

Java lists, maps and not only...

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Outline

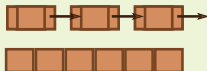
- 1 Hello
- 2 Algorithmics
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- 3 List
 - ArrayList
 - Generics
 - LinkedList
- 4 Maps
 - HashMap
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- 6 Summary
 - What have we learned today?

Introduction

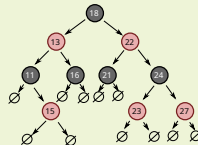
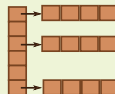
- Java collections are sophisticated implementations of algorithms for storing data (see also Cormen's "Introduction to algorithms")
- All collections implement iteration
- Lists are implemented as sequential collections backed by arrays and linked lists
- Sets, besides iteration, are also allowing checking if the item is present.
- Maps are associative containers which allow quick access by object-key (any object, also other than integer).
- Trees are underlying some implementations of Maps and Sets
- Lists maintain item order while it is not guaranteed in Sets and Maps.

The plan

Lists



Maps



$$\sum_{n=1}^{\infty} 0 < \left| \lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} \right| < \infty$$

Algorithm analysis

The Big O, Ω , Θ notation

When there exists $n_0 > 0$ and $c > 0$ for which we can say that:

$$\forall n > n_0 \quad f(n) \leq c \cdot g(n)$$

we say that f is big-O g and write $f(x) = O(g(x))$

When there exists $n_0 > 0$ and $c > 0$ for which we can say that:

$$\forall n > n_0 \quad f(n) \geq c \cdot g(n)$$

we say that f is big- Ω g and write $f(x) = \Omega(g(x))$

$$f(x) = \Theta(g(x))$$

Examples

- Given $g(x) = x^2$ and $f(x) = 3x^2 + 4$ we can say that $f(x) = O(g(x)) \dots$
- Given $g(x) = x^3$ and $f(x) = 100x^2 + 4x^3$ we can say that $f(x) = O(g(x)) \dots$
- Given $g(x) = 2^x$ and $f(x) = 2^x + x^{64}$ we can say that $f(x) = O(g(x)) \dots$

We can also easily say that $x^2 = O(2^x)$, and $x^2 = \Omega(1)$ but it is not very interesting. This observation is very weak. We are rather interested in finding the closest matches.

Θ sufficient condition

The most interesting is finding the simple in form, but exact, match of the function. This match is symbolised by Θ .

$$f(x) = O(g(x)) \wedge f(x) = o(g(x)) \implies f(x) = \Theta(g(x))$$

If $f(x) = \Theta(g(x))$, it means that f grows asymptotically as fast as g .

Sufficient condition for $f(x) = \Theta(g(x))$ is:

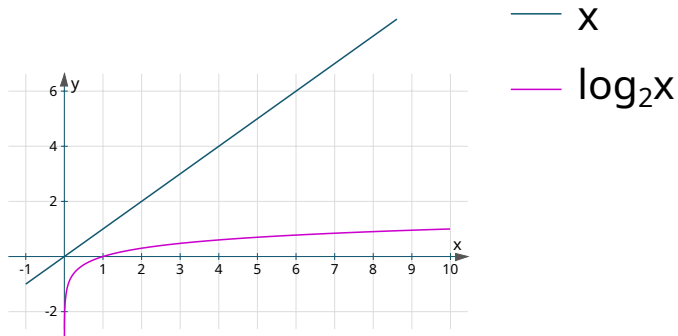
$$0 < \lim_{x \rightarrow \infty} \left| \frac{f(x)}{g(x)} \right| < \infty$$

Some things that programmers usually know

We use given notation extensively to describe algorithm time of execution (and memory consumption). Hence for us (generally):

- x^2 is worse than x
- x is worse than $\log x$
- 2^x is worse than x^{16} and it's bad
- $x!$ – it's so much, we don't distinguish between $x!$ and 2^x . Those are equally BAD.
- We often don't distinguish between logarithm and exponential bases

Examples



Worst case scenario

Programmers also consider worst case (pessimistic) scenarios of algorithms and their complexity. Some algorithms work very slow on some special cases of input data. E.g. **quicksort**, despite of being $O(n \cdot \log n)$ runs at n^2 time, when in every partitioning selected pivot divides data to lengths: 1 and rest.

Having that in mind, we can select, for example, **heapsort**, which has worst case running time still $n \cdot \log n$

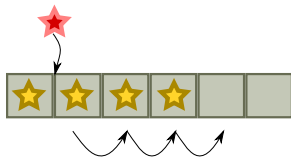
- Allocated as a one block in memory (more “array-ish” than “list-ish”)
- Quick “please get me an element at position n ” (further referred to as **get**): $\Theta(1)$
- Slow “please insert element at position n , moving all following elements to the right” (further referred to as **insert**): $O(n)$
- Slow “please delete element at position n , moving all following elements to the left” (further referred to as **delete**): $O(n)$

ArrayList

- Quick **append** - a special-case of the insert at the last place, i.e. when for n equal number of elements: $\Theta(1)$, but only when the underlying array size is $N > n + 1$
- When ArrayList is full, to perform insert we need to expand underlying array. We do it by increasing size by twice the current size.
- Is really the **append** operation $\Theta(1)$???

ArrayList pros and cons

- Insert at an arbitrary place costs n operations for $n = s - i$ where s is number of items and i is the position we insert at.



Amortised cost. Amortised analysis

- We analyse amortised cost of operations amortised by their number.
- Amortised cost for given $f(n)$ is $F(n)$ that:

$$F(n) = \frac{T(n)}{n}$$

where

$$T(n) = \sum_{x=0}^n f(x)$$

Amortised cost. Amortised analysis

- For array initialised to 2, in the **append** operation we have cost:

$$1, 1, n_1 = 2, 1, 1, n_2 = 4, 1, 1, 1, 1, n_3 = 8$$

where n_x is n-th resize cost equals array's size at the moment of resize.

- If we use an accounting method to tell that every operation, we put additional 2 operations' time on a special account, for a later use, we can reuse that time at critical sections of resizing. Our account state is then:

Amortised cost. Amortised analysis

Operation	Account state
+2	2
+2	4
-2	2
+2	4
+2	6
-4	2
+2	4
+2	6
+2	8
+2	10
-8	2

Table 1: Amortised analysis using accounting method

Amortised cost. Amortised analysis

SVG Animation for the amortised cost accounting method

Amortised cost. Amortised analysis

- In this case we say that operation **append** on ArrayList is $O(1)$ and keep in mind that sometimes it can stop our system for a very long time. So if we want to get overall computation time fit, we can allow us to expand a very large ArrayList, but when we are low-latency needers, we might need to search for a better solution.

Generics and wildcards

```
public void addToList(List<? extends Number> lst) {  
    Number item1 = lst.get(0);  
    lst.add(3);  
}
```

```
List<Float> floatList = ...;  
addToList(floatList);
```

Generics and wildcards

```
public void addToList(List<? extends Number> lst) {  
    Number item1 = lst.get(0); // ok  
    lst.add(3); // error  
}
```

```
List<Float> floatList = ...;  
addToList(floatList);
```

Generics and wildcards

```
public void addToList(List<? super Integer>) {  
    Integer i = lst.get(0);  
    lst.add(4);  
}
```

```
List<Number> numberList = ...;  
addToList(numberList);
```

Generics and wildcards

```
public void addToList(List<? super Integer>) {  
    Integer i = lst.get(0); // error  
    lst.add(4); // ok  
}
```

```
List<Number> numberList = ...;  
addToList(numberList);
```

Generics and wildcards

```
// https://stackoverflow.com/questions/4343202

public class Collections {
    public static <T> void copy(List<? super T> dest,
List<? extends T> src) {
        for (int i = 0; i < src.size(); i++)
            dest.set(i, src.get(i));
    }
}
```

Type erasure

```

1 public class Node<T> {
2     private T data;
3     private Node<T> next;
4
5     public Node(T data, Node<T> next) {
6         this.data = data;
7         this.next = next;
8     }
9
10    public void setData(T data) {
11        this.data = data;
12    }
13
14    public T getData() {
15        return data;
16    }
17 }
18
19 public class Main {
20     public static void main(String[] args) {
21         Node<Integer> node = new Node<>(10, null);
22         Integer i = node.getData();
23         node.setData(20);
24     }
25 }
26

```

```

1 public class Node {
2     private Object data;
3     private Node next;
4
5     public Node(Object data, Node next) {
6         this.data = data;
7         this.next = next;
8     }
9
10    public void setData(Object data) {
11        this.data = data;
12    }
13
14    public Object getData() {
15        return this.data;
16    }
17 }
18
19 public class Main {
20     public Main() {
21     }
22
23     public static void main(String[] args) {
24         Node node = new Node(10, (Node)null);
25         Integer i = (Integer)node.getData();
26         node.setData(20);
27     }
28 }
29

```

Figure 5: Type erasure

Type erasure

- Replace unbound type E with Object
- Add typecasting to all method calls
- Object has no information of type it was created with¹
- `ClassCastException` will be called if constraints are violated

¹except for class that inherits type-specified generic

LinkedList

- Allocated as linked list of many objects
- Every inserted object needs a wrapper object, so the number of objects in memory are at least **twice** the number of elements in list.
- Fast **append**: $\Theta(1)$
- Fast **insert** but only when given a **preceeding node** in the list.
- Fast **delete** but under the same conditions as **insert**
- Slow **get**: $O(n)$
- Does not need to grow - **append** is not lagging from time to time.

LinkedList - pros

- LinkedList implements Deque which gives us nice stack and fifo methods: pollFirst, pollLast, peekFirst, peekLast

LinkedList - cons

- Every item in LinkedList is of a class LinkedList.Item which has much overhead.

mark hash gcAge tLock tId	klassPtr ref to class	length length ptr	padding
4-8 bytes	4-8 bytes	4 bytes	up to mod arch

Remove all strings containing letter a

```
List<String> lst = new ArrayList<>();  
lst.addAll(Arrays.asList("marek", "basia", "kotek", "paw"  
    ));  
// remove all words containing letter a
```

Remove all strings containing letter a

```
List<String> lst = new ArrayList<>();
lst.addAll(Arrays.asList("marek", "basia", "kotek", "paw"
    ));
// remove all words containing letter a
for (String l : lst) {
    if (l.contains("a")) {
        lst.remove(l);
    }
}
```

Remove all strings containing letter a

```
List<String> lst = new ArrayList<>();
lst.addAll(Arrays.asList("marek", "basia", "kotek", "paw"
    ));
// remove all words containing letter a
for (String l : lst) { //
    ConcurrentModificationException
    if (l.contains("a")) {
        lst.remove(l);
    }
}
```

Hello
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Algorithmics
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List
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Maps
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Excercises
oooooooo

Summary
ooooo

[LinkedList](#)

Listlterator

How much do the object weight in Java

```
class A {}
```

```
class B extends A {}
```

```
class C {  
    boolean a;  
}
```

How much do the object weight in Java - compressed pointers

```
java -XX:+UseCompressedOops
```

```
sizeof A = 16
```

```
sizeof B = 16
```

```
sizeof C = 16
```

```
sizeof Boolean = 16
```

```
sizeof Char = 16
```

```
sizeof Integer = 16
```

```
sizeof Long = 24
```

```
sizeof class java.util.LinkedList\$$Node = 24
```

How much do the object weight in Java - *not* compressed pointers

java -XX:-UseCompressedOops

sizeof A = 16

sizeof B = 16

sizeof C = 24

sizeof Boolean = 24

sizeof Char = 24

sizeof Integer = 24

sizeof Long = 24

sizeof class java.util.LinkedList\ \$Node = 40

How much do the object weight in Java

```
public class C1 {  
    boolean a;  
    int b;  
} // size 24
```

```
public class C2 {  
    boolean a;  
    int b;  
    int c;  
} // size 24
```

```
public class C3 {  
    boolean a;  
    int b;  
    int c;  
} // size 32
```

HashMap

- Fast finding element by the key (further referred to as: **lookup**): $O(1)$.²
- Fast putting a key-value pair (further referred to as: **insert**): ($O(1)$, pessimitically $n = \text{size}$, when map needs to grow).
- Fast remove: $O(1)$ (HashMap does not shrink).
- Values with the same hash stored in *LinkedList*, collisions may occur.

²When time is longer, in properly setup map (i.e. number of buckets $>$ size), it means that collisions occur. This can be the sign of incorrect hash function.

HashMap

$x.\text{equals}(y) \implies x.\text{hashCode}() == y.\text{hashCode}()$

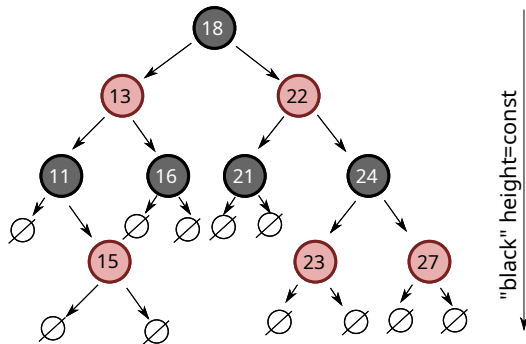
$x.\text{hashCode}() == y.\text{hashCode}() \not\Rightarrow x.\text{equals}(y)$

TreeMap

- Implemented by Red-Black tree.
- Keys are sorted.
- Quite good lookup: $O(\log n)$
- Quite good insert: $O(\log n)$
- Quite good delete: $O(\log n)$
- Memory consumption: $O(n)$.

Red-black tree

For each path count of black nodes ("black length") is constant



Algorithms visualised

<https://www.cs.usfca.edu/~galles/visualization/java/visualization.html>

How to write tests

- Create test that checks constraint on empty or broken code
- Run the test to confirm that it fails (if it passes, the checks are incorrect)
- During the test compilation it might be needed to create more empty methods - this is the Development part in TDD
- Fix the code and fill the empty methods
- Tests should pass

Excercise 1

- Open the project, and enter the t1 package
- in `src/main/java/t1` we have code that we need to fix
- in `src/test/java/t1` we have tests.

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 - ① Run `FIFOQueueImplTest` - it fails.

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 - 2 Implement `FIFOQueue` using appropriate Java's list implementation.

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 - 3 Run `LIFOQueueImplTest` - it should pass. How long does it run?
- Run the efficiency test `PerformanceTestsForT1` - all should

Excercise 2

- see `src/main/java/t2` and `src/test/java/t2`
- First focus on `first_exercise_testGBPtoPLN` - right click on the test name and *run testname*.

Excercise 2

- see `src/main/java/t2` and `src/test/java/t2`
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- why there is no such currency conversion, providing that we have added it in the map in `initStatic` method?

Excercise 2

- see `src/main/java/t2` and `src/test/java/t2`
- First focus on `first_exercise_testGBPtoPLN` - right click on the test name and *run testname*.
- why there is no such currency conversion, providing that we have added it in the map in `initStatic` method?
- Next focus on `second_excercise_testPLNtoGBP`

Excercise 2

- see `src/main/java/t2` and `src/test/java/t2`
- First focus on `first_exercise_testGBPtoPLN` - right click on the test name and *run testname*.
- why there is no such currency conversion, providing that we have added it in the map in `initStatic` method?
- Next focus on `second_excercise_testPLNtoGBP`
- Why the test is failing? What is the printed result?

Exercise 3

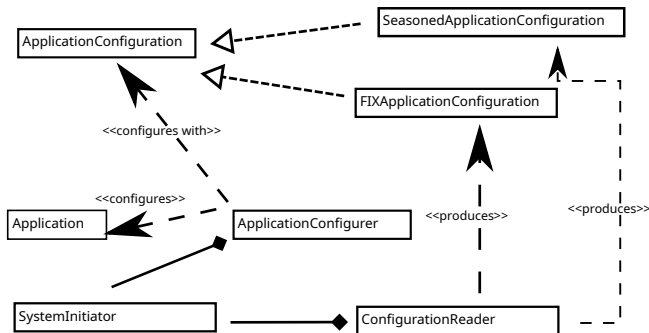
- see `src/main/java/t3` and `src/test/java/t3`
- List must not grow, must throw proper exceptions.
- List must contain the java's Array as storage for the items
- How to initialize the storing array for any T? Why `new T[]` is not working?

Excercise 4

- see `src/main/java/t4` and `src/test/java/t4`
- Why the `assertNull` does not work?

Excercise 5

- see `src/main/java/t5` and `src/test/java/t5`
- Uncomment code in the test
- Why the code does not compile? How can we fix it?



Excercise 6

- see `src/main/java/t6` and `src/test/java/t6`
- `testRunningTimeEasy` — focus on time of execution, why it is so slow

Excercise 6

- see `src/main/java/t6` and `src/test/java/t6`
- `testRunningTimeEasy` — focus on time of execution, why it is so slow
- `testMemEfficiencyEasy` — try to experiment with memory and drop the required by test below 42. You will see the test fails.

Excercise 6

- see `src/main/java/t6` and `src/test/java/t6`
- `testRunningTimeEasy` — focus on time of execution, why it is so slow
- `testMemEfficiencyEasy` — try to experiment with memory and drop the required by test below 42. You will see the test fails.
- `testMemEfficiencyIntermediate` — try to achieve the 32 MB memory by using list other than `LinkedList`

Excercise 6

- see `src/main/java/t6` and `src/test/java/t6`
- `testRunningTimeEasy` — focus on time of execution, why it is so slow
- `testMemEfficiencyEasy` — try to experiment with memory and drop the required by test below 42. You will see the test fails.
- `testMemEfficiencyIntermediate` — try to achieve the 32 MB memory by using list other than `LinkedList`
- `testMemEfficiencyExpert` — try to achieve 26 MB by initializing the list with the number of items known a priori.

Excercise 6

- see `src/main/java/t6` and `src/test/java/t6`
- `testRunningTimeEasy` — focus on time of execution, why it is so slow
- `testMemEfficiencyEasy` — try to experiment with memory and drop the required by test below 42. You will see the test fails.
- `testMemEfficiencyIntermediate` — try to achieve the 32 MB memory by using list other than `LinkedList`
- `testMemEfficiencyExpert` — try to achieve 26 MB by initializing the list with the number of items known a priori.
- `testMemEfficiencyMaster` — try to achieve 12 MB by writing your own suited `int`-based collection.

Hello
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Algorithmics
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List
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Maps
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Excercises
oooooooo

Summary
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What have we learned today?

Summary

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Test Driven Development

- Always start with test

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- Even though, using incorrect implementation, we can break things up

What have we learned today?

Complexity as a gift from Java

- Big O notation, small o notation, theta θ notation
- Java's *sophisticated* collections give us access to best algorithm implementation
- Even though, using incorrect implementation, we can break things up
- Even though, sometimes we may need to have our own implementation (read: **use libraries**)

What have we learned today?

Generics

- One code for many usages
- ? extends and ? super
- Type erasure will clear typed info

Hello
oo

Algorithmics
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List
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Maps
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Excercises
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Summary
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What have we learned today?

Thank you