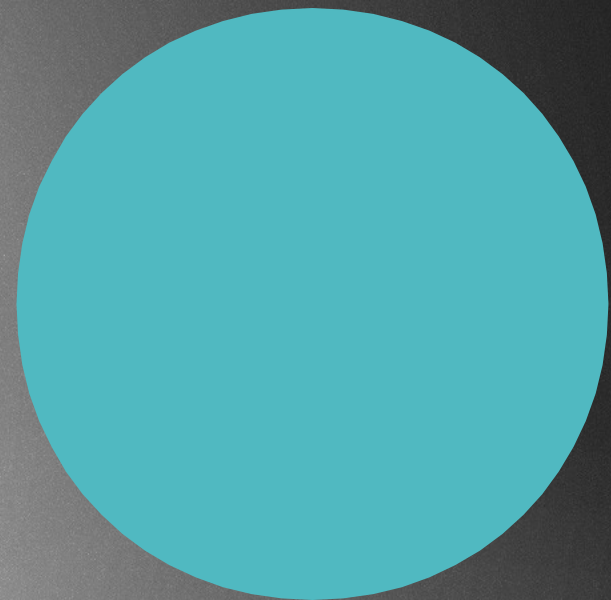
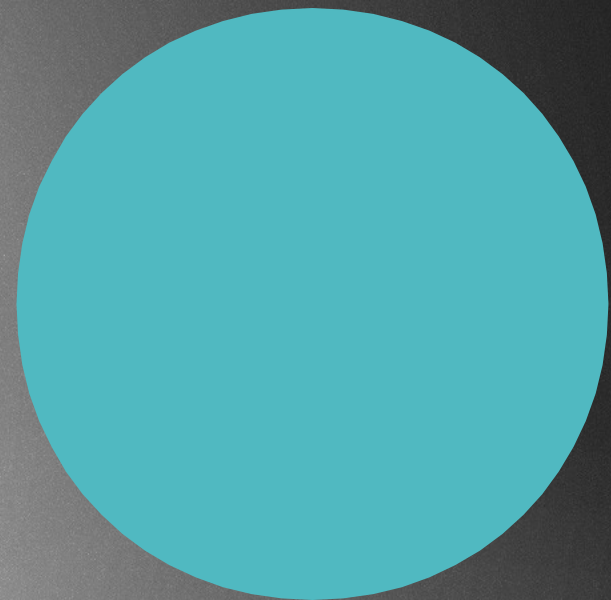
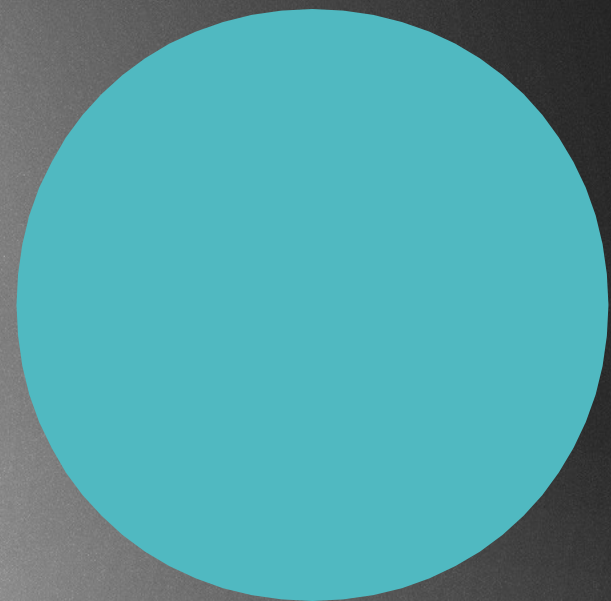


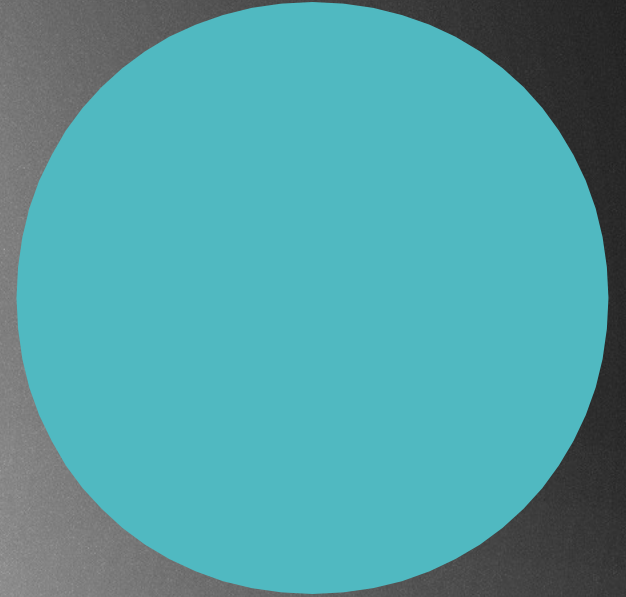
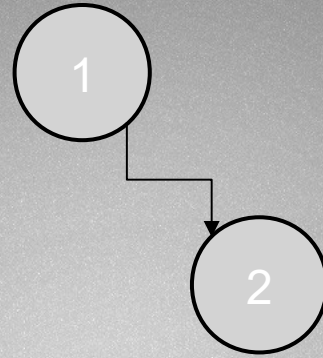
RED-BLACK TREES

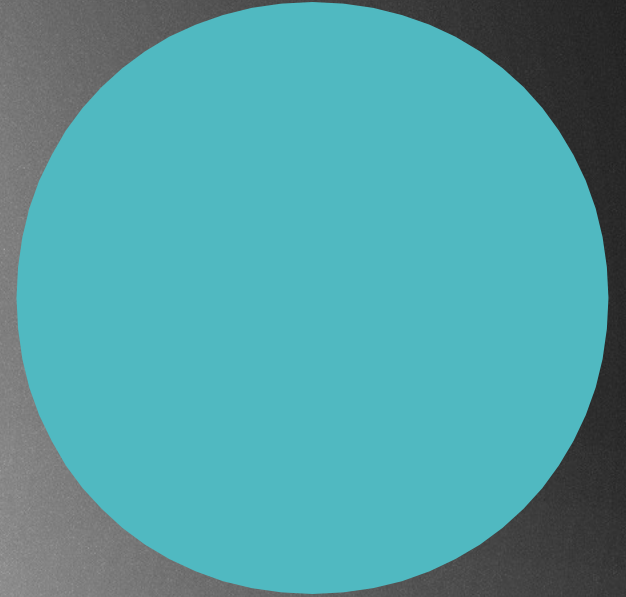
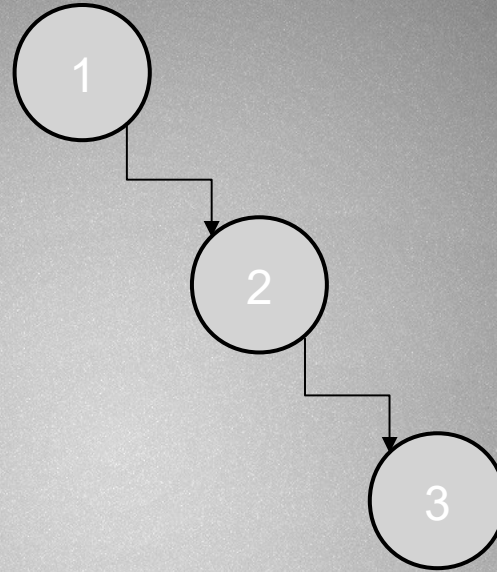


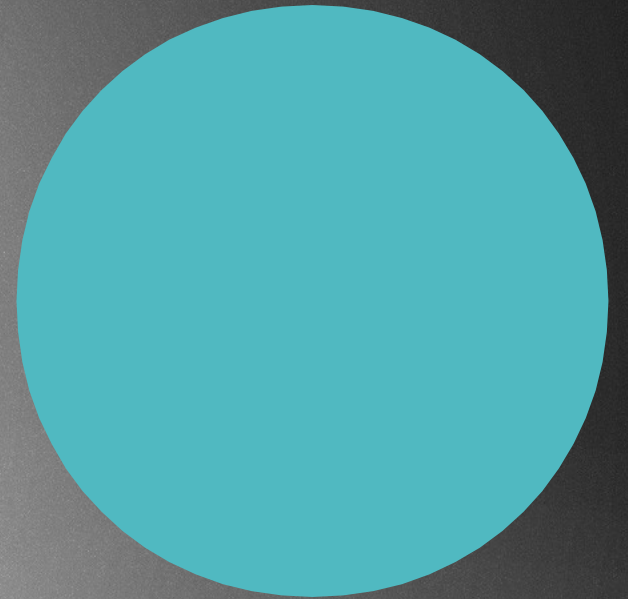
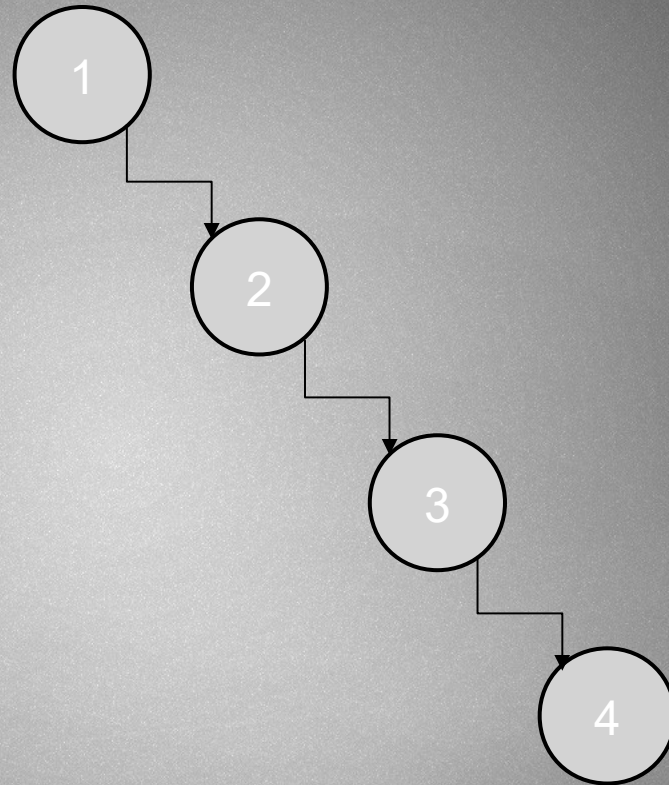
Construct a BST from a sorted array
[1,2,3,4]






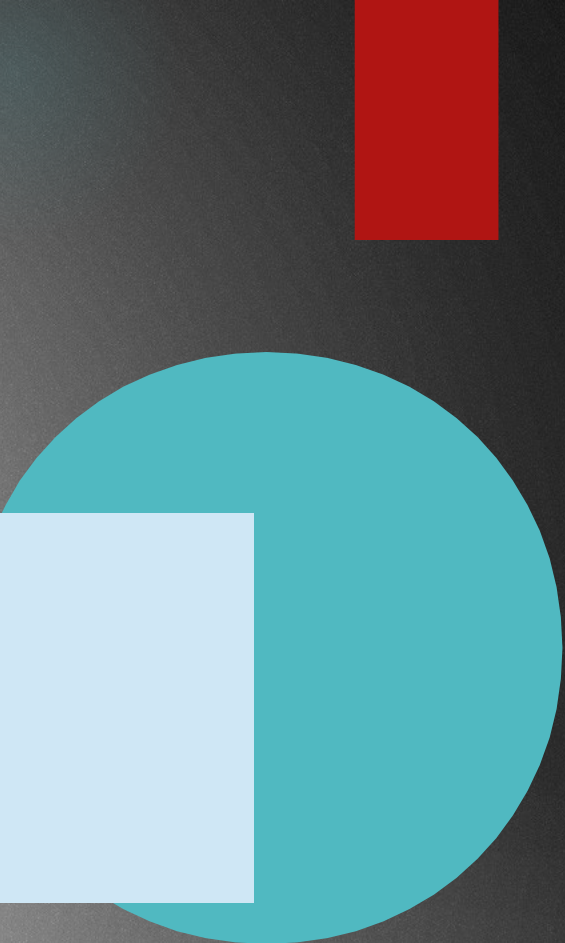






It is a linked list: **$O(N)$** !!!

- 
- ▶ The running time of BST operations depends on the height of the binary search tree: we should keep the tree balanced in order to get the best performance
 - ▶ In an AVL tree, the heights of the two child subtrees of any node differ by at most one
 - ▶ Another solution to the problem is a red-black trees
 - ▶ AVL trees are faster than red-black trees because they are more rigidly balanced BUT needs more work
 - ▶ Operating Systems relies heavily on these data structures !!!

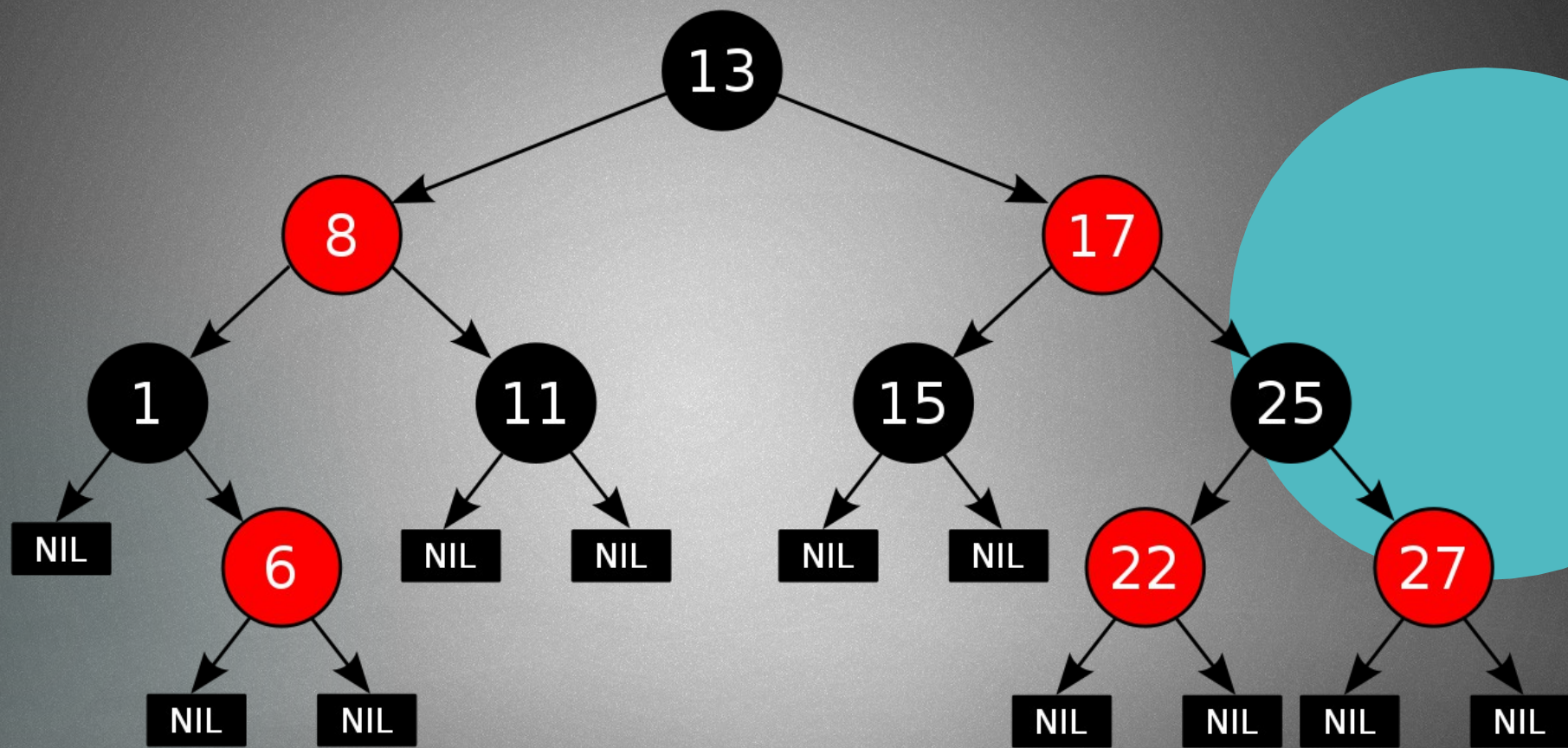


	Average case	Worst case
Space	$O(n)$	$O(n)$
Insert	$O(\log n)$	$O(\log n)$
Delete	$O(\log n)$	$O(\log n)$
Search	$O(\log n)$	$O(\log n)$

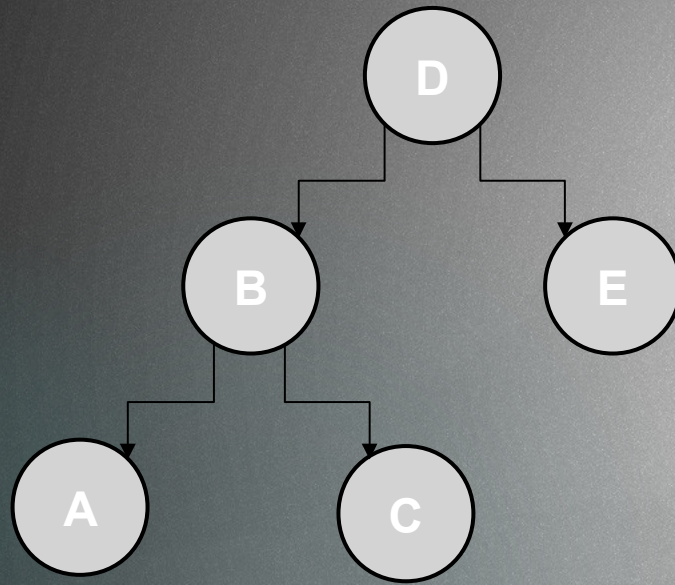
Red-black properties:

- ▶ Each node is either red or black
- ▶ The root node is always **black**
- ▶ All leaves (NIL or NULL) are black
- ▶ Every red node must have two black child nodes and no other children → it must have a black parent
- ▶ Every path from a given node to any of its descendant NIL/NULL nodes contains the same number of black nodes





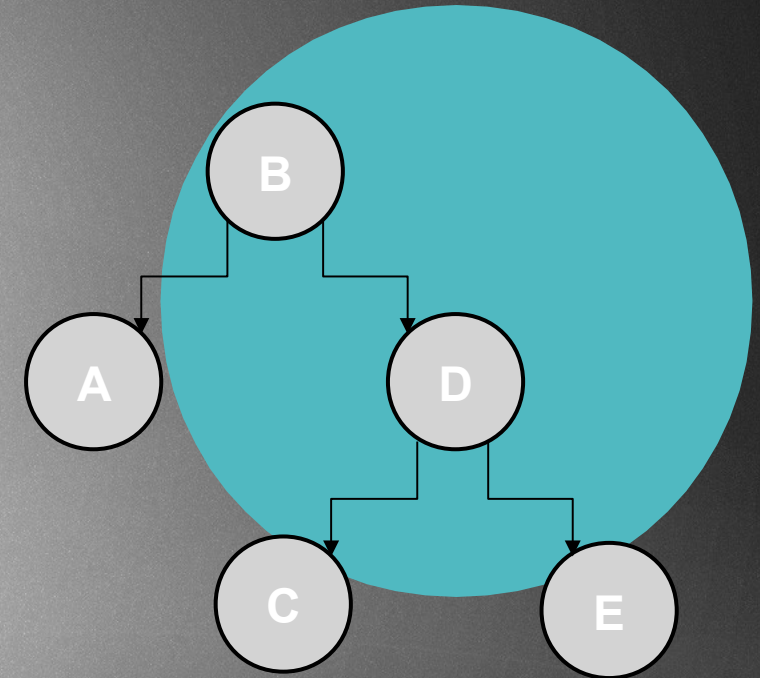
Rotations



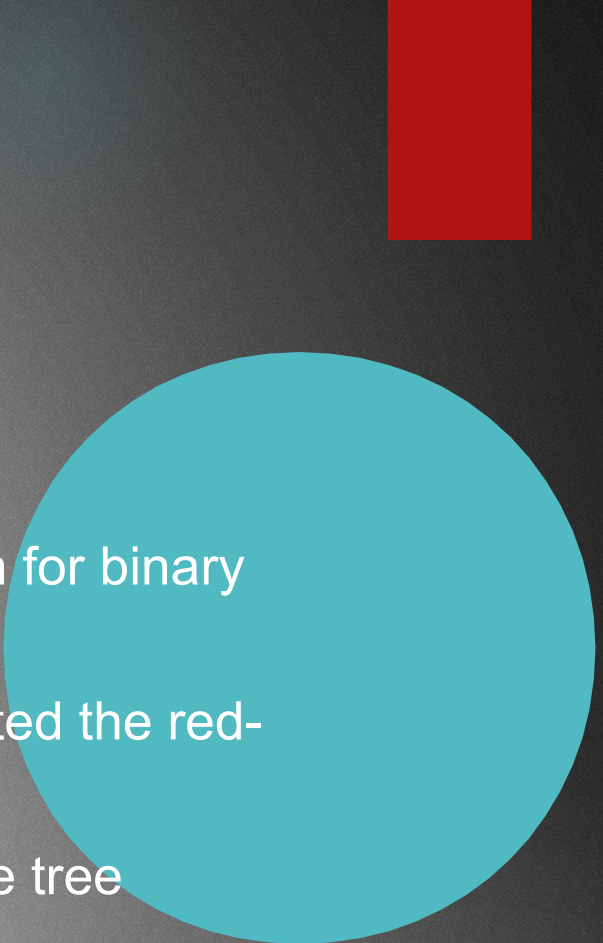
rightRotate(D)

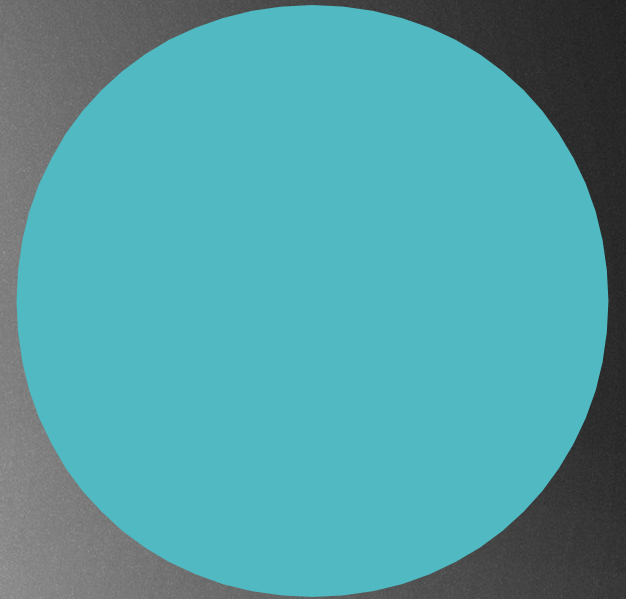


leftRotate(B)

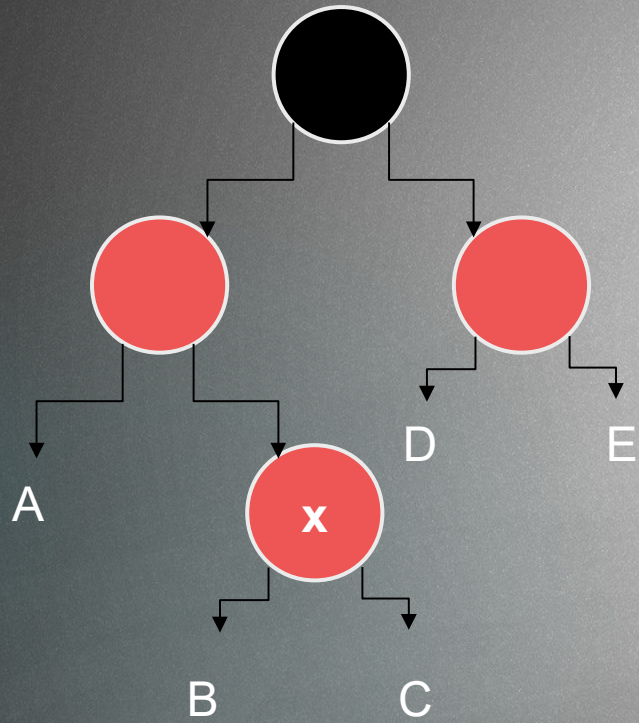


We just have to update the references which can be done in **O(1)** time complexity !!! (the in-order traversal is the same)

- 
- ▶ Every new node is red by default
 - ▶ We keep inserting new node in the same way as we have seen for binary search trees (or AVL trees)
 - ▶ On every insertion → we have to check whether we have violated the red-black properties or not
 - ▶ If we have violated the RB properties: we have to rebalance the tree
 - ▶ 1.) make rotations
 - ▶ 2.) OR just recolor the nodes

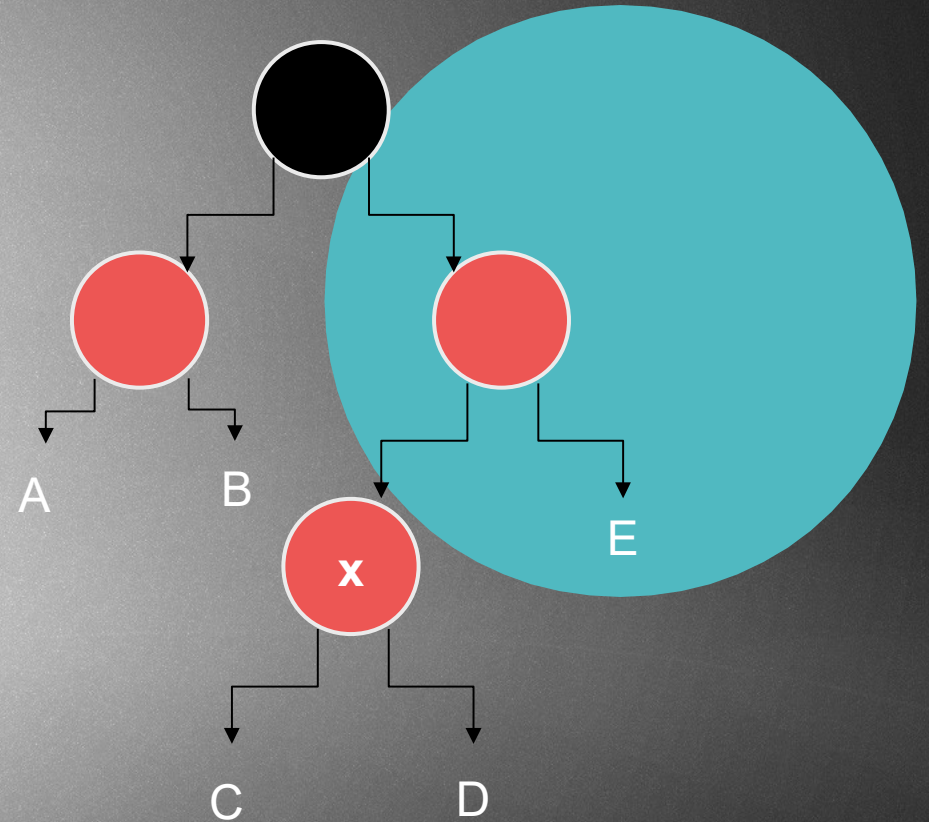
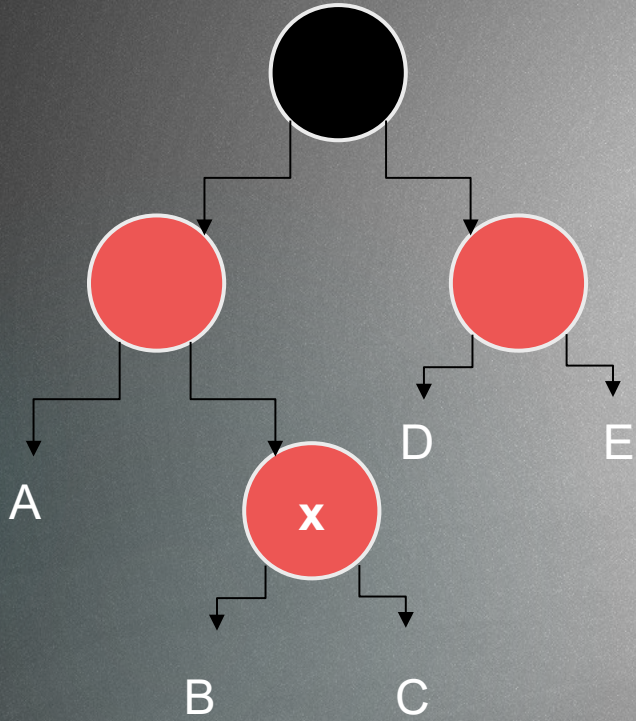


Case 1:



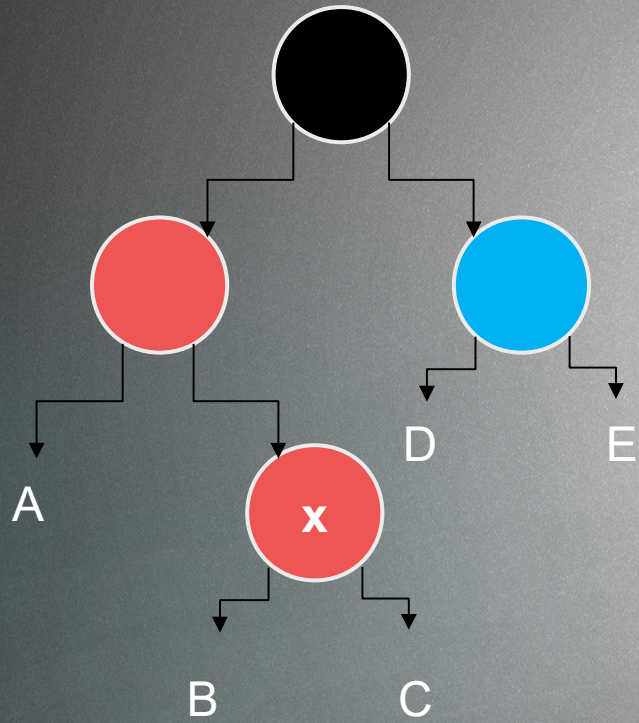
Case 1:

This problem is symmetric !!!



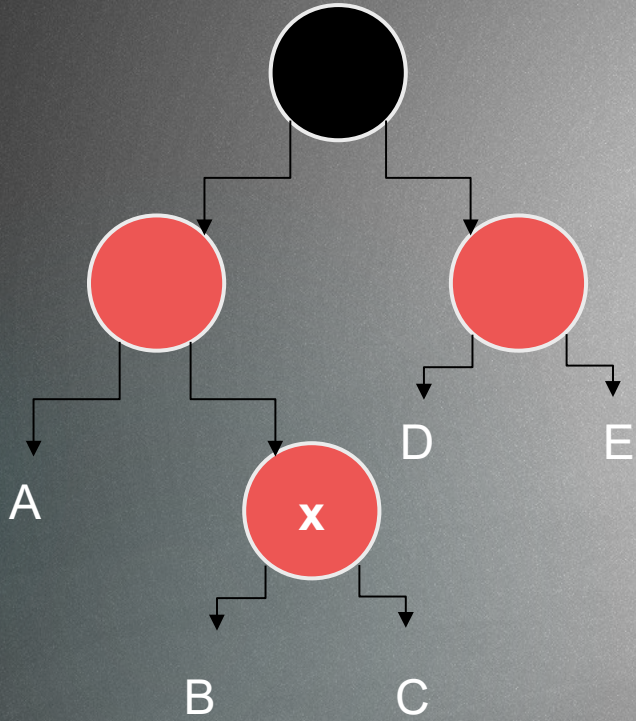
Case 1:

The „uncle” of x is red too !!!



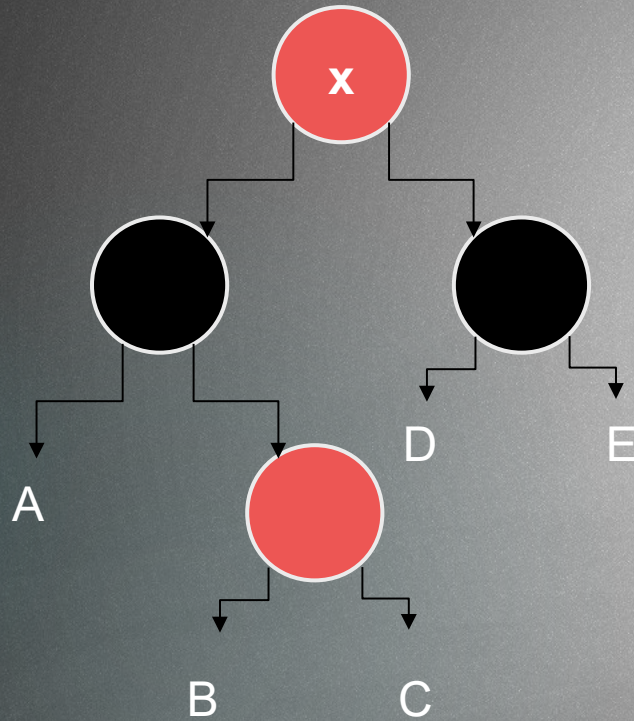
Case 1:

We just have to recolor some nodes, quite easy case



Case 1:

We just have to recolor some nodes, quite easy case + the **x** will be the root node in this case

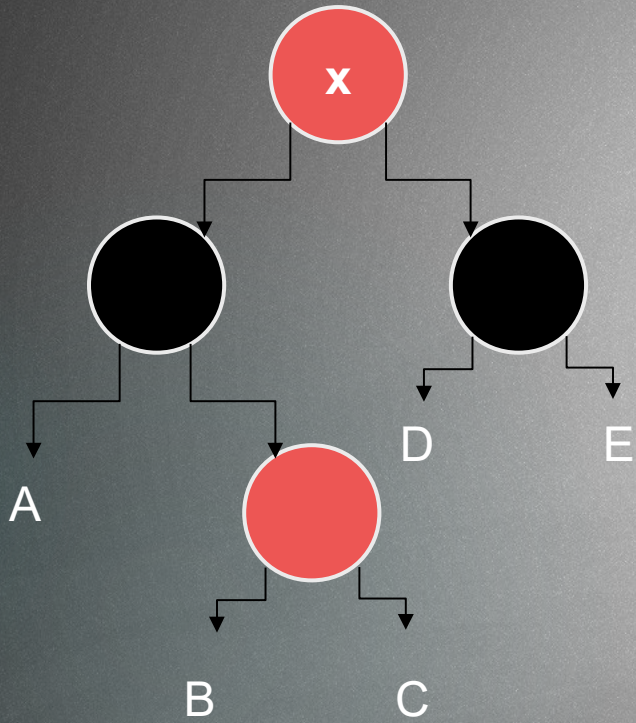


IMPORTANT: with this recoloring, maybe we violate the red-black properties in other parts of the tree !!!

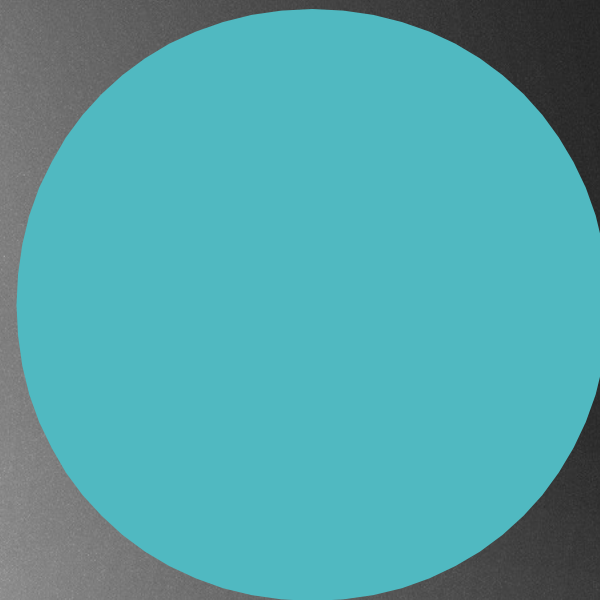
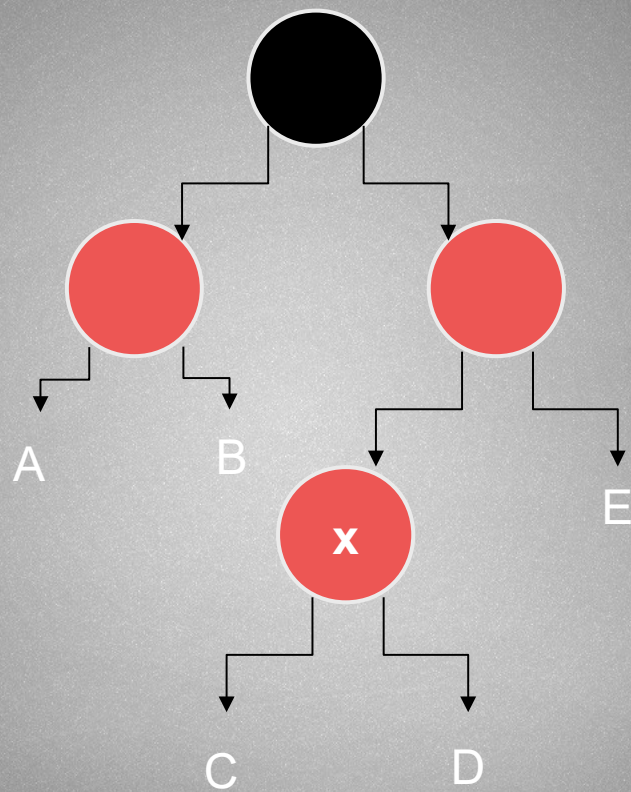
~ have to check recursively on the whole tree

Case 1:

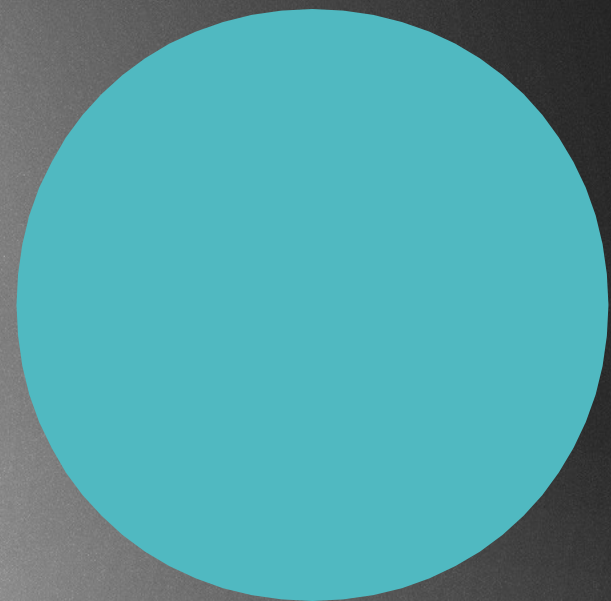
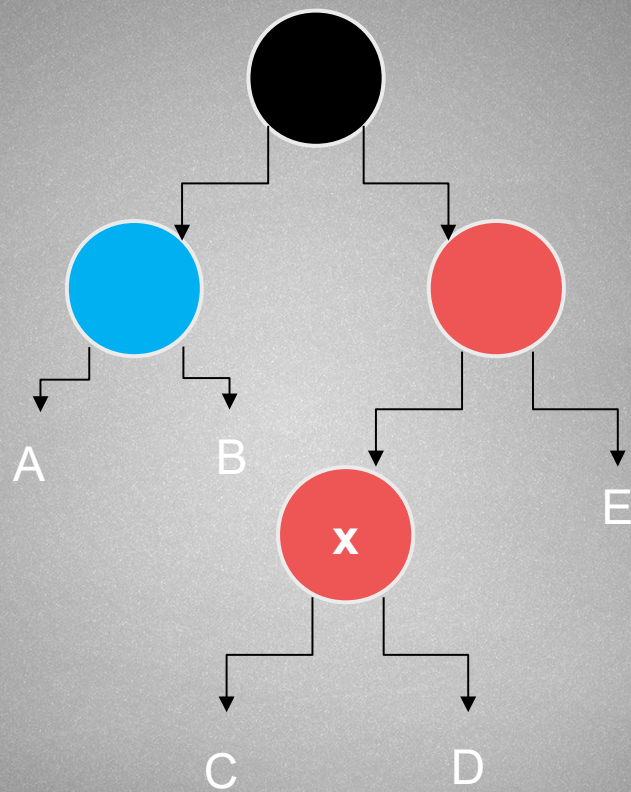
We have to check recursively (from bottom to top) whether the red-black properties are violated or not



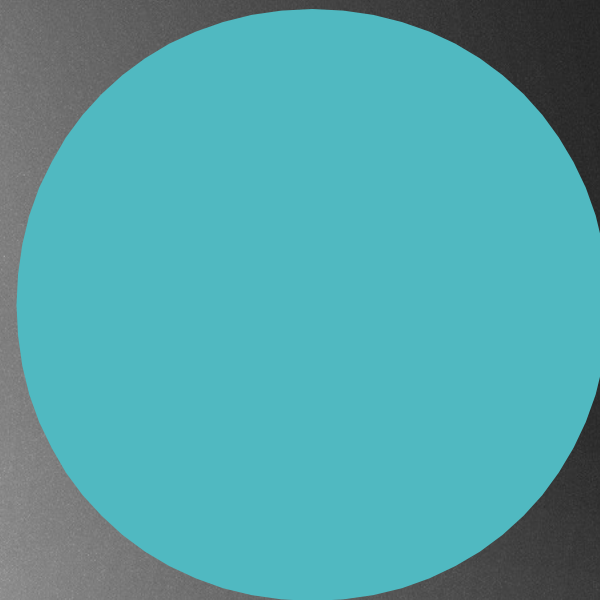
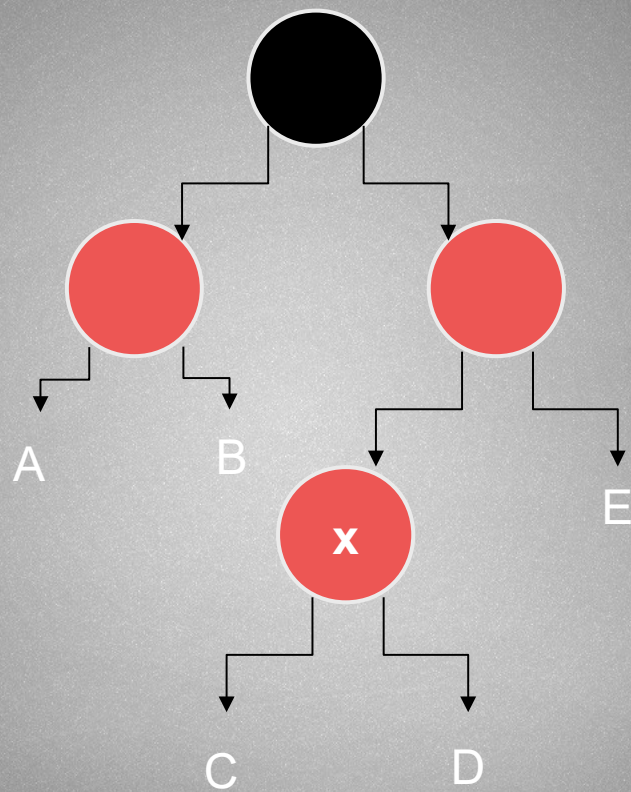
Case 1:



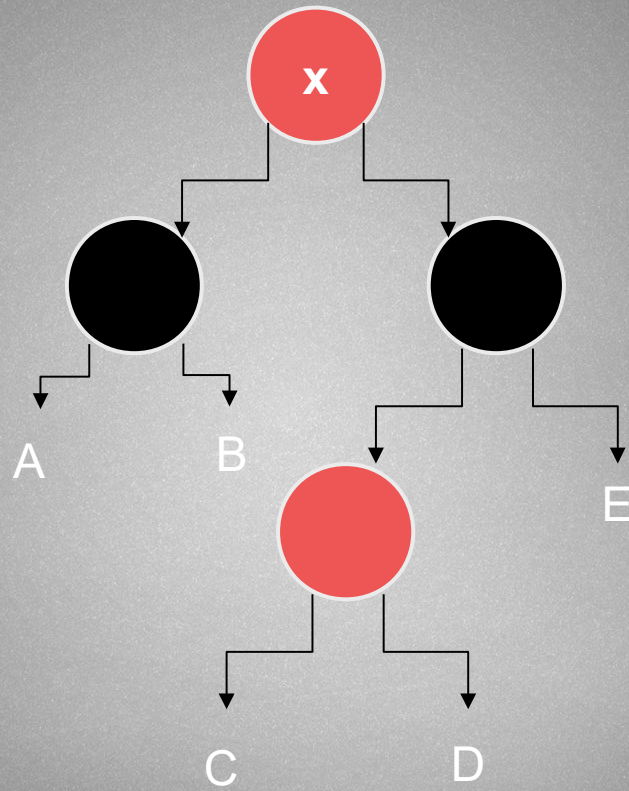
Case 1:

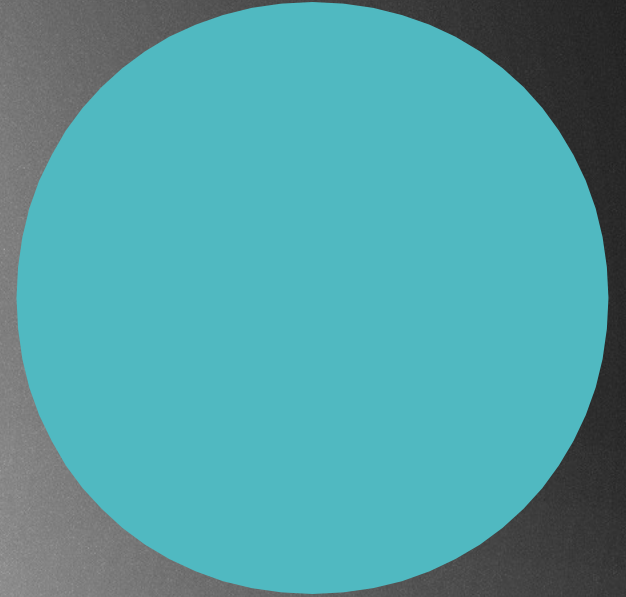


Case 1:

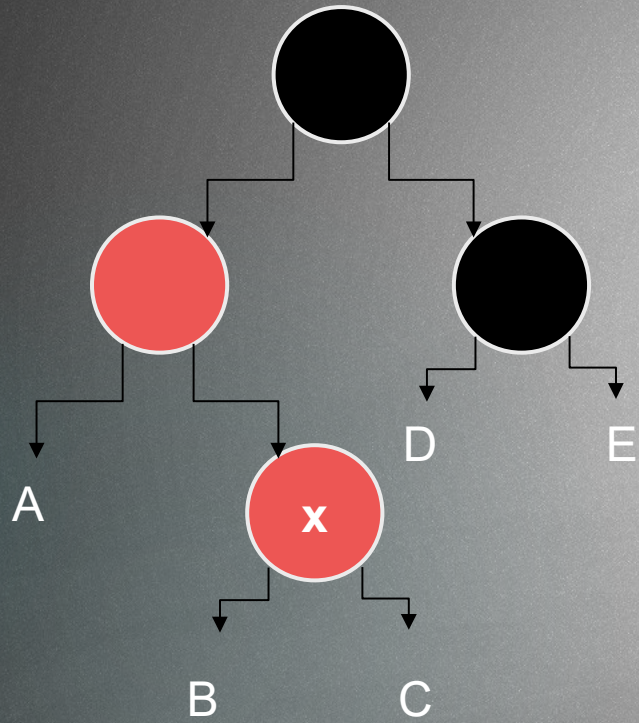


Case 1:



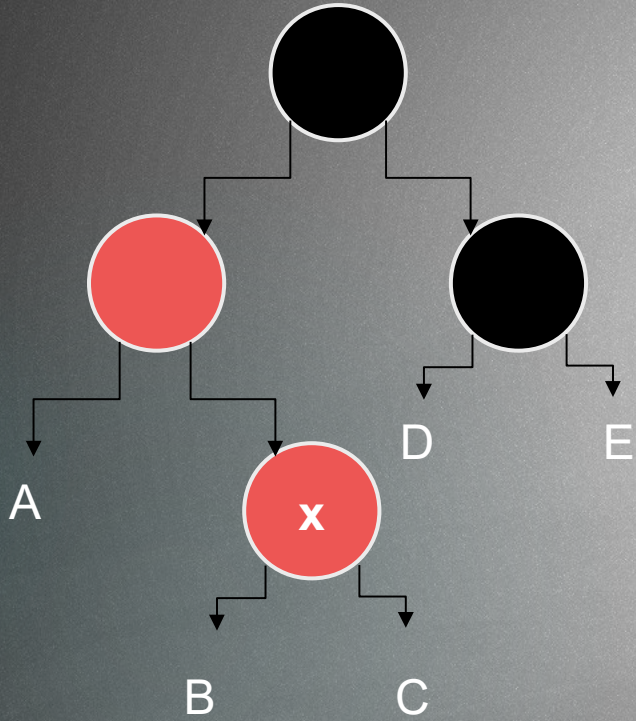


Case 2:



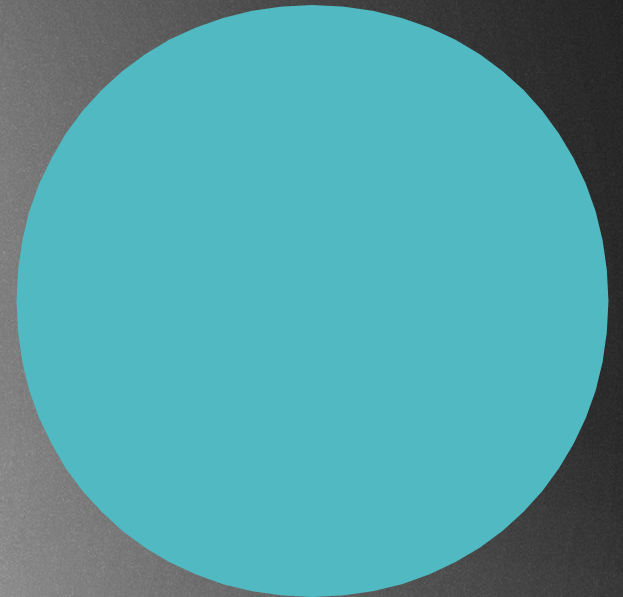
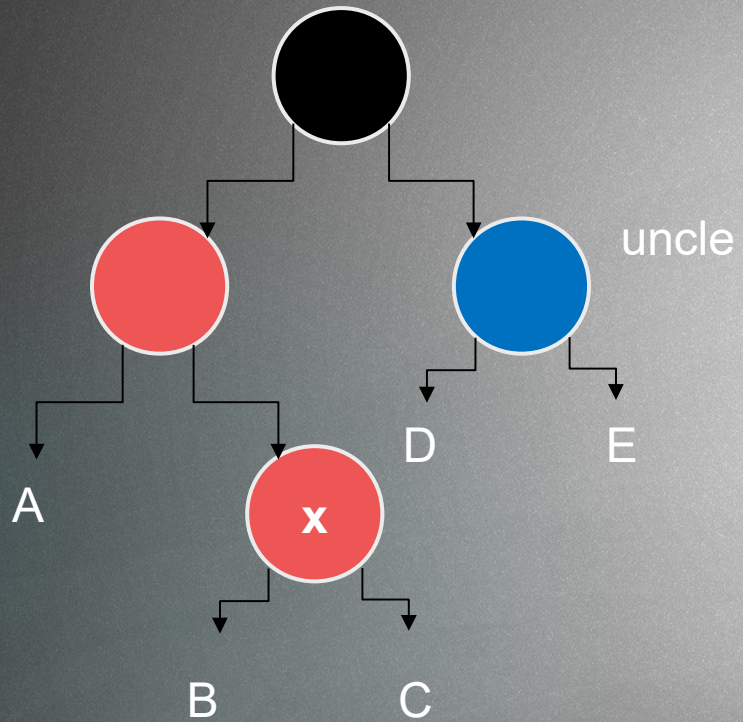
Case 2:

The uncle of node **x** is a black node + node **x** is a right child

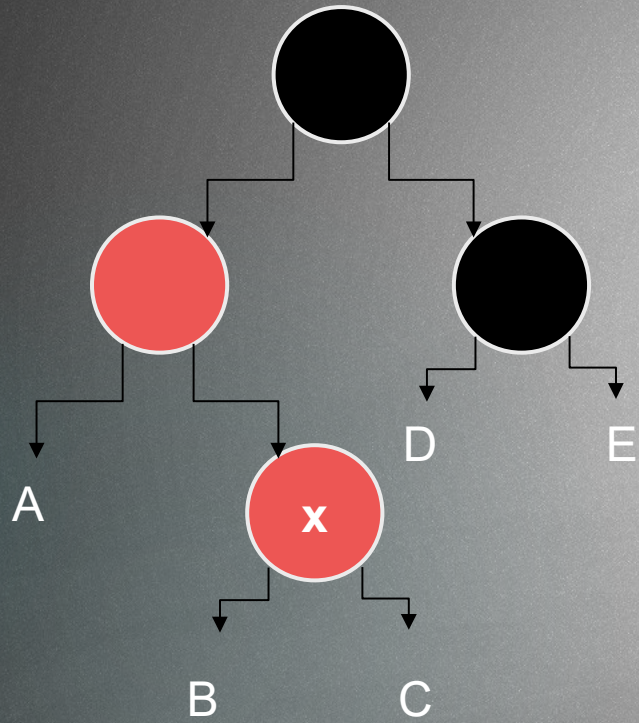


Case 2:

The uncle of node x is a black node + node x is a right child

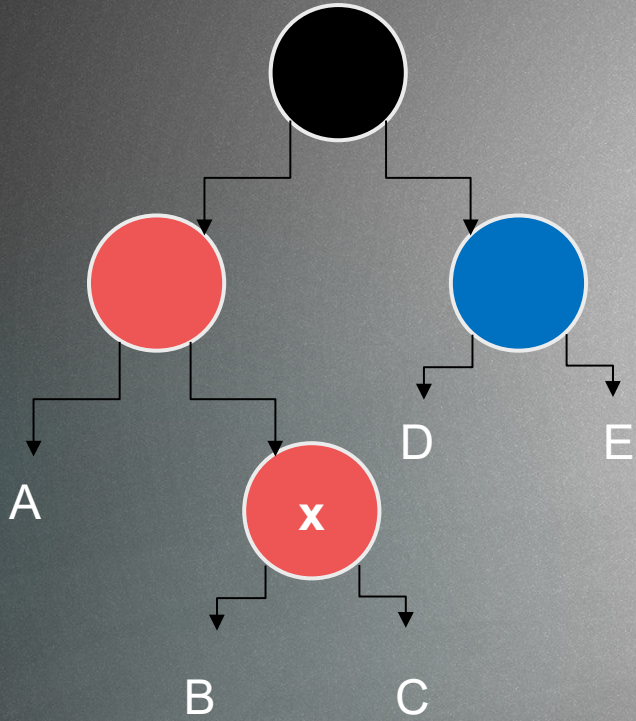


Case 2:

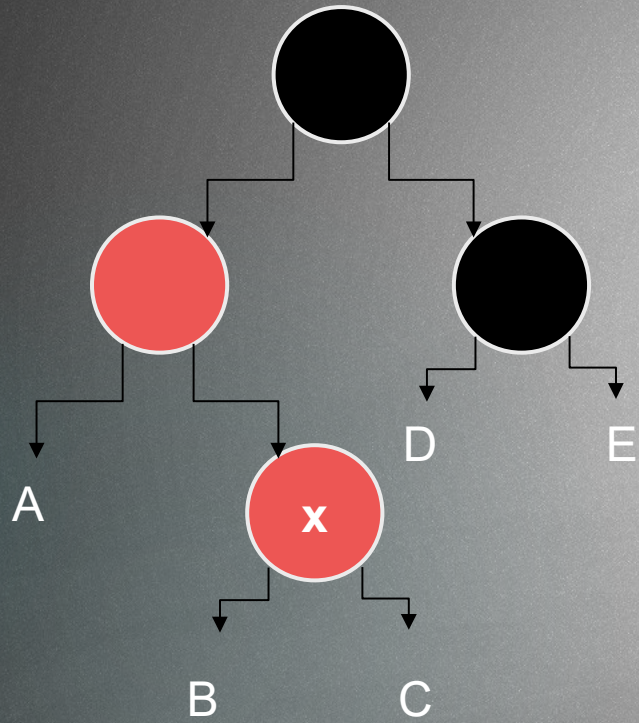


Case 2:

The uncle of node **x** is a black node + node **x** is a right child
We have to make a left rotation on the node **x**'s parent

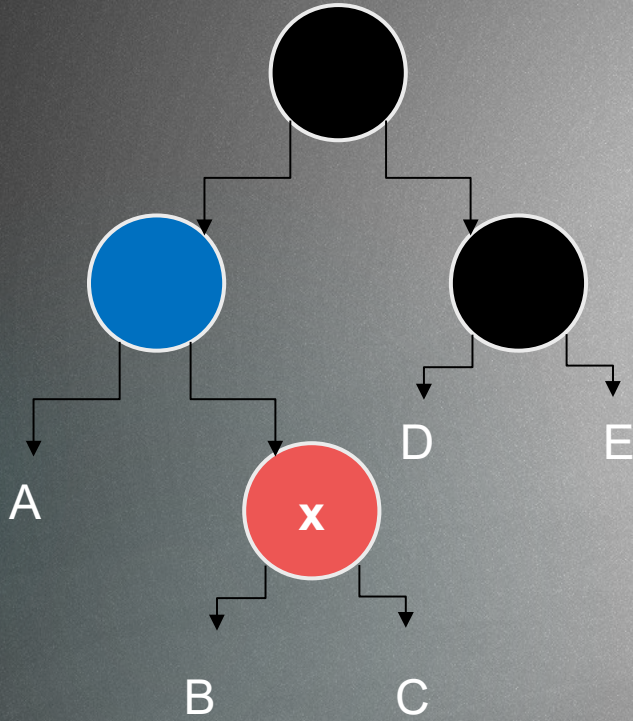


Case 2:

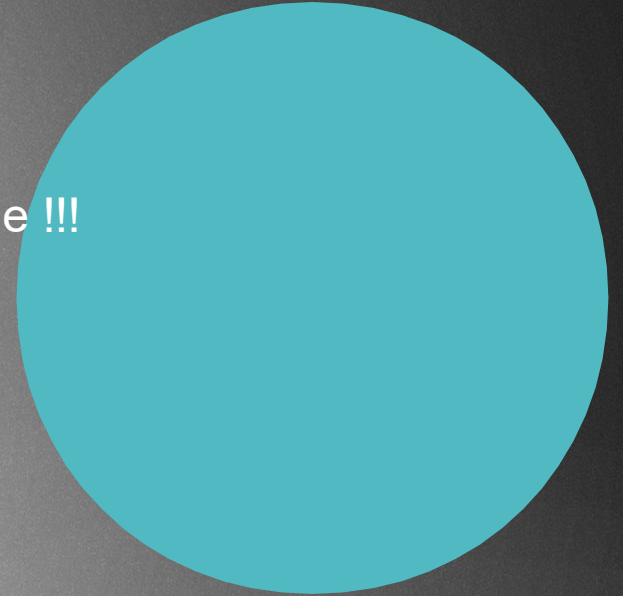


Case 2:

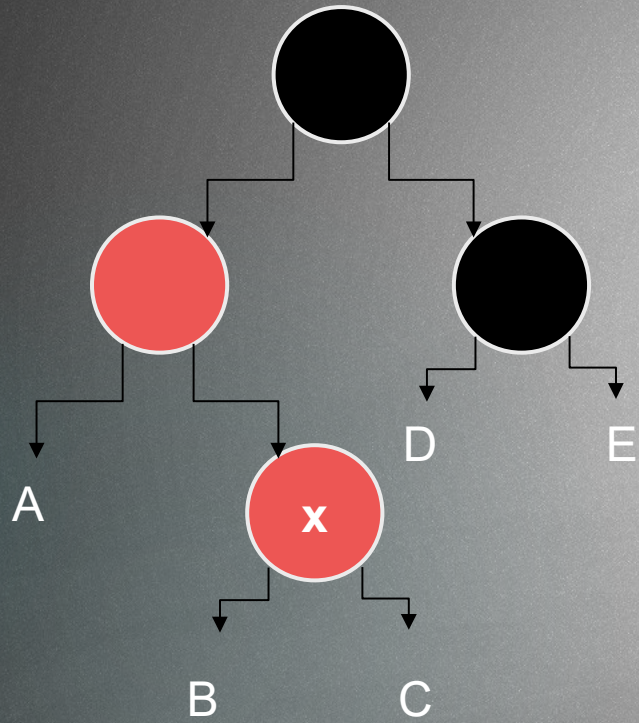
The uncle of node **x** is a black node + node **x** is a right child
We have to make a left rotation on the node **x**'s parent



We have to rotate the blue node !!!

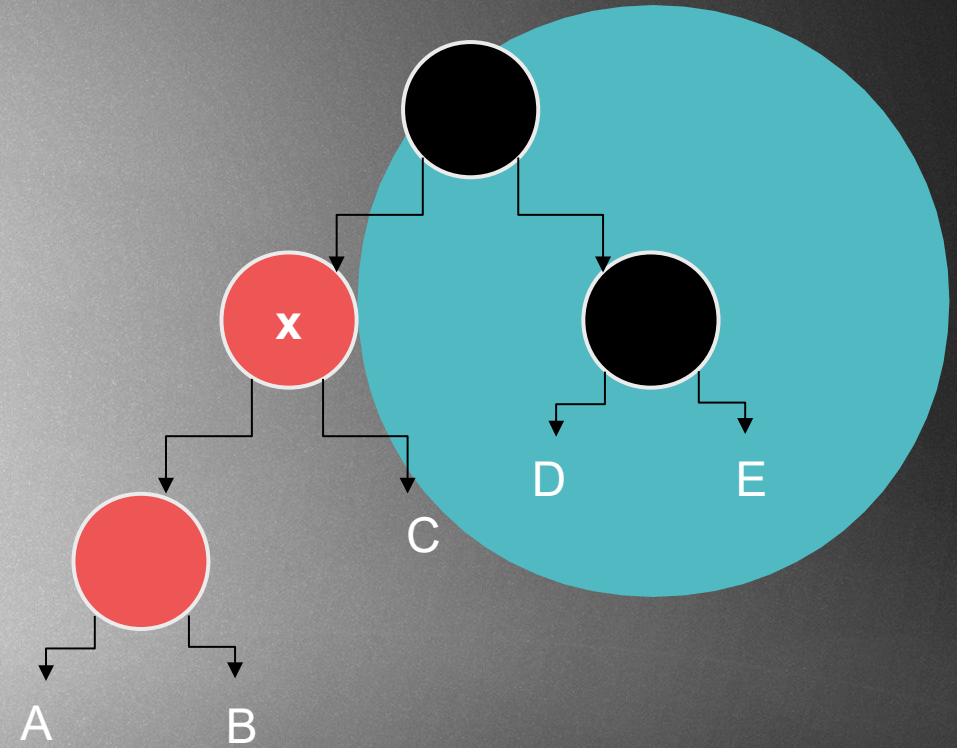
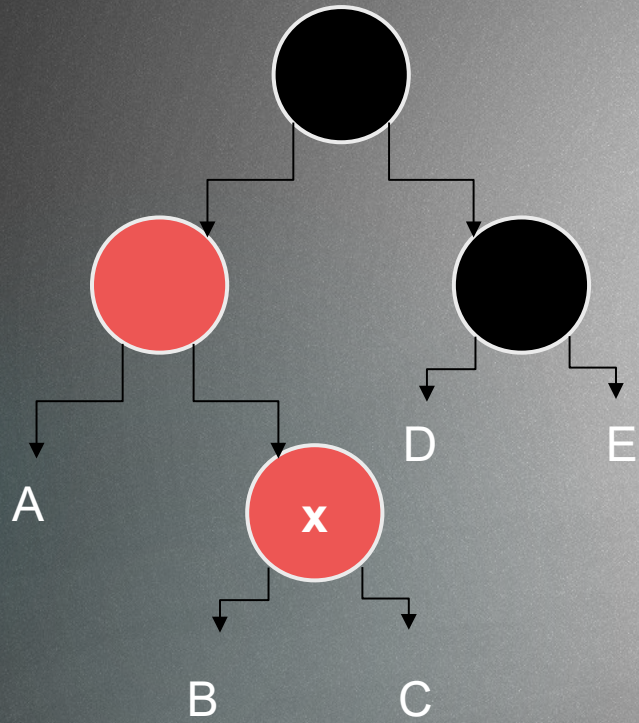


Case 2:



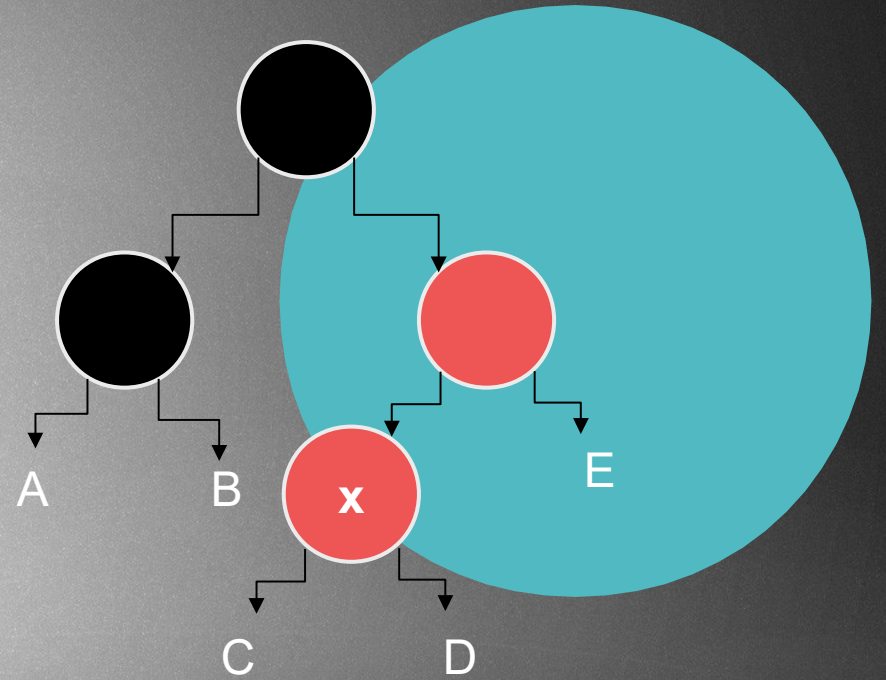
Case 2:

The uncle of node **x** is a black node + node **x** is a right child
We have to make a left rotation on the node **x**'s parent



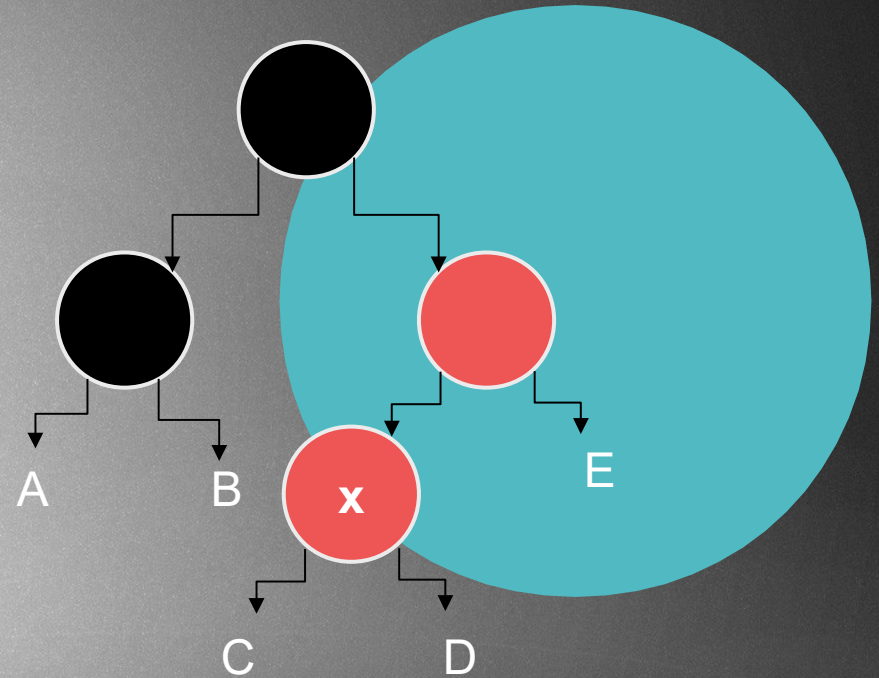
Case 2:

same problem basically but the symmetric version



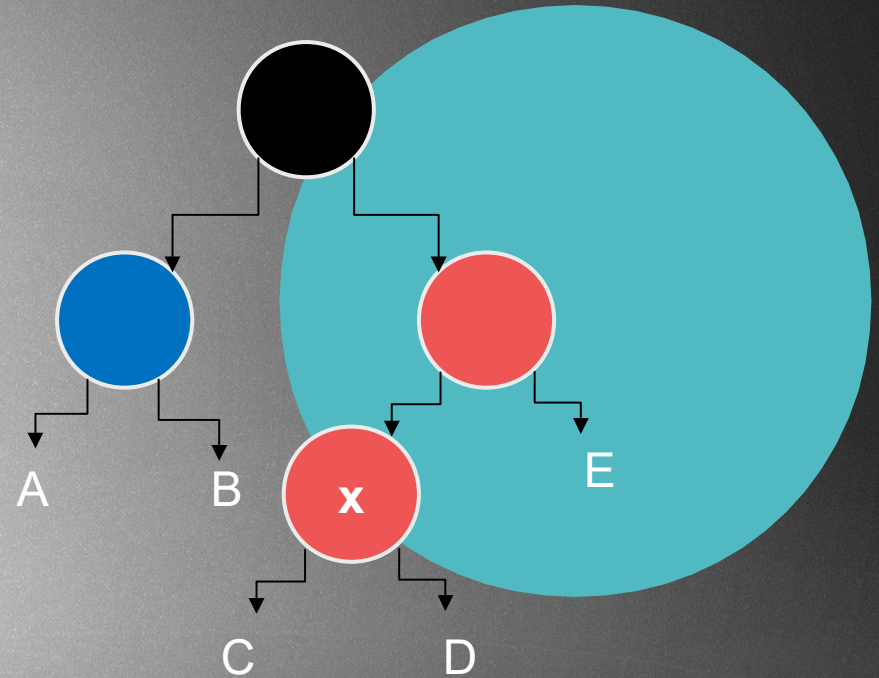
Case 2:

The given **x** node is a left child
+ the parent is red
+ the uncle is black



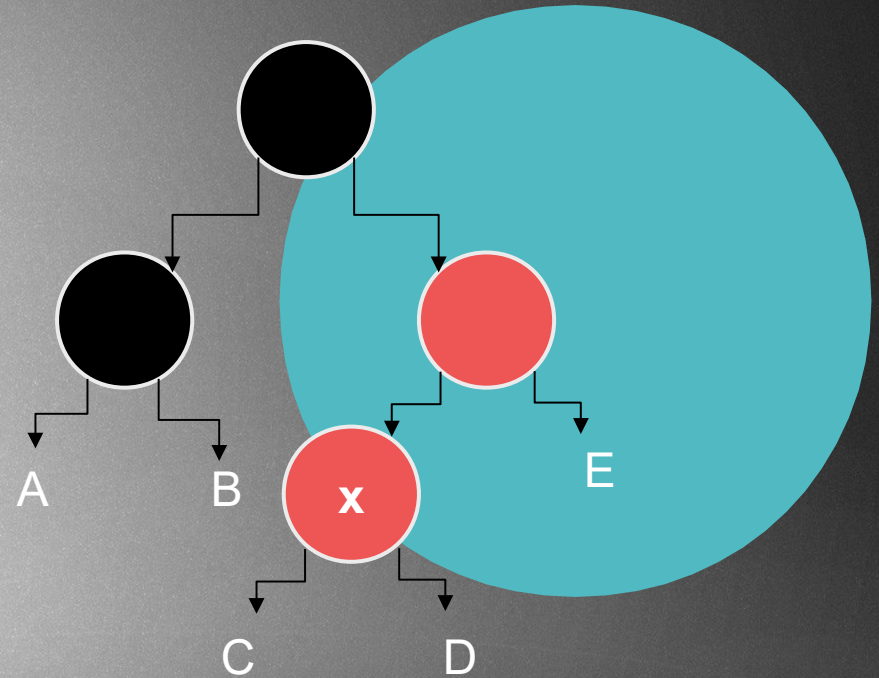
Case 2:

The given **x** node is a left child
+ the parent is red
+ the uncle is black



Case 2:

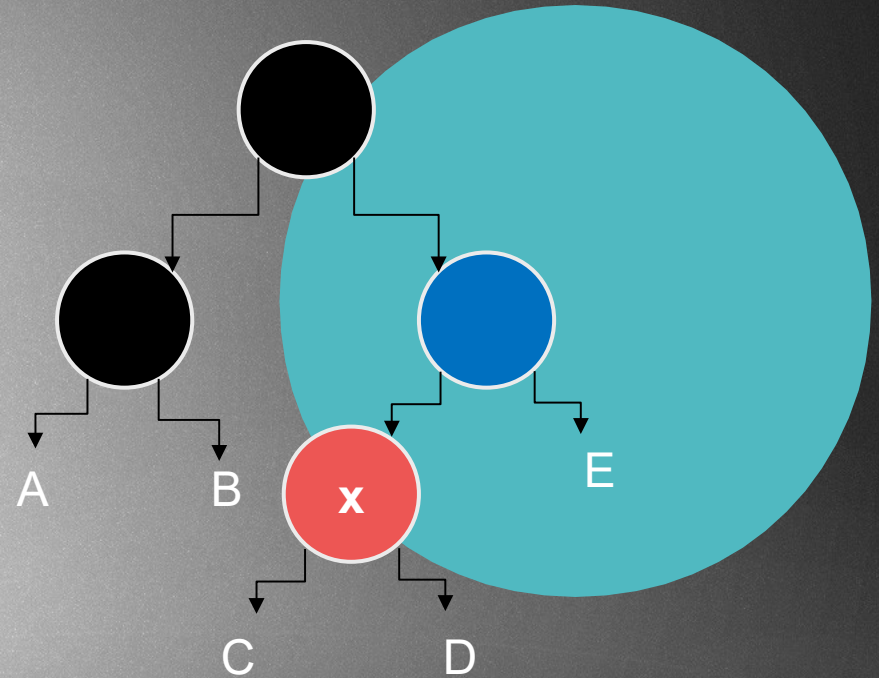
The given **x** node is a left child
+ the parent is red
+ the uncle is black



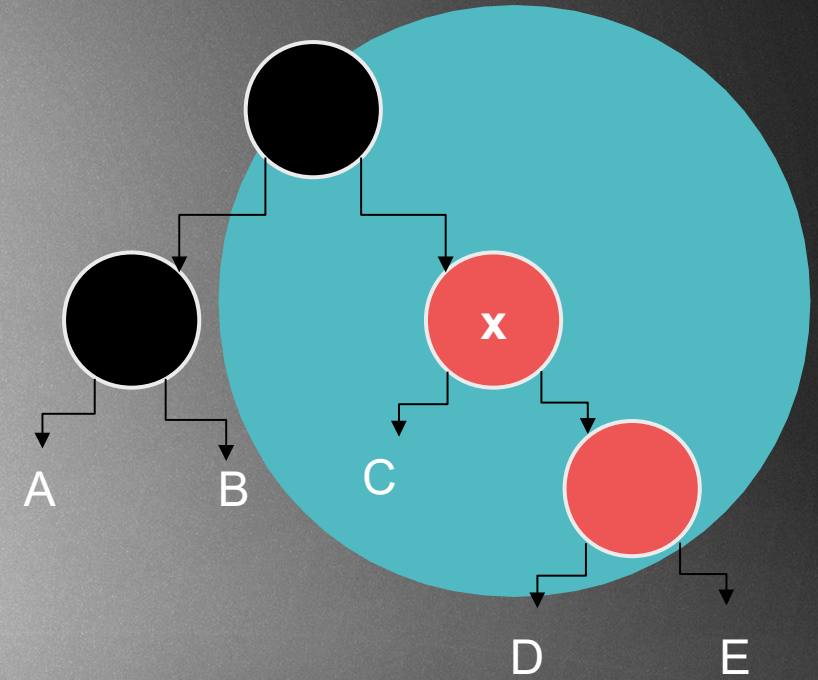
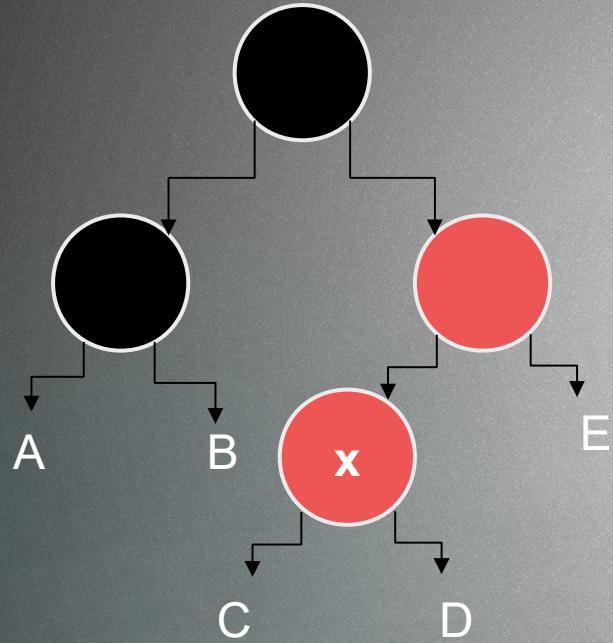
Case 2:

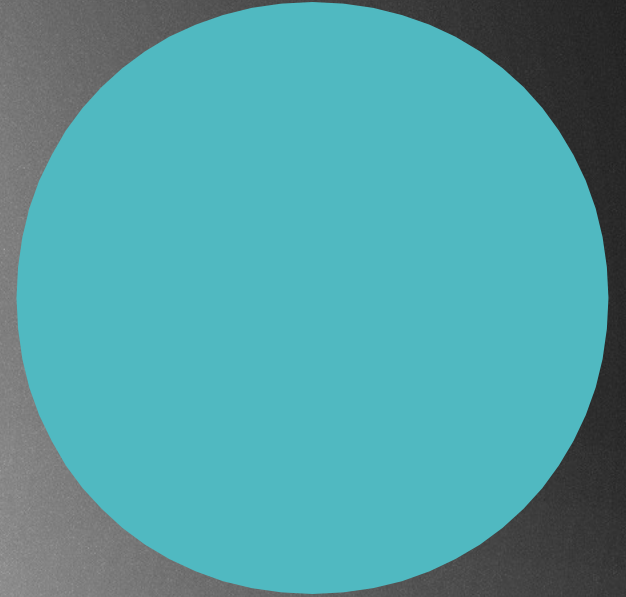
The given x node is a left child
+ the parent is red
+ the uncle is black

We just have to make a right rotation
on the blue node !!!



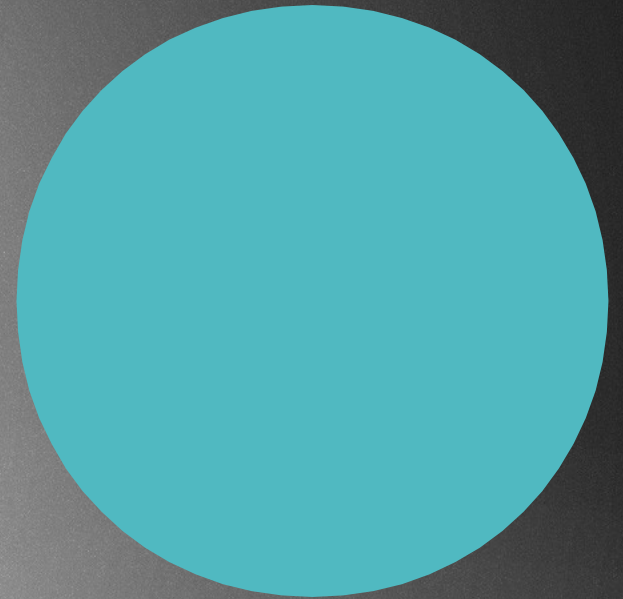
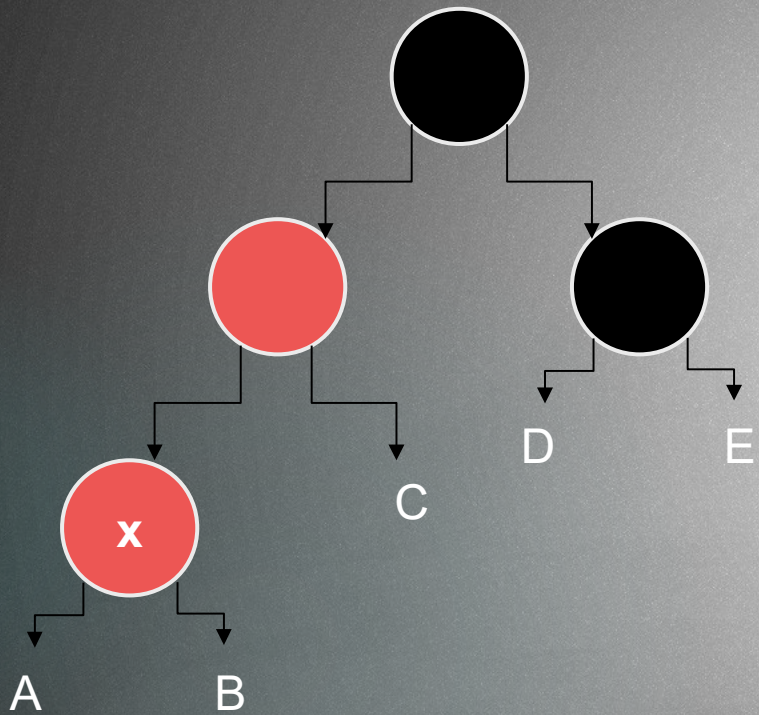
Case 2:





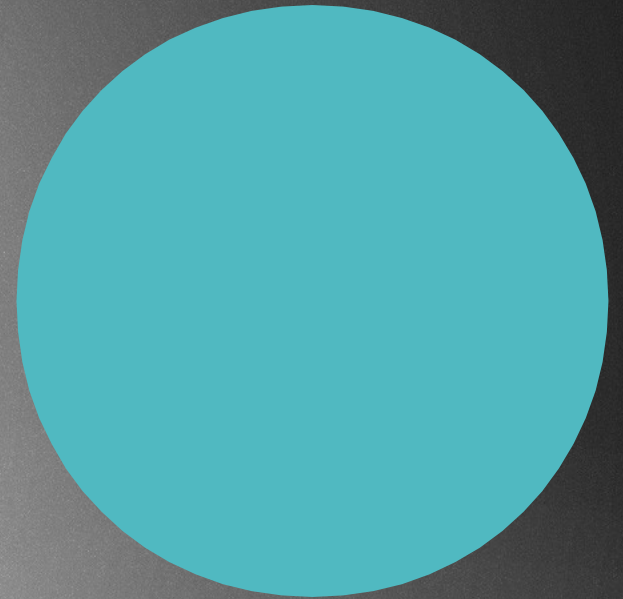
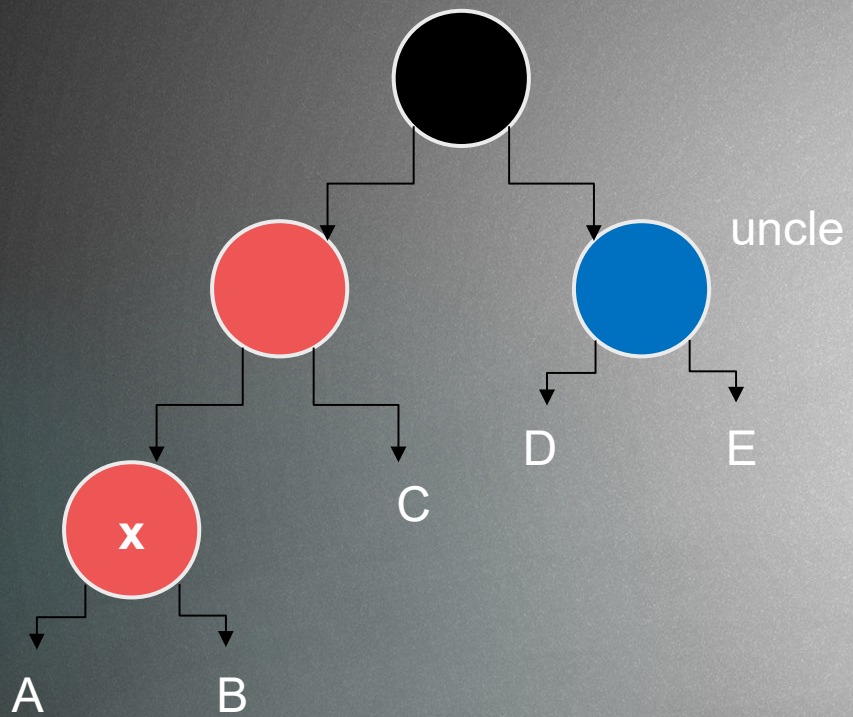
Case 3:

Uncle of node **x** is black and node **x** is a left child



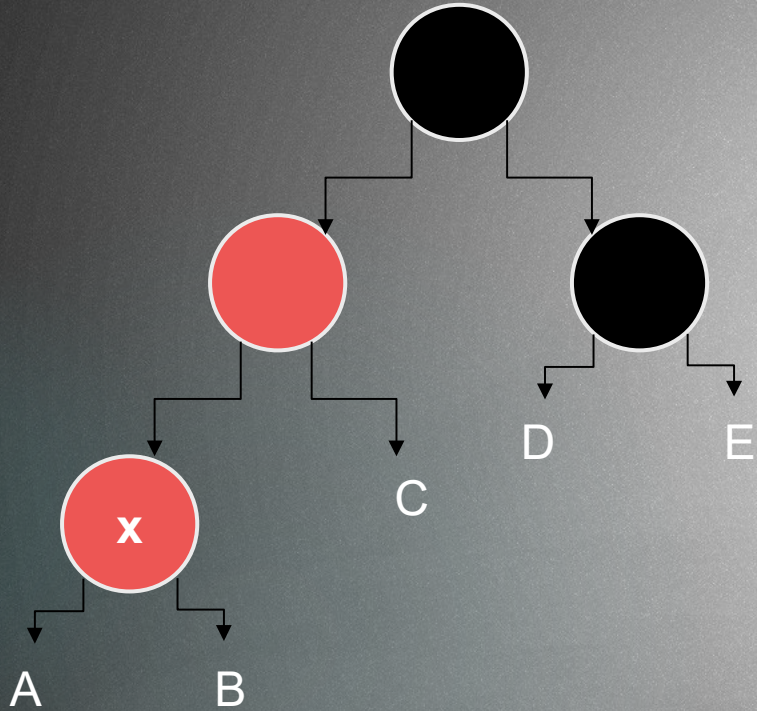
Case 3:

Uncle of node **x** is black and node **x** is a left child



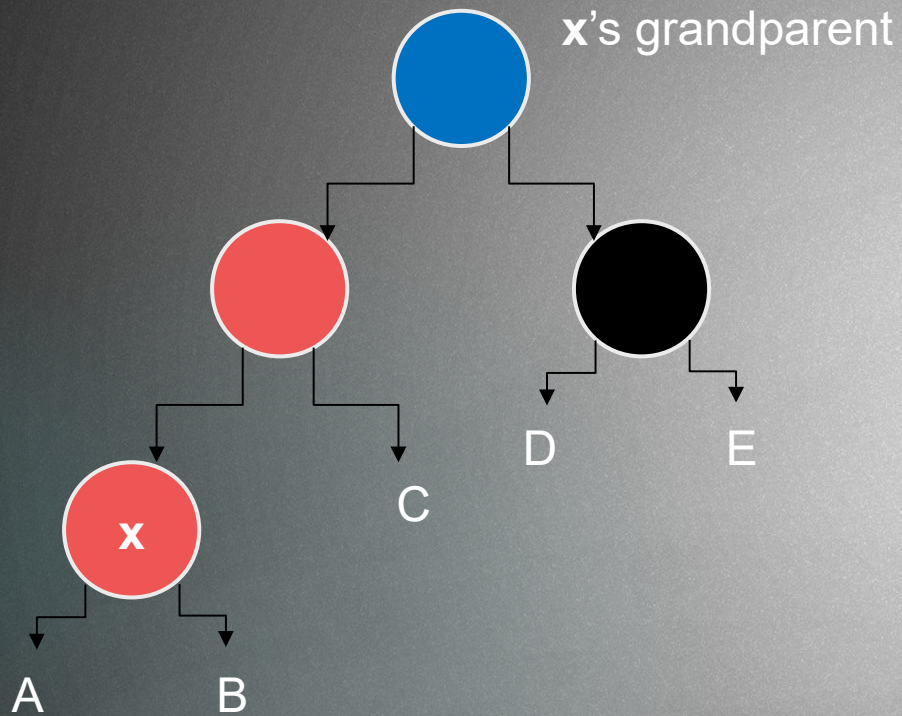
Case 3:

Uncle of node **x** is black and node **x** is a left child
We have to make a rotation on **x** node's grandparent



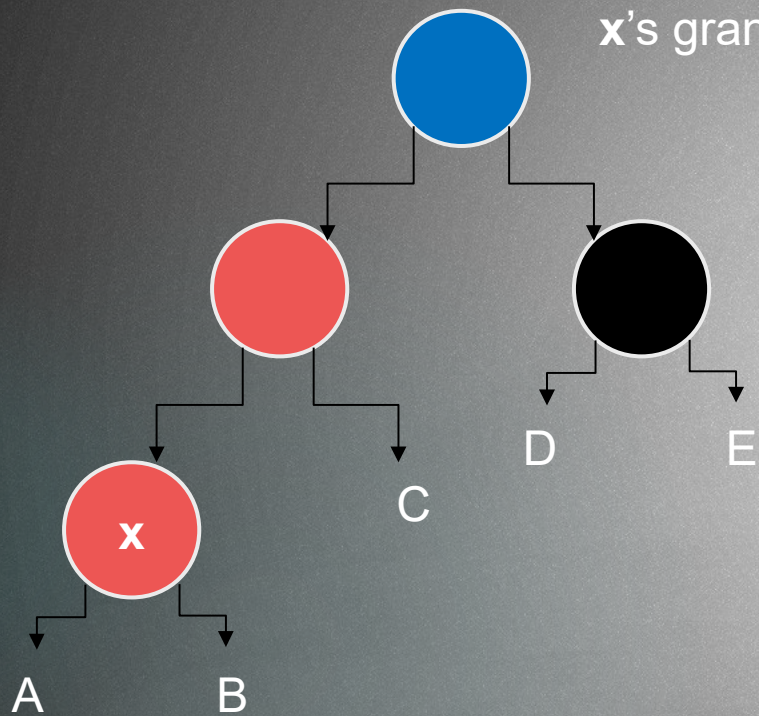
Case 3:

Uncle of node **x** is black and node **x** is a left child
We have to make a rotation on **x** node's grandparent



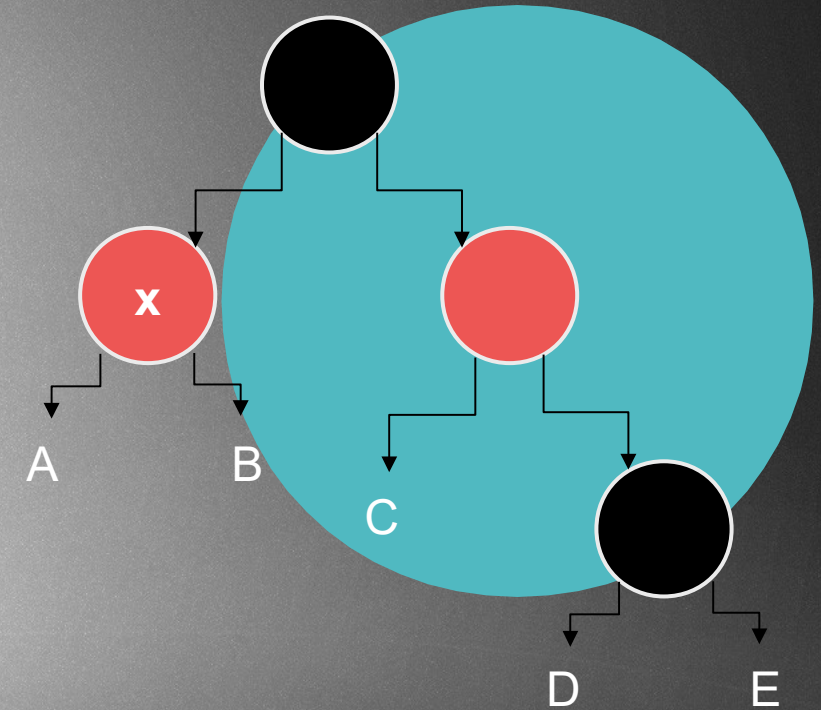
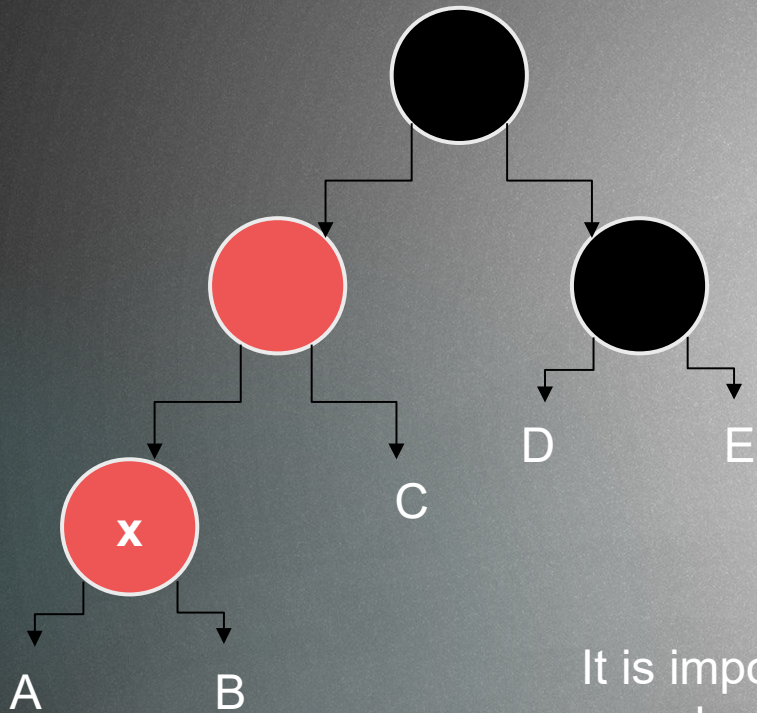
Case 3:

Uncle of node **x** is black and node **x** is a left child
We have to make a rotation on **x** node's grandparent



Case 3:

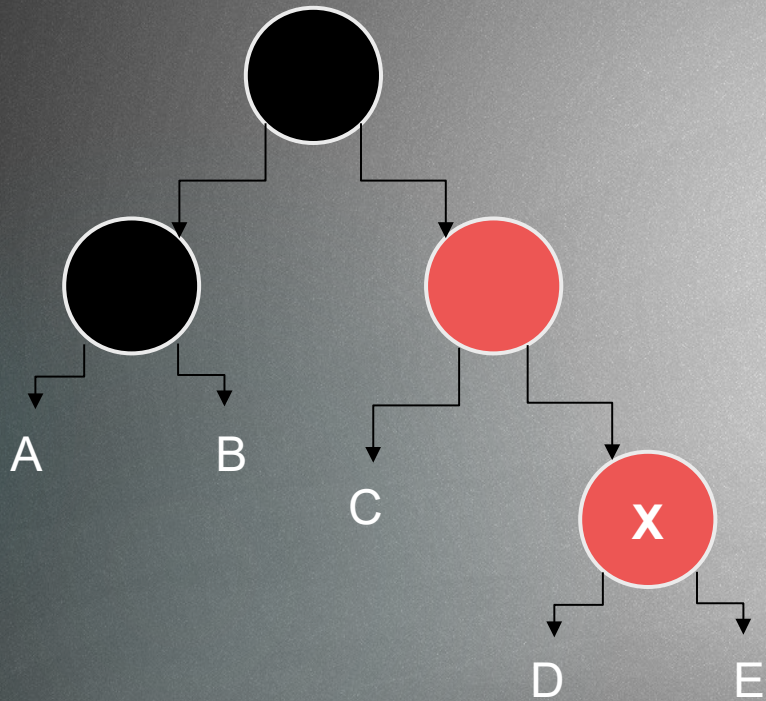
Uncle of node **x** is black and node **x** is a left child
We have to make a rotation on **x** node's grandparent



It is important that after the rotations
we have to make some recoloring as well !!!

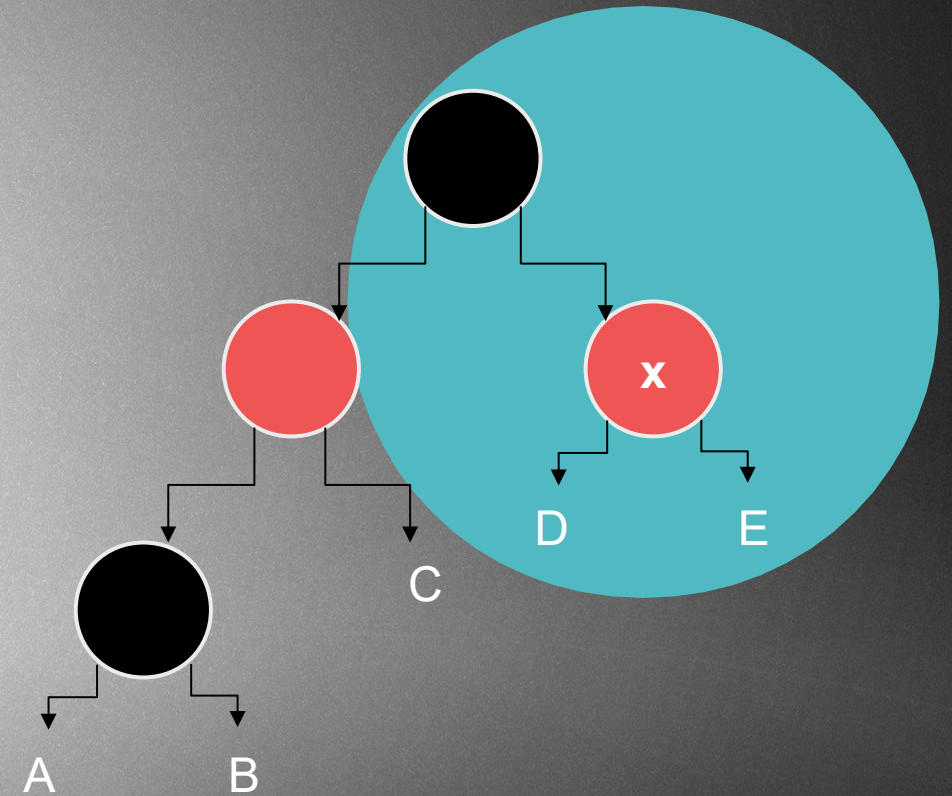
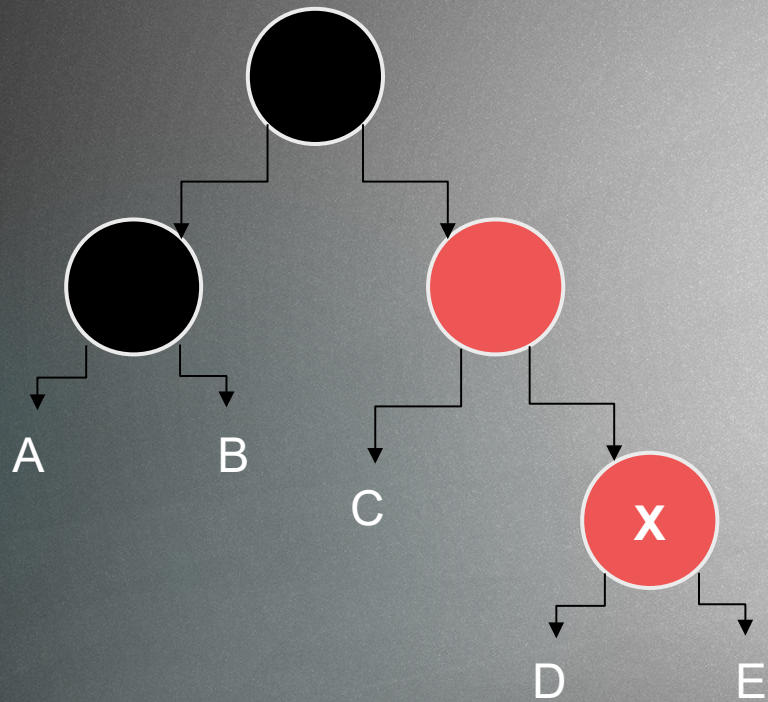
Case 3:

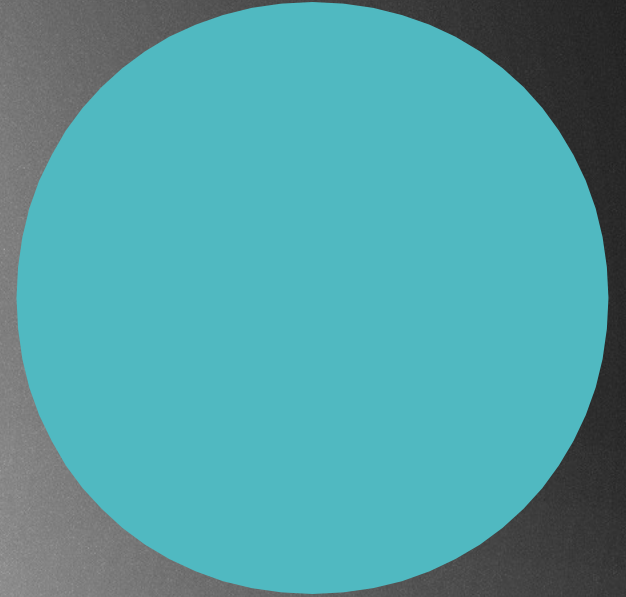
The symmetric version is basically the same just rotate in the opposite direction



Case 3:

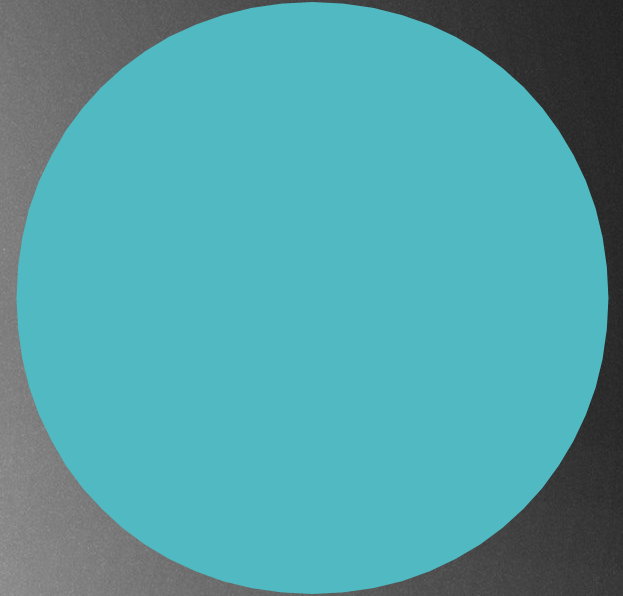
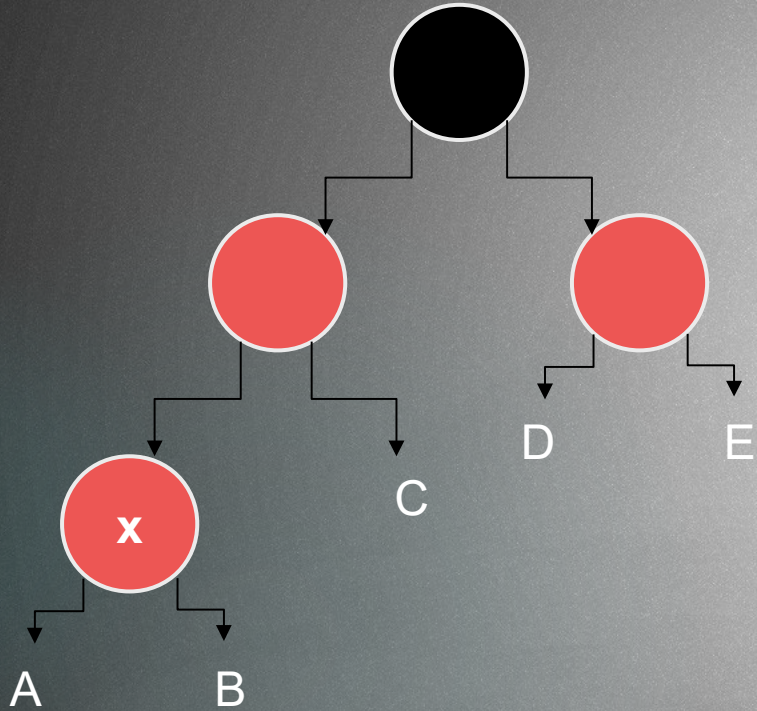
The symmetric version is basically the same just rotate in the opposite direction





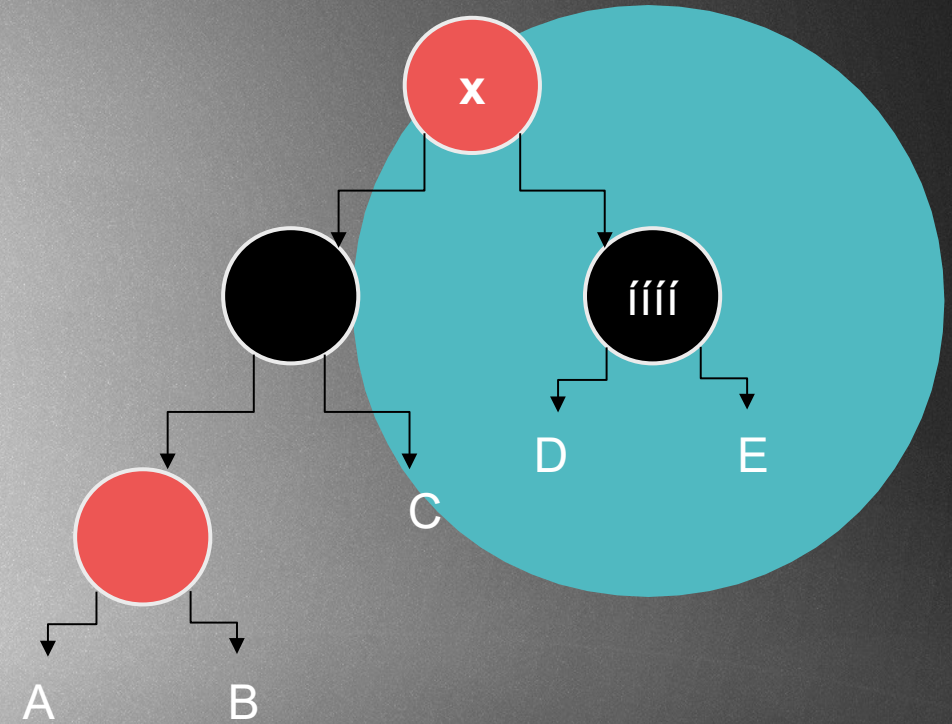
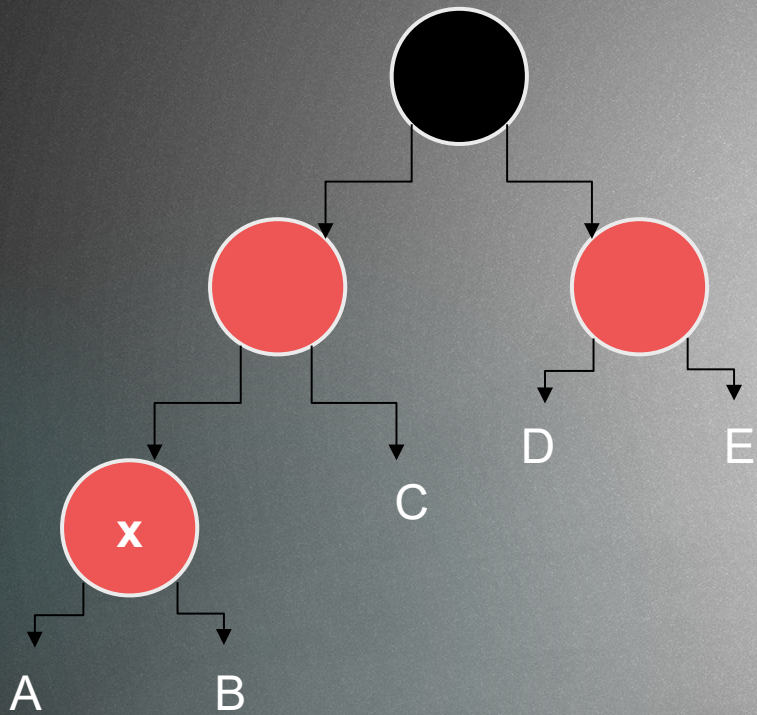
Case 4:

Uncle of node **x** is black and node **x** is a left child
+ uncle is red here in this case
We have to recolor some nodes !!!



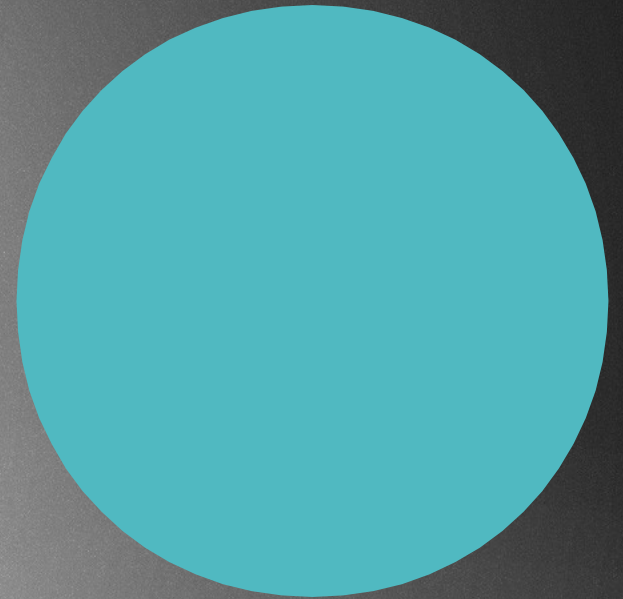
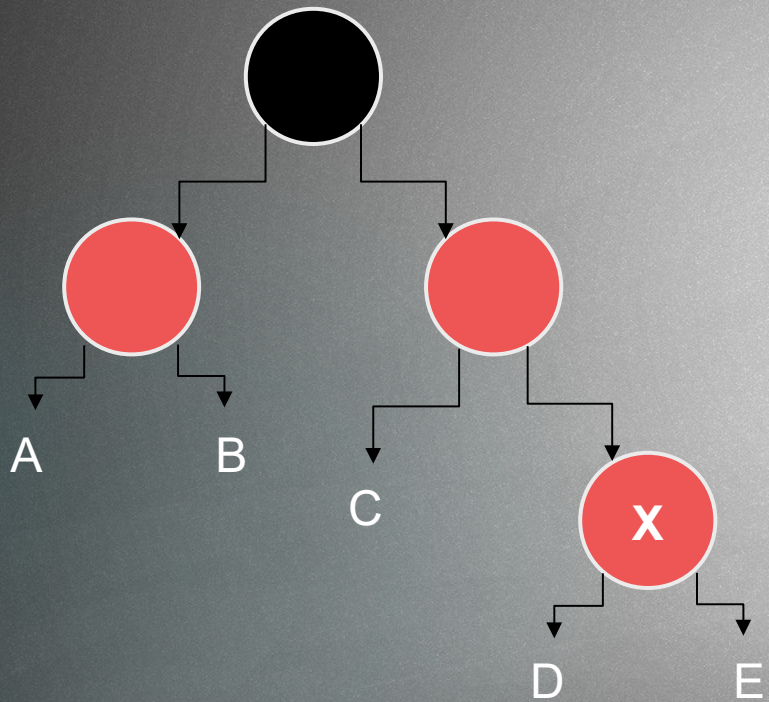
Case 4:

Uncle of node **x** is black and node **x** is a left child
We have to recolor some nodes



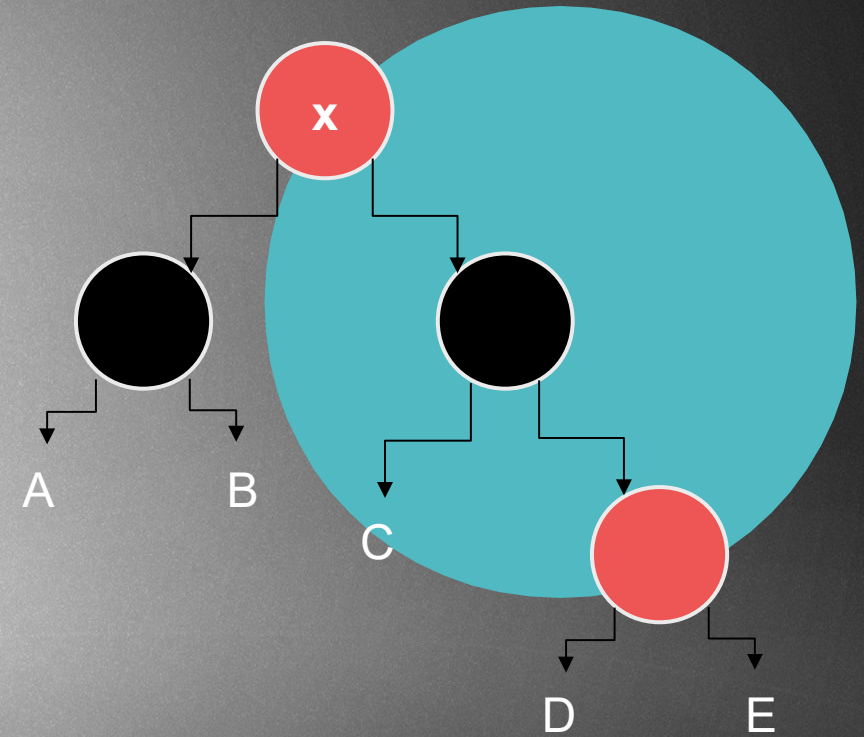
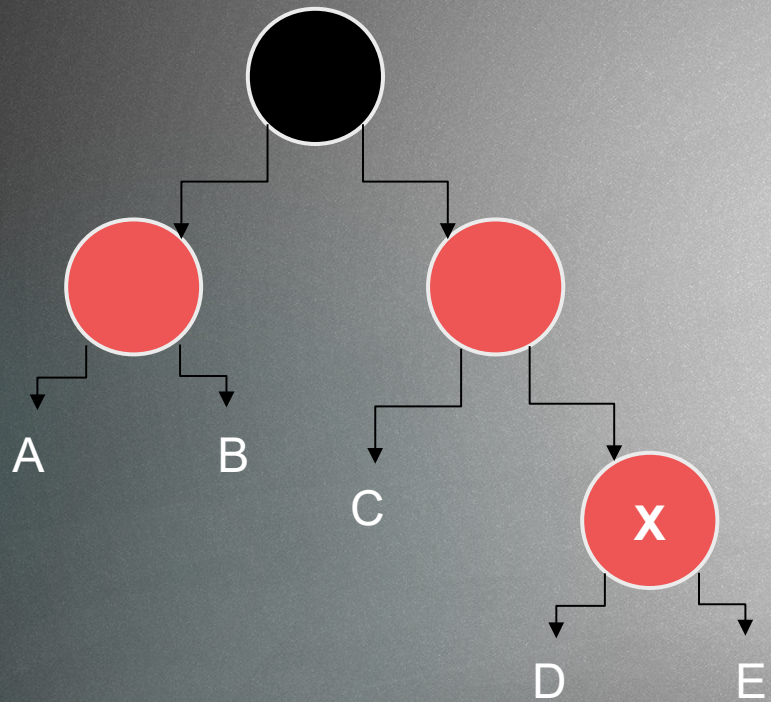
Case 4:

This is the same problem but the symmetric version !!!

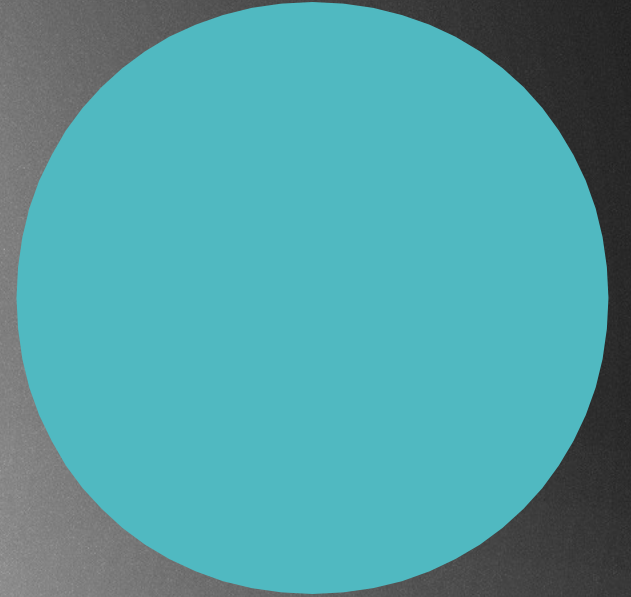


Case 4:

This is the same problem but the symmetric version !!!

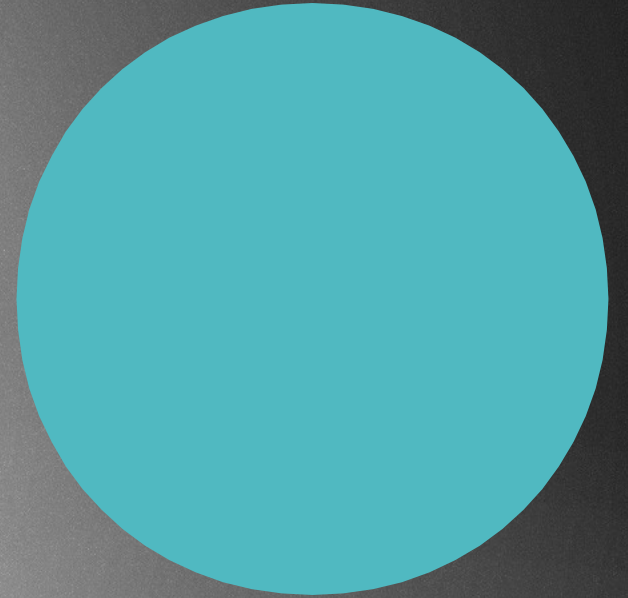


Example:



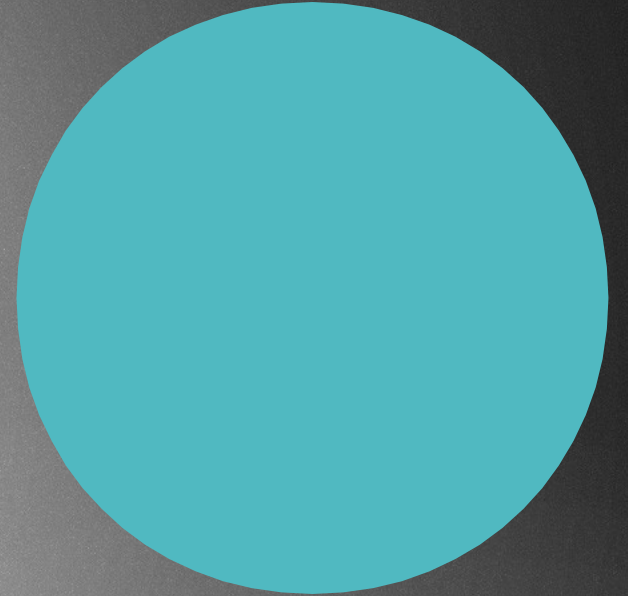
Example:

insert(1)



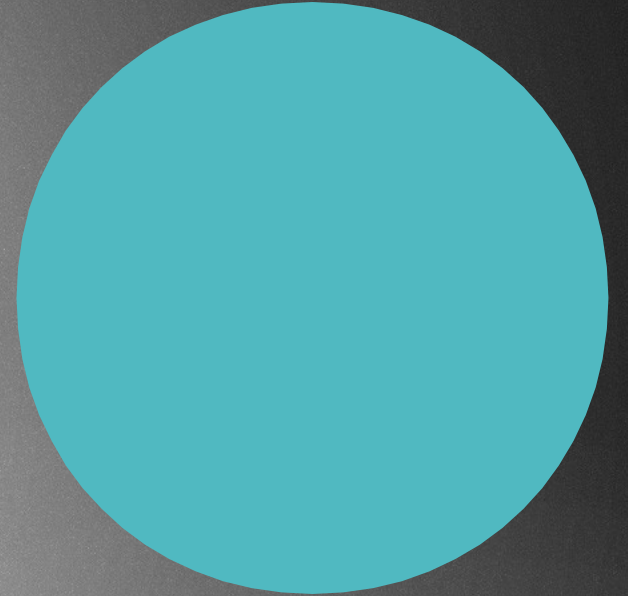
Example:

insert(1)



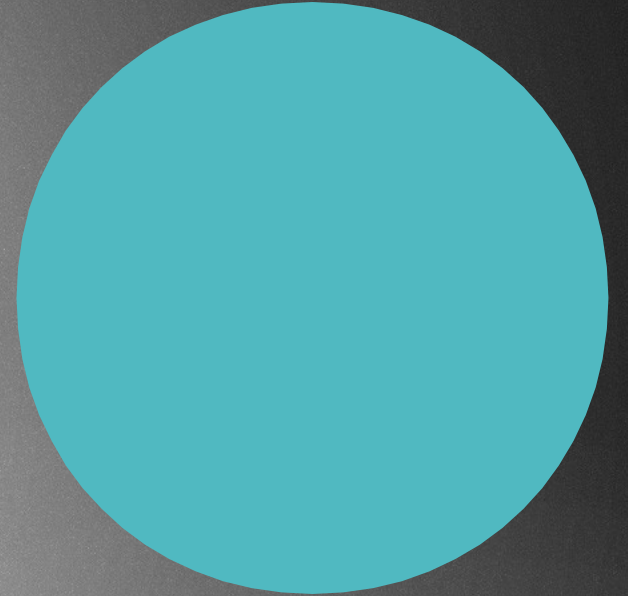
Example:

insert(1)



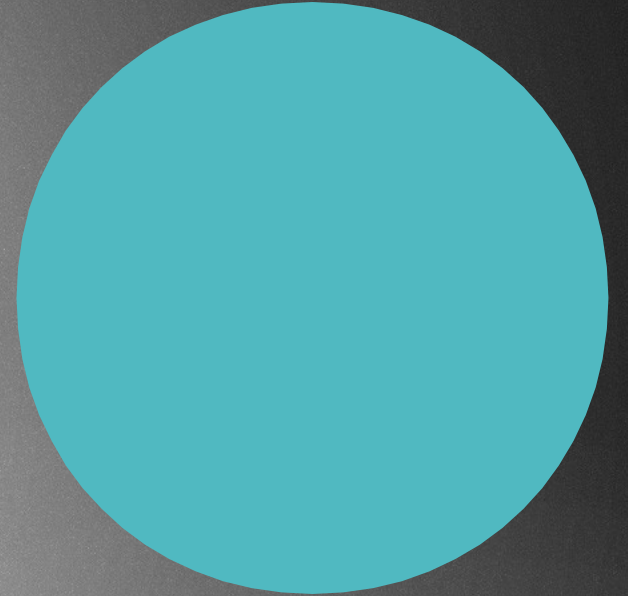
Example:

insert(2)



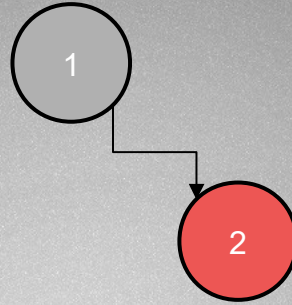
Example:

insert(2)



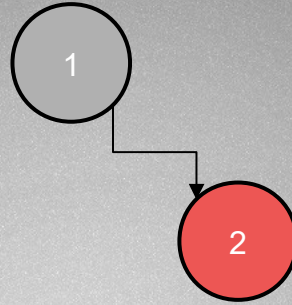
Example:

insert(2)



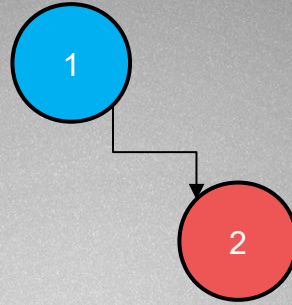
Example:

insert(3)



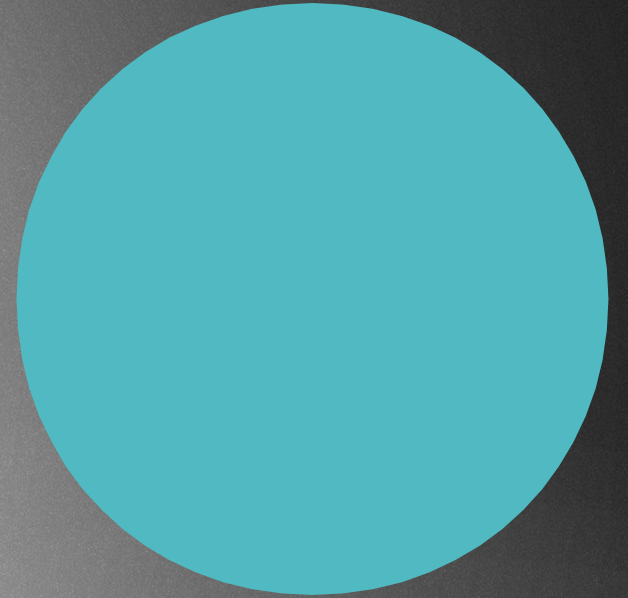
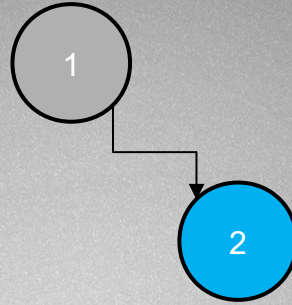
Example:

insert(3)

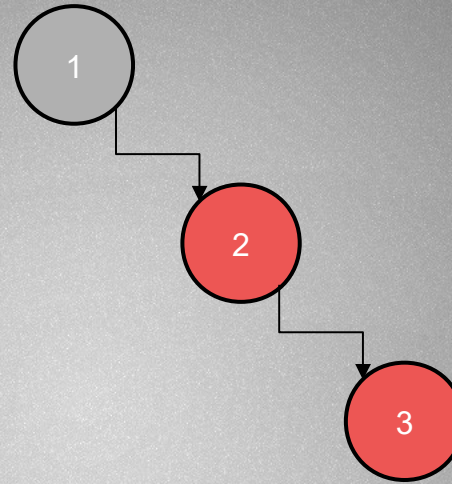


Example:

insert(3)



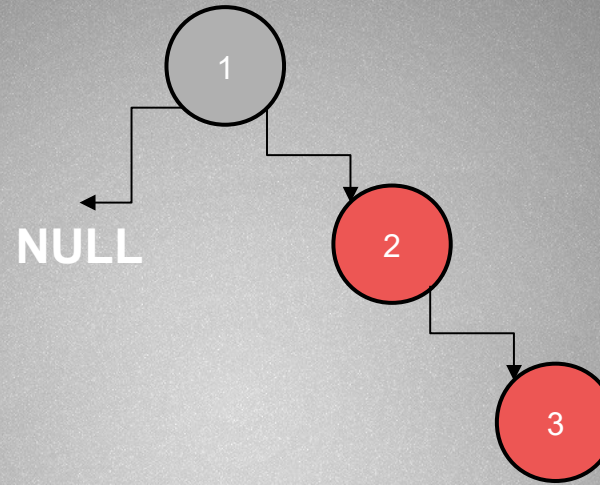
Example:



It is the Case 3: because the NULL is considered to be a black node

- the uncle of Node 3 is black
- Have to make a rotation + recolor the nodes if necessary

Example:

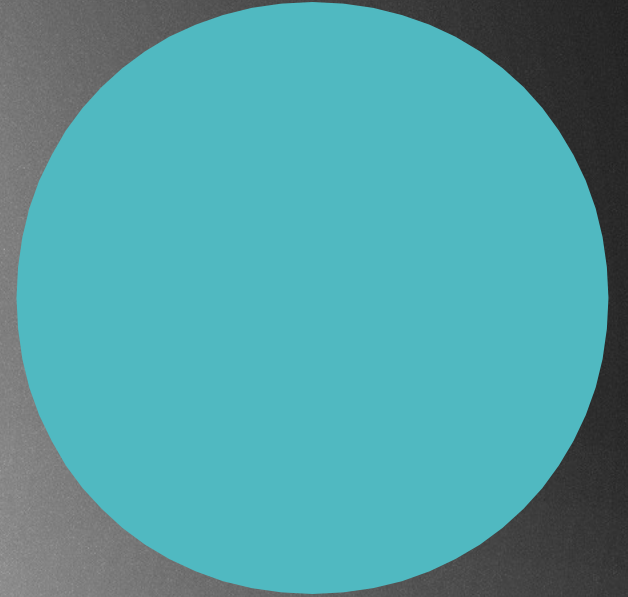
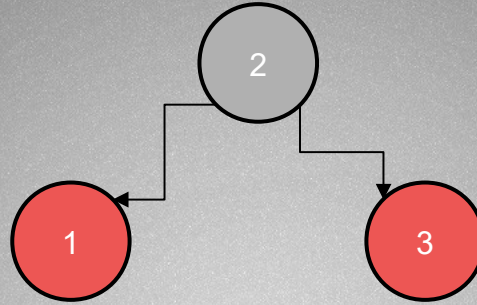


It is the Case 3: because the NULL is considered to be a black node

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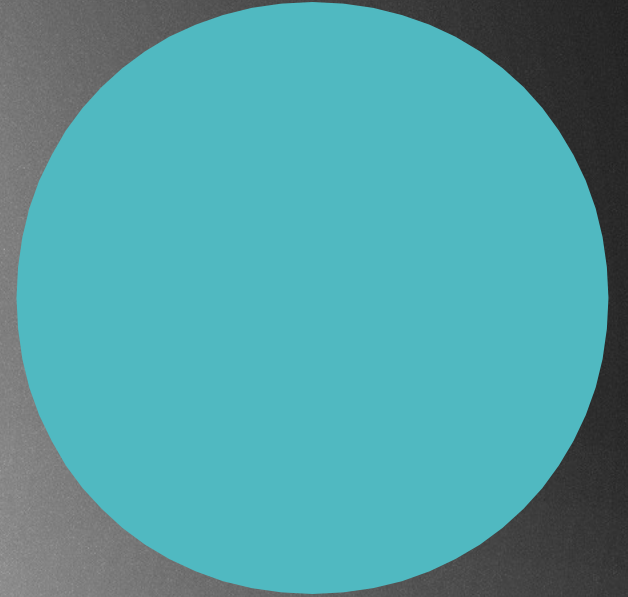
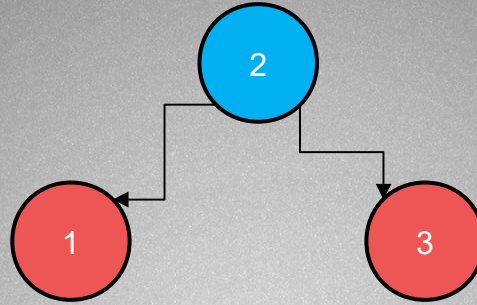
Example:

insert(4)



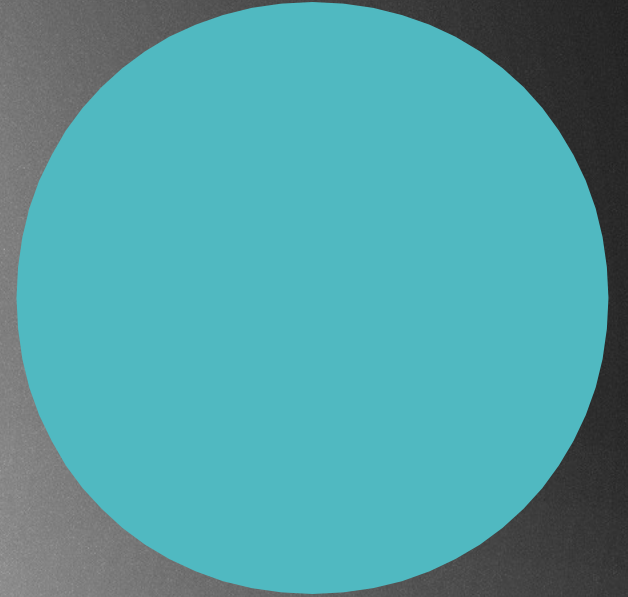
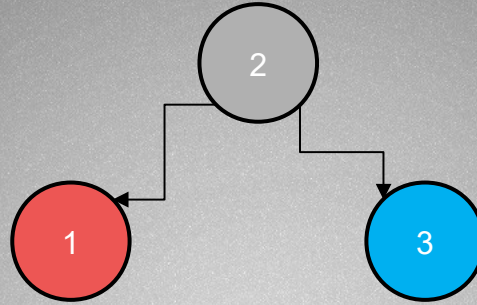
Example:

insert(4)

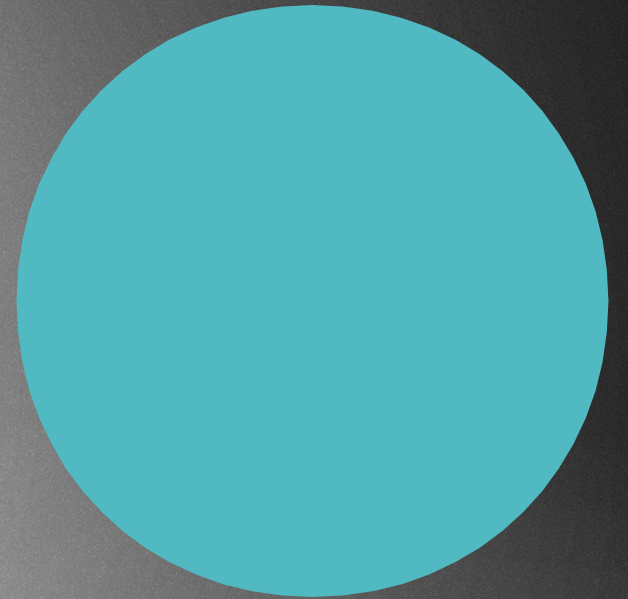
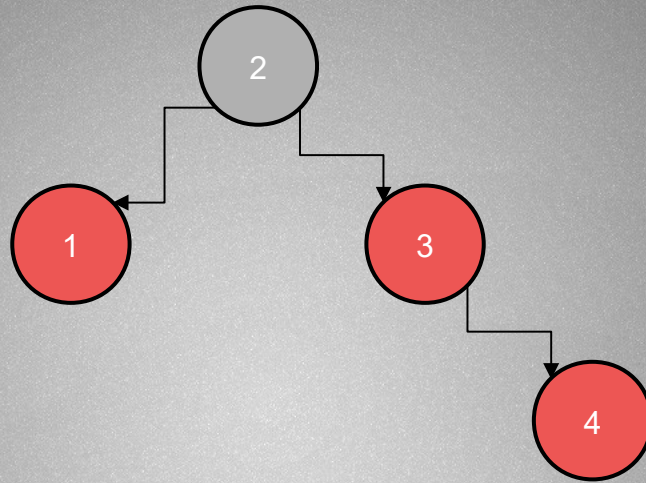


Example:

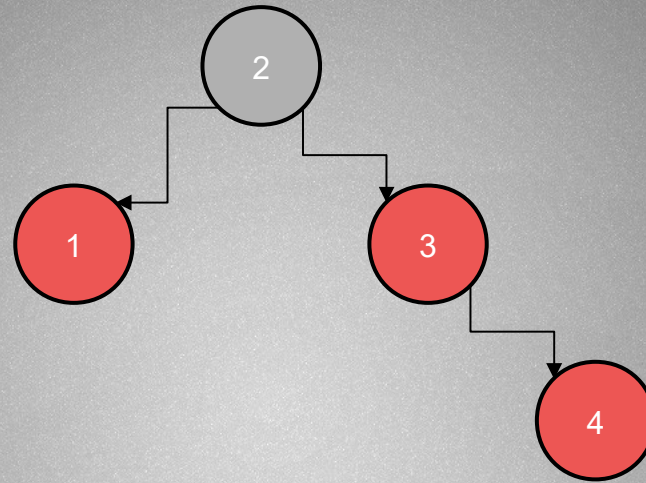
insert(4)



Example:



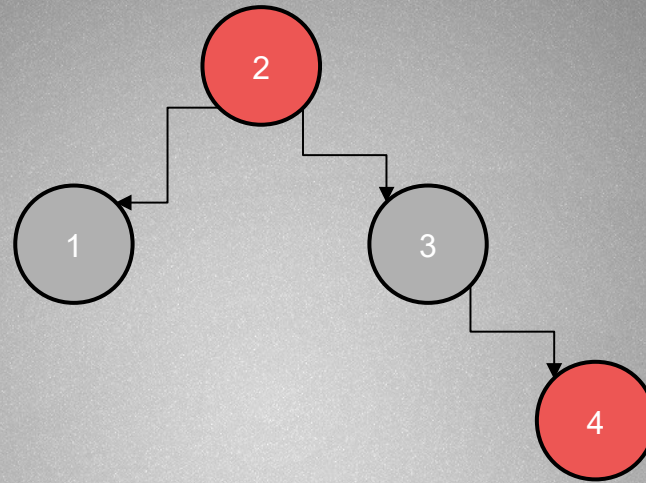
Example:



It is the Case 1: the given node 4 and the parent are both red + uncle is red
Color uncle + parent to be black



Example:

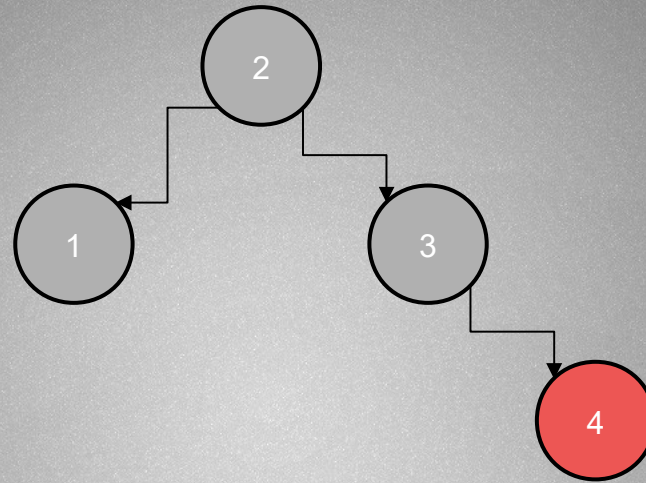


It is the Case 1: the given node 4 and the parent are both red + uncle is red

Color uncle + parent to be black



Example:



It is the Case 1: the given node 4 and the parent are both red + uncle is red

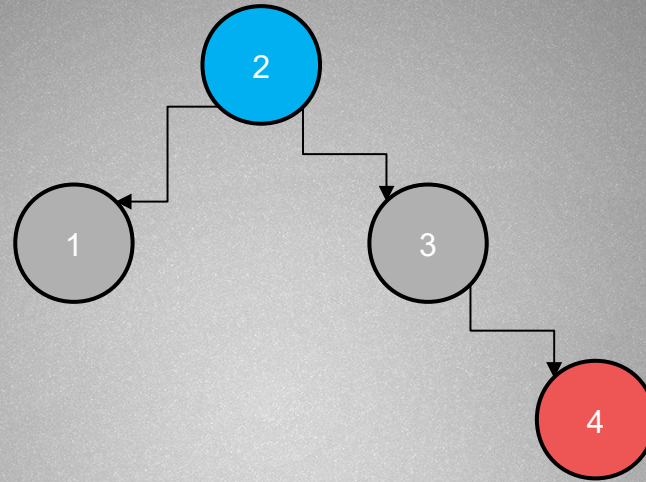
Color uncle + parent to be black

VIOLATES THE PROPERTIES AGAIN

Root has to be black as well

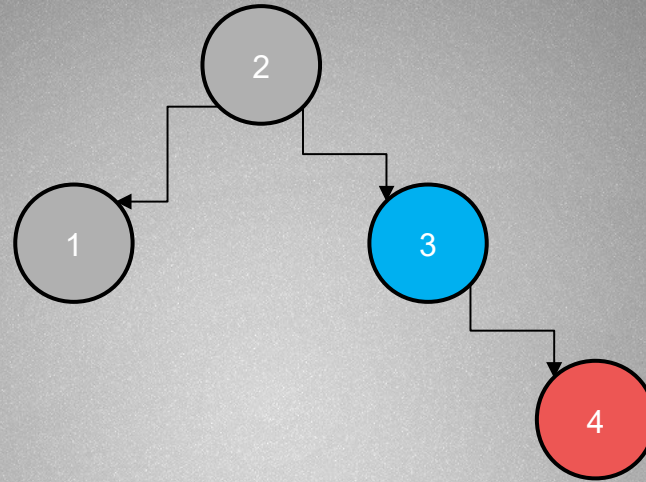
Example:

insert(5)



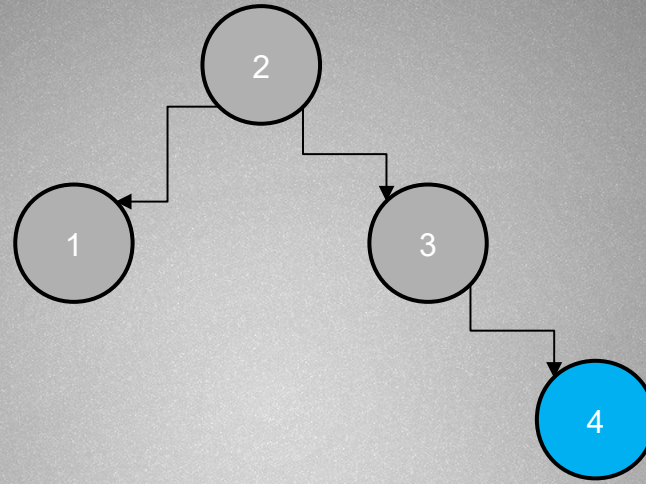
Example:

insert(5)



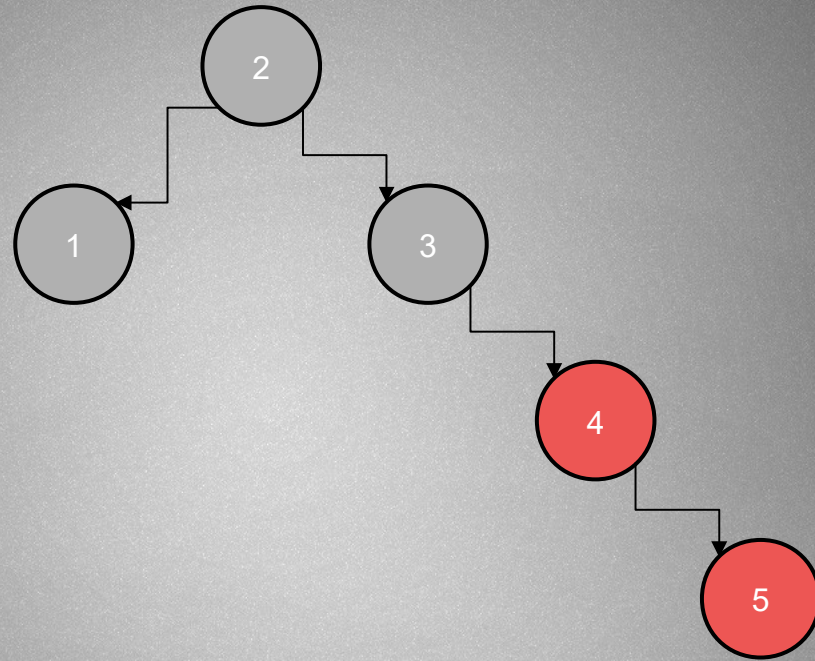
Example:

insert(5)



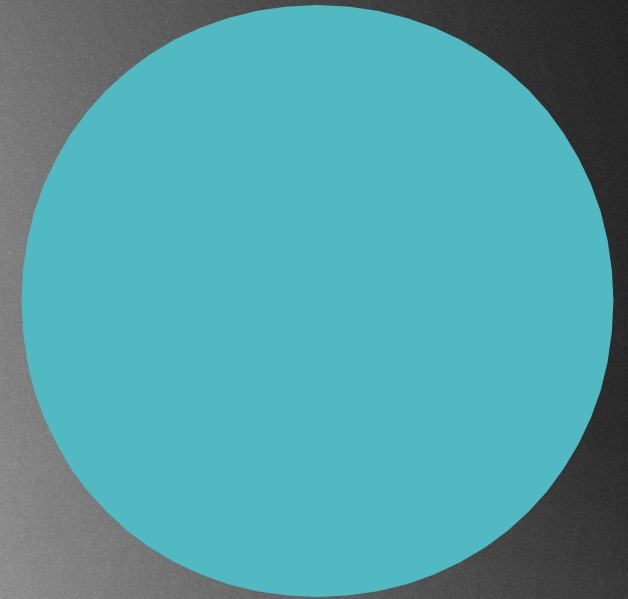
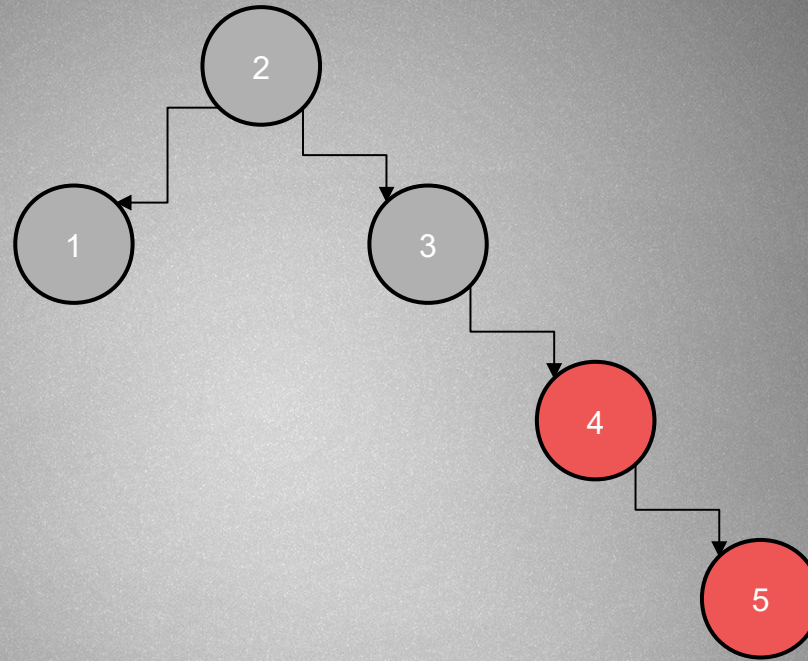
Example:

insert(5)

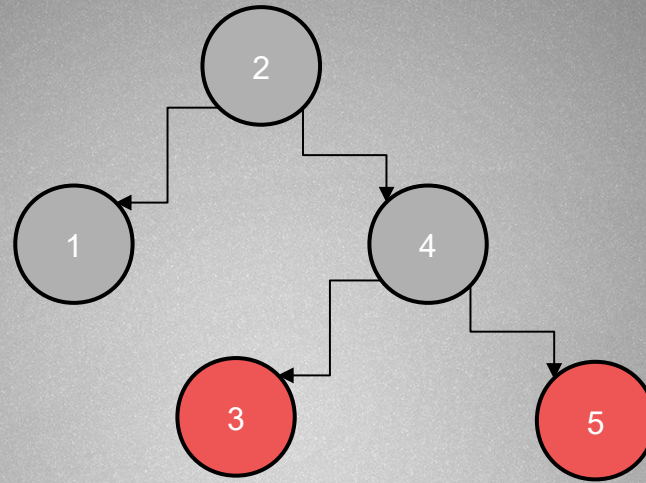


Example:

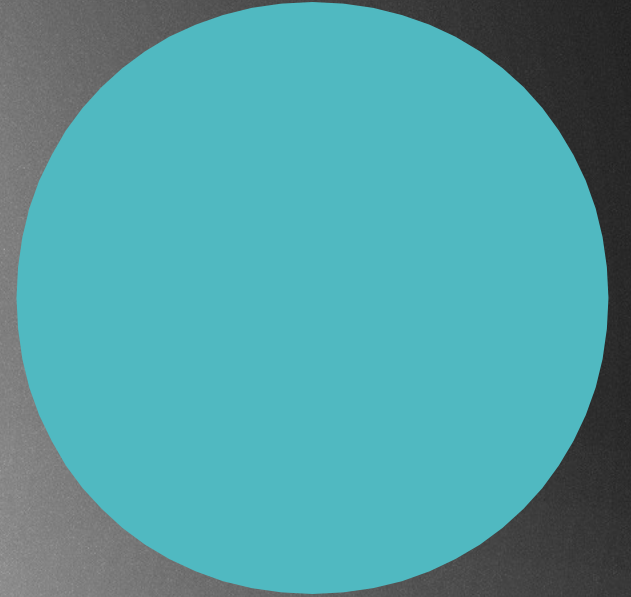
This is the Case 3 situation, we have to make some rotations !!!



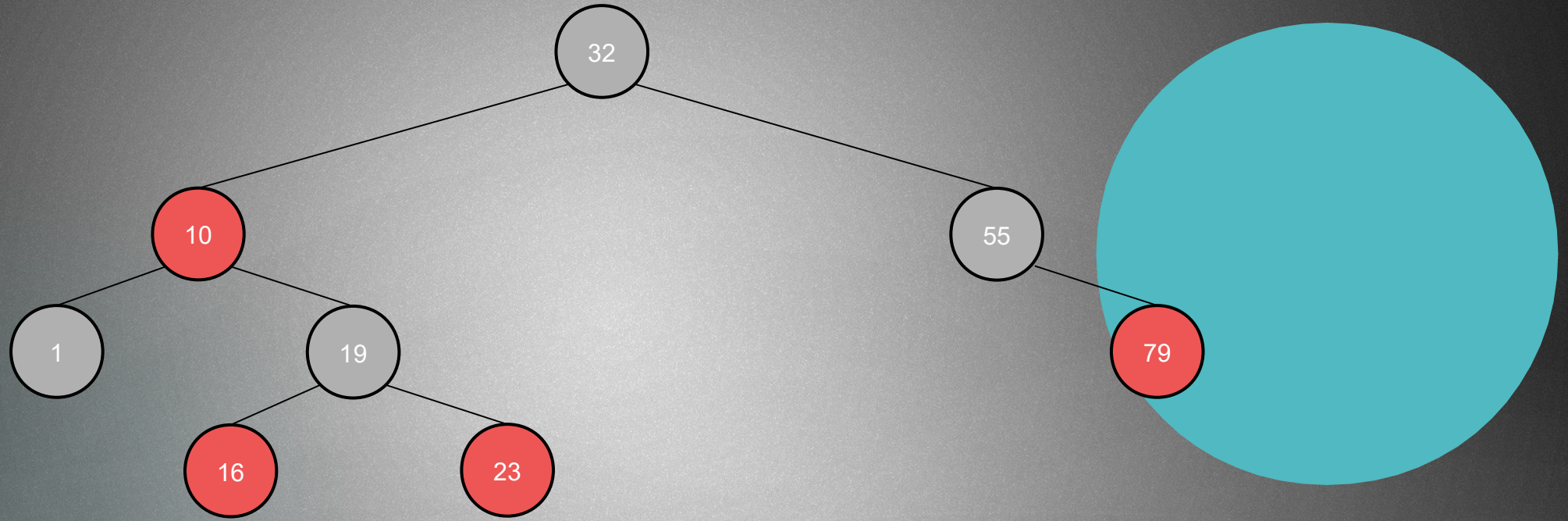
Example:



Example:

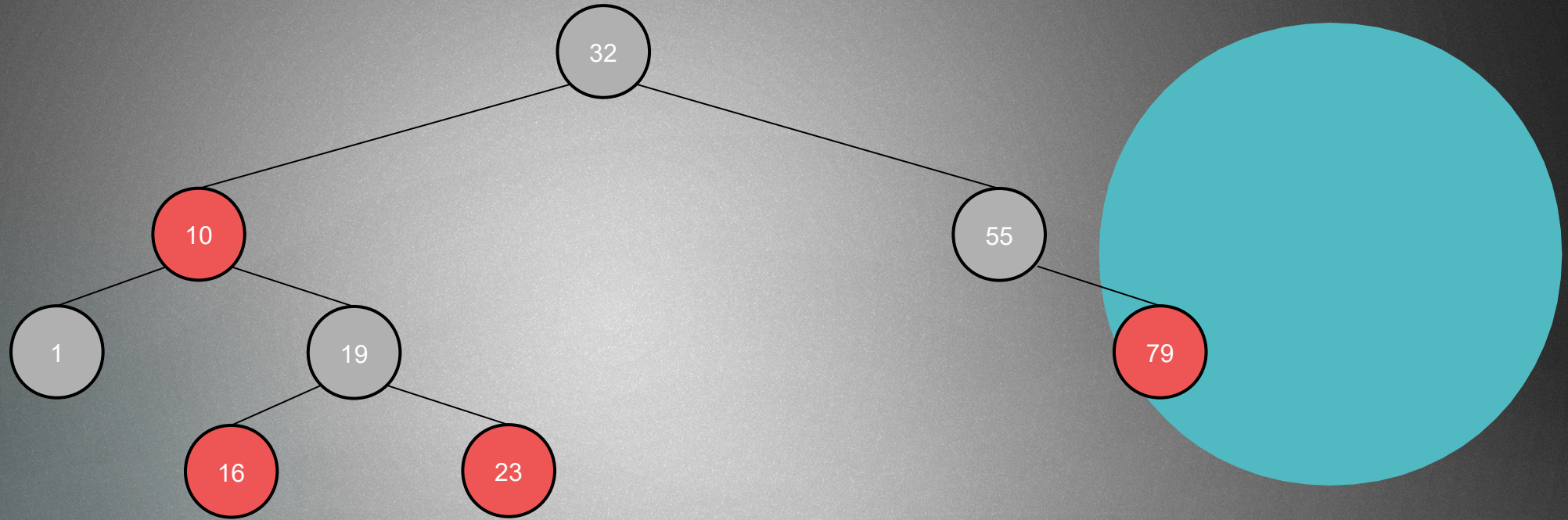


Example:



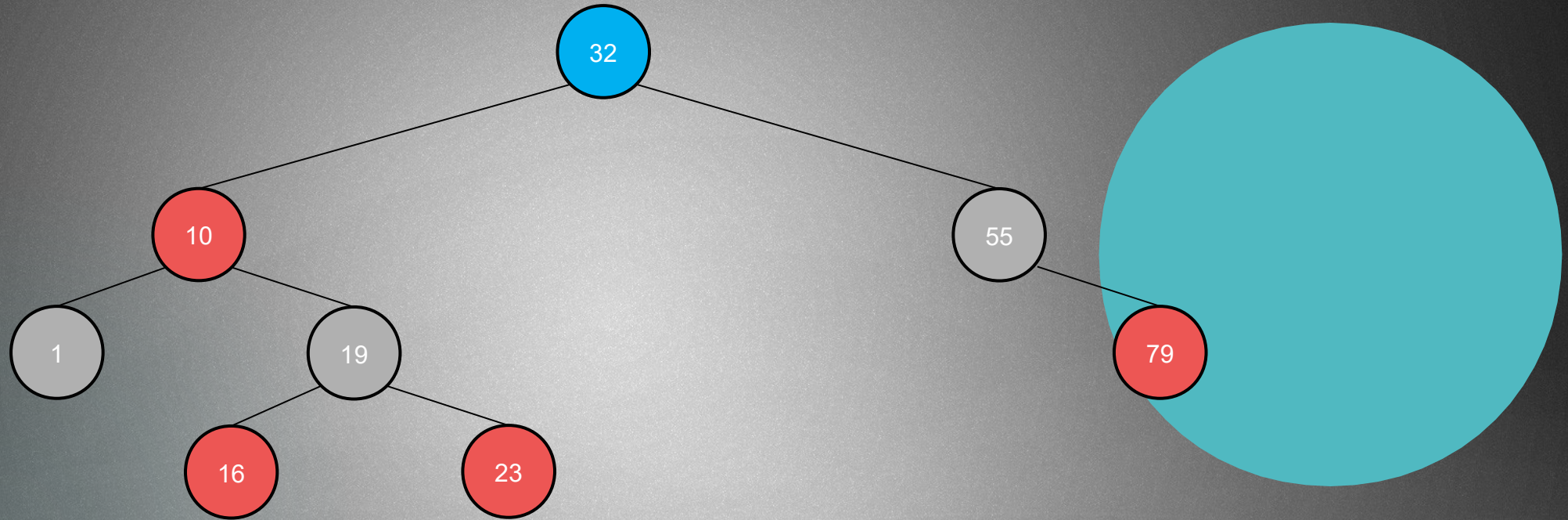
Example:

We want to insert 12



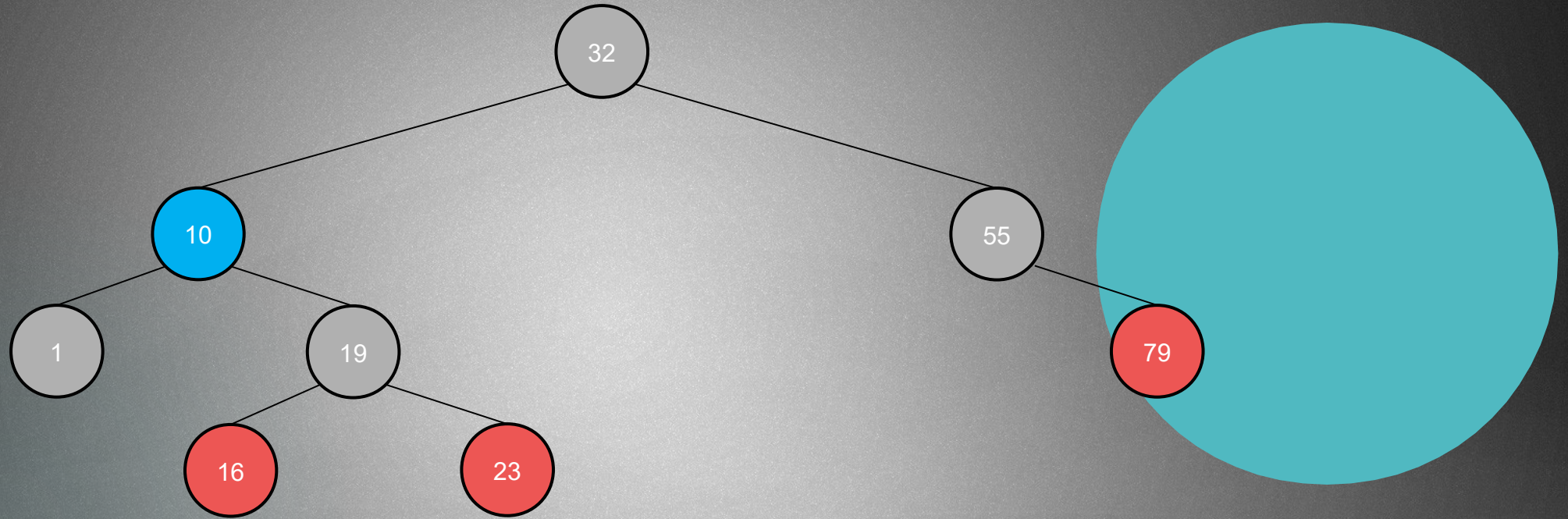
Example:

We want to insert 12



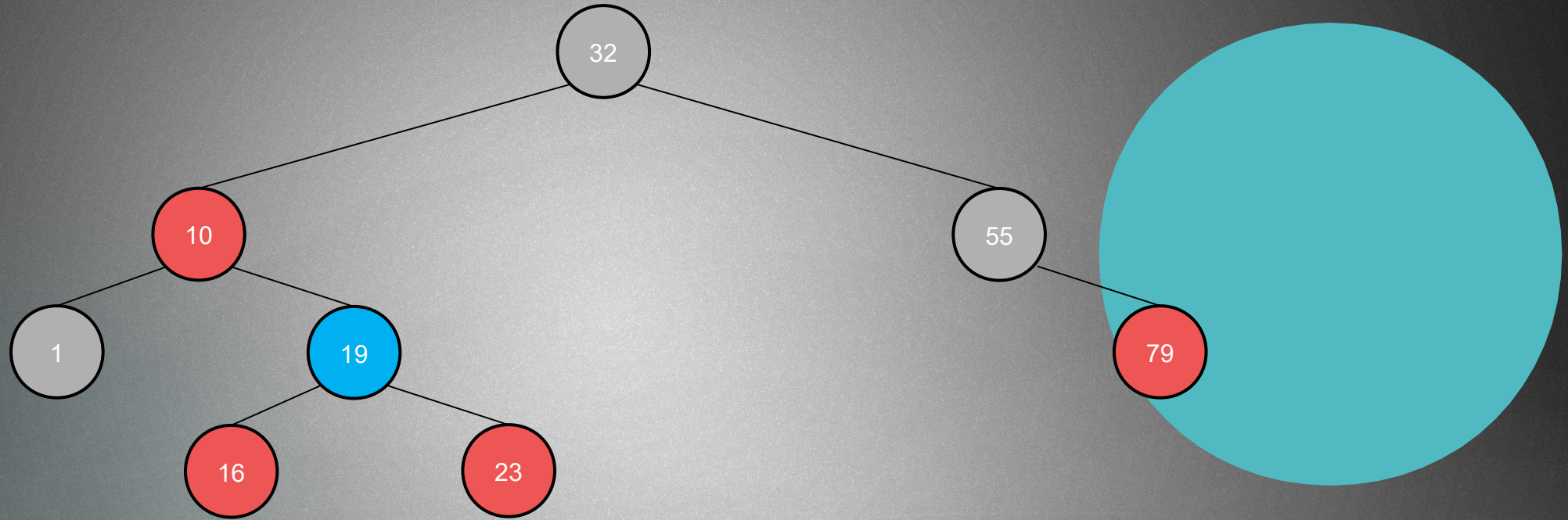
Example:

We want to insert 12



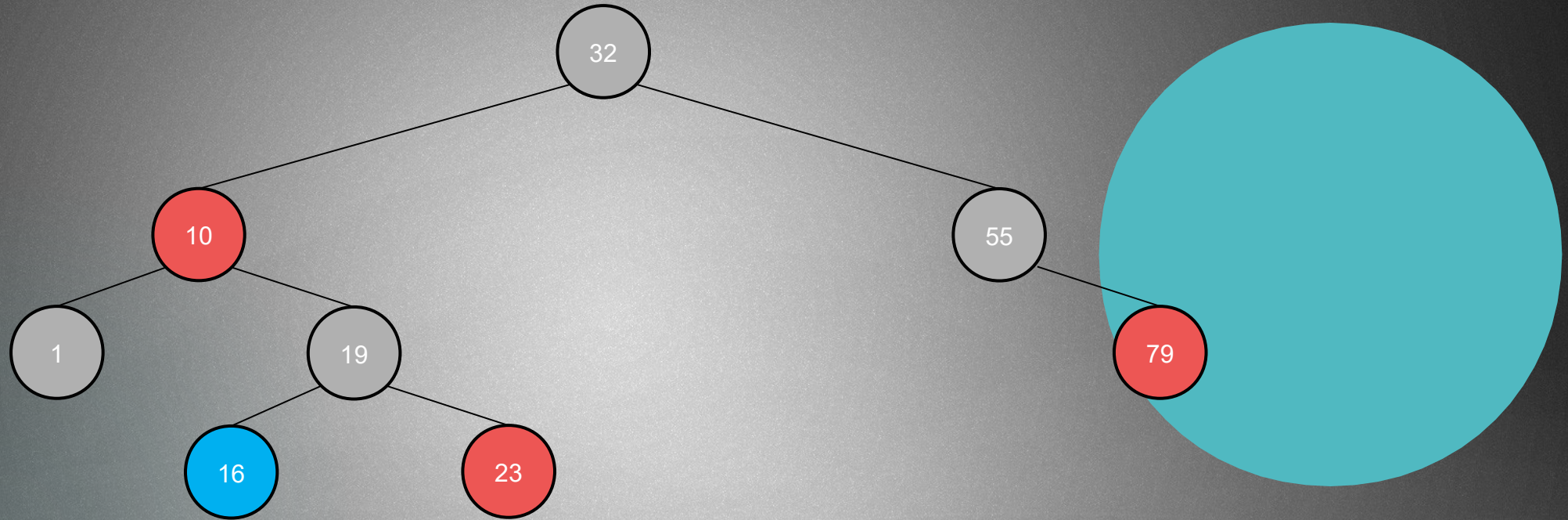
Example:

We want to insert 12

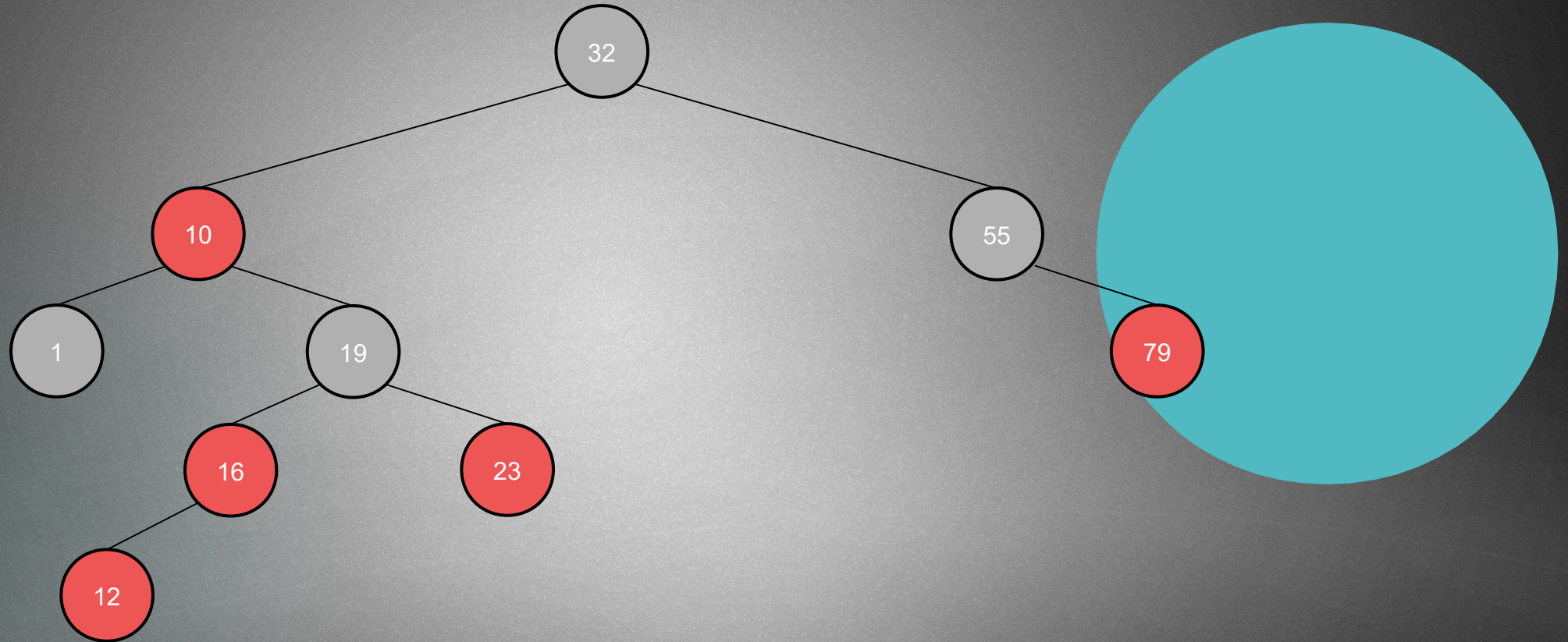


Example:

We want to insert 12

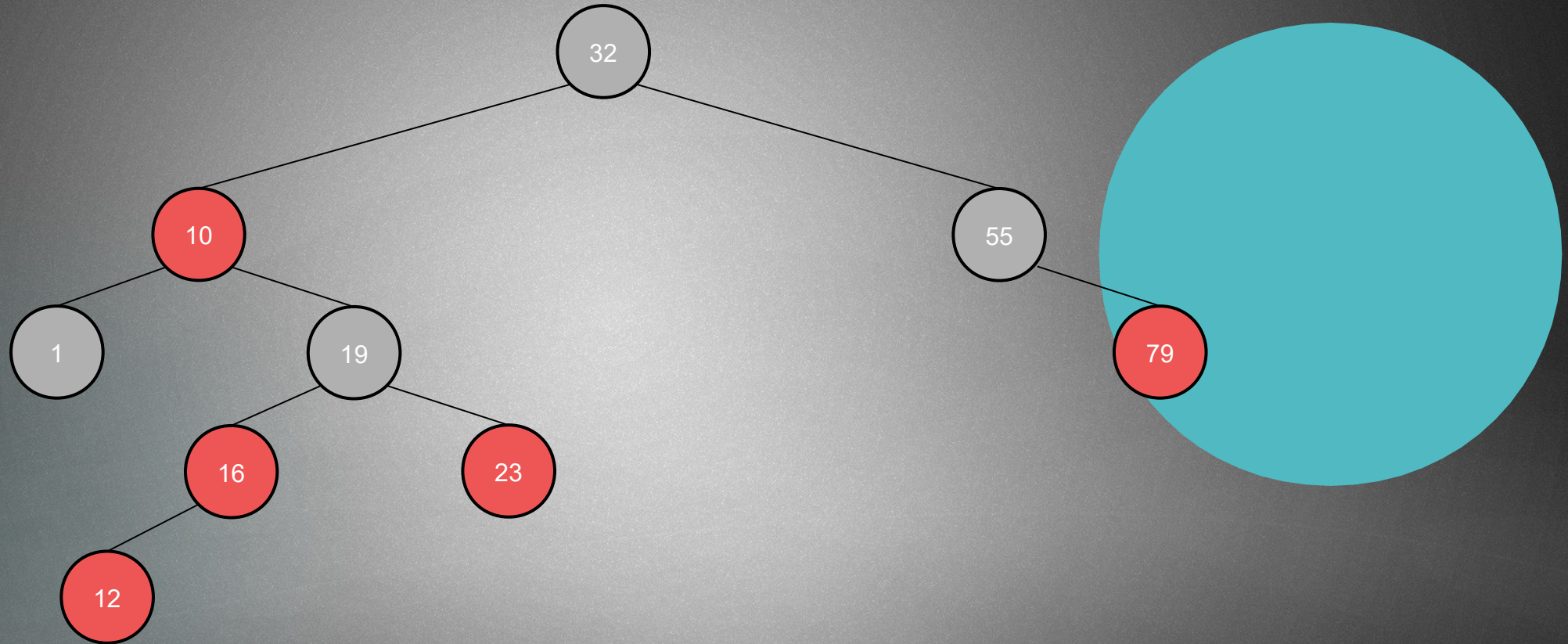


Example:



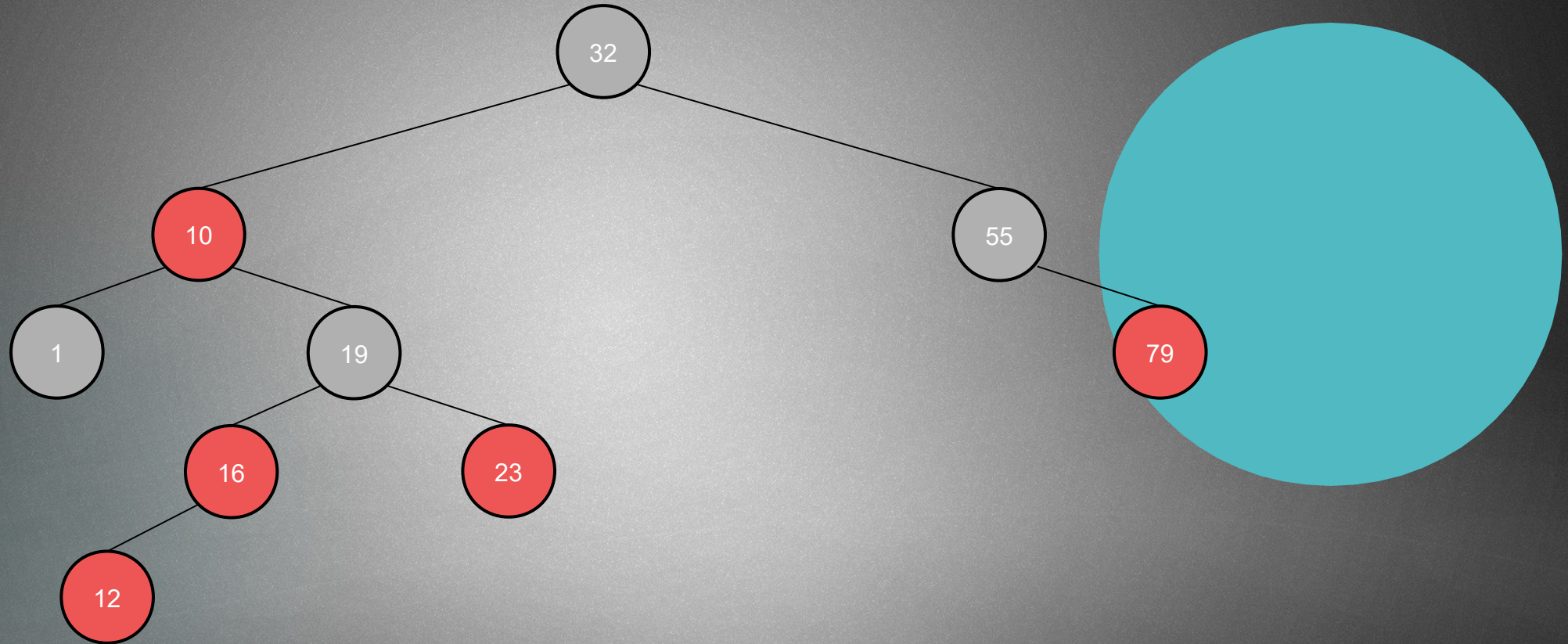
Example:

On every insertion we have to check whether the red black tree properties are violated or not !!!



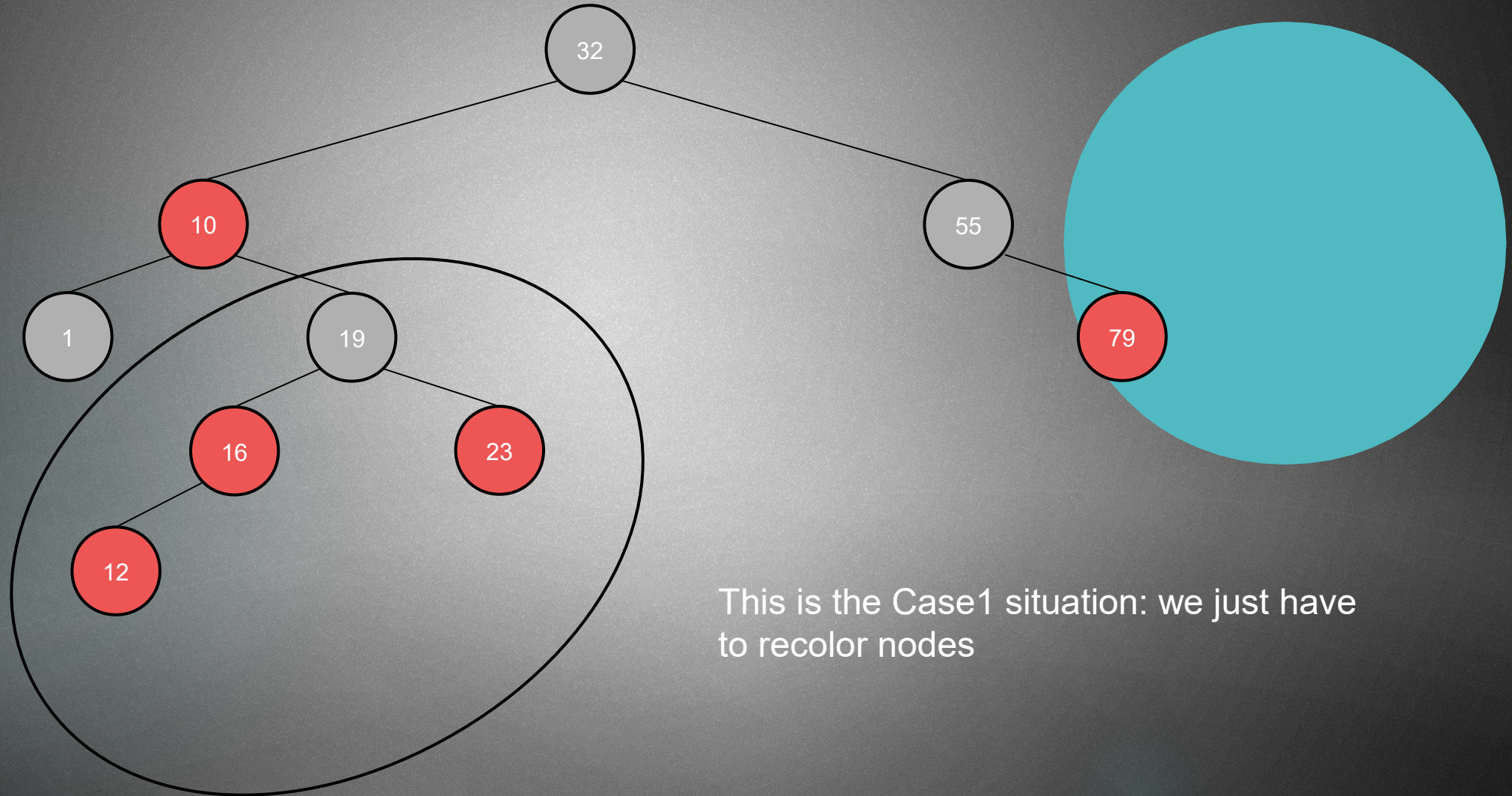
Example:

On every insertion we have to check whether the red black tree properties are violated or not !!!

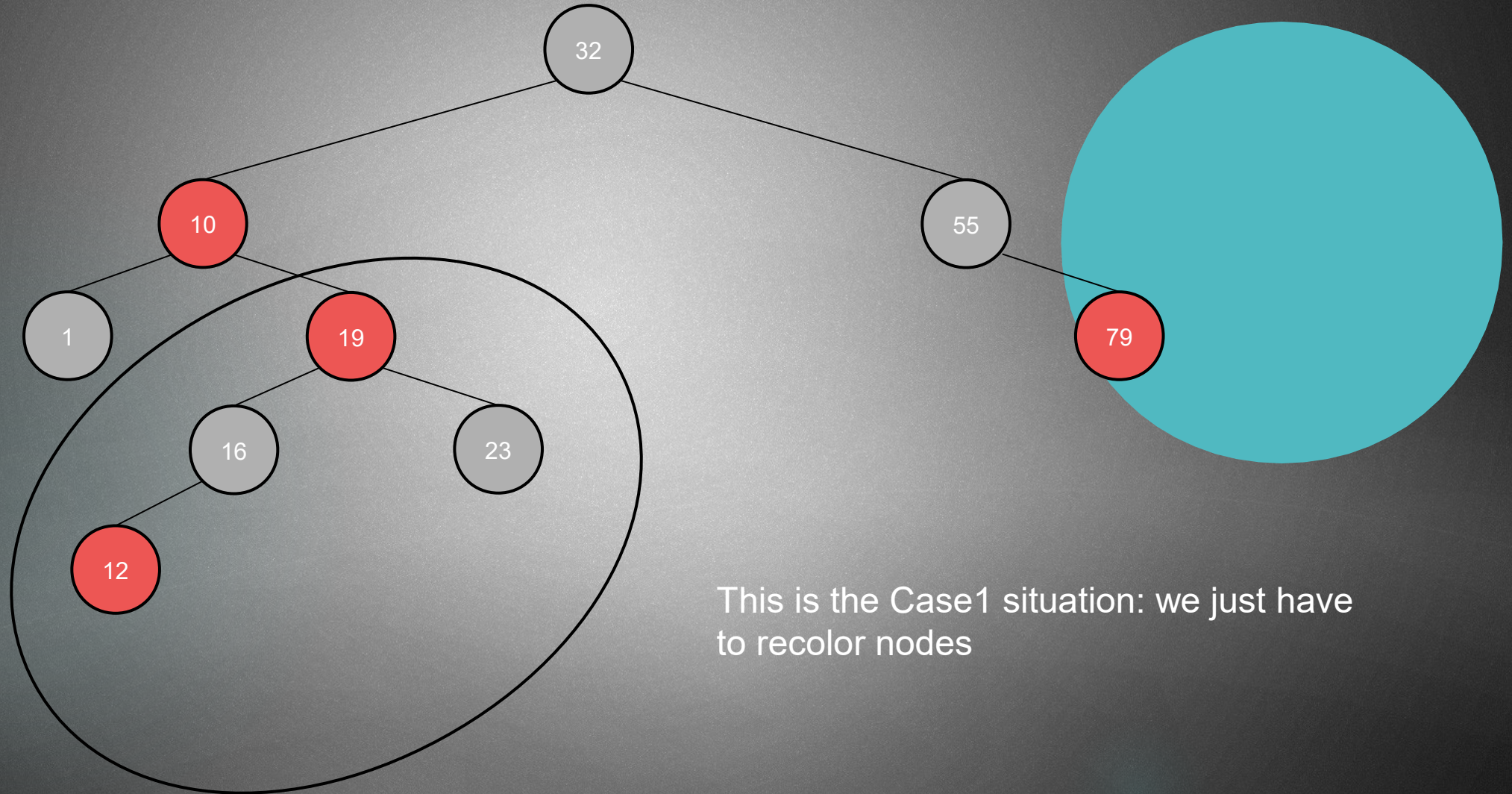


It is violated because red node has a single red child

Example:

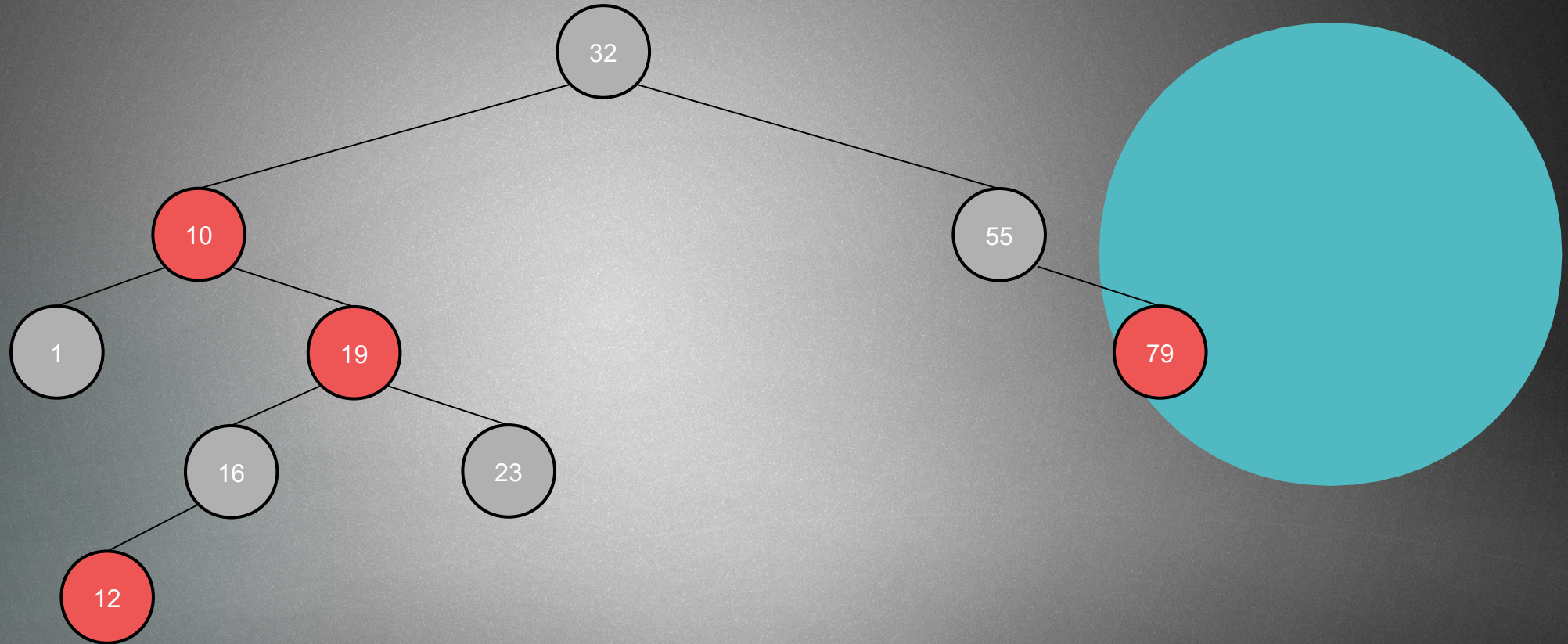


Example:

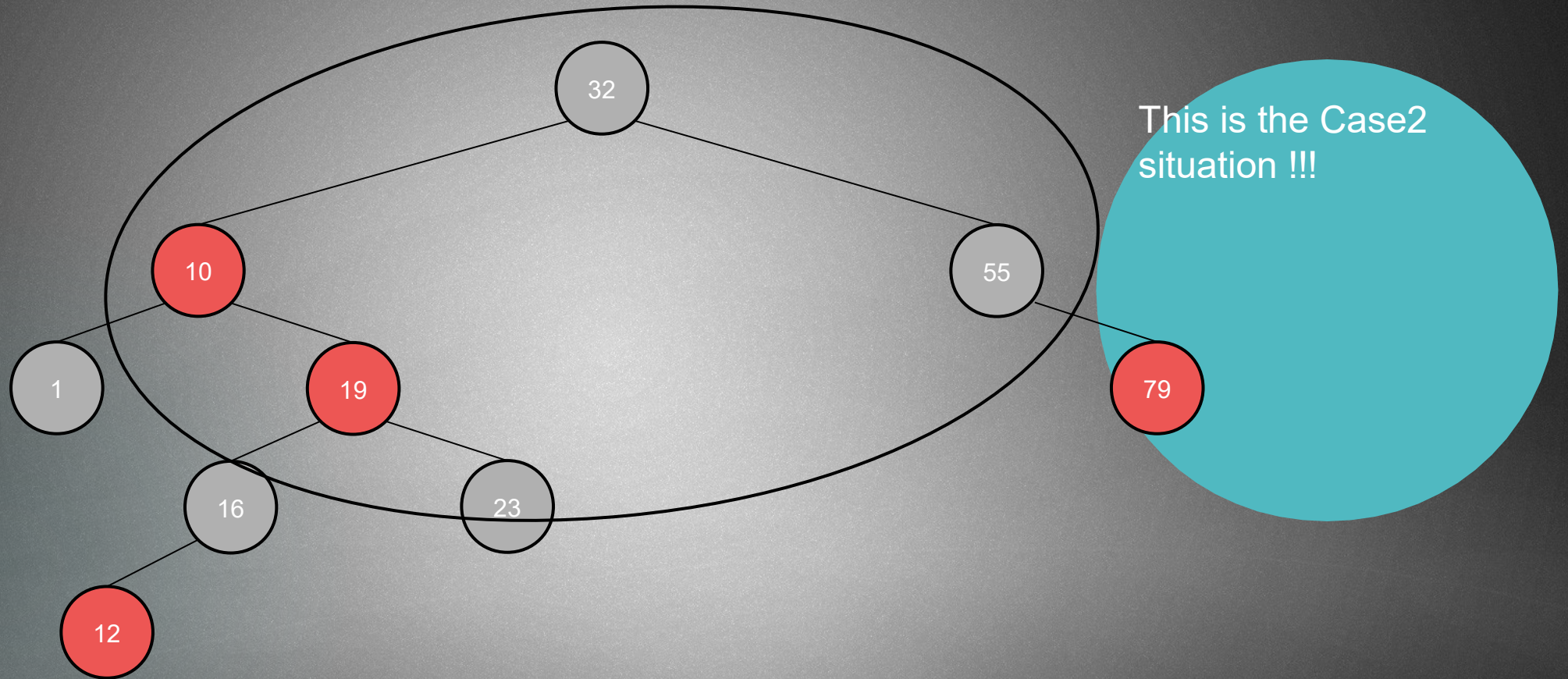


Example:

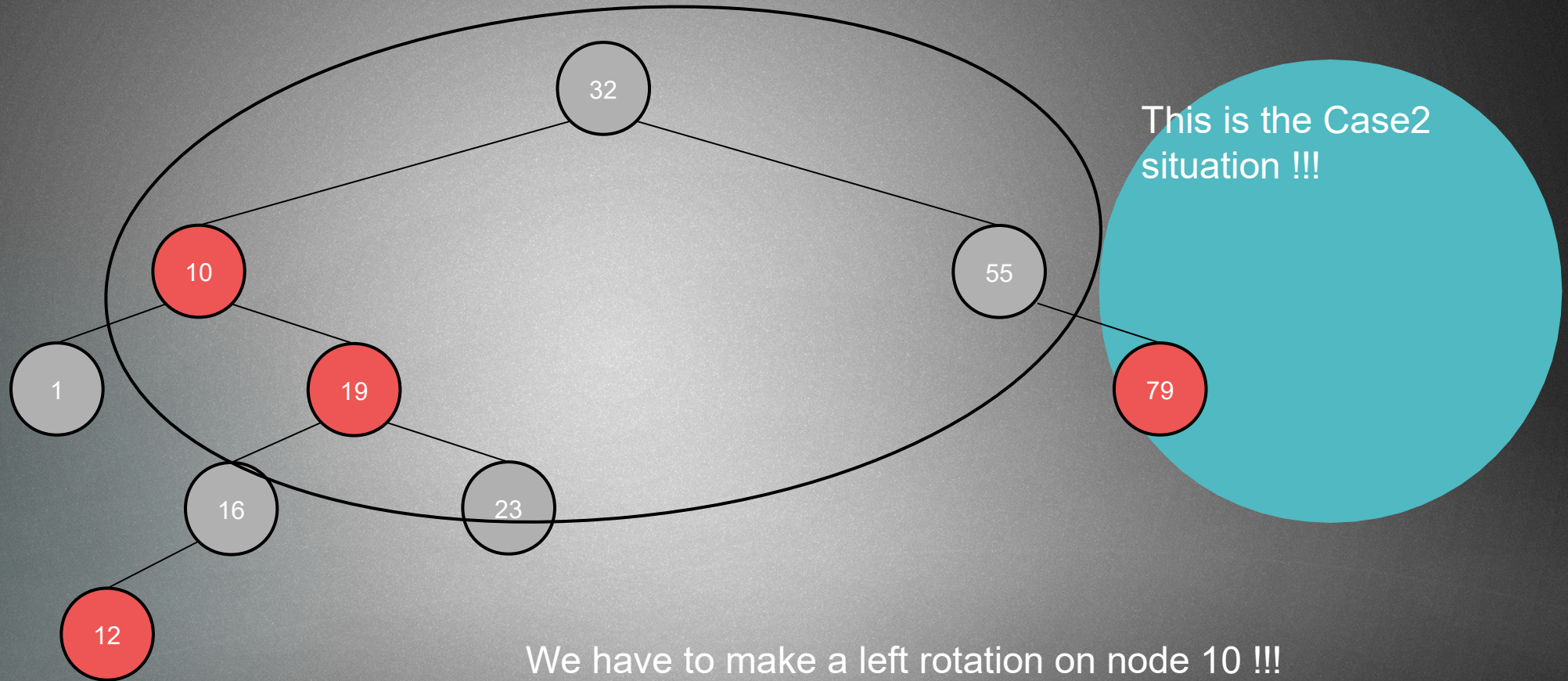
We have to check whether
the red black tree properties are violated or not !!!



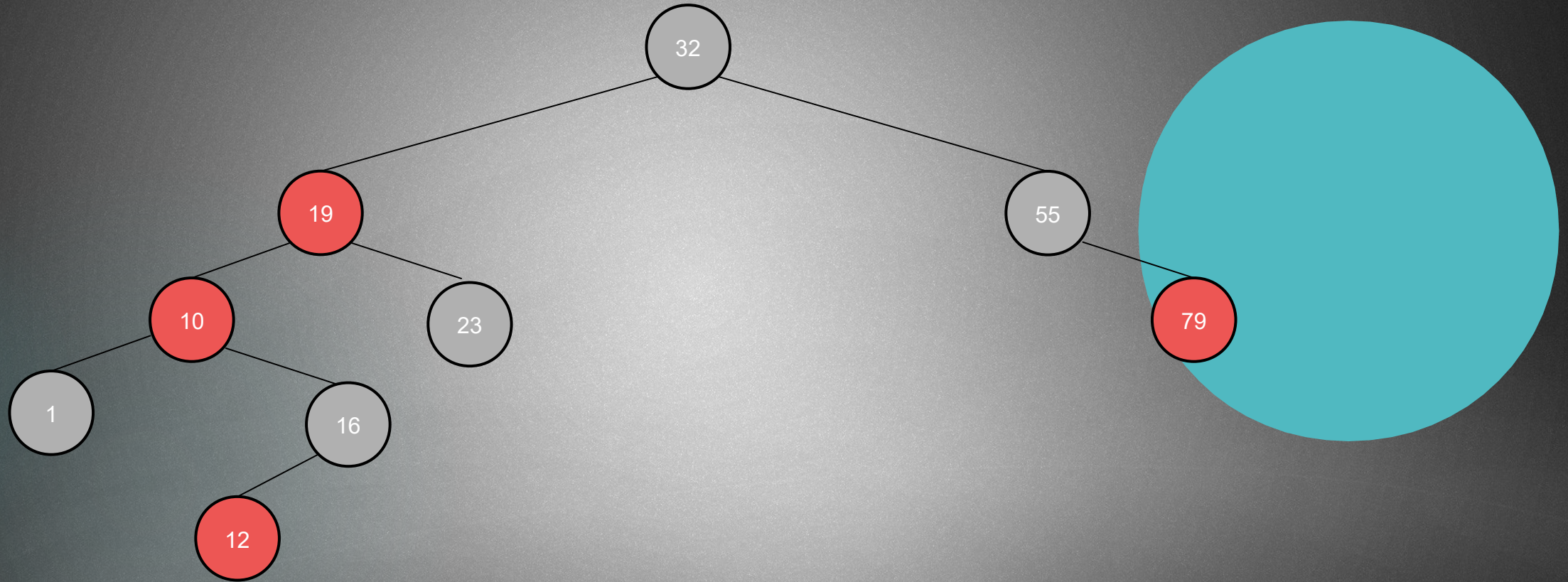
Example:



Example:

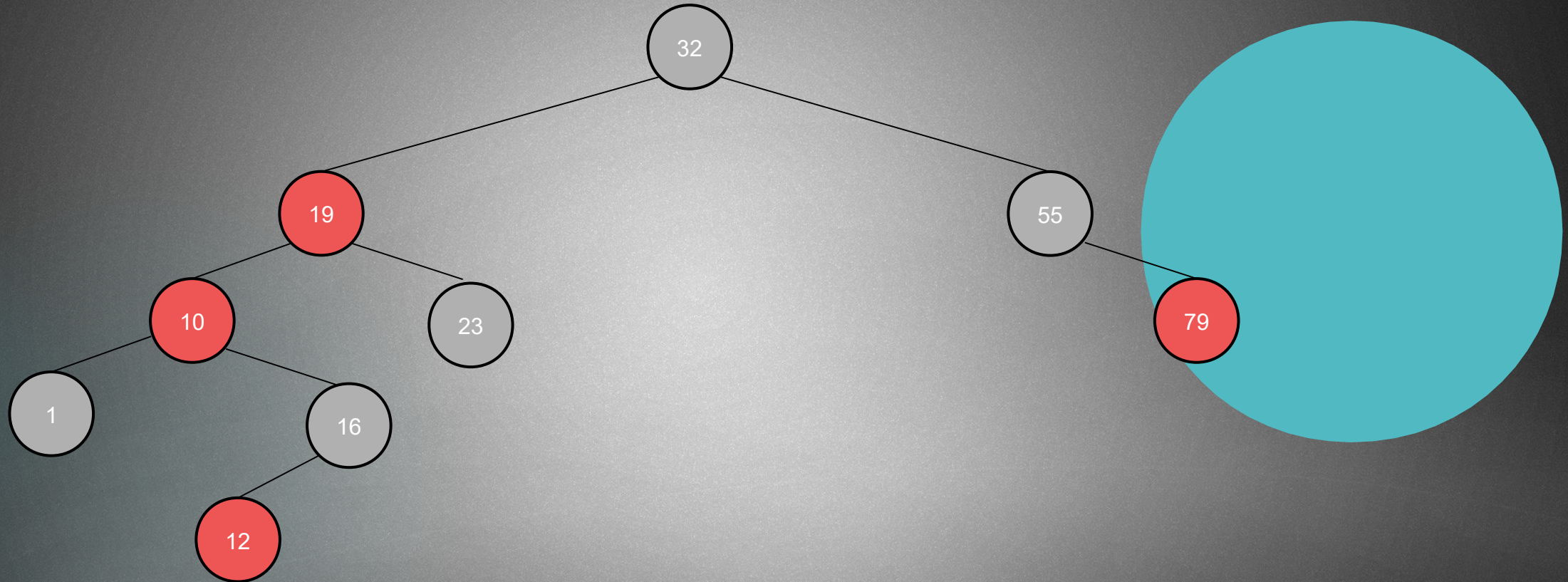


Example:



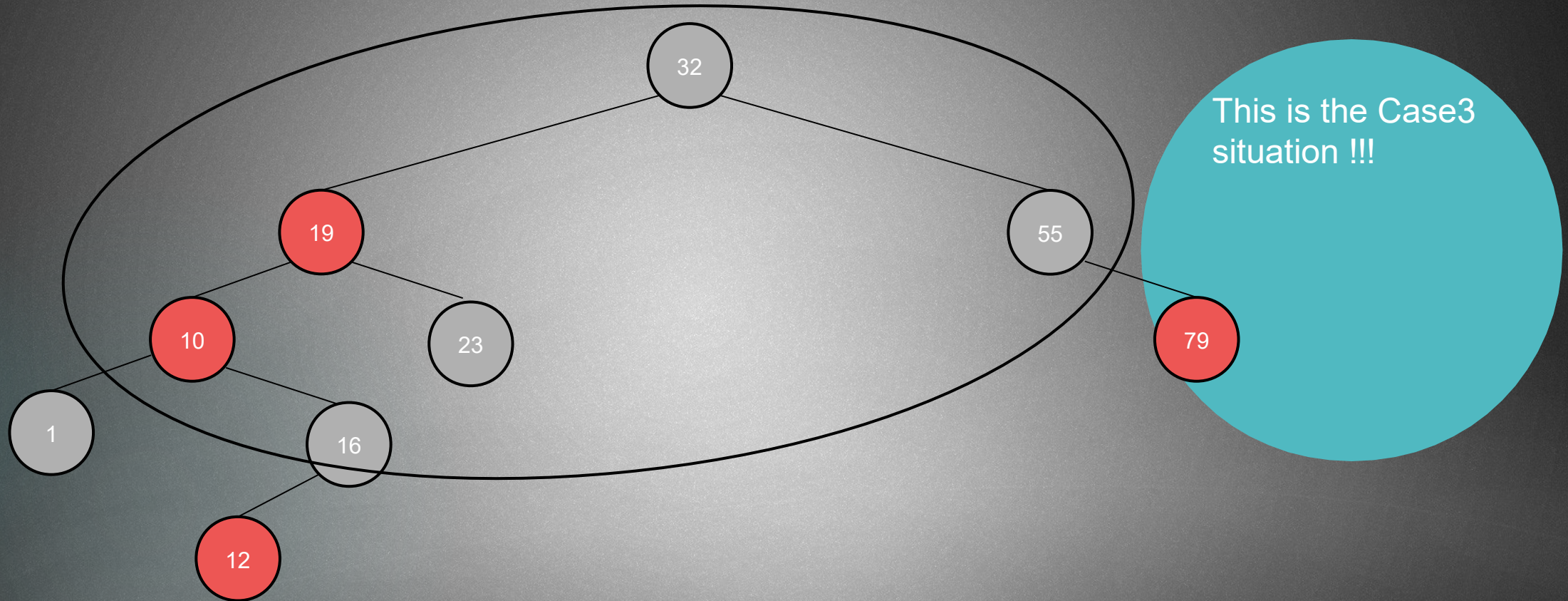
Example:

We have to check whether
the red black tree properties are violated or not !!!



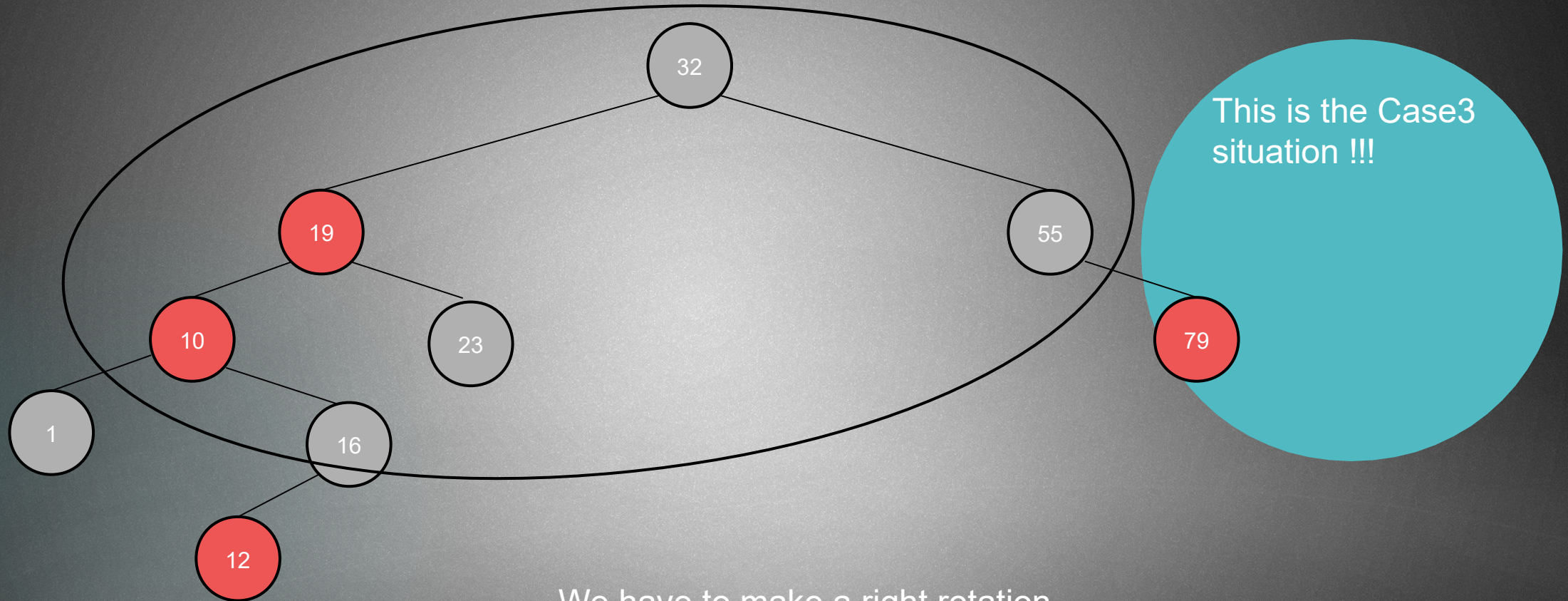
Example:

We have to check whether
the red black tree properties are violated or not !!!



Example:

We have to check whether
the red black tree properties are violated or not !!!



We have to make a right rotation
on the root node !!!

Example:

