CS540
Uninformed Search

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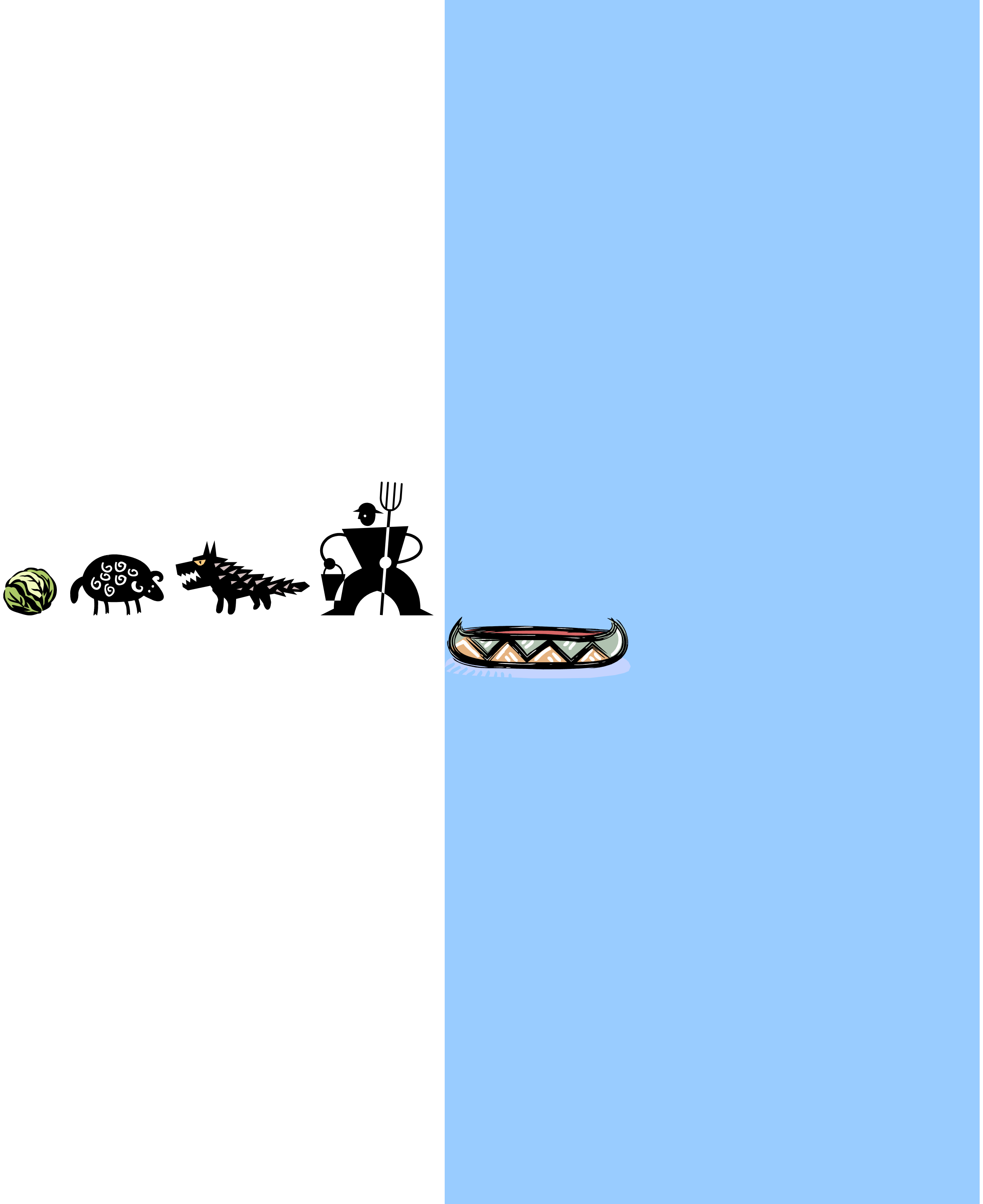
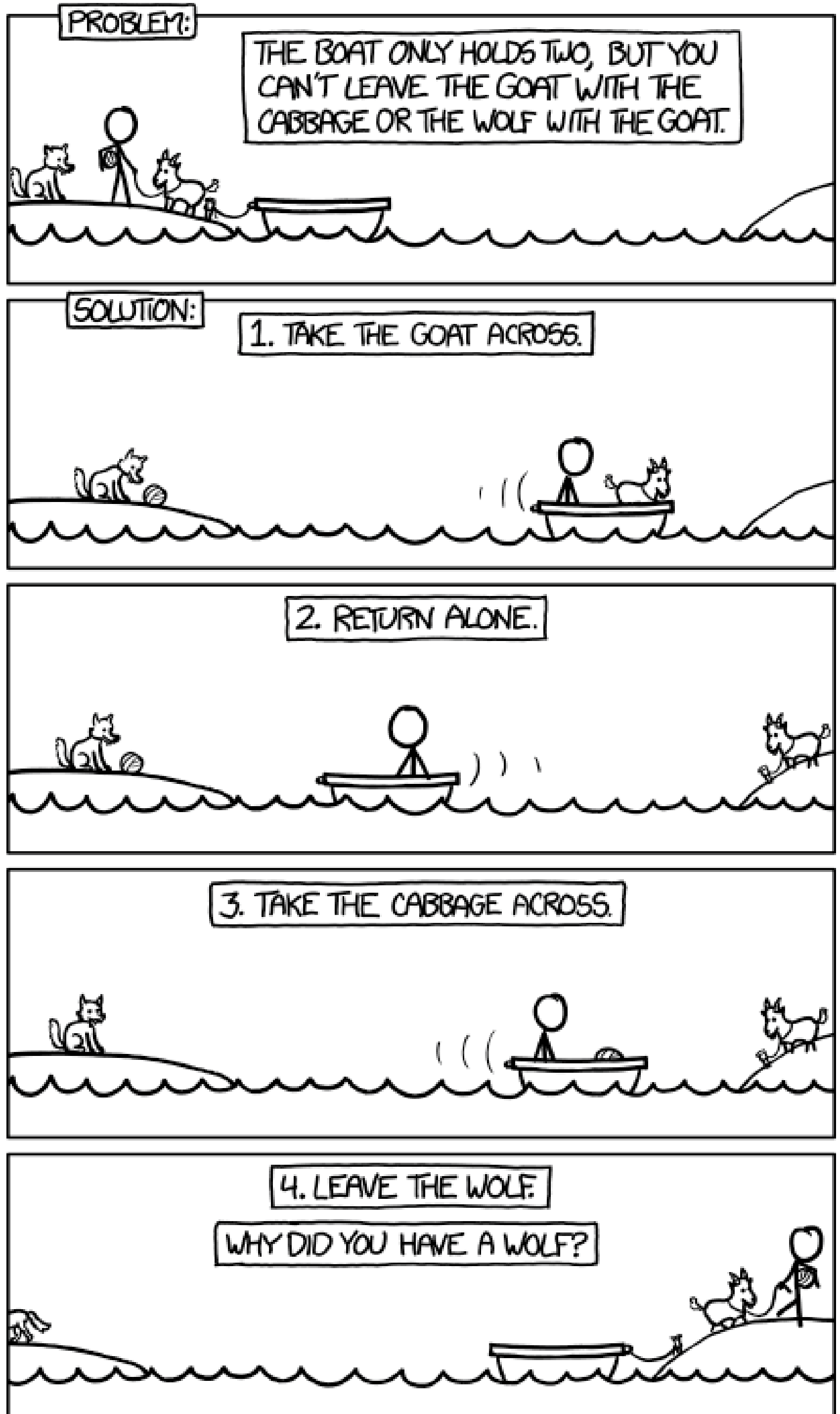
Main messages

- Many AI problems can be formulated as search.
- Iterative deepening is good when you don’t know much.





http://xkcd.com/1134/

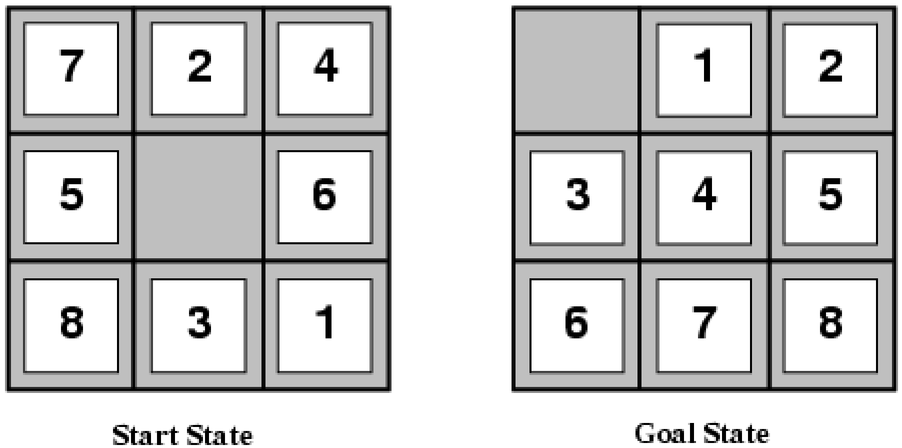
The search problem

- State space S : all valid configurations
- Initial states (nodes) I={(CSDF,)} $\subseteq S$
  - Where's the boat?
- Goal states G={(,CSDF)} $\subseteq S$

- Successor function $succs$(s)$\subseteq S$: states reachable in one step (one arc) from s
  - $succs$((CSDF,)) = {(CD, SF)}
  - $succs$((CDF,S)) = {(CD,FS), (D,CFS), (C, DFS)}
- Cost(s,s')=1 for all arcs. (weighted later)
- The search problem: find a solution path from a state in I to a state in G.
  - Optionally minimize the cost of the solution.

Search examples

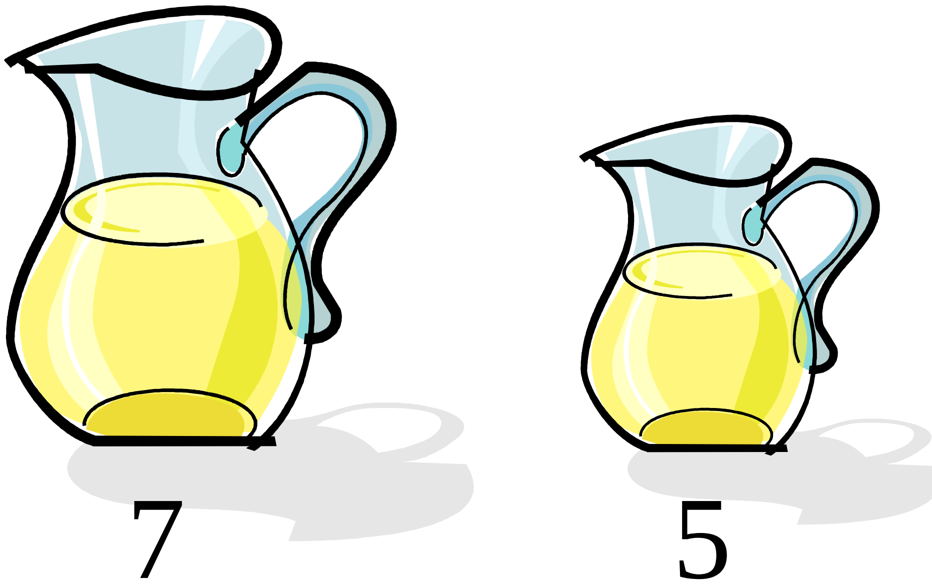
- 8-puzzle



- States = configurations
- successor function = up to 4 kinds of movement
- Cost = 1 for each move

Search examples

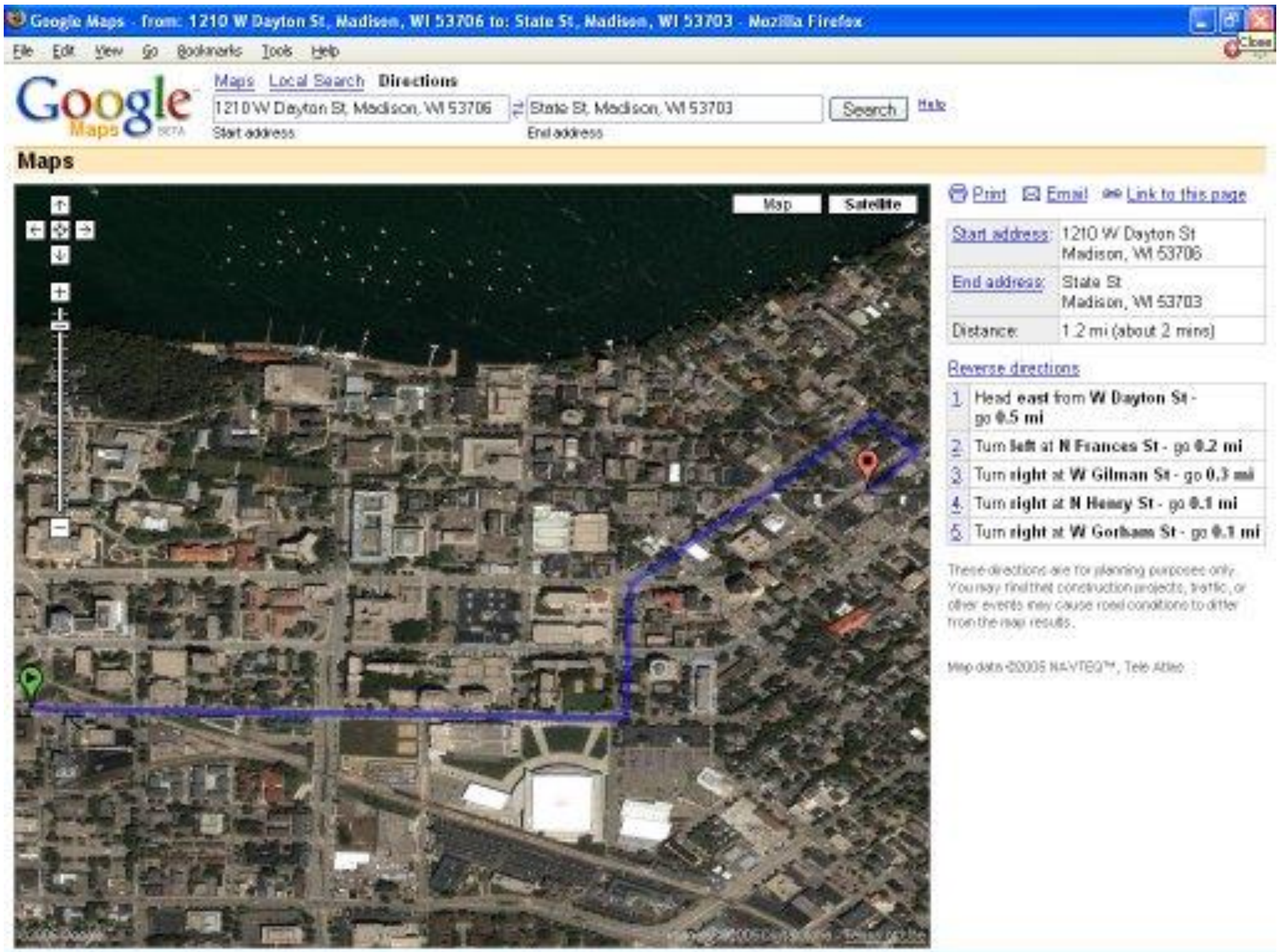
- Water jugs: how to get 1?



- Goal? (How many goal states?)
- Successor function: fill up (from tap or other jug), empty (to ground or other jug)

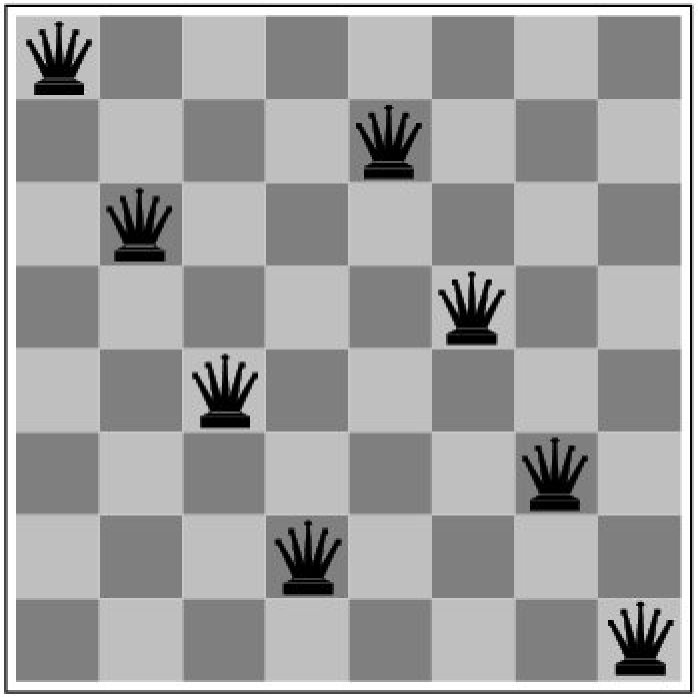
Search examples

- Route finding (state? Successors? Cost weighted)



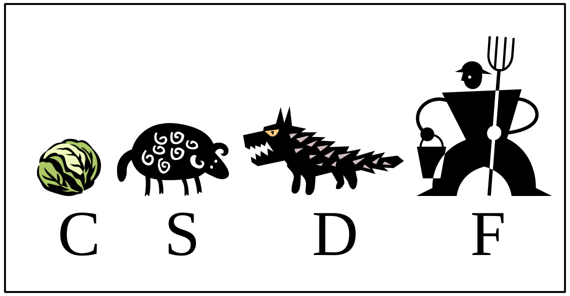
8-queens

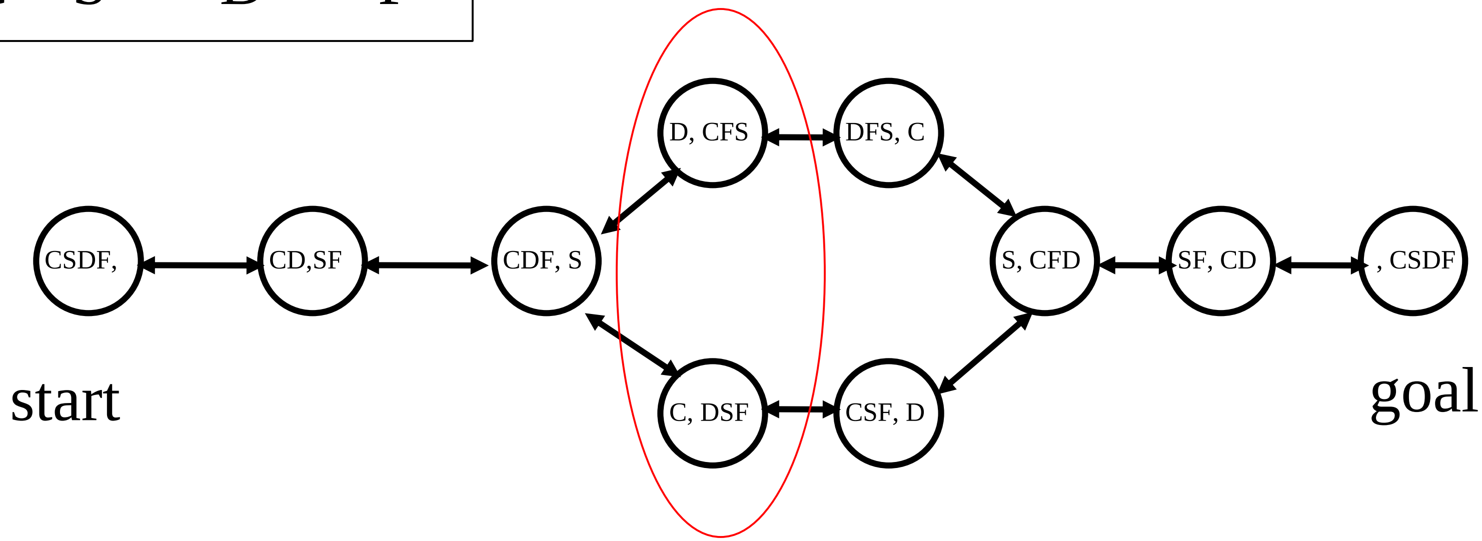
- State: complete configuration vs. column-by-column



- Tree instead of graph

A directed graph in state space

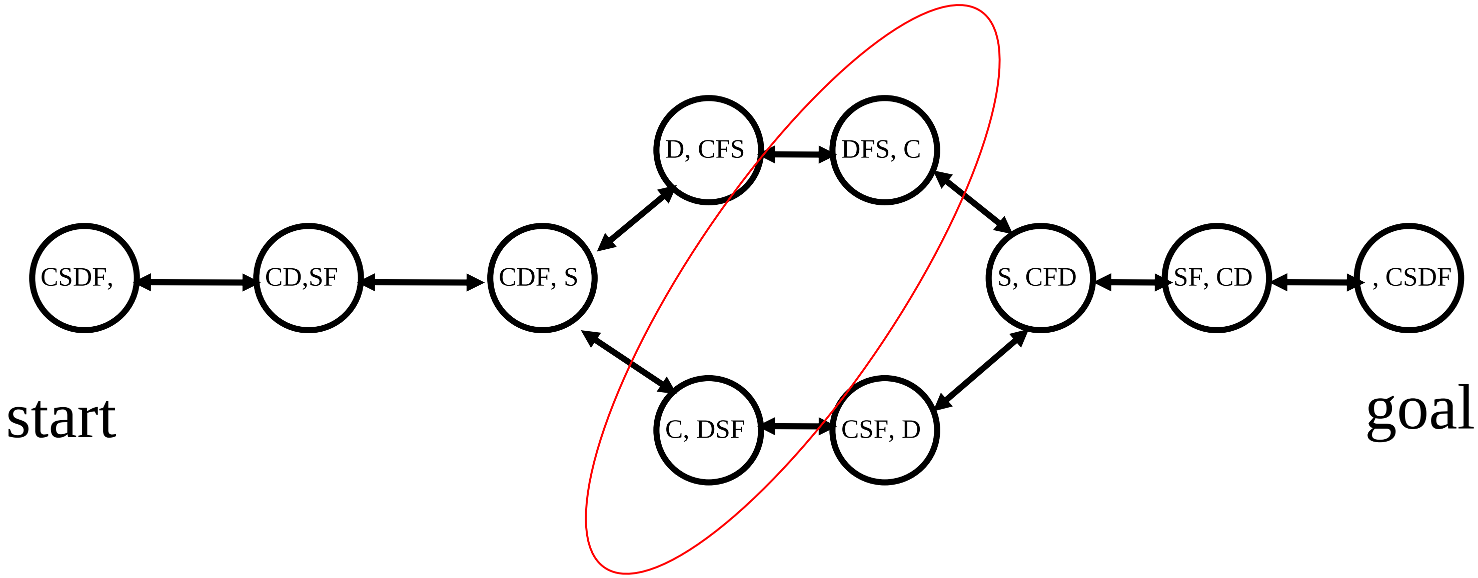




- In general there will be many generated, but un-expanded states at any given time
- One has to choose which one to expand next

Different search strategies

- The generated, but not yet expanded states form the fringe (OPEN).
- The essential difference is which one to expand first.
- Deep or shallow?



Uninformed search on trees

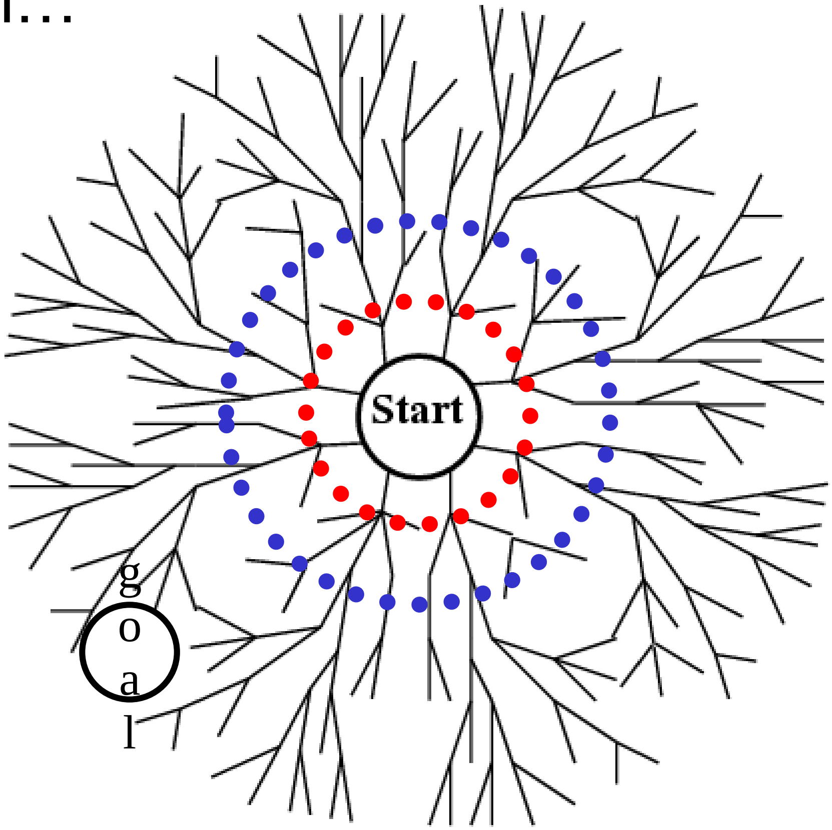
- Uninformed means we only know:
  - The goal test
  - The ***succs***() function
- But not which non-goal states are better: that would be informed search (next lecture).
- For now, we also assume ***succs***() graph is a tree.
  - Won't encounter repeated states.
  - We will relax it later.
- Search strategies: BFS, UCS, DFS, IDS, BIBFS
- Differ by what un-expanded nodes to expand

Breadth-first search (BFS)

Expand the shallowest node first

- Examine states **one** step away from the initial states
- Examine states **two** steps away from the initial states
- and so on…

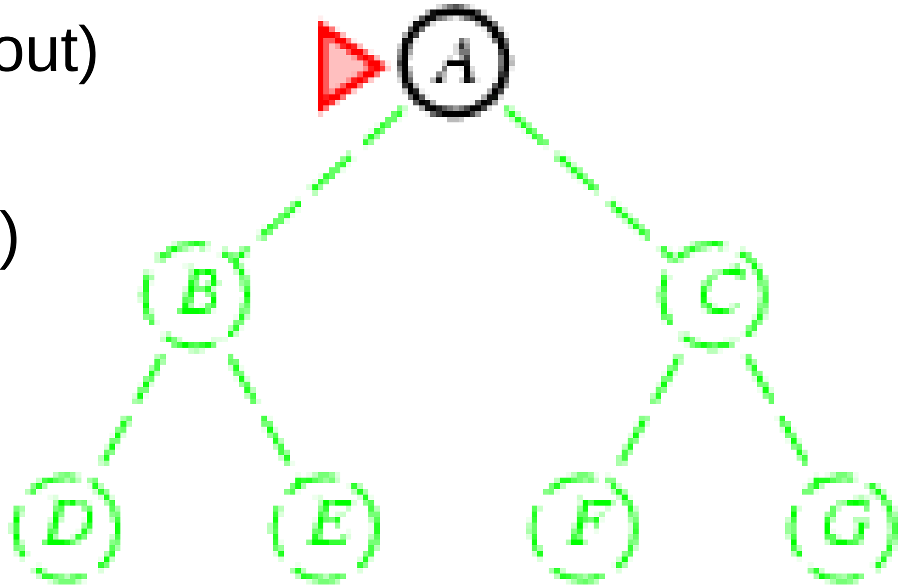
ripple



Breadth-first search (BFS)

Use a queue (First-in First-out)

1. en_queue(Initial states)
2. While (queue not empty)
3.   s = de_queue()
4.   if (s==goal) success!
5.   T = succs(s)
6.   en_queue(T)
7. endWhile



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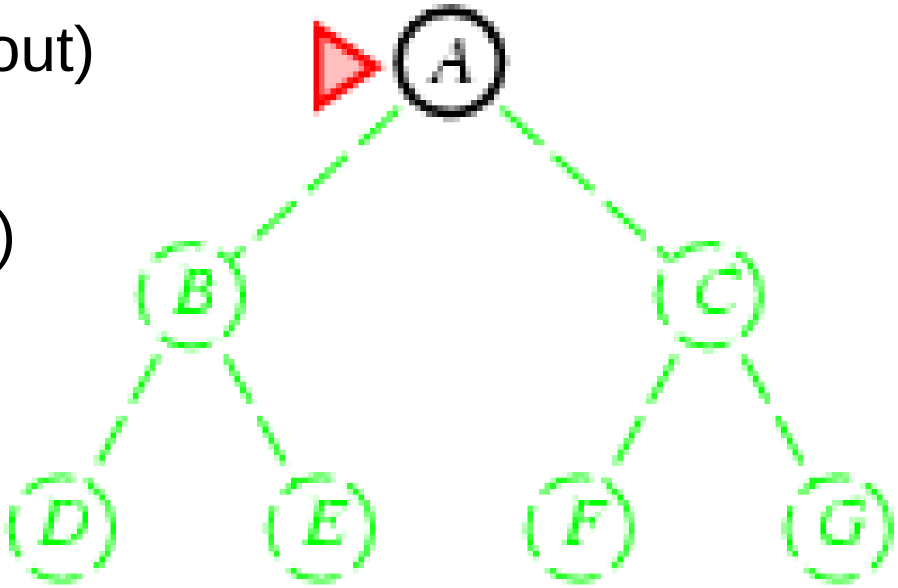

queue (fringe, OPEN)

→ [A] →

Breadth-first search (BFS)

Use a queue (First-in First-out)

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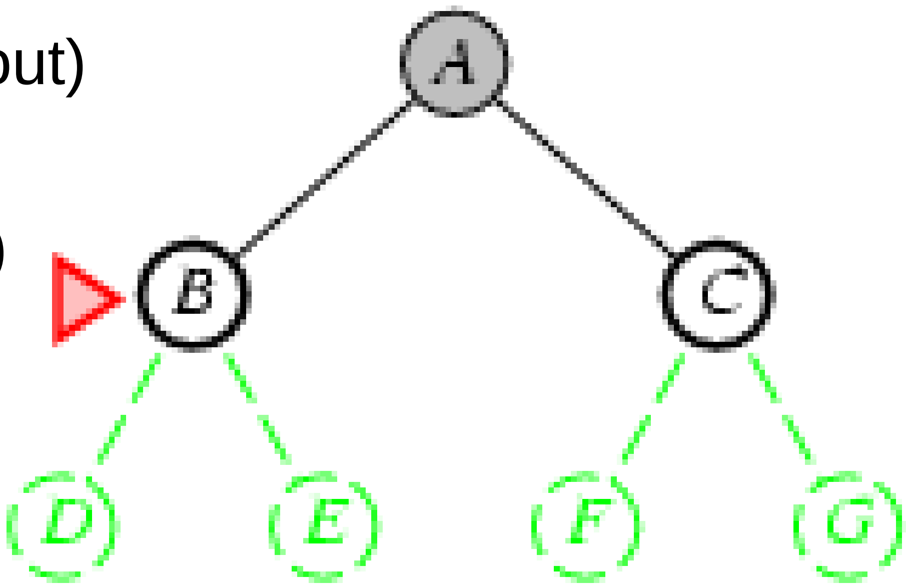

queue (fringe, OPEN)
→ [CB] → A

Breadth-first search (BFS)

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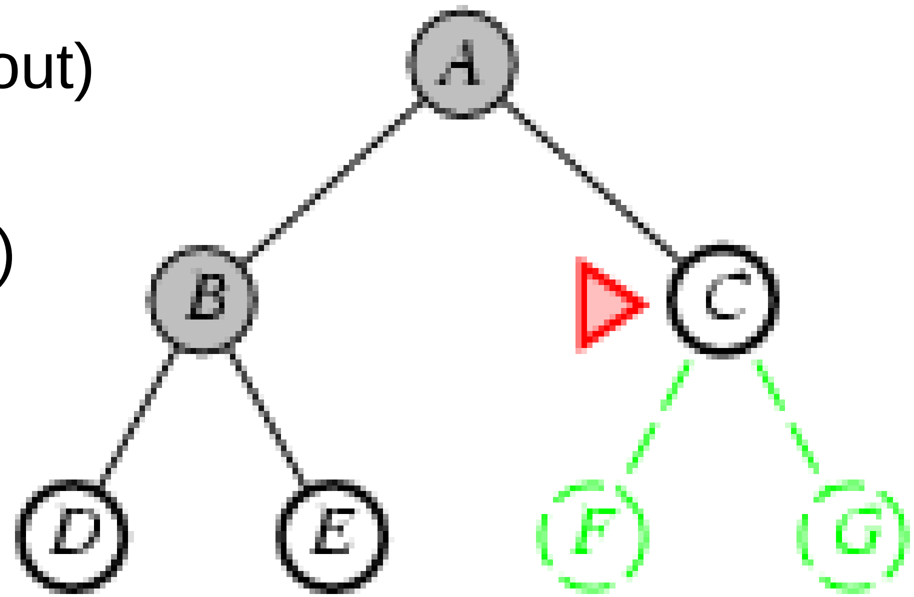

queue (fringe, OPEN)
→ [EDC] → B

Breadth-first search (BFS)

Use a queue (First-in First-out)

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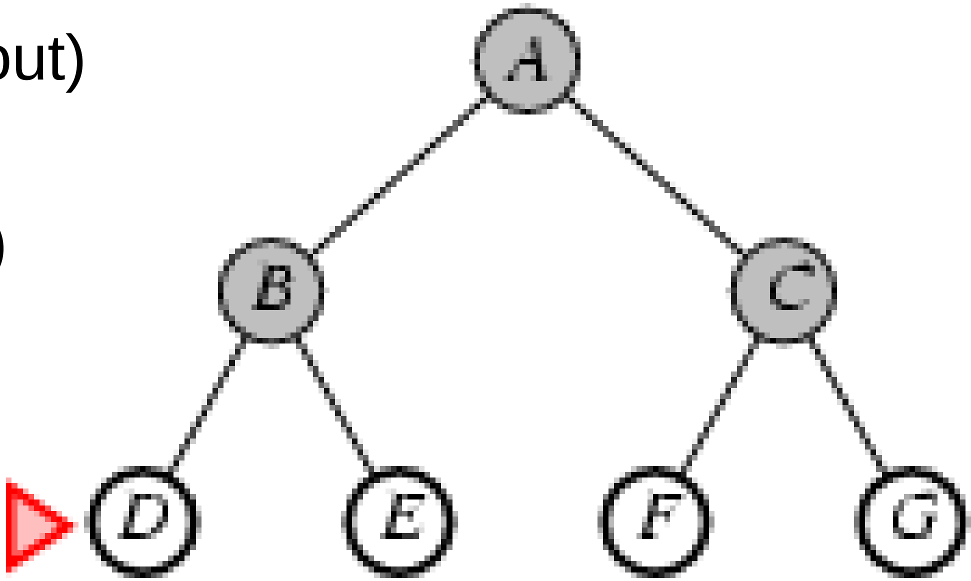


queue (fringe , OPEN)

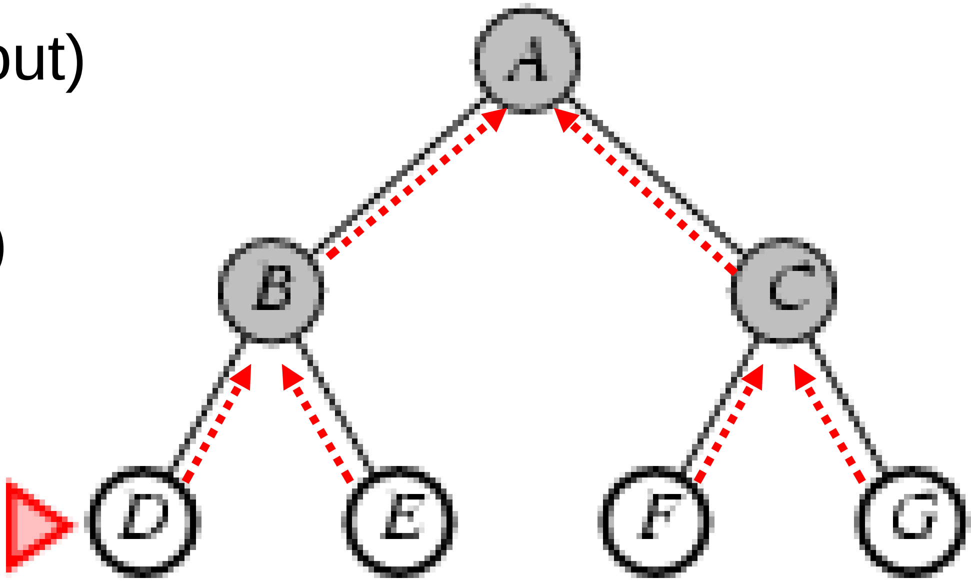
→[GFED] → C

If G is a goal, we've seen it, but we don't stop!

Breadth-first search (BFS)

Use a queue (First-in First-out)

- en_queue(Initial states)
- While (queue not empty)
-     s = de_queue()
-     if (s==goal) success!
-     T = succs(s)
-     for t in T: t.prev=s
-     en_queue(T)
- endWhile



queue
→[] →G

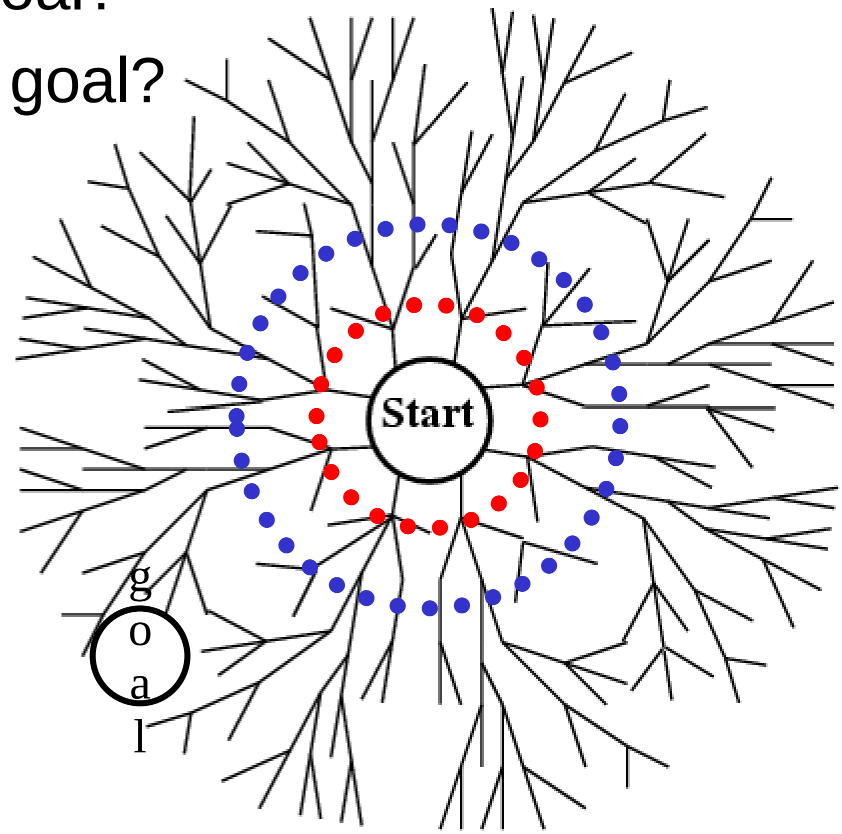
... until much later we pop G.

We need back pointers to recover the solution path.

Looking stupid? Indeed. But let's be consistent…

Performance of BFS

- Assume:
  - the graph may be infinite.
  - Goal(s) exists and is only finite steps away.
- Will BFS find at least one goal?
- Will BFS find the least cost goal?
- Time complexity?
  - # states generated
  - Goal ***d*** edges away
  - Branching factor ***b***
- Space complexity?
  - # states stored



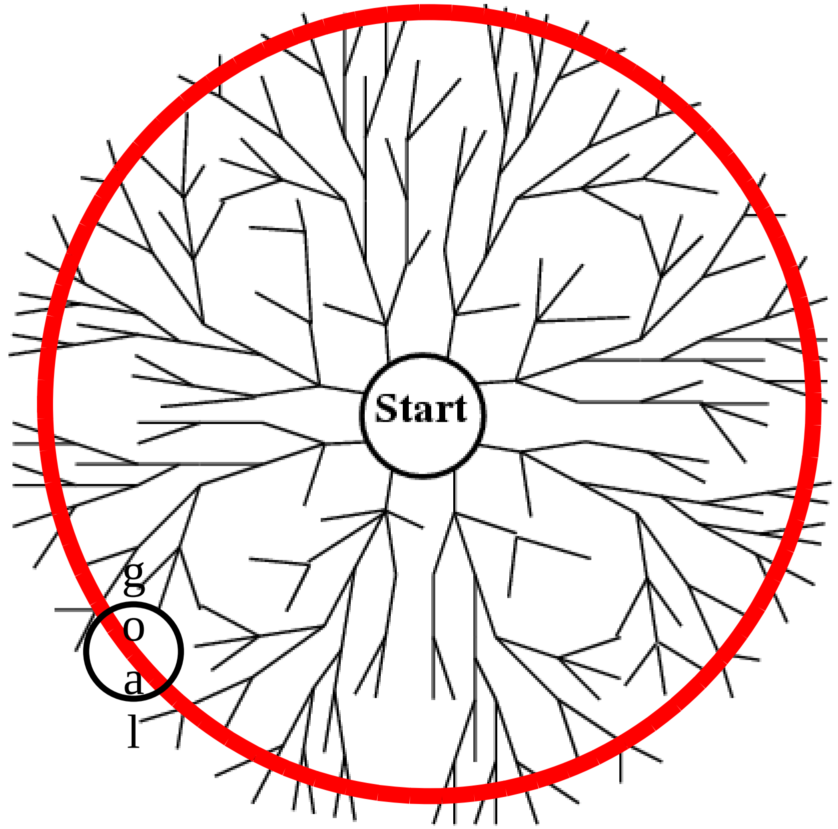
Performance of BFS

Four measures of search algorithms:

- Completeness (not finding all goals): yes, BFS will find a goal.
- Optimality: yes if edges cost 1 (more generally positive non-decreasing in depth), no otherwise.
- Time complexity (worst case): goal is the last node at radius d.
  - Have to generate all nodes at radius d.
  - $b + b^2 + \ldots + b^d \sim O(b^d)$
- Space complexity (bad)
  - Back pointers for all generated nodes $O(b^d)$
  - The queue / fringe (smaller, but still $O(b^d)$)

What's in the fringe (queue) for BFS?

- Convince yourself this is $O(b^d)$



Performance of search algorithms on trees

b: branching factor (assume finite)   d: goal depth

	Complete	optimal	time	space
Breadth-first search	Y	Y, if [1]	$O(b^d)$	$O(b^d)$

1. Edge cost constant, or positive non-decreasing in depth

Performance of BFS

Solution: Uniform-cost search

Four measures of search algorithms:

- Completeness (not finding all goals): find a goal.
- Optimality: yes if edges cost 1 (more generally positive non-decreasing with depth), no otherwise.
- Time complexity (worst case): goal is the last node at radius d.
  - Have to generate all nodes at radius d.
  - $b + b^2 + \ldots + b^d \sim O(b^d)$
- Space complexity (bad, Figure 3.11)
  - Back points for all generated nodes $O(b^d)$
  - The queue (smaller, but still $O(b^d)$)

Uniform-cost search

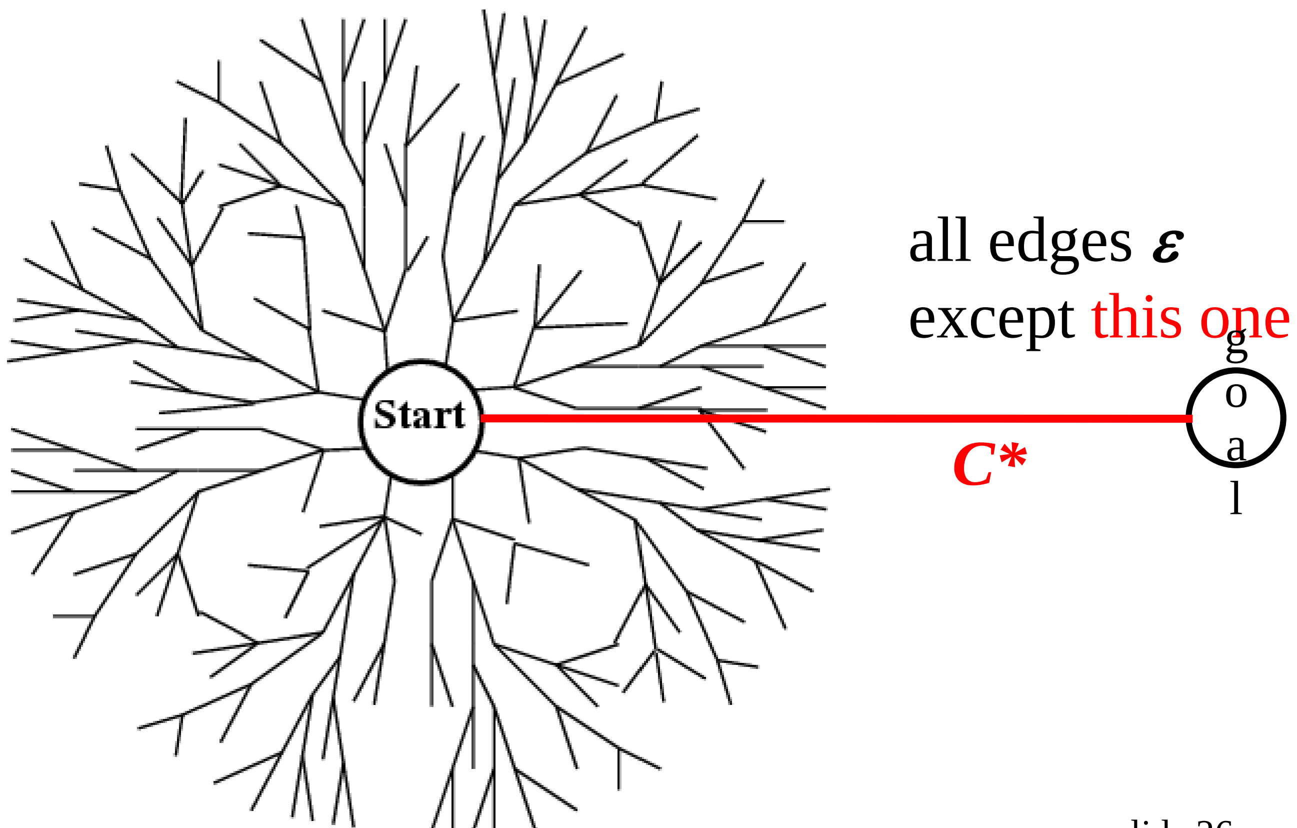
- Find the least-cost goal
- Each node has a path cost from start (= sum of edge costs along the path). Expand the least cost node first.
- Use a priority queue instead of a normal queue
  - Always take out the least cost item
  - Remember *heap*? time ***O***(***log***(#items in heap))

That's it*

* Complications on graphs (instead of trees). Later.

Uniform-cost search (UCS)

- Complete and optimal (if edge costs $\geq \varepsilon > 0$)
- Time and space: can be much worse than BFS
  - Let C^* be the cost of the least-cost goal
  - $O(b^{C^*/\varepsilon})$, possibly $C^*/\varepsilon >> d$



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1. edge cost constant, or positive non-decreasing in depth
- edge costs $\geq \varepsilon > 0$. C* is the best goal path cost.

General State-Space Search Algorithm

```
function general-search(problem, QUEUEING-FUNCTION)
  ;; problem describes the start state, operators, goal test, and
  ;;   operator costs
  ;; queueing-function is a comparator function that ranks two states
  ;; general-search returns either a goal node or "failure"

  nodes = MAKE-QUEUE(MAKE-NODE(problem.INITIAL-STATE))
  loop
    if EMPTY(nodes) then return "failure"
    node = REMOVE-FRONT(nodes)
    if problem.GOAL-TEST(node.STATE) succeeds
then return node
    nodes = QUEUEING-FUNCTION(nodes, EXPAND(node,
                                  problem.OPERATORS))
    ;; succ(s)=EXPAND(s, OPERATORS)
    ;; Note: The goal test is NOT done when nodes are generated
    ;; Note: This algorithm does not detect loops
  end
```

Recall the bad space complexity of BFS

Solution: Uniform-cost search

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- Completeness (not finding all goals): find a goal.
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Solution: Depth-first search

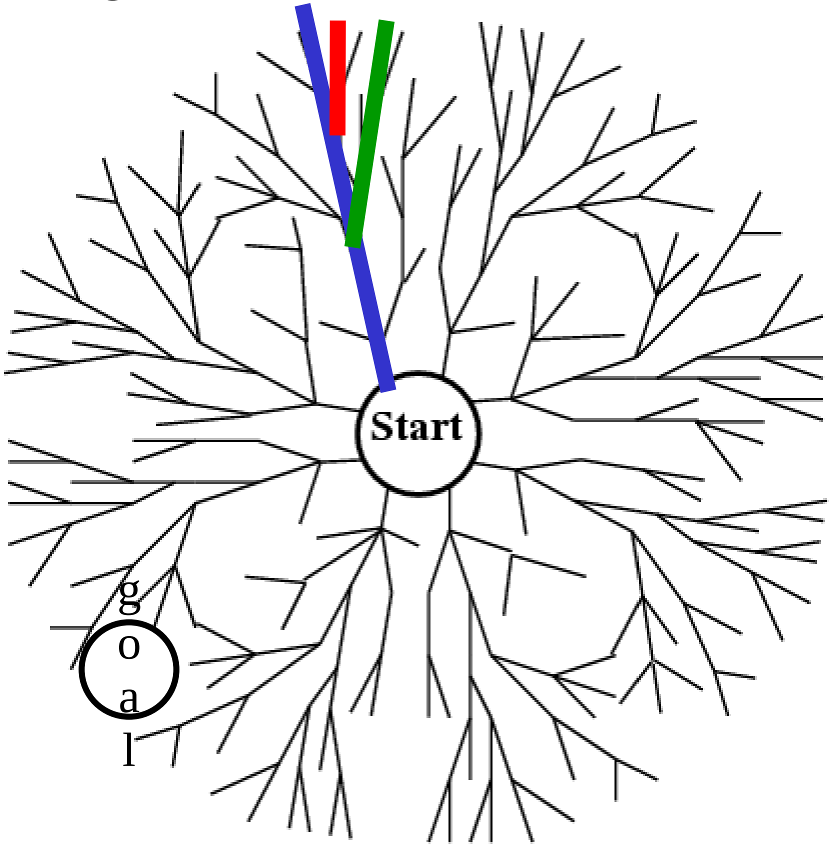
- Space complexity (bad, Figure 3.11)
  - Back points for all generated nodes $O(b^d)$
  - The queue (smaller, but still $O(b^d)$)

Depth-first search

Expand the deepest node first

1. Select a direction, go deep to the end
2. Slightly change the end
3. Slightly change the end some more…

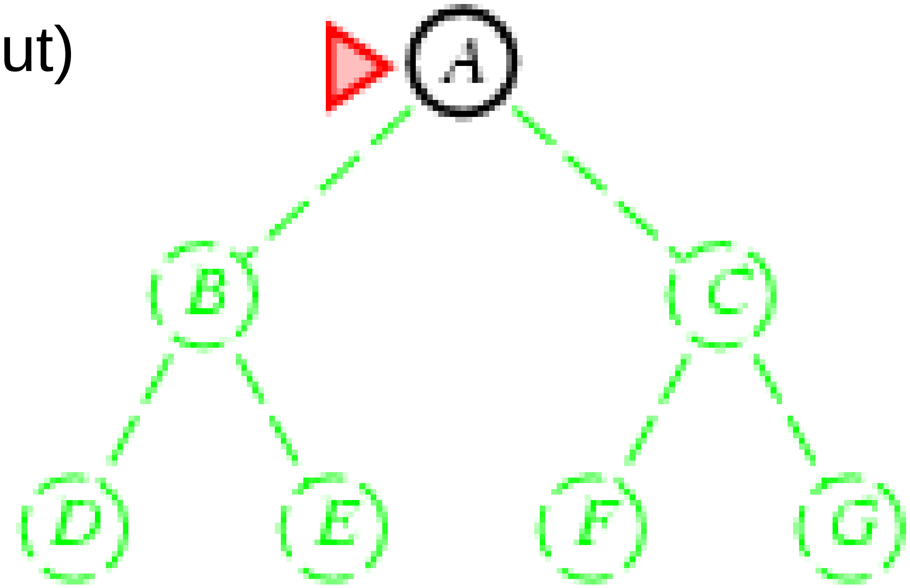
fan



Depth-first search (DFS)

Use a stack (First-in Last-out)

1. push(Initial states)
2. While (stack not empty)
3.   s = pop()
4.   if (s==goal) success!
5.   T = succs(s)
6.   push(T)
7. endWhile



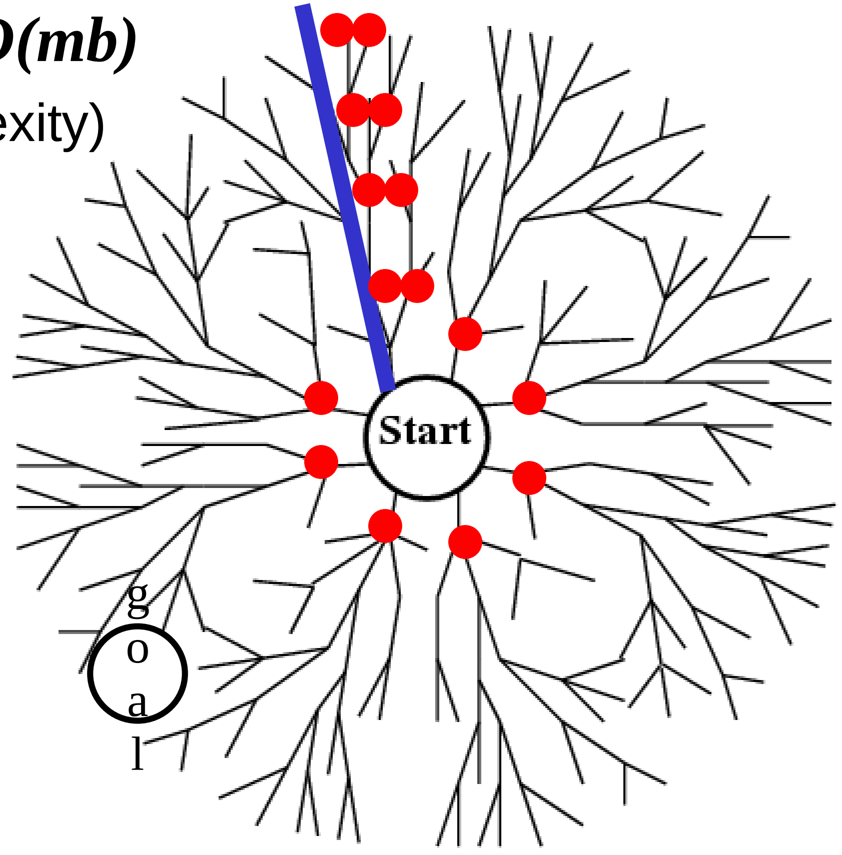
stack (fringe)

[] ⇔

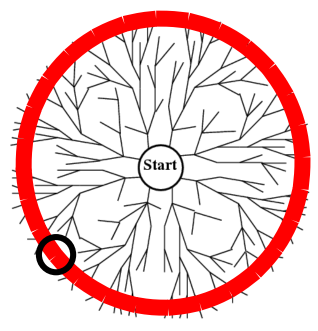
What's in the fringe for DFS?

- m = maximum depth of graph from start
- ***m(b-1) ~ O(mb)***

(Space complexity)



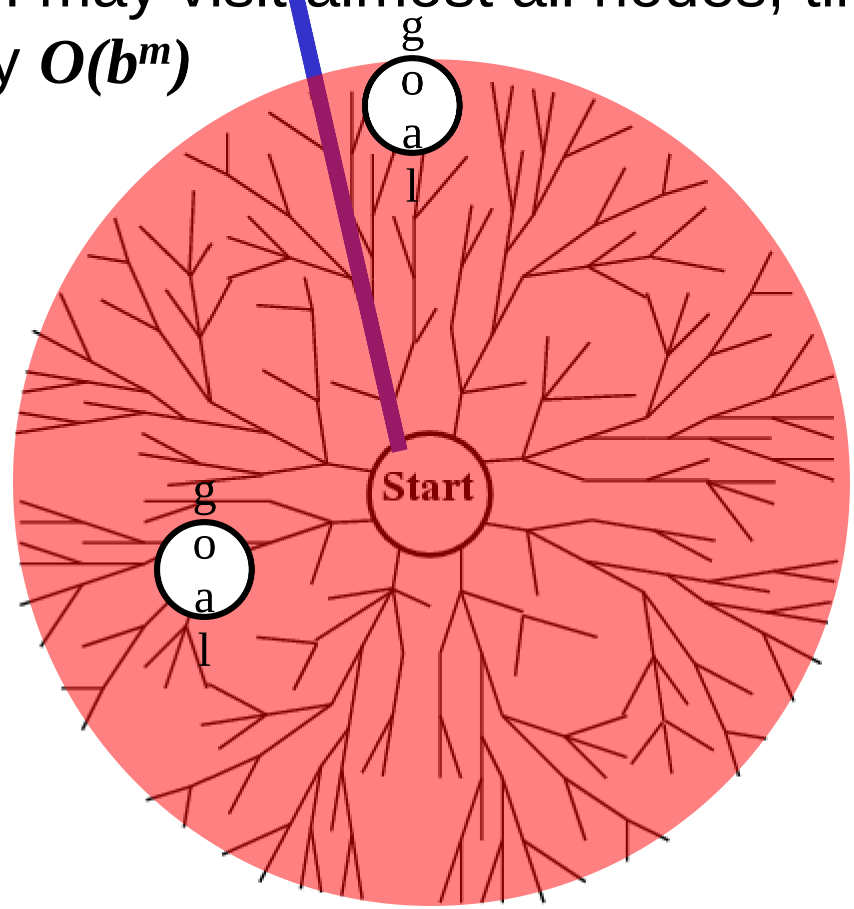
c.f. BFS $\boldsymbol{O(b^d)}$



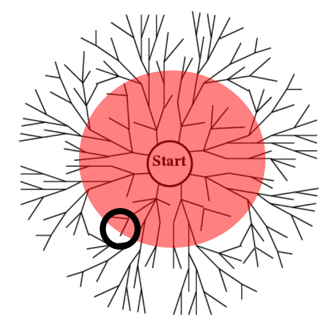
- "backtracking search" even less space
  - generate siblings (if applicable)

What's wrong with DFS?

- Infinite tree: may not find goal (incomplete)
- May not be optimal
- Finite tree: may visit almost all nodes, time complexity $O(b^m)$



c.f. BFS $O(b^d)$



Performance of search algorithms on trees

b: branching factor (assume finite) d: goal depth m: graph depth

	Complete	optimal	time	space
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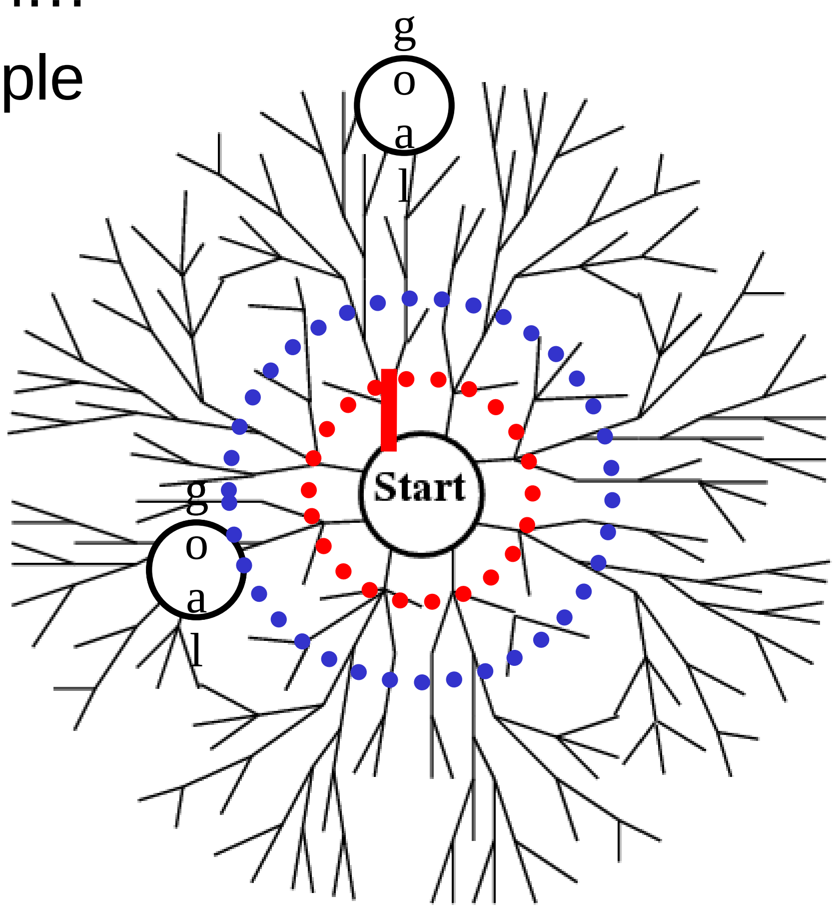
Uniform-cost search[2]	Y	Y	$O(b^{C^*/\varepsilon})$	$O(b^{C^*/\varepsilon})$
Depth-first search	N	N	$O(b^m)$	$O(bm)$

1. edge cost constant, or positive non-decreasing in depth
- edge costs $\geq \varepsilon > 0$. C* is the best goal path cost.

How about this?

1. DFS, but stop if path length > 1.
2. If goal not found, repeat DFS, stop if path length >2.
3. And so on…

fan within ripple



Iterative deepening

- Search proceeds like BFS, but fringe is like DFS
  - Complete, optimal like BFS
  - Small space complexity like DFS
- A huge waste?
  - Each deepening repeats DFS from the beginning
  - No! $db+(d-1)b^2+(d-2)b^3+\ldots+b^d \sim O(b^d)$
  - Time complexity like BFS
- Preferred uninformed search method

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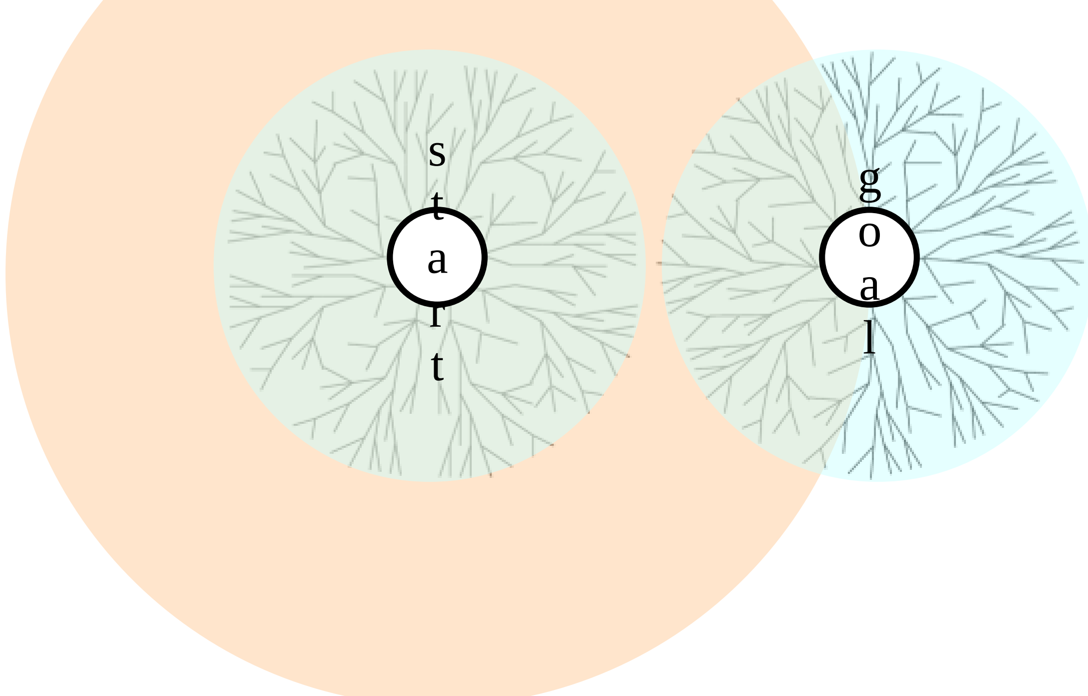
Breadth-first			$O(b^d)$	$O(b^d)$
			$O(b^{C^*/\varepsilon})$	$O(b^{C^*/\varepsilon})$
			$O(b^m)$	$O(bm)$
It d			$O(b^d)$	$O(bd)$

How to reduce the number of states we have to generate?

1. edge cost constant, or positive non-decreasing in depth
- edge costs $\geq \varepsilon > 0$. C^* is the best goal path cost.

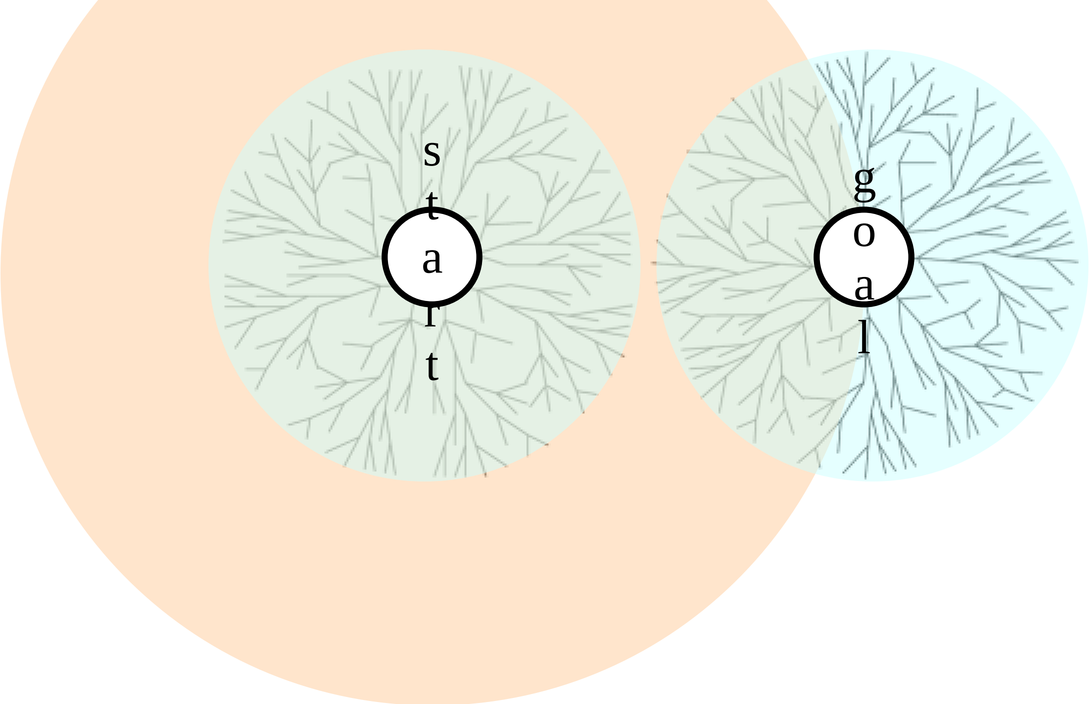
Bidirectional search

- Breadth-first search from both start and goal
- Fringes meet
- Generates $O(b^{d/2})$ instead of $O(b^d)$ nodes



Bidirectional search

- But
  - The fringes are $O(b^{d/2})$
  - How do you start from the 8-queens goals?



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	Complete	optimal	time	space
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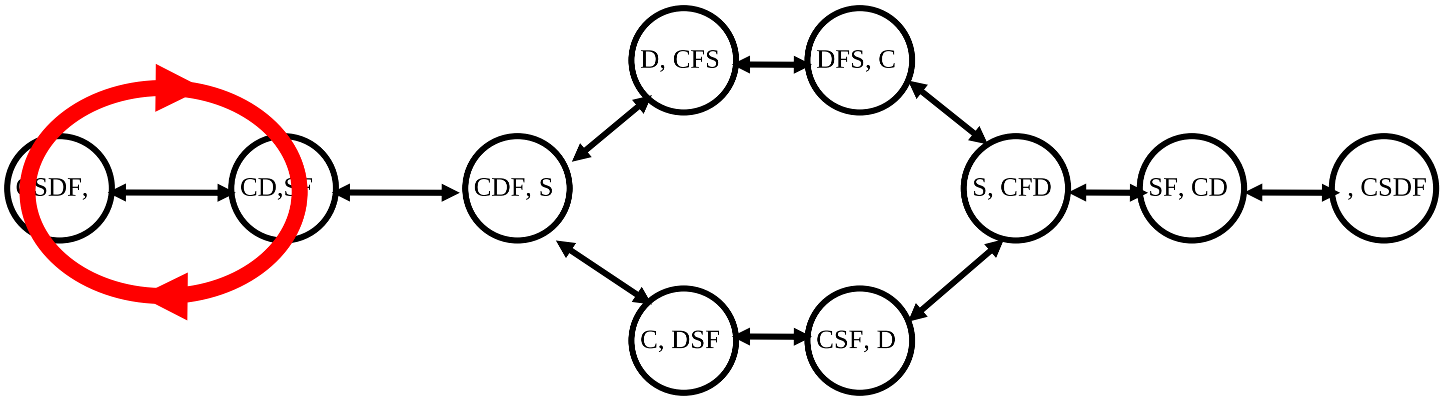
Uniform-cost search[2]	Y	Y	$O(b^{C^*/\varepsilon})$	$O(b^{C^*/\varepsilon})$
Depth-first search	N	N	$O(b^m)$	$O(bm)$
Iterative deepening	Y	Y, if [1]	$O(b^d)$	$O(bd)$
Bidirectional search[3]	Y	Y, if [1]	$O(b^{d/2})$	$O(b^{d/2})$

1. edge cost constant, or positive non-decreasing in depth
- edge costs $\geq \varepsilon > 0$. C^* is the best goal path cost.
- both directions BFS; not always feasible.

If state space graph is not a tree

- The problem: repeated states



- Ignore the danger of repeated states: wasteful (BFS) or impossible (DFS). Can you see why?
- How to prevent it?

If state space graph is not a tree

- We have to remember already-expanded states (CLOSED).
- When we take out a state from the fringe (OPEN), check whether it is in CLOSED (already expanded).
  - If yes, throw it away.
  - If no, expand it (add successors to OPEN), and move it to CLOSED.

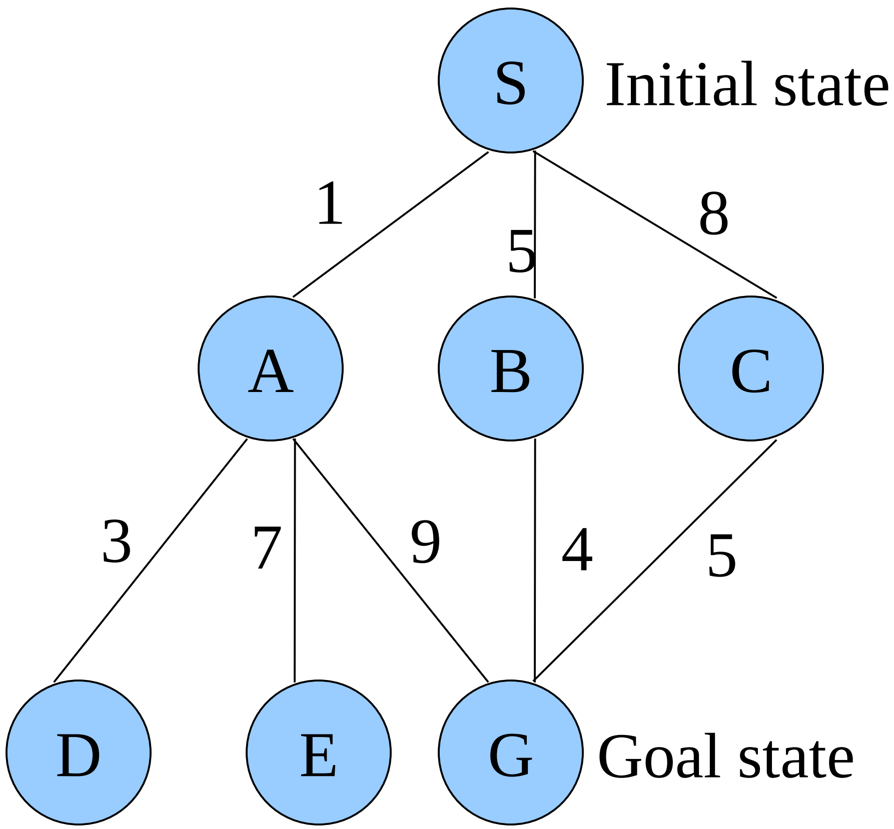
If state space graph is not a tree

- BFS:
  - Still $O(b^d)$ space complexity, not worse
- DFS:
  - Known as Memorizing DFS (MEMDFS)
    - Space and time now $O(min(N, b^M))$ – much worse!
    - N: number of states in problem
    - M: length of longest cycle-free path from start to anywhere
  - Alternative: Path Check DFS (PCDFS): remember only expanded states on current path (from start to the current node)
    - Space O(M)
    - Time $O(b^M)$

Path Checking DFS

1. Maintain a “prefix” path from root to current node, initially empty.
2. Pop a state s. If s in prefix, skip to next pop
3. Goal-checking s.
4. s comes with a backpointer to its parent p. The prefix should contain p somewhere as in initial, ..., p, ...
5. Remove everything after p and put s there, so prefix is now initial, ..., p, s.
6. When you generate a successor t of s, check if t is in prefix or stack. If no, push t to the stack; if yes, do not push it.

Example



(All edges are directed, pointing downwards)

Nodes expanded by:

- Depth-First Search: S A D E G

  Solution found: S A G

- Breadth-First Search: S A B C D E G

  Solution found: S A G

- Uniform-Cost Search: S A D B C E G

  Solution found: S B G (This is the only uninformed search that worries about costs.)

- Iterative-Deepening Search: S A B C S A D E G

  Solution found: S A G

Depth-First Search

```
expanded
node          nodes list
----          ----------
              { S }
  S           { A B C }
  A           { D E G B C}
  D           { E G B C }
  E           { G B C }
  G           { B C }

Solution path found is S A G  <-- this G has cost 10
Number of nodes expanded (including goal node) = 5
```

Breadth-First Search

```
expanded
node          nodes list
----          ----------
              { S }
  S           { A B C }
  A           { B C D E G }
  B           { C D E G G' }
  C           { D E G G' G" }
  D           { E G G' G" }
  E           { G G' G" }
  G           { G' G" }

Solution path found is S A G  <-- this G also has cost 10
Number of nodes expanded (including goal node) = 7
```

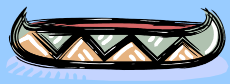
Uniform-Cost Search

```
expanded
node          nodes list
----          ----------
              { S }
  S           { A(1) B(5) C(8) }
  A           { D(4) B(5) C(8) E(8) G(10) }    (note, we don't return G)
  D           { B(5) C(8) E(8) G(10) }
  B           { C(8) E(8) G(9) G(10) }
  C           { E(8) G(9) G(10) G(13) }
  E           { G(9) G(10) G(13) }
  G           { }

Solution path found is S B G  <-- this G has cost 9, not 10
Number of nodes expanded (including goal node) = 7
```

What you should know

- Problem solving as search: state, successors, goal test
- Uninformed search
  - Breadth-first search
    - Uniform-cost search
  - Depth-first search
  - Iterative deepening
  - Bidirectional search



- Can you unify them (except bidirectional) using the same algorithm, with different priority functions?
- Performance measures
  - Completeness, optimality, time complexity, space complexity