

Figure 1: This is a caption for this figure.

### Abstract

This is the abstract for my master's thesis on Certified Security by Design (CSBD) and Access-control logic (ACL) because ACL is so cool and effective...because effective is cool. Here is some text for formatting.

At least 18 people have been killed and dozens trapped in the Indian city of Varanasi after a flyover collapsed, crushing vehicles beneath it. The flyover was still being built when portions of its cement structure fell on the road being used under it. Officials from the National Disaster Response Force said 18 bodies had been recovered so far. A rescue operation is continuing for those believed to still be trapped, but their number and condition is unknown. Photographs and video from the scene showed cars and a bus crushed beneath the weight of the concrete, many of which still held people inside. Local media reported that a handful of people had been successfully rescued, as seven cranes attempted to lift the concrete pillar. A large crowd also gathered at the scene. One eyewitness told reporters they were nearby when the collapse happened. "At least four cars, an auto-rickshaw and a minibus were crushed under it," they said.

India's NDTV also reported that many of those trapped are believed to be construction workers who had been building the flyover. The cause of the collapse is not yet known, and an inquiry has been ordered, NDTV added. Major collapses of buildings and other infrastructure are not uncommon in India, where the enforcement of construction standards is weaker than many Western countries. In September, 33 people died when a six-storey Mumbai building toppled and more than 20 people died in 2016 when a flyover collapsed in Kolkata. Other collapses with smaller death tolls are frequent. Varanasi is the home constituency of India's Prime Minister Narendra Modi, who said he was "extremely saddened by the loss of lives due to the collapse". "I pray that the injured recover soon. Spoke to officials and asked them to ensure all possible support to those affected," he tweeted.

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# List of Acronyms

 $\mathbf{ACL}$  access-control logic.

 $\mathbf{CSBD}$  Certified Security by Design.

 $\mathbf{WFF}$  well-formed formula.

# Introduction

Some text here.[1] testing citations from the references.

### 1.1 Motivation

abelsec:intro:motivation

### 1.1.1 Systems Are Everywhere

### 1.1.2 CIA: Confidentiality, Integrity, and Accountability

### 1.2 This Master Thesis

This master thesis describes a method for designing secure systems. The method is called Certified Security by Design (CSBD). CSBD has been successfully demonstrated

on automated systems such as ... and .... But, until this research, it has not been demonstrated on non-automated, human-centered systems.

Systems span the range of fully automated to fully non-automated. This master thesis focuses on one end of this range: non-automated, human-centered systems.

The question addressed in this master thesis is whether CSBD could be successfully applied to non-automated, human-centered systems. An example of a non-automated, human-centered system is the patrol base operations defined in the United States Army Ranger Handbook[2]. Patrol base operations exemplify a non-automated, human-centered system wherein security is critical to mission success. In this master thesis, the results of applying CSBD to patrol base operations is discussed.

The patrol base operations are also an example of a predefined system. This means that this thesis also addresses the question of whether CSBD could be successfully applied to a pre-designed, non-automated, human-centered system. This is important because many such systems in use today are already designed and implemented. This means that a method for verifying and documenting the security properties of current, in-use systems is needed.

The successful application of CSBD to patrol base operations also suggests its use in combining automation with human-centered systems. The approach employed by this master thesis involves describing the patrol base operations as a hierarchy of secure state machines. This hierarchy has the property that it is easy to demonstrate security properties of the system, which is the goal of CSBD. But, it also has the property that it describes the patrol base operations as a system that is amiable to automation. Such automations of pre-defined non-automated, human-centered systems could include, for example, accountability systems for tracking supplies and personel. In the not-so-distant future, the miliatary, in particular, will most likely seek tracking and accountability systems for pre-existing, non-automated military operations. These

systems, like all security-sensitive military systems, should be designed according to NIST 800-160 standards. These standards require the formal verification and documentation provided by CSBD and demonstrated in this thesis.

# Background

Formal Methods

**Functional Programming** 

Higher Order Logic (HOL) Interactive Theorem Prover

Other Interactive Theorem Provers

# Systems Security Engineering

### 3.1 Systems

## 3.2 Systems Engineering

This is some text with figures. There is a figure here and text referencing figure?. More text before the figure call.

### 3.2.1 ISO/IEC/IEEE 15288

### 3.3 Systems Security Engineering

### 3.3.1 NIST Special Publication 800-160

### 3.3.1.1 Systems Security Engineering Framework

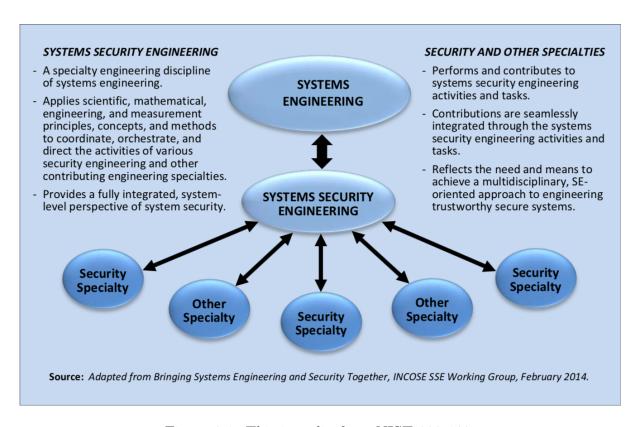


Figure 3.1: This is a clip from NIST 800-160.

This is some text after the figure.

3.3.2	Trustworthiness
3.3.2.1	Complete Mediation
3.3.3	Verification
3.3.4	Documentation
3.3.5	Reproducibility
3.4	Verification & Documentation
3.5	Principle of Complete Mediation
3.5.1	Formal Verification Using Computer-Aided Reasoning

Certified Security by Design (CSBD) &

# Access-Control Logic (ACL)

- 4.1 Certified Security by Design (CSBD)
- 4.1.1 The Principle of Complete Mediation
- 4.2 Access-Control Logic (ACL)
- 4.2.1 Motivation for Using ACL
- 4.2.2 ACL: A Command and Control (C2) Calculus

This section discusses the access-control logic in sufficient detail to understand the research reported in this master thesis. The material is adapted from *Access Control*, *Security*, and *Trust: A Logical Approach*[1]. For more indepth coverage of the ACL, read the text. Any references to "the text" in this section refer to the aforementioned text book.

### 4.2.3 Principals

Principals should be thought of as actors in the access-control logic. Principals can make statements or requests. They can be assigned privileges or authority over objects or actions. The text defines allowable principals using the identifier **Princ**:

### Princ ::= PName / Princ & Princ / Princ | Princ

This is a recursive definition. **PName** refers to the name of a principal (i.e., Jane, PlatoonLeader, sensor1). **Princ & Princ** is read "Princ with Princ" or "Princ and Princ" (i.e., Principal1 with Principal2). **Princ** | **Princ** is read as "Princ quoting Princ" (i.e., Principal1 quoting Principal2).

# 4.2.4 Propositional Variables, Requests, Authority, and Jurisdiction

To reason about access-control and trust, the ACL uses propositional variables, requests, authority, and jurisdiction to make statements.

Propositions in logic are assertions that are either true or false. For example, "I am reading this master thesis" is a proposition because either you are or you are not reading this. Propositional variables are just place holders for propositions. For example, "I am reading something", where the propositional variable "something" is what you are reading.

Principals can make requests. In the ACL, principals make requests using the says operator. Requests have the form P says  $\varphi$ , where P represents some principal and  $\varphi$  represents some assertion. For example, PlatoonLeader says platoonHalt. In this example, the Platoon Leader is issuing a command (or request) for the platoon to halt.

Principals can have authority over assertions. In the ACL, authority is conveyed using the *controls* operator. Statements of authority have the form P controls  $\varphi$ , where P represents some principal and  $\varphi$  represents some assertion. For example, PlatoonLeader controls platoonHalt. This example states that the Platoon Leader has the authority to issue the command (or request) for the platoon to halt.

Principals can also have jurisdiction over assertions. Both authority and jurisdiction use the *controls* operator. Statements of jurisdiction have the same form as statements of authority. Statements of authority are typically defined in an organization's policy. Statements of jurisdiction are statements that are readily believed given the context. For example, *PresidentOfUS controls (PlatoonLeader controls platoonHalt)*. In this example, the President of the United States, per the U.S. Constitution, has jurisdiction over the authority invested in the Platoon Leader. In particular, the President of the United States has the jurisdiction to give the Platoon Leader the authority to command her platoon to halt.

In addition, principals can speak for other principals. Principals do this using the *speaks* for operator. The ACL represents the *speaks* for operator with the symbol  $\Rightarrow$ . These types of statements have the form  $P \Rightarrow Q$ , where both P and Q are principals. For example,  $PlatoonLeader \Rightarrow PresidentOfUS$ . This example states that the Platoon Leader speaks for the President of the United States.

### 4.2.5 Well-formed Formulas

Well-formed formulas (WFFs) valid statements or expressions in the ACL. All statements or requests must be a WFF. The text book defines the set of WFFs using the identifier **Form**:

Form ::= PropVar / 
$$\neg$$
 Form / (Form  $\lor$  Form) / (Form  $\land$  Form) / (Form  $\supset$  Form) / (Form  $\equiv$  Form) / (Princ  $\Rightarrow$  Princ) / (Princ says Princ) / (Princ controls Form)

### 4.2.6 Kripke Structure

A Kripke structure deals with three things: worlds, propositions, and principals. The worlds can be thought of as possible states or configurations of some system.

Propositions are just statements that are either true or false. And, principals are just actors. A Kripke structure  $\mathcal{M} = \langle W, I, J \rangle$  is defined as a three-tuple consisting of a set of worlds W, a function I that maps propositions to worlds, and a functions J that maps principals to relations on worlds. A more formal definition is definition 2.1 in the text:

- W is a nonempty set, whose elements are called worlds.
- $I: \mathbf{PropVar} \to \mathcal{P}(W)$  is an interpretation function that maps each propositional variable to a set of worlds.
- $J: \mathbf{PName} \to \mathcal{P}((W \times W) \text{ is a function that maps each principal name to a } relation on worlds.$

### **4.2.6.1** satisies

### 4.2.6.2 soundness

### 4.2.7 Inference Rules

The inference rules for the access-control logic (ACL) are shown in figure 4.1. All the inference rules are sound. Details of proofs of soundness can be found in *Access Control*, *Security, and Trust: A Logical Approach*[1].

$$P \ \text{controls} \ \varphi \ \stackrel{\text{def}}{=} \ (P \ \text{says} \ \varphi) \supset \varphi \qquad P \ \text{reps} \ Q \ \text{on} \ \varphi \stackrel{\text{def}}{=} P \mid Q \ \text{says} \ \varphi \supset Q \ \text{says} \ \varphi$$
 
$$Modus \ Ponens \ \frac{\varphi \quad \varphi \supset \varphi'}{\varphi'} \qquad Says \ \frac{\varphi}{P \ \text{says} \ \varphi} \qquad Controls \ \frac{P \ \text{controls} \ \varphi}{\varphi} \qquad P \ \text{says} \ \varphi}$$
 
$$Derived \ Speaks \ For \ \frac{P \Rightarrow Q \quad P \ \text{says} \ \varphi}{Q \ \text{says} \ \varphi} \qquad Reps \qquad \frac{Q \ \text{controls} \ \varphi}{\varphi} \qquad P \ \text{reps} \ Q \ \text{on} \ \varphi \qquad P \mid Q \ \text{says} \ \varphi}{\varphi}$$
 
$$\& \ Says \ (1) \quad \frac{P \ \& \ Q \ \text{says} \ \varphi}{P \ \text{says} \ \varphi \land Q \ \text{says} \ \varphi} \qquad \& \ Says \ (2) \quad \frac{P \ \text{says} \ \varphi \land Q \ \text{says} \ \varphi}{P \ \& \ Q \ \text{says} \ \varphi}$$
 
$$Quoting \ (1) \quad \frac{P \mid Q \ \text{says} \ \varphi}{P \ \text{says} \ Q \ \text{says} \ \varphi} \qquad Quoting \ (2) \quad \frac{P \ \text{says} \ Q \ \text{says} \ \varphi}{P \mid Q \ \text{says} \ \varphi}$$
 
$$Idempotency \ of \Rightarrow \qquad \frac{P \Rightarrow P}{P \Rightarrow P} \qquad Monotonicity \ of \Rightarrow \qquad \frac{P' \Rightarrow P \quad Q' \Rightarrow Q}{P' \mid Q' \Rightarrow P \mid Q}$$

Figure 4.1: The ACL inference rules. Image taken from  $Access\ Control$ , Security, and Trust:  $A\ Logical\ Approach[1]$ 

### 4.2.8 Complete mediation

Fundamental to this work is the concept of complete mediation (discussed in section 4.1.1). In the ACL, this means that each principal must be authenticated and authorized on each request. ACL does this primarily by the *Controls* inference rule in figure 4.1 and shown again here in figure 4.2.

$$\begin{array}{ccc} Controls & \frac{P \; \text{controls} \; \varphi & P \; \text{says} \; \varphi}{\varphi} \end{array}$$

Figure 4.2: The Controls inference rule. Image taken from Access Control, Security, and Trust: A Logical Approach[1]

ACL refers to the left statement as an authorization<sup>1</sup>. The principal P controls (is authorized on) some action  $\varphi$ . ACL refers to the right statement in this inference rule as a request <sup>2</sup>. The principal P requests some action  $\varphi$ . The conjunction of the authorization and the request of P on  $\varphi$  results in the action  $\varphi$ . That is, if P controls  $\varphi$  and P says  $\varphi$  then  $\varphi$  is true.

 $<sup>^{1}</sup>$ or a control in the C2 calculus

<sup>&</sup>lt;sup>2</sup>or a *command* in the C2 calculus

The Reps rule also demonstrates complete mediation. It follows a similar logic.

However, the *Reps* rule is not used in this master thesis.

### 4.2.9 Well formed statements

### 4.2.10 Inference Rules

### 4.3 ACL in HOL

The equivalence of the ACL formulas implemented in HOL are shown in figure 4.3.

Access-Control Logic Formula	HOL Syntax
$\langle jump \rangle$	prop jump
$\neg \langle jump \rangle$	notf (prop jump)
$\langle run \rangle \wedge \langle jump \rangle$	prop run andf prop jump
$\langle run \rangle \vee \langle stop \rangle$	prop run orf prop stop
$\langle run \rangle \supset \langle jump \rangle$	prop run impf prop jump
$\langle walk \rangle \equiv \langle stop \rangle$	prop walk eqf prop stop
$Alice$ says $\langle jump \rangle$	Name Alice says prop jump
$Alice \& Bob $ says $\langle stop \rangle$	Name Alice meet Name Bob says prop stop
$Bob \mid Carol \text{ says } \langle run \rangle$	Name Bob quoting Name Carol says prop run
$Bob$ controls $\langle walk \rangle$	Name Bob controls prop walk
$Bob$ reps $Alice$ on $\langle jump \rangle$	reps (Name Bob) (Name Alice) (prop jump)
$Carol \Rightarrow Bob$	Name Carol speaks_for Name Bob

Figure 4.3: The ACL formulas in HOL. Image taken from  $Access\ Control,\ Security,\ and\ Trust:\ A\ Logical\ Approach[1]$ 

Using this syntax, an ACL request of the form P says  $\varphi$  would have the form

Name P says prop  $\varphi$ 

- 4.3.1 Complete Mediation
- 4.3.2 satList

# Patrol Base Operations

### 5.1 Motivation

The patrol base operations described in the patrol base Ranger Handbook

### 5.2 Ranger Handbook Description

# 5.3 Modeling the Patrol Base Operations from the Ranger Handbook

Modeling a system requires the knowledge of an expert on the system. This is necessary because only someone who is familiar with the system, especially with regards to security, can detail its nuances. For this reason, a subject matter expert from the United States Army (Jesse Nathaniel Hall) is employed to develop a model of the patrol base operations.

The model of the patrol base operations needs to be amiable to complete mediation and verification using an access-control logic (section 4.2). This is necessary to prove security properties of the patrol base operations. To do this, the patrol base operations are abstracted from the Ranger Manual and modeled in Visio<sup>1</sup>. The result of doing this is a hierarchy of secure state machines (SSMs). (SSMs are described in section 6.2.)

Each level of the hierarchy of SSMs represents a level of abstraction of the patrol base operations. The most abstract level of the hierarchy is the top level SSM. A diagram of this most abstract level is shown in figure 5.1.

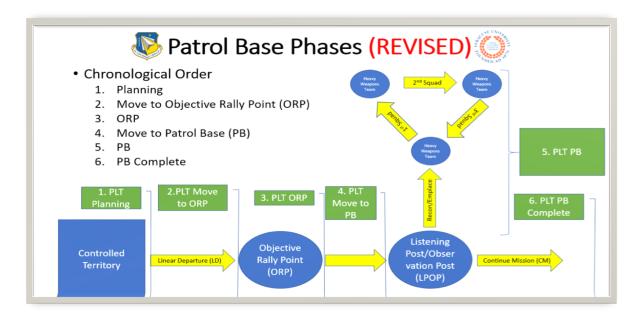


Figure 5.1: A diagram of the most abstract level in the hierarchy of secure state machines.

The diagram describes a chronological order of abstract phases (modeled as states) of the patrol base operations. The operations begin with the planning phase (1). Next, they move to the objective rally point (ORP) (2). At the ORP, operations commence (3). When these are complete, the patrol base operations move to the actual patrol base (4). At the patrol base, operations proceed (5). Finally, the patrol base operations are

<sup>&</sup>lt;sup>1</sup>This work began as a collaboration between Jesse Nathaniel Hall and the author. Once the hierarchy of secure state machines was decided upon, the abstraction of the Ranger Handbook was done by Jesse Nathaniel Hall with only structural consultation with the author. Concurrently, the author focused on proving the properties of complete mediation in the ACL using HOL. Thus, there was a great deal of separation of work. Jesse's work is described here because it is necessary to put the entire system into context for this master thesis. This means that Jesse's work provided the model of the system for which the principle of complete mediation was proved, verified and documented.

complete (6). These are the six states in the top level SSM.

The next level of abstraction in the hierarchy of SSMs represents a horizontal slice through the patrol base operations. This is the second level of the hierarchical description of the patrol base operations. It is referred to as the sub-level. In this documentation, SSMs at this level are referred to as the sub-level, sublevel, or sub-level SSMs. This slice describes the patrol base operations at a lower level of abstraction. It expands each of the states in the top level (except for the last state PB Complete). For example, the planning phase (1) in figure 5.1 is expanded into an SSM of its own. This is called ssmPlanPB. It consists of several states (see section 5.4.5.1) which detail activities conducted during the planning phase of the patrol base operations. Each state in the top level (except for PB Complete) has it's own SSM (see the next section).

At yet another lower level of abstraction is the sub-sub (3rd) level. In this documentation, SSMs at this level are referred to as the sub-sub-level, subsublevel, or subsubLevel SSMs. This level expands upon the states in the sub level (one level above) in the same manner that the sub level expands upon the states in the top level SSM. In this manner, each level is a lower level of abstraction than the level above it.

A vertical slice through the diagram is also modeled. This slice models the patrol base operations from the top level down to the most detailed level (level 8). This vertical slice consists of a series of SSMs. Each SSM expands upon only one state in the level above it. This differs from the horizontal slice which expands upon all states in the level above it. Expanding upon only one state focuses on a vertical slice through all 8 states of the hierarchy of SSMs.

The vertical slice begins at the top level SSM. Next, it expands upon one state at this level, the move to ORP state (2). This results in a sub level SSM named ssmMoveToORP. The vertical slice progresses in this manner, by expanding one state at each level into a new SSM. From ssmMoveToORP, the state secure halt is expanded to

ssmSecureHalt. From within this SSM, the state *ORP Recon* is expanded into ssmORPRecon. From within this SSM, the state *Move to ORP 4L* (fourth level move to ORP state) is expanded into ssmMoveToORP4L. Finally, from within this SSM, the state *Form RT* is expanded into ssmFormRT.

The vertical slice spans the all eight levels. However, not all levels are represented with an SSM. The last SSM in the vertical slice, ssmFormRT, is actually at the 5th level of the hierarchy. This SSM, consists of three states. These three states reside at the 7th level because ssmFormRT does not have states at the 6th level (it skips the 6th level). Furthermore, the 7th level states are not expanded into an SSM because each of these 7th level states expand into only one state at the 8th level.

In addition to the horizontal and vertical slices, an escape level is also modeled. Actions in the escape level are reachable from any phase of the patrol base operations. These actions are the unacceptable circumstances that require the patrol base operations to abort. For example, if the patrol base contacts the enemy in any phase of the operations, then the command *react to combat* is issued. The patrol base operations are subsequently aborted.

Excluding the escape level, there are eight levels of the hierarchy of SSMs.

Note that, the purpose of this master thesis is to demonstrate that the properties of complete mediation could be applied and verified on a non-automated, human-centered systems. Thus, it is sufficient to demonstrate this on a horizontal and vertical slice of the patrol base operations.

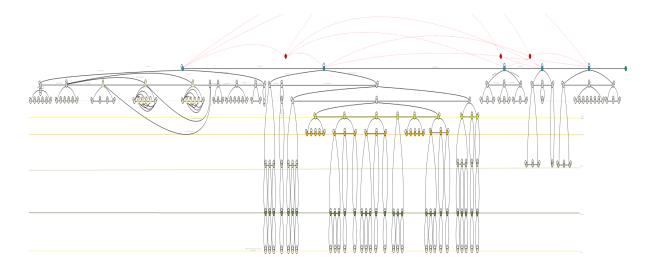


Figure 5.2: Diagrammatic description of patrol base operations as a hierarchy of secure state machines. (Generated by Jesse Nathaniel Hall.)

### 5.4 Hierarchy of Secure State Machines

### 5.4.1 Diagrammatic Description in Visio

The enormity of the hierarchy of SSMs is evident in figure 5.2. This is a squashed version of the Visio diagram for the hierarchy of SSMs. The diagram is included as a Visio file with the files for this project (LaTeX/figures/diagram.vis).

The straight, colored lines that span the diagram in figure 5.2 delineate levels of the hierarchy of SSMs. The top lines are obscured by the size of this squashed version of the diagram. The most visible bright yellow line delineates the sub-sub-sub (4th) level of the hierarchy, for example.

The small, colored dots in figure 5.2 represent states (phases) of the patrol base operations. The red dots are an exception. The labels for these states are not readable in this diagram. The dots are color coded. The colors correspond to the level of those states. For example, the dots at the top level (level below the red dots) are all dark blue.

In figure 5.2, the red dots at the top of the diagram represent the escape level SSM.

But, they do not represent states in the SSM. This is because the escape level SSM is accessible by all phases of the patrol base operations. This means that if can not be ascribed to any one level of the hierarchy.

The dots representing states are connected to each other by lines. These lines represent allowable transitions from one state to another. The escape level is again an exception. If no line connects one state to another then no transition is allowed.

The red dots representing the escape level SSM are best thought of as multiple copies of a floating SSM. The escape level SSM acts as a sub-SSM for all SSMs in the hierarchy. During patrol base operations, abortion of the patrol base operations can occur at any action from any state at any level. This means that any state at any level can be terminated by the escape level SSM. Drawing lines from all states to the escape level SSM and drawing really long lines clutters the diagram. Therefore, only lines at the top level are drawn and the red dots are duplicated.

The lines in the diagram are annotated by SSM requests. (Annotations are visible in the original Visio diagram, but not in this squashed version.) For example, a line connecting the top level state PLAN\_PB is annotated with the request *PlatoonLeader says crossLD*. crossLD is an abbreviation for "cross the line of discrimination" and it is the command to transition to the MOVE\_TO\_ORP state. Lines are not annotated beyond the sub level.

Details of each level follow in the next section.

# 5.4.2 OMNI-Level5.4.3 Escape5.4.4 Top Level

Horizontal Slice

 $5.4.5.1 \quad ssmPlanPB$ 

5.4.5

5.4.5.2 ssmMoveToORP

labels ssec: ssmMoveToORP

5.4.5.3 ssmConductORP

labels ssec:ssmConductORP

5.4.5.4 ssmMoveToPB

labels ssec:ssmMoveToPB

### 5.4.5.5 ssmConductPB

### 5.4.6 Vertical Slice

### 5.4.6.1 ssmSecureHalt

### 5.4.6.2 ssmORPRecon

### 5.4.6.3 ssmMoveToORP4L

### 5.4.6.4 ssmFormRT

secure state machine (SSM)-glesentryfull

secure state machines (SSMs)-glesentryfullpl

SSM-glsentrytext

SSMs-glsentry shortpl

# Secure State Machine Model

C 1	$\mathbf{C}\mathbf{A} - \mathbf{A} =$	<b>N</b> / <b>C</b> - 1-	•
6.1	State	lviacr	nnes

- 6.1.1 Next-state Function
- 6.1.2 Next-output Function
- 6.1.3 Transition Commands

### 6.2 Secure State Machines

- 6.2.1 State Machine Versus Secure State Machine
- 6.2.2 Transition Types
- 6.2.3 Authentication

# Chapter 7

# Patrol Base Operations as Secure State Machines

- 7.1 ssmPB: An Example from the Hierarchy
- 7.1.1 Principals
- **7.1.2** States
- 7.1.3 Commands
- 7.1.4 Next-State Function
- 7.1.5 Next-Output Function
- 7.1.6 Authentication

# Chapter 8

## Discussion

- 8.1 Recap
- 8.2 Mission Accomplished
- 8.3 Stop-Gaps, Lessons Learned, & Advice
- 8.4 Other Verifiable Theories
- 8.4.1 Platoon Theory, Soldier Theory, Squad Theory, etc.
- 8.4.2 Soldiers in Roles

# Chapter 9

# Future Work & Implications

This is the future works section. But, as I am typing this, it is the current working section for LaTeXThe point here is to get the margins in order. This means that there must be text of sufficient length to visually verify that the text meets LORI's standards. LORI is complying with SU standards for the senior thesis. Therefore, meeting LORI's standards is synonymous with meeting SU's standards. Resistance will only degrade you.

### 9.1 The Devil Is in The Details

Of course, there are top margins and bottom margins. This means that we'll need more text. You know, the best wait to generate text is to just cut-n-paste some random stuff. Perinton, N.Y. – The FBI conducted a search of Morgan Management LLC's offices in Monroe County Monday as part of an ongoing investigation into the development company's business practices, according to Rochester area media reports.

Agents were seen carrying boxes in and out of the company's headquarters at 1080

Pittsford Victor Road in the town of Perinton, according to the reports.

An FBI spokeswoman confirmed that agents conducted "court-authorized activity at 1080 Pittsford Victor Road," the Democrat & Chronicle reported. The company's founder, developer Robert Morgan, was in the office as agents conducted the search, the newspaper said.

The newspaper reported in September that a federal investigation is focused on bank loans to Morgan's real estate portfolio, which, according to the company's website, has grown to 140 properties and more than 34,000 apartment units across 14 states since the company's founding in 1979.

The investigation is centered largely on Buffalo-region apartment complexes purchased by Morgan's companies and whether the information the company gave lenders to obtain the loans was accurate, according to the newspaper.

However, the Buffalo News reported in March that the investigation includes a look at Morgan's purchase of the Rugby Square apartment complex on Dorchester Avenue in Syracuse. One of Morgan's companies borrowed \$5.56 million to buy the apartment complex in a distress sale in 2012, then obtained a new \$9 million mortgage on the property just 10 months later after reporting a major turnaround of the complex, the newspaper said.

Morgan has said his companies have done nothing illegal to obtain financing. No charges have been filed in connection with the investigation.

According to the company's website, Morgan operates 13 apartment complexes in the Syracuse area.

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### 9.2 Accountability Systems

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### 9.3 Applicability

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# Appendices

# Appendix A

Access Control Logic Theories: Pretty-Printed Theories

### Contents

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#### 1 aclfoundation Theory

Built: 25 February 2018

Parent Theories: indexedLists, patternMatches

#### 1.1 Datatypes

```
Form =
    TT
  | FF
  | prop 'aavar
  | notf (('aavar, 'apn, 'il, 'sl) Form)
  | (andf) (('aavar, 'apn, 'il, 'sl) Form)
           (('aavar, 'apn, 'il, 'sl) Form)
  | (orf) (('aavar, 'apn, 'il, 'sl) Form)
          (('aavar, 'apn, 'il, 'sl) Form)
  | (impf) (('aavar, 'apn, 'il, 'sl) Form)
           (('aavar, 'apn, 'il, 'sl) Form)
  | (eqf) (('aavar, 'apn, 'il, 'sl) Form)
          (('aavar, 'apn, 'il, 'sl) Form)
  | (says) ('apn Princ) (('aavar, 'apn, 'il, 'sl) Form)
  | (speaks_for) ('apn Princ) ('apn Princ)
  | (controls) ('apn Princ) (('aavar, 'apn, 'il, 'sl) Form)
  | reps ('apn Princ) ('apn Princ)
         (('aavar, 'apn, 'il, 'sl) Form)
  | (domi) (('apn, 'il) IntLevel) (('apn, 'il) IntLevel)
  | (eqi) (('apn, 'il) IntLevel) (('apn, 'il) IntLevel)
  | (doms) (('apn, 'sl) SecLevel) (('apn, 'sl) SecLevel)
  | (eqs) (('apn, 'sl) SecLevel) (('apn, 'sl) SecLevel)
  | (eqn) num num
  | (lte) num num
  | (lt) num num
Kripke =
    KS ('aavar -> 'aaworld -> bool)
       ('apn -> 'aaworld -> 'aaworld -> bool) ('apn -> 'il)
       ('apn -> 'sl)
Princ =
   Name 'apn
  | (meet) ('apn Princ) ('apn Princ)
  | (quoting) ('apn Princ) ('apn Princ);
IntLevel = iLab 'il | il 'apn ;
SecLevel = sLab 'sl | sl 'apn
```

#### 1.2 Definitions

```
[imapKS_def]
  \vdash \forall \mathit{Intp} \ \mathit{Jfn} \ \mathit{ilmap} \ \mathit{slmap}.
        imapKS (KS Intp Jfn ilmap slmap) = ilmap
[intpKS_def]
  \vdash \ \forall \mathit{Intp} \ \mathit{Jfn} \ \mathit{ilmap} \ \mathit{slmap}.
        intpKS (KS Intp Jfn ilmap slmap) = Intp
[jKS_def]
 \vdash \forall Intp \ Jfn \ ilmap \ slmap. jKS (KS Intp \ Jfn \ ilmap \ slmap) = Jfn
[01_def]
 ⊢ 01 = PO one_weakorder
[one_weakorder_def]
 \vdash \forall x \ y. \ \text{one\_weakorder} \ x \ y \iff \mathtt{T}
[po_TY_DEF]
 \vdash \exists \mathit{rep}. TYPE_DEFINITION WeakOrder \mathit{rep}
[po_tybij]
 \vdash (\forall a. PO (repPO a) = a) \land
    \forall r. WeakOrder r \iff (repPO (PO r) = r)
[prod_PO_def]
  \vdash \forall PO_1 \ PO_2.
        prod_PO PO_1 PO_2 = PO (RPROD (repPO PO_1) (repPO PO_2))
[smapKS_def]
  \vdash \forall Intp \ Jfn \ ilmap \ slmap.
        smapKS (KS Intp Jfn ilmap slmap) = slmap
[Subset_PO_def]
 \vdash Subset_P0 = P0 (\subseteq)
1.3 Theorems
[abs_po11]
  \vdash \forall r \ r'.
        \texttt{WeakOrder} \ r \ \Rightarrow \ \texttt{WeakOrder} \ r' \ \Rightarrow \ \texttt{((PO} \ r \ \texttt{=} \ \texttt{PO} \ r') \ \Longleftrightarrow \ (r \ \texttt{=} \ r'))
[absPO_fn_onto]
 \vdash \forall a. \exists r. (a = PO r) \land WeakOrder r
```

```
[antisym_prod_antisym]
 \vdash \forall r \ s.
       antisymmetric r \wedge \text{antisymmetric } s \Rightarrow
       antisymmetric (RPROD r s)
[EQ_WeakOrder]
 ⊢ WeakOrder (=)
[KS_bij]
 \vdash \forall M. M = KS \text{ (intpKS } M) \text{ (jKS } M) \text{ (imapKS } M) \text{ (smapKS } M)
[one_weakorder_WO]
 ⊢ WeakOrder one_weakorder
[onto_po]
 \vdash \ \forall \, r. WeakOrder r \iff \exists \, a. r = repPO a
[po_bij]
 \vdash (\forall a. PO (repPO a) = a) \land
    \forall r. WeakOrder r \iff (repPO (PO r) = r)
[PO_repPO]
 \vdash \forall a. \ PO \ (repPO \ a) = a
[refl_prod_refl]
 \vdash \ \forall \, r \ s. reflexive r \ \land reflexive s \ \Rightarrow reflexive (RPROD r \ s)
[repPO_iPO_partial_order]
 \vdash (\forall x. repPO iPO x x) \land
     (\forall x \ y. \ \texttt{repPO} \ iPO \ x \ y \ \land \ \texttt{repPO} \ iPO \ y \ x \ \Rightarrow \ (x = y)) \ \land
    \forall x \ y \ z. repPO iPO \ x \ y \ \land repPO iPO \ y \ z \Rightarrow repPO iPO \ x \ z
[repP0_01]
 ⊢ repPO 01 = one_weakorder
[repPO_prod_PO]
 \vdash \forall po_1 po_2.
       repPO (prod_PO po_1 po_2) = RPROD (repPO po_1) (repPO po_2)
[repPO_Subset_PO]
 \vdash repPO Subset_PO = (\subseteq)
[RPROD_THM]
 \vdash \forall r \ s \ a \ b.
       RPROD r s a b \iff r (FST a) (FST b) \wedge s (SND a) (SND b)
```

```
[SUBSET\_WO] \\ \vdash WeakOrder (\subseteq) \\ [trans\_prod\_trans] \\ \vdash \forall r \ s. \ transitive \ r \land transitive \ s \Rightarrow transitive \ (RPROD \ r \ s) \\ [WeakOrder\_Exists] \\ \vdash \exists R. \ WeakOrder \ R \\ [WO\_prod\_WO] \\ \vdash \forall r \ s. \ WeakOrder \ r \land WeakOrder \ s \Rightarrow WeakOrder \ (RPROD \ r \ s) \\ [WO\_repPO] \\ \vdash \forall r. \ WeakOrder \ r \iff (repPO \ (PO \ r) = r) \\
```

#### 2 aclsemantics Theory

Built: 25 February 2018

Parent Theories: aclfoundation

#### 2.1 Definitions

```
[Efn_def]
 \vdash (\forall Oi \ Os \ M. Efn Oi \ Os \ M TT = \mathcal{U}(:,v)) \land
     (\forall Oi \ Os \ M. \ Efn \ Oi \ Os \ M \ FF = \{\}) \land
     (\forall~Oi~Os~M~p. Efn Oi~Os~M (prop p) = intpKS M~p) \land
     (\forall Oi \ Os \ M \ f.
         Efn Oi\ Os\ M (notf f) = \mathcal{U}(:'v) DIFF Efn Oi\ Os\ M f) \land
     (\forall Oi \ Os \ M \ f_1 \ f_2.
         Efn Oi Os M (f_1 and f_2) =
         Efn Oi Os M f_1 \cap Efn Oi Os M f_2) \wedge
     (\forall Oi \ Os \ M \ f_1 \ f_2.
         Efn Oi Os M (f_1 orf f_2) =
         Efn Oi Os M f_1 \cup Efn Oi Os M f_2) \wedge
     (\forall Oi \ Os \ M \ f_1 \ f_2.
         Efn Oi Os M (f_1 \text{ impf } f_2) =
         \mathcal{U}(: \text{'v}) DIFF Efn Oi\ Os\ M\ f_1\ \cup Efn Oi\ Os\ M\ f_2)\ \wedge
     (\forall Oi \ Os \ M \ f_1 \ f_2.
         Efn Oi Os M (f_1 eqf f_2) =
         (\mathcal{U}(:\,\,{}^{\backprime}\mathtt{v}) DIFF Efn Oi\ Os\ M\ f_1\ \cup Efn Oi\ Os\ M\ f_2) \cap
         (\mathcal{U}(:, v) DIFF Efn Oi\ Os\ M\ f_2\ \cup Efn Oi\ Os\ M\ f_1)) \wedge
     (\forall Oi \ Os \ M \ P \ f.
         Efn Oi\ Os\ M\ (P\ says\ f) =
         \{w \mid \text{Jext (jKS } M) \mid P \mid w \subseteq \text{Efn } Oi \mid Os \mid M \mid f\}) \land
     (\forall Oi \ Os \ M \ P \ Q.
         Efn Oi \ Os \ M (P speaks_for Q) =
```

```
if Jext (jKS M) Q RSUBSET Jext (jKS M) P then \mathcal{U}(:'v)
         else { }) \ \
     (\forall Oi \ Os \ M \ P \ f.
         Efn Oi\ Os\ M (P controls f) =
         \mathcal{U}(: `v) DIFF \{w \mid \text{Jext (jKS } M) \mid P \mid w \subseteq \text{Efn } Oi \mid Os \mid M \mid f\} \cup \mathcal{U}(: `v)
         Efn Oi Os M f) \land
     (\forall Oi \ Os \ M \ P \ Q \ f.
         Efn Oi\ Os\ M (reps P\ Q\ f) =
         \mathcal{U}(:,v) DIFF
         \{w \mid \text{Jext (jKS } M) \mid (P \text{ quoting } Q) \mid w \subseteq \text{Efn } Oi \mid Os \mid M \mid f\} \cup G
         \{w \mid \text{Jext (jKS } M) \mid Q \mid w \subseteq \text{Efn } Oi \mid Os \mid M \mid f\} \}
     (\forall Oi \ Os \ M \ intl_1 \ intl_2.
         Efn Oi \ Os \ M \ (intl_1 \ domi \ intl_2) =
         if repPO Oi (Lifn M intl_2) (Lifn M intl_1) then \mathcal{U}(:,v)
         else { }) \ \
     (\forall Oi \ Os \ M \ intl_2 \ intl_1.
         Efn Oi \ Os \ M \ (intl_2 \ eqi \ intl_1) =
         (if repPO Oi (Lifn M intl_2) (Lifn M intl_1) then \mathcal{U}(:,v)
          else { }) ∩
         if repPO Oi (Lifn M intl_1) (Lifn M intl_2) then \mathcal{U}(:,v)
         else { }) \
     (\forall Oi \ Os \ M \ secl_1 \ secl_2.
         Efn Oi \ Os \ M \ (secl_1 \ doms \ secl_2) =
         if repPO Os (Lsfn M secl_2) (Lsfn M secl_1) then \mathcal{U}(:,v)
         else { }) \ \
     (\forall Oi \ Os \ M \ secl_2 \ secl_1.
         Efn Oi\ Os\ M\ (secl_2\ eqs\ secl_1) =
         (if repPO Os (Lsfn M secl_2) (Lsfn M secl_1) then \mathcal{U}(:,v)
          else { }) ∩
         if repPO Os (Lsfn M secl_1) (Lsfn M secl_2) then \mathcal{U}(:,v)
         else { }) \ \
     (\forall Oi \ Os \ M \ numExp_1 \ numExp_2.
         Efn Oi\ Os\ M\ (numExp_1\ eqn\ numExp_2) =
         if numExp_1 = numExp_2 then \mathcal{U}(:,v) else \{\}) \land
     (\forall Oi \ Os \ M \ numExp_1 \ numExp_2.
         Efn Oi\ Os\ M\ (numExp_1\ lte\ numExp_2) =
         if numExp_1 \leq numExp_2 then \mathcal{U}(:'v) else \{\}) \land
     \forall Oi \ Os \ M \ numExp_1 \ numExp_2.
       Efn Oi Os M (numExp_1 lt numExp_2) =
       if numExp_1 < numExp_2 then \mathcal{U}(:,v) else \{\}
[Jext_def]
 \vdash (\forall J \ s. Jext J (Name s) = J \ s) \land
     (\forall J P_1 P_2.
         Jext J (P_1 meet P_2) = Jext J P_1 RUNION Jext J P_2) \wedge
     \forall J \ P_1 \ P_2. Jext J \ (P_1 \ \text{quoting} \ P_2) = Jext J \ P_2 O Jext J \ P_1
[Lifn_def]
 \vdash (\forall M \ l. Lifn M (iLab l) = l) \land
    \forall M \ name. Lifn M (il name) = imapKS M name
```

```
[Lsfn_def]
 \vdash (\forall M \ l. Lsfn M (sLab l) = l) \land
    \forall\,M name. Lsfn M (sl name) = smapKS M name
2.2
       Theorems
[andf_def]
 \vdash \ \forall \ Oi \ Os \ M \ f_1 \ f_2.
       Efn Oi Os M (f_1 and f_2) = Efn Oi Os M f_1 \cap Efn Oi Os M f_2
[controls_def]
 \vdash \forall Oi \ Os \ M \ P \ f.
       Efn Oi \ Os \ M \ (P \ controls \ f) =
       \mathcal{U}(:"v) DIFF \{w \mid \text{Jext (jKS } M) \mid P \mid w \subseteq \text{Efn } Oi \mid Os \mid M \mid f\} \cup \mathcal{U}(:"v)
       Efn Oi Os M f
[controls_says]
 \vdash \forall M \ P \ f.
       Efn Oi\ Os\ M (P controls f) = Efn Oi\ Os\ M (P says f impf f)
[domi_def]
 \vdash \ \forall \ Oi \ Os \ M \ intl_1 \ intl_2.
       Efn Oi Os M (intl_1 domi intl_2) =
       if repPO Oi (Lifn M intl_2) (Lifn M intl_1) then \mathcal{U}(:'v)
       else { }
[doms_def]
 \vdash \ \forall \ Oi \ Os \ M \ secl_1 \ secl_2.
       Efn Oi \ Os \ M \ (secl_1 \ doms \ secl_2) =
       if repPO \mathit{Os} (Lsfn \mathit{M} \mathit{secl}_2) (Lsfn \mathit{M} \mathit{secl}_1) then \mathcal{U}(:\, \mbox{`v})
       else { }
[eqf_def]
 \vdash \ \forall \ Oi \ Os \ M \ f_1 \ f_2.
       Efn Oi Os M (f_1 \text{ eqf } f_2) =
       (\mathcal{U}(:, v) DIFF Efn Oi\ Os\ M\ f_2\ \cup Efn Oi\ Os\ M\ f_1)
[eqf_impf]
 \vdash \ \forall M \ f_1 \ f_2.
       Efn Oi Os M (f_1 eqf f_2) =
       Efn Oi Os M ((f_1 impf f_2) andf (f_2 impf f_1))
```

```
[eqi_def]
 \vdash \ \forall \ Oi \ \ Os \ \ M \ \ intl_2 \ \ intl_1 \, .
       Efn Oi\ Os\ M\ (intl_2\ eqi\ intl_1) =
       (if repPO Oi (Lifn M intl_2) (Lifn M intl_1) then \mathcal{U}(:,v)
        else { }) ∩
       if repPO Oi (Lifn M intl_1) (Lifn M intl_2) then \mathcal{U}(:,v)
       else { }
[eqi_domi]
 \vdash \ \forall M \ intL_1 \ intL_2.
       Efn Oi \ Os \ M \ (intL_1 \ eqi \ intL_2) =
       Efn Oi Os M (intL_2 domi intL_1 and intL_1 domi intL_2)
eqn_def
 \vdash \ \forall \ Oi \ Os \ M \ numExp_1 \ numExp_2.
       Efn Oi\ Os\ M (numExp_1 eqn numExp_2) =
       if numExp_1 = numExp_2 then \mathcal{U}(:,v) else \{\}
[eqs_def]
 \vdash \forall Oi \ Os \ M \ secl_2 \ secl_1.
       Efn Oi \ Os \ M \ (secl_2 \ eqs \ secl_1) =
       (if repPO Os (Lsfn M secl_2) (Lsfn M secl_1) then \mathcal{U}(:,v)
        else { }) ∩
       if repPO Os (Lsfn M secl_1) (Lsfn M secl_2) then \mathcal{U}(:,v)
       else { }
[eqs_doms]
 \vdash \forall M \ secL_1 \ secL_2.
       Efn Oi\ Os\ M\ (secL_1\ eqs\ secL_2) =
       Efn Oi Os M (secL_2 doms secL_1 and secL_1 doms secL_2)
[FF_def]
 \vdash \forall Oi \ Os \ M. Efn Oi \ Os \ M FF = {}
[impf_def]
 \vdash \ \forall \ Oi \ Os \ M \ f_1 \ f_2.
       Efn Oi \ Os \ M \ (f_1 \ \text{impf} \ f_2) =
       \mathcal{U}(:, v) DIFF Efn Oi Os M f_1 \cup Efn Oi Os M f_2
[lt_def]
 \vdash \ \forall \ Oi \ Os \ M \ numExp_1 \ numExp_2.
       Efn Oi\ Os\ M\ (numExp_1\ lt\ numExp_2) =
       if numExp_1 < numExp_2 then \mathcal{U}(:,v) else \{\}
[lte_def]
 \vdash \forall Oi \ Os \ M \ numExp_1 \ numExp_2.
       Efn Oi\ Os\ M (numExp_1 lte numExp_2) =
       if numExp_1 \leq numExp_2 then \mathcal{U}(:,v) else \{\}
```

```
[meet_def]
 \vdash \forall J \ P_1 \ P_2. Jext J \ (P_1 \ \text{meet} \ P_2) = Jext J \ P_1 RUNION Jext J \ P_2
[name_def]
 \vdash \forall J \ s. \ \texttt{Jext} \ J \ (\texttt{Name} \ s) = J \ s
[notf_def]
 \vdash \forall Oi \ Os \ M \ f. \ \texttt{Efn} \ Oi \ Os \ M \ (\texttt{notf} \ f) = \mathcal{U}(:'\texttt{v}) \ \texttt{DIFF} \ \texttt{Efn} \ Oi \ Os \ M \ f
[orf_def]
 \vdash \forall Oi \ Os \ M \ f_1 \ f_2.
        Efn Oi Os M (f_1 orf f_2) = Efn Oi Os M f_1 \cup Efn Oi Os M f_2
[prop_def]
 \vdash \ \forall \ Oi \ Os \ M \ p. Efn Oi \ Os \ M (prop p) = intpKS M p
[quoting_def]
 \vdash \forall J \ P_1 \ P_2. Jext J (P_1 quoting P_2) = Jext J P_2 O Jext J P_1
[reps_def]
  \vdash \ \forall \ Oi \ Os \ M \ P \ Q \ f.
        Efn Oi\ Os\ M (reps P\ Q\ f) =
        \mathcal{U}(:,v) DIFF
        \{w \mid \mathsf{Jext} \ (\mathsf{jKS} \ M) \ (P \ \mathsf{quoting} \ Q) \ w \subseteq \mathsf{Efn} \ \mathit{Oi} \ \mathit{Os} \ M \ f\} \ \cup
        \{w \mid \text{Jext (jKS } M) \mid Q \mid w \subseteq \text{Efn } Oi \mid Os \mid M \mid f\}
[says_def]
 \vdash \forall Oi \ Os \ M \ P \ f.
        Efn Oi \ Os \ M \ (P \ \text{says} \ f) =
        \{w \mid \text{Jext (jKS } M) \mid P \mid w \subseteq \text{Efn } Oi \mid Os \mid M \mid f\}
[speaks_for_def]
  \vdash \forall Oi \ Os \ M \ P \ Q.
        Efn Oi\ Os\ M (P speaks_for Q) =
        if Jext (jKS M) Q RSUBSET Jext (jKS M) P then \mathcal{U}(:'v)
[TT_def]
 \vdash \ \forall \ Oi \ Os \ M . Efn Oi \ Os \ M TT = \mathcal{U}(: `v)
```

#### 3 aclrules Theory

Built: 25 February 2018

Parent Theories: aclsemantics

#### 3.1 Definitions

```
[sat_def]
 \vdash \forall M \ Oi \ Os \ f. \ (M,Oi,Os) \ \text{sat} \ f \iff (\text{Efn} \ Oi \ Os \ M \ f = \mathcal{U}(:'world))
3.2
       Theorems
[And_Says]
 \vdash \forall M \ Oi \ Os \ P \ Q \ f.
       (M,Oi,Os) sat P meet Q says f eqf P says f and Q says f
[And_Says_Eq]
 \vdash (M, Oi, Os) sat P meet Q says f \iff
     (M,Oi,Os) sat P says f and Q says f
[and_says_lemma]
 \vdash \forall M \ Oi \ Os \ P \ Q \ f.
       (M,Oi,Os) sat P meet Q says f impf P says f and f says f
[Controls_Eq]
 \vdash \forall M \ Oi \ Os \ P \ f.
       (M,Oi,Os) sat P controls f\iff (M,Oi,Os) sat P says f impf f
[DIFF_UNIV_SUBSET]
 \vdash (\mathcal{U}(:'a) DIFF s \cup t = \mathcal{U}(:'a)) \iff s \subseteq t
[domi_antisymmetric]
 \vdash \ \forall M \ Oi \ Os \ l_1 \ l_2.
       (M,Oi,Os) sat l_1 domi l_2 \Rightarrow
       (M, Oi, Os) sat l_2 domi l_1 \Rightarrow
       (M,Oi,Os) sat l_1 eqi l_2
[domi_reflexive]
 \vdash \ \forall \, M \ Oi \ Os \ l. \ (M,Oi,Os) \ {\it sat} \ l \ {\it domi} \ l
[domi_transitive]
 \vdash \ \forall M \ Oi \ Os \ l_1 \ l_2 \ l_3.
       (M, Oi, Os) sat l_1 domi l_2 \Rightarrow
       (M, Oi, Os) sat l_2 domi l_3 \Rightarrow
       (M,Oi,Os) sat l_1 domi l_3
[doms_antisymmetric]
 \vdash \ \forall M \ Oi \ Os \ l_1 \ l_2.
       (M,Oi,Os) sat l_1 doms l_2 \Rightarrow
       (M,Oi,Os) sat l_2 doms l_1 \Rightarrow
       (M,Oi,Os) sat l_1 eqs l_2
```

```
[doms_reflexive]
 \vdash \forall \, M \ Oi \ Os \ l. \ (M,Oi,Os) \ {\it sat} \ l \ {\it doms} \ l
[doms_transitive]
 \vdash \ \forall M \ Oi \ Os \ l_1 \ l_2 \ l_3.
        (M, Oi, Os) sat l_1 doms l_2 \Rightarrow
        (M,Oi,Os) sat l_2 doms l_3 \Rightarrow
        (M, Oi, Os) sat l_1 doms l_3
[eqf_and_impf]
 \vdash \ \forall M \ Oi \ Os \ f_1 \ f_2.
        (M,Oi,Os) sat f_1 eqf f_2 \iff
        (M,Oi,Os) sat (f_1 \text{ impf } f_2) and (f_2 \text{ impf } f_1)
[eqf_andf1]
 \vdash \ \forall \, M \ \ Oi \ \ Os \ f \ \ f' \ \ g \, .
        (M, Oi, Os) sat f \neq f' \Rightarrow
        (M,Oi,Os) sat f and g \Rightarrow
        (M, Oi, Os) sat f' and g
[eqf_andf2]
 \vdash \ \forall \, M \ \ Oi \ \ Os \ f \ \ f' \ \ g \, .
        (M,Oi,Os) sat f eqf f' \Rightarrow
        (M,Oi,Os) sat g and f \Rightarrow
        (M,Oi,Os) sat g and f'
eqf_controls
 \vdash \forall M \ Oi \ Os \ P \ f \ f'.
        (M,Oi,Os) sat f \neq f' \Rightarrow
        (M,Oi,Os) sat P controls f \Rightarrow
        (M, Oi, Os) sat P controls f'
[eqf_eq]
 \vdash (Efn Oi\ Os\ M\ (f_1\ \mathsf{eqf}\ f_2) = \mathcal{U}(:\ \mathsf{'b})) \iff
     (Efn Oi Os M f_1 = Efn Oi Os M f_2)
[eqf_eqf1]
 \vdash \ \forall M \ Oi \ Os \ f \ f' \ g.
        (M,Oi,Os) sat f eqf f' \Rightarrow
        (M, Oi, Os) sat f \neq g \Rightarrow
        (M,Oi,Os) sat f' eqf g
[eqf_eqf2]
 \vdash \ \forall \, M \ \ Oi \ \ Os \ f \ f' \ g \, .
        (M,Oi,Os) sat f \neq f' \Rightarrow
        (M,Oi,Os) sat g eqf f \Rightarrow
        (M,Oi,Os) sat g eqf f'
```

```
[eqf_impf1]
 \vdash \forall M \ Oi \ Os \ f \ f' \ g.
       (M,Oi,Os) sat f eqf f' \Rightarrow
       (M,Oi,Os) sat f impf g \Rightarrow
       (M,Oi,Os) sat f' impf g
[eqf_impf2]
 \vdash \forall M \ Oi \ Os \ f \ f' \ g.
       (M,Oi,Os) sat f eqf f' \Rightarrow
       (M,Oi,Os) sat g impf f \Rightarrow
       (M,Oi,Os) sat g impf f'
[eqf_notf]
 \vdash \forall M \ Oi \ Os \ f \ f'.
       (M,Oi,Os) sat f \neq f' \Rightarrow
       (M,Oi,Os) sat notf f \Rightarrow
       (M,Oi,Os) sat notf f'
[eqf_orf1]
 \vdash \ \forall M \ Oi \ Os \ f \ f' \ g.
       (M,Oi,Os) sat f eqf f' \Rightarrow
       (M,Oi,Os) sat f orf g \Rightarrow
       (M,Oi,Os) sat f' orf g
[eqf_orf2]
 \vdash \forall M \ Oi \ Os \ f \ f' \ g.
       (M,Oi,Os) sat f eqf f' \Rightarrow
       (M,Oi,Os) sat g orf f \Rightarrow
       (M,Oi,Os) sat g orf f'
[eqf_reps]
 \vdash \forall M \ Oi \ Os \ P \ Q \ f \ f'.
       (M,Oi,Os) sat f eqf f' \Rightarrow
       (M,Oi,Os) sat reps P Q f \Rightarrow
       (M,Oi,Os) sat reps P Q f'
[eqf_sat]
 \vdash \forall M \ Oi \ Os \ f_1 \ f_2.
       (M,Oi,Os) sat f_1 eqf f_2 \Rightarrow
       ((M,Oi,Os) \text{ sat } f_1 \iff (M,Oi,Os) \text{ sat } f_2)
[eqf_says]
 \vdash \forall M \ Oi \ Os \ P \ f \ f'.
       (M,Oi,Os) sat f eqf f' \Rightarrow
       (M,Oi,Os) sat P says f \Rightarrow
       (M,Oi,Os) sat P says f'
```

```
[eqi_Eq]
 \vdash \forall M \ Oi \ Os \ l_1 \ l_2.
        (M,Oi,Os) sat l_1 eqi l_2 \iff
        (M,Oi,Os) sat l_2 domi l_1 andf l_1 domi l_2
[eqs_Eq]
 \vdash \ \forall M \ Oi \ Os \ l_1 \ l_2.
        (M,Oi,Os) sat l_1 eqs l_2 \iff
        (M,Oi,Os) sat l_2 doms l_1 and l_1 doms l_2
[Idemp_Speaks_For]
 \vdash \ \forall M \ Oi \ Os \ P. (M,Oi,Os) sat P speaks_for P
[Image_cmp]
 \vdash \forall R_1 \ R_2 \ R_3 \ u. (R_1 \ \mathsf{O} \ R_2) u \subseteq R_3 \iff R_2 \ u \subseteq \{y \mid R_1 \ y \subseteq R_3\}
[Image_SUBSET]
 \vdash \ \forall \, R_1 \ R_2 \,. \ R_2 \ \text{RSUBSET} \ R_1 \ \Rightarrow \ \forall \, w \,. \ R_2 \ w \ \subseteq \ R_1 \ w
[Image_UNION]
 \vdash \forall R_1 R_2 w. (R_1 RUNION R_2) w = R_1 w \cup R_2 w
[INTER_EQ_UNIV]
 \vdash (s \cap t = \mathcal{U}(:'a)) \iff (s = \mathcal{U}(:'a)) \land (t = \mathcal{U}(:'a))
[Modus_Ponens]
 \vdash \ \forall M \ Oi \ Os \ f_1 \ f_2.
        (M, Oi, Os) sat f_1 \Rightarrow
        (M, Oi, Os) sat f_1 impf f_2 \Rightarrow
        (M,Oi,Os) sat f_2
[Mono_speaks_for]
 \vdash \ \forall M \ Oi \ Os \ P \ P' \ Q \ Q'.
        (M,Oi,Os) sat P speaks_for P' \Rightarrow (M,Oi,Os) sat Q speaks_for Q' \Rightarrow
        (M,Oi,Os) sat P quoting Q speaks_for P' quoting Q'
[MP_Says]
 \vdash \forall M \ Oi \ Os \ P \ f_1 \ f_2.
        (M,Oi,Os) sat
        P says (f_1 \text{ impf } f_2) impf P says f_1 impf P says f_2
Quoting
 \vdash \forall M \ Oi \ Os \ P \ Q \ f.
        (M,Oi,Os) sat P quoting Q says f eqf P says Q says f
```

```
[Quoting_Eq]
 \vdash \ \forall M \ Oi \ Os \ P \ Q \ f.
        (M, Oi, Os) sat P quoting Q says f \iff
       (M,Oi,Os) sat P says Q says f
[reps_def_lemma]
 \vdash \ \forall M \ Oi \ Os \ P \ Q \ f.
       Efn Oi\ Os\ M (reps P\ Q\ f) =
       Efn Oi Os M (P quoting Q says f impf Q says f)
[Reps_Eq]
 \vdash \forall M \ Oi \ Os \ P \ Q \ f.
        (M,Oi,Os) sat reps P Q f \iff
       (M,Oi,Os) sat P quoting Q says f impf Q says f
[sat_allworld]
 \vdash \ \forall \ M \ f. \ (M,Oi,Os) \ {\sf sat} \ f \iff \forall \ w. \ w \in {\sf Efn} \ Oi \ Os \ M \ f
[sat_andf_eq_and_sat]
 \vdash (M, Oi, Os) sat f_1 and f_2 \iff
     (M,Oi,Os) sat f_1 \wedge (M,Oi,Os) sat f_2
sat_TT
 \vdash (M, Oi, Os) sat TT
[Says]
 \vdash \ \forall M \ Oi \ Os \ P \ f. \ (M,Oi,Os) \ {\tt sat} \ f \ \Rightarrow \ (M,Oi,Os) \ {\tt sat} \ P \ {\tt says} \ f
[says_and_lemma]
 \vdash \ \forall \, M \ Oi \ Os \ P \ Q \ f \, .
       (M,Oi,Os) sat P says f and f says f impf P meet f says f
[Speaks_For]
 \vdash \forall M \ Oi \ Os \ P \ Q \ f.
        (M,Oi,Os) sat P speaks_for Q impf P says f impf Q says f
[speaks_for_SUBSET]
 \vdash \forall R_3 \ R_2 \ R_1.
       R_2 RSUBSET R_1 \Rightarrow \forall w. \{w \mid R_1 \mid w \subseteq R_3\} \subseteq \{w \mid R_2 \mid w \subseteq R_3\}
[SUBSET_Image_SUBSET]
 \vdash \ \forall R_1 \ R_2 \ R_3.
       (\forall w_1. R_2 w_1 \subseteq R_1 w_1) \Rightarrow
       \forall w. \{w \mid R_1 \ w \subseteq R_3\} \subseteq \{w \mid R_2 \ w \subseteq R_3\}
```

```
Trans_Speaks_For
  \vdash \forall M \ Oi \ Os \ P \ Q \ R.
           (M, Oi, Os) sat P speaks_for Q \Rightarrow
           (M, Oi, Os) sat Q speaks_for R \Rightarrow
           (M,Oi,Os) sat P speaks_for R
[UNIV_DIFF_SUBSET]
  \vdash \forall R_1 \ R_2. \ R_1 \subseteq R_2 \Rightarrow (\mathcal{U}(:\ 'a)\ \mathtt{DIFF}\ R_1 \cup R_2 = \mathcal{U}(:\ 'a))
[world_and]
  \vdash \forall M \ Oi \ Os \ f_1 \ f_2 \ w.
           w \in \mathsf{Efn}\ \mathit{Oi}\ \mathit{Os}\ \mathit{M}\ (\mathit{f}_1\ \mathsf{andf}\ \mathit{f}_2) \iff
           w \in \text{Efn } Oi \ Os \ M \ f_1 \ \land \ w \in \text{Efn } Oi \ Os \ M \ f_2
[world_eq]
  \vdash \ \forall M \ Oi \ Os \ f_1 \ f_2 \ w.
           w \in \text{Efn } Oi \ Os \ M \ (f_1 \ \text{eqf} \ f_2) \iff
           (w \in \mathsf{Efn}\ \mathit{Oi}\ \mathit{Os}\ \mathit{M}\ \mathit{f}_1 \iff w \in \mathsf{Efn}\ \mathit{Oi}\ \mathit{Os}\ \mathit{M}\ \mathit{f}_2)
world_eqn
  \vdash \forall M \ Oi \ Os \ n_1 \ n_2 \ w. \ w \in \texttt{Efn} \ Oi \ Os \ m \ (n_1 \ \texttt{eqn} \ n_2) \iff (n_1 \ \texttt{=} \ n_2)
[world_F]
  \vdash \ \forall \, M \ Oi \ Os \ w \, . \ w \, \notin \, \mathtt{Efn} \ Oi \ Os \ M \ \mathtt{FF}
world_imp
  \vdash \forall M \ Oi \ Os \ f_1 \ f_2 \ w.
           w \in \text{Efn } Oi \ Os \ M \ (f_1 \ \text{impf} \ f_2) \iff
           w \in \texttt{Efn} \ Oi \ Os \ M \ f_1 \ \Rightarrow \ w \in \texttt{Efn} \ Oi \ Os \ M \ f_2
[world_lt]
  \vdash \ orall \ \mathit{M} \ \mathit{Oi} \ \mathit{Os} \ \mathit{n}_1 \ \mathit{n}_2 \ \mathit{w} . \ \mathit{w} \in \mathsf{Efn} \ \mathit{Oi} \ \mathit{Os} \ \mathit{m} \ (\mathit{n}_1 \ \mathsf{lt} \ \mathit{n}_2) \iff \mathit{n}_1 < \mathit{n}_2
[world_lte]
  \vdash \ orall \ \mathit{M} \ \mathit{Oi} \ \mathit{Os} \ \mathit{n}_1 \ \mathit{n}_2 \ \mathit{w} . \ \mathit{w} \ \in \ \mathsf{Efn} \ \mathit{Oi} \ \mathit{Os} \ \mathit{m} \ (\mathit{n}_1 \ \mathsf{lte} \ \mathit{n}_2) \ \Longleftrightarrow \ \mathit{n}_1 \ \leq \ \mathit{n}_2
[world_not]
  [world_or]
  \vdash \forall M \ f_1 \ f_2 \ w.
           w \in \mathsf{Efn} \ \mathit{Oi} \ \mathit{Os} \ \mathit{M} \ (\mathit{f}_1 \ \mathsf{orf} \ \mathit{f}_2) \iff
           w \in \mathtt{Efn}\ Oi\ Os\ M\ f_1\ \lor\ w \in \mathtt{Efn}\ Oi\ Os\ M\ f_2
[world_says]
  \vdash \ \forall M \ Oi \ Os \ P \ f \ w.
           w \in \text{Efn } Oi \ Os \ M \ (P \ \text{says} \ f) \iff
           \forall v. v \in \text{Jext (jKS } M) \ P \ w \Rightarrow v \in \text{Efn } Oi \ Os \ M \ f
world_T
  \vdash \forall M \ Oi \ Os \ w. \ w \in \texttt{Efn} \ Oi \ Os \ M \ \texttt{TT}
```

### 4 aclDrules Theory

Built: 25 February 2018

Parent Theories: aclrules

#### 4.1 Theorems

```
[Conjunction]
 \vdash \ \forall M \ Oi \ Os \ f_1 \ f_2.
       (M,Oi,Os) sat f_1 \Rightarrow
       (M,Oi,Os) sat f_2 \Rightarrow
       (M, Oi, Os) sat f_1 and f_2
[Controls]
 \vdash \ \forall M \ Oi \ Os \ P \ f.
       (M,Oi,Os) sat P says f \Rightarrow
       (M, Oi, Os) sat P controls f \Rightarrow
       (M,Oi,Os) sat f
[Derived_Controls]
 \vdash \forall M \ Oi \ Os \ P \ Q \ f.
       (M,Oi,Os) sat P speaks_for Q \Rightarrow
       (M,Oi,Os) sat Q controls f \Rightarrow
       (M,Oi,Os) sat P controls f
[Derived_Speaks_For]
 \vdash \ \forall M \ Oi \ Os \ P \ Q \ f.
       (M,Oi,Os) sat P speaks_for Q \Rightarrow
       (M,Oi,Os) sat P says f \Rightarrow
       (M,Oi,Os) sat Q says f
[Disjunction1]
 \vdash \forall M \ Oi \ Os \ f_1 \ f_2. \ (M,Oi,Os) \ sat \ f_1 \Rightarrow (M,Oi,Os) \ sat \ f_1 \ orf \ f_2
[Disjunction2]
 \vdash \forall M \ Oi \ Os \ f_1 \ f_2. (M,Oi,Os) sat f_2 \Rightarrow (M,Oi,Os) sat f_1 orf f_2
[Disjunctive_Syllogism]
 \vdash \ \forall M \ Oi \ Os \ f_1 \ f_2.
       (M,Oi,Os) sat f_1 orf f_2 \Rightarrow
       (M,Oi,Os) sat notf f_1 \Rightarrow
       (M,Oi,Os) sat f_2
[Double_Negation]
 \vdash \forall M \ Oi \ Os \ f. \ (M,Oi,Os) \ \text{sat notf (notf } f) \Rightarrow (M,Oi,Os) \ \text{sat } f
```

```
[eqn_eqn]
 \vdash (M, Oi, Os) sat c_1 eqn n_1 \Rightarrow
     (M,Oi,Os) sat c_2 eqn n_2 \Rightarrow
     (M,Oi,Os) sat n_1 eqn n_2 \Rightarrow
     (M, Oi, Os) sat c_1 eqn c_2
[eqn_lt]
 \vdash (M, Oi, Os) sat c_1 eqn n_1 \Rightarrow
     (M,Oi,Os) sat c_2 eqn n_2 \Rightarrow
     (M,Oi,Os) sat n_1 lt n_2 \Rightarrow
     (M, Oi, Os) sat c_1 lt c_2
[eqn_lte]
 \vdash (M, Oi, Os) sat c_1 eqn n_1 \Rightarrow
     (M,Oi,Os) sat c_2 eqn n_2 \Rightarrow
     (M,Oi,Os) sat n_1 lte n_2 \Rightarrow
     (M,Oi,Os) sat c_1 lte c_2
[Hypothetical_Syllogism]
 \vdash \ \forall M \ Oi \ Os \ f_1 \ f_2 \ f_3.
        (M, Oi, Os) sat f_1 impf f_2 \Rightarrow
        (M,Oi,Os) sat f_2 impf f_3 \Rightarrow
        (M, Oi, Os) sat f_1 impf f_3
[il_domi]
 \vdash \ \forall \, M \ \ Oi \ \ Os \ \ P \ \ Q \ \ l_1 \ \ l_2 \, .
        (M, Oi, Os) sat il P eqi l_1 \Rightarrow
        (M, Oi, Os) sat il Q eqi l_2 \Rightarrow
        (M,Oi,Os) sat l_2 domi l_1 \Rightarrow
        (M,Oi,Os) sat il Q domi il P
[INTER_EQ_UNIV]
 \vdash \forall s_1 \ s_2. \ (s_1 \cap s_2 = \mathcal{U}(:'a)) \iff (s_1 = \mathcal{U}(:'a)) \land (s_2 = \mathcal{U}(:'a))
[Modus_Tollens]
 \vdash \ \forall M \ Oi \ Os \ f_1 \ f_2.
        (M,Oi,Os) sat f_1 impf f_2 \Rightarrow
        (M,Oi,Os) sat notf f_2 \Rightarrow
        (M, Oi, Os) sat notf f_1
[Rep_Controls_Eq]
 \vdash \forall M \ Oi \ Os \ A \ B \ f.
        (M,Oi,Os) sat reps A B f \iff
        (M, Oi, Os) sat A controls B says f
```

```
[Rep_Says]
 \vdash \ \forall M \ Oi \ Os \ P \ Q \ f.
        (M,Oi,Os) sat reps P Q f \Rightarrow
        (M,Oi,Os) sat P quoting Q says f \Rightarrow
        (M,Oi,Os) sat Q says f
[Reps]
 \vdash \ \forall \, M \ \ Oi \ \ Os \ \ P \ \ Q \ \ f \, .
        (M,Oi,Os) sat reps P Q f \Rightarrow
        (M,Oi,Os) sat P quoting Q says f \Rightarrow
        (M,Oi,Os) sat Q controls f \Rightarrow
        (M, Oi, Os) sat f
[Says_Simplification1]
 \vdash \ \forall M \ Oi \ Os \ P \ f_1 \ f_2.
        (M,Oi,Os) sat P says (f_1 \text{ andf } f_2) \Rightarrow (M,Oi,Os) sat P says f_1
[Says_Simplification2]
 \vdash \forall M \ Oi \ Os \ P \ f_1 \ f_2.
        (M,Oi,Os) sat P says (f_1 \text{ andf } f_2) \Rightarrow (M,Oi,Os) sat P says f_2
[Simplification1]
 \vdash \ \forall \ M \ Oi \ Os \ f_1 \ f_2. (M,Oi,Os) sat f_1 andf f_2 \Rightarrow (M,Oi,Os) sat f_1
[Simplification2]
 \vdash \forall \, M \;\; Oi \;\; Os \;\; f_1 \;\; f_2 . (M, Oi, Os) sat f_1 andf f_2 \;\Rightarrow\; (M, Oi, Os) sat f_2
[sl_doms]
 \vdash \ \forall \, M \ \ Oi \ \ Os \ \ P \ \ Q \ \ l_1 \ \ l_2 \, .
        (M, Oi, Os) sat sl P eqs l_1 \Rightarrow
        (M,Oi,Os) sat sl Q eqs l_2 \Rightarrow
        (M,Oi,Os) sat l_2 doms l_1 \Rightarrow
        (M,Oi,Os) sat sl Q doms sl P
```

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# Appendix B

Secure State Machine & Patrol Base Operations: Pretty-Printed Theories

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# 1 OMNIType Theory

**Built:** 13 May 2018

Parent Theories: indexedLists, patternMatches

### 1.1 Datatypes

```
command = ESCc escCommand | SLc 'slCommand
escCommand = returnToBase | changeMission | resupply
               | reactToContact
escOutput = ReturnToBase | ChangeMission | Resupply
             | ReactToContact
escState = RTB | CM | RESUPPLY | RTC
output = ESCo escOutput | SLo 'slOutput
principal = SR 'stateRole
state = ESCs escState | SLs 'slState
1.2
       Theorems
[command_distinct_clauses]
 \vdash \ \forall \ a' \ a. ESCc a \neq \operatorname{SLc} \ a'
[command_one_one]
 \vdash (\forall a \ a'. (ESCc a = ESCc \ a') \iff (a = a')) \land
    \forall a \ a'. (SLc a = SLc \ a') \iff (a = a')
[escCommand_distinct_clauses]
 \vdash returnToBase \neq changeMission \land returnToBase \neq resupply \land
    returnToBase \neq reactToContact \land changeMission \neq resupply \land
    \texttt{changeMission} \neq \texttt{reactToContact} \ \land \ \texttt{resupply} \neq \texttt{reactToContact}
[escOutput_distinct_clauses]
 \vdash ReturnToBase \neq ChangeMission \land ReturnToBase \neq Resupply \land
    \texttt{ReturnToBase} \neq \texttt{ReactToContact} \ \land \ \texttt{ChangeMission} \neq \texttt{Resupply} \ \land
    \texttt{ChangeMission} \neq \texttt{ReactToContact} \ \land \ \texttt{Resupply} \neq \texttt{ReactToContact}
[escState_distinct_clauses]
 \vdash RTB \neq CM \land RTB \neq RESUPPLY \land RTB \neq RTC \land CM \neq RESUPPLY \land
    \mathtt{CM} \neq \mathtt{RTC} \wedge \mathtt{RESUPPLY} \neq \mathtt{RTC}
```

```
[output_distinct_clauses]
\vdash \forall a' \ a. \ ESCo \ a \neq SLo \ a'
[output_one_one]
\vdash (\forall a \ a'. \ (ESCo \ a = ESCo \ a') \iff (a = a')) \land \\ \forall a \ a'. \ (SLo \ a = SLo \ a') \iff (a = a')
[principal_one_one]
\vdash \forall a \ a'. \ (SR \ a = SR \ a') \iff (a = a')
[state_distinct_clauses]
\vdash \forall a' \ a. \ ESCs \ a \neq SLs \ a'
[state_one_one]
\vdash (\forall a \ a'. \ (ESCs \ a = ESCs \ a') \iff (a = a')) \land \\ \forall a \ a'. \ (SLs \ a = SLs \ a') \iff (a = a')
```

## 2 ssm11 Theory

**Built:** 13 May 2018

Parent Theories: satList

### 2.1 Datatypes

```
configuration =
    CFG (('command order, 'principal, 'd, 'e) Form -> bool)
        ('state -> ('command order, 'principal, 'd, 'e) Form)
        (('command order, 'principal, 'd, 'e) Form list)
        (('command order, 'principal, 'd, 'e) Form list) 'state
        ('output list)

order = SOME 'command | NONE

trType = discard 'command | trap 'command | exec 'command
```

### 2.2 Definitions

```
\begin{array}{l} [\mathsf{TR\_def}] \\ \vdash \mathsf{TR} = \\ (\lambda \, a_0 \ a_1 \ a_2 \ a_3 \, . \\ \forall \, \mathit{TR'} \, . \\ (\forall \, a_0 \ a_1 \ a_2 \ a_3 \, . \\ (\exists \, \mathit{authenticationTest} \, P \, \mathit{NS} \, \mathit{M} \, \mathit{Oi} \, \mathit{Os} \, \mathit{Out} \, s \\ securityContext \, \, \mathit{stateInterp} \, \mathit{cmd} \, \mathit{ins} \, \mathit{outs} \, . \\ (a_0 = (M, \mathit{Oi}, \mathit{Os})) \, \land \, (a_1 = \mathsf{exec} \, \mathit{cmd}) \, \land \\ (a_2 = (a_2 + a_3)) & (a_3 = (a_3 + a_3)) & (a_4 = (a_4 + a_3)) & (a_4 + a_4)) & (a_4 = (a_4 + a_4)) & (a_4 + a_4)) & (a_4 = (a_4 + a_4)) & (a_4 = (a_4 + a_4)) & (a_4 + a_4)) & (a_4 = (a_4 + a_4)) & (a_4
```

SSM11 THEORY Theorems

```
{\tt CFG} authentication Test stateInterp
                     securityContext (P says prop (SOME cmd)::ins) s
                     outs) \land
                  (a_3 =
                   CFG authentication Test stateInterp
                     securityContext ins (NS s (exec cmd))
                      (Out \ s \ (exec \ cmd)::outs)) \land
                  authenticationTest (P says prop (SOME cmd)) \land
                  CFGInterpret (M, Oi, Os)
                    (CFG authentication Test stateInterp
                        securityContext (P says prop (SOME cmd)::ins)
                        s outs)) \vee
              (\exists authenticationTest\ P\ NS\ M\ Oi\ Os\ Out\ s
                   security Context\ state Interp\ cmd\ ins\ outs .
                  (a_0 = (M, Oi, Os)) \land (a_1 = trap \ cmd) \land
                   {\tt CFG} authentication Test stateInterp
                     securityContext (P says prop (SOME cmd)::ins) s
                     outs) \land
                  (a_3 =
                   CFG authentication Test stateInterp
                     securityContext ins (NS s (trap cmd))
                      (Out \ s \ (trap \ cmd) :: outs)) \land
                  authenticationTest (P says prop (SOME cmd)) \land
                  CFGInterpret (M, Oi, Os)
                    (CFG authenticationTest\ stateInterp
                        securityContext (P says prop (SOME cmd)::ins)
                        s outs)) \vee
              (\exists authentication Test\ NS\ M\ Oi\ Os\ Out\ s\ security Context
                   stateInterp\ cmd\ x\ ins\ outs.
                  (a_0 = (M, Oi, Os)) \land (a_1 = discard \ cmd) \land
                  (a_2 =
                   CFG authentication Test stateInterp
                     securityContext (x::ins) s outs) \land
                  (a_3 =
                   {\tt CFG} authentication Test stateInterp
                     securityContext\ ins\ (NS\ s\ ({\tt discard}\ cmd))
                      (Out \ s \ (discard \ cmd)::outs)) \ \land
                  \neg authentication Test x) \Rightarrow
              TR' a_0 a_1 a_2 a_3) \Rightarrow
          TR' a_0 a_1 a_2 a_3)
      Theorems
[CFGInterpret_def]
 \vdash CFGInterpret (M, Oi, Os)
       (CFG \ authentication Test \ stateInterp \ security Context
```

(input::ins) state  $outputStream) \iff$ 

2.3

SSM11 THEORY Theorems

```
(M,Oi,Os) satList securityContext \land (M,Oi,Os) sat input \land
     (M,Oi,Os) sat stateInterp state
[CFGInterpret_ind]
 \vdash \ \forall P.
       (\forall M \ Oi \ Os \ authentication Test \ stateInterp \ security Context
             input ins state outputStream.
            P (M, Oi, Os)
               (CFG \ authentication Test \ stateInterp \ security Context
                   (input::ins) state outputStream)) <math>\land
       (\forall v_{15} \ v_{10} \ v_{11} \ v_{12} \ v_{13} \ v_{14}.
            P \ v_{15} \ (CFG \ v_{10} \ v_{11} \ v_{12} \ [] \ v_{13} \ v_{14})) \ \Rightarrow
       \forall v \ v_1 \ v_2 \ v_3. P \ (v, v_1, v_2) \ v_3
[configuration_one_one]
 \vdash \forall a_0 \ a_1 \ a_2 \ a_3 \ a_4 \ a_5 \ a_0' \ a_1' \ a_2' \ a_3' \ a_4' \ a_5'.
       (CFG a_0 a_1 a_2 a_3 a_4 a_5 = CFG a_0' a_1' a_2' a_3' a_4' a_5') \iff
       (a_0 = a_0') \wedge (a_1 = a_1') \wedge (a_2 = a_2') \wedge (a_3 = a_3') \wedge
       (a_4 = a_4') \wedge (a_5 = a_5')
[order_distinct_clauses]
 \vdash \ \forall \, a. SOME a \neq \mathtt{NONE}
[order_one_one]
 \vdash \forall a \ a'. (SOME a = \text{SOME } a') \iff (a = a')
[TR_cases]
 \vdash \forall a_0 \ a_1 \ a_2 \ a_3.
       TR a_0 a_1 a_2 a_3 \iff
       (\exists authenticationTest\ P\ NS\ M\ Oi\ Os\ Out\ s\ securityContext
             stateInterp cmd ins outs.
            (a_0 = (M, Oi, Os)) \land (a_1 = exec \ cmd) \land
            (a_2 =
             {\tt CFG} \ \ authentication Test \ \ state Interp \ \ security Context
                (P says prop (SOME cmd)::ins) s outs) \land
            (a_3 =
             {\tt CFG} authentication Test state Interp security Context ins
                (NS s (exec cmd)) (Out s (exec cmd)::outs)) \land
            authenticationTest (P says prop (SOME cmd)) \land
           CFGInterpret (M, Oi, Os)
               (CFG \ authentication Test \ stateInterp \ security Context
                   (P says prop (SOME cmd)::ins) s outs)) \lor
       (\exists authenticationTest\ P\ NS\ M\ Oi\ Os\ Out\ s\ securityContext
             stateInterp cmd ins outs.
            (a_0 = (M, Oi, Os)) \land (a_1 = trap \ cmd) \land
            (a_2 =
             {\tt CFG} authentication Test state Interp security Context
                (P says prop (SOME cmd)::ins) s outs) \wedge
```

Theorems SSM11 THEORY

```
{\tt CFG} authentication Test state Interp security Context ins
              (NS s (trap cmd)) (Out s (trap cmd)::outs)) \wedge
          authenticationTest (P says prop (SOME cmd)) \land
          CFGInterpret (M, Oi, Os)
             (CFG \ authentication Test \ stateInterp \ security Context
                (P \text{ says prop (SOME } cmd)::ins) \ s \ outs)) \ \lor
      \exists authenticationTest NS M Oi Os Out s securityContext
          stateInterp\ cmd\ x\ ins\ outs.
         (a_0 = (M, Oi, Os)) \land (a_1 = discard \ cmd) \land
         (a_2 =
          {\tt CFG} authentication Test state Interp security Context
             (x::ins) s outs) \wedge
         (a_3 =
          {\tt CFG} authentication Test state Interp security Context ins
             (NS \ s \ (discard \ cmd)) (Out \ s \ (discard \ cmd)::outs)) \land
         \neg authentication Test x
[TR_discard_cmd_rule]
 \vdash TR (M, Oi, Os) (discard cmd)
       (CFG authenticationTest stateInterp securityContext
          (x::ins) s outs)
       (\mathtt{CFG}\ authenticationTest\ stateInterp\ securityContext\ ins
          (NS \ s \ (discard \ cmd)) \ (Out \ s \ (discard \ cmd)::outs)) \iff
    \neg authenticationTest x
[TR_EQ_rules_thm]
 \vdash (TR (M, Oi, Os) (exec cmd)
        (CFG authentication Test stateInterp securityContext
            (P says prop (SOME cmd)::ins) s outs)
        (\mathtt{CFG}\ authenticationTest\ stateInterp\ securityContext\ ins
           (NS \ s \ (exec \ cmd)) \ (Out \ s \ (exec \ cmd)::outs)) \iff
     authenticationTest (P says prop (SOME cmd)) \land
     CFGInterpret (M, Oi, Os)
        (CFG authentication Test stateInterp securityContext
            (P says prop (SOME cmd)::ins) s outs)) \land
    (TR (M, Oi, Os) (trap cmd)
        (CFG \ authentication Test \ stateInterp \ security Context
            (P says prop (SOME cmd)::ins) s outs)
        (\mathtt{CFG}\ authenticationTest\ stateInterp\ securityContext\ ins
            (NS \ s \ (trap \ cmd)) \ (Out \ s \ (trap \ cmd)::outs)) \iff
     authenticationTest (P says prop (SOME cmd)) \land
     CFGInterpret (M, Oi, Os)
        (CFG \ authentication Test \ stateInterp \ security Context
            (P says prop (SOME cmd)::ins) s outs)) \land
    (TR (M,Oi,Os) (discard cmd)
        ({\tt CFG}\ \ authentication Test\ \ state Interp\ \ security Context
            (x::ins) s outs)
        (CFG authenticationTest stateInterp securityContext ins
```

SSM11 THEORY Theorems

```
(NS \ s \ (discard \ cmd)) \ (Out \ s \ (discard \ cmd)::outs)) \iff
     \neg authentication Test x)
[TR_exec_cmd_rule]
 \vdash \ \forall \ authenticationTest \ \ securityContext \ \ stateInterp \ \ P \ \ cmd \ \ ins \ \ s
       (\forall M \ Oi \ Os.
          CFGInterpret (M, Oi, Os)
             (CFG \ authentication Test \ stateInterp \ security Context
                 (P \text{ says prop (SOME } cmd)::ins) \ s \ outs) \Rightarrow
           (M,Oi,Os) sat prop (SOME cmd)) \Rightarrow
      \forall NS \ Out \ M \ Oi \ Os.
         TR (M, Oi, Os) (exec cmd)
            (CFG authenticationTest stateInterp securityContext
                (P says prop (SOME cmd)::ins) s outs)
            (\mathtt{CFG}\ authentication\ Test\ stateInterp\ security\ Context\ ins
                (NS \ s \ (exec \ cmd)) \ (Out \ s \ (exec \ cmd)::outs)) \iff
         authenticationTest (P says prop (SOME cmd)) \land
         CFGInterpret (M, Oi, Os)
            (CFG \ authentication Test \ stateInterp \ security Context
                (P \text{ says prop } (SOME \ cmd) :: ins) \ s \ outs) \ \land
         (M, Oi, Os) sat prop (SOME cmd)
[TR_ind]
 \vdash \forall TR'.
       (\forall authentication Test P NS M Oi Os Out s security Context
            stateInterp cmd ins outs.
           authenticationTest (P says prop (SOME cmd)) \land
          CFGInterpret (M, Oi, Os)
             (CFG \ authentication Test \ stateInterp \ security Context
                 (P \text{ says prop (SOME } cmd)::ins) \ s \ outs) \Rightarrow
           TR' (M, Oi, Os) (exec cmd)
             (CFG \ authentication Test \ stateInterp \ security Context
                 (P says prop (SOME cmd)::ins) s outs)
             ({\tt CFG}\ authentication Test\ state Interp\ security Context
                 ins \ (NS \ s \ (exec \ cmd)) \ (Out \ s \ (exec \ cmd)::outs))) \ \land
       (\forall authentication Test\ P\ NS\ M\ Oi\ Os\ Out\ s\ security Context
            stateInterp cmd ins outs.
           authenticationTest (P says prop (SOME cmd)) \land
          CFGInterpret (M, Oi, Os)
             (CFG authenticationTest stateInterp securityContext
                 (P \text{ says prop (SOME } cmd)::ins) \ s \ outs) \Rightarrow
           TR' (M, Oi, Os) (trap cmd)
             (CFG \ authentication Test \ stateInterp \ security Context
                 (P says prop (SOME cmd)::ins) s outs)
             (\mathtt{CFG}\ authenticationTest\ stateInterp\ securityContext)
                 ins \ (NS \ s \ (trap \ cmd)) \ (Out \ s \ (trap \ cmd)::outs))) \ \land
       (\forall\, authentication Test\ NS\ M\ Oi\ Os\ Out\ s\ security Context
            stateInterp\ cmd\ x\ ins\ outs.
```

Theorems SSM11 THEORY

```
\neg authentication Test \ x \Rightarrow
           TR' (M, Oi, Os) (discard cmd)
              (\mathtt{CFG}\ authentication\ Test\ stateInterp\ security\ Context
                  (x::ins) s outs)
              (CFG \ authentication Test \ stateInterp \ security Context
                 ins (NS \ s \ (discard \ cmd))
                  (Out \ s \ (discard \ cmd)::outs))) \Rightarrow
       \forall a_0 \ a_1 \ a_2 \ a_3. TR a_0 \ a_1 \ a_2 \ a_3 \Rightarrow TR' \ a_0 \ a_1 \ a_2 \ a_3
[TR_rules]
 \vdash (\forall authenticationTest\ P\ NS\ M\ Oi\ Os\ Out\ s\ securityContext
         stateInterp cmd ins outs.
        authenticationTest (P says prop (SOME cmd)) \land
        CFGInterpret (M, Oi, Os)
           (\mathtt{CFG}\ authentication\ Test\ stateInterp\ securityContext
               (P \text{ says prop } (SOME \ cmd) :: ins) \ s \ outs) \Rightarrow
        TR (M, Oi, Os) (exec cmd)
           (CFG authenticationTest stateInterp securityContext
               (P \text{ says prop } (SOME \ cmd)::ins) \ s \ outs)
           (CFG \ authentication Test \ stateInterp \ security Context \ ins
               (NS \ s \ (exec \ cmd)) \ (Out \ s \ (exec \ cmd)::outs))) \ \land
     (\forall authenticationTest\ P\ NS\ M\ Oi\ Os\ Out\ s\ securityContext
          stateInterp cmd ins outs.
        authenticationTest (P says prop (SOME cmd)) \land
        CFGInterpret (M, Oi, Os)
           (CFG \ authentication Test \ stateInterp \ security Context
               (P \text{ says prop (SOME } cmd)::ins) \ s \ outs) \Rightarrow
        TR (M, Oi, Os) (trap cmd)
           (CFG authenticationTest stateInterp securityContext
               (P says prop (SOME cmd)::ins) s outs)
           (CFG\ authentication Test\ state Interp\ security Context\ ins
               (NS s (trap cmd)) (Out s (trap cmd)::outs))) \land
    \forall authenticationTest \ NS \ M \ Oi \ Os \ Out \ s \ securityContext
        stateInterp\ cmd\ x\ ins\ outs.
       \neg authenticationTest \ x \Rightarrow
       TR (M, Oi, Os) (discard cmd)
          (\mathtt{CFG}\ authentication\ Test\ stateInterp\ securityContext
              (x::ins) s outs)
          (CFG \ authentication Test \ state Interp \ security Context \ ins
              (NS \ s \ (discard \ cmd)) \ (Out \ s \ (discard \ cmd)::outs))
[TR_strongind]
 \vdash \forall TR'.
       (\forall authentication Test\ P\ NS\ M\ Oi\ Os\ Out\ s\ security Context
            stateInterp cmd ins outs.
           authenticationTest (P says prop (SOME cmd)) \land
           CFGInterpret (M, Oi, Os)
              (CFG \ authentication Test \ stateInterp \ security Context
                  (P \text{ says prop } (SOME \ cmd)::ins) \ s \ outs) \Rightarrow
```

SSM11 THEORY Theorems

```
TR' (M, Oi, Os) (exec cmd)
             (CFG \ authentication Test \ stateInterp \ security Context
                 (P says prop (SOME cmd)::ins) s outs)
             (CFG authenticationTest stateInterp securityContext
                 ins (NS s (exec cmd)) (Out s (exec cmd)::outs))) \land
       (\forall authentication Test\ P\ NS\ M\ Oi\ Os\ Out\ s\ security Context
            stateInterp cmd ins outs.
           authenticationTest (P says prop (SOME cmd)) \land
          CFGInterpret (M, Oi, Os)
             (CFG \ authentication Test \ stateInterp \ security Context
                 (P \text{ says prop (SOME } cmd)::ins) \ s \ outs) \Rightarrow
           TR' (M, Oi, Os) (trap cmd)
             (CFG authenticationTest stateInterp securityContext
                 (P says prop (SOME cmd)::ins) s outs)
             ({\tt CFG}\ authentication Test\ state Interp\ security Context
                 ins (NS s (trap cmd)) (Out s (trap cmd)::outs))) \land
       (\forall authentication Test\ NS\ M\ Oi\ Os\ Out\ s\ security Context
            stateInterp\ cmd\ x\ ins\ outs.
           \neg authentication Test \ x \Rightarrow
           TR' (M, Oi, Os) (discard cmd)
             (CFG \ authentication Test \ stateInterp \ security Context
                 (x::ins) s outs)
             (CFG \ authentication Test \ stateInterp \ security Context
                 ins (NS \ s \ (discard \ cmd))
                 (Out \ s \ (discard \ cmd)::outs))) \Rightarrow
      \forall a_0 \ a_1 \ a_2 \ a_3. TR a_0 \ a_1 \ a_2 \ a_3 \Rightarrow TR' \ a_0 \ a_1 \ a_2 \ a_3
[TR_trap_cmd_rule]
 \vdash \ \forall \ authenticationTest \ \ stateInterp \ \ securityContext \ \ P \ \ cmd \ \ ins \ \ s
        outs.
       (\forall M \ Oi \ Os.
          CFGInterpret (M, Oi, Os)
             (CFG authenticationTest stateInterp securityContext
                 (P \text{ says prop } (SOME \ cmd) :: ins) \ s \ outs) \Rightarrow
           (M,Oi,Os) sat prop NONE) \Rightarrow
      \forall NS \ Out \ M \ Oi \ Os.
         TR (M, Oi, Os) (trap cmd)
            (\mathtt{CFG}\ authentication} \mathit{Test}\ stateInterp\ securityContext
                (P says prop (SOME cmd)::ins) s outs)
            (CFG\ authentication Test\ state Interp\ security Context\ ins
                (NS \ s \ (trap \ cmd)) \ (Out \ s \ (trap \ cmd)::outs)) \iff
         authenticationTest (P says prop (SOME cmd)) \land
         CFGInterpret (M, Oi, Os)
            (CFG authenticationTest stateInterp securityContext
                (P says prop (SOME cmd)::ins) s outs) \land
         (M, Oi, Os) sat prop NONE
[TRrule0]
 \vdash TR (M, Oi, Os) (exec cmd)
       (CFG authentication Test stateInterp security Context
```

```
(P says prop (SOME cmd)::ins) s outs)
       (CFG \ authentication Test \ stateInterp \ security Context \ ins
          (NS \ s \ (exec \ cmd)) \ (Out \ s \ (exec \ cmd)::outs)) \iff
    authenticationTest (P says prop (SOME cmd)) \land
    CFGInterpret (M, Oi, Os)
       (CFG \ authentication Test \ stateInterp \ security Context
          (P says prop (SOME cmd)::ins) s outs)
[TRrule1]
 \vdash TR (M, Oi, Os) (trap cmd)
       (CFG \ authentication Test \ stateInterp \ security Context
          (P says prop (SOME cmd)::ins) s outs)
       (\mathtt{CFG}\ authenticationTest\ stateInterp\ securityContext\ ins
          (NS \ s \ (trap \ cmd)) \ (Out \ s \ (trap \ cmd)::outs)) \iff
    authenticationTest (P says prop (SOME cmd)) \land
    CFGInterpret (M, Oi, Os)
       (\mathtt{CFG}\ authentication\ Test\ stateInterp\ security\ Context
          (P \text{ says prop } (SOME \ cmd) :: ins) \ s \ outs)
[trType_distinct_clauses]
 \vdash (\forall a' \ a. \ discard \ a \neq trap \ a') \land (\forall a' \ a. \ discard \ a \neq exec \ a') \land
    \forall a' \ a. \ \mathsf{trap} \ a \neq \mathsf{exec} \ a'
[trType_one_one]
 \vdash (\forall a \ a'. (discard a =  discard a') \iff (a = a')) \land
    (\forall a \ a'. \ (\text{trap} \ a = \text{trap} \ a') \iff (a = a')) \land
    \forall a \ a'. (exec a = \text{exec } a') \iff (a = a')
3
     ssm Theory
Built: 13 May 2018
Parent Theories: satList
3.1
       Datatypes
configuration =
     CFG (('command option, 'principal, 'd, 'e) Form -> bool)
          ('state ->
           ('command option, 'principal, 'd, 'e) Form list ->
           ('command option, 'principal, 'd, 'e) Form list)
          (('command option, 'principal, 'd, 'e) Form list ->
           ('command option, 'principal, 'd, 'e) Form list)
          (('command option, 'principal, 'd, 'e) Form list list)
```

'state ('output list)

trType = discard 'cmdlist | trap 'cmdlist | exec 'cmdlist

SSM THEORY Definitions

#### 3.2 Definitions

```
[authenticationTest_def]
 \vdash \forall elementTest x.
       \verb|authenticationTest|| elementTest|| x \iff
       FOLDR (\lambda p \ q. \ p \land q) T (MAP elementTest \ x)
[commandList_def]
 \vdash \ \forall \, x \,. commandList x = MAP extractCommand x
[inputList_def]
 \vdash \ \forall \, xs. inputList xs = MAP extractInput xs
[propCommandList_def]
 \vdash \ \forall \, x. propCommandList x = MAP extractPropCommand x
[TR_def]
 \vdash TR =
    (\lambda \ a_0 \ a_1 \ a_2 \ a_3.
        \forall TR'.
           (\forall a_0 \ a_1 \ a_2 \ a_3.
               (\exists elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x
                   (a_0 = (M, Oi, Os)) \land (a_1 = exec (inputList x)) \land
                   (a_2 =
                    CFG elementTest stateInterp context (x::ins) s
                       outs) \wedge
                   (a_3 =
                    {\tt CFG} elementTest stateInterp context ins
                       (NS \ s \ (exec \ (inputList \ x)))
                       (Out \ s (exec (inputList x))::outs)) \land
                  \verb|authenticationTest|| elementTest|| x \ \land
                  CFGInterpret (M, Oi, Os)
                     (CFG elementTest stateInterp context (x::ins) s
                         outs)) \
               (\exists elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x
                   (a_0 = (M, Oi, Os)) \land (a_1 = trap (inputList x)) \land
                    CFG elementTest stateInterp context (x::ins) s
                       outs) \wedge
                   (a_3 =
                    {\tt CFG} \ elementTest \ stateInterp \ context \ ins
                       (NS \ s \ (trap \ (inputList \ x)))
                       (Out \ s \ (trap \ (inputList \ x))::outs)) \ \land
                  \verb|authenticationTest|| elementTest||x| \wedge
                  CFGInterpret (M, Oi, Os)
                     (CFG elementTest stateInterp context (x::ins) s
```

Theorems SSM THEORY

 $outs)) \lor$ 

```
(\exists elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x
                     (a_0 = (M, Oi, Os)) \land (a_1 = discard (inputList x)) \land
                     (a_2 =
                      CFG elementTest stateInterp context (x::ins) s
                         outs) \land
                     (a_3 =
                      CFG elementTest stateInterp context ins
                         (NS \ s \ (discard \ (inputList \ x)))
                         (Out \ s (discard (inputList x))::outs)) \land
                     \negauthenticationTest elementTest x) \Rightarrow
                 TR' a_0 a_1 a_2 a_3) \Rightarrow
            TR' a_0 a_1 a_2 a_3)
3.3
        Theorems
[CFGInterpret_def]
 \vdash CFGInterpret (M, Oi, Os)
        (CFG elementTest stateInterp context (x::ins) state
            outStream) \iff
     (M,Oi,Os) satList context \ x \land (M,Oi,Os) satList x \land (M,Oi,Os)
     (M,Oi,Os) satList stateInterp state x
[CFGInterpret_ind]
 \vdash \forall P.
        (\forall M \ Oi \ Os \ elementTest \ stateInterp \ context \ x \ ins \ state
             outStream.
            P (M, Oi, Os)
               (CFG elementTest stateInterp context (x::ins) state
                   outStream)) \land
        (\forall v_{15} \ v_{10} \ v_{11} \ v_{12} \ v_{13} \ v_{14}.
            P \ v_{15} \ (\text{CFG} \ v_{10} \ v_{11} \ v_{12} \ [] \ v_{13} \ v_{14})) \ \Rightarrow
       \forall v \ v_1 \ v_2 \ v_3. P \ (v, v_1, v_2) \ v_3
[configuration_one_one]
 \vdash \forall a_0 \ a_1 \ a_2 \ a_3 \ a_4 \ a_5 \ a_0' \ a_1' \ a_2' \ a_3' \ a_4' \ a_5'.
        (CFG a_0 a_1 a_2 a_3 a_4 a_5 = CFG a_0' a_1' a_2' a_3' a_4' a_5') \iff
        (a_0 = a_0') \wedge (a_1 = a_1') \wedge (a_2 = a_2') \wedge (a_3 = a_3') \wedge
        (a_4 = a'_4) \wedge (a_5 = a'_5)
[extractCommand_def]
 \vdash extractCommand (P says prop (SOME cmd)) = cmd
[extractCommand_ind]
 \vdash \forall P'.
        (\forall P \ cmd. \ P' \ (P \ \text{says prop (SOME} \ cmd))) \ \land \ P' \ \text{TT} \ \land \ P' \ \text{FF} \ \land
        (\forall v_1. P' \text{ (prop } v_1)) \land (\forall v_3. P' \text{ (notf } v_3)) \land
```

SSM THEORY Theorems

```
(\forall v_6 \ v_7. \ P' \ (v_6 \ \text{andf} \ v_7)) \ \land \ (\forall v_{10} \ v_{11}. \ P' \ (v_{10} \ \text{orf} \ v_{11})) \ \land
           (\forall v_{14} \ v_{15}. \ P' \ (v_{14} \ \text{impf} \ v_{15})) \ \land
           (\forall v_{18} \ v_{19}. \ P' \ (v_{18} \ \text{eqf} \ v_{19})) \ \land \ (\forall v_{129}. \ P' \ (v_{129} \ \text{says} \ \text{TT})) \ \land
           (\forall v130. P' (v130 \text{ says FF})) \land
           (\forall v132. P' (v132 \text{ says prop NONE})) \land
           (\forall v133 v_{66}. P' (v133 says notf v_{66})) \wedge
           (\forall v134\ v_{69}\ v_{70}. P' (v134 says (v_{69} andf v_{70}))) \land
           (\forall v135 \ v_{73} \ v_{74}. \ P' \ (v135 \ \text{says} \ (v_{73} \ \text{orf} \ v_{74}))) \land
           (\forall v136 \ v_{77} \ v_{78}. \ P' \ (v136 \ \text{says} \ (v_{77} \ \text{impf} \ v_{78}))) \land
           (\forall v137 \ v_{81} \ v_{82}. \ P' \ (v137 \ \text{says} \ (v_{81} \ \text{eqf} \ v_{82}))) \ \land
           (\forall v138 \ v_{85} \ v_{86}. \ P' \ (v138 \ \text{says} \ v_{85} \ \text{says} \ v_{86})) \ \land
           (\forall v139 \ v_{89} \ v_{90}. \ P' \ (v139 \ \text{says} \ v_{89} \ \text{speaks\_for} \ v_{90})) \ \land
           (\forall v140 \ v_{93} \ v_{94}. \ P' \ (v140 \ \text{says} \ v_{93} \ \text{controls} \ v_{94})) \ \land
           (\forall v141 \ v_{98} \ v_{99} \ v100. \ P' \ (v141 \ {\tt says \ reps} \ v_{98} \ v_{99} \ v100)) \ \land
           (\forall v142 v103 v104. P' (v142 says v103 domi v104)) \land
           (\forall v143 \ v107 \ v108. \ P' \ (v143 \ \text{says} \ v107 \ \text{eqi} \ v108)) \ \land
           (\forall v144 \ v111 \ v112. \ P' \ (v144 \ \text{says} \ v111 \ \text{doms} \ v112)) \ \land
           (\forall v145 \ v115 \ v116. \ P' \ (v145 \ \text{says} \ v115 \ \text{eqs} \ v116)) \ \land
           (\forall v146 \ v119 \ v120. \ P' \ (v146 \ \text{says} \ v119 \ \text{eqn} \ v120)) \ \land
           (\forall v147 \ v123 \ v124. \ P' \ (v147 \ \text{says} \ v123 \ \text{lte} \ v124)) \ \land
           (\forall v148 \ v127 \ v128. \ P' \ (v148 \ \text{says} \ v127 \ \text{lt} \ v128)) \ \land
           (\forall v_{24} \ v_{25}. \ P' \ (v_{24} \ \text{speaks\_for} \ v_{25})) \ \land
           (\forall v_{28} \ v_{29}. P' (v_{28} controls v_{29})) \land
           (\forall v_{33} \ v_{34} \ v_{35}. \ P' \ (reps \ v_{33} \ v_{34} \ v_{35})) \ \land
           (\forall v_{38} \ v_{39}. \ P' \ (v_{38} \ \text{domi} \ v_{39})) \land
           (\forall v_{42} \ v_{43}. \ P' \ (v_{42} \ \text{eqi} \ v_{43})) \ \land
           (\forall v_{46} \ v_{47}. \ P' \ (v_{46} \ \text{doms} \ v_{47})) \ \land
           (\forall \, v_{50} \ v_{51}. P' (v_{50} eqs v_{51})) \wedge
           (\forall v_{54} \ v_{55}. \ P' \ (v_{54} \ \text{eqn} \ v_{55})) \ \land
           (\forall v_{58} \ v_{59}. \ P' \ (v_{58} \ \text{lte} \ v_{59})) \ \land
           (\forall v_{62} \ v_{63}. \ P' \ (v_{62} \ \text{lt} \ v_{63})) \Rightarrow
          \forall v. P' v
[extractInput_def]
  \vdash extractInput (P says prop x) = x
[extractInput_ind]
  \vdash \forall P'.
           (\forall P \ x. \ P' \ (P \ \text{says prop} \ x)) \ \land \ P' \ \text{TT} \ \land \ P' \ \text{FF} \ \land
           (\forall v_1.\ P'\ (\texttt{prop}\ v_1))\ \land\ (\forall v_3.\ P'\ (\texttt{notf}\ v_3))\ \land
           (\forall v_6 \ v_7. \ P' \ (v_6 \ \text{andf} \ v_7)) \land (\forall v_{10} \ v_{11}. \ P' \ (v_{10} \ \text{orf} \ v_{11})) \land
           (\forall v_{14} \ v_{15}. \ P' \ (v_{14} \ \text{impf} \ v_{15})) \ \land
           (\forall \, v_{18} \ v_{19}. P' (v_{18} eqf v_{19})) \land (\forall \, v129. P' (v129 says TT)) \land
           (\forall v130. P' (v130 \text{ says FF})) \land
           (\forall\,v131\ v_{66}. P' (v131 says notf v_{66})) \land
           (\forall\,v132\ v_{69}\ v_{70}. P' (v132\ \mathrm{says} (v_{69}\ \mathrm{andf}\ v_{70}))) \wedge
           (\forall v133 \ v_{73} \ v_{74}. \ P' \ (v133 \ \text{says} \ (v_{73} \ \text{orf} \ v_{74}))) \ \land
           (\forall v134 \ v_{77} \ v_{78}. \ P' \ (v134 \ \text{says} \ (v_{77} \ \text{impf} \ v_{78}))) \ \land
           (\forall v135 \ v_{81} \ v_{82}. \ P' \ (v135 \ \text{says} \ (v_{81} \ \text{eqf} \ v_{82}))) \land
```

Theorems SSM THEORY

```
(\forall\,v136\ v_{85}\ v_{86}. P' (v136\ \mathrm{says}\ v_{85}\ \mathrm{says}\ v_{86})) \wedge
          (\forall\,v137\ v_{89}\ v_{90}. P' (v137 says v_{89} speaks_for v_{90})) \wedge
          (\forall v138 \ v_{93} \ v_{94}. \ P' \ (v138 \ \text{says} \ v_{93} \ \text{controls} \ v_{94})) \ \land
          (\forall v139 \ v_{98} \ v_{99} \ v100. \ P' \ (v139 \ \text{says reps} \ v_{98} \ v_{99} \ v100)) \ \land
          (\forall \, v140 \ v103 \ v104 \,. \ P' \ (v140 \ {\tt says} \ v103 \ {\tt domi} \ v104)) \ \land
          (\forall\,v141\ v107\ v108. P' (v141\ \mathrm{says}\ v107\ \mathrm{eqi}\ v108)) \land
          ( \forall\,v142\ v111\ v112 . P' ( v142\ {\rm says}\ v111\ {\rm doms}\ v112 ) \land
          (\forall\,v143\ v115\ v116. P' (v143\ \mathrm{says}\ v115\ \mathrm{eqs}\ v116)) \land
          (\forall v144 v119 v120. P' (v144 says v119 eqn v120)) \land
          (\forall v145 \ v123 \ v124 . P' \ (v145 \ \text{says} \ v123 \ \text{lte} \ v124)) \land
          (\forall v146 \ v127 \ v128. \ P' \ (v146 \ \text{says} \ v127 \ \text{lt} \ v128)) \ \land
          (\forall v_{24} \ v_{25}. \ P' \ (v_{24} \ \text{speaks\_for} \ v_{25})) \land
          (\forall v_{28} \ v_{29}. P' (v_{28} controls v_{29})) \land
          (\forall v_{33} \ v_{34} \ v_{35}. \ P' \ (\text{reps} \ v_{33} \ v_{34} \ v_{35})) \ \land
          (\forall \, v_{38} \ v_{39}. P' (v_{38} domi v_{39})) \wedge
          (\forall v_{42} \ v_{43}. \ P' \ (v_{42} \ \text{eqi} \ v_{43})) \ \land
          (\forall v_{46} \ v_{47}. \ P' \ (v_{46} \ \text{doms} \ v_{47})) \land
          (\forall v_{50} \ v_{51}. \ P' \ (v_{50} \ \text{eqs} \ v_{51})) \ \land
          (\forall v_{54} \ v_{55}. \ P' \ (v_{54} \ \text{eqn} \ v_{55})) \ \land
          (\forall v_{58} \ v_{59}. \ P' \ (v_{58} \ \text{lte} \ v_{59})) \land
          (\forall v_{62} \ v_{63}. \ P' \ (v_{62} \ \text{lt} \ v_{63})) \Rightarrow
         \forall v. P' v
[extractPropCommand_def]
  \vdash extractPropCommand (P says prop (SOME cmd)) = prop (SOME cmd)
[extractPropCommand_ind]
  \vdash \forall P'.
          (\forall P \ cmd. \ P' \ (P \ \text{says prop (SOME} \ cmd))) \ \land \ P' \ \text{TT} \ \land \ P' \ \text{FF} \ \land
          (\forall v_{14} \ v_{15}. P' (v_{14} impf v_{15})) \land
          (\forall v_{18} \ v_{19}. \ P' \ (v_{18} \ \text{eqf} \ v_{19})) \ \land \ (\forall v_{129}. \ P' \ (v_{129} \ \text{says} \ \text{TT})) \ \land
          (\forall\,v130. P' (v130 says FF)) \land
          (\forall v132. P' (v132 says prop NONE)) \land
          (\forall v133 \ v_{66}. \ P' \ (v133 \ \text{says notf} \ v_{66})) \ \land
          (\forall v134 \ v_{69} \ v_{70}. \ P' \ (v134 \ \text{says} \ (v_{69} \ \text{andf} \ v_{70}))) \land
          (\forall v135 \ v_{73} \ v_{74}. \ P' \ (v135 \ \text{says} \ (v_{73} \ \text{orf} \ v_{74}))) \ \land
          (\forall\,v136\ v_{77}\ v_{78}.\ P' (v136\ \mathrm{says} (v_{77}\ \mathrm{impf}\ v_{78}))) \wedge
          (\forall v137 \ v_{81} \ v_{82}. \ P' \ (v137 \ \text{says} \ (v_{81} \ \text{eqf} \ v_{82}))) \ \land
          (\forall v138 \ v_{85} \ v_{86}. \ P' \ (v138 \ \text{says} \ v_{85} \ \text{says} \ v_{86})) \ \land
          (\forall v139 \ v_{89} \ v_{90}. \ P' \ (v139 \ \text{says} \ v_{89} \ \text{speaks\_for} \ v_{90})) \ \land
          (\forall v140 \ v_{93} \ v_{94}. P' (v140 says v_{93} controls v_{94})) \land
          (\forall v141 \ v_{98} \ v_{99} \ v100. \ P' \ (v141 \ \text{says reps} \ v_{98} \ v_{99} \ v100)) \ \land
          (\forall\,v142\ v103\ v104. P' (v142\ {\rm says}\ v103\ {\rm domi}\ v104)) \land
          (\forall\,v143\ v107\ v108. P' (v143\ \mathrm{says}\ v107\ \mathrm{eqi}\ v108)) \land
          (\forall v114 v111 v112. P' (v144 says v111 doms v112)) \land
          (\forall\,v145\ v115\ v116. P' (v145\ \mathrm{says}\ v115\ \mathrm{eqs}\ v116)) \land
          (\forall\,v146\ v119\ v120. P' (v146\ \mathrm{says}\ v119\ \mathrm{eqn}\ v120)) \wedge
```

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```
(\forall v147 \ v123 \ v124. \ P' \ (v147 \ \text{says} \ v123 \ \text{lte} \ v124)) \ \land
        (\forall v148 \ v127 \ v128. P' (v148 \ \mathrm{says} \ v127 \ \mathrm{lt} \ v128)) \wedge
        (\forall v_{24} \ v_{25}. \ P' \ (v_{24} \ \text{speaks\_for} \ v_{25})) \ \land
        (\forall \, v_{28} \ v_{29}. P' (v_{28} controls v_{29})) \land
        (\forall v_{33} v_{34} v_{35}. P' (reps v_{33} v_{34} v_{35})) \wedge
        (\forall\,v_{38}\;\;v_{39}. P' (v_{38}\; domi v_{39})) \wedge
        (\forall v_{42} \ v_{43}. \ P' \ (v_{42} \ \text{eqi} \ v_{43})) \ \land
        (\forall v_{46} \ v_{47}. \ P' \ (v_{46} \ \text{doms} \ v_{47})) \ \land
        (\forall v_{50} \ v_{51}. \ P' \ (v_{50} \ \text{eqs} \ v_{51})) \ \land
        (\forall v_{54} \ v_{55}. \ P' \ (v_{54} \ \text{eqn} \ v_{55})) \ \land
        (\forall v_{58} \ v_{59}. P' (v_{58} lte v_{59})) \wedge
        (\forall v_{62} \ v_{63}. \ P' \ (v_{62} \ \text{lt} \ v_{63})) \Rightarrow
        \forall v. P'v
[TR_cases]
 \vdash \forall a_0 \ a_1 \ a_2 \ a_3.
        TR a_0 a_1 a_2 a_3 \iff
        (\exists elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
             (a_0 = (M, Oi, Os)) \land (a_1 = exec (inputList x)) \land
             (a_2 =
              CFG elementTest stateInterp context (x::ins) s outs) \land
             (a_3 =
              {\tt CFG} \ elementTest \ stateInterp \ context \ ins
                 (NS \ s \ (exec \ (inputList \ x)))
                 (Out \ s \ (exec \ (inputList \ x))::outs)) \land
            authenticationTest elementTest x \land
            CFGInterpret (M, Oi, Os)
                (CFG elementTest stateInterp context (x::ins) s
                    outs)) ∨
        (\exists elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
             (a_0 = (M, Oi, Os)) \land (a_1 = trap (inputList x)) \land
             (a_2 =
              CFG elementTest stateInterp context (x::ins) s outs) \land
             (a_3 =
              CFG elementTest stateInterp context ins
                 (NS \ s \ (trap \ (inputList \ x)))
                 (Out \ s \ (trap \ (inputList \ x))::outs)) \ \land
            authenticationTest elementTest x \land
             CFGInterpret (M, Oi, Os)
                (CFG elementTest stateInterp context (x::ins) s
                    outs)) \lor
        \exists \ elementTest \ NS \ M \ Oi \ Os \ Out \ s \ context \ stateInterp \ x \ ins
           (a_0 = (M, Oi, Os)) \land (a_1 = discard (inputList x)) \land
           (a_2 =
            CFG elementTest stateInterp context (x::ins) s outs) \land
           (a_3 =
```

Theorems SSM THEORY

```
CFG elementTest stateInterp context ins
            (NS \ s \ (discard \ (inputList \ x)))
            (Out \ s (discard (inputList x))::outs)) \land
         \negauthenticationTest elementTest x
[TR_discard_cmd_rule]
 \vdash TR (M, Oi, Os) (discard (inputList x))
      (CFG elementTest stateInterp context (x::ins) s outs)
      (CFG elementTest stateInterp context ins
          (NS \ s \ (discard \ (inputList \ x)))
          (Out \ s \ (discard \ (inputList \ x))::outs)) \iff
    \negauthenticationTest elementTest x
[TR_EQ_rules_thm]
 \vdash (TR (M, Oi, Os) (exec (inputList x))
       (CFG elementTest stateInterp context (x::ins) s outs)
        (CFG elementTest stateInterp context ins
           (NS \ s \ (exec \ (inputList \ x)))
           (Out \ s \ (exec \ (inputList \ x))::outs)) \iff
     authenticationTest elementTest \ x \ \land
     CFGInterpret (M, Oi, Os)
        (CFG elementTest stateInterp context (x::ins) s outs)) \land
    (TR (M, Oi, Os) (trap (inputList x))
       (CFG elementTest stateInterp context (x::ins) s outs)
        (CFG elementTest stateInterp context ins
           (NS \ s \ (trap \ (inputList \ x)))
           (Out \ s \ (trap \ (inputList \ x))::outs)) \iff
     \verb|authenticationTest|| elementTest|| x \ \land
     CFGInterpret (M, Oi, Os)
        (CFG elementTest stateInterp context (x::ins) s outs)) \land
    (TR (M, Oi, Os) (discard (inputList x))
       (CFG elementTest stateInterp context (x::ins) s outs)
        (CFG elementTest stateInterp context ins
           (NS \ s \ (discard \ (inputList \ x)))
           (Out \ s \ (discard \ (inputList \ x))::outs)) \iff
     \negauthenticationTest elementTest x)
[TR_exec_cmd_rule]
 \vdash \forall elementTest \ context \ stateInterp \ x \ ins \ s \ outs.
      (\forall M \ Oi \ Os.
         CFGInterpret (M, Oi, Os)
            (CFG elementTest stateInterp context (x::ins) s
                outs) \Rightarrow
          (M, Oi, Os) satList propCommandList x) \Rightarrow
      \forall NS \ Out \ M \ Oi \ Os.
        TR (M, Oi, Os) (exec (inputList x))
           (CFG elementTest stateInterp context (x::ins) s outs)
           (CFG elementTest stateInterp context ins
```

SSM THEORY Theorems

```
(NS \ s \ (exec \ (inputList \ x)))
               (Out \ s \ (exec \ (inputList \ x))::outs)) \iff
         authenticationTest elementTest \ x \ \land
         CFGInterpret (M, Oi, Os)
            (CFG elementTest stateInterp context (x::ins) s outs) \land
         (M,Oi,Os) satList propCommandList x
[TR_ind]
 \vdash \ \forall \ TR'.
       (\forall elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
            outs.
          \verb|authenticationTest|| elementTest|| x \ \land
          CFGInterpret (M, Oi, Os)
             (CFG elementTest stateInterp context (x::ins) s
                 outs) \Rightarrow
          TR' (M, Oi, Os) (exec (inputList x))
             (CFG elementTest stateInterp context (x::ins) s outs)
             (CFG elementTest stateInterp context ins
                (NS \ s \ (exec \ (inputList \ x)))
                (Out \ s \ (exec \ (inputList \ x))::outs))) \land
       (\forall elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
          authenticationTest elementTest x \land
          CFGInterpret (M, Oi, Os)
             (CFG elementTest stateInterp context (x::ins) s
                outs) \Rightarrow
          TR' (M, Oi, Os) (trap (inputList x))
             (CFG elementTest stateInterp context (x::ins) s outs)
             (CFG elementTest stateInterp context ins
                 (NS \ s \ (trap \ (inputList \ x)))
                (Out \ s \ (trap \ (inputList \ x))::outs))) \land
       (\forall elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
          \negauthenticationTest elementTest x \Rightarrow
          TR' (M, Oi, Os) (discard (inputList x))
             (CFG elementTest stateInterp context (x::ins) s outs)
             (CFG elementTest stateInterp context ins
                (NS \ s \ (discard \ (inputList \ x)))
                (Out \ s \ (discard \ (inputList \ x))::outs))) \Rightarrow
      \forall a_0 \ a_1 \ a_2 \ a_3. TR a_0 \ a_1 \ a_2 \ a_3 \Rightarrow TR' \ a_0 \ a_1 \ a_2 \ a_3
[TR_rules]
 \vdash (\forall elementTest NS M Oi Os Out s context stateInterp x ins
         outs.
        \verb|authenticationTest|| elementTest||x| \wedge
        CFGInterpret (M, Oi, Os)
          (CFG elementTest stateInterp context (x::ins) s outs) \Rightarrow
        TR (M, Oi, Os) (exec (inputList x))
          (CFG elementTest stateInterp context (x::ins) s outs)
```

Theorems SSM THEORY

```
(CFG elementTest stateInterp context ins
              (NS \ s \ (exec \ (inputList \ x)))
              (Out \ s \ (exec \ (inputList \ x))::outs))) \land
    (\forall elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
       \verb|authenticationTest|| elementTest|| x \ \land
       CFGInterpret (M, Oi, Os)
          (CFG elementTest\ stateInterp\ context\ (x::ins)\ s\ outs) \Rightarrow
       TR (M, Oi, Os) (trap (inputList x))
          (CFG elementTest stateInterp context (x::ins) s outs)
          (CFG elementTest stateInterp context ins
              (NS \ s \ (trap \ (inputList \ x)))
              (Out \ s \ (trap \ (inputList \ x))::outs))) \land
    \forall elementTest NS M Oi Os Out s context stateInterp x ins outs.
       \negauthenticationTest elementTest x \Rightarrow
      TR (M, Oi, Os) (discard (inputList x))
         (CFG elementTest stateInterp context (x::ins) s outs)
         (CFG elementTest stateInterp context ins
            (NS \ s \ (discard \ (inputList \ x)))
            (Out s (discard (inputList x))::outs))
[TR_strongind]
 \vdash \forall TR'.
       (\forall elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
           outs.
          authenticationTest elementTest x \land
          CFGInterpret (M, Oi, Os)
            (CFG elementTest stateInterp context (x::ins) s
                outs) \Rightarrow
          TR' (M,Oi,Os) (exec (inputList x))
            (CFG elementTest stateInterp context (x::ins) s outs)
            (CFG elementTest\ stateInterp\ context\ ins
                (NS \ s \ (exec \ (inputList \ x)))
                (Out \ s \ (exec \ (inputList \ x))::outs))) \land
       (\forall elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
          authenticationTest elementTest x \land
          CFGInterpret (M, Oi, Os)
            (CFG elementTest stateInterp context (x::ins) s
                outs) \Rightarrow
          TR' (M, Oi, Os) (trap (inputList x))
            (CFG elementTest stateInterp context (x::ins) s outs)
            (CFG elementTest\ stateInterp\ context\ ins
                (NS \ s \ (trap \ (inputList \ x)))
                (Out \ s \ (trap \ (inputList \ x))::outs))) \land
       (\forall elementTest\ NS\ M\ Oi\ Os\ Out\ s\ context\ stateInterp\ x\ ins
          \negauthenticationTest elementTest x \Rightarrow
          TR' (M, Oi, Os) (discard (inputList x))
```

SSM THEORY Theorems

```
(CFG elementTest stateInterp context (x::ins) s outs)
             (CFG elementTest stateInterp context ins
                 (NS \ s \ (discard \ (inputList \ x)))
                 (Out \ s \ (discard \ (inputList \ x))::outs))) \Rightarrow
      \forall a_0 \ a_1 \ a_2 \ a_3. TR a_0 \ a_1 \ a_2 \ a_3 \Rightarrow TR' \ a_0 \ a_1 \ a_2 \ a_3
[TR_trap_cmd_rule]
 \vdash \ \forall \ elementTest \ \ context \ \ stateInterp \ \ x \ \ ins \ \ s \ \ outs.
       (\forall M \ Oi \ Os.
          CFGInterpret (M, Oi, Os)
             (CFG elementTest stateInterp context (x::ins) s
                 outs) \Rightarrow
           (M,Oi,Os) sat prop NONE) \Rightarrow
      \forall NS \ Out \ M \ Oi \ Os.
         TR (M, Oi, Os) (trap (inputList x))
            (CFG elementTest stateInterp context (x::ins) s outs)
            (CFG elementTest stateInterp context ins
                (NS \ s \ (trap \ (inputList \ x)))
                (Out \ s \ (trap \ (inputList \ x))::outs)) \iff
         authenticationTest elementTest x \land
         CFGInterpret (M, Oi, Os)
            (CFG elementTest stateInterp context (x::ins) s outs) \land
         (M,Oi,Os) sat prop NONE
[TRrule0]
 \vdash TR (M, Oi, Os) (exec (inputList x))
       (CFG elementTest stateInterp context (x::ins) s outs)
       (CFG elementTest stateInterp context ins
           (NS \ s \ (exec \ (inputList \ x)))
           (Out \ s \ (exec \ (inputList \ x))::outs)) \iff
    \verb|authenticationTest|| elementTest|| x \ \land
    CFGInterpret (M, Oi, Os)
       (CFG elementTest stateInterp context (x::ins) s outs)
[TRrule1]
 \vdash TR (M, Oi, Os) (trap (inputList x))
       (CFG elementTest stateInterp context (x::ins) s outs)
       (CFG elementTest stateInterp context ins
           (NS \ s \ (trap \ (inputList \ x)))
           (Out \ s \ (trap \ (inputList \ x))::outs)) \iff
    authenticationTest elementTest x \land
    CFGInterpret (M, Oi, Os)
       (CFG elementTest stateInterp context (x::ins) s outs)
[trType_distinct_clauses]
 \vdash (\forall a' \ a. discard a \neq \text{trap } a') \land (\forall a' \ a. discard a \neq \text{exec } a') \land
    \forall a' \ a. \ \mathsf{trap} \ a \neq \mathsf{exec} \ a'
```

```
[trType_one_one] 
 \vdash (\forall a \ a') (\text{discard } a = \text{discard } a') \iff (a = a')) \land (\forall a \ a') (\text{trap } a = \text{trap } a') \iff (a = a')) \land \forall a \ a'. (\text{exec } a = \text{exec } a') \iff (a = a')
```

## 4 satList Theory

**Built:** 13 May 2018

Parent Theories: aclDrules

#### 4.1 Definitions

## 5 ssmPB Theory

**Built:** 13 May 2018

Parent Theories: PBType, ssm11, OMNIType

## 5.1 Definitions

SSMPB THEORY Theorems

#### 5.2 Theorems

```
[authenticationTest_cmd_reject_lemma]
 \vdash \forall cmd. \neg authenticationTest (prop (SOME cmd))
authenticationTest_def
 \vdash (authenticationTest (Name PlatoonLeader says prop cmd) \iff
     T) \land (authenticationTest TT \iff F) \land
    (authenticationTest FF \iff F) \land
    (authenticationTest (prop v) \iff F) \land
    (authenticationTest (notf v_1) \iff F) \land
    (authenticationTest (v_2 andf v_3) \iff F) \wedge
    (authenticationTest (v_4 orf v_5) \iff F) \land
    (authenticationTest (v_6 impf v_7) \iff F) \land
    (authenticationTest (v_8 eqf v_9) \iff F) \land
    (authenticationTest (v_{10} says TT) \iff F) \wedge
    (authenticationTest (v_{10} says FF) \iff F) \wedge
    (authenticationTest (v133 meet v134 says prop v_{66}) \iff F) \land
    (authenticationTest (v135 quoting v136 says prop v_{66}) \iff F) \land
    (authenticationTest (v_{10} says notf v_{67}) \iff F) \land
    (authenticationTest (v_{10} says (v_{68} andf v_{69})) \iff F) \land
    (authenticationTest (v_{10} says (v_{70} orf v_{71})) \iff F) \land
    (authenticationTest (v_{10} says (v_{72} impf v_{73})) \iff F) \wedge
    (authenticationTest (v_{10} says (v_{74} eqf v_{75})) \iff F) \wedge
    (authenticationTest (v_{10} says v_{76} says v_{77}) \iff F) \land
    (authenticationTest (v_{10} says v_{78} speaks_for v_{79}) \iff F) \land
    (authenticationTest (v_{10} says v_{80} controls v_{81}) \iff F) \wedge
    (authenticationTest (v_{10} says reps v_{82} v_{83} v_{84}) \iff F) \land
    (authenticationTest (v_{10} says v_{85} domi v_{86}) \iff F) \wedge
    (authenticationTest (v_{10} says v_{87} eqi v_{88}) \iff F) \land
    (authenticationTest (v_{10} says v_{89} doms v_{90}) \iff F) \wedge
    (authenticationTest (v_{10} says v_{91} eqs v_{92}) \iff F) \wedge
    (authenticationTest (v_{10} says v_{93} eqn v_{94}) \iff F) \wedge
    (authenticationTest (v_{10} says v_{95} lte v_{96}) \iff F) \wedge
    (authenticationTest (v_{10} says v_{97} lt v_{98}) \iff F) \land
    (authenticationTest (v_{12} speaks_for v_{13}) \iff F) \wedge
    (authenticationTest (v_{14} controls v_{15}) \iff F) \wedge
    (authenticationTest (reps v_{16} v_{17} v_{18}) \iff F) \wedge
    (authenticationTest (v_{19} domi v_{20}) \iff F) \wedge
    (authenticationTest (v_{21} eqi v_{22}) \iff F) \land
    (authenticationTest (v_{23} doms v_{24}) \iff F) \wedge
    (authenticationTest (v_{25} eqs v_{26}) \iff F) \wedge
    (authenticationTest (v_{27} eqn v_{28}) \iff F) \wedge
    (authenticationTest (v_{29} lte v_{30}) \iff F) \land
    (authenticationTest (v_{31} lt v_{32}) \iff F)
authenticationTest_ind
 \vdash \forall P.
      (\forall cmd. P \text{ (Name PlatoonLeader says prop } cmd)) \land P \text{ TT } \land
```

Theorems SSMPB THEORY

```
P \text{ FF } \wedge (\forall v. P \text{ (prop } v)) \wedge (\forall v_1. P \text{ (notf } v_1)) \wedge
          (\forall v_2 \ v_3. \ P \ (v_2 \ \text{andf} \ v_3)) \ \land \ (\forall v_4 \ v_5. \ P \ (v_4 \ \text{orf} \ v_5)) \ \land
          (\forall v_6 \ v_7. \ P \ (v_6 \ \text{impf} \ v_7)) \ \land \ (\forall v_8 \ v_9. \ P \ (v_8 \ \text{eqf} \ v_9)) \ \land
          (\forall v_{10}. \ P \ (v_{10} \ \text{says TT})) \ \land \ (\forall v_{10}. \ P \ (v_{10} \ \text{says FF})) \ \land
          (\forall v133 \ v134 \ v_{66}. \ P \ (v133 \ \text{meet} \ v134 \ \text{says prop} \ v_{66})) \ \land
          (\forall\,v135\ v136\ v_{66}. P (v135 quoting v136 says prop v_{66})) \wedge
          (\forall v_{10} \ v_{67}. P (v_{10} says notf v_{67})) \land
          (\forall v_{10} \ v_{68} \ v_{69}. \ P \ (v_{10} \ \text{says} \ (v_{68} \ \text{andf} \ v_{69}))) \ \land
          (\forall v_{10} \ v_{70} \ v_{71}. \ P \ (v_{10} \ \text{says} \ (v_{70} \ \text{orf} \ v_{71}))) \land
          (\forall v_{10} \ v_{72} \ v_{73}. \ P \ (v_{10} \ \text{says} \ (v_{72} \ \text{impf} \ v_{73}))) \ \land
          (\forall v_{10} \ v_{74} \ v_{75}. \ P \ (v_{10} \ \text{says} \ (v_{74} \ \text{eqf} \ v_{75}))) \land
          (\forall v_{10} \ v_{76} \ v_{77}. \ P \ (v_{10} \ \text{says} \ v_{76} \ \text{says} \ v_{77})) \ \land
          (\forall\,v_{10}\,v_{78}\,v_{79}. P (v_{10} says v_{78} speaks_for v_{79})) \wedge
          (\forall v_{10} \ v_{80} \ v_{81}. \ P \ (v_{10} \ \text{says} \ v_{80} \ \text{controls} \ v_{81})) \ \land
          (\forall v_{10} \ v_{82} \ v_{83} \ v_{84}. P (v_{10} says reps v_{82} \ v_{83} \ v_{84})) \land
          (\forall v_{10} \ v_{85} \ v_{86}. \ P \ (v_{10} \ \text{says} \ v_{85} \ \text{domi} \ v_{86})) \ \land
          (\forall v_{10} \ v_{89} \ v_{90}. \ P \ (v_{10} \ {\tt says} \ v_{89} \ {\tt doms} \ v_{90})) \ \land
          (\forall \, v_{10} \ v_{91} \ v_{92}. P (v_{10} says v_{91} eqs v_{92})) \wedge
          (\forall v_{10} \ v_{93} \ v_{94}. \ P \ (v_{10} \ \text{says} \ v_{93} \ \text{eqn} \ v_{94})) \ \land
          (\forall v_{10} \ v_{95} \ v_{96}. P (v_{10} says v_{95} lte v_{96})) \land
          (\forall v_{10} \ v_{97} \ v_{98}. \ P \ (v_{10} \ \text{says} \ v_{97} \ \text{lt} \ v_{98})) \ \land
          (\forall v_{12} \ v_{13}. \ P \ (v_{12} \ \text{speaks\_for} \ v_{13})) \land
          (\forall v_{14} \ v_{15}. \ P \ (v_{14} \ \text{controls} \ v_{15})) \land
          (\forall v_{16} \ v_{17} \ v_{18}. \ P \ (\text{reps} \ v_{16} \ v_{17} \ v_{18})) \ \land
          (\forall v_{19} \ v_{20}. \ P \ (v_{19} \ \text{domi} \ v_{20})) \ \land
          (\forall v_{21} \ v_{22}. \ P \ (v_{21} \ \text{eqi} \ v_{22})) \ \land
          (\forall v_{23} \ v_{24}. \ P \ (v_{23} \ \text{doms} \ v_{24})) \ \land
          (\forall v_{25} \ v_{26}. \ P \ (v_{25} \ \text{eqs} \ v_{26})) \ \land \ (\forall v_{27} \ v_{28}. \ P \ (v_{27} \ \text{eqn} \ v_{28})) \ \land
          (\forall v_{29} \ v_{30}. \ P \ (v_{29} \ \text{lte} \ v_{30})) \land (\forall v_{31} \ v_{32}. \ P \ (v_{31} \ \text{lt} \ v_{32})) \Rightarrow
         \forall v. P v
[PBNS_def]
  ⊢ (PBNS PLAN_PB (exec (SLc crossLD)) = MOVE_TO_ORP) ∧
      (PBNS PLAN_PB (exec (SLc incomplete)) = PLAN_PB) \(\lambda\)
      (PBNS MOVE_TO_ORP (exec (SLc conductORP)) = CONDUCT_ORP) \(\lambda\)
      (PBNS MOVE_TO_ORP (exec (SLc incomplete)) = MOVE_TO_ORP) \(\lambda\)
      (PBNS CONDUCT_ORP (exec (SLc moveToPB)) = MOVE_TO_PB) \(\lambda\)
      (PBNS CONDUCT_ORP (exec (SLc incomplete)) = CONDUCT_ORP) \(\lambda\)
      (PBNS MOVE_TO_PB (exec (SLc conductPB)) = CONDUCT_PB) \( \)
      (PBNS MOVE_TO_PB (exec (SLc incomplete)) = MOVE_TO_PB) \(\lambda\)
      (PBNS CONDUCT_PB (exec (SLc completePB)) = COMPLETE_PB) \land
      (PBNS CONDUCT_PB (exec (SLc incomplete)) = CONDUCT_PB) \( \)
      (PBNS s (trap (SLc cmd)) = s) \land
      (PBNS s (discard (SLc cmd)) = s)
[PBNS_ind]
  \vdash \forall P.
          P PLAN_PB (exec (SLc crossLD)) \wedge
```

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```
P PLAN_PB (exec (SLc incomplete)) \wedge
       P MOVE_TO_ORP (exec (SLc conductORP)) \wedge
      P MOVE_TO_ORP (exec (SLc incomplete)) \wedge
      P CONDUCT_ORP (exec (SLc moveToPB)) \wedge
      P CONDUCT_ORP (exec (SLc incomplete)) \wedge
      P MOVE_TO_PB (exec (SLc conductPB)) \wedge
       P MOVE_TO_PB (exec (SLc incomplete)) \wedge
       P CONDUCT_PB (exec (SLc completePB)) \wedge
      P CONDUCT_PB (exec (SLc incomplete)) \wedge
       (\forall s \ cmd. \ P \ s \ (trap \ (SLc \ cmd))) \land
       (\forall s \ v_6. \ P \ s \ (\texttt{discard} \ (\texttt{ESCc} \ v_6))) \ \land
       (\forall s \ v_9. \ P \ s \ (trap \ (ESCc \ v_9))) \ \land
       (\forall v_{12}. P PLAN_PB (exec (ESCc v_{12}))) \land
       P PLAN_PB (exec (SLc conductORP)) \wedge
       P PLAN_PB (exec (SLc moveToPB)) \wedge
       P PLAN_PB (exec (SLc conductPB)) \wedge
       P PLAN_PB (exec (SLc completePB)) \wedge
       (\forall v_{15}. \ P \ \texttt{MOVE\_TO\_ORP} \ (\texttt{exec} \ (\texttt{ESCc} \ v_{15}))) \ \land
       P MOVE_TO_ORP (exec (SLc crossLD)) \wedge
       P MOVE_TO_ORP (exec (SLc moveToPB)) \wedge
       P MOVE_TO_ORP (exec (SLc conductPB)) \wedge
       P MOVE_TO_ORP (exec (SLc completePB)) \wedge
       (\forall v_{18}. \ P \ \texttt{CONDUCT\_ORP} \ (\texttt{exec} \ (\texttt{ESCc} \ v_{18}))) \ \land
       P CONDUCT_ORP (exec (SLc crossLD)) \wedge
      P CONDUCT_ORP (exec (SLc conductORP)) \wedge
      P CONDUCT_ORP (exec (SLc conductPB)) \wedge
      P CONDUCT_ORP (exec (SLc completePB)) \wedge
       (\forall v_{21}. P \text{ MOVE\_TO\_PB (exec (ESCc } v_{21}))) \land
       P MOVE_TO_PB (exec (SLc crossLD)) \wedge
       P \text{ MOVE\_TO\_PB (exec (SLc conductORP))} \land
       P MOVE_TO_PB (exec (SLc moveToPB)) \land
       P MOVE_TO_PB (exec (SLc completePB)) \wedge
       (\forall v_{24}.\ P CONDUCT_PB (exec (ESCc v_{24}))) \land
       P CONDUCT_PB (exec (SLc crossLD)) \wedge
      P CONDUCT_PB (exec (SLc conductORP)) \wedge
       P CONDUCT_PB (exec (SLc moveToPB)) \wedge
       P CONDUCT_PB (exec (SLc conductPB)) \wedge
       (\forall v_{26}. \ P \ \texttt{COMPLETE\_PB} \ (\texttt{exec} \ v_{26})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
[PBOut_def]
 ⊢ (PBOut PLAN_PB (exec (SLc crossLD)) = MoveToORP) ∧
    (PBOut PLAN_PB (exec (SLc incomplete)) = PlanPB) \(\lambda\)
    (PBOut MOVE_TO_ORP (exec (SLc conductORP)) = ConductORP) \( \)
    (PBOut MOVE_TO_ORP (exec (SLc incomplete)) = MoveToORP) \(\lambda\)
    (PBOut CONDUCT_ORP (exec (SLc moveToPB)) = MoveToPB) \(\lambda\)
    (PBOut CONDUCT_ORP (exec (SLc incomplete)) = ConductORP) \(\lambda\)
    (PBOut MOVE_TO_PB (exec (SLc conductPB)) = ConductPB) \(\lambda\)
```

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```
(PBOut MOVE_TO_PB (exec (SLc incomplete)) = MoveToPB) \(\lambda\)
    (PBOut CONDUCT_PB (exec (SLc completePB)) = CompletePB) \(\lambda\)
    (PBOut CONDUCT_PB (exec (SLc incomplete)) = ConductPB) \land
    (PBOut s (trap (SLc cmd)) = unAuthorized) \land
    (PBOut s (discard (SLc cmd)) = unAuthenticated)
[PBOut_ind]
 \vdash \forall P.
       P PLAN_PB (exec (SLc crossLD)) \wedge
       P PLAN_PB (exec (SLc incomplete)) \wedge
       P MOVE_TO_ORP (exec (SLc conductORP)) \wedge
       P MOVE_TO_ORP (exec (SLc incomplete)) \wedge
       P CONDUCT_ORP (exec (SLc moveToPB)) \wedge
       P CONDUCT_ORP (exec (SLc incomplete)) \wedge
       P MOVE_TO_PB (exec (SLc conductPB)) \wedge
       P MOVE_TO_PB (exec (SLc incomplete)) \wedge
       P CONDUCT_PB (exec (SLc completePB)) \wedge
       P CONDUCT_PB (exec (SLc incomplete)) \wedge
       (\forall s \ cmd. \ P \ s \ (trap \ (SLc \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (discard \ (SLc \ cmd))) \ \land
       (\forall s \ v_6. \ P \ s \ (discard \ (ESCc \ v_6))) \ \land
       (\forall s \ v_9. \ P \ s \ (trap \ (ESCc \ v_9))) \ \land
       (\forall v_{12}. P PLAN_PB (exec (ESCc v_{12}))) \land
       P PLAN_PB (exec (SLc conductORP)) \wedge
       P PLAN_PB (exec (SLc moveToPB)) \wedge
       P PLAN_PB (exec (SLc conductPB)) \wedge
       P PLAN_PB (exec (SLc completePB)) \wedge
       (\forall v_{15}. \ P \ \texttt{MOVE\_TO\_ORP} \ (\texttt{exec} \ (\texttt{ESCc} \ v_{15}))) \ \land
       P MOVE_TO_ORP (exec (SLc crossLD)) \wedge
       P MOVE_TO_ORP (exec (SLc moveToPB)) \wedge
       P MOVE_TO_ORP (exec (SLc conductPB)) \land
       P MOVE_TO_ORP (exec (SLc completePB)) \wedge
       (\forall v_{18}. \ P \ \texttt{CONDUCT\_ORP} \ (\texttt{exec} \ (\texttt{ESCc} \ v_{18}))) \ \land
       P CONDUCT_ORP (exec (SLc crossLD)) \wedge
       P CONDUCT_ORP (exec (SLc conductORP)) \wedge
       P CONDUCT_ORP (exec (SLc conductPB)) \wedge
       P CONDUCT_ORP (exec (SLc completePB)) \wedge
       (\forall v_{21}.\ P MOVE_TO_PB (exec (ESCc v_{21}))) \land
       P MOVE_TO_PB (exec (SLc crossLD)) \wedge
       P MOVE_TO_PB (exec (SLc conductORP)) \wedge
       P MOVE_TO_PB (exec (SLc moveToPB)) \wedge
       P MOVE_TO_PB (exec (SLc completePB)) \wedge
       (\forall v_{24}. \ P \ \texttt{CONDUCT\_PB} \ (\texttt{exec} \ (\texttt{ESCc} \ v_{24}))) \ \land
       P CONDUCT_PB (exec (SLc crossLD)) \wedge
       P CONDUCT_PB (exec (SLc conductORP)) \wedge
       P CONDUCT_PB (exec (SLc moveToPB)) \wedge
       P CONDUCT_PB (exec (SLc conductPB)) \wedge
       (\forall v_{26}. \ P \ \texttt{COMPLETE\_PB} \ (\texttt{exec} \ v_{26})) \Rightarrow
       \forall v \ v_1 . \ P \ v \ v_1
```

```
[PlatoonLeader_exec_slCommand_justified_thm]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os) (exec (SLc slCommand))
        (CFG authenticationTest ssmPBStateInterp
           (secContext slCommand)
           (Name PlatoonLeader says prop (SOME (SLc slCommand))::
                ins) s outs)
        (CFG authenticationTest ssmPBStateInterp
           (secContext slCommand) ins
           (NS \ s \ (exec \ (SLc \ slCommand)))
           (Out \ s \ (exec \ (SLc \ slCommand))::outs)) \iff
     authenticationTest
        (Name PlatoonLeader says prop (SOME (SLc slCommand))) \land
     CFGInterpret (M, Oi, Os)
        (CFG authenticationTest ssmPBStateInterp
           (secContext slCommand)
           (Name PlatoonLeader says prop (SOME (SLc slCommand))::
                ins) s outs) \wedge
      (M, Oi, Os) sat prop (SOME (SLc slCommand))
[PlatoonLeader_slCommand_lemma]
 \vdash CFGInterpret (M, Oi, Os)
      (CFG authenticationTest ssmPBStateInterp
         (secContext slCommand)
         (Name PlatoonLeader says prop (SOME (SLc slCommand))::
              ins) s outs) \Rightarrow
    (M, Oi, Os) sat prop (SOME (SLc slCommand))
```

# 6 PBTypeIntegrated Theory

**Built:** 13 May 2018

Parent Theories: OMNIType

### 6.1 Datatypes

```
| CONDUCT_PB | COMPLETE_PB
stateRole = PlatoonLeader | Omni
6.2
        Theorems
[omniCommand_distinct_clauses]
 \vdash ssmPlanPBComplete \neq ssmMoveToORPComplete \land
     {\tt ssmPlanPBComplete} \, \neq \, {\tt ssmConductORPComplete} \, \, \land \, \,
     {\tt ssmPlanPBComplete} \, \neq \, {\tt ssmMoveToPBComplete} \, \, \land \, \,
     {\tt ssmPlanPBComplete} \, \neq \, {\tt ssmConductPBComplete} \, \, \wedge \,
     ssmPlanPBComplete \neq invalidOmniCommand \land
     {\tt ssmMoveToORPComplete} \, \neq \, {\tt ssmConductORPComplete} \, \, \wedge \,
     {\tt ssmMoveToORPComplete} \, \neq \, {\tt ssmMoveToPBComplete} \, \, \wedge \,
     ssmMoveToORPComplete \neq ssmConductPBComplete \land
     ssmMoveToORPComplete \neq invalidOmniCommand \land
     {\tt ssmConductORPComplete} \, \neq \, {\tt ssmMoveToPBComplete} \, \, \wedge \,
     {\tt ssmConductORPComplete} \, \neq \, {\tt ssmConductPBComplete} \, \, \wedge \,
     {\tt ssmConductORPComplete} \neq {\tt invalidOmniCommand} \ \land \\
     {\tt ssmMoveToPBComplete} \, \neq \, {\tt ssmConductPBComplete} \, \, \wedge \,
     {\tt ssmMoveToPBComplete} \, \neq \, {\tt invalidOmniCommand} \, \, \land \, \,
     ssmConductPBComplete \neq invalidOmniCommand
[plCommand_distinct_clauses]
 \vdash crossLD \neq conductORP \land crossLD \neq moveToPB \land
     \texttt{crossLD} \neq \texttt{conductPB} \ \land \ \texttt{crossLD} \neq \texttt{completePB} \ \land
     \verb|crossLD| \neq \verb|incomplete| \land \verb|conductORP| \neq \verb|moveToPB| \land
     conductORP \neq conductPB \land conductORP \neq completePB \land
     \mathtt{conductORP} \, \neq \, \mathtt{incomplete} \, \, \land \, \, \mathtt{moveToPB} \, \neq \, \mathtt{conductPB} \, \, \land \, \,
     moveToPB \neq completePB \land moveToPB \neq incomplete \land
     conductPB \neq completePB \land conductPB \neq incomplete \land
     completePB \neq incomplete
[slCommand_distinct_clauses]
 \vdash \forall a' \ a. \ PL \ a \neq OMNI \ a'
[slCommand_one_one]
 \vdash (\forall a \ a'. (PL \ a = PL \ a') \iff (a = a')) \land
     \forall a \ a'. (OMNI a = OMNI \ a') \iff (a = a')
[slOutput_distinct_clauses]
 \vdash PlanPB \neq MoveToORP \land PlanPB \neq ConductORP \land
     {\tt PlanPB} \neq {\tt MoveToPB} \ \land \ {\tt PlanPB} \neq {\tt ConductPB} \ \land \\
     PlanPB \neq CompletePB \land PlanPB \neq unAuthenticated \land
     {\tt PlanPB} \, \neq \, {\tt unAuthorized} \, \wedge \, {\tt MoveToORP} \, \neq \, {\tt ConductORP} \, \wedge \,
```

 $\texttt{MoveToORP} \, \neq \, \texttt{MoveToPB} \, \wedge \, \texttt{MoveToORP} \, \neq \, \texttt{ConductPB} \, \wedge \,$ 

 $MoveToORP \neq CompletePB \land MoveToORP \neq unAuthenticated \land$ 

 $slState = PLAN_PB \mid MOVE_TO_ORP \mid CONDUCT_ORP \mid MOVE_TO_PB$ 

```
MoveToORP \neq unAuthorized \land ConductORP \neq MoveToPB \land
          \texttt{ConductORP} \, \neq \, \texttt{ConductPB} \, \wedge \, \texttt{ConductORP} \, \neq \, \texttt{CompletePB} \, \wedge \,
          {\tt ConductORP} \neq {\tt unAuthenticated} \ \land \ {\tt ConductORP} \neq {\tt unAuthorized} \ \land
          \texttt{MoveToPB} \neq \texttt{ConductPB} \ \land \ \texttt{MoveToPB} \neq \texttt{CompletePB} \ \land \\
          	exttt{MoveToPB} 
eq 	exttt{unAuthenticated} \land 	exttt{MoveToPB} 
eq 	exttt{unAuthorized} \land
          {\tt ConductPB} \, \neq \, {\tt CompletePB} \, \wedge \, {\tt ConductPB} \, \neq \, {\tt unAuthenticated} \, \wedge \,
          {\tt ConductPB} \, \neq \, {\tt unAuthorized} \, \, \land \, \, {\tt CompletePB} \, \neq \, {\tt unAuthenticated} \, \, \land \, \,
          {\tt CompletePB} \neq {\tt unAuthorized} \ \land \ {\tt unAuthenticated} \ \neq \ {\tt unAuthorized}
[slState_distinct_clauses]
   \vdash PLAN_PB \neq MOVE_TO_ORP \land PLAN_PB \neq CONDUCT_ORP \land
          PLAN_PB \neq MOVE_TO_PB \wedge PLAN_PB \neq CONDUCT_PB \wedge
          {\tt PLAN\_PB} \ \neq \ {\tt COMPLETE\_PB} \ \land \ {\tt MOVE\_TO\_ORP} \ \neq \ {\tt CONDUCT\_ORP} \ \land \\
          \texttt{MOVE\_TO\_ORP} \ \neq \ \texttt{MOVE\_TO\_PB} \ \land \ \texttt{MOVE\_TO\_ORP} \ \neq \ \texttt{CONDUCT\_PB} \ \land \\
          MOVE_TO_ORP ≠ COMPLETE_PB ∧ CONDUCT_ORP ≠ MOVE_TO_PB ∧
          {\tt CONDUCT\_ORP} \ \neq \ {\tt CONDUCT\_PB} \ \land \ {\tt CONDUCT\_ORP} \ \neq \ {\tt COMPLETE\_PB} \ \land \\
          	exttt{MOVE\_TO\_PB} 
eq 	exttt{CONDUCT\_PB} 
ightharpoonup MOVE\_TO\_PB 
eq 	exttt{COMPLETE\_PB} 
eq 	exttt{COMPLET
          CONDUCT_PB \neq COMPLETE_PB
[stateRole_distinct_clauses]
   \vdash PlatoonLeader \neq Omni
             PBIntegratedDef Theory
Built: 13 May 2018
Parent Theories: PBTypeIntegrated, aclfoundation
                Definitions
7.1
[secAuthorization_def]
   \vdash \forall xs. secAuthorization xs = secHelper (getOmniCommand xs)
[secHelper_def]
   \vdash \forall cmd.
                secHelper \ cmd =
                [Name Omni controls prop (SOME (SLc (OMNI cmd)))]
7.2
                Theorems
[getOmniCommand_def]
   ├ (get0mniCommand [] = invalid0mniCommand) ∧
          (\forall xs \ cmd.
                   get0mniCommand
                         (Name Omni controls prop (SOME (SLc (OMNI cmd)))::xs) =
                   cmd) \wedge
           (\forall xs. \text{ getOmniCommand } (TT::xs) = \text{getOmniCommand } xs) \land
```

```
(\forall xs. \text{ getOmniCommand } (FF::xs) = \text{getOmniCommand } xs) \land
(\forall xs \ v_2. \ \text{getOmniCommand (prop } v_2::xs) = \text{getOmniCommand } xs) \land
(\forall xs \ v_3. \ \text{getOmniCommand (notf} \ v_3::xs) = \text{getOmniCommand} \ xs) \ \land
(\forall xs \ v_5 \ v_4.
   getOmniCommand (v_4 andf v_5::x_5) = getOmniCommand x_5) \land
(\forall xs \ v_7 \ v_6.
   getOmniCommand (v_6 orf v_7::x_8) = getOmniCommand x_8) \land
(\forall xs \ v_9 \ v_8.
   getOmniCommand (v_8 impf v_9::x_s) = getOmniCommand x_s) \land
(\forall xs \ v_{11} \ v_{10}.
   getOmniCommand (v_{10} eqf v_{11}::xs) = getOmniCommand xs) \land
(\forall xs \ v_{13} \ v_{12}.
   getOmniCommand (v_{12} says v_{13}::xs) = getOmniCommand xs) \land
(\forall xs \ v_{15} \ v_{14}.
   getOmniCommand (v_{14} speaks_for v_{15}::xs) =
   getOmniCommand xs) \land
(\forall xs \ v_{16}.
   getOmniCommand (v_{16} controls TT::xs) =
   getOmniCommand xs) \land
(\forall xs \ v_{16}.
   getOmniCommand (v_{16} controls FF::xs) =
   getOmniCommand xs) \land
(\forall xs \ v134.
   getOmniCommand (Name v134 controls prop NONE::xs) =
   get0mniCommand xs) \land
(\forall xs \ v144.
   get0mniCommand
      (Name PlatoonLeader controls prop (SOME v144)::xs) =
   getOmniCommand xs) \land
(\forall xs \ v146.
   get0mniCommand
      (Name Omni controls prop (SOME (ESCc v146))::xs) =
   getOmniCommand xs) \land
(\forall xs \ v150.
   get0mniCommand
      (Name Omni controls prop (SOME (SLc (PL v150)))::xs) =
   getOmniCommand xs) \land
(\forall xs \ v_{68} \ v136 \ v135.
   getOmniCommand (v135 meet v136 controls prop v_{68}::xs) =
   getOmniCommand xs) \wedge
(\forall xs \ v_{68} \ v_{138} \ v_{137}.
   getOmniCommand (v137 quoting v138 controls prop v_{68}::xs) =
   getOmniCommand xs) \land
(\forall xs \ v_{69} \ v_{16}.
   getOmniCommand (v_{16} controls notf v_{69}::xs) =
   getOmniCommand xs) \land
(\forall xs \ v_{71} \ v_{70} \ v_{16}.
   getOmniCommand (v_{16} controls (v_{70} andf v_{71})::xs) =
   getOmniCommand xs) \land
```

```
(\forall xs \ v_{73} \ v_{72} \ v_{16}.
    getOmniCommand (v_{16} controls (v_{72} orf v_{73})::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{75} \ v_{74} \ v_{16}.
    getOmniCommand (v_{16} controls (v_{74} impf v_{75})::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{77} \ v_{76} \ v_{16}.
    getOmniCommand (v_{16} controls (v_{76} eqf v_{77})::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{79} \ v_{78} \ v_{16}.
    getOmniCommand (v_{16} controls v_{78} says v_{79}::xs) =
    getOmniCommand xs) \wedge
(\forall xs \ v_{81} \ v_{80} \ v_{16}.
    getOmniCommand (v_{16} controls v_{80} speaks_for v_{81}::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{83} \ v_{82} \ v_{16}.
    getOmniCommand (v_{16} controls v_{82} controls v_{83}::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{86} \ v_{85} \ v_{84} \ v_{16}.
    getOmniCommand (v_{16} controls reps v_{84} v_{85} v_{86}::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{88} \ v_{87} \ v_{16}.
    getOmniCommand (v_{16} controls v_{87} domi v_{88}::xs) =
    getOmniCommand xs) \wedge
(\forall xs \ v_{90} \ v_{89} \ v_{16}.
    getOmniCommand (v_{16} controls v_{89} eqi v_{90}::xs) =
    getOmniCommand xs) \wedge
(\forall xs \ v_{92} \ v_{91} \ v_{16}.
    getOmniCommand (v_{16} controls v_{91} doms v_{92}::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{94} \ v_{93} \ v_{16}.
    getOmniCommand (v_{16} controls v_{93} eqs v_{94}::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{96} \ v_{95} \ v_{16}.
    getOmniCommand (v_{16} controls v_{95} eqn v_{96}::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{98} \ v_{97} \ v_{16}.
    getOmniCommand (v_{16} controls v_{97} lte v_{98}::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{99} \ v_{16} \ v_{100}).
    getOmniCommand (v_{16} controls v_{99} lt v100::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{20} \ v_{19} \ v_{18}.
    getOmniCommand (reps v_{18} v_{19} v_{20}::xs) =
    getOmniCommand xs) \land
(\forall xs \ v_{22} \ v_{21}.
    getOmniCommand (v_{21} domi v_{22}::xs) = getOmniCommand xs) \land
(\forall xs \ v_{24} \ v_{23}.
    getOmniCommand (v_{23} eqi v_{24}::xs) = getOmniCommand xs) \land
```

```
(\forall xs \ v_{26} \ v_{25}.
           getOmniCommand (v_{25} doms v_{26}::xs) = getOmniCommand xs) \land
      (\forall xs \ v_{28} \ v_{27}.
           \verb"getOmniCommand" (v_{27} \verb" eqs" v_{28}{::}xs) = \verb"getOmniCommand" xs) \ \land
      (\forall xs \ v_{30} \ v_{29}.
           getOmniCommand (v_{29} eqn v_{30}::xs) = getOmniCommand xs) \land
      (\forall xs \ v_{32} \ v_{31}.
           getOmniCommand (v_{31} lte v_{32}::xs) = getOmniCommand xs) \land
     \forall xs \ v_{34} \ v_{33}.
         getOmniCommand (v_{33} lt v_{34}::xs) = getOmniCommand xs
[getOmniCommand_ind]
  \vdash \forall P.
         P [] \land
         (\forall cmd xs.
                  (Name Omni controls prop (SOME (SLc (OMNI cmd)))::
                          xs)) \land (\forall xs. P xs \Rightarrow P (TT::xs)) \land
         (\forall xs. P xs \Rightarrow P (FF::xs)) \land
         (\forall v_2 \ xs. \ P \ xs \Rightarrow P \ (prop \ v_2::xs)) \ \land
         (\forall v_3 \ xs. \ P \ xs \Rightarrow P \ (\texttt{notf} \ v_3::xs)) \ \land
         (\forall v_4 \ v_5 \ xs. \ P \ xs \Rightarrow P \ (v_4 \ \text{andf} \ v_5::xs)) \ \land
         (\forall v_6 \ v_7 \ xs. \ P \ xs \Rightarrow P \ (v_6 \ \text{orf} \ v_7::xs)) \ \land
         (\forall v_8 \ v_9 \ xs. \ P \ xs \Rightarrow P \ (v_8 \ \text{impf} \ v_9::xs)) \land
         (\forall v_{10} \ v_{11} \ xs. \ P \ xs \Rightarrow P \ (v_{10} \ \mathsf{eqf} \ v_{11} :: xs)) \ \land
         (\forall v_{12} \ v_{13} \ xs. \ P \ xs \Rightarrow P \ (v_{12} \ \text{says} \ v_{13}::xs)) \ \land
         (\forall v_{14} \ v_{15} \ xs. \ P \ xs \Rightarrow P \ (v_{14} \ \text{speaks\_for} \ v_{15}::xs)) \land
         (\forall v_{16} \ xs. \ P \ xs \Rightarrow P \ (v_{16} \ \text{controls TT::} xs)) \ \land
         (\forall v_{16} \ xs. \ P \ xs \Rightarrow P \ (v_{16} \ \text{controls FF}::xs)) \land
         (\forall v134 \ xs. \ P \ xs \Rightarrow P \ (Name \ v134 \ controls \ prop \ NONE::xs)) \land
         (\forall v144 xs.
              P xs \Rightarrow
              P (Name PlatoonLeader controls prop (SOME v144)::xs)) \land
         (\forall v146 \ xs.
              P xs \Rightarrow
              P (Name Omni controls prop (SOME (ESCc v146))::xs)) \land
         (\forall v150 xs.
              P xs \Rightarrow
              P
                  (Name Omni controls prop (SOME (SLc (PL v150)))::
                          xs)) \wedge
         (\forall v135 \ v136 \ v_{68} \ xs.)
              P xs \Rightarrow P (v135 \text{ meet } v136 \text{ controls prop } v_{68}::xs)) \land
         (\forall v137 \ v138 \ v_{68} \ xs.
              P xs \Rightarrow P (v137 \text{ quoting } v138 \text{ controls prop } v_{68}::xs)) \land
         (\forall v_{16} \ v_{69} \ xs. \ P \ xs \Rightarrow P \ (v_{16} \ \text{controls notf} \ v_{69}::xs)) \land
         (\forall v_{16} \ v_{70} \ v_{71} \ xs.
              P xs \Rightarrow P (v_{16} \text{ controls } (v_{70} \text{ andf } v_{71})::xs)) \land
         (\forall v_{16} \ v_{72} \ v_{73} \ xs.
```

```
P xs \Rightarrow P (v_{16} \text{ controls } (v_{72} \text{ orf } v_{73})::xs)) \land
          (\forall v_{16} \ v_{74} \ v_{75} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } (v_{74} \text{ impf } v_{75})::xs)) \land
          (\forall v_{16} \ v_{76} \ v_{77} \ xs.)
                P \ xs \Rightarrow P \ (v_{16} \ {\tt controls} \ (v_{76} \ {\tt eqf} \ v_{77})::xs)) \ \land
          (\forall v_{16} \ v_{78} \ v_{79} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } v_{78} \text{ says } v_{79} :: xs)) \land
          (\forall v_{16} \ v_{80} \ v_{81} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } v_{80} \text{ speaks\_for } v_{81} :: xs)) \land
          (\forall v_{16} \ v_{82} \ v_{83} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } v_{82} \text{ controls } v_{83} :: xs)) \land
          (\forall v_{16} \ v_{84} \ v_{85} \ v_{86} \ xs.
                P \ xs \Rightarrow P \ (v_{16} \ {\tt controls} \ {\tt reps} \ v_{84} \ v_{85} \ v_{86}{\tt ::} xs)) \ \land
          (\forall v_{16} \ v_{87} \ v_{88} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } v_{87} \text{ domi } v_{88}::xs)) \land
          (\forall v_{16} \ v_{89} \ v_{90} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } v_{89} \text{ eqi } v_{90} :: xs)) \land
          (\forall v_{16} \ v_{91} \ v_{92} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } v_{91} \text{ doms } v_{92} :: xs)) \land
          (\forall v_{16} \ v_{93} \ v_{94} \ xs.
                P \ xs \Rightarrow P \ (v_{16} \ \text{controls} \ v_{93} \ \text{eqs} \ v_{94}\!::\!xs)) \ \land
          (\forall v_{16} \ v_{95} \ v_{96} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } v_{95} \text{ eqn } v_{96}::xs)) \land
          (\forall v_{16} \ v_{97} \ v_{98} \ xs.
                P xs \Rightarrow P (v_{16} \text{ controls } v_{97} \text{ lte } v_{98}::xs)) \land
          (\forall v_{16} \ v_{99} \ v_{100} \ xs.
                P \ xs \Rightarrow P \ (v_{16} \ {\tt controls} \ v_{99} \ {\tt lt} \ v100::xs)) \ \land
          (\forall v_{18} \ v_{19} \ v_{20} \ xs. \ P \ xs \Rightarrow P \ (\texttt{reps} \ v_{18} \ v_{19} \ v_{20} :: xs)) \ \land
          (\forall v_{21} \ v_{22} \ xs. \ P \ xs \Rightarrow P \ (v_{21} \ \text{domi} \ v_{22}\!::\!xs)) \ \land
          (\forall v_{23} \ v_{24} \ xs. \ P \ xs \Rightarrow P \ (v_{23} \ \text{eqi} \ v_{24}::xs)) \land
          (\forall v_{25} \ v_{26} \ xs. \ P \ xs \Rightarrow P \ (v_{25} \ \text{doms} \ v_{26}::xs)) \land
          (\forall v_{27} \ v_{28} \ xs. \ P \ xs \Rightarrow P \ (v_{27} \ \text{eqs} \ v_{28}::xs)) \land
          (\forall v_{29} \ v_{30} \ xs. \ P \ xs \Rightarrow P \ (v_{29} \ \text{eqn} \ v_{30}{::}xs)) \ \land
          (\forall v_{31} \ v_{32} \ xs. \ P \ xs \Rightarrow P \ (v_{31} \ \text{lte} \ v_{32} :: xs)) \ \land
          (\forall v_{33} \ v_{34} \ xs. \ P \ xs \Rightarrow P \ (v_{33} \ \text{lt} \ v_{34}\!:\!:\!xs)) \Rightarrow
          \forall v. P v
[secContext_def]
  \vdash (secContext PLAN_PB (x::xs) =
        [prop (SOME (SLc (OMNI ssmPlanPBComplete))) impf
          Name PlatoonLeader controls
          prop (SOME (SLc (PL crossLD)))]) \cap 
       (secContext MOVE_TO_ORP (x::xs) =
        [prop (SOME (SLc (OMNI ssmMoveToORPComplete))) impf
          Name PlatoonLeader controls
          prop (SOME (SLc (PL conductORP)))]) \cap \big|
       (secContext CONDUCT_ORP (x::xs) =
        [prop (SOME (SLc (OMNI ssmConductORPComplete))) impf
         Name PlatoonLeader controls
```

```
prop (SOME (SLc (PL moveToPB)))]) \capsum
    (secContext MOVE_TO_PB (x::xs) =
     [prop (SOME (SLc (OMNI ssmMoveToPBComplete))) impf
      Name PlatoonLeader controls
      prop (SOME (SLc (PL conductPB)))]) \capses
    (secContext CONDUCT_PB (x::xs) =
     [prop (SOME (SLc (OMNI ssmConductPBComplete))) impf
      Name PlatoonLeader controls
      prop (SOME (SLc (PL completePB)))])
[secContext_ind]
 \vdash \forall P.
       (\forall x \ xs. \ P \ \mathtt{PLAN\_PB} \ (x::xs)) \ \land
       (\forall x \ xs. \ P \ MOVE\_TO\_ORP \ (x::xs)) \land
       (\forall x \ xs. \ P \ \texttt{CONDUCT\_ORP} \ (x::xs)) \ \land
       (\forall x \ xs. \ P \ \texttt{MOVE\_TO\_PB} \ (x::xs)) \ \land
       (\forall x \ xs. \ P \ \texttt{CONDUCT\_PB} \ (x \colon : xs)) \ \land \ (\forall \, v_4 \ . \ P \ v_4 \ []) \ \land
       (\forall v_5 \ v_6. \ P \ COMPLETE\_PB \ (v_5::v_6)) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
     ssmConductORP Theory
Built: 13 May 2018
Parent Theories: ConductORPType, ssm11, OMNIType
8.1
       Definitions
[secContextConductORP_def]
 \vdash \forall plcmd psgcmd incomplete.
      {\tt secContextConductORP}\ plcmd\ psgcmd\ incomplete =
       [Name PlatoonLeader controls prop (SOME (SLc (PL plcmd)));
       Name PlatoonSergeant controls
       prop (SOME (SLc (PSG psgcmd)));
       Name PlatoonLeader says
       prop (SOME (SLc (PSG psgcmd))) impf prop NONE;
       Name PlatoonSergeant says
       prop (SOME (SLc (PL plcmd))) impf prop NONE]
[ssmConductORPStateInterp_def]
 \vdash \forall slState. ssmConductORPStateInterp slState = TT
8.2
       Theorems
[authTestConductORP_cmd_reject_lemma]
```

 $\vdash \forall cmd$ .  $\neg authTestConductORP$  (prop (SOME cmd))

```
[authTestConductORP_def]
 \vdash (authTestConductORP (Name PlatoonLeader says prop cmd) \iff
    (authTestConductORP (Name PlatoonSergeant says prop cmd) \iff
     T) \land (authTestConductORP TT \iff F) \land
    (authTestConductORP FF \iff F) \land
    (authTestConductORP (prop v) \iff F) \land
    (authTestConductORP (notf v_1) \iff F) \land
    (authTestConductORP (v_2 andf v_3) \iff F) \wedge
    (authTestConductORP (v_4 orf v_5) \iff F) \wedge
    (authTestConductORP (v_6 impf v_7) \iff F) \land
    (authTestConductORP (v_8 eqf v_9) \iff F) \land
    (authTestConductORP (v_{10} says TT) \iff F) \wedge
    (authTestConductORP (v_{10} says FF) \iff F) \wedge
    (authTestConductORP (v133 meet v134 says prop v_{66}) \iff F) \land
    (authTestConductORP (v135 quoting v136 says prop v_{66}) \iff F) \land
    (authTestConductORP (v_{10} says notf v_{67}) \iff F) \wedge
    (authTestConductORP (v_{10} says (v_{68} andf v_{69})) \iff F) \wedge
    (authTestConductORP (v_{10} says (v_{70} orf v_{71})) \iff F) \land
    (authTestConductORP (v_{10} says (v_{72} impf v_{73})) \iff F) \wedge
    (authTestConductORP (v_{10} says (v_{74} eqf v_{75})) \iff F) \wedge
    (authTestConductORP (v_{10} says v_{76} says v_{77}) \iff F) \wedge
    (authTestConductORP (v_{10} says v_{78} speaks_for v_{79}) \iff F) \wedge
    (authTestConductORP (v_{10} says v_{80} controls v_{81}) \iff F) \wedge
    (authTestConductORP (v_{10} says reps v_{82} v_{83} v_{84}) \iff F) \land
    (authTestConductORP (v_{10} says v_{85} domi v_{86}) \iff F) \wedge
    (authTestConductORP (v_{10} says v_{87} eqi v_{88}) \iff F) \wedge
    (authTestConductORP (v_{10} says v_{89} doms v_{90}) \iff F) \wedge
    (authTestConductORP (v_{10} says v_{91} eqs v_{92}) \iff F) \land
    (authTestConductORP (v_{10} says v_{93} eqn v_{94}) \iff F) \wedge
    (authTestConductORP (v_{10} says v_{95} lte v_{96}) \iff F) \wedge
    (authTestConductORP (v_{10} says v_{97} lt v_{98}) \iff F) \wedge
    (authTestConductORP (v_{12} speaks_for v_{13}) \iff F) \wedge
    (authTestConductORP (v_{14} controls v_{15}) \iff F) \wedge
    (authTestConductORP (reps v_{16} v_{17} v_{18}) \iff F) \wedge
    (authTestConductORP (v_{19} domi v_{20}) \iff F) \land
    (authTestConductORP (v_{21} eqi v_{22}) \iff F) \wedge
    (authTestConductORP (v_{23} doms v_{24}) \iff F) \wedge
    (authTestConductORP (v_{25} eqs v_{26}) \iff F) \wedge
    (authTestConductORP (v_{27} eqn v_{28}) \iff F) \wedge
    (authTestConductORP (v_{29} lte v_{30}) \iff F) \wedge
    (authTestConductORP (v_{31} lt v_{32}) \iff F)
[authTestConductORP_ind]
 \vdash \forall P.
       (\forall \, cmd \, . \, P \, (\texttt{Name PlatoonLeader says prop} \, \, cmd)) \, \land \,
       (\forall \, cmd \, . \, P \, \, ({\tt Name \, PlatoonSergeant \, says \, prop \, } \, cmd)) \, \wedge \, P \, \, {\tt TT} \, \, \wedge \,
       P FF \land (\forall v. P (prop v)) \land (\forall v_1. P (notf v_1)) \land
       (\forall v_2 \ v_3. \ P \ (v_2 \ \text{andf} \ v_3)) \land (\forall v_4 \ v_5. \ P \ (v_4 \ \text{orf} \ v_5)) \land
```

```
(\forall v_6 \ v_7. \ P \ (v_6 \ \text{impf} \ v_7)) \land (\forall v_8 \ v_9. \ P \ (v_8 \ \text{eqf} \ v_9)) \land
          (\forall v_{10}. P (v_{10} \text{ says TT})) \land (\forall v_{10}. P (v_{10} \text{ says FF})) \land
          (\forall v133 \ v134 \ v_{66}. \ P \ (v133 \ \text{meet} \ v134 \ \text{says prop} \ v_{66})) \ \land
          (\forall v135 \ v136 \ v_{66}. \ P \ (v135 \ {
m quoting} \ v136 \ {
m says} \ {
m prop} \ v_{66})) \ \land
          (\forall v_{10} \ v_{67}. P (v_{10} says notf v_{67})) \land
          (\forall\,v_{10}\,v_{68}\,v_{69}. P (v_{10} says (v_{68} andf v_{69}))) \wedge
          (\forall v_{10} \ v_{70} \ v_{71}. \ P \ (v_{10} \ \text{says} \ (v_{70} \ \text{orf} \ v_{71}))) \land
          (\forall v_{10} \ v_{72} \ v_{73}. \ P \ (v_{10} \ \text{says} \ (v_{72} \ \text{impf} \ v_{73}))) \land
          (\forall v_{10} \ v_{74} \ v_{75}. \ P \ (v_{10} \ \text{says} \ (v_{74} \ \text{eqf} \ v_{75}))) \land
          (\forall v_{10} \ v_{76} \ v_{77}. \ P \ (v_{10} \ \text{says} \ v_{76} \ \text{says} \ v_{77})) \ \land
          (\forall v_{10} \ v_{78} \ v_{79}. \ P \ (v_{10} \ \text{says} \ v_{78} \ \text{speaks\_for} \ v_{79})) \ \land
          (\forall v_{10} \ v_{80} \ v_{81}. \ P \ (v_{10} \ \text{says} \ v_{80} \ \text{controls} \ v_{81})) \ \land
          (\forall v_{10} \ v_{82} \ v_{83} \ v_{84}. \ P \ (v_{10} \ \text{says reps} \ v_{82} \ v_{83} \ v_{84})) \ \land
          (\forall v_{10} \ v_{85} \ v_{86}. P (v_{10} says v_{85} domi v_{86})) \land
          (\forall v_{10} \ v_{87} \ v_{88}. P (v_{10} says v_{87} eqi v_{88})) \land
          (\forall v_{10} \ v_{89} \ v_{90}. \ P \ (v_{10} \ \text{says} \ v_{89} \ \text{doms} \ v_{90})) \ \land
          (\forall v_{10} \ v_{91} \ v_{92}. \ P \ (v_{10} \ \text{says} \ v_{91} \ \text{eqs} \ v_{92})) \ \land
          (\forall v_{10} \ v_{93} \ v_{94}. P (v_{10} says v_{93} eqn v_{94})) \land
          (\forall v_{10} \ v_{95} \ v_{96}. \ P \ (v_{10} \ \text{says} \ v_{95} \ \text{lte} \ v_{96})) \ \land
          (\forall v_{10} \ v_{97} \ v_{98}. \ P \ (v_{10} \ \text{says} \ v_{97} \ \text{lt} \ v_{98})) \ \land
          (\forall v_{12} \ v_{13}. \ P \ (v_{12} \ \text{speaks\_for} \ v_{13})) \ \land
          (\forall v_{14} \ v_{15}. P (v_{14} controls v_{15})) \land
          (\forall v_{16} \ v_{17} \ v_{18}. \ P \ (reps \ v_{16} \ v_{17} \ v_{18})) \ \land
          (\forall v_{19} \ v_{20}. \ P \ (v_{19} \ \text{domi} \ v_{20})) \land
          (\forall v_{21} \ v_{22}. \ P \ (v_{21} \ \text{eqi} \ v_{22})) \ \land
          (\forall v_{23} \ v_{24}. \ P \ (v_{23} \ \text{doms} \ v_{24})) \ \land
          (\forall v_{25} \ v_{26}. \ P \ (v_{25} \ \text{eqs} \ v_{26})) \ \land \ (\forall v_{27} \ v_{28}. \ P \ (v_{27} \ \text{eqn} \ v_{28})) \ \land
          (\forall v_{29} \ v_{30}. \ P \ (v_{29} \ \text{lte} \ v_{30})) \land (\forall v_{31} \ v_{32}. \ P \ (v_{31} \ \text{lt} \ v_{32})) \Rightarrow
         \forall v. P v
[conductORPNS_def]
  ├ (conductORPNS CONDUCT_ORP (exec (PL secure)) = SECURE) ∧
      (conductORPNS CONDUCT_ORP (exec (PL plIncomplete)) =
        CONDUCT_ORP) \
      (conductORPNS SECURE (exec (PSG actionsIn)) = ACTIONS_IN) \land
      (conductORPNS SECURE (exec (PSG psgIncomplete)) = SECURE) \land
       (conductORPNS ACTIONS_IN (exec (PL withdraw)) = WITHDRAW) \( \)
      (conductORPNS ACTIONS_IN (exec (PL plIncomplete)) =
       ACTIONS_IN) \
      (conductORPNS WITHDRAW (exec (PL complete)) = COMPLETE) \(\lambda\)
      (conductORPNS WITHDRAW (exec (PL plIncomplete)) = WITHDRAW) \land
      (conductORPNS s (trap (PL cmd')) = s) \land
      (conductORPNS s (trap (PSG cmd)) = s) \land
      (conductORPNS s (discard (PL cmd')) = s) \land
      (conductORPNS \ s \ (discard \ (PSG \ cmd)) = s)
[conductORPNS_ind]
  \vdash \forall P.
         P CONDUCT_ORP (exec (PL secure)) \wedge
```

```
P CONDUCT_ORP (exec (PL plIncomplete)) \wedge
      P SECURE (exec (PSG actionsIn)) \wedge
      P SECURE (exec (PSG psgIncomplete)) \wedge
      P ACTIONS_IN (exec (PL withdraw)) \wedge
      P ACTIONS_IN (exec (PL plIncomplete)) \land
      P WITHDRAW (exec (PL complete)) \wedge
      P WITHDRAW (exec (PL plIncomplete)) \wedge
      (\forall s \ cmd. \ P \ s \ (trap \ (PL \ cmd))) \land
      (\forall s \ cmd. \ P \ s \ (trap \ (PSG \ cmd))) \land
      (\forall s \ cmd. \ P \ s \ (discard \ (PL \ cmd))) \land
      (\forall s \ cmd. \ P \ s \ (\texttt{discard} \ (\texttt{PSG} \ cmd))) \ \land
      P CONDUCT_ORP (exec (PL withdraw)) \wedge
      P CONDUCT_ORP (exec (PL complete)) \wedge
      (\forall v_{11}. P CONDUCT_ORP (exec (PSG v_{11}))) \land
      (\forall v_{13}. P SECURE (exec (PL v_{13}))) \land
      P ACTIONS_IN (exec (PL secure)) \wedge
      P ACTIONS_IN (exec (PL complete)) \wedge
      (\forall v_{17}. \ P \ ACTIONS_IN \ (exec \ (PSG \ v_{17}))) \land
      P WITHDRAW (exec (PL secure)) \wedge
      P WITHDRAW (exec (PL withdraw)) \wedge
      (\forall v_{20}. \ P \ \text{WITHDRAW} \ (\text{exec (PSG} \ v_{20}))) \ \land
      (\forall v_{21}. \ P \ \texttt{COMPLETE} \ (\texttt{exec} \ v_{21})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
[conductORPOut_def]
 \vdash (conductORPOut CONDUCT_ORP (exec (PL secure)) = Secure) \land
    (conductORPOut CONDUCT_ORP (exec (PL plincomplete)) =
     ConductORP) ∧
    (conductORPOut SECURE (exec (PSG actionsIn)) = ActionsIn) \( \)
    (conductORPOut SECURE (exec (PSG psgIncomplete)) = Secure) \( \lambda \)
    (conductORPOut ACTIONS_IN (exec (PL withdraw)) = Withdraw) \( \lambda \)
    (conductORPOut ACTIONS_IN (exec (PL plIncomplete)) =
     ActionsIn) \( \)
    (conductORPOut WITHDRAW (exec (PL complete)) = Complete) \( \)
    (conductORPOut WITHDRAW (exec (PL plIncomplete)) =
    (conductORPOut s (trap (PL cmd')) = unAuthorized) \land
    (conductORPOut s (trap (PSG cmd)) = unAuthorized) \land
    (conductORPOut s (discard (PL cmd')) = unAuthenticated) \land
    (conductORPOut s (discard (PSG cmd)) = unAuthenticated)
[conductORPOut_ind]
 \vdash \forall P.
      P CONDUCT_ORP (exec (PL secure)) \wedge
      P CONDUCT_ORP (exec (PL plIncomplete)) \wedge
      P SECURE (exec (PSG actionsIn)) \wedge
      P SECURE (exec (PSG psgIncomplete)) \wedge
      P ACTIONS_IN (exec (PL withdraw)) \wedge
      P ACTIONS_IN (exec (PL plIncomplete)) \wedge
```

```
P WITHDRAW (exec (PL complete)) \wedge
      P WITHDRAW (exec (PL plIncomplete)) \wedge
      (\forall s \ cmd. \ P \ s \ (trap \ (PL \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (trap \ (PSG \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (\texttt{discard} \ (\texttt{PL} \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (\texttt{discard} \ (\texttt{PSG} \ cmd))) \ \land
      P CONDUCT_ORP (exec (PL withdraw)) \wedge
      P CONDUCT_ORP (exec (PL complete)) \wedge
       (\forall v_{11}. P CONDUCT_ORP (exec (PSG <math>v_{11}))) \land
       (\forall v_{13}. P SECURE (exec (PL <math>v_{13}))) \land
      P ACTIONS_IN (exec (PL secure)) \wedge
      P ACTIONS_IN (exec (PL complete)) \wedge
       (\forall v_{17}. P ACTIONS_{IN} (exec (PSG <math>v_{17}))) \land
       P WITHDRAW (exec (PL secure)) \wedge
      P WITHDRAW (exec (PL withdraw)) \wedge
      (\forall v_{20}. P WITHDRAW (exec (PSG v_{20}))) \wedge
       (\forall v_{21}. \ P \ \texttt{COMPLETE} \ (\texttt{exec} \ v_{21})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
[PlatoonLeader_exec_plCommand_justified_thm]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
      TR (M, Oi, Os) (exec (SLc (PL plCommand)))
         (CFG authTestConductORP ssmConductORPStateInterp
             (secContextConductORP plCommand psqCommand incomplete)
             (Name PlatoonLeader says
              prop (SOME (SLc (PL plCommand)))::ins) s outs)
         ({\tt CFG}\ auth{\tt TestConductORP}\ {\tt ssmConductORPStateInterp}
             (secContextConductORP \ plCommand \ psgCommand \ incomplete)
             ins (NS s (exec (SLc (PL plCommand))))
             (Out \ s \ (exec \ (SLc \ (PL \ plCommand)))::outs)) \iff
      authTestConductORP
         (Name PlatoonLeader says
          prop (SOME (SLc (PL plCommand)))) \cap

      CFGInterpret (M, Oi, Os)
         (CFG authTestConductORP ssmConductORPStateInterp
             (secContextConductORP plCommand psgCommand incomplete)
             (Name PlatoonLeader says
              prop (SOME (SLc (PL plCommand)))::ins) s outs) \land
       (M,Oi,Os) sat prop (SOME (SLc (PL plCommand)))
[PlatoonLeader_plCommand_lemma]
 \vdash CFGInterpret (M, Oi, Os)
       (CFG authTestConductORP ssmConductORPStateInterp
          (secContextConductORP plCommand psgCommand incomplete)
          (Name PlatoonLeader says
           prop (SOME (SLc (PL plCommand)))::ins) s outs) \Rightarrow
    (M, Oi, Os) sat prop (SOME (SLc (PL plCommand)))
```

```
PlatoonSergeant_exec_psgCommand_justified_thm
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os) (exec (SLc (PSG psgCommand)))
        (CFG authTestConductORP ssmConductORPStateInterp
           (secContextConductORP plCommand psgCommand incomplete)
           (Name PlatoonSergeant says
           prop (SOME (SLc (PSG psgCommand)))::ins) s outs)
        (CFG authTestConductORP ssmConductORPStateInterp
           (\verb"secContextConductORP" plCommand psgCommand incomplete)
           ins (NS s (exec (SLc (PSG psgCommand))))
           (Out \ s \ (exec \ (SLc \ (PSG \ psgCommand)))::outs)) \iff
     authTestConductORP
        (Name PlatoonSergeant says
        prop (SOME (SLc (PSG psgCommand)))) \land
     CFGInterpret (M, Oi, Os)
        (CFG authTestConductORP ssmConductORPStateInterp
           (\verb|secContextConductORP|| plCommand|| psgCommand|| incomplete)
           (Name PlatoonSergeant says
            prop (SOME (SLc (PSG psgCommand)))::ins) s outs) \land
     (M, Oi, Os) sat prop (SOME (SLc (PSG psgCommand)))
[PlatoonSergeant_psgCommand_lemma]
 \vdash CFGInterpret (M, Oi, Os)
     (CFG authTestConductORP ssmConductORPStateInterp
         (secContextConductORP plCommand psgCommand incomplete)
         (Name PlatoonSergeant says
         prop (SOME (SLc (PSG psgCommand)))::ins) s outs) \Rightarrow
   (M,Oi,Os) sat prop (SOME (SLc (PSG psgCommand)))
```

## 9 ConductORPType Theory

**Built:** 13 May 2018

Parent Theories: indexedLists, patternMatches

#### 9.1 Datatypes

```
plCommand = secure | withdraw | complete | plIncomplete

psgCommand = actionsIn | psgIncomplete

slCommand =
    PL ConductORPType$plCommand
    | PSG ConductORPType$psgCommand

slOutput = ConductORP | Secure | ActionsIn | Withdraw | Complete | unAuthenticated | unAuthorized

slState = CONDUCT_ORP | SECURE | ACTIONS_IN | WITHDRAW | COMPLETE

stateRole = PlatoonLeader | PlatoonSergeant
```

#### 9.2 Theorems

```
[plCommand_distinct_clauses]
 \vdash secure \neq withdraw \land secure \neq complete \land
     secure \neq plIncomplete \land withdraw \neq complete \land
     withdraw \neq plIncomplete \wedge complete \neq plIncomplete
[psgCommand_distinct_clauses]
 \vdash actionsIn \neq psgIncomplete
[slCommand_distinct_clauses]
 \vdash \forall a' \ a. \ PL \ a \neq PSG \ a'
[slCommand_one_one]
 \vdash (\forall a \ a'. (PL a = PL \ a') \iff (a = a')) \land
    \forall a \ a'. (PSG a = PSG \ a') \iff (a = a')
[slOutput_distinct_clauses]
 \vdash ConductORP \neq Secure \land ConductORP \neq ActionsIn \land
     \texttt{ConductORP} \neq \texttt{Withdraw} \ \land \ \texttt{ConductORP} \neq \texttt{Complete} \ \land
     {\tt ConductORP} \, \neq \, {\tt unAuthenticated} \, \wedge \, {\tt ConductORP} \, \neq \, {\tt unAuthorized} \, \wedge \,
     Secure \neq ActionsIn \wedge Secure \neq Withdraw \wedge Secure \neq Complete \wedge
     Secure \neq unAuthenticated \wedge Secure \neq unAuthorized \wedge
     ActionsIn \neq Withdraw \wedge ActionsIn \neq Complete \wedge
     ActionsIn \neq unAuthenticated \wedge ActionsIn \neq unAuthorized \wedge
     Withdraw \neq Complete \wedge Withdraw \neq unAuthenticated \wedge
     Withdraw \neq unAuthorized \wedge Complete \neq unAuthenticated \wedge
     {\tt Complete} \neq {\tt unAuthorized} \ \land \ {\tt unAuthenticated} \neq {\tt unAuthorized}
[slRole_distinct_clauses]
 ⊢ PlatoonLeader ≠ PlatoonSergeant
[slState_distinct_clauses]
 \vdash CONDUCT_ORP \neq SECURE \land CONDUCT_ORP \neq ACTIONS_IN \land
     {\tt CONDUCT\_ORP} \ \neq \ {\tt WITHDRAW} \ \land \ {\tt CONDUCT\_ORP} \ \neq \ {\tt COMPLETE} \ \land
     \mathtt{SECURE} \neq \mathtt{ACTIONS\_IN} \ \land \ \mathtt{SECURE} \neq \mathtt{WITHDRAW} \ \land \ \mathtt{SECURE} \neq \mathtt{COMPLETE} \ \land
     ACTIONS_IN \neq WITHDRAW \wedge ACTIONS_IN \neq COMPLETE \wedge
     WITHDRAW \neq COMPLETE
```

## 10 ssmConductPB Theory

**Built:** 13 May 2018

Parent Theories: ConductPBType, ssm11, OMNIType

#### 10.1 Definitions

```
[secContextConductPB_def]
 \vdash \forall plcmd psgcmd incomplete.
      {\tt secContextConductPB}\ plcmd\ psgcmd\ incomplete =
      [Name PlatoonLeader controls prop (SOME (SLc (PL plcmd)));
       Name PlatoonSergeant controls
       prop (SOME (SLc (PSG psgcmd)));
       Name PlatoonLeader says
       prop (SOME (SLc (PSG psgcmd))) impf prop NONE;
       Name PlatoonSergeant says
       prop (SOME (SLc (PL plcmd))) impf prop NONE]
[ssmConductPBStateInterp_def]
 \vdash \forall slState. ssmConductPBStateInterp slState = TT
10.2
       Theorems
[authTestConductPB_cmd_reject_lemma]
 \vdash \forall \, cmd. \neg authTestConductPB (prop (SOME cmd))
[authTestConductPB_def]
 \vdash (authTestConductPB (Name PlatoonLeader says prop cmd) \iff T) \land
    (authTestConductPB (Name PlatoonSergeant says prop cmd) \iff
    T) \land (authTestConductPB TT \iff F) \land
    (authTestConductPB FF \iff F) \land
    (authTestConductPB (prop v) \iff F) \land
    (authTestConductPB (notf v_1) \iff F) \land
    (authTestConductPB (v_2 andf v_3) \iff F) \wedge
    (authTestConductPB (v_4 orf v_5) \iff F) \land
    (authTestConductPB (v_6 impf v_7) \iff F) \wedge
    (authTestConductPB (v_8 eqf v_9) \iff F) \land
    (authTestConductPB (v_{10} says TT) \iff F) \wedge
    (authTestConductPB (v_{10} says FF) \iff F) \wedge
    (authTestConductPB (v133 meet v134 says prop v_{66}) \iff F) \land
    (authTestConductPB (v135 quoting v136 says prop v_{66}) \iff F) \land
    (authTestConductPB (v_{10} says notf v_{67}) \iff F) \wedge
    (authTestConductPB (v_{10} says (v_{68} andf v_{69})) \iff F) \land
    (authTestConductPB (v_{10} says (v_{70} orf v_{71})) \iff F) \land
    (authTestConductPB (v_{10} says (v_{72} impf v_{73})) \iff F) \land
    (authTestConductPB (v_{10} says (v_{74} eqf v_{75})) \iff F) \land
    (authTestConductPB (v_{10} says v_{76} says v_{77}) \iff F) \wedge
    (authTestConductPB (v_{10} says v_{78} speaks_for v_{79}) \iff F) \wedge
    (authTestConductPB (v_{10} says v_{80} controls v_{81}) \iff F) \wedge
    (authTestConductPB (v_{10} says reps v_{82} v_{83} v_{84}) \iff F) \wedge
    (authTestConductPB (v_{10} says v_{85} domi v_{86}) \iff F) \wedge
    (authTestConductPB (v_{10} says v_{87} eqi v_{88}) \iff F) \land
    (authTestConductPB (v_{10} says v_{89} doms v_{90}) \iff F) \land
```

```
(authTestConductPB (v_{10} says v_{91} eqs v_{92}) \iff F) \wedge
       (authTestConductPB (v_{10} says v_{93} eqn v_{94}) \iff F) \wedge
       (authTestConductPB (v_{10} says v_{95} lte v_{96}) \iff F) \wedge
       (authTestConductPB (v_{10} says v_{97} lt v_{98}) \iff F) \land
       (authTestConductPB (v_{12} speaks_for v_{13}) \iff F) \land
       (authTestConductPB (v_{14} controls v_{15}) \iff F) \wedge
      (authTestConductPB (reps v_{16} v_{17} v_{18}) \iff F) \wedge
      (authTestConductPB (v_{19} domi v_{20}) \iff F) \wedge
      (authTestConductPB (v_{21} eqi v_{22}) \iff F) \land
       (authTestConductPB (v_{23} doms v_{24}) \iff F) \land
       (authTestConductPB (v_{25} eqs v_{26}) \iff F) \land
      (authTestConductPB (v_{27} eqn v_{28}) \iff F) \land (authTestConductPB (v_{29} lte v_{30}) \iff F) \land
       (authTestConductPB (v_{31} lt v_{32}) \iff F)
[authTestConductPB_ind]
  \vdash \forall P.
          (\forall \, cmd \, . \, P \, \, ({\tt Name \, PlatoonLeader \, says \, prop \, } \, cmd)) \, \, \wedge \, \,
           (\forall \, cmd \, . \, P \, (\texttt{Name PlatoonSergeant says prop} \, cmd)) \, \land \, P \, \texttt{TT} \, \land \,
          P FF \land (\forall v. P (prop v)) \land (\forall v_1. P (notf v_1)) \land
           (\forall v_2 \ v_3. \ P \ (v_2 \ \text{andf} \ v_3)) \ \land \ (\forall v_4 \ v_5. \ P \ (v_4 \ \text{orf} \ v_5)) \ \land
           (\forall v_6 \ v_7. \ P \ (v_6 \ \text{impf} \ v_7)) \ \land \ (\forall v_8 \ v_9. \ P \ (v_8 \ \text{eqf} \ v_9)) \ \land
           (\forall v_{10}. \ P \ (v_{10} \ \text{says TT})) \ \land \ (\forall v_{10}. \ P \ (v_{10} \ \text{says FF})) \ \land
           (\forall\,v133\ v134\ v_{66}. P (v133\ \mathrm{meet}\ v134\ \mathrm{says}\ \mathrm{prop}\ v_{66})) \wedge
           (\forall v135 \ v136 \ v_{66}. \ P \ (v135 \ \text{quoting} \ v136 \ \text{says prop} \ v_{66})) \ \land
           (\forall v_{10} \ v_{67}. \ P \ (v_{10} \ \text{says notf} \ v_{67})) \land
           (\forall v_{10} \ v_{68} \ v_{69}. \ P \ (v_{10} \ \text{says} \ (v_{68} \ \text{andf} \ v_{69}))) \land
           (\forall v_{10} \ v_{70} \ v_{71}. \ P \ (v_{10} \ \text{says} \ (v_{70} \ \text{orf} \ v_{71}))) \land
           (\forall v_{10} \ v_{72} \ v_{73}. \ P \ (v_{10} \ \text{says} \ (v_{72} \ \text{impf} \ v_{73}))) \ \land
           (\forall v_{10} \ v_{74} \ v_{75}. \ P \ (v_{10} \ \text{says} \ (v_{74} \ \text{eqf} \ v_{75}))) \ \land
           (\forall v_{10} \ v_{76} \ v_{77}. \ P \ (v_{10} \ \text{says} \ v_{76} \ \text{says} \ v_{77})) \ \land
           (\forall v_{10} \ v_{78} \ v_{79}. P (v_{10} says v_{78} speaks_for v_{79})) \wedge
           (\forall v_{10} \ v_{80} \ v_{81}. \ P \ (v_{10} \ \text{says} \ v_{80} \ \text{controls} \ v_{81})) \ \land
           (\forall v_{10} \ v_{82} \ v_{83} \ v_{84}. \ P \ (v_{10} \ {\tt says \ reps} \ v_{82} \ v_{83} \ v_{84})) \ \land
           (\forall v_{10} \ v_{85} \ v_{86}. \ P \ (v_{10} \ {\tt says} \ v_{85} \ {\tt domi} \ v_{86})) \ \land
           (\forall v_{10} \ v_{89} \ v_{90}. \ P \ (v_{10} \ \text{says} \ v_{89} \ \text{doms} \ v_{90})) \ \land
           (\forall v_{10} \ v_{91} \ v_{92}. \ P \ (v_{10} \ {\tt says} \ v_{91} \ {\tt eqs} \ v_{92})) \ \land
           (\forall v_{10} \ v_{93} \ v_{94}. \ P \ (v_{10} \ \text{says} \ v_{93} \ \text{eqn} \ v_{94})) \ \land
           (\forall v_{10} \ v_{95} \ v_{96}. \ P \ (v_{10} \ \text{says} \ v_{95} \ \text{lte} \ v_{96})) \ \land
           (\forall v_{10} \ v_{97} \ v_{98}. \ P \ (v_{10} \ \text{says} \ v_{97} \ \text{lt} \ v_{98})) \land
           (\forall v_{12} \ v_{13}. P (v_{12} speaks_for v_{13})) \land
           (\forall v_{14} \ v_{15}. P (v_{14} controls v_{15})) \land
           (\forall v_{16} \ v_{17} \ v_{18}. \ P \ (reps \ v_{16} \ v_{17} \ v_{18})) \ \land
           (\forall v_{19} \ v_{20}. P (v_{19} domi v_{20})) \land
           (\forall v_{21} \ v_{22}. \ P \ (v_{21} \ \text{eqi} \ v_{22})) \ \land
           (\forall v_{23} \ v_{24}. \ P \ (v_{23} \ \text{doms} \ v_{24})) \land
           (\forall v_{25} \ v_{26}. \ P \ (v_{25} \ \text{eqs} \ v_{26})) \land (\forall v_{27} \ v_{28}. \ P \ (v_{27} \ \text{eqn} \ v_{28})) \land
           (\forall v_{29} \ v_{30}. \ P \ (v_{29} \ \text{lte} \ v_{30})) \land (\forall v_{31} \ v_{32}. \ P \ (v_{31} \ \text{lt} \ v_{32})) \Rightarrow
```

```
\forall v. P v
[conductPBNS_def]
 \vdash (conductPBNS CONDUCT_PB (exec (PL securePB)) = SECURE_PB) \land
    (conductPBNS CONDUCT_PB (exec (PL plIncompletePB)) =
     CONDUCT_PB) \
    (conductPBNS SECURE_PB (exec (PSG actionsInPB)) =
     ACTIONS_IN_PB) ∧
    (conductPBNS SECURE_PB (exec (PSG psgIncompletePB)) =
     SECURE_PB) \
    (conductPBNS ACTIONS_IN_PB (exec (PL withdrawPB)) =
     WITHDRAW_PB) \
    (conductPBNS ACTIONS_IN_PB (exec (PL plIncompletePB)) =
     ACTIONS_IN_PB) ∧
    (conductPBNS WITHDRAW_PB (exec (PL completePB)) =
     COMPLETE_PB) ∧
    (conductPBNS WITHDRAW_PB (exec (PL plIncompletePB)) =
     WITHDRAW_PB) \land (conductPBNS s (trap (PL cmd')) = s) \land
    (conductPBNS s (trap (PSG cmd)) = s) \land
    (conductPBNS s (discard (PL cmd')) = s) \land
    (conductPBNS s (discard (PSG cmd)) = s)
[conductPBNS_ind]
 \vdash \forall P.
      P CONDUCT_PB (exec (PL securePB)) \wedge
      P CONDUCT_PB (exec (PL plIncompletePB)) \wedge
      P SECURE_PB (exec (PSG actionsInPB)) \wedge
      P SECURE_PB (exec (PSG psgIncompletePB)) \wedge
      P ACTIONS_IN_PB (exec (PL withdrawPB)) \wedge
      P ACTIONS_IN_PB (exec (PL plIncompletePB)) \wedge
      P WITHDRAW_PB (exec (PL completePB)) \wedge
      P WITHDRAW_PB (exec (PL plIncompletePB)) \wedge
      (\forall s \ cmd. \ P \ s \ (trap \ (PL \ cmd))) \ \land
      (\forall s \ cmd. \ P \ s \ (trap \ (PSG \ cmd))) \land
      (\forall s \ cmd. \ P \ s \ (\texttt{discard} \ (\texttt{PL} \ cmd))) \ \land
      (\forall s \ cmd. \ P \ s \ (\texttt{discard} \ (\texttt{PSG} \ cmd))) \ \land
      P CONDUCT_PB (exec (PL withdrawPB)) \wedge
      P CONDUCT_PB (exec (PL completePB)) \wedge
      (\forall v_{11}. P CONDUCT_{PB} (exec (PSG v_{11}))) \land
      (\forall v_{13}. P SECURE_PB (exec (PL v_{13}))) \land
      P ACTIONS_IN_PB (exec (PL securePB)) \wedge
      P ACTIONS_IN_PB (exec (PL completePB)) \wedge
      (\forall v_{17}.\ P ACTIONS_IN_PB (exec (PSG v_{17}))) \land
      P WITHDRAW_PB (exec (PL securePB)) \wedge
      P WITHDRAW_PB (exec (PL withdrawPB)) \wedge
      (\forall v_{20}. \ P \ \text{WITHDRAW\_PB (exec (PSG} \ v_{20}))) \ \land
      (\forall v_{21}. \ P \ \texttt{COMPLETE\_PB} \ (\texttt{exec} \ v_{21})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
```

```
[conductPBOut_def]
 ⊢ (conductPBOut CONDUCT_PB (exec (PL securePB)) = ConductPB) ∧
    (conductPBOut CONDUCT_PB (exec (PL plIncompletePB)) =
     ConductPB) ∧
    (conductPBOut SECURE_PB (exec (PSG actionsInPB)) =
     SecurePB) ∧
    (conductPBOut SECURE_PB (exec (PSG psgIncompletePB)) =
     SecurePB) ∧
    (conductPBOut ACTIONS_IN_PB (exec (PL withdrawPB)) =
     ActionsInPB) \( \)
    (conductPBOut ACTIONS_IN_PB (exec (PL plIncompletePB)) =
     ActionsInPB) \( \)
    (conductPBOut WITHDRAW_PB (exec (PL completePB)) =
     WithdrawPB) ∧
    (conductPBOut WITHDRAW_PB (exec (PL plIncompletePB)) =
     WithdrawPB) ∧
    (conductPBOut s (trap (PL cmd')) = unAuthorized) \land
    (conductPBOut s (trap (PSG cmd)) = unAuthorized) \land
    (conductPBOut s (discard (PL cmd')) = unAuthenticated) \land
    (conductPBOut s (discard (PSG cmd)) = unAuthenticated)
[conductPBOut_ind]
 \vdash \forall P.
      P CONDUCT_PB (exec (PL securePB)) \wedge
      P CONDUCT_PB (exec (PL plIncompletePB)) \wedge
      P SECURE_PB (exec (PSG actionsInPB)) \wedge
      P SECURE_PB (exec (PSG psgIncompletePB)) \wedge
      P ACTIONS_IN_PB (exec (PL withdrawPB)) \wedge
      P ACTIONS_IN_PB (exec (PL plIncompletePB)) \wedge
      P WITHDRAW_PB (exec (PL completePB)) \wedge
      P WITHDRAW_PB (exec (PL plIncompletePB)) \wedge
       (\forall s \ cmd. \ P \ s \ (trap \ (PL \ cmd))) \land
       (\forall s \ cmd. \ P \ s \ (trap \ (PSG \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (discard \ (PL \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (discard \ (PSG \ cmd))) \ \land
       P CONDUCT_PB (exec (PL withdrawPB)) \wedge
      P CONDUCT_PB (exec (PL completePB)) \wedge
       (\forall v_{11}. \ P \ \texttt{CONDUCT\_PB} \ (\texttt{exec} \ (\texttt{PSG} \ v_{11}))) \ \land
       (\forall v_{13}. \ P \ \texttt{SECURE\_PB} \ (\texttt{exec} \ (\texttt{PL} \ v_{13}))) \ \land
       P ACTIONS_IN_PB (exec (PL securePB)) \wedge
       P ACTIONS_IN_PB (exec (PL completePB)) \land
       (\forall v_{17}. \ P \ ACTIONS_IN_PB \ (exec \ (PSG \ v_{17}))) \land
       P WITHDRAW_PB (exec (PL securePB)) \wedge
      P WITHDRAW_PB (exec (PL withdrawPB)) \wedge
       (\forall v_{20}. \ P \ \texttt{WITHDRAW\_PB} \ (\texttt{exec} \ (\texttt{PSG} \ v_{20}))) \ \land
       (\forall v_{21}. \ P \ \texttt{COMPLETE\_PB} \ (\texttt{exec} \ v_{21})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
```

```
PlatoonLeader_exec_plCommandPB_justified_thm
 \vdash \forall NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os) (exec (SLc (PL plCommand)))
       (CFG authTestConductPB ssmConductPBStateInterp
           (secContextConductPB plCommand psgCommand incomplete)
           (Name PlatoonLeader says
           prop (SOME (SLc (PL plCommand)))::ins) s outs)
       (CFG authTestConductPB ssmConductPBStateInterp
           (secContextConductPB plCommand psgCommand incomplete)
           ins (NS s (exec (SLc (PL plCommand))))
           (Out \ s \ (exec \ (SLc \ (PL \ plCommand)))::outs)) \iff
     authTestConductPB
       (Name PlatoonLeader says
        prop (SOME (SLc (PL plCommand)))) ∧
     CFGInterpret (M, Oi, Os)
       (CFG authTestConductPB ssmConductPBStateInterp
           (secContextConductPB plCommand psqCommand incomplete)
           (Name PlatoonLeader says
           prop (SOME (SLc (PL plCommand)))::ins) s outs) \land
     (M, Oi, Os) sat prop (SOME (SLc (PL plCommand)))
[PlatoonLeader_plCommandPB_lemma]
 \vdash CFGInterpret (M, Oi, Os)
     (CFG authTestConductPB ssmConductPBStateInterp
         (secContextConductPB plCommand psqCommand incomplete)
         (Name PlatoonLeader says
         prop (SOME (SLc (PL plCommand)))::ins) s outs) \Rightarrow
   (M, Oi, Os) sat prop (SOME (SLc (PL plCommand)))
[PlatoonSergeant_exec_psgCommandPB_justified_thm]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os) (exec (SLc (PSG psgCommand)))
       (CFG authTestConductPB ssmConductPBStateInterp
           (secContextConductPB plCommand psgCommand incomplete)
           (Name PlatoonSergeant says
           prop (SOME (SLc (PSG psgCommand)))::ins) s outs)
       (CFG authTestConductPB ssmConductPBStateInterp
           (secContextConductPB plCommand psgCommand incomplete)
           ins (NS s (exec (SLc (PSG psgCommand))))
           (Out \ s \ (exec \ (SLc \ (PSG \ psgCommand)))::outs)) \iff
     authTestConductPB
       (Name PlatoonSergeant says
        prop (SOME (SLc (PSG psgCommand)))) \cap \langle 
     CFGInterpret (M, Oi, Os)
       (CFG authTestConductPB ssmConductPBStateInterp
           (secContextConductPB plCommand psgCommand incomplete)
           (Name PlatoonSergeant says
           prop (SOME (SLc (PSG psgCommand)))::ins) s outs) \land
     (M, Oi, Os) sat prop (SOME (SLc (PSG psgCommand)))
```

```
[PlatoonSergeant_psgCommandPB_lemma]
 \vdash CFGInterpret (M, Oi, Os)
      (CFG authTestConductPB ssmConductPBStateInterp
          (secContextConductPB plCommand psgCommand incomplete)
         (Name PlatoonSergeant says
          prop (SOME (SLc (PSG psgCommand)))::ins) s outs) \Rightarrow
    (M, Oi, Os) sat prop (SOME (SLc (PSG psgCommand)))
11
       ConductPBType Theory
Built: 13 May 2018
Parent Theories: indexedLists, patternMatches
11.1 Datatypes
plCommandPB = securePB | withdrawPB | completePB
              | plIncompletePB
psgCommandPB = actionsInPB | psgIncompletePB
slCommand = PL plCommandPB | PSG psgCommandPB
slOutput = ConductPB | SecurePB | ActionsInPB | WithdrawPB
          | CompletePB | unAuthenticated | unAuthorized
slState = {\tt CONDUCT\_PB} \mid {\tt SECURE\_PB} \mid {\tt ACTIONS\_IN\_PB} \mid {\tt WITHDRAW\_PB}
         | COMPLETE_PB
stateRole = PlatoonLeader | PlatoonSergeant
11.2
       Theorems
[plCommandPB_distinct_clauses]
 \vdash securePB \neq withdrawPB \land securePB \neq completePB \land
    \texttt{securePB} \neq \texttt{plIncompletePB} \ \land \ \texttt{withdrawPB} \neq \texttt{completePB} \ \land \\
    withdrawPB \neq plIncompletePB \wedge completePB \neq plIncompletePB
[psgCommandPB_distinct_clauses]
 \vdash actionsInPB \neq psgIncompletePB
[slCommand_distinct_clauses]
 \vdash \ \forall \ a' \ a. PL a \neq \ \mathtt{PSG} \ a'
[slCommand_one_one]
```

 $\vdash$  ( $\forall a \ a'$ . (PL  $a = PL \ a'$ )  $\iff$  (a = a'))  $\land$   $\forall a \ a'$ . (PSG  $a = PSG \ a'$ )  $\iff$  (a = a')

[slOutput\_distinct\_clauses]

```
\vdash ConductPB \neq SecurePB \land ConductPB \neq ActionsInPB \land
     ConductPB \neq WithdrawPB \land ConductPB \neq CompletePB \land
     {\tt ConductPB} \, \neq \, {\tt unAuthenticated} \, \wedge \, {\tt ConductPB} \, \neq \, {\tt unAuthorized} \, \wedge \,
     \texttt{SecurePB} \neq \texttt{ActionsInPB} \ \land \ \texttt{SecurePB} \neq \texttt{WithdrawPB} \ \land \\
     SecurePB \neq CompletePB \wedge SecurePB \neq unAuthenticated \wedge
     \texttt{SecurePB} \neq \texttt{unAuthorized} \ \land \ \texttt{ActionsInPB} \neq \texttt{WithdrawPB} \ \land
     {\tt ActionsInPB} \neq {\tt CompletePB} \ \land \ {\tt ActionsInPB} \neq {\tt unAuthenticated} \ \land \\
     {\tt ActionsInPB} \neq {\tt unAuthorized} \ \land \ {\tt WithdrawPB} \neq {\tt CompletePB} \ \land \\
     WithdrawPB \neq unAuthenticated \wedge WithdrawPB \neq unAuthorized \wedge
     {\tt CompletePB} \neq {\tt unAuthenticated} \ \land \ {\tt CompletePB} \neq {\tt unAuthorized} \ \land \\
     unAuthenticated \neq unAuthorized
[slRole_distinct_clauses]
 ⊢ PlatoonLeader ≠ PlatoonSergeant
[slState_distinct_clauses]
 \vdash CONDUCT_PB \neq SECURE_PB \land CONDUCT_PB \neq ACTIONS_IN_PB \land
     {\tt CONDUCT\_PB} \ \neq \ {\tt WITHDRAW\_PB} \ \land \ {\tt CONDUCT\_PB} \ \neq \ {\tt COMPLETE\_PB} \ \land
     {\tt SECURE\_PB} \ \neq \ {\tt ACTIONS\_IN\_PB} \ \land \ {\tt SECURE\_PB} \ \neq \ {\tt WITHDRAW\_PB} \ \land \\
     {\tt SECURE\_PB} \ \neq \ {\tt COMPLETE\_PB} \ \land \ {\tt ACTIONS\_IN\_PB} \ \neq \ {\tt WITHDRAW\_PB} \ \land \\
     {\tt ACTIONS\_IN\_PB} \ \neq \ {\tt COMPLETE\_PB} \ \land \ {\tt WITHDRAW\_PB} \ \neq \ {\tt COMPLETE\_PB}
12
         ssmMoveToORP Theory
Built: 13 May 2018
Parent Theories: MoveToORPType, ssm11, OMNIType
12.1
          Definitions
[secContextMoveToORP_def]
 \vdash \forall cmd.
        {\tt secContextMoveToORP}\ cmd =
        [Name PlatoonLeader controls prop (SOME (SLc cmd))]
[ssmMoveToORPStateInterp_def]
 \vdash \ \forall \, state. ssmMoveToORPStateInterp state = TT
12.2
          Theorems
[authTestMoveToORP_cmd_reject_lemma]
 \vdash \ \forall \ cmd. \neg \texttt{authTestMoveToORP} (prop (SOME cmd))
```

```
[authTestMoveToORP_def]
 \vdash (authTestMoveToORP (Name PlatoonLeader says prop cmd) \iff T) \land
     (authTestMoveToORP TT \iff F) \land (authTestMoveToORP FF \iff F) \land
     (authTestMoveToORP (prop v) \iff F) \land
     (authTestMoveToORP (notf v_1) \iff F) \land
     (authTestMoveToORP (v_2 andf v_3) \iff F) \wedge
     (authTestMoveToORP (v_4 orf v_5) \iff F) \land
     (authTestMoveToORP (v_6 impf v_7) \iff F) \land
     (authTestMoveToORP (v_8 eqf v_9) \iff F) \land
    (authTestMoveToORP (v_{10} says TT) \iff F) \wedge
    (authTestMoveToORP (v_{10} says FF) \iff F) \land
    (authTestMoveToORP (v133 meet v134 says prop v_{66}) \iff F) \land
    (authTestMoveToORP (v135 quoting v136 says prop v_{66}) \iff F) \land
     (authTestMoveToORP (v_{10} says notf v_{67}) \iff F) \wedge
     (authTestMoveToORP (v_{10} says (v_{68} andf v_{69})) \iff F) \land
     (authTestMoveToORP (v_{10} says (v_{70} orf v_{71})) \iff F) \land
     (authTestMoveToORP (v_{10} says (v_{72} impf v_{73})) \iff F) \wedge
     (authTestMoveToORP (v_{10} says (v_{74} eqf v_{75})) \iff F) \wedge
     (authTestMoveToORP (v_{10} says v_{76} says v_{77}) \iff F) \land
     (authTestMoveToORP (v_{10} says v_{78} speaks_for v_{79}) \iff F) \wedge
     (authTestMoveToORP (v_{10} says v_{80} controls v_{81}) \iff F) \wedge
     (authTestMoveToORP (v_{10} says reps v_{82} v_{83} v_{84}) \iff F) \wedge
     (authTestMoveToORP (v_{10} says v_{85} domi v_{86}) \iff F) \wedge
     (authTestMoveToORP (v_{10} says v_{87} eqi v_{88}) \iff F) \wedge
     (authTestMoveToORP (v_{10} says v_{89} doms v_{90}) \iff F) \wedge
    (authTestMoveToORP (v_{10} says v_{91} eqs v_{92}) \iff F) \land
     (authTestMoveToORP (v_{10} says v_{93} eqn v_{94}) \iff F) \wedge
     (authTestMoveToORP (v_{10} says v_{95} lte v_{96}) \iff F) \wedge
     (authTestMoveToORP (v_{10} says v_{97} lt v_{98}) \iff F) \wedge
     (authTestMoveToORP (v_{12} speaks_for v_{13}) \iff F) \land
     (authTestMoveToORP (v_{14} controls v_{15}) \iff F) \wedge
     (authTestMoveToORP (reps v_{16} v_{17} v_{18}) \iff F) \wedge
    (authTestMoveToORP (v_{19} domi v_{20}) \iff F) \wedge
     (authTestMoveToORP (v_{21} eqi v_{22}) \iff F) \wedge
     (authTestMoveToORP (v_{23} doms v_{24}) \iff F) \wedge
     (authTestMoveToORP (v_{25} eqs v_{26}) \iff F) \land
     (authTestMoveToORP (v_{27} eqn v_{28}) \iff F) \wedge
     (authTestMoveToORP (v_{29} lte v_{30}) \iff F) \wedge
     (authTestMoveToORP (v_{31} lt v_{32}) \iff F)
[authTestMoveToORP_ind]
 \vdash \forall P.
       (\forall \, cmd \, . \, P \, \, (\text{Name PlatoonLeader says prop} \, \, cmd)) \, \wedge \, P \, \, \text{TT} \, \wedge \,
       P FF \wedge (\forall v. P (prop v)) \wedge (\forall v_1. P (notf v_1)) \wedge
       (\forall v_2 \ v_3. \ P \ (v_2 \ \text{andf} \ v_3)) \land (\forall v_4 \ v_5. \ P \ (v_4 \ \text{orf} \ v_5)) \land
       (\forall v_6 \ v_7. \ P \ (v_6 \ \text{impf} \ v_7)) \ \land \ (\forall v_8 \ v_9. \ P \ (v_8 \ \text{eqf} \ v_9)) \ \land
       (\forall v_{10}. \ P \ (v_{10} \ \text{says TT})) \land (\forall v_{10}. \ P \ (v_{10} \ \text{says FF})) \land
       (\forall v133 \ v134 \ v_{66}. \ P \ (v133 \ \text{meet} \ v134 \ \text{says prop} \ v_{66})) \ \land
       (\forall v135 \ v136 \ v_{66}. \ P \ (v135 \ \text{quoting} \ v136 \ \text{says prop} \ v_{66})) \ \land
```

```
(\forall v_{10} \ v_{67}. \ P \ (v_{10} \ \text{says notf} \ v_{67})) \land
         (\forall v_{10} \ v_{68} \ v_{69}. \ P \ (v_{10} \ \text{says} \ (v_{68} \ \text{andf} \ v_{69}))) \ \land
         (\forall v_{10} \ v_{70} \ v_{71}. \ P \ (v_{10} \ \text{says} \ (v_{70} \ \text{orf} \ v_{71}))) \ \land
         (\forall v_{10} \ v_{72} \ v_{73}. \ P \ (v_{10} \ {\tt says} \ (v_{72} \ {\tt impf} \ v_{73}))) \ \land
         (\forall v_{10} \ v_{74} \ v_{75}. \ P \ (v_{10} \ \text{says} \ (v_{74} \ \text{eqf} \ v_{75}))) \ \land
         (\forall v_{10} \ v_{76} \ v_{77}. P (v_{10} says v_{76} says v_{77})) \land
         (\forall v_{10} \ v_{78} \ v_{79}. \ P \ (v_{10} \ \text{says} \ v_{78} \ \text{speaks\_for} \ v_{79})) \ \land
         (\forall v_{10} \ v_{80} \ v_{81}. \ P \ (v_{10} \ \text{says} \ v_{80} \ \text{controls} \ v_{81})) \ \land
         (\forall v_{10} \ v_{82} \ v_{83} \ v_{84}. \ P \ (v_{10} \ \text{says reps} \ v_{82} \ v_{83} \ v_{84})) \ \land
         (\forall v_{10} \ v_{85} \ v_{86}. \ P \ (v_{10} \ \text{says} \ v_{85} \ \text{domi} \ v_{86})) \ \land
         (\forall v_{10} \ v_{87} \ v_{88}. \ P \ (v_{10} \ \text{says} \ v_{87} \ \text{eqi} \ v_{88})) \ \land
         (\forall v_{10} \ v_{89} \ v_{90}. \ P \ (v_{10} \ \text{says} \ v_{89} \ \text{doms} \ v_{90})) \ \land
         (\forall v_{10} \ v_{91} \ v_{92}. \ P \ (v_{10} \ {\tt says} \ v_{91} \ {\tt eqs} \ v_{92})) \ \land
         (\forall v_{10} \ v_{93} \ v_{94}. P (v_{10} says v_{93} eqn v_{94})) \land
         (\forall\,v_{10}\ v_{95}\ v_{96}. P (v_{10} says v_{95} lte v_{96})) \wedge
         (\forall v_{10} \ v_{97} \ v_{98}. \ P \ (v_{10} \ \text{says} \ v_{97} \ \text{lt} \ v_{98})) \land
         (\forall v_{12} \ v_{13}. \ P \ (v_{12} \ \text{speaks\_for} \ v_{13})) \ \land
         (\forall \, v_{14} \ v_{15}. P (v_{14} controls v_{15})) \wedge
         (\forall v_{16} \ v_{17} \ v_{18}. \ P \ (\text{reps} \ v_{16} \ v_{17} \ v_{18})) \ \land
         (\forall v_{19} \ v_{20}. \ P \ (v_{19} \ \text{domi} \ v_{20})) \land
         (\forall v_{21} \ v_{22}. \ P \ (v_{21} \ \text{eqi} \ v_{22})) \ \land
         (\forall v_{23} \ v_{24}. \ P \ (v_{23} \ \text{doms} \ v_{24})) \ \land
         (\forall v_{25} \ v_{26}. \ P \ (v_{25} \ \text{eqs} \ v_{26})) \land (\forall v_{27} \ v_{28}. \ P \ (v_{27} \ \text{eqn} \ v_{28})) \land
         (\forall v_{29} \ v_{30}. \ P \ (v_{29} \ \text{lte} \ v_{30})) \land (\forall v_{31} \ v_{32}. \ P \ (v_{31} \ \text{lt} \ v_{32})) \Rightarrow
         \forall v. P v
[moveToORPNS_def]
  ⊢ (moveToORPNS MOVE_TO_ORP (exec (SLc pltForm)) = PLT_FORM) ∧
      (moveToORPNS MOVE_TO_ORP (exec (SLc incomplete)) =
       MOVE TO ORP) ∧
      (moveToORPNS PLT_FORM (exec (SLc pltMove)) = PLT_MOVE) \land
      (moveToORPNS PLT_FORM (exec (SLc incomplete)) = PLT_FORM) \(\lambda\)
      (moveToORPNS PLT_MOVE (exec (SLc pltSecureHalt)) =
       PLT_SECURE_HALT) \( \tag{ }
      (moveToORPNS PLT_MOVE (exec (SLc incomplete)) = PLT_MOVE) \(\lambda\)
      (moveToORPNS PLT_SECURE_HALT (exec (SLc complete)) =
       COMPLETE) ∧
      (moveToORPNS PLT_SECURE_HALT (exec (SLc incomplete)) =
       PLT_SECURE_HALT) \land (moveToORPNS s (trap (SLc cmd)) = s) \land
      (moveToORPNS s (discard (SLc cmd)) = s)
[moveToORPNS_ind]
  \vdash \forall P.
         P MOVE_TO_ORP (exec (SLc pltForm)) \wedge
         P MOVE_TO_ORP (exec (SLc incomplete)) \land
         P PLT_FORM (exec (SLc pltMove)) \wedge
         P PLT_FORM (exec (SLc incomplete)) \wedge
         P PLT_MOVE (exec (SLc pltSecureHalt)) \wedge
         P PLT_MOVE (exec (SLc incomplete)) \wedge
```

```
P PLT_SECURE_HALT (exec (SLc complete)) \wedge
       P PLT_SECURE_HALT (exec (SLc incomplete)) \wedge
       (\forall s \ cmd. \ P \ s \ (trap \ (SLc \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (\texttt{discard} \ (\texttt{SLc} \ cmd))) \ \land
       (\forall s \ v_6. \ P \ s \ (\texttt{discard} \ (\texttt{ESCc} \ v_6))) \ \land
       (\forall s \ v_9. \ P \ s \ (trap \ (ESCc \ v_9))) \land
       (\forall v_{12}. P MOVE_TO_ORP (exec (ESCc v_{12}))) \land
       P MOVE_TO_ORP (exec (SLc pltMove)) \wedge
      P MOVE_TO_ORP (exec (SLc pltSecureHalt)) \wedge
      P MOVE_TO_ORP (exec (SLc complete)) \wedge
       (\forall v_{15}.\ P\ \mathtt{PLT\_FORM}\ (\mathtt{exec}\ (\mathtt{ESCc}\ v_{15})))\ \land
      P PLT_FORM (exec (SLc pltForm)) \wedge
      P PLT_FORM (exec (SLc pltSecureHalt)) \wedge
      P PLT_FORM (exec (SLc complete)) \wedge
       (\forall v_{18}.\ P PLT_MOVE (exec (ESCc v_{18}))) \land
       P PLT_MOVE (exec (SLc pltForm)) \wedge
      P PLT_MOVE (exec (SLc pltMove)) \wedge
      P PLT_MOVE (exec (SLc complete)) \wedge
      (\forall v_{21}. \ P \ \text{PLT\_SECURE\_HALT (exec (ESCc} \ v_{21}))) \ \land
      P PLT_SECURE_HALT (exec (SLc pltForm)) \wedge
      P PLT_SECURE_HALT (exec (SLc pltMove)) \wedge
      P PLT_SECURE_HALT (exec (SLc pltSecureHalt)) \wedge
      (\forall v_{23}. \ P \ \texttt{COMPLETE} \ (\texttt{exec} \ v_{23})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
[moveToORPOut_def]
 ⊢ (moveToORPOut MOVE_TO_ORP (exec (SLc pltForm)) = PLTForm) ∧
    (moveToORPOut MOVE_TO_ORP (exec (SLc incomplete)) =
     MoveToORP) ∧
    (moveToORPOut PLT_FORM (exec (SLc pltMove)) = PLTMove) \cap \)
    (moveToORPOut PLT_FORM (exec (SLc incomplete)) = PLTForm) \( \)
    (moveToORPOut PLT_MOVE (exec (SLc pltSecureHalt)) =
     PLTSecureHalt) ∧
    (moveToORPOut PLT_MOVE (exec (SLc incomplete)) = PLTMove) \( \)
    (moveToORPOut PLT_SECURE_HALT (exec (SLc complete)) =
     Complete) \( \)
    (moveToORPOut PLT_SECURE_HALT (exec (SLc incomplete)) =
     PLTSecureHalt) \wedge
    (moveToORPOut s (trap (SLc cmd)) = unAuthorized) \land
    (moveToORPOut s (discard (SLc cmd)) = unAuthenticated)
[moveToORPOut_ind]
 \vdash \forall P.
      P MOVE_TO_ORP (exec (SLc pltForm)) \wedge
      P MOVE_TO_ORP (exec (SLc incomplete)) \wedge
      P PLT_FORM (exec (SLc pltMove)) \wedge
      P PLT_FORM (exec (SLc incomplete)) \wedge
      P PLT_MOVE (exec (SLc pltSecureHalt)) \wedge
      P PLT_MOVE (exec (SLc incomplete)) \wedge
```

```
P PLT_SECURE_HALT (exec (SLc complete)) \wedge
       P PLT_SECURE_HALT (exec (SLc incomplete)) \wedge
       (\forall s \ cmd. \ P \ s \ (trap \ (SLc \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (\texttt{discard} \ (\texttt{SLc} \ cmd))) \ \land
       (\forall s \ v_6. \ P \ s \ (discard \ (ESCc \ v_6))) \ \land
       (\forall s \ v_9. \ P \ s \ (trap \ (ESCc \ v_9))) \ \land
       (\forall v_{12}.\ P MOVE_TO_ORP (exec (ESCc v_{12}))) \land
       P MOVE_TO_ORP (exec (SLc pltMove)) \wedge
      P MOVE_TO_ORP (exec (SLc pltSecureHalt)) \wedge
      P MOVE_TO_ORP (exec (SLc complete)) \wedge
       (\forall v_{15}. \ P \ \mathtt{PLT\_FORM} \ (\mathtt{exec} \ (\mathtt{ESCc} \ v_{15}))) \ \land
       P PLT_FORM (exec (SLc pltForm)) \wedge
       P PLT_FORM (exec (SLc pltSecureHalt)) \( \)
       P PLT_FORM (exec (SLc complete)) \wedge
       (\forall v_{18}. \ P \ \text{PLT\_MOVE} \ (\text{exec (ESCc} \ v_{18}))) \ \land
       P PLT_MOVE (exec (SLc pltForm)) \wedge
       P PLT_MOVE (exec (SLc pltMove)) \wedge
       P PLT_MOVE (exec (SLc complete)) \wedge
       (\forall v_{21}. \ P \ \text{PLT\_SECURE\_HALT (exec (ESCc} \ v_{21}))) \ \land
       P PLT_SECURE_HALT (exec (SLc pltForm)) \wedge
       P PLT_SECURE_HALT (exec (SLc pltMove)) \wedge
       P PLT_SECURE_HALT (exec (SLc pltSecureHalt)) \wedge
       (\forall v_{23}. \ P \ \texttt{COMPLETE} \ (\texttt{exec} \ v_{23})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
[PlatoonLeader_exec_slCommand_justified_thm]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
      TR (M, Oi, Os) (exec (SLc slCommand))
         (CFG authTestMoveToORP ssmMoveToORPStateInterp
             (secContextMoveToORP slCommand)
             (Name PlatoonLeader says prop (SOME (SLc slCommand))::
                   ins) s outs)
         (CFG authTestMoveToORP ssmMoveToORPStateInterp
             (secContextMoveToORP slCommand) ins
             (NS \ s (exec (SLc \ slCommand)))
             (Out \ s \ (exec \ (SLc \ slCommand))::outs)) \iff
       authTestMoveToORP
         (Name PlatoonLeader says prop (SOME (SLc slCommand))) \land
      CFGInterpret (M, Oi, Os)
         (CFG authTestMoveToORP ssmMoveToORPStateInterp
             (secContextMoveToORP slCommand)
             (Name PlatoonLeader says prop (SOME (SLc slCommand))::
                   ins) s outs) \wedge
       (M, Oi, Os) sat prop (SOME (SLc slCommand))
[PlatoonLeader_slCommand_lemma]
 \vdash CFGInterpret (M, Oi, Os)
       (CFG authTestMoveToORP ssmMoveToORPStateInterp
           (secContextMoveToORP slCommand)
```

```
(Name PlatoonLeader says prop (SOME (SLc slCommand))::
          ins) s outs) \Rightarrow
(M, Oi, Os) sat prop (SOME (SLc slCommand))
```

#### 13 MoveToORPType Theory

**Built:** 13 May 2018

Parent Theories: indexedLists, patternMatches

## 13.1 Datatypes

```
slCommand = pltForm | pltMove | pltSecureHalt | complete
            | incomplete
slOutput = MoveToORP | PLTForm | PLTMove | PLTSecureHalt
           | Complete | unAuthorized | unAuthenticated
slState = \texttt{MOVE\_TO\_ORP} \mid \texttt{PLT\_FORM} \mid \texttt{PLT\_MOVE} \mid \texttt{PLT\_SECURE\_HALT}
          | COMPLETE
stateRole = PlatoonLeader
13.2
        Theorems
[slCommand_distinct_clauses]
```

```
\vdash pltForm \neq pltMove \land pltForm \neq pltSecureHalt \land
           pltForm \neq complete \land pltForm \neq incomplete \land
           pltMove \neq pltSecureHalt \land pltMove \neq complete \land
           pltMove \neq incomplete \land pltSecureHalt \neq complete \land
           pltSecureHalt \neq incomplete \land complete \neq incomplete
[slOutput_distinct_clauses]
    \vdash MoveToORP \neq PLTForm \land MoveToORP \neq PLTMove \land
           \texttt{MoveToORP} \neq \texttt{PLTSecureHalt} \ \land \ \texttt{MoveToORP} \neq \texttt{Complete} \ \land
           {	t MoveToORP} 
eq {	t unAuthorized} \land {	t MoveToORP} 
eq {	t unAuthenticated} \land
           {\tt PLTForm} \neq {\tt PLTMove} \ \land \ {\tt PLTForm} \neq {\tt PLTSecureHalt} \ \land
           PLTForm \neq Complete \land PLTForm \neq unAuthorized \land
           {\tt PLTForm} \neq {\tt unAuthenticated} \ \land \ {\tt PLTMove} \neq {\tt PLTSecureHalt} \ \land
           {\tt PLTMove} \, \neq \, {\tt Complete} \, \wedge \, {\tt PLTMove} \, \neq \, {\tt unAuthorized} \, \wedge \,
          PLTMove \neq unAuthenticated \wedge PLTSecureHalt \neq Complete \wedge
          PLTSecureHalt \neq unAuthorized \land
          {\tt PLTSecureHalt} \, \neq \, {\tt unAuthenticated} \, \wedge \, {\tt Complete} \, \neq \, {\tt unAuthorized} \, \wedge \,
          {\tt Complete} \, \neq \, {\tt unAuthenticated} \, \wedge \, {\tt unAuthorized} \, \neq \, {\tt unAuthenticated}
slState_distinct_clauses
    \vdash MOVE_TO_ORP \neq PLT_FORM \land MOVE_TO_ORP \neq PLT_MOVE \land
           	exttt{MOVE\_TO\_ORP} 
eq 	exttt{PLT\_SECURE\_HALT} 
which wove\_TO\_ORP 
eq 	exttt{COMPLETE} 
which wove\_TO_ORP 
eq 	exttt{COMPLE
           {\tt PLT\_FORM} \ \neq \ {\tt PLT\_MOVE} \ \land \ {\tt PLT\_FORM} \ \neq \ {\tt PLT\_SECURE\_HALT} \ \land \\
           PLT\_FORM \neq COMPLETE \land PLT\_MOVE \neq PLT\_SECURE\_HALT \land
           PLT_MOVE ≠ COMPLETE ∧ PLT_SECURE_HALT ≠ COMPLETE
```

**Built:** 13 May 2018

## 14 ssmMoveToPB Theory

Parent Theories: MoveToPBType, ssm11, OMNIType

```
14.1
        Definitions
[secContextMoveToPB_def]
 \vdash \forall cmd.
      secContextMoveToPB \ cmd =
      [Name PlatoonLeader controls prop (SOME (SLc cmd))]
[ssmMoveToPBStateInterp_def]
 \vdash \ \forall \, state. ssmMoveToPBStateInterp state = TT
        Theorems
14.2
[authTestMoveToPB_cmd_reject_lemma]
 \vdash \forall \, cmd. \neg authTestMoveToPB (prop (SOME cmd))
[authTestMoveToPB_def]
 \vdash (authTestMoveToPB (Name PlatoonLeader says prop cmd) \iff T) \land
    (authTestMoveToPB TT \iff F) \land (authTestMoveToPB FF \iff F) \land
    (authTestMoveToPB (prop v) \iff F) \land
    (authTestMoveToPB (notf v_1) \iff F) \land
    (authTestMoveToPB (v_2 andf v_3) \iff F) \wedge
    (authTestMoveToPB (v_4 orf v_5) \iff F) \wedge
    (authTestMoveToPB (v_6 impf v_7) \iff F) \land
    (authTestMoveToPB (v_8 eqf v_9) \iff F) \land
    (authTestMoveToPB (v_{10} says TT) \iff F) \wedge
    (authTestMoveToPB (v_{10} says FF) \iff F) \wedge
    (authTestMoveToPB (v133 meet v134 says prop v_{66}) \iff F) \land
    (authTestMoveToPB (v135 quoting v136 says prop v_{66}) \iff F) \wedge
    (authTestMoveToPB (v_{10} says notf v_{67}) \iff F) \wedge
    (authTestMoveToPB (v_{10} says (v_{68} andf v_{69})) \iff F) \land
    (authTestMoveToPB (v_{10} says (v_{70} orf v_{71})) \iff F) \land
    (authTestMoveToPB (v_{10} says (v_{72} impf v_{73})) \iff F) \land
    (authTestMoveToPB (v_{10} says (v_{74} eqf v_{75})) \iff F) \land
    (authTestMoveToPB (v_{10} says v_{76} says v_{77}) \iff F) \land
    (authTestMoveToPB (v_{10} says v_{78} speaks_for v_{79}) \iff F) \wedge
    (authTestMoveToPB (v_{10} says v_{80} controls v_{81}) \iff F) \wedge
    (authTestMoveToPB (v_{10} says reps v_{82} v_{83} v_{84}) \iff F) \wedge
    (authTestMoveToPB (v_{10} says v_{85} domi v_{86}) \iff F) \wedge
    (authTestMoveToPB (v_{10} says v_{87} eqi v_{88}) \iff F) \wedge
    (authTestMoveToPB (v_{10} says v_{89} doms v_{90}) \iff F) \land
    (authTestMoveToPB (v_{10} says v_{91} eqs v_{92}) \iff F) \land
```

(authTestMoveToPB ( $v_{10}$  says  $v_{93}$  eqn  $v_{94}$ )  $\iff$  F)  $\wedge$ 

```
(authTestMoveToPB (v_{10} says v_{95} lte v_{96}) \iff F) \wedge
       (authTestMoveToPB (v_{10} says v_{97} lt v_{98}) \iff F) \wedge
       (authTestMoveToPB (v_{12} speaks_for v_{13}) \iff F) \wedge
       (authTestMoveToPB (v_{14} controls v_{15}) \iff F) \wedge
       (authTestMoveToPB (reps v_{16} v_{17} v_{18}) \iff F) \land
       (authTestMoveToPB (v_{19} domi v_{20}) \iff F) \wedge
      (authTestMoveToPB (v_{21} eqi v_{22}) \iff F) \wedge
      (authTestMoveToPB (v_{23} doms v_{24}) \iff F) \wedge
      (authTestMoveToPB (v_{25} eqs v_{26}) \iff F) \land
       (authTestMoveToPB (v_{27} eqn v_{28}) \iff F) \wedge
       (authTestMoveToPB (v_{29} lte v_{30}) \iff F) \wedge
       (authTestMoveToPB (v_{31} lt v_{32}) \iff F)
[authTestMoveToPB_ind]
  \vdash \forall P.
          (\forall \, cmd \,.\,\, P (Name PlatoonLeader says prop cmd)) \land \,\, P TT \land
          P \text{ FF } \wedge (\forall v. P \text{ (prop } v)) \wedge (\forall v_1. P \text{ (notf } v_1)) \wedge
           (\forall \ v_2 \ v_3 . P (v_2 andf v_3)) \wedge (\forall \ v_4 \ v_5 . P (v_4 orf v_5)) \wedge
           (\forall v_6 \ v_7. \ P \ (v_6 \ \text{impf} \ v_7)) \land (\forall v_8 \ v_9. \ P \ (v_8 \ \text{eqf} \ v_9)) \land
           (\forall v_{10}. \ P \ (v_{10} \ \text{says TT})) \land (\forall v_{10}. \ P \ (v_{10} \ \text{says FF})) \land
           (\forall v133 \ v134 \ v_{66}. \ P \ (v133 \ \text{meet} \ v134 \ \text{says prop} \ v_{66})) \ \land
           (\forall\,v135\ v136\ v_{66}. P (v135 quoting v136 says prop v_{66})) \wedge
           (\forall v_{10} \ v_{67}. \ P \ (v_{10} \ \text{says notf} \ v_{67})) \land
           (\forall v_{10} \ v_{68} \ v_{69}. \ P \ (v_{10} \ \text{says} \ (v_{68} \ \text{andf} \ v_{69}))) \land
           (\forall v_{10} \ v_{70} \ v_{71}. \ P \ (v_{10} \ \text{says} \ (v_{70} \ \text{orf} \ v_{71}))) \land
           (\forall v_{10} \ v_{72} \ v_{73}. \ P \ (v_{10} \ \text{says} \ (v_{72} \ \text{impf} \ v_{73}))) \ \land
           (\forall v_{10} \ v_{74} \ v_{75}. \ P \ (v_{10} \ \text{says} \ (v_{74} \ \text{eqf} \ v_{75}))) \ \land
           (\forall v_{10} \ v_{76} \ v_{77}. \ P \ (v_{10} \ \text{says} \ v_{76} \ \text{says} \ v_{77})) \ \land
           (\forall \, v_{10} \ v_{78} \ v_{79}. P (v_{10} says v_{78} speaks_for v_{79})) \wedge
           (\forall \, v_{10} \ v_{80} \ v_{81}. P (v_{10} says v_{80} controls v_{81})) \wedge
           (\forall v_{10} \ v_{82} \ v_{83} \ v_{84}. \ P \ (v_{10} \ \text{says reps} \ v_{82} \ v_{83} \ v_{84})) \ \land
           (\forall v_{10} \ v_{85} \ v_{86}. P (v_{10} says v_{85} domi v_{86})) \land
           (\forall v_{10} \ v_{87} \ v_{88}. P (v_{10} says v_{87} eqi v_{88})) \land
           (\forall v_{10} \ v_{89} \ v_{90}. \ P \ (v_{10} \ {\tt says} \ v_{89} \ {\tt doms} \ v_{90})) \ \land
           (\forall v_{10} \ v_{91} \ v_{92}. P (v_{10} says v_{91} eqs v_{92})) \land
           (\forall v_{10} \ v_{93} \ v_{94}. \ P \ (v_{10} \ \text{says} \ v_{93} \ \text{eqn} \ v_{94})) \ \land
           (\forall v_{10} \ v_{95} \ v_{96}. \ P \ (v_{10} \ {\tt says} \ v_{95} \ {\tt lte} \ v_{96})) \ \land
           (\forall v_{10} \ v_{97} \ v_{98}. \ P \ (v_{10} \ {\tt says} \ v_{97} \ {\tt lt} \ v_{98})) \ \land \ 
           (\forall v_{12} \ v_{13}. \ P \ (v_{12} \ \text{speaks\_for} \ v_{13})) \land
           (\forall v_{14} \ v_{15}. \ P \ (v_{14} \ \text{controls} \ v_{15})) \land
           (\forall v_{16} \ v_{17} \ v_{18}. \ P \ (reps \ v_{16} \ v_{17} \ v_{18})) \ \land
           (\forall v_{19} \ v_{20}. P (v_{19} domi v_{20})) \land
           (\forall v_{21} \ v_{22}. \ P \ (v_{21} \ \mathsf{eqi} \ v_{22})) \ \land
           (\forall v_{23} \ v_{24}. \ P \ (v_{23} \ \text{doms} \ v_{24})) \land
           (\forall v_{25} \ v_{26}. \ P \ (v_{25} \ \text{eqs} \ v_{26})) \ \land \ (\forall v_{27} \ v_{28}. \ P \ (v_{27} \ \text{eqn} \ v_{28})) \ \land
           (\forall v_{29} \ v_{30}. \ P \ (v_{29} \ \text{lte} \ v_{30})) \land (\forall v_{31} \ v_{32}. \ P \ (v_{31} \ \text{lt} \ v_{32})) \Rightarrow
          \forall v. P v
```

[moveToPBNS\_def]

```
⊢ (moveToPBNS MOVE_TO_PB (exec (SLc pltForm)) = PLT_FORM) ∧
    (moveToPBNS MOVE_TO_PB (exec (SLc incomplete)) =
     MOVE_TO_PB) ∧
    (moveToPBNS PLT_FORM (exec (SLc pltMove)) = PLT_MOVE) \cap \)
    (moveToPBNS PLT_FORM (exec (SLc incomplete)) = PLT_FORM) \(\lambda\)
    (moveToPBNS PLT_MOVE (exec (SLc pltHalt)) = PLT_HALT) \cap 
    (moveToPBNS PLT_MOVE (exec (SLc incomplete)) = PLT_MOVE) \cap \( \)
    (moveToPBNS PLT_HALT (exec (SLc complete)) = COMPLETE) \land
    (moveToPBNS PLT_HALT (exec (SLc incomplete)) = PLT_HALT) \cap \( \)
    (moveToPBNS s (trap (SLc cmd)) = s) \land
    (moveToPBNS s (discard (SLc cmd)) = s)
[moveToPBNS_ind]
 \vdash \forall P.
       P MOVE_TO_PB (exec (SLc pltForm)) \wedge
      P MOVE_TO_PB (exec (SLc incomplete)) \land
      P PLT_FORM (exec (SLc pltMove)) \wedge
      P PLT_FORM (exec (SLc incomplete)) \wedge
      P PLT_MOVE (exec (SLc pltHalt)) \wedge
      P PLT_MOVE (exec (SLc incomplete)) \wedge
      P PLT_HALT (exec (SLc complete)) \wedge
      P PLT_HALT (exec (SLc incomplete)) \wedge
       (\forall s \ cmd. \ P \ s \ (trap \ (SLc \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (\texttt{discard} \ (\texttt{SLc} \ cmd))) \ \land
       (\forall s \ v_6. \ P \ s \ (discard \ (ESCc \ v_6))) \land
       (\forall s \ v_9. \ P \ s \ (trap \ (ESCc \ v_9))) \ \land
       (\forall v_{12}. P MOVE_TO_PB (exec (ESCc v_{12}))) \land
       P MOVE_TO_PB (exec (SLc pltMove)) \wedge
      P MOVE_TO_PB (exec (SLc pltHalt)) \wedge
       P MOVE_TO_PB (exec (SLc complete)) \wedge
       (\forall v_{15}. \ P \ \mathtt{PLT\_FORM} \ (\mathtt{exec} \ (\mathtt{ESCc} \ v_{15}))) \ \land
       P PLT_FORM (exec (SLc pltForm)) \wedge
      P PLT_FORM (exec (SLc pltHalt)) \wedge
      P PLT_FORM (exec (SLc complete)) \wedge
       (\forall v_{18}.\ P PLT_MOVE (exec (ESCc v_{18}))) \land
       P PLT_MOVE (exec (SLc pltForm)) \wedge
      P PLT_MOVE (exec (SLc pltMove)) \wedge
      P PLT_MOVE (exec (SLc complete)) \wedge
       (\forall v_{21}. \ P \ PLT\_HALT \ (exec \ (ESCc \ v_{21}))) \land
      P PLT_HALT (exec (SLc pltForm)) \wedge
      P PLT_HALT (exec (SLc pltMove)) \wedge
       P PLT_HALT (exec (SLc pltHalt)) \wedge
       (\forall v_{23}. \ P \ \texttt{COMPLETE} \ (\texttt{exec} \ v_{23})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
[moveToPBOut_def]
 ⊢ (moveToPBOut MOVE_TO_PB (exec (SLc pltForm)) = PLTForm) ∧
    (moveToPBOut MOVE_TO_PB (exec (SLc incomplete)) = MoveToPB) \(\lambda\)
    (moveToPBOut PLT_FORM (exec (SLc pltMove)) = PLTMove) \(\lambda\)
```

```
(moveToPBOut PLT_FORM (exec (SLc incomplete)) = PLTForm) \( \)
    (moveToPBOut PLT_MOVE (exec (SLc pltHalt)) = PLTHalt) \(\lambda\)
    (moveToPBOut PLT_MOVE (exec (SLc incomplete)) = PLTMove) \( \lambda \)
    (moveToPBOut PLT_HALT (exec (SLc complete)) = Complete) ^
    (moveToPBOut PLT_HALT (exec (SLc incomplete)) = PLTHalt) \(\lambda\)
    (moveToPBOut s (trap (SLc cmd)) = unAuthorized) \land
    (moveToPBOut s (discard (SLc cmd)) = unAuthenticated)
[moveToPBOut_ind]
 \vdash \forall P.
      P MOVE_TO_PB (exec (SLc pltForm)) \wedge
      P MOVE_TO_PB (exec (SLc incomplete)) \wedge
      P PLT_FORM (exec (SLc pltMove)) \wedge
      P PLT_FORM (exec (SLc incomplete)) \wedge
      P PLT_MOVE (exec (SLc pltHalt)) \wedge
      P PLT_MOVE (exec (SLc incomplete)) \wedge
      P PLT_HALT (exec (SLc complete)) \wedge
      P PLT_HALT (exec (SLc incomplete)) \wedge
       (\forall s \ cmd. \ P \ s \ (trap \ (SLc \ cmd))) \ \land
       (\forall s \ cmd. \ P \ s \ (discard \ (SLc \ cmd))) \ \land
       (\forall s \ v_6. \ P \ s \ (discard \ (ESCc \ v_6))) \land
       (\forall s \ v_9. \ P \ s \ (trap \ (ESCc \ v_9))) \land
       (\forall v_{12}. \ P \ \texttt{MOVE\_TO\_PB} \ (\texttt{exec} \ (\texttt{ESCc} \ v_{12}))) \ \land
      P MOVE_TO_PB (exec (SLc pltMove)) \wedge
      P MOVE_TO_PB (exec (SLc pltHalt)) \wedge
      P MOVE_TO_PB (exec (SLc complete)) \wedge
       (\forall v_{15}.\ P PLT_FORM (exec (ESCc v_{15}))) \land
       P PLT_FORM (exec (SLc pltForm)) \wedge
      P PLT_FORM (exec (SLc pltHalt)) \wedge
      P PLT_FORM (exec (SLc complete)) \wedge
       (\forall v_{18}.\ P PLT_MOVE (exec (ESCc v_{18}))) \land
       P PLT_MOVE (exec (SLc pltForm)) \wedge
      P PLT_MOVE (exec (SLc pltMove)) \wedge
      P PLT_MOVE (exec (SLc complete)) \wedge
       (\forall v_{21}. \ P \ \text{PLT\_HALT} \ (\text{exec (ESCc } v_{21}))) \ \land
      P PLT_HALT (exec (SLc pltForm)) \wedge
      P PLT_HALT (exec (SLc pltMove)) \wedge
      P PLT_HALT (exec (SLc pltHalt)) \wedge
       (\forall v_{23}. \ P \ \texttt{COMPLETE} \ (\texttt{exec} \ v_{23})) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
[PlatoonLeader_exec_slCommand_justified_thm]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
      TR (M, Oi, Os) (exec (SLc slCommand))
         (CFG authTestMoveToPB ssmMoveToPBStateInterp
             (secContextMoveToPB slCommand)
             (Name PlatoonLeader says prop (SOME (SLc slCommand))::
                   ins) s outs)
         (CFG authTestMoveToPB ssmMoveToPBStateInterp
```

```
(secContextMoveToPB slCommand) ins
           (NS \ s \ (exec \ (SLc \ slCommand)))
           (Out \ s \ (exec \ (SLc \ slCommand))::outs)) \iff
     authTestMoveToPB
        (Name PlatoonLeader says prop (SOME (SLc slCommand))) \land
     CFGInterpret (M, Oi, Os)
        (CFG authTestMoveToPB ssmMoveToPBStateInterp
           (secContextMoveToPB slCommand)
           (Name PlatoonLeader says prop (SOME (SLc slCommand))::
                ins) s outs) \wedge
      (M, Oi, Os) sat prop (SOME (SLc slCommand))
[PlatoonLeader_slCommand_lemma]
 \vdash CFGInterpret (M, Oi, Os)
      (CFG authTestMoveToPB ssmMoveToPBStateInterp
         ({\tt secContextMoveToPB}\ \mathit{slCommand})
         (Name PlatoonLeader says prop (SOME (SLc slCommand))::
              ins) s outs) \Rightarrow
    (M, Oi, Os) sat prop (SOME (SLc slCommand))
```

## 15 MoveToPBType Theory

**Built:** 13 May 2018

Parent Theories: indexedLists, patternMatches

## 15.1 Datatypes

## 15.2 Theorems

```
[slCommand_distinct_clauses]

⊢ pltForm ≠ pltMove ∧ pltForm ≠ pltHalt ∧ pltForm ≠ complete ∧
    pltForm ≠ incomplete ∧ pltMove ≠ pltHalt ∧
    pltMove ≠ complete ∧ pltMove ≠ incomplete ∧
    pltHalt ≠ complete ∧ pltHalt ≠ incomplete ∧
    complete ≠ incomplete
```

```
[slOutput_distinct_clauses]
  \vdash MoveToPB \neq PLTForm \land MoveToPB \neq PLTMove \land
     MoveToPB \neq PLTHalt \land MoveToPB \neq Complete \land
     	exttt{MoveToPB} 
eq 	exttt{unAuthorized} 
abla 	exttt{MoveToPB} 
eq 	exttt{unAuthenticated} 
abla 
     {\tt PLTForm} \neq {\tt PLTMove} \ \land \ {\tt PLTForm} \neq {\tt PLTHalt} \ \land \ {\tt PLTForm} \neq {\tt Complete} \ \land
     {\tt PLTForm} \neq {\tt unAuthorized} \ \land \ {\tt PLTForm} \neq {\tt unAuthenticated} \ \land
     {\tt PLTMove} \neq {\tt PLTHalt} \ \land \ {\tt PLTMove} \neq {\tt Complete} \ \land
     {\tt PLTMove} \, \neq \, {\tt unAuthorized} \, \wedge \, {\tt PLTMove} \, \neq \, {\tt unAuthenticated} \, \wedge \,
     {\tt PLTHalt} \neq {\tt Complete} \ \land \ {\tt PLTHalt} \neq {\tt unAuthorized} \ \land \\
     PLTHalt \neq unAuthenticated \wedge Complete \neq unAuthorized \wedge
     {\tt Complete} \neq {\tt unAuthenticated} \ \land \ {\tt unAuthorized} \neq {\tt unAuthenticated}
[slState_distinct_clauses]
 \vdash MOVE_TO_PB \neq PLT_FORM \land MOVE_TO_PB \neq PLT_MOVE \land
     	exttt{MOVE\_TO\_PB} \neq 	exttt{PLT\_HALT} \land 	exttt{MOVE\_TO\_PB} \neq 	exttt{COMPLETE} \land
     {\tt PLT\_FORM} \, \neq \, {\tt PLT\_MOVE} \, \wedge \, {\tt PLT\_FORM} \, \neq \, {\tt PLT\_HALT} \, \wedge \,
     \mathtt{PLT\_FORM} \neq \mathtt{COMPLETE} \ \land \ \mathtt{PLT\_MOVE} \neq \mathtt{PLT\_HALT} \ \land
     PLT\_MOVE \neq COMPLETE \land PLT\_HALT \neq COMPLETE
```

## 16 ssmPlanPB Theory

**Built:** 13 May 2018

Parent Theories: PlanPBDef, ssm

#### 16.1 Theorems

```
[inputOK_def]
```

```
\vdash (inputOK (Name PlatoonLeader says prop cmd) \iff T) \land
   (inputOK (Name PlatoonSergeant says prop cmd) \iff T) \land
   (inputOK TT \iff F) \land (inputOK FF \iff F) \land
   (inputOK (prop v) \iff F) \land (inputOK (notf v_1) \iff F) \land
   (inputOK (v_2 andf v_3) \iff F) \wedge (inputOK (v_4 orf v_5) \iff F) \wedge
   (inputOK (v_6 impf v_7) \iff F) \land (inputOK (v_8 eqf v_9) \iff F) \land
   (inputOK (v_{10} says TT) \iff F) \wedge (inputOK (v_{10} says FF) \iff F) \wedge
   (inputOK (v133 meet v134 says prop v_{66}) \iff F) \land
   (inputOK (v135 quoting v136 says prop v_{66}) \iff F) \land
   (inputOK (v_{10} says notf v_{67}) \iff F) \wedge
   (inputOK (v_{10} says (v_{68} andf v_{69})) \iff F) \land
   (inputOK (v_{10} says (v_{70} orf v_{71})) \iff F) \land
   (inputOK (v_{10} says (v_{72} impf v_{73})) \iff F) \wedge
   (inputOK (v_{10} says (v_{74} eqf v_{75})) \iff F) \land
   (inputOK (v_{10} says v_{76} says v_{77}) \iff F) \wedge
   (inputOK (v_{10} says v_{78} speaks_for v_{79}) \iff F) \wedge
   (inputOK (v_{10} says v_{80} controls v_{81}) \iff F) \wedge
   (inputOK (v_{10} says reps v_{82} v_{83} v_{84}) \iff F) \land
   (inputOK (v_{10} says v_{85} domi v_{86}) \iff F) \wedge
   (inputOK (v_{10} says v_{87} eqi v_{88}) \iff F) \wedge
```

```
(inputOK (v_{10} says v_{89} doms v_{90}) \iff F) \land
       (inputOK (v_{10} says v_{91} eqs v_{92}) \iff F) \land
       (inputOK (v_{10} says v_{93} eqn v_{94}) \iff F) \wedge
       (inputOK (v_{10} says v_{95} lte v_{96}) \iff F) \wedge
       (inputOK (v_{10} says v_{97} lt v_{98}) \iff F) \wedge
       (inputOK (v_{12} speaks_for v_{13}) \iff F) \wedge
       (inputOK (v_{14} controls v_{15}) \iff F) \wedge
       (inputOK (reps v_{16} v_{17} v_{18}) \iff F) \land
       (inputOK (v_{19} domi v_{20}) \iff F) \land
       (inputOK (v_{21} eqi v_{22}) \iff F) \wedge
       (inputOK (v_{23} doms v_{24}) \iff F) \wedge
       (inputOK (v_{25} eqs v_{26}) \iff F) \wedge (inputOK (v_{27} eqn v_{28}) \iff F) \wedge (inputOK (v_{29} lte v_{30}) \iff F) \wedge (inputOK (v_{31} lt v_{32}) \iff F)
[inputOK_ind]
  \vdash \forall P.
           (\forall \, cmd \, . \, \, P \, \, ({\tt Name \, \, PlatoonLeader \, \, says \, \, prop \, \, } cmd)) \, \, \wedge \, \,
           (\forall \, cmd \,.\,\, P (Name PlatoonSergeant says prop cmd)) \wedge \,\, P TT \wedge
           P \text{ FF } \wedge (\forall v. P \text{ (prop } v)) \wedge (\forall v_1. P \text{ (notf } v_1)) \wedge
           (\forall v_2 \ v_3. \ P \ (v_2 \ \text{andf} \ v_3)) \land (\forall v_4 \ v_5. \ P \ (v_4 \ \text{orf} \ v_5)) \land
           (\forall v_6 \ v_7. \ P \ (v_6 \ \text{impf} \ v_7)) \land (\forall v_8 \ v_9. \ P \ (v_8 \ \text{eqf} \ v_9)) \land
           (\forall v_{10}. \ P \ (v_{10} \ \text{says TT})) \ \land \ (\forall v_{10}. \ P \ (v_{10} \ \text{says FF})) \ \land
           (\forall\,v133\ v134\ v_{66}. P (v133 meet v134 says prop v_{66})) \wedge
           (\forall v135 \ v136 \ v_{66}. P (v135 quoting v136 says prop v_{66})) \land
           (\forall v_{10} \ v_{67}. P (v_{10} says notf v_{67})) \land
           (\forall v_{10} \ v_{68} \ v_{69}. \ P \ (v_{10} \ \text{says} \ (v_{68} \ \text{andf} \ v_{69}))) \land
           (\forall v_{10} \ v_{70} \ v_{71}. \ P \ (v_{10} \ \text{says} \ (v_{70} \ \text{orf} \ v_{71}))) \land
           (\forall v_{10} \ v_{72} \ v_{73}. \ P \ (v_{10} \ \text{says} \ (v_{72} \ \text{impf} \ v_{73}))) \ \land
           (\forall v_{10} \ v_{74} \ v_{75}. \ P \ (v_{10} \ {\tt says} \ (v_{74} \ {\tt eqf} \ v_{75}))) \ \land
           (\forall v_{10} \ v_{76} \ v_{77}. \ P \ (v_{10} \ \text{says} \ v_{76} \ \text{says} \ v_{77})) \ \land
           (\forall \, v_{10} \ v_{78} \ v_{79}. P (v_{10} says v_{78} speaks_for v_{79})) \wedge
           (\forall v_{10} \ v_{80} \ v_{81}. \ P \ (v_{10} \ \text{says} \ v_{80} \ \text{controls} \ v_{81})) \ \land
           (\forall v_{10} \ v_{82} \ v_{83} \ v_{84}. \ P \ (v_{10} \ {\tt says \ reps} \ v_{82} \ v_{83} \ v_{84})) \ \land
           (\forall v_{10} v_{85} v_{86}. P (v_{10} says v_{85} domi v_{86})) \wedge
           (\forall v_{10} \ v_{87} \ v_{88}. \ P \ (v_{10} \ {\tt says} \ v_{87} \ {\tt eqi} \ v_{88})) \ \land
           (\forall v_{10} \ v_{89} \ v_{90}. \ P \ (v_{10} \ \text{says} \ v_{89} \ \text{doms} \ v_{90})) \ \land
           (\forall v_{10} \ v_{91} \ v_{92}. \ P \ (v_{10} \ \text{says} \ v_{91} \ \text{eqs} \ v_{92})) \ \land
           (\forall v_{10} \ v_{93} \ v_{94}. \ P \ (v_{10} \ \text{says} \ v_{93} \ \text{eqn} \ v_{94})) \ \land
           (\forall v_{10} \ v_{95} \ v_{96}. P (v_{10} says v_{95} lte v_{96})) \wedge
           (\forall v_{10} \ v_{97} \ v_{98}. \ P \ (v_{10} \ \text{says} \ v_{97} \ \text{lt} \ v_{98})) \ \land
           (\forall v_{12} \ v_{13}. \ P \ (v_{12} \ \text{speaks\_for} \ v_{13})) \land
           (\forall v_{14} \ v_{15}. \ P \ (v_{14} \ \text{controls} \ v_{15})) \ \land
           (\forall v_{16} \ v_{17} \ v_{18}. \ P \ (reps \ v_{16} \ v_{17} \ v_{18})) \ \land
           (\forall v_{19} \ v_{20}. \ P \ (v_{19} \ \text{domi} \ v_{20})) \ \land
           (\forall v_{21} \ v_{22}. \ P \ (v_{21} \ \text{eqi} \ v_{22})) \ \land
           (\forall v_{23} \ v_{24}. \ P \ (v_{23} \ \text{doms} \ v_{24})) \land
           (\forall v_{25} \ v_{26}. \ P \ (v_{25} \ \text{eqs} \ v_{26})) \land (\forall v_{27} \ v_{28}. \ P \ (v_{27} \ \text{eqn} \ v_{28})) \land
           (\forall v_{29} \ v_{30}. \ P \ (v_{29} \ \text{lte} \ v_{30})) \land (\forall v_{31} \ v_{32}. \ P \ (v_{31} \ \text{lt} \ v_{32})) \Rightarrow
          \forall v. P v
```

```
[planPBNS_def]
 \vdash (planPBNS WARNO (exec x) =
     if
        (getRecon x = [SOME (SLc (PL recon))]) \land
        (getTenativePlan x = [SOME (SLc (PL tentativePlan))]) \land
        (getReport x = [SOME (SLc (PL report1))]) \land
        (getInitMove x = [SOME (SLc (PSG initiateMovement))])
     then
        REPORT1
     else WARNO) ∧
     (planPBNS PLAN_PB (exec x) =
     if getPlCom x = receiveMission then RECEIVE_MISSION
     else PLAN_PB) ∧
     (planPBNS RECEIVE_MISSION (exec x) =
     if getPlCom x = warno then WARNO else RECEIVE_MISSION) \wedge
     (planPBNS REPORT1 (exec x) =
     if getPlCom x = completePlan then COMPLETE_PLAN
     else REPORT1) ∧
     (planPBNS COMPLETE_PLAN (exec x) =
     if getPlCom x = opoid then OPOID else COMPLETE_PLAN) \wedge
    (planPBNS OPOID (exec x) =
     if getPlCom x = supervise then SUPERVISE else OPOID) \wedge
    (planPBNS SUPERVISE (exec x) =
     if getPlCom x = report2 then REPORT2 else SUPERVISE) \wedge
     (planPBNS REPORT2 (exec x) =
     if getPlCom x = complete then COMPLETE else REPORT2) \wedge
     (planPBNS s (trap v_0) = s) \wedge (planPBNS s (discard v_1) = s)
[planPBNS_ind]
 \vdash \forall P.
       (\forall x. \ P \ \text{WARNO (exec} \ x)) \ \land \ (\forall x. \ P \ \text{PLAN\_PB (exec} \ x)) \ \land
       (\forall x. P RECEIVE\_MISSION (exec x)) \land
       (\forall x. \ P \ \text{REPORT1 (exec} \ x)) \land (\forall x. \ P \ \text{COMPLETE\_PLAN (exec} \ x)) \land
       (\forall x. \ P \ \text{OPOID (exec} \ x)) \ \land \ (\forall x. \ P \ \text{SUPERVISE (exec} \ x)) \ \land
       (\forall x. \ P \ \texttt{REPORT2} \ (\texttt{exec} \ x)) \ \land \ (\forall s \ v_0. \ P \ s \ (\texttt{trap} \ v_0)) \ \land
       (\forall s \ v_1. \ P \ s \ (\texttt{discard} \ v_1)) \ \land
       (\forall v_6. P TENTATIVE_PLAN (exec v_6)) \land
       (\forall v_7. P INITIATE_MOVEMENT (exec v_7)) \land
       (\forall v_8. \ P \ \text{RECON (exec} \ v_8)) \land (\forall v_9. \ P \ \text{COMPLETE (exec} \ v_9)) \Rightarrow
      \forall v \ v_1 . \ P \ v \ v_1
[planPBOut_def]
 \vdash (planPBOut WARNO (exec x) =
     if
        (getRecon x = [SOME (SLc (PL recon))]) \land
        (getTenativePlan x = [SOME (SLc (PL tentativePlan))]) \land
        (getReport x = [SOME (SLc (PL report1))]) \land
        (getInitMove x = [SOME (SLc (PSG initiateMovement))])
```

```
then
        Report1
     else unAuthorized) \wedge
    (planPBOut PLAN_PB (exec x) =
     if getPlCom x = receiveMission then ReceiveMission
     else unAuthorized) ∧
    (planPBOut RECEIVE_MISSION (exec x) =
     if getPlCom x = warno then Warno else unAuthorized) \land
    (planPBOut REPORT1 (exec x) =
     if getPlCom x = completePlan then CompletePlan
     \textbf{else} \text{ unAuthorized) } \land \\
    (planPBOut COMPLETE_PLAN (exec x) =
     if getPlCom x = opoid then Opoid else unAuthorized) \wedge
    (planPBOut OPOID (exec x) =
     if getPlCom x = supervise then Supervise
     else unAuthorized) ∧
    (planPBOut SUPERVISE (exec x) =
     if getPlCom x = report2 then Report2 else unAuthorized) \land
    (planPBOut REPORT2 (exec x) =
     if getPlCom x = complete then Complete else unAuthorized) \wedge
    (planPBOut s (trap v_0) = unAuthorized) \wedge
    (planPBOut s (discard v_1) = unAuthenticated)
[planPBOut_ind]
 \vdash \forall P.
       (\forall x. \ P \ \text{WARNO (exec} \ x)) \ \land \ (\forall x. \ P \ \text{PLAN\_PB (exec} \ x)) \ \land
       (\forall x.\ P\ \texttt{RECEIVE\_MISSION} (exec x)) \land
       (\forall x.\ P\ \text{REPORT1}\ (\text{exec}\ x))\ \land\ (\forall x.\ P\ \text{COMPLETE\_PLAN}\ (\text{exec}\ x))\ \land
       (\forall x. P \text{ OPOID (exec } x)) \land (\forall x. P \text{ SUPERVISE (exec } x)) \land
       (\forall x.\ P\ \text{REPORT2 (exec }x))\ \land\ (\forall s\ v_0.\ P\ s\ (\text{trap }v_0))\ \land
       (\forall s \ v_1. \ P \ s \ (discard \ v_1)) \ \land
       (\forall \, v_6 \,.\,\,P TENTATIVE_PLAN (exec v_6)) \land
       (\forall v_7. P \text{ INITIATE\_MOVEMENT (exec } v_7)) \land
       (\forall v_8. \ P \ \text{RECON (exec} \ v_8)) \land (\forall v_9. \ P \ \text{COMPLETE (exec} \ v_9)) \Rightarrow
       \forall v \ v_1 . \ P \ v \ v_1
[PlatoonLeader_notWARNO_notreport1_exec_plCommand_justified_lemma]
 \vdash s \neq \mathtt{WARNO} \Rightarrow
    plCommand \neq invalidPlCommand \Rightarrow
    plCommand \neq report1 \Rightarrow
    \forall NS \ Out \ M \ Oi \ Os.
       TR (M, Oi, Os)
          (exec
              (inputList
                  [Name PlatoonLeader says
                   prop (SOME (SLc (PL plCommand)))]))
          (CFG inputOK secContext secContextNull
              ([Name PlatoonLeader says
                prop (SOME (SLc (PL plCommand)))]::ins) s outs)
```

```
(CFG inputOK secContext secContextNull ins
           (NS \ s
               (exec
                  (inputList
                     [Name PlatoonLeader says
                      prop (SOME (SLc (PL plCommand)))])))
           (Out s
               (exec
                  (inputList
                     [Name PlatoonLeader says
                      prop (SOME (SLc (PL plCommand)))]))::
                 outs)) \iff
      authenticationTest inputOK
        [Name PlatoonLeader says
         prop (SOME (SLc (PL plCommand)))] \land
     CFGInterpret (M, Oi, Os)
        (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
             prop (SOME (SLc (PL plCommand)))]::ins) s outs) \land
      (M,Oi,Os) satList
     propCommandList
        [Name PlatoonLeader says
         prop (SOME (SLc (PL plCommand)))]
[PlatoonLeader_notWARNO_notreport1_exec_plCommand_justified_thm]
 \vdash s \neq \mathtt{WARNO} \Rightarrow
   plCommand \neq invalidPlCommand \Rightarrow
   plCommand \neq report1 \Rightarrow
   \forall NS \ Out \ M \ Oi \ Os.
      TR (M, Oi, Os) (exec [SOME (SLc (PL plCommand))])
        (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
             prop (SOME (SLc (PL plCommand)))]::ins) s outs)
        (CFG inputOK secContext secContextNull ins
           (NS \ s \ (exec \ [SOME \ (SLc \ (PL \ plCommand))]))
           (Out \ s \ (exec \ [SOME \ (SLc \ (PL \ plCommand))])::outs)) \iff
      authenticationTest inputOK
        [Name PlatoonLeader says
         prop (SOME (SLc (PL plCommand)))] \land
     CFGInterpret (M, Oi, Os)
        (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
             prop (SOME (SLc (PL plCommand)))]::ins) s outs) \land
      (M, Oi, Os) satList [prop (SOME (SLc (PL plCommand)))]
[PlatoonLeader_notWARNO_notreport1_exec_plCommand_lemma]
 \vdash s \neq WARNO \Rightarrow
   plCommand \neq invalidPlCommand \Rightarrow
   plCommand \neq report1 \Rightarrow
```

```
\forall M \ Oi \ Os.
     CFGInterpret (M, Oi, Os)
        (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
             prop (SOME (SLc (PL plCommand)))]::ins) s outs) \Rightarrow
      (M,Oi,Os) satList
     propCommandList
        [Name PlatoonLeader says
         prop (SOME (SLc (PL plCommand)))]
[PlatoonLeader_psgCommand_notDiscard_thm]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
      \neg \texttt{TR} \ (M, Oi, Os)
         (discard
            (inputList
               [Name PlatoonLeader says
                prop (SOME (SLc (PSG psgCommand)))]))
         (CFG inputOK secContext secContextNull
            ([Name PlatoonLeader says
              prop (SOME (SLc (PSG psgCommand)))]::ins) s outs)
         (CFG inputOK secContext secContextNull ins
            (NS \ s
               (discard
                   (inputList
                      [Name PlatoonLeader says
                       prop (SOME (SLc (PSG psgCommand)))])))
            (Out s
               (discard
                   (inputList
                      [Name PlatoonLeader says
                       prop (SOME (SLc (PSG psgCommand)))]))::
                  outs))
[PlatoonLeader_trap_psgCommand_justified_lemma]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os)
        (trap
           (inputList
              [Name PlatoonLeader says
               prop (SOME (SLc (PSG psgCommand)))]))
        (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
             \verb|prop (SOME (SLc (PSG | psgCommand)))]:: ins) | s | outs)|
        (CFG inputOK secContext secContextNull ins
           (NS s
              (trap
                  (inputList
                     [Name PlatoonLeader says
                      prop (SOME (SLc (PSG psgCommand)))])))
```

```
(Out \ s
              (trap
                 (inputList
                    [Name PlatoonLeader says
                     prop (SOME (SLc (PSG psgCommand)))]))::
                outs)) \iff
     authenticationTest inputOK
        [Name PlatoonLeader says
        prop (SOME (SLc (PSG psqCommand)))] \land
     CFGInterpret (M, Oi, Os)
        (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
             prop (SOME (SLc (PSG psgCommand)))]::ins) s outs) \land
      (M,Oi,Os) sat prop NONE
[PlatoonLeader_trap_psgCommand_lemma]
 \vdash \forall M \ Oi \ Os.
     CFGInterpret (M, Oi, Os)
        (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
             prop (SOME (SLc (PSG psgCommand)))]::ins) s outs) \Rightarrow
      (M,Oi,Os) sat prop NONE
[PlatoonLeader_WARNO_exec_report1_justified_lemma]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os)
        (exec
           (inputList
              [Name PlatoonLeader says
               prop (SOME (SLc (PL recon)));
               Name PlatoonLeader says
               prop (SOME (SLc (PL tentativePlan)));
               Name PlatoonSergeant says
               prop (SOME (SLc (PSG initiateMovement)));
               Name PlatoonLeader says
               prop (SOME (SLc (PL report1)))]))
        (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
             prop (SOME (SLc (PL recon)));
             Name PlatoonLeader says
             prop (SOME (SLc (PL tentativePlan)));
             Name PlatoonSergeant says
             prop (SOME (SLc (PSG initiateMovement)));
             Name PlatoonLeader says
             prop (SOME (SLc (PL report1)))]::ins) WARNO outs)
        (CFG inputOK secContext secContextNull ins
           (NS WARNO
              (exec
                 (inputList
```

```
[Name PlatoonLeader says
                     prop (SOME (SLc (PL recon)));
                     Name PlatoonLeader says
                     prop (SOME (SLc (PL tentativePlan)));
                     Name PlatoonSergeant says
                     prop (SOME (SLc (PSG initiateMovement)));
                     Name PlatoonLeader says
                     prop (SOME (SLc (PL report1)))])))
           (Out WARNO
              (exec
                 (inputList
                    [Name PlatoonLeader says
                     prop (SOME (SLc (PL recon)));
                     Name PlatoonLeader says
                     prop (SOME (SLc (PL tentativePlan)));
                     Name PlatoonSergeant says
                     prop (SOME (SLc (PSG initiateMovement)));
                     Name PlatoonLeader says
                     prop (SOME (SLc (PL report1)))]))::outs)) <=>
     authenticationTest inputOK
        [Name PlatoonLeader says prop (SOME (SLc (PL recon)));
        Name PlatoonLeader says
        prop (SOME (SLc (PL tentativePlan)));
        Name PlatoonSergeant says
        prop (SOME (SLc (PSG initiateMovement)));
        Name PlatoonLeader says
        prop (SOME (SLc (PL report1)))] \cap \)
     CFGInterpret (M, Oi, Os)
       (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
            prop (SOME (SLc (PL recon)));
            Name PlatoonLeader says
            prop (SOME (SLc (PL tentativePlan)));
            Name PlatoonSergeant says
            prop (SOME (SLc (PSG initiateMovement)));
            Name PlatoonLeader says
            prop (SOME (SLc (PL report1)))]::ins) WARNO outs) \land
     (M,Oi,Os) satList
     propCommandList
       [Name PlatoonLeader says prop (SOME (SLc (PL recon)));
        Name PlatoonLeader says
        prop (SOME (SLc (PL tentativePlan)));
        Name PlatoonSergeant says
        prop (SOME (SLc (PSG initiateMovement)));
        Name PlatoonLeader says prop (SOME (SLc (PL report1)))]
[PlatoonLeader_WARNO_exec_report1_justified_thm]
 \vdash \ \forall \, NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os)
```

```
(exec
     [SOME (SLc (PL recon)); SOME (SLc (PL tentativePlan));
      SOME (SLc (PSG initiateMovement));
      SOME (SLc (PL report1))])
  (CFG inputOK secContext secContextNull
     ([Name PlatoonLeader says
       prop (SOME (SLc (PL recon)));
       Name PlatoonLeader says
       prop (SOME (SLc (PL tentativePlan)));
       Name PlatoonSergeant says
       prop (SOME (SLc (PSG initiateMovement)));
       Name PlatoonLeader says
       prop (SOME (SLc (PL report1)))]::ins) WARNO outs)
  (CFG inputOK secContext secContextNull ins
     (NS WARNO
        (exec
           [SOME (SLc (PL recon));
            SOME (SLc (PL tentativePlan));
            SOME (SLc (PSG initiateMovement));
            SOME (SLc (PL report1))]))
     (Out WARNO
        (exec
           [SOME (SLc (PL recon));
            SOME (SLc (PL tentativePlan));
            SOME (SLc (PSG initiateMovement));
            SOME (SLc (PL report1))])::outs)) \iff
authenticationTest inputOK
  [Name PlatoonLeader says prop (SOME (SLc (PL recon)));
   Name PlatoonLeader says
   prop (SOME (SLc (PL tentativePlan)));
   Name PlatoonSergeant says
   prop (SOME (SLc (PSG initiateMovement)));
   Name PlatoonLeader says
   prop (SOME (SLc (PL report1)))] \cap \)
CFGInterpret (M, Oi, Os)
  (CFG inputOK secContext secContextNull
     ([Name PlatoonLeader says
       prop (SOME (SLc (PL recon)));
       Name PlatoonLeader says
       prop (SOME (SLc (PL tentativePlan)));
       Name PlatoonSergeant says
       prop (SOME (SLc (PSG initiateMovement)));
       Name PlatoonLeader says
       prop (SOME (SLc (PL report1)))]::ins) WARNO outs) \land
(M,Oi,Os) satList
[prop (SOME (SLc (PL recon)));
 prop (SOME (SLc (PL tentativePlan)));
 prop (SOME (SLc (PSG initiateMovement)));
 prop (SOME (SLc (PL report1)))]
```

```
PlatoonLeader_WARNO_exec_report1_lemma
 \vdash \forall M \ Oi \ Os.
     CFGInterpret (M, Oi, Os)
       (CFG inputOK secContext secContextNull
           ([Name PlatoonLeader says
            prop (SOME (SLc (PL recon)));
            Name PlatoonLeader says
            prop (SOME (SLc (PL tentativePlan)));
            Name PlatoonSergeant says
            prop (SOME (SLc (PSG initiateMovement)));
            Name PlatoonLeader says
            prop (SOME (SLc (PL report1)))]::ins) WARNO outs) \Rightarrow
     (M,Oi,Os) satList
     propCommandList
        [Name PlatoonLeader says prop (SOME (SLc (PL recon)));
        Name PlatoonLeader says
        prop (SOME (SLc (PL tentativePlan)));
        Name PlatoonSergeant says
        prop (SOME (SLc (PSG initiateMovement)));
        Name PlatoonLeader says prop (SOME (SLc (PL report1)))]
[PlatoonSergeant_trap_plCommand_justified_lemma]
 \vdash \forall NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os)
       (trap
           (inputList
              [Name PlatoonSergeant says
               prop (SOME (SLc (PL plCommand)))]))
       (CFG inputOK secContext secContextNull
           ([Name PlatoonSergeant says
            prop (SOME (SLc (PL plCommand)))]::ins) s outs)
       (CFG inputOK secContext secContextNull ins
           (NS s
              (trap
                 (inputList
                    [Name PlatoonSergeant says
                     prop (SOME (SLc (PL plCommand)))])))
           (Out s
              (trap
                 (inputList
                    [Name PlatoonSergeant says
                     prop (SOME (SLc (PL plCommand)))]))::
                outs)) \iff
     authenticationTest inputOK
        [Name PlatoonSergeant says
        prop (SOME (SLc (PL plCommand)))] \land
     CFGInterpret (M, Oi, Os)
       (CFG inputOK secContext secContextNull
           ([Name PlatoonSergeant says
```

```
prop (SOME (SLc (PL plCommand)))]::ins) s outs) \land
      (M,Oi,Os) sat prop NONE
[PlatoonSergeant_trap_plCommand_justified_thm]
 \vdash \ \forall NS \ Out \ M \ Oi \ Os.
     TR (M, Oi, Os) (trap [SOME (SLc (PL plCommand))])
        (CFG inputOK secContext secContextNull
           ([Name PlatoonSergeant says
             prop (SOME (SLc (PL plCommand)))]::ins) s outs)
        (CFG inputOK secContext secContextNull ins
           (NS \ s \ (trap \ [SOME \ (SLc \ (PL \ plCommand))]))
           (Out \ s \ (trap \ [SOME \ (SLc \ (PL \ plCommand))])::outs)) \iff
     authenticationTest inputOK
        [Name PlatoonSergeant says
         prop (SOME (SLc (PL plCommand)))] \land
     CFGInterpret (M, Oi, Os)
        (CFG inputOK secContext secContextNull
           ([Name PlatoonSergeant says
             prop (SOME (SLc (PL plCommand)))]::ins) s outs) \land
      (M,Oi,Os) sat prop NONE
[PlatoonSergeant_trap_plCommand_lemma]
 \vdash \forall M \ Oi \ Os.
     CFGInterpret (M, Oi, Os)
        (CFG inputOK secContext secContextNull
           ([Name PlatoonSergeant says
             prop (SOME (SLc (PL plCommand)))]::ins) s outs) \Rightarrow
      (M,Oi,Os) sat prop NONE
```

## 17 PlanPBType Theory

**Built:** 13 May 2018

Parent Theories: indexedLists, patternMatches

#### 17.1 Datatypes

### 17.2 Theorems

```
[plCommand_distinct_clauses]
 \vdash receiveMission \neq warno \land receiveMission \neq tentativePlan \land
    receiveMission \neq recon \land receiveMission \neq report1 \land
    {\tt receiveMission} \neq {\tt completePlan} \ \land \ {\tt receiveMission} \neq {\tt opoid} \ \land
    \verb|receiveMission| \neq \verb|supervise| \land \verb|receiveMission| \neq \verb|report2| \land
    receiveMission \neq complete \land receiveMission \neq plIncomplete \land
    receiveMission \neq invalidPlCommand \land warno \neq tentativePlan \land
    	ext{warno} 
eq 	ext{recon} \land 	ext{warno} 
eq 	ext{report1} \land 	ext{warno} 
eq 	ext{completePlan} \land
    warno \neq opoid \wedge warno \neq supervise \wedge warno \neq report2 \wedge
    warno \neq complete \wedge warno \neq plIncomplete \wedge
    warno \neq invalidPlCommand \wedge tentativePlan \neq recon \wedge
    tentativePlan \neq report1 \land tentativePlan \neq completePlan \land
    tentativePlan \neq opoid \land tentativePlan \neq supervise \land
    tentativePlan \neq report2 \land tentativePlan \neq complete \land
    \texttt{tentativePlan} \neq \texttt{plIncomplete} \ \land
    \texttt{tentativePlan} \neq \texttt{invalidPlCommand} \ \land \ \texttt{recon} \neq \texttt{report1} \ \land
    recon \neq completePlan \land recon \neq opoid \land recon \neq supervise \land
    recon \neq report2 \land recon \neq complete \land recon \neq plIncomplete \land
    recon \neq invalidPlCommand \land report1 \neq completePlan \land
    report1 \neq opoid \land report1 \neq supervise \land report1 \neq report2 \land
    \texttt{report1} \neq \texttt{complete} \ \land \ \texttt{report1} \neq \texttt{plIncomplete} \ \land \\
    report1 \neq invalidPlCommand \wedge completePlan \neq opoid \wedge
    {\tt completePlan} \, \neq \, {\tt supervise} \, \wedge \, {\tt completePlan} \, \neq \, {\tt report2} \, \wedge \,
    {\tt completePlan} \neq {\tt complete} \ \land \ {\tt completePlan} \neq {\tt plIncomplete} \ \land
    completePlan \neq invalidPlCommand \land opoid \neq supervise \land
    opoid \neq report2 \wedge opoid \neq complete \wedge opoid \neq plIncomplete \wedge
    opoid \neq invalidPlCommand \land supervise \neq report2 \land
    supervise \neq complete \land supervise \neq plIncomplete \land
    supervise \neq invalidPlCommand \land report2 \neq complete \land
    report2 \neq plIncomplete \land report2 \neq invalidPlCommand \land
    complete \neq plIncomplete \land complete \neq invalidPlCommand \land
    plIncomplete \neq invalidPlCommand
[psgCommand_distinct_clauses]
 \vdash initiateMovement \neq psgIncomplete \land
    initiateMovement \neq invalidPsgCommand \land
    psgIncomplete \neq invalidPsgCommand
[slCommand_distinct_clauses]
```

 $\vdash \forall a' \ a. \ PL \ a \neq PSG \ a'$ 

```
[slCommand_one_one]
  \vdash (\forall a \ a'. (PL a = PL \ a') \iff (a = a')) \land
     \forall a \ a'. (PSG a = PSG \ a') \iff (a = a')
[slOutput_distinct_clauses]
  \vdash PlanPB \neq ReceiveMission \land PlanPB \neq Warno \land
     PlanPB \neq TentativePlan \land PlanPB \neq InitiateMovement \land
     {\tt PlanPB} \neq {\tt Recon} \ \land \ {\tt PlanPB} \neq {\tt Report1} \ \land \ {\tt PlanPB} \neq {\tt CompletePlan} \ \land \\
     {\tt PlanPB} \neq {\tt Opoid} \ \land \ {\tt PlanPB} \neq {\tt Supervise} \ \land \ {\tt PlanPB} \neq {\tt Report2} \ \land
     PlanPB \neq Complete \land PlanPB \neq unAuthenticated \land
     PlanPB \neq unAuthorized \land ReceiveMission \neq Warno \land
     {\tt Receive Mission} \neq {\tt Tentative Plan} \ \land
     \texttt{ReceiveMission} \neq \texttt{InitiateMovement} \ \land \ \texttt{ReceiveMission} \neq \texttt{Recon} \ \land
     ReceiveMission \neq Report1 \wedge ReceiveMission \neq CompletePlan \wedge
     \texttt{ReceiveMission} \neq \texttt{Opoid} \ \land \ \texttt{ReceiveMission} \neq \texttt{Supervise} \ \land
     \texttt{ReceiveMission} \neq \texttt{Report2} \ \land \ \texttt{ReceiveMission} \neq \texttt{Complete} \ \land
     {\tt ReceiveMission} \, \neq \, {\tt unAuthenticated} \, \, \wedge \,
     ReceiveMission \neq unAuthorized \land Warno \neq TentativePlan \land
     	exttt{Warno} 
eq 	exttt{InitiateMovement} 
\wedge 	exttt{Warno} 
eq 	exttt{Recon} 
\wedge 	exttt{Warno} 
eq 	exttt{Report1} 
\wedge
     \texttt{Warno} \neq \texttt{Report2} \ \land \ \texttt{Warno} \neq \texttt{Complete} \ \land
     Warno \neq unAuthenticated \wedge Warno \neq unAuthorized \wedge
     \texttt{TentativePlan} \neq \texttt{InitiateMovement} \ \land \ \texttt{TentativePlan} \neq \texttt{Recon} \ \land \\
     {\tt TentativePlan} \, \neq \, {\tt Report1} \, \wedge \, {\tt TentativePlan} \, \neq \, {\tt CompletePlan} \, \wedge \,
     \texttt{TentativePlan} \neq \texttt{Opoid} \ \land \ \texttt{TentativePlan} \neq \texttt{Supervise} \ \land
     TentativePlan \neq Report2 \wedge TentativePlan \neq Complete \wedge
     TentativePlan \neq unAuthenticated \wedge
     \texttt{TentativePlan} \neq \texttt{unAuthorized} \ \land \ \texttt{InitiateMovement} \neq \texttt{Recon} \ \land \\
     {\tt InitiateMovement} \, \neq \, {\tt Report1} \, \, \wedge \,
     InitiateMovement \neq CompletePlan \land InitiateMovement \neq Opoid \land
     {\tt InitiateMovement} \neq {\tt Supervise} \ \land \ {\tt InitiateMovement} \neq {\tt Report2} \ \land \\
     {\tt InitiateMovement} \, \neq \, {\tt Complete} \, \, \wedge \,
     InitiateMovement \neq unAuthenticated \land
     InitiateMovement \neq unAuthorized \wedge Recon \neq Report1 \wedge
     \texttt{Recon} \neq \texttt{CompletePlan} \ \land \ \texttt{Recon} \neq \texttt{Opoid} \ \land \ \texttt{Recon} \neq \texttt{Supervise} \ \land
     \texttt{Recon} \neq \texttt{Report2} \ \land \ \texttt{Recon} \neq \texttt{Complete} \ \land
     Recon \neq unAuthenticated \land Recon \neq unAuthorized \land
     \texttt{Report1} \neq \texttt{CompletePlan} \ \land \ \texttt{Report1} \neq \texttt{Opoid} \ \land
     \texttt{Report1} \neq \texttt{Supervise} \ \land \ \texttt{Report1} \neq \texttt{Report2} \ \land
     \texttt{Report1} \neq \texttt{Complete} \ \land \ \texttt{Report1} \neq \texttt{unAuthenticated} \ \land
     Report1 \neq unAuthorized \land CompletePlan \neq Opoid \land
     {\tt CompletePlan} \neq {\tt Supervise} \ \land \ {\tt CompletePlan} \neq {\tt Report2} \ \land
     {\tt CompletePlan} \neq {\tt Complete} \ \land \ {\tt CompletePlan} \neq {\tt unAuthenticated} \ \land \\
     {\tt CompletePlan} \neq {\tt unAuthorized} \ \land \ {\tt Opoid} \neq {\tt Supervise} \ \land \\
     Opoid \neq Report2 \wedge Opoid \neq Complete \wedge
     Opoid \neq unAuthenticated \wedge Opoid \neq unAuthorized \wedge
     Supervise \neq Report2 \wedge Supervise \neq Complete \wedge
     Supervise \neq unAuthenticated \wedge Supervise \neq unAuthorized \wedge
     Report2 \neq Complete \wedge Report2 \neq unAuthenticated \wedge
```

```
Report2 \neq unAuthorized \land Complete \neq unAuthenticated \land
        Complete \neq unAuthorized \wedge unAuthenticated \neq unAuthorized
[slRole_distinct_clauses]
  ⊢ PlatoonLeader ≠ PlatoonSergeant
[slState_distinct_clauses]
  \vdash PLAN_PB \neq RECEIVE_MISSION \land PLAN_PB \neq WARNO \land
        PLAN_PB \neq TENTATIVE_PLAN \land PLAN_PB \neq INITIATE_MOVEMENT \land
        {\tt PLAN\_PB} \, \neq \, {\tt RECON} \, \wedge \, {\tt PLAN\_PB} \, \neq \, {\tt REPORT1} \, \wedge \,
        {\tt PLAN\_PB} \, \neq \, {\tt COMPLETE\_PLAN} \, \wedge \, {\tt PLAN\_PB} \, \neq \, {\tt OPOID} \, \wedge \,
        PLAN_PB \neq SUPERVISE \wedge PLAN_PB \neq REPORT2 \wedge
        PLAN_PB \neq COMPLETE \wedge RECEIVE_MISSION \neq WARNO \wedge
        RECEIVE_MISSION \neq TENTATIVE_PLAN \wedge
        RECEIVE_MISSION \neq INITIATE_MOVEMENT \wedge
        \mathtt{RECEIVE\_MISSION} \neq \mathtt{RECON} \land \mathtt{RECEIVe\_MISSION} \neq \mathtt{REPORT1} \land
        RECEIVE_MISSION ≠ COMPLETE_PLAN ∧ RECEIVE_MISSION ≠ OPOID ∧
        RECEIVE_MISSION \neq SUPERVISE \wedge RECEIVE_MISSION \neq REPORT2 \wedge
        RECEIVE_MISSION \neq COMPLETE \wedge WARNO \neq TENTATIVE_PLAN \wedge
        	extsf{warno} 
eq 	extsf{initiate_movement} 
warno 
eq 	extsf{recon} 
warno 
eq 	extsf{recon} 
otag 	extsf{warno} 
eq 	extsf{recon} 
otag 	extsf
        WARNO \neq COMPLETE_PLAN \wedge WARNO \neq OPOID \wedge WARNO \neq SUPERVISE \wedge
        WARNO \neq REPORT2 \land WARNO \neq COMPLETE \land
        TENTATIVE_PLAN \neq INITIATE_MOVEMENT \wedge TENTATIVE_PLAN \neq RECON \wedge
        TENTATIVE_PLAN \neq REPORT1 \wedge TENTATIVE_PLAN \neq COMPLETE_PLAN \wedge
        TENTATIVE_PLAN \neq OPOID \wedge TENTATIVE_PLAN \neq SUPERVISE \wedge
        TENTATIVE_PLAN \neq REPORT2 \wedge TENTATIVE_PLAN \neq COMPLETE \wedge
        INITIATE_MOVEMENT \neq RECON \wedge INITIATE_MOVEMENT \neq REPORT1 \wedge
        \texttt{INITIATE\_MOVEMENT} \neq \texttt{COMPLETE\_PLAN} \ \land
        {\tt INITIATE\_MOVEMENT} \neq {\tt OPOID} \ \land \ {\tt INITIATE\_MOVEMENT} \neq {\tt SUPERVISE} \ \land
        {\tt INITIATE\_MOVEMENT} \neq {\tt REPORT2} \ \land \ {\tt INITIATE\_MOVEMENT} \neq {\tt COMPLETE} \ \land \\
        \mathtt{RECON} \neq \mathtt{REPORT1} \ \land \ \mathtt{RECON} \neq \mathtt{COMPLETE\_PLAN} \ \land \ \mathtt{RECON} \neq \mathtt{OPOID} \ \land
        {\tt RECON} \neq {\tt SUPERVISE} \land {\tt RECON} \neq {\tt REPORT2} \land {\tt RECON} \neq {\tt COMPLETE} \land
        REPORT1 \neq COMPLETE_PLAN \wedge REPORT1 \neq OPOID \wedge
        \texttt{REPORT1} \neq \texttt{SUPERVISE} \ \land \ \texttt{REPORT1} \neq \texttt{REPORT2} \ \land
        REPORT1 \neq COMPLETE \wedge COMPLETE_PLAN \neq OPOID \wedge
        COMPLETE_PLAN \neq SUPERVISE \wedge COMPLETE_PLAN \neq REPORT2 \wedge
        \mathtt{COMPLETE\_PLAN} \neq \mathtt{COMPLETE} \land \mathtt{OPOID} \neq \mathtt{SUPERVISE} \land
        OPOID \neq REPORT2 \wedge OPOID \neq COMPLETE \wedge SUPERVISE \neq REPORT2 \wedge
        	ext{SUPERVISE} 
eq 	ext{COMPLETE} \wedge 	ext{REPORT2} 
eq 	ext{COMPLETE}
```

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# Appendix C

# Secure State Machine Theories: HOL Script Files

#### C.1 ssm

```
(* Secure State Machine Theory: authentication, authorization, and state
(* interpretation.
(* Author: Shiu-Kai Chin
(* Date: 27 November 2015
structure ssmScript = struct
(* = = = Interactive \ mode = = = =
app load ["TypeBase", "ssminfRules", "listTheory", "optionTheory", "acl_infRules", "satListTheory", "ssmTheory"];
app\ load\ ["TypeBase",\ "ssminfRules"," list Theory"," option Theory"," acl\_infRules",
         "satListTheory"];
open TypeBase listTheory ssminfRules optionTheory acl_infRules satListTheory
    ssmTheory
==== end interactive mode ===== *)
open HolKernel boolLib Parse bossLib
open TypeBase listTheory optionTheory ssminfRules acl_infRules satListTheory
(********************************
(* create a new theory *)
(*****************************
val _ = new_theory "ssm";
(*\ Define\ the\ type\ of\ transition:\ discard\ ,\ execute\ ,\ or\ trap\ . We discard from
(* the input stream those inputs that are not of the form P says command. We
(* execute commands that users and supervisors are authorized for. We trap
(* commands that users are not authorized to execute.
(* In keeping with virtual machine design principles as described by Popek
(* and Goldberg, we add a TRAP instruction to the commands by users.
                                                                           *)
(* In effect, we are LIFTING the commands available to users to include the
(* TRAP instruction used by the state machine to handle authorization errors.
```

```
val =
Datatype
trTvpe =
  discard 'cmdlist | trap 'cmdlist | exec 'cmdlist '
val trType_distinct_clauses = distinct_of '': 'cmdlist trType''
val _ = save_thm("trType_distinct_clauses", trType_distinct_clauses)
val trType_one_one = one_one_of '': 'cmdlist trType''
val _ = save_thm("trType_one_one", trType_one_one)
(* Define configuration to include the security context within which the
(*\ inputs\ are\ evaluated . The components are as follows: (1) the authentication*)
(* function, (2) the interretation of the state, (3) the security context,
(* (4) the input stream, (5) the state, and (6) the output stream.
(* -
val _ =
Datatype
'configuration =
CFG
  (('command option, 'principal, 'd, 'e)Form -> bool)
  (('state -> ('command option, 'principal, 'd, 'e)Form list ->
('command option, 'principal, 'd, 'e)Form list))
  ((('command option, 'principal, 'd, 'e)Form list) ->
((('command option, 'principal, 'd, 'e)Form list))
(((('command option, 'principal, 'd, 'e)Form list))
(((('command option, 'principal, 'd, 'e)Form) list) list)
    state)
  ('output list)'
(* Prove one-to-one properties of configuration
val configuration_one_one =
    one_one_of '': ('command option, 'd, 'e, 'output, 'principal, 'state) configuration ''
val _ = save_thm("configuration_one_one", configuration_one_one)
(* The interpretation of configuration is the conjunction of the formulas in
                                                                                           *)
(* the context and the first element of a non-empty input stream.
val CFGInterpret_def =
Define
`CFGInterpret
 ((M: ('command option, 'b, 'principal, 'd, 'e) Kripke), Oi: 'd po, Os: 'e po)
 (CFG
  (elementTest:('command option,'principal,'d,'e)Form -> bool)
  (stateInterp: state -> (('command option, 'principal, 'd, 'e)Form list) -> (('command option, 'principal, 'd, 'e)Form list))
  (context:(('command option, 'principal, 'd, 'e)Form list) ->
  (('command option, 'principal, 'd, 'e)Form list))
((x:('command option, 'principal, 'd, 'e)Form list)::ins)
  (state: 'state)
  (outStream: 'output list))
  ((M,Oi,Os) satList (context x)) /\
  ((M,Oi,Os) satList x) /\
  ((M,Oi,Os) satList (stateInterp state x))'
(* In the following definitions of authenticationTest, extractCommand, and
(* commandList, we implicitly assume that the only authenticated inputs are
                                                                                           * )
(* of the form P says phi, i.e., we know who is making statement phi.
val authenticationTest_def =
Define
`authentication Test\\
 (elementTest:('command option,'principal,'d,'e)Form -> bool)
(x:('command option,'principal,'d,'e)Form list) =
FOLDR (\p q.p /\ q) T (MAP elementTest x)';
```

```
val extractCommand_def =
Define
'extractCommand (P says (prop (SOME cmd)): ('command option, 'principal, 'd, 'e)Form) =
   cmd ';
val commandList_def =
Define
'commandList (x:('command option, 'principal, 'd, 'e)Form list) =
MAP extractCommand x';
val extractPropCommand_def =
Define
'(extractPropCommand (P says (prop (SOME cmd)):('command option,'principal,'d,'e)Form) =
   ((prop (SOME cmd)): ('command option, 'principal, 'd, 'e)Form)) ';
\mathbf{val} \ \mathbf{propCommandList\_def} =
'propCommandList (x:('command option,'principal,'d,'e)Form list) =
MAP extractPropCommand x';
val extractInput_def =
Define
'extractInput (P says (prop x):('command option,'principal,'d,'e)Form) = x';
val inputList_def =
Define
'inputList (xs:('command option,'principal,'d,'e)Form list) =
MAP extractInput xs';
(* Define transition relation among configurations. This definition is
(* parameterized in terms of next-state transition function and output
(* function.
(*
val (TR_rules, TR_ind, TR_cases) =
Hol_reln
'(!(elementTest:('command option,'principal,'d,'e)Form -> bool)
    (NS: 'state -> ('command option list) trType -> 'state) M Oi Os Out (s:'state)
   (context:(('command option, 'principal, 'd, 'e)Form list) -> (('command option, 'principal, 'd, 'e)Form list))
    (stateInterp:'state -> ('command option, 'principal, 'd, 'e)Form list ->
    ('command option, 'principal, 'd, 'e)Form list)
(x:('command option, 'principal, 'd, 'e)Form list)
    (ins:('command option,'principal,'d,'e)Form list list)
    (outs: 'output list).
 (authenticationTest elementTest x) /\
 (CFGInterpret (M, Oi, Os)
  (CFG elementTest stateInterp context (x::ins) s outs)) =>>
 (TR
  ((M:('command option,'b,'principal,'d,'e)Kripke),Oi:'d po,Os:'e po)
  (exec (inputList x))
  (CFG elementTest stateInterp context (x::ins) s outs)
  (CFG elementTest stateInterp context ins
        (NS s (exec (inputList x)))
 ((Out s (exec (inputList x)))::outs)))) /\
(!(elementTest:('command option, 'principal, 'd, 'e)Form -> bool)
(NS: 'state -> ('command option list) trType -> 'state) M Oi Os Out (s:'state)
    (context:(('command option, 'principal, 'd, 'e)Form list) ->
   (('command option, 'principal, 'd, 'e)Form list))
    (stateInterp:'state -> ('command option, 'principal, 'd, 'e)Form list ->
    ('command option, 'principal, 'd, 'e)Form list)
(x:('command option, 'principal, 'd, 'e)Form list)
    (ins:('command option,'principal,'d,'e)Form list list)
    (outs: 'output list).
 (authenticationTest elementTest x)
 (CFGInterpret (M, Oi, Os)
  (CFG elementTest stateInterp context (x::ins) s outs)) ==>
 (TR.
  ((M: ('command option, 'b, 'principal, 'd, 'e) Kripke), Oi: 'd po, Os: 'e po)
  (trap (inputList x))
 (CFG elementTest stateInterp context (x::ins) s outs)
 (CFG elementTest stateInterp context ins
```

```
(NS s (trap (inputList x)))
 ((Out s (trap (inputList x)))::outs)))) /\
(!(elementTest:('command option,'principal,'d,'e)Form -> bool)
     (NS: 'state \rightarrow ('command option list') trType \rightarrow 'state) M Oi Os Out (s:'state)
   (context:(('command option, 'principal, 'd, 'e)Form list) ->
(('command option, 'principal, 'd, 'e)Form list))
     (stateInterp:'state -> ('command option, 'principal, 'd, 'e)Form list ->
    ('command option, 'principal, 'd, 'e)Form list)
(x:('command option, 'principal, 'd, 'e)Form list)
(ins:('command option, 'principal, 'd, 'e)Form list list)
     (outs: 'output list).
 ~(authenticationTest elementTest x) ==>
 (TR
  ((M: ('command option, 'b, 'principal, 'd, 'e) Kripke), Oi: 'd po, Os: 'e po)
  (discard (inputList x))
 (CFG elementTest stateInterp context (x::ins) s outs)
 (CFG elementTest stateInterp context ins
       (NS s (discard (inputList x)))
       ((Out s (discard (inputList x)))::outs))))'
(* Split up TR_rules into individual clauses
(* -
val [rule0 , rule1 , rule2] = CONJUNCTS TR_rules
(* Prove the converse of rule0, rule1, and rule2
val TR_lemma0 =
TAC_PROOF(([], flip_TR_rules rule0),
DISCH_TAC THEN
IMP_RES_TAC TR_cases THEN
PAT_ASSUM
 ``exec\ cmd\ =\ y``
 (fn th => ASSUME_TAC(REWRITE_RULE[trType_one_one, trType_distinct_clauses]th)) THEN
PROVE.TAC [\ configuration\_one\_one\ , list\_11\ , trType\_distinct\_clauses\ ])
val TR_lemma1 =
TAC_PROOF(([], flip_TR_rules rule1),
DISCH_TAC THEN
IMP_RES_TAC TR_cases THEN
PAT_ASSUM
  "trap cmd = y"
 (\ fn\ th\ \Longrightarrow ASSUME\_TAC(REWRITE\_RULE[\ trType\_one\_one\ , trType\_distinct\_clauses\ ]\ th\ ))\ THEN
 PROVE_TAC[configuration_one_one, list_11, trType_distinct_clauses])
val TR_lemma2 =
\begin{array}{l} {\rm TACPROOF(([],flip\_TR\_rules\ rule2)}\,,\\ {\rm DISCH\_TAC\ THEN} \end{array}
IMP_RES_TAC TR_cases THEN
PAT_ASSUM
 "discard (inputList x)= y"
 (fn th => ASSUME_TAC(REWRITE_RULE[trType_one_one, trType_distinct_clauses]th)) THEN
PROVE_TAC[configuration_one_one , list_11 , trType_distinct_clauses])
val TR rules converse =
TAC_PROOF(([], flip_TR_rules TR_rules),
REWRITE_TAC[TR_lemma0, TR_lemma1, TR_lemma2])
val TR_EQ_rules_thm = TR_EQ_rules TR_rules_TR_rules_converse
val _ = save_thm("TR_EQ_rules_thm",TR_EQ_rules_thm)
val [TRrule0, TRrule1, TR_discard_cmd_rule] = CONJUNCTS TR_EQ_rules_thm
val _ = save_thm("TRrule0",TRrule0)
val _ = save_thm("TRrule1", TRrule1)
val _ = save_thm("TR_discard_cmd_rule", TR_discard_cmd_rule)
(* If (CFGInterpret
```

```
(*
(*
       (M, Oi, Os)
       (CFG\ elementTest\ stateInterpret\ certList
             ((P \ says \ (prop \ (CMD \ cmd)))::ins) \ s \ outs) ==>
(*
      ((M, Oi, Os) \ sat \ (prop \ (CMD \ cmd))))
(*
(* is a valid inference rule, then executing cmd the exec (CMD cmd) transition
(* occurs if and only if prop (CMD cmd), elementTest, and
(* CFGInterpret (M, Oi, Os)
(* (CFG elementTest stateInterpret certList (P says prop (CMD cmd)::ins) s outs)
(* are true.
                                                                                  * )
(*
val TR_exec_cmd_rule =
TAC_PROOF(([],
''!elementTest context stateInterp (x:('command option, 'principal, 'd, 'e)Form list)
   ins s outs.
 (!M Oi Os.
 (CFGInterpret
  ((M:('command option, 'b, 'principal, 'd, 'e) Kripke),(Oi:'d po), (Os:'e po))
  (CFG element Test
       (stateInterp:'state -> ('command option, 'principal, 'd, 'e)Form list ->
        ('command option, 'principal, 'd, 'e) Form list) context
       (x::ins)
       (s:'state) (outs:'output list))) ==>
  (M, Oi, Os) satList (propCommandList (x:('command option, 'principal, 'd, 'e)Form list))) =>
(!NS Out M Oi Os.
  ((M:('command option, 'b, 'principal, 'd, 'e) Kripke),(Oi:'d po),
   (Os:'e po)) (exec (inputList x))
  ('command option, 'principal, 'd, 'e)Form list)
(context : ('command option, 'principal, 'd, 'e) Form list ->
            ('command option, 'principal, 'd, 'e) Form list)
            (x::ins)
           (s:'state) (outs:'output list))
   (CFG elementTest stateInterp context ins
            ((NS: 'state -> 'command option list trType -> 'state) s (exec (inputList x)))
            (Out s (exec (inputList x))::outs)) <=>
  (authenticationTest elementTest x) /\
  (CFGInterpret (M, Oi, Os)
   (CFG elementTest stateInterp context (x::ins) s outs)) /\
(M,Oi,Os) satList (propCommandList x))''), REWRITE_TAC[TRrule0] THEN
REPEAT STRIP_TAC THEN
EQ_TAC THEN
REPEAT STRIP_TAC THEN
PROVE_TAC[])
val _ = save_thm("TR_exec_cmd_rule", TR_exec_cmd_rule)
(* If (CFGInterpret
(*
       (M, Oi, Os)
(*
       (CFG elementTest stateInterpret certList
(*
            ((P \ says \ (prop \ (CMD \ cmd)))::ins) \ s \ outs) \Longrightarrow
      ((M, Oi, Os) \ sat \ (prop \ TRAP)))
(*
(* is a valid inference rule, then executing cmd the trap(CMD cmd) transition
(* occurs if and only if prop TRAP, elementTest, and
(* CFGInterpret (M, Oi, Os)
    (CFG elementTest stateInterpret certList (P says prop (CMD cmd)::ins)
(*
         s outs) are true.
(* -
val TR_trap_cmd_rule =
TAC_PROOF(
([], ''!elementTest context stateInterp (x:('command option, 'principal, 'd, 'e)Form list)
   ins s outs.
 (!M Oi Os.
 (CFGInterpret
  ((M:('command option, 'b, 'principal, 'd, 'e) Kripke),(Oi:'d po), (Os:'e po))
  (CFG element Test
       (stateInterp:'state -> ('command option, 'principal, 'd, 'e)Form list ->
        ('command option, 'principal, 'd, 'e)Form list) context
       (x :: ins)
       (s:'state) (outs:'output list))) ==>
```

```
(M,Oi,Os) sat (prop NONE)) ==>
(!NS Out M Oi Os.
 TR.
  ((M:('command option, 'b, 'principal, 'd, 'e) Kripke),(Oi:'d po),
    (Os : 'e po)) (trap (inputList x))
  (CFG (elementTest : ('command option, 'principal, 'd, 'e) Form -> bool)

(stateInterp:'state -> ('command option, 'principal, 'd, 'e) Form list ->
             ('command option, 'principal, 'd, 'e)Form list)
(context : ('command option, 'principal, 'd, 'e) Form list ->
              ('command option, 'principal, 'd, 'e) Form list)
             (x::ins)
             (s:'state) (outs:'output list))
   (CFG elementTest stateInterp context ins
             ((NS: 'state -> 'command option list trType -> 'state) s (trap (inputList x)))
             (Out s (trap (inputList x))::outs)) <=>
  (authenticationTest elementTest x) /\
  (CFGInterpret (M, Oi, Os)
    (CFG elementTest stateInterp context (x::ins) s outs)) / 
    (M, Oi, Os) sat (prop NONE)) ''),
REWRITE_TAC[TRrule1] THEN
REPEAT STRIP_TAC THEN EQ_TAC THEN
REPEAT STRIP_TAC THEN
PROVE_TAC[])
val _ = save_thm("TR_trap_cmd_rule",TR_trap_cmd_rule)
(* = = start here =
==== end here ==== *)
val _ = export_theory ();
val _ = print_theory "-";
end (* structure *)
```

#### C.2 satList

```
(*\ Definition\ of\ satList\ for\ conjunctions\ of\ ACL\ formulas
                                                                           *)
(* Author: Shiu-Kai Chin
(* Date: 24 July 2014
structure satListScript = struct
(* interactive mode
 app load
 ["TypeBase", "listTheory", "acl_infRules"];
open HolKernel boolLib Parse bossLib
open TypeBase acl_infRules listTheory
(*******
* create a new theory
********
val _ = new_theory "satList";
(*********************************
(* Configurations and policies are represented by lists
(* of formulas in the access-control logic.
(*\ Previously\ ,\ for\ a\ formula\ f\ in\ the\ access-control\ logic\ ,
(* we ultimately interpreted it within the context of a
(* Kripke structure M and partial orders Oi: 'Int po and
(* Os: 'Sec po. This is represented as (M, Oi, Os) sat f.
(* The natural extension is to interpret a list of formulas *)
(* [f0; ...; fn] as a conjunction:
(* (M, Oi, Os) \ sat \ fo \land \ldots \land (M, Oi, Os) \ sat \ fn
val _ set_fixity "satList" (Infixr 540);
```

```
val satList_def =
Define
'((M:('prop,'world,'pName,'Int,'Sec)Kripke),(Oi:'Int po),(Os:'Sec po))
 satList
 formList =
 FOLDR
 (\x y. x /\ y) T
 (MAP
  (\ (f:('prop,'pName,'Int,'Sec)Form).
((M:('prop,'world,'pName,'Int,'Sec)Kripke),
    Oi: 'Int po, Os: 'Sec po) sat f) formList) ';
(* Properties of satList *)
(*********************************
val satList_nil =
TAC_PROOF(
([], ''((M:('prop,'world,'pName,'Int,'Sec)Kripke),(Oi:'Int po),(Os:'Sec po)) satList []''),
REWRITE_TAC[satList_def,FOLDR,MAP])
val _ = save_thm("satList_nil", satList_nil)
val satList_conj =
TAC_PROOF(
([]]
'!ll l2 M Oi Os.(((M:('prop, 'world, 'pName, 'Int, 'Sec)Kripke),(Oi:'Int po),(Os:'Sec po))
   satList l1) /\
  (((M:('prop,'world,'pName,'Int,'Sec)Kripke),(Oi:'Int po),(Os:'Sec po))
   satList 12) =
  (((M:('prop, 'world, 'pName, 'Int, 'Sec)Kripke),(Oi:'Int po),(Os:'Sec po))
   satList (l1 ++ l2))''),
Induct THEN
REWRITE_TAC[APPEND, satList_nil] THEN
REWRITE_TAC[satList_def,MAP] THEN
CONV_TAC(DEPTH_CONV BETA_CONV) THEN
REWRITE_TAC[FOLDR] THEN
CONV_TAC(DEPTH_CONV BETA_CONV) THEN
REWRITE_TAC [GSYM satList_def] THEN
PROVE_TAC[])
val _ = save_thm("satList_conj", satList_conj)
val satList_CONS =
TACPROOF(([],
''!h t M Oi Os.(((M:('prop,'world,'pName,'Int,'Sec)Kripke),(Oi:'Int po),(Os:'Sec po))
   satList (h::t)) =
  (((M,Oi,Os) sat h) /\
(((M:('prop,'world,'pName,'Int,'Sec)Kripke),(Oi:'Int po),(Os:'Sec po))
satList t))''),
REPEAT STRIP_TAC THEN
REWRITE_TAC[satList_def,MAP] THEN
CONV_TAC(DEPTH_CONV BETA_CONV) THEN
REWRITE_TAC[FOLDR] THEN
CONV.TAC(DEPTH.CONV BETA.CONV) THEN
REWRITE_TAC[])
val _ = save_thm("satList_CONS",satList_CONS)
val _ = export_theory ();
val _ = print_theory "-";
end (* structure *)
```

## Appendix D

# Secure State Machine Theories Applied to Patrol Base Operations: HOL Script Files

#### D.1 OMNILevel

```
(* OMNIScript
(* Author: Lori Pickering
(* Date: 10 May 2018
(*\ This\ file\ is\ intended\ to\ allow\ for\ integration\ among\ the\ ssms.\ The\ idea
(* is to provide an OMNI-level integrating theory, in the sense of a super-
(* conscious that knows when each ssm is complete and provides that info to
(* higher-level state machines.
structure OMNIScript = struct
(* ==== Interactive Mode ======
app\ load\ ["TypeBase"," list Theory", "option Theory",
        "OMNIType Theory",\\
"acl\_infRules", "aclDrulesTheory", "aclrulesTheory"]; open TypeBase listTheory optionTheory
    OMNIType\ Theory
    acl\_infRules aclDrulesTheory aclrulesTheory
=== End Interactive Mode ==== *)
open HolKernel Parse boolLib bossLib;
open TypeBase listTheory optionTheory
open OMNITypeTheory
\mathbf{open} \ \mathbf{acl\_infRules} \ \mathbf{aclDrulesTheory} \ \mathbf{aclrulesTheory}
val _ = new_theory "OMNI";
(* Define slCommands for OMNI.
(* ==== Area 52 ======
val =
Datatype 'stateRole = Omni'
```

```
v\,a\,l _=
Datatype \quad `omniCommand = ssmPlanPBComplete
                     ssmMoveToORPComplete\\
                      ssmConductORPComplete
                      ssmMoveToPBComplete\\
                      ssmConductPBComplete
val \ omniCommand\_distinct\_clauses = distinct\_of ``:omniCommand``
val = save_thm("omniCommand_distinct_clauses",
                omniCommand\_distinct\_clauses)
val =
Datatype 'slCommand = OMNI \ omniCommand'
val\ omniAuthentication\_def =
Define
`(omniAuthentication"
        (\textit{Name Omni says prop } (\textit{cmd:} ((\textit{slCommand command}) \ \textit{option}))
        :((slCommand\ command)\ option\ ,\ stateRole\ ,\ 'd\ ,'e)Form)\ =\ T)\ /
 (omniAuthentication = F)
val\ omniAuthorization\_def =
Define
(omniAuthorization
        (Name\ Omni\ controls\ prop\ (cmd:((slCommand\ command)\ option))
        : ((slCommand\ command)\ option\ ,\ stateRole\ ,\ 'd\ , 'e)Form)\ =\ T)\ / \ 
 (omniAuthorization = F)
This may not be necessary...But, it is interesting. Save for a later time.
(* Prove that
(* Omni \ says \ omniCommand \Longrightarrow omniCommand
set\_goal([],
     (Name Omni says prop (cmd:((slCommand command) option))
        :((slCommand command) option, stateRole, 'd, 'e)Form) ==>
                    prop (cmd:((slCommand command) option))'')
val \ th \ 1 \ = \textit{ASSUME''(Name Omni says prop (cmd:((slCommand command) option))}
       :((slCommand\ command)\ option\ ,\ stateRole\ ,\ 'd\ ,'e)Form)=TT'
val\ th2 = REWRITE\_RULE[omniAuthentication\_def]th1
===== End Area 52 ==== *)
val = export_theory();
end
```

#### D.2 TopLevel

#### D.2.1 PBTypeIntegrated Theory: Type Definitions

```
(* ==== Interactive Mode ====
app load ["TypeBase"]
open\ TypeBase
 ==== end Interactive Mode ===== *)
open HolKernel Parse boolLib bossLib;
open TypeBase OMNITypeTheory
val _= new_theory "PBTypeIntegrated";
(* Define types
val _=
{\tt Datatype \ `plCommand = crossLD \ } \ \ (* \ \textit{Move to MOVE\_TO\_ORP state *})
                      conductORP
                      moveToPB
                      {\tt conductPB}
                      completePB
                    | incomplete '
val plCommand_distinct_clauses = distinct_of '':plCommand''
val _= save_thm("plCommand_distinct_clauses",
                 plCommand_distinct_clauses)
val =
Datatype 'omniCommand = ssmPlanPBComplete
                      ssmMoveToORPComplete
                      {\tt ssmConductORPComplete}
                      ssmMoveToPBComplete\\
                      ssmConductPBComplete
                      invalid Omni Command\ ``
\mathbf{val} \ \mathtt{omniCommand\_distinct\_clauses} \ = \ \mathtt{distinct\_of} \ ``:\mathtt{omniCommand}``
val _= save_thm("omniCommand_distinct_clauses",
                 omniCommand_distinct_clauses)
val _=
Datatype 'slCommand = PL plCommand
                    | OMNI omniCommand '
val slCommand_distinct_clauses = distinct_of '':slCommand''
val _= save_thm("slCommand_distinct_clauses",
                 slCommand_distinct_clauses)
val slCommand_one_one = one_one_of ' ':slCommand' '
val _= save_thm("slCommand_one_one", slCommand_one_one)
val _=
Datatype 'stateRole = PlatoonLeader | Omni'
val stateRole_distinct_clauses = distinct_of '': stateRole''
val _= save_thm("stateRole_distinct_clauses",
                 stateRole_distinct_clauses)
val _=
{\tt Datatype 'slState = PLAN\_PB}
                   MOVE_TO_ORP
                    CONDUCT_ORP
                    MOVE_TO_PB
                    CONDUCT_PB
                   COMPLETE PB
val slState_distinct_clauses = distinct_of '':slState''
val _ = save_thm("slState_distinct_clauses", slState_distinct_clauses)
```

# D.2.2 PBIntegratedDef Theory: Authentication & Authorization Definitions

```
(*\ PBIntegratedDefTheory
(* Author: Lori Pickering
(* Date: 7 May 2018
(* Definitions for ssmPBIntegrated Theory.
structure PBIntegratedDefScript = struct
(* ==== Interactive Mode ==
           ["TypeBase", "listTheory", "optionTheory",
"uavUtilities",
"OMNITypeTheory",
app load
            "PBIntegrated Def Theory", "PBType Integrated Theory"];\\
open \ Type Base \ list Theory \ option Theory
      a clse mantics \, Theory \quad a \, cl foundation \, Theory
      u\,a\,v\,U\,t\,i\,l\,i\,t\,i\,e\,s
      OMNIType\ Theory
      PBIntegrated Def Theory \ PBType Integrated Theory
     = end Interactive Mode ==== *)
open HolKernel Parse boolLib bossLib;
open TypeBase listTheory optionTheory
open uavUtilities
open OMNITypeTheory PBTypeIntegratedTheory
val _ = new_theory "PBIntegratedDef";
(* state Interpretation function
(* This function doesn't do anything but is necessary to specialize other
val secContext_def = Define '
    secContext (x:((slCommand command)option, stateRole, 'd,'e)Form list) =
         [(TT:((slCommand command)option, stateRole, 'd, 'e)Form)]
val secHelper =
Define '
  (secHelper (cmd:omniCommand) =
      [(Name Omni) controls prop (SOME (SLc (OMNI (cmd:omniCommand)))))]) '
\mathbf{val}\ \mathrm{getOmniCommand\_def} =
Define '
  (getOmniCommand \ ([]:((slCommand \ command) \ option \ , \ stateRole \ , \ 'd \ , 'e)Form \ list)
                          = invalidOmniCommand:omniCommand) /
  (\texttt{getOmniCommand} \ (((\texttt{Name Omni}) \ \texttt{controls prop} \ (\texttt{SOME} \ (\texttt{SLc} \ (\texttt{OMNI} \ \texttt{cmd}))))) :: xs)
                          = (cmd:omniCommand)) / 
  (getOmniCommand ((x:((slCommand command)option, stateRole, 'd,'e)Form)::xs)
                          = (getOmniCommand xs)) '
```

```
val secAuthorization_def =
Define '
  (sec Authorization \ (xs: ((slCommand\ command)\ option\ ,\ stateRole\ , 'd, 'e) Form\ list\ )
                     = secHelper (getOmniCommand xs))
val secContext_def =
Define '
 (secContext (PLAN_PB) ((x:((slCommand command)option, stateRole, 'd,'e)Form)::xs) =
         [(prop (SOME (SLc (OMNI (ssmPlanPBComplete))))
           : ((slCommand\ command)\ option\ ,\ stateRole\ ,\ 'd\ ,'e)Form)\ impf
           (Name PlatoonLeader) controls prop (SOME (SLc (PL crossLD)))
           :((slCommand command)option, stateRole, 'd, 'e)Form])
 (secContext (MOVE_TO_ORP) ((x:((slCommand command) option, stateRole, 'd, 'e)Form)::xs) =
         [prop (SOME (SLc (OMNI (ssmMoveToORPComplete)))) impf
          (Name PlatoonLeader) controls prop (SOME (SLc (PL conductORP)))]) /\
  (secContext \ (CONDUCT\_ORP) \ ((x:((slCommand \ command) \ option \ , \ stateRole \ , \ 'd \ , 'e)Form)::xs) = [prop \ (SOME \ (SLc \ (OMNI \ (ssmConductORPComplete)))) \ impf 
           (Name PlatoonLeader) controls prop (SOME (SLc (PL moveToPB)))]) /
 (secContext (MOVE_TO_PB) ((x:((slCommand command)option, stateRole, ''d, 'e)Form)::xs) = [prop (SOME (SLc (OMNI (ssmMoveToPBComplete)))) impf
          (Name PlatoonLeader) controls prop (SOME (SLc (PL conductPB)))]) /\
 (secContext (CONDUCT.PB) ((x:((slCommand command)option, stateRole, 'd,'e)Form)::xs) =
    [prop (SOME (SLc (OMNI (ssmConductPBComplete)))) impf
           (Name PlatoonLeader) controls prop (SOME (SLc (PL completePB)))])'
(* ==== Area 52 =
==== End Area 52 ==== *)
val = export_theory();
```

#### D.2.3 ssmPlanPBIntegrated Theory: Theorems

```
(* ssmPBIntegratedTheory)
(* Author: Lori Pickering
(* Date: 7 May 2018
(* This theory aims to integrate the topLevel ssm with the sublevel ssms.
(* does this by adding a condition to the security context. In particular,
(* it requires that the "COMPLETE" state in the subLevel ssm must preceede
(* transition to the next state at the topLeve. I.e.,
           planPBComplete \Longrightarrow
           PlatoonLeader\ controls\ crossLD .
(* In the ssmPlanPB ssm, the last state is COMPLETE. This is reached when the *)
(* the appropriate authority says complete and the transition is made.
(* Note that following the ACL, if P says x and P controls x, then x.
(* Therefore, it is not necessary for anyone to say x at the topLevel, because*)
(* it is already proved at the lower level.
(* However, indicating that at the topLevel remains something to workout.
structure ssmPBIntegratedScript = struct
(* ==== Interactive Mode ====
                     ["TypeBase", "listTheory", "optionTheory", "listSyntax",
app load
                         'acl_infRules", "aclDrulesTheory", "aclrulesTheory",
                        "aclsemantics Theory", "aclfoundation Theory",
                       "satListTheory", "ssmTheory", "ssminfRules", "uavUtilities", "OMNITypeTheory", "PBTypeIntegratedTheory", "PBIntegratedDefTheory", "psintegratedDefTheory", "psintegratedD
                        "ssmPBIntegratedTheory"];
open \ TypeBase \ listTheory \ optionTheory \ listSyntax
            acl_infRules aclDrulesTheory aclrulesTheory
           a clse mantics Theory \ a clfoundation Theory
            satListTheory \ ssmTheory \ ssminfRules \ uavUtilities
           OMNITy pe Theory \ PBTy pe Integrated Theory \ PBIntegrated Def Theory
           ssmPBIntegratedTheory
```

```
==== end Interactive Mode ===== *)
open HolKernel Parse boolLib bossLib;
open TypeBase listTheory optionTheory
{\bf open} \ \ {\bf acl\_infRules} \ \ {\bf aclDrulesTheory} \ \ {\bf aclrulesTheory}
open satListTheory ssmTheory ssminfRules uavUtilities
open OMNITypeTheory PBTypeIntegratedTheory PBIntegratedDefTheory
val _ = new_theory "ssmPBIntegrated";
(* Define next-state and next-output functions
_
val PBNS_def =
Define '
(PBNS PLAN_PB (exec [SOME (SLc (PL crossLD))]) = MOVE_TO_ORP) /\
(PBNS MOVE_TO_ORP (exec [SOME (SLc (PL conductORP))]) = CONDUCT_ORP) /\
(PBNS CONDUCT_ORP (exec [SOME (SLc (PL conductPB))]) = MOVE_TO_PB) /\
(PBNS CONDUCT_ORP (exec [SOME (SLc (PL conductPB))]) = CONDUCT_PB) /\
(PBNS CONDUCT_PB (exec [SOME (SLc (PL completePB))]) = COMPLETE_PB) /\
(PBNS (s:slState) (trap _) = s) /\
(PBNS (s:slState) (discard _{-}) = s)'
val PBOut_def =
Define '
(PBOut PLAN_PB
                  (exec [SOME (SLc (PL crossLD))])
                                                      = MoveToORP)
(PBOut MOVE_TO_ORP (exec [SOME (SLc (PL conductORP))]) = ConductORP) /\
(PBOut CONDUCT_ORP (exec [SOME (SLc (PL moveToPB))]) = MoveToPB) /\
(PBOut MOVE_TO_PB (exec [SOME (SLc (PL conductPB))]) = ConductPB) /\
(PBOut CONDUCT_PB (exec [SOME (SLc (PL conductPB))]) = CompletePB) /\
(PBOut (s:slState) (trap _) = unAuthorized) /
(PBOut (s:slState) (discard _) = unAuthenticated)'
(* Define authentication function
val inputOK_def =
Define '
(inputOK (((Name PlatoonLeader) says prop (cmd:((slCommand command)option)))
           :((slCommand command) option, stateRole, 'd, 'e)Form) = T) /
                               says prop (cmd:((slCommand command)option)))
(inputOK (((Name Omni)
          :((slCommand command)option, stateRole, 'd, 'e)Form) = T) /\
(inputOK _ = F)
(* Prove that commands are rejected unless that are requested by a properly
(* authenticated principal.
val inputOK_cmd_reject_lemma =
Q. prove ('!cmd. ~(inputOK
                 ((prop (SOME cmd))))',
                              (PROVE_TAC[inputOK_def]))
(* ==== Just playing around with this ====
val\ inputOK\_not\_reject\_lemma =
Q. prove ( '! cmd. ~(
          (input OK \ (((Name\ Platoon Leader)\ says\ prop\ (cmd:((slCommand\ command)\ option)))
           : ((slCommand\ command)\ option\ ,\ stateRole\ ,\ 'd\ ,\ 'e)Form))\ \setminus /
          (inputOK\ (((Name\ Omni)\ says\ prop\ (cmd:((slCommand\ command)\ option)))
           :((slCommand command) option, stateRole, 'd, 'e)Form)))
==== OK, done fooling around ==== *)
val = export_theory();
```

 $\mathbf{end}$ 

#### D.3 Horizontal Slice

D.3.1	ssmPlanPB
D.3.1.1	PlanPBType Theory: Type Definitions
D.3.1.2	PlanPBDef Theory: Authentication & Authorization Definitions
D.3.1.3	ssmPlanPB Theory: Theorems
D.3.2	${\bf ssmMoveToORP}$
D.3.2.1	MoveToORPType Theory: Type Definitions
D.3.2.2	MoveToORPDef Theory: Authentication & Authorization Definitions
D.3.2.3	ssmMoveToORP Theory: Theorems
D.3.3	$\operatorname{ssmConductORP}$
D.3.3.1	ConductORPType Theory: Type Definitions
D.3.3.2	ConductORPDef Theory: Authentication & Authorization Definitions
D.3.3.3	ssmConductORP Theory: Theorems
D.3.4	ssmMoveToPB
D.3.4.1	MoveToPBType Theory: Type Definitions
D.3.4.2	MoveToPBDef Theory: Authentication & Authorization Definitions

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 $D.3.5 \quad ssmConductPB$ 

D.3.4.3 ssmMoveToPB Theory: Theorems

# Appendix E

Map of The File Folder Structure

### References

- [1] Shiu-Kai Chin and Susan Beth Older. Access Control, Security, and Trust: A Logical Approach. Chapman & Hall: CRC Cryptography and Network Security Series. Chapman and Hall/CRC, July 2010.
- [2] United States Army Ranger School, ATTN: ATSH-RB, 10850 Schneider Rd, Bldg 5024, Ft Benning, GA 31905. Ranger handbook, April 2017.