# Calculus: Homework #2

Due on February 12, 2014 at 3:10 pm

 $Professor\ Is a a c\ Newton\ Section\ A$ 

Josh Davis

Give an appropriate positive constant c such that  $f(n) \le c \cdot g(n)$  for all n > 1.

1. 
$$f(n) = n^2 + n + 1$$
,  $g(n) = 2n^3$ 

2. 
$$f(n) = n\sqrt{n} + n^2$$
,  $g(n) = n^2$ 

3. 
$$f(n) = n^2 - n + 1$$
,  $g(n) = n^2/2$ 

	one	two	three	or	more	premisses	
and	any	nu	mber	of	conclusions	s as	well

#### Solution

We solve each solution algebraically to determine a possible constant c.

#### Part One

$$n^{2} + n + 1 =$$
 $\leq n^{2} + n^{2} + n^{2}$ 
 $= 3n^{2}$ 
 $\leq c \cdot 2n^{3}$ 

Thus a valid c could be when c = 2.

#### Part Two

$$n^{2} + n\sqrt{n} =$$

$$= n^{2} + n^{3/2}$$

$$\leq n^{2} + n^{4/2}$$

$$= n^{2} + n^{2}$$

$$= 2n^{2}$$

$$\leq c \cdot n^{2}$$

Thus a valid c is c = 2.

### Part Three

$$n^2 - n + 1 =$$

$$\leq n^2$$

$$\leq c \cdot n^2/2$$

Thus a valid c is c = 2.

Let  $\Sigma = \{0, 1\}$ . Construct a DFA A that recognizes the language that consists of all binary numbers that can be divided by 5.

Let the state  $q_k$  indicate the remainder of k divided by 5. For example, the remainder of 2 would correlate to state  $q_2$  because 7 mod 5 = 2.

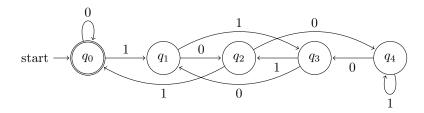


Figure 1: DFA, A, this is really beautiful, ya know?

#### Justification

Take a given binary number, x. Since there are only two inputs to our state machine, x can either become x0 or x1. When a 0 comes into the state machine, it is the same as taking the binary number and multiplying it by two. When a 1 comes into the machine, it is the same as multiplying by two and adding one.

Using this knowledge, we can construct a transition table that tell us where to go:

	$x \mod 5 = 0$	$x \mod 5 = 1$	$x \mod 5 = 2$	$x \mod 5 = 3$	$x \mod 5 = 4$
x0	0	2	4	1	3
x1	1	3	0	2	4

Therefore on state  $q_0$  or  $(x \mod 5 = 0)$ , a transition line should go to state  $q_0$  for the input 0 and a line should go to state  $q_1$  for input 1. Continuing this gives us the Figure 1.

# Problem 3

Write part of Quick-Sort(list, start, end)

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1: function QUICK-SORT(list, start, end)
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- 2: **if**  $start \ge end$  **then**
- 3: return
- 4: end if
- 5:  $mid \leftarrow Partition(list, start, end)$
- 6: Quick-Sort(list, start, mid 1)
- 7: Quick-Sort(list, mid + 1, end)
- 8: end function

Algorithm 1: Start of QuickSort

Suppose we would like to fit a straight line through the origin, i.e.,  $Y_i = \beta_1 x_i + e_i$  with i = 1, ..., n,  $E[e_i] = 0$ , and  $Var[e_i] = \sigma_e^2$  and  $Cov[e_i, e_j] = 0$ ,  $\forall i \neq j$ .

#### Part A

Find the least squares esimator for  $\hat{\beta}_1$  for the slope  $\beta_1$ .

#### Solution

To find the least squares estimator, we should minimize our Residual Sum of Squares, RSS:

$$RSS = \sum_{i=1}^{n} (Y_i - \hat{Y}_i)^2$$
$$= \sum_{i=1}^{n} (Y_i - \hat{\beta}_1 x_i)^2$$

By taking the partial derivative in respect to  $\hat{\beta}_1$ , we get:

$$\frac{\partial}{\partial \hat{\beta}_1}(RSS) = -2\sum_{i=1}^n x_i(Y_i - \hat{\beta}_1 x_i) = 0$$

This gives us:

$$\sum_{i=1}^{n} x_i (Y_i - \hat{\beta}_1 x_i) = \sum_{i=1}^{n} x_i Y_i - \sum_{i=1}^{n} \hat{\beta}_1 x_i^2$$
$$= \sum_{i=1}^{n} x_i Y_i - \hat{\beta}_1 \sum_{i=1}^{n} x_i^2$$

Solving for  $\hat{\beta}_1$  gives the final estimator for  $\beta_1$ :

$$\hat{\beta_1} = \frac{\sum x_i Y_i}{\sum x_i^2}$$

## Part B

Calculate the bias and the variance for the estimated slope  $\hat{\beta}_1$ .

## Solution

For the bias, we need to calculate the expected value  $E[\hat{\beta}_1]$ :

$$\begin{split} \mathbf{E}[\hat{\beta}_1] &= \mathbf{E}\left[\frac{\sum x_i Y_i}{\sum x_i^2}\right] \\ &= \frac{\sum x_i \mathbf{E}[Y_i]}{\sum x_i^2} \\ &= \frac{\sum x_i (\beta_1 x_i)}{\sum x_i^2} \\ &= \frac{\sum x_i^2 \beta_1}{\sum x_i^2} \\ &= \beta_1 \frac{\sum x_i^2 \beta_1}{\sum x_i^2} \\ &= \beta_1 \end{split}$$

Thus since our estimator's expected value is  $\beta_1$ , we can conclude that the bias of our estimator is 0.

For the variance:

$$\begin{aligned} \operatorname{Var}[\hat{\beta}_1] &= \operatorname{Var}\left[\frac{\sum x_i Y_i}{\sum x_i^2}\right] \\ &= \frac{\sum x_i^2}{\sum x_i^2 \sum x_i^2} \operatorname{Var}[Y_i] \\ &= \frac{\sum x_i^2}{\sum x_i^2 \sum x_i^2} \operatorname{Var}[Y_i] \\ &= \frac{1}{\sum x_i^2} \operatorname{Var}[Y_i] \\ &= \frac{1}{\sum x_i^2} \sigma^2 \\ &= \frac{\sigma^2}{\sum x_i^2} \end{aligned}$$

# Problem 5

Prove a polynomial of degree k,  $a_k n^k + a_{k-1} n^{k-1} + \ldots + a_1 n^1 + a_0 n^0$  is a member of  $\Theta(n^k)$  where  $a_k \ldots a_0$  are nonnegative constants.

*Proof.* To prove that  $a_k n^k + a_{k-1} n^{k-1} + \ldots + a_1 n^1 + a_0 n^0$ , we must show the following:

$$\exists c_1 \exists c_2 \forall n \ge n_0, \ c_1 \cdot g(n) \le f(n) \le c_2 \cdot g(n)$$

For the first inequality, it is easy to see that it holds because no matter what the constants are,  $n^k \le a_k n^k + a_{k-1} n^{k-1} + \ldots + a_1 n^1 + a_0 n^0$  even if  $c_1 = 1$  and  $n_0 = 1$ . This is because  $n^k \le c_1 \cdot a_k n^k$  for any nonnegative constant,  $c_1$  and  $a_k$ .

Taking the second inequality, we prove it in the following way. By summation,  $\sum_{i=0}^{k} a_i$  will give us a new constant, A. By taking this value of A, we can then do the following:

$$a_k n^k + a_{k-1} n^{k-1} + \ldots + a_1 n^1 + a_0 n^0 =$$

$$\leq (a_k + a_{k-1} \ldots a_1 + a_0) \cdot n^k$$

$$= A \cdot n^k$$

$$< c_2 \cdot n^k$$

where  $n_0 = 1$  and  $c_2 = A$ .  $c_2$  is just a constant. Thus the proof is complete.

Evaluate  $\sum_{k=1}^{5} k^2$  and  $\sum_{k=1}^{5} (k-1)^2$ .

# Problem 19

Find the derivative of  $f(x) = x^4 + 3x^2 - 2$ 

# Problem 6

Evaluate the integrals  $\int_0^1 (1-x^2) \mathrm{d}x$  and  $\int_1^\infty \frac{1}{x^2} \mathrm{d}x$ .