

## Lecture 14 Outline

Wednesday, October 9, 2019 3:22 PM

### Midterm Exam

Wed, October 16

In-class ( $\sim 1$  hr, 20 minutes)

Closed book, closed notes; no calculators,  
phones or computers

All material through Lecture 10 on linear  
transformations and HW 1-4

HW 4 Review

Orthogonal + Orthonormal Basis

# Orthogonal & Orthonormal Basis

Sunday, October 6, 2019 10:42 PM

Recall that vectors are orthogonal, if

$$\underline{q}_0 \cdot \underline{q}_1 = 0$$

An orthogonal basis is one where all the vectors of the basis are orthogonal to each other

Example:  $B = \{ \underline{q}_0, \underline{q}_1, \underline{q}_2 \}, \underline{q}_i \cdot \underline{q}_j = 0, i \neq j$

An orthonormal basis is one where

$$\underline{q}_i \cdot \underline{q}_j = 0 \quad i \neq j$$

$$\underline{q}_i \cdot \underline{q}_i = 1$$

Now, consider a matrix with orthonormal columns

$$\underline{Q} = \begin{bmatrix} | & | & | & \dots & | \\ \underline{q}_1 & \underline{q}_2 & \underline{q}_3 & \dots & \underline{q}_n \\ | & | & | & \dots & | \end{bmatrix} \quad \begin{array}{ll} \underline{q}_1 \cdot \underline{q}_1 = 1 & \underline{q}_1 \cdot \underline{q}_2 = 0 \\ \vdots & \\ \underline{q}_n \cdot \underline{q}_n = 1 & \underline{q}_n \cdot \underline{q}_{n-1} = 0 \end{array}$$

Look at  $\underline{Q}^T \underline{Q}$

Look at  $\underline{Q} \quad \underline{Q}$

$$\begin{bmatrix} \underline{q}_1^T \\ \underline{q}_2^T \\ \vdots \\ \underline{q}_n^T \end{bmatrix} \begin{bmatrix} \underline{q}_1 & \underline{q}_2 & \cdots & \underline{q}_n \end{bmatrix} = \begin{bmatrix} 1 & 0 & & \\ 0 & 1 & & \\ & & \ddots & \\ & & & 0 & 1 \end{bmatrix}$$

$$\Rightarrow \underline{Q}^T \underline{Q} = \underline{I}$$

If  $\underline{Q}$  is square, then  $\underline{Q} \underline{Q}^T = \underline{I}$

$$\Rightarrow \underline{Q}^T = \underline{Q}^{-1} \text{ for square } \underline{Q} \text{ and is}$$

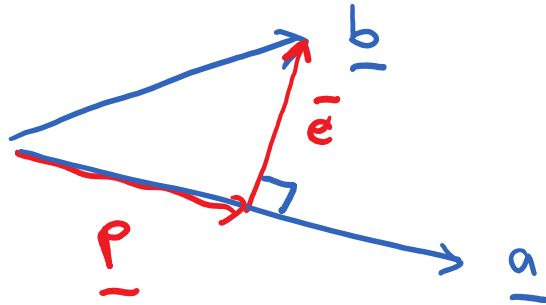
called a unitary matrix

How can  $\underline{Q}$  be constructed?

### Orthonormal Basis Construction

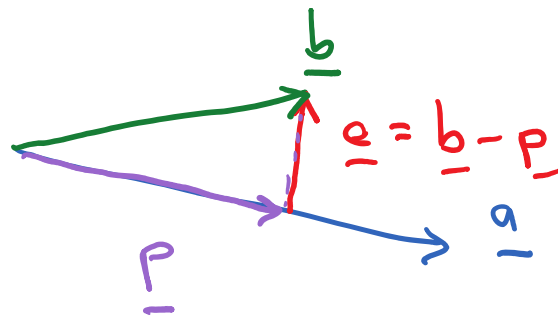
Given a set of vectors that span a subspace,  
find an orthonormal basis that also spans  
that subspace

Example: In  $\mathbb{R}^2$ , let  $\underline{a} \neq \underline{b}$  be non-parallel vectors



$\underline{a} \neq \underline{b}$  span  $\mathbb{R}^2$ , but are not orthogonal

Recall the projection of  $\underline{b}$  onto  $\underline{a}$



$$\underline{e} = \underline{b} - \underline{p} = \underline{b} - \frac{\underline{a}^T \underline{b}}{\underline{a}^T \underline{a}} \underline{a} ; \underline{e} \perp \underline{a}$$

Thus, an orthonormal basis is

$$\underline{q}_1 = \frac{\underline{a}}{\|\underline{a}\|_2}, \quad \underline{q}_2 = \frac{\underline{e}}{\|\underline{e}\|_2} = \frac{\underline{b} - (\underline{a}^T \underline{b} / \underline{a}^T \underline{a}) \underline{a}}{\|\underline{e}\|_2}$$

Are these vectors  $\underline{q}_1$  &  $\underline{q}_2$  unique?

No! For example, one could project  $\underline{a}$  onto  $\underline{b}$  instead

Now, consider a matrix and find an orthonormal basis to a column space  $C(A)$ , with

$$\underline{A} = \begin{bmatrix} a_1 & b_1 & c_1 \\ a_2 & b_2 & c_2 \\ a_3 & b_3 & c_3 \\ a_4 & b_4 & c_4 \end{bmatrix} = [\underline{a} \quad \underline{b} \quad \underline{c}]$$

Step 1: Set  $\underline{t}_1 = \underline{a}$

Step 2: Project onto  $\perp$  space of  $\underline{a}$

$$\underline{t}_2 = \underline{b} - \frac{\underline{t}_1^T \underline{b}}{\underline{t}_1^T \underline{t}_1} \underline{t}_1 \quad (\text{i.e. } \underline{I} - \underline{A})$$

$$\Rightarrow \underline{t}_1 \cdot \underline{t}_2 = 0$$

Step 3: Project onto  $\perp$  space of  $\underline{t}_1$  &  $\underline{t}_2$

Step 3: Project onto  $\perp$  space of  $\underline{t}_1$  &  $\underline{t}_2$

$$\underline{t}_3 = \underline{c} - \frac{\underline{t}_1^T \underline{c}}{\underline{t}_1^T \underline{t}_1} \underline{t}_1 - \frac{\underline{t}_2^T \underline{c}}{\underline{t}_2^T \underline{t}_2} \underline{t}_2$$

$$\Rightarrow \underline{t}_2 \cdot \underline{t}_3 = 0 \quad \& \quad \underline{t}_1 \cdot \underline{t}_3 = 0$$

Step 4: Normalize

$$\underline{q}_1 = \frac{\underline{t}_1}{\|\underline{t}_1\|}, \quad \underline{q}_2 = \frac{\underline{t}_2}{\|\underline{t}_2\|}, \quad \underline{q}_3 = \frac{\underline{t}_3}{\|\underline{t}_3\|}$$

This process is called Gram-Schmidt (G-S) orthonormalization

The result is an orthonormal basis to  $C(\underline{A})$

How do  $\underline{Q}$  and  $\underline{A}$  relate?

Recall that G-S stated that

$$\underline{q}_1 = \frac{\underline{a}}{\|\underline{a}\|} = \frac{\underline{a}}{r_{11}} \Rightarrow \underline{a} = r_{11} \underline{q}_1$$

$$\underline{q}_2 = \frac{\underline{b} - r_{21} \underline{q}_1}{\|\underline{b} - r_{21} \underline{q}_1\|} = \frac{\underline{b} - r_{21} \underline{q}_1}{r_{22}}$$

$$= \frac{1}{r_{22}} \underline{b} - \frac{r_{21}}{r_{22}} \underline{q}_1$$

$$\underline{b} = r_{21} \underline{q}_1 + r_{22} \underline{q}_2$$

$$\text{Similarly, } \underline{c} = r_{31} \underline{q}_1 + r_{32} \underline{q}_2 + r_{33} \underline{q}_3$$

$$\therefore \underline{A} = [\underline{a} \ \underline{b} \ \underline{c}] = \begin{bmatrix} \underline{q}_1 & \underline{q}_2 & \underline{q}_3 \end{bmatrix} \begin{bmatrix} r_{11} & r_{12} & r_{13} \\ 0 & r_{22} & r_{23} \\ 0 & 0 & r_{33} \end{bmatrix}$$

$$\underline{A} = \underline{Q} \underline{R} \quad \text{QR decomposition}$$

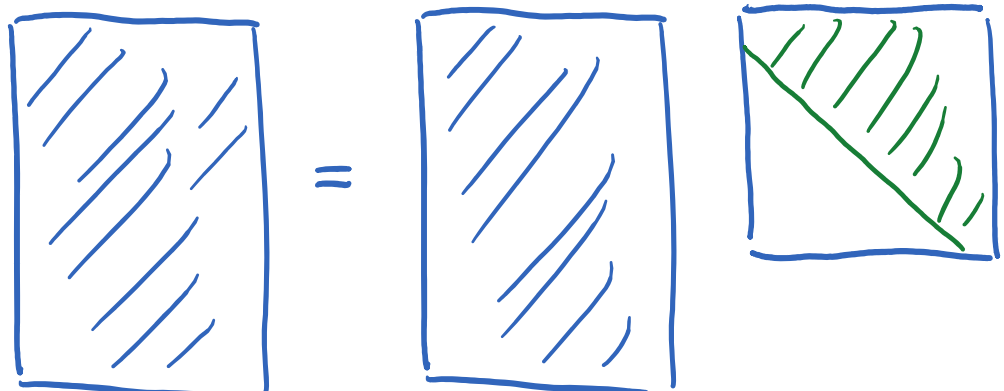
↙
↘

orthonormal matrix      upper triangular

This is actually called

# Reduced QR Decomposition

typically written as  $\underline{\hat{A}} = \underline{\hat{Q}} \underline{\hat{R}}$



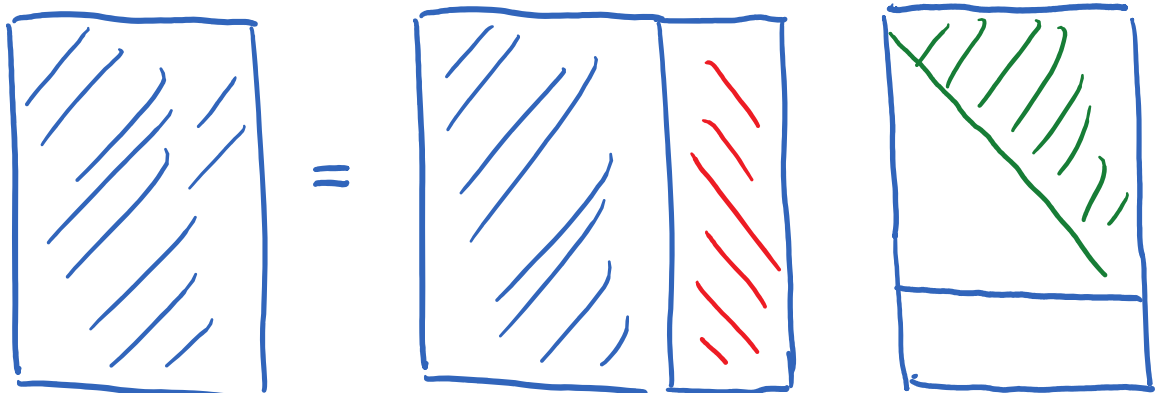
$$\underline{\hat{A}}_{m \times n} = \underline{\hat{Q}}_{m \times n} \underline{\hat{R}}_{n \times n}$$

Note:  $\underline{\hat{R}} = \begin{bmatrix} \underline{q}_1^T \underline{a} & \underline{q}_1^T \underline{b} & \underline{q}_1^T \underline{c} \\ 0 & \underline{q}_2^T \underline{b} & \underline{q}_2^T \underline{c} \\ 0 & 0 & \underline{q}_3^T \underline{c} \end{bmatrix}$

Also, one can determine a Full QR Factorization by appending columns to  $\underline{\hat{Q}}$  to make it  $m \times m$



(typically with  $m \geq n$ )


$$\begin{array}{ccc} \begin{array}{c} \underline{A} \\ m \times n \end{array} & = & \begin{array}{c} \underline{Q} \\ m \times m \end{array} \begin{array}{c} \underline{R} \\ m \times n \end{array} \end{array}$$

The columns  $\underline{q}_j$  for  $j > n$  must be orthogonal to the range( $\underline{A}$ ).

If  $\text{rank}(\underline{A}) = n$ , then these columns are the orthonormal basis to  $\text{null}(\underline{A}^T)$

Why is this useful?

Theorem: Every  $\underline{A} \in \mathbb{R}^{m \times n}$  with  $m \geq n$  has a full QR factorization + a reduced

## QR factorization

Theorem: Each  $\underline{A} \in \mathbb{R}^{m \times n}$  with  $m \geq n$  of full rank ( $\text{rank}(\underline{A}) = n$ ) has a unique reduced QR with  $r_{jj} > 0$

All diagonals of  $\underline{R}$  are positive

$\Rightarrow \underline{R}^{-1}$  exists and so does  $\underline{R}^{-T}$

Now, let's look at  $\underline{A}\underline{x} = \underline{b}$ ,  $\underline{A}$  is full rank

Solve via  $\hat{\underline{Q}} \hat{\underline{R}}$

① Decompose:  $\underline{A} = \hat{\underline{Q}} \hat{\underline{R}}$  ✓ expensive

$$\textcircled{2} \quad \hat{\underline{Q}} \hat{\underline{R}} \underline{x} = \underline{b} \Rightarrow \hat{\underline{Q}}^T \hat{\underline{Q}} \hat{\underline{R}} \underline{x} = \hat{\underline{Q}}^T \underline{b}$$

$$\textcircled{3} \quad \hat{\underline{R}} \underline{x} = \hat{\underline{Q}}^T \underline{b} \quad \checkmark \text{ Solve (cheap)}$$

If  $\underline{b}$  changes (new right-hand side), but  $\underline{A}$  does not, then it is cheap to solve for new solution  $\underline{x}$

Note: Drop  $\wedge$

Return to consider  $\underline{A}^T \underline{A} \underline{x} = \underline{A}^T \underline{b}$

(Least squares formulation)

$$\underline{A}^T \underline{A} \underline{x} = \underline{A}^T \underline{b} \quad \text{Find } \underline{A} = \underline{Q} \underline{R}$$

$$(\underline{Q} \underline{R})^T (\underline{Q} \underline{R}) \underline{x} = (\underline{Q} \underline{R})^T \underline{b}$$

$$\underline{R}^T \underline{Q}^T \underline{Q} \underline{R} \underline{x} = \underline{R}^T \underline{Q}^T \underline{b}$$

$$\underline{R}^T \underline{R} \underline{x} = \underline{R}^T \underline{Q}^T \underline{b}, \text{ where } \underline{R}^{-T} \text{ exists}$$

$$\Rightarrow \underline{R} \underline{x} = \underline{Q}^T \underline{b}$$

If  $m > n$  for  $\underline{A} \in \mathbb{R}^{m \times n}$ , then

solving  $\underline{R}\underline{x} = \underline{Q}^T \underline{b}$  is the solution  
that minimizes the error

Another advantage: Solving  $\underline{R}\underline{x} = \underline{Q}^T \underline{b}$   
is much more stable than

$$\underline{x} = (\underline{A}^T \underline{A})^{-1} \underline{A}^T \underline{b}$$

↳ Remember condition  
number for this!

### Classical Gram-Schmidt Algorithm

One algorithm for reduced QR of  $\underline{A}$

$$\text{Let } \underline{A} = [\underline{a}_1 \ \underline{a}_2 \ \dots \ \underline{a}_n]$$

$$\text{Recall that } \underline{q}_1 = \frac{\underline{a}_1}{\|\underline{a}_1\|} = \frac{\underline{a}_1}{\|\underline{a}_1\|}$$

$$- \| \underline{q} \| \quad r_{11}$$

$$\underline{q}_2 = \frac{\underline{q}_2 - r_{12} \underline{q}_1}{r_{22}}$$

⋮

$$\text{where } r_{ij} = \underline{q}_i^T \underline{q}_j \text{ for } i \neq j$$

$$\text{and } |r_{jj}| = \| \underline{q}_j - \sum_{i=1}^{j-1} r_{ij} \underline{q}_i \|_2$$

Note:  $r_{jj}$  can be either + or -,

choose  $\oplus$  value

Algorithm: Classical G-S

for  $j = 1 : n$

$$\underline{v}_j = \underline{q}_j$$

for  $i = 1 : j-1$

$$r_{ij} = \underline{q}_i^T \underline{q}_j$$

$$\underline{Q} = \begin{bmatrix} \underline{q}_1 & \underline{q}_2 & \dots & \underline{q}_n \end{bmatrix}$$

$$r_{ij} = \underline{q}_i^T \underline{a}_j$$

$$\underline{v}_j = \underline{v}_j - r_{ij} \underline{q}_i$$

end

$$r_{jj} = \|\underline{v}_j\|_2$$

$$\underline{q}_j = \underline{v}_j / r_{jj}$$

end

Operation Count for G-S

Most expensive operation  $r_{ij} = \underline{q}_i^T \underline{v}_j$

$$+ \underline{v}_j = \underline{v}_j - r_{ij} \underline{q}_i$$

$$\Rightarrow \sum_{i=1}^n \sum_{j=i+1}^n 4m \sim \sum_{i=1}^n 4mi \sim 2mn^2 //$$

However, Classical G-S is not numerically stable  $\Rightarrow$  Roundoff errors cause issues

(We will not prove, because this requires complicated stability + error analysis. There will be a related HW problem.)

Consequently, a better method is needed:

Modified Gram-Schmidt