



University of Pittsburgh

# CS 1541 Introduction

Instructor Introduction

Course Introduction

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Department of Computer Science

School of Computing and Information





# *Instructor Introduction*



# My Technical Background

## ■ Wonsun Ahn

- First name is pronounced *one-sun* (if you can manage)
- Or you can just call me Dr. Ahn (rhymes with *naan*)

## ■ PhD in CPU Design and Compilers

- University of Illinois at Urbana Champaign

## ■ Industry Experience

- Bluebird Corporation (70-person startup company)
  - Manufactures industrial hand-held devices
  - Me: Built software stack based on Windows Embedded
- IBM Research (thousands of people)
  - Does next-gen stuff like carbon nanotubes, quantum computers
  - Me: Designed supercomputers for ease of parallel programming



# My World View

## ■ Everything is connected

- Pandemic: If my neighbors catch the virus, so will I
- Environment: If my neighbors pollute, I will feel the effects
- Economy: Think of how the subprime mortgage crisis spread

## ■ Zero-sum thinking (old way of thinking)

- “If you get a larger slice of the pie, I get a smaller slice.”
- Therefore, if you lose, I win (and vice versa)

## ■ Zero-sum thinking no longer works

- If you catch the virus, do I become safer from the virus?

## ■ Collaboration is replacing competition



# Collaboration is Replacing Competition

- Is happening in all spheres of life
- Collaboration is also happening in the IT industry
  - The *open source* movement
  - Increasing importance of the software/hardware *ecosystem*
  - Increasing importance of the developer *community*
- Collaboration is also important for learning
  - During my undergrad years, what do I remember best?
  - Stuff that I explained to my classmates
  - Stuff that my classmates taught me



# Supporting Collaborative Learning

- *I do not grade on a curve*
  - You will not be competing against your classmates
  - You are graded on your own work on an absolute scale
  
- You are a member of a Team
  - You are already a member of the class on Microsoft Teams
  - I encourage you to be on Teams at most times (I will too)
    - You can install app on both laptop and cell phone
  - If you have a question, you can ask in the Team “Posts” tab
    - Either your classmate or your instructor will answer
  - You can chat with any individual on the Team
    - “Manage Team” item in the “...” Team context menu



# Supporting Collaborative Learning

## ■ You are a member of a Study Group

1. Right now, form a group of 4~6 members with neighbors
2. Create a Teams chat group comprised of these members
3. Invite your instructor to chat group (ID: wahn@pitt.edu)
4. Send at least one message in chat group so I can see it
  - Can discuss TopHat questions/answers before submitting
  - Goal: develop a basis of knowledge for homeworks/exams

## ■ You are a member of a Project Group

- Find one partner with whom you want to do projects
- On GradeScope, submit to “Partnership Contract”  
(Please do one submission per group by adding partner)



# *Course Introduction*



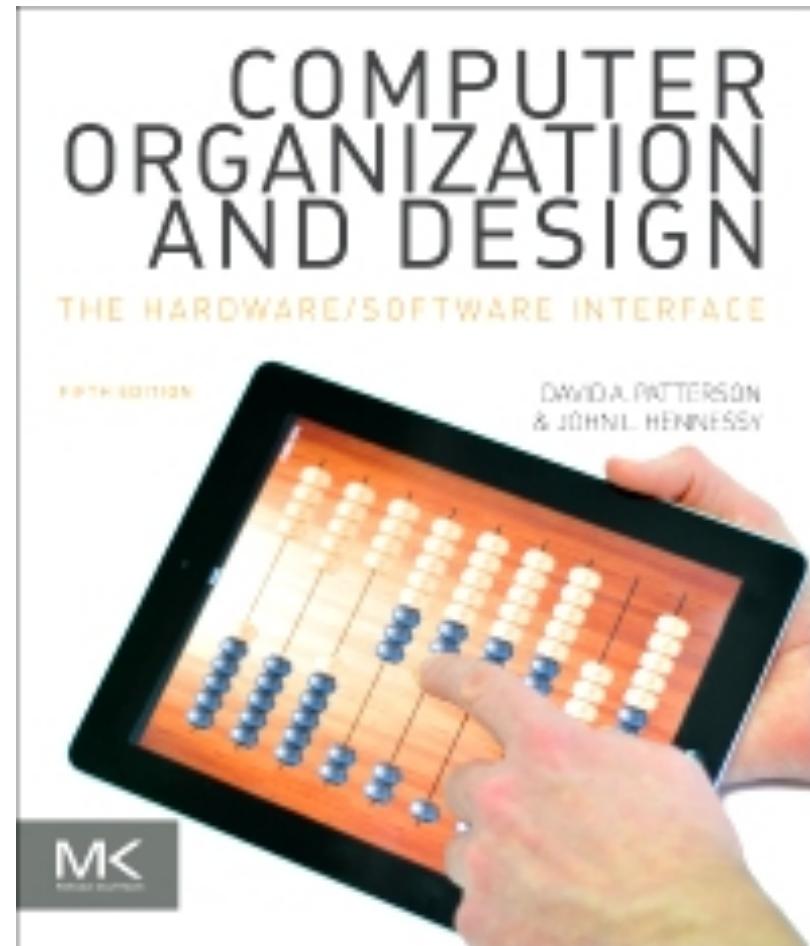
# Structure of the Course

- (45% of grade) Two midterms
- (20% of grade) Two projects
  - Implementing a CPU simulator using C programming language
- (20% of grade) Four homeworks
- (15% of grade) Participation
  - Attendance, TopHat lecture questions, Teams participation
- Class resources:
  - Canvas: announcements, Zoom meetings, recorded lectures
  - GitHub: syllabus, lectures, homeworks, projects
  - Tophat: online lecture questions
  - GradeScope: homework / projects submission, grading and feedback
  - Microsoft Teams: all out-of-class communication



# Textbook (You Probably Have it)

- “Computer Organization and Design - The Hardware/Software Interface” by David Patterson and John Hennessy Fifth Edition - Morgan & Kaufmann.





# For More Details

- Please refer to the syllabus page:

[https://github.com/wonsunahn/CS1541\\_Spring2023  
/blob/main/syllabus.md](https://github.com/wonsunahn/CS1541_Spring2023/blob/main/syllabus.md)

- Please follow the course schedule:

[https://github.com/wonsunahn/CS1541\\_Spring2023  
/blob/main/schedule.md](https://github.com/wonsunahn/CS1541_Spring2023/blob/main/schedule.md)

# Computer Architecture.

## What is it?

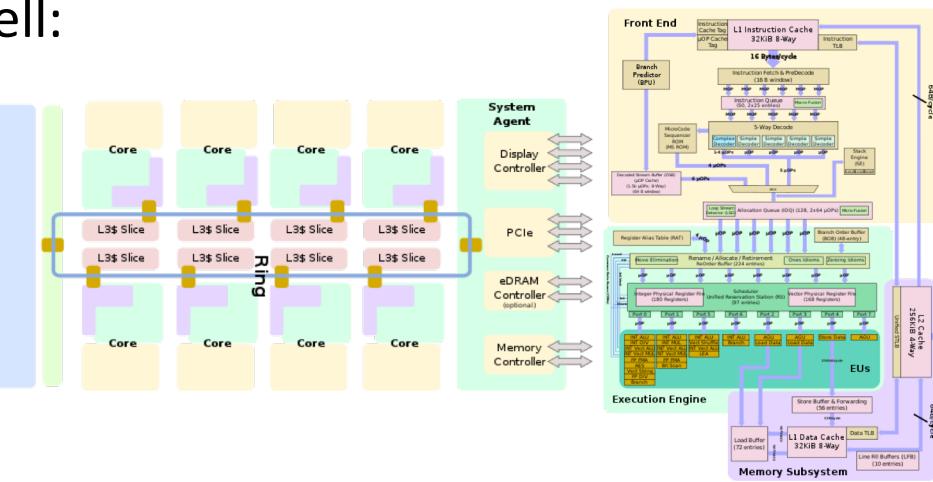
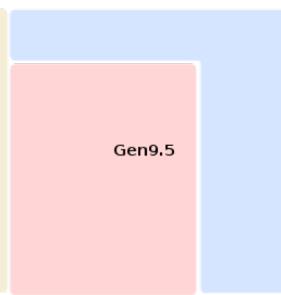
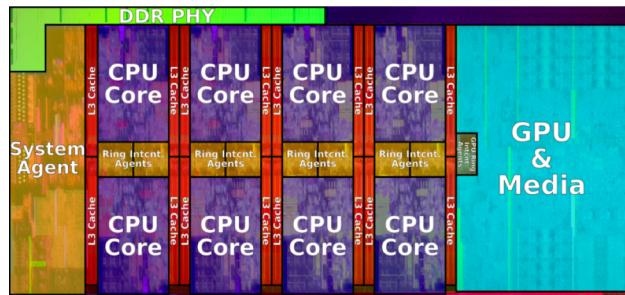
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# Computer Architecture: the architecture of a CPU

- Every building has an architecture:



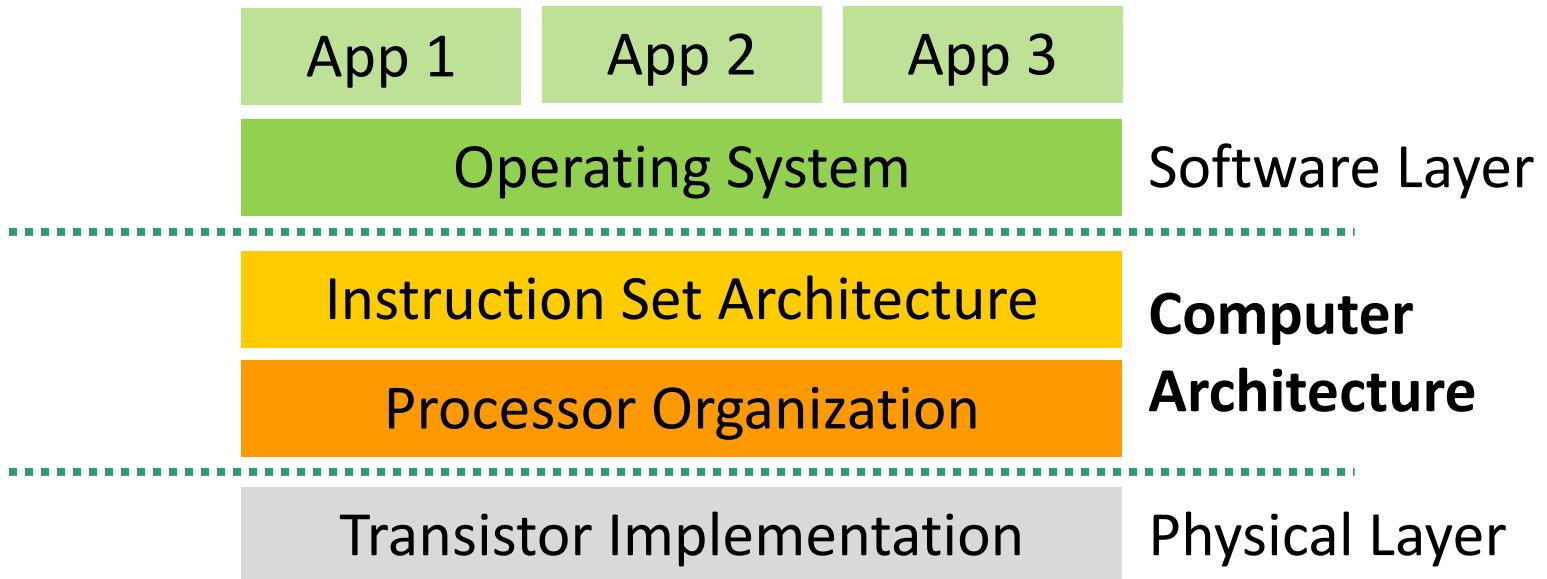
- Every CPU has an architecture as well:



# ISA: the architecture of an instruction set

- At a high-level: how a computer is built
  - Computer here meaning the processor (CPU)
- You probably heard of a similar term before: ISA
  - ISA (Instruction Set Architecture)
- Review: what is defined by an ISA?
  - *Set of instructions* usable by the computer
  - *Set of registers* available in the computer
  - Other functional attributes
- What is *not* defined by an ISA?
  - *Speed* of computer
  - *Energy efficiency* of computer
  - *Reliability* of computer
  - Other performance attributes

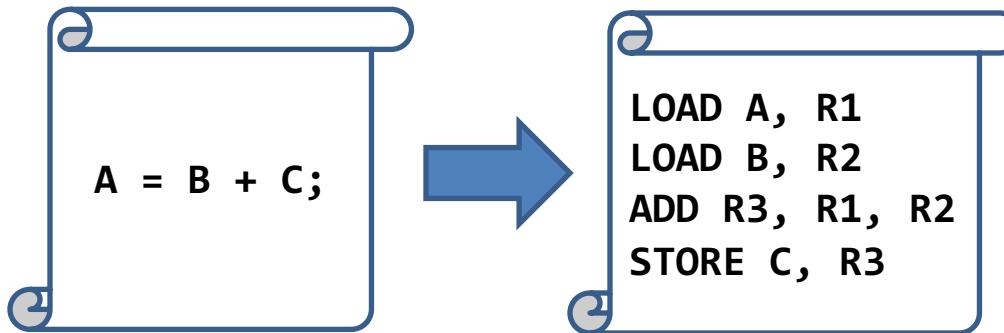
# Computer Architecture = ISA+Processor Organization



- **Computer Architecture = ISA + Processor Organization**
  - Processor organization is also called *Microarchitecture*
- Given an ISA, performance is determined by:
  - Processor organization (internal design of the processor)
  - Transistor implementation (semiconductor technology)

# A CPU is like a factory

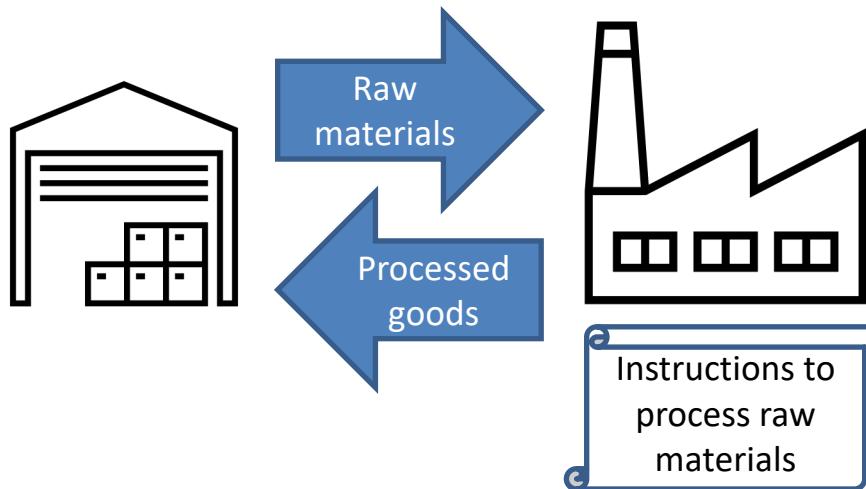
- A program written in Python, Java, or C is converted to instructions



- A CPU processes data according to instructions, just like a factory



# Factory efficiency depends on factory design

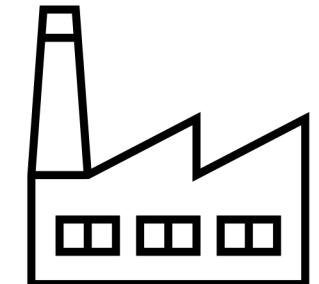
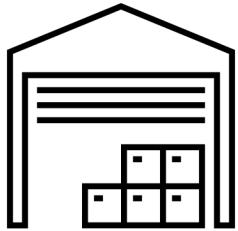


1. How efficiently goods are moved between warehouse and factory
2. How efficiently workers process goods in factory

# How to design a shirt factory

## Part 1: Logistics

# You cannot succeed without good logistics

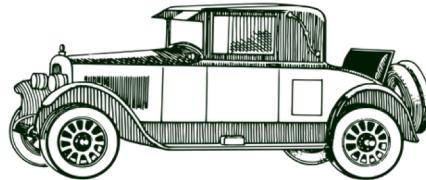


Instructions to process raw materials

- If you can't move goods in a timely manner your output will suffer
- Factory efficiency doesn't matter if logistics is the bottleneck
  - If deliveries are slow (high memory latency)
  - Or do not arrive in sufficient quantities (low memory bandwidth)

# Memory Latency vs. Memory Bandwidth

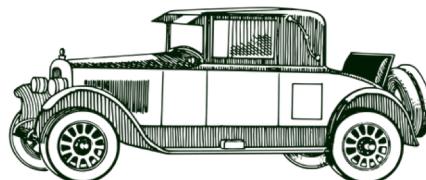
- Memory bottleneck comes from two sources:
  - **Memory latency**: hours to handle a single delivery



vs.



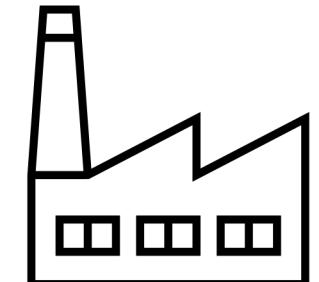
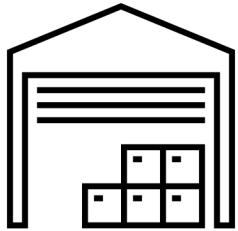
- **Memory bandwidth**: maximum volume handled per hour



vs.



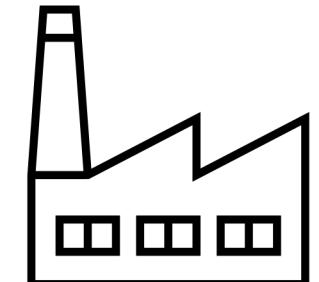
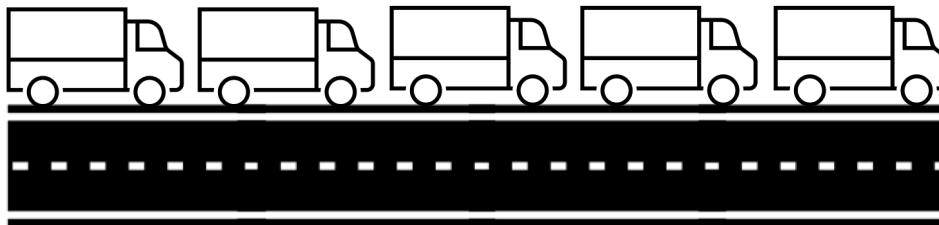
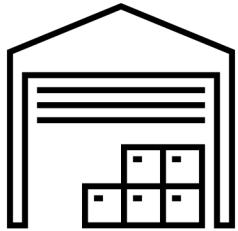
# Memory Latency vs. Memory Bandwidth



Instructions to process raw materials

- Race cars are rarely used for logistics. Trucks are. Why?
  1. Can place orders ahead of time to arrive just in time (*prefetching*)
  2. Or the factory can produce something else while waiting (*scheduling*)
- There are ways to circumvent latency so it does not impact performance.
  - Except when neither prefetching or scheduling works

# Memory bandwidth puts a cap on performance



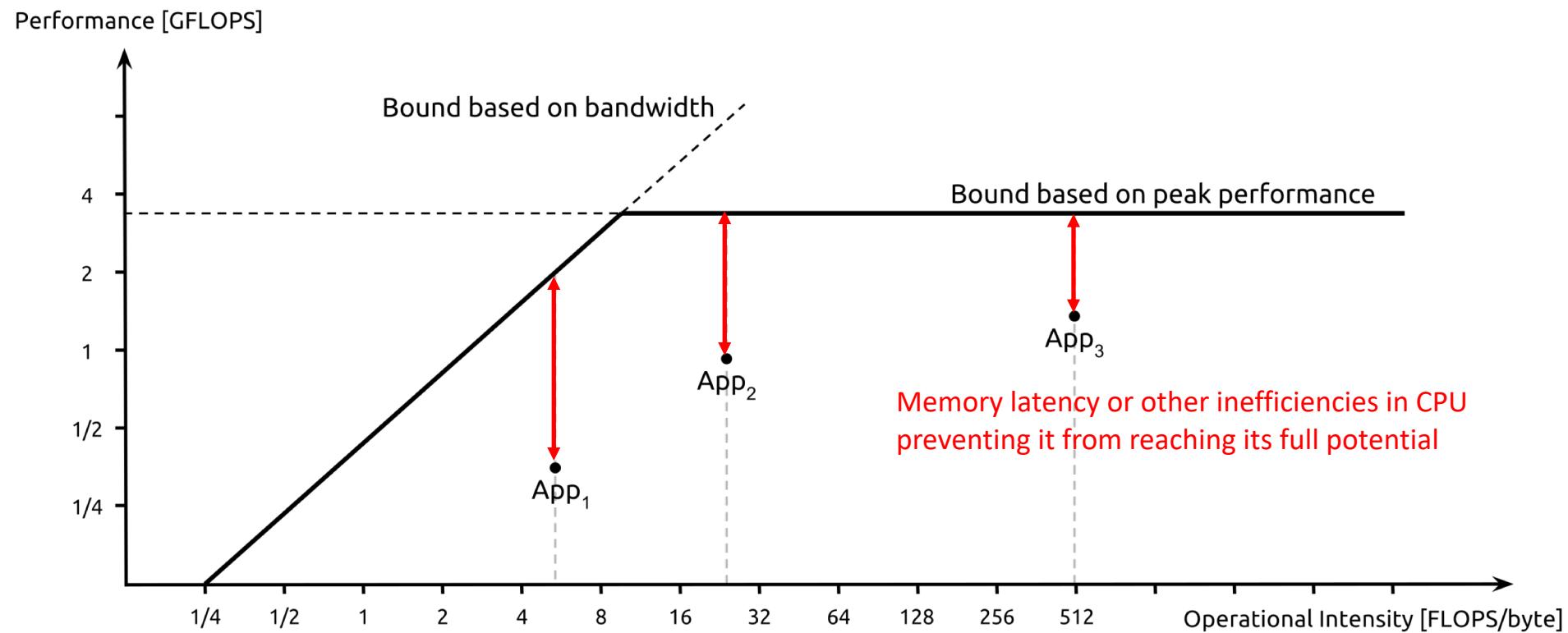
- If the road is fully utilized but still can't keep up with demand...
  - Factory needs 5 trucks / minute but road handles 1 truck / minute
- No amount of pre-ordering or scheduling can alleviate this problem  
→ Unlike latency, bandwidth puts a hard limit on performance!

# Memory bandwidth puts a cap on performance

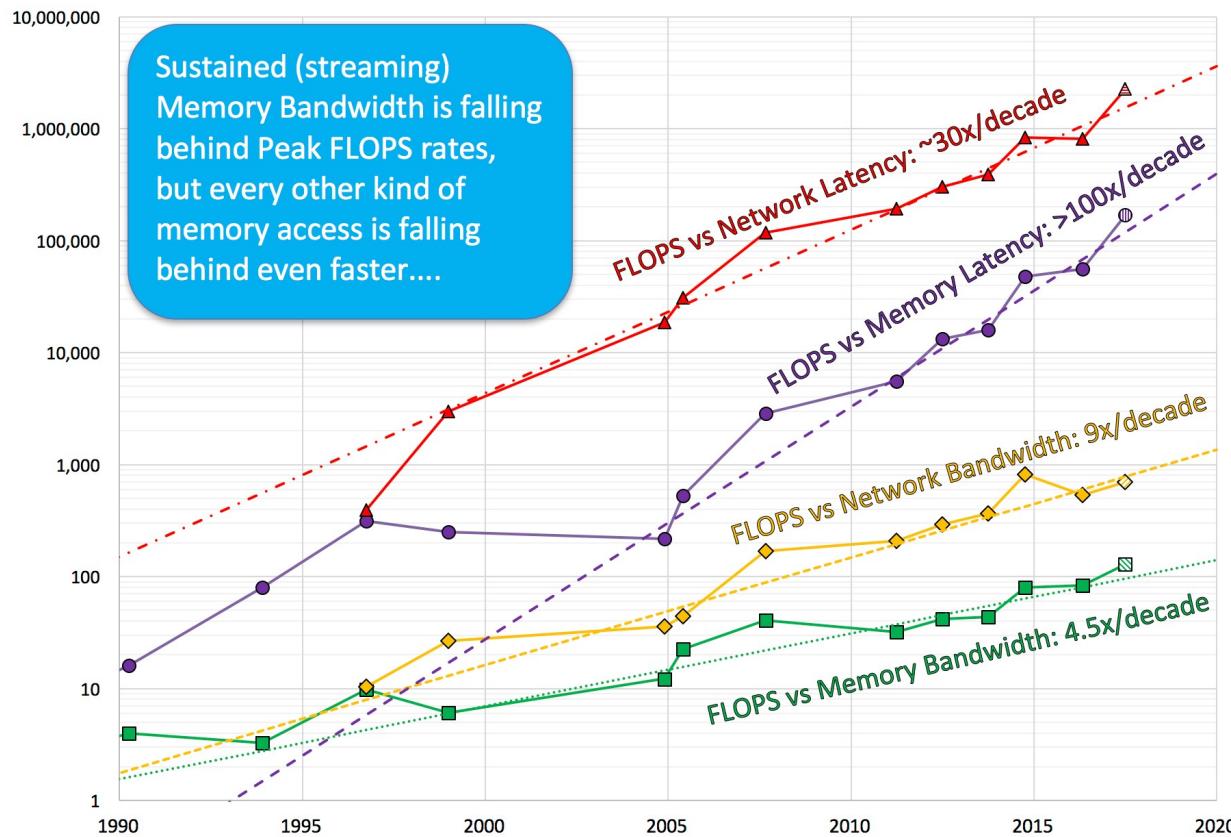
- When you get a traffic jam, performance is dictated by bandwidth
    - How quickly you pull data in, rather than how fast you process it
    - **Operational intensity** = work (FLOP) per memory access (byte)
    - Performance = work / second  
= work / byte \* byte / second  
= operational intensity \* memory bandwidth
- Linear relationship between performance and operational intensity

# The Roofline Model: A bird's eye view of performance

- **Roofline**: theoretical performance achievable given an intensity
  - Formed by memory bandwidth bounds + peak CPU performance



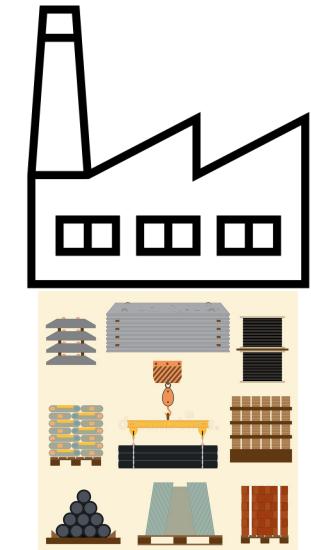
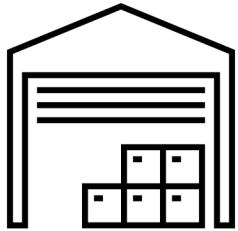
# Memory Wall: widening gap between memory and CPU



Source: SC16 Invited Talk “Memory Bandwidth and System Balance in HPC Systems” by John D. McCaughan

- FLOPS = floating point operations per second (CPU speed)

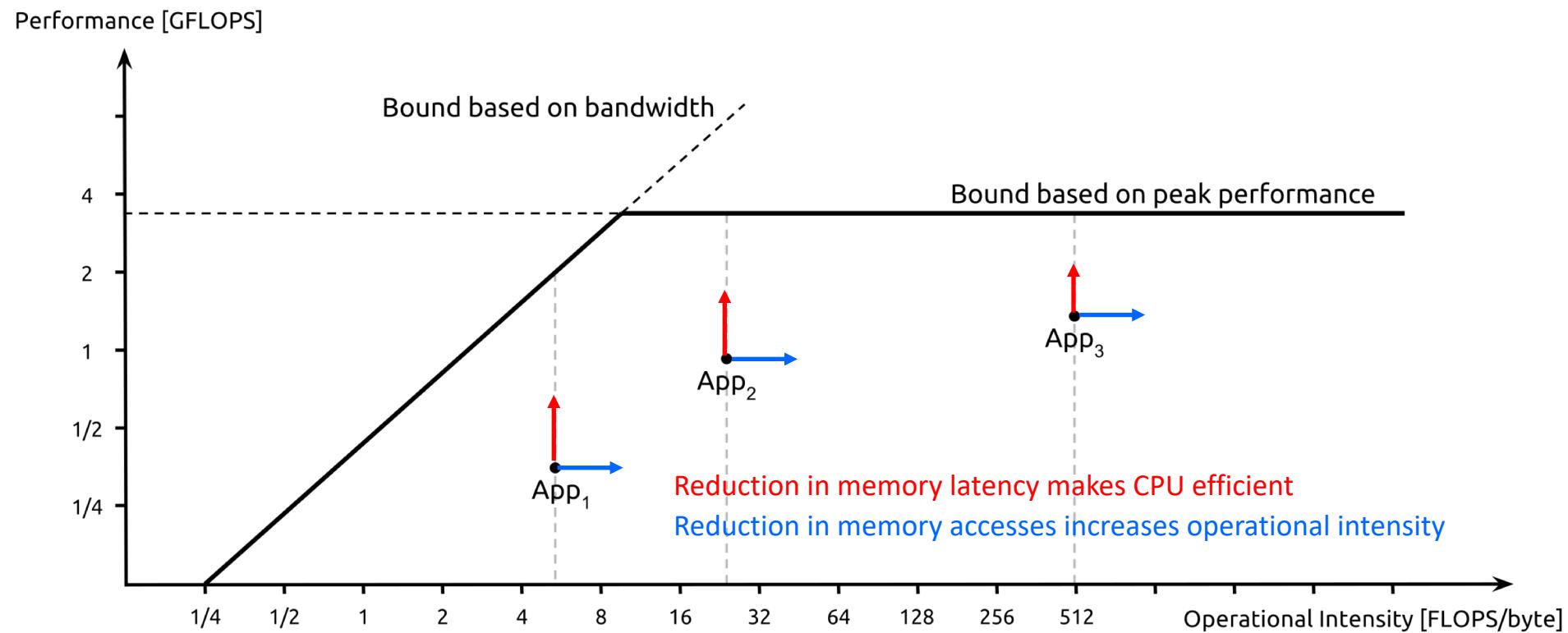
# Cache: storage in CPU for frequently accessed data



- What if we had a storage **cache** in factory premises?
  - Bring in a truckload of material on an order (more than needed)
  - Store the rest of the material in cache for future use
- Then, both memory latency and bandwidth problems are alleviated
  - Latency: access is quick when material is found in the cache
  - Bandwidth: less need to go back and forth to the warehouse

# Effect of Cache on Roofline Model

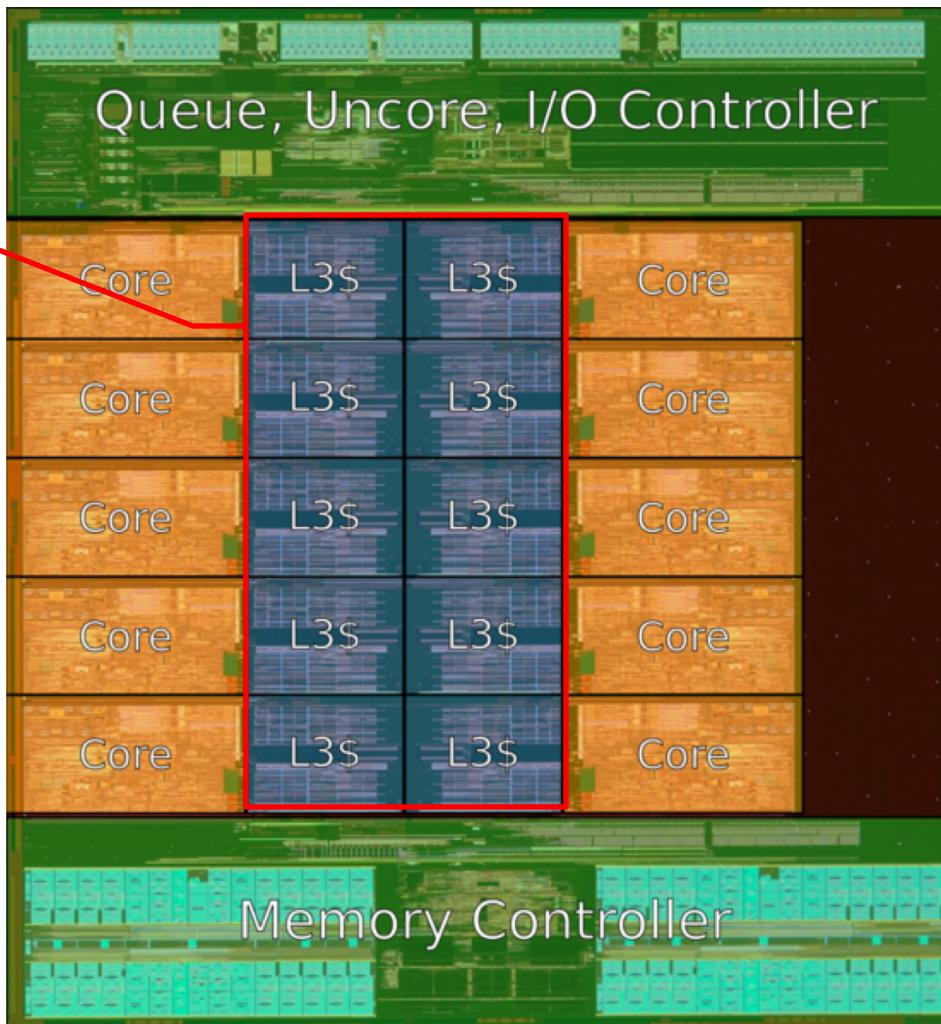
- With caching, apps shift upwards and rightwards



# Caches have become bigger and bigger

- As the memory wall got higher and higher, caches got bigger and bigger
- On the right is a diagram of the Intel Xeon Broadwell CPU
  - Caches take up almost as much real estate as cores!
  - A memory access is that painful (10~100X slower)

Caches



# How to design a shirt factory

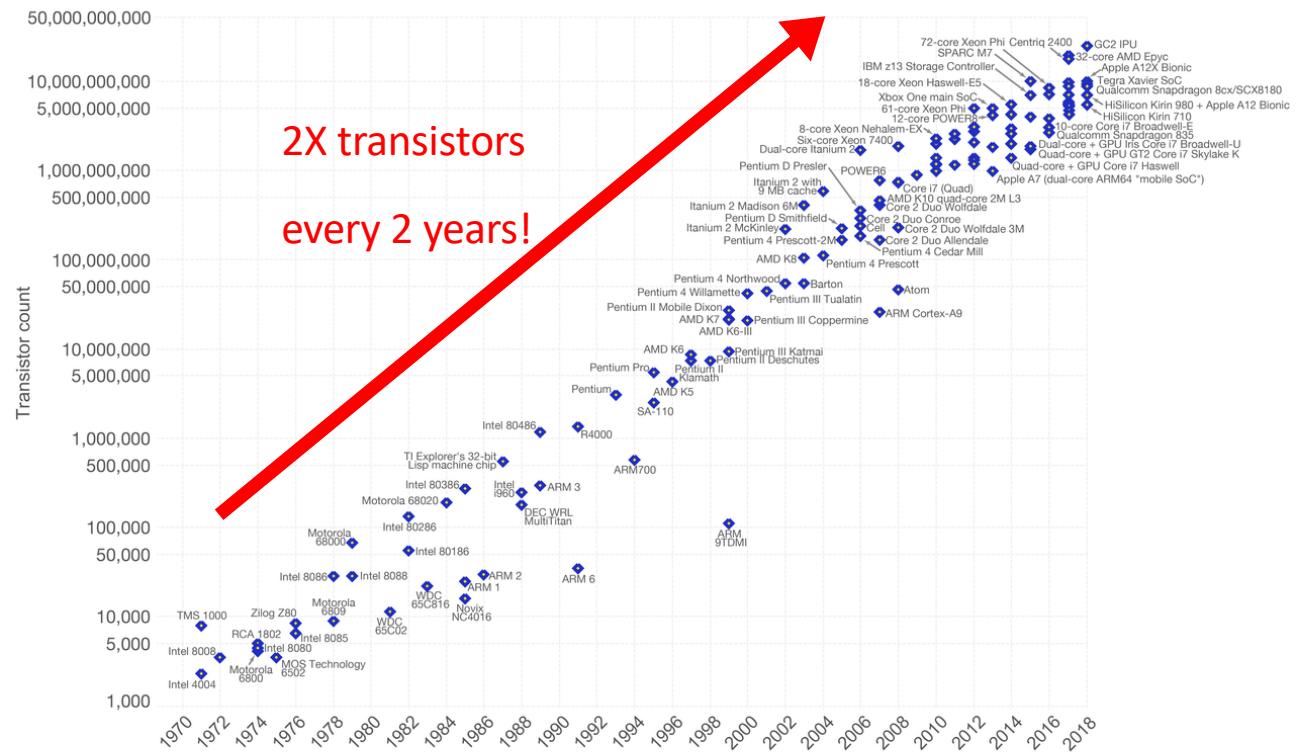
## Part 2: Worker Utilization

# Moore's Law

## Moore's Law – The number of transistors on integrated circuit chips (1971-2018)

Moore's law describes the empirical regularity that the number of transistors on integrated circuits doubles approximately every two years. This advancement is important as other aspects of technological progress – such as processing speed or the price of electronic products – are linked to Moore's law.

OurWorld  
in Data

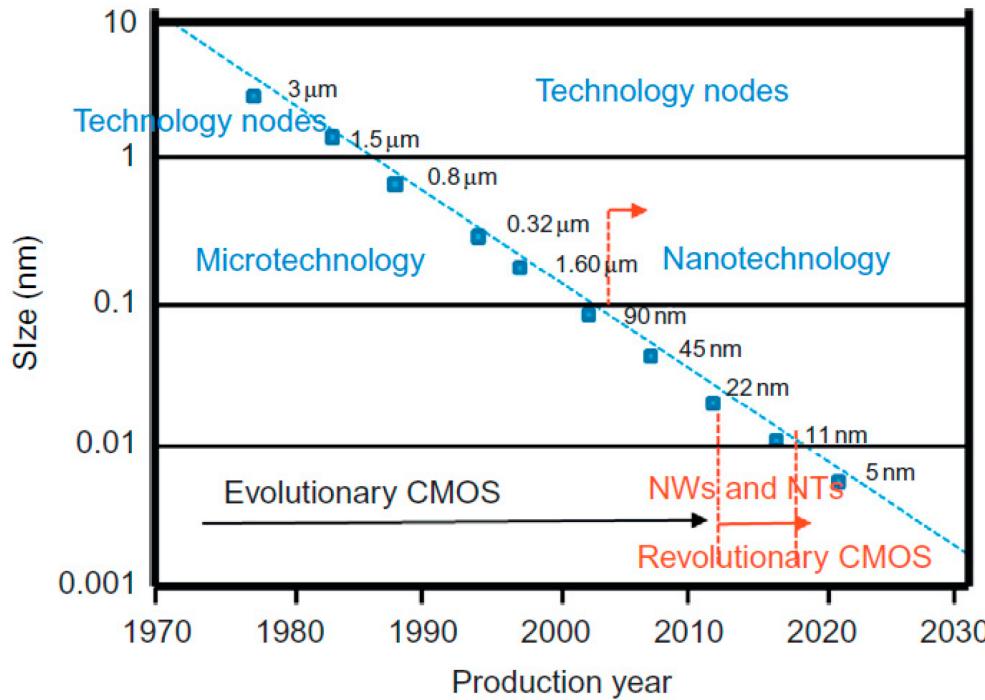


Data source: Wikipedia ([https://en.wikipedia.org/wiki/Transistor\\_count](https://en.wikipedia.org/wiki/Transistor_count))

The data visualization is available at [OurWorldInData.org](http://OurWorldInData.org). There you find more visualizations and research on this topic.

Licensed under CC-BY-SA by the author Max Roser.

# Miniaturization of Transistors



Data source: Radamson, H.H.; He, X.; Zhang, Q.; Liu, J.; Cui, H.; Xiang, J.; Kong, Z.; Xiong, W.; Li, J.; Gao, J.; Yang, H.; Gu, S.; Zhao, X.; Du, Y.; Yu, J.; Wang, G. Miniaturization of CMOS. *Micromachines* **2019**, *10*, 293.

- Moore's Law has been driven by transistor miniaturization

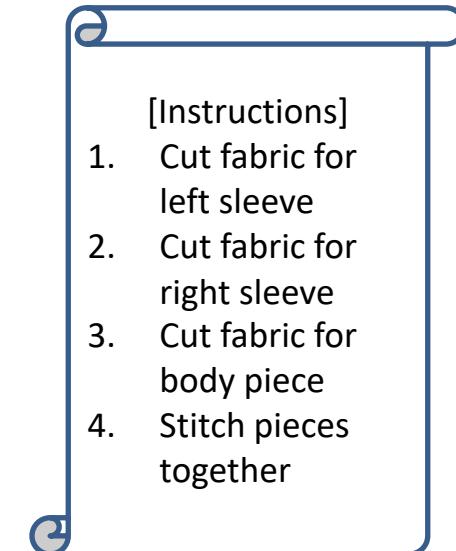
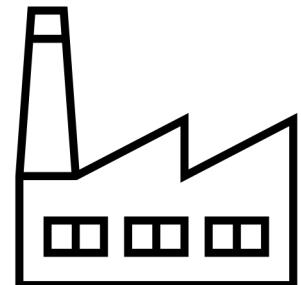
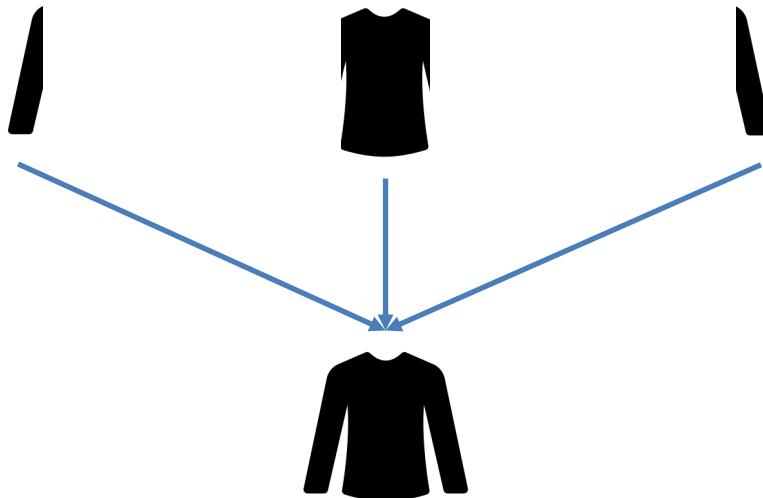
# Our factory is bursting with workers!

- Transistors are cobbled together to create **Functional Units**
  - Things like Add Unit, Multiply Unit, Load Unit, Store Unit
- With billions of transistors, CPUs can contain many Functional Units
- What is the problem with a factory with a lot of workers?
  - It is hard to keep all the workers productive!
  - Workers always find an excuse to slack...
  - Same goes for Functional Units



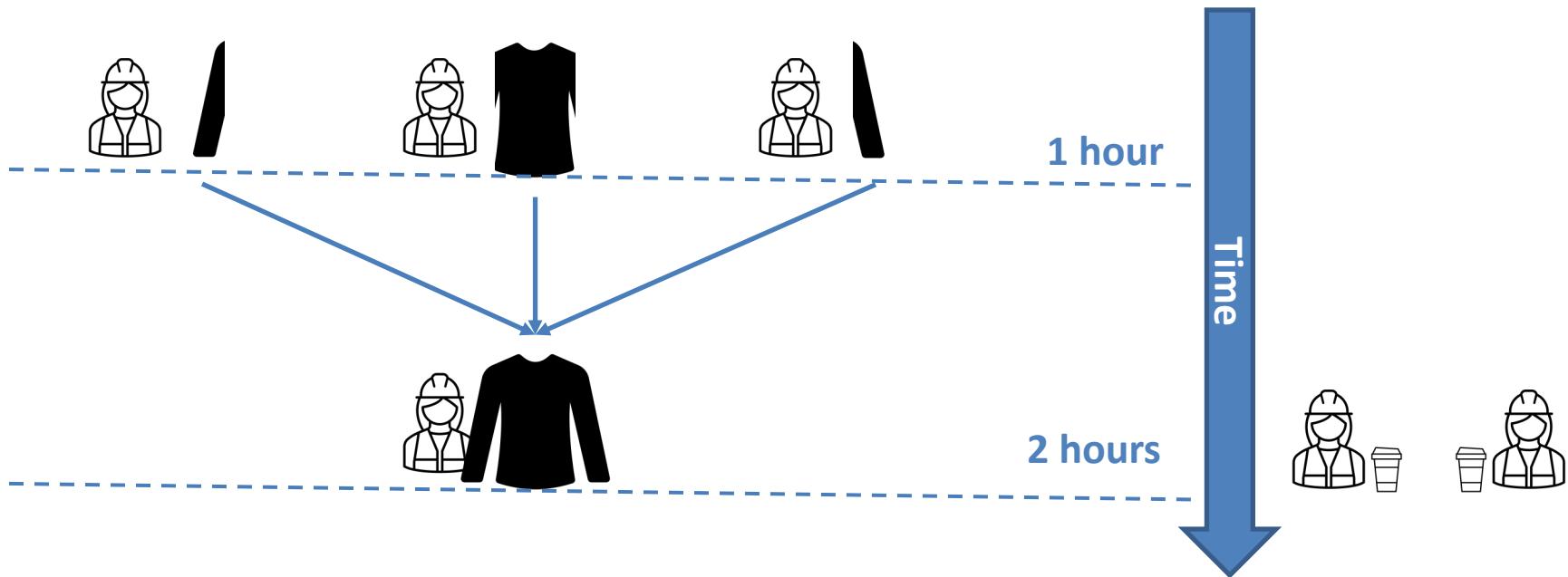
# Shirt manufacturing instructions

- Let's assume one man hour for each instruction.
- How many workers would you hire?
- **Data Dependence Graph** for shirt manufacturing



# 3 workers = fastest but low utilization

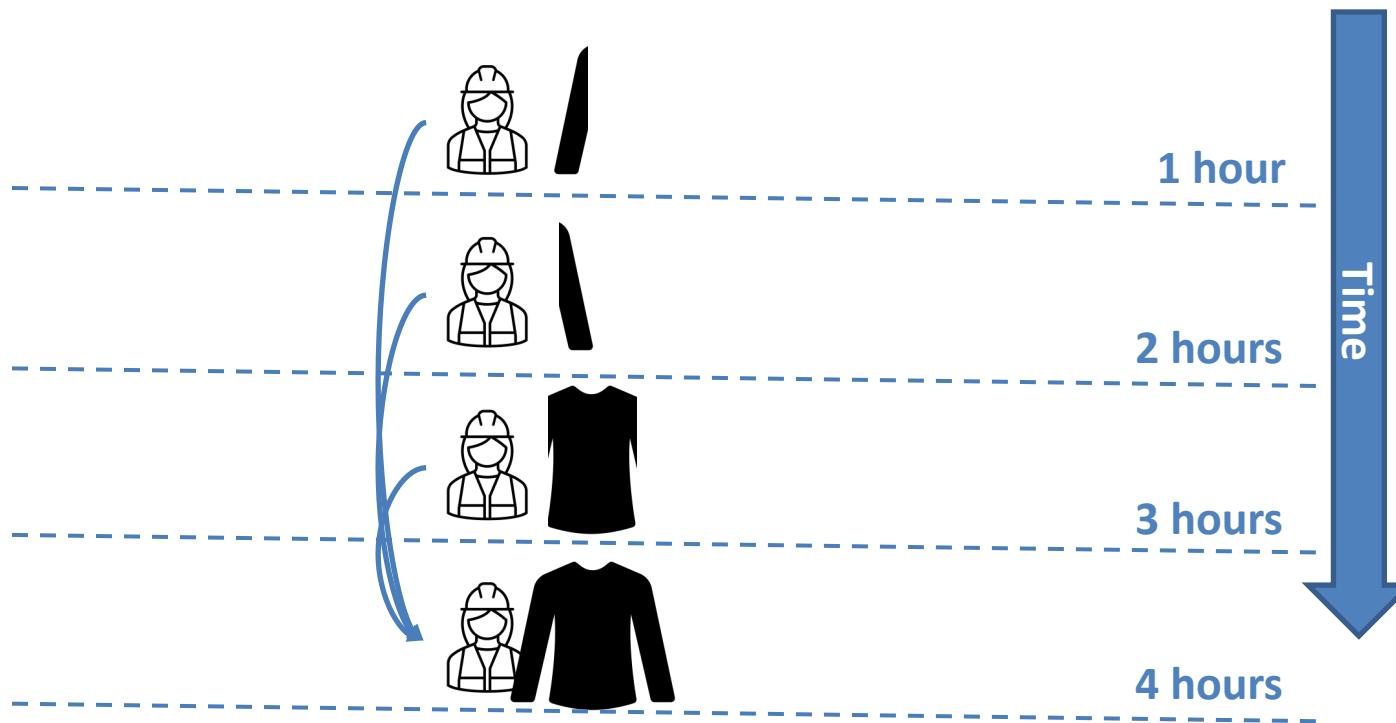
- Can produce a shirt in just 2 hours:



- During the 2<sup>nd</sup> hour, 2 workers were slacking over coffee. Oh no!
- Utilization = 4 steps / 6 worker-hours \* 100 ≈ 67%

# 1 worker = slow but highest utilization

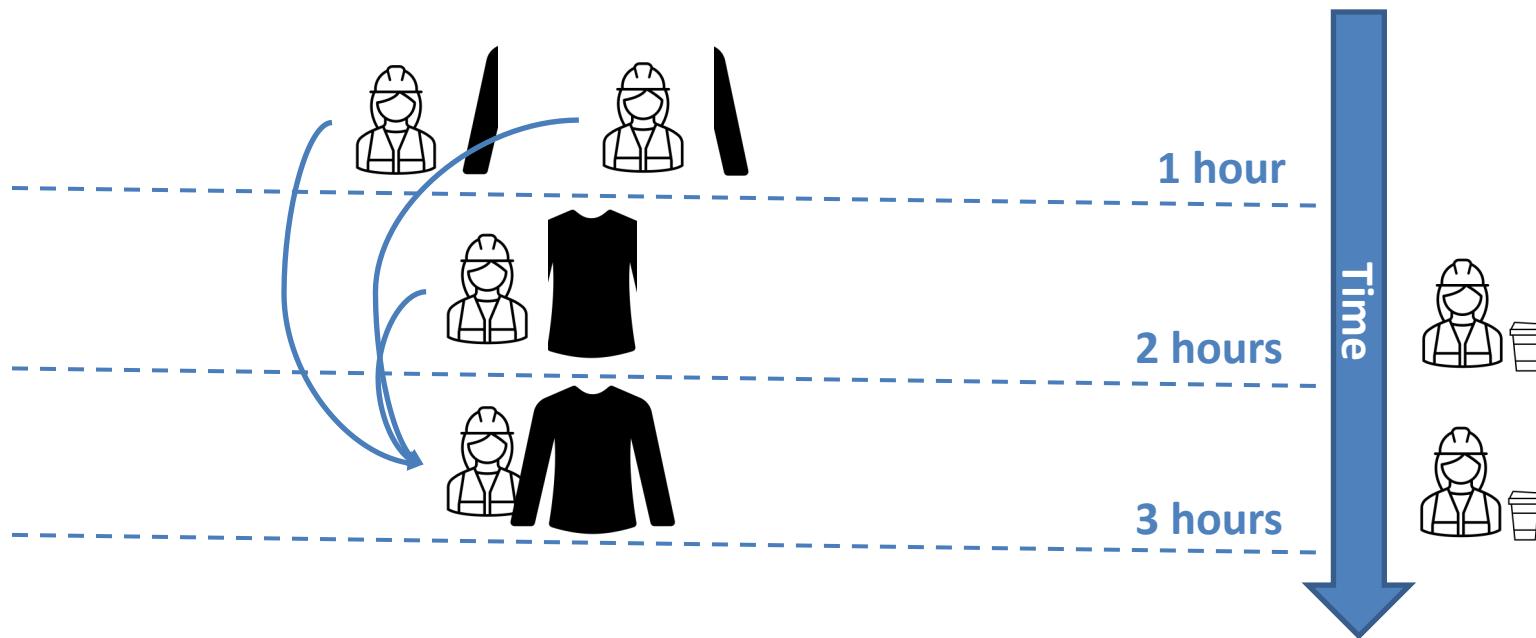
- Needs 4 hours to produce a shirt, but no slacking!



- Utilization = 4 steps / 4 worker-hours \* 100 = 100%

# 2 workers = mediocre speed and low utilization

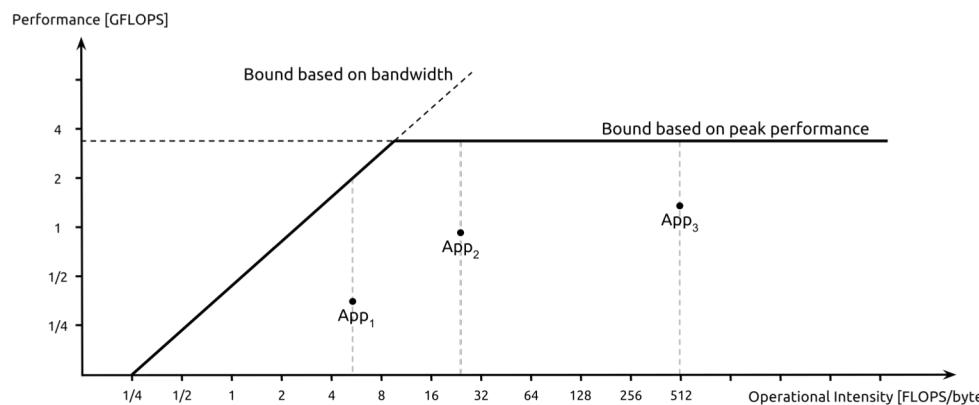
- Needs 3 hours with some slacking:



- Utilization = 4 steps / 6 worker-hours \* 100 ≈ 67%
- Same utilization as 3 workers but just slower. **Fail!**

# What are the implications of low utilization?

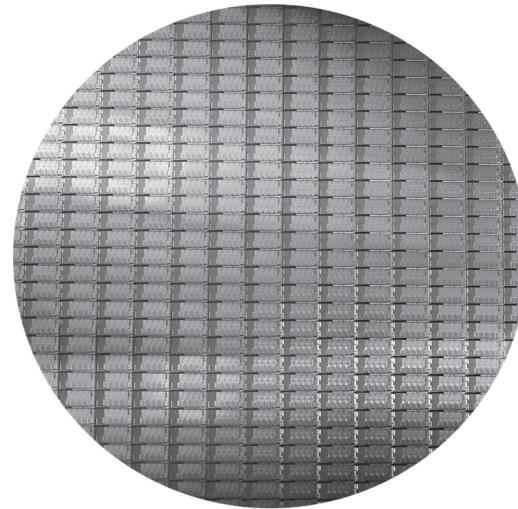
- Factory: leads to high cost for producing the same shirt
- CPU: leads to **high energy use** for executing the same instructions
  - Functional units still use energy while slacking off
- So, should we go with 3 workers or 1 worker? It depends!
  - 3 workers = Gaming CPU that runs fast but uses lots of power
  - 1 worker = Mobile CPU that runs slower but uses less battery life
- Low utilization is what prevents apps from reaching full potential





# Scope of Class

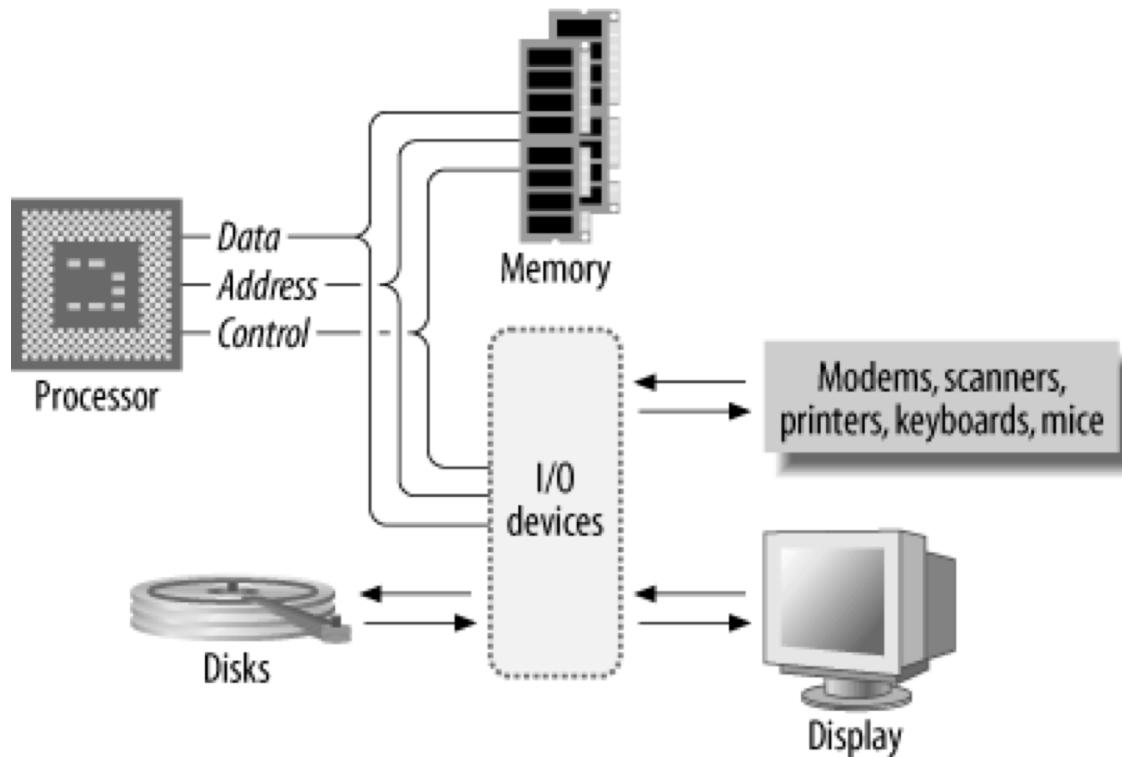
- Physical layer is beyond the scope of the class
- We will focus mostly on processor organization
  - And how performance goals are achieved





# Scope of Class

- Computer architecture is part of system architecture

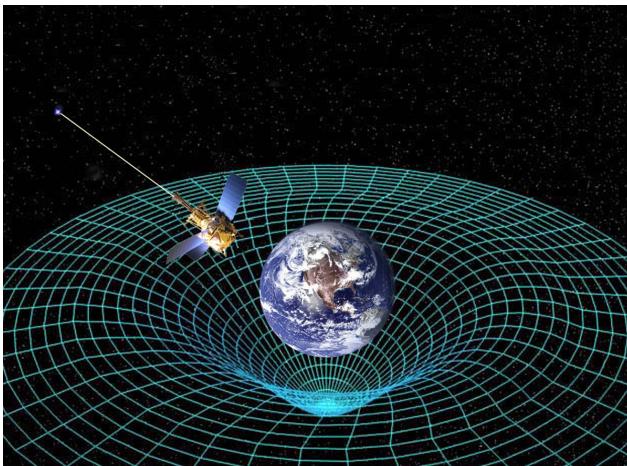


- Other components beside processor is beyond the scope



# Two Forces on Computer Architecture

## 1. Application pull



## 2. Technology push



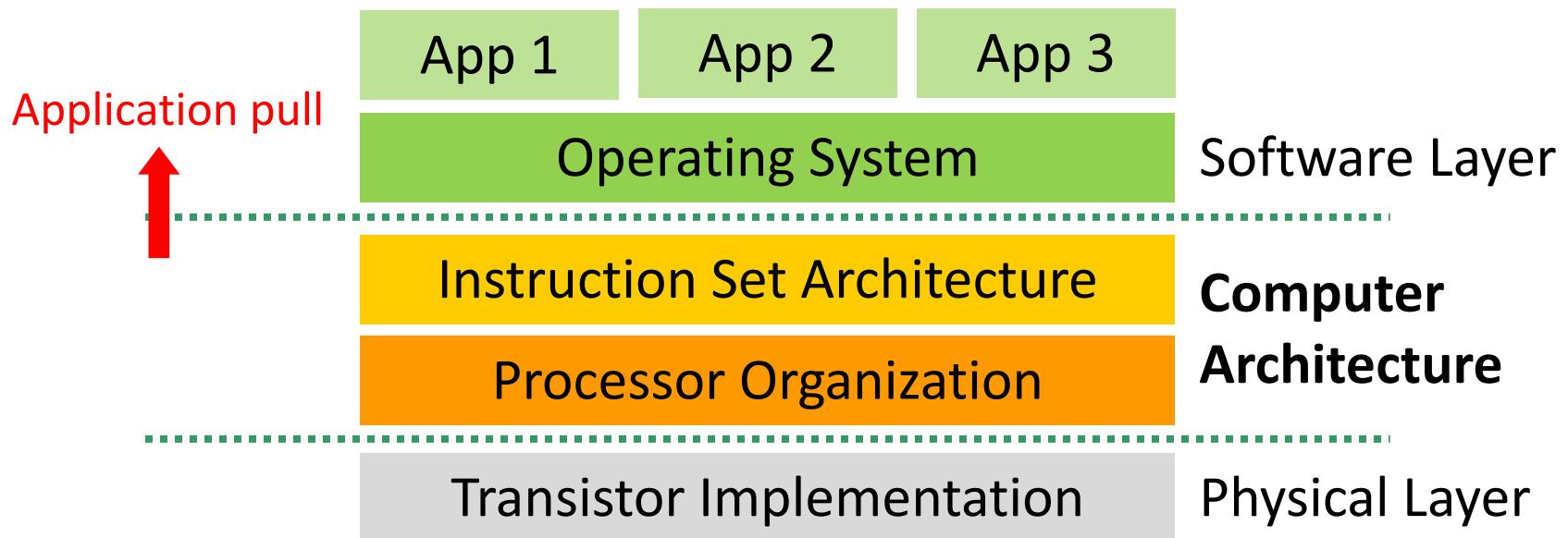
- Market forces pull architecture towards popular applications

- Advances in silicon technology push architecture to change



# Application Pull

- Different applications pull in different directions



- Real-time app (e.g. Game): *Short latency*
  - Server app: *High throughput*
  - Mobile app: *High energy-efficiency* (battery life)
  - Mission critical app: *High reliability*
- An app typically has multiple goals that are important



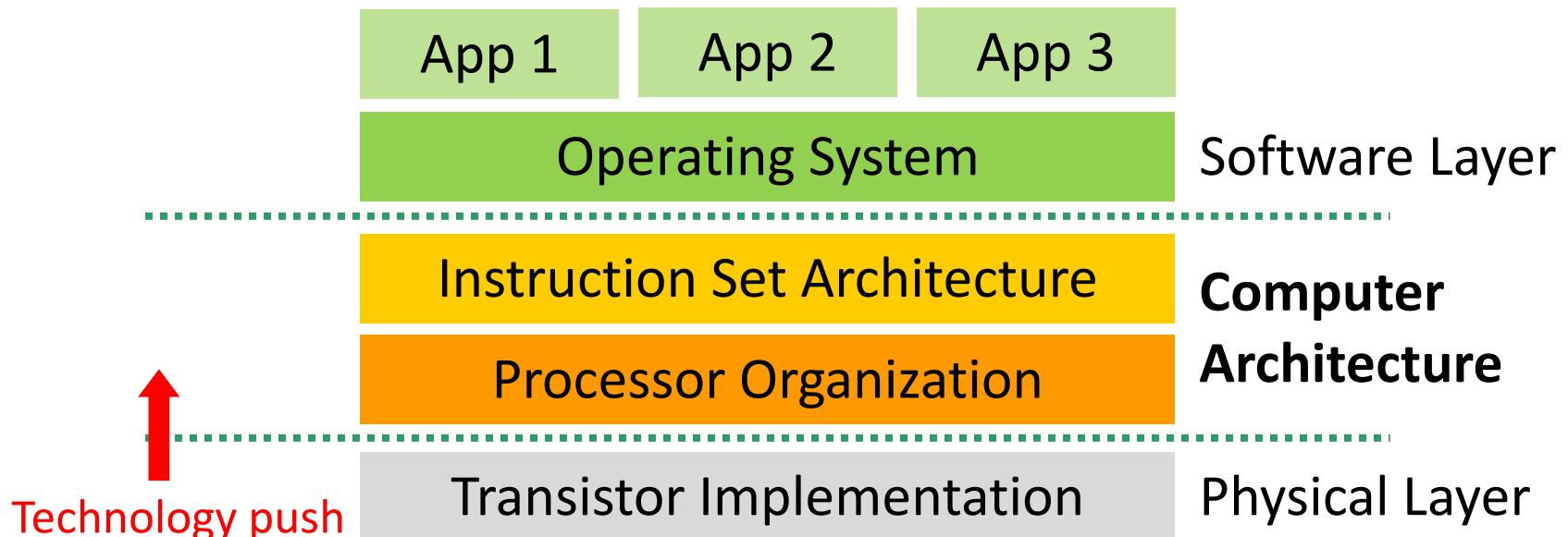
# Application Pull

- Some goals can be incompatible
  - E.g. Speed and energy-efficiency are incompatible
    - Running is faster than walking but uses more energy
    - A Ferrari is faster than a Prius but has worse fuel efficiency
  - E.g. Reliability is incompatible with many other goals
    - If you use redundancy, you use twice the amount of energy
- Even when sharing a goal, apps have unique needs
  - Scientific apps need lots of *floating point units* to go fast
  - Database apps need lots of *memory cache* to go fast
- An architecture is a **compromise** among all the apps
  - When app achieves market critical mass, designs diverge (Mobile chips / Server chips / GPUs / TPUs diverged)
  - Sometimes even ISAs diverge (GPUs and TPUs)



# Technology Push

- Trends in technology pushes architecture too



- Trends can be *advances* in technology
  - Trends can be *constraints* technology couldn't overcome
- \* “Technology” in CPU design refers to the physical layer
- Manufacturing technology used for transistor implementation