Compare and Swap > non-backing algorithms which Provides

The Pechnique used when Designing Concurrent Algorithms. Basically

(As compares an expected value to the Convite value of a

Variable, and if the Congrate value of the variable is cavale

to the expected value, Swaps the value of the variable for a

New Manable.

Single Proposed Exposed Expose:

Shorted Proposed Forked Exposed:

14 (1) locked) &

locked = true:

greture face:

CAS as Atomic Oberation:

Public Static Cass Mylock &

Private Atomic Boolean locked = new Atomic

Boolean (False):

roctor locked. Containe And Sct (face true)

of there it comfares the value of locked to false and it it is false it lets the new value of Atomic Boolean to true.

her Compare and Set () method rectured the It the value is

As is Jenneally much felter than locking, but it does defend on the degree of Contention. Because the may force of certain it the name Changes between recading and Companing, a thread can get Stock in a busy-wait if the namable in a content is that hand by many other threads.

Synchronised Convertions We consumer Contestions.

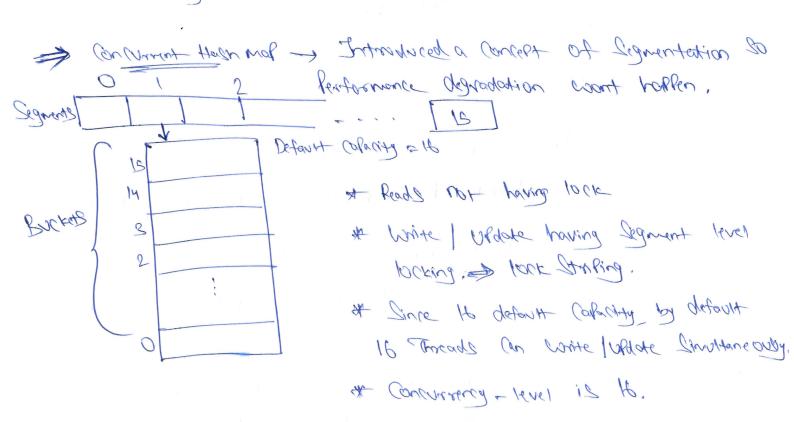
* Collections. Synchronized mas () and Collections. Synchronized List ()

ferrides a basic Conditionally thread - Sate inhementation of

nos and 1124. (Fail - Fast Steredor)

Single Collection readle lock is an impediment to Indubility and it often becames necessary to lock a Collection for a Considerable.

Three during Heration to Request Converent Modification Exception.



* Hash Table all methods including Set is Gracemonized.

* (HM - Pet & regimen use look Striking > And now cound.

- Chet: wave is volodite, next is final - Here in Stake reads

Sync Chesing Corning hollend-betone Mutual-exclusion (Resinistic book)

Vonatile X

Atomic (CAS)

- of Oftimistic lock Read 15 allowed arrows thread. write 15
- * Pesimilare lock > no read write till lock is opposed.

I sues with Symphonization:

- O are condition for look
- (2) NO lock folling, i.e Thread which does not get a lock lan't do conthing else. [trylock () method)
- (2) NO Third lak waits. [try lock (hardtime Thre voit)]
- (i) No Internited lock waits.
- (E) Bronk Structured look (Irm is doing an instruction look.

 Ex! Colling an importment look)

World () | Notify ():

- O I Jua, the Object is the entity that is showed between thrull conich allows them to communicate with each other. The threeds have no knowledge of each other and they sun Asynchronously.
- (2) In Java, you wantly on a Particular instance of an Object Called the monitor. Notify () will notify other through waiting on Same Object.

- there is and notify () was on a Thread, took thread would have to know the Status of every other thread, then would thread I know that thread 2 was bailting for access to Ponticipa resource?
- of Jun has a hidden boot avenes, which Puts all threads in a amove waiting for Same Object and then notify () will deavene & Put to remains State.
 - notify () -> Printer example

 notify An() -> Approarion Corne George -> Obstate wake of Call

 Ex: Parallel harbage Conector.

=> Frankler Queve

- * Extends Blocking Queve.
- * transfer co method is like a Sychronous Putco method.
- # try Francher () returns booken

Je Maybe continuée dos 40 contine un ne not à

Inked Frankler Ougre! Inflements Frankler Quere

Dod Queve or Did blaking Queve hence Colled

Procking many Debends on Consumer

halvorting continue (but 1) -> rectum blocking Court

A factor blocking Queve (JDK 5 & 6)

-> linked Brocking Queve (JDK 7)