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CML Interpreter Design Document

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0.2	06-03-2014	Anders Kael Malmos	Added introduction and domain description
0.3	26-03-2014	Anders Kael Malmos	Added draft of static and dynamic structure for the core interpreter

¹⁸ **Abstract**

¹⁹ This document describes the overall design of the CML interpreter and provides an
²⁰ overview of the code structure targeting developers. It assume a basic knowledge of
²¹ CML.

22 **Contents**

23 1 Introduction	6
24 1.1 Problem Domain	6
25 1.2 Definitions	8
26 2 Software Layers	9
27 3 Layer design and Implementation	9
28 3.1 The Core Layer	10
29 3.2 The Dynamic Model	16
30 3.3 The IDE Layer	18

³¹ 1 Introduction

³² This document is targeted at developers and describes the overall design of the CML
³³ simulator, it is not a detailed description of every part of the source code. This kind of
³⁴ documentation is done in Javadoc and can be generated automatically from the code.
³⁵ It is assumed that common design patterns are known like ?? and a basic understanding
³⁶ of CML.

³⁷ 1.1 Problem Domain

³⁸ The goal of the interpreter is to enable simulation/animation of a given CML [?] model
³⁹ and be able to visualize this in the Eclipse IDE Debugger. CML has a UTP semantics
⁴⁰ defined in [?] which dictates the interpretation. Therefore, the overall goal of the CML
⁴¹ interpreter is to adhere to the semantic rules defined in those documents and to visualize
⁴² this in the Eclipse Debugger.

⁴³ In order to get a high level understanding of how CML is interpreted without knowing
⁴⁴ all the details, a short illustration of how the interpreter represents and evolves a CML
⁴⁵ model is given below.

⁴⁶ In Listing 1 a CML model consisting of three CML processes is given. It has a R
⁴⁷ (Reader) process which reads a value from the inp channel and writes it on the out
⁴⁸ channel. The W (Writer) process writes the value 1 to the inp channel and finishes.
⁴⁹ The S (System) process is a parallel composition of these two processes where they
⁵⁰ must synchronize all events on the inp channel.

```
51 channels
52 inp : int
53 out : int
54
55 process W =
56 begin
57   @ inp!1 -> Skip
58 end
59
60 process R =
61 begin
62   @ inp?x -> out!x -> Skip
63 end
64
65 process S = W || ($inp$) || R
```

Listing 1: A process S composed of a parallel composition of a reader and writer process

⁶⁶ The interpretation of a CML model is done through a series of steps/transitions starting
⁶⁷ from a given entry point. In figure 1 the first step in the interpretation of the model
⁶⁸ is shown, it is assumed that the S process is given as the starting point. Process are
⁶⁹ represented as a circle along with its current position in the model. Each step of the
⁷⁰ execution is split up in two phases, the inspection phase and the execution phase. The
⁷¹ dashed lines represents the environment (another actor that invokes the operation e.g a
⁷² human user or another process) initiating the phase.

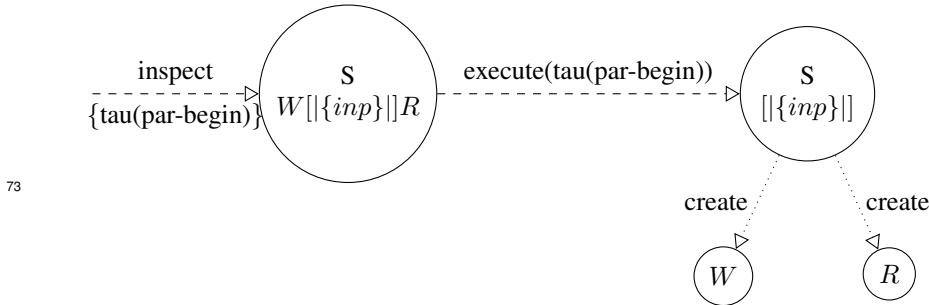


Figure 1: Initial step of Listing 1 with process S as entry point.

74 The inspection phase determines the possible transitions that are available in the next
 75 step of execution. The result of the inspection is shown as a set of transitions below
 76 “inspect”. As seen on figure Figure 1 process P starts out by pointing to the parallel
 77 composition constructs, this construct has a semantic begin rule which does the initial-
 78 ization needed. In the figure Figure 1 that rule is named tau(par-begin) and is therefore
 79 returned from the inspection. The reason for the name tau(..) is that transitions can
 80 be either observable or silent, so in principle any tau transition is not observable from
 81 the outside of the process. However, in the interpreter all transitions flows out of the
 82 inspection phase. When the inspection phase has completed, the execution phase be-
 83 gins. The execution phase executes one of the transitions returned from the inspection
 84 phase. In this case, only a single transition is available so the tau(par-begin)) is ex-
 85 ecuted which creates the two child processes. The result of each of the shown steps
 86 are the first configuration shown in the next step. So in this case the resulting process
 87 configuration of Figure 1 is shown in figure Figure 2.

88 The second step on Figure 2 has a more interesting inspection phase. According to the
 89 parallel composition rule, we have that any event on the *inp* channel must be synchro-
 90 nized, meaning that W and R must only perform transition that involves *inp* channel
 91 events synchronously.

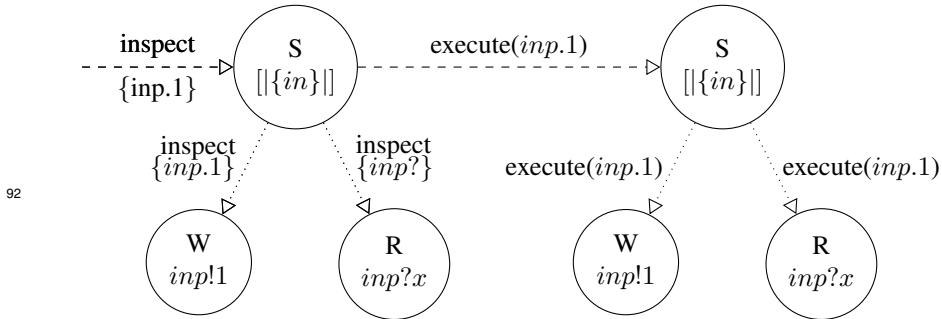


Figure 2: Second step of Listing 1 with S as entry point.

93 Therefore, when P is inspected it must inspect its child processes to determine the
 94 possible transitions. In this case W can perform the *inp*.1 event and R can perform
 95 any event on *inp* and therefore, the only possible transition is the one that performs the
 96 *inp*.1 event. This is then given to the execution phase which result in the *inp*.1 event
 97 and moves both child processes into their next state.

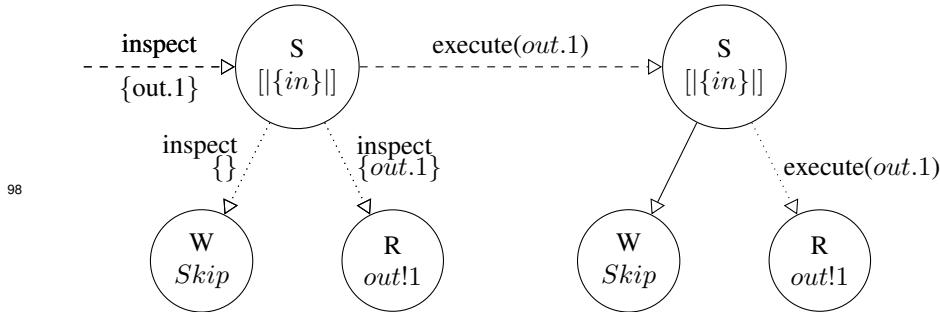


Figure 3: Third step of Listing 1 with S as entry point

99 In the third step on figure Figure 3 W is now Skip which means that it is successfully
 100 terminated. The inspection for W therefore results in an empty set of possible transitions.
 101 R is now waiting for the *out.1* event after 1 was writting to *x* in the last step
 102 and therefore returns this transition. The execution phase is a little different and S now
 103 knows only to execute R.

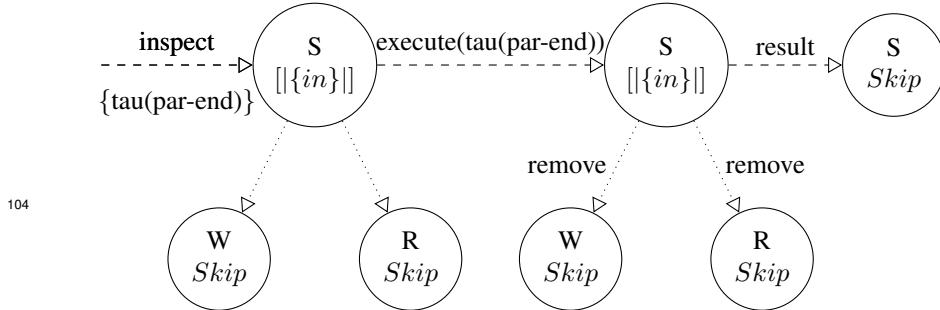


Figure 4: Final step of Listing 1 where the parallel composition collapses onto a Skip process

105 The fourth and final step shown in Figure 4 of the interpretation starts out with both W
 106 and R as Skip, this triggeres the parallel end rules, which evolves into Skip. S therefore
 107 returns the silent transition the triggers this end rule.

108 1.2 Definitions

109 **Animation** Animation is when the user are involved in taking the decisions when in-
 110 terpreting the CML model

111 **CML** Compass Modelling Language

112 **UTP** Unified Theory of Programming (a semantic framework)

113 **Simulation** Simulation is when the interpreter runs without any form of user interac-
 114 tion other than starting and stoppping.

115 **trace** A sequence of observable events performed by a behavior.

116 2 Software Layers

117 This section describes the layers of the CML interpreter. As depicted in figure 5 two
 118 highlevel layers exists.

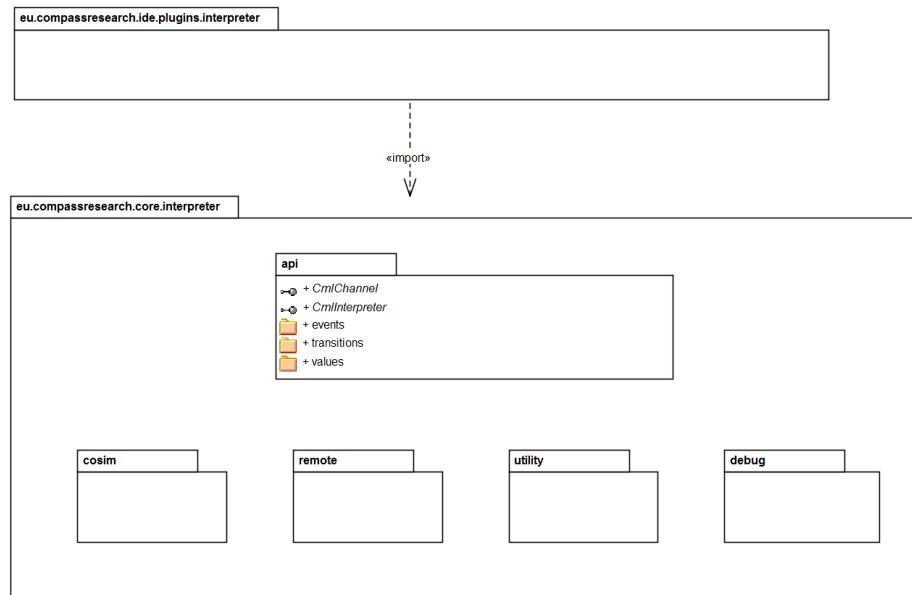


Figure 5: The layers of the CML Interpreter

119 Each of these components will be described in further detail in the following sections.
 120 The major reason behind this layering is that the implementation of the semantics
 121 should be independent of the view showing the results.

122 **Core Layer** This layer has the overall responsibility of interpreting a CML model as
 123 described in the operational semantics that are defined in [?] and is located in the
 124 package *eu.compassresearch.core.interpreter*

125 **IDE Layer** Has the overall responsibility of visualizing the outputs of a running inter-
 126 pretation a CML model in the Eclipse Debugger. It is located in the *eu.compassresearch.ide.plugins.interpreter*
 127 package. The IDE part is integrating the interpreter into Eclipse, enabling CML
 128 models to be debugged through the Eclipse debugger.

129 3 Layer design and Implementation

130 This section describes the static and dynamic structure of the components involved in
 131 interpreting a CML model.

132 3.1 The Core Layer

133 The core layer is responsible for the overall interpretation of a given CML model.
134 To understand some of the choice made, the design philosophy needs a short word.
135 The design philosophy of the top-level structure is to encapsulate all the classes and
136 interfaces (hence make elements package accessible only when appropriate) that makes
137 up the implementation of the core functionality and only expose those that are needed
138 to utilize the interpreter. This provides a clean separation between the implementation
139 and interface and makes it clear for both the users, which not necessarily wants to
140 know about the implementation details, and developers which parts they need to work
141 with.

142 In the following section both the static and dynamic model will be described in more
143 details.

144 3.1.1 The Static Model

145 Packages

146 The following packages defines the top level structure of the core:

147 **eu.compassresearch.core.interpreter** This package contains all the internal classes
148 and interfaces that defines the core functionality of the interpreter. There is one
149 important public class in the package, namely the **VanillaInterpreterFactory** fac-
150 tory class, that any user of the interpreter must invoke to use the interpreter. This
151 can creates instances of the **CmlInterpreter** interface. Furthermore, this pack-
152 age is split into two seperate source folders, each representing a different logical
153 component. The following folders are present:

154 **src/main/java** This folder contains all public classes and interfaces as described
155 above.

156 **src/main/behavior** This folder contains all the internal classes and interfaces
157 that the default interpreter implementation is comprised of. This will be
158 described in more details in Subsection 3.1.1.

159 **eu.compassresearch.core.interpreter.api** This package and sub-packages contains all
160 the public classes and interfaces that defines the API of the interpreter. Some of
161 the most important entities of this package includes the main interpreter interface
162 **CmlInterpreter** along with the **CmlBehaviour** interface that represents a CML
163 process or action. It corresponds to the circles in the figures of Subsection 1.1.

164 **eu.compassresearch.core.interpreter.api.events** This package contains all the public
165 components that enable users of the interpreter to subscribe to multiple events
166 (this it not CML channel events) from both **CmlIntepreter** and **CmlBehaviour**
167 instances.

168 **eu.compassresearch.core.interpreter.api.transitions** This package contains all the
169 possible types of transitions that a **CmlBehaviour** instance can make. This will
170 be explained in more detail in section 3.1.1.

171 **eu.compassresearch.core.interpreter.api.values** This package contains all the val-
 172 ues used by the CML interpreter. They represent the values of variables and
 173 constants in a context.

174 **eu.compassresearch.core.interpreter.cosim** Has the responsibility of running a co-
 175 simulation. A co-simulation can be either between multiple instances of the
 176 CML intepreter co-simulating a CML model, or a CML intepreter instance co-
 177 simulating a CML model with a real live system.

178 **eu.compassresearch.core.interpreter.remote** This has the responsiblilty of exposing
 179 the CML interpreter to be remote controlled.

180 **eu.compassresearch.core.interpreter.debug** Has the responsibility of controlling a
 181 debugging sessions, which only includes the Eclipse debugger at this point.

182 **eu.compassresearch.core.interpreter.utility** The utility packages contains reusable
 183 classes and interfaces that are use across packages.

184 **The Top Level Elements**

185 The top level interfaces and classes of the interpreter structure is depicted in Figure 6,
 186 followed by a short description of each the depicted components.

187 Before going into detais with each element on figure 6 a few things needs mentioning.
 188 First of all, any CML model has a top level Process. Because of this, the interpreter
 189 need only to interact with the top level CmlBehaviour instance. This explains the one-
 190 to-one correspondence between the CmlInterpreter and the CMLBehaviour. However,
 191 the behavior of top level CmlBehaviour is determined by the binary tree of CmlBe-
 192 haviour instances that itself and it's child behaviours defines. So in effect, the CmlIn-
 193 terpreter along with the selection strategy controls every observable transition that any
 194 CmlBehaviour makes.

195 **CmlInterpreter** The interface exposing the functionality of the interpreter compo-
 196 nent. This interface has the overall responsibility of interpreting. It exposes
 197 methods to inspect and execute and it is implemented by the **VanillaCmlInter-
 198 preter** class in the default simulation settings.

199 **CmlBehavior** Interface that represents a behavior specified by either a CML process
 200 or action. Most importantly it exposes the two methods: *inspect* which calcu-
 201 lates the immediate set of possible transitions that it currently allows and *execute*
 202 which takes one of the possible transitions determined by it's supervisor. This
 203 process is described in Subsection 1.1 where a CmlBehavior is represented as a
 204 circle in the figures. As seen both in Subsection 1.1 and Figure 6 associations
 205 between CmlBehavior instances are structured as a binary tree, where a parent
 206 supervises its child behaviors. In this context supervises means that they control
 207 the flow of possible transitions and determines when to execute them. The reason
 208 for this is that is corresponds nicely to the structure of the CML semantics.

209 **SelectionStrategy** This interface has the responsibility of choosing a CmlTransition
 210 from a given CmlTransitionSet. This could be seen as the last chain in the super-
 211 visor hierarchy, since this is where all the possible transitions flows to and the
 212 decision of which one to execute next is taken here. The purpose of this interface

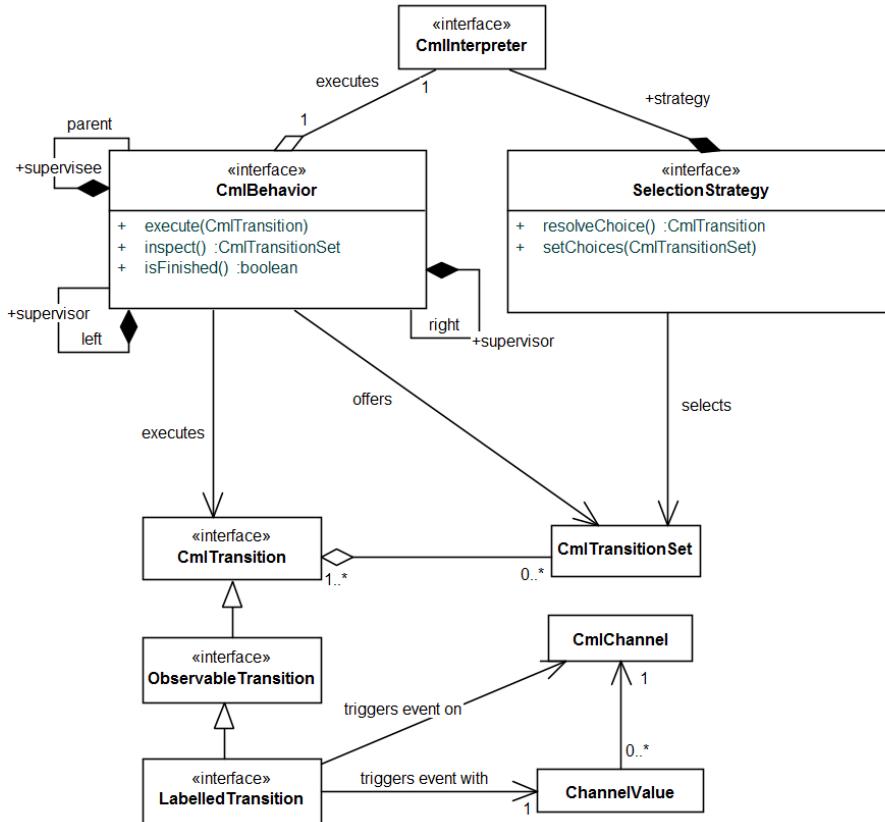


Figure 6: The high level classes and interfaces of the interpreter core component

213 is to allow different kinds of strategies for choosing the next transition. e.g there
 214 is a strategy that picks one at random and another that enables a user to pick.

215 **CmlTransition** Interface that represents any kind of transition that a CmlBehavior
 216 can make. They are not all depicted here and will be described in greater details
 217 in ???. But overall, only transitions that implements the ObservableTransition
 218 interface can produce an observable trace of a behavior.

219 **CmlTransitionSet** This is an immutable set of CmlTransition objects and is the return
 220 value of the inspect method on a CmlBehavior. The reason for it being immutable
 221 is to ensure that calculations never change the input sets.

222 The Transitions Model

223 As described in the previous sections a CML model is represented by a binary tree of
 224 CmlBehaviour instances and each of these has a set of possible transitions that they can
 225 make. A class diagram of all the classes and interfaces that makes up transitions are
 226 shown in figure ???, followed by a description of each of the elements.

227 A transition taken by a CmlBehavior is represented by a CmlTransition. This represent

- 228 a possible next step in the model which can be either observable or silent (also called a
229 tau transition).
- 230 An observable transition represents either that time passes or that a communication/syn-
231 chronization event takes place on a given channel. All of these transitions are captured
232 in the ObservableTransition interface. A silent transitions is captured by the TauTrans-
233 ition and HiddenTransition class and can respectively marks the occurrence of a an
234 internal transition of a behavior or a hidden channel transition.

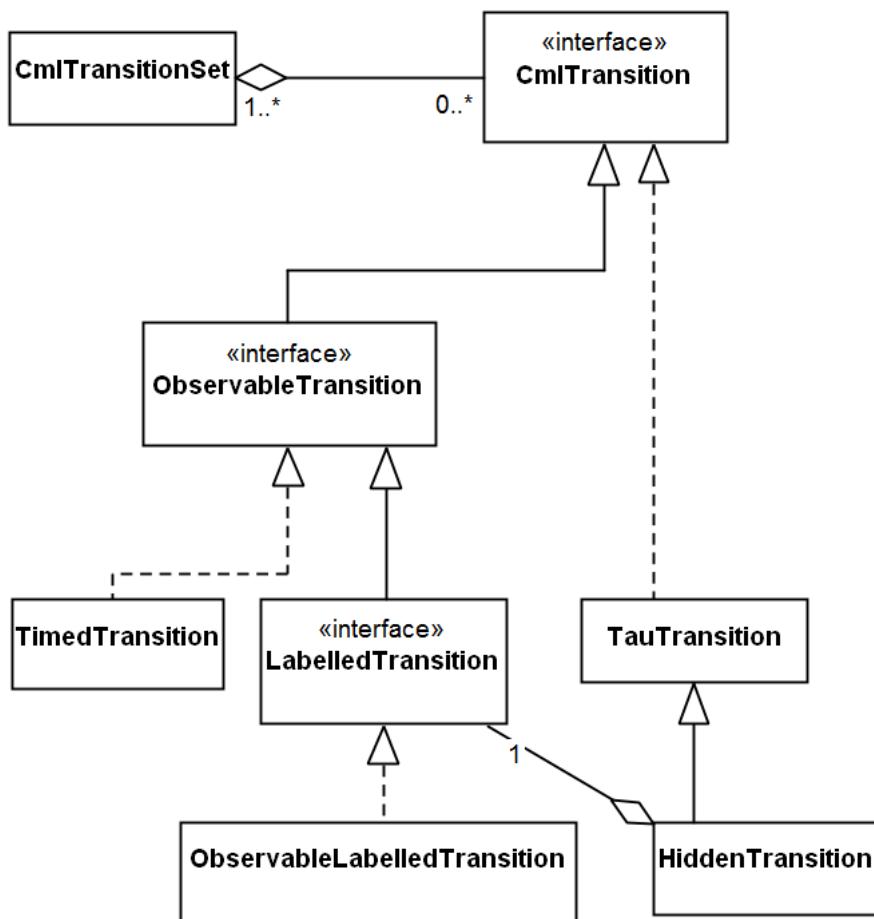


Figure 7: The classes and interfaces that defines transitions

- 235 **CmlTransition** Represents any possible transition.
- 236 **CmlTransitionSet** Represents a set of CmlTransition objects.
- 237 **ObservableTransition** This represents any observable transition.
- 238 **LabelledTransition** This represents any transition that results in a observable channel
239 event
- 240 **TimedTransition** This represents a tock event marking the passage of a time unit.

- 241 **ObservableLabelledTransition** This represents the occurrence of a observable channel event which can be either a communication event or a synchronization event.
- 242 **TauTransition** This represents any non-observable transitions that can be taken in a behavior.
- 243 **HiddenEvent** This represents the occurrence of a hidden channel event in the form of a tau transition.

247 The Default CmlBehavior Implementation

- 248 Actions and processes are both represented by the CmlBehaviour interface. A class diagram of the important classes that implements this interface is shown in Figure 8 When

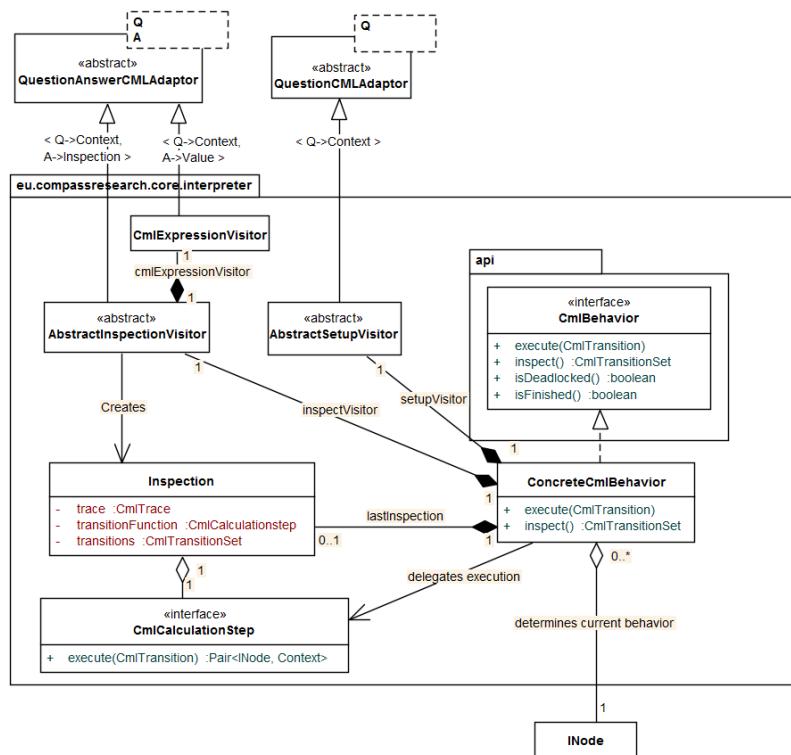


Figure 8: The classes and interfaces making up the default implementation the CmlBehavior interface

- 249 the interpreter runs in the default operation mode, meaning where only a single interpreter instance runs (opposed to the co-simulation modes where multiple instances of the interpreter might run or connected to an externally running system). Then all CmlBehavior instances will be in the form of the ConcreteCmlBehavior class. As described above a CmlBehavior has the responsibility to behave as a given action or process. However, as shown in Figure 8 the ConcreteCmlBehavior class delegates a large part of its responsibility to other classes. The actual behavior of a ConcreteCmlBehavior instance is decided by its current INode instance, so when a ConcreteCmlBehavior instance is created a INode instance must be given. The INode interface is implemented

259 by all the CML AST nodes and can therefore be any CML process or action. The
 260 actual implementation of the behavior of any process/action is delegated to internal
 261 visitor classes as depicted in Figure 8. The used visitors are all extending generated
 262 abstract visitors that have the infrastructure to visit any CML AST node. The reason for
 263 this structure is to be able to utilize the already generated visitors by the AST-creator
 264 [] that enables traversing of CML AST's.

265 Here a brief description of each new element depicted in Figure 8:

266 **CmlExpressionVisitor** This has the responsibility to evaluate CML expressions given
 267 a Context.

268 **AbstractSetupVisitor** This has the responsibility of performing any required setup
 269 for a behavior. This visitor is invoked whenever a new INode instance is loaded.

270 **AbstractAlphabetVisitor** This has the responsibility of creating an Inspection object
 271 given the current state of the behavior, which is represented by a INode and a
 272 Context object.

273 **Inspection** Contains the next possible transitions (in a CmlTransitionSet) along with
 274 a transition function in the form of a CmlCalculationStep.

275 **CmlCalculationStep** Responsible for executing the actual behavior that occurs in a
 276 transition from one state to another. This is where the actual implementation of
 277 the semantics is.

278 The Visitors

279 In figure 9 a more detailed look at the inspection visitor structure is given.

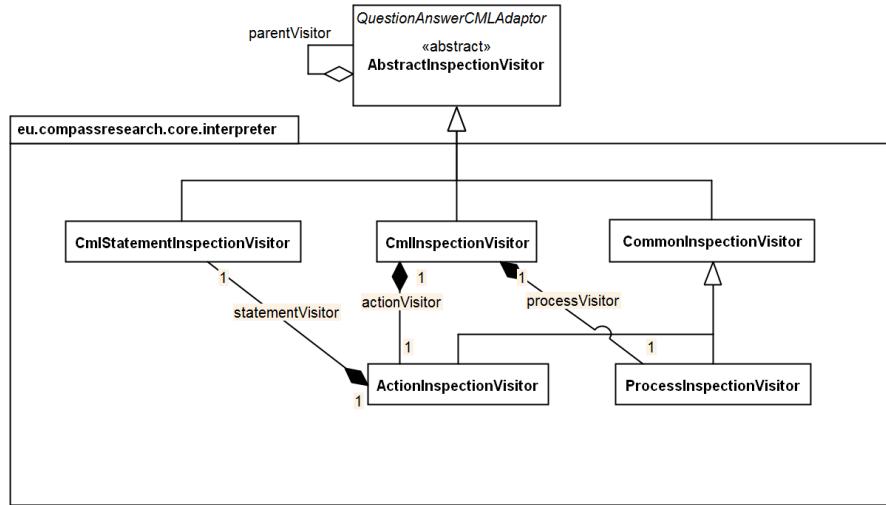


Figure 9: Visitor structure

280 As depicted the visitors are split into several visitors that handle different parts of the
 281 CML language. The sole reason for doing this is to avoid having one large visitor

282 that handles all the cases. At run-time the visitors are setup in a tree structure. For
 283 the inspection the top most visitor is the CmlInspectionVisitor which then delegates
 284 to either ActionInspectionVisitor or ProcessEvaluationVisitor depending on the given
 285 INode. This structures resembles the structure of the setup visitors.

286 The CmlExpressionVisitor is however a little different from the others. It takes care
 287 of all CML expressions, but delegates the entire subset of VDM expression constructs
 288 that are contained in CML to the DelegateExpressionEvaluator overture class, which
 289 can evaluate VDM expressions. The reason for doing this is of course reuse.

290 **3.2 The Dynamic Model**

291 This section will describe the high-level dynamic model. First of all, in the default oper-
 292 ation mode (as mentioned above a single running instance of the interpreter) the entire
 293 CML interpreter runs in a single thread. This is mainly due to the inherent complex-
 294 ity of concurrent programming. You could argue that since a large part of COMPASS
 295 is about modelling complex concurrent systems, we also need a concurrent interpre-
 296 tation of the models. However, the semantics is perfectly implementable in a single
 297 thread which makes a multi-threaded interpreter optional. There are of course benefits
 298 to a multi-threaded interpreter, but for matters such as the testing and deterministic be-
 299 haviour a single threaded interpreter is much easier to handle and comprehend.

300 **The Top Execution Loop**

301 To start a simulation/animation of a CML model, you first of all need an instance of
 302 the CmlInterpreter interface. This is created through the VanillaInterpreterFactory by
 303 invoking the newInterpreter method with a typechecked AST of the CML model. The
 304 default returned instance is the VanillaCmlInterpreter class. Once a CmlInterpreter is
 305 instantiated the interpretation of the CML model is started by invoking the execute
 306 method.

307 In figure 10 a sequence diagram of the execute method on the VanillaCmlInterpreter
 308 class is depicted.

309 As seen in the figure the execution continues until the top level process is either suc-
 310 cessfully terminated or deadlocked. Each round taken in this loop is one step taken in
 311 the model, where the meaning of a step is explained in Subsection 1.1 with an inspec-
 312 tion and execution phase. The actual decision of which transition to be taken next is
 313 decided by the given SelectionStrategy instance to the execute method. This decision
 314 is delegated to the two methods setChoices and resolveChoice.

315 **Dynamics of the ConcreteCmlBehavior**

316 As mentioned multiple times the ConcreteCmlBehavior class is the default realization
 317 of the CmlBehavior interface and is the only one of them explained in details in this
 318 report. To understand the dynamic model we need to see what happens in the inspect
 319 and execute methods, as these together determines the possible transitions at the top
 320 level shown in the last section.

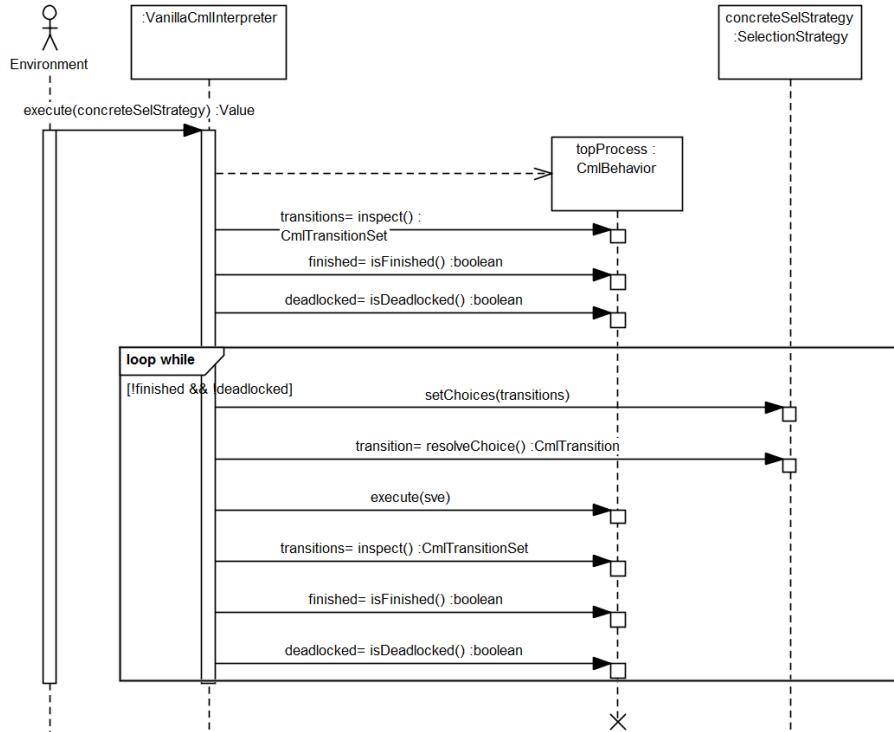


Figure 10: The top level dynamics

321 In Figure 11 the general inspect dynamics is depicted. When the inspect method is
 322 called on a ConcreteCmlBehavior it uses its nextNode (in the java source nextNode
 323 and nextContext is actually a Pair₁INode, Context₂) to delegate the actual inspection
 324 to the CmlInspectionVisitor. The CmlInspectionVisitor contains a method case₁INode
 325 instance name₂ for every CML AST node. So e.g. if nextNode is a AInterleavingAction
 326 then the visitor method caseAInterleavingAction(..) method is called with the
 next

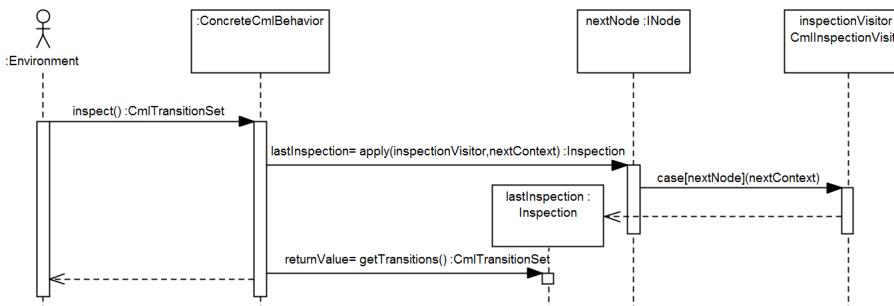


Figure 11: The general dynamics of the inspect method

327 possible transitions and a transition function to be called if the execute method is to be
 328 invoked. The last call in Figure 11 just grabs the CmlTransitionSet from the returned
 329

330 Inspection object and return this as the result of the inspect call.

331 The execute method, shown in Figure 12, will execute the given transition and must
 332 only be called if one the returned transitions from the inspect method has been chosen
 333 for execution. The actual execution is delegated to the CmlCalculationStep instance
 334 contained in the last calculated Inspection object (lastInspection in the figure) in the
 inspect method. The instance of a CmlCalculationStep is an anonymous class created

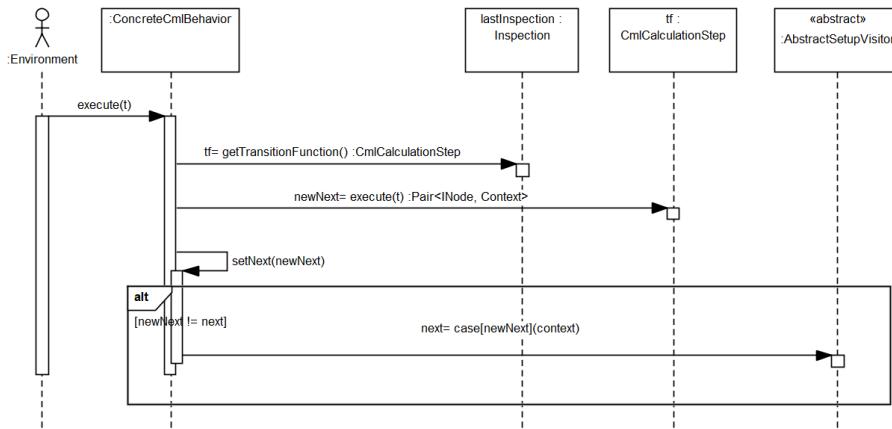


Figure 12: The general dynamics of the execute method

335 in a inspection visitor case, so the behavior of the execute method is entirely dependent
 336 on the current node contained in the next pair. The result of the execute method is the
 337 next node and context. As seen in Figure 12 the setup visitor is called if the newly
 338 returned pair is different from the current one. This enables any case specific setup
 339 behavior to be implemented here. E.g in some cases the context needs to be updated
 340 before the inspection phase is commenced.
 341

3.3 The IDE Layer

3.3.1 Deployment Model

344 In order to get the big picture of how the IDE layer works together with the Core, a
 345 deployment view of the IDE is shown in Figure 13.

346 An Eclipse debugging session involves two JVM instances, the one that the Eclipse
 347 platform is executing in and one where only the Core executes in. All communication
 348 between them is done via JSON through a TCP connection.

349 The JSON protocol need to be defined

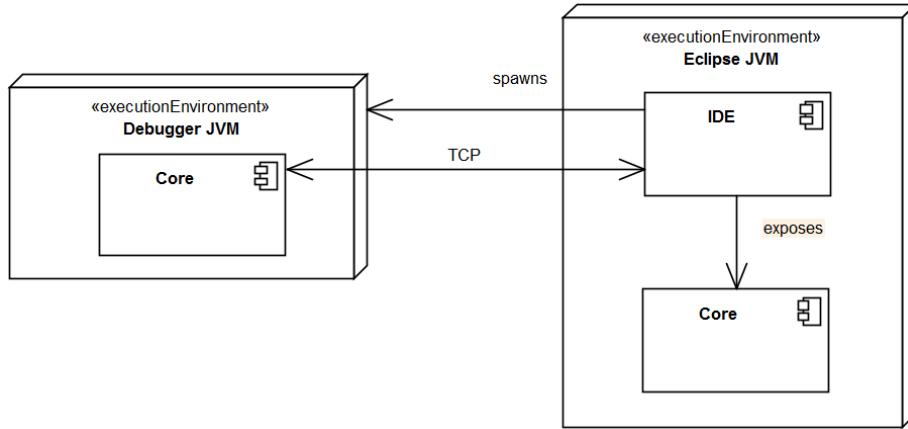


Figure 13: Deployment diagram of the debugger

3.3.2 Static Model

351 Packages

- 352 The following packages defines the top level structure of the IDE:
- 353 **eu.compassresearch.ide.interpreter.model** Contains all the classes that implements the Eclipse debug model [?]
- 354 **eu.compassresearch.ide.interpreter.launching** Classes that deals with launching a debugging session.
- 355 **eu.compassresearch.ide.interpreter.protocol** Classes that deals with the communication between the Eclipse CML debugger and a CML interpreter instance.
- 356 **eu.compassresearch.ide.interpreter.view** Contains the custom views of the Eclipse CML debugger.
- 361 Before explaining the steps involved in a debugging session, there are two important classes worth mentioning:
- 363 **CmlDebugger** Interface with the responsibility of controlling the CmlInterpreter execution in a debugging session.
- 365 **SocketServerDebugger** Realization of the CmlDebugger interface that enables controlling the debugging session over a tcp connection.
- 367 **DebugMain** Class that contains the main method that initializes the core component on the debugger JVM side. This involves
- 369 **CmlDebugTarget** This class is part of the Eclipse debugging model. It has the responsibility of representing a running interpreter on the Eclipse side. All communications to and from the Eclipse debugger are handled in this class.
- 372 A debugging session has the following steps:
- 373 1. The user launches a debug session

- 374 2. On the Eclipse JVM a CmlDebugTarget instance is created, which listens for an
375 incomming TCP connection.
- 376 3. A Debugger JVM is spawned with the main method in the DebugMain class as
377 starting point.
- 378 4. A SocketServerDebugger instance is created and tries to connect to the created
379 connection from step 2.
- 380 5. When the connection is established, the SocketServerDebugger will send a START-
381 ING status message along with additional details
- 382 6. The **CmlDebugTarget** updates the GUI accordingly.
- 383 7. When the interpreter is running, status messages will be sent from SocketServerDe-
384 bugger and commands and request messages are sent from **CmlDebugTarget**
- 385 8. This continues until either the CML model successfully terminates or the user
386 stops.