

原文：<https://www.zerodayinitiative.com/blog/2018/6/19/analyzing-an-integer-overflow-in-bitdefender-av-part-1-the-vulnerability>

在众多的软件缺陷中，安全软件中出现的漏洞被认为比其他漏洞尤为严重。这是因为，安全软件是帮助用户抵御攻击者的靠山，所以，如果这些软件中出现安全问题，不仅会影响Internet

Security中的安全漏洞的原因之一，如果该漏洞被利用，这款安全防护软件就会被攻击者用于远程执行代码。实际上，这个漏洞是由整数溢出引起的，尽管现在已经修复，但

当前，尽管系统已经提供了DEP和ASLR等防御机制，但是，人们仍然可以对现代软件中的整数溢出进行有效的利用。在本文中，我们将对这个运行攻击者在目标系统上面远

Bitdefender概述

Bitdefender提供了多种具有防病毒（AV）功能的产品，并且所有这些AV产品都安装了SYSTEM服务vsserv.exe。该服务的作用是处理AV扫描请求。就像各种AV引擎一样，Files\Bitdefender\Bitdefender Threat

Scanner\Antivirus_XXX\Plugins\目录中。除此之外，该引擎还有一个鲜为人知的功能，即模拟遇到的可执行代码，以此检测未知或经过混淆处理的病毒。

遇到x86可移植可执行文件（PE）时，Bitdefender实际上是通过保存在cevakrnl.xmd和ceva_emu.cvd文件中高度复杂的代码仿真器来虚拟执行PE的。这些仿真器会创建一API的实现，并为频繁执行的地址创建相应的JIT（Just-In-Time，JIT）代码。

杀毒原理简介

当仿真器检测到函数调用时，就会调用cevakrnl.xmd!sub_366F4D0()。该函数处理被仿真函数中的第一条指令，实际上就是在一个常量表中搜索匹配项。在处理后面的指令

```
cevakrnl.xmd:0366F4D0 sub_366F4D0      proc near          ; CODE XREF: sub_363F400+Fp
cevakrnl.xmd:0366F4D0
...
cevakrnl.xmd:0366F530      mov     edi, [edx+esi*4] ; offset
cevakrnl.xmd:0366F533      mov     bl, [ebx+eax]   ; first byte in code
cevakrnl.xmd:0366F536      add     edi, ecx
cevakrnl.xmd:0366F538      inc     eax
cevakrnl.xmd:0366F539      mov     [ebp+var_218], eax
cevakrnl.xmd:0366F53F      movzx   eax, byte ptr [edi]
cevakrnl.xmd:0366F542      and     eax, 0FFFFFFFh
cevakrnl.xmd:0366F547      dec     eax              ; interpreter state
cevakrnl.xmd:0366F548      cmp     eax, 3           ; switch 4 cases
cevakrnl.xmd:0366F54B      ja      loc_366F5E2      ; jumptable 0366F551 default case
cevakrnl.xmd:0366F551      jmp     ds:off_366F7F4[eax*4] ; switch jump
...
cevakrnl.xmd:0366F58F      movzx   edx, byte ptr [edi+1] ; search distance
cevakrnl.xmd:0366F593      xor     esi, esi
cevakrnl.xmd:0366F595      test    edx, edx
cevakrnl.xmd:0366F597      jz      @out
cevakrnl.xmd:0366F59D      lea     ecx, [ecx+0]
cevakrnl.xmd:0366F5A0      @locatelopcode:
cevakrnl.xmd:0366F5A0      lea     eax, [edx+esi]      ; CODE XREF: sub_366F4D0+EBj
cevakrnl.xmd:0366F5A3      shr     eax, 1
cevakrnl.xmd:0366F5A5      lea     ecx, [edi+eax*2] ; edi + eax*3 + 2 = opcode
cevakrnl.xmd:0366F5A8      mov     cl, [eax+ecx+2]
cevakrnl.xmd:0366F5AC      cmp     bl, cl
cevakrnl.xmd:0366F5AE      jnb     short loc_366F5B4
cevakrnl.xmd:0366F5B0      mov     edx, eax
cevakrnl.xmd:0366F5B2      jmp     short loc_366F5B9
cevakrnl.xmd:0366F5B4 ;
-----
cevakrnl.xmd:0366F5B4      loc_366F5B4:
cevakrnl.xmd:0366F5B4      jbe     short @opcodematchfound ; CODE XREF: sub_366F4D0+DEj
cevakrnl.xmd:0366F5B6      lea     esi, [eax+1]
cevakrnl.xmd:0366F5B9      loc_366F5B9:
cevakrnl.xmd:0366F5B9      cmp     esi, edx          ; CODE XREF: sub_366F4D0+E2j
cevakrnl.xmd:0366F5BB      jb      short @locatelopcode
cevakrnl.xmd:0366F5BD      jmp     @out
cevakrnl.xmd:0366F5C2 ;
-----
```

当函数的前16个字节被仿真器找到相应的匹配项时，或者当指令序列是未知的指令序列时，初始搜索就宣告结束。

常量表最后的匹配项会提供函数起始位置之前一些已知字节的内容，以及函数起始位置后16字节的内容。

[X bytes before]

FUNCTIONSTART:

16 bytes already matched

[Y bytes after]

然后，该引擎会读取并检查这些字节的内容，以确定是否与已知的AV签名相匹配。如果匹配，则调用检测到的代码签名对应的处理程序做进一步处理。

漏洞分析

遇到Themida加壳软件代码序列时，该杀毒软件就会调用sub_36906D0()函数以便对匹配的序列进行代码解释。

```
cevakrn1.xmd:0368BAA0      push     0A69h
cevakrn1.xmd:0368BAA5      push     97Fh
cevakrn1.xmd:0368BAAA      push     16Fh
cevakrn1.xmd:0368BAAF      push     931h
cevakrn1.xmd:0368BAB4      push     ebx
cevakrn1.xmd:0368BAB5      push     edi
cevakrn1.xmd:0368BAB6      call     sub_36906D0
```

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被解释的代码中，唯一的变量为ebp偏移量“X”，并且ebx的值就是通过它读取的。该函数从代码流中提取偏移量，并使用它从仿真堆栈中获取ebx的值。

```
cevakrn1.xmd:036906D0 sub_36906D0      proc near
...
cevakrn1.xmd:036906D0      push     ebp
cevakrn1.xmd:036906D1      mov      ebp, esp
cevakrn1.xmd:036906D3      sub      esp, 84h
cevakrn1.xmd:036906D9      mov      eax, ds:dword_369A5C4
cevakrn1.xmd:036906DE      xor      eax, ebp
cevakrn1.xmd:036906E0      mov      [ebp+var_4], eax
cevakrn1.xmd:036906E3      push     ebx
cevakrn1.xmd:036906E4      mov      ebx, [ebp+arg_4]
cevakrn1.xmd:036906E7      push     edi
cevakrn1.xmd:036906E8      mov      edi, [ebp+arg_0]
cevakrn1.xmd:036906EB      push     0 ; int
cevakrn1.xmd:036906ED      mov      eax, [ebx+88h] ; virtual function address
cevakrn1.xmd:036906F3      sub      eax, 0Fh
cevakrn1.xmd:036906F6      mov      [ebp+var_60], ebx
cevakrn1.xmd:036906F9      push     eax ; addr
cevakrn1.xmd:036906FA      lea      eax, [ebp+put] ; mov ebx, [ebp+0xdeadbeef]
cevakrn1.xmd:036906FD      push     eax ; output
cevakrn1.xmd:036906FE      push     dword ptr [edi+0Ch] ; int
cevakrn1.xmd:03690701      call     get_virtual_dword ; read ebp offset x
cevakrn1.xmd:03690706      add      esp, 10h
cevakrn1.xmd:03690709      test     eax, eax
cevakrn1.xmd:0369070B      jnz      loc_3690C64
cevakrn1.xmd:03690711      mov      eax, [ebx+28h] ; 0x12ffc0 (ebp)
cevakrn1.xmd:03690711      ; put has to be dataaddr-0x12ffc0
cevakrn1.xmd:03690711      ; dataaddr holds first dword after
standard code
cevakrn1.xmd:03690711      ; which can be controlled
cevakrn1.xmd:03690711      ; functionstart+0x35C = dataaddr
cevakrn1.xmd:03690714      add      eax, [ebp+put] ; X => [ebp+X]
cevakrn1.xmd:03690717      push     0 ; int
cevakrn1.xmd:03690719      push     eax ; addr
cevakrn1.xmd:0369071A      lea      eax, [ebp+readebx] ; ebx
cevakrn1.xmd:0369071D      push     eax ; output
cevakrn1.xmd:0369071E      push     dword ptr [edi+0Ch] ; int
cevakrn1.xmd:03690721      call     get_virtual_dword ; read ebx from [ebp+X]
```

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接下来，会解释mov ecx, [ebx]指令，ecx的值是从保存在ebx中的模拟地址中提取的。

```

cevakrn1.xmd:036906D0 sub_36906D0 proc near
...
cevakrn1.xmd:036906D0 push ebp
cevakrn1.xmd:036906D1 mov ebp, esp
cevakrn1.xmd:036906D3 sub esp, 84h
cevakrn1.xmd:036906D9 mov eax, ds:dword_369A5C4
cevakrn1.xmd:036906DE xor eax, ebp
cevakrn1.xmd:036906E0 mov [ebp+var_4], eax
cevakrn1.xmd:036906E3 push ebx
cevakrn1.xmd:036906E4 mov ebx, [ebp+arg_4]
cevakrn1.xmd:036906E7 push edi
cevakrn1.xmd:036906E8 mov edi, [ebp+arg_0]
cevakrn1.xmd:036906EB push 0 ; int
cevakrn1.xmd:036906ED mov eax, [ebx+88h] ; virtual function address
cevakrn1.xmd:036906F3 sub eax, 0Fh
cevakrn1.xmd:036906F6 mov [ebp+var_60], ebx
cevakrn1.xmd:036906F9 push eax ; addr
cevakrn1.xmd:036906FA lea eax, [ebp+put] ; mov ebx, [ebp+0xdeadbeef]
cevakrn1.xmd:036906FD push eax ; output
cevakrn1.xmd:036906FE push dword ptr [edi+0Ch] ; int
cevakrn1.xmd:03690701 call get_virtual_dword ; read ebp offset x
cevakrn1.xmd:03690706 add esp, 10h
cevakrn1.xmd:03690709 test eax, eax
cevakrn1.xmd:0369070B jnz loc_3690C64
cevakrn1.xmd:03690711 mov eax, [ebx+28h] ; 0x12ffc0 (ebp)
cevakrn1.xmd:03690711 ; put has to be dataaddr-0x12ffc0
cevakrn1.xmd:03690711 ; dataaddr holds first dword after
standard code
cevakrn1.xmd:03690711 ; which can be controlled
cevakrn1.xmd:03690711 ; functionstart+0x35C = dataaddr
cevakrn1.xmd:03690714 add eax, [ebp+put] ; X => [ebp+X]
cevakrn1.xmd:03690717 push 0 ; int
cevakrn1.xmd:03690719 push eax ; addr
cevakrn1.xmd:0369071A lea eax, [ebp+readebx] ; ebx
cevakrn1.xmd:0369071D push eax ; output
cevakrn1.xmd:0369071E push dword ptr [edi+0Ch] ; int
cevakrn1.xmd:03690721 call get_virtual_dword ; read ebx from [ebp+X]

```

接下来解释下面的代码序列：

```

cevakrn1.xmd:03690806 push [ebp+readebx] ; addr
cevakrn1.xmd:03690809 mov [ebp+var_68], eax
cevakrn1.xmd:0369080C lea eax, [ebp+readecx]
cevakrn1.xmd:0369080F push eax ; output
cevakrn1.xmd:03690810 push dword ptr [edi+0Ch] ; int
cevakrn1.xmd:03690813 call get_virtual_dword

```

这里，会根据 $ecx + \text{dword}[ecx + 0x78 + \text{word}[ecx+0x3c]]$ 计算出仿真esi的值。

```

.data:000000A5 push eax
.data:000000A6 push ecx
.data:000000A7 pusha
.data:000000A8 xor eax, eax
.data:000000AA mov [ebp+0], eax
.data:000000B0 mov esi, 3Ch
.data:000000B5 BYTESAFTER:
.data:000000B5 add esi, [esp+20h] ; saved ecx offset 0x3C
.data:000000B9 lodsw
.data:000000BB add eax, [esp+20h] ; saved ecx
.data:000000BF mov esi, [eax+78h] ; dword [ecx + word [ecx+0x3C]+0x78]
.data:000000C2 add esi, [esp+20h] ; ecx + dword [ecx + word [ecx+0x3C]+0x78]

```

接下来从仿真的esi中总共提取0x28字节，并读取偏移0x18处的dword（对应仿真代码的mov edi, [esi + 18h]指令）。


```

cevakrn1.xmd:0369082E      add     eax, 3Ch
cevakrn1.xmd:03690831      mov     [ebp+var_30], 0
cevakrn1.xmd:03690838      push    0
cevakrn1.xmd:0369083A      push    eax                ; ecx+0x3c
cevakrn1.xmd:0369083B      lea     eax, [ebp+pout]
cevakrn1.xmd:0369083E      push    eax
cevakrn1.xmd:0369083F      push    dword ptr [edi+0Ch]
cevakrn1.xmd:03690842      call    get_virtual_word ; word ptr [ecx+0x3C]
cevakrn1.xmd:03690847      add     esp, 10h
cevakrn1.xmd:0369084A      test    eax, eax
cevakrn1.xmd:0369084C      jnz     loc_3690C63
cevakrn1.xmd:03690852      mov     eax, [ebp+readecx]
cevakrn1.xmd:03690855      movzx   ecx, [ebp+pout] ; word[ecx+0x3C]
cevakrn1.xmd:03690859      add     eax, 78h
cevakrn1.xmd:0369085C      add     eax, ecx            ; ecx + 0x78 + word[ecx+0x3C]
cevakrn1.xmd:0369085E      push    0                  ; int
cevakrn1.xmd:03690860      push    eax                ; addr
cevakrn1.xmd:03690861      lea     eax, [ebp+output]
cevakrn1.xmd:03690864      push    eax                ; output
cevakrn1.xmd:03690865      push    dword ptr [edi+0Ch] ; int
cevakrn1.xmd:03690868      call    get_virtual_dword
cevakrn1.xmd:0369086D      add     esp, 10h
cevakrn1.xmd:03690870      test    eax, eax
cevakrn1.xmd:03690872      jnz     loc_3690C63
cevakrn1.xmd:03690878      mov     edx, [ebp+output] ; dword[ecx + 0x78 + word[ecx+0x3C]]
cevakrn1.xmd:0369087B      lea     esi, [ebp+var_2C]
cevakrn1.xmd:0369087E      mov     ecx, [edi+0Ch]
cevakrn1.xmd:03690881      add     edx, [ebp+readecx] ; ecx + dword[ecx + 0x78 +
word[ecx+0x3C]] =
cevakrn1.xmd:03690881                                     ; emulated esi

```

对dword (N) 乘以4的时候，没有进行整数边界检查，然后调用了malloc()。这样的话，将导致分配的缓冲区过小。

```

cevakrn1.xmd:03690884      push    esi
cevakrn1.xmd:03690885      push    0
cevakrn1.xmd:03690887      push    28h
cevakrn1.xmd:03690889      mov     [ebp+output], edx
cevakrn1.xmd:0369088C      mov     eax, [ecx]
cevakrn1.xmd:0369088E      push    0
cevakrn1.xmd:03690890      push    edx
cevakrn1.xmd:03690891      call    dword ptr [eax+0C0h]
cevakrn1.xmd:03690897      test    eax, eax
cevakrn1.xmd:03690899      jnz     loc_3690C63
cevakrn1.xmd:0369089F      mov     eax, dword ptr [ebp+structure+18h] ; at offset 0x18

```

接下来进入一个循环，在该循环中，将会把通过位于仿真堆栈中的偏移量找到的N个字符串的CRC32校验和填充到这个缓冲区中。当字符串的偏移量过大的时候，循环将会中

```

cevakrn1.xmd:036908A6      shl     eax, 2              ; integer overflow
cevakrn1.xmd:036908A6                                     ; 0x40 00 00 00
cevakrn1.xmd:036908A9      push    eax
cevakrn1.xmd:036908AA      call    near ptr 15D08D0h ; malloc
cevakrn1.xmd:036908AF      mov     ebx, eax
cevakrn1.xmd:036908B1      add     esp, 4
cevakrn1.xmd:036908B4      mov     [ebp+buffer], ebx

```

这提供了执行代码时所需的一切。首先，我们需要能够用自己的内容覆盖任意数量的字节。这就需要用到一些其CRC值符合特定要求的字符串。当然，这里的CRC32算法是

```

cevakrn1.xmd:036908C6 @fillbuffer:          ; CODE XREF: sub_36906D0+245j
cevakrn1.xmd:036908C6      mov     eax, dword ptr [ebp+structure+20h] ; at offset 0x18 +
8                                     ; (8 bytes after large num)
cevakrn1.xmd:036908C6                                     ; offset to a table of 32-bit offsets
cevakrn1.xmd:036908C6                                     ; to data
cevakrn1.xmd:036908C9      push    0 ; int
cevakrn1.xmd:036908CB      lea     eax, [eax+esi*4]
cevakrn1.xmd:036908CE      add     eax, [ebp+readecx]
cevakrn1.xmd:036908D1      push    eax ; addr
cevakrn1.xmd:036908D2      lea     eax, [ebp+dataoffset]
cevakrn1.xmd:036908D5      push    eax ; output
cevakrn1.xmd:036908D6      push    dword ptr [edi+0Ch] ; int
cevakrn1.xmd:036908D9      call    get_virtual_dword
cevakrn1.xmd:036908DE      mov     ebx, eax
cevakrn1.xmd:036908E0      add     esp, 10h
cevakrn1.xmd:036908E3      test    ebx, ebx
cevakrn1.xmd:036908E5      jnz     @error ; in case of error 32-bit zero is
written
cevakrn1.xmd:036908EB      mov     eax, [ebp+buffer]
cevakrn1.xmd:036908EE      lea     eax, [eax+esi*4]
cevakrn1.xmd:036908F1      push    eax
cevakrn1.xmd:036908F2      mov     eax, [ebp+dataoffset]
cevakrn1.xmd:036908F5      add     eax, [ebp+readecx]
cevakrn1.xmd:036908F8      push    0EDB88320h ; xorval
cevakrn1.xmd:036908FD      push    eax ; srcvirt
cevakrn1.xmd:036908FE      push    edi
cevakrn1.xmd:036908FF      call    crc32b ; unsigned int crc32b(unsigned char
*message) {
cevakrn1.xmd:036908FF      ; int i, j;
cevakrn1.xmd:036908FF      ; unsigned int byte, crc, mask;
cevakrn1.xmd:036908FF      ;
cevakrn1.xmd:036908FF      ; i = 0;
cevakrn1.xmd:036908FF      ; crc = 0xFFFFFFFF;
cevakrn1.xmd:036908FF      ; while (message[i] != 0) {
cevakrn1.xmd:036908FF      ;     byte = message[i];
cevakrn1.xmd:036908FF      ;
cevakrn1.xmd:036908FF      // Get next byte.
cevakrn1.xmd:036908FF      ;     crc = crc ^ byte;
cevakrn1.xmd:036908FF      ;     for (j = 7; j >= 0; j--) {
cevakrn1.xmd:036908FF      ;         mask = -(crc & 1);
cevakrn1.xmd:036908FF      ;         crc = (crc >> 1) ^
(0xEDB88320 & mask);
cevakrn1.xmd:036908FF      ;     }
cevakrn1.xmd:036908FF      ;     i = i + 1;
cevakrn1.xmd:036908FF      ; }
cevakrn1.xmd:036908FF      ; return ~crc;
cevakrn1.xmd:036908FF      ; }
cevakrn1.xmd:03690904      mov     ebx, eax
cevakrn1.xmd:03690906      add     esp, 10h
cevakrn1.xmd:03690909      test    ebx, ebx
cevakrn1.xmd:0369090B      jnz     @error
cevakrn1.xmd:03690911      inc     esi
cevakrn1.xmd:03690912      cmp     esi, dword ptr [ebp+structure+18h]
cevakrn1.xmd:03690915      jb      short @fillbuffer

```

当匹配的函数被调用时，漏洞被触发，即利用任意内容覆盖特定数量的内存空间。

小结

通过深入分析，我们可以看到，攻击者可以含有整数溢出漏洞的Bitdefender版本上执行任意代码。目前，软件供应商已经通过更新修复了该漏洞（以及其他漏洞）。因此，

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