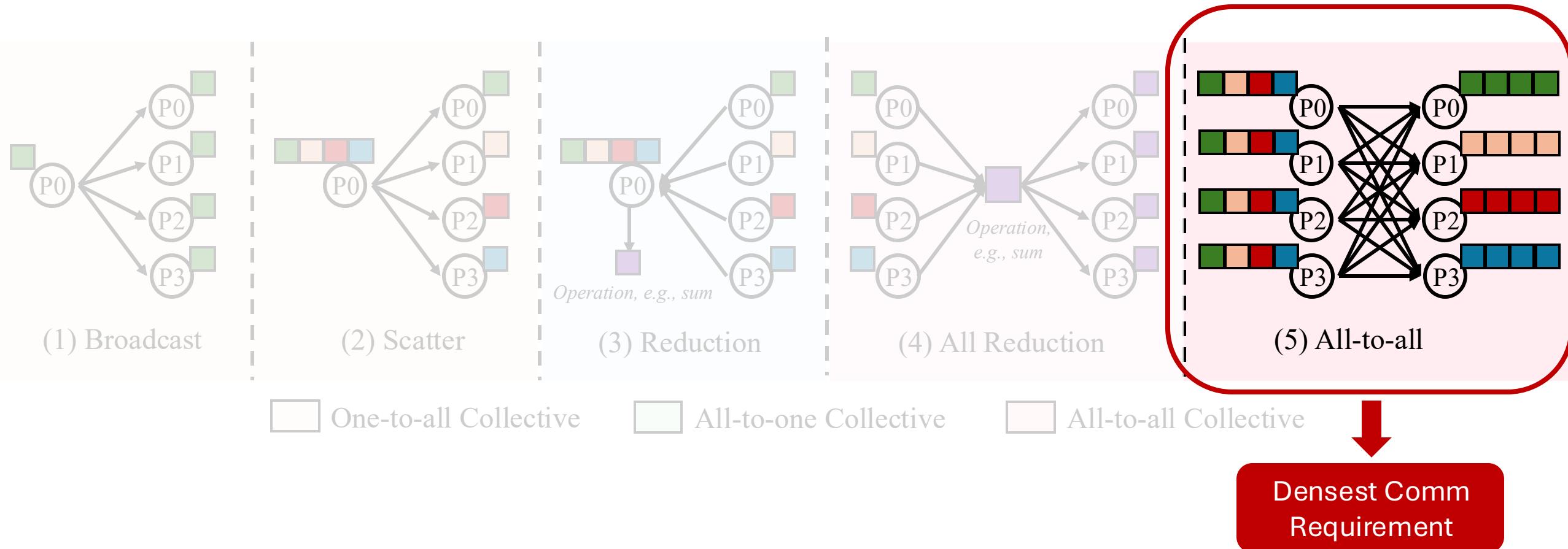
Collective Involves Communication Among All Processes



Some Widely Used Collective Primitives



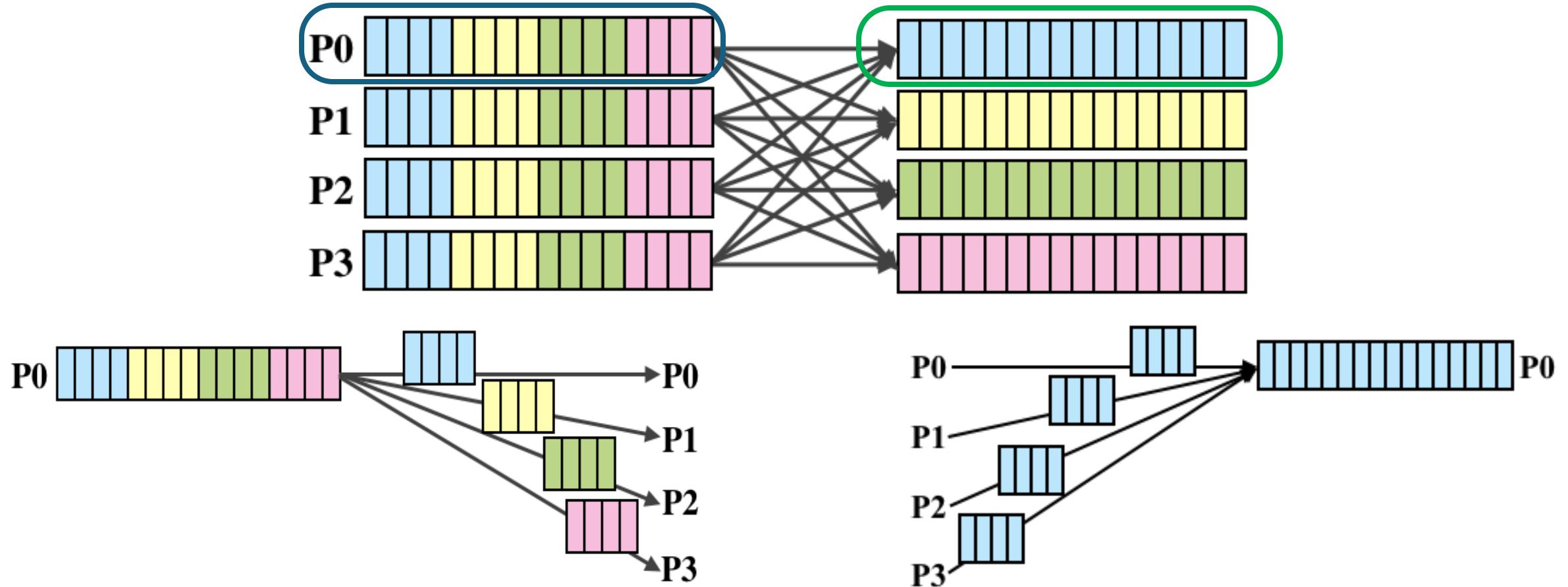
All-to-all Data Exchange is the Most Challenging to Scale and Optimize



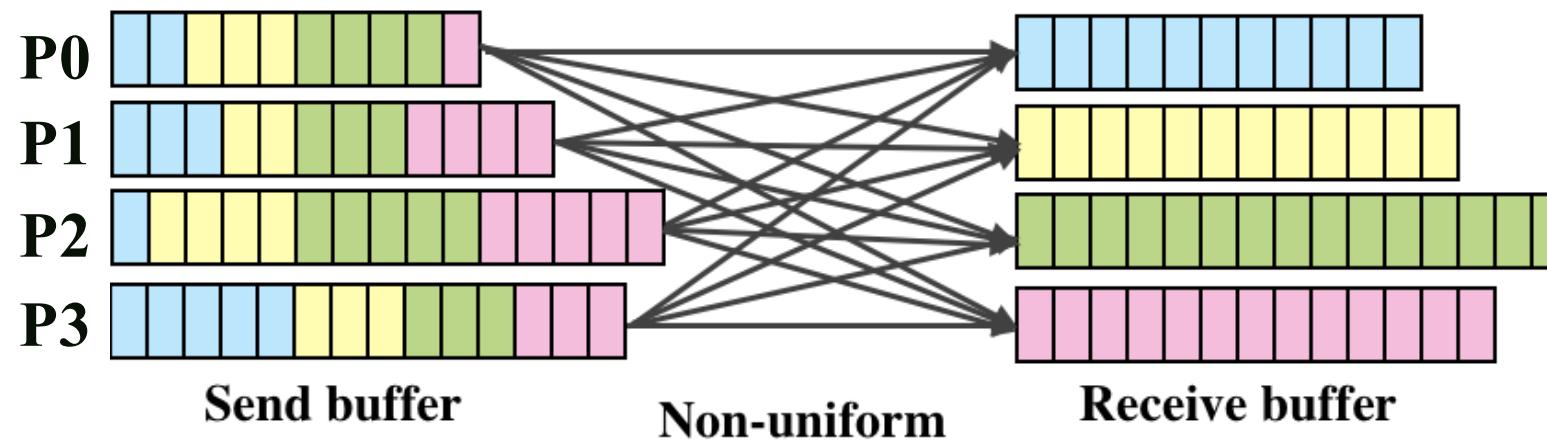
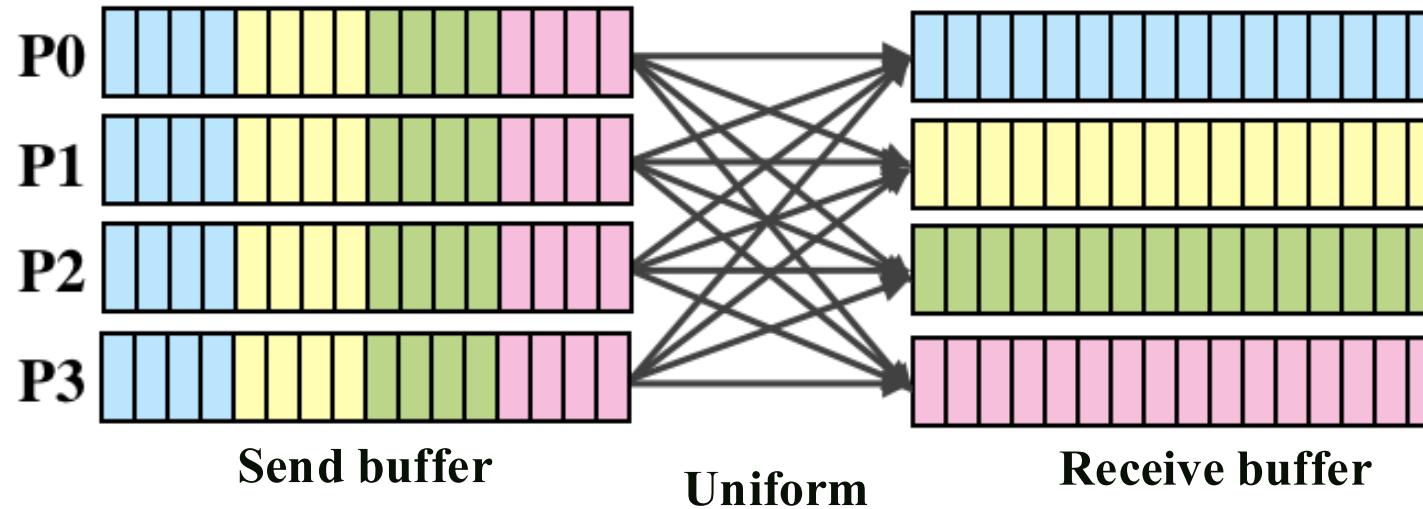
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All-to-all – Every process Sends and receives data from every other process



Uniform and Non-uniform All-to-all



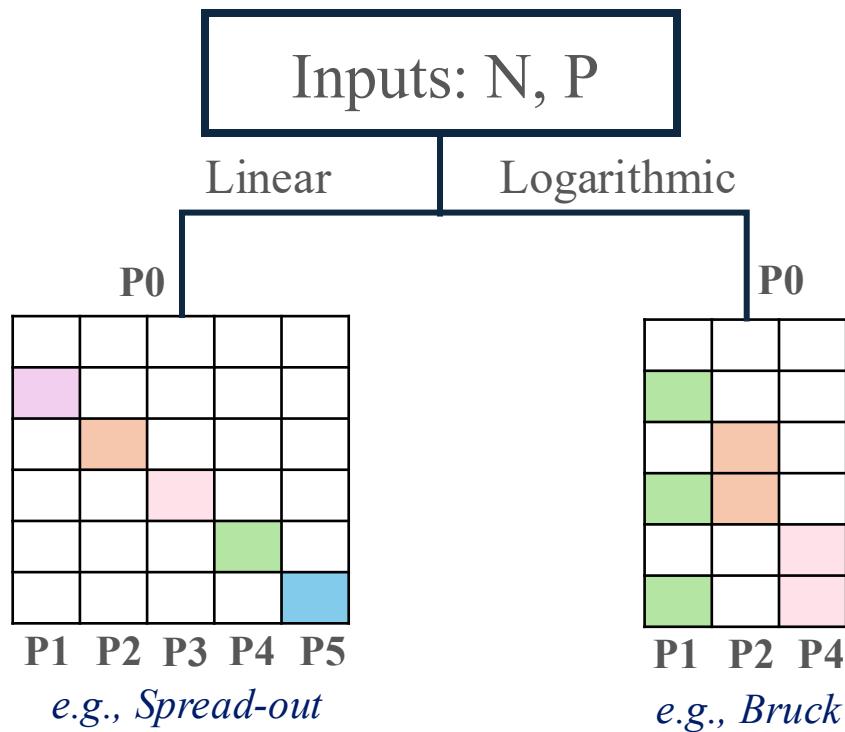
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Standard MPI All-to-all Implementations

N : Size of data-block

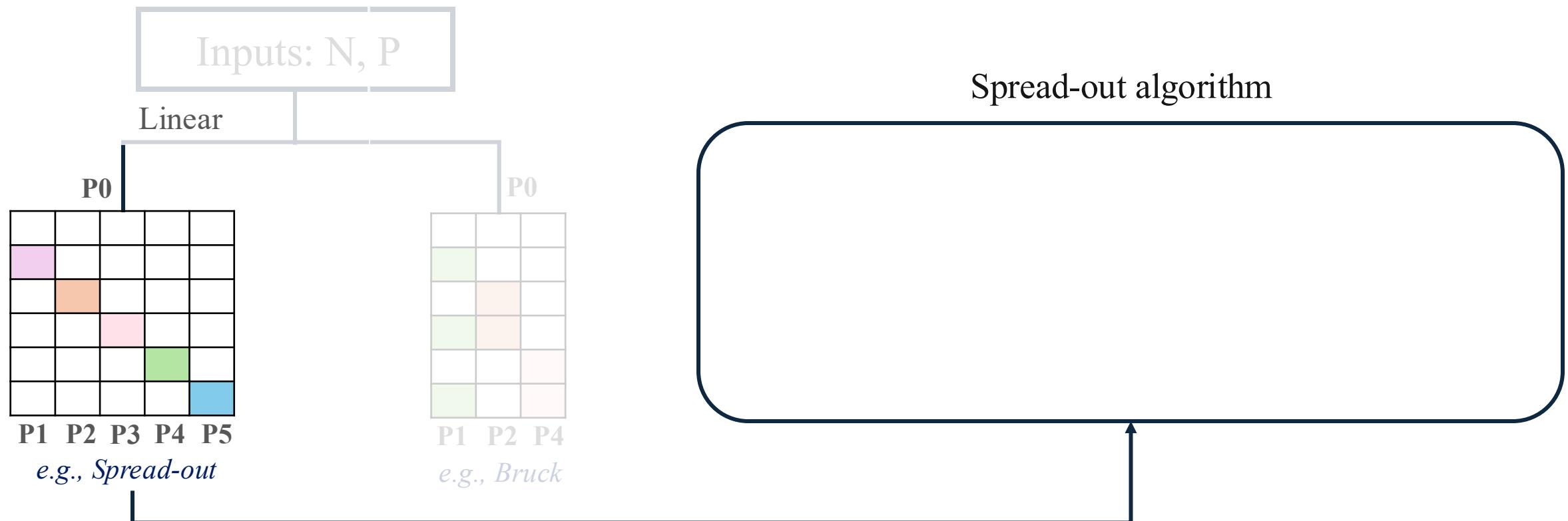
P : Process count



Standard MPI All-to-all Implementations: Spread-out

N: Size of data-block

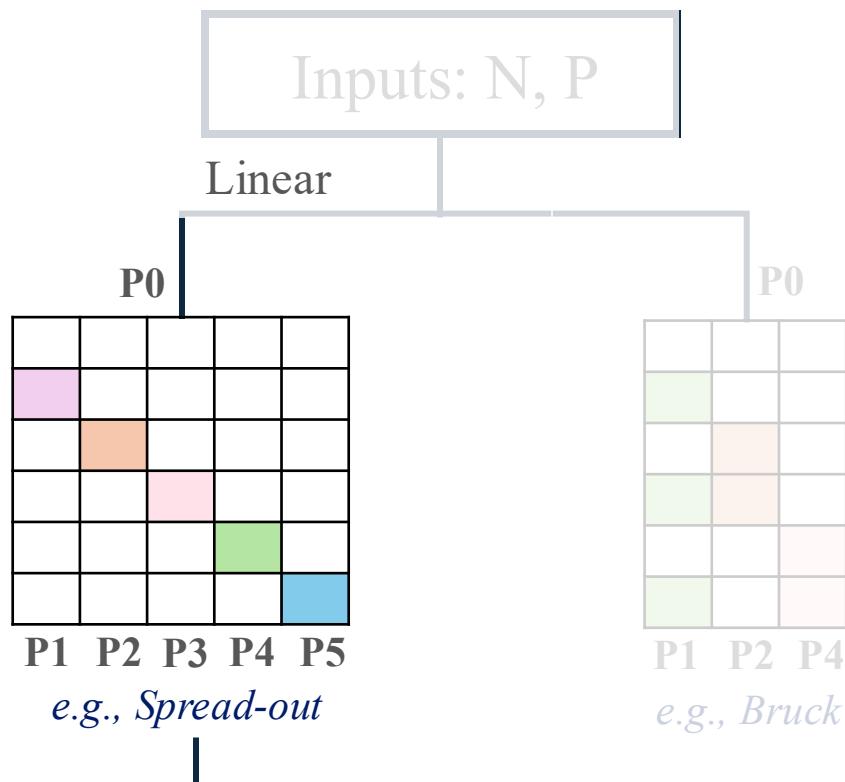
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Standard MPI All-to-all Implementations: Spread-out

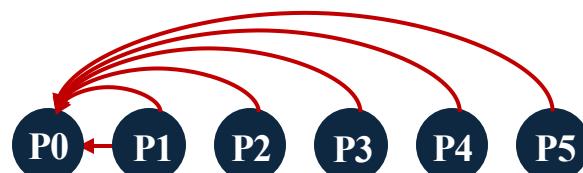
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Spread-out algorithm

**Pose receive requests
using `MPI_Irecv`**



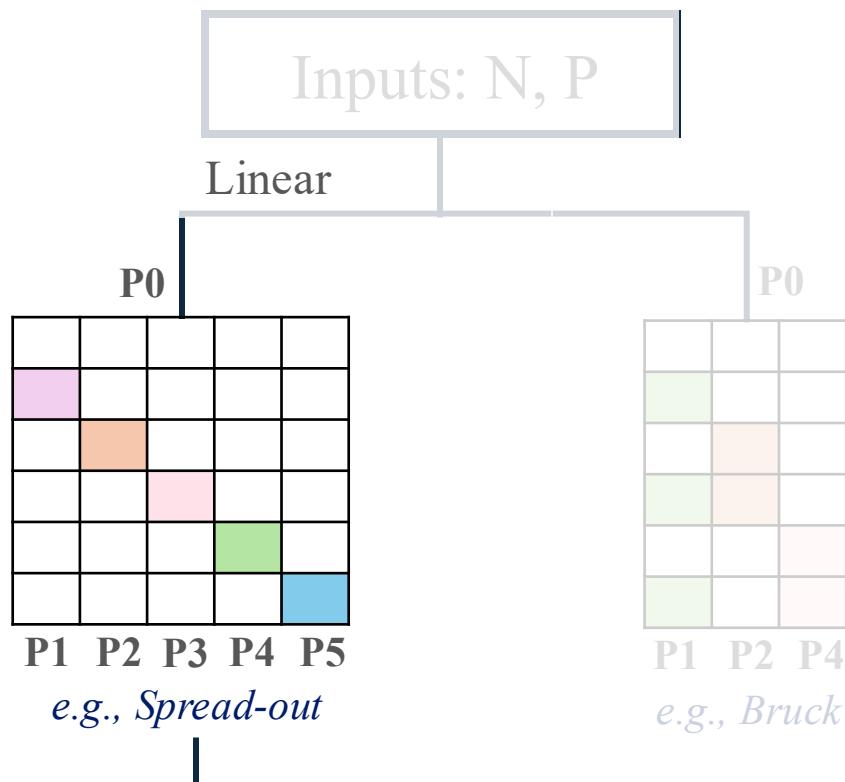
$(rank + i) \% P$

Taking P_0 as example

Standard MPI All-to-all Implementations: Spread-out

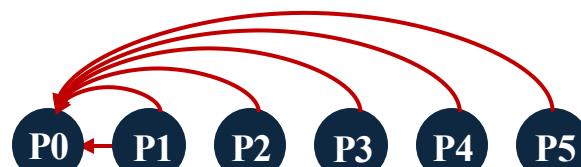
N: Size of data-block

P: Process count

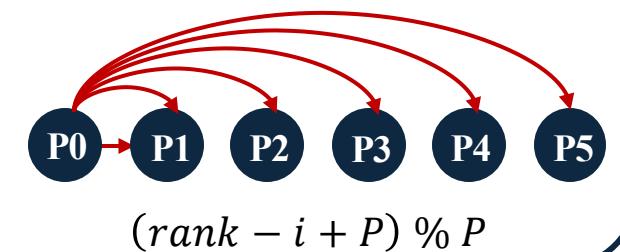


Spread-out algorithm

**Pose receive requests
using MPI_Irecv**



**Pose send requests
using MPI_Isend**

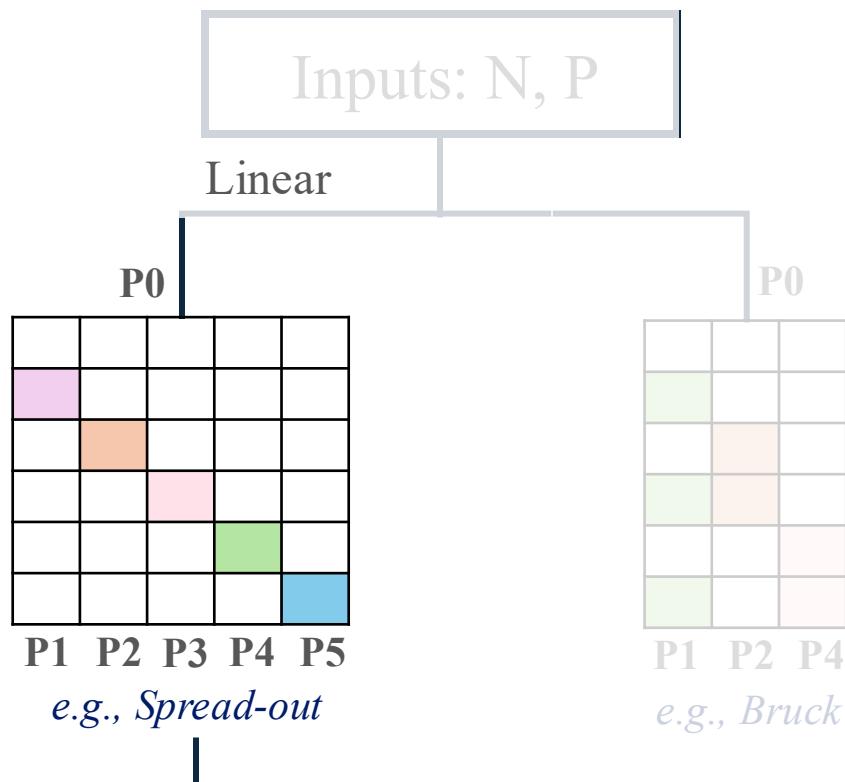


Taking P0 as example

Standard MPI All-to-all Implementations: Spread-out

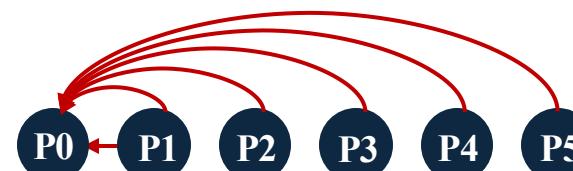
N : Size of data-block

P : Process count

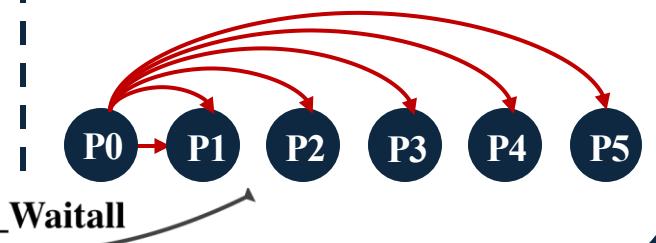


Spread-out algorithm

Pose receive requests
using **`MPI_Irecv`**



Pose send requests
using **`MPI_Isend`**

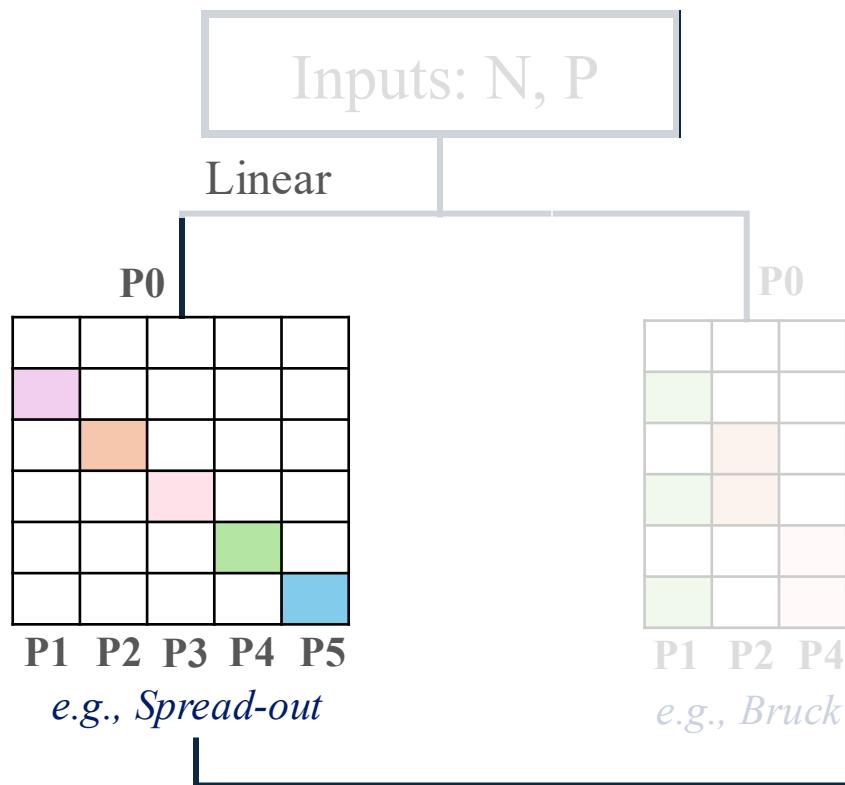


Taking $P0$ as example

Standard MPI All-to-all Implementations: Spread-out

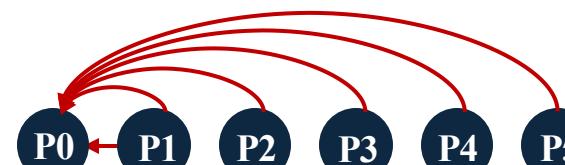
N: Size of data-block

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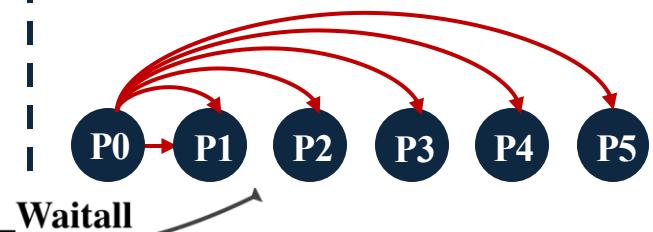
Spread-out algorithm

**Pose receive requests
using MPI_Irecv**



MPI_Waitall

**Pose send requests
using MPI_Isend**

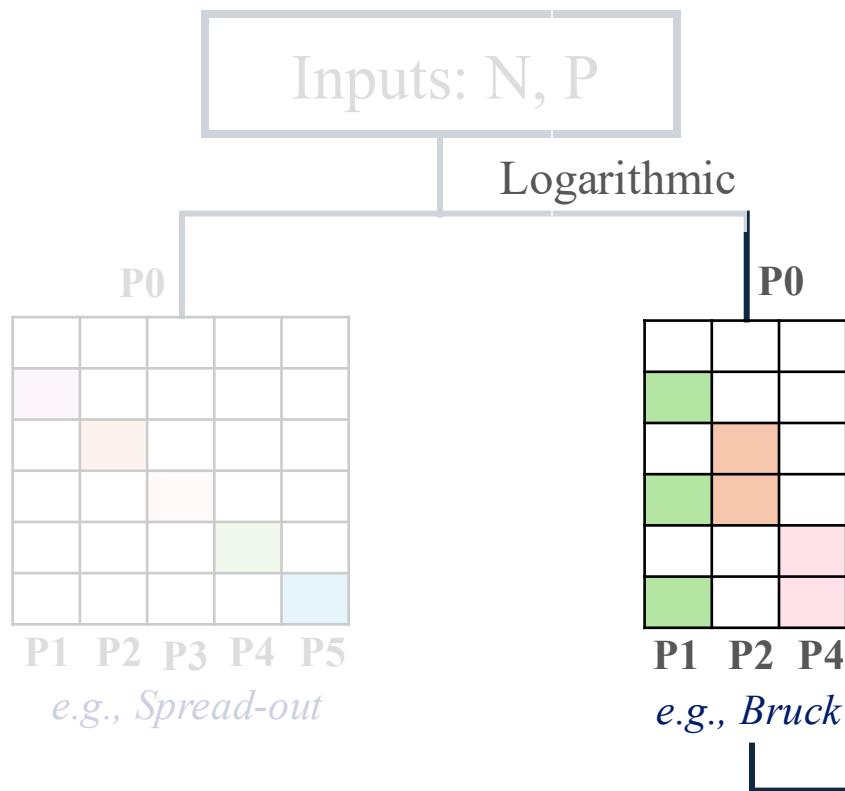


(P - 1) communication steps.

Standard MPI All-to-all Implementations: Bruck

N : Size of data-block

P : Process count



Bruck algorithm

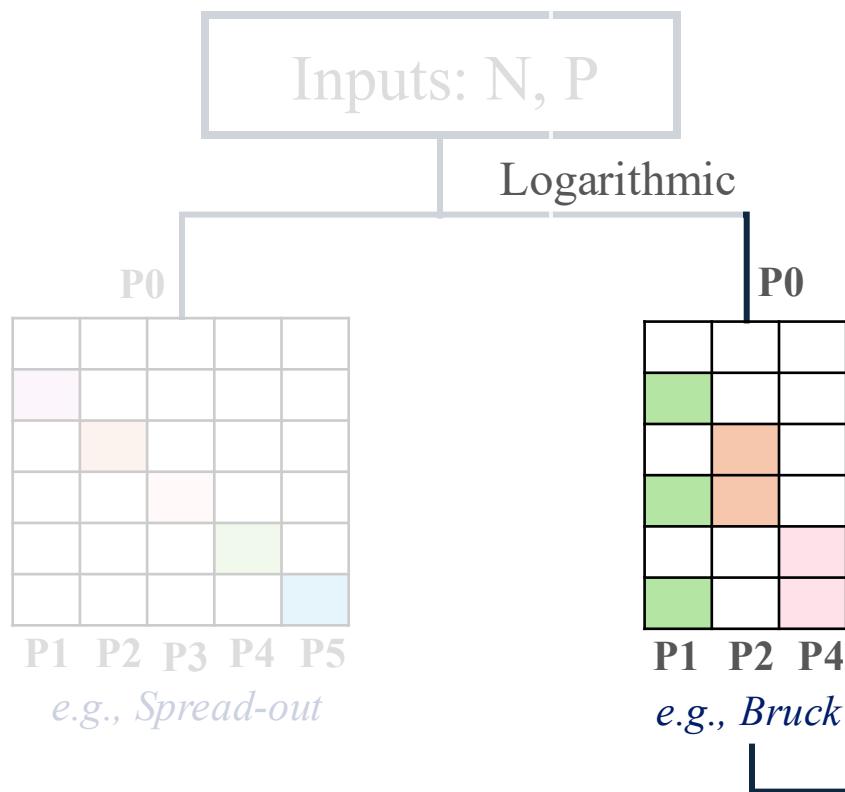


$\log(P)$ communication steps

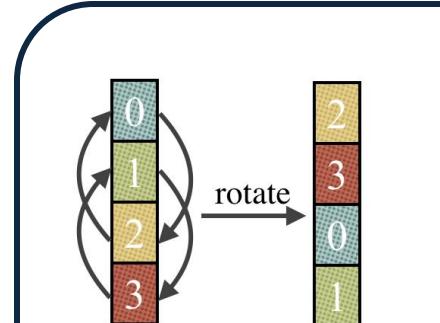
Standard MPI All-to-all Implementations: Bruck

N : Size of data-block

P : Process count



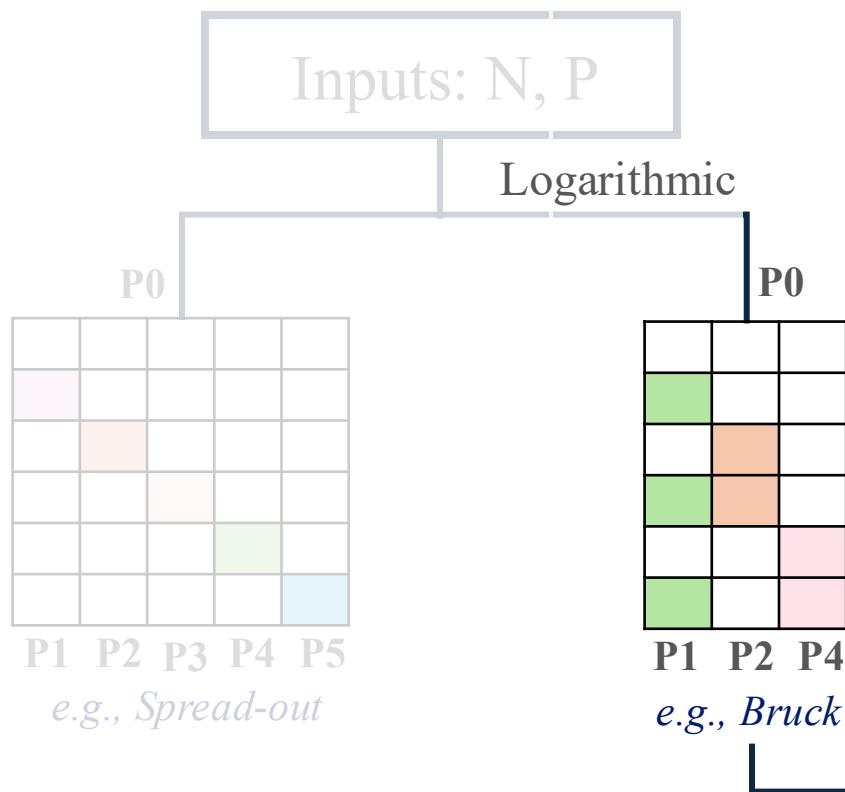
Bruck algorithm



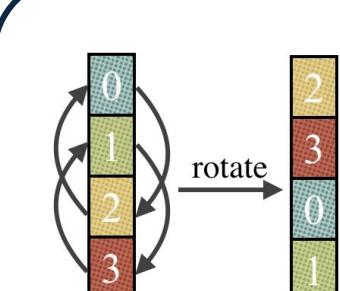
Standard MPI All-to-all Implementations: Bruck

N: Size of data-block

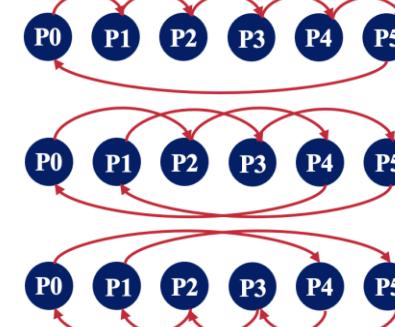
P: Process count



Bruck algorithm



2. $\log(P)$ communication steps.

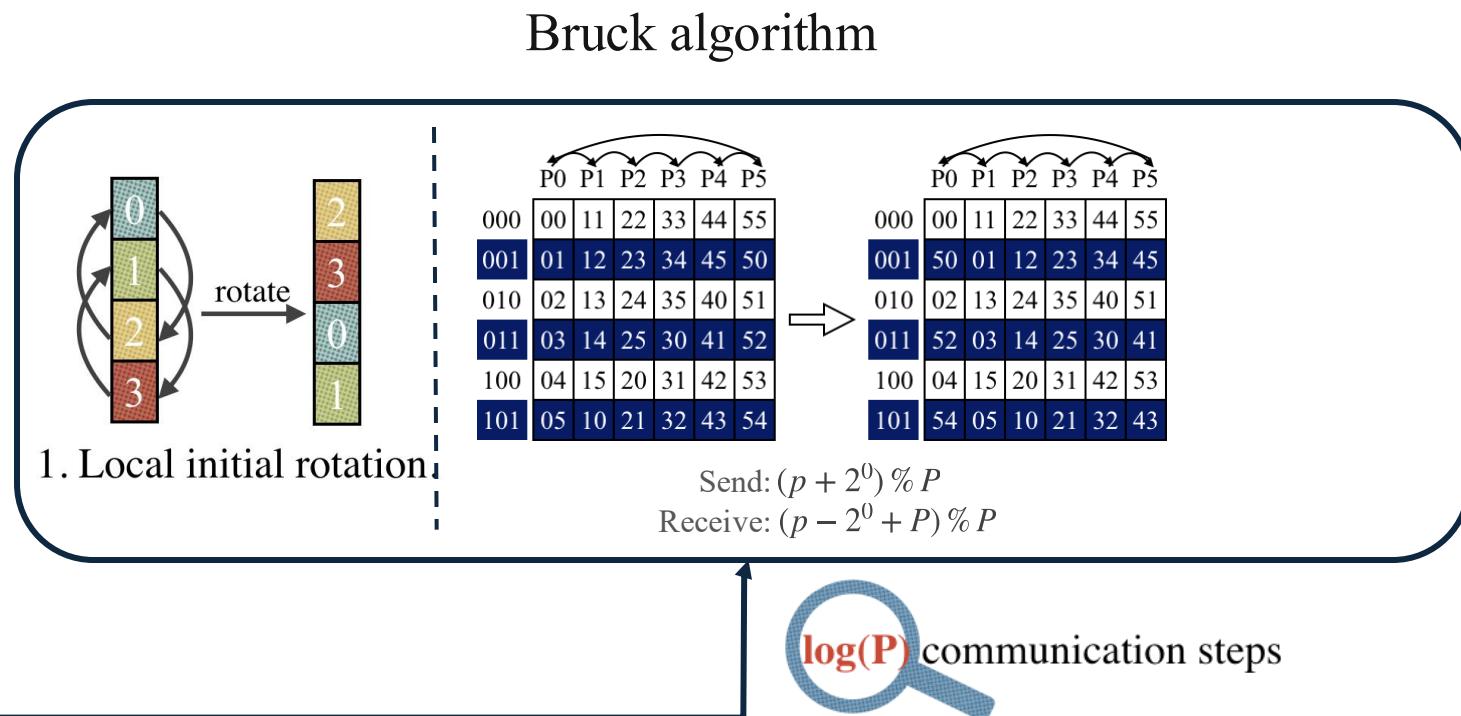
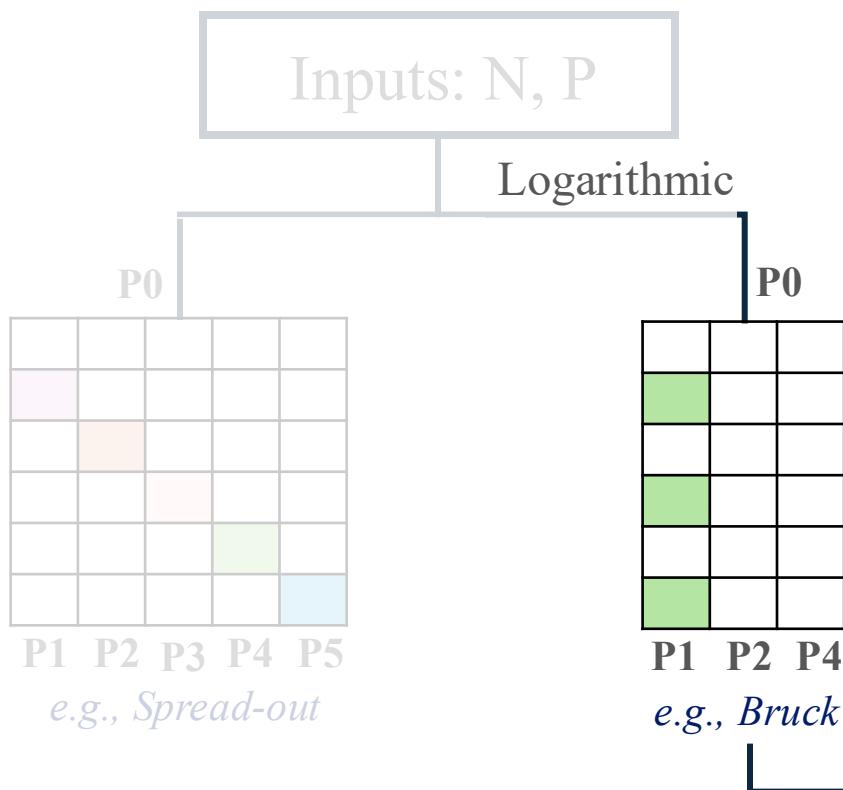


$\log(P)$ communication steps

Standard MPI All-to-all Implementations: Comm-0

N: Size of data-block

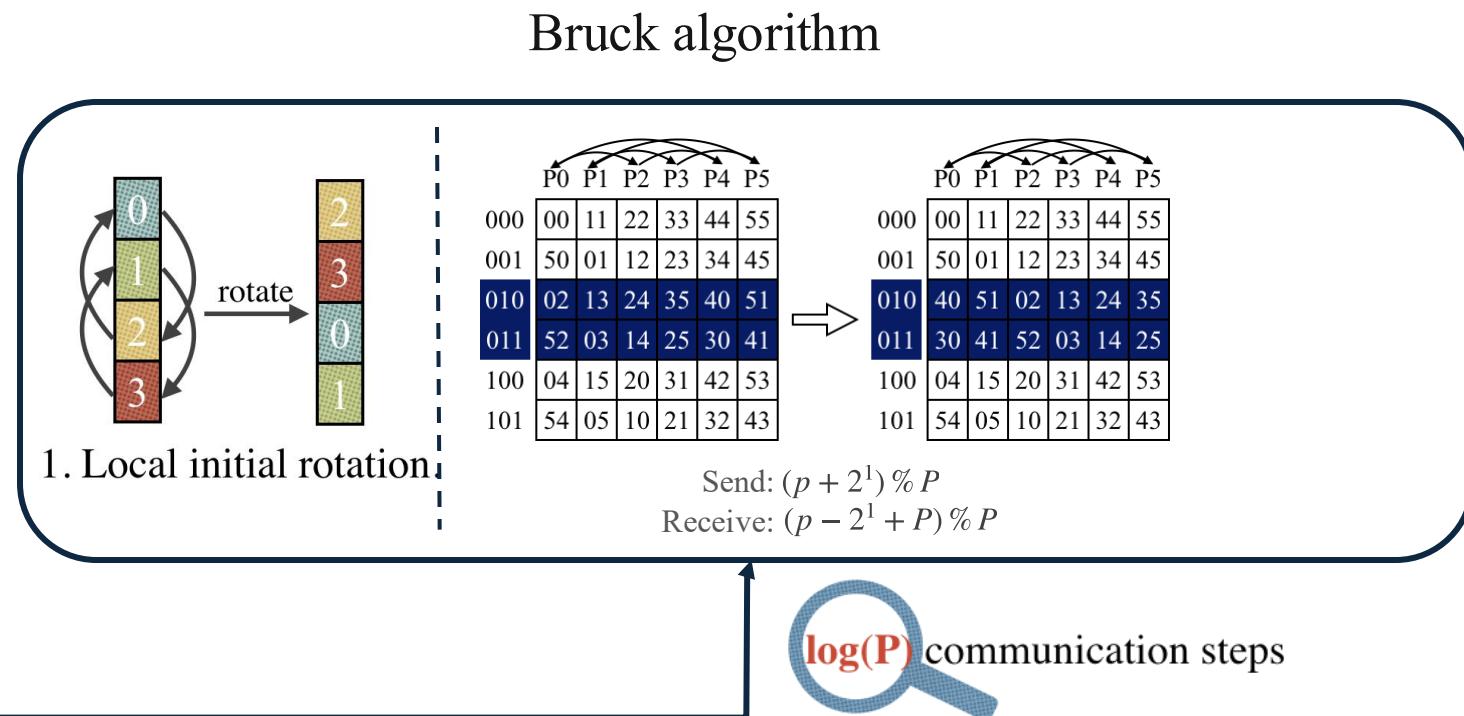
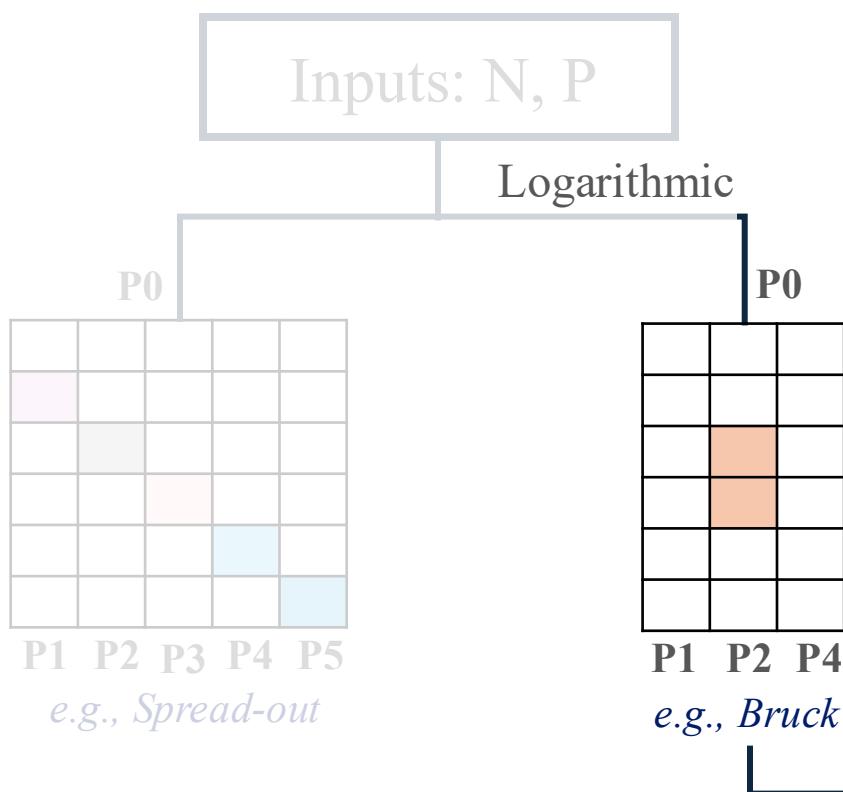
P: Process count



Standard MPI All-to-all Implementations: Comm-1

N: Size of data-block

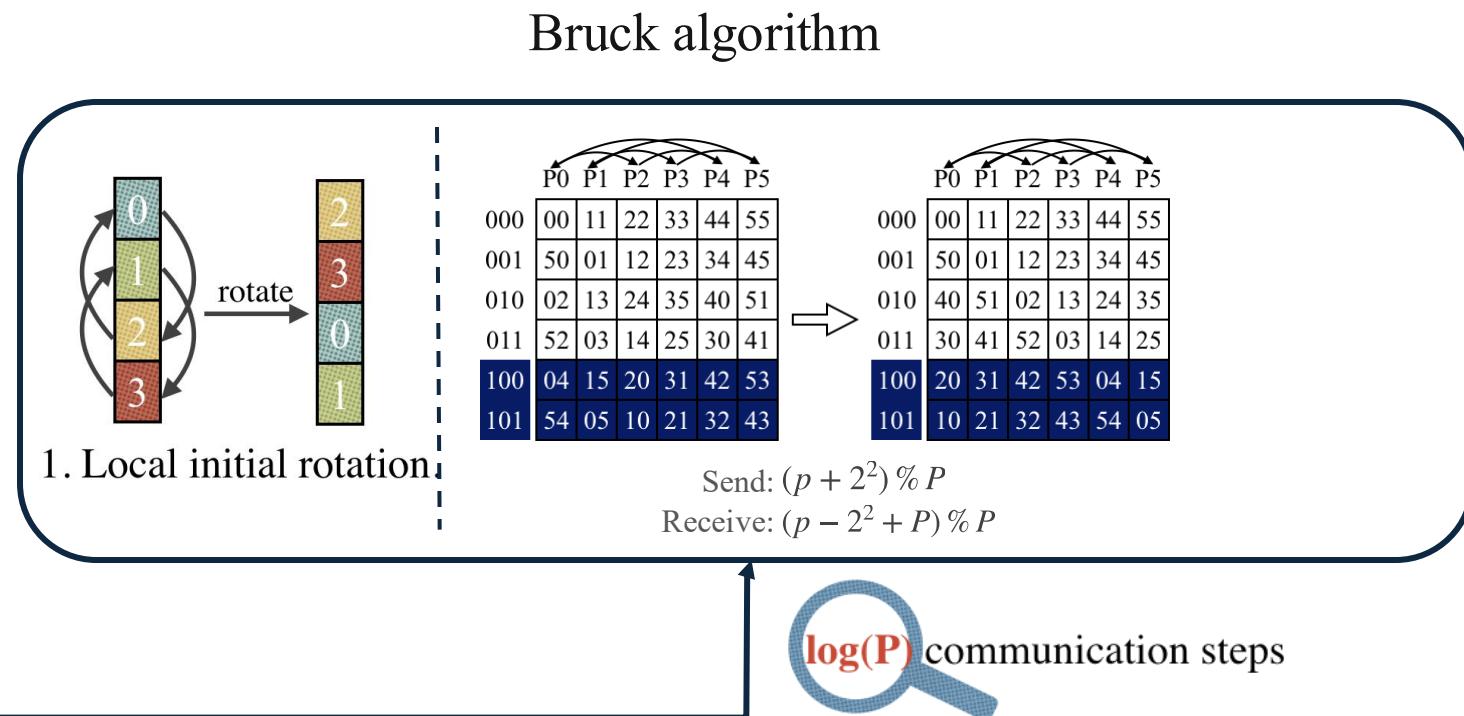
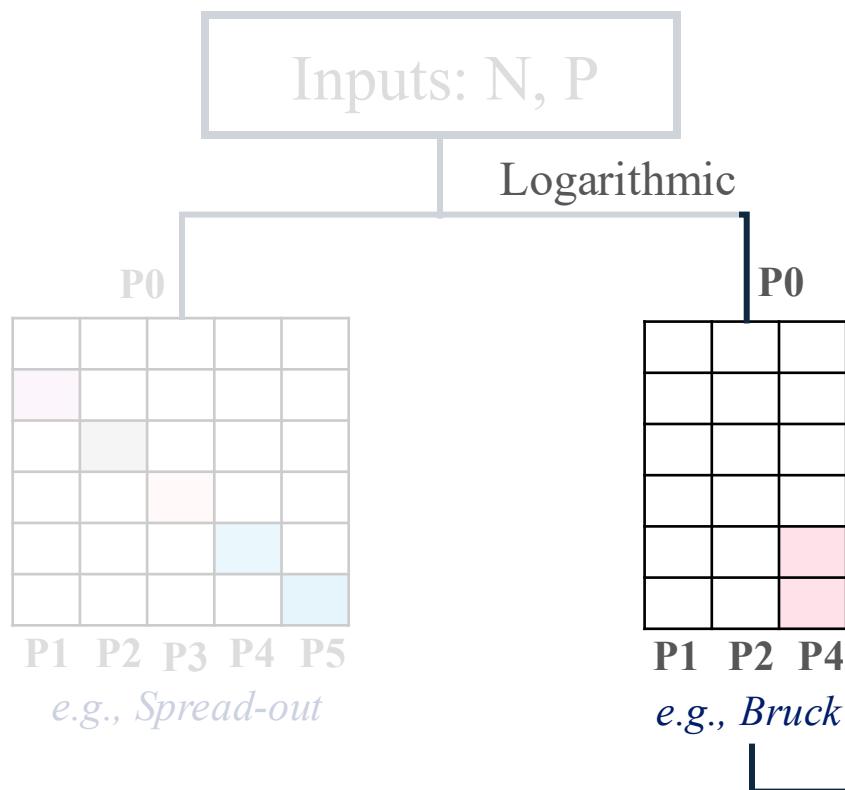
P: Process count



Standard MPI All-to-all Implementations: Comm-2

N: Size of data-block

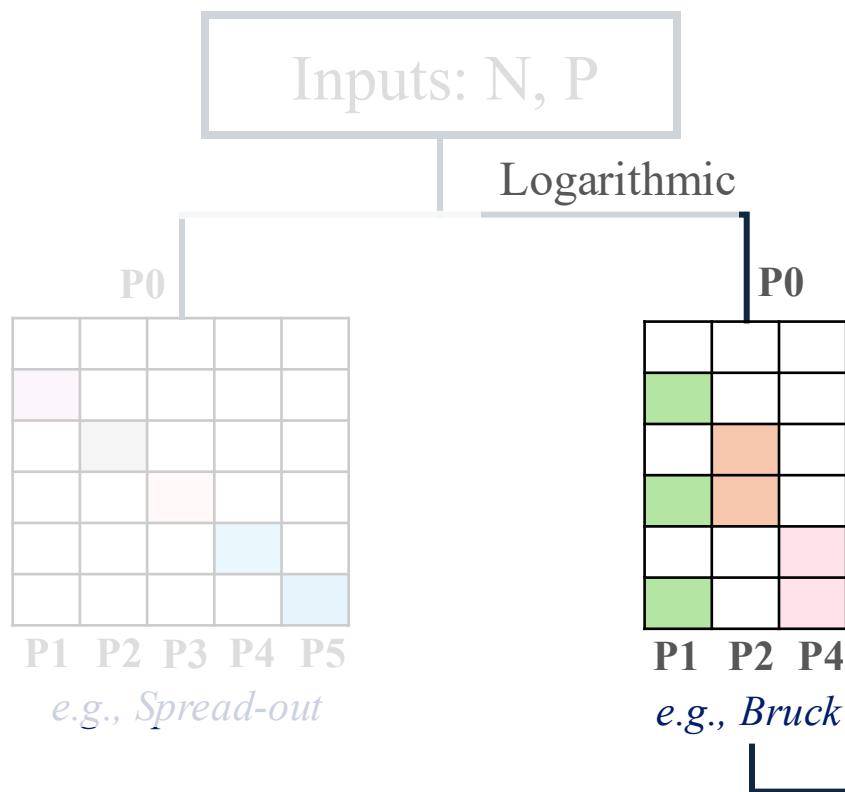
P: Process count



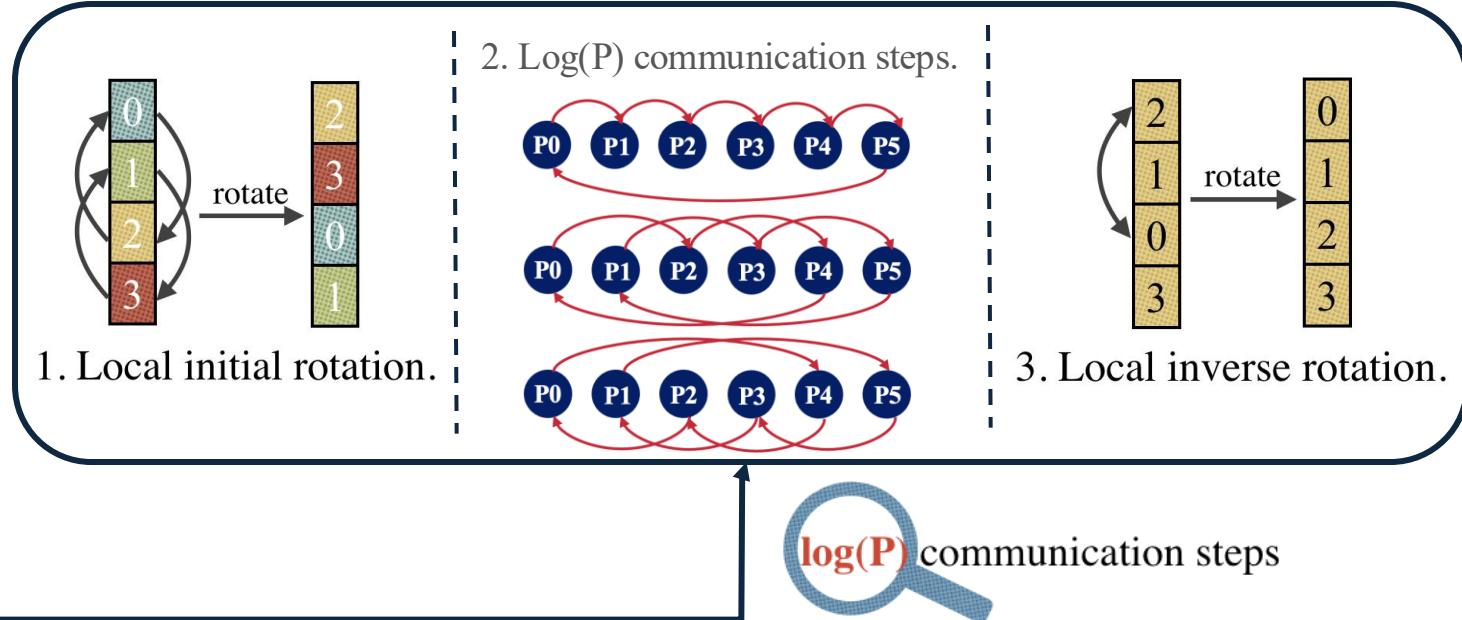
Standard MPI All-to-all Implementations: Bruck

N : Size of data-block

P : Process count



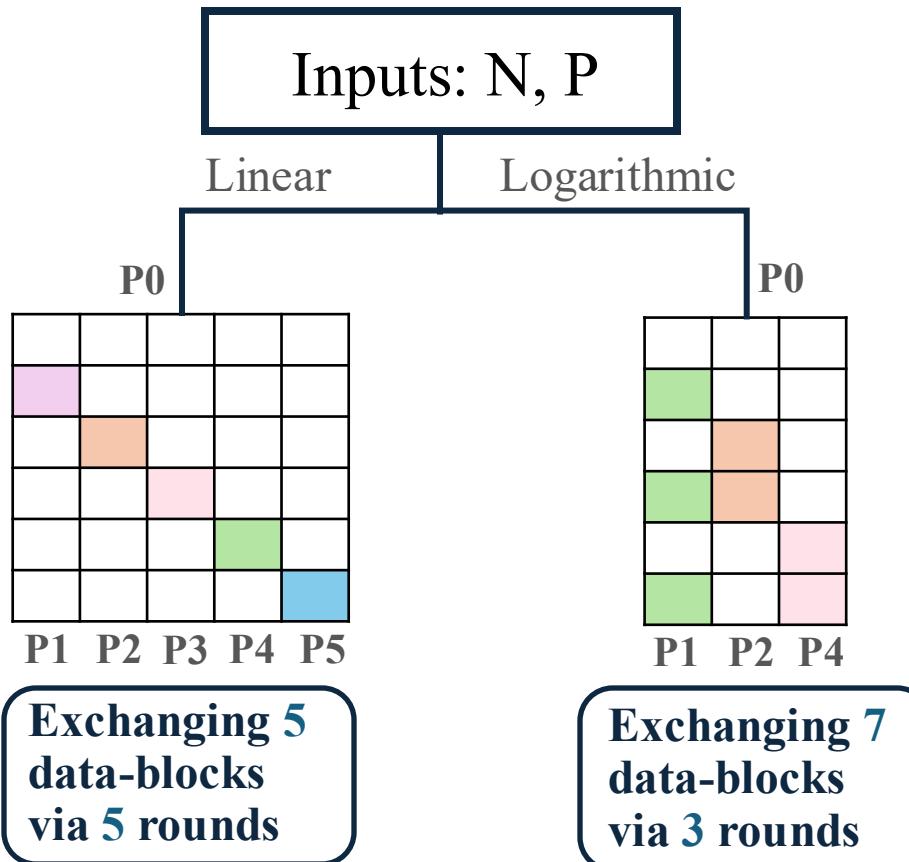
Bruck algorithm



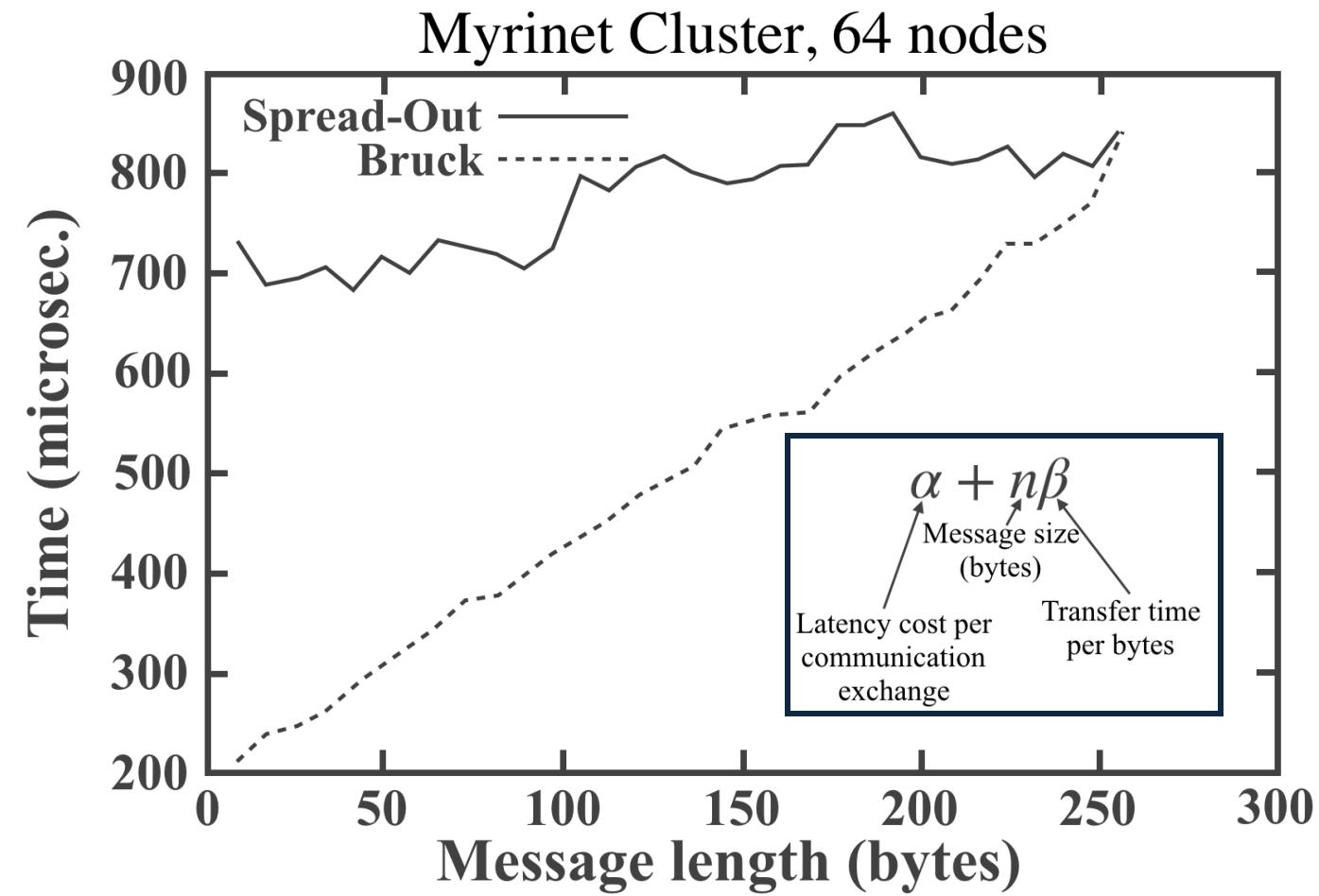
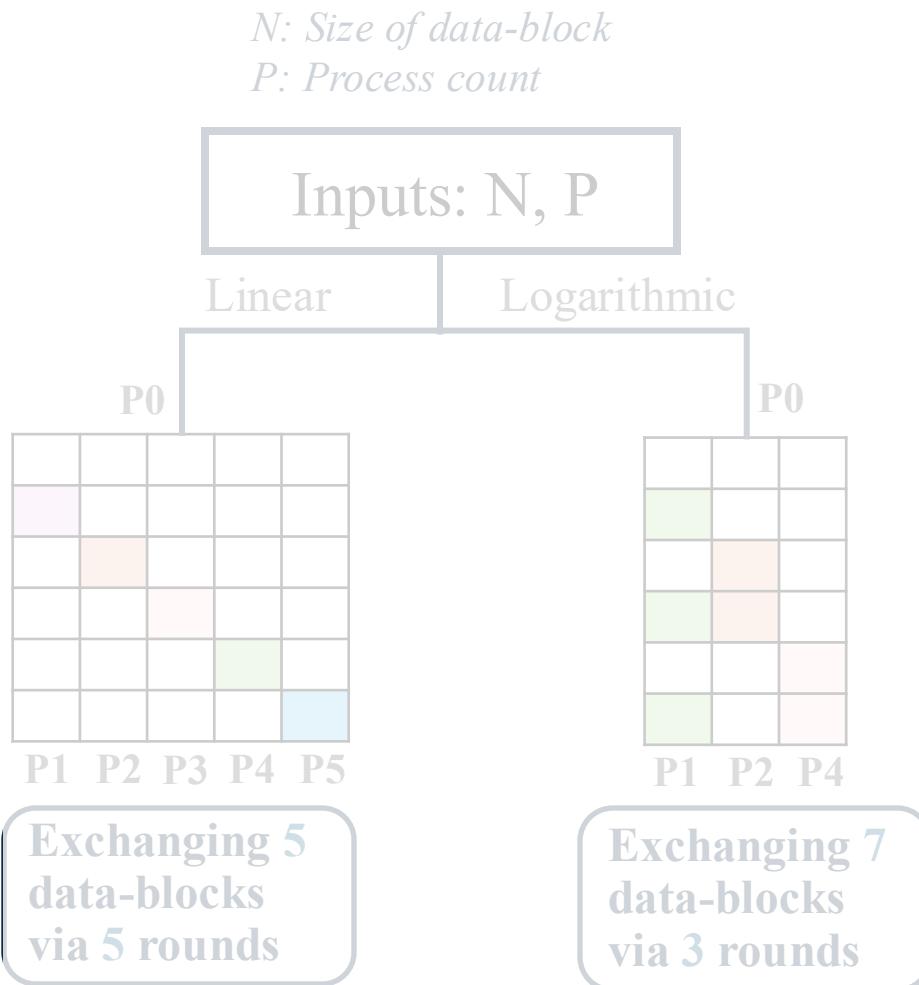
Comparison of These Two Algorithms

N : Size of data-block

P : Process count



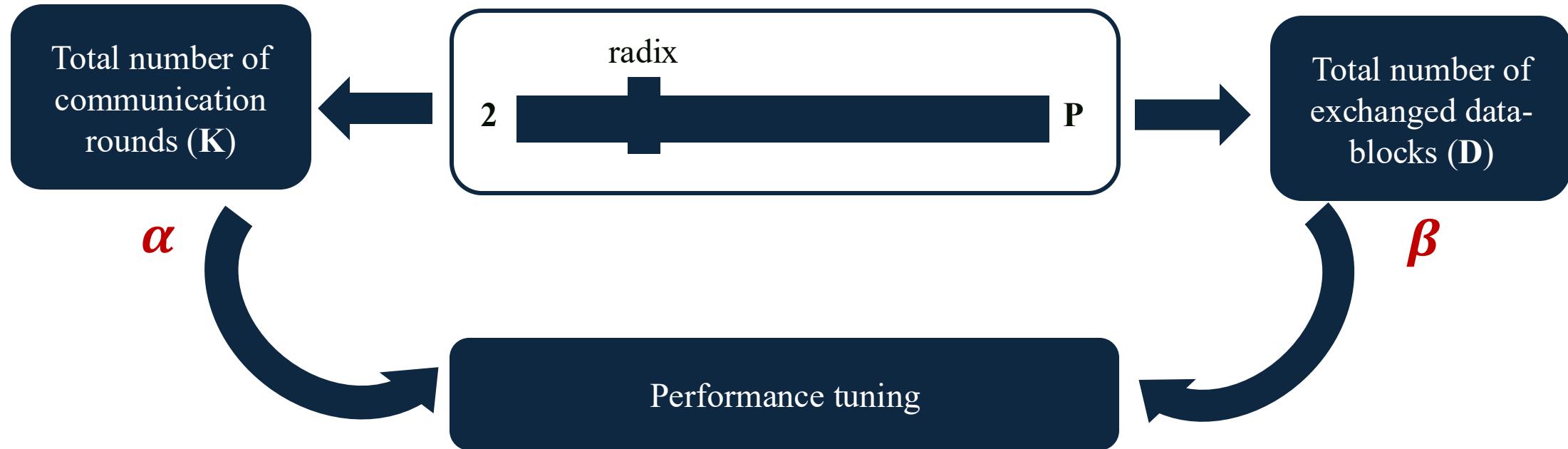
Comparison of These Two Algorithms



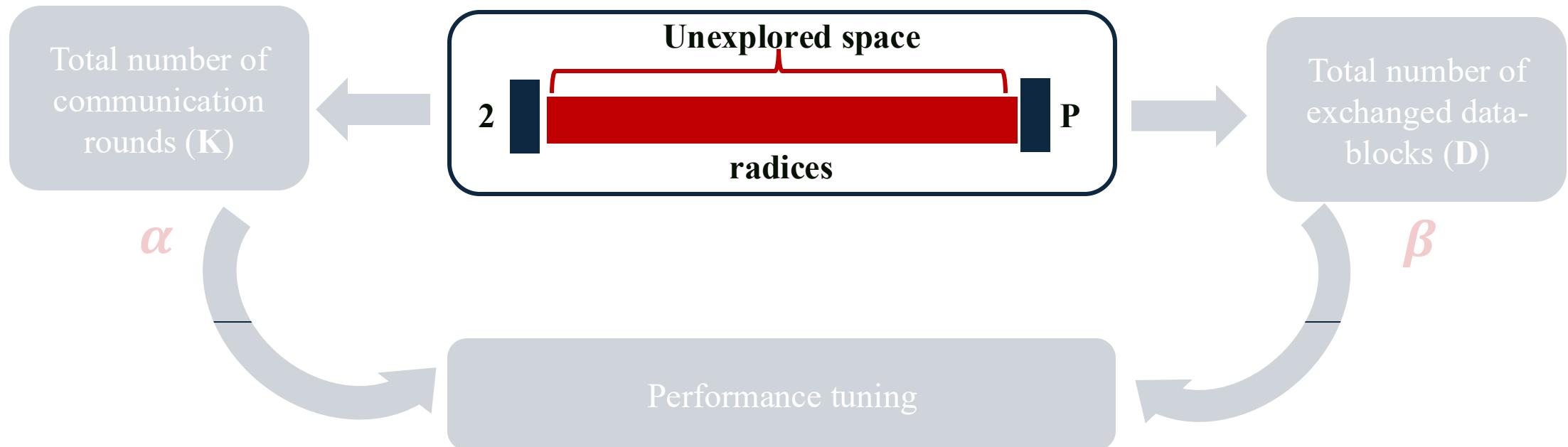
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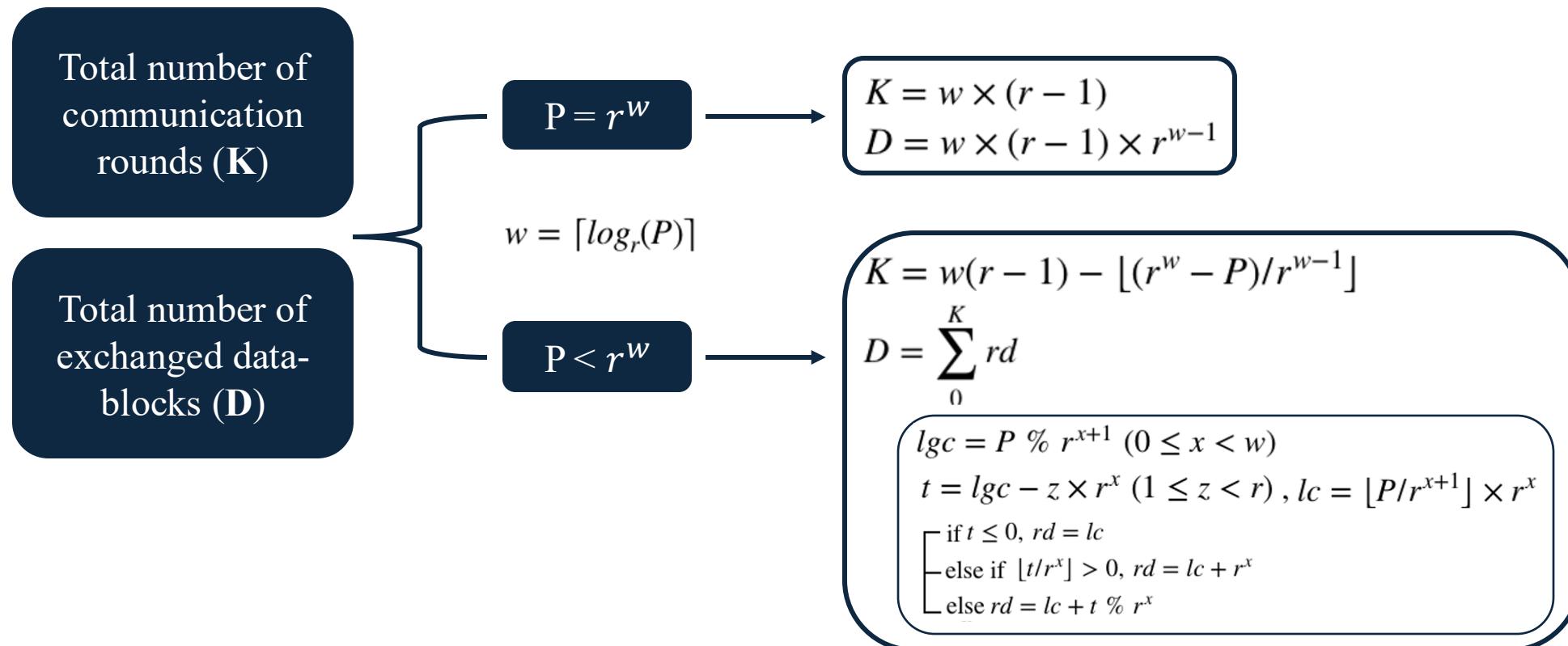
Tuning Radices Enables Performance Tuning



Standard MPI All-to-all Implementation Leaves a Large Unexplored Space



Mathematics of Tunable Radix All-to-all



Tunable Radix All-to-all Example

$$P = 8, r = 3$$

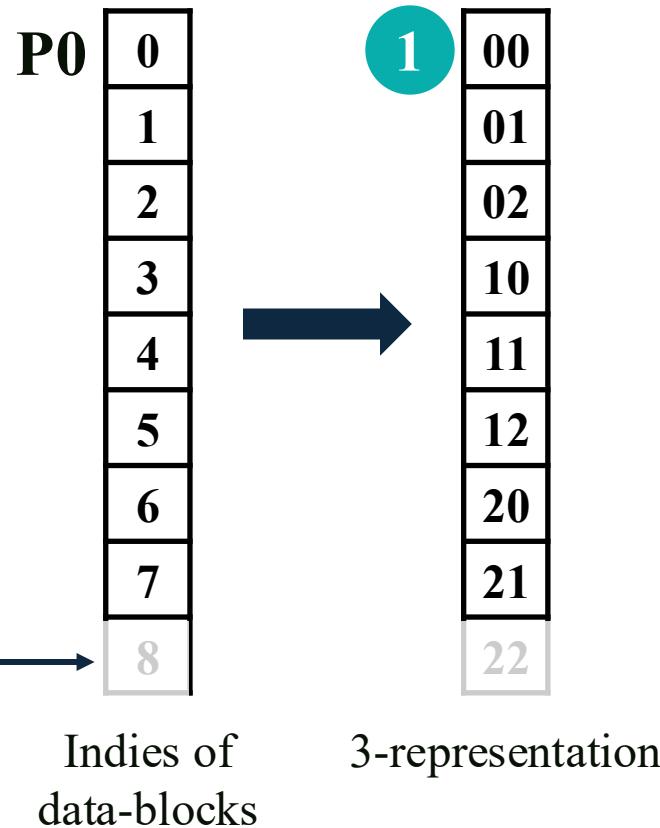
P0
0
1
2
3
4
5
6
7
8

if $\log_r(P)$
is integer

Indies of
data-blocks

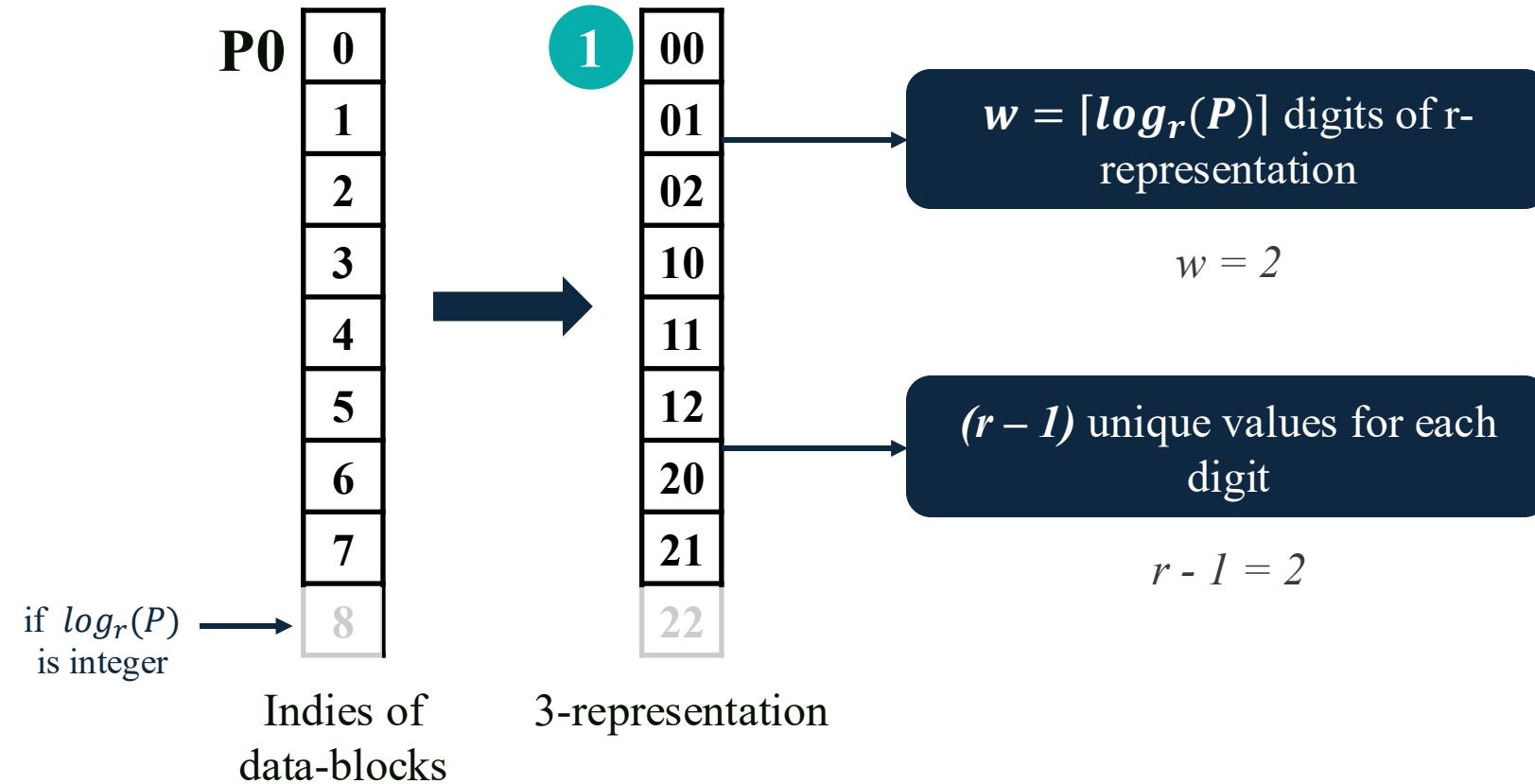
Tunable Radix All-to-all Example: r-representation

$$P = 8, r = 3$$



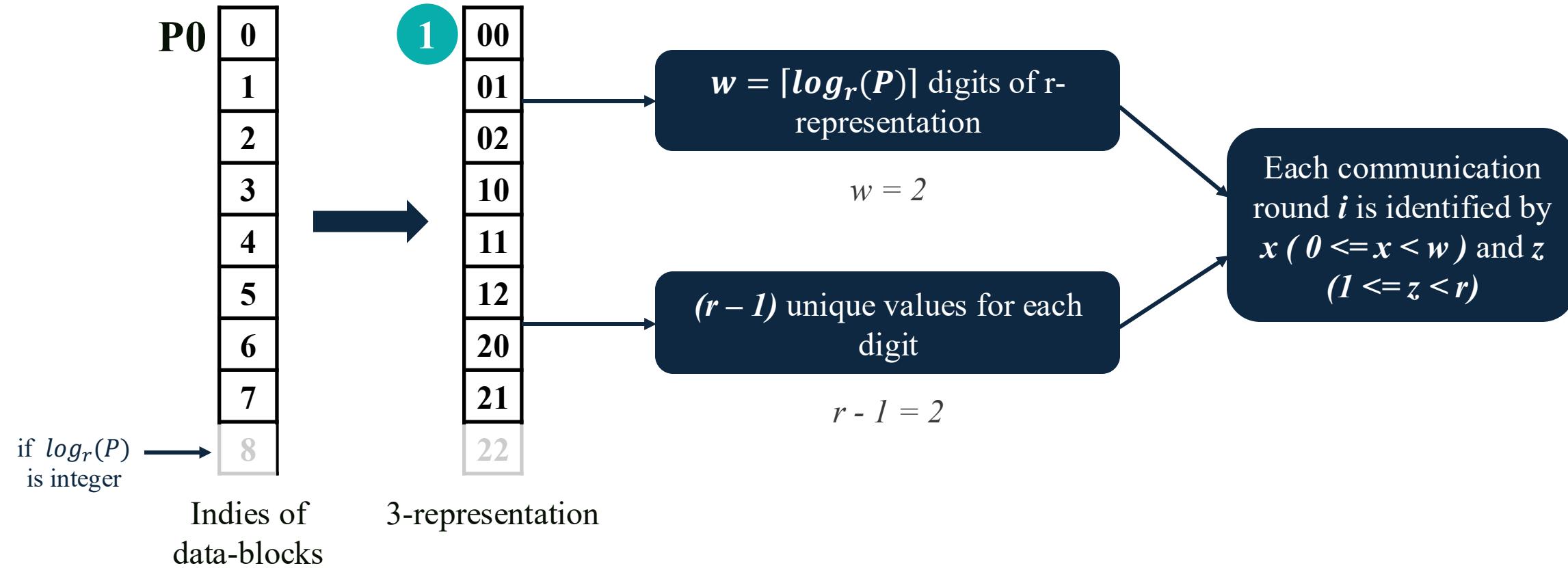
Tunable Radix All-to-all Example: r-representation

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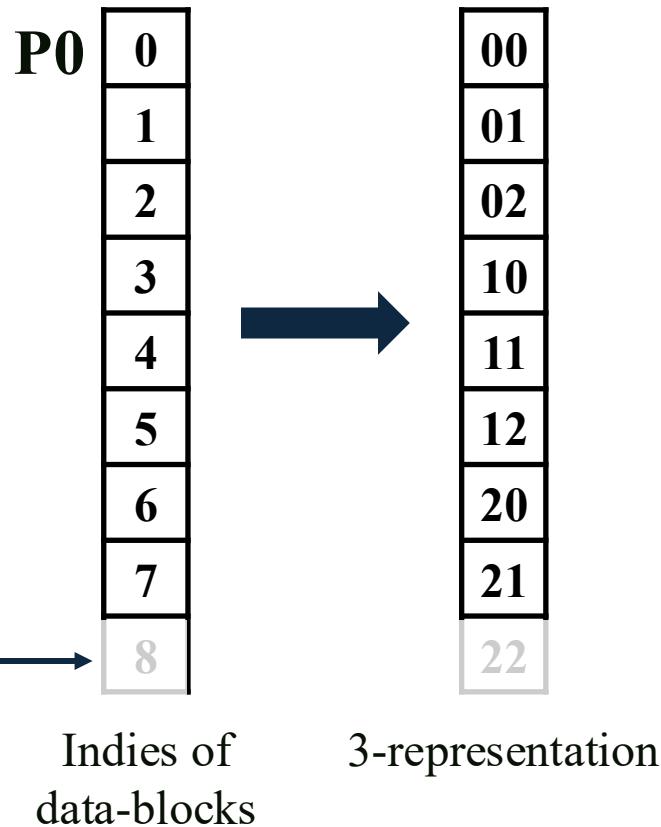
Tunable Radix All-to-all Example: r-representation

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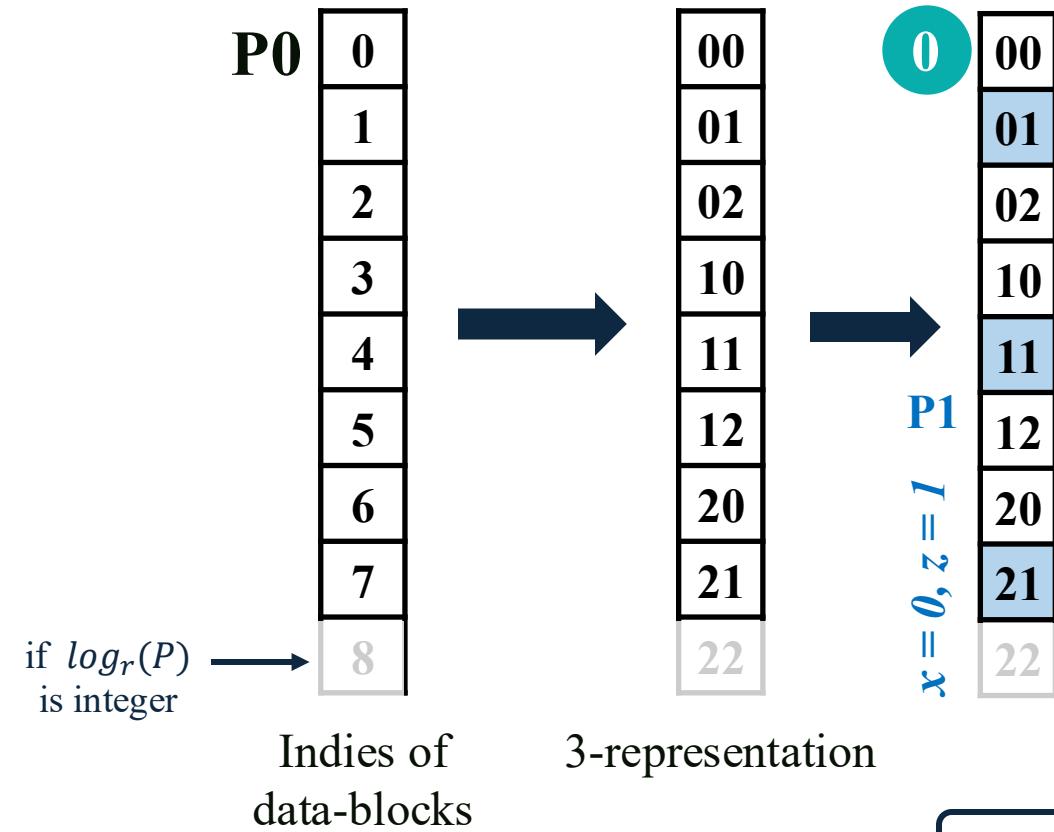
Tunable Radix All-to-all Example: Communication

$$P = 8, r = 3$$



Tunable Radix All-to-all Example: Comm-round-0

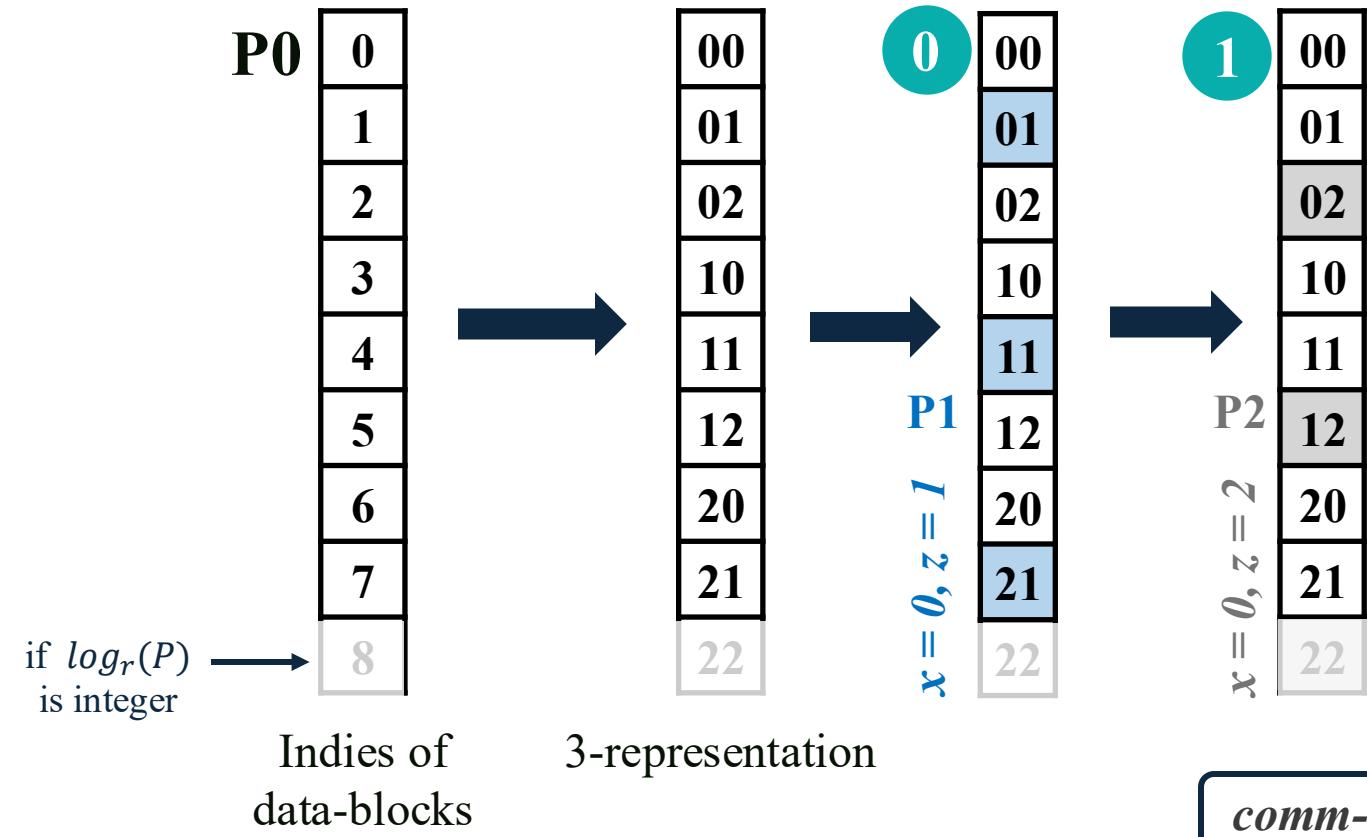
$$P = 8, r = 3$$



comm-round-0: $x = 0, z = 1, P_0$ sends to P_1

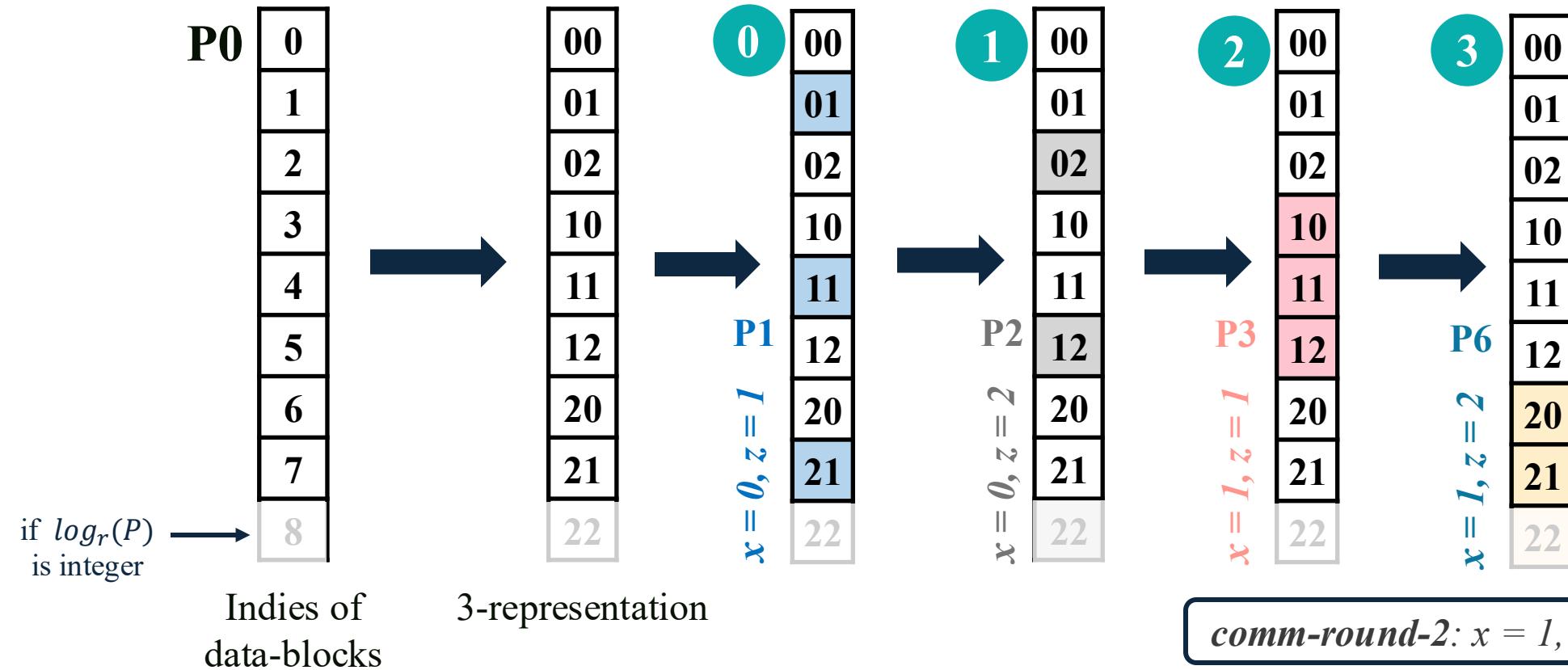
Tunable Radix All-to-all Example: Comm-round-1

$$P = 8, r = 3$$



Tunable Radix All-to-all Example: Comm-round-2,3

$$P = 8, r = 3$$



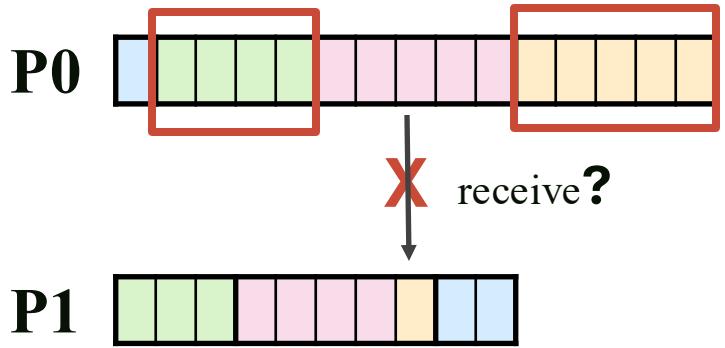
comm-round-2: $x = 1, z = 1$, P0 sends to P3

comm-round-3: $x = 1, z = 2$, P0 sends to P6

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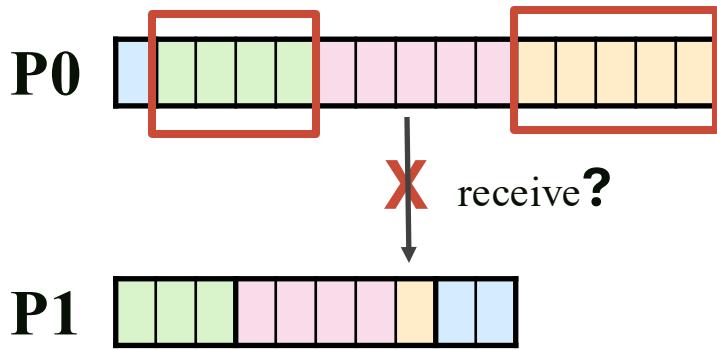
Challenges of Applying the Logarithmic Algorithm to Non-uniform All-to-all Data Exchange



Each process is unaware of how much data to expect during each intermediate communication round.

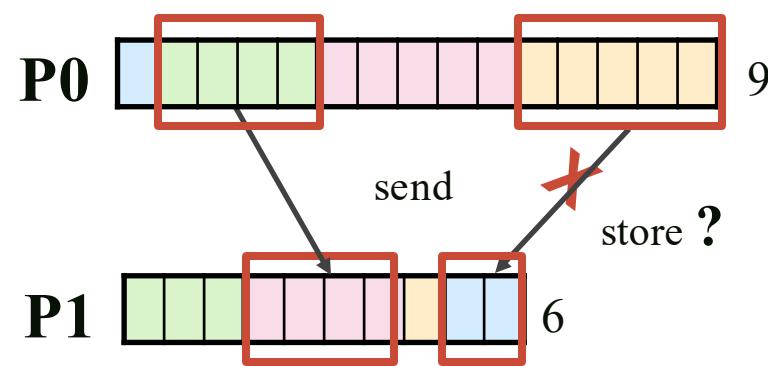
Non-uniform
workloads

Challenges of Applying the Logarithmic Algorithm to Non-uniform All-to-all Data Exchange



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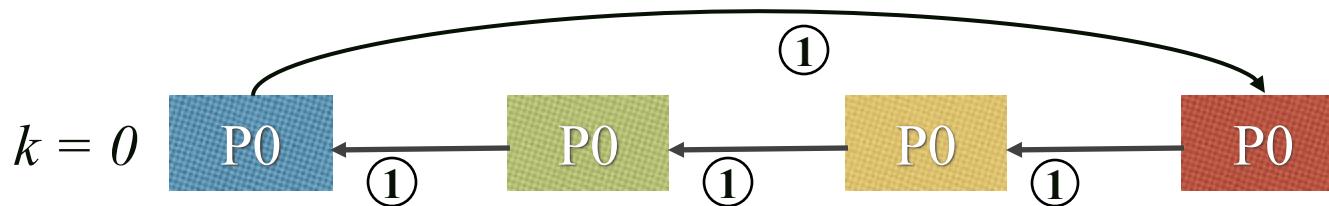
Non-uniform workloads



The received data elements could be too large to fit into the segment in the send buffer.

Store-and-forward

Two-phase Communication Scheme

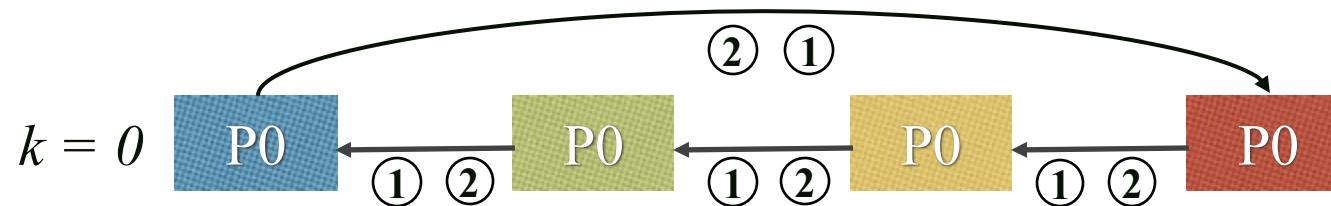


① *Metadata Exchange*

a	b
---	---

 Size of each block

Two-phase Communication Scheme



① *Metadata Exchange*

a	b
---	---

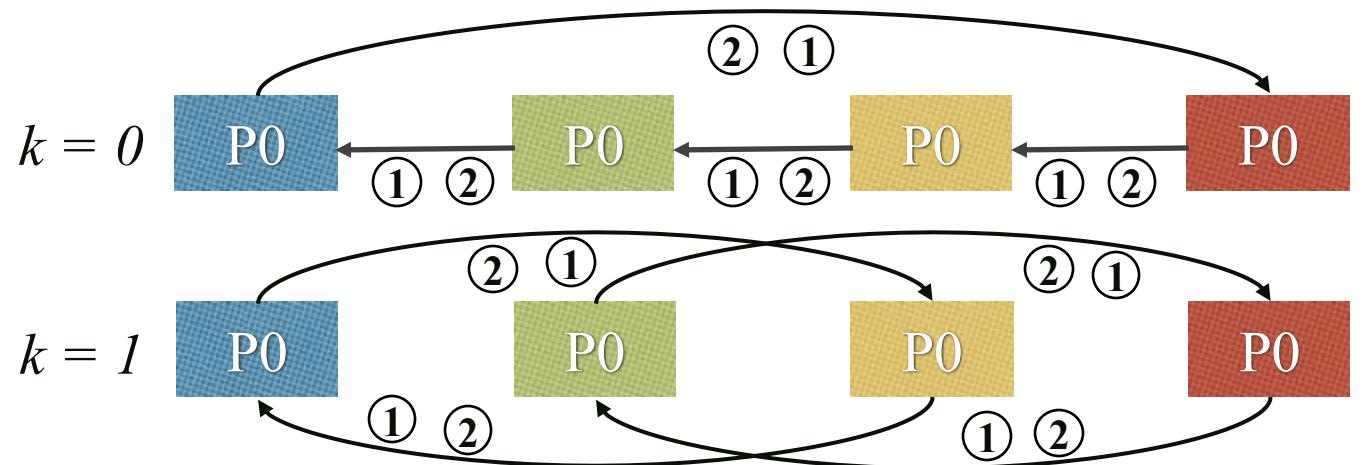
 Size of each block

② *Data Exchange*

Block 1	Block 3
---------	---------

 Actual data-blocks

Two-phase Communication Scheme



① *Metadata Exchange*

a	b
---	---

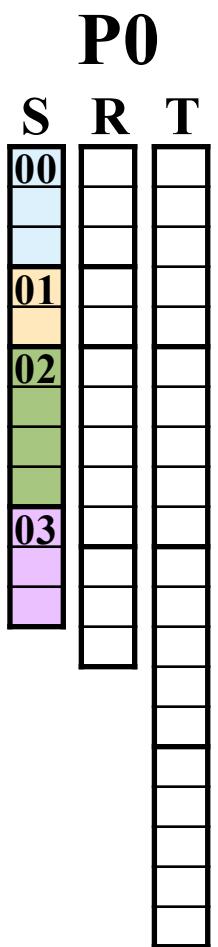
 Size of each block

② *Data Exchange*

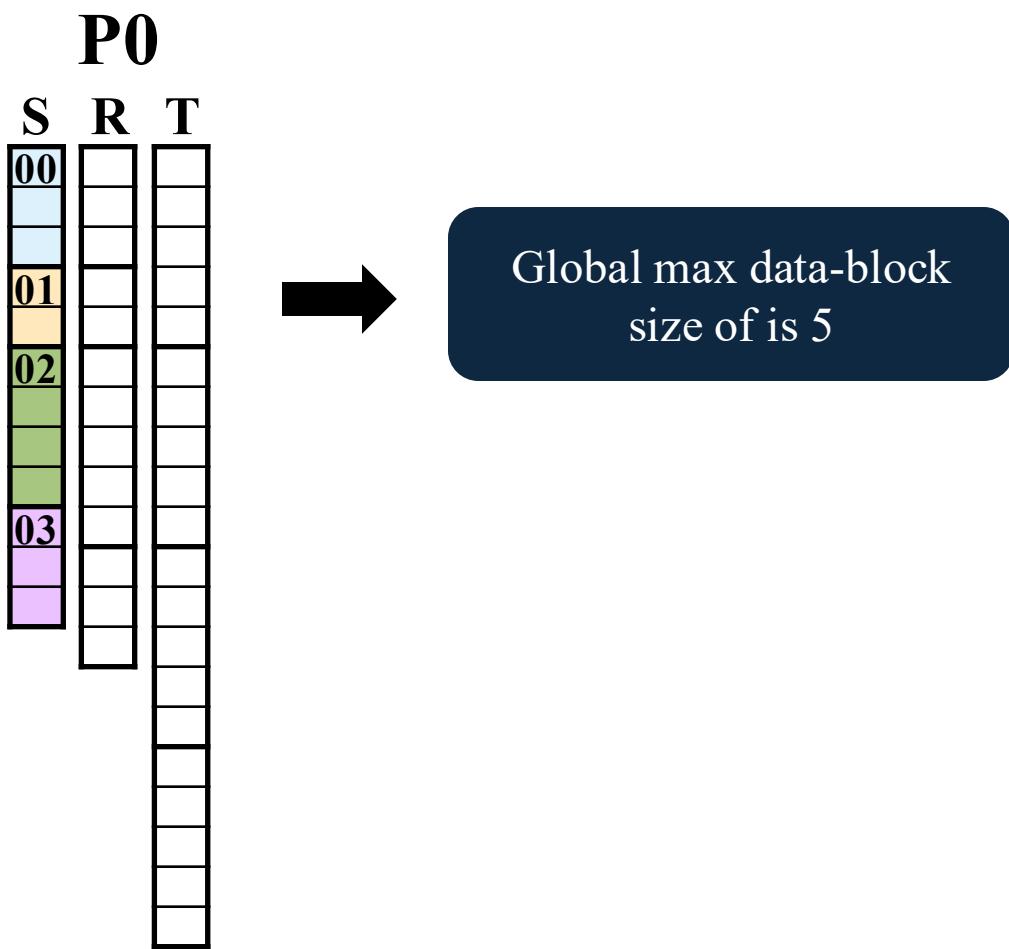
Block 1	Block 3
---------	---------

 Actual data-blocks

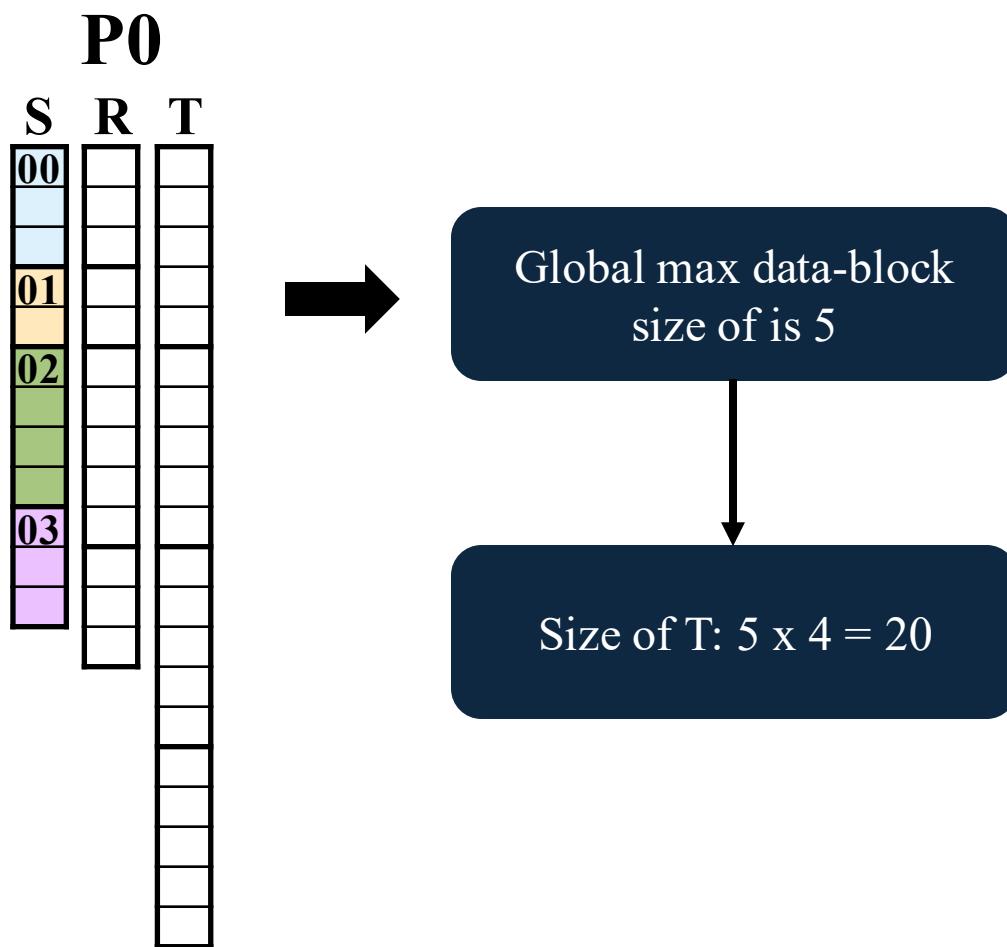
Using Temporary Buffer to Solve Challenge 2



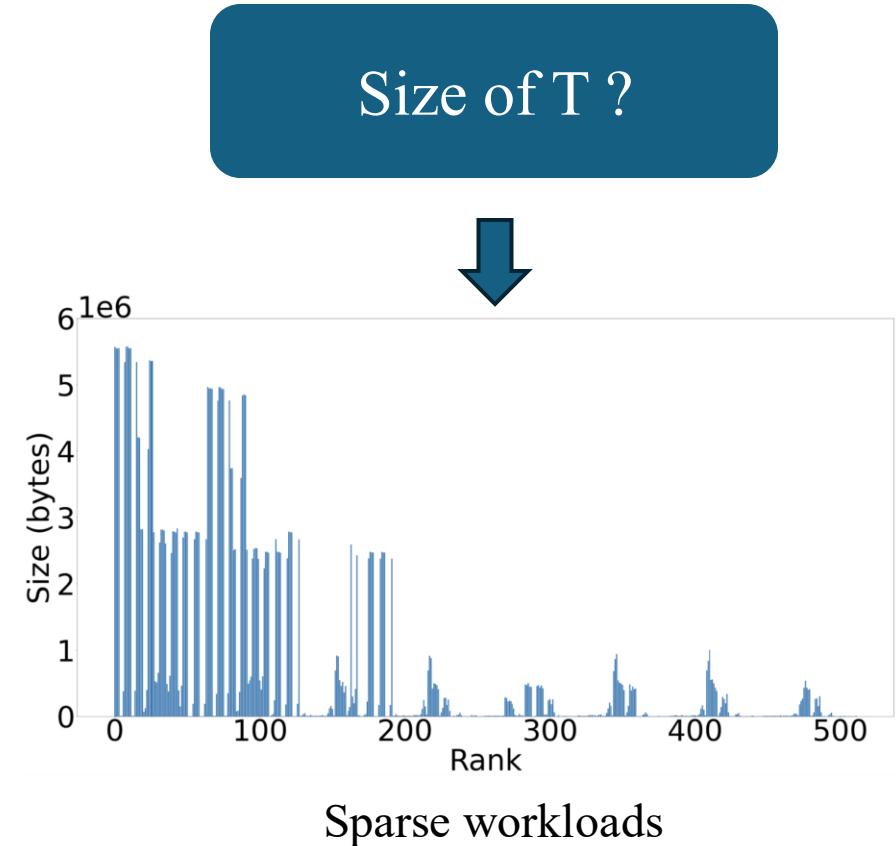
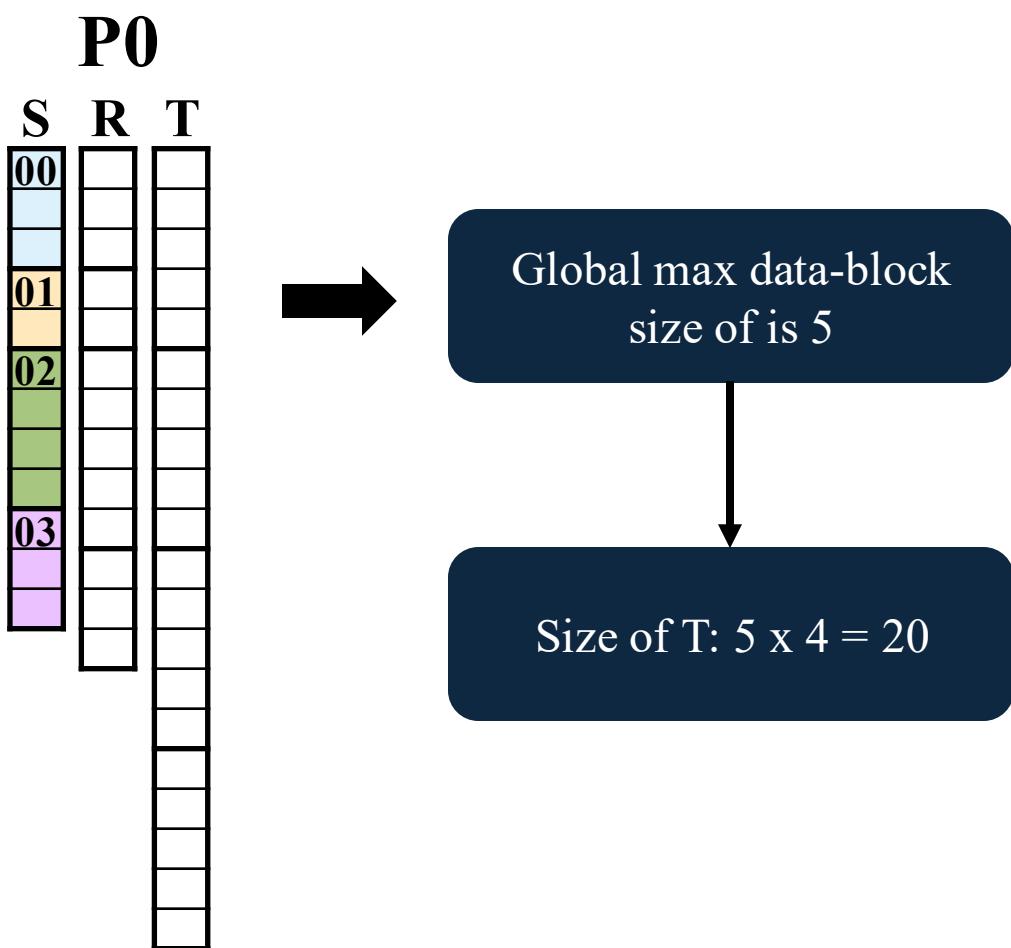
Estimating Temporary Buffer Size



Estimating Temporary Buffer Size

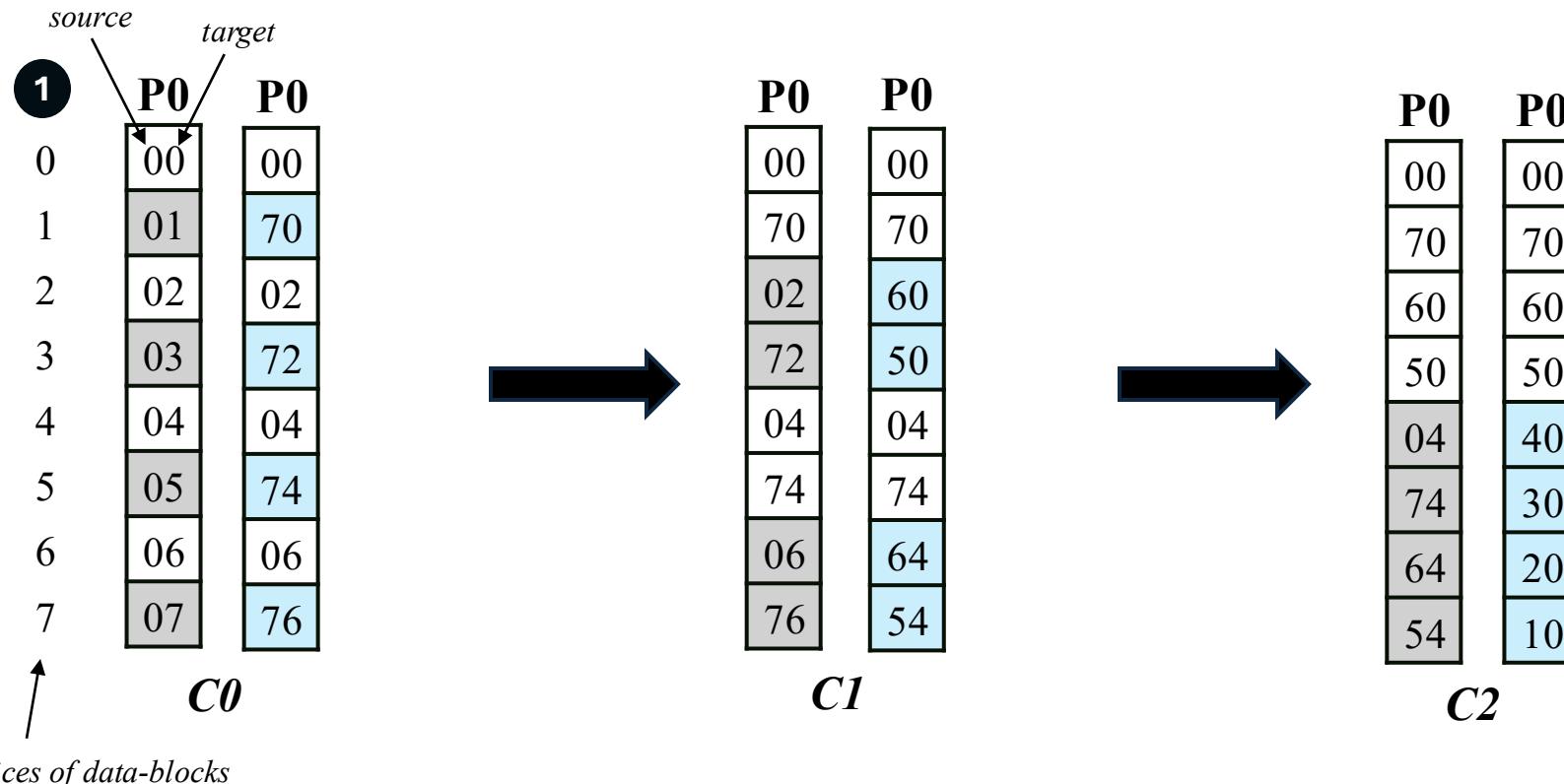


Temporary Buffer Size Can Be Very Large



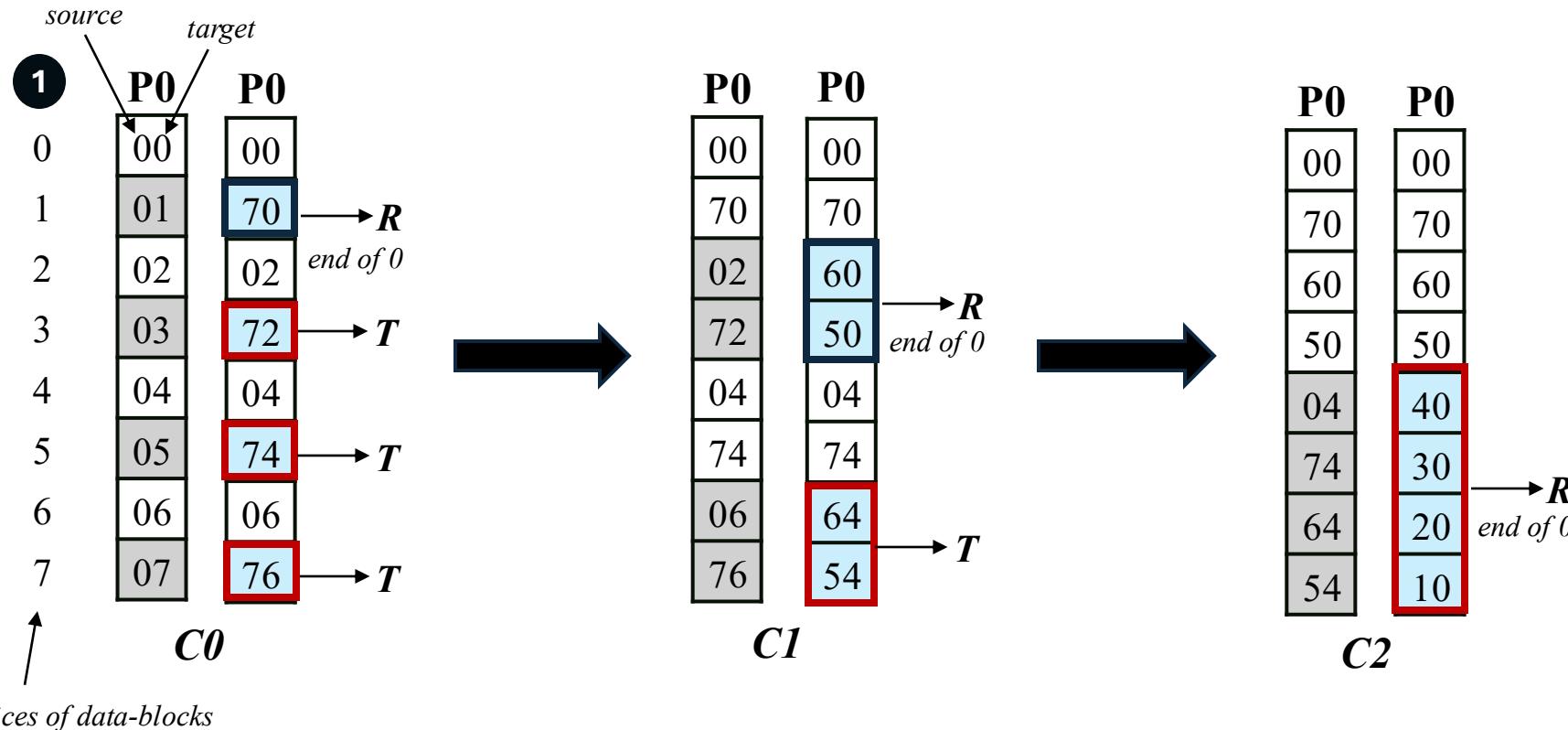
Temporary Buffer Size Analysis

$$P = 8, r = 2$$



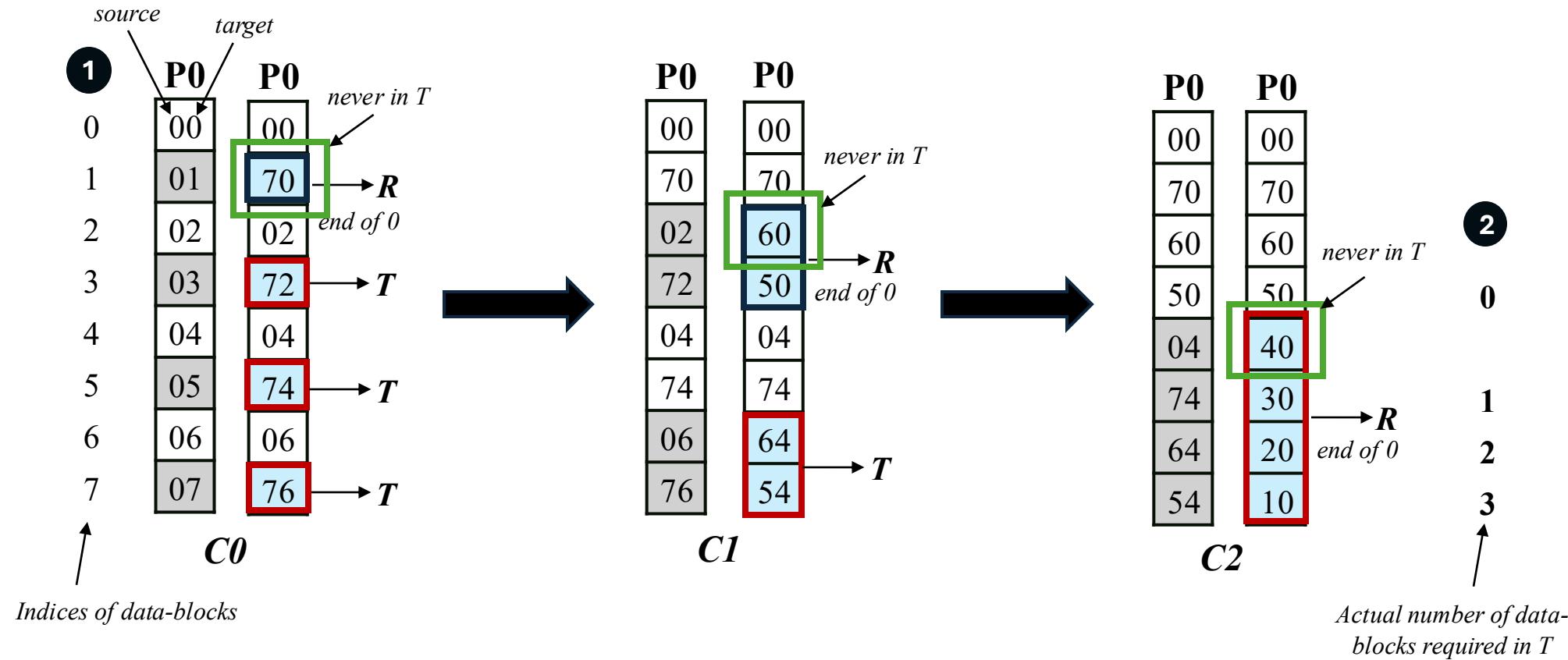
Temporary Buffer Size Analysis

$$P = 8, r = 2$$

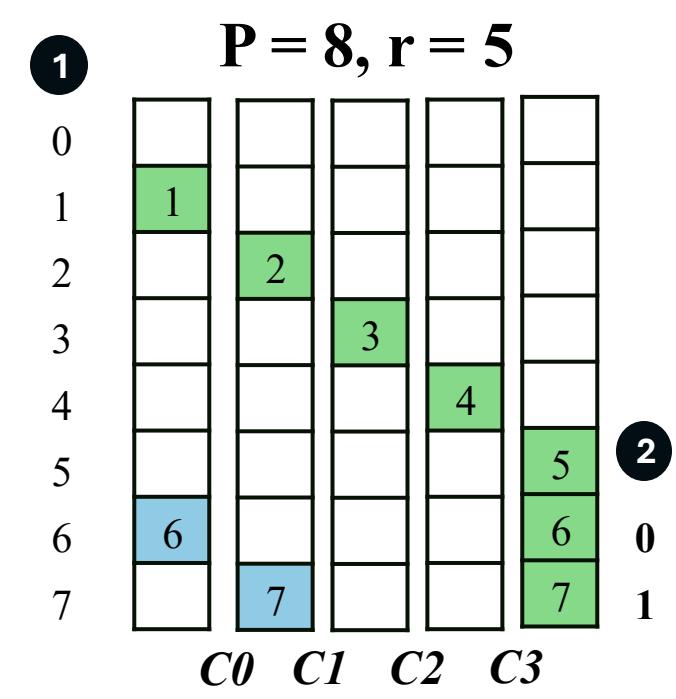
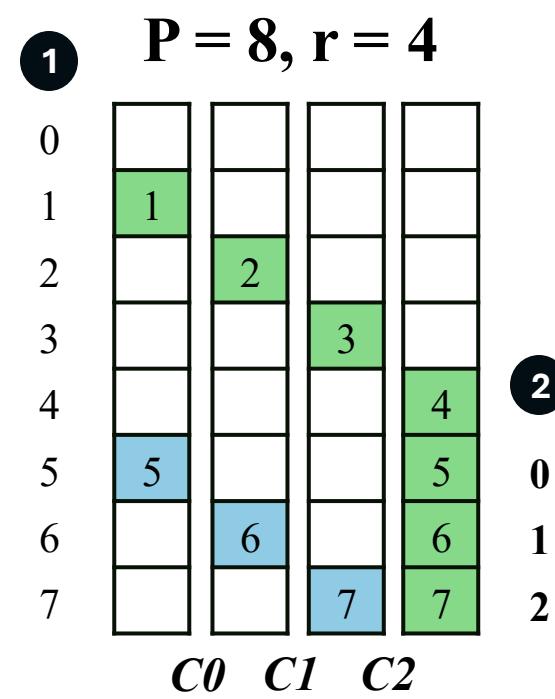
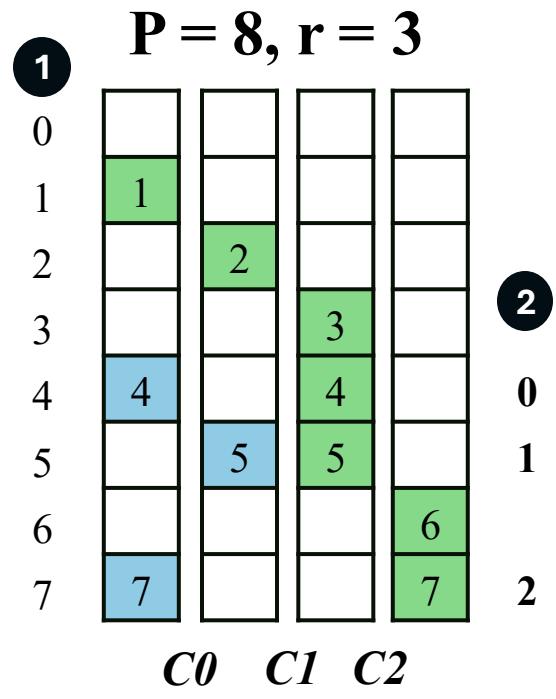


Temporary Buffer Size Analysis

$$P = 8, r = 2$$



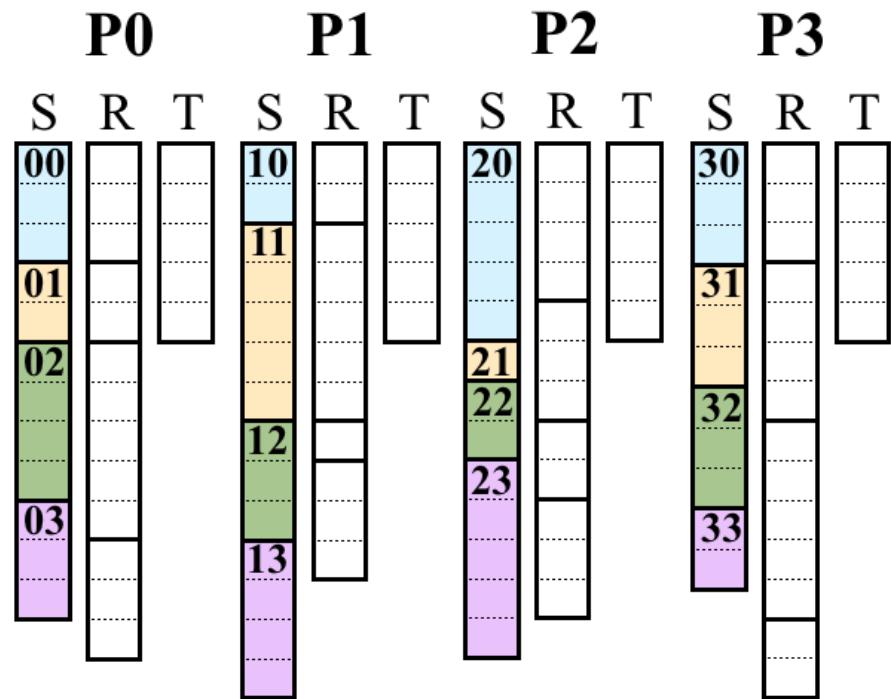
Temporary Buffer Size Analysis



Agenda

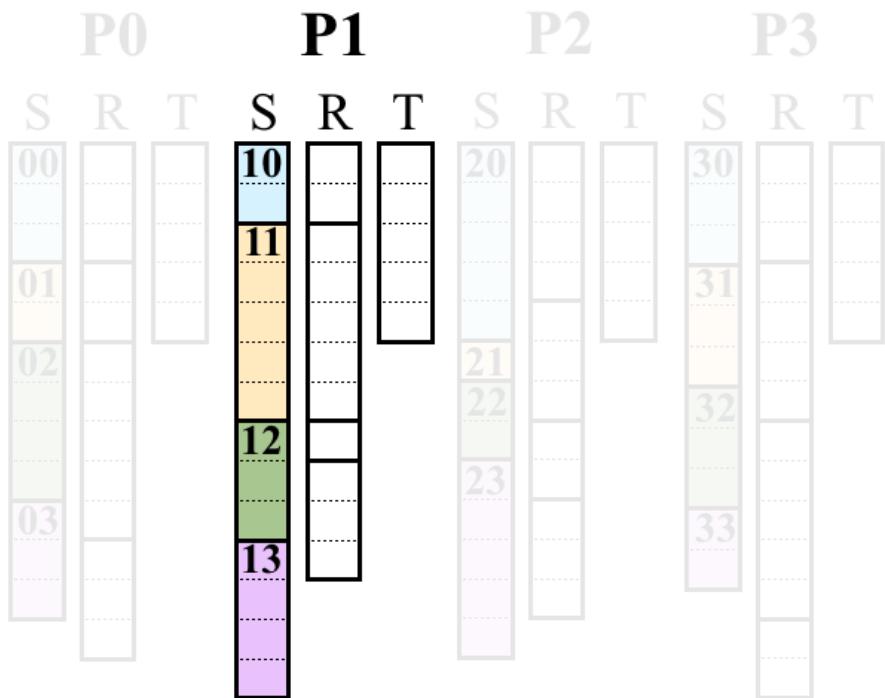
- Introduction and Background
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- **Parameterized Logarithmic Algorithm**
- Parameterized Linear Algorithm
- Evaluation
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- Conclusion

Example of the ParLogNa with $P = 4$ and $r = 2$.



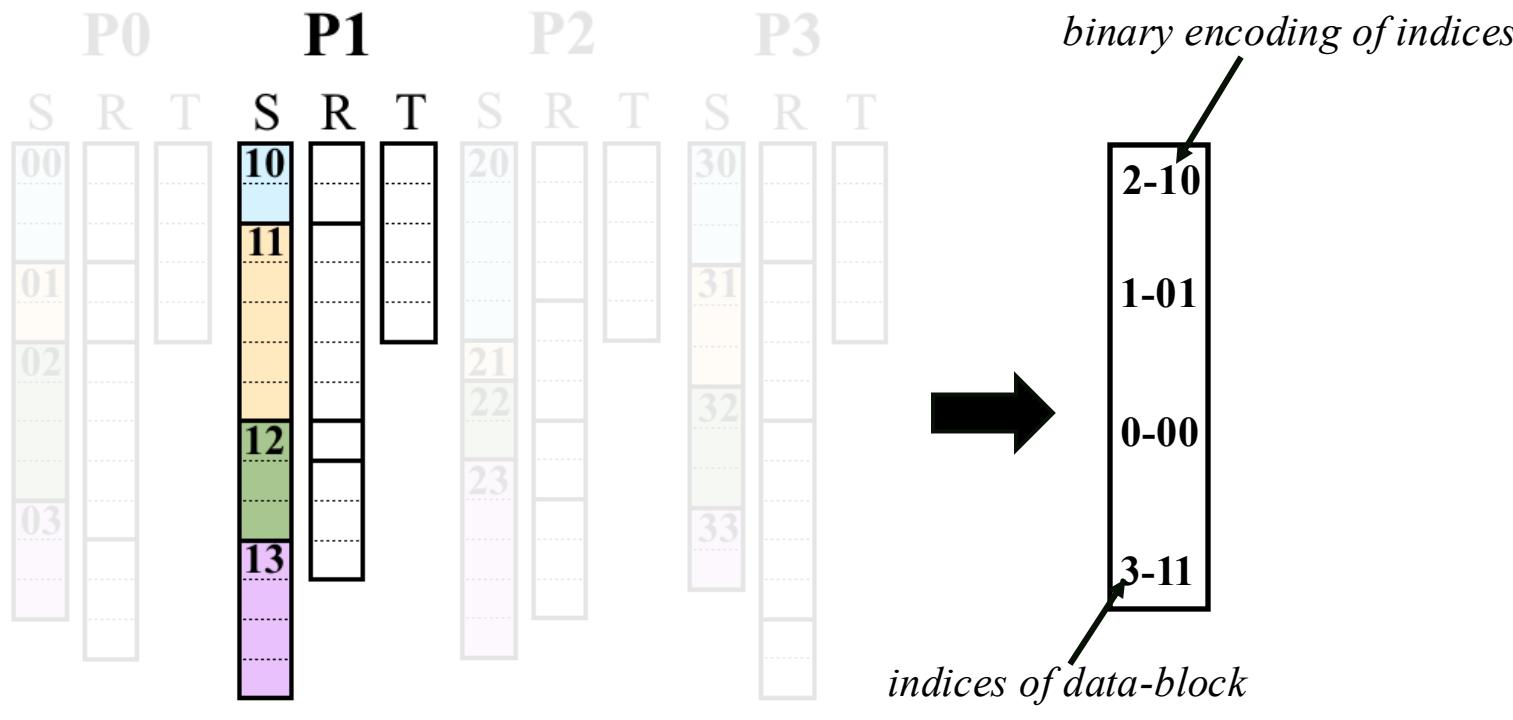
Initial data for 4 process

Example of the ParLogNa with $P = 4$ and $r = 2$.



Taking P1 for example

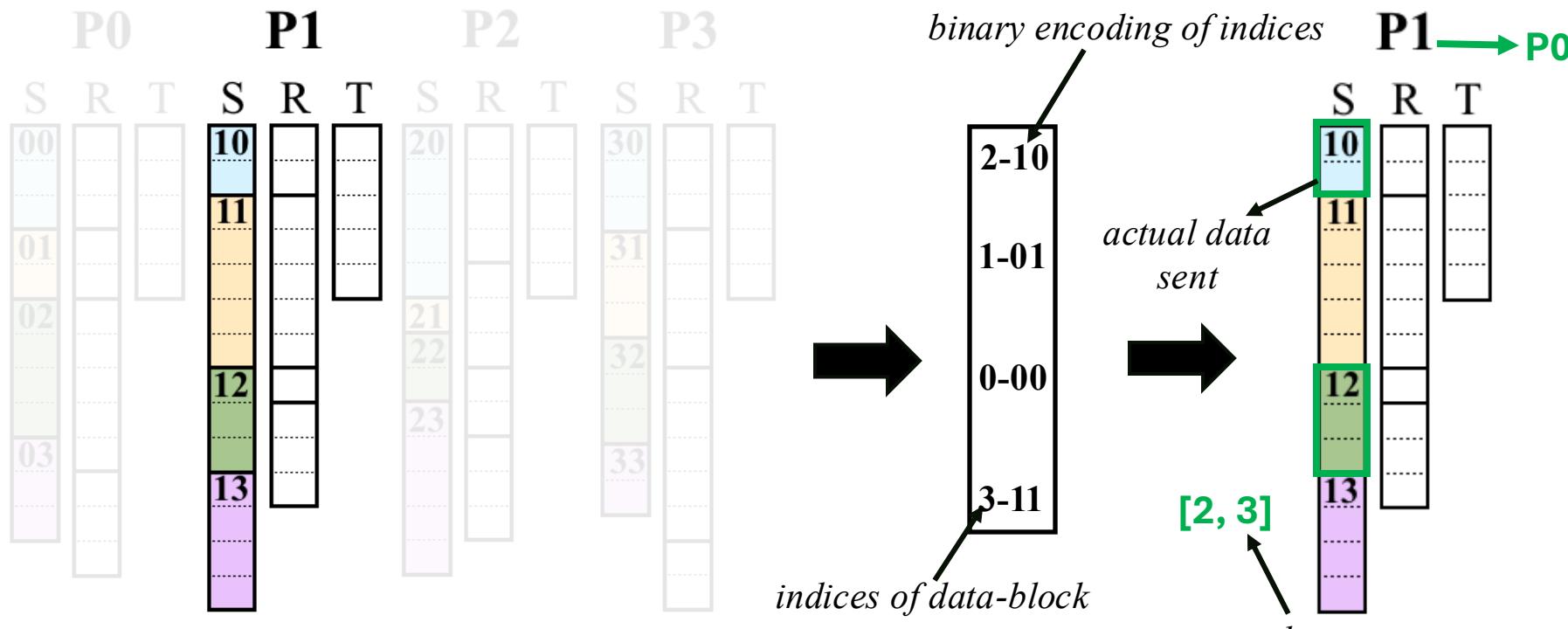
Example of the ParLogNa: Rotation Array and Binary Encoding



Taking P1 for example

Rotation array

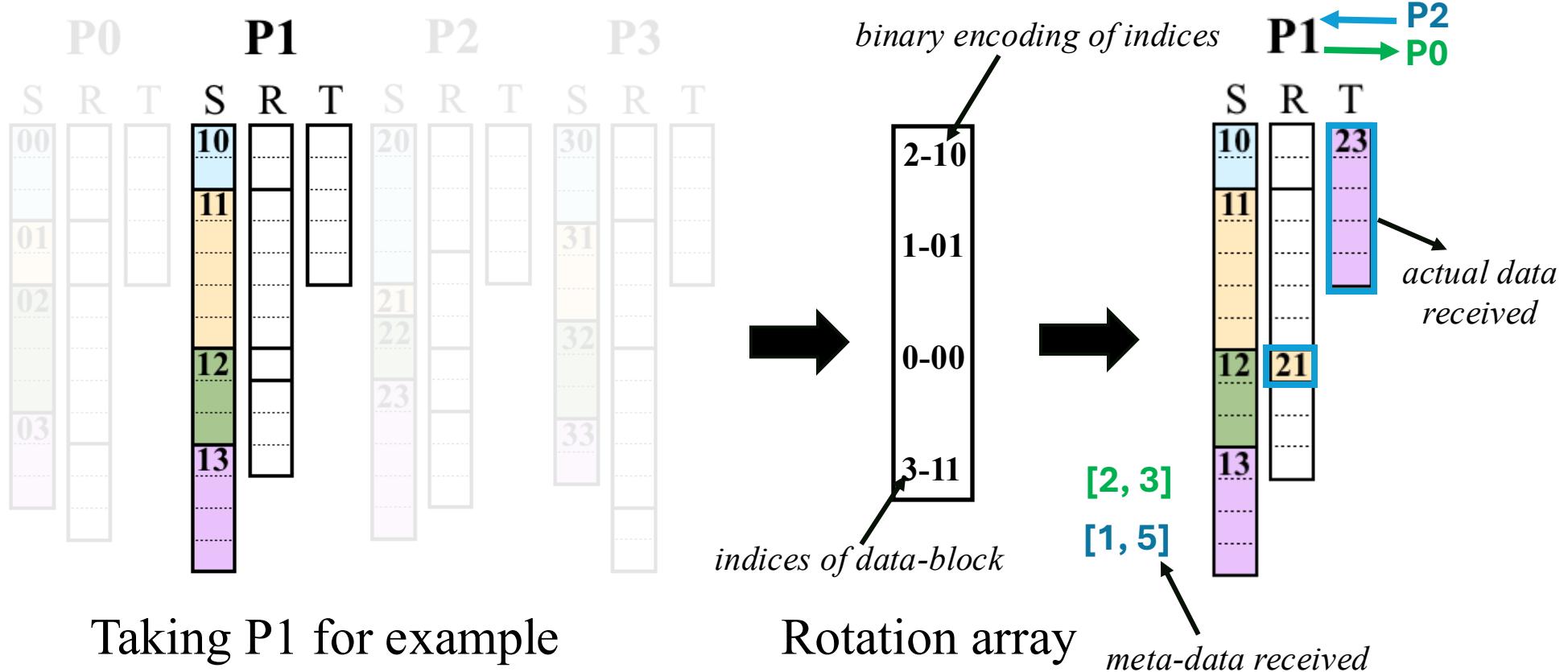
Example of the ParLogNa: Comm Round 0



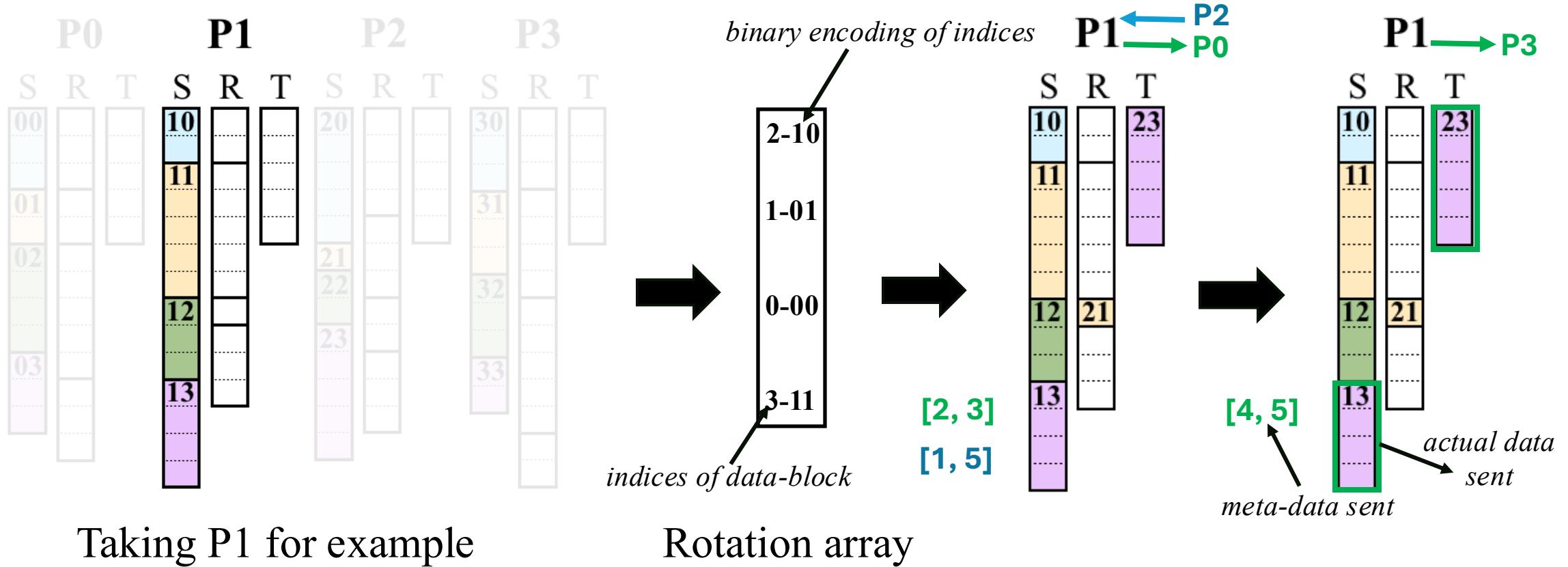
Taking P1 for example

Rotation array

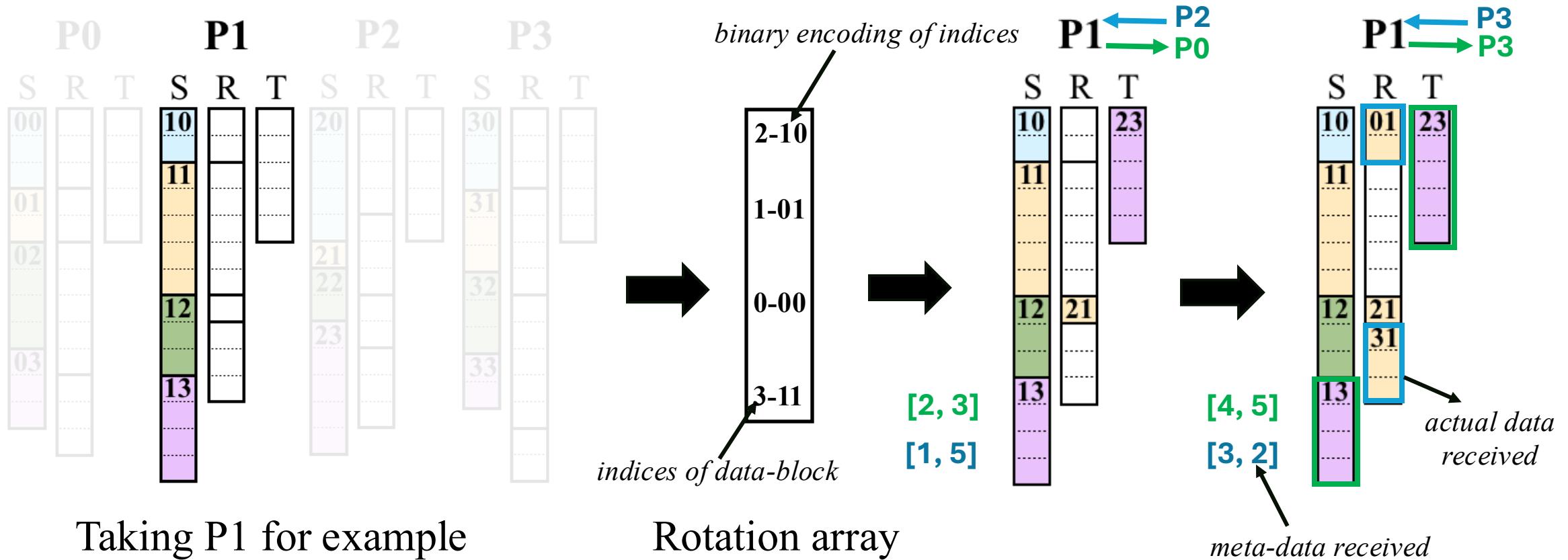
Example of the ParLogNa: Comm Round 0



Example of the ParLogNa: Comm Round 1



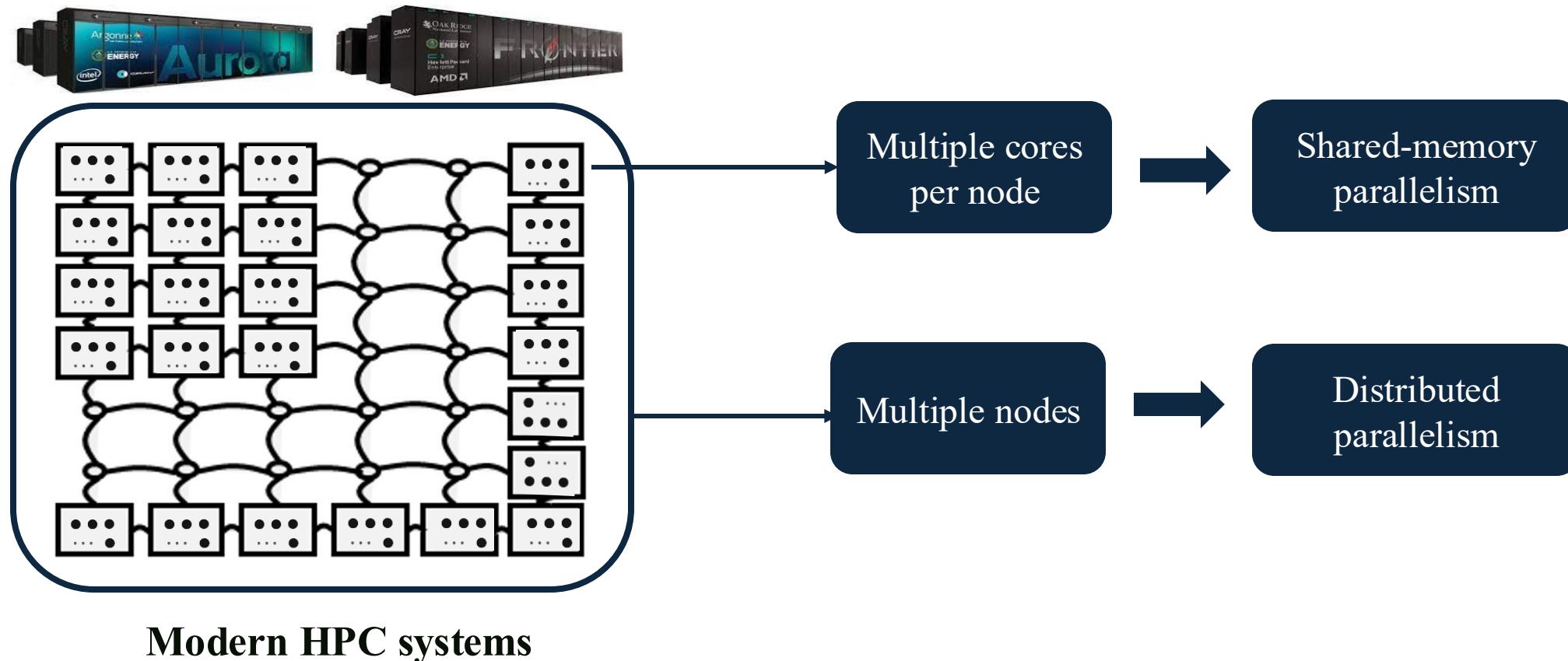
Example of the ParLogNa: Comm Round 1



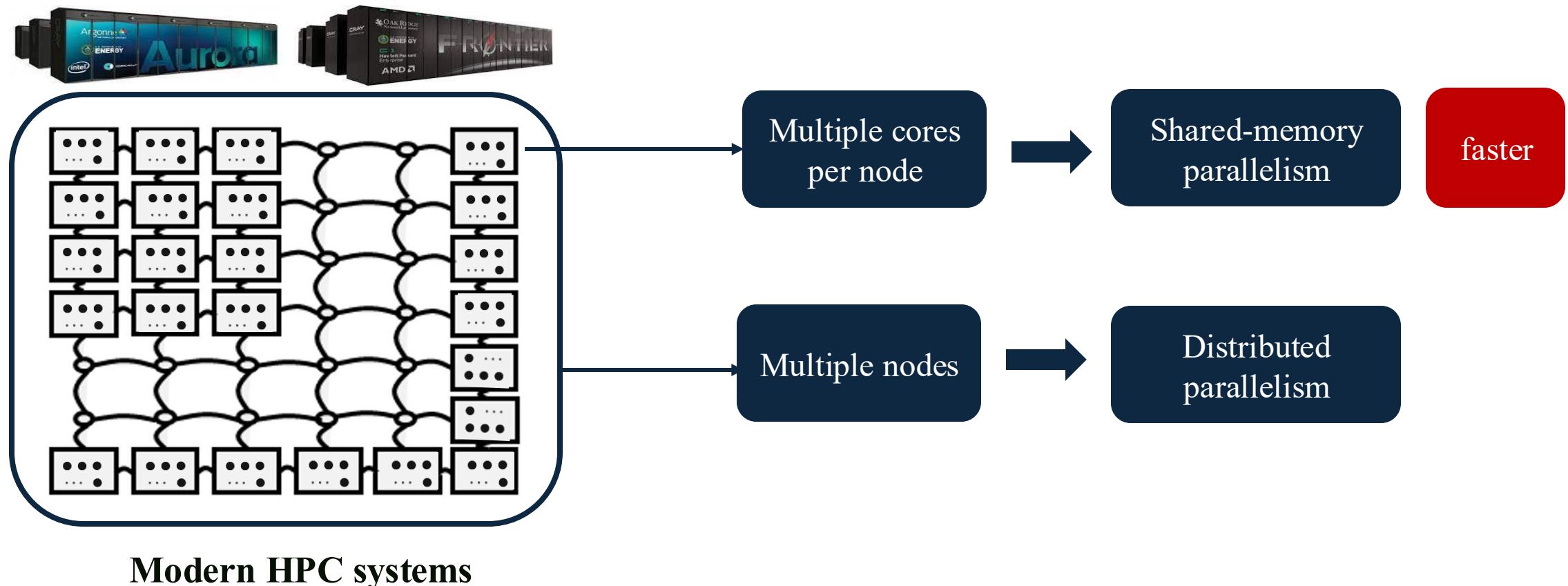
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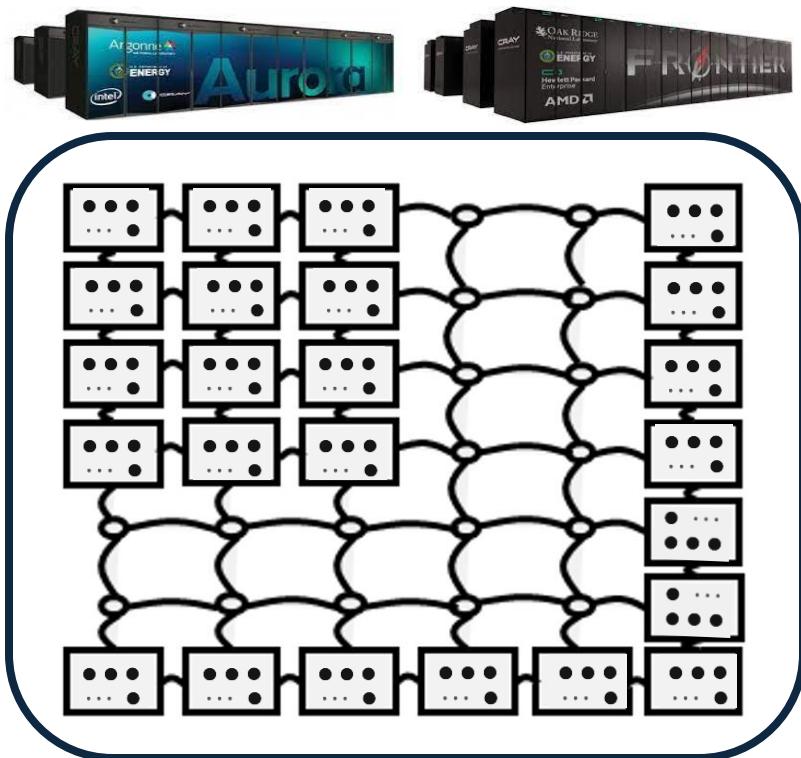
Architecture of Modern HPC Systems



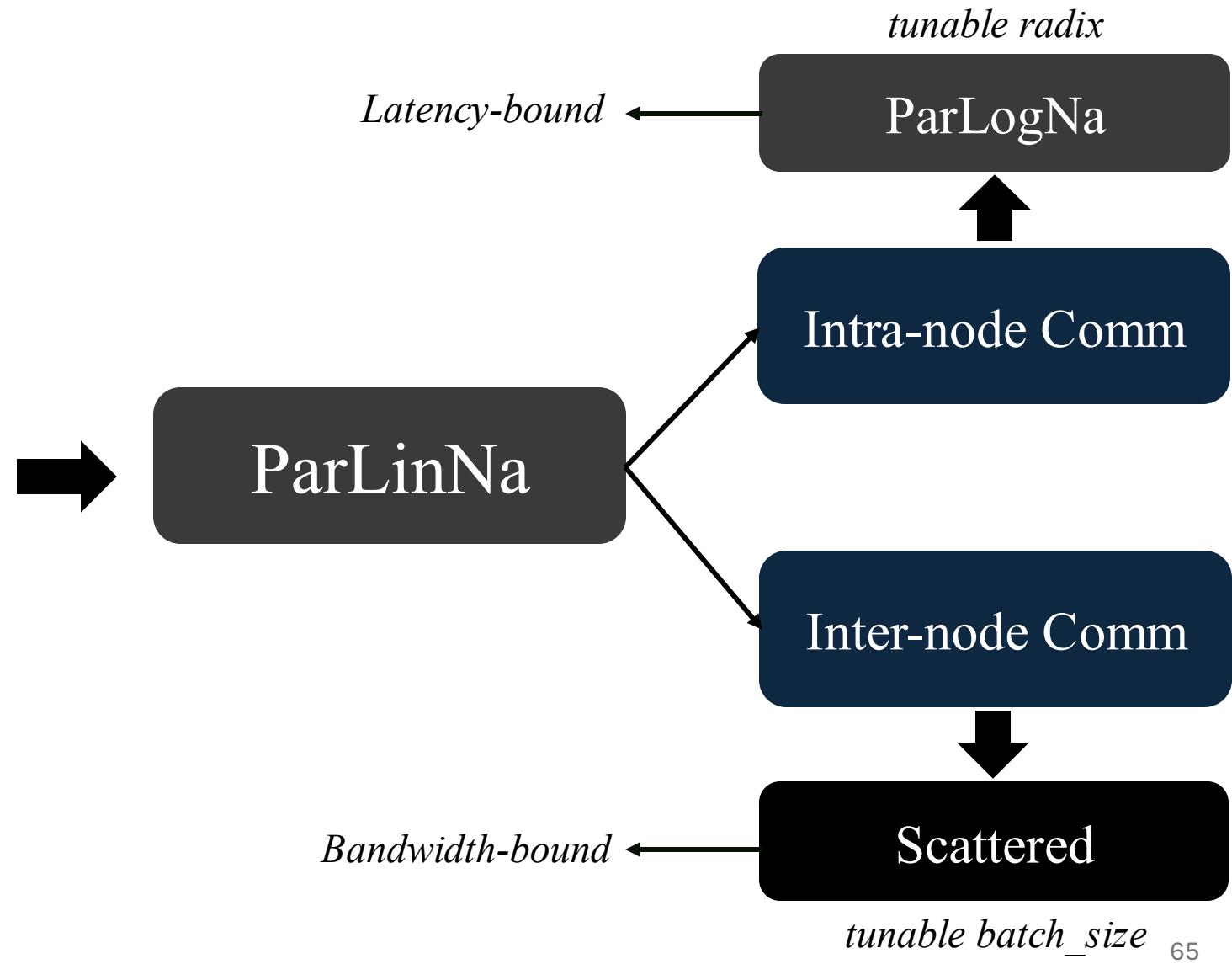
Architecture of Modern HPC Systems



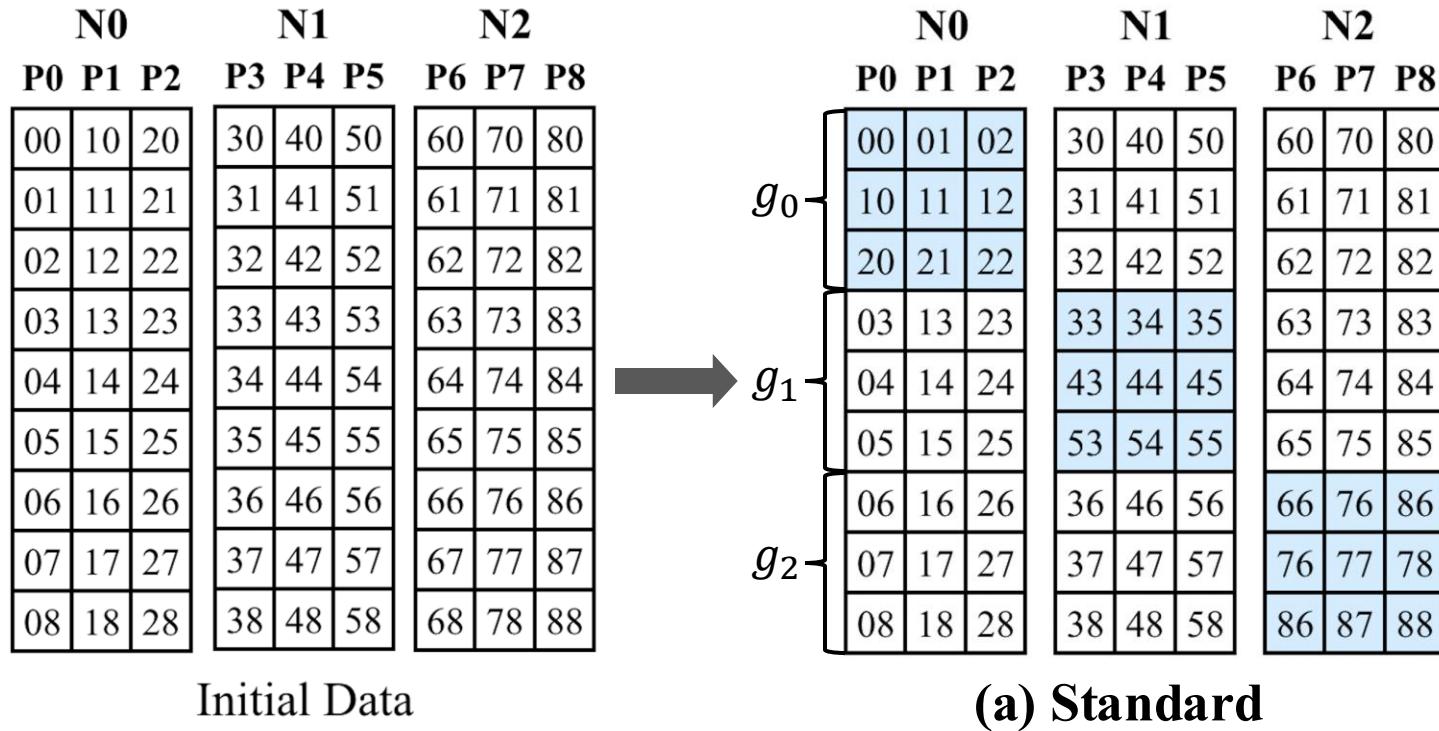
Parameterized Linear Non-uniform All-to-all (ParLinNa)



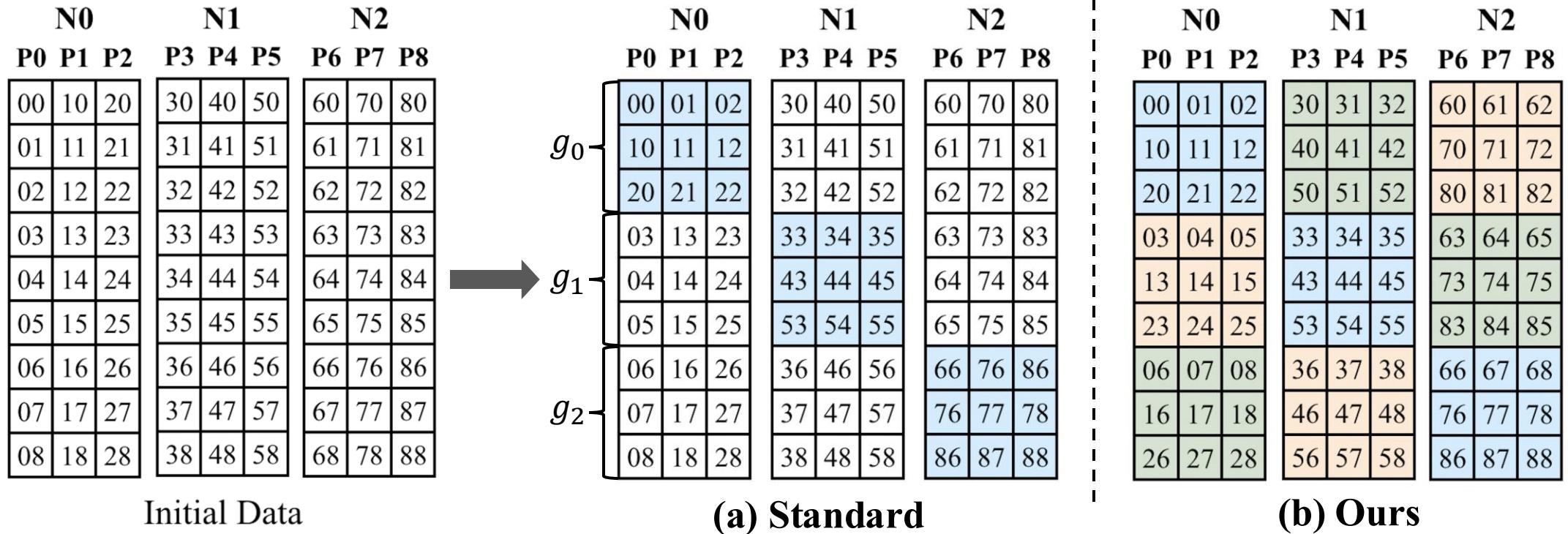
Modern HPC systems



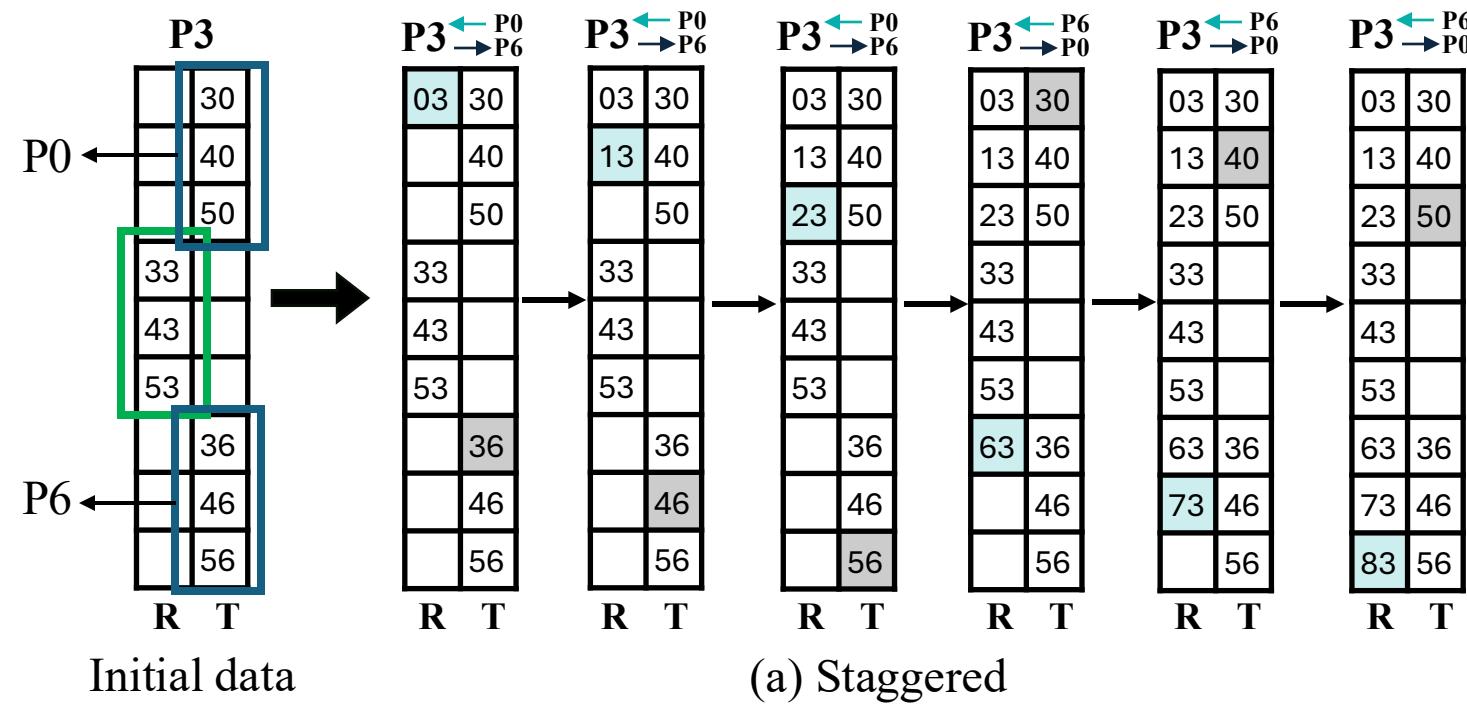
Intra-node Communication: Two Strategies



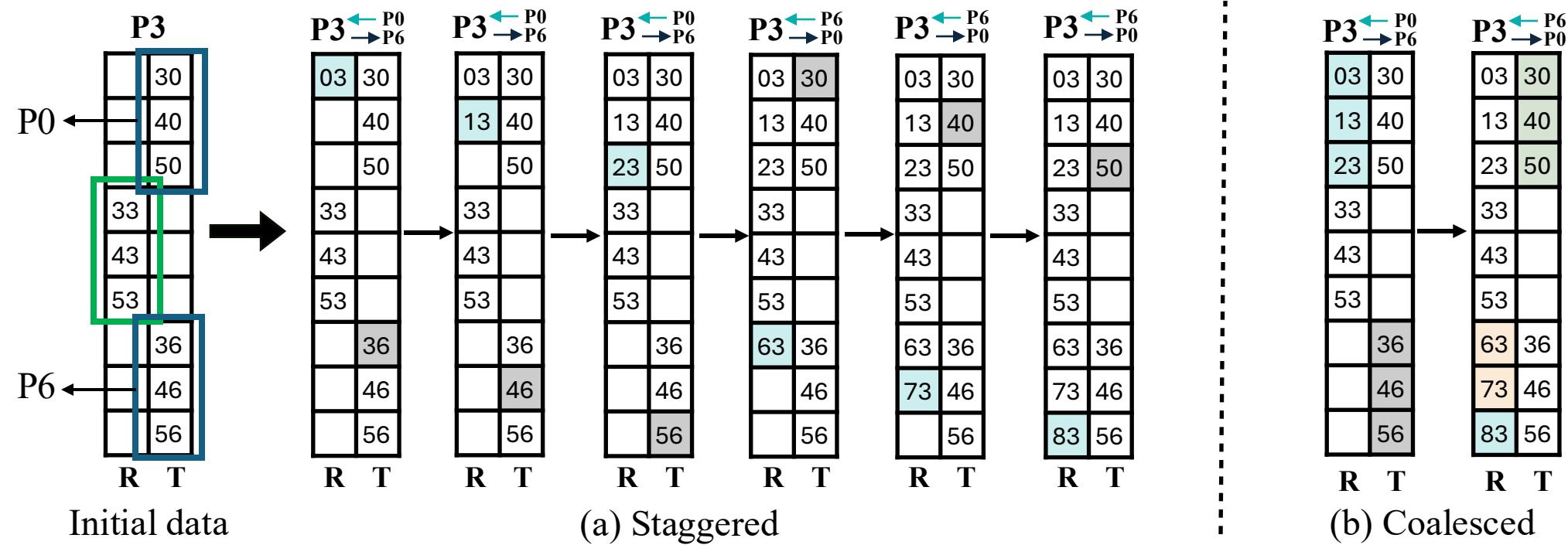
Intra-node Communication: Two Strategies



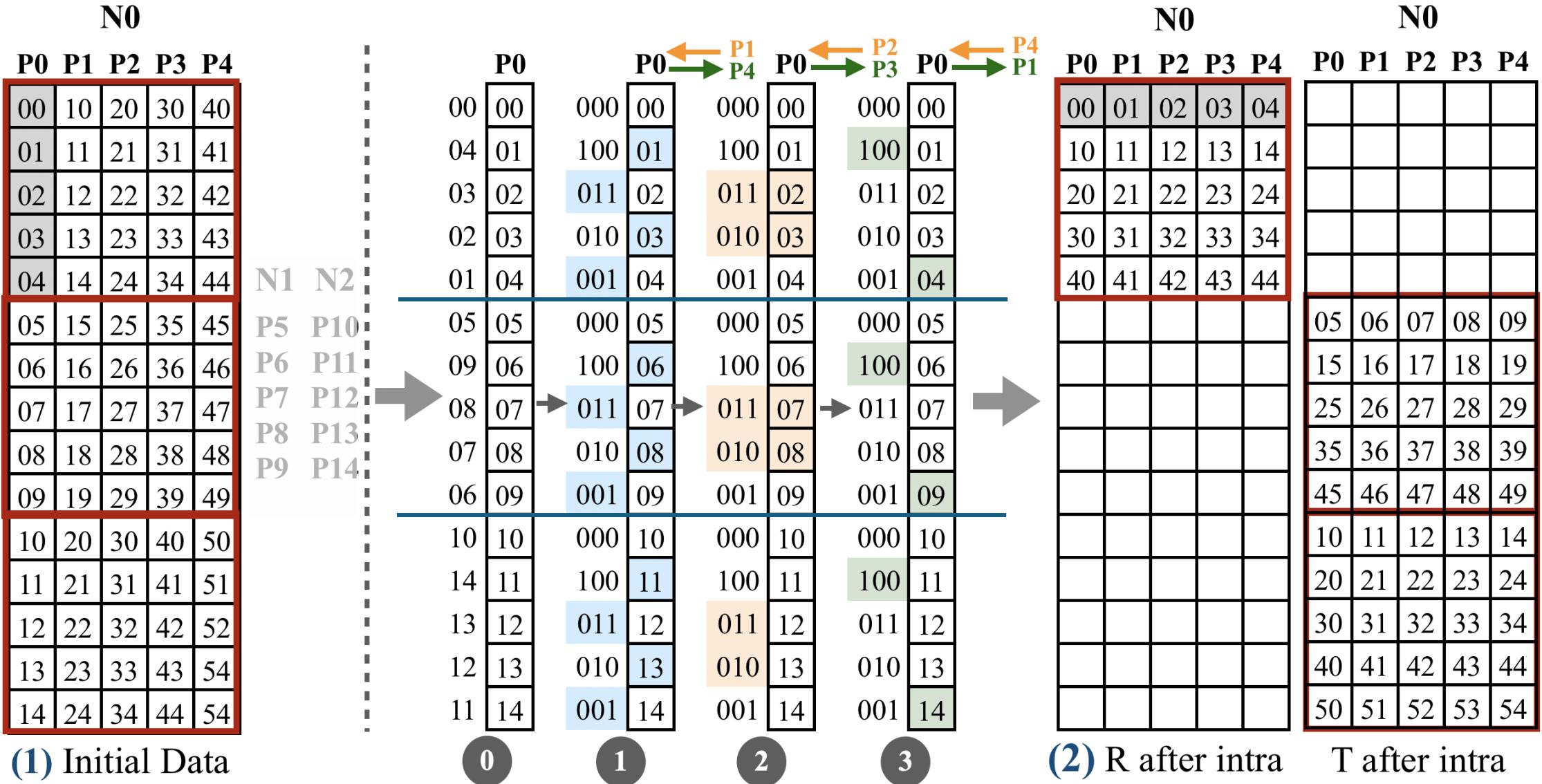
Inter-node Communication: Two Patterns



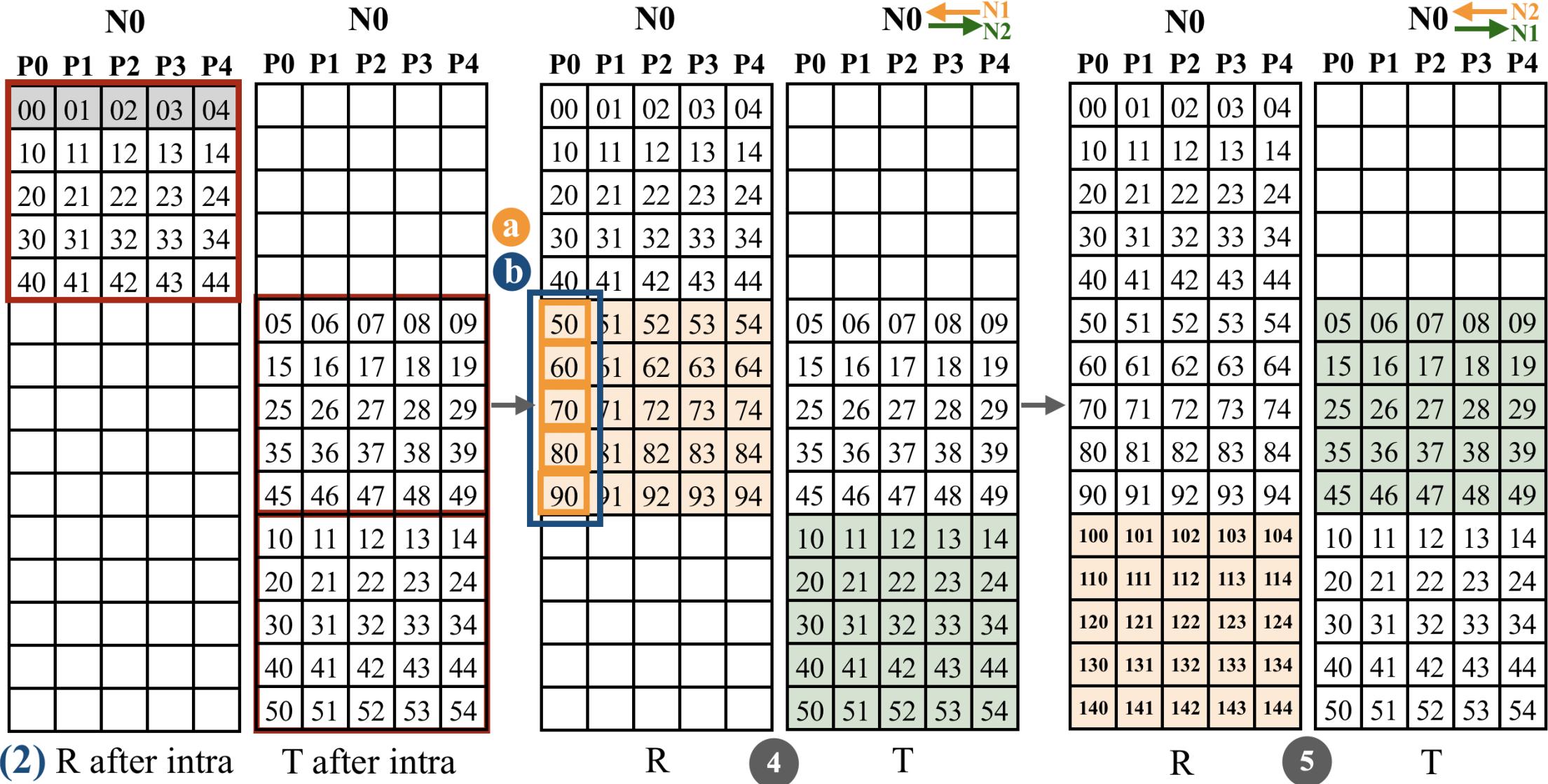
Inter-node Communication: Two Patterns



Intra-node Communication Example



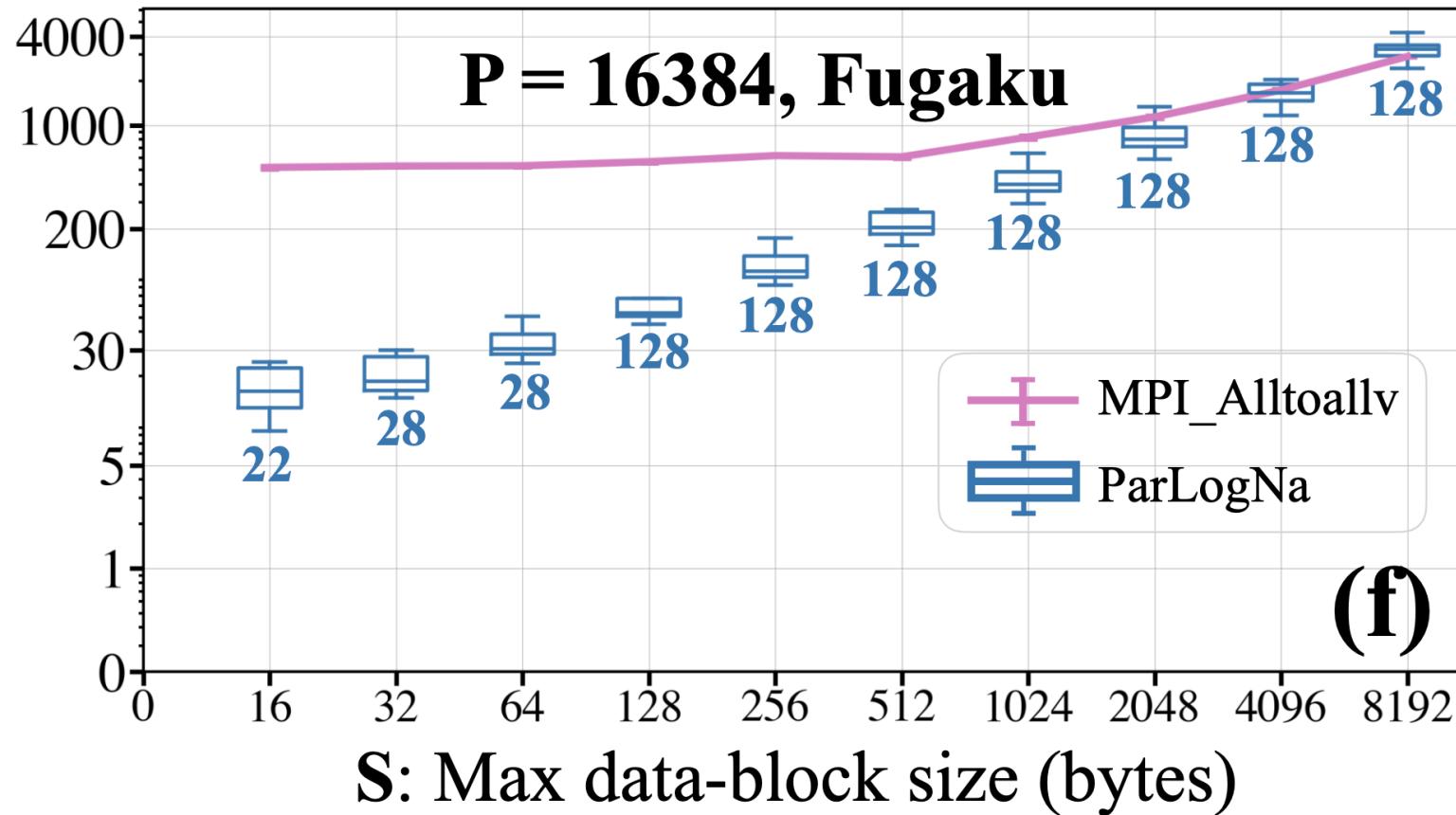
Inter-node Communication Example



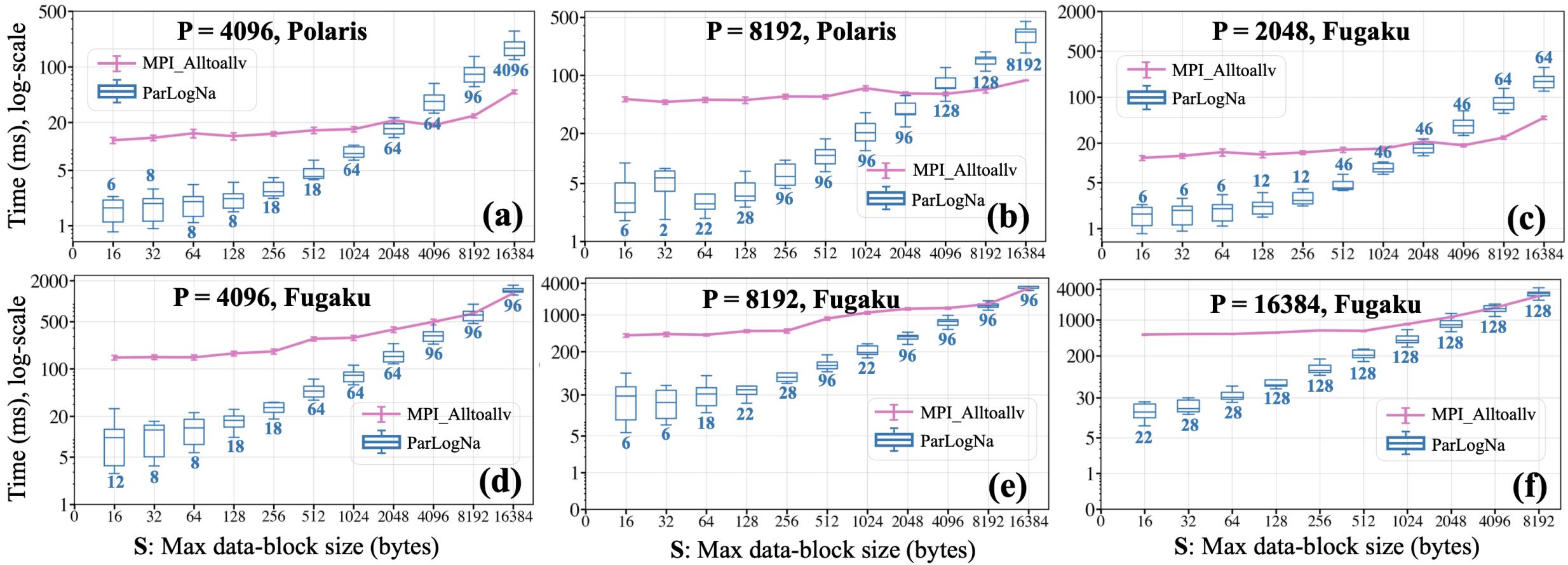
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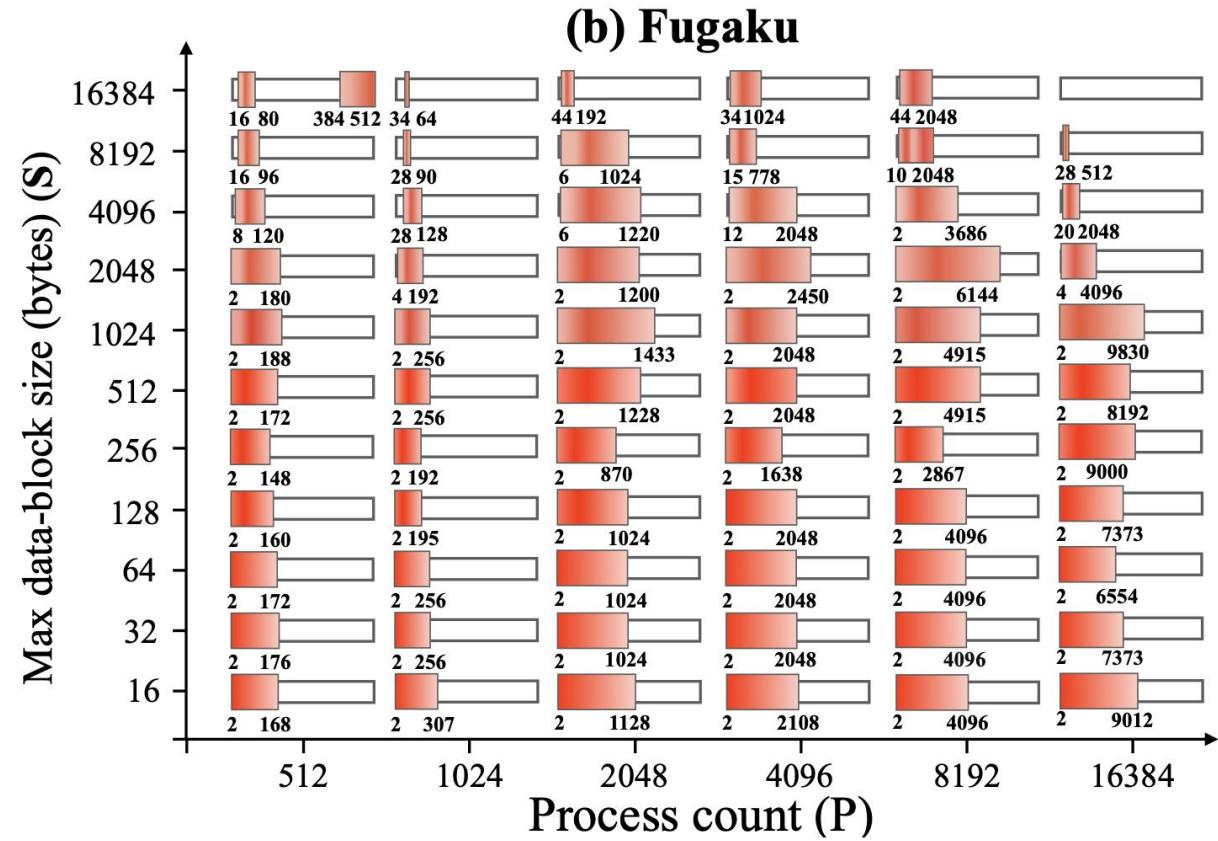
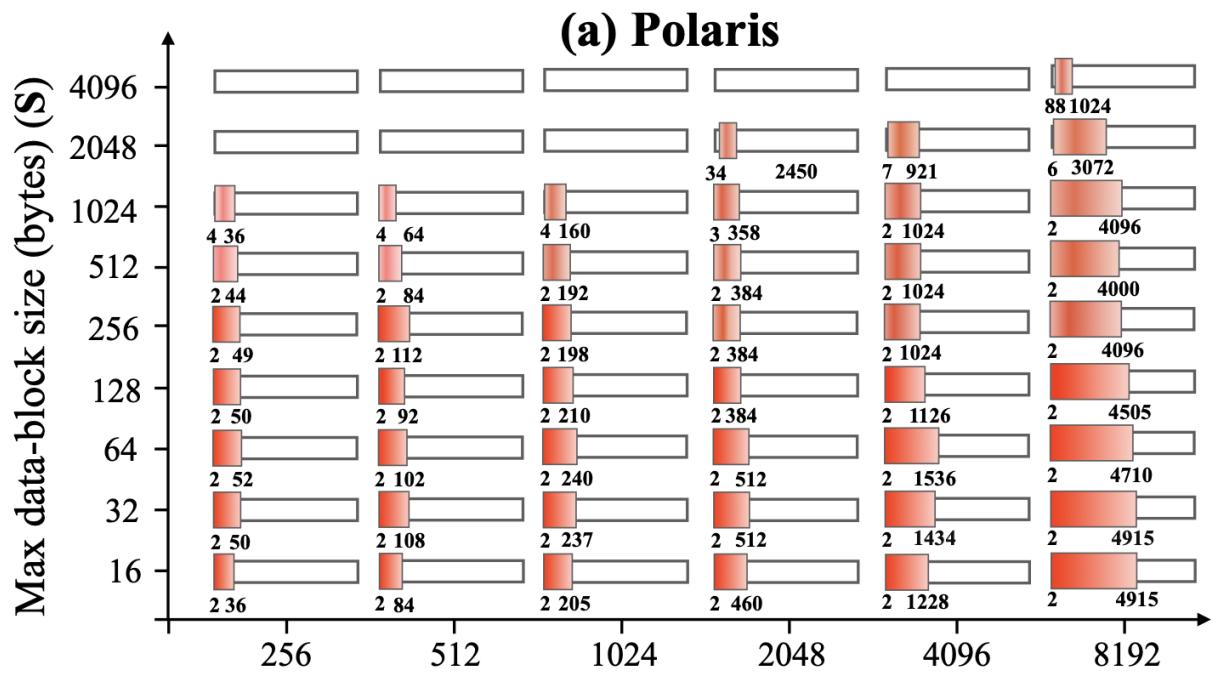
Comparing ParLogNa with MPI_Alltoallv



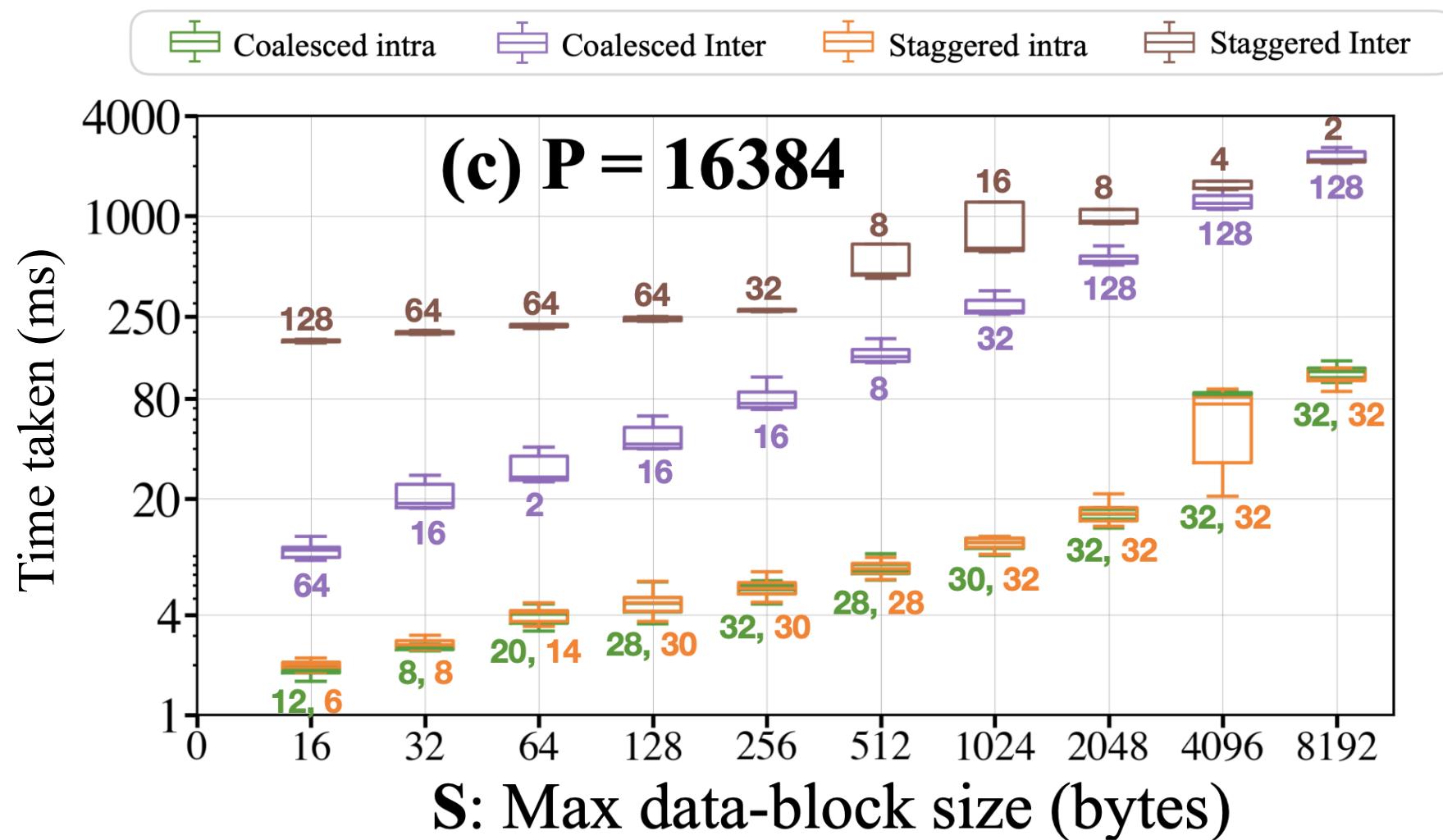
Comparing ParLogNa with MPI_Alltoallv



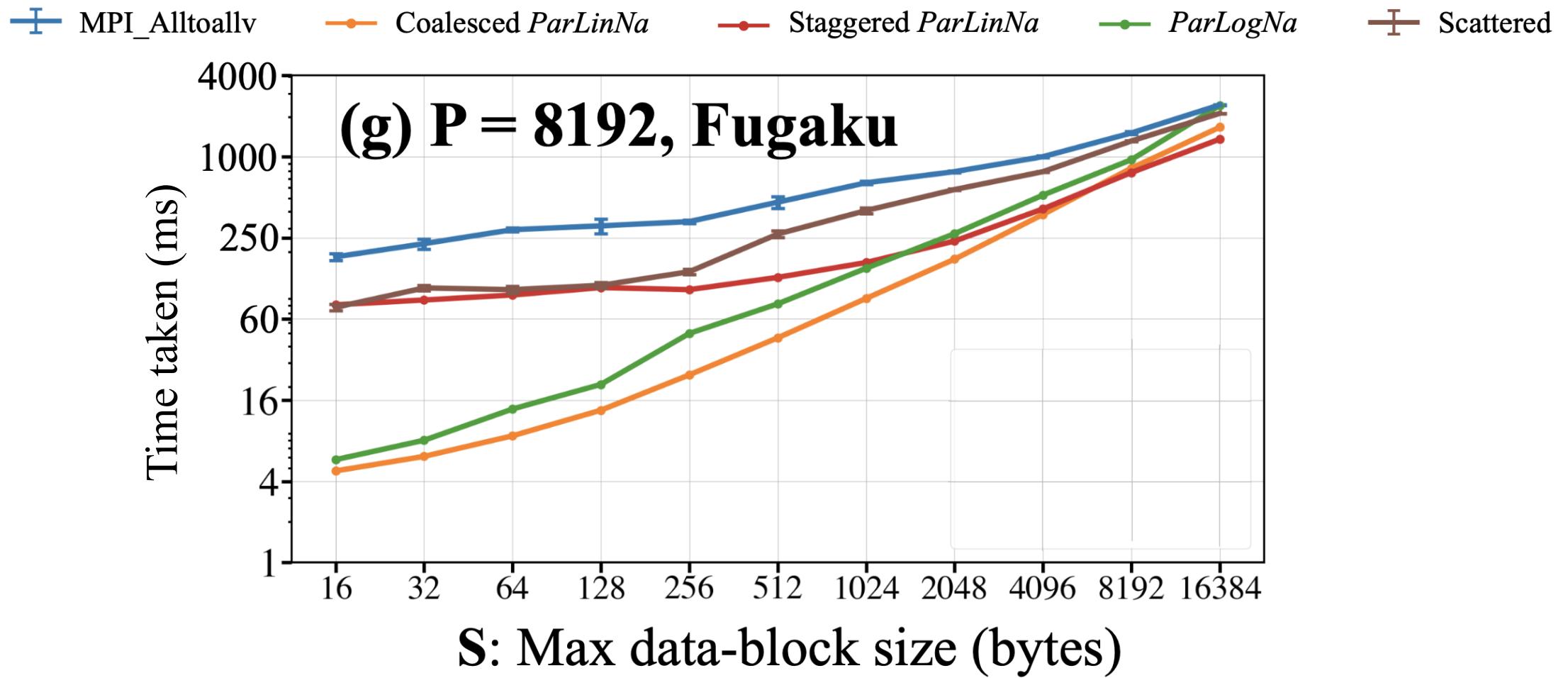
Ranges of radix where ParLogNa outperforms MPI_Alltoallv



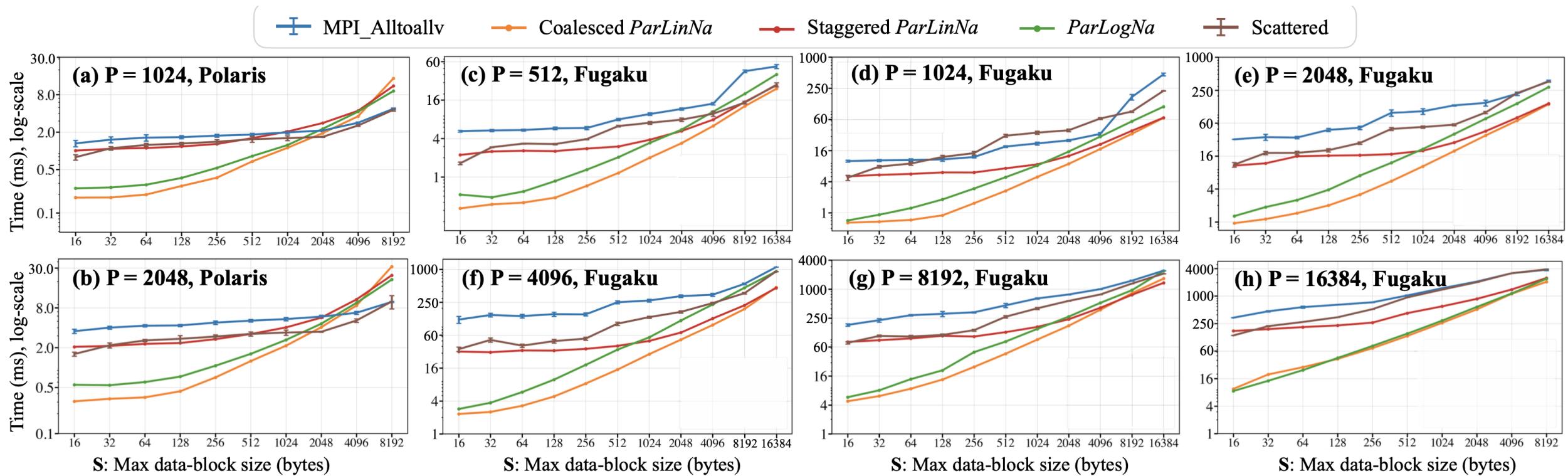
Comparing Coalesced and Staggered ParLinNa



Comparing Proposed Algorithms with the Top-performing Benchmark



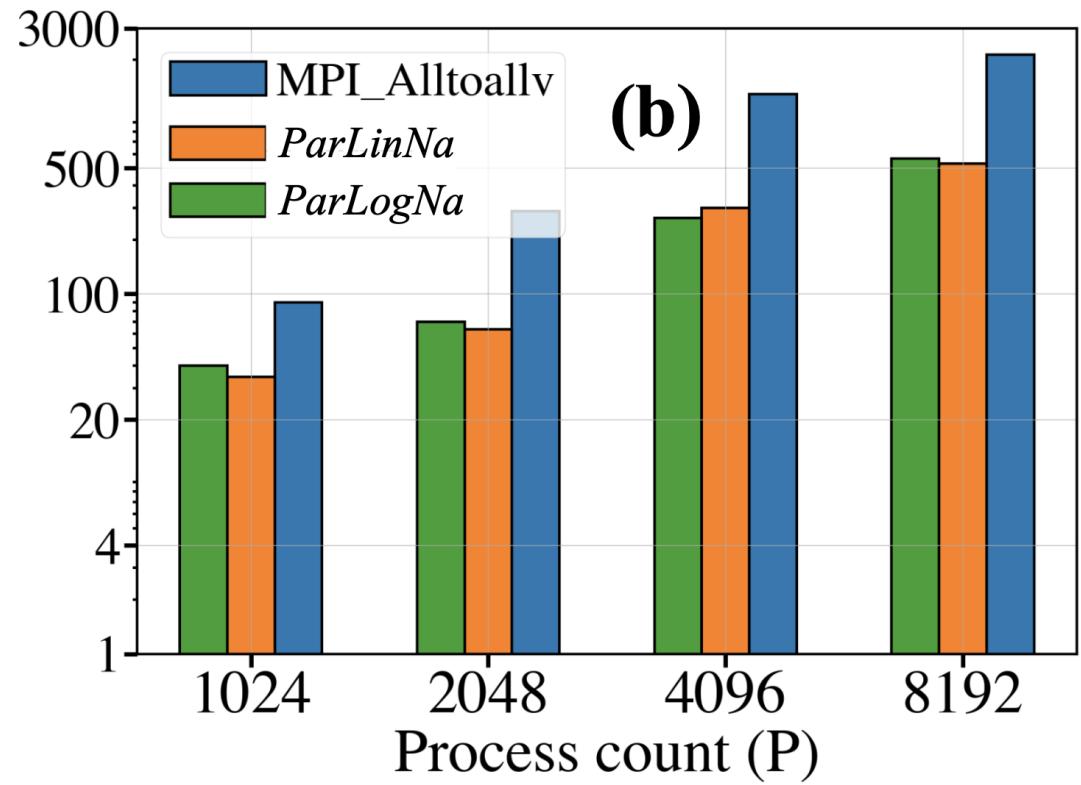
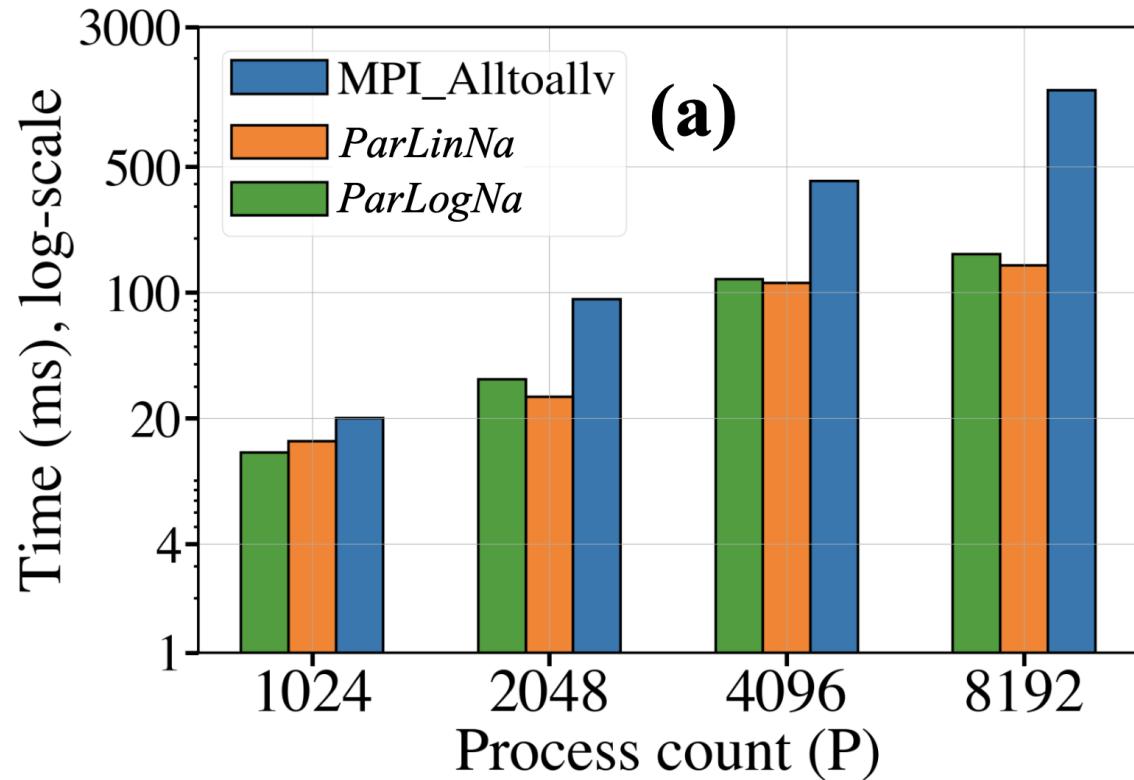
Comparing Proposed Algorithms with the Top-performing Benchmark



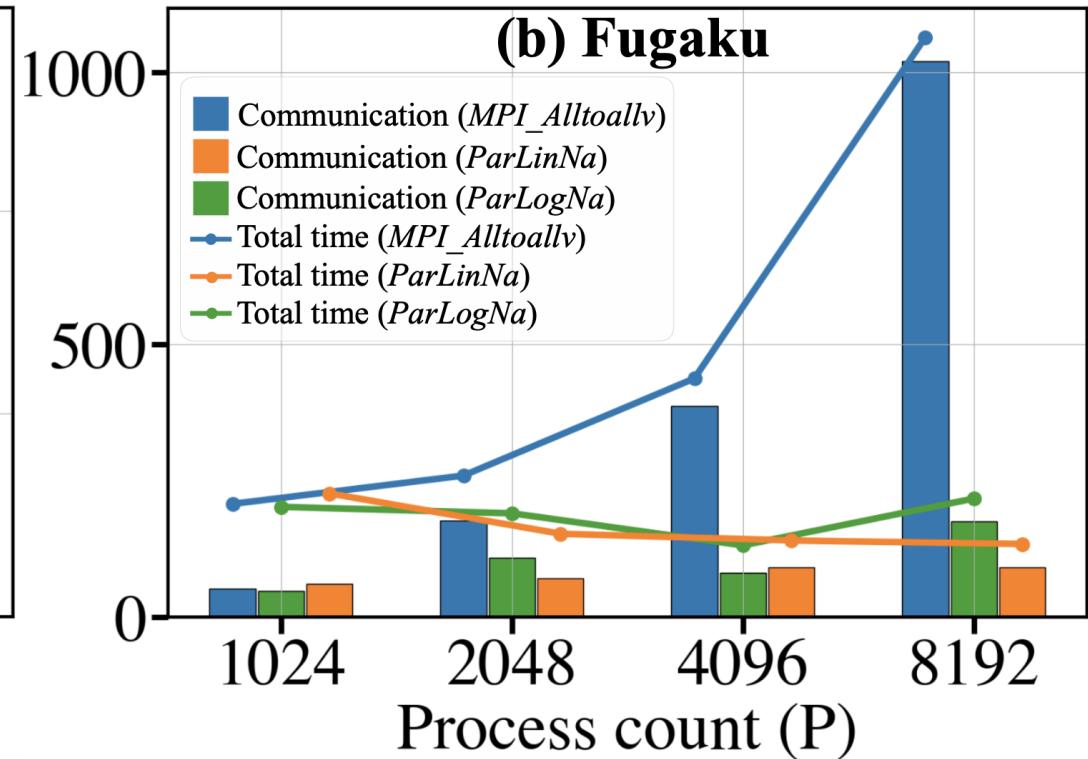
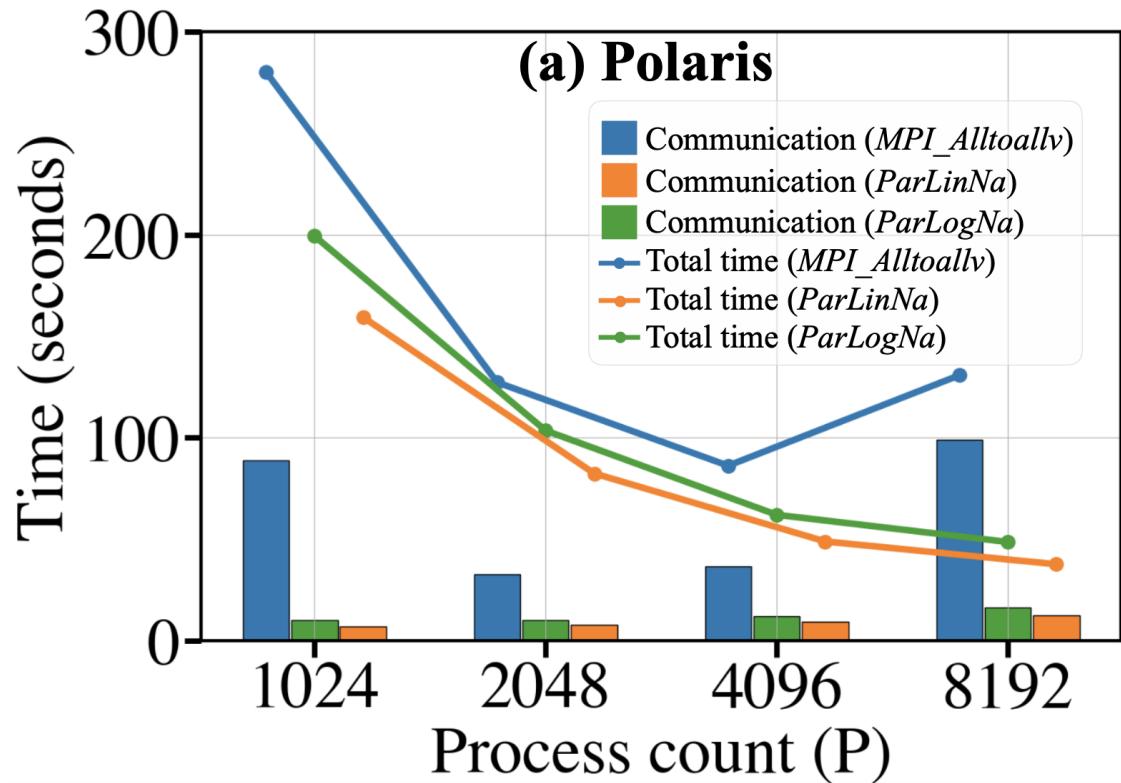
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Performance of Applying Our Algorithms to FFT



Performance of Applying Our Algorithms to Path Finding



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Conclusion

- We have presented algorithms (ParLogNa and ParLinNa) that can improve an important class of collective functions.
- The work can help vendors further optimize their collective routines.
- The work can have a direct impact on a range of applications that uses all to all communication.

Thank you!