```
목차
                                                        pii(a.fi+b.fi, a.se+b.se);}
                                                        pii operator - (pii &a,pii &b){return pii(a.fi-
헤더 (1)
                                                        b.fi, a.se-b.se);}
                                                        pll operator + (pll &a,pll &b){return
Math (1-6)
                                                        pll(a.fi+b.fi, a.se+b.se);}
                                                        pll operator - (pll &a,pll &b){return pll(a.fi-
Geometry (6-8)
                                                        b.fi, a.se-b.se);}
Flow (8-11)
                                                        int main(){
Graph (11-14)
                                                            fastio;
String (14-18)
Tree (18-19)
Data Structure (19~25)
                                                        서용헤더
                                                        # pragma GCC optimize ("03")
민규헤더
                                                        # pragma GCC optimize ("Ofast")
                                                        # pragma GCC optimize ("unroll-loops")
#include<bits/stdc++.h>
                                                        # pragma GCC
#include <bits/extc++.h>
                                                        target("sse,sse2,sse3,sse4,avx,avx2")
#include<ext/rope>
                                                        #include <bits/stdc++.h>
using namespace std;
                                                        using namespace std;
using namespace gnu cxx;
                                                        typedef long long 11;
using namespace __gnu_pbds;
                                                        typedef long double ld;
#define fi first
                                                        const 11 \mod = 1e9+7;
#define se second
                                                        const int MxN = 2210;
#define fastio
                                                        const int inf = INT MAX;
ios base::sync with stdio(false);cin.tie(0)
                                                        #define REP0(i,n) for(int i=0; i< n; i++)
#define fopen freopen("input.txt", "r", stdin)
                                                        #define REP1(i,n) for(int i=1; i<=n; i++)</pre>
#define eb emplace back
                                                        #define REP(i,a,b) for(int i=a; i<=b; i++)</pre>
#define em emplace
                                                        #define sz(v) ((int)(v).size())
#define prec(a) cout<<fixed;cout.precision(a);</pre>
                                                        #define all(v) (v).begin(), (v).end()
#define all(a) (a).begin(), (a).end()
                                                        #define compress(v) sort(all(v));
typedef long long ll;typedef long double ld;typedef
                                                        v.erase( unique(all(v)) , v.end() )
unsigned long long ul; typedef unsigned int
                                                        #define reset(X) memset(X, 0, sizeof(X))
ui;typedef pair<int,int> pii;typedef pair<ll,ll>
                                                        #define pb push back
pll;typedef complex<double> cpx;
                                                        #define fi first
typedef tuple<int,int,int> tiii;
                                                        #define se second
template<class T>
                                                        #define pii pair<int, int>
using ordered set = tree<T, null type, less<T>,
                                                        #define pll pair<ll, 11>
rb tree tag, tree order statistics node update>;
                                                        #define vint vector<int>
template<class T>
                                                        #define vll vector<ll>
void pr(T t) {cerr << t << " ";}</pre>
                                                        #define vpii vector<pair<int, int>>
template<class T, class ...Args>
                                                        #define vpll vector<pair<11,11>>
void pr(T a, Args ...args) {cerr << a << "</pre>
                                                        #define endl "\n"
";pr(args...);}
                                                        #define LOG 18
template<class ...Args>
void prl(Args ...args) {pr(args...);cerr << endl;}</pre>
const 11 INF = 2e18;
                                                        희승헤더
const int inf = 2e9;
const double PI = acos(-1);
                                                        #include <bits/stdc++.h>
```

pii operator + (pii &a,pii &b){return

```
#define eb emplace back
#define em emplace
#define all(v) v.begin(), v.end()
#define fi first
#define se second
using namespace std;
typedef long long 11;
typedef pair <int, int> pii;
typedef pair <ll, 11> pll;
const int MAX = 2020;
const int INF = 1e9;
const 11 LINF = 9e18;
Math-1. Fastest FFT
const double PI = acos(-1);
typedef complex<double> cpx;
void FFT(vector<cpx> &a, bool inv) {
   int n = a.size();
   // bit reversal
   for (int i = 1, j = 0; i < n; ++i) {
       int bit = n >> 1;
       while (!((j ^= bit) & bit)) bit >>= 1;
       if (i<j) swap(a[i], a[j]);</pre>
   for (int i = 1; i<n; i <<= 1) {
       double x = inv ? PI / i : -PI / i;
       cpx w = cpx(cos(x), sin(x));
       for (int j = 0; j<n; j += i << 1) {
           cpx p = cpx(1, 0);
           for (int k = 0; k < i; ++k) {
               cpx tmp = a[i + j + k] * p;
               a[i + j + k] = a[j + k] - tmp;
               a[j + k] += tmp;
               p *= w;
   if (inv) {
       for (int i = 0; i < n; ++i) a[i] /= n;
vector<ll> multiply(vector<ll> const& a, vector<ll>
const& b) {
```

```
vector<cpx> fa(a.begin(), a.end()),
fb(b.begin(), b.end());
   // 가장 가까운 2<sup>k</sup> 꼴의 n 으로 크기를 정해줌.
   int n = 1;
   while (n < a.size() + b.size())</pre>
       n <<= 1;
   fa.resize(n);
   fb.resize(n);
   // fft : coefficient rep to value rep :
O(nLogn)
   FFT(fa, 0);
   FFT(fb, 0);
   // value rep 에서 곱하기 : O(n)
   for (int i = 0; i < n; i++)</pre>
       fa[i] *= fb[i];
   //Inverse FFT : value rep to coefficeint
rep :0(nlogn)
   FFT(fa, 1);
   vector<ll> result(n);
   for (int i = 0; i < n; i++)</pre>
       result[i] = (ll)round(fa[i].real());
   return result;
```

Math-2. High-precision FFT

```
const long double PI = acos(-1);
typedef complex<long double> cpx;
void FFT(vector<cpx> &a, bool inv) {
   int n = a.size();
   // bit reversal
   for (int i = 1, j = 0; i < n; ++i) {
       int bit = n >> 1;
       while (!((j ^= bit) & bit)) bit >>= 1;
       if (i<j) swap(a[i], a[j]);</pre>
   for (int i = 1; i<n; i <<= 1) {
       long double x = inv ? PI / i : -PI / i;
       cpx w = cpx(cos(x), sin(x));
       for (int j = 0; j < n; j += i << 1) {
           cpx p = cpx(1, 0);
           for (int k = 0; k < i; ++k) {
               cpx tmp = a[i + j + k] * p;
               a[i + j + k] = a[j + k] - tmp;
```

```
a[j + k] += tmp;
               p *= w;
    if (inv) {
       for (int i = 0; i<n; ++i) a[i] /= n;</pre>
vector<ll> multiply_high_precision(vector<ll>
const& a, vector<ll> const& b) {
   // 1023 = 2^10-1 : select this value to be near
sqrt(Max possible coefficient)
   // 가장 가까운 2<sup>k</sup> 꼴의 n 으로 크기를 정해줌.
   int n = 1; while (n < a.size() + b.size()) {n
   // a(i) = amsb(i) * cut + alsb(i)
   vector<cpx> amsb(n), alsb(n), bmsb(n), blsb(n);
    for(int i=0; i<sz(a);i++){amsb[i] =</pre>
a[i]>>10;alsb[i]=a[i]&1023;}
   for(int i=0; i<sz(b);i++){bmsb[i] =</pre>
b[i]>>10;blsb[i]=b[i]&1023;}
   // fft : coefficient rep to value rep :
O(nLogn)
    FFT(amsb,0); FFT(alsb,0); FFT(bmsb,0);
FFT(blsb,0):
   vector<cpx> high(n), middle(n), low(n);
   // value rep 에서 곱하기 : O(n) amsb*bmsb*2^20 +
(amsb*blsb+alsb*bmsb)*2^10 + alsb*blsb
    for (int i = 0; i < n; i++)</pre>
       high[i] = amsb[i]*bmsb[i];
       middle[i] =
amsb[i]*blsb[i]+alsb[i]*bmsb[i];
       low[i] = alsb[i]*blsb[i];
   //Inverse FFT : value rep to coefficeint
rep :O(nlogn)
    FFT(high, 1);FFT(middle, 1);FFT(low, 1);
    //Summation ->Result
    vector<ll> result(n);
    for (int i = 0; i < n; i++)
       11 h = (11)round(high[i].real());
       11 m = (11)round(middle[i].real());
       11 1 = (11)round(low[i].real());
       result[i] = (h << 20) + (m << 10) + 1;
```

```
return result;
```

Math-3. Miller-Rabin

```
#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
const 11 \mod = 1e9+7;
const int MxN = 1e5+1;
const int MxK = 1e9;
const int inf = INT MAX;
#define REP0(i,n) for(int i=0; i< n; i++)
#define REP1(i,n) for(int i=1; i<=n; i++)
#define REP(i,a,b) for(int i=a; i<=b; i++)
#define pb push back
#define fi first
#define se second
#define reset(X) memset(X, ∅, sizeof(X))
#define pii pair<int, int>
#define endl "\n"
#define vint vector<int>
#define pll pair<11, 11>
#define vll vector<ll>
#define LOG 18 //enough for 10^5
//O(numrep * (Logn)^3).
typedef int128 111;
11 modpow(int x, 11 y, 11 mod)
   if (y == 0) return 1;
   111 ans = modpow(x,y/2,mod);
   ans = (ans*ans)%mod;
   if (y\%2 == 1) ans = (ans*x)\%mod;
   return ans;
bool miillerTest(ll d, ll n, int a)
   // Compute a^d % n
   111 \times = modpow(a, d, n);
   if (x == 1 | x == n-1)
      return true;
   while (d != n-1)
```

```
x = (x * x) % n;
       d *= 2;
       if (x == 1)
                       return false;
       if (x == n-1)
                       return true;
    return false;
bool Miller_Rabin(ll n, int numrep)
   if (n <= 1 \mid | n == 4) return false;
   if (n <= 3) return true;</pre>
   11 d = n - 1;
   while (d \% 2 == 0)d /= 2;
   // 하나라도 false 가 뜨면 바로 합성수. 합성수들을
걸러내는 작업으로 보면됨.
   int a[]={2, 3, 5, 7, 11, 13, 17, 19, 23, 29,
31, 37}; //n 이 작은 LongLong 일때 결정적으로
파병함.
   for (int i = 0; i < numrep; i++)</pre>
       if(n==a[i])return true;
       if (!miillerTest(d, n,a[i]))return false;
   return true;
int main()
   for(int i=1; i<100; i++)
       if(Miller Rabin(i,12))cout<<i<<endl;</pre>
Math-4. NTT
민규 헤더 생략
*/
class NTT{
   //2^23 크기까지 가능
   const 11 mod = 998244353, w = 3;
public:
```

```
11 Pow(11 x, 11 y) {
       11 r = 1;
       while (v) {
            if (y \& 1) r = r * x % mod;
           x = x * x % mod;
           y >>= 1;
       return r;
   void FFT(vector<11> &a, bool f) {
       int N = a.size();
       ll i, j=0, k, x, y, z;
       for (i = 1; i < N; i++) {
           for (k = N >> 1; j >= k; k >>= 1) j -=
k;
            j += k;
            if (i < j) swap(a[i], a[j]);</pre>
       for (i = 1; i < N; i <<= 1) {
           x = Pow(f ? Pow(w, mod - 2) : w, mod /
i >> 1);
           for (j = 0; j < N; j += i << 1) {
               \vee = 1;
               for (k = 0; k < i; k++) {
                   z = a[i \mid j \mid k] * y % mod;
                   a[i | j | k] = a[j | k] - z;
                    if (a[i | j | k] < 0) a[i | j |
k] += mod;
                   a[j \mid k] += z;
                   if (a[j \mid k] >= mod) a[j \mid k] -=
mod;
                   y = y * x % mod;
       if (f) {
           j = Pow(a.size(), mod - 2);
           for (i = 0; i < N; i++) a[i] = a[i] *
j % mod;
   vector<ll> mult(vector<ll> a, vector<ll> b){
       int n = 1;
       while(n < a.size() + b.size() - 1) n<<=1;</pre>
       a.resize(n);
       b.resize(n);
       vector<ll> c(n);
       FFT(a, false);
       FFT(b, false);
       for(int i=0; i<n; i++) c[i] = a[i] * b[i] %</pre>
```

```
FFT(c, true);
    return c;
}
}Ntt;
int main() {
    vector<ll> a(3,1), b(3,1);
    vector<ll> c = Ntt.mult(a,b);
    for(auto i:c) cout<<i<<" "; //1 2 3 2 1 0 0 0
}</pre>
```

Math-5. Pollard's rho algorithm

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 \mod = 1e9+7;
const int MxN = 1e5;
const int MxK = 1e9;
const ll inf = INT MAX;
#define REP0(i,n) for(int i=0; i<n; i++)
#define REP1(i,n) for(int i=1; i<=n; i++)</pre>
#define REP(i,a,b) for(int i=a; i<=b; i++)
#define pb push back
#define fi first
#define se second
#define reset(X) memset(X, 0, sizeof(X))
#define pii pair<int, int>
#define endl "\n"
#define vint vector<int>
#define pll pair<ll, ll>
#define vll vector<ll>
#define LOG 18 //enough for 10^5
typedef int128 lll;
//Find a divisor that divides n in under than
O(sqrt(minP)) <= O(N^{1/4})
//code works for Long Long n
long long gcd(long long a, long long b)
    if(b==0)return a;
    return gcd(b, a%b);
long long f(long long x, long long c, long long
mod)
   return ((111)x*x%mod+c)%mod;
```

```
long long PollardRho(long long n)
   /* initialize random seed */
   //srand ((unsigned int)time(NULL));
   /* no prime divisor for 1 */
   if (n==1) return n;
   /* even number means one of the divisors is 2
   if (n \% 2 == 0) return 2;
   long long x = (rand()\%(n-2))+2;
   long long y = x;
   long long c = (rand()\%20)+1;
   /* Initialize candidate divisor (or result) */
   long long d = 1;
   /* until the prime factor isn't obtained.
       If n is prime, return n */
   while (d==1)
       x=f(x,c,n);
       y=f(y,c,n);
       y=f(y,c,n);
       /* check qcd of |x-y| and n */
       d = gcd(abs(x-y),n);
       if (d==n)
            x = (rand()\%(n-2))+2;
           \vee = \times;
            c = (rand()\%20)+1;
            d = 1;
   return d;
int main()
   11 x = mod*mod;
   cout<<x<<endl;</pre>
   cout<<PollardRho(x)<<endl;</pre>
```

Math-6. Basic Math

```
-density of primes : O(N/LnN)
 -goldbach's conjecture : even number = sum of two
primes
 -twin prime conjecture : p,p+2 prime
 -legendre's conjecture : prime exists between n^2,
-Lagrange's theorem:positive int = sum of four
squares
-Zeckendorf's theorem :positive int : unique sum
of Fibonacci where no numbers are same and no
numbers are consecutive Fibonacci
-Pythagorean triple : 0 < m < n, (n, m) = 1, n or m is
even: (n^2-m^2, 2nm, n^2+m^2) produces all
pythaaorean triple.
-Fermat's theorem : if p: prime, (x,p)=1,
then x^{(p-1)} \% p = 1
-Euler's theorem : if (x,m)=1, x^{(phi(m))} % m = 1
 -wilson's theorem : p is prime \langle - \rangle (p-1)! %p = p-1
-Burnside Lemma :
-The number of orbits are equal to the average
number of fixed points int the group action(taken
over all the group elements)
*/
//modular inverse : if (x,m)=1, inverse of x
modulo\ m = x^{(phi(m)-1)}
//O(sgrt mod)
int inverse(int x, int mod)
    if(gcd(x,mod)!=1)return 0; //doesn't exist.
   return modpow(x, phi(mod)-1, mod);
   // return modpow(x, mod-2, mod) if mod is
prime.
//general egn = particular sol + \{kb/q, -ka/q\},
particular sol = Extended * c/q.
//Extended euclidean algorithm : solve ax+by =
qcd(a,b);
//O(Log(min(a,b)))
pii Extended(int a, int b)
    if (a == 0) return {0,1};
   int x,y;
```

```
pii p = Extended(b%a, a);
   x=p.first; y=p.second;
   return {y - (b/a) * x, x};
//find modular inverse in O(log mod)
int inverse_Euclidean(int x, int mod)
   return Extended(x, mod).fi;
//Chinese remainder theorem : O(NlogM) if we use
extended euclidean algorithm for calculating
inverse, O(NsgrtM) if we use phi function for
inverse.
// find x st. a1 \pmod{m1} & a2 \pmod{m2} &....&
an(mod mn). mi's are coprime
int Chinese(vector<int> modular, vector<int>
remainder)
   // Compute product of all numbers
   int total product = 1;
   int n = (int)modular.size();
   for (int i = 0; i < n; i++)
       total product *= modular[i];
   // Initialize result
   int result = 0;
   // Apply formula
   for (int i = 0; i < n; i++) {
       int pp = total product / modular[i];
       result += remainder[i] * Extended(pp,
modular[i]).fi * pp;
       result %= total product;
   return result % total_product; //result +
total product*k 꼴은 모두 해임
/* nCr calculation */
//Solution 1. < just dp : O(n^2)
long long nCr(int n, int r, int mod)
   if(n<r)return ∅;
   if(n==0 && r==0)return 1;
```

```
long long dp[n+1][n+1];
   for(int i=1; i<=n; i++){dp[i][0]=1;</pre>
dp[i][i]=1;}
   for(int i=2; i<=n; i++)</pre>
       for(int j=1; j<i; j++)</pre>
           dp[i][j]= dp[i-1][j-1]+dp[i-1][j];
           dp[i][j]%=mod;
   return dp[n][r];
//Solution 2. <페르마 소정리. inverse.O(n)>
long long modpow(int x, int y, int mod)
   if (y == 0) return 1;
   long long ans = modpow(x,y/2,mod);
   ans = (ans*ans)%mod;
   if (y\%2 == 1) ans = (ans*x)\%mod;
   return ans;
long long inverse(int x, int mod) {return modpow(x,
mod-2, mod);}
long long nCr Fermat(int n, int r, int mod) //mod
is prime!!
   if (n < r) return 0;</pre>
   if (r == 0)return 1;
   long long fac[n + 1]; fac[\theta] = 1;
   for (int i = 1; i <= n; i++) fac[i] = (fac[i -
1] * i) % mod;
   return (fac[n] * inverse((int)fac[r], mod) %
mod * inverse((int)fac[n - r], mod) % mod) % mod;
//Solution 3. <Lucas Theorem. 나누는 수가 작을때
사용.0(p^2logp) >
// nCr%p 는 n 과 r 을 p 진법으로 바꾸고 각각의
niCri 들의 곱과 같다.
long long nCr_Lucas(int n, int r, int p)
   if (r==0) return 1;
   // Compute last digits of n and r in base p
   int ni = n%p, ri = r%p;
```

```
return nCr Lucas(n/p, r/p, p) * nCr(ni, ri, p)%
р;
Math-7. More Math
typedef long long C;
typedef complex<C> Point;
#define X real()
#define Y imag()
C cross product(Point a, Point b)
   return (conj(a)*b).Y;
int orientation(Point 11, Point 12, Point p)
   C val = cross_product(p-l1, p-l2); //양수면 l1-
>L2 선분에 대해 왼쪽, 음수면 오른쪽, 0이면 선분위
   if(val>0)return 1;
   if(val==0)return 0;
   return -1;
bool online(Point 11, Point 12, Point
p) //collinear 한 세점이 주어졌을때 p 는 l1, l2
사이에 있는가?
   C = 11.X; C = 12.X;
   C down =11.Y; C up =12.Y;
   if(left>right)swap(left, right);
   if(down>up)swap(down, up);
   return (left<=p.X && p.X<=right && down<=p.Y &&
p.Y <= up);
bool intersect(Point p1, Point q1, Point p2, Point
q2) //does line segments intersect with each
other??
   if(orientation(p1,q1,p2)!=orientation(p1,q1,q2)
&&
orientation(p2,q2,p1)!=orientation(p2,q2,q1))return
true;
   else{
       if(orientation(p1,q1,p2)==0 \&\&
orientation(p1,q1,q2)==0)//네점 collinear
```

```
if (online(p1, q1, p2) ||
online(p1,q1,q2) | online(p2,q2,p1) |
online(p2,q2,q1))return true; //무한히 많은 교점
       return false;
// Finding inverse = making RREF = Solving Linear
egn ...
//int -> double & remove p& modulo op for ordinary
void GaussJordan(vector<vint> & A, vector<vint> &
invA, int n, int p, vint & inverse)
   int curr=0:
   for(int j=0; j<n; j++)</pre>
       //find first row with nonzero component in
ith column
       int index=-1;
       for(int i=curr;
i < n; i++ if(A[i][j]!=0) \{index = i; break; \}
       if(index==-1)continue;
       //swap that row
       swap(A[curr], A[index]);
       swap(invA[curr] , invA[index]);
       //make Leading 1
       int divide = inverse[A[curr][j]];
       for(int k=j;k<n;k++)A[curr][k] = A[curr][k]</pre>
* divide % p;
       for(int k=0; k<n;k++)invA[curr][k] =</pre>
invA[curr][k] * divide % p;
       //update the upper & Lower rows
       for(int i=0; i<n;i++)</pre>
           if(i!=curr){
          int mult = A[i][j];
          for(int k = j; k < n; k++) A[i][k] =
((A[i][k] - A[curr][k] *mult)%p+p)%p;
          for(int k=0;k<n;k++)invA[i][k] =
((invA[i][k] - invA[curr][k] *mult) %p+p)%p;
       curr++;
```

```
//finding inverse in modulo p
vector<vint> inverse(vector<vint> & A , int p)
   int n = sz(A);
   // identity matrix invA
   vector<vint> invA;
   for(int rep=0; rep<n;rep++)</pre>
       vint temp(n);
       invA.pb(temp);
   REP0(i,n)invA[i][i]=1;
   //preprocess inverse in moulo p
   vint inverse(p);
   inverse[1] = 1;
   for(int i=2; i<p;i++)inverse[i] = p -</pre>
(p/i)*inverse[p%i]%p;
   //Gauss Jordan elimination
   GaussJordan(A, invA,n,p, inverse);
   return invA;
희승 헤더 생략
*/
```

Geometry-1. Bulldozer

```
struct Point {
   int x, y;
   ll val;
   bool operator < (const Point &p) const {</pre>
       return pii(x, y) < pii(p.x, p.y);
};
struct Evt{
   int dx, dy, idx1, idx2;
   Evt(int dx, int dy, int idx1, int idx2) :
dx(dx), dy(dy), idx1(idx1), idx2(idx2) {}
   bool operator < (const Evt &s)const{</pre>
   11 k = 111 * dx * s.dy - 111 * dy * s.dx;
   if(k == 0) return pii(idx1, idx2) < pii(s.idx1,
s.idx2);
       return k > 0;
};
11 ccw(Evt a, Evt b) {
return (11)a.dx * b.dy - (11)b.dx * a.dy;
```

```
int main() {
    ios::sync with stdio(false); cin.tie(nullptr);
    int n;
    cin >> n;
    vector <Point> a(n);
    vector <int> pos(n);
    for(int i = 0; i < n; i++) {</pre>
        cin >> a[i].x >> a[i].y >> a[i].val;
    if(n == 1) {
    // 예외처리
        return 0;
    sort(all(a));
    for(int i = 0; i < n; i++) {</pre>
        pos[i] = i;
    vector <Evt> line;
    for(int i = 0; i < n; i++) {</pre>
        for(int j = i + 1; j < n; j++) {
            line.eb(a[j].x - a[i].x, a[j].y -
a[i].y, i, j);
    sort(all(line));
    // modify data structures
    11 \text{ ans} = 0;
    for(int i = 0; i < line.size(); i++) {</pre>
        int rx = pos[line[i].idx1], ry =
pos[line[i].idx2];
        swap(a[rx], a[ry]);
        swap(pos[line[i].idx1], pos[line[i].idx2]);
    // modify data structures
        if(i + 1 < line.size() && ccw(line[i],</pre>
line[i+1]) == 0) continue;
       //modify ans
    cout << ans << endl;</pre>
```

Geometry-2. ConvexHull

#include<bits/stdc++.h>

```
using namespace std;
#define fi first
#define se second
typedef long long 11;
typedef pair<11,11> pll;
//coding by red1108
11 cross(pll a, pll b){return a.fi*b.se-a.se*b.fi;}
vector<pll> convex hull(vector<pll> feed){
   sort(feed.begin(), feed.end());
    sort(feed.begin()+1, feed.end(), [&](pll a, pll
b){
       11 t = cross({a.fi-feed[0].fi,a.se-
feed[0].se},{b.fi-feed[0].fi,b.se-feed[0].se});
       if(t==0) return a.fi<b.fi;
       return t>0;
   });
   vector<pll> ret;
   for(pll i:feed){
       while(ret.size()>1&&cross({ret[ret.size()-
1].fi-ret[ret.size()-2].fi,ret[ret.size()-1].se-
ret[ret.size()-2].se},{i.fi-ret[ret.size()-
1].fi,i.se-ret[ret.size()-1].se})<=0)
ret.pop_back();
       ret.push back(i);
   pll i=feed[0];
   while(ret.size()>2&&cross({ret[ret.size()-
1].fi-ret[ret.size()-2].fi,ret[ret.size()-1].se-
ret[ret.size()-2].se},{i.fi-ret[ret.size()-
1].fi,i.se-ret[ret.size()-1].se})<=0)
ret.pop back();
   return ret;
// 민규야... 컨벡스헐 이렇게 드롭게 짜니 내 코드 보셈
int main(){
   vector<pll> test;
   test.emplace back(1,1);
   test.emplace back(0,0);
   test.emplace back(0,2);
   test.emplace back(2,0);
   test.emplace back(2,2);
   test = convex hull(test);
   for(auto i:test) cout<<i.first<<"</pre>
"<<i.second<<endl;</pre>
```

Geometry-3. Rotating_Callipus

```
희승 헤더 생략
*/
const int MAX = 101010;
const int INF = 1e9;
const 11 LINF = 9e18;
11 ccw(pll a, pll b, pll c) {
   return (b.fi - a.fi) * (c.se - a.se) - (b.se -
a.se) * (c.fi - a.fi);
11 ccw2(pll a, pll b, pll c, pll d) {
   return (b.fi - a.fi) * (d.se - c.se) - (b.se -
a.se) * (d.fi - c.fi);
void make convex hull(vector <pll> &A, vector <pll>
&CH) {
   sort(all(A), [&](pll a, pll b) {
       return a < b;
   sort(next(A.begin()), A.end(), [&A](pll a, pll
       return ccw(A[0], a, b) > 0;
   });
   for(auto i : A) {
       while(CH.size() >= 2 && ccw(CH[CH.size()-
2], CH[CH.size()-1], i) <= 0) CH.pop back();
       CH.eb(i);
double dist(pll a, pll b) {
   return sqrt((a.fi - b.fi) * (a.fi - b.fi) +
(a.se - b.se) * (a.se - b.se));
double ldist(pll a, pll b1, pll b2) {
   return abs((ll)(a.fi - b1.fi) * (b2.se - b1.se)
- (ll)(a.se - b1.se) * (b2.fi - b1.fi)) / dist(b1,
b2);
double farthest(vector <pll> CH) {
   int idx = 0;
   double ret = 0.0;
   for(int i = 0; i < CH.size(); i++) {</pre>
```

```
while(idx + 1 < CH.size() \&\& ccw2(CH[i]),
CH[i+1], CH[idx], CH[idx+1]) >= 0) {
            ret = max(ret, dist(CH[i], CH[idx]));
       ret = max(ret, dist(CH[i], CH[idx]));
   return ret;
int main() {
   ios::sync with stdio(false); cin.tie(nullptr);
    cout.precision(9); cout << fixed;</pre>
    int n;
    cin >> n;
   vector <pll> A(n);
   for(int i = 0; i < n; i++) {</pre>
        cin >> A[i].fi >> A[i].se;
   vector <pll> CH;
    make convex hull(A, CH);
    cout << farthest(CH) << endl;</pre>
```

Flow-0. Flow summary

```
Dinic: min (EF, V^2E) //very pessimistic
// 바깥쪽 bfs 가 O(V) 번 돌고, 주어진 Level graph 에서
O(E) 번 dfs 가 돌면서 각 Level graph 는 O(VE) 안에
처리됨.
//in unit capacity : min(V^{(2/3)} * E, E^{(3/2)})
//in unit network : O(EV^1/2) == Hopcroft-Karp
//adj 는 양방향으로 연결, residual 은 진짜 간선있는
곳에만 update, adj/residual testcase 마다 초기화,
MxN 값 vertex 보다 크게 설정.
// source sink 가 아닌 두 점 사이에서 flow 를 새로
흘리면 우회 경로를 찾을 수 있음
//Maxflow = Mincut // residual 이 0 인 놈들 기준으로
node 를 나누면 됨, 이후 실제로 있는 간선들 중에서
도달가능-도달 불가능 페어를 찾으면 됨.
//disjoint path( #of edge-disjoint path = maxflow,
vertex->vertex 선분들을 만든 후, #of vertex-disjoint
path = maxflow)
-Bipartite matching
```

// Find maximum matching in bipartite graph

```
// Hall's theorem : perfect matching <-> for
arbitary set X, |X| <= |f(X)|
//Konig's theorem : Minimum node cover = Maximum
matching
// Maximum independent set + Minimum node cover = V
//find minimum Node-disjoint path cover & find
minimum general path cover
 -Mincost Maxflow : find the minimum cost among the
maxflows.
//Maxcost Maxflow
//cost 는 양방향으로(양수, 음수) 입력,
adj,residual,cost 초기화.
//양방향 mcmf(두 노드 사이의 edge 가 cost 다른 경우) :
 정점분할 필수
 -Global Mincut, all pairs maximum flow.
 -Circulation problem
//Flow with demands : 각 정점이 demand(공급 -
 방출량)을 가지는 경우.
//1. 모든 vertex 에 대해 demand 의 합은 0 임을 확인
//2. demand 가 음수인 놈 에 newsource 가 공급,
demand 가 양수인 놈은 newsink 에 공급 후 maxflow 값이
demand 합과 같은지 확인
//유량의 하한이 존재하는 경우
//sink->source INF 간선 만들고 진행.
//u->v [L,R] : 1.change capacity to R-L, demand(u)
+=L, demand(v) -=L
//복원시에는 L 값을 다시 더해줘야함.
//maxflow 구하고 싶으면 flow with demands 풀고 원래
source MM sink \neq 3 \Rightarrow 4 M \Rightarrow
```

Flow-1. Dinic

```
/*
서용 해더 생략
MxN = 2210
*/
int residual[MxN][MxN];
vint adj[MxN];
int level[MxN];
int savepoint[MxN];
```

```
bool bfs(int n, int source, int sink)
                                          //flow
더 보내기 가능?
   REP1(i,n)level[i]=-1;
   level[source]=0;
   queue<int> q;
   q.push(source);
   while(!q.empty())
       int curr = q.front(); q.pop();
       for (auto next : adj[curr])
           if (residual[curr][next]>0 &&
level[next]==-1)
              level[next]=level[curr]+1;
              q.push(next);
   return (level[sink]!=-1); // is it possible to
reach the sink??
int dfs(int curr, int destination, int
flow) //curr -> destination 까지 flow 를 넘기는
하나의 path 찾기. 그 과정에서 residual update
   if(curr==destination)return flow;
   for (int &i= savepoint[curr];
i<adj[curr].size(); i++) //look for all adjacent</pre>
edges
       int next = adj[curr][i];
       if(level[next]==level[curr]+1 &&
residual[curr][next]>0) //Level 이 1증가하고
이동가능(residual>0) 하면. 즉 Level graph 상에서
edge 가 있으면 가본다.
          int possibleflow = dfs(next,
destination, min(flow, residual[curr][next]));
          if(possibleflow>0) //경로를
찾았다!!!!!!!
```

```
residual[curr][next] -=
                 //residual graph update
possibleflow;
               residual[next][curr] +=
possibleflow;
              return possibleflow;
                                      //한번
dfs(source, sink)부를때마다 하나의 path 를 따라서
flow 를 보낸다. 이과정을 하나의 Level graph 에 대해
불가능할때까지 반복.
   return ∅; //no possible flow from source to
sink in current level graph.
int Dinic(int n, int source, int sink)
   int maxflow=0;
   while(bfs(n, source, sink)) //there exits
path that can flow
       reset(savepoint); //we are going to
continuously work on a single level graph, so we
need to store the edges that we have traveled
before.
       while(true)
                       //for a single level graph
           int flow = dfs(source, sink, inf);
           maxflow += flow;
           if(flow==0)break;
   return maxflow;
int demand[MxN];
void solve()
   int m,n,p, pp; cin>>m>>n>>p>>pp;
   pii q[n+1];
   REP1(i,n)
       int x,y; cin>>x>>y; q[i] = \{x,y\};
   vector<pair<int, pii> > r[m+1];
   int totalk =0;
```

REP1(i,m)

```
int k; cin>>k;
       totalk += k;
       while(k--)
           int d,rp,rpp;cin>>d>>rp>>rpp;
           r[i].pb({d,{rp,rpp}});
   int source = n+m+totalk+1;
   int sink = n+m+totalk+2;
   int newsource = n+m+totalk+3;
   int newsink =n+m+totalk+4;
   //sink -> source : inf
   adj[sink].pb(source);
   adj[source].pb(sink);
   residual[sink][source]=inf;
   //source -> members
   REP1(i,m)
       adj[source].pb(i);adj[i].pb(source);residua
l[source][i]=pp-p;
       demand[source] += n-pp;
       demand[i] -= (n-pp);
   //days -> sink
   REP(i, m+1, m+n)
       adj[i].pb(sink); adj[sink].pb(i);
residual[i][sink]=q[i-m].se-q[i-m].fi;
       demand[i] += m-q[i-m].se;
       demand[sink] -= m-q[i-m].se;
   //middle nodes
   int curr = m+n+1;
   REP1(i,m)
       for(auto x : r[i])
           adj[i].pb(curr);
adj[curr].pb(curr);residual[i][curr] = x.se.se-
x.se.fi+1-x.fi;
           demand[i] += x.fi;
           demand[curr] -=x.fi;
           REP(point, x.se.fi,x.se.se)
               adj[curr].pb(point+m);adj[point+m].
pb(curr); residual[curr][point+m]=1;
```

```
curr++;
   int sumofdemand =0;
   REP1(i,n+m+totalk+2)
        if(demand[i]>0)
            sumofdemand += demand[i];
            adj[i].pb(newsink);adj[newsink].pb(i);r
esidual[i][newsink]=demand[i];
       else{
            adj[newsource].pb(i);adj[i].pb(newsourc
e);residual[newsource][i]=-demand[i];
   curr=n+m+1;
    if( sumofdemand == Dinic(n+m+totalk+4,
newsource, newsink) )
        cout<<1<<endl;</pre>
       REP1(i,m)
            vint v;
            for(auto x : r[i])
                REP(point, x.se.fi, x.se.se)
                    if(residual[curr][m+point]==0)
                        v.pb(point);
                curr++;
            cout<<sz(v)<<" ";
            for(auto x: v)cout<<x<<" ";</pre>
            cout<<endl;</pre>
   else cout<<-1<<endl;</pre>
```

```
int main()
{
    ios::sync_with_stdio(0);
    cin.tie(0);

    int t; t=1;
    while(t--)
    {
        solve();
    }
}
```

Flow-2. MCMF

```
/*
서용 헤더 생략
MxN = 2020
int residual[MxN][MxN];
vint adj[MxN];
int cost[MxN][MxN];
int dist[MxN];
int parent[MxN];
bool inQ[MxN];
bool SPFA(int n, int source, int sink)
   REP1(i,n)
       parent[i]=-1;
       dist[i]=inf;
       inQ[i]=false;
   dist[source]=0;
   parent[source]=0;
   queue<int> q;
   q.push(source);
   inQ[source]=true;
   while(!q.empty())
       int curr = q.front(); q.pop();
inQ[curr]=false;
       for (auto next : adj[curr])
           if (residual[curr][next]>0 &&
dist[next]>dist[curr]+cost[curr][next]) //if
```

```
there is path(in adj& residual>0) and distance can
be updated
               dist[next]=dist[curr]+cost[curr][ne
xt];
        //update dist
               parent[next]=curr;
      //update parent
               if(!inO[next]) //if next is not
in Q we add it.
                   q.push(next);
                   inQ[next]=true;
   return (parent[sink]!=-1); //경로가 있는지
없는지 return
pii MCMF(int n, int source, int sink)
   int maxflow=0;
   int mincost=0;
   while(SPFA(n, source, sink)) //there exits
path that can flow
       //calculate flowofpath
       int flowofpath=INT MAX;
       int curr = sink;
       while(curr!=source)
           flowofpath = min(flowofpath,
residual[parent[curr]][curr]);
           curr = parent[curr];
       //update residual graph, maxflow, mincost
       curr=sink;
       while(curr!=source)
           mincost +=
cost[parent[curr]][curr]*flowofpath;
           residual[parent[curr]][curr] -=
flowofpath;
           residual[curr][parent[curr]] +=
flowofpath;
           curr = parent[curr];
       //update maxflow
       maxflow += flowofpath;
```

```
return {mincost, maxflow}; //pair of mincost,
maxflow
void solve()
   int e,v;
   while(cin>>v>>e)
       int source = 2*v+1;
       int sink = 2*v+2;
       //source to 1
       adi[source].pb(1);
       adj[1].pb(source);
       residual[source][1] =2;
       cost[source][1] =0;
       cost[1][source]=0;
       // 2*v to sink
       adj[2*v].pb(sink);
       adj[sink].pb(2*v);
       residual[2*v][sink]=2;
       cost[2*v][sink]=0;
       cost[sink][2*v] = 0;
       //given edges
       REP1(i,e){
           int from,to, w; cin>>from>>to>>w;
           adj[from +v].pb(to);
           adj[to].pb(from+v);
           residual[from+v][to]=1;
           cost[from+v][to] = w;
           cost[to][from+v] = -w;
       //node to node
       REP(i, 1, v)
           adj[i].pb(i+v);
           adj[i+v].pb(i);
           if(i==1 || i==v)residual[i][i+v]=2;
           else residual[i][i+v]=1;
           cost[i][i+v]= 0;
           cost[i+v][i]=0;
       cout<<MCMF(2*v+2, source ,sink).fi<<endl;</pre>
```

```
REP1(i, 2*v+2)
           adj[i].clear();
           REP1(j, 2*v+2)
               residual[i][j] = 0;
               cost[i][j]=0;
int main()
   ios::sync with stdio(0);
   cin.tie(0);
   int t; t=1;
   while(t--)
       solve();
```

Flow-3. Mincut

```
if (residual[curr][next]>0 &&
canvisit[next]==false)
           dfs findcut(next, canvisit);
vector<pii> MinCut(int n, int source)
   vector<pii> cuts;
   bool canvisit[n+1];
   reset(canvisit);
   dfs findcut(source, canvisit);
   REP1(startvertex,n){
       for (auto endvertex :
realadi[startvertex])
                              //실제 간선들 중에서
print 해야함.
(canvisit[startvertex]&& !canvisit[endvertex])
cuts.pb({startvertex, endvertex});
   return cuts;
```

Flow-4. Edmonds-Karp algo

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 \mod = 1e9+7;
const int MxN = 410;
const int MxK = 1e9;
const ll inf = INT MAX;
#define REP0(i,n) for(int i=0; i< n; i++)
#define REP1(i,n) for(int i=1; i<=n; i++)</pre>
#define REP(i,a,b) for(int i=a; i<=b; i++)</pre>
#define pb push back
#define fi first
#define se second
#define reset(X) memset(X, 0, sizeof(X))
#define pii pair<int, int>
#define endl "\n"
#define vint vector<int>
```

```
#define pll pair<11, 11>
#define vll vector<ll>
#define LOG 18 //enough for 10^5
//Edmonds Karp algorithm : O(min(E^2 V),
min(EF)) : a improved version of ford-fulkerson
int residual[MxN][MxN];
vint adj[MxN];
bool visited[MxN];
int parent[MxN];
bool bfs(int source, int sink)
   reset(visited);
   visited[source]=true;
   parent[source]=0;
   queue<int> q;
   q.push(source);
   while(!q.empty())
       int curr = q.front(); q.pop();
       for (auto next : adj[curr])
           if (residual[curr][next]>0 &&
visited[next]==false)
               visited[next]=true;
               parent[next]=curr;
               q.push(next);
               if(next==sink) return
        //possible to reach sink;
true;
   return false; //no path to reach sink
int Edmonds Karp(int source, int sink)
   int maxflow=0;
   while(bfs(source, sink)) //there exits path
that can flow
   {
       //calculate flowofpath
       int flowofpath=INT MAX;
       int curr = sink;
       while(curr!=source)
```

```
flowofpath = min(flowofpath,
residual[parent[curr]][curr]);
           curr = parent[curr];
       //update residual graph
       curr=sink;
       while(curr!=source)
           residual[parent[curr]][curr] -=
flowofpath:
           residual[curr][parent[curr]] +=
flowofpath;
           curr = parent[curr];
       //update maxflow
       maxflow += flowofpath;
   return maxflow;
void solve()
   //initialize adj, and residual graph.
   int n,m; cin>>n>m;
   int source=1; int sink=2;
   reset(residual);
   REP0(i,m)
       int u, v; cin>>u>>v;
       adj[u].pb(v);
       adj[v].pb(u); //even if it is directed
graph!!!
       residual[u][v]++;
   cout<<Edmonds Karp(source, sink)<<endl;</pre>
int main()
   ios::sync with stdio(0);
   cin.tie(0);
   int t; t=1;
   while(t--)
       solve();
```

Graph-1. BCC

```
#include <bits/stdc++.h>
#define eb emplace back
#define em emplace
#define all(v) v.begin(), v.end()
#define fi first
#define se second
using namespace std;
typedef long long 11;
typedef pair <int, int> pii;
typedef pair <ll, ll> pll;
const int MAX = 101010;
const int INF = (int)1e9;
const 11 LINF = (11)1e18;
vector <int> g[MAX], stk, ng[MAX];
int ord[MAX], low[MAX], cnt, c, BCC[MAX];
void dfs(int x, int pa) {
   low[x] = ord[x] = ++cnt;
   stk.eb(x):
   bool flag = false;
   for(auto i : g[x]) {
       if(i == pa && !flag) {
           flag = true;
           continue;
       if(ord[i]) low[x] = min(low[x], ord[i]);
       else {
           dfs(i, x);
           low[x] = min(low[x], low[i]);
   if(low[x] == ord[x]) {
       while(stk.back() != x) {
           BCC[stk.back()] = c;
           stk.pop_back();
       BCC[x] = c;
       stk.pop back();
```

```
int main() {
   ios::sync with stdio(false); cin.tie(nullptr);
   int n, m;
   cin >> n >> m;
   for(int i = 1; i <= m; i++) {
       int u, v;
       cin >> u >> v;
       g[u].eb(v), g[v].eb(u);
   dfs(1, 0);
   for(int i = 1; i <= n; i++) {</pre>
       for(auto j : g[i]) {
           if(BCC[i] != BCC[j])
ng[BCC[i]].eb(BCC[j]);
```

Graph-2. SCC

```
#include <bits/stdc++.h>
#define eb emplace back
#define em emplace
#define all(v) v.begin(), v.end()
#define fi first
#define se second
using namespace std;
typedef long long ll;
typedef pair <int, int> pii;
typedef pair <ll, 11> pll;
const int MAX = 101010;
const int INF = (int)1e9;
const 11 LINF = (11)1e63;
int sz, low[505050], num[505050], scc[505050];
stack <int> stk;
vector <int> g[505050];
bool chk[505050];
int cnt;
void dfs(int x) {
   chk[x] = true;
   low[x] = num[x] = ++cnt;
   stk.push(x);
   for(auto i : g[x]) {
```

```
if(num[i] == 0) {
            dfs(i);
            low[x] = min(low[x], low[i]);
       else if(chk[i]) {
            low[x] = min(low[x], num[i]);
    if(low[x] == num[x]) {
       SZ++;
       while(!stk.empty()) {
            int u = stk.top();
            stk.pop();
            chk[u] = false;
            scc[u] = sz;
            if(x == u) break;
int main() {
    ios::sync with stdio(false); cin.tie(0);
   int n, m;
   cin >> n >> m;
   for(int i = 0; i < m; i++) {</pre>
       int u, v; cin >> u >> v;
       g[u].eb(v);
   for(int i = 1; i <= n; i++) {
       if(num[i] == 0)dfs(i);
```

Graph-3. 2-SAT

```
서용 헤더 생략
vint adj[MxN+1];
vint rev[MxN+1];
vint component[MxN+1];
vint family[MxN+1];
vint topology;
int value[MxN+1];
bool visited[MxN+1];
int parent[MxN+1];
int n,m;
```

```
int NOT(int x){return (x>n)?(x-n):(x+n);}
void dfs visit(int curr)
   visited[curr]=true;
   for(auto
next :adj[curr])if(visited[next]==false)dfs_visit(n
ext);
   topology.pb(curr);
void dfs rev(int curr, int p)
   family[p].pb(curr);
   visited[curr]=true;
   parent[curr]=p;
   for(auto next :
rev[curr])if(visited[next]==false)dfs rev(next,p);
void build scc()
   REP1(i,2*n)
       if(visited[i]==false)dfs_visit(i);
   reverse(all(topology));
   reset(visited);
   for(auto i : topology)
       if(visited[i]==false)dfs rev(i,i);
void dfs comp(int curr)
   visited[curr]=true;
   for(auto next : component[curr])
       if(visited[next]==false)dfs comp(next);
   topology.pb(curr);
void sort component()
```

```
topology.clear();
   reset(visited);
   REP1(i, 2*n)
       if(parent[i]==i && visited[i]==false)
           dfs_comp(i);
   reverse(all(topology));
void solve()
   cin>>n>>m;
   while(m--)
       int a,b; cin>>a>>b;
       if(a<0)a = NOT(-a);
       if(b<0)b=NOT(-b);
       adj[NOT(a)].pb(b);
       adj[NOT(b)].pb(a);
       rev[b].pb(NOT(a));
       rev[a].pb(NOT(b));
   build scc();
   for(int i=1; i<=2*n; i++)</pre>
       for(auto x: adj[i])
           if(parent[i]!=parent[x])
               component[parent[i]].pb(parent[x]);
   sort component();
   reverse(all(topology));
   memset(value, -1, sizeof(value));
   REP0(i,topology.size())
       for(auto x : family[topology[i]])
           if(value[x]!=-1)continue;
           if(value[NOT(x)]!=-1)
               value[x] = (value[NOT(x)]^1);
               continue;
```

```
value[x] =1;
            for(auto next : adj[x])
                if(value[next]==0)value[x]=0;
   bool pos=true;
   REP1(i,n)
       if(parent[i] == parent[NOT(i)])pos=false;
    cout<<pos<<endl;</pre>
    if(pos)REP1(i,n)cout<<value[i]<<" ";</pre>
int main()
   ios::sync with stdio(0);
   cin.tie(0);
   int t; t=1;
   while(t--)
       solve();
```

Graph-4. Offline_DynamicConnectivity

```
#include <bits/stdc++.h>
#define eb emplace_back
#define em emplace
#define all(v) v.begin(), v.end()
#define fi first
#define se second

using namespace std;
```

```
typedef long long ll;
typedef pair <int, int> pii;
typedef pair <ll, ll> pll;
const int MAX = 101010;
const int INF = (int)1e9;
const 11 LINF = (11)1e63;
struct Query {
   int t, a, b;
};
Query q[MAX];
vector <pii> E[4*MAX];
vector <int> ans;
void update(int node, int s, int e, int l, int r,
pii k) {
   if(s > r \mid\mid e < 1) return;
   if(s >= 1 \&\& e <= r) {
       E[node].eb(k);
       return;
   int m = (s + e) / 2;
   update(node*2, s, m, 1, r, k);
   update(node*2+1, m+1, e, l, r, k);
struct UFT {
   int p[MAX], rnk[MAX];
   vector <int> bp[4*MAX], br[4*MAX];
   void init(int n) {
       for(int i = 1; i <= n; i++) {
           p[i] = i;
   int find(int x) {
       if(x == p[x]) return x;
       return find(p[x]);
   void uni(int x, int y, int node) {
       x = find(x), y = find(y);
       if(x == y) return;
       if(rnk[x] < rnk[y]) swap(x, y);
       p[y] = x;
       bp[node].eb(y);
       if(rnk[x] == rnk[y]) {
           rnk[x]++;
           br[node].eb(x);
```

```
void rollback(int node) {
       for(auto i : br[node]) rnk[i]--;
       for(auto i : bp[node]) p[i] = i;
} tree;
void dnc(int node, int s, int e) {
   for(auto i : E[node]) {
       tree.uni(i.fi, i.se, node);
   if(s == e) {
       if(q[s].t == 3) {
           ans.eb(tree.find(q[s].a) ==
tree.find(q[s].b));
   else {
       int m = (s + e) / 2;
       dnc(node*2, s, m);
       dnc(node*2+1, m+1, e);
   tree.rollback(node);
int main() {
    ios::sync with stdio(false); cin.tie(0);
   int n, m;
   cin >> n >> m;
   tree.init(n);
   map <pii, int> mp;
   for(int i = 1; i <= m; i++) {</pre>
       cin >> q[i].t >> q[i].a >> q[i].b;
       if(q[i].a > q[i].b) swap(q[i].a, q[i].b);
       if(q[i].t == 1) {
           mp[{q[i].a, q[i].b}] = i;
       if(q[i].t == 2) {
           int val = mp[{q[i].a, q[i].b}];
           update(1, 1, m, val, i, {q[i].a,
q[i].b});
           mp.erase({q[i].a, q[i].b});
   for(auto i : mp) {
       update(1, 1, m, i.se, m, i.fi);
   dnc(1, 1, m);
```

```
for(int i : ans) {
    cout << i << '\n';
}</pre>
```

String-1. KMP

```
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const 11 \mod = 1e9+7;
const int MxN = 1e5;
const int MxK = 1e9;
const ll inf = INT MAX;
#define REP0(i,n) for(int i=0; i< n; i++)
#define REP1(i,n) for(int i=1; i<=n; i++)</pre>
#define REP(i,a,b) for(int i=a; i<=b; i++)
#define pb push back
#define fi first
#define se second
#define reset(X) memset(X, 0, sizeof(X))
#define pii pair<int, int>
#define endl "\n"
#define vint vector<int>
//vector<int> adj[MxN+1];
//LL dp[MxN];
//given a text and a pattern we want to print out
all the occurences!
// Time complexity O(n) in the worst case
(preprocessing -failure ftn takes O(m) time)
vector<int> failure ftn(const string& pat)
   int m = (int)pat.size();
   vector<int> f(m, ∅);
   int begin=1; int j=0;
   while(begin+j<m)</pre>
       if (pat[begin+j]==pat[j])
           j++;
           f[begin+j-1]=j;
```

```
if (j==0)
               begin++;
           else{
               begin += j-f[j-1];
               j=f[j-1];
   return f;
vector<int> KMP(const string& text, const string&
pat)
   int n = (int)text.size();
   int m = (int)pat.size();
   vector<int> f =
failure ftn(pat); //prerpocessed vector
   vector<int> find; //where we will store the
starting positions
   int begin=0; int j=0; //begin indicates the
starting position of text, j indicates the position
of pat
   while(begin <= n-m) // length condition</pre>
should be satisfied
       if (j<m && text[begin+j]==pat[j]){</pre>
           j++; // move on to the next character
           if
(j==m)find.push back(begin); //found the
matching position!
                  //mismatched
       else{
           if(j==0) begin++;
           eLse{
               begin = begin+j-f[j-1];
               j=f[j-1];
   return find;
```

```
void solve()
    string text, pat; cin>>text>>pat;
   vint v;
   v = KMP(text, pat);
   REP0(i,v.size()) cout<< v[i]<<" ";</pre>
                                           //zero
based indexing
int main()
   ios::sync with stdio(0);
   cin.tie(0);
   int t; t=1;
   while(t--)
       solve();
String-2. LCP
서용 헤더 생략
*/
void counting_sort(vint & p, vint & c)
   int n = sz(p);
   vint cnt(n);
   vint pos(n);
   for(auto x : c)cnt[x]++;
   pos[0]=0;
   for(int i=1; i<n; i++)pos[i] = pos[i-1]+cnt[i-</pre>
1];
   vint p_update(n);
   for(auto x: p)
       p_update[pos[c[x]]] = x;
```

```
pos[c[x]]++;
   p = p_update;
pair<vint, vint> suffix lcp array(string
s) //build suffix array in O(nlogn), lcp in O(n)
   s += "$";
   int n = sz(s);
   vint p(n), c(n);
   vector<pair<char, int>> temp(n);
   REPO(i,n)temp[i] = {s[i], i};
   sort(all(temp));
   REP0(i,n)p[i]=temp[i].se;
   c[p[0]]=0;
   REP1(i,n-1)
       if(temp[i].fi!=temp[i-1].fi)c[p[i]]=c[p[i-1].fi)
1]]+1;
       else c[p[i]]=c[p[i-1]];
   int k=0;
   while((1 << k) < n)
       REP0(i,n)p[i] = (p[i]-(1<< k)+n)%n;
       counting_sort(p,c);
       vint c update(n);
       c update[p[0]]=0;
       REP1(i,n-1)
           pii prev = \{c[p[i-1]], c[(p[i-1] +
(1<<k))%n]};
           pii curr = \{c[p[i]], c[(p[i] +
(1<<k))%n]};
           if(prev!=curr)c_update[p[i]]=c_update[p
[i-1]]+1;
           else c_update[p[i]]=c_update[p[i-1]];
       c = c_update;
       k ++;
   vint lcp(n);
   int skip = 0;
```

```
for(int i=0; i<n-1; i++)</pre>
        int pi = c[i];
        int j = p[pi - 1];
        while(s[i+skip]==s[j+skip])skip++;
        lcp[pi] = skip;
        skip=max(skip-1,0);
   return {p, lcp};
   // suffix array : p
   //lcp array : lcp
void solve()
   string s; cin>>s;
   vint sa,lcp;
   pair<vint, vint> ans = suffix_lcp_array(s);
   sa = ans.fi; lcp = ans.se;
   //printing out the suffix array and lcp array.
   REP1(i, sz(sa)-1)cout<<sa[i]+1<< " "; //problem</pre>
uses 1 based indexing
   cout<<endl;
   cout<<"x"<<" ";
   REP(i,2, sz(lcp)-1)cout<<lcp[i]<<" ";</pre>
int main()
   ios::sync with stdio(∅);
   cin.tie(∅);
   int t; t=1;
   while(t--)
        solve();
```

String-3. TRIE

```
민규 헤더 생략
알파벳 개수 = 26
대문자 기준 -> 'A' (기본값)
소문자 기준 -> 'a'
string 의 경우 c str()로 바꿔서 함수호출.
struct Node{
   Node* childs[26] = {NULL};
   bool finished=false;
   Node(){}
};
struct Trie{
   Node* root = new Node();
   void insert(const char* str){
       Node* cur = root;
       for (int i=0;i<strlen(str);i++)</pre>
           if (cur->childs[str[i] - 'A']==
NULL)cur->childs[str[i] - 'A'] = new Node();
           cur = cur->childs[str[i] - 'A'];
       cur->finished = true;
   bool find(const char* str){
       Node* cur = root;
       for(int i=0;i<strlen(str);i++){</pre>
           if(cur->childs[str[i]-'A'] == NULL)
return false;
           cur = cur -> childs[str[i]-'A'];
       return cur->finished;
}trie;
int main()
   trie.insert("ABC");
   trie.insert("ABCDE");
   trie.insert("ABCD");
   cout<<trie.find("ABCD");</pre>
```

String-4. Manacher

```
/*
서용 헤더 생략
//find palindrome in O(n)
string manipulate(string s)
   string changed="#";
   for(int i=0; i<(int)s.size(); i++)</pre>
       changed+=s[i];
       changed+="#";
   return changed;
vector<int> Manacher(string s) //returns longest
palindrome substring lps array.
   vector<int> lps((int)s.size()); //original
string size:N -> manipulated string size:2N+1;
   int r=0; int c=0; // [2*c-r, r] is palindrome
and r is the maximum untill now
   for(int i=0; i<(int)s.size(); i++)</pre>
       if(i \leftarrow r) lps[i] = min(r-i, lps[2*c-i]);
       else lps[i]=0;
       while((i-lps[i]-
1)>=0 && (i+lps[i]+1)<s.size() && s[i+lps[i]+1]
==s[i-lps[i]-1] )lps[i]++;
       if(r<lps[i]+i)
           r=lps[i]+i;
           c=i;
   return lps;
void solve()
   vint v= Manacher(manipulate("aaaaa"));
   REP0(i,v.size())cout<<v[i]<<endl;</pre>
```

```
int main()
   ios::sync with stdio(0);
    cin.tie(∅);
   int t; t=1;
   while(t--)
       solve();
String-5. Aho corasick
서용 헤더 생략
*/
//KMP : O(nk + Sigma(mi))
//Ahocorasick : O(n+ Sigma(mi) )
//given string s, multiple patterns m1,
m2, m3, \ldots, mk
//we decide if the pattern exists inside the string
or not
const int MAXN = 100005, MAXC = 26;
struct aho corasick{
   int trie[MAXN][MAXC], piv; // trie
   int fail[MAXN]; // failure link
   int term[MAXN]; // output check
   void init(vector<string> &v){
       memset(trie, 0, sizeof(trie));
       memset(fail, 0, sizeof(fail));
       memset(term, 0, sizeof(term));
       piv = 0;
       for(auto &i : v){
           int p = 0;
           for(auto &j : i){
               if(!trie[p][j-'a']) trie[p][j-'a']
= ++piv;
               p = trie[p][j-'a'];
           term[p] = 1;
       queue<int> que;
       for(int i=0; i<MAXC; i++){</pre>
```

if(trie[0][i]) que.push(trie[0][i]);

while(!que.empty()){

int x = que.front();

```
que.pop();
            for(int i=0; i<MAXC; i++){</pre>
               if(trie[x][i]){
                   int p = fail[x];
                   while(p && !trie[p][i]) p =
fail[p];
                   p = trie[p][i];
                   fail[trie[x][i]] = p;
                   if(term[p]) term[trie[x][i]] =
1;
                   que.push(trie[x][i]);
   bool query(string &s){
       int p = 0;
       for(auto &i : s){
            while(p && !trie[p][i-'a']) p =
fail[p];
            p = trie[p][i-'a'];
            if(term[p]) return 1;
       return 0;
}aho_corasick;
void solve()
   int n;cin>>n;
   vector<string> v;
   REP0(i,n){string s; cin>>s;v.pb(s);}
   aho corasick.init(v);
   int q; cin>>q;
   while(q--)
       string s; cin>>s;
       if(aho corasick.query(s))cout<<"YES"<<endl;</pre>
       else cout<<"NO"<<endl;</pre>
int main()
   ios::sync with stdio(0);
   cin.tie(0);
   int t; t=1;
   while(t--)
```

```
solve();
}
```

Tree-1. Centroid_Decomposition

```
#include <bits/stdc++.h>
#define eb emplace back
#define em emplace
#define all(v) v.begin(), v.end()
#define fi first
#define se second
using namespace std;
typedef long long ll;
typedef pair <int, int> pii;
typedef pair <ll, ll> pll;
const int MAX = 101010;
const int INF = (int)1e9;
const 11 LINF = (11)1e18;
vector <int> g[MAX];
int sz[MAX], crt, ctp[MAX], p[20][MAX], dep[MAX],
c[MAX];
bool chk[MAX];
void dfs(int x, int pa) {
   SZ[X] = 1;
   p[0][x] = pa;
   dep[x] = dep[pa] + 1;
   for(int i = 1; i <= 17; i++) {</pre>
       p[i][x] = p[i-1][p[i-1][x]];
   for(auto i : g[x]) {
       if(i == pa) continue;
       dfs(i, x);
       sz[x] += sz[i];
int find cen(int x) {
   int csz = 1, mx = 0, y = 0;
   for(auto i : g[x]) {
       if(chk[i]) continue;
       CSZ += SZ[i];
       if(sz[i] > mx) {
           mx = sz[i];
```

```
if(mx \le csz / 2) return x;
    sz[x] = csz - mx;
    return find cen(y);
int make ct(int x) {
    x = find cen(x);
    chk[x] = true;
    for(auto i : g[x]) {
        if(chk[i]) continue;
        ctp[make ct(i)] = x;
    return x;
int main() {
    ios::sync with stdio(false);
    cin.tie(0);
    int n;
    cin >> n;
    for(int i = 1; i < n; i++) {</pre>
        int u, v;
        cin >> u >> v;
        g[u].eb(v);
        g[v].eb(u);
    dfs(1, 0);
    crt = make ct(1);
```

Tree-2, HLD

```
#include <bits/stdc++.h>
#define eb emplace_back
#define em emplace
#define all(v) v.begin(), v.end()
#define fi first
#define se second

using namespace std;

typedef long long ll;
typedef pair <int, int> pii;
typedef pair <ll, ll> pll;

const int MAX = 101010;
```

```
const int INF = (int)1e9;
const 11 LINF = (11)1e63;
vector <pii> g[MAX];
int n:
int sz[MAX], in[MAX], nxt[MAX], p[MAX], dep[MAX],
val[MAX], po[MAX], cnt;
int tree[4*MAX];
void dfs(int x, int pa) {
   SZ[X] = 1;
   p[x] = pa;
   dep[x] = dep[pa] + 1;
   for(int i = 0; i < g[x].size(); i++) {</pre>
       int y = g[x][i].fi;
       if(y == pa) continue;
       dfs(y, x);
       SZ[X] += SZ[Y];
       po[g[x][i].se] = y;
       if(sz[y] > sz[g[x][0].fi] || g[x][0].fi ==
pa) {
           swap(g[x][i], g[x][0]);
void hld(int x) {
   in[x] = ++cnt;
   for(int i = 0; i < g[x].size(); i++) {
       int y = g[x][i].fi;
       if(y == p[x]) continue;
       if(i == 0) nxt[y] = nxt[x];
       else nxt[y] = y;
       hld(y);
void update(int node, int s, int e, int k, int x) {
   if(s == e) {
       tree[node] = x;
       return;
   int m = (s + e) / 2;
   if(k \le m) \text{ update(node * 2, s, m, k, x)};
   else update(node * 2 + 1, m + 1, e, k, x);
   tree[node] = max(tree[node * 2], tree[node * 2
+ 1]);
int cal(int node, int s, int e, int l, int r) {
   if(s > r \mid\mid e < 1) return 0;
if(s >= 1 \&\& e <= r) return tree[node];
```

```
int m = (s + e) / 2;
    int ret = max(cal(node * 2, s, m, l, r),
cal(node * 2 + 1, m + 1, e, l, r));
    return ret;
int maxquery(int u, int v) {
   int ret = 0:
   while(nxt[u] != nxt[v]) {
       if(dep[nxt[u]] > dep[nxt[v]]) swap(u, v);
       ret = max(ret, cal(1, 1, n, in[nxt[v]],
in[v]));
       v = p[nxt[v]];
    if(dep[u] > dep[v]) swap(u, v);
    if(u != v) ret = max(ret, cal(1, 1, n, in[u] +
1, in[v]));
   return ret;
int main() {
    ios::sync with stdio(false); cin.tie(0);
    cin >> n:
   for(int i = 1; i < n; i++) {</pre>
       int u, v, w;
       cin >> u >> v >> w;
       g[u].eb(v, i);
       g[v].eb(u, i);
       val[i] = w;
   dfs(1, 0);
    nxt[1] = 1;
   hld(1);
   for(int i = 1; i < n; i++) {</pre>
       update(1, 1, n, in[po[i]], val[i]);
    int q;
    cin >> q;
   while(q--) {
       int t, u, v;
       cin >> t >> u >> v;
       if(t == 1) {
           val[u] = v;
            update(1, 1, n, in[po[u]], v);
       if(t == 2) {
            cout << maxquery(u, v) << '\n';</pre>
```

DataStructure-1. LichaoTree in SegmentTree

```
#include <bits/stdc++.h>
#define eb emplace back
#define em emplace
#define all(v) v.begin(), v.end()
#define fi first
#define se second
using namespace std;
typedef long long ll;
typedef pair <int, ll> pii;
typedef pair <ll, 11> pll;
const int MAX = 303030;
const int INF = (int)1e9;
const 11 LINF = (11)9e18;
class LiChaoSegmentTree {
   struct Node {
       Node *1 = nullptr, *r = nullptr;
       pii val;
   } *tree;
public :
   11 cal(pii 1, int x) {
       return (11)1.fi * x + 1.se;
   void add line(Node *node, pii val, int s = -
INF, int e = INF + 1) {
       int m = s + e \gg 1;
       bool left = cal(val, s) > cal(node->val,
s);
       bool mid = cal(val, m) > cal(node->val, m);
       if(mid) swap(val, node->val);
       if(e == s + 1) return;
       if(left != mid) {
           if(!node->1) {
               node->1 = new Node;
               node->l->val = val;
           else add line(node->1, val, s, m);
       else {
           if(!node->r) {
               node->r = new Node;
```

```
node->r->val = val;
           else add line(node->r, val, m, e);
   void Add line(pii val) {
       if(!tree) {
           tree = new Node;
           tree->val = val;
           return;
       else add line(tree, val);
   11 cal max(Node *node, int x, int s = -INF, int
e = INF + 1) {
       int m = s + e \gg 1;
       if(e == s + 1) return cal(node->val, x);
       if(x < m) {
           if(node->1) return max(cal(node->val,
x), cal max(node->1, x, s, m));
           else return cal(node->val, x);
       else {
           if(node->r) return max(cal(node->val,
x), cal max(node->r, x, m, e);
           else return cal(node->val, x);
   11 Cal max(int x) {
       if(!tree) return -LINF;
       return cal max(tree, x);
};
pii line[MAX];
bool chk[MAX];
vector <pii> pt;
int main() {
   ios::sync with stdio(false); cin.tie(nullptr);
   int n:
   cin >> n;
   ST.Init(n);
   for(int i = 1; i <= n; i++) {
       int t;
       cin >> t;
       if(t == 1) {
```

```
cin >> line[i].fi >> line[i].se;
        chk[i] = true;
    if(t == 2) {
        int x;
        cin >> x;
        ST.Update(x, i, line[x]);
        chk[x] = false;
    if(t == 3) {
        int x;
        cin >> x;
        pt.eb(i, x);
for(int i = 1; i <= n; i++) {
    if(chk[i]) ST.Update(i, n, line[i]);
for(auto i : pt) {
    11 ans = ST.Cal_max(i.fi, i.se);
    if(ans == -LINF) cout << "EMPTY" << '\n';</pre>
    else cout << ans << '\n';</pre>
```

DataStructure-2. LineContainer

```
/**
* Author: Simon Lindholm
* Date: 2017-04-20
* License: CC0
* Source: own work
* Description: Container where you can add lines
of the form kx+m, and query maximum values at
points x.
* Useful for dynamic programming (``convex hull
trick'').
* Time: O(\Log N)
* Status: stress-tested
 */
#pragma once
struct Line {
   mutable 11 k, m, p;
   bool operator<(const Line& o) const { return k</pre>
く o.k; } // >로 바꾸면 min query
   bool operator<(ll x) const { return p < x; }</pre>
};
```

```
struct LineContainer : multiset<Line, less<>>> {
    // (for doubles, use inf = 1/.0, div(a,b) =
a/b)
    static const 11 inf = LLONG MAX;
    11 div(11 a, 11 b) { // floored division
        return a / b - ((a ^ b) < 0 && a % b); }
    bool isect(iterator x, iterator y) {
        if (y == end()) return x \rightarrow p = inf, 0;
        if (x->k == y->k) x->p = x->m > y->m ?
inf : -inf; //<로 바꾸면 min query
        else x \rightarrow p = div(y \rightarrow m - x \rightarrow m, x \rightarrow k - y \rightarrow k);
        return x \rightarrow p >= y \rightarrow p;
    void add(ll k, ll m) {
        auto z = insert(\{k, m, \emptyset\}), y = z++, x = y;
        while (isect(y, z)) z = erase(z);
        if (x != begin() && isect(--x, y)) isect(x,
y = erase(y));
        while ((y = x) != begin() && (--x)->p >= y-
>p)
            isect(x, erase(y));
    11 querv(ll x) {
        assert(!emptv());
        auto 1 = *lower bound(x);
        return 1.k * x + 1.m;
};
```

DataStructure-3. Mo's algorithm

```
/*
서용 해더 생략
*/

/* Mo's Algo
1.offline 2.static 3. processing
insertion/deletion in active range takes f(n).
Mo's algorithm Time complexity:
0( (Q+N)sqrt(N)*F + QlogQ )
i~j 에 존재하는 서로 다른 수의 개수 (online:
Mergesorttree or pst)
i~j 에 가장 많이 등장한 수
i~j 에 가장 많이 등장한 수
*/
```

```
query
                  //
int q;
int sqrtQ;
struct Query
   int index, left, right;
   bool operator< (const Query & q1)</pre>
       if (left/sqrtQ!=q1.left/sqrtQ)return
left/sqrtQ<q1.left/sqrtQ;</pre>
       return right<q1.right;</pre>
vector<Query> queries;
vint ansforquery;
void init()
   sqrtQ = sqrt(q);
   ansforquery.resize(q);
   sort(all(queries));
int ar[MxN+1];
vint store;
int occurrence[MxN];
int pos[MxN+1];
int getindex(int x) {return
lower bound(all(store),x) - store.begin();}
void solve()
   int n; cin>>n;
   REP1(i,n){cin>>ar[i]; store.pb(ar[i]);}
   compress(store);
   REP1(i,n)pos[i] = getindex(ar[i]);
   cin>>q;
   REP0(i,q)
       int s,e; cin>>s>>e;
       queries.pb({i,s,e});
   init();
```

```
int activeleft = queries[0].left; int
activeright = queries[0].left-1; int ans=0;
   REP0(i,q)
       int left = queries[i].left; int right =
queries[i].right; int index = queries[i].index;
       while(activeright < right)</pre>
            activeright++;
            if(occurrence[pos[activeright]]
==0 )ans++;
            occurrence[pos[activeright]] ++;
       while(left < activeleft)</pre>
            activeleft --;
            if(occurrence[pos[activeleft]]==0)ans
++;
            occurrence[pos[activeleft]]++;
       while(activeright>right)
            occurrence[pos[activeright]] -- ;
            if(occurrence[pos[activeright]]==0)ans
            activeright--;
       while(activeleft<left)</pre>
            occurrence[pos[activeleft]]--;
            if(occurrence[pos[activeleft]]==0)ans-
-;
            activeleft++;
       ansforquery[index] = ans;
   REP0(i, q)cout<<ansforquery[i]<<endl;</pre>
int main()
   ios::sync with stdio(0);
   cin.tie(∅);
   int t; t=1;
```

```
while(t--)
{
          solve();
}
```

DataStructure-4. Pbds

```
#include<bits/stdc++.h>
#include <bits/extc++.h>
using namespace std;
using namespace gnu pbds;
template<class T>
using ordered set = tree<T, null type, less<T>,
rb_tree_tag,tree_order_statistics_node_update>;
ordered set<int> X;
int main()
   X.insert(5);
   X.insert(3);
   X.insert(10);
   X.insert(15);
   cout<<*X.find by order(1)<<endl; //10 (0-based)</pre>
   X.erase(X.find(10));
    cout<<*X.find by order(2)<<endl; //15 (0-based)</pre>
   for(auto i:X){ // 3 5 15
        cout<<i<<" ":
```

DataStructure-5. SegmentTreeBeats

```
#include <bits/stdc++.h>
#define eb emplace_back
#define em emplace
#define all(v) v.begin(), v.end()
#define fi first
#define se second

using namespace std;
typedef long long ll;
```

```
typedef pair <int, int> pii;
typedef pair <ll, 1l> pll;
const int MAX = 101010;
const int INF = 1e9:
const 11 LINF = 1e18;
// 수열과 쿼리 19 정답 코드임. 어차피 세그 비츠는 어떤
특성값 C 를 정해서 세그먼트 트리 노드 전체의 C 값에 대한
고찰이 필요함.
// 단지 이런 느낌으로 코딩하면 된다는 참고용
class SegmentTree {
   struct Node {
       11 sum, mn, mx;
       Node(11 \text{ sum} = 0, 11 \text{ mn} = \text{LINF}, 11 \text{ mx} = -
LINF) : sum(sum), mn(mn), mx(mx) {}
       Node operator + (Node b) {
           Node ret:
           ret.sum = sum + b.sum;
           ret.mx = max(mx, b.mx);
           ret.mn = min(mn, b.mn);
           return ret;
   };
   vector <Node> tree;
   vector <pll> lazy;
   int n;
public:
   11 div(11 a, 11 d) {
       if(a >= 0) return a / d;
       else return (a - (d - 1)) / d;
   void update lazy(int node, int s, int e) {
       tree[node].sum = tree[node].sum *
lazy[node].fi + lazy[node].se * (e-s+1);
       tree[node].mx = tree[node].mx *
lazy[node].fi + lazy[node].se;
       tree[node].mn = tree[node].mn *
lazy[node].fi + lazy[node].se;
       if(s != e) {
           lazy[node<<1].fi = lazy[node<<1].fi *</pre>
lazy[node].fi;
           lazy[node<<1].se = lazy[node<<1].se *</pre>
lazy[node].fi + lazy[node].se;
           lazy[node<<1|1].fi = lazy[node<<1|1].fi</pre>
* lazv[node].fi;
           lazy[node<<1|1].se = lazy[node<<1|1].se</pre>
* lazy[node].fi + lazy[node].se;
```

```
lazy[node] = \{1, 0\};
   void Init(int N, vector <ll> &A) {
       n = N;
       tree.resize(4*n+10);
       lazy.resize(4*n+10);
       init(1, 0, n-1, A);
   void init(int node, int s, int e, vector <11>
&A) {
       lazy[node] = \{1, 0\};
       if(s == e) {
            tree[node].sum = tree[node].mn =
tree[node].mx = A[s];
            return;
       int m = s + e \gg 1;
       init(node<<1, s, m, A);</pre>
       init(node<<1 | 1, m+1, e, A);
       tree[node] = tree[node<<1] +</pre>
tree[node<<1|1]:
   void divide(int 1, int r, ll d) {
       update d(1, 0, n-1, 1, r, d);
   void update d(int node, int s, int e, int l,
int r, ll d) {
       update lazy(node, s, e);
       if(s > r \mid\mid e < 1) return;
       if(s >= 1 \&\& e <= r \&\& div(tree[node].mx,
d) == div(tree[node].mn, d)) {
            lazy[node].fi = 0;
            lazy[node].se = div(tree[node].mx, d);
            update lazy(node, s, e);
            return:
       if(s >= 1 && e <= r && tree[node].mx -</pre>
tree[node].mn == 1) {
            lazy[node].se = div(tree[node].mn, d) -
tree[node].mn;
            update_lazy(node, s, e);
            return;
       int m = s + e \gg 1;
       update d(node<<1, s, m, l, r, d);
       update d(node<<1|1, m+1, e, l, r, d);
```

```
tree[node] = tree[node<<1] +</pre>
tree[node<<1|1];
    void add(int 1, int r, ll s) {
        update s(1, 0, n-1, 1, r, s);
    void update s(int node, int s, int e, int l,
int r, 11 c) {
        update lazy(node, s, e);
        if(s > r \mid\mid e < 1) return;
        if(s >= 1 \&\& e <= r) {
            lazv[node].se = c;
            update lazy(node, s, e);
            return;
        int m = s + e \gg 1;
        update s(node <<1, s, m, l, r, c);
        update s(node<<1|1, m+1, e, l, r, c);
        tree[node] = tree[node<<1] +</pre>
tree[node<<1 | 1];
    11 sum(int 1, int r) {
        return cal(1, 0, n-1, 1, r).sum;
    11 mn(int 1, int r) {
        return cal(1, 0, n-1, 1, r).mn;
    Node cal(int node, int s, int e, int l, int r)
        update lazy(node, s, e);
        if(s > r \mid\mid e < 1) return Node();
        if(s >= 1 \&\& e <= r) return tree[node];
        int m = s + e \gg 1;
        return cal(node<<1, s, m, l, r) +</pre>
cal(node << 1 | 1, m+1, e, l, r);
} ST;
int main() {
    ios::sync with stdio(false); cin.tie(0);
    int n, q;
    cin >> n >> q;
    vector <1l> A(n);
    for(int i = 0; i < n; i++) {
        cin >> A[i];
```

```
ST.Init(n, A);
while(q--) {
   int t, 1, r;
   cin >> t >> 1 >> r;
   if(t == 1) {
        11 c;
        cin >> c;
        ST.add(1, r, c);
   if(t == 2) {
        11 d;
        cin >> d;
        ST.divide(1, r, d);
   if(t == 3) {
        cout << ST.mn(1, r) << '\n';
   if(t == 4) {
        cout << ST.sum(1, r) << '\n';</pre>
```

DataStructure-6. SplayTree

```
희승 헤더 생략
*/
template <
    class T,
    class Container = vector <T>,
    class Compare = less <T>
> class SplayTree {
    struct Node{
        Node *1, *r, *p;
        int c;
        T val;
    } *tree;
public :
   void upd(Node *x) {
        x->c = 1;
        if(x\rightarrow 1) x\rightarrow c += x\rightarrow 1\rightarrow c;
        if(x->r) x->c += x->r->c;
```

```
void rotate(Node *x) {
        Node *p = x \rightarrow p;
        Node *m;
        if(x == p->1) {
             p->1 = m = x->r;
             x->r = p;
        else {
             p->r = m = x->1;
             x - > 1 = p;
        if(m) \quad m->p = p;
        x->p = p->p;
        p \rightarrow p = x;
        if(x->p) {
             if(p == x->p->1) x->p->1 = x;
             else x \rightarrow p \rightarrow r = x;
        else tree = x;
        upd(p);
        upd(x);
    void splay(Node *x) {
        while(x->p) {
             Node *p = x->p, *pp = p->p;
             if(pp) {
                  if((pp->1 == p) == (p->1 == x))
rotate(p);
                  else rotate(x);
             rotate(x);
    void splay k(int k) {
        Node *x = tree;
        while(1) {
             while(x - > 1 && x - > 1 - > c > k) x = x - > 1;
             if(x\rightarrow 1) k \rightarrow x\rightarrow 1\rightarrow c;
             if(!k) break;
             k--;
             x = x - r;
             splay(x);
    void splay itv(int s, int e) {
        splay k(s - 1);
        Node *tmp = tree;
        tree = tmp->r;
```

```
tree->p = nullptr;
        splay k(e - s + 1);
        tmp->r = tree;
        tree->p = tmp;
        tree = tmp;
    void insert(T val) {
        Node *p = tree;
        Node **pp;
        if(!p) {
            Node *x = new Node;
            tree = x;
            x->1 = x->r = x->p = NULL;
            x \rightarrow val = val;
            return;
        while(1) {
            //if(val == p->val) return; //중복 포함
or 미포함
            if(Compare(val, p->val)) {
                 if(!p->1) {
                    pp = &p->1;
                    break;
                p = p - > 1;
            else {
                 if(!p->r) {
                    pp = &p->r;
                    break;
                p = p - r;
        Node *x = new Node;
        *pp = x;
        x->1 = x->r = NULL;
        x \rightarrow p = p, x \rightarrow val = val;
        splay(x);
    bool find(T val) {
        Node *p = tree;
        if(!p) return false;
        while(p) {
            if(val == p->val) break;
            if(compare(val, p->val)) {
                if(!p->1) break;
                p = p \rightarrow 1;
```

```
p = p - > r;
        splay(p);
        return val == p->val;
    void delete(int key){
        if(!find(key)) return;
        node* p = tree;
        if(p->1 && p->r) {
             tree = p->1; tree->p = NULL;
             node* x = tree;
            while(x->r) x = x->r;
            x->r = p->r; p->r->p = x;
            delete p;
            return;
        if(p\rightarrow 1) {
            tree = p->1; tree->p = NULL;
            delete p;
             return;
        if(p->r) {
            tree = p->r; tree->p = NULL;
             delete p;
            return;
        delete p; tree = NULL;
    void init(Container V) {
        int n = V.size();
        v.push back((T)0);
        Node *x = new Node;
        tree = x;
        x->1 = x->r = x->p = nullptr;
        x->c = n + 2:
        x \rightarrow val = (T)0;
        for(int i = 0; i <= n; i++) {</pre>
            x->r = new Node;
            x \rightarrow r \rightarrow p = x;
            x = x->r;
             x \rightarrow 1 = x \rightarrow r = nullptr:
            x - > c = n + 1 - i;
            x \rightarrow val = V[i];
};
```

if(!p->r) break;

```
int main() {
    ios::sync_with_stdio(false);

    string s;
    cin >> s;
    SplayTree <int> ST;
}
```

DataStructure-7. Dominator tree

```
서용 헤더 생략
struct dominator{
   vint dtree[MxN+2]; //dominator tree
   vint adj[MxN+2]; //original graph
   vint reverse[MxN+2]; //reverse graph (in
index. everything below is in index)
   vint nodes having semidom[MxN+2];
   int dom[MxN+2],
sdom[MxN+2],parent[MxN+2];
                                 //idom, sdom,
parent in dfs tree
   int link[MxN+2],label[MxN+2]; //about dsu ,
label[i] stores vertex with min sdom Lying on
i~root
   int indexof[MxN+2], valueof[MxN+2];
//mapping (i <-> index in dfs tree )
   int T=0; //used for counting index of dfstree
   dominator()
       REP0(i,MxN+2)
           dom[i] =
0;sdom[i]=0;parent[i]=0;link[i]=0;
           label[i]=0; indexof[i]=0;valueof[i]=0;
   void build()
       initialize(0); //build dfstree from
source(node 0)
       calcsdom();
       calcidom();
```

```
int Find(int u,int x=0) //return minimum
sdom Lying on path v to r
   {
       if(u==link[u])return x?-1:u;
       int v = Find(link[u], x+1);
       if(v<0)return u;
       if(sdom[label[link[u]]]<sdom[label[u]])</pre>
           label[u] = label[link[u]];
       link[u] = v;
       return x?v:label[u];
   void Union(int u,int v){ //Merge dsu tree where
u and v is placed
       link[v]=u;
   void initialize(int node)
       T++;
       indexof[node] =T; valueof[T] = node;
       label[T]=T;
       sdom[T] = T;
       link[T] =T;
       for (auto next : adj[node])
            if(indexof[next]==0)
               initialize(next);
               parent[indexof[next]] =
indexof[node]:
           reverse[indexof[next]].pb(indexof[node]
);
   void calcsdom()
       int n=T;
       for(int i=n;i>=1;i--)
           for(int
j=0;j<sz(reverse[i]);j++)sdom[i] =</pre>
min(sdom[i],sdom[Find(reverse[i][j])]);
           if(i>1)nodes having semidom[sdom[i]].pb
(i);
           for(int
j=0;j<sz(nodes having semidom[i]);j++)</pre>
               int w =
nodes having semidom[i][j],v = Find(w);
```

```
if(sdom[v]==sdom[w])
dom[w] = sdom[w];
               else dom[w] = v;
           if(i>1)Union(parent[i],i);
   void calcidom()
       int n=T;
       REP(i,2,n)
           if(dom[i] !=sdom[i])dom[i] =
dom[dom[i]];
           //make dfstree
           dtree[valueof[i]].pb(valueof[dom[i]]);
           dtree[valueof[dom[i]]].pb(valueof[i]);
};
void solve()
   int n,m,k; cin>>n>>m>>k;
   //build graph. don't forget source node
   dominator D;
   REP0(i,k){int u,v; cin>>u>>v; D.adj[u].pb(v);}
   int source =0;
   REP1(i,m)
       D.adj[source].pb(i);
   D.build();
   cout<<sz(D.dtree[source])<<endl;</pre>
int main()
   ios::sync with stdio(0);
   cin.tie(∅);
   int t; t=1;
   while(t--)
```

```
solve();
}
```

Optimization-1. DnC Opt

```
서용헤더 생략
Knuth Optimization
Recurrence : dp[i][j] =
min[i <= k < j](dp[i][k]+dp[k+1][j]+C[i][j])
Condition. #1. C[i][j] is Monge array
#2.C[a][d]>=C[b][c] (a<=b<=c<=d)
Following Property holds. opt[i][j-
1]<=opt[i][j]<=opt[i+1][j].
Using the property, we calculate dp[i][j] in the
order of increasing j-i, and we limit the range of
This leads to naive O(n^3) \rightarrow O(n^2)
Divide and Conquer Optimization: O(mn^2) -> O(m
nlogn)
Recursion: dp[i][j] = Min[k] (dp[i-1][k]
+C[k][j])
Condition : (고정된 i 에서) j 일때 최소값을 주는 k 를
opt(j) 라 할때 optj 가 monotonic.
각각의 dp[i][1~n]를 O(nlogn)에 calculate.
만약 k<=j & C[k][j]가 Monge array 라면 optj 가
monotonic 임을 보일 수 있음.
*/
ll val[8001];
ll sum[8001];
ll dp[801][8001];
long long C(int k, int j) //C[k][j] 값을 계산
   return (j-k)*(sum[j]-sum[k]);
void compute(int 1, int r, int opt1, int optr, int
i) //compute dp[i][l] ~ dp[i][r]을 찾는다. 가능한
if(1>r)return;
```

```
int mid = (1+r)/2;
   pair<ll, int> best = {LLONG MAX, -
     //min 을 구하는 것이므로 LLONG MAX
   for (int k = optl; k <=min(mid-1,optr); k++)</pre>
       best = min(best, \{dp[i-1][k] + C(k, mid),
k});
   dp[i][mid] = best.first;
   int opt = best.second;
   compute(l, mid - 1, optl, opt,i);
   compute(mid + 1, r, opt, optr,i);
void solve()
   int n,m; cin>>n>>m;
   REP1(i,n)cin>>val[i];
   REP1(i,n)sum[i] = sum[i-1]+val[i];
   //initialize:
   REP0(i, m+1)dp[i][0]=0;
   REP0(i,n+1)dp[1][i]=i*sum[i];
   //fill the dp array
   for(int i=2; i<=m; i++)</pre>
       for(int j=1; j<i; j++)</pre>
           dp[i][j] = sum[j];
       compute(i,n,1, n-1,i); //fill dp[i][1~n]
   cout<<dp[m][n]<<endl;</pre>
int main()
   ios::sync with stdio(∅);
   cin.tie(0);
   int t; t=1;
   while(t--)
       solve();
```