Wizards

Standard Code Library

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1. 数论

1.1 $O(m^2 \log n)$ 线性递推

```
Given a_0, a_1, \ldots, a_{m-1}

a_n = c_0 \times a_{n-m} + \cdots + c_{m-1} \times a_{n-1}

Solve for a_n = v_0 \times a_0 + v_1 \times a_1 + \cdots + v_{m-1} \times a_{m-1}
```

```
1
   void linear_recurrence(long long n, int m, int a[], int
      \hookrightarrow \texttt{c[], int p) } \{
 2
     long long v[M] = \{1 \% p\}, u[M << 1], msk = !!n;
     for(long long i(n); i > 1; i >>= 1) {
 3
 4
       msk <<= 1:
 5
 6
     for(long long x(0); msk; msk >>= 1, x <<= 1) {
 7
       fill_n(u, m << 1, 0);
       int b(!!(n & msk));
 8
9
       x \mid = b;
10
        if(x < m) {
11
          u[x] = 1 \% p;
12
        }else {
13
          for(int i(0); i < m; i++) {</pre>
14
            for(int j(0), t(i + b); j < m; j++, t++) {
15
              u[t] = (u[t] + v[i] * v[j]) % p;
            }
16
17
          for(int i((m << 1) - 1); i >= m; i--) {
18
19
            for(int j(0), t(i - m); j < m; j++, t++) {
              u[t] = (u[t] + c[j] * u[i]) % p;
20
21
22
          }
23
        }
24
        copy(u, u + m, v);
     }
25
     //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] *
26
        \hookrightarrow a[m - 1].
27
     for(int i(m); i < 2 * m; i++) {</pre>
28
        a[i] = 0:
29
        for(int j(0); j < m; j++) {</pre>
30
          a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31
     }
32
33
     for(int j(0); j < m; j++) {
       b[j] = 0;
        for(int i(0); i < m; i++) {</pre>
35
36
          b[j] = (b[j] + v[i] * a[i + j]) % p;
37
38
     }
39
     for(int j(0); j < m; j++) {
40
        a[j] = b[j];
41
  }
42
```

1.2 求逆元

```
void ex_gcd(long long a, long long b, long long &x, long
      \hookrightarrow long &y) {
     if (b == 0) {
       x = 1;
3
       y = 0;
4
5
       return;
7
     long long xx, yy;
8
     ex_gcd(b, a % b, xx, yy);
9
     y = xx - a / b * yy;
10
     x = yy;
11
   }
12
   long long inv(long long x, long long MODN) {
13
14
     long long inv_x, y;
     ex_gcd(x, MODN, inv_x, y);
15
     return (inv_x % MODN + MODN) % MODN;
16
```

17 }

1.3 中国剩余定理

```
//返回 (ans, M), 其中 ans 是模 M 意义下的解
   std::pair<long long, long long> CRT(const std::vector<long
     \hookrightarrow long>& m, const std::vector<long long>& a) {
3
    long long M = 1, ans = 0;
    int n = m.size();
4
5
    for (int i = 0; i < n; i++) M *= m[i];</pre>
6
    for (int i = 0; i < n; i++) {
       ans = (ans + (M / m[i]) * a[i] % M * inv(M / m[i],
         → m[i])) % M; // 可能需要大整数相乘取模
8
9
    return std::make_pair(ans, M);
10 }
```

1.4 素性测试

```
int strong_pseudo_primetest(long long n,int base) {
2
       long long n2=n-1,res;
3
        int s=0:
 4
        while(n2\%2==0) n2>>=1,s++;
 5
        res=powmod(base,n2,n);
 6
        if((res==1)||(res==n-1)) return 1;
 7
        s--;
 8
        while(s>=0) {
 g
            res=mulmod(res,res,n);
10
            if(res==n-1) return 1;
11
12
       }
13
        return 0; // n is not a strong pseudo prime
14
15
   int isprime(long long n) {
16
     static LL testNum[]={2,3,5,7,11,13,17,19,23,29,31,37};
17
        \rightarrow \lim[]=\{4,0,1373653LL,25326001LL,25000000000LL,21523028987\}
         \rightarrow 3474749660383LL,341550071728321LL,0,0,0,0);
     if(n<2||n==3215031751LL) return 0;
18
     for(int i=0;i<12;++i){</pre>
19
20
        if(n<lim[i]) return 1;</pre>
        if(strong_pseudo_primetest(n,testNum[i])==0) return 0;
21
22
     }
23
     return 1;
24
```

1.5 质因数分解

```
int ansn; LL ans[1000];
   LL func(LL x,LL n){ return(mod_mul(x,x,n)+1)%n; }
   LL Pollard(LL n){
     LL i,x,y,p;
     if(Rabin_Miller(n)) return n;
     if(!(n&1)) return 2;
6
 7
     for(i=1;i<20;i++){</pre>
8
       x=i; y=func(x,n); p=gcd(y-x,n);
9
        while(p==1) {x=func(x,n); y=func(func(y,n),n);
           \hookrightarrow p=\gcd((y-x+n)%n,n)%n;
10
        if(p==0||p==n) continue;
11
        return p;
12
13
14
   void factor(LL n){
     LL x;
16
     x=Pollard(n);
17
     if(x==n){ ans[ansn++]=x; return; }
     factor(x), factor(n/x);
18
19
   }
```

1.6 佩尔方程

```
import java.math.BigInteger;
   import java.util.Scanner;
2
3/a[n]=(g[n]+a[0])/h[n]
4 //g[n] = a[n-1] *h[n-1] -g[n-1]
5 / h[n] = (N-g[n]*g[n])/h[n-1]
6 //p[n]=a[n-1]*p[n-1]+p[n-2]
7 / q[n] = a[n-1] * q[n-1] + q[n-2]
8 //so:
   //p[n]*q[n-1]-p[n-1]*q[n]=(-1)^(n+1);
9
  //p[n]^2-N*q[n]^2=(-1)^(n+1)*h[n+1];
10
   public class Main {
11
12
       public static BigInteger p, q;
       public static void solve(int n) {
13
14
           BigInteger N, p1, p2, q1, q2, a0, a1, a2, g1, g2,
              \hookrightarrow h1. h2:
            g1 = q2 = p1 = BigInteger.ZERO;
15
16
           h1 = q1 = p2 = BigInteger.ONE;
17
            a0 = a1 =

→ BigInteger.valueOf((long)Math.sqrt(1.0*n));
18
           N = BigInteger.valueOf(n);
           while (true) {
19
                g2 = a1.multiply(h1).subtract(g1);
20
                   \hookrightarrow //g2=a1*h1-g1
21
                h2 = N.subtract(g2.pow(2)).divide(h1);
                   \hookrightarrow //h2=(n-g2^2)/h1
                a2 = g2.add(a0).divide(h2);
22
                   \hookrightarrow //a2=(g2+a0)/h2
                p = a1.multiply(p2).add(p1);
23
                   \hookrightarrow //p=a1*p2+p1
                q = a1.multiply(q2).add(q1);
24
                   \leftrightarrow //q=a1*q2+q1
25
                   \Rightarrow == 0) return;//p^2-n*q^2=1
                g1 = g2;h1 = h2;a1 = a2;
26
27
                p1 = p2; p2 = p;
28
                q1 = q2; q2 = q;
29
           }
30
31
32
       public static void main(String[] args) {
           Scanner cin = new Scanner(System.in);
33
34
            int t=cin.nextInt();
            while (t--!=0) {
35
                solve(cin.nextInt());
36
                System.out.println(p + " " + q);
37
38
           }
39
       }
40
  }
```

1.7 二次剩余

```
1 // x^2 = a (mod p), 0 <= a < p, 返回 true or false 代表
     →是否存在解
  // p 必须是质数, 若是多个单次质数的乘积, 可以分别
     →求解再用 CRT 合并
  // 复杂度为 D(log n)
3
  void multiply(ll &c, ll &d, ll a, ll b, ll w) {
      int cc = (a * c + b * d % MOD * w) % MOD;
5
6
      int dd = (a * d + b * c) \% MOD;
      c = cc, d = dd;
7
  }
8
g
10 bool solve(int n, int &x) {
      if (MOD == 2) return x = 1, true;
```

```
12
       if (power(n, MOD / 2, MOD) == MOD - 1) return false;
13
       11 c = 1, d = 0, b = 1, a, w;
14
       // finding a such that a^2 - n is not a square
       do { a = rand() % MOD;
15
           w = (a * a - n + MOD) \% MOD;
16
17
           if (w == 0) return x = a, true;
       } while (power(w, MOD / 2, MOD) != MOD - 1);
18
19
       for (int times = (MOD + 1) / 2; times; times >>= 1) {
20
           if (times & 1) multiply(c, d, a, b, w);
21
           multiply(a, b, a, b, w);
22
23
       // x = (a + sqrt(w)) ^ ((p + 1) / 2)
24
       return x = c, true;
25
```

1.8 一元三次方程

```
double a(p[3]), b(p[2]), c(p[1]), d(p[0]);
   double k(b / a), m(c / a), n(d / a);
   double p(-k * k / 3. + m);
   double q(2. * k * k * k / 27 - k * m / 3. + n);
   Complex omega[3] = \{Complex(1, 0), Complex(-0.5, 0.5 *
      \hookrightarrow sqrt(3)), Complex(-0.5, -0.5 * sqrt(3))};
   Complex r1, r2;
   double delta(q * q / 4 + p * p * p / 27);
8
   if (delta > 0) {
       r1 = cubrt(-q / 2. + sqrt(delta));
10
       r2 = cubrt(-q / 2. - sqrt(delta));
11
   } else {
       r1 = pow(-q / 2. + pow(Complex(delta), 0.5), 1. / 3);
12
       r2 = pow(-q / 2. - pow(Complex(delta), 0.5), 1. / 3);
13
14
   }
15
   for(int _(0); _ < 3; _++) {
16
       Complex x = -k / 3. + r1 * omega[_ * 1] + r2 * omega[_
          \hookrightarrow * 2 % 3];
17
```

1.9 线下整点

```
// \sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor , n,m,a,b>0
  LL solve(LL n,LL a,LL b,LL m){
    if(b==0) return n*(a/m);
     if(a>=m) return n*(a/m)+solve(n,a%m,b,m);
    if(b>=m) return (n-1)*n/2*(b/m)+solve(n,a,b%m,m);
5
6
     return solve((a+b*n)/m,(a+b*n)%m,m,b);
7
```

1.10 线性同余不等式

```
1 // Find the minimal non-negtive solutions for
      \rightarrow l < d \cdot x \mod m < r
   // 0 \le d, l, r \le m; l \le r, O(\log n)
   11 cal(l1 m, l1 d, l1 l, l1 r) {
       if (1 == 0) return 0;
       if (d == 0) return MXL; // 无解
 5
       if (d * 2 > m) return cal(m, m - d, m - r, m - 1);
 6
 7
       if ((1 - 1) / d < r / d) return (1 - 1) / d + 1;
       ll k = cal(d, (-m % d + d) % d, 1 % d, r % d);
 8
       return k == MXL ? MXL : (k * m + 1 - 1) / d + 1; // 无
 9
          →解 2
10 }
11
   // return all x satisfying l1<=x<=r1 and
      \rightarrow 12<=(x*mul+add)%LIM<=r2
   // here LIM = 2^32 so we use UI instead of "%".
13
   // O(\log p + \#solutions)
15
   struct Jump {
16
       UI val, step;
       Jump(UI val, UI step) : val(val), step(step) { }
17
18
       Jump operator + (const Jump & b) const {
```

```
19
            return Jump(val + b.val, step + b.step); }
20
        Jump operator - (const Jump & b) const {
21
            return Jump(val - b.val, step + b.step);
22
   inline Jump operator * (UI x, const Jump & a) {
23
       return Jump(x * a.val, x * a.step);
24
25 }
   vector<UI> solve(UI 11, UI r1, UI 12, UI r2, pair<UI, UI>
26
      \hookrightarrow muladd) {
       UI mul = muladd.first, add = muladd.second, w = r2 -
          Jump up(mul, 1), dn(-mul, 1);
28
29
       UI s(11 * mul + add);
        Jump lo(r2 - s, 0), hi(s - 12, 0);
30
31
       function<void(Jump &, Jump &)> sub = [&](Jump & a,
           \hookrightarrow Jump & b) {
            if (a.val > w) {
32
33
                UI t(((long long)a.val - max(011, w + 111 -
                   \hookrightarrow b.val)) / b.val);
34
                a = a - t * b:
            }
35
36
       };
       sub(lo, up), sub(hi, dn);
37
38
        while (up.val > w || dn.val > w) {
39
            sub(up, dn); sub(lo, up);
            sub(dn, up); sub(hi, dn); }
40
41
        assert(up.val + dn.val > w);
       vector<UI> res;
42
        Jump bg(s + mul * min(lo.step, hi.step), min(lo.step,
43
          \hookrightarrow hi.step));
        while (bg.step <= r1 - 11) {
            if (12 <= bg.val && bg.val <= r2)</pre>
45
46
                res.push_back(bg.step + 11);
47
            if (12 <= bg.val - dn.val && bg.val - dn.val <=
               \hookrightarrow r2) {
48
                bg = bg - dn;
49
            } else bg = bg + up;
50
       } return res;
51
  }
```

1.11 组合数取模

```
1 LL prod=1,P;
   pair<LL,LL> comput(LL n,LL p,LL k){
2
       if(n<=1)return make_pair(0,1);</pre>
3
4
       LL ans=1,cnt=0;
5
       ans=pow(prod,n/P,P);
6
       cnt=n/p;
       pair<LL,LL>res=comput(n/p,p,k);
7
8
       cnt+=res.first;
9
       ans=ans*res.second%P;
10
       for(int i=n-n%P+1;i<=n;i++)if(i%p){</pre>
11
12
            ans=ans*i%P;
13
14
       return make_pair(cnt,ans);
15
   pair<LL,LL> calc(LL n,LL p,LL k){
16
17
       prod=1;P=pow(p,k,1e18);
       for(int i=1;i<P;i++)if(i%p)prod=prod*i%P;</pre>
18
19
       pair<LL,LL> res=comput(n,p,k);
  // res.second=res.second*pow(p,res.first%k,P)%P;
20
  // res.first-=res.first%k:
22
       return res;
23 | }
  LL calc(LL n,LL m,LL p,LL k){
24
       pair<LL,LL>A,B,C;
25
26
       LL P=pow(p,k,1e18);
27
       A=calc(n,p,k);
28
       B=calc(m,p,k);
29
       C=calc(n-m,p,k);
       LL ans=1:
30
       ans=pow(p,A.first-B.first-C.first,P);
31
```

1.12 Schreier-Sims

```
struct Perm{
 2
     vector<int> P; Perm() {} Perm(int n) { P.resize(n); }
 3
     Perm inv()const{
       Perm ret(P.size()):
4
 5
       for(int i = 0; i < int(P.size()); ++i) ret.P[P[i]] =
 6
7
     int &operator [](const int &dn){ return P[dn]; }
8
     void resize(const size_t &sz){ P.resize(sz); }
9
10
     size_t size()const{ return P.size(); }
11
     const int &operator [](const int &dn)const{ return
        \hookrightarrow P[dn]; }
12
   };
   Perm operator *(const Perm &a, const Perm &b){
13
     Perm ret(a.size());
14
     for(int i = 0; i < (int)a.size(); ++i) ret[i] = b[a[i]];</pre>
15
16
     return ret;
17
   }
18
   typedef vector<Perm> Bucket;
19
   typedef vector<int> Table;
20
   typedef pair<int,int> PII;
21
   int n, m;
   vector<Bucket> buckets, bucketsInv; vector<Table>
23
   int fastFilter(const Perm &g, bool addToGroup = true) {
     int n = buckets.size();
24
25
     Perm p(g);
26
     for(int i = 0; i < n; ++i){
27
       int res = lookupTable[i][p[i]];
28
       if(res == -1){
29
         if (addToGroup) {
30
           buckets[i].push_back(p);

→ bucketsInv[i].push_back(p.inv());
           lookupTable[i][p[i]] = (int)buckets[i].size() - 1;
31
32
33
         return i;
34
35
         = p * bucketsInv[i][res];
36
37
     return -1:
38
   }
   long long calcTotalSize(){
40
     long long ret = 1;
41
     for(int i = 0; i < n; ++i) ret *= buckets[i].size();</pre>
42
     return ret:
43
   bool inGroup(const Perm &g){ return fastFilter(g, false)
   void solve(const Bucket &gen,int _n){// m perm[0..n - 1]s
46
     n = _n, m = gen.size();
     {//clear all
47
48
       vector<Bucket> _buckets(n); swap(buckets, _buckets);
       vector<Bucket> _bucketsInv(n); swap(bucketsInv,

→ bucketsInv):
       vector<Table> _lookupTable(n); swap(lookupTable,
50
          \hookrightarrow _lookupTable);
     }
51
52
     for(int i = 0; i < n; ++i){
53
       lookupTable[i].resize(n);
       fill(lookupTable[i].begin(), lookupTable[i].end(),
           \hookrightarrow -1);
55
56
     Perm id(n):
     for(int i = 0; i < n; ++i) id[i] = i;</pre>
```

```
for(int i = 0; i < n; ++i){
58
       buckets[i].push_back(id); bucketsInv[i].push_back(id);
59
60
       lookupTable[i][i] = 0;
61
     for(int i = 0; i < m; ++i) fastFilter(gen[i]);</pre>
62
     queue<pair<PII,PII> > toUpdate;
63
     for(int i = 0; i < n; ++i)
       for(int j = i; j < n; ++j)
65
66
         for(int k = 0; k < (int)buckets[i].size(); ++k)</pre>
67
            for(int 1 = 0; 1 < (int)buckets[j].size(); ++1)</pre>
68
              toUpdate.push(make_pair(PII(i,k), PII(j,l)));
69
     while(!toUpdate.emptv()){
70
       PII a = toUpdate.front().first, b =
          71
       toUpdate.pop();
       int res = fastFilter(buckets[a.first][a.second] *
72

    buckets[b.first][b.second]):
73
       if(res==-1) continue;
74
       PII newPair(res, (int)buckets[res].size() - 1);
       for(int i = 0; i < n; ++i)
75
76
         for(int j = 0; j < (int)buckets[i].size(); ++j){</pre>
77
           if(i <= res) toUpdate.push(make_pair(PII(i, j),</pre>
              → newPair)):
78
           if(res <= i) toUpdate.push(make_pair(newPair,</pre>
              \hookrightarrow PII(i, j));
79
80
     }
81 }
```

1.13 分治 FFT

```
1 struct complex
2 | {
3
     double x , yi;
 4
5
     complex(double x = 0, double yi = 0): x(x), yi(yi) {}
 6
 7
     friend complex operator + (const complex a, const complex
        \hookrightarrow b)
8
     {
9
        return complex(a.x + b.x, a.yi + b.yi);
10
11
     friend complex operator - (const complex a, const complex
12
     {
13
       return complex(a.x - b.x, a.yi - b.yi);
14
15
     friend complex operator * (const complex a, const complex
16
17
       return complex(a.x * b.x - a.yi * b.yi , a.x * b.yi +
           \hookrightarrow a.yi * b.x);
18
     }
19
     friend complex operator / (const complex a, const double
20
21
        return complex(a.x / b, a.yi / b);
22
23
   ∣}:
24
   void FFT(complex *X,int n,int flag)
25
26
     for(int i = 0; i < n; i++)</pre>
27
       int p = 0, t = i;
       for(int j = 1; j < n; j <<= 1)
29
30
          p <<= 1, p |= (t & 1), t >>= 1;
31
        if(i < p) std::swap(X[i], X[p]);</pre>
32
     }
33
34
     for(int m = 2; m <= n; m <<= 1)
35
36
        complex wm = complex(cos((double) 2 * pi * flag / m),
           \hookrightarrow \sin((double)2 * pi * flag / m));
```

```
38
        for(int i = 0; i < n; i += m)
39
            complex wk = complex(1, 0);
40
41
42
          for(int j = 0; j < (m >> 1); wk = wk * wm, j++)
43
            complex u = X[i + j], t = wk * X[i + j + (m >>
44
               \hookrightarrow 1)];
45
46
            X[i + j] = u + t, X[i + j + (m >> 1)] = u - t;
47
48
        }
49
     }
50
      if(flag == -1) for(int i = 0; i < n; i++) X[i] = X[i] /
51

¬ n:

52
   void solve(int l,int r)
54
55
56
     if(1 == r) return;
57
     static complex A[maxn], B[maxn];
58
59
     int mid = (1 + r) >> 1;
60
     int len = 1;
61
     solve(1, mid);
62
63
     while(len < (r - 1 + 1)) len <<= 1;
64
     len <<= 1;
66
67
      for(int i = 0; i < len; i++) A[i] = B[i] = complex(0,
68
     for(int i = 1; i <= r - 1; i++) A[i] = complex(a[i], 0);</pre>
69
70
     for(int i = 1; i <= mid; i++) B[i - 1] = complex(f[i],</pre>
        \hookrightarrow 0);
71
     FFT(A, len, 1);
     FFT(B, len, 1);
72
     for(int i = 0; i < len; i++) A[i] = A[i] * B[i];</pre>
73
74
     FFT(A, len, -1);
     for(int i = mid + 1; i <= r; i++) (f[i] += round(A[i -</pre>
         \hookrightarrow 1].x)) %= mod;
76
77
     solve(mid + 1, r);
78
```

2. 代数

2.1 快速傅里叶变换

```
// n 必须是 2 的次幂
   void fft(Complex a[], int n, int f) {
     for (int i = 0; i < n; ++i)
3
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
7
       Complex w = Complex(1, 0);
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
       for (int p = i \ll 1, j = 0; j \ll n; j \neq p) {
10
         for (int k = 0; k < i; ++k) {
11
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
15
    }
16
  }
```

2.2 分治卷积

```
1 // n 必须是 2 的次幂
  void fft(Complex a[], int n, int f) {
     for (int i = 0; i < n; ++i)
3
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
7
       Complex w = Complex(1, 0);
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
       for (int p = i \iff 1, j = 0; j \iff n; j += p) {
10
         for (int k = 0; k < i; ++k) {
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
11
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
       }
14
15
     }
16 }
```

2.3 快速数论变换

```
1 // n 必须是 2 的次幂
   void fft(Complex a[], int n, int f) {
    for (int i = 0; i < n; ++i)
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
    for (int i = 1, h = 0; i < n; i <<= 1, h++) {
5
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
6
       Complex w = Complex(1, 0);
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
       for (int p = i \iff 1, j = 0; j \iff n; j += p) {
10
         for (int k = 0; k < i; ++k) {
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
11
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
15
     }
  }
16
```

2.4 光速数论变换

```
#define meminit(A, 1, r) memset(A + (1), 0, sizeof(*A) *
      \hookrightarrow ((r) - (1)))
   #define memcopy(B, A, 1, r) memcpy(B, A + (1), sizeof(*A)
      \hookrightarrow * ((r) - (1)))
   void DFT(int *a, int n, int f) { //f=1 逆 DFT
     for (register int i = 0, j = 0; i < n; i++) {
       if (i > j) std::swap(a[i], a[j]);
5
       for (register int t = n \gg 1; (j ^= t) < t; t \gg = 1);
6
    }
7
     for (register int i = 2; i <= n; i <<= 1) {
9
       static int exp[MAXN];
10
       exp[0] = 1; exp[1] = fpm(PRT, (MOD - 1) / i, MOD);
       if (f == 1) exp[1] = fpm(exp[1], MOD - 2, MOD);
11
       for (register int k = 2; k < (i >> 1); k++) {
12
13
         exp[k] = 111 * exp[k - 1] * exp[1] % MOD;
14
       for (register int j = 0; j < n; j += i) {
15
16
         for (register int k = 0; k < (i >> 1); k++) {
           register int &pA = a[j + k], &pB = a[j + k + (i >>
17
18
           register long long B = 111 * pB * exp[k];
           pB = (pA - B) \% MOD;
19
           pA = (pA + B) \% MOD;
20
21
       }
22
     }
23
24
     if (f == 1) {
       register int rev = fpm(n, MOD - 2, MOD);
25
       for (register int i = 0; i < n; i++) {
26
         a[i] = 111 * a[i] * rev % MOD;
27
         if (a[i] < 0) { a[i] += MOD; }</pre>
29
```

```
31
  // 在不写高精度的情况下合并 FFT 所得结果对 MOD
     →取模后的答案
  // 值得注意的是,这个东西不能最后再合并,而是应该
     →每做一次多项式乘法就 CRT 一次
   int CRT(int *a) {
    static int x[3];
    for (int i = 0; i < 3; i++) {
36
      x[i] = a[i];
37
38
      for (int j = 0; j < i; j++) {
        int t = (x[i] - x[j] + FFT[i] -> MOD) % FFT[i] ->
40
        if (t < 0) t += FFT[i] -> MOD;
        x[i] = 1LL * t * inv[j][i] % FFT[i] -> MOD;
41
42
43
    }
44
    int sum = 1, ret = x[0] % MOD;
    for (int i = 1; i < 3; i ++) {
46
      sum = 1LL * sum * FFT[i - 1] -> MOD % MOD;
47
      ret += 1LL * x[i] * sum % MOD;
48
      if(ret >= MOD) ret -= MOD;
49
    }
50
    return ret:
51
  for (int i = 0; i < 3; i++) // inv 数组的预处理过程,
     \rightarrow inverse(x, p) 表示求 x 在 p 下逆元
53
    for (int j = 0; j < 3; j++)
      inv[i][j] = inverse(FFT[i] -> MOD, FFT[j] -> MOD);
```

2.5 多项式除法

```
void divide(int n, int m, int *a, int *b, int *d, int *r)
      \rightarrow { // n、m 分别为多项式 A (被除数) 和 B (除数)
      →的指数 + 1
     static int M, tA[MAXN], tB[MAXN], inv[MAXN], tD[MAXN];
 2
     for (; n > 0 && a[n - 1] == 0; n--);
 3
4
     for (; m > 0 \&\& b[m - 1] == 0; m--);
     for (int i = 0; i < n; i++) tA[i] = a[n - i - 1];</pre>
5
     for (int i = 0; i < m; i++) tB[i] = b[m - i - 1];
     for (M = 1; M <= n - m + 1; M <<= 1);</pre>
7
8
     if (m < M) meminit(tB, m, M);</pre>
9
     getInv(tB, inv, M):
10
     for (M = 1; M \le 2 * (n - m + 1); M \le 1);
     meminit(inv, n - m + 1, M);
11
     meminit(tA, n - m + 1, M);
12
13
     DFT(inv, M, 0);
     DFT(tA, M, 0);
14
     for (int i = 0; i < M; i++) {</pre>
15
16
       d[i] = 111 * inv[i] * tA[i] % MOD;
17
     DFT(d, M, 1);
18
19
     std::reverse(d, d + n - m + 1);
     for (M = 1; M <= n; M <<= 1);</pre>
20
     memcopy(tB, b, 0, m);
21
     if (m < M) meminit(tB, m, M);</pre>
22
23
     memcopy(tD, d, 0, n - m + 1);
24
     meminit(tD, n - m + 1, M);
25
     DFT(tD, M, 0);
26
     DFT(tB, M, 0);
27
     for (int i = 0; i < M; i++) {
28
       r[i] = 111 * tD[i] * tB[i] % MOD;
29
30
     DFT(r, M, 1);
31
     meminit(r, n, M);
     for (int i = 0; i < n; i++) {
32
33
       r[i] = (a[i] - r[i] + MOD) % MOD;
34
     }
35
   }
```

2.6 多项式求逆

```
void getInv(int *a, int *b, int n) {
2
     static int tmp[MAXN];
     b[0] = fpm(a[0], MOD - 2, MOD);
3
     for (int c = 2, M = 1; c < (n << 1); c <<= 1) {
       for (; M \le 3 * (c - 1); M \le 1);
5
       meminit(b, c, M);
       meminit(tmp, c, M);
7
8
       memcopy(tmp, a, 0, c);
g
       DFT(tmp, M, 0);
10
       DFT(b, M, 0);
       for (int i = 0; i < M; i++) {</pre>
11
12
         b[i] = 111 * b[i] * (211 - 111 * tmp[i] * b[i] % MOD
             \hookrightarrow + MOD) % MOD;
13
14
       DFT(b, M, 1);
       meminit(b, c, M);
15
16
17
  ۱,
```

2.7 多项式取对数

```
// n 必须是 2 的次幂
   void fft(Complex a[], int n, int f) {
     for (int i = 0; i < n; ++i)
3
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
5
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
       Complex w = Complex(1, 0);
7
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
9
10
        for (int k = 0; k < i; ++k) {
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
11
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
       }
     }
15
16 }
```

2.8 快速沃尔什变换

```
void FWT(LL a[],int n,int ty){
 2
     for(int d=1;d<n;d<<=1){</pre>
 3
        for(int m=(d<<1),i=0;i<n;i+=m){</pre>
 4
          if(ty==1){
            for(int j=0; j<d; j++){</pre>
 5
              LL x=a[i+j], y=a[i+j+d];
 6
               a[i+j]=x+y;
 8
               a[i+j+d]=x-y;
q
                        //xor:a[i+j]=x+y,a[i+j+d]=x-y;
10
                        //and:a[i+j]=x+y;
11
                        //or:a[i+j+d]=x+y;
12
            }
13
            for(int j=0; j< d; j++){
14
15
               LL x=a[i+j], y=a[i+j+d];
               a[i+j]=(x+y)/2;
16
17
               a[i+j+d]=(x-y)/2;
18
                        //xor:a[i+j]=(x+y)/2,a[i+j+d]=(x-y)/2;
19
                        //and:a[i+j]=x-y;
20
                        //or:a[i+j+d]=y-x;
21
22
23
        }
24
25
   }
26
     FWT(a, 1 << n, 1);
     FWT(b.1<<n.1):
27
     for(int i=0;i<(1<<n);i++)</pre>
28
```

```
29 c[i]=a[i]*b[i];
30 FWT(c,1<<n,-1);
```

2.9 自适应辛普森积分

```
namespace adaptive_simpson {
 2
     template<typename function>
 3
     inline double area(function f, const double &left, const
        double mid = (left + right) / 2;
 4
        return (right - left) * (f(left) + 4 * f(mid) +
 5
           \hookrightarrow f(right)) / 6;
6
 7
 8
     template<typename function>
 9
     inline double simpson(function f, const double &left,
        \hookrightarrow \texttt{const} double &right, const double &eps, const

    double &area_sum) {
10
        double mid = (left + right) / 2;
11
        double area_left = area(f, left, mid);
12
        double area_right = area(f, mid, right);
        double area_total = area_left + area_right;
13
        if (fabs(area_total - area_sum) <= 15 * eps) {</pre>
14
          return area_total + (area_total - area_sum) / 15;
15
       }
16
17
       return simpson(f, left, right, eps / 2, area_left) +
           \hookrightarrow simpson(f, mid, right, eps / 2, area_right);
18
19
20
     template<typename function>
21
     inline double simpson(function f, const double &left,
        \hookrightarrow const double &right, const double &eps) {
        return simpson(f, left, right, eps, area(f, left,
22

    right));
23
24
   }
```

2.10 单纯形

```
const double eps = 1e-8;
   // max{c * x | Ax <= b, x >= 0} 的解, 无解返回空的
      → vector, 否则就是解.
   vector<double> simplex(vector<vector<double> > &A,
 3
      \hookrightarrow vector<double> b, vector<double> c) {
     int n = A.size(), m = A[0].size() + 1, r = n, s = m - 1;
 5
     vector<vector<double> > D(n + 2, vector<double>(m + 1));
     vector<int> ix(n + m);
6
     for(int i = 0; i < n + m; i++) {
7
8
       ix[i] = i;
9
10
     for(int i = 0; i < n; i++) {
11
       for(int j = 0; j < m - 1; j++) {
12
         D[i][j] = -A[i][j];
13
14
       D[i][m - 1] = 1;
15
       D[i][m] = b[i]:
       if (D[r][m] > D[i][m]) {
16
17
         r = i;
18
       }
19
20
     for(int j = 0; j < m - 1; j++) {
21
22
       D[n][j] = c[j];
23
     D[n + 1][m - 1] = -1;
24
25
     for(double d; ;) {
26
       if (r < n) {</pre>
27
         swap(ix[s], ix[r + m]);
28
         D[r][s] = 1. / D[r][s];
         for(int j = 0; j \le m; j++) {
29
           if (j != s) {
30
             D[r][j] *= -D[r][s];
31
```

```
32
            }
33
          }
34
          for(int i = 0; i <= n + 1; i++) {
35
            if (i != r) {
              for(int j = 0; j \le m; j++) {
36
                 if (j != s) {
37
                   D[i][j] += D[r][j] * D[i][s];
38
39
40
              }
41
              D[i][s] *= D[r][s];
42
43
          }
44
45
        r = -1, s = -1;
46
        for(int j = 0; j < m; j++) {
          if (s < 0 || ix[s] > ix[j]) {
47
            if (D[n + 1][j] > eps || D[n + 1][j] > -eps &&
48
               \hookrightarrow D[n][j] > eps) {
49
               s = i;
            }
50
          }
51
52
        }
        if (s < 0) {
53
          break;
54
55
56
        for(int i = 0; i < n; i++) {</pre>
57
          if (D[i][s] < -eps) {</pre>
            if (r < 0 \mid \mid (d = D[r][m] / D[r][s] - D[i][m] /
58
               \hookrightarrow D[i][s]) < -eps
               || d < eps && ix[r + m] > ix[i + m]) {
59
60
              r = i;
61
62
            }
63
          }
64
65
66
        if (r < 0) {
          return vector<double> ();
67
68
69
     }
70
     if (D[n + 1][m] < -eps) {
71
        return vector<double> ();
72
73
74
     vector<double> x(m - 1);
75
     for(int i = m; i < n + m; i++) {</pre>
       if (ix[i] < m - 1) {</pre>
76
77
          x[ix[i]] = D[i - m][m];
78
79
     }
80
     return x;
81 }
```

3. 计算几何

3.1 二维

3.1.1 点类

```
1 int sign(DB x) {
2
    return (x > eps) - (x < -eps);
3 }
4 DB msqrt(DB x) {
5
    return sign(x) > 0 ? sqrt(x) : 0;
6
  ۱,
8
   struct Point {
    DB x, y;
9
    Point rotate(DB ang) const { // 逆时针旋转 ang 弧度
10
      return Point(cos(ang) * x - sin(ang) * y,
11
          cos(ang) * y + sin(ang) * x);
12
13
    Point turn90() const { // 逆时针旋转 90 度
14
```

```
15
              return Point(-y, x);
16
17
         Point unit() const {
18
             return *this / len();
19
20
     }:
     DB dot(const Point& a, const Point& b) {
         return a.x * b.x + a.y * b.y;
23
     }
24
     DB det(const Point& a, const Point& b) {
25
       return a.x * b.y - a.y * b.x;
26
27
      #define cross(p1,p2,p3)
           \hookrightarrow ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
28
      #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
     bool isLL(const Line& 11, const Line& 12, Point& p) { //
           →直线与直线交点
          DB s1 = det(12.b - 12.a, 11.a - 12.a),
30
               s2 = -det(12.b - 12.a, 11.b - 12.a);
31
          if (!sign(s1 + s2)) return false;
32
33
         p = (11.a * s2 + 11.b * s1) / (s1 + s2);
34
         return true:
     }
35
     bool onSeg(const Line& 1, const Point& p) { // 点在线段
36
          return sign(det(p - 1.a, 1.b - 1.a)) == 0 && sign(dot(p
37
               \hookrightarrow - l.a, p - l.b)) <= 0;
38
     Point projection(const Line & 1, const Point& p) {
39
         return 1.a + (1.b - 1.a) * (dot(p - 1.a, 1.b - 1.a) /
               \hookrightarrow (1.b - 1.a).len2());
     }
41
     DB disToLine(const Line& 1, const Point& p) { // 点到 *
42
           →直线 * 距离
         return fabs(det(p - 1.a, 1.b - 1.a) / (1.b -
43
               \hookrightarrow 1.a).len());
44
     DB disToSeg(const Line& 1, const Point& p) { // 点到线段
45
           →距离
         return sign(dot(p - 1.a, 1.b - 1.a)) * sign(dot(p - 1.b, a)) * sign(dot(p - 
               \rightarrow l.a - l.b)) == 1 ? disToLine(l, p) : std::min((p -
               \hookrightarrow l.a).len(), (p - l.b).len());
47
      // 圆与直线交点
48
      bool isCL(Circle a, Line 1, Point& p1, Point& p2) {
49
         DB x = dot(1.a - a.o, 1.b - 1.a),
               y = (1.b - 1.a).len2(),
52
                d = x * x - y * ((1.a - a.o).len2() - a.r * a.r);
          if (sign(d) < 0) return false;</pre>
53
         Point p = 1.a - ((1.b - 1.a) * (x / y)), delta = (1.b - 1.a)
54
               \hookrightarrow 1.a) * (msqrt(d) / y);
          p1 = p + delta; p2 = p - delta;
56
         return true;
57
     }
      //圆与圆的交面积
58
     DB areaCC(const Circle& c1, const Circle& c2) {
59
         DB d = (c1.o - c2.o).len();
60
          if (sign(d - (c1.r + c2.r)) >= 0) return 0;
61
62
          if (sign(d - std::abs(c1.r - c2.r)) <= 0) {</pre>
63
              DB r = std::min(c1.r, c2.r);
64
              return r * r * PI;
65
         DB x = (d * d + c1.r * c1.r - c2.r * c2.r) / (2 * d),
66
             t1 = acos(x / c1.r), t2 = acos((d - x) / c2.r);
67
68
          return c1.r * c1.r * t1 + c2.r * c2.r * t2 - d * c1.r *
               \hookrightarrow \sin(t1):
69
     }
70 // 圆与圆交点
71
     bool isCC(Circle a, Circle b, P& p1, P& p2) {
         DB s1 = (a.o - b.o).len();
72
73
         if (sign(s1 - a.r - b.r) > 0 || sign(s1 - std::abs(a.r -
               \hookrightarrow b.r)) < 0) return false;
```

```
DB s2 = (a.r * a.r - b.r * b.r) / s1;
      DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
75
76
     P \circ = (b.o - a.o) * (aa / (aa + bb)) + a.o;
     P delta = (b.o - a.o).unit().turn90() * msqrt(a.r * a.r
77
        \hookrightarrow - aa * aa):
     p1 = o + delta, p2 = o - delta;
78
     return true;
80 }
81 // 求点到圆的切点,按关于点的顺时针方向返回两个点
   | bool tanCP(const Circle &c, const Point &p0, Point &p1,
      → Point &p2) {
     double x = (p0 - c.o).len2(), d = x - c.r * c.r;
83
     if (d < eps) return false; // 点在圆上认为没有切点
84
     Point p = (p0 - c.o) * (c.r * c.r / x);
85
     Point delta = ((p0 - c.o) * (-c.r * sqrt(d) /
86
        \hookrightarrow x)).turn90();
      p1 = c.o + p + delta;
87
88
      p2 = c.o + p - delta;
89
      return true;
90 }
   // 求圆到圆的外共切线,按关于 c1.o 的顺时针方向返
91
    vector<Line> extanCC(const Circle &c1, const Circle &c2) {
93
     vector<Line> ret;
     if (sign(c1.r - c2.r) == 0) {
94
       Point dir = c2.o - c1.o;
95
       dir = (dir * (c1.r / dir.len())).turn90();
96
       ret.push_back(Line(c1.o + dir, c2.o + dir));
       ret.push_back(Line(c1.o - dir, c2.o - dir));
98
99
     } else {
100
       Point p = (c1.0 * -c2.r + c2.o * c1.r) / (c1.r - c2.r + c2.o * c1.r) / (c1.r - c2.r + c2.o * c1.r)
          \hookrightarrow c2.r):
101
       Point p1, p2, q1, q2;
102
       if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) {
103
          if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);</pre>
104
         ret.push_back(Line(p1, q1));
105
          ret.push_back(Line(p2, q2));
106
107
     }
108
     return ret;
109 }
   // 求圆到圆的内共切线,按关于 c1.o 的顺时针方向返
      →回两条线
111 std::vector<Line> intanCC(const Circle &c1, const Circle
      std::vector<Line> ret;
112
     Point p = (c1.0 * c2.r + c2.o * c1.r) / (c1.r + c2.r);
114
     Point p1, p2, q1, q2;
     if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { //
115
        →两圆相切认为没有切线
116
       ret.push_back(Line(p1, q1));
117
       ret.push_back(Line(p2, q2));
     }
118
119
     return ret:
120
   |bool contain(vector<Point> polygon, Point p) { // 判断点
121
      →p 是否被多边形包含,包括落在边界上
      int ret = 0, n = polygon.size();
122
     for(int i = 0; i < n; ++ i) {
123
124
       Point u = polygon[i], v = polygon[(i + 1) % n];
        if (onSeg(Line(u, v), p)) return true; // Here I
        if (sign(u.y - v.y) \le 0) swap(u, v);
126
        if (sign(p.y - u.y) > 0 \mid \mid sign(p.y - v.y) \le 0)
127
          128
       ret += sign(det(p, v, u)) > 0;
129
130
      return ret & 1;
131 }
132 // 用半平面 (q1,q2) 的逆时针方向去切凸多边形
133 | std::vector<Point> convexCut(const std::vector<Point>&ps,
      \hookrightarrow Point q1, Point q2) {
```

```
std::vector<Point> qs; int n = ps.size();
134
135
      for (int i = 0; i < n; ++i) {
136
        Point p1 = ps[i], p2 = ps[(i + 1) % n];
        int d1 = crossOp(q1,q2,p1), d2 = crossOp(q1,q2,p2);
137
        if (d1 >= 0) qs.push_back(p1);
138
        if (d1 * d2 < 0) qs.push_back(isSS(p1, p2, q1, q2));</pre>
139
      }
140
141
      return qs;
142 }
   // 求凸包
    std::vector<Point> convexHull(std::vector<Point> ps) {
145
      int n = ps.size(); if (n <= 1) return ps;</pre>
146
      std::sort(ps.begin(), ps.end());
147
      std::vector<Point> qs;
      for (int i = 0; i < n; qs.push_back(ps[i ++]))</pre>
        while (qs.size() > 1 && sign(det(qs[qs.size() - 2],
149
           \hookrightarrow qs.back(), ps[i])) <= 0)
150
          qs.pop_back();
      for (int i = n - 2, t = qs.size(); i \ge 0;
151
         while ((int)qs.size() > t && sign(det(qs[qs.size() -
           \hookrightarrow 2], qs.back(), ps[i])) <= 0)
153
          qs.pop_back();
      return qs;
```

3.1.2 凸包

```
// 凸包中的点按逆时针方向
 2
   struct Convex {
 3
     int n:
4
     std::vector<Point> a, upper, lower;
     void make_shell(const std::vector<Point>& p,
         std::vector<Point>& shell) { // p needs to be
 6
            \hookrightarrow sorted.
 7
       clear(shell); int n = p.size();
8
       for (int i = 0, j = 0; i < n; i++, j++) {
         for (; j \ge 2 \&\& sign(det(shell[j-1] - shell[j-2],
 9
10
                 p[i] - shell[j-2])) \le 0; --j)
                    11
         shell.push_back(p[i]);
       }
12
13
     }
14
     void make_convex() {
       std::sort(a.begin(), a.end());
16
       make_shell(a, lower);
17
       std::reverse(a.begin(), a.end());
18
       make_shell(a, upper);
19
       a = lower; a.pop_back();
20
       a.insert(a.end(), upper.begin(), upper.end());
       if ((int)a.size() >= 2) a.pop_back();
21
22
       n = a.size();
23
     }
24
     void init(const std::vector<Point>& _a) {
       clear(a); a = _a; n = a.size();
25
26
       make_convex();
27
28
     void read(int _n) { // Won't make convex.
29
       clear(a); n = _n; a.resize(n);
       for (int i = 0; i < n; i++)
30
31
         a[i].read();
32
     std::pair<DB, int> get_tangent(
33
         const std::vector<Point>& convex, const Point& vec)
           ← {
       int l = 0, r = (int)convex.size() - 2;
35
       assert(r >= 0):
36
37
       for (; 1 + 1 < r; ) {
         int mid = (1 + r) / 2;
38
39
         if (sign(det(convex[mid + 1] - convex[mid], vec)) >
           r = mid:
40
         else 1 = mid;
41
```

```
42
      return std::max(std::make_pair(det(vec, convex[r]),
43
          std::make_pair(det(vec, convex[0]), 0));
44
     }
45
     int binary_search(Point u, Point v, int 1, int r) {
46
      int s1 = sign(det(v - u, a[1 % n] - u));
47
      for (; 1 + 1 < r; ) {
48
49
        int mid = (1 + r) / 2;
50
        int smid = sign(det(v - u, a[mid % n] - u));
        if (smid == s1) 1 = mid;
51
52
        else r = mid:
53
54
      return 1 % n;
     }
55
     // 求凸包上和向量 vec 叉积最大的点, 返回编号, 共
56
       → 线的多个切点返回任意一个
57
     int get_tangent(Point vec) {
58
      std::pair<DB, int> ret = get_tangent(upper, vec);
59
      ret.second = (ret.second + (int)lower.size() - 1) % n;
60
      ret = std::max(ret, get_tangent(lower, vec));
61
      return ret.second;
    }
62
     // 求凸包和直线 u, v 的交点, 如果不相交返回 false,
63
       →如果有则是和 (i, next(i)) 的交点, 交在点上不
       →确定返回前后两条边其中之一
     bool get_intersection(Point u, Point v, int &i0, int
64
       int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
66
      if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p0] - u))
         \hookrightarrow a[p1] - u)) \le 0) 
        if (p0 > p1) std::swap(p0, p1);
67
68
        i0 = binary_search(u, v, p0, p1);
        i1 = binary_search(u, v, p1, p0 + n);
69
70
        return true:
71
72
      else return false;
73
    }
74 \ \ \ \ ;
```

3.1.3 凸包最近点对

```
#include<cstdio>
   #include<cmath>
3 | #include<cstring>
4 #include<iostream>
5 #include<algorithm>
6 | #include < cstdlib>
  #include<queue>
7
8
  #include<map>
9
   #include<stack>
10 #include<set>
   #define e exp(1.0); //2.718281828
11
   #define mod 100000007
12
   #define INF 0x7fffffff
13
   #define inf 0x3f3f3f3f
14
   typedef long long LL;
15
16
   using namespace std;
17
   #define zero(x) (((x)>0?(x):(-x))<eps)
19 const double eps=1e-8;
20
21 //判断数 k 的符号 -1 负数 1 正数 0 零
22 | int dcmp(double k) {
23
       return k<-eps?-1:k>eps?1:0;
24 }
25
   inline double sqr(double x) {
26
27
       return x*x;
28 }
   struct point {
29
30
       double x,y;
```

```
31
       point() {};
32
       point(double a,double b):x(a),y(b) {};
33
       void input() {
34
           scanf("%lf %lf",&x,&y);
35
36
       friend point operator + (const point &a, const point
           return point(a.x+b.x,a.y+b.y);
37
38
       }
39
       friend point operator - (const point &a, const point
          40
           return point(a.x-b.x,a.y-b.y);
41
42
       friend bool operator == (const point &a,const point
          43
           return dcmp(a.x-b.x)==0\&\&dcmp(a.y-b.y)==0;
44
45
       friend point operator * (const point &a, const double
           return point(a.x*b,a.y*b);
46
47
48
       friend point operator * (const double &a, const point
         49
           return point(a*b.x,a*b.y);
50
51
       friend point operator / (const point &a,const double
52
           return point(a.x/b,a.y/b);
       }
53
54
       friend bool operator < (const point &a, const point
55
           return a.x < b.x || (a.x == b.x && a.y < b.y);
56
       }
57
       double norm() {
58
           return sqrt(sqr(x)+sqr(y));
59
60
   //计算两个向量的叉积
61
62
   double cross(const point &a,const point &b) {
63
       return a.x*b.y-a.y*b.x;
64
65
   double cross3(point A,point B,point C) { //叉乘
66
       return (B.x-A.x)*(C.y-A.y)-(B.y-A.y)*(C.x-A.x);
67
   }
   //计算两个点的点积
68
69
   double dot(const point &a,const point &b) {
70
       return a.x*b.x+a.y*b.y;
71
  }
72
   double dot3(point A, point B, point C) { //点乘
73
       return (C.x-A.x)*(B.x-A.x)+(C.y-A.y)*(B.y-A.y);
74
75
   //向量长度
76
77
   double length(const point &a) {
78
      return sqrt(dot(a,a));
79
   //两个向量的角度
80
81
   double angle(const point &a,const point &b) {
82
       return acos(dot(a,b)/length(a)/length(b));
83
   //计算两个点的距离
84
85
   double dist(const point &a,const point &b) {
86
       return (a-b).norm();
87
  }
   //op 沿远点逆时针旋转角度 A
88
   point rotate_point(const point &p,double A) {
89
       double tx=p.x,ty=p.y;
90
91
       return point(tx*cos(A)-ty*sin(A),tx*sin(A)+ty*cos(A));
92
  }
   double TriArea(const point &a, const point &b, const point
94
       return fabs( cross( b - a, c - a ) ) / 2;
95
```

```
point Normal(const point &a) {
97
        double L = length(a);
                                                                      157
                                                                                  res[m++] = P[i];
98
        return point(-a.y/L, a.x/L);
                                                                      158
                                                                              }
   ۱,
                                                                              int k = m:
99
                                                                      159
    //求两条直线的交点, p 和 q 分别为两条直线上的点, v
                                                                              for (int i = cnt - 2; i >= 0; i--) {
                                                                      160
100
                                                                                  while (m > k \&\& cross(res[m - 1] - res[m - 2],
                                                                      161
       →和 w 分别为直线的方向向量
                                                                                     \hookrightarrow P[i] - res[m - 2]) \le 0
    point GetLineIntersection(point p, point v, point q, point
                                                                      162
       \hookrightarrow w) {
                                                                      163
                                                                                  res[m++] = P[i];
102
        point u = p - q;
                                                                              }
                                                                      164
        double t = cross(w, u) / cross(v, w);
103
                                                                      165
                                                                              if (cnt > 1) m--;
        return p + v * t;
104
                                                                      166
                                                                              return m;
105 }
                                                                      167
    //求点 p 到直线 ab 的距离
106
                                                                      168
    double DistanceToLine(point p, point a, point b) {
                                                                          //判断点是否在多边形内
        point v1 = b - a, v2 = p - a;
108
                                                                      170
                                                                          int isPointInPolygon(point p, point *a, int n) {
109
        return fabs(cross(v1,v2)) / length(v1);
                                                                      171
                                                                              int cnt = 0:
110
   | }
                                                                              for(int i=0; i<n; ++i) {</pre>
                                                                      172
    //求点 p 到线段 ab 的距离
111
                                                                      173
                                                                                  if(OnSegment(p, a[i], a[(i+1)%n])) return -1;
    double DistanceToSegment(point p, point a, point b) {
112
                                                                      174
                                                                                  double k = cross(a[(i+1)%n]-a[i], p-a[i]);
        if(a==b) return length(p - a);
113
                                                                      175
                                                                                  double d1 = a[i].y - p.y;
114
        point v1 = b - a, v2 = p - a, v3 = p - b;
                                                                              double d2 = a[(i+1)].y - p.y;
                                                                      176
        if(dcmp(dot(v1,v2)) < 0) return length(v2);</pre>
115
                                                                                  if(k>0 &&d1<=0 &&d2>0)//点在线段的左侧
                                                                      177
        else if(dcmp(dot(v1,v3)) > 0) return length(v3);
116
                                                                      178
117
        else return fabs(cross(v1,v2)) / length(v1);
                                                                                  if(k<0 &&d2<=0 &&d1>0)//点在线段的右侧
118 }
    //判断直线 a1a2 和直线 b1b2 是否规范相交
                                                                      180
                                                                                      cnt++;
119
                                                                                  //k==0, 点和线段共线的情况不考虑
    bool SegmentProperIntersection(point a1, point a2, point
                                                                      181
120
                                                                      182
      \hookrightarrow b1, point b2) {
                                                                      183
                                                                              if(cnt&1)return 1;
121
        double c1 = cross(a2-a1,b1-a1), c2 = cross(a2-a1,b1-a1)
                                                                              return 0;
                                                                      184
           \hookrightarrow b2-a1):
                                                                      185
122
        double c3 = cross(b2-b1, a1-b1), c4 = cross(b2-b1,
                                                                          //判断凸包是否相离
                                                                      186
           \rightarrow a2-b1):
                                                                          bool two_getaway_ConvexHull(point *cha, int n1, point
        return dcmp(c1) * dcmp(c2) <0 && dcmp(c3) * dcmp(c4) <
                                                                      187
123
                                                                             \hookrightarrow *chb, int m1) {
124
    }
                                                                      188
                                                                              if(n1==1 && m1==1) {
                                                                                  if(cha[0] == chb[0])
125
                                                                      189
    //判断点 p 是否在直线 a1a2 上
                                                                      190
                                                                                      return false;
126
                                                                              } else if(n1==1 && m1==2) {
                                                                      191
    bool OnSegment(point p, point a1, point a2) {
                                                                      192
                                                                                  if(OnSegment(cha[0], chb[0], chb[1]))
128
        return dcmp(cross(a1-p,a2-p)) ==0 &&
                                                                      193
                                                                                      return false;
           \hookrightarrow dcmp(dot(a1-p,a2-p))<0;
                                                                      194
                                                                              } else if(n1==2 && m1==1) {
129 }
    //判断线段 a1a2 和线段 b1b2 是否相交, 可以在端点处
                                                                      195
                                                                                  if(OnSegment(chb[0], cha[0], cha[1]))
130
                                                                      196
                                                                                       return false;
       →相交
                                                                      197
                                                                              } else if(n1==2 && m1==2) {
131
    bool SegmentIntersection(point a1, point a2, point b1,
                                                                      198
                                                                                  if(SegmentIntersection(cha[0], cha[1], chb[0],
       \hookrightarrow point b2) {
        return SegmentProperIntersection(a1, a2, b1, b2) ||
132
                                                                      199
                                                                                      return false;
           } else if(n1==2) {
                                                                      200
133 }
                                                                      201
                                                                                  for(int i=0; i<n1; ++i)</pre>
                                                                      202
                                                                                       if(isPointInPolygon(cha[i], chb, m1))
    double SegmentToSegment(point a1, point a2, point b1,
                                                                      203
                                                                                          return false;
       \hookrightarrow point b2) {
                                                                              } else if(m1==2) {
        //线段间的最短距离分为四种情况
136
                                                                                  for(int i=0; i<m1; ++i)</pre>
                                                                      205
137
        double t1 = DistanceToSegment(b1, a1, a2);
                                                                      206
                                                                                       if(isPointInPolygon(chb[i], cha, n1))
        double t2 = DistanceToSegment(b2, a1, a2);
138
                                                                      207
                                                                                          return false;
        double t3 = DistanceToSegment(a1, b1, b2);
139
                                                                      208
                                                                              } else {
140
        double t4 = DistanceToSegment(a2, b1, b2);
                                                                      209
                                                                                  for(int i=0; i<n1; ++i) {</pre>
        return min(t1,min(t2,min(t3,t4)));
141
                                                                      210
                                                                                      for(int j=0; j<m1; ++j) {</pre>
142
   ۱,
                                                                                           if (SegmentIntersection(cha[i],
                                                                      211
    //使点集逆时针转
143
                                                                                             \hookrightarrow cha[(i+1)%n1], chb[j],
144
    void antiClockSort(point *ch, int n) {
                                                                                             \hookrightarrow chb[(j+1)\%m1]))
        double res = cross(ch[1] - ch[0], ch[2] - ch[0]);
145
                                                                      212
                                                                                               return false;
        if(dcmp(res) >= 0) return;
146
                                                                      213
                                                                                      }
        reverse(ch, ch+n);
147
                                                                      214
148
   ۱,
                                                                      215
                                                                                  for(int i=0; i<n1; ++i)</pre>
149
                                                                      216
                                                                                       if(isPointInPolygon(cha[i], chb, m1))
    int ConvexHull(point* P, int cnt, point* res) {
150
                                                                      217
                                                                                          return false;
        sort(P, P + cnt);
151
                                                                      218
                                                                                  for(int i=0; i<m1; ++i)</pre>
        cnt = (int) (unique(P, P + cnt) - P);
152
                                                                      219
                                                                                       if(isPointInPolygon(chb[i], cha, n1))
153
        int m = 0;
                                                                      220
                                                                                           return false;
154
        for (int i = 0; i < cnt; i++) {</pre>
                                                                      221
            while (m > 1 && cross(res[m - 1] - res[m - 2],
155
                                                                      222
                                                                              return true;
               \hookrightarrow P[i] - res[m - 2]) \le 0
                                                                      223 }
```

```
224 //旋转卡壳求两个凸包最近距离
225
    double solve(point *P, point *Q, int n, int m) {
226
        if(n==1 && m==1) {
             return length(P[0] - Q[0]);
227
        else if(n==1 \&\& m==2) {
228
             return DistanceToSegment(P[0], Q[0], Q[1]);
229
230
        } else if(n==2 && m==1) {
             return DistanceToSegment(Q[0], P[0], P[1]);
231
232
        } else if(n==2 && m==2) {
             return SegmentToSegment(P[0], P[1], Q[0], Q[1]);
233
234
235
236
        int yminP = 0, ymaxQ = 0;
         for(int i=0; i<n; ++i) if(P[i].y < P[yminP].y) yminP =</pre>
237
        for(int i=0; i<m; ++i) if(Q[i].y > Q[ymaxQ].y) ymaxQ =
238
           239
        P[n] = P[0];
240
        Q[n] = Q[0];
        double INF2 = 1e100;
241
242
        double arg, ans = INF2;
243
        for(int i=0; i<n; ++i) {</pre>
244
             //当叉积负正转正时,说明点 ymaxQ 就是对踵点
245
             while((arg=cross(P[yminP] - P[yminP+1],Q[ymaxQ+1]
246
                \hookrightarrow - Q[ymaxQ])) < -eps)
                 ymaxQ = (ymaxQ+1)%m;
248
             double ret;
249
             if(arg > eps) { //卡住第二个凸包上的点。
250
                 ret = DistanceToSegment(Q[ymaxQ], P[yminP],
251
                    ans = min(ans,ret);
252
             } else { //arg==0, 卡住第二个凸包的边
253
                 ret =
254
                    \hookrightarrow \texttt{SegmentToSegment(P[yminP],P[yminP+1],Q[ymaxQ],Q[ymaxQ+1])}, \&\& \texttt{ sign(x) >= 0);} \\
                 ans = min(ans,ret);
255
             }
256
257
             yminP = (yminP+1)%n;
258
259
        return ans;
260
    double mindis_twotubao(point *P, point *Q, int n, int m){
261
        //尼玛, hdu2823 要判是否分离, poj3608 不判
262
         //return min(solve(P, Q, n, m),solve(Q,P,m,n));
263
        //判断凸包是不是相离,如果不是,输出 o
264
265
         if(two_getaway_ConvexHull(P,n,Q,m)==true) return
           \hookrightarrow \texttt{min}(\texttt{solve}(\texttt{P}, \ \texttt{Q}, \ \texttt{n}, \ \texttt{m}), \texttt{solve}(\texttt{Q}, \texttt{P}, \texttt{m}, \texttt{n}))\,;
        else return 0.0;
266
267
   }
268
    const int N=10005;
269
270 point a[N],b[N];
271
    point cha[N],chb[N];
    int main() {
272
        int n.m:
273
274
        while(scanf("%d%d",&n,&m)!=EOF){
275
             for(int i=0;i<n;++i)</pre>
                \hookrightarrow scanf("%lf%lf",&a[i].x,&a[i].y);
276
             for(int i=0;i<m;++i)</pre>
                \hookrightarrow scanf("%lf%lf",&b[i].x,&b[i].y);
             //先求凸包
277
             int n1 = ConvexHull(a, n, cha);
278
             int m1 = ConvexHull(b, m, chb);
279
280
             printf("%.4f\n",mindis_twotubao(cha,chb,n1,m1));
281
282
         return 0;
283
    }
```

3.1.4 三角形的心

```
Point inCenter(const Point &A, const Point &B, const Point
      → &C) { // 内心
     double a = (B - C).len(), b = (C - A).len(), c = (A -
 2
        \hookrightarrow B).len(),
       s = fabs(det(B - A, C - A)),
 4
       r = s / p;
 5
     return (A * a + B * b + C * c) / (a + b + c);
   }
6
 7
   Point circumCenter(const Point &a, const Point &b, const
      → Point &c) { // 外心
     Point bb = b - a, cc = c - a;
 8
     double db = bb.len2(), dc = cc.len2(), d = 2 * det(bb)
9
     return a - Point(bb.y * dc - cc.y * db, cc.x * db - bb.x
        \rightarrow * dc) / d;
11
   }
   Point othroCenter(const Point &a, const Point &b, const
12
      → Point &c) { // 垂心
13
     Point ba = b - a, ca = c - a, bc = b - c;
     double Y = ba.y * ca.y * bc.y,
14
          A = ca.x * ba.y - ba.x * ca.y,
15
16
          x0 = (Y + ca.x * ba.y * b.x - ba.x * ca.y * c.x) /
          y0 = -ba.x * (x0 - c.x) / ba.y + ca.y;
17
18
     return Point(x0, y0);
19
```

3.1.5 半平面交

```
struct Point {
     int quad() const { return sign(y) == 1 || (sign(y) == 0
4
   struct Line {
 5
     bool include(const Point &p) const { return sign(det(b -
        \hookrightarrow a, p - a)) > 0; }
     Line push() const{ // 将半平面向外推 eps
 6
       const double eps = 1e-6;
 7
8
       Point delta = (b - a).turn90().norm() * eps;
9
        return Line(a - delta, b - delta);
10
     }
11
   };
12
   bool sameDir(const Line &10, const Line &11) { return
      \hookrightarrow parallel(10, 11) && sign(dot(10.b - 10.a, 11.b -
      \hookrightarrow 11.a)) == 1; }
13
   bool operator < (const Point &a, const Point &b) {</pre>
     if (a.quad() != b.quad()) {
15
       return a.quad() < b.quad();</pre>
16
     } else {
17
       return sign(det(a, b)) > 0;
18
19
   }
20
   bool operator < (const Line &10, const Line &11) {</pre>
21
     if (sameDir(10, 11)) {
22
       return 11.include(10.a);
23
     } else {
24
        return (10.b - 10.a) < (11.b - 11.a);
25
26
   }
   bool check(const Line &u, const Line &v, const Line &w) {

    return w.include(intersect(u, v)): }

   vector<Point> intersection(vector<Line> &1) {
28
29
     sort(1.begin(), 1.end());
     deque<Line> q;
     for (int i = 0; i < (int)1.size(); ++i) {</pre>
       if (i && sameDir(l[i], l[i - 1])) {
32
33
          continue:
34
35
        while (q.size() > 1 && !check(q[q.size() - 2],
           \hookrightarrow q[q.size() - 1], l[i])) q.pop_back();
```

```
while (q.size() > 1 && !check(q[1], q[0], 1[i]))
36

    q.pop_front();
37
         q.push_back(l[i]);
      7
38
      while (q.size() > 2 && !check(q[q.size() - 2],
39
         \label{eq:q0} \leftarrow \texttt{q[q.size() - 1], q[0])) \ q.pop\_back();}
      while (q.size() > 2 && !check(q[1], q[0], q[q.size() -
40
         \hookrightarrow 1])) q.pop_front();
41
      vector<Point> ret;
42
      for (int i = 0; i < (int)q.size(); ++i)</pre>
         \hookrightarrow ret.push_back(intersect(q[i], q[(i + 1) %
          \hookrightarrow q.size()]));
43
      return ret;
   }
44
```

3.1.6 最大空凸包

```
#include <iostream>
1
   #include <cmath>
2
3
   #include <cstdio>
   #include <algorithm>
   using namespace std;
   typedef double type_p;
   const double eps = 1e-6;
8 const int maxn = 510;
   double dp[maxn][maxn];
10
   inline double eq(double x, double y)
11
   {
12
       return fabs(x-y)<eps;</pre>
13
  | }
14
   inline int eq(int x, int y)
15
   {
16
       return x==y;
17
   }
18
   struct point
19
   {
20
       type_p x,y;
21
  | }:
22
   type_p xmult(point a, point b, point o)
23
   {
24
        return (a.x-o.x)*(o.y-b.y)-(a.y-o.y)*(o.x-b.x);//b at
          \hookrightarrow ao left if negative, at right if positive
   }
25
26
   type_p dist(point a, point b)
27
28
       return (a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y);
  }
29
   point o;
30
   bool cmp_angle(point a,point b)
31
32
   {
33
        if(eq(xmult(a,b,o),0.0))
34
35
            return dist(a,o) < dist(b,o);</pre>
36
37
       return xmult(a.o.b)>0:
38
   }
39
40
   Input: p: Point set
41
            pn: size of the point set
42
43
   Output: the area of the largest empty convex
44
45
   double empty_convex(point *p, int pn)
46
   {
47
       double ans=0;
       for(int i=0; i<pn; i++)</pre>
48
49
50
            for(int j=0; j<pn; j++)
51
52
                dp[i][j]=0;
53
       }
54
55
```

```
56
         for(int i=0; i<pn; i++)</pre>
 57
 58
              int j = i-1;
 59
             while(j \ge 0 && eq(xmult(p[i], p[j],
                \hookrightarrow o),0.0))j--;//coline
 60
             bool flag= j==i-1;
 62
 63
             while(j>=0)
 64
             {
 65
                  int k = j-1;
 66
                  while(k >= 0 && xmult(p[i],p[k],p[j])>0)k--;
 67
                  double area = fabs(xmult(p[i],p[j],o))/2;
 68
                  if(k >= 0)area+=dp[j][k];
 69
                  if(flag) dp[i][j]=area;
 70
                  ans=max(ans, area);
 71
                  j=k;
 72
             }
 73
             if(flag)
 74
             {
 75
                  for(int j=1; j<i; j++)</pre>
 76
                  {
                      dp[i][j] = max(dp[i][j],dp[i][j-1]);
 77
 78
 79
             }
 80
         }
 81
         return ans;
 82
 83
    double largest_empty_convex(point *p, int pn)
 84
 85
         point data[maxn];
         double ans=0;
 86
 87
         for(int i=0; i<pn; i++)</pre>
 88
         {
 89
             o=p[i];
 90
             int dn=0;
 91
             for(int j=0; j<pn; j++)</pre>
 92
 93
                  if(p[j].y>o.y||(p[j].y==o.y&&p[j].x>=o.x))
 94
 95
                      data[dn++]=p[j];
                  }
 96
 98
             sort(data, data+dn, cmp_angle);
 99
             ans=max(ans, empty_convex(data, dn));
         7
100
101
         return ans;
102
103
    int main()
104
105
         point p[110];
106
         int t;
         scanf("%d",&t);
107
         while(t--)
108
109
         {
110
             int pn;
111
             scanf("%d",&pn);
112
             for(int i=0; i<pn; i++)</pre>
113
             {
                  scanf("%lf%lf",&p[i].x,&p[i].y);
114
115
116
             printf("%.1f\n",largest_empty_convex(p,pn));
117
118
         return 0;
119
```

3.1.7 平面最近点对

```
#include <iostream>
#include <cstdio>
#include <cstdlib>
#include <cstring>
```

```
5 | #include <algorithm>
   #include <cmath>
6
8 using namespace std;
9
10 const double eps = 1e-8;
11 const int INF = 0x7ffffffff;
12 int n;
13
   struct Point
14
15 | €
16
     double x.v:
17
     Point(double x=0, double y=0):x(x),y(y) {}
18
     bool operator < (const Point& p) const
19
20
       if(x != p.x) return x < p.x;
21
       else return y < p.y;</pre>
22
23
  }p[200000+5],temp[200000+5];
25 bool cmpy(Point a, Point b)
26 {
27
    return a.y < b.y;</pre>
28
  | }
29
30
   double Dis(Point a, Point b)
31
    return sqrt((a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y));
32
33 }
34
35
  double Closest_Pair(int left, int right)
36 {
37
     double d = INF;
38
    if(left == right)
      return d;
39
40
     if(left +1 == right)
41
      return Dis(p[left],p[right]);
42
     int mid = (left+right)>>1;
43
     double d1 = Closest_Pair(left,mid);
     double d2 = Closest_Pair(mid,right);
44
45
     d = min(d1.d2):
     int k = 0;
46
47
     for(int i = left; i <= right; i++)</pre>
48
49
       if(fabs(p[mid].x - p[i].x) \le d)
50
         temp[k++] = p[i];
     }
51
52
     sort(temp,temp+k,cmpy);
53
     for(int i = 0; i < k; i++)</pre>
55
       for(int j = i+1; j < k && temp[j].y - temp[i].y < d;
          \hookrightarrow j++)
56
         double d3 = Dis(temp[i],temp[j]);
57
58
         d = min(d,d3);
59
60
     }
61
     return d;
62
  ۱ }
63
64
   int main()
65
66
     cin>>n;
67
     for(int i=0; i<n; i++)</pre>
68
69
       double a.b:
       scanf("%lf%lf",&a,&b);
70
71
       p[i] = Point(a,b);
72
73
     sort(p,p+n);
74
     printf("%.3f",Closest_Pair(0,n-1));
75
```

3.1.8 最小覆盖圆

```
#include<cmath>
2
   #include<cstdio>
3
   #include<algorithm>
   using namespace std;
4
5
   const double eps=1e-6;
6
   struct couple
7
8
     double x, y;
     couple(){}
9
10
     couple(const double &xx, const double &yy)
11
12
       x = xx; y = yy;
13
14
   } a[100001];
15
   int n;
   bool operator < (const couple & a, const couple & b)
16
17
18
    return a.x < b.x - eps or (abs(a.x - b.x) < eps and a.y
        \hookrightarrow < b.y - eps);
19
   }
   bool operator == (const couple & a, const couple & b)
20
21
22
    return !(a < b) and !(b < a);
23
   inline couple operator - (const couple &a, const couple
25
   {
26
     return couple(a.x-b.x, a.y-b.y);
27
28
   inline couple operator + (const couple &a, const couple
29
30
     return couple(a.x+b.x, a.y+b.y);
31
   }
32
   inline couple operator * (const couple &a, const double
33
34
     return couple(a.x*b, a.y*b);
35
36
   inline couple operator / (const couple &a, const double
      37
38
     return a*(1/b);
39
40
   inline double operator * (const couple &a, const couple
41
42
     return a.x*b.y-a.y*b.x;
43
44
   inline double len(const couple &a)
45
   {
46
    return a.x*a.x+a.y*a.y;
47
48
   inline double di2(const couple &a, const couple &b)
49
50
    return (a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y);
51
   inline double dis(const couple &a, const couple &b)
52
53
54
    return sqrt((a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y));
   }
55
56
   struct circle
57
58
    double r; couple c;
59
   } cir;
60
   inline bool inside(const couple & x)
62
     return di2(x, cir.c) < cir.r*cir.r+eps;</pre>
63
64
   inline void p2c(int x, int y)
65 {
```

```
cir.c.x = (a[x].x+a[y].x)/2;
      cir.c.y = (a[x].y+a[y].y)/2;
67
68
      cir.r = dis(cir.c, a[x]);
69 }
70 | inline void p3c(int i, int j, int k)
71 | {
      couple x = a[i], y = a[j], z = a[k];
72
73
         \rightarrow sqrt(di2(x,y)*di2(y,z)*di2(z,x))/fabs(x*y+y*z+z*x)/2;
74
      couple t1((x-y).x, (y-z).x), t2((x-y).y, (y-z).y),
         \hookrightarrow t3((len(x)-len(y))/2, (len(y)-len(z))/2);
      cir.c = couple(t3*t2, t1*t3)/(t1*t2);
75
76
77
    inline circle mi()
78
      sort(a + 1, a + 1 + n);
79
      n = unique(a + 1, a + 1 + n) - a - 1;
80
81
      if(n == 1)
82
        cir.c = a[1];
83
84
        cir.r = 0:
        return cir;
85
86
87
      random_shuffle(a + 1, a + 1 + n);
88
      p2c(1, 2);
89
      for(int i = 3; i <= n; i++)
90
        if(!inside(a[i]))
91
          p2c(1, i);
92
          for(int j = 2; j < i; j++)
93
             if(!inside(a[j]))
95
96
               p2c(i, j);
97
               for(int k = 1; k < j; k++)
98
                 if(!inside(a[k]))
99
                   p3c(i,j, k);
100
101
102
      return cir;
103 }
```

3.1.9 多边形内部可视

```
#include <bits/stdc++.h>
3
   using namespace std;
   const int N = 510:
5
   const double eps = 1e-3;
6
7
8 struct Point {
9
    double x, y;
10
    Point() {}
11
    Point(double x, double y): x(x), y(y) {}
    void read() {
12
13
      scanf("%lf %lf", &x, &y);
14
15
     void print() const {
16
       printf("%.10f %.10f\n", x, y);
17
18 }:
19
20 Point p[N];
21 Point A, B;
22 int n, dfn;
23 int g[N][N], vis[N][N], f[N][N], v[N][N];
24
25 | Point operator + (const Point & a, const Point & b) {
    return Point(a.x + b.x, a.y + b.y);
26
27 }
28
29 Point operator - (const Point & a, const Point & b) {
30 return Point(a.x - b.x, a.y - b.y);
```

```
32
33 Point operator * (const Point & a, double p) {
34
    return Point(a.x * p, a.y * p);
35
36
   Point operator / (const Point & a, double p) {
    return Point(a.x / p, a.y / p);
38
39
40
41
   double Cross(const Point & a, const Point & b) {
42
     return a.x * b.y - a.y * b.x;
43
44
   double Dot(const Point & a, const Point & b) {
45
46
     return a.x * b.x + a.y * b.y;
47
48
   int dcmp(double x) {
     if (fabs(x) < eps) return 0;</pre>
51
     return x < 0 ? -1 : 1;
52
53
   Point Get(const Point & P, const Point & v, const Point &
      \hookrightarrow \mathbb{Q}, const Point & w) {
     Point u = P - Q;
56
     double t = Cross(w, u) / Cross(v, w);
     return P + v * t;
57
58
59
   int OnLine(const Point & a, const Point & b, const Point &
     return dcmp(Cross(b - a, b - c)) == 0 && dcmp(Dot(b - a,
61
        \hookrightarrow b - c)) < 0;
62
63
   int C(const Point & P, const Point & A, const Point & Q,
      \hookrightarrow \texttt{const} Point & B) {
65
     Point C = Get(P, A - P, Q, Q - B);
66
     return OnLine(Q, C, B);
67
68
   int Onleft(const Point & a, const Point &b, const Point &
70
     return dcmp(Cross(b - c, a - c)) > 0;
71
   }
72
73
   int visible(int x, int y) {
     int P = (x + n - 1) \% n, Q = (x + 1) \% n;
     Point u = p[y] - p[x], v = p[x] - p[P], w = p[x] - p[Q];
76
     if (Onleft(p[Q], p[x], p[P])) {
77
       return dcmp(Cross(v, u)) > 0 && dcmp(Cross(w, u)) < 0;
78
       return !(dcmp(Cross(v, u)) < 0 && dcmp(Cross(w, u)) >
79
          \hookrightarrow 0);
80
     }
81
   }
82
83
   int solve(int x, int y) {
     if (vis[x][y] == dfn) return g[x][y];
84
85
     vis[x][y] = dfn;
     if (x == y || y == x + 1) return g[x][y] = 1;
87
     for (int i = x; i + 1 \le y; i++) {
88
       if (C(p[x], p[y], p[i], p[i + 1])) return g[x][y] = 0;
89
     for (int i = x + 1; i < y; i++) {
90
91
       if (OnLine(p[x], p[i], p[y])) {
         return g[x][y] = solve(x, i) && solve(i, y);
92
93
       }
     }
94
     if (!visible(x, y) || !visible(y, x)) return g[x][y] =
95
     return g[x][y] = 1;
```

```
97 }
 98
 99
    void DP(int x, int y) {
      if (v[x][y] == dfn || x > y) return;
100
      v[x][y] = dfn;
101
      if (x == y) {
102
103
        f[x][y] = 1;
104
        return;
105
      }
106
      DP(x + 1, y);
107
      DP(x, y - 1);
      f[x][y] = max(f[x][y - 1], f[x + 1][y]);
108
109
      if (g[x][y] == 0) {
110
        int z = x;
111
        while(!g[z][y] && z < y) ++z;
112
        DP(x, z - 1);
113
        DP(z + 1, y);
114
        f[x][y] = max(f[x][y], f[x][z - 1] + f[z + 1][y]);
115
      }
116 }
117
118
    vector<int> ans;
119
    void DFS(int x, int y) {
120
121
      if (x > y) return;
122
      if (x == y) {
123
         ans.push_back(x);
124
        return;
      }
125
      if (f[x][y] == f[x][y - 1]) {
126
127
        DFS(x, y - 1);
      } else if (f[x][y] == f[x + 1][y]) {
128
129
        DFS(x + 1, y);
130
      } else {
131
         int z = x;
         while (!g[z][y] \&\& z < y) ++z;
132
133
        DFS(x, z - 1);
134
         DFS(z + 1, y);
135
    ۱,
136
137
138 int main() {
      freopen("hide.in", "r", stdin);
139
      freopen("hide.out", "w", stdout);
140
141
      while (scanf("%d", &n) && n) {
142
        ++dfn:
        for (int i = 0; i < n; i++) {
143
144
           p[i].read();
145
146
        for (int i = 1; i < n; i++) {
147
           for (int j = i; j < n; j++) {
148
             g[i][j] = solve(i, j);
149
         }
150
151
         DP(1, n - 1);
         cout << f[1][n - 1] << endl;</pre>
152
         ans.clear();
153
154
         DFS(1, n - 1);
         for (int i = 0; i < ans.size(); i++) {</pre>
           printf("%d%c", ans[i] + 1, i + 1 < ans.size() ? ' '
156
              \hookrightarrow : '\n');
157
158
159
      return 0;
160 }
```

3.1.10 V 图

```
const int AIX = 5;
const int MAXM = AIX * MAXN;

struct point {
   double x, y;
}
```

```
6
     int index:
     struct Edge *in;
8
     point(double _x = 0, double _y = 0) : x(_x), y(_y) {}
9
   inline bool operator< (const point &a, const point &b) {</pre>
10
    return a.x < b.x \mid \mid (sgn(a.x - b.x) == 0 && a.y < b.y);
11
12 }
   inline double cross(const point &a, const point &b, const
13
      \hookrightarrow point &c) { return det
14
       (b - a, c - a); }
15
   struct Edge {
     point *Org, *Dest;
16
17
     Edge *Onext, *Oprev, *Dnext, *Dprev;
18
19
   inline point* Other(const Edge *e, const point *p) {
      \hookrightarrow return e->Org == p ?
        e->Dest : e->Org; }
20
21
   inline Edge* Next(const Edge *e, const point *p) { return
      \hookrightarrow e->0rg == p ? e->0next
        : e->Dnext: }
   inline Edge* Prev(const Edge *e, const point *p) { return
      \hookrightarrow e->0rg == p ? e->0prev
        : e->Dprev; }
24
25
   struct gEdge {
26
     int u, v;
27
     double w;
28
     gEdge() {}
29
     gEdge(int _u, int _v, double _w) : u(_u), v(_v), w(_w)
        → {}
30 };
31
   inline bool operator< (const gEdge &a, const gEdge &b) {
      32
   point p[MAXN], *Q[MAXN];
33
   Edge mem[AIX * MAXN], *elist[AIX * MAXN];
34
   static int nfree;
35
   //Alloc memory
36
   inline void Alloc_Memory(const int &n) {
     nfree = AIX * n;
37
38
     Edge *e = mem;
     for (int i = 0; i < nfree; ++i)</pre>
39
       elist[i] = e++;
40
41
42
   //Add an edge to a ring of edges
   inline void Splice(Edge *a, Edge *b, point *v) {
     Edge *next;
45
     if (a->0rg == v)
46
       next = a->0next, a->0next = b;
47
     else
48
       next = a->Dnext, a->Dnext = b;
49
     if (next->Org == v)
50
       next->Oprev = b;
51
     else
52
       next->Dprev = b;
53
     if (b->0rg == v)
       b->Onext = next, b->Oprev = a;
54
55
56
       b->Dnext = next, b->Dprev = a;
57
   //Initialise a new edge
   inline Edge *MakeEdge(point *u, point *v) {
59
60
     Edge *e = elist[--nfree];
     e->Onext = e->Oprev = e->Dnext = e->Dprev = e;
62
     e->0rg = u, e->Dest = v;
63
     if (!u->in)
       u->in = e;
64
     if (!v->in)
65
66
       v->in = e;
67
     return e;
68
69
   //Creates a new edge and adds it to two rings of edges.
70 inline Edge *Join(Edge *a, point *u, Edge *b, point *v,
      → int side) {
     Edge *e = MakeEdge(u, v);
```

```
if (side == 1) {
                                                                                Edge *El = Next(B, 0), *Er = Prev(B, D), *next, *prev;
72
73
        if (a->0rg == u)
                                                                                point *l = Other(El, O), *r = Other(Er, D);
                                                                       144
74
          Splice(a->Oprev, e, u);
                                                                       145
                                                                                11 = *0 - *1, 12 = *D - *1, r1 = *0 - *r, r2 = *D -
75
        else
                                                                                  → *r:
          Splice(a->Dprev, e, u);
                                                                                double cl = det(11, 12), cr = det(r1, r2);
76
                                                                       146
                                                                                bool BL = cl > EPS, BR = cr > EPS;
77
        Splice(b, e, v);
                                                                       147
      }
78
                                                                                if (!BL && !BR)
79
      else {
                                                                       149
                                                                                  break;
80
        Splice(a, e, u);
                                                                       150
                                                                               if (BL) {
81
        if (b->0rg == v)
                                                                       151
                                                                                 double dl = dot(11, 12);
          Splice(b->Oprev, e, v);
                                                                                  cot_L = dl / cl;
82
                                                                       152
83
                                                                       153
                                                                                  do {
84
          Splice(b->Dprev, e, v);
                                                                       154
                                                                                   next = Next(E1, 0);
85
      }
                                                                       155
                                                                                    uu = *0 - *0ther(next, 0);
                                                                                   vv = *D - *Other(next, 0);
86
      return e;
                                                                       156
87 }
                                                                       157
                                                                                   N1 = det(uu, vv);
                                                                                   if (!(N1 > EPS))
    //Remove an edge
                                                                       158
88
89 inline void Remove(Edge *e) {
                                                                       159
                                                                                     break;
90
     point *u = e->Org, *v = e->Dest;
                                                                       160
                                                                                    cot_N = dot(uu, vv) / N1;
     if (u->in == e)
                                                                                   if (cot_N > cot_L)
                                                                       161
92
        u->in = e->0next;
                                                                       162
                                                                                     break:
     if (v->in == e)
                                                                                    Remove(E1);
93
                                                                       163
        v->in = e->Dnext;
                                                                                    El = next:
                                                                       164
94
      if (e->0next->0rg == u)
                                                                       165
                                                                                    cot_L = cot_N;
95
96
        e->Onext->Oprev = e->Oprev;
                                                                       166
97
                                                                       167
                                                                                  while (true);
98
        e->Onext->Dprev = e->Oprev;
                                                                       168
                                                                                if (BR) {
99
      if (e->0prev->0rg == u)
                                                                       169
        e->Oprev->Onext = e->Onext;
                                                                                  double dr = dot(r1, r2);
100
                                                                       170
                                                                                  cot_R = dr / cr;
101
                                                                       171
      else
        e->Oprev->Dnext = e->Onext;
102
                                                                       172
                                                                                  do {
      if (e->Dnext->Org == v)
                                                                                   prev = Prev(Er, D);
103
                                                                       173
104
        e->Dnext->Oprev = e->Dprev;
                                                                       174
                                                                                    uu = *0 - *0ther(prev, D);
105
      else
                                                                       175
                                                                                    vv = *D - *Other(prev, D);
106
        e->Dnext->Dprev = e->Dprev;
                                                                       176
                                                                                   P1 = det(uu, vv);
107
      if (e->Dprev->Org == v)
                                                                       177
                                                                                   if (!(P1 > EPS))
108
        e->Dprev->Onext = e->Dnext;
                                                                       178
                                                                                     break;
109
                                                                       179
                                                                                    cot_P = dot(uu, vv) / P1;
110
        e->Dprev->Dnext = e->Dnext;
                                                                       180
                                                                                    if (cot_P > cot_R)
111
      elist[nfree++] = e;
                                                                       181
                                                                                     break:
112 }
                                                                       182
                                                                                    Remove(Er):
113 //Determines the lower tangent of two triangulations
                                                                       183
                                                                                    Er = prev;
114 inline void Low_tangent(Edge *e_1, point *o_1, Edge *e_r,
                                                                       184
                                                                                    cot_R = cot_P;
                                                                                  }
       \hookrightarrow point *o_r, Edge
115
         **1_low, point **0L, Edge **r_low, point **0R) {
                                                                       186
                                                                                  while (true);
                                                                               }
116
      point *d_1 = Other(e_1, o_1), *d_r = Other(e_r, o_r);
                                                                       187
      while (true) {
                                                                               1 = Other(E1, 0); r = Other(Er, D);
117
                                                                       188
118
       if (cross(*o_1, *o_r, *d_1) < -EPS) {</pre>
                                                                       189
                                                                                if (!BL || (BL && BR && cot_R < cot_L)) {</pre>
119
          e_l = Prev(e_l, d_l);
                                                                       190
                                                                                  B = Join(B, 0, Er, r, 0);
120
          o_1 = d_1;
                                                                       191
                                                                                  D = r;
                                                                               }
121
          d_1 = Other(e_1, o_1);
                                                                       192
122
                                                                       193
                                                                                else {
123
        else if (cross(*o_1, *o_r, *d_r) < -EPS) {
                                                                       194
                                                                                  B = Join(El, 1, B, D, 0);
          e_r = Next(e_r, d_r);
                                                                       195
                                                                                  0 = 1:
124
125
          o_r = d_r;
                                                                       196
          d_r = Other(e_r, o_r);
                                                                             }
126
                                                                       197
127
        }
                                                                       198
                                                                             while (true);
128
        else
                                                                       199
129
          break;
                                                                       200
                                                                           inline void Divide(int s, int t, Edge **L, Edge **R) {
130
                                                                       201
                                                                             Edge *a, *b, *c, *ll, *lr, *rl, *rr, *tangent;
131
      *OL = o_1, *OR = o_r;
                                                                       202
                                                                              int n = t - s + 1;
                                                                              if (n == 2)
132
      *l_low = e_l, *r_low = e_r;
                                                                       203
133 }
                                                                       204
                                                                                *L = *R = MakeEdge(Q[s], Q[t]);
                                                                              else if (n == 3) {
134
    inline void Merge(Edge *lr, point *s, Edge *rl, point *u,
                                                                       205
                                                                       206
                                                                               a = MakeEdge(Q[s], Q[s + 1]);
       double cot_L, cot_R, N1, cot_N, P1, cot_P;
135
                                                                       207
                                                                               b = MakeEdge(Q[s + 1], Q[t]);
      point 11, 12, r1, r2, uu, vv;
                                                                       208
                                                                               Splice(a, b, Q[s + 1]);
136
      point *0, *D, *OR, *OL;
                                                                                double v = cross(*Q[s], *Q[s + 1], *Q[t]);
                                                                       209
138
      Edge *B, *L, *R;
                                                                       210
                                                                               if (v > EPS) {
139
      Low_tangent(lr, s, rl, u, &L, &OL, &R, &OR);
                                                                       211
                                                                                  c = Join(a, Q[s], b, Q[t], 0);
      *tangent = B = Join(L, OL, R, OR, O);
                                                                       212
                                                                                  *L = a, *R = b;
140
141
      0 = OL, D = OR;
                                                                       213
142
      do {
                                                                       214
                                                                               else if (v < -EPS) {
```

```
c = Join(a, Q[s], b, Q[t], 1);
215
216
          *L = c, *R = c;
217
        }
218
        else
219
          *L = a, *R = b;
220
      else if(n > 3) {
221
        int split = (s + t) / 2;
222
223
        Divide(s, split, &ll, &lr);
224
        Divide(split + 1, t, &rl, &rr);
225
        Merge(lr, Q[split], rl, Q[split + 1], &tangent);
226
        if (tangent->Org == Q[s])
227
          11 = tangent;
        if (tangent->Dest == Q[t])
228
229
          rr = tangent;
230
        *L = 11; *R = rr;
231
      }
232 }
233 int task, n, m, k, root[MAXN];
234 gEdge E[MAXM], MST[MAXN];
235 | inline int Make_Graph() {
236
      Edge *start, *e;
      int M = 0:
237
238
      point *u, *v;
239
      for(int i = 0; i < n; ++i) {
240
        u = p + i;
241
        start = e = u->in;
        do {
242
          v = Other(e, u);
243
244
          if (u < v)
245
            E[M++] = gEdge(u - p + 1, v - p + 1, dis(*u, *v));
246
          e = Next(e, u);
247
        }
248
        while(e != start);
249
      }
250
      return M;
251
    int find_root(const int &x) { return root[x] ? root[x] =
252
        find_root(root[x]) : x;
253
         }
254
    inline bool merge(const int &x, const int &y) {
      int p = find_root(x), q = find_root(y);
255
      if (p != q) {
256
        root[p] = q;
258
        return true;
      }
259
260
      else
261
        return false;
262
263
    inline void kruskal(gEdge *E, int m, int n, gEdge* MST) {
264
      for (int i = 1; i <= n; ++i)
265
       root[i] = 0;
      sort(E, E + m);
266
      int tot = 0:
267
268
      for (int i = 0; i < m; ++i)
269
        if (merge(E[i].u, E[i].v))
270
          MST[tot++] = E[i];
271 | }
272
   inline void MinimumEuclideanSpaningTree(point* p, int n,
      273
      Alloc_Memory(n);
274
      sort(p, p + n);
275
      for (int i = 0; i < n; ++i)
276
        Q[i] = p + i;
277
      Edge *L, *R;
      Divide(0, n - 1, &L, &R);
278
279
      m = Make_Graph();
280
      kruskal(E, m, n, MST);
281 }
282 | int main() {
     for (scanf("%d", &task); task--; ) {
283
        scanf("%d", &k);
284
285
        for (n = 0; scanf("%lf", &p[n].x) == 1 && p[n].x !=
           \hookrightarrow -1; ++n) {
```

```
286
           scanf("%lf", &p[n].y);
          p[n].in = NULL;
287
288
          p[n].index = n;
289
        if (n == 1) {
290
          printf("0\n");
291
          continue:
293
294
        MinimumEuclideanSpaningTree(p, n, MST);
295
        printf("%d\n", int(ceil(k > n ? 0 : MST[n - k - 1].w))

→ + EPS)):
296
297
```

3.2 三维

3.2.1 三维点类

```
// 三维绕轴旋转, 大拇指指向 axis 向量方向, 四指弯曲
     →方向转 w 弧度
 2
  Point rotate(const Point& s, const Point& axis, DB w) {
    DB x = axis.x, y = axis.y, z = axis.z;
4
    DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
5
        cosw = cos(w), sinw = sin(w);
6
     DB a[4][4];
     memset(a, 0, sizeof a);
 7
8
     a[3][3] = 1;
     a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
9
     a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
     a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
11
12
     a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
     a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
13
14
     a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
15
     a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
     a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
16
     a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
17
18
     DB ans [4] = \{0, 0, 0, 0\}, c[4] = \{s.x, s.y, s.z, 1\};
19
     for (int i = 0; i < 4; ++ i)
20
      for (int j = 0; j < 4; ++ j)
21
         ans[i] += a[j][i] * c[j];
22
     return Point(ans[0], ans[1], ans[2]);
23
```

3.2.2 凸包

```
__inline P cross(const P& a, const P& b) {
     return P(
 3
         a.y * b.z - a.z * b.y,
         a.z * b.x - a.x * b.z,
 5
         a.x * b.y - a.y * b.x
6
           );
7
   }
8
9
    _inline DB mix(const P& a, const P& b, const P& c) {
10
    return dot(cross(a, b), c);
11
12
   __inline DB volume(const P& a, const P& b, const P& c,
     14
    return mix(b - a, c - a, d - a);
15
16
17
   struct Face {
18
    int a, b, c;
     __inline Face() {}
19
     __inline Face(int _a, int _b, int _c):
20
       a(_a), b(_b), c(_c) {}
22
     __inline DB area() const {
23
       return 0.5 * cross(p[b] - p[a], p[c] - p[a]).len();
24
25
     __inline P normal() const {
```

```
return cross(p[b] - p[a], p[c] - p[a]).unit();
26
27
28
     __inline DB dis(const P& p0) const {
29
       return dot(normal(), p0 - p[a]);
     }
30
   };
31
32
   std::vector<Face> face, tmp; // Should be O(n).
33
   int mark[N][N], Time, n;
35
   __inline void add(int v) {
36
37
     ++ Time:
38
     clear(tmp);
39
     for (int i = 0; i < (int)face.size(); ++ i) {</pre>
40
       int a = face[i].a, b = face[i].b, c = face[i].c;
       if (sign(volume(p[v], p[a], p[b], p[c])) > 0) {
41
         mark[a][b] = mark[b][a] = mark[a][c] =
42
43
            mark[c][a] = mark[b][c] = mark[c][b] = Time;
44
45
46
         tmp.push_back(face[i]);
47
       }
48
     }
49
     clear(face); face = tmp;
50
     for (int i = 0; i < (int)tmp.size(); ++ i) {</pre>
51
       int a = face[i].a, b = face[i].b, c = face[i].c;
52
       if (mark[a][b] == Time) face.emplace_back(v, b, a);
       if (mark[b][c] == Time) face.emplace_back(v, c, b);
53
       if (mark[c][a] == Time) face.emplace_back(v, a, c);
54
       assert(face.size() < 500u);</pre>
55
56
   }
57
58
59
   void reorder() {
     for (int i = 2; i < n; ++ i) {
60
61
       P \text{ tmp} = cross(p[i] - p[0], p[i] - p[1]);
62
       if (sign(tmp.len())) {
63
          std::swap(p[i], p[2]);
         for (int j = 3; j < n; ++ j)
64
65
            if (sign(volume(p[0], p[1], p[2], p[j]))) {
              std::swap(p[j], p[3]);
66
67
              return;
68
            }
69
70
     }
71
   }
72
73
   void build_convex() {
74
     reorder();
75
     clear(face);
76
     face.emplace_back(0, 1, 2);
77
     face.emplace_back(0, 2, 1);
78
     for (int i = 3; i < n; ++ i)
79
       add(i):
80 }
```

3.2.3 最小覆盖球

```
#include<iostream>
2
   #include<cstring>
3 #include<algorithm>
4 #include<cstdio>
  #include<cmath>
7
   using namespace std;
8
9
   const int eps = 1e-8;
10
11
   struct Tpoint
12
13
     double x, y, z;
14
  | }:
15
```

```
16
   int npoint, nouter;
17
18
   Tpoint pt[200000], outer[4],res;
   double radius,tmp;
19
   inline double dist(Tpoint p1, Tpoint p2) {
20
     double dx=p1.x-p2.x, dy=p1.y-p2.y, dz=p1.z-p2.z;
21
22
     return ( dx*dx + dy*dy + dz*dz );
23
24
   inline double dot(Tpoint p1, Tpoint p2) {
25
     return p1.x*p2.x + p1.y*p2.y + p1.z*p2.z;
26
27
   void ball() {
28
     Tpoint q[3]; double m[3][3], sol[3], L[3], det;
29
30
     res.x = res.y = res.z = radius = 0;
     switch ( nouter ) {
31
32
       case 1: res=outer[0]: break:
33
           res.x=(outer[0].x+outer[1].x)/2;
            res.y=(outer[0].y+outer[1].y)/2;
35
36
           res.z=(outer[0].z+outer[1].z)/2;
           radius=dist(res, outer[0]);
37
38
           break:
39
       case 3:
40
           for (i=0; i<2; ++i ) {
41
              q[i].x=outer[i+1].x-outer[0].x;
42
              q[i].y=outer[i+1].y-outer[0].y;
43
              q[i].z=outer[i+1].z-outer[0].z;
44
45
           for (i=0; i<2; ++i) for(j=0; j<2; ++j)
46
              m[i][j]=dot(q[i], q[j])*2;
47
            for (i=0; i<2; ++i ) sol[i]=dot(q[i], q[i]);</pre>
48
            if (fabs(det=m[0][0]*m[1][1]-m[0][1]*m[1][0]) < eps)
49
              return;
50
           L[0]=(sol[0]*m[1][1]-sol[1]*m[0][1])/det:
51
           L[1]=(sol[1]*m[0][0]-sol[0]*m[1][0])/det;
52
            res.x=outer[0].x+q[0].x*L[0]+q[1].x*L[1];
53
            res.y=outer[0].y+q[0].y*L[0]+q[1].y*L[1];
54
            res.z=outer[0].z+q[0].z*L[0]+q[1].z*L[1];
           radius=dist(res, outer[0]);
55
56
           break:
57
       case 4:
           for (i=0; i<3; ++i) {
58
              q[i].x=outer[i+1].x-outer[0].x;
59
60
              q[i].y=outer[i+1].y-outer[0].y;
61
              q[i].z=outer[i+1].z-outer[0].z;
62
              sol[i]=dot(q[i], q[i]);
63
64
            for (i=0;i<3;++i)</pre>
65
              for(j=0;j<3;++j) m[i][j]=dot(q[i],q[j])*2;</pre>
66
            det= m[0][0]*m[1][1]*m[2][2]
67
             + m[0][1]*m[1][2]*m[2][0]
68
              + m[0][2]*m[2][1]*m[1][0]
              - m[0][2]*m[1][1]*m[2][0]
69
70
              - m[0][1]*m[1][0]*m[2][2]
              - m[0][0]*m[1][2]*m[2][1];
71
72
            if ( fabs(det) < eps ) return;</pre>
73
            for (j=0; j<3; ++j) {
74
              for (i=0; i<3; ++i) m[i][j]=sol[i];</pre>
75
              L[i] = (m[0][0]*m[1][1]*m[2][2]
76
                  + m[0][1]*m[1][2]*m[2][0]
                  + m[0][2]*m[2][1]*m[1][0]
77
78
                  - m[0][2]*m[1][1]*m[2][0]
79
                  - m[0][1]*m[1][0]*m[2][2]
80
                  - m[0][0]*m[1][2]*m[2][1]
                 ) / det;
81
82
              for (i=0; i<3; ++i)
                m[i][j]=dot(q[i], q[j])*2;
83
84
           }
85
           res=outer[0];
           for (i=0; i<3; ++i ) {
86
87
             res.x += q[i].x * L[i];
88
             res.y += q[i].y * L[i];
```

```
res.z += q[i].z * L[i];
 89
 90
 91
             radius=dist(res, outer[0]);
 92
 93 }
    void minball(int n) {
 94
      ball():
      //printf("(%.3lf,%.3lf,%.3lf) %.3lf\n",
 96
          \hookrightarrow res.x,res.y,res.z,radius);
 97
      if ( nouter<4 )
 98
        for (int i=0; i<n; ++i)
           if (dist(res, pt[i])-radius>eps) {
 99
100
             outer[nouter]=pt[i];
101
             ++nouter;
102
             minball(i);
103
             --nouter;
             if (i>0) {
104
105
               Tpoint Tt = pt[i];
106
               memmove(&pt[1], &pt[0], sizeof(Tpoint)*i);
               pt[0]=Tt;
107
108
             }
109
110
    | }
111
    void solve()
112
113
      for (int i=0;i<npoint;i++)</pre>
          \hookrightarrow scanf("%lf%lf",&pt[i].x,&pt[i].y,&pt[i].z);
114
      random_shuffle(pt, pt + npoint);
      radius=-1:
115
      for (int i=0;i<npoint;i++){</pre>
116
117
        if (dist(res,pt[i])-radius>eps){
118
119
           outer[0] = pt[i];
120
           minball(i);
121
122
123
      printf("%.5f\n",sqrt(radius));
124
125
    int main(){
      for( ; cin >> npoint && npoint; )
126
127
        solve():
128
      return 0;
129 }
```

4. 字符串

4.1 AC 自动机

```
1 int newnode()
2 {
3
4
     memset(ch[tot], 0, sizeof(ch[tot]));
5
     fail[tot] = 0;
6
     dep[tot] = 0;
7
     par[tot] = 0;
8
9
     return tot;
10
   }
11
   void insert(char *s,int x)
12
     if(*s == '\0') return;
13
14
     else
15
       int &y = ch[x][*s - 'a'];
16
17
       if(y == 0)
18
19
         y = newnode();
20
         par[y] = x;
21
22
         dep[y] = dep[x] + 1;
23
24
25
       insert(s + 1, y);
```

```
27
28
   void build()
29
30
     int line[maxn];
     int f = 0, r = 0;
31
32
33
      fail[root] = root;
34
35
     for(int i = 0; i < alpha; i++)</pre>
36
     {
37
        if(ch[root][i])
38
        {
39
          fail[ch[root][i]] = root;
40
          line[r++] = ch[root][i];
        7
41
42
        else
43
        {
          ch[root][i] = root;
45
        }
     }
46
47
     while(f != r)
48
49
50
        int x = line[f++];
51
52
        for(int i = 0; i < alpha; i++)</pre>
53
          if(ch[x][i])
54
55
          {
            fail[ch[x][i]] = ch[fail[x]][i];
57
            line[r++] = ch[x][i];
58
          }
59
          else
60
          {
            ch[x][i] = ch[fail[x]][i];
61
62
63
64
     }
65
```

4.2 后缀数组

```
const int MAXN = MAXL * 2 + 1;
   int a[MAXN], x[MAXN], y[MAXN], c[MAXN], sa[MAXN],
      \hookrightarrow rank[MAXN], height[MAXN];
   void calc_sa(int n) {
     int m = alphabet, k = 1;
     memset(c, 0, sizeof(*c) * (m + 1));
     for (int i = 1; i <= n; ++i) c[x[i] = a[i]]++;</pre>
     for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
8
     for (int i = n; i; --i) sa[c[x[i]]--] = i;
9
     for (; k <= n; k <<= 1) {
10
       int tot = k:
11
       for (int i = n - k + 1; i \le n; ++i) y[i - n + k] = i;
12
       for (int i = 1; i <= n; ++i)
13
         if (sa[i] > k) y[++tot] = sa[i] - k;
14
       memset(c, 0, sizeof(*c) * (m + 1));
       for (int i = 1; i \le n; ++i) c[x[i]]++;
15
16
       for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
17
       for (int i = n; i; --i) sa[c[x[y[i]]]--] = y[i];
       for (int i = 1; i <= n; ++i) y[i] = x[i];</pre>
19
       tot = 1; x[sa[1]] = 1;
20
       for (int i = 2; i <= n; ++i) {
         if (max(sa[i], sa[i - 1]) + k > n || y[sa[i]] !=
21
             \hookrightarrow y[sa[i - 1]] || y[sa[i] + k] != y[sa[i - 1] +
             \hookrightarrow k]) ++tot;
         x[sa[i]] = tot;
22
23
24
       if (tot == n) break; else m = tot;
25
26 }
```

```
27
   void calc_height(int n) {
     for (int i = 1; i <= n; ++i) rank[sa[i]] = i;</pre>
28
29
     for (int i = 1; i <= n; ++i) {
       \label{eq:height[rank[i]] = max(0, height[rank[i - 1]] - 1);} \\
30
        if (rank[i] == 1) continue;
        int j = sa[rank[i] - 1];
32
        while (max(i, j) + height[rank[i]] <= n && a[i +</pre>
33

    height[rank[i]]] == a[j + height[rank[i]]])

→ ++height[rank[i]];

34
     }
35
   | }
```

4.3 后缀自动机

```
1 static const int MAXL = MAXN * 2; // MAXN is original
   static const int alphabet = 26; // sometimes need
      \hookrightarrow \texttt{changing}
   int 1, last, cnt, trans[MAXL][alphabet], par[MAXL],
      \hookrightarrow \text{sum}\,[\text{MAXL}]\,,\,\,\text{seq}\,[\text{MAXL}]\,,\,\,\text{mxl}\,[\text{MAXL}]\,,\,\,\text{size}\,[\text{MAXL}]\,;\,\,\,\,\,//\,\,\,\text{mxl}
      \hookrightarrow is maxlength, size is the size of right
   char str[MAXL];
   inline void init() {
     l = strlen(str + 1); cnt = last = 1;
 6
     for (int i = 0; i \le 1 * 2; ++i) memset(trans[i], 0,
        ⇔ sizeof(trans[i])):
8
     memset(par, 0, sizeof(*par) * (1 * 2 + 1));
      memset(mxl, 0, sizeof(*mxl) * (1 * 2 + 1));
      memset(size, 0, sizeof(*size) * (1 * 2 + 1));
10
11 | }
12 inline void extend(int pos, int c) {
     int p = last, np = last = ++cnt;
     mxl[np] = mxl[p] + 1; size[np] = 1;
14
15
      for (; p && !trans[p][c]; p = par[p]) trans[p][c] = np;
16
      if (!p) par[np] = 1;
17
      else {
        int q = trans[p][c];
18
        if (mxl[p] + 1 == mxl[q]) par[np] = q;
19
20
        else {
21
          int nq = ++cnt;
          mxl[nq] = mxl[p] + 1;
23
          memcpy(trans[nq], trans[q], sizeof(trans[nq]));
24
          par[nq] = par[q];
25
          par[np] = par[q] = nq;
          for (; trans[p][c] == q; p = par[p]) trans[p][c] =
26
27
28
     }
29 }
   inline void buildsam() {
30
     for (int i = 1; i <= 1; ++i) extend(i, str[i] - 'a');
31
     memset(sum, 0, sizeof(*sum) * (1 * 2 + 1));
32
     for (int i = 1; i <= cnt; ++i) sum[mx1[i]]++;</pre>
     for (int i = 1; i <= 1; ++i) sum[i] += sum[i - 1];</pre>
34
     for (int i = cnt; i; --i) seq[sum[mxl[i]]--] = i;
35
     for (int i = cnt; i; --i) size[par[seq[i]]] +=
36
         \hookrightarrow size[seq[i]];
   }
37
```

4.4 广义后缀自动机

```
for (; lastnode && c[lastnode][x] == nownode;
              → lastnode = par[lastnode]) {
11
            c[lastnode][x] = auxnode;
          }
12
13
          last = auxnode;
       }
14
15
     } else {
16
       int newnode = ++cnt; l[newnode] = l[lastnode] + 1;
17
        for (; lastnode && !c[lastnode][x]; lastnode =

→ par[lastnode]) c[lastnode][x] = newnode;
18
        if (!lastnode) par[newnode] = 1;
19
        else {
20
          int nownode = c[lastnode][x];
          if (l[lastnode] + 1 == l[nownode]) par[newnode] =
             \hookrightarrow nownode:
22
          else {
            int auxnode = ++cnt; 1[auxnode] = 1[lastnode] + 1;
23
24
            for (int i = 0; i < alphabet; ++i) c[auxnode][i] =</pre>
               \hookrightarrow c[nownode][i];
            par[auxnode] = par[nownode]; par[nownode] =
               → par[newnode] = auxnode;
            for (; lastnode && c[lastnode][x] == nownode;
26
               \hookrightarrow lastnode = par[lastnode]) {
27
              c[lastnode][x] = auxnode;
28
29
30
31
        last = newnode;
32
33 }
```

4.5 manacher

```
1
   void Manacher(std::string s,int p[])
 2
   {
 3
        string t = "$#";
 4
 5
        for (int i = 0; i < s.size(); i++)</pre>
 6
 7
            t += s[i];
 8
            t += "#";
 9
10
11
        std::vector<int> p(t.size(), 0);
12
13
        int mx = 0, id = 0;
14
15
        for (int i = 1; i < t.size(); i++)</pre>
16
            p[i] = mx > i ? min(p[2 * id - i], mx - i) : 1;
17
            while (t[i + p[i]] == t[i - p[i]]) ++p[i];
18
19
20
            if (mx < i + p[i])
21
            {
22
                 mx = i + p[i];
23
                 id = i;
24
25
```

4.6 回文自动机

```
int nT, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN],
     \hookrightarrow 1[MAXN], s[MAXN];
2
  int allocate(int len) {
3
    l[nT] = len;
4
    r[nT] = 0;
    fail[nT] = 0:
    memset(c[nT], 0, sizeof(c[nT]));
6
7
    return nT++;
8
  }
9 void init() {
```

```
10
     nT = nStr = 0:
     int newE = allocate(0);
11
12
     int new0 = allocate(-1);
     last = newE:
13
     fail[newE] = newO:
14
     fail[new0] = newE;
15
     s[0] = -1;
16
17 }
18
  void add(int x) {
     s[++nStr] = x;
19
20
     int now = last:
21
     while (s[nStr - l[now] - 1] != s[nStr]) now = fail[now];
22
     if (!c[now][x]) {
       int newnode = allocate(l[now] + 2), &newfail =
23
          \hookrightarrow fail[newnode];
24
       newfail = fail[now];
       while (s[nStr - l[newfail] - 1] != s[nStr]) newfail =
25
          → fail[newfail]:
26
       newfail = c[newfail][x];
       c[now][x] = newnode;
27
28
29
     last = c[now][x];
     r[last]++:
30
31
32
   void count() {
33
     for (int i = nT - 1; i >= 0; i--) {
34
       r[fail[i]] += r[i];
35
36 }
```

4.7 循环串的最小表示

```
1 // n 必须是 2 的次幂
  void fft(Complex a[], int n, int f) {
    for (int i = 0; i < n; ++i)
3
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
7
       Complex w = Complex(1, 0);
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
8
9
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
10
         for (int k = 0; k < i; ++k) {
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
11
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
       }
     }
15
  }
16
```

5. 数据结构

5.1 可并堆

```
1 int merge(int x,int y)
2
  //p[i] 结点 i 的权值,这里是维护大根堆
  //d[i] 在 i 的子树中, i 到右叶子结点的最远距离.
5
6
      if(!x) return v:
7
      if(!y) return x;
8
9
      if(p[x] < p[y]) std::swap(x, y);
10
      r[x] = merge(r[x], y);
11
      if(r[x]) fa[r[x]] = x;
12
13
      if(d[l[x]] < d[r[x]]) std::swap(l[x], r[x]);//调整树
14
        →的结构,使其满足左偏性质
15
      d[x] = d[r[x]] + 1;
16
17
      return x;
```

18 }

5.2 KD-Tree

```
long long norm(const long long &x) {
 2
            For manhattan distance
       //
 3
        return std::abs(x);
 4
            For euclid distance
 5
        return x * x;
6
 8
   struct Point {
       int x, y, id;
10
11
        const int& operator [] (int index) const {
            if (index == 0) {
12
13
                return x;
            } else {
14
15
                return y;
16
17
18
19
        friend long long dist(const Point &a, const Point &b)
20
            long long result = 0;
            for (int i = 0; i < 2; ++i) {
21
22
                result += norm(a[i] - b[i]);
23
            }
24
            return result;
        }
25
26
   } point[N];
28
   struct Rectangle {
        int min[2], max[2];
29
30
31
        Rectangle() {
            min[0] = min[1] = INT_MAX; // sometimes int is
               \hookrightarrow not enough
33
            max[0] = max[1] = INT_MIN;
34
       }
35
        void add(const Point &p) {
36
37
            for (int i = 0; i < 2; ++i) {
38
                min[i] = std::min(min[i], p[i]);
39
                max[i] = std::max(max[i], p[i]);
            }
40
       }
41
42
        long long dist(const Point &p) {
43
            long long result = 0;
45
            for (int i = 0; i < 2; ++i) {
46
                //
                     For minimum distance
47
                result += norm(std::min(std::max(p[i],
                   \hookrightarrow \min[i]), \max[i]) - p[i]);
48
                      For maximum distance
                result += std::max(norm(max[i] - p[i]),
49
                   \hookrightarrow norm(min[i] - p[i]));
            7
50
51
            return result;
       }
52
53
   };
55
   struct Node {
56
       Point seperator;
57
        Rectangle rectangle;
58
       int child[2];
59
        void reset(const Point &p) {
            seperator = p;
61
62
            rectangle = Rectangle();
63
            rectangle.add(p);
            child[0] = child[1] = 0;
64
```

```
65
        }
    } tree[N << 1];
66
67
68
    int size, pivot;
69
    bool compare(const Point &a, const Point &b) {
70
        if (a[pivot] != b[pivot]) {
71
72
            return a[pivot] < b[pivot];</pre>
73
74
        return a.id < b.id;</pre>
   }
75
76
    // 左閉右開: build(1, n + 1)
77
78
    int build(int 1, int r, int type = 1) {
        pivot = type;
79
80
        if (1 >= r) {
81
            return 0;
82
        int x = ++size:
83
84
        int mid = 1 + r >> 1;
        std::nth_element(point + 1, point + mid, point + r,
85
86
        tree[x].reset(point[mid]);
        for (int i = 1; i < r; ++i) {
87
            tree[x].rectangle.add(point[i]);
88
89
90
        tree[x].child[0] = build(1, mid, type ^ 1);
91
        tree[x].child[1] = build(mid + 1, r, type ^ 1);
92
        return x;
93 }
94
    int insert(int x, const Point &p, int type = 1) {
95
96
        pivot = type;
97
        if (x == 0) {
98
            tree[++size].reset(p);
99
            return size;
100
101
        tree[x].rectangle.add(p);
102
        if (compare(p, tree[x].seperator)) {
103
             tree[x].child[0] = insert(tree[x].child[0], p,
               \hookrightarrow type ^ 1);
        } else {
104
            tree[x].child[1] = insert(tree[x].child[1], p,
105
                \hookrightarrow type ^ 1);
106
107
        return x;
108
   | }
109
    // For minimum distance
110
    // For maximum: 下面递归 query 时 0, 1 换顺序;< and
111

→ >; min and max

112
    void query(int x, const Point &p, std::pair<long long,</pre>
       \hookrightarrow int> &answer, int type = 1) {
113
        pivot = type;
        if (x == 0 || tree[x].rectangle.dist(p) >
114
           \hookrightarrow answer.first) {
115
            return;
116
117
        answer = std::min(answer,
118
                  std::make_pair(dist(tree[x].seperator, p),
                     if (compare(p, tree[x].seperator)) {
119
             query(tree[x].child[0], p, answer, type ^{1});
120
121
             query(tree[x].child[1], p, answer, type ^ 1);
122
        } else {
123
             query(tree[x].child[1], p, answer, type ^ 1);
             query(tree[x].child[0], p, answer, type ^ 1);
124
125
   }
126
127
128
    std::priority_queue<std::pair<long long, int> > answer;
129
130 void query(int x, const Point &p, int k, int type = 1) {
```

```
pivot = type;
        if (x == 0 || (int)answer.size() == k &&
132
          133
       7
134
        answer.push(std::make_pair(dist(tree[x].seperator, p),
135
          \hookrightarrow tree[x].seperator.id));
        if ((int)answer.size() > k) {
136
137
           answer.pop();
138
        }
139
        if (compare(p, tree[x].seperator)) {
140
           query(tree[x].child[0], p, k, type ^ 1);
           query(tree[x].child[1], p, k, type ^ 1);
141
142
143
           query(tree[x].child[1], p, k, type ^ 1);
           query(tree[x].child[0], p, k, type ^ 1);
144
        }
145
146
   }
```

5.3 Treap

```
struct Node{
 1
 2
     int mn, key, size, tag;
     bool rev;
 3
 4
     Node* ch[2]:
     Node(int mn, int key, int size): mn(mn), key(key),
        \hookrightarrow \texttt{size}(\texttt{size})\,,\,\,\texttt{rev}(0)\,,\,\,\texttt{tag}(0)\{\}
 6
     void downtag();
 7
     Node* update(){
8
       mn = min(ch[0] \rightarrow mn, min(key, ch[1] \rightarrow mn));
9
        size = ch[0] -> size + 1 + ch[1] -> size;
10
        return this;
11
12
   };
13
   typedef pair<Node*, Node*> Pair;
14
   Node *null. *root:
   void Node::downtag(){
     if(rev){
17
        for(int i = 0; i < 2; i++)
18
          if(ch[i] != null){
19
            ch[i] -> rev ^= 1;
20
            swap(ch[i] -> ch[0], ch[i] -> ch[1]);
21
          }
22
       rev = 0;
23
24
      if(tag){
25
        for(int i = 0; i < 2; i++)</pre>
          if(ch[i] != null){
26
27
            ch[i] -> key += tag;
            ch[i] -> mn += tag;
29
            ch[i] -> tag += tag;
30
          }
31
        tag = 0;
32
     }
33
34
35
     static int s = 3023192386;
36
     return (s += (s << 3) + 1) & (~0u >> 1);
37
   }
38
   bool random(int x, int y){
39
     return r() \% (x + y) < x;
40
   Node* merge(Node *p, Node *q){
41
42
     if(p == null) return q;
43
     if(q == null) return p;
     p -> downtag();
44
45
     q -> downtag();
     if(random(p -> size, q -> size)){
46
47
        p \to ch[1] = merge(p \to ch[1], q);
48
        return p -> update();
49
     }else{
50
        q -> ch[0] = merge(p, q -> ch[0]);
```

```
51
        return q -> update();
52
53 }
54 Pair split(Node *x, int n){
     if(x == null) return make_pair(null, null);
     x -> downtag();
56
     if(n \le x \rightarrow ch[0] \rightarrow size){
57
       Pair ret = split(x -> ch[0], n);
58
59
        x \rightarrow ch[0] = ret.second;
60
       return make_pair(ret.first, x -> update());
61
     }
62
     Pair ret = split(x \rightarrow ch[1], n - x \rightarrow ch[0] \rightarrow size -
         \hookrightarrow 1);
63
      x \rightarrow ch[1] = ret.first;
64
      return make_pair(x -> update(), ret.second);
65 }
   pair<Node*, Pair> get_segment(int 1, int r){
66
67
     Pair ret = split(root, l - 1);
68
     return make_pair(ret.first, split(ret.second, r - 1 +
69 }
70 | int main(){
     null = new Node(INF, INF, 0);
71
72
     null \rightarrow ch[0] = null \rightarrow ch[1] = null;
73
     root = null;
74 }
```

5.4 Splay

```
template<class T>void checkmin(T &x,T y)
1
2
  {
3
     if(y < x) x = y;
4
   }
5
   struct Node
6
     Node *c[2], *fa;
7
8
     int size, rev;
10
     LL val, add, min;
11
12
     Node *init(LL v)
13
       val = min = v:
14
       add = rev = 0;
15
16
       c[0] = c[1] = fa = NULL;
17
       size = 1;
18
19
       return this;
     }
20
21
     void rvs()
22
23
       std::swap(c[0], c[1]);
24
       rev ^= 1;
25
     }
     void inc(LL x)
26
27
28
       val += x:
29
       add += x;
30
       min += x;
     }
31
32
     void pushdown()
33
     {
       if(rev)
34
35
         if(c[0]) c[0]->rvs();
36
         if(c[1]) c[1]->rvs();
37
         rev = 0;
38
39
       }
40
       if(add)
41
42
         if(c[0]) c[0]->inc(add);
         if(c[1]) c[1]->inc(add);
43
         add = 0;
44
```

```
46
47
      void update()
48
49
        min = val:
        if(c[0]) checkmin(min, c[0]->min);
50
        if(c[1]) checkmin(min, c[1]->min);
51
52
53
        size = 1;
        if(c[0]) size += c[0]->size;
54
55
        if(c[1]) size += c[1]->size;
56
57
58
    } *root;
59
    Node* newnode(LL x)
60
61
62
      static Node pool[maxs], *p = pool;
63
      return (++p)->init(x);
65
    }
66
    void setc(Node *x,int t,Node *y)
67
68
      x->c[t] = y;
69
      if(y) y \rightarrow fa = x;
 70
 71
    Node *find(int k)
 72
73
      Node *now = root:
74
 75
      while(true)
76
77
        now->pushdown();
78
        int t = (now->c[0] ? now->c[0]->size : 0) + 1;
79
80
        if(t == k) break;
81
82
83
        if(t > k) now = now->c[0];
84
        else now = now->c[1], k -= t;
85
86
87
      return now;
88
89
    void rotate(Node *x,Node* &k)
90
      Node *y = x->fa, *z = y->fa;
91
92
93
      if(y != k) z->c[z->c[1] == y] = x;
94
      else k = x;
95
96
      x->fa = z;
97
      int i = (y->c[1] == x);
98
100
      setc(y, i, x->c[i ^ 1]);
101
      setc(x, i ^ 1, y);
102
103
      y->update(), x->update();
104
    void spaly(Node *x,Node* &k)
105
106
107
      static Node *st[maxs];
108
      int top = 0;
109
      Node *y, *z;
110
111
      y = x;
      while(y != k) st[++top] = y, y = y->fa;
112
113
      st[++top] = y;
114
      while(top) st[top]->pushdown(), top--;
115
116
```

```
while(x != k)
117
118
119
        y = x->fa, z = y->fa;
120
        if(y != k)
121
122
           if((y == z->c[1]) ^ (x == y->c[1])) rotate(x, k);
124
           else rotate(y, k);
125
126
127
        rotate(x, k);
128
129
130
    Node *subtree(int 1,int r)
131
      assert((++1) <= (++r));
132
      spaly(find(l - 1), root);
133
134
      spaly(find(r + 1), root->c[1]);
135
      return root->c[1]->c[0];
136
137 }
138
    void ins(int pos,int v)
139 {
140
      pos++;
141
      spaly(find(pos), root);
142
      spaly(find(pos + 1), root->c[1]);
143
      setc(root->c[1], 0, newnode(v));
      root->c[1]->update();
144
      root->update();
145
146 }
    void del(int pos)
147
148 {
149
150
      spaly(find(pos - 1), root);
      spaly(find(pos + 1), root->c[1]);
151
      root->c[1]->c[0] = NULL;
152
153
      root->c[1]->update();
      root->update();
154
155 }
156
    void init()
157
158
     root = newnode(0);
159
      setc(root, 1, newnode(0));
160
      root->update();
161 }
```

5.5 Link cut Tree

```
1 inline void reverse(int x) {
2
     tr[x].rev ^= 1; swap(tr[x].c[0], tr[x].c[1]);
3 }
4
5
   inline void rotate(int x, int k) {
6
     int y = tr[x].fa, z = tr[y].fa;
       tr[x].fa = z; tr[z].c[tr[z].c[1] == y] = x;
7
8
       tr[tr[x].c[k ^ 1]].fa = y; tr[y].c[k] = tr[x].c[k ^ 1]
9
       tr[x].c[k ^ 1] = y; tr[y].fa = x;
  ۱,
10
11
   inline void splay(int x, int w) {
12
     int z = x; pushdown(x);
13
     while (tr[x].fa != w) {
       int y = tr[x].fa; z = tr[y].fa;
15
16
       if (z == w) {
         pushdown(z = y); pushdown(x);
17
18
         rotate(x, tr[y].c[1] == x);
         update(y); update(x);
19
20
21
         pushdown(z); pushdown(y); pushdown(x);
22
         int t1 = tr[y].c[1] == x, t2 = tr[z].c[1] == y;
         if (t1 == t2) rotate(y, t2), rotate(x, t1);
23
24
         else rotate(x, t1), rotate(x, t2);
```

```
update(z); update(y); update(x);
26
       }
27
     }
28
     update(x);
29
     if (x != z) par[x] = par[z], par[z] = 0;
30
31
   inline void access(int x) {
33
     for (int y = 0; x; y = x, x = par[x]) {
34
       splay(x, 0);
35
       if (tr[x].c[1]) par[tr[x].c[1]] = x, tr[tr[x].c[1]].fa
36
       tr[x].c[1] = y; par[y] = 0; tr[y].fa = x; update(x);
37
38
39
40
   inline void makeroot(int x) {
41
     access(x); splay(x, 0); reverse(x);
42
43
44
   inline void link(int x, int y) {
45
     makeroot(x); par[x] = y;
46
47
48
   inline void cut(int x, int y) {
     access(x); splay(y, 0);
50
     if (par[y] != x) swap(x, y), access(x), splay(y, 0);
51
     par[y] = 0;
52
53
   inline void split(int x, int y) { // x will be the root
      \hookrightarrow of the tree
55
     makeroot(y); access(x); splay(x, 0);
56
```

5.6 树上莫队

```
void dfs(int u)
2
 3
     dep[u] = dep[fa[u][0]] + 1;
4
     for(int i = 1; i < logn; i++)</pre>
       fa[u][i] = fa[fa[u][i - 1]][i - 1];
5
6
     stk.push(u);
 8
     for(int i = 0; i < vec[u].size(); i++)</pre>
9
       int v = vec[u][i];
10
11
12
       if(v == fa[u][0]) continue;
13
14
        fa[v][0] = u, dfs(v);
15
        size[u] += size[v];
16
17
18
        if(size[u] >= bufsize)
19
20
          ++bcnt;
21
22
          while(stk.top() != u)
23
24
            block[stk.top()] = bcnt;
25
            stk.pop();
26
          }
27
28
          size[u] = 0;
29
       }
30
32
     size[u]++;
33
34
   void prework()
35 {
```

```
36
     dfs(1);
37
38
     ++bcnt;
39
     while(!stk.empty())
40
       block[stk.top()] = bcnt;
41
42
        stk.pop();
43
44
   }
   void rev(int u)
45
46
   {
47
     now -= (cnt[val[u]] > 0);
48
49
     if(used[u])
50
       cnt[val[u]]--;
51
       used[u] = false;
52
53
54
     else
55
       cnt[val[u]]++;
56
57
       used[u] = true;
58
59
60
     now += (cnt[val[u]] > 0);
61 }
62
   void move(int &x,int y,int z)
63 {
     int fwd = y;
64
65
66
     rev(getlca(x, z));
     rev(getlca(y, z));
67
68
     while(x != y)
69
70
71
       if(dep[x] < dep[y]) std::swap(x, y);</pre>
72
73
       rev(x), x = fa[x][0];
74
75
76
     x = fwd:
77 }
78
   void solve()
79
80
     std::sort(query + 1, query + m + 1);
81
82
     int L = 1, R = 1;
83
     rev(1);
84
85
     for(int i = 1; i <= m; i++)</pre>
86
87
       int 1 = query[i].u;
88
       int r = query[i].v;
89
90
       move(L, 1, R);
       move(R, r, L);
91
92
93
        ans[query[i].t] = now;
94
     }
95 }
```

5.7 CDQ 分治

```
struct Node

int x, y, z, idx;

friend bool operator == (const Node &a,const Node &b)

freturn a.x == b.x && a.y == b.y && a.z == b.z;

friend bool operator < (const Node &a,const Node &b)

friend bool operator < (const Node &a,const Node &b)

{</pre>
```

```
11
       return a.y < b.y;</pre>
12
13
14
   } triple[maxn];
15
   bool cmpx(const Node &a,const Node &b)
16
17
18
     if(a.x != b.x) return a.x < b.x;</pre>
19
     if(a.y != b.y) return a.y < b.y;</pre>
20
     return a.z < b.z;</pre>
21
22
   void solve(int l,int r)
23
24
     if(1 == r) return;
25
26
     int mid = (1 + r) >> 1;
27
28
     solve(1, mid);
29
30
     static std::pair<Node,int> Lt[maxn], Rt[maxn];
31
     int Ls = 0, Rs = 0;
32
33
34
     for(int i = 1; i <= mid; i++)</pre>
35
       Lt[++Ls] = std::make_pair(triple[i], i);
36
     for(int i = mid + 1; i <= r; i++)
37
       Rt[++Rs] = std::make_pair(triple[i], i);
38
     int pos = 1;
39
40
     std::sort(Lt + 1, Lt + Ls + 1);
41
42
     std::sort(Rt + 1, Rt + Rs + 1);
43
44
     backup.clear();
     for(int i = 1; i <= Rs; i++)</pre>
45
46
       while(pos <= Ls && !(Rt[i].first < Lt[pos].first))</pre>
47
48
49
          insert(Lt[pos].first.z, 1);
50
51
          pos++;
52
53
54
       f[Rt[i].second] += query(Rt[i].first.z);
55
56
57
     for(int i = 0; i < backup.size(); i++) pre[backup[i]] =</pre>
        58
59
     solve(mid + 1, r);
60
```

5.8 整体二分

```
1
   void solve(int l,int r,std::vector<int> q)
 2
 3
     if(1 == r || q.empty())
 4
 5
        for(int i = 0; i < q.size(); i++)</pre>
 6
 7
          ans[q[i]] = 1;
 8
        }
 9
     }
10
     else
11
       int mid = (1 + r) >> 1;
12
13
        backup.clear();
14
15
        for(int i = 1; i <= mid; i++)</pre>
16
17
          Event e = event[i];
18
```

```
19
20
          if(e.1 <= e.r)
21
22
            add(e.1, e.v);
            add(e.r + 1, -e.v);
23
24
25
26
          {
27
            add(1, e.v);
28
            add(e.r + 1, -e.v);
29
            add(e.1, e.v);
30
31
32
33
        std::vector<int> qL, qR;
34
        for(int i = 0; i < q.size(); i++)</pre>
35
36
37
          LL val = 0;
38
          for(int j = 0; j < vec[q[i]].size(); j++)</pre>
39
40
            val += count(vec[q[i]][j]);
41
42
43
            if(val >= p[q[i]]) break;
44
45
46
          if(cnt[q[i]] + val >= p[q[i]])
47
            qL.push_back(q[i]);
48
          }
49
50
          else
51
          {
52
            cnt[q[i]] += val;
53
            qR.push_back(q[i]);
54
55
56
57
        for(int i = 0; i < backup.size(); i++) sum[backup[i]]</pre>
           \hookrightarrow = 0:
58
        solve(1, mid, qL);
59
        solve(mid + 1, r, qR);
60
61 }
```

6. 图论

6.1 2-SAT

```
1 // n 必须是 2 的次幂
  void fft(Complex a[], int n, int f) {
    for (int i = 0; i < n; ++i)
3
4
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
5
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
7
       Complex w = Complex(1, 0);
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
10
         for (int k = 0; k < i; ++k) {
11
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
           a[j + k] = x + y; a[j + k + i] = x - y;
12
13
14
       }
     }
15
16 }
```

6.2 2-SAT (tarjan)

```
template < class TAT > void checkmin(TAT &x,TAT y)

if (y < x) x = y;

}
</pre>
```

```
5
   void tarjan(int u)
6
7
     dfn[u] = low[u] = ++dt;
     flag[u] = true;
8
9
     stk.push(u);
10
     for(int i = 0; i < vec[u].size(); i++)</pre>
11
12
13
       int v = vec[u][i];
14
15
       if(!dfn[v])
16
17
          tarjan(v);
18
          checkmin(low[u], low[v]);
19
20
        else if(flag[v])
21
22
          checkmin(low[u], dfn[v]);
23
24
     }
25
     if(low[u] == dfn[u])
26
27
28
        ++bcnt;
29
        while(stk.top() != u)
30
31
          block[stk.top()] = bcnt;
          flag[stk.top()] = false;
32
33
          stk.pop();
34
36
        block[u] = bcnt;
37
        flag[u] = false;
38
        stk.pop();
39
40
41
   bool solve()
42
43
        for(int i = 1; i <= 2 * n; i++)
         if(!dfn[i]) tarjan(i);
44
45
46
       bool ans = true;
        for(int i = 1; i <= n; i++)
48
49
          if(block[2 * i] == block[2 * i - 1])
50
          {
51
            ans = false;
52
            break;
53
54
55
        return ans;
56
```

6.3 KM

```
struct KM {
    // Truly O(n^3)
    // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时
      → 边权取负, 但不能连的还是 -INF, 使用时先对 1
      → -> n 调用 hungary() ,再 get_ans() 求值
    int w[N][N];
    int lx[N], ly[N], match[N], way[N], slack[N];
6
    bool used[N];
7
    void init() {
      for (int i = 1; i \le n; i++) {
8
9
        match[i] = 0;
10
        lx[i] = 0;
        ly[i] = 0;
11
12
        way[i] = 0;
13
    }
14
15
    void hungary(int x) {
```

```
16
       match[0] = x:
17
        int j0 = 0;
18
       for (int j = 0; j \le n; j++) {
         slack[j] = INF;
19
         used[j] = false;
20
21
22
23
       do {
24
         used[j0] = true;
25
         int i0 = match[j0], delta = INF, j1 = 0;
26
         for (int j = 1; j \le n; j++) {
            if (used[j] == false) {
27
28
              int cur = -w[i0][j] - lx[i0] - ly[j];
29
              if (cur < slack[j]) {</pre>
30
                slack[j] = cur;
31
                way[j] = j0;
              }
32
33
              if (slack[j] < delta) {</pre>
34
                delta = slack[j];
35
                j1 = j;
              }
36
           }
37
         }
38
39
          for (int j = 0; j \le n; j++) {
40
            if (used[j]) {
41
              lx[match[j]] += delta;
42
              ly[j] -= delta;
43
            else slack[j] -= delta;
44
45
46
         j0 = j1;
       } while (match[j0] != 0);
47
48
49
       do {
         int j1 = way[j0];
50
51
         match[j0] = match[j1];
52
          j0 = j1;
53
       } while (j0);
54
     }
55
56
     int get_ans() {
57
       int sum = 0;
58
       for(int i = 1; i <= n; i++) {
         if (w[match[i]][i] == -INF); // 无解
59
          if (match[i] > 0) sum += w[match[i]][i];
60
       }
61
62
       return sum;
     }
63
64
  } km;
```

6.4 点双连通分量

```
const bool BCC_VERTEX = 0, BCC_EDGE = 1;
2
   struct BCC { // N = NO + MO. Remember to call

→ init(&raw_graph).
3
     Graph *g, forest; // g is raw graph ptr.
4
     int dfn[N], DFN, low[N];
5
     int stack[N], top;
6
     int expand_to[N];
                            // Where edge i is expanded to in
        \hookrightarrow expaned graph.
7
     // Vertex i expaned to i.
8
     int compress_to[N]; // Where vertex i is compressed to.
     bool vertex_type[N], cut[N], compress_cut[N], branch[M];
9
     //std::vector<int> BCC_component[N]; // Cut vertex
        \hookrightarrow \texttt{belongs} to none.
11
     __inline void init(Graph *raw_graph) {
12
       g = raw_graph;
13
     }
     void DFS(int u, int pe) {
14
15
       dfn[u] = low[u] = ++DFN; cut[u] = false;
16
       if (!~g->adj[u]) {
         cut[u] = 1;
17
18
         compress_to[u] = forest.new_node();
```

```
19
         compress_cut[compress_to[u]] = 1;
20
21
       for (int e = g->adj[u]; ~e; e = g->nxt[e]) {
22
         int v = g - v[e];
         if ((e ^ pe) > 1 && dfn[v] > 0 && dfn[v] < dfn[u]) {
23
24
           stack[top++] = e;
           low[u] = std::min(low[u], dfn[v]);
25
26
27
         else if (!dfn[v]) {
28
           stack[top++] = e; branch[e] = 1;
29
           DFS(v, e);
30
           low[u] = std::min(low[v], low[u]);
31
           if (low[v] >= dfn[u]) {
32
              if (!cut[u]) {
33
                cut[u] = 1;
34
                compress_to[u] = forest.new_node();
35
                compress_cut[compress_to[u]] = 1;
36
37
             int cc = forest.new_node();
38
             forest.bi ins(compress to[u], cc):
39
             compress_cut[cc] = 0;
40
              //BCC_component[cc].clear();
41
             do {
42
                int cur_e = stack[--top];
43
                compress_to[expand_to[cur_e]] = cc;
44
                compress_to[expand_to[cur_e^1]] = cc;
45
                if (branch[cur_e]) {
                  int v = g->v[cur_e];
46
                  if (cut[v])
47
48
                    forest.bi_ins(cc, compress_to[v]);
49
                    //BCC_component[cc].push_back(v);
50
51
                    compress_to[v] = cc;
52
                  }
53
54
             } while (stack[top] != e);
55
56
57
       }
58
59
     void solve() {
60
       forest.init(g->base);
       int n = g->n;
61
       for (int i = 0; i < g->e; i++) {
63
         expand_to[i] = g->new_node();
64
65
       memset(branch, 0, sizeof(*branch) * g->e);
66
       memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0;
67
       for (int i = 0; i < n; i++)
68
         if (!dfn[i + g->base]) {
69
           top = 0;
70
           DFS(i + g->base, -1);
         }
71
     }
72
73
   } bcc;
74
75
   bcc.init(&raw_graph);
76
   bcc.solve();
   // Do something with bcc.forest ...
```

6.5 边双连诵分量

```
struct BCC {
    Graph *g, forest;
    int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N],
3
       4
    // tot[] is the size of each BCC, belong[] is the BCC
       \hookrightarrow that each node belongs to
5
    pair<int, int > ori[M]; // bridge in raw_graph(raw node)
6
    bool is_bridge[M];
7
     __inline void init(Graph *raw_graph) {
8
      g = raw_graph;
```

```
9
        memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
10
        memset(vis + g \rightarrow base, 0, sizeof(*vis) * g \rightarrow n);
11
     }
12
     void tarjan(int u, int from) {
       dfn[u] = low[u] = ++dfs_clock; vis[u] = 1;
13

    stack[++top] = u;
14
        for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
          if ((p ^ 1) == from) continue;
15
16
          int v = g -> v[p];
17
          if (vis[v]) {
18
            if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
19
          } else {
20
            tarjan(v, p);
21
            low[u] = min(low[u], low[v]);
22
            if (low[v] > dfn[u]) is_bridge[p / 2] = true;
23
        }
24
25
        if (dfn[u] != low[u]) return;
26
        tot[forest.new_node()] = 0;
27
28
          belong[stack[top]] = forest.n;
29
          vis[stack[top]] = 2;
          tot[forest.n]++;
30
31
          --top;
32
       } while (stack[top + 1] != u);
33
     }
34
     void solve() {
35
       forest.init(g -> base);
        int n = g \rightarrow n;
36
        for (int i = 0; i < n; ++i)
37
38
          if (!vis[i + g -> base]) {
            top = dfs_clock = 0;
39
40
            tarjan(i + g \rightarrow base, -1);
41
       for (int i = 0; i < g -> e / 2; ++i)
42
          if (is_bridge[i]) {
43
44
            int e = forest.e;
            forest.bi_ins(belong[g -> v[i * 2]], belong[g ->
45
               \hookrightarrow v[i * 2 + 1]], g \rightarrow w[i * 2]);
            ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i *
46
            ori[e + 1] = make_pair(g -> v[i * 2], g -> v[i * 2]
47
               \hookrightarrow + 1]);
48
49
     }
50
   } bcc;
```

6.6 最小树形图

```
const int MAXN,INF;// INF >= sum( W_ij )
   int from[MAXN + 10][MAXN * 2 + 10],n,m,edge[MAXN +
      \hookrightarrow 10] [MAXN * 2 + 10];
 3 int sel[MAXN * 2 + 10], fa[MAXN * 2 + 10], vis[MAXN * 2 +
   int getfa(int x){if(x == fa[x]) return x; return fa[x] =
      \hookrightarrow getfa(fa[x]);}
 5
   void liuzhu(){ // 1-base: root is 1, answer = (sel[i], i)
      \hookrightarrow for i in [2..n]
     fa[1] = 1;
 6
 7
     for(int i = 2; i <= n; ++i){</pre>
 8
        sel[i] = 1; fa[i] = i;
        for(int j = 1; j \le n; ++j) if(fa[j] != i)
 9
10
           if(from[j][i] = i, edge[sel[i]][i] > edge[j][i])
             \hookrightarrow sel[i] = j;
11
      int limit = n;
12
13
        int prelimit = limit; memset(vis, 0, sizeof(vis));
14
           \hookrightarrow vis[1] = 1;
        for(int i = 2; i <= prelimit; ++i) if(fa[i] == i &&</pre>
15

    !vis[i]){
16
          int j = i; while(!vis[j]) vis[j] = i, j =
             \hookrightarrow \text{getfa(sel[j])};
```

```
if(j == 1 || vis[j] != i) continue; vector<int> C;
17
             \hookrightarrow int k = j;
18
          do C.push_back(k), k = getfa(sel[k]); while(k != j);
19
          ++limit:
20
          for(int i = 1; i <= n; ++i){</pre>
            edge[i][limit] = INF, from[i][limit] = limit;
21
22
23
          fa[limit] = vis[limit] = limit;
24
          for(int i = 0; i < int(C.size()); ++i){</pre>
25
            int x = C[i], fa[x] = limit;
26
            for(int j = 1; j \le n; ++j)
27
              if(edge[j][x] != INF && edge[j][limit] >
                 \hookrightarrow \text{edge[j][x] - edge[sel[x]][x])}
28
                 edge[j][limit] = edge[j][x] - edge[sel[x]][x];
29
                from[j][limit] = x;
30
31
32
          for(int j=1; j<=n;++j) if(getfa(j)==limit)</pre>
             sel[limit] = 1;
33
34
          for(int j = 1; j \le n; ++j)
35
            if(edge[sel[limit]][limit] > edge[j][limit])

    sel[limit] = i:
36
37
        if(prelimit == limit) break;
38
39
     for(int i = limit; i > 1; --i) sel[from[sel[i]][i]] =
        \hookrightarrow sel[i];
40
```

6.7 带花树

```
vector<int> link[maxn];
   int n,match[maxn],Queue[maxn],head,tail;
   int pred[maxn],base[maxn],start,finish,newbase;
   bool InQueue[maxn], InBlossom[maxn];
   void push(int u){ Queue[tail++]=u;InQueue[u]=true; }
   int pop(){ return Queue[head++]; }
   int FindCommonAncestor(int u,int v){
     bool InPath[maxn];
9
     for(int i=0;i<n;i++) InPath[i]=0;</pre>
10
     while(true){ u=base[u];InPath[u]=true;if(u==start)
        ⇔ break;u=pred[match[u]]; }
     while(true){ v=base[v];if(InPath[v])
        ⇔ break; v=pred[match[v]]; }
12
     return v;
   }
13
   void ResetTrace(int u){
14
15
     while(base[u]!=newbase){
16
17
       v=match[u];
18
       InBlossom[base[u]]=InBlossom[base[v]]=true;
19
       u=pred[v]:
20
       if(base[u]!=newbase) pred[u]=v;
21
22
   void BlossomContract(int u,int v){
23
24
     newbase=FindCommonAncestor(u,v);
     for (int i=0;i<n;i++)</pre>
25
26
     InBlossom[i]=0:
27
     ResetTrace(u);ResetTrace(v);
     if(base[u]!=newbase) pred[u]=v;
     if(base[v]!=newbase) pred[v]=u;
29
30
     for(int i=0;i<n;++i)</pre>
31
     if(InBlossom[base[i]]){
32
       base[i]=newbase:
33
       if(!InQueue[i]) push(i):
34
35
36
   bool FindAugmentingPath(int u){
37
     bool found=false:
     for(int i=0;i<n;++i) pred[i]=-1,base[i]=i;</pre>
38
```

```
39
     for (int i=0;i<n;i++) InQueue[i]=0;</pre>
     start=u;finish=-1; head=tail=0; push(start);
40
41
     while(head<tail){
42
       int u=pop();
       for(int i=link[u].size()-1;i>=0;i--){
43
         int v=link[u][i];
44
         if (base[u]!=base[v]&&match[u]!=v)
45
           if(v==start||(match[v]>=0&&pred[match[v]]>=0))
46
47
             BlossomContract(u,v);
48
           else if(pred[v]==-1){
49
             pred[v]=u;
             if(match[v]>=0) push(match[v]);
50
51
             else{ finish=v; return true; }
52
53
       }
     7
54
     return found:
55
56 }
57
   void AugmentPath(){
     int u=finish,v,w;
59
     while(u>=0){

    v=pred[u]; w=match[v]; match[v]=u; match[u]=v; u=w; }

60 }
61
   void FindMaxMatching(){
62
    for(int i=0;i<n;++i) match[i]=-1;</pre>
     for(int i=0;i<n;++i) if(match[i]==-1)</pre>
        64 }
```

6.8 支配树

```
vector<int> prec[N], succ[N];
   vector<int> ord;
3
   int stamp, vis[N];
4 int num[N];
5 int fa[N];
6 void dfs(int u) {
7
    vis[u] = stamp;
    num[u] = ord.size();
9
     ord.push_back(u);
10
     for (int i = 0; i < (int)succ[u].size(); ++i) {</pre>
11
       int v = succ[u][i];
       if (vis[v] != stamp) {
12
13
         fa[v] = u;
14
          dfs(v);
15
16
     }
  1
17
   int fs[N], mins[N], dom[N], sem[N];
18
  int find(int u) {
19
20
     if (u != fs[u]) {
21
       int v = fs[u];
22
       fs[u] = find(fs[u]);
       if (mins[v] != -1 && num[sem[mins[v]]] <</pre>
23
          \hookrightarrow num[sem[mins[u]]]) {
24
         mins[u] = mins[v];
25
     }
26
27
     return fs[u];
28 }
   void merge(int u, int v) { fs[u] = v; }
29
30 | vector<int> buf[N];
31 int buf2[N];
32 void mark(int source) {
33
    ord.clear();
     ++stamp:
34
     dfs(source):
35
     for (int i = 0; i < (int)ord.size(); ++i) {</pre>
36
37
       int u = ord[i]:
38
       fs[u] = u, mins[u] = -1, buf2[u] = -1;
39
     for (int i = (int) \text{ ord. size}() - 1; i > 0; --i) {
40
41
       int u = ord[i], p = fa[u];
```

```
42
        sem[u] = p;
43
        for (int j = 0; j < (int)prec[u].size(); ++j) {</pre>
44
          int v = prec[u][j];
45
          if (use[v] != stamp) continue;
46
          if (num[v] > num[u]) {
47
            find(v); v = sem[mins[v]];
          }
48
49
          if (num[v] < num[sem[u]]) {</pre>
50
            sem[u] = v:
51
          }
        }
52
53
        buf[sem[u]].push_back(u);
54
       mins[u] = u;
55
        merge(u, p);
56
        while (buf[p].size()) {
          int v = buf[p].back();
57
58
          buf[p].pop_back();
59
          find(v);
          if (sem[v] == sem[mins[v]]) {
            dom[v] = sem[v];
62
          } else {
            buf2[v] = mins[v];
63
64
          }
65
       }
66
67
     dom[ord[0]] = ord[0];
68
     for (int i = 0; i < (int)ord.size(); ++i) {</pre>
       int u = ord[i];
69
       if (~buf2[u]) {
70
71
          dom[u] = dom[buf2[u]];
72
73
     }
74
```

6.9 无向图最小割

```
int cost[maxn] [maxn], seq[maxn], len[maxn], n, m, pop, ans;
 2
   bool used[maxn];
 3
   void Init(){
     int i,j,a,b,c;
5
     for(i=0;i<n;i++) for(j=0;j<n;j++) cost[i][j]=0;</pre>
6
     for(i=0:i<m:i++){</pre>
 7
       scanf("%d %d %d",&a,&b,&c); cost[a][b]+=c;
           \hookrightarrow cost[b][a]+=c;
8
9
     pop=n; for(i=0;i<n;i++) seq[i]=i;
   }
10
   void Work(){
11
     ans=inf; int i,j,k,l,mm,sum,pk;
12
     while(pop > 1){
14
       for(i=1;i<pop;i++) used[seq[i]]=0; used[seq[0]]=1;</pre>
15
       for(i=1;i<pop;i++) len[seq[i]]=cost[seq[0]][seq[i]];</pre>
16
        pk=0; mm=-inf; k=-1;
        for(i=1;i<pop;i++) if(len[seq[i]] > mm){
17
           18
        for(i=1;i<pop;i++){</pre>
          used[seq[l=k]]=1;
19
20
          if(i==pop-2) pk=k;
21
         if(i==pop-1) break;
22
          mm=-inf:
23
          for(j=1;j<pop;j++) if(!used[seq[j]])</pre>
24
            if((len[seq[j]]+=cost[seq[1]][seq[j]]) > mm)
25
              mm=len[seq[j]], k=j;
26
        s11m=0:
27
       for(i=0;i<pop;i++) if(i != k)</pre>
28

    sum+=cost[seq[k]][seq[i]];

29
        ans=min(ans.sum):
30
        for(i=0;i<pop;i++)</pre>
          cost[seq[k]][seq[i]]=cost[seq[i]][seq[k]]+=cost[seq[pk]][s
31
32
        seq[pk]=seq[--pop];
33
```

```
34 | printf("%d\n",ans);
35 |}
```

6.10 最大团搜索

```
1
   const. int. N = 1000 + 7:
   vector<vector<bool> > adj;
2
3
   class MaxClique {
       const vector<vector<bool> > adi:
5
       const int n;
6
       vector<int> result, cur_res;
7
       vector<vector<int> > color set:
8
       const double t_limit; // MAGIC
9
     int para, level;
     vector<pair<int, int> > steps;
10
11
   public:
12
       class Vertex {
       public:
13
            int i, d;
14
15
            Vertex(int i, int d = 0) : i(i), d(d) {}
16
17
       void reorder(vector<Vertex> &p) {
18
            for (auto &u : p) {
                u.d = 0:
19
20
                for (auto v : p) u.d += adj[v.i][u.i];
21
22
            sort(p.begin(), p.end(), [&](const Vertex &a,
               \hookrightarrow const Vertex &b) { return a.d > b.d; } );
23
       }
24
     // reuse p[i].d to denote the maximum possible clique
        \hookrightarrow for first i vertices.
25
       void init_color(vector<Vertex> &p) {
26
            int maxd = p[0].d;
            for (int i = 0; i < p.size(); i++) p[i].d = min(i,</pre>
27
               \hookrightarrow maxd) + 1;
28
29
       bool bridge(const vector<int> &s, int x) {
30
            for (auto v : s) if (adj[v][x]) return true;
31
            return false;
32
33
     // approximate estimate the p[i].d
     // Do not care about first mink color class (For better
34
        \hookrightarrow result, we must get some vertex in some color class
        → larger than mink )
35
       void color_sort(vector<Vertex> &cur) {
36
            int totc = 0, ptr = 0, mink =
               for (int i = 0; i < cur.size(); i++) {</pre>
37
                int x = cur[i].i, k = 0;
38
39
                while (k < totc && bridge(color_set[k], x))</pre>
                   40
                if (k == totc) color_set[totc++].clear();
41
                color_set[k].push_back(x);
42
                if (k < mink) cur[ptr++].i = x;</pre>
43
            }
            if (ptr) cur[ptr - 1].d = 0;
            for (int i = mink; i < totc; i ++) {</pre>
45
46
                for (auto v : color_set[i]) {
                     cur[ptr++] = Vertex(v, i + 1);
47
                }
48
            }
49
50
        void expand(vector<Vertex> &cur) {
51
52
        steps[level].second = steps[level].second -
           \hookrightarrow \mathtt{steps}\, \texttt{[level]}. \texttt{first} \,\, + \,\, \mathtt{steps}\, \texttt{[level - 1]}. \texttt{first};
        steps[level].first = steps[level - 1].second;
53
            while (cur.size()) {
54
                if (cur_res.size() + cur.back().d <=</pre>
55

    result.size()) return ;

56
                int x = cur.back().i:
                cur_res.push_back(x); cur.pop_back();
57
58
                vector<Vertex> remain;
```

```
for (auto v : cur) {
60
                     if (adj[v.i][x]) remain.push_back(v.i);
61
                if (remain.size() == 0) {
62
                     if (cur_res.size() > result.size()) result
63
                        } else {
65
            // Magic ballance.
66
            if (1. * steps[level].second / ++para < t_limit)</pre>

→ reorder(remain);
67
                     color_sort(remain);
68
            steps[level++].second++:
69
                     expand(remain);
70
            level--;
71
                }
72
                cur_res.pop_back();
            }
73
74
       }
75
   public:
       MaxClique(const vector<vector<bool> > &_adj, int n,
          \rightarrow double tt = 0.025) : adj(_adj), n(n), t_limit(tt)
            result.clear():
77
78
            cur_res.clear();
79
            color_set.resize(n);
80
        steps.resize(n + 1);
81
        fill(steps.begin(), steps.end(), make_pair(0, 0));
        level = 1;
82
       para = 0;
83
84
        vector<int> solve() {
            vector<Vertex> p;
86
            for (int i = 0; i < n; i++)
               \hookrightarrow \texttt{p.push\_back(Vertex(i));}
88
            reorder(p);
89
            init_color(p);
90
            expand(p);
91
            return result;
92
93
   };
```

6.11 弦图判定

```
// n 必须是 2 的次幂
   void fft(Complex a[], int n, int f) {
 2
 3
     for (int i = 0; i < n; ++i)
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
 5
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
 6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
       Complex w = Complex(1, 0);
7
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
10
         for (int k = 0; k < i; ++k) {
11
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
       }
15
     }
16
```

6.12 斯坦纳树

```
void SPFA(int *dist)
{
    static int line[maxn + 5];
    static bool hash[maxn + 5];
    int f = 0, r = 0;

for(int i = 1; i <= N; i++)
    if(dist[i] < inf)
}</pre>
```

```
10
                 line[r] = i;
11
                hash[i] = true;
12
                r = (r + 1) \% (N + 1);
13
14
        while(f != r)
15
16
17
            int t = line[f];
18
            hash[t] = false;
            f = (f + 1) \% (N + 1);
19
20
21
            for(int i = head[t]; i ; i = edge[i].next)
22
23
                 int v = edge[i].v, dt = dist[t] + edge[i].w;
24
                 if(dt < dist[v])</pre>
25
26
                     dist[v] = dt;
27
28
                     if(!hash[v])
29
30
                         if(dist[v] < dist[line[f]])</pre>
31
32
                          {
33
                              f = (f + N) \% (N + 1);
34
                              line[f] = v;
35
                         }
36
                          else
37
                          {
                              line[r] = v:
38
                              r = (r + 1) \% (N + 1);
39
40
41
42
                         hash[v] = true;
                     }
43
                }
44
45
            }
46
47
   }
48
   void solve()
49
        for(int i = 1; i <= S; i++)
50
51
52
            for(int j = 1; j \le N; j++)
                for(int k = (i - 1) & i; k ; k = (k - 1) & i)
53
54
                     G[i][j] = std::min(G[i][j], G[k][j] + G[k]
                        → ^ i][j]);
55
56
            SPFA(G[i]);
57
58
   }
```

6.13 虚树

```
1
  bool cmp(const int lhs,const int rhs)
2
  1
3
    return dfn[lhs] < dfn[rhs];</pre>
4
   }
5
   void build()
6
7
     std::sort(h + 1, h + 1 + m, cmp);
8
9
     int top = 0;
10
11
     for (int i = 1; i <= m; i++)
12
       if (!top) father[st[++top] = h[i]] = 0;
13
14
       else
15
       {
            int p = h[i], lca = LCA(h[i],st[top]);
16
17
            while(d[st[top]] > d[lca])
18
19
                if (d[st[top - 1]] <= d[lca])</pre>
20
```

```
father[st[top]] = lca;
21
22
23
                top--;
            }
24
25
26
            if (st[top] != lca)
27
28
                t[++tot] = lca;
29
                father[lca] = st[top];
30
                st[++top] = lca;
            }
31
32
33
            father[p] = lca;
34
            st[++top] = p;
35
36
   }
37
```

6.14 点分治

```
template<class TAT>void checkmax(TAT &x,TAT y)
2
3
     if(x < y) x = y;
4
   }
   template<class TAT>void checkmin(TAT &x,TAT y)
5
 6
7
     if(y < x) x = y;
8
   }
9
   void getsize(int u,int fa)
10
     size[u] = 1;
11
12
     smax[u] = 0;
13
     for(int i = 0; i < G[u].size(); i++)</pre>
14
15
16
       int v = G[u][i];
17
18
       if(v == fa || ban[v]) continue;
19
20
       getsize(v, u);
21
       size[u] += size[v];
22
23
       checkmax(smax[u], size[v]);
24
25
26
   int getroot(int u,int ts,int fa)
27
     checkmax(smax[u], ts - size[u]);
28
29
30
     int res = u;
31
32
     for(int i = 0; i < G[u].size(); i++)</pre>
33
34
       int v = G[u][i];
35
       if(v == fa || ban[v]) continue;
36
37
38
       int w = getroot(v, ts, u);
39
40
       if(smax[w] < smax[res]) res = w;</pre>
     }
41
42
43
     return res;
44
45
   void solve()
46
47
     static int line[maxn];
     static std::vector<int> vec;
48
49
     int f = 0, r = 0;
50
     line[r++] = 1;
51
52
```

```
while(f != r)
53
54
55
        int u = line[f++];
56
        getsize(u, 0);
57
        u = getroot(u, size[u], 0);
58
59
60
        ban[u] = true;
61
        vec.clear();
62
        for(int i = 0; i < G[u].size(); i++)</pre>
63
          if(!ban[G[u][i]]) vec.push_back(G[u][i]);
64
65
66
67
        do something you like...
68
69
70
71
72
        for(int i = 0; i < vec.size(); i++)</pre>
          line[r++] = vec[i];
73
74
     }
75 }
```

6.15 最小割最大流

```
bool BFS()
 2
   {
3
        for(int i = 1; i <= ind; i++) dep[i] = 0;</pre>
 4
5
        dep[S] = 1, line.push(S);
 6
 7
        while(!line.empty())
 8
            int now = line.front();
g
10
            line.pop();
11
            for(int i = head[now], p; i ; i = edge[i].next)
12
13
                if(edge[i].cap && !dep[p = edge[i].v])
14
                     dep[p] = dep[now] + 1, line.push(p);
15
       }
16
        if (dep[T])
17
18
        {
19
            for(int i = 1; i <= ind; i++)</pre>
20
                cur[i] = head[i];
21
            return true;
        }
22
23
        else
24
            return false;
25 }
26
   int DFS(int a,int flow)
27
   {
28
        if(a == T) return flow;
29
30
        int ret = 0;
31
32
        for(int &i = cur[a], p; i ; i = edge[i].next)
            if(dep[p = edge[i].v] == dep[a] + 1 &&
33
               \hookrightarrow \texttt{edge[i].cap)}
34
35
                int ff = DFS(p, std::min(flow, edge[i].cap));
36
37
                flow -= ff, edge[i].cap -= ff;
38
                ret += ff, edge[i ^ 1].cap += ff;
39
40
                if(!flow) break:
41
            }
42
43
            return ret;
44 }
45
   int solve()
46 {
```

```
47    int totflow = 0;
48
49    while(BFS())
50    {
51        totflow += DFS(S, INF);
52    }
53
54    return totflow;
55 }
```

6.16 最小费用流

```
bool SPFA()
 2
   {
 3
        static int line[maxv];
        static bool hash[maxv];
 4
        register int f = 0, r = 0;
 5
6
 7
     for(int i = 1; i <= ind; i++)</pre>
8
9
          dist[i] = inf;
          from[i] = 0;
10
11
12
        dist[S] = 0, line[r] = S, r = (r + 1) \% maxv;
13
14
        hash[S] = true;
15
16
        while(f != r)
17
18
            int x = line[f];
19
20
            line[f] = 0, f = (f + 1) % maxv;
21
            hash[x] = false;
22
23
            for(int i = head[x]; i; i = edge[i].next)
24
                 if(edge[i].cap)
26
                     int v = edge[i].v;
27
                     int w = dist[x] + edge[i].cost;
28
29
                     if(w < dist[v])</pre>
30
                     {
31
                          dist[v] = w;
32
                          from[v] = i;
33
                          if(!hash[v])
34
35
                               if(f != r && dist[v] <=</pre>
36
                                  \hookrightarrow dist[line[f]])
                                   f = (f - 1 + maxv) \% maxv,
37
                                      \hookrightarrow line[f] = v;
38
                               else
39
                                   line[r] = v, r = (r + 1) %
                                      \hookrightarrow maxv:
40
                               hash[v] = true;
41
42
                          }
                     }
43
44
                 }
45
46
47
        return from[T];
48
   }
49
50
   int back(int x,int flow)
51
52
     if(from[x])
53
        flow = back(edge[from[x] ^ 1].v, std::min(flow,
54

    edge[from[x]].cap));
55
56
        edge[from[x]].cap -= flow;
```

```
edge[from[x] ^ 1].cap += flow;
57
58
59
60
     return flow;
61 }
   int solve()
62
64
       int mincost = 0, maxflow = 0;
65
       while(SPFA())
66
67
       {
            int flow = back(T, inf);
68
69
70
            mincost += dist[T] * flow;
71
            maxflow += flow:
72
73
74
       return mincost;
75 }
```

6.17 zkw 费用流

```
1 int S, T, totFlow, totCost;
2
3 int dis[N], slack[N], visit[N];
4
5
   int modlable () {
6
       int delta = INF;
7
       for (int i = 1; i <= T; i++) {
8
            if (!visit[i] && slack[i] < delta) delta =</pre>
              → slack[i]:
9
           slack[i] = INF;
10
11
       if (delta == INF) return 1;
       for (int i = 1; i <= T; i++)
12
           if (visit[i]) dis[i] += delta;
13
14
       return 0:
15 }
16
17
   int dfs (int x, int flow) {
18
       if (x == T) {
           totFlow += flow;
19
           totCost += flow * (dis[S] - dis[T]);
20
21
           return flow;
23
       visit[x] = 1;
24
       int left = flow;
       for (int i = e.last[x]; ~i; i = e.succ[i])
25
           if (e.cap[i] > 0 && !visit[e.other[i]]) {
26
27
                int y = e.other[i];
                if (dis[y] + e.cost[i] == dis[x]) {
28
29
                    int delta = dfs (y, min (left, e.cap[i]));
30
                    e.cap[i] -= delta;
                    e.cap[i ^ 1] += delta;
31
                    left -= delta;
32
33
                    if (!left) { visit[x] = 0; return flow; }
34
35
                    slack[y] = min (slack[y], dis[y] +
                       \hookrightarrow e.cost[i] - dis[x]);
                }
36
           }
37
38
       return flow - left;
  }
39
40
41
   pair <int, int> minCost () {
       totFlow = 0; totCost = 0;
42
43
       fill (dis + 1, dis + T + 1, 0);
44
       do {
45
46
                fill (visit + 1, visit + T + 1, 0);
           } while (dfs (S, INF));
47
       } while (!modlable ()):
48
       return make_pair (totFlow, totCost);
49
```

```
50 }
```

6.18 最小割树

```
#include<iostream>
   #include<cstdio>
   #include<cstdlib>
   #include<cstring>
   #include<algorithm>
   #include<queue>
   #define inf 0x3f3f3f3f
   #define N 155
   using namespace std;
10
11
      \hookrightarrow cnt,n,m,dis[N],last[N],a[N],tmp[N],ans[N][N],s,t,mark[N];
12
   struct edge{int to,c,next;}e[N*200];
   queue <int> q;
13
14
15
   void addedge(int u,int v,int c)
16
17
           \rightarrow e[++cnt].to=v;e[cnt].c=c;e[cnt].next=last[u];last[\mu]=cn
18
           \rightarrow e[++cnt].to=u;e[cnt].c=c;e[cnt].next=last[v];last[v]=cn
19
20
   bool bfs()
21
22
   {
23
        memset(dis.0.sizeof(dis)):
24
        dis[s]=2;
25
        while (!q.empty()) q.pop();
26
        q.push(s);
27
        while (!q.empty())
28
29
            int u=q.front();
            q.pop();
31
            for (int i=last[u];i;i=e[i].next)
32
                if (e[i].c&&!dis[e[i].to])
33
                {
                     dis[e[i].to]=dis[u]+1;
34
35
                     if (e[i].to==t) return 1;
36
                     q.push(e[i].to);
37
38
39
        return 0;
40
41
42
   int dfs(int x,int maxf)
43
44
        if (x==t||!maxf) return maxf;
45
        int ret=0;
46
        for (int i=last[x];i;i=e[i].next)
47
            if (e[i].c&&dis[e[i].to]==dis[x]+1)
48
                int f=dfs(e[i].to,min(e[i].c,maxf-ret));
49
50
                e[i].c-=f;
51
                e[i<sup>1</sup>].c+=f;
52
                ret+=f:
53
                if (ret==maxf) break;
55
        if (!ret) dis[x]=0;
56
        return ret;
57
58
59
   void dfs(int x)
60
   {
        mark[x]=1;
61
62
        for (int i=last[x];i;i=e[i].next)
63
            if (e[i].c&&!mark[e[i].to]) dfs(e[i].to);
   }
64
65
```

```
void solve(int l,int r)
 67
    {
 68
         if (l==r) return;
 69
         s=a[1]:t=a[r]:
         for (int i=2;i<=cnt;i+=2)</pre>
 70
             e[i].c=e[i^1].c=(e[i].c+e[i^1].c)/2;
 71
 72
 73
         while (bfs()) flow+=dfs(s,inf);
 74
         memset(mark,0,sizeof(mark));
 75
         dfs(s);
 76
         for (int i=1;i<=n;i++)</pre>
             if (mark[i])
 77
 78
                  for (int j=1; j<=n; j++)
 79
                       if (!mark[j])
 80
                              \rightarrow ans[i][j]=ans[j][i]=min(ans[i][j],flow);
 81
         int i=1.i=r:
 82
         for (int k=1;k<=r;k++)</pre>
 83
             if (mark[a[k]]) tmp[i++]=a[k];
             else tmp[j--]=a[k];
 85
         for (int k=1;k<=r;k++)</pre>
 86
             a[k]=tmp[k];
         solve(1,i-1);
 87
 88
         solve(j+1,r);
 89
    }
 90
 91
    int main()
    {
 92
 93
         int cas:
         scanf("%d",&cas);
 94
 95
         while (cas--)
 96
 97
             scanf("%d%d",&n,&m);
 98
             cnt=1;
             for (int i=1:i<=n:i++)
 99
100
                  a[i]=i;
101
             memset(last,0,sizeof(last));
             memset(ans,inf,sizeof(ans));
102
103
              for (int i=1;i<=m;i++)</pre>
104
              {
105
                  int x,y,z;
                  scanf("%d%d%d",&x,&y,&z);
106
107
                  addedge(x,y,z);
             }
108
109
              solve(1,n);
110
             int q;
             scanf("%d",&q);
111
112
             for (int i=1;i<=q;i++)</pre>
113
114
                  int x,tot=0;
115
                  scanf("%d",&x);
116
                  for (int i=1;i<n;i++)</pre>
117
                       for (int j=i+1; j<=n; j++)</pre>
                           if (ans[i][j]<=x) tot++;</pre>
118
119
                  printf("%d\n",tot);
              }
120
121
              cout << end1;
122
123
         return 0;
124
```

6.19 上下界网络流建图

B(u,v) 表示边 (u,v) 流量的下界,C(u,v) 表示边 (u,v) 流量的上界,F(u,v) 表示边 (u,v) 的流量。设 G(u,v)=F(u,v)-B(u,v),显然有

$$0 \le G(u, v) \le C(u, v) - B(u, v)$$

6.19.1 无源汇的上下界可行流

建立超级源点 S^* 和超级汇点 T^* ,对于原图每条边 (u,v) 在新网络中连如下三条边: $S^* \to v$,容量为 B(u,v);

 $u \to T^*$, 容量为 B(u,v); $u \to v$, 容量为 C(u,v) - B(u,v)。最后求新网络的最大流,判断从超级源点 S^* 出发的边是否都满流即可,边 (u,v) 的最终解中的实际流量为 G(u,v) + B(u,v)。

6.19.2 有源汇的上下界可行流

从汇点 T 到源点 S 连一条上界为 ∞ ,下界为 0 的边。按照**无源汇的上下界可行流**一样做即可,流量即为 $T \to S$ 边上的流量。

6.19.3 有源汇的上下界最大流

- 1. 在**有源汇的上下界可行流**中,从汇点 T 到源点 S 的 边改为连一条上界为 ∞ ,下届为 x 的边。x 满足二分性质,找到最大的 x 使得新网络存在**无源汇的上下界可行流**即为原图的最大流。
- 2. 从汇点 T 到源点 S 连一条上界为 ∞ ,下界为 0 的 边,变成无源汇的网络。按照**无源汇的上下界可行流**的方法,建立超级源点 S^* 和超级汇点 T^* ,求一遍 $S^* \to T^*$ 的最大流,再将从汇点 T 到源点 S 的这条边拆掉,求一次 $S \to T$ 的最大流即可。

6.19.4 有源汇的上下界最小流

- 1. 在**有源汇的上下界可行流**中,从汇点 T 到源点 S 的 边改为连一条上界为 x,下界为 0 的边。x 满足二分性质,找到最小的 x 使得新网络存在**无源汇的上下界可行流**即为原图的最小流。
- 2. 按照**无源汇的上下界可行流**的方法,建立超级源点 S^* 与超级汇点 T^* ,求一遍 $S^* \to T^*$ 的最大流,但是注意这一次不加上汇点 T 到源点 S 的这条边,即不使之改为无源汇的网络去求解。求完后,再加上那条汇点 T 到源点 S 上界 ∞ 的边。因为这条边下界为0,所以 S^* , T^* 无影响,再直接求一次 $S^* \to T^*$ 的最大流。若超级源点 S^* 出发的边全部满流,则 $T \to S$ 边上的流量即为原图的最小流,否则无解。

7. 其他

7.1 Dancing Links

7.1.1 精确覆盖

```
#pragma comment(linker, "/STACK:1024000000,1024000000")
   #include<iostream>
   #include<cstdio>
   #include<cstring>
   #include<algorithm>
   #include<map>
   #include<queue>
8
   #include<set>
Q
   #include<cmath>
10
   #include<bitset>
   #define mem(a,b) memset(a,b,sizeof(a))
   #define lson i<<1,1,mid
12
13
   #define rson i<<1|1,mid+1,r</pre>
14
   #define llson j<<1,1,mid
15
   #define rrson j<<1|1,mid+1,r</pre>
   #define INF 0x7fffffff
   #define maxn 1000005
   typedef long long 11;
   typedef unsigned long long ull;
20
   using namespace std;
   int head.sz:
21
22
   int U[maxn],D[maxn],L[maxn],R[maxn];//上下左右链表指针
23
   int H[maxn],ROW[maxn],C[maxn],S[maxn],O[maxn];
   void remove(int c)
24
25
   {
26
       L[R[c]]=L[c]:
       R[L[c]]=R[c]:
27
28
       for(int i=D[c]; i!=c; i=D[i])
```

```
29
             for(int j=R[i]; j!=i; j=R[j])
30
31
                 U[D[j]]=U[j];
                 D[U[j]]=D[j];
32
                 --S[C[j]];
33
34
35
   }
 36
    void resume(int c)
37
        for(int i=U[c]; i!=c; i=U[i])
38
39
40
             for(int j=L[i]; j!=i; j=L[j])
41
42
                 ++S[C[j]];
43
                 U[D[j]]=j;
                 D[U[j]]=j;
44
45
46
        L[R[c]]=c;
47
        R[L[c]]=c;
48
   }
49
    void init(int m)//m 是列
50
    ł
51
        head=0;//头指针为 0
52
53
        for(int i=0; i<=m; i++)</pre>
54
55
             U[i]=i;
             D[i]=i;//建立双向十字链表
56
57
             L[i]=i-1;
58
             R[i]=i+1;
59
             S[i]=0;
60
        }
        R[m] = 0:
61
        L[0]=m;
62
63
        S[0]=INF+1;
        sz=m+1;
65
        memset(H,0,sizeof(H));
66
   | }
67
    void insert(int i, int j)
68
    {
         if(H[i])
69
70
             L[sz] = L[H[i]];
71
             R[sz] = H[i];
72
73
             L[R[sz]] = sz;
             R[L[sz]] = sz;
74
75
76
        else
77
         {
78
             L[sz] = sz;
79
             R[sz] = sz;
             H[i] = sz;
80
81
        }
        U[sz] = U[j];
82
        D[sz] = j;
83
84
        U[D[sz]] = sz;
        D[U[sz]] = sz;
85
        C[sz] = j;
86
87
        ROW[sz] = i;
88
         ++S[j];
89
         ++sz;
90
   }
91
    bool dfs(int k,int len)
92
    {
93
        if (R[head] ==head)
94
95
             sort(0,0+len*len);
96
             int p=0;
             for(int i=0; i<len; i++)</pre>
97
98
99
                 for(int j=0; j<len; j++)</pre>
100
                      int num=0[p++];
101
```

```
num=num-(i*len+j)*len;
102
103
                     printf("%d",num);
104
                 7
                 puts("");
105
            }
106
107
            return true:
        }
108
109
        int s=INF,c;
110
        for (int t=R[head]; t!=head; t=R[t])
111
            if (S[t] <s) s=S[t],c=t;</pre>
112
        remove(c):
113
        for(int i=D[c]; i!=c; i=D[i])
114
        {
115
            0[k]=ROW[i];
            for(int j=R[i]; j!=i; j=R[j])
116
117
                 remove(C[i]);
            if(dfs(k+1.len))
118
119
                 return true;
120
            for(int j=L[i]; j!=i; j=L[j])
121
                 resume(C[j]);
        }
122
123
        resume(c);
124
        return false;
125
126
    void calc(int i,int j,int k,int len)
127
128
        int r=(i*len+j-1)*len+k;
        int base=sqrt(len);
129
        //第 i 行有数字 k
130
        insert(r,i*len+k);
131
        //第 j 列有数字 k
132
        insert(r,len*len+(j-1)*len+k);
133
        //第 k 块有数字 k
134
        int block=(j-1)/base*base+i/base;
135
        insert(r,len*len*2+block*len+k);
136
        //第 i 行 j 列有一个数字 (限制一个出格子只填一
137
           →个数)
        insert(r,len*len*3+i*len+j);
138
   }
139
    void build(char s[][10],int len)//len 表示是几宫数独
140
141
142
        int i,j,k;
143
        init(len*len*4);
        for(i=0; i<len; i++)</pre>
144
145
            for(j=1; j<=len; j++)</pre>
146
            {
147
                 if(s[i][j-1]=='0')
                     for(k=1; k<=len; k++)</pre>
148
149
                         calc(i,j,k,len);
150
                 else calc(i,j,s[i][j-1]-'0',len);
151
            }
152
    }
153
    int main()
    {
154
        //freopen("1.txt","r",stdin);
155
156
        int t;
157
        cin>>t;
        while(t--)
158
159
160
            char s[10][10];
161
            for(int i=0; i<9; i++)</pre>
162
                 scanf("%s",s[i]);
163
            build(s,9);
            dfs(0,9);//从根开始搜
164
        }
165
166
        return 0;
167
```

7.1.2 重复覆盖

```
Problem : 2295 ( Radar ) Judge Status : Accepted
RunId : 4355553 Language : G++ Author : zhuyawei
```

```
3 | Code Render Status : Rendered By HDOJ G++ Code Render
      # include<stdio.h>
5
  # include<math.h>
6 # include<string.h>
7
  # define eps 1e-8
8 # define N 55
9 # define V 3600
10 int n.m.K:
11 int L[V],R[V];
12 int D[V],U[V];
13
  int C[V];
   int S[N],H[N];
14
15
   int ak,size;
16
   double dis(double x1, double y1, double x2, double y2)
17
       return sqrt((x2-x1)*(x2-x1) + (y2-y1)*(y2-y1));
18
19 }
   void Link(int r,int c)
20
21
       S[c]++;C[size]=c;
22
       U[size]=U[c];D[U[c]]=size;
23
       D[size]=c;U[c]=size;
24
25
       if(H[r]==-1) H[r]=L[size]=R[size]=size;
26
27
28
           L[size]=L[H[r]];R[L[H[r]]]=size;
           R[size]=H[r];L[H[r]]=size;
29
       }
30
       size++;
31
  }
32
   void remove(int c)
34
   {
35
       int i;
       for(i=D[c];i!=c;i=D[i])
36
37
           L[R[i]]=L[i],R[L[i]]=R[i];
38
   }
39
   void resume(int c)
40
41
       int i;
       for(i=U[c];i!=c;i=U[i])
42
43
           L[R[i]]=R[L[i]]=i;
  }
44
45
   int h()
46
   {
       int i,j,k,count=0;
47
48
       bool visit[N];
49
       memset(visit,0,sizeof(visit));
50
       for(i=R[0];i;i=R[i])
51
52
            if(visit[i]) continue;
53
           count++:
54
           visit[i]=1;
           for(j=D[i];j!=i;j=D[j])
55
56
                for(k=R[j];k!=j;k=R[k])
57
58
                    visit[C[k]]=1;
59
            }
60
61
       return count:
62
   }
   void Dance(int k)
63
64
   {
65
       int i,j,c,Min,ans;
       ans=h():
66
       if(k+ans>K || k+ans>=ak) return;
67
68
       if(!R[0])
69
70
            if(k<ak) ak=k;</pre>
71
           return;
72
73
       for(Min=N,i=R[0];i;i=R[i])
74
           if(S[i]<Min) Min=S[i],c=i;</pre>
```

```
for(i=D[c];i!=c;i=D[i])
 75
 76
 77
             remove(i);
             for(j=R[i];j!=i;j=R[j])
 78
 79
                 remove(j);
 80
             Dance(k+1):
             for(j=L[i];j!=i;j=L[j])
 81
 82
                  resume(j);
 83
             resume(i);
 84
         7
 85
         return:
 86
 87
    int main()
 88
 89
         int i,j,ncase;
         double x[N],y[N],x1[N],y1[N];
 90
 91
         double left, right, ans, mid;
 92
         scanf("%d",&ncase);
         while(ncase--)
 94
             scanf("%d%d%d",&n,&m,&K);
 95
 96
             for(i=1;i<=n;i++)
 97
                  scanf("%lf%lf",&x[i],&y[i]);
 98
             for(i=1;i<=m;i++)</pre>
 99
                  scanf("%lf%lf",&x1[i],&y1[i]);
100
             left=0;
101
             right=1416.0;
102
             ans=right;
103
             while(right>=left)
104
                  for(i=0;i<=n;i++)</pre>
105
106
107
                       S[i]=0;
108
                       U[i]=D[i]=i;
                       L[i+1]=i;R[i]=i+1;
109
110
                  R[n]=0;
111
                  memset(H,-1,sizeof(H));
112
                  size=n+1;
113
                  mid=(left+right)/2;
114
                  for(i=1;i<=m;i++)</pre>
115
                       for(j=1; j<=n; j++)</pre>
116
117
                           if (mid>=dis(x1[i],y1[i],x[j],y[j]))
                              \hookrightarrow Link(i,j);
                  }
118
                  ak=N:
119
                  Dance(0);
120
121
                  if(ak<=K) {ans=mid<ans?mid:ans;right=mid-eps;}</pre>
122
                  else left=mid+eps;
123
124
             printf("%.61f\n",ans);
         }
125
126
         return 0;
127
```

7.1.3 斜率优化

```
#include<set>
   #include<map>
 3
  #include<ctime>
  #include<queue>
  #include<cmath>
  #include<cstdio>
   #include<vector>
8
  #include<cstring>
9
   #include<cstdlib>
10
   #include<iostream>
11
   #include<algorithm>
   #define inf 900000000000000000LL
12
13
   #define mp make_pair
14
   #define pa pair<ll,int>
   #define ll long long
15
16 using namespace std;
```

```
17 int read()
18
   {
19
      int x=0,f=1;char ch=getchar();
      while(ch<'0'||ch>'9'){if(ch=='-')f=-1;ch=getchar();}
20
     \label{lem:while(ch>='0'&&ch<='9')} $$ x=x*10+ch-'0'; ch=getchar(); $$
21
     return x*f:
22
23 }
24 int n,K;
25 | 11 sum[100005],f[100005],g[100005];
26
   int a[100005],q[100005];
27
   double cal(int j,int k)
28 | {
29
         \hookrightarrow (double)(sum[k]*sum[k]-sum[j]*sum[j]+g[j]-g[k])/(double)(\hat{s}u
   }
30
   void tran(int x)
31
   {
32
33
      int head=1,tail=0;
34
     for(int i=x;i<=n;i++)</pre>
35
        while(head<tail&&cal(q[tail-1],q[tail])>cal(q[tail],i-1))tail4+yoid tran(int x)
36
37
        q[++tail]=i-1;
38

→ while (head < tail & & cal (q [head], q [head+1]) < sum [i]) head++47
</p>
39
        int t=q[head];
40
        f[i]=g[t]+(sum[i]-sum[t])*sum[t];
41
      for(int i=x;i<=n;i++)swap(f[i],g[i]);</pre>
42
   ۱,
43
   void dp()
44
45
46
        for(int i=1;i<=K;i++)</pre>
47
        tran(i);
     printf("%lld\n",g[n]);
48
49
   }
50
   int main()
51
   {
52
     n=read();K=read();
53
     for(int i=1;i<=n;i++)a[i]=read();</pre>
54
     int top=0;
     for(int i=1;i<=n;i++)if(a[i]!=0)a[++top]=a[i];</pre>
55
56
57
     for(int i=1;i<=n;i++)</pre>
58
       sum[i]=sum[i-1]+a[i];
59
      dp();
60
      return 0;
61 }
```

7.1.4 决策单调性

```
#include<set>
  #include<map>
3 | #include<ctime>
   #include<queue>
5
   #include<cmath>
6
   #include<cstdio>
   #include<vector>
   #include<cstring>
8
   #include<cstdlib>
10 #include<iostream>
11 #include <algorithm>
12 #define inf 900000000000000000LL
13 #define mp make_pair
14 | #define pa pair<11,int>
15 #define 11 long long
16 using namespace std;
17
   int read()
18
   {
19
       int x=0,f=1;char ch=getchar();
20
       while(ch<'0'||ch>'9'){if(ch=='-')f=-1;ch=getchar();}
21
       while(ch>='0'&&ch<='9'){x=x*10+ch-'0'; ch=getchar();}
       return x*f:
22
23 }
```

```
24 int n.K:
25
   ll a[100005],sum[100005],f[100005],g[100005];
   struct data{
        int 1,r,p;
   }a[100005]:
28
   ll cal(int i,int j)
29
30
31
        return g[i]+(sum[j]-sum[i])*sum[i];
32
   }
33
   int find(data t,int q)
34
   {
35
        int l=t.1,r=t.r,mid;
36
        while(1<=r)
   [k]-s(um[j]);
38
            mid=(1+r)>>1;
39
            if(cal(q,mid)>cal(t.p,mid))r=mid-1;
            else l=mid+1;
40
41
        return 1;
   }
43
45
46
        int head=1.tail=0:
        q[++tail]=(data)\{0,n,x-1\};
48
        for(int i=x;i<=n;i++)</pre>
49
50
            if(i>q[head].r)head++;
51
            f[i]=cal(q[head].p,i);
            if(tail<head||cal(i,n)>cal(q[tail].p,n))
52
53

→ while(head<=tail&&cal(i,q[tail].1)>cal(q[tail].
55
                     tail--;
                 if (head<=tail)</pre>
56
57
                 {
58
                     int t=find(q[tail],i);
59
                     q[tail].r=t-1;
                     q[++tail]=(data){t,n,i};
60
61
62
                 else q[++tail]=(data){i,n,i};
            }
63
64
65
        for(int i=x;i<=n;i++)swap(f[i],g[i]);</pre>
66
   }
67
   void dp()
68
   {
69
        for(int i=1;i<=K;i++)</pre>
70
            tran(i);
71
        printf("%lld\n",g[n]);
72
73
   int main()
74
75
        n=read(); K=read();
        for(int i=1;i<=n;i++)a[i]=read(),sum[i]=sum[i-1]+a[i];</pre>
76
77
        dp();
78
        return 0;
79
```

7.2 蔡勒公式

```
int zeller(int y,int m,int d) {
   if (m<=2) y--,m+=12; int c=y/100; y%=100;
   int w=((c>>2)-(c<<1)+y+(y>>2)+(13*(m+1)/5)+d-1)%7;
   if (w<0) w+=7; return(w);
}</pre>
```

7.3 五边形数定理

$$p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k-1} p(n - \frac{k(3k-1)}{2})$$

```
#include<iostream>
   #include<cstdio>
3
   using namespace std;
4 #define LL __int64
5 const int N=100005;
6 const int MOD=1000000007;
7
  LL dp[N],fi[N];
8 LL five(LL x) { return (3*x*x-x)/2; }
   //五边形数
Q
   void wbxs(){
10
       dp[0]=1;
11
       int t=1000; //其实可以等于 sqrt(N)
12
       for(int i=-t;i<=t;++i)</pre>
13
14
           fi[i+t]=five(i);
       for(int i=1;i<=100000;++i){
15
16
           int flag=1;
17
           for(int j=1;;++j){
                LL a=fi[j+t],b=fi[-j+t];
18
                if(a>i && b>i) break;
19
20
                if(a<=i) dp[i]=(dp[i]+dp[i-a]*flag+MOD)%MOD;</pre>
21
                if(b<=i) dp[i]=(dp[i]+dp[i-b]*flag+MOD)%MOD;</pre>
22
                flag*=-1;
           }
23
       }
24
25
   }
26
   int main(){
27
       wbxs();
       int T,n;
28
       scanf("%d",&T);
29
       while(T--){
30
           scanf("%d",&n);
31
           printf("%I64d\n",dp[n]);
32
33
       }
34
       return 0;
35
  | }
```

7.4 凸包闵可夫斯基和

```
// cv[0..1] 为两个顺时针凸包, 其中起点等于终点, 求
    →出的闵可夫斯基和不一定是严格凸包
  int i[2] = {0, 0}, len[2] = {(int)cv[0].size() - 1,
    3
 vector<P> mnk;
 mnk.push_back(cv[0][0] + cv[1][0]);
4
5
 do {
6
   int d((cv[0][i[0] + 1] - cv[0][i[0]]) * (cv[1][i[1] + 1]
      \hookrightarrow - cv[1][i[1]]) >= 0);
   mnk.push_back(cv[d][i[d] + 1] - cv[d][i[d]] +
      \rightarrow mnk.back());
   i[d] = (i[d] + 1) \% len[d];
 } while(i[0] || i[1]);
```

8. 技巧

8.1 STL 归还空间

```
template <typename T>
__inline void clear(T& container) {
  container.clear(); // 或者删除了一堆元素
  T(container).swap(container);
}
```

8.2 大整数取模

```
4 return t < 0 ? t + MODN : t;
5 }</pre>
```

8.3 读入优化

```
// getchar() 读入优化 << 关同步 cin << 此优化
   // 用 isdigit() 会小幅变慢
   // 返回 false 表示读到文件尾
   namespace Reader {
       const int L = (1 << 15) + 5;
       char buffer[L], *S, *T;
       __inline bool getchar(char &ch) {
           if (S == T) {
8
9
               T = (S = buffer) + fread(buffer, 1, L, stdin);
10
               if (S == T) {
           ch = EOF;
11
12
           return false;
         }
13
           }
14
15
       ch = *S++;
16
       return true;
17
18
       __inline bool getint(int &x) {
19
       char ch; bool neg = 0;
       for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^=
20
          \hookrightarrow ch == '-';
       if (ch == EOF) return false;
       x = ch - '0';
23
       for (; getchar(ch), ch >= '0' && ch <= '9'; )
         x = x * 10 + ch - '0';
24
       if (neg) x = -x;
25
26
       return true;
27
28
```

8.4 二次随机法

```
#include <random>
int main() {
    std::mt19937 g(seed); // std::mt19937_64
    std::cout << g() << std::endl;
}</pre>
```

8.5 vimrc

```
set number
   set smartindent
   set autoindent
   set tabstop=4
   set softtabstop=4
   set shiftwidth=4
   set hlsearch
   set incsearch
   set autoread
11
   set backspace=2
12
   set mouse=a
13
14
   syntax on
   nmap <C-A> ggVG
16
17
   vmap <C-C> "+v
18
  filetype plugin indent on
   autocmd FileType cpp set cindent
   autocmd FileType cpp map <F9> :!g++ % -o %< -g -std=c++11
      \hookrightarrow -Wall -Wextra -Wconversion && size %< <CR>
   autocmd FileType cpp map <C-F9> :!g++ % -o %< -std=c++11
      \hookrightarrow -02 && size %< <CR>
```

```
autocmd FileType cpp map <F8> :!time ./%< < %<.in <CR>
   autocmd FileType cpp map <F5> :!time ./%< <CR>
26
  map <F3> :vnew %<.in <CR>
27
28 m
```

控制 cout 输出实数精度

1 | std::cout << std::fixed << std::setprecision(5);</pre>

让 make 支持 c++11

export CXXFLAGS='-std=c++11 -Wall'

8.8 tuple 相关

```
//
mytuple = std::make_tuple (10, 2.6, 'a');
   \hookrightarrow packing values into tuple
std::tie (myint, std::ignore, mychar) = mytuple;
   \hookrightarrow unpacking tuple into variables
std::get<I>(mytuple) = 20;
                                                             // get
std::cout << std::get<I>(mytuple) << std::endl;</pre>
   \hookrightarrow the Ith(const) element
```

9. 提示

9.1 线性规划转对偶

maximize $\mathbf{c}^T\mathbf{x}$ $\begin{array}{l} \text{maximize } \mathbf{c}^T \mathbf{x} \\ \text{subject to } \mathbf{A} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0 \end{array} \Longleftrightarrow \begin{array}{l} \text{minimize } \mathbf{y}^T \mathbf{b} \\ \text{subject to } \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T, \mathbf{y} \geq 0 \end{array}$

NTT 素数及其原根

Prime	Primitive root
1053818881	7
1051721729	6
1045430273	3
1012924417	5
1007681537	3

9.3 积分表

9.3.1
$$ax^2 + bx + c(a > 0)$$

1.
$$\int \frac{\mathrm{d}x}{ax^2 + bx + c} = \begin{cases} \frac{2}{\sqrt{4ac - b^2}} \arctan \frac{2ax + b}{\sqrt{4ac - b^2}} + C & (b^2 < 4ac) \\ \frac{1}{\sqrt{b^2 - 4ac}} \ln \left| \frac{2ax + b - \sqrt{b^2 - 4ac}}{2ax + b + \sqrt{b^2 - 4ac}} \right| + C & (b^2 > 4ac) \end{cases}$$

2.
$$\int \frac{x}{ax^2 + bx + c} dx = \frac{1}{2a} \ln|ax^2 + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^2 + bx + c}$$

9.3.2 $\sqrt{\pm ax^2 + bx + c}$ (a > 0)

1.
$$\int \frac{\mathrm{d}x}{\sqrt{ax^2 + bx + c}} = \frac{1}{\sqrt{a}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

2.
$$\int \sqrt{ax^2 + bx + c} dx = \frac{\frac{2ax+b}{4a}\sqrt{ax^2 + bx + c}}{\frac{4ac-b^2}{8\sqrt{a^3}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C} + C$$

3.
$$\int \frac{x}{\sqrt{ax^2 + bx + c}} dx = \frac{1}{a} \sqrt{ax^2 + bx + c} - \frac{b}{2\sqrt{a^3}} \ln|2ax + b + 2\sqrt{a} \sqrt{ax^2 + bx + c}| + C$$

4.
$$\int \frac{\mathrm{d}x}{\sqrt{c+bx-ax^2}} = -\frac{1}{\sqrt{a}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

5.
$$\int \sqrt{c + bx - ax^2} dx = \frac{2ax - b}{4a} \sqrt{c + bx - ax^2} + \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$

6.
$$\int \frac{x}{\sqrt{c+bx-ax^2}} dx = -\frac{1}{a} \sqrt{c+bx-ax^2} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{a^3}} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{b^3+4ac}} + \frac{b}{2\sqrt{a^3}} +$$

9.3.3
$$\sqrt{\pm \frac{x-a}{x-b}}$$
 或 $\sqrt{(x-a)(x-b)}$

1.
$$\int \frac{\mathrm{d}x}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C \ (a < b)$$

$$\int \sqrt{(x-a)(b-x)} dx = \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C, (a < b) \quad (1)$$

9.3.4 三角函数的积分

- 1. $\int \tan x dx = -\ln|\cos x| + C$
- 2. $\int \cot x dx = \ln|\sin x| + C$
- 3. $\int \sec x dx = \ln \left| \tan \left(\frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$
- 4. $\int \csc x dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x \cot x \right| + C$
- 5. $\int \sec^2 x dx = \tan x + C$
- 6. $\int \csc^2 x dx = -\cot x + C$
- 7. $\int \sec x \tan x dx = \sec x + C$
- 8. $\int \csc x \cot x dx = -\csc x + C$
- 9. $\int \sin^2 x dx = \frac{x}{2} \frac{1}{4} \sin 2x + C$
- 10. $\int \cos^2 x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$
- 11. $\int \sin^n x dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x dx$
- 12. $\int \cos^n x dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x dx$
- 13. $\int \frac{\mathrm{d}x}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{\mathrm{d}x}{\sin^{n-2} x}$
- 14. $\int \frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$

$$\int \cos^m x \sin^n x dx$$

$$= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x dx$$

$$= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x dx$$

16.
$$\int \frac{\mathrm{d}x}{a+b\sin x} = \begin{cases} \frac{2}{\sqrt{a^2 - b^2}} \arctan \frac{a\tan \frac{x}{2} + b}{\sqrt{a^2 - b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2 - a^2}} \ln \left| \frac{a\tan \frac{x}{2} + b - \sqrt{b^2 - a^2}}{a\tan \frac{x}{2} + b + \sqrt{b^2 - a^2}} \right| + C & (a^2 < b^2) \end{cases}$$
17.
$$\int \frac{\mathrm{d}x}{a+b\cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan \left(\sqrt{\frac{a-b}{a+b}} \tan \frac{x}{2}\right) + C & (a^2 > b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln \left| \frac{\tan \frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan \frac{x}{2} - \sqrt{\frac{a+b}{b-a}}} \right| + C & (a^2 < b^2) \end{cases}$$

7.
$$\int \frac{dx}{a+b\cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan\left(\sqrt{\frac{a-b}{a+b}} \tan\frac{x}{2}\right) + C & (a^2 > b) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln\left|\frac{\tan\frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan\frac{x}{2} - \sqrt{\frac{a+b}{b-a}}}\right| + C & (a^2 < b) \end{cases}$$

- 18. $\int \frac{\mathrm{d}x}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan\left(\frac{b}{a} \tan x\right) + C$
- 19. $\int \frac{dx}{a^2 \cos^2 x b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x a} \right| + C$
- 20. $\int x \sin ax dx = \frac{1}{a^2} \sin ax \frac{1}{a} x \cos ax + C$
- 21. $\int x^2 \sin ax dx = -\frac{1}{a}x^2 \cos ax + \frac{2}{a^2}x \sin ax + \frac{2}{a^3}\cos ax + C$
- 22. $\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{1}{a} x \sin ax + C$
- 23. $\int x^2 \cos ax dx = \frac{1}{a}x^2 \sin ax + \frac{2}{a^2}x \cos ax \frac{2}{a^3} \sin ax + C$

9.3.5 反三角函数的积分 (其中 a > 0)

- 1. $\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 x^2} + C$
- 2. $\int x \arcsin \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 x^2} + C$
- 3. $\int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{2} \arcsin \frac{x}{a} + \frac{1}{9}(x^2 + 2a^2)\sqrt{a^2 x^2} + C$
- 4. $\int \arccos \frac{x}{a} dx = x \ \arccos \frac{x}{a} \sqrt{a^2 x^2} + C$
- 5. $\int x \arccos \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arccos \frac{x}{a} \frac{x}{4} \sqrt{a^2 x^2} + C$
- 6. $\int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} \frac{1}{9}(x^2 + 2a^2)\sqrt{a^2 x^2} + C$
- 7. $\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} \frac{a}{2} \ln(a^2 + x^2) + C$
- 8. $\int x \arctan \frac{x}{a} dx = \frac{1}{2}(a^2 + x^2) \arctan \frac{x}{a} \frac{a}{2}x + C$
- 9. $\int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} \frac{a}{6}x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C$

9.3.6 指数函数的积分

$$1. \int a^x dx = \frac{1}{\ln a} a^x + C$$

1.
$$\int a \, \mathrm{d}x = \frac{1}{\ln a} a + C$$

2.
$$\int e^{ax} \, \mathrm{d}x = \frac{1}{a} a^{ax} + C$$

3.
$$\int_{0}^{a} x e^{ax} dx = \frac{1}{a^{2}} (ax - 1)a^{ax} + C$$

4.
$$\int x^n e^{ax} dx = \frac{1}{a} x^n e^{ax} - \frac{n}{a} \int x^{n-1} e^{ax} dx$$

5.
$$\int xa^x dx = \frac{x}{\ln a}a^x - \frac{1}{(\ln a)^2}a^x + C$$

6.
$$\int x^n a^x dx = \frac{1}{\ln a} x^n a^x - \frac{n}{\ln a} \int x^{n-1} a^x dx$$

7.
$$\int e^{ax} \sin bx dx = \frac{1}{a^2 + b^2} e^{ax} (a \sin bx - b \cos bx) + C$$

8.
$$\int e^{ax} \cos bx dx = \frac{1}{a^2 + b^2} e^{ax} (b \sin bx + a \cos bx) + C$$

9.
$$\int e^{ax} \sin^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \sin^{n-1} bx (a \sin bx - nb \cos bx) + \frac{n(n-1)b^2}{a^2 + b^2 n^2} \int e^{ax} \sin^{n-2} bx dx$$

10.
$$\int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) + \frac{n(n-1)b^2}{a^2 + b^2 n^2} \int e^{ax} \cos^{n-2} bx dx$$

9.3.7 对数函数的积分

1.
$$\int \ln x dx = x \ln x - x + C$$

$$2. \int \frac{\mathrm{d}x}{x \ln x} = \ln \left| \ln x \right| + C$$

3.
$$\int x^n \ln x dx = \frac{1}{n+1} x^{n+1} (\ln x - \frac{1}{n+1}) + C$$

4.
$$\int (\ln x)^n dx = x(\ln x)^n - n \int (\ln x)^{n-1} dx$$

5.
$$\int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n - \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$$