

# MysticAndCandies

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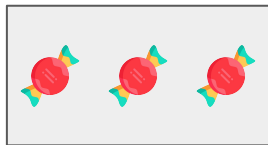
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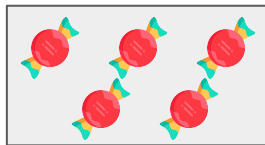
# The Problem

- mystic\_tc wants to eat a certain number of candies
- Many closed boxes with random # of candies in each
- Open minimum # of boxes
- Must eat all candies in every box opened
- The catch: only know ranges in each box

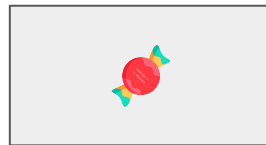
2 - 3



5 - 10



1 - 4



# The Method

X: # candies to eat  
C: total candies  
N: number of boxes

- Class name: MysticAndCandies
- Method name: minBoxes
- Parameters:

C	int	Total # of candies	$\text{Sum}(\text{Lows}) \leq C \leq \text{Sum}(\text{Highs}), 1 \leq C \leq 50$
X	int	# of candies to eat	$1 \leq X \leq C$
low	int[]	Lower bounds	$\text{low}[i] \leq \text{high}[i], \text{size}(\text{high}) = \text{size}(\text{low})$
high	int[]	Upper bounds	$\dots 1 \leq \text{size}(\text{high}) \leq 10,000,000$

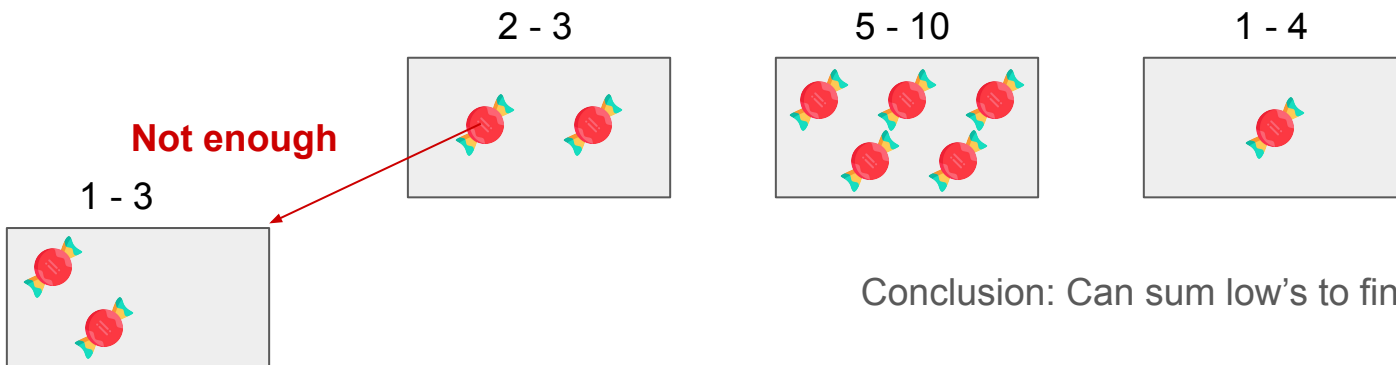
- Returns: int
- Signature: int MysticAndCandies::minBoxes(int C, int X, int[] low, int[] high)

# The Solution

X: # candies to eat  
C: total candies  
N: number of boxes

- Consider:
  - mystic\_tc chooses a solution
  - # candies in solution  $\geq X$
  - To guarantee the above, one of two cases  $\rightarrow$

1. Boxes in solution are at their min capacity



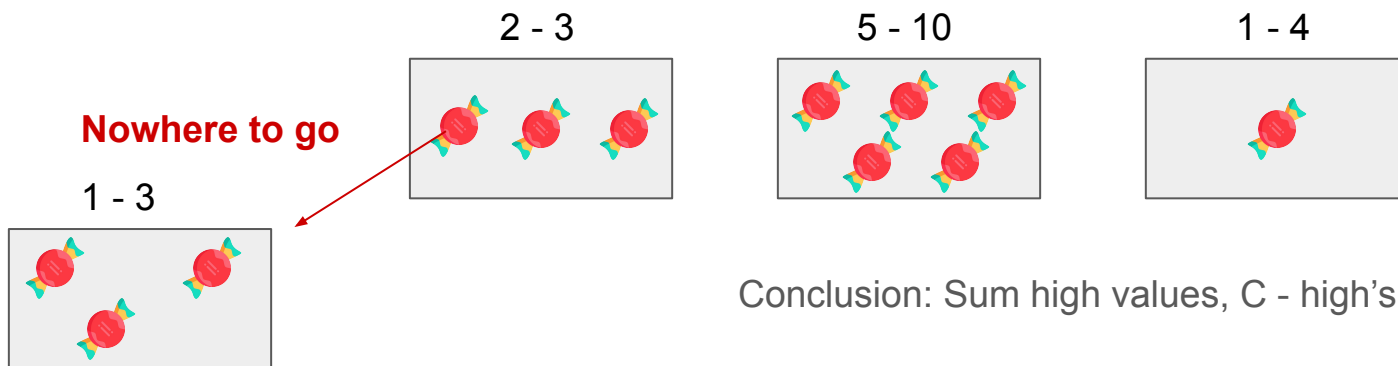
Conclusion: Can sum low's to find # boxes

# The Solution

X: # candies to eat  
C: total candies  
N: number of boxes

- Consider:
  - mystic\_tc chooses a solution
  - # candies in solution  $\geq X$
  - To guarantee the above, one of two cases  $\rightarrow$

2. Boxes outside solution are at their max capacity



Conclusion: Sum high values, C - high's not in the solution

# The Solution

X: # candies to eat  
C: total candies  
N: number of boxes

Sort low's in descending order

Sort high's in ascending order

For i from 0 to N:

sum += low[i]

If  $\text{sum} \geq X$ , minBoxes = i+1

Subset of boxes at  
minimum capacity

For i from 0 to N:

sum += high[i]

If  $C - \text{sum} \leq X$ , minBoxes = N-i-1

Rest of boxes at  
maximum capacity

Return smaller minBoxes between the two

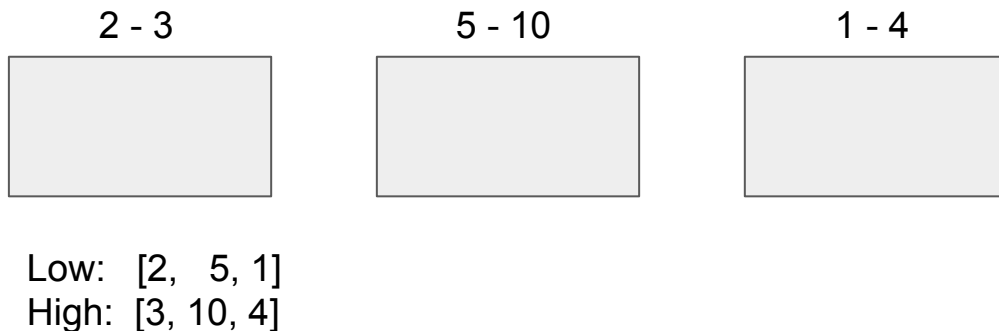
# An Example

X: # candies to eat  
C: total candies  
N: number of boxes

- **X = 5, C = 15**

1. Sort low: [2, 5, 1]  $\rightarrow$  [5, 2, 1]
2. Sort high: [3, 10, 4]  $\rightarrow$  [3, 4, 10]
3. Sum low's: 5,  $5 + 2 = 7$ ,  $5 + 2 + 1 = 8 \rightarrow$  [5, 7, 8]

$5 \geq X \rightarrow$  **minBoxes = 1** (i = 0)



# An Example

X: # candies to eat  
C: total candies  
N: number of boxes

- **X = 9, C = 15**

1. Sort low:  $\rightarrow [5, 2, 1]$

2. Sort high:  $\rightarrow [3, 4, 10]$

3. Sum low's:  $\rightarrow 5, 5 + 2 = 7, 5 + 2 + 1 = 8 \rightarrow [5, 7, 8] \rightarrow$  **None are enough!**

4. Sum high's:  $15 - 3 = 12, 15 - (3 + 4) = 8, 15 - (3 + 4 + 10) = -2 \rightarrow [12, 8, -2]$

$8 \leq X \rightarrow$  **minBoxes = 1** ( $i = 1$ )

2 - 3



5 - 10



1 - 4



Low: [2, 5, 1]

High: [3, 10, 4]



# The Complexity

X: # candies to eat  
C: total candies  
N: number of boxes

Sort low's in descending order

Sort high's in ascending order

For i from 0 to N:

Sum += low[i]

If sum  $\geq$  X, minBoxes = i+1

For i from 0 to N:

Sum += high[i]

If C-sum  $\leq$  X, minBoxes = N-i-1

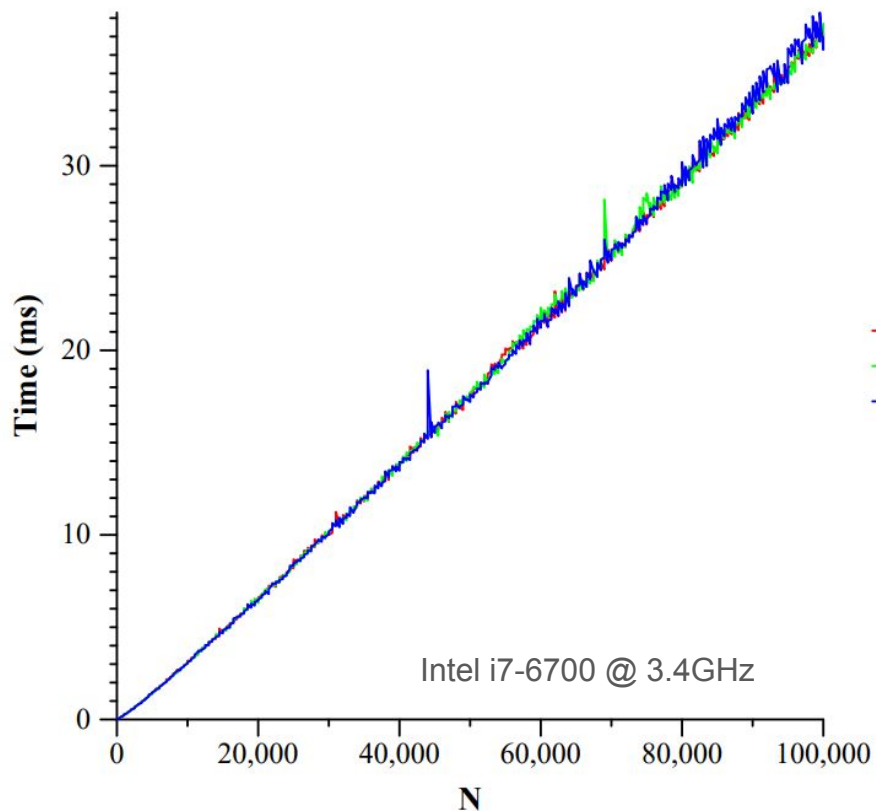
Return smaller minBoxes between the two

std::sort  $\rightarrow O(N \log N)$  average case

for i from 0 to N  $\rightarrow O(N)$

$\rightarrow O(2N + 2N \log N) \rightarrow O(N \log N)$

# The Data



X: # candies to eat  
C: total candies  
N: number of boxes

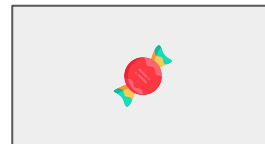
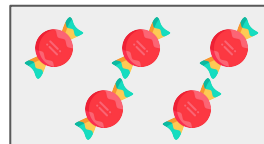
→  **$O(N \log N)$**

Making it faster:

1. Use a priority queue instead of sorting  
→ `std::priority_queue`  $O(N)$  from vector
2. Base case  $C = X$   
→ Gets rid of worse case

# The Statistics

- 742 people opened
- 501 people submitted a solution (67.5%)
- 73.05% of submissions were correct
- Best score: 293.0
- Average score: 199.05
- Best time: 4:40 (5:39 in C++)
- Average time: 26:27



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