

Experiment 01 : First () and Follow() Set

Learning Objective: Student should be able to Compute First () and Follow () set of given grammar.

Tools: Jdk1.8, Turbo C/C++, Python, Notepad++

Theory:

1. Algorithm to Compute FIRST as follows:

- Let a be a string of terminals and non-terminals.
- First (a) is the set of all terminals that can begin strings derived from a .

Compute FIRST(X) as follows:

- a) if X is a terminal, then $\text{FIRST}(X) = \{X\}$
- b) if $X \rightarrow \epsilon$ is a production, then add ϵ to FIRST(X)
- c) if X is a non-terminal and $X \rightarrow Y_1 Y_2 \dots Y_n$ is a production, add FIRST(Y_i) to FIRST(X) if the preceding Y_j s contain ϵ in their FIRSTs

2. Algorithm to Compute FOLLOW as follows:

- a) FOLLOW(S) contains EOF
- b) For productions $A \rightarrow \alpha B \beta$, everything in FIRST (β) except ϵ goes into FOLLOW (B)
- c) For productions $A \rightarrow \alpha B$ or $A \rightarrow \alpha B \beta$ where FIRST (β) contains ϵ , FOLLOW(B) contains everything that is in FOLLOW(A)

Original grammar:

$$E \rightarrow E + E$$
$$E \rightarrow E * E$$
$$E \rightarrow (E)$$
$$E \rightarrow id$$

This grammar is left-recursive, ambiguous and requires left-factoring. It needs to be modified before we build a predictive parser for it:

Step 1: Remove Ambiguity.

$$E \rightarrow E + T$$
$$T \rightarrow T * F$$
$$F \rightarrow (E)$$
$$F \rightarrow id$$

Grammar is left recursive hence Remove left recursion:

$$E \rightarrow TE'$$
$$E' \rightarrow +TE' | \epsilon$$
$$T \rightarrow FT'$$
$$T' \rightarrow *FT' | \epsilon$$
$$F \rightarrow (E)$$
$$F \rightarrow id$$

Step 2: Grammar is already left factored.

Step 3: Find First & Follow set to construct predictive parser table:-

$$\text{FIRST}(E) = \text{FIRST}(T) = \text{FIRST}(F) = \{ (, id \}$$
$$\text{FIRST}(E') = \{ +, \epsilon \}$$
$$\text{FIRST}(T') = \{ *, \epsilon \}$$
$$\text{FOLLOW}(E) = \text{FOLLOW}(E') = \{ \$,) \}$$

$$\text{FOLLOW}(T) = \text{FOLLOW}(T') = \{+, \$,)\}$$

$$\text{FOLLOW}(F) = \{*, +, \$,)\}$$

Example:

$$E \rightarrow TE'$$

$$E' \rightarrow +TE' | \epsilon$$

$$T \rightarrow FT'$$

$$T' \rightarrow *FT' | \epsilon$$

$$F \rightarrow (E)$$

$$F \rightarrow \text{id}$$

$$\text{FIRST}(E) = \text{FIRST}(T) = \text{FIRST}(F) = \{ (, \text{id} \}$$

$$\text{FIRST}(E') = \{ +, \epsilon \}$$

$$\text{FIRST}(T') = \{ *, \epsilon \}$$

$$\text{FOLLOW}(E) = \text{FOLLOW}(E') = \{ \$,) \}$$

$$\text{FOLLOW}(T) = \text{FOLLOW}(T') = \{ +, \$,) \}$$

$$\text{FOLLOW}(F) = \{ *, +, \$,) \}$$

Application: To design Top Down and Bottom up Parsers.

Design:

Result and Discussion:

Learning Outcomes: The student should have the ability to

LO1: Identify type of grammar G.

LO2: Define First () and Follow () sets.

LO3: Find First () and Follow () sets for given grammar G.

LO4: Apply First () and Follow () sets for designing Top Down and Bottom up Parsers

Course Outcomes: Upon completion of the course students will be able to analyze the analysis and synthesis phase of compiler for writing application programs and construct different parsers for given context free grammars.

Conclusion:

For Faculty Use

Correction Parameters	Formative Assessment [40%]	Timely completion of Practical [40%]	Attendance / Learning Attitude [20%]	
Marks Obtained				