

# Numerical sheet music analysis, L3 intership (CNAM / INRIA) 27/05/24 - 02/08/24

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## I. INTRODUCTION

We present here some results regarding the analysis of tempo curve of musical performances, with score-based and scoreless approaches extending previously existing models.

The [Music Information Retrieval \(MIR\)](#) community focuses on three ways to compute musical information. The first one is raw audio, either recorded or generated, encoded in .wav or .mp3 files. The computation is based on a physical understanding of signals, using audio frames and spectrum, and represents the most common and accessible type of data. The second is a more musically-informed format, that indicates mainly two parameters : pitch (i.e., the note that the listener hear) and duration, encoded within a .mid (or MIDI) file. Such a file can be displayed as a piano roll, that is a graph which x-axis is time and y-axis is pitch (hence, the y-axis is discrete). The last way to encode musical information is the computed counterpart of sheet music. A sheet music is a way to write down a musical score, that is usually computed as a .music\_xml file, mainly for display purposes. It comes with a symbolic and abstract notation for time, that only describes the length of events in relation to a specific abstract unit, called a [beat](#), and the pitch of each event. This kind of data is actually the least common and accessible.

To actually play a sheet music, one needs a given [tempo](#), usually indicated as the amount of beat per minute (BPM). Therefore, the notion of tempo allows to translate symbolic notation (expressed in musical unit, e.g., beats) to real time events (expressed in real time unit, e.g., seconds). We will discuss later on a formal definition of tempo. However, tempo itself is insufficient to describe an actual performance of a sheet music, i.e., the sequence of real time events. Indeed, S. D. Peter, C. E. Cancino-Chacón, E. Karystinaios, and G. Widmer [1] present four parameters, among which tempo and [articulation](#) appear the most salient in contrast with [velocity](#) and timing. The latter represents the delay between the theoretical real time onset according to the current tempo, and the actual onset heard in the performance. Even though such a delay is inevitable for neurological and biological reasons, those timings are usually overemphasized and

understood as part of the musical expressivity of the performance.

In this study, we shall focus mainly on tempo estimation for a given performance recorded as a MIDI file, on both a local and global level.

## II. STATE OF ART

Even though the community studies the four parameters, the hierarchy exposed by [1] embodies quite well their relative priority within litterature. O. F. B. Katerina Kosta Rafael Ramírez and E. Chew [2] present results pointing that, although velocities don't help to meaningfully estimate tempo, the latter allows to marginally upgrade velocity-related predictions. Actually, velocity appears to be more of a score parameter rather than a performance one : automatic learning methods trained on performances of a single piece showed much better results when asked to predict velocities employed by another performer on the same piece than when trained on other performances of the same performer.

Tempo and related works actually hold a prominent place in litterature. Direct tempo estimation was first computed based on probabilistic models (C. Raphael [3], E. Nakamura, N. Ono, S. Sagayama, and K. Watanabe [4], [5]), and physical / neurological models (E. W. Large and M. R. Jones [6], H.-H. Schulze, A. Cordes, and D. Vorberg [7]) ; before the community tried neural network models [2] and hybrids approaches (K. Shibata, E. Nakamura, and K. Yoshii [8]). As the majority of previous examples, we shall focus here on mathematically and/or musically explainable methods.

Since tempo needs a symbolic representation to be meaningful, one can consider transcription as a tempo-related work. We will keep this discussion for Section IV.D and Section VI.

However, note-alignment, that is matching each note of a performance with those indicated by a given score is a very useful preprocessing technique, especially for direct tempo estimation and further analysis, such as [9]–[11]. Two main methods are to be found in litterature : a dynamic programming algorithm, equivalent to finding a shortest path (M. Müller [12]), that can works on raw audio (.wav files) ; and

a Hidden Markov Model (E. Nakamura, K. Yoshii, and H. Katayose [13]) that needs more formatted data, such as MIDI files.

In this report, we will present the following contributions :

- a justified proposition for a formal definition of tempo based on C. Raphael [3], [9] and P. Hu and G. Widmer [11] (III.A); and some immediate consequences (III.B)
- a revision of E. W. Large and M. R. Jones [6] and H.-H. Schulze, A. Cordes, and D. Vorberg [7] to fit a score-based approach (III.C)
- a theoretical framework for scoreless tempo estimation (IV.B)
- an extension of G. Romero-García, C. Guichaoua, and E. Chew [14], to fit tempo estimation, and some new theoretical results (IV.D, IV.E and Appendix C)
- generated data based on F. Foscarin, A. Mcleod, P. Rigaux, F. Jacquemard, and M. Sakai [15] and S. D. Peter *et al.* [16] (Section V.A)

### III. SCORE-BASED APPROACHES

#### A. Formal considerations

Since we chose to focus on MIDI files, we will represent a performance as a strictly increasing sequence of events  $(t_n)_{n \in \mathbb{N}}$ , each element of whose indicates the onset of the corresponding performance event. Such a definition is very close to an actual MIDI representation.

For practical considerations, we will stack together all events whose distance in time is smaller than  $\varepsilon = 20$  ms. This order of magnitude, calculated by E. Nakamura, T. Nakamura, Y. Saito, N. Ono, and S. Sagayama [5] represents the limits of human ability to tell two rhythmic events apart, and is widely used within the field [8]–[11], [13]–[17].

Likewise, a sheet music will be represented as a strictly increasing sequence of events  $(b_n)_{n \in \mathbb{N}}$ . In both of those definition, the terms of the sequence do not indicate the nature of the event (chord, single note, rest...). Moreover, in terms of units,  $(t_n)$  corresponds to real onset, thus expressed in seconds, whereas  $(b_n)$  corresponds to theorical or symbolic onsets, expressed in beats.

With those definitions, let us formally define tempo  $T(t)$  so that, for all  $n \in \mathbb{N}$ ,  $\int_{t_0}^{t_n} T(t) dt = b_n - b_0$ .

Appendix A shows that this definition is equivalent to :  $\forall n \in \mathbb{N}$ ,  $\int_{t_n}^{t_{n+1}} T(t) dt = b_{n+1} - b_n$ . However, tempo is only tangible (or observable) between two events *a priori*. We will then define the canonical tempo  $T^*(t)$  so that :

$$\forall x \in \mathbb{R}^+, \forall n \in \mathbb{N}, x \in [t_n, t_{n+1}[ \Rightarrow T^*(x) = \frac{b_{n+1} - b_n}{t_{n+1} - t_n}.$$

The reader can verify that this function is a formal tempo according to the previous definition. From now on, we will consider the convention :  $t_0 = 0$  (s) et  $b_0 = 0$  (beat).

Even though there is a general consensus in the field as for the interest and informal definition of tempo, several formal definitions coexist within litterature : K. Shibata, E. Nakamura, and K. Yoshii [8] and E. Nakamura, N. Ono, S. Sagayama, and K. Watanabe [4] take  $\frac{1}{T^*}$  as definition ; C. Raphael [3], [9] et P. Hu and G. Widmer [11] choose similar definitions than the one given here (approximated at the scale of a *measure* or a section for instance).

$T^*$  has the advantage to coincide with the tempo actually indicated on traditional sheet music (and therefore on .music\_xml format), hence allowing a simpler and more direct interpretation of results.

#### B. Naive use of formalism

As said in introduction, the more formatted data, the less accessible it is ; and the field contains only a few datasets containing both sheet music and corresponding audio, more or less anoted with various labels [9]–[11], [15], [16].

In our study, we chose to rely on the (n)-ASAP dataset [16] that presents a vast amount of performances, with over 1000 different pieces of classical music, all note-aligned with the corresponding score. From there, we can easily compute our definition of tempo. Figure 1 presents the results for a specific piece of the (n)-ASAP dataset with a logarithmic y-scale, that contains a few brutal tempo change, whilst maintaining a rather stable tempo value in-between.

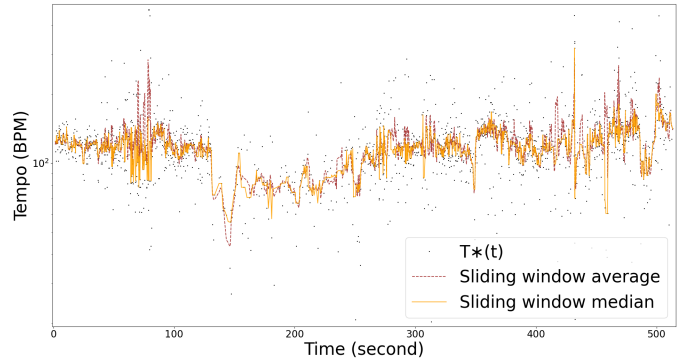


Figure 1: Tempo curve for a performance of Islamey, Op.18, M. Balakirev, with naive algorithm

In this graph, one can notice how  $T^*$  (plotted as little dots) appears noisy over time; even though allowing to distinguish a tempo change at  $t_1 = 130$  s and  $t_2 = 270$  s. Both the sliding window average (dotted line) and median (full line) of  $T^*$  seem unstable, presenting undesirable peaks, whereas the “feeled” tempo is quite constant for the listener, although the median line is a bit more stable than the average line, as expected. There are two explanation for those results. First, fast events are harder to play exactly on time, and the very definition being a ratio with a small theorical value as the denominator explains the deviation and absurd immediate tempo plotted. In fact, we can read that about 10 points are

plotted over 400 BPM (keep in mind that usual tempo are in the range 40 - 250 BPM). Second, the notion of timing and tempo are mixed together in this computation, hence giving results that do not match the listener feeling of a stable tempo. Actually, timing can be seen as little modifications to the “official” score, and using the resulting score would allow for curves that fit better the listener feeling, though needing an actual transcription of the performance first.

### C. Physical models

Among the tasks needing tempo estimation, the problem of real time estimation to allow a dedicated machine to play an accompagnement by following at least one real musician has been tackled by various approaches in litterature. C. Raphael [3] started with a probabilistic model, but those methods have found themselves replaced by a more physical understanding of tempo *via* the notion of internal pulse, as explained by E. W. Large and M. R. Jones [6]. In fact, their method has recently been developped to a commercial form<sup>1</sup>, based on an a previous adaption of the latters by A. Cont [18].

The approach developped by E. W. Large and M. R. Jones [6] consider a simplified neurological model, where listening is a fundamentally active process, implying a synchronization between external events (those of the performance) and an internal oscillator, whose complexity depends of hypothesis on the shape of the first ones. The model consists of two equations for the internal parameters:

$$\begin{aligned}\Phi_{n+1} &= \left[ \Phi_n + \frac{t_{n+1} - t_n}{p_n} - \eta_\Phi F(\Phi_n) \right] \bmod_{[-0.5, 0.5]} 1 \quad (1) \\ p_{n+1} &= p_n (1 + \eta_p F(\Phi_n)) \quad (2)\end{aligned}$$

Here,  $(\Phi_n)$  corresponds to the phase, or rather the phase shift between the oscillator and the external events, and  $(p_n)$  embodies its period. Finally,  $\eta_p$  and  $\eta_\Phi$  are both constant parameters. This initial model is then modified to consider a notion of attending *via* the  $\kappa$  parameter, whose value change over time according to other equations. The new model contains the same formulas, with the following definition for  $F$

$$F : \Phi, \kappa \rightarrow \frac{\exp(\kappa \cos(2\pi\Phi)) \sin(2\pi\Phi)}{\exp(\kappa)} \cdot \frac{1}{2\pi}.$$

Even though this model shows pretty good results, has been validated through some experiments in [6], and is still used in the previously presented version (E. W. Large *et al.* [19]), a theoretical study of the system behavior remains quite complex, even in simplified theoretical cases [7], notably because of the function  $F$  expression.

In order to simplify the previous model, H.-H. Schulze, A. Cordes, and D. Vorberg [7] present *TimeKeeper*, that can be

seen as a linearization of the previous approach, valid in the theoretical framework of a metronome presenting small tempo variations. In fact, there is a strong analogy between the two models, that are almost equivalent under specific circumstances, as shown by J. D. Loehr, E. W. Large, and C. Palmer [20]. Here, we used the derandomised version presented in [20], where  $M_i = 0$  and  $T_i = \tau_i$  for all  $i \in \mathbb{N}$ .

None of those models have an inherent comprehension of musical score information, since the both rely on a rather stable metronome. In the version displayed hereafter, they were modified to consider score information, in order to to create a more stable and precise value of tempo than the naive approach previously presented. Those modifications are detailed in the following paragraph (OR IN APPENDIX ?), and were made in order to keep consistency with the original models in their initial theoretical framework of validity. Let  $\min_{\text{abs}} : a, b \mapsto \begin{cases} a & \text{if } |a| < |b| \\ b & \text{otherwise} \end{cases}$

We first modify the E. W. Large and M. R. Jones [6] equations accordingly :

$$\Phi_{n+1} = \Phi_n + \frac{t_{n+1} - t_n}{p_n} - \eta_\Phi F(\Psi_n, \kappa_n) \quad (3)$$

$$p_{n+1} = p_n (1 + \eta_p F(\Psi_n, \kappa_n)) \quad (4)$$

$$\Psi_n = -\min_{\text{abs}}(k + b_n - \Phi_n, k + 1 + b_n - \Phi_n)$$

$$k = \lfloor \Phi_n - b_n \rfloor$$

Using the analogy presented in [20], we then obtain the following equations for *TimeKeeper* :

$$A_{i+1} = K_i (1 - \alpha) + \tau_i - (t_{i+1} - t_i) \quad (5)$$

$$\tau_{i+1} = \tau_i - \beta * (K_i \bmod_{[-0.5, 0.5]} 1) \quad (6)$$

$$K_i = -\min_{\text{abs}}(k\tau + b_i - A_i, (k+1)\tau + b_i - A_i)$$

$$k = \left\lfloor \frac{A_i - b_i}{\tau_i} \right\rfloor$$

In both of those extensions, the  $\min_{\text{abs}}$  function is used in order to represent a choice between two corrections. The first argument can be interpreted as a correction with respect to the most recent passed beat time occurring exactly on a actual beat, i.e.,  $a_1 = \max_{n \in \mathbb{N} : b_n \leq B_i} \lfloor b_n \rfloor$  where  $B_i$  represents the internal value at time  $i$  acting as a beat unit ( $\Phi_i$  for E. W. Large and M. R. Jones [6] and  $A_i$  for H.-H. Schulze, A. Cordes, and D. Vorberg [7]). The second argument embodies the correction according to the following beat,  $a_2 = a_1 + 1$ .

One can notice that the phase is actually always used modulo 1 in E. W. Large and M. R. Jones [6], since it appears only multiplied by  $2\pi$  in either cos or sin functions. Using this remark, one can verify that, in the initial presentation of the model with a metronome, i.e.,  $\forall n \in \mathbb{N}, b_n = 0 \bmod 1$ , the extension proposed here is equivalent to the original approach,

<sup>1</sup><https://metronautapp.com/>

i.e., (1), (2)  $\Leftrightarrow$  (3), (4), hence justifying the designation “extension”.

Figure 2 displays the results of those two models, in regards with the canonical, or immediate tempo. One can notice that E. W. Large and M. R. Jones [6] model is less stable than *TimeKeeper*, although faster to converge.

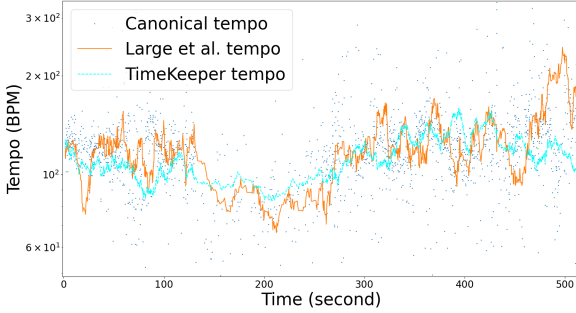


Figure 2: Tempo curve for a performance of Islamey, Op.18, M. Balakirev, according to various models

Figure 3 (below) exposes the visible difference in tempo initialization of the two models, starting both here with the initial tempo of 70 BPM ( $\text{♩} = 70$ , i.e., the *beat* unit here is a quarter note). *TimeKeeper* does not manage to converge to any significant tempo. Such a behavior was to be expected, considering the theoretical framework for *TimeKeeper*, that is small tempo variation, and correct initialization. However, Large et al model manages to converge to a meaningful result. In fact, in the range 9 to 70 seconds, the estimated tempo according to Large is exactly half of the actual tempo hinted by the blue dots (canonical tempo).

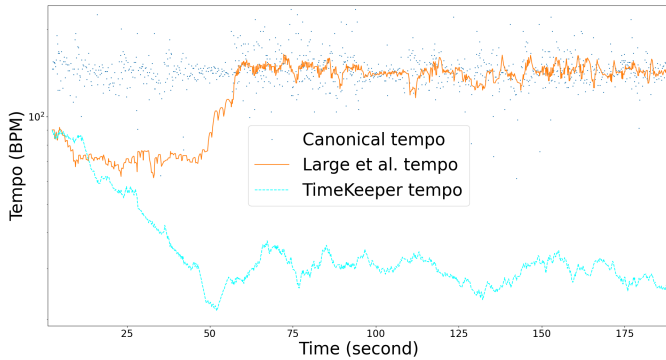


Figure 3: Tempo curve for a performance of Piano Sonata No. 11 in A Major (K. 331: III), W.A Mozart, according to the previous models with irrelevant initialization

## IV. SCORELESS APPROACHES

### A. Motivations

There are three main issues with the previous models, apart from the necessary knowledge of the sheet music, that

are : salient sensibility to tempo initialization (cf. Figure 3), unstability that requires some time to (possibly) converge (cf. Figure 2 and Figure 3), and difficulty to accurately estimate relevant values of the constant internal parameters. According to our implementation, E. W. Large and M. R. Jones [6] is a particularly chaotic model regarding the latter.

We will present here two models focusing on tackling mainly the first two issues previously presented. Those rely on a specific musical property of division : in symbolic notations of music, every single event can be comprehended as a multiple of a certain unit called a *tatum*, usually expressed in beat unit. Therefore, the real events of a performance, or rather their duration, can be interpreted as multiple of this tatum. However, considering a non-constant tempo, the real value (i.e., real duration in seconds) of this tatum may evolve through time, whereas the symbolic value remains constant anyway. Actually, detecting the tatum is equivalent to transcribe the performance to sheet music, which is a rather more complicated task than tempo estimation. For instance, there are several ways to write down sheet musics that are undistinguishable when performed. We call this ambiguity *tempo octaves* (cf Appendix A).

We considered here two quantization methods extracted from literature : D. Murphy [17] and G. Romero-García, C. Guichaoua, and E. Chew [14]. Both of these papers present a theory of approximate division, that is a way to find a rational number, interpreted as the ratio of two symbolic events expressed in arbitrary unit (for instance in tatum) from the real events durations in seconds. Such a link is equivalent to defining a tempo with the formalism presented in III. Although [17] provides an algorithm to find candidate rationals, they do not include a way to compare those candidates, thus leaving no choice but an exhaustive approach (top-down in the paper). With another formalism, G. Romero-García, C. Guichaoua, and E. Chew [14] define a graph, restrained only to consistent values of tatum. We choose to adapt the latter, although their presentation is clearly user-oriented rather than automatic, since the introduced graph allows to mathematically (and therefore automatically) define a “good” and a “best” choice among all possible found tatum values.

### B. Introduction of an estimator based approach

Given a sequence  $(u_n)_{n \in \mathbb{N}}$ , we now introduce the notation  $(\Delta u_n)_{n \in \mathbb{N}} := (u_{n+1} - u_n)_{n \in \mathbb{N}}$  for the next sections.

The reason why the previous models have to converge is because they both try to find an exact value of tempo, and therefore sudden and huge tempo changes will lead to an uncertain period for the resulting tempo estimation, that is in this way the exact same problem as tempo initialization. When doing tempo estimation, we are in fact much more interested in a local tempo variation, relative to the previ-



ous estimation, rather than an absolute value, especially on a local time scale (where we can often assume tempo to be almost constant). Using the formalism presented in III, we first present the following result since  $T_n^* > 0$ :

$$T_{n+1}^* = T_n^* \frac{T_{n+1}^*}{T_n^*} = T_n^* \frac{\Delta t_n}{\Delta t_{n+1}} \frac{\Delta b_{n+1}}{\Delta b_n} = T_n^* \frac{\Delta t_n}{\Delta t_{n+1}} \times \frac{T_{n+1}^* \Delta t_{n+1}}{T_n^* \Delta t_n}$$

Let  $T_n$  be an estimation of  $T_n^*$  by a certain given model and  $\alpha_n = \frac{T_n}{T_n^*}$ . We obtain :

$$\alpha_n T_{n+1}^* = \underbrace{\alpha_n T_n^*}_{T_n} \frac{\Delta t_n}{\Delta t_{n+1}} \times \frac{\Delta b_{n+1}}{\Delta b_n}$$

In the above formula, the only value to actually estimate is therefore  $\frac{b_{n+2}-b_{n+1}}{b_{n+1}-b_n}$ , where both the numerator and denominator are translation permissive (i.e., we can afford a locally constant shift in both of our estimation), hence the resulting value is invariant by translation, or constant multiplication of our estimation. Furthermore, the value to be estimated only deals with symbolic units, meaning that we can use musical properties to find a consistent result. As a result, we can obtain a tempo estimation with the same multiplicative shift as the previous estimation  $T_n$ , thus, by using the formula recursively, we obtain a model that can track tempo variations over time without any need of convergence, and that is robust to tempo initialization, while using only local methods (in other words, the resulting model is [online](#)). As noticed in the beginning of this section, symbolic value have some bindings that actually help to determine their values. Moreover, we find :  $\frac{\Delta b_{n+1}}{\Delta b_n} = \frac{T_{n+1}^* \Delta t_{n+1}}{T_n^* \Delta t_n} = \frac{\frac{1}{\alpha_n} T_{n+1}}{\frac{1}{\alpha_n} T_n} \times \frac{\Delta t_{n+1}}{\Delta t_n}$ ,

$$\text{hence } \frac{\Delta b_{n+1}}{\Delta b_n} = \frac{T_{n+1}}{T_n} \times \frac{\Delta t_{n+1}}{\Delta t_n}.$$

Let us then write the actual formula of the model :

$$\frac{T_{n+1}}{T_n} = \frac{\Delta t_n}{\Delta t_{n+1}} E \left( \underbrace{\frac{T_{n+1}}{T_n} \times \frac{\Delta t_{n+1}}{\Delta t_n}}_{\frac{\Delta b_{n+1}}{\Delta b_n}} \right) \quad (7)$$

where  $E$  is a function-like object (closer to a *object* in computer science than an actual mathematical function), designated by *estimator*. This function is supposed to act, on a theoretical ground, as an oracle that returns the correct value of the symbolic  $\frac{\Delta b_{n+1}}{\Delta b_n}$  from the given real values indicated in (7), therefore supposed to rarely match the theoretical values.

Given an estimator  $E$ , the tempo value defined as  $T_{n+1}$ , computed from both  $T_n$  and local data, is obtained *via* the following equation, where  $x$  represents  $\frac{T_{n+1}}{T_n}$  in (7) :

$$T_{n+1} = T_n \underset{x \in [\frac{2}{3}T_n, \frac{4}{3}T_n]}{\operatorname{argmin}} d \left( x, \frac{\Delta t_n}{\Delta t_{n+1}} E \left( x \frac{\Delta t_{n+1}}{\Delta t_n} \right) \right) \quad (8)$$

where  $d : a, b \mapsto k_* |\log(\frac{a}{b})|$ ,  $k_* \in \mathbb{R}_+$ , is a logarithmic distance, chosen since an absolute distance would have favor small values by triangle inequality in the hereunder process.

In the implementation presented here, the estimator role is more to quantify the ratio in order to output a musically relevant value. In our test, we limited these quantification to

accept only regular division (i.e., powers of 2). Furthermore, the numerical resolution for the previous equation was done by a logarithmically evenly spaced search and favor  $x$  values closer to 1 (i.e.,  $T_{n+1}$  closer to  $T_n$ ) in case of distance equality.

Such a research allows for a musically explainable result : the current estimation is the nearest most probable tempo, and both halving and doubling the previous tempo is considered as improbable, and as further going from the initial tempo. Appendix B gives further explanation about (8).

### C. Study of the model

Since this approach fundamentally search to estimate tempo variation rather than actual values, it is not easy to visualize the relevance of the result by naive means. We choose here to define  $(\alpha_n := \frac{T_n}{T_n^*})_{n \in [1, N]}$  and  $(\tilde{\alpha}_n := \exp(\ln(\frac{T_n}{T_n^*}) - \lfloor \log_2(\frac{T_n}{T_n^*}) \rfloor \ln(2)))_{n \in [1, N]}$ .  $(\tilde{\alpha}_n)_{n \in [1, N]}$  is then called the *normalized* sequence of ratio, where each value is uniquely determined within the range  $[1, 2[$ . Such a choice allows for merging together the tempo octaves, as explained in Appendix A. In this representation, we actually have  $\tilde{2} = \tilde{1}$ , and adding  $\tilde{1}$  is equivalent to multiply the initial value by 2. We will now define a *spectrum*  $S = (\tilde{\alpha}_n)_{n \in [1, N]}$ . We then call  $|S|$  the value  $N \in \mathbb{N}$  and  $\mathcal{C}$  the range  $[1, 2[$  seen as a circle according to the following application :  $c : [1, 2[ \rightarrow \mathcal{C}(0, 1)$

cation :  $x \mapsto (\cos(2\pi x), \sin(2\pi x))$ , so that  $c(1) = c(2)$

We then define the *measure* of a spectrum  $S$ , that embodies a standard deviation on  $\mathcal{C}$ , so that :

$$m(S, \Delta) = \max_{\tilde{d} \in \mathcal{C}} \frac{|\{n \in [1, |S|] : d(\tilde{\alpha}_n, \tilde{d}) \leq \Delta\}|}{|S|} \quad (9)$$

Where  $\Delta \in [0, \frac{1}{2}]$  embodies the measure accuracy, and  $d$  is still a logarithmic distance, slightly modified on  $\mathcal{C}$  to be consistent with  $d(\tilde{1}, \tilde{2}) = 0$ . Actually, it can be shown that on  $\mathcal{C}$ ,  $d : \tilde{a}, \tilde{b} \mapsto \min(|\tilde{a} - \tilde{b}|, 1 - |\tilde{a} - \tilde{b}|)$ .

The reader can verify that this measure is invariant with respect to spectrum rotation by any  $\lambda \in R_+$  (i.e.,  $m(S, \Delta) = m((\tilde{\lambda}\alpha_n)_{n \in [1, N]}, \Delta)$ ), and does not depend on the normalisation interval (here  $[1, 2[$ , but actually  $[\lambda, 2\lambda[$  would work just as well, with a different expression of  $d$  on  $\mathcal{C}$ ). Finally,  $0 \leq m(S, \Delta) \leq 1$ , and  $m(S, \Delta) = 0 \Leftrightarrow |S| = 0$ ,  $m(S, \Delta) = 1$  iff  $S$  only contains values within a  $2\Delta$  range.

This measure allows to quantify the quality of this model, without considering tempo octaves, or equivalently to quantify the quality of the estimator. Appendix B presents the C++ implementation of this measure used to obtain the following figures.

résultats sur (n-)Asap selon les périodes, les compositeurs, etc..., éventuellement en annexe ?

#### D. Towards a quantified approach

The previous model supposes the existence of an oracle, more or less correct in its predictions, that is actually a (partial) transcriber. In this section, we will focus on this transcribing part, by extending G. Romero-García, C. Guichaoua, and E. Chew [14] model with the previous formalism. In fact, we extend the previous approach by considering the estimator as our central model and then extracting tempo values rather than the opposite.

Let  $n \in \mathbb{N}^*$  and  $D \subset (R^+)^n$  be a set of some durations of real time events. [14] defined the continuous function  $\varepsilon_D$  as :

$$\varepsilon_D : a \mapsto \max_{d \in D} \min_{m \in \mathbb{Z}} |d - ma| \quad (10)$$

This function is called the *transcription error*, and can be interpreted as maximum error (in real time unit) between all real events  $d \in D$  and theoretical real duration of  $ma$ , where  $m$  is a symbolic notation expressed in arbitrary symbolic unit, and  $a$  a real time value corresponding to a *tatum* at a given tempo. We prove in Appendix C that the set of all local maxima of  $\varepsilon_D$ , except those that also are minima, is :

$$\begin{aligned} M_D &= \left\{ \frac{d}{k + \frac{1}{2}}, d \in D, k \in \mathbb{N} \right\} \\ &= \bigcup_{d \in D} \left\{ \frac{d}{k + \frac{1}{2}}, k \in \mathbb{N} \right\} \end{aligned} \quad (11)$$

In fact, each of these local maxima corresponds to a change of the  $m$  giving the minimum in the expression of  $\varepsilon_D$ , hence the following result : in-between two such successive local maxima, the quantification remains the same, i.e.,

**Proposition IV.D.1** Let  $m_1, m_2$  be two successive local maxima of  $\varepsilon_D$ ,  $a_1 \in ]m_1, m_2[$ ,  $a_2 \in [m_1, m_2]$ ,  $d \in D$  and  $m \in \mathbb{Z}$ . Then  $m \in \operatorname{argmin}_{k \in \mathbb{Z}} |d - ka_1| \Rightarrow m \in \operatorname{argmin}_{k \in \mathbb{Z}} |d - ka_2|$ .

With this property, we can then choose to consider only local minima of  $\varepsilon_D$  as in [14], since there is exactly one local minima in-between two such successive local maxima, and choosing any other value in this range would result in the exact same transcription, with a higher error by definition of a local maxima (that is global on the considered interval). The correctness of the following algorithm to find all local minima within a given interval is proven in Appendix C.

---

FindLocalMinima( $D \neq \emptyset$ , start, end) :

---

```

1  $M \leftarrow \left\{ \frac{d}{k + \frac{1}{2}}, d \in D, k \in \left[ \left\lfloor \frac{d}{\text{end}}, \frac{d}{\text{start}} \right\rfloor \right] \right\}$ 
2  $P \leftarrow \left\{ \frac{d_1 + d_2}{k}, (d_1, d_2) \in D^2, k \in \left[ \left\lfloor \frac{d_1 + d_2}{\text{end}}, \frac{d_1 + d_2}{\text{start}} \right\rfloor \right] \right\}$ 
3 localMinima  $\leftarrow \emptyset$ 
4 for  $(m_1, m_2) \in M^2$  two successive maxima :
5   | in_range  $\leftarrow \{p \in P : m_1 \leq p < m_2\}$ 
6   | localMinima  $\leftarrow \text{localMinima} \cup \{\min(\text{in\_range})\}$ 
7 return localMinima
```

---

Algorithm 1: Returns all local minima within [start, end]

G. Romero-García, C. Guichaoua, and E. Chew [14] then defined  $G = (V, E)$  a graph which vertices are the local minima of  $\varepsilon_D$  with  $D$  a sliding window, or *frame*, on a given performance, and which edges are so that they can guarantee a *consistency property*, explained hereafter.

The *consistency property* for two tatums  $a_1, a_2$  specifies that, if  $F_\cap$  is the set of all values in common between two successive frame, for all  $d \in F_\cap$ ,  $d$  is quantified the same way according to the tatum  $a_1$  and  $a_2$ , i.e., the symbolic value of  $d$  is the same when considering either  $a_1$  or  $a_2$  as the duration of the same given tatum at some tempo (respectively  $\frac{1}{a_1}$  and  $\frac{1}{a_2}$  as shown in Section IV.E). From these definitions, we can now define a *tempo curve* as a path in  $G$ . In fact, [14] call such a path a *transcription* rather than a tempo curve, but since an exact tempo curve would be  $(T_n^*)$ , those two problems are actually equivalent.

Actually, the consistency property is not that restrictive when considering tempo curves. Let  $F_1, F_2$  be two successive frames,  $F_\cap = F_1 \cap F_2$ ,  $d \in F_\cap$ , and  $p$  a path in  $G$  containing a local minima  $a_1$  of  $\varepsilon_{F_1}$ . According to (11), we can divide the set of all local minima in two, with those “caused” by  $F_\cap$ ,  $M_{F_\cap}$ , and the others. Let then  $m_1, m_2 \in M_{F_\cap}$ . Those are local maxima for both  $\varepsilon_{F_1}$  and  $\varepsilon_{F_2}$  by (11) since  $F_\cap \subset F_2$ , and therefore there is at least one local minima within the range  $]m_1, m_2[$  for both of these functions. However, thanks to Proposition IV.D.1, we know that both these local minima will quantify the elements of  $D$  the same way. Hence, by defining :

- $m_1 = \max\{m \in M_{F_\cap} : m < a_1\}$
- $m_2 = \min\{m \in M_{F_\cap} : m > a_1\}$
- $a_2$  a local minima of  $\varepsilon_{F_2}$  in the range  $]m_1, m_2[$ , which exists since  $m_1$  and  $m_2$  are local maxima (that are not local minima).

We obtain :  $(a_1, a_2)$  is *consistent* according to the consistency property.

**Corollary IV.D.2** The *consistency property* only implies restrictions relative to the interval of research. In other words,

any given strictly partial path  $p$  in  $G$  can be extended, even if it means considering a bigger interval, for any given performance, and any given frame length for defining  $G$ .

However, this restriction on  $G$  appears to have some interest. Indeed, let  $p$  a path in  $G$  locally inconsistent, i.e., such that  $a_1, a_2 \in p$  so that  $d \in D$  is quantified differently according to  $a_1$  and  $a_2$ , with  $a_1$  and  $a_2$  local minima of successive frames. We therefore have a two partial transcriptions of  $d$  being either :  $m_1$  at tempo  $\frac{1}{a_1}$  and  $m_2$  at tempo  $\frac{1}{a_2}$ ,  $m_1, m_2$  expressed in tatum unit, with  $m_1 \neq m_2$ . WHAT IS THE POINT ?

#### E. Quantification revised

Let us define from now our tatum  $\varepsilon = \frac{1}{60} \downarrow$ , which correspond to an sixteenth note wrapped within a triplet within a quintuplet, and has the property that  $1 \varepsilon / s = 1 \downarrow / m$ .

With our tatum defined, we can now choose to express all our symbolic units as multiple of this tatum, hence the unit for symbolic values is now  $\varepsilon$ . We then have :  $T = \frac{\Delta b}{\Delta t} = \frac{1 \varepsilon}{a}$ , where  $a$  is the theoretical duration of  $\varepsilon$  at tempo  $T$ . From there, we can define  $\sigma_D : a \mapsto \frac{1}{a} \varepsilon_D(a)$ , the *normalized error*, or *symbolic error*, since it embodies the error between a transcription of  $d \in D$  as  $m$  expressed in tatum, hence a quantified and valid transcription, and  $d \times \frac{1}{a} = d \times T$ , which is the expression of the symbolic duration of  $d$  at tempo  $T$  according to the [definition of canonical tempo](#).

- LR
- bidi (2 passes: LR + RL) : justification (en annexe) : retour à la définition formelle de Tempo : valide dans les deux sens, d'où la possibilité de le faire en bidirectionnel + parler rapidement d'une application à Large
- RT : avec valeur initiale de tempo

#### F. résultats évaluation (comparaison avec 3)

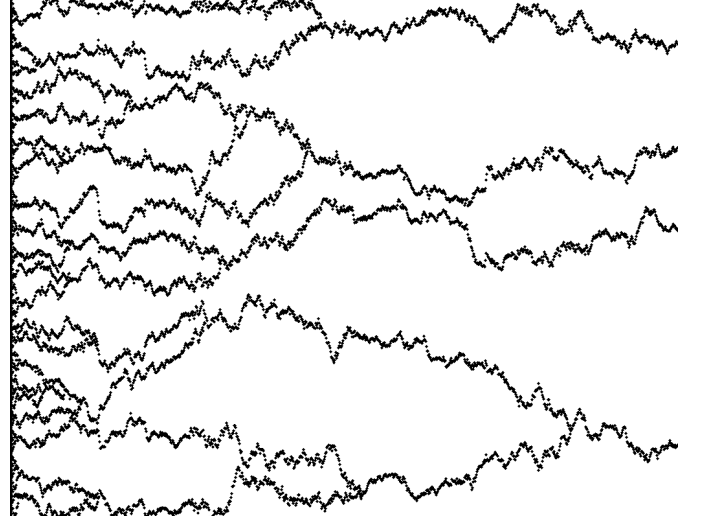


Figure 4: All possible tempo curves for a performance of Piano Sonata No. 11 in A Major (K. 331: III), W.A. Mozart. The tempo scale is linear between  $\downarrow = 40$  (bottom) and  $\downarrow = 240$  (top)

## V. APPLICATIONS

### A. Data extension

Let  $(t_n)_{n \in \mathbb{N}}, (T_n^*)_{n \in \mathbb{N}}, (T_n)_{n \in \mathbb{N}}$  be respectively a performance, a canonical tempo for this performance, an estimated tempo curve (supposed to be flattened with respect to the canonical tempo) and  $T_c \in R_+^*$  a given tempo value. We define  $(s_n := \Delta t_n (T_n - T_n^*))_{n \in \mathbb{N}}$  the symbolic shift of the performance according to  $(T_n^*)$  and  $(T_n)$ . The *normalized performance to tempo  $T_c$*  of  $(t_n)$  is then defined as :

$$\hat{t}_0 = t_0, \forall n \in \mathbb{N}, \hat{t}_{n+1} = \hat{t}_n + \underbrace{\Delta t_n \frac{T_n}{T_c}}_{\alpha_n} + \underbrace{\frac{s_n}{T_c}}_{\beta_n} \quad (12)$$

Where  $\alpha_n$  represents the new duration of  $\Delta t_n$  at tempo  $T_c$ , since  $\alpha_n = \frac{\Delta b_n}{T_c}$ , and  $\beta_n$  embodies the actual time shift at tempo  $T_c$ .

Proposition V.A.1  $(\hat{t}_n)_{n \in \mathbb{N}}$  is a performance as defined in Section III.A, and  $(\hat{T}_n^*)_{n \in \mathbb{N}} = \left( \frac{1}{1 + \frac{s_n}{\Delta b_n}} T_c \right)_{n \in \mathbb{N}}$

Proof For all  $n \in \mathbb{N}^*$ ,  $\Delta \hat{t}_n = \alpha_n + \beta_n = \Delta \frac{t_n}{T_c} (T_n^* + T_n - T_n^*) = \Delta t_n \frac{T_n}{T_c} > 0$

Furthermore,  $\hat{T}_n^* = \frac{\Delta b_n}{\Delta \hat{t}_n} = \frac{\Delta b_n}{\Delta \frac{t_n}{T_c} + \frac{s_n}{T_c}} = \frac{1}{1 + \frac{s_n}{\Delta b_n}} T_c \quad \square$

Depending on the use of this data, one can choose to adapt the definition of  $(s_n)$ , for instance by normalizing its values or cutting all shifts that represent over a certain portion of their duration. [Proposition V.A.1](#) exposes how such definitions of  $(s_n)$  allow for guarantee on tempo change and deviation of the canonical tempo for the *normalized performance* with respect to  $T_c$ .

- génération de données “performance” : pour data augmentation ou test robustesse (fuzz testing) aplanissement de tempo démo MIDI?
- transcription MIDI par parsing : pre-processing d’évaluation tempo (approche partie 4)
- analyse “musicologique” quantitative de performances humaines de réf. (à la Mazurka BL) données quantitatives de tempo et time-shifts

## VI. CONCLUSION & PERSPECTIVES

- intégration pour couplage avec transcription par parsing (+ plus court chemin multi-critère)
- Section IV.E presented a model which appear to share some similarities with E. W. Large and M. R. Jones [6] as a score-based approach. Therefore, studying the formalism for a quantifier might allow to obtain some theoretical results for the previous model, regarding for instance convergence guarantee and meaningfulness of the result.

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## APPENDIX A : SOME FORMAL CONSIDERATIONS

### A. Equivalence of tempo formal definitions

Let  $n \in \mathbb{N}$ .  $\int_{t_0}^{t_n} T(t) dt = \sum_{i=0}^{n-1} \int_{t_i}^{t_{i+1}} T(t) dt$   
 Furthermore,  $\int_{t_0}^{t_{n+1}} T(t) dt = \int_{t_0}^{t_{n+1}} T(t) dt + \int_{t_n}^{t_0} T(t) dt = \int_{t_0}^{t_{n+1}} T(t) dt - \int_{t_0}^{t_n} T(t) dt$ .

Let  $T$  be a formal tempo according to the first definition. For all  $n \in \mathbb{N}$ , we then have :

$$\begin{aligned} \int_{t_n}^{t_{n+1}} T(t) dt &= \int_{t_0}^{t_{n+1}} T(t) dt - \int_{t_0}^{t_n} T(t) dt \\ &= b_{n+1} - b_0 - (b_n - b_0) \\ &= b_{n+1} - b_n \end{aligned}$$

Let  $T$  be a formal tempo according to the second definition.

For all  $n \in \mathbb{N}$ :

$$\begin{aligned} \int_{t_0}^{t_n} T(t) dt &= \sum_{i=0}^{n-1} \int_{t_i}^{t_{i+1}} T(t) dt \\ &= \sum_{i=0}^{n-1} b_{i+1} - b_i \\ &= b_n - b_0 \end{aligned}$$

We thus obtain the two implications, hence the equivalence.

### B. The tempo octave problem

When estimating tempo, or transcribing a performance, there always exist several equivalent possibilities. For instance, given a “correct” transcription  $(b_n)$  of a performance  $(t_n)$ , one can choose to define its own transcription as  $t = \left(\frac{b_n}{2}\right)_{n \in \mathbb{N}}$ .

Then, the canonical tempo according to  $t$ , called  $(T_1^*)$ , and the one according to  $(b_n)$ , called  $(T_2^*)$  verify :

$\forall n \in \mathbb{N}, T_{1,n}^* = \frac{\frac{b_{n+1}}{2} - \frac{b_n}{2}}{t_{n+1} - t_n} = \frac{1}{2} T_{2,n}^*$ . Actually, the  $t$  transcription corrections to  $(b_n)$  where all durations are indicated doubled, but played twice faster, hence giving the exact same theoretical performance. Unfortunately, there is no absolute way to decide which of those two transcription is better than the other. This problem is here known as the tempo octave problem, and should be kept in mind when transcribing or estimating tempo. We present in Section IV.B a model resistant to these tempo octaves, as well as other kind of octaves not discussed here (for instance multiplying the tempo by 3 by using [triplet](#)).

### C. Tempo conservation when reversing time

First, we want to insist on the fact that none of the sequence  $(b_n)$  and  $(t_n)$  are infinite, but in order to simplify the notation, we chose to indicate them as usual infinite sequences, or rather only consider them on a finite number of indexes, called  $|(b_n)|$  and  $|(t_n)|$  respectively, with  $|(b_n)| = |(t_n)|$ . Let us then define the reversed sequence of  $(u_n)_{n \in \mathbb{N}}$  as  $r((u_n)_{n \in \mathbb{N}}) := (\overline{u_n} = u_{|(u_n)| - n})_{n \in [0, |(u_n)|]}$ . Both  $\bar{b} = r((b_n)_{n \in \mathbb{N}})$

and  $\bar{t} = r((t_n)_{n \in \mathbb{N}})$  are correct representations of a sheet music and performance respectively, as defined in Section I-II.A.

Let  $t^* = t_{|(t_n|)}$  and  $q = |(t_n)|$ ,  $T$  a formal tempo with respect to  $(t_n)$  and  $(b_n)$ , i.e.,  $\forall n \in \mathbb{N}$ ,  $\int_{t_n}^{t_{n+1}} T(t) dt = b_{n+1} - b_n$ ,  $n \in \llbracket 0, q-1 \rrbracket$ , and  $T_r : t \mapsto T(t^* - t)$ .

$$\begin{aligned} \int_{t_n}^{t_{n+1}} T_r(t) dt &= \int_{t^*-t_{q-n-1}}^{t^*-t_{q-n}} T(t^* - t) dt = \int_{t_{q-n}}^{t_{q-n-1}} -T(x) dx \\ &= \int_{t_{q-n-1}}^{t_{q-n}} T(t) dt \\ &= b_{q-n} - b_{q-n-1} \\ &= (b_{q-n} - b_q) - (b_{q-n-1} - b_q) \\ &= -\bar{b}_n + \bar{b}_{n+1} \\ &= \bar{b}_{n+1} - \bar{b}_n \end{aligned}$$

Hence  $T_r$  is a formal tempo with respect to  $(\bar{t}_n)$  and  $(\bar{b}_n)$ .

#### D. Musical explication of the choice of a tempo distance

In terms of tempo, halving and doubling are considered as far as each other from the initial value. Therefore a usual absolute distance does not fit this notion, and we will rather use a logarithmic distance when comparing tempi.

### APPENDIX B : ESTIMATOR MODEL

- d est une distance
- argmin : d est continue (distance), donc bornée et atteint ses bornes sur l'intervalle

Remarque sur la distance log (pas d'importance de  $k_*$ ).

Pourquoi 2/3 et 4/3 :

- candidat unique pour chaque changement de tempo : résistant aux octaves
- easier to search ( $|\frac{2}{3} - 1| = |\frac{4}{3} - 1|$ ) : solution of  $2x - 1 = 1 - x \Leftrightarrow 3x = 2$  : on veut un tempo unique, DONC (à montrer éventuellement) entre  $x$  et  $2x$ , et on prend  $x$  "centré" en 1.
- si l'estimateur est l'identité : par hypothèse d'oracle, le tempo est constant, et joué parfaitement (i.e., fichier midi), et le résultat est le bon
- l'estimateur n'est pas une fonction : l'output à un instant donné peut dépendre des outputs précédents (cas extrême : transcription en temps réel) : très bonne courbe, mais estimateur très complexe, or le problème ici est justement relaxé : ajouter un exemple de fausse transcription et de correcte avec les courbes de tempo correspondantes.

The C++ code for the measure of  $S, \Delta$  is presented hereafter :

```
#measure.hpp :
using couple = std::pair<int, int>;

bool couple_eq(couple c1, couple c2) {
    auto [a, b] = c1;
    auto [c, d] = c2;
    return (a == c) && (b == d);
}

template <typename T>
class Circle {
    std::vector<T> t;
    std::vector<int> count;
    double x;
    int len;
    couple start;

public:
    Circle(std::vector<T> spectre, double x, int
i) : t(std::vector<T>()), count(std::vector<int>()),
x(x), len(spectre.size()), start(couple(i, 0)) {
        count.push_back(1);
        t.push_back(spectre[0]);
        for (unsigned int ind = 0; ind + 1 <
spectre.size(); ind++) {
            if (spectre[ind] == spectre[ind + 1]) {
                count[count.size() - 1]++;
            } else {
                count.push_back(1);
                t.push_back(spectre[ind+1]);
            }
        }
    }

    int size() {
        return t.size();
    }

    bool is_start(couple end) {
        auto [i, k] = start;
        auto [j, l] = end;
        return i == j;
    }

    couple get_start() {
        return start;
    }

    double get(couple a) {
        auto [i, k] = a;
        return t[i] + x*k;
    }

    couple get_next(couple a) {
        auto [i, k] = a;
        if (i + 1 == t.size()) {
            return couple(0, k+1);
        }
        return couple(i+1, k);
    }
}
```

```

couple get_previous(couple a) {
    auto [i, k] = a;
    if (i == 0) {
        return couple(t.size() - 1, k-1);
    }
    return couple(i-1, k);
}

int get_count(couple a) {
    auto [i, _] = a;
    return count[i];
}

int get_cardinal(couple a1, couple a2) {
    auto [i1, k1] = a1;
    auto [i2, k2] = a2;
    int s = 0;
    int delta = i2 - i1 + (k2 - k1) * t.size();
    int q = delta / t.size();

    for (int i = 0; i < t.size(); i++) {
        int c = q;
        if (i1 > i2) {
            c += (i >= i1 || i <= i2);
        } else {
            c += (i >= i1 && i <= i2);
        }
        s += count[i] * c;
    }
    return s;
}

int length() {
    return len;
}
};

```

#measure.cpp :

```

double distance(double a, double b, double x) {
    return fmin(fabs(a - b), x - fabs(a - b));
}

double get_measure(std::vector<double> spectre,
const double delta, const double x, int i) {
    std::sort(spectre.begin(), spectre.end());
    Circle<double> c = Circle<double>(spectre, x, i);
    couple ind_end = c.get_start();
    couple ind_start = ind_end;
    double d = c.get(ind_end) - delta;
    for (int i = 0; i < c.size() - 1; i++) {
        ind_start = c.get_previous(ind_start);
        if (distance(c.get(ind_start), d, x) > delta)
        {
            ind_start = c.get_next(ind_start);
            break;
        }
    }
    int p = c.get_cardinal(ind_start, ind_end);
    int S = p;
}

```

```

int h = 0;
while (h+1 < c.size()) {
    double d1 = delta - distance(c.get(ind_start),
d, x); // Next start index change
    if (c.get(ind_start) > d) {
        d1 = -d1 + 2*delta;
    }

    double d2 =
distance(c.get(c.get_next(ind_end)), d, x) -
delta; // Next end index change

    if (d1 <= d2) {
        p -= c.get_count(ind_start);
    }
    if (d2 <= d1) {
        h++; // A new point enters the delta area
        p += c.get_count(ind_end);
        S = fmax(S, p);
    }
    d += fmin(d1, d2);

    if (d1 <= d2) {
        ind_start = c.get_next(ind_start);
        if (couple_eq(ind_start, ind_end)) {
            ind_end = ind_start;
        }
    }

    if (d2 <= d1) {
        ind_end = c.get_next(ind_end);
    }
}
return (double)S / (double)(c.length());
}

```

## APPENDIX C : QUANTIFIED MODEL

In this section we use the notation introduced by G. Romero-García, C. Guichaoua, and E. Chew [14], that essentially replace  $D$  by  $T$ .

Let us first define  $g : x \mapsto \min(x - \lfloor x \rfloor, 1 + \lfloor x \rfloor - x)$ . One can make sure that  $g : x \mapsto \begin{cases} x - \lfloor x \rfloor & \text{si } x - \lfloor x \rfloor \leq \frac{1}{2} \\ 1 - (x - \lfloor x \rfloor) & \text{sinon} \end{cases}$  and that  $g$  is 1-periodic, continuous on  $\mathbb{R}$ .

Then, by definition :

$$\begin{aligned}
\varepsilon_T(a) &= \max_{t \in T} \left( \min_{m \in \mathbb{Z}} |t - ma| \right) \\
&= \max_{t \in T} \min \left( t - \left\lfloor \frac{t}{a} \right\rfloor a, \left( \left\lfloor \frac{t}{a} \right\rfloor + 1 \right) a - t \right) \\
&= a \max_{t \in T} \min \left( \underbrace{\frac{t}{a} - \left\lfloor \frac{t}{a} \right\rfloor, \left\lfloor \frac{t}{a} \right\rfloor + 1 - \frac{t}{a}}_{g\left(\frac{t}{a}\right)} \right) \\
&= a \max_{t \in T} g\left(\frac{t}{a}\right)
\end{aligned}$$

hence we have proved  $\varepsilon_T$  to be continuous on  $R_+^*$ .

Furthermore, for  $n \in \mathbb{N}^*$ ,  $T \subset (\mathbb{R}_+^*)^n$ ,  $a \in R_+^*$ ,  
 $\varepsilon_T(a) = a \max_{t \in T} g\left(\frac{t}{a}\right) = a \times 1 \max_{t \in T} g\left(\frac{t}{a} 1^{-1}\right) = a \varepsilon_{T/a}(1)$ .  
Hence the intuitive following result : the smaller the tatum,  
the smaller the bound of the error.

*A. Characterization of local maxima that happen not be a local minima*

1) *First implication:*

Let  $a$  be a local maxima (and not a local minima) of  $\varepsilon_T$ ,  $a > 0$ . By definition, there is a  $a > \varepsilon > 0$  so that :

$$\forall \delta \in ]-\varepsilon, \varepsilon[, \varepsilon_T(a) \geq \varepsilon_T(a + \delta).$$

Let  $t \in \arg\max_{t' \in T} g\left(\frac{t'}{a}\right)$ , hence  $\varepsilon_T(a) = ag\left(\frac{t}{a}\right)$ .

For all  $\delta \in ]-\varepsilon, \varepsilon[, \varepsilon_T(a) = ag\left(\frac{t}{a}\right) \geq \varepsilon_T(a + \delta)$ ,  
and  $\varepsilon_T(a + \delta) = (a + \delta) \max_{t' \in T} g\left(\frac{t'}{a + \delta}\right) \geq (a + \delta)g\left(\frac{t}{a + \delta}\right)$ .  
For  $\delta \geq 0$ ,  $a + \delta \geq a$ , so  $ag\left(\frac{t}{a}\right) \geq (a + \delta)g\left(\frac{t}{a + \delta}\right) \geq ag\left(\frac{t}{a + \delta}\right)$ .  
Hence,  $g\left(\frac{t}{a}\right) \geq g\left(\frac{t}{a + \delta}\right)$  since  $a > 0$ , for all  $\delta \in [0, \varepsilon]$ .

Therefore,  $g$  increases monotonically in the range  $]\frac{t}{a + \delta}, \frac{t}{a}[,$   
since  $g$  has a unique local maxima (modulo 1), considering  
the previous range as a neighbourhood of  $\frac{t}{a}$ .

Hence,  $g = x \mapsto x - \lfloor x \rfloor$  within the considered range, and  
 $g\left(\frac{t}{a}\right) = \frac{t}{a} - \lfloor \frac{t}{a} \rfloor, \varepsilon_T(a) = t - a \lfloor \frac{t}{a} \rfloor$ .

The function  $\varepsilon_T$  is the maximum of a finite set of continuous  
functions, with a countable set of  $A$  of intersection, i.e.,  $A =$   
 $\{x \in \mathbb{R}_+^* : \exists (t_1, t_2) \in T^2 : t_1 \neq t_2 \wedge xg\left(\frac{t_1}{x}\right) = xg\left(\frac{t_2}{x}\right)\}$ . In-

deed,  
 $x \in A \Leftrightarrow \exists (t_1, t_2) \in T^2 : t_1 \neq t_2 \wedge g\left(\frac{t_1}{x}\right) = g\left(\frac{t_2}{x}\right)$

$$\Leftrightarrow \exists (t_1, t_2) \in T^2 : t_1 \neq t_2 \wedge \frac{t_1}{x} = \pm \frac{t_2}{x} \text{ mod } 1$$

$$\Leftrightarrow \exists (t_1, t_2) \in T^2 : t_1 \neq t_2 \wedge x = \frac{t_1 \mp t_2}{n}, n \in \mathbb{Z}^*$$

Hence  $A \subset \left\{ \frac{t_1 \mp t_2}{n}, (t_1, t_2) \in T^2, n \in \mathbb{Z}^* \right\}$ , because  $T \subset$   
 $(R_+^*)^{|T|}$ . Therefore, there is a countable set of closed convex  
intervals, which union is  $R_+^*$  so that on each of these inter-  
vals,  $\varepsilon_T$  is equal to  $f_t : a \mapsto ag\left(\frac{t}{a}\right)$  for a  $t \in T$ . Let then  $t$  be  
so that for all  $x \in ]a - \delta', a[, \varepsilon_T(x) = f_{t(x)}$ , where  $]a - \delta', a[$   
is included in one the previous intervals. Since  $f_t$  and  $\varepsilon_T$   
are both continuous on  $[a - \delta', a]$ ,  $f_t(a) = \varepsilon_T(a)$  and there-  
fore,  $t \in \arg\max_{t' \in T} g\left(\frac{t'}{a}\right)$ . The previous paragraph showed that  
 $g$  is increasing on a left neighbourhood of  $\frac{t}{a}$ . Therefore, on a  
right neighbourhood of  $\frac{t}{a}$ ,  $g$  is either increasing or decreasing  
by its definition.

- if  $g$  is increasing on this neighbourhood, called  $N\left(\frac{t}{a}\right)^+$   
in the following, the previous expression of  $g$  re-  
mains valid, i.e.,  $\forall x \in N\left(\frac{t}{a}\right)^+, g(x) = x - \lfloor x \rfloor$ . More-  
over,  $x \mapsto \lfloor x \rfloor$  is right-continuous, hence by restrict-  
ing  $N\left(\frac{t}{a}\right)^+$ , we can assure for all  $x \in N\left(\frac{t}{a}\right)^+, \lfloor x \rfloor =$   
 $\lfloor \frac{t}{a} \rfloor$ . Let then  $y = a - \frac{t}{x}$  so that  $x = \frac{t}{a - y}$ , we then  
have  $\varepsilon_T(a - y) = t - (a - y) \lfloor \frac{t}{a} \rfloor \leq \varepsilon_T(a) = t - a \lfloor \frac{t}{a} \rfloor$   
because  $a$  is a local maxima of  $\varepsilon_T$  and  $a - y$  is within a

(left) neighbourhood of  $a$ , even if it means restricting  
 $\delta'$  or  $N\left(\frac{t}{a}\right)^+$ . Hence,  $t - \lfloor \frac{t}{a} \rfloor a \geq t - (a - y) \lfloor \frac{t}{a} \rfloor$  i.e.,  
 $a \lfloor \frac{t}{a} \rfloor \leq (a - y) \lfloor \frac{t}{a} \rfloor \Leftrightarrow 0 \leq -y \lfloor \frac{t}{a} \rfloor$  i.e.,  $\lfloor \frac{t}{a} \rfloor \leq 0$  i.e.,  
 $\lfloor \frac{t}{a} \rfloor = 0$ , since  $y, t$  and  $a$  are all positive values. Then,  
 $a > t$  and therefore  $\varepsilon_T(a - y) = t = \varepsilon_T(a)$  for  $\lfloor \frac{t}{a} \rfloor =$   
 $0$ . This interval where  $\varepsilon_T$  is constant is then either  
going on infinitely on the right of  $a$ , or else  $\varepsilon_T$  will  
reach a value greater than  $\varepsilon_T(a) = t$ , since  $\varepsilon_T$  can then  
be rewritten as  $x \mapsto \max\left(\max_{t' \in T \setminus \{t\}} xg\left(\frac{t'}{x}\right), \varepsilon_T(a)\right)$  on  
 $[a, +\infty[$ . Hence  $a$  is a local minima on the right, and  
since  $\varepsilon_T$  is constant on a left neighbourhood of  $a$ ,  $a$  is  
also a local minima on the left. Finally,  $a$  is a local min-  
ima, which is absurd by definition.

- else,  $g$  is decreasing on  $N\left(\frac{t}{a}\right)^+$ ,  $\frac{t}{a}$  is by definition a lo-  
cal maxima of  $g$ . However,  $g$  only has a unique local  
maxima modulo 1, that is  $\frac{1}{2}$ . Hence,  $\frac{t}{a} = \frac{1}{2} \text{ mod } 1$ , i.e.,  
 $\frac{t}{a} = \frac{1}{2} + k, k \in \mathbb{Z}$ , or  $a = \frac{t}{k + \frac{1}{2}}, k \in \mathbb{N}$ , since  $a > 0$ .

2) *Second implication:*

Let  $(t, k) \in T \times \mathbb{N}$ ,  $a = \frac{t}{k + \frac{1}{2}}$ .

By definition :  $g\left(\frac{t}{a}\right) = g\left(\frac{t}{\frac{t}{k + \frac{1}{2}}}\right) = g\left(\frac{1}{2}\right) = \frac{1}{2} = \max_{\mathbb{R}} g$ .

Therefore,  $\varepsilon_T(a) = a \max_{t' \in T} g\left(\frac{t'}{a}\right) = ag\left(\frac{t}{a}\right) = \frac{a}{2}$ .

For all  $x \in ]0, a[, \varepsilon_T(x) = x \max_{t' \in T} g\left(\frac{t'}{x}\right) \leq \frac{x}{2} < \frac{a}{2} =$   
 $\varepsilon_T(a)$ , i.e.,  $\varepsilon_T(a) > \varepsilon_T(x)$ .

Let  $T^* = \{t' \in T : g\left(\frac{t'}{a}\right) = \frac{1}{2}\}$ . Since  $t \in T^*, |T^*| > 1$ .

Let  $t^* \in T^*$ . For all  $t' \in T \setminus T^*, g\left(\frac{t'}{a}\right) < g\left(\frac{t^*}{a}\right)$ .

Since  $h_{t'} : x \mapsto g\left(\frac{t'}{x}\right) - g\left(\frac{t^*}{x}\right)$  is continuous in a neighbour-  
hood of  $a > 0$ , we have the existence of  $\varepsilon_{t'} > 0$  so that  $h_{t'}$  is  
strictly positive within  $[a, a + \varepsilon_{t'}]$ .

Let  $\varepsilon_{t^*} = \min_{t' \in T} \varepsilon_{t'}$  and finally  $\varepsilon_1 = \min_{t^* \in T^*} \varepsilon_{t^*}$ .

Let  $(t_1, t_2) \in (T^*)^2$ .

In the following,  $N(a)^+$  is a right neighbourhood of  $a$  such  
that  $a \notin N(a)^+$ .

Let  $\text{tmp} : x \mapsto g\left(\frac{t_1}{x}\right) - g\left(\frac{t_2}{x}\right)$  be a continous function on  
 $N(a)^+$  and  $A$  be the set of all  $x^* \in N(a)^+$  so that  $\text{tmp}(x^*) =$   
 $0 \Leftrightarrow g\left(\frac{t_1}{x^*}\right) = g\left(\frac{t_2}{x^*}\right)$ .

We have for all  $x^* \in A, g\left(\frac{t_1}{x^*}\right) = g\left(\frac{t_2}{x^*}\right)$  by definition. Con-  
sidering the expression of  $g$ , we then find :  $\frac{t_1}{x^*} = \pm \frac{t_2}{x^*} \text{ mod } 1$ .  
Moreover, since  $g$  only reach  $g\left(\frac{t_1}{a}\right) = \frac{1}{2}$  once per period, we  
have  $\frac{t_1}{a} = \frac{t_2}{a} \text{ mod } 1$ , i.e.,  $|\frac{t_1}{a} - \frac{t_2}{a}| = k_a \in \mathbb{N}$ .

Then,  $\frac{t_1}{x^*} = \pm \frac{t_2}{x^*} \text{ mod } 1$  i.e.,  $|\frac{t_1}{x^*} \mp \frac{t_2}{x^*}| = k_* \in \mathbb{N}$ , and there-  
fore  $|t_1 \mp t_2| = ak_* = x^* k_*$ , and  $x^* > a$  implies  $k_a > k_* \geq$   
 $0$ . However,  $x^* = \frac{|t_1 \mp t_2|}{k_*}$ , hence  $A$  is finite if  $A \neq \emptyset$ , and  $\emptyset$  is  
a finite set. Finally,  $A$  is a finite set, i.e.,  $|A| \in \mathbb{N}$ .

Let then  $x_{t_1, t_2} = \begin{cases} \min A \text{ if } A \neq \emptyset \\ x \in N(a)^+ \setminus \{a\} \text{ otherwise} \end{cases}$  WLOG  $g\left(\frac{t_1}{x}\right) \geq$

$g\left(\frac{t_2}{x}\right) \forall x \in [a, x_{t_1, t_2}]$

Let  $a_2 = \min_{(t_1, t_2) \in T^*{}^2} \underbrace{x_{t_1, t_2}}_{> a}$  and  $a_1 \in ]a, a_2[$ ,



let  $t^* = \operatorname{argmax}_{t' \in T^*} g\left(\frac{t'}{a_1}\right)$  We finally have  $\forall x \in ]a, a_2[, g\left(\frac{t^*}{x}\right) \geq g\left(\frac{t'}{x}\right), \forall t' \in T^*$ .

Let then  $\tilde{a} = \min(a + \varepsilon_1, a_2)$  so that for all  $x \in ]a, \tilde{a}[$ ,  $t' \in T$ ,  $g\left(\frac{t'}{x}\right) \geq g\left(\frac{t'}{\tilde{a}}\right)$ , hence  $\varepsilon_T(x) = xg\left(\frac{t^*}{x}\right)$ .  
Let  $f : x \mapsto g\left(\frac{t^*}{x}\right)$ ,  $f(a) = g\left(\frac{t^*}{a}\right) = \frac{1}{2}$  because  $t^* \in T^*$  hence  $f$  is increasing on a right neighbourhood of  $a$ ,  $N(a)^+$ , since  $\frac{1}{2}$  is a global maxima of  $g$ , therefore  $g$  is increasing on  $N\left(\frac{t^*}{a}\right)$  a left neighbourhood of  $\frac{t^*}{a}$ , i.e.,  $f$  is increasing on  $N(a)^+$ . Therefore, we know that  $g\left(\frac{t^*}{x}\right) = \frac{t^*}{x} - \left\lfloor \frac{t^*}{x} \right\rfloor$ , since the only other possible expression for  $g$  would imply a decreasing function on  $N\left(\frac{t^*}{a}\right)$ .

Hence,  $\varepsilon_T(x) = xg\left(\frac{t^*}{x}\right) = x\left(\frac{t^*}{x} - \left\lfloor \frac{t^*}{x} \right\rfloor\right) = t^* - x\left\lfloor \frac{t^*}{x} \right\rfloor$  and  $\varepsilon_T(a) = t^* - a\left\lfloor \frac{t^*}{a} \right\rfloor$  since  $\varepsilon_T$  is continuous on  $\mathbb{R}_+^*$ .

By definition :  $\left\lfloor \frac{t^*}{a} \right\rfloor \leq \frac{t^*}{a} < \left\lfloor \frac{t^*}{a} \right\rfloor + 1$ .

However,  $f(a) = \frac{1}{2} = \frac{t^*}{a} - \left\lfloor \frac{t^*}{a} \right\rfloor$ , therefore  $\left\lfloor \frac{t^*}{a} \right\rfloor < \frac{t^*}{a}$ .

Then, there is  $\alpha \in \mathbb{R}_+^*$  so that  $\left\lfloor \frac{t^*}{a} \right\rfloor < \alpha < \frac{t^*}{a}$ , let  $y = \frac{t^*}{\alpha}$ , i.e.,  $\alpha = \frac{t^*}{y}$ , with  $y > a$ .

Let  $a' = \min(y, \tilde{a})$  and  $k = \left\lfloor \frac{t^*}{a} \right\rfloor$ . For all  $x \in ]a, a'[,$

- $x \leq y = \frac{t^*}{\alpha}$  hence  $\left\lfloor \frac{t^*}{x} \right\rfloor \leq \alpha \leq \frac{t^*}{x}$
- $x \geq a$  hence  $\frac{t^*}{x} \leq \frac{t^*}{a} < \left\lfloor \frac{t^*}{a} \right\rfloor + 1$

In the end,  $\left\lfloor \frac{t^*}{x} \right\rfloor = \left\lfloor \frac{t^*}{a} \right\rfloor = k$  by definition.

Then,  $\varepsilon_T(x) = t^* - xk$  and  $\varepsilon_T(a) = t^* - xa$ , with  $a < x$ .

Finally,  $\varepsilon_T(a) > \varepsilon_T(x)$ .

To conclude, for all  $x \in ]0, a'[,$

- if  $x \leq a$ ,  $\varepsilon_T(a) \geq \varepsilon_T(x)$
- if  $x \geq a$ ,  $x \in [a, a'[,$  and  $\varepsilon_T(a) \geq \varepsilon_T(x)$

Hence  $a$  is a local maxima of  $\varepsilon_T$  and not a local minima, and then the set of all such local maxima of  $\varepsilon_T$  is  $M_T = \left\{ \frac{t}{k+\frac{1}{2}}, t \in T, k \in \mathbb{N} \right\}$  and therefore with the notations introduced in Section IV.D, we proved (11) ■

## B. Necessary condition to be a local minima

Let  $a$  be a local minima of  $\varepsilon_T$ , i.e.,

$\varepsilon_T$  is constant on a neighbourhood of  $+\infty$ , donc on va d'abord trouver un maximum local qui n'est pas un minimum local, puis le reste de la preuve fonctionne

Par continuité de  $\varepsilon_T$ , on est assuré de l'existence d'exactement un unique minimum local entre deux maximums locaux, qui est alors global sur cet intervalle.

Par la condition nécessaire précédente, il suffit donc, pour déterminer ce minimum local, de déterminer le plus petit élément parmi les points obtenus, contenus dans l'intervalle. On en déduit ainsi un algorithme en  $\mathcal{O}\left(|T|^2 \frac{t^*}{\tau} \log\left(|T| \frac{t^*}{\tau}\right)\right)$  permettant de déterminer tous les minimums locaux accordés par le seuil  $\tau$  fixé, sur l'intervalle  $]2\tau, t_* + \tau[$

On en déduit la correction de Algorithm 1.

## APPENDIX D : MUSICAL GLOSSARY

### A. Acronyms

*MIR* – Music Information Retrieval. 1

### B. Definitions

**articulation:** describes how a specific note is played by the performer. For instance, *staccato* means the note shall not be maintained, and instead last only a few musical units, depending on the context. On the other hand, a fermata (*point d'orgue* in French) indicates that the note should stay longer than indicated, to the performer's discretion. 1

**beat:**

Unité de temps d'une partition, le beat est défini par une signature temps, ou division temporelle. Bien que sa valeur ne soit *a priori* pas fixe d'une partition à une autre, ni même sur une même partition, la notion de beat est en général l'unité la plus pratique quant à la description d'un passage rythmique, lorsque la signature temps est adéquatement définie. 1, 13

**chord:** A chord is by definition the simultaneous production of at least three musical events with different pitches 2

**measure:** Une mesure est une unité de temps musicale, contenant un certain nombre (entier) de beat. Ce nombre est indiqué par la *time signature* 2

**online:** Définition d'une méthode / d'un algo online 5

**rest:** A symbolic notation for silence, following the same rules as actual note notations. 2

**tatum:** Résolution minimal d'une unité musicale, exprimé en beat. Bien que de nombreuses valeurs soit possible, la définition formelle d'un tatum serait la suivante :  $\sup\{r \mid \forall n \in \mathbb{N}, \exists k \in \mathbb{N} : b_n = kr, r \in \mathbb{R}_+^*\}$ . Pour des raisons pratiques, il arrive que le tatum soit un élément plus petit que la définition donnée, en particulier si cet élément est plus facilement expressible dans une partition, ou a plus de sens d'un point de vue musical. On notera dans la définition de l'ensemble donnée,  $k$  n'a pas d'unité, ce qui montre clairement que le tatum s'exprime en beat comme dit précédemment. 4, 6

**tempo:**

Formally defined in Section III by :  $T_n^* = \frac{b_{n+1} - b_n}{t_{n+1} - t_n}$ , tempo is a measure of the immediate speed of a performance, usually written on the score. It can be seen as a ratio between the symbolic speed indicated by the score, and the actual speed of a performance. Tempo is often expressed in *beat per minute*, or *bpm* 1

**time signature.** 13

**triplet:** A triplet is a musical symbol indicating to play a third of the indicated duration 9

*velocity*: The velocity describes how loud a sound shall be played, or is actually played. <sup>1</sup>