

# **Tugas 2 : Kopleksitas Waktu**

## **Praktikum Analisis Algoritma**



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## 1. Mencari nilai max

### - Algoritma

```
procedure CariMaks(input  $x_1, x_2, \dots, x_n$ : integer, output maks: integer)
{ Mencari elemen terbesar dari sekumpulan elemen larik integer  $x_1, x_2, \dots, x_n$ . Elemen
  terbesar akan disimpan di dalam maks
  Input:  $x_1, x_2, \dots, x_n$ 
  Output: maks (nilai terbesar)
}
```

#### Deklarasi

$i$  : integer

#### Algoritma

```
maks  $\leftarrow x_1$ 
 $i \leftarrow 2$ 
while  $i \leq n$  do
  if  $x_i > \text{maks}$  then
    maks  $\leftarrow x_i$ 
  endif
   $i \leftarrow i + 1$ 
endwhile
{ $i > n$ }
```

### - Code

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int maks = x[0];
    int i = 1;
    while (i <= n)
    {
        if (x[i] > maks)
            maks = x[i];
    }
}
```

```

        i++;
    }
    cout << "Maksimum Number : " << maks << endl;

    return 0;
}

```

#### - Kompleksitas waktu

$\text{maks} \leftarrow x_1$       1 kali

$i \leftarrow 2$       1 kali

$\text{maks} \leftarrow x_i$       n kali

$i \leftarrow i + 1$       n kali

$$T(n) = 1 + 1 + n + n = 2n + 2$$

## 2. Sequential Search

### - Algoritma

```
procedure SequentialSearch(input  $x_1, x_2, \dots, x_n$  : integer, y : integer, output idx : integer)
{   Mencari y di dalam elemen  $x_1, x_2, \dots, x_n$ . Lokasi (indeks elemen) tempat y ditemukan
    diisi ke dalam idx. Jika y tidak ditemukan, maka idx diisi dengan 0.
    Input:  $x_1, x_2, \dots, x_n$ 
    Output: idx
}
```

#### Deklarasi

i : integer

found : boolean { bernilai true jika y ditemukan atau false jika y tidak ditemukan }

#### Algoritma

i  $\leftarrow$  1

found  $\leftarrow$  false

while (i  $\leq$  n) and (not found) do

if  $x_i = y$  then

found  $\leftarrow$  true

else

i  $\leftarrow$  i + 1

endif

endwhile

{ i < n or found }

If found then { y ditemukan }

idx  $\leftarrow$  i

else

idx  $\leftarrow$  0 { y tidak ditemukan }

endif

### - Code

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
```

```

        cin >> x[i];
    }

    int y;
    cout << "Masukkan yang dicari : ";
    cin >> y;

    int i = 0;
    bool found = false;
    int idx;
    while ((i < n) && (!found))
    {
        if (x[i] == y)
            found = true;
        else
            i++;
    }
    if (found)
        idx = i+1;
    else
        idx = 0;

    cout << "Yang dicari berada di urutan : " << idx << endl;

    return 0;
}

```

### - Kompleksitas waktu

- Best Case :

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow 1$	1 kali

$$T_{\min}(n) = 1 + 1 + 1 + 1 = 4$$

- Average Case :

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$i \leftarrow i + 1$	$\frac{1}{2} n$ kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow 1$	1 kali

$$T_{\text{avg}}(n) = 1 + 1 + \frac{1}{2} n + 1 + 1 = \frac{1}{2} n + 4$$

- Worst Case :
 

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$i \leftarrow i + 1$	$n$ kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow l$	1 kali

$$T_{\max}(n) = 1 + 1 + n + 1 + 1 = n + 4$$

### 3. Binary Search

#### - Algoritma

```
procedure BinarySearch(input  $x_1, x_2, \dots, x_n$  : integer, x : integer, output : idx : integer)
{   Mencari y di dalam elemen  $x_1, x_2, \dots, x_n$ . Lokasi (indeks elemen) tempat y ditemukan
    diisi ke dalam idx. Jika y tidak ditemukan maka dx diisi dengan 0.
    Input:  $x_1, x_2, \dots, x_n$ 
    Output: idx
}
Deklarasi
    i, j, mid : integer
    found : Boolean
Algoritma
    i  $\leftarrow$  1
    j  $\leftarrow$  n
    found  $\leftarrow$  false
    while (not found) and (i  $\leq$  j) do
        mid  $\leftarrow$  (i + j) div 2
        if  $x_{mid} = y$  then
            found  $\leftarrow$  true
        else
            if  $x_{mid} < y$  then {mencari di bagian kanan}
                i  $\leftarrow$  mid + 1
            else {mencari di bagian kiri}
                j  $\leftarrow$  mid - 1
            endif
        endif
    endwhile
    {found or i > j}

    If found then
        Idx  $\leftarrow$  mid
    else
        Idx  $\leftarrow$  0
    endif
```

#### - Code

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
```

```

    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int y;
    cout << "Masukkan yang dicari : ";
    cin >> y;

    int i = 0;
    int j = n-1;
    bool found = false;
    int idx;
    int mid;
    while ((i <= j) && (!found))
    {
        mid = (i + j)/2;
        if (x[mid] == y)
            found = true;
        else
        {
            if (x[mid] < y)
                i = mid + 1;
            else
                j = mid - 1;
        }
    }

    if (found)
        idx = mid+1;
    else
        idx = 0;

    cout << "Yang dicari berada di urutan : " << idx << endl;

    return 0;
}

```



## - Kompleksitas waktu

- Best Case :

$i \leftarrow 1$	1 kali
$j \leftarrow n$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{mid} \leftarrow (i + j) \text{ div } 2$	1 kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{Idx} \leftarrow \text{mid}$	1 kali

$$T_{\min}(n) = 1 + 1 + 1 + 1 + 1 + 1 = 6$$

- Average Case :

$i \leftarrow 1$	1 kali
$j \leftarrow n$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{mid} \leftarrow (i + j) \text{ div } 2$	$\frac{1}{2}n + 1$ kali
$i \leftarrow \text{mid} + 1 \text{ or } j \leftarrow \text{mid} - 1$	$\frac{1}{2}n$ kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{Idx} \leftarrow \text{mid}$	1 kali

$$T_{\text{avg}}(n) = 1 + 1 + 1 + \frac{1}{2}n + 1 + \frac{1}{2}n + 1 + 1 = n + 6$$

- Worst Case :

$i \leftarrow 1$	1 kali
$j \leftarrow n$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{mid} \leftarrow (i + j) \text{ div } 2$	$n + 1$ kali
$i \leftarrow \text{mid} + 1 \text{ or } j \leftarrow \text{mid} - 1$	$n$ kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{Idx} \leftarrow \text{mid}$	1 kali

$$T_{\max}(n) = 1 + 1 + 1 + n + 1 + n + 1 + 1 = 2n + 6$$

## 4. Insertion Sort

### - Algoritma

```
procedure InsertionSort(input/output  $x_1, x_2, \dots, x_n$  : integer)
{
    Mengurutkan elemen-elemen  $x_1, x_2, \dots, x_n$  dengan metode insertion sort.
    Input:  $x_1, x_2, \dots, x_n$ 
    Output:  $x_1, x_2, \dots, x_n$  (sudah terurut menaik)
}
Deklarasi
    i, j, insert : integer
Algoritma
    for i  $\leftarrow$  2 to n do
        insert  $\leftarrow$   $x_i$ 
        j  $\leftarrow$  i
        while (j < i) and ( $x[j-i]$  > insert) do
             $x[j] \leftarrow x[j-1]$ 
            j  $\leftarrow$  j-1
        endwhile
         $x[j] =$  insert
    endfor
```

### - Code

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }
    cout << "Data Sebelum di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";
    cout << endl;

    int insert;
```

```

int j;
for (int i = 1; i < n; i++)
{
    insert = x[i];
    j = i-1;
    while ((j >= 0) && (x[j] > insert))
    {
        x[j+1] = x[j];
        j--;
    }
    x[j+1] = insert;
}

cout << "Data setelah di Sorting : ";
for (int i = 0; i < n; i++)
    cout << x[i] << " ";

return 0;
}

```

### - Kompleksitas waktu

- Best Case :
 

fori $\leftarrow$ 2 to n do	1 kali
insert $\leftarrow$ xi	n kali
j $\leftarrow$ i	n kali
x[j] = insert	n kali

$$T_{min}(n) = 1 + n + n + n = 3n + 1$$

- Average Case :
 

fori $\leftarrow$ 2 to n do	1 kali
insert $\leftarrow$ xi	n kali
j $\leftarrow$ i	n kali
x[j] $\leftarrow$ x[j-1]	n * ½ n kali
j $\leftarrow$ j-1	n * ½ n kali
x[j] = insert	n kali

$$T_{avg}(n) = 1 + n + n + \frac{1}{2} n^2 + \frac{1}{2} n^2 + n = n^2 + 3n + 1$$

- Worst Case :
 

fori $\leftarrow$ 2 to n do	1 kali
insert $\leftarrow$ xi	n kali
j $\leftarrow$ i	n kali

$x[j] \leftarrow x[j-1]$	$n * n$ kali
$j \leftarrow j-1$	$n * n$ kali
$x[j] = \text{insert}$	$n$ kali

$$T_{\max}(n) = 1 + n + n + n^2 + n^2 + n = 2n^2 + 3n + 1$$

## 5. Selection Sort

### - Algoritma

```
procedure SelectionSort(input/output  $x_1, x_2, \dots x_n$  : integer)
{  Mengurutkan elemen-elemen  $x_1, x_2, \dots x_n$  dengan metode selection sort.
  Input:  $x_1, x_2, \dots x_n$ 
  Output:  $x_1, x_2, \dots x_n$  (sudah terurut menaik)
}
Deklarasi
  i, j, imaks, temp : integer
Algoritma
  for i  $\leftarrow$  n downto 2 do {pass sebanyak n-1 kali}
    imaks  $\leftarrow$  1
    for j  $\leftarrow$  2 to i do
      if  $x_j > x_{imaks}$  then
        imaks  $\leftarrow$  j
      endif
    endfor
    {pertukarkan  $x_{imaks}$  dengan  $x_i$ }
    temp  $\leftarrow$   $x_i$ 
     $x_i \leftarrow x_{imaks}$ 
     $x_{imaks} \leftarrow$  temp
  endfor
```

### - Code

```
#include <iostream>
using namespace std;

int main()
{
  int n;
  int x[10];
  cout << "Masukkan Jumlah Data : ";
  cin >> n;
  for (int i = 0; i < n; i++)
  {
    cout << "Masukkan Data ke - " << i+1 << " : ";
    cin >> x[i];
  }
  cout << "Data Sebelum di Sorting : ";
  for (int i = 0; i < n; i++)
    cout << x[i] << " ";
  cout << endl;

  int imaks;
```

```

int temp;
for (int i = n-1; i >= 1; i--)
{
    imaks = 0;
    for (int j = 1; j <= i; j++)
    {
        if (x[j] > x[imaks])
            imaks = j;
    }
    temp = x[i];
    x[i] = x[imaks];
    x[imaks] = temp;
}

cout << "Data setelah di Sorting : ";
for (int i = 0; i < n; i++)
    cout << x[i] << " ";

return 0;
}

```

#### - Kompleksitas waktu

- Best Case :
 

fori $\leftarrow$ n downto 2 do	1 kali
imaks $\leftarrow$ 1	n kali
forj $\leftarrow$ 2 to i do	n kali
imaks $\leftarrow$ j	n * 1 kali
temp $\leftarrow$ xi	n kali
xi $\leftarrow$ ximaks	n kali
ximaks $\leftarrow$ temp	n kali

$$T_{min}(n) = 1 + n + n + n * 1 + n + n + n = 6n + 1$$

- Average Case :
 

fori $\leftarrow$ n downto 2 do	1 kali
imaks $\leftarrow$ 1	n kali
forj $\leftarrow$ 2 to i do	n kali
imaks $\leftarrow$ j	n * $\frac{1}{2}$ n kali
temp $\leftarrow$ xi	n kali
xi $\leftarrow$ ximaks	n kali
ximaks $\leftarrow$ temp	n kali

$$T_{avg}(n) = 1 + n + n + \frac{1}{2}n^2 + n + n + n = \frac{1}{2}n^2 + 5n + 1$$

- Worst Case :
 

fori $\leftarrow$ n downto 2 do	1 kali
imaks $\leftarrow$ 1	n kali
forj $\leftarrow$ 2 to i do	n kali
imaks $\leftarrow$ j	n * n kali
temp $\leftarrow$ xi	n kali
xi $\leftarrow$ ximaks	n kali
ximaks $\leftarrow$ temp	n kali

$$T_{max}(n) = 1 + n + n + n^2 + n + n + n = n^2 + 5n + 1$$