Lab 7: Graph Search Functions

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This paper aims to demonstrate implementations of a disjoint set forest and graph search functionality to solve a randomly produced maze. The three different types of graph search function utilized and analyzed in this lab is the Breath First implementation, Depth First based on stacks instead of queses, and a recursive Depth First Search. By timing the performance of each function, a conclusion is drawn on which of these tree implementations is the proffered method.

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I. Introduction

A graph is a set of vertices with a specific amount of edges which connects particular vertices together. It can be represented in three different forms: Matrix Form, Adjacency List, and an Edge Lists. All three methods capitalizes on existing edges, thus representing connected vertices. This data structure is commonly found when representing individual objects or items with existing relations between them. A notable example is how a friends list from an individual will be represented on a broader scale.

In this lab, we will implement and demonstrate the three different approaches of search function through the graph. The maze will be randomly generated using a disjoint set forest approach which is similar to the previous lab. The graph built for this lab will be based on the number of cells within the maze. Each cell represents a vertices within the graph. Each edge which connects the vertices represents a path between two cells which is not obstructed by an existing wall. The three search functions will solve the maze and compared according to time performances. The resulting path will be printed along the maze to verify the results.

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II. Design

This lab is based upon constructing a maze with a given set forest and solving it using a graph. The size of the graph depends on the number of cells within the maze. The size of the maze and the number of walls to be removed is up for user discretion, but it also affects of the number of paths, if they exits, that the program provides. If the number of wall to be removed is less than the number cells -1, then there is no guaranteed path form the cell 0 to cell n-1. If the number of walls to removed is exactly n-1, then there is a guaranteed unique path. Anything over produces more than one path. Of course, the basic design of the maze is based upon the previous lab.

A graph is then produced based on the exiting lists containing the wall within the maze. It is built in adjacency form and marks edges between vertices given that it is not located within the walls lists. With the graph produces, different methods can be utilized in order to solve the maze creating a path. The following parts showcases three different methods.

A. Part 1: Breadth First Search (BFS)

The Breadth First Search algorithm utilizes queues to transverses through the graph. The design follows similar approach to any to the basic BFS functions. This methods takes $\theta(-v-+E-)$ and is predicted to share the same execution time as the DFS function.

B. Parts 2: Depth First Search Function (DFS) - Stack

The Depth First Search algorithm utilizes the stack method in order to transverse through the graph. The design follows the same logic as in part 1, except this time the search algorithm searches the path in question until it reaches the final cell or cannot go further due to obstructed walls. In terms of time performance, it is predicted they it will have the same, if not similar outcomes at in part 1 since the time complexity of this function is $\theta(-v-+E-)$.

C. Parts 3: Depth First Search Function (DFS) - Recursion

The recursion approach to the depth first function was a bit more challenging than the previous parts. In this method although, the function passes the array in which contains the path cells at each recursions call. This enables to only append to the specific list only when the path is found. The resulting list will only contain the cell numbers which the final path takes effectively making the drawing of the path easier. This is why the maze is only drawn with this approach.

III. Results

After the implementations of both disjoint set forest methods and all three graph search functions were successful, each functions were executed and recorded for time performance. The time values were used to evaluate each program and compares each data structure with one another. Below are their time executions and over all statistical results.

Figure 1: Standard Maze

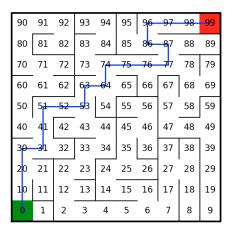
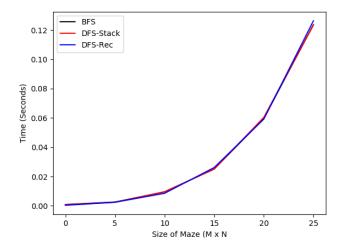


Figure 2: Run Times



IV. Discussion

Overall, implementing all three methods to to search a given constructed maze results in a successful path computation. There is a a small difference between all there methods, but the recursion method turned out to take the most time. From Figure 3, we can conclude that the discrepancy that both BFS and DFS do share the same time performance. A side note that must be written along the results is that the time results do not include the time it takes to draw the resulting path within the maze.

V. Source Code

```
1 """
<sup>2</sup> Author: Steven J. Robles
3 Class: CS 2302 Data Structures III
4 Instructor: Olac Fuentes
5 TA: Anindita Nath And Maliheh Zargaran
6 Last Modified: 04/29/2019
7 Discreption: Lab 7:
           This program is desinged to act as the main file of this project. It produces the
      main menu and recieves
           user input to call functions. It also records the time it takes for each maze to be
      built.
11
  ,, ,, ,,
12
13
14 from BuildMaze import initiateMazeC
15 import matplotlib.pyplot as plt
  import numpy as np
17 import timeit
18
20 loop = True #commences the loop
21
22 #the is the while loop which pompts the user.
while loop:
    print ("1. Build and Graph Maze\n2. Time Trail")
24
    number = input("3. Exit \n")
25
    print("***************")
    #trys converting the input into an int. if it fails, the pormpt runs again
27
28
    try:
29
      choice = int(number)
    except:
30
31
      choice = -1
32
    #The fist if statement bulds the maze by asking of its dimensions first
33
    if choice == 1:
34
      cont = True
35
36
       while cont:
         rows = input ("Enter value for rows : \n")
37
         cols = input ("Enter value for columns : \n")
38
         remove = input ("Enter number of rows to be removed : \n")
39
40
        #the following converts the input into ints if it's possible
41
        try:
          rows = int (rows)
42
43
           cols = int(cols)
           remove = int(remove)
44
           if remove < 0:
45
            print("Try Again")
46
           else :
47
            cont = False
48
49
50
         except:
          print("Try Again")
51
52
       print()
       if remove < rows*cols -1:
         print ('A path from source to destination is not guaranteed to exist')
54
55
       elif remove = rows*cols -1:
         print('The is a unique path from source to destination')
56
57
         initiateMazeC (rows, cols, start, end, True)
58
        print('There is at least one path from source to destination')
59
60
    #Choice number two times the preformance of the functions with a determined
61
    #set of sisze mazes.
62
    elif choice == 2:
63
       dimensions = [[5,5],[10,10],[15,15],[20,20],[25,25],[30,30]]
64
      times = [[],[],[]]
```

```
for i in range(len(dimensions)):
 66
                   #times the BFS
 67
                   start = timeit.default_timer() # starts timer
 68
                   initiate Maze C \, (\, dimensions \, [\, i\, ] \, [\, 0\, ] \,\,, \,\,\, dimensions \, [\, i\, ] \, [\, 1\, ] \,\,, \,\,\, 0\,, \,\,\, (\, dimensions \, [\, i\, ] \, [\, 0\, ] \, * \, dimensions \, [\, i\, ] \, [\, 0\, ] \,
 69
               [1] -1, False, 1, True
                   stop = timeit.default_timer() # ends timers
 70
 71
                   times [0]. append (stop-start)
                   #times the DFS _ Stack
 72
                   start = timeit.default_timer() # starts timer
 73
                   initiate Maze C (dimensions [i][0] \;,\; dimensions [i][1] \;,\; 0 \;,\; (dimensions [i][0] * dimensions [i][0] * dimensions [i][0] \;,\; dimensi
 74
               ][1])-1, False, 2, True)
                   stop = timeit.default_timer() # ends timers
 75
                   times [1].append(stop-start)
 76
                   #times hte DFs - Recursion
 77
                   start = timeit.default_timer() # starts timer
 78
                   initiate Maze C (dimensions [i][0], \ dimensions [i][1], \ 0, \ (dimensions [i][0]* dimensions [i][0] \\
 79
               [1]) -1, False, 3, True)
                   stop = timeit.default_timer() # ends timers
 80
                   times [2].append(stop-start)
 81
 82
               fig, ax = plt.subplots()
 83
 84
                   #proceeds to plot the time results
               plt.xlabel('Size of Maze (M x N')
plt.ylabel('Time (Seconds)')
 85
 86
              x = np.arange(0,30,5)
 87
              plt.plot(x, times[0], 'k', label= 'BFS')
plt.plot(x, times[1], 'r', label='DFS-Stack')
plt.plot(x, times[2], 'b', label='DFS-Rec')
 88
 89
 90
               plt.legend()
 91
               plt.savefig('RunTimes')
 92
              plt.show()
 93
 94
          #program exits
           elif choice == 3:
 95
               print("Good Bye!")
 96
               loop = False
 97
 98
               print("Try Again")
 99
          100
  1
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  3 Class: CS 2302 Data Structures III
  4 Instructor: Olac Fuentes
  5 TA: Anindita Nath And Maliheh Zargaran
  6 Last Modified: 04/29/2019
  7 Discreption: Lab 7:
                       The purpose of this program is to build construct a adjecancy lists for a graph
              based on the
                       maze matrix. This graph will be utilized to solve the mazes in three different
               approcahes:
                       BFS, DFS - Stacked, and DFS - Recursive
 10
 11
      ,, ,, ,,
 12
 13 import math
 14
 15 #bulds the graph giaven the walls lists
      def buildGraph (walls, rows, columns):
 16
          G = [[] for i in range(rows*columns)]
 17
          v = 0
 18
          #checks if there is a wall between to cells before making an edge
 19
 20
          while v <rows*columns:
               if (v + 1) %columns != 0 and [v, v+1] not in walls and [v+1, v] not in walls:
 21
 22
                  G[v]. append (v+1)
               if v %columns != 0 and [v, v-1] not in walls and [v-1, v] not in walls:
 23
 24
                  G[v]. append (v-1)
               if v-columns >= 0 and [v, v-columns] not in walls and [v-columns, v] not in walls:
 25
                  G[v].append(v-columns)
 26
               if v +columns < rows*columns and [v, v+columns] not in walls and [v+columns, v] not in
 27
```

walls:

G[v]. append (v+columns)

```
v+=1
    return G
30
31
32 #BFS method approach
  def BFS(G, v):
    Q = []
34
35
    visited = [False]*(len(G))
    prev = [-1] * (len(G))
36
    Q. append (v)
37
    visited\,[\,v\,]\,\,=\,\, True
38
     while len(Q) > 0:
39
      u = Q. pop(0)
40
      for t in G[u]:
41
         if visited [t]== False:
42
           visited[t] = True
43
           prev[t] = u
44
45
           Q. append(t)
    return prev
46
47
  #DFS stack approach
48
  def DFS_Stack(G, v):
49
    Q = []
    visited = [False]*(len(G))
51
    prev = [-1] * (len(G))
52
    Q. append (v)
53
    visited[v] = True
54
    while len(Q) > 0:
55
56
      u = Q. pop()
      for t in G[u]:
57
        if visited [t]== False:
58
           visited[t] = True
59
60
           prev[t] = u
           Q. append(t)
61
62
    return prev
64 #DFS recursion approach
def DFS_Rec(G, source, visited, prev, final):
66
    visited [source] = True
67
    if source == final:
      return prev, True
68
     for t in G[source]:
69
      if visited[t] == False:
70
        prev, add = DFS_Rec(G, t, visited, prev, final)
71
         if add:
72
           prev.append(t)
73
74
           return prev, True
    return prev, False
75
76
77 #builds the graph in an adjecancy list for the other methods
def mazeGraph (walls, rows, columns):
    G = buildGraph (walls, rows, columns)
   return G
80
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  Discreption: Lab 7:
          The purpose of this program is to build a maze with any given number of cells and
      make them all
           one complex cell. This will is done by created a disjoint set forest making all of
      the single cells
           belong to one tree. Resutls are shown and the preformances are timed. A solution
       path is also calculated
           by sovling it int terms of grapsh using BSF and DSF methods. They are ploted wihtin
11
      the maze.
12
13 """
```

```
15 #from Graph import DFS
16 from Graph import mazeGraph
17 from Graph import DFS_Rec
  from Graph import DFS_Stack
  from Graph import BFS
  import matplotlib.pyplot as plt
  import numpy as np
21
  import random
22
24 #the following function is the intial call for the iterative function which solves the maze
  def solveMaze(walls, rows, columns, start, finish, trackNum = 0):
26
       G = maxeGraph(walls, rows, columns)
27
       print ("Adjectancy List: \n", G)
28
       print ('Final Path For Breath Frist Search: \n', BFS(G, 0))
29
       print('Final Path For Detph First Search - Stack: \n', DFS_Stack(G, 0))
30
       final Path, dud = DFS\_Rec(G, 0, [False]*(len(G)), [], (rows*columns) -1)
31
32
       finalPath.append(0)
       print ('Final Path For Detph First Search - Recursion: \n', finalPath)
33
34
       convertedPath = []
35
       #the solution path are convereted to x and y coordinates to plot the blue lines for the
       path
       for i in range(len(finalPath)):
36
            if i > 0:
37
                c1 = finalPath[i] %columns
38
                r1 = (finalPath[i]-c1) / columns
39
                c2 = finalPath [i-1]\% columns
40
                r2 = (finalPath[i-1]-c2) / columns
41
                convertedPath.append([[c1 + .5, c2 + .5], [r1 + .5, r2 + .5]])
42
43
       return convertedPath
44
45
  #plots the maze as provided in class
  def draw_maze(walls, maze_rows, maze_cols, trackNum=0, start=0, finish=0, cell_nums=True):
47
       finalPath = solveMaze(walls, maze_rows, maze_cols, start, finish, trackNum)
       fig , ax = plt.subplots()
49
50
       for w in walls:
51
            if w[1]-w[0] ==1: #vertical wall
                x0 = (w[1]\% maze\_cols)
                x1 = x0
53
                y0 = (w[1]//maze\_cols)
55
                y1 = y0+1
            else:#horizontal wall
56
                x0 = (w[0]\% maze\_cols)
57
                x1 = x0+1
58
                y0 = (w[1]//maze\_cols)
60
                y1 = y0
            ax.plot([x0,x1],[y0,y1],linewidth=1,color='k')
61
62
       sx = maze\_cols
       sy = maze_rows
63
       ax.\,plot\,(\,[\,0\;,0\;,sx\;,sx\;,0\,]\;,[\,0\;,sy\;,sy\;,0\;,0\,]\;,\\ linewidt\,h\,=\,2,color\,=\,\,\,'\,k\,\,'\,)
64
65
       c1 = start%maze_cols
       r1 = (start-c1) / maze\_cols
66
67
       c2 = finish\%maze\_cols
       r2 = (finish - c2) / maze\_cols
68
       #fills the start and end cells with green adn red coolors
69
       \begin{array}{l} \text{ax. fill ([c1,\ c1+1,\ c1+1,\ c1,\ c1],\ [r1,\ r1,\ r1+1,\ r1+1,\ r1],\ color = 'g')} \\ \text{ax. fill ([c2,\ c2+1,\ c2+1,\ c2,\ c2],\ [r2,\ r2,\ r2+1,\ r2+1,\ r2],\ color = 'r')} \end{array}
70
71
       ax.axis('off')
72
73
       ax.set_aspect(1.0)
       for i in range(len(finalPath)): #plots the solution path
74
            ax.plot(finalPath[i][0], finalPath[i][1], color = 'b')
75
       if cell_nums:
76
            for r in range(maze_rows):
77
                for c in range(maze_cols):
78
                     cell = c + r*maze\_cols
79
80
                     ax.text((c+.5),(r+.5), str(cell), size=10,
                              ha="center", va="center")
81
       ax.axis('off')
```

```
83
        ax.set_aspect(1.0)
        plt.show()
84
85
86 #bulds the wall lists as porvided
   def wall_list(maze_rows, maze_cols):
87
        # Creates a list with all the walls in the maze
88
        w = []
89
        for r in range(maze_rows):
90
            for c in range(maze_cols):
91
                 cell = c + r*maze\_cols
92
                 if c!=maze\_cols-1:
93
                     w.append([cell,cell+1])
94
                 if r!=maze\_rows-1:
95
                     w.append([cell,cell+maze_cols])
96
97
        return w
98
99 #initializes the disjoint set forest list
   def DisjointSetForest(size):
100
        return np. zeros (size, dtype=np.int)-1
101
   #comppressed find fucntion where all of the cells point
103
   #directty to its root
   def findC(S,i):
        if S[i] < 0:
106
            return i
107
        S[i] = findC(S, S[i])
108
109
        return S[i]
   #joins trees together depending on thier existing sizes
   def unionBySize(S, i, j):
112
        ri = findC(S, i)
113
114
        rj = findC(S, j)
        if ri!= rj:
            if S[ri] < S[rj]:
116
                 S[ri] += S[rj]
117
                 S[rj] = ri
118
            else:
119
120
                 S[rj] += S[ri]
                 S[ri] = rj
   #returns a boolean if two cells belong to the same tree or not for compressed
   def checkPathC(S, walls, d):
124
        \begin{array}{ll} if \ \ findC(S, \ walls [d][0]) == findC(S, \ walls [d][1]): \\ return \ \ False \end{array}
125
126
        return True
127
128
#buids a maze based upon the compressed method
   def initiateMazeC(maze_rows, maze_cols, start=0, end=0 ,draw=False, trackNum = 0, time =
130
        False):
        #initializes the components to build the mazw
        walls = wall_list (maze_rows, maze_cols)
        T = DisjointSetForest (maze_rows * maze_cols)
134
        setCount2 = len(T)
        #randomly removes a wall until all cells are connected
136
        while setCount2 > 1:
            d = random.randint(0, len(walls)-1)
137
            if checkPathC(T, walls, d):
138
                 unionBySize\left(T,\ walls\left[d\right]\left[0\right],\ walls\left[d\right]\left[1\right]\right)
139
                 walls.pop(d)
140
                 setCount2 -=1
141
        if draw:
142
            draw_maze(walls, maze_rows, maze_cols, trackNum, start, end)
143
144
        if time:
            G = mazeGraph(walls, maze_rows, maze_cols)
145
            if trackNum == 1:
146
                 BFS(G, 0)
147
148
                 return
            if trackNum == 2:
149
                 DFS_Stack(G, 0)
```

```
if trackNum == 3:

DFS_Rec(G, 0, [False]*(len(G)), [], (maze_rows*maze_cols) -1)

return
```

VI. Academic Dishonesty

Scholastic Dishonesty

Any student who commits an act of scholastic dishonesty is subject to discipline. Scholastic dishonesty includes, but not limited to cheating, plagiarism, collusion, the submission for credit of any work or materials that are attributable to another person.

Cheating

- Copying form the test paper of another student
- o Communicating with another student during a test
- o Giving or seeking aid from another student during a test
- o Possession and/or use of unauthorized materials during tests (i.e. Crib notes, class notes, books, etc)
- Substituting for another person to take a test
- o Falsifying research data, reports, academic work offered for credit

• Plagiarism

- Using someone's work in your assignments without the proper citations
- Submitting the same paper or assignment from a different course, without direct permission of instructors

Collusion

o Unauthorized collaboration with another person in preparing academic assignments

Sign:	The Molle	Date:	04/30/2019
·			