Computing Large Scale Linear Least Squares with SKiLLS

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1 Overview

SKiLLS (SKetchIng-Linear-Least-Sqaures) is a C++ package for finding solutions to over-determined linear least square problems. SKiLLS uses a modern dimensionality reduction technique called sketching, and is particularly suited for large scale linear least squares where the number of measurements/observations is far greater than the number of variables.

Mathematically, SkiLLS solves

$$\min_{x \in \mathbb{R}^d} f(x) := \|Ax - b\|_2^2,\tag{1}$$

where $A \in \mathbb{R}^{n \times d}$ and $b \in \mathbb{R}^n$ given. The matrix A is allowed to be rank-deficient or nearly rank-deficient.

1.1 When to use SKiLLS

If A in (1) is dense, the state-of-the-art sketching solver is Blendenpik [2]. Comparing to solvers in the classical state-of-the-art numerical package LAPACK [1], Blendenpik is two times faster on matrices of size $20,000 \times 500$ and four times faster on matrices of size $100,000 \times 2500$. Comparing to classical iterative solver LSQR [7], Blendenpik is 80 times faster on matrices of size $20,000 \times 1,000$ with condition number 100.

Blendenpik only solves (1) when the matrix A has full numerical rank, as illustrated in Table 1. Two solvers are included in the library SKiLLS for dense A. The robust version,

 $ls_dense_hashing_blendenpik$ solves problems 1 when A is numecially rank-deficient but is as fast as Blendenpik when A has full rank (takes 70% to 130% time on matrices of different sizes). The fast version, $ls_dense_hashing_blendenpik_noCPQR$ is 1.6 times faster than Blendenpik on coherent matrices¹ of size $40,000 \times 4,000$ and $90,000 \times 2250$, and slightly faster than Blendenpik on other types of input (about 1.1 times faster).

If A in (1) is sparse, the state-of-the-art solvers are SPQR [3] and HSL [8], which uses sparse QR factorization of A and LSQR with an incomplete Cholesky preconditioner, respectively. Our library has a single sparse solver, ls_sparse_spqr . It significantly outperforms both state-of-the-art sparse solvers on large random sparse ill-conditioned matrices. It is 10 times faster than HSL and 7 times faster than SPQR on $120,000 \times 5,000$ random sparse matrices with 1% non-zero entries and condition number equals to 10^6 . And it is 3 times faster than both HSL and SPQR on $40,000 \times 2,000$ matrices with 1% non-zero entries and condition number equals to 10^6 .

¹Matrices of the form $A \in \mathbb{R}^{n \times d} = \binom{I_{d \times d}}{0} + 10^{-8} J_{n,d}$ where $J_{n,d} \in \mathbb{R}^{n \times d}$ is a matrix of all ones.

Our sparse solver also performs extremely well on highly over-determined sparse inputs from real applications. It is the fastest solver on more than 75% of problems in the Florida Matrix Collection [5] with the matrix A defined in (1) having $n \geq 30d$ and $n \geq 20,000$. It is also the fastest solver on more than 90% of problems with the matrix A having $n \geq 10d$, $n \geq 20,000$ but additionally having more than 1% of non-zero entries.

Because sketching exploits redundancy of rows, the performance gain compared to non-sketching solvers² will be larger on A with larger n/d radio.

For further details of computational advantages, see our paper. Computing Large Scale Linear Least Squares with SKiLLS, C. Cartis and Z. Shao.

2 Installing SKiLLS

- 2.1 Dependency
- 2.2 Installation instruction
- 2.3 Test the installation

3 Using SKiLLS

This is a library of three routines for solving problem (1).

- If the matrix A is dense, $ls_dense_hashing_blendenpik$ uses LSQR with a preconditioner built from a sketch of the matrix A using Hashed-Randomised-Hadamard-Transform (HRHT) and Randomised-Column-Pivoted-QR (RCPQR) to solve (1).
- If A is dense and has full numerical rank, $ls_dense_hashing_blendenpik_noCPQR$ uses LSQR with a preconditioner built from a sketch of the matrix A using HRHT and Column-Pivoted-QR (CPQR) to solve (1).
- If A is sparse, ls_sparse_spqr uses LSQR with a preconditioner built from a sketch of the matrix A using s—hashing and sparse QR (SPQR) to solve (1).

3.1 Data structure

Dense Vector and Matrix SKiLLS uses C++ class Vec for dense vectors and C++ class Mat for dense matrices. To create a $n \times 1$ vector from already existed data, take a pointer to the data and call Vec(n, *data). To create a $n \times d$ matrix from already existed data, take a pointer to the data and call Mat(n, d, *data). Basic matrix operations such as return a specific row, column, or matrix-matrix and matrix vector multiplications are implemented. These classes are defined in include/Vec.hpp and include/impl/Mat.hpp respectively.

Sparse Matrix SKiLLS uses Compressed Column Data C structure for sparse matrices and vectors. The sparse data structure is from SuiteSparse [4]. The cs_dl and cholmod_sparse structures are defined in cs.h and cholmod_core.h respectively the in SuiteSparse/include folder.

²LAPACK, HSL, SPQR

3.2 Dense problems

To compute a solution of a linear least square problem where the matrix A is dense, a call of the following form should be made.

ls_dense_hashing_blendenpik(A, b, x, rank, flag, it, gamma, k, abs_tol, rcond, it_tol, max_it, debug, wisdom):

• Inputs

- 1. **A** is type Mat d containing a $\mathbb{R}^{m \times n}$ matrix defined in (1).
- 2. **b** is type Vec d containing a \mathbb{R}^m vector defined in (1).
- Main Outputs
 - 1. **x** is type Vec_d containing a \mathbb{R}^n vector which is a solution of (1).
- Auxillary Outputs
 - 1. \mathbf{rank} is a scalar containing the detected numerical rank of the matrix A.
 - 2. **flag** is a scaler. flag=0 indicates LSQR has converged. flag=1 indicates LSQR has not converged.
 - 3. it is a scaler indicating the number of LSQR iterations taken.

• Parameters

- 1. **gamma** is the over-sampling ratio used in sketching. Bigger gamma typically gives higher quality preconditioner but slower running time. The default value of gamma is 1.7.
- 2. \mathbf{k} is number of non-zeros per column in the hashing matrix as part of sketching. Bigger k typically gives higher quality preconditioner but slower running time. The default value of k is 1.
- 3. **abs_tol** specifies an absolute residual tolerance for the solution x of 1. If the algorithm finds an x satisfies $||Ax b||_2 \le abs_t ol$, it terminates and returns x. The default value is 10^{-8}
- 4. it _tol specifies the relative residual tolerance for LSQR convergence, see [7]. The default value is 10⁻⁶.
- 5. max it specifies the maximum iteration of LSQR, the default value is 10,000.
- 6. **debug** is a flag. If debug=1, the solver prints additional outputs for diagosis.
- 7. **wisdom** is a flag. If wisdom=1, the macro variable FFTW_WISDOM_FILE in inclde/Config.hpp must be defined as the path to a FFTW wisdom file, see fftw documentation at fftw.org. If wisdom=0, the solver does not use FFTW wisdom and may run slower.

If the matrix A is known to be of full numerical rank, then a different routine can be used which is slightly faster:

ls_dense_hashing_blendenpik_noCPQR(A, b, x, rank, flag, it, gamma, k, it_tol, max_it, debug, wisdom). The arguments usage is the same as above, but we don't need the rank-detection parameter roond, and the shortcut related to abs_tol is not implemented.

Example The following program solves an example of problem (1) with given A and b.

```
#include <iostream>
#include "Config.hpp"
#include "SpMat.hpp"
#include "Random.hpp"
#include <math.h>
#include "RandSolver.hpp"
#include "cs.h"
#include <stdio.h>
#include "SuiteSparseQR.hpp"
#include "cholmod.h"
#include "lsex_all_in_c.h"
#include "IterSolver.hpp"
#include <sys/timeb.h>
#include "bench config.hpp"
#include "blendenpik.hpp"
int main(int argc, char **argv)
        // create data
        long m = 4;
        long n = 2;
        double data A[8] = \{0.306051, 1.53078, 1.64493, 1.61322,
         0.2829, 0.474476, 0.586278, 0.610202;
        double data b[4] = \{0.0649338, 0.845946, 0.0164085, 0.247119\};
        Mat_d A(m, n, data A);
        Vec d b(m, data b);
        // parameters
        long k = NNZ PER COLUMN;
        double gamma = OVER SAMPLING RATIO;
        long max it = MAX \overline{\Pi};
        double it tol = IT TOL;
        double rcond = 1e 10;
        int wisdom = 0;
        double abs tol = ABS TOL;
        // storage
        long it;
        int flag;
        long rank;
        double residual;
        Vec d x(A.n());
        // solve
        ls dense hashing blendenpik (
        A, b, x, rank, flag, it, gamma, k, abs tol,
        rcond, it tol, max it, debug, wisdom);
        std::cout << "A_is:_"<< A << std::endl;
        std::cout << "b_is:_"<< b << std::endl;
        std::cout << "solution_is:_"<< x << std::endl;
        // compute residual
        A.mv('n', 1, x, 1, b);
```

```
std::cout << "Residual_of_ls_blendenpik_hashing_is:_"<< nrm2(b) << std::endl;
std::cout << "Iteration_of_ls_blendenpik_hashing_is:_" << it << std::endl;
}</pre>
```

3.3 Sparse problems

To compute a solution of a linear least square problem where the matrix A is sparse, a call of the following form should be made.

ls_sparse_spqr(*A, b, x, rank, flag, it, gamma, k, abs_tol, ordering, it_tol, max_it, rcond, peturb, debug):

- Inputs
 - 1. A is pointer to type cs dl containing a $\mathbb{R}^{m \times n}$ sparse matrix defined in (1).
 - 2. **b** is type Vec_d containing a \mathbb{R}^m vector defined in (1).
- Main Outputs
 - 1. **x** is type Vec_d containing a \mathbb{R}^n vector which is a solution of (1).
- Auxillary Outputs
 - 1. \mathbf{rank} is a scalar containing the detected numerical rank of the matrix A.
 - 2. **flag** is a scaler. flag=0 indicates LSQR has converged. flag=1 indicates LSQR has not converged.
 - 3. it is a scaler indicating the number of LSQR iterations taken.
- Parameters
 - 1. **gamma** is the over-sampling ratio used in sketching. Bigger gamma typically gives higher quality preconditioner but slower running time. The default value of gamma is 1.7.
 - 2. \mathbf{k} is number of non-zeros per column in the hashing matrix as part of sketching. Bigger k typically gives higher quality preconditioner but slower running time. The default value of k is 1.
 - 3. **abs_tol** specifies an absolute residual tolerance for the solution x of 1. If the algorithm finds an x satisfies $||Ax b||_2 \le abs_t ol$, it terminates and returns x. The default value is 10^{-8}
 - 4. **ordering** is an integer specifying a fill-reduced ordering for sparse QR factorization. The default value is SPQR_ORDERING_DEFAULT. See SuiteSparseQR documentation for more information.
 - 5. **it_tol** specifies the relative residual tolerance for LSQR convergence, see [7]. The default value is 10^{-6} .
 - 6. max it specifies the maximum iteration of LSQR, the default value is 10,000.
 - 7. **rcond** is numerical-ill-conditioning tolerance for the preconditioner returned by sparse QR. If the condition number of the preconditioner is larger than 1/rcond, a warning will be printed. The default value of rcond is 10^{-12}
 - 8. **perturb** is a real number. It has no-effect on the algorithm and is part of an ongoing development.
 - 9. **debug** is a flag. If debug=1, the solver prints additional outputs for diagosis.

Example The following program solves an example of problem (1) with given sparse A and dense b.

```
#include <iostream>
#include "Config.hpp"
#include "SpMat.hpp"
#include "Random.hpp"
#include <math.h>
#include "RandSolver.hpp"
#include <time.h>
#include "cs.h"
#include <stdio.h>
#include "SuiteSparseQR.hpp"
#include "cholmod.h"
#include "lsex all in c.h"
#include "IterSolver.hpp"
#include <sys/timeb.h>
#include "bench config.hpp"
int main(int argc, char const *argv[])
{
         // Create a sparse matrix in compressed column format
     long m = 4;
     long n = 3;
     long nnz = 5;
     \begin{array}{lll} \textbf{long} & col\left[4\right] = \{0\,, & 1\,, & 3\,, & 5\}; \\ \textbf{long} & row\left[5\right] = \{0\,, & 0\,, & 2\,, & 1\,, & 3\}; \\ \textbf{double} & val\left[5\right] = \{2.0\,, & 3.0\,, & 1.0\,, & 4.0\,, & 5.0\}; \end{array}
     cholmod common Common, *cc;
     cc = \&Common;
     cholmod_l_start(cc);
     cc > print = 4;
     cholmod sparse* A;
     A = cholmod l allocate sparse(
     m, n, nnz, true, true, 0, CHOLMOD REAL, cc);
     A > p = col;
     A > i = row;
     A > x = val;
     cholmod_l_print_sparse( A, "A", cc);
     cholmod l finish(cc);
          //input
          cs_dl *A;
     A = to cs dl(Acholmod);
     double *b data = (double*) malloc((A>m)*sizeof(double));
     for (long i=0; i < A > m; i++){
          b data[i] = 1;
     Vec d b(A>m, b data);
     // parameters
     long k = NNZ PER COLUMN;
     double gamma = OVER SAMPLING RATIO;
```

```
long max it = MAX IT;
double it tol = IT TOL;
double abs tol = ABS TOL;
double roond = 1e 12;
int ordering = SPQR ORDERING;
// storage
long it;
int flag;
long rank;
double t_start;
double t finish;
double residual;
Vec d x ls qr(A > n);
// solve
ls sparse spqr(*A, b, x ls qr,
rank, flag, it, gamma, k, abs_tol, ordering, it_tol, max_it, RCOND THRESHOLD, PERTURB, debug);
CSC Mat d A CSC(A > m, A > n, A > nzmax,
    (\mathbf{long}*)A > i, (\mathbf{long}*)A > p, (\mathbf{double}*)A > x);
A CSC.mv('n', 1, x ls qr, 1, b);
std::cout << "Residual_of_ls sparse spqr_is:_"<< nrm2(b) << std::endl;
    return 0:
```

References

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- [2] H. Avron, P. Maymounkov, and S. Toledo. Blendenpik: Supercharging LAPACK's Least-Squares Solver. SIAM Journal on Scientific Computing, 32(3):1217–1236, Jan. 2010.
- [3] T. A. Davis. Algorithm 915, SuiteSparseQR: Multifrontal multithreaded rank-revealing sparse QR factorization. ACM Transactions on Mathematical Software (TOMS), 38(1):8, Nov. 2011.
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A Numerical Results

	$lp_ship12l$	Franz1	GL7d26	cis-n4c6-b2	$lp_modszk1$	rel5	ch5-5-b1
SVD	18.336	26.503	50.875	6.1E-14	33.236	14.020	7.3194
Blendenpk	NaN	9730.700	NaN	$3.0E{+}02$	NaN	NaN	340.9200
Ski-LLS-dense	18.336	26.503	50.875	5.3E-14	33.236	14.020	7.3194
	n3c5-b2	ch4-4-b1	n3c5-b1	n3c4-b1	connectus	landmark	cis-n4c6-b3
SVD	9.0E-15	4.2328	3.4641	1.8257	282.67	1.1E-05	30.996
Blendenpk	$1.3E{+}02$	66.9330	409.8000	8.9443	NaN	NaN	3756.200
Diendenpn							

Table 1: Residual of solvers for a range of approximate rank-deficient problems taken from the Florida matrix collection [6]. We see while Blendenpik struggles, SKiLLS achives the same residual accuracy as SVD method.