### Approximate Query Processing

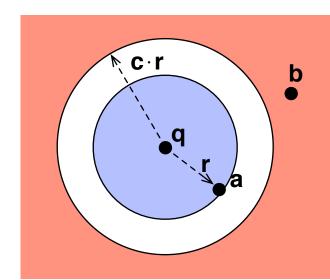
- Space Partitioning-based
  - Tree
  - Encoding
  - Locality Sensitive Hashing
- Graph-based Methods

#### Notes:

- Recent works mainly in the Database area
- Prefer ease of exposition over rigor
- Categorization is not fixed/unique

- □ From the perspective of collision probability
  - (Ordinary) hash function h:
    - $Pr[h(x) = h(y)] = \varepsilon, \text{ if } x \neq y$

c.f., Cryptographic hash functions Pr [ h(x) = h(y) ] =  $2^{-m}$ , if Hamming(x, y) = 1

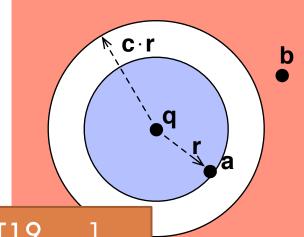


- □ From the perspective of collision probability
  - (Ordinary) hash function h:
    - $Pr[h(x) = h(y)] = \varepsilon, \text{ if } x \neq y$
  - LSH
    - Pr[h(x) = h(y)] increases with locality
    - $\blacksquare$  Randomness comes from r.v.  $h \in H$

```
(r_1, r_2, p_1, p_2)-sensitive [IM98]
```

- $Pr[h(x) = h(y)] \ge p_1$ , if  $dist(x, y) \le r_1$
- $Pr[h(x) = h(y)] \le p_2$ , if  $dist(x, y) \ge r_2$

$$Pr[h(x) = h(y)] = sim(x, y) [C02] too narrow$$



c.f., Cryptographic hash functions

Pr  $[h(x) = h(y)] = 2^{-m}$ , if

Hamming(x, y) = 1

Can be generalized [SWQZ+14, ACPS18, CKPT19, ...]

- Equality search
  - Index: store o into bucket h(o)
  - Query:
    - retrieve every o<sub>i</sub> in the bucket h(q)
    - verify if o<sub>i</sub> = q

$$Pr[h(q) = h(o)] = 1/B$$

- Equality search
  - Index: store o into bucket h(o)

Pr[h(q) = h(o)] = 1/B

- Query:
  - retrieve every o<sub>i</sub> in the bucket h(q)
  - verify if o<sub>i</sub> = q
- □ LSH

c.f., [PIM12] for the <u>rigorous</u> QP procedure

- $\square \forall h \in LSH$ -family,  $\Pr[Q(h(q)) = Q(h(o))] = f(Dist(q, o))$ 
  - Q(): quantization (not essential)
  - "Near-by" points have more chance of colliding with q than "far-away" points
- □ Similar index & query procedures, with a weak
  probabilistic guarantee
  → Repeat to boost the guarantee

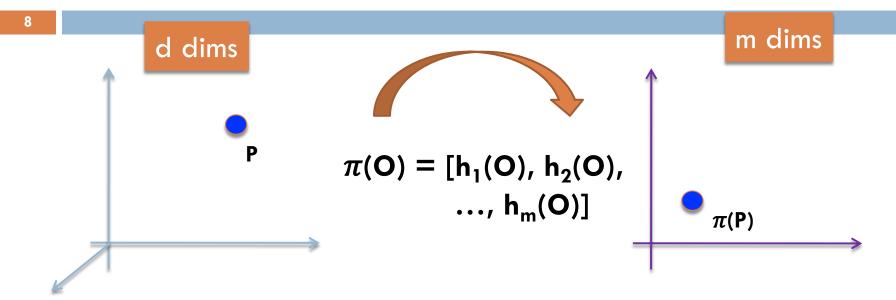
#### LSH Families

- Many are known
  - Arr L<sub>p</sub> (0 \leq 2): use p-stable distribution to generate the projection vector
    - For L<sub>2</sub>, just use random Gaussian vector
    - Other families exists, e.g., sparse random projection
  - Angular distance (arccos): SimHash
  - Jaccard: minhash (based on random permutation)
  - Hamming:
    - random projection
    - covering LSH

#### Comments

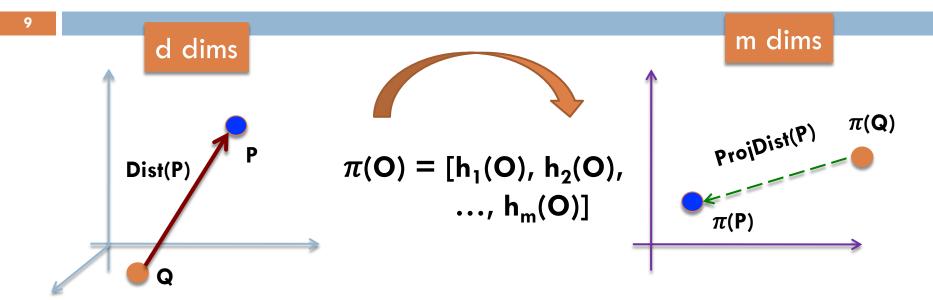
- New queries can be reduced to known LSH cases
  - Maximum inner product search (MIPS)
  - Set containment
  - Group aggregated query
- Related to various distortion-bounded embedding
  - Edit distance: CGK-embedding to Hamming with O(K) distortion

#### Probabilistic Mapping



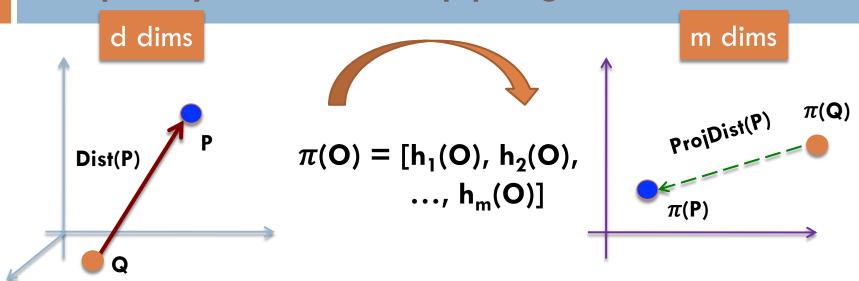
 Probabilistic, linear mapping from the original space to the projected space

#### Probabilistic Mapping



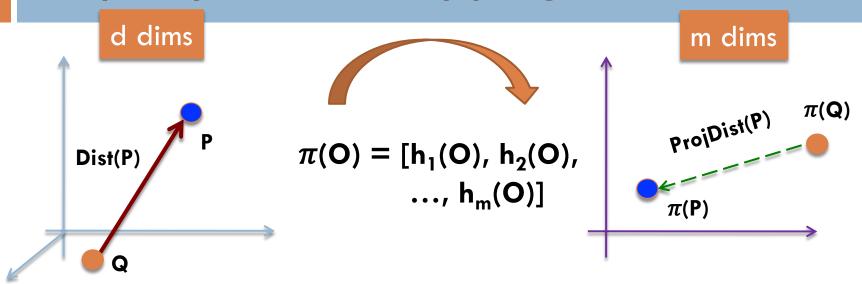
- Probabilistic, linear mapping from the original space to the projected space
- What about the distances (wrt Q or  $\pi(Q)$ ) in these two spaces?

# Probabilistic Distance Tracking Property of the Mapping



- □ ProjDist(P)<sup>2</sup> ~ Dist(P)<sup>2</sup> \*  $\chi^2_m$  [SWQZ+14]
  - ProjDist(P)<sup>2</sup> can be computed (incrementally) from  $h_i(P)$  and  $h_i(Q)$  due to the linearity of the hash function
  - Can be generalized to other p-stable LSH functions

# Probabilistic Distance Tracking Property of the Mapping



LSH provides a probabilistic distance-preserving mapping between the two spaces

Johnson & Lindenstrauss Lemma [JL84] only works for L2 and induces a method that requires more space than LSH [AIR18]

#### Roadmap

- Roadmap
  - Practical LSH methods (i.e., linear index complexity)
  - Data-dependent LSH methods

#### Practical Variations of LSH

- Easy to relax the LSH method in practice at the cost of no worst-case guarantees
  - E2LSH: use fewer number of random projections
  - Multiprobe LSH (entropyLSH and other variants): space-time tradeoff
  - LSH in practice: use empirically tuned parameters (k, l)
  - SRS: a space-saving LSH index with in-memory or disk implementations
  - QALSH: use multiple B-trees for the index and use collision counting strategy based on LSH
  - HD-index: space filling curves as pseudo-LSH functions
  - SK-LSH: Replace LSB-tree/forest by a dimension-wise linear mapping

#### Data-sensitive Hashing

- LSH is data-insensitive
  - Indexing hyper-parameters determined by the shape of the data only
  - Indexing parameters are randomly generated
- Efforts to make data-sensitive, LSH-like methods
  - □ [AR15]
    - lacktriangle Aim: break the lower bounds of ho

c.f., [AIR18]

- DSH
- OPFA / NeOPFA
- Learning-to-hash methods
  - NSH [PCM15]
  - [LYZX+18] and many in the ML/CV communities

#### DSH [GJLO14]

Learn a family of (hash) functions, H, that preserves
 kNN of queries

- 1. Training data:  $\mathbf{W}_{ij} = \begin{cases} 1 & \text{, if } o_j \in kNN(q_i) \\ -1 & \text{, if } o_j \not\in kNN(q_i) \land o_j \text{ is sampled} \\ 0 & \text{otherwise.} \end{cases}$ 
  - their k-NN objects (+ve)
  - samples non-c\*k-NN objects (-ve)
- 2. Function family:
  - Thresholded linear functions  $h(\mathbf{x}; \mathbf{a}) = \operatorname{sgn}(\mathbf{a}^{\top} \mathbf{x})$
- 3. Learn one hash function

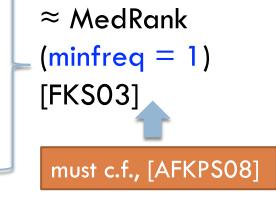
$$\arg\min_{h} \sum_{i} \sum_{j} \ell(q_i, o_j) \mathbf{W}_{ij} \quad \text{, where } \ell(q, o) = (h(q) - h(o))^2$$

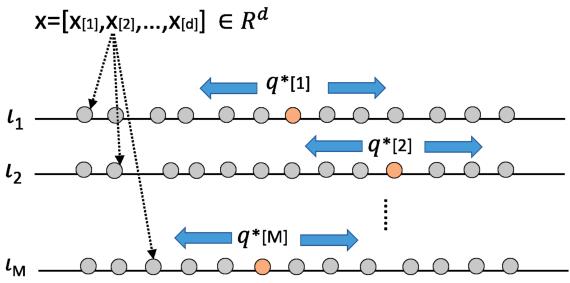
#### DSH [GJLO14]

- Learn a family of (hash) functions, H, that preserves
  kNN of queries
  - 4. Learn multiple hash functions
    - Multiplicative updates on W<sub>ij</sub>
      - Increase W<sub>ij</sub> if incorrectly classified
      - Decrease W<sub>ii</sub> if correctly classified
    - (Under some assumptions) obtain H that satisfies the (k, ck,  $p_1$ ,  $p_2$ )-sensitive property for the training data
      - if  $o \in NN(q, k)$ , then collision probability from  $H \ge p_1$
      - if o  $\notin$  NN(q, ck), then collision probability from H  $\leq$  p<sub>2</sub>

#### Learned ANN Index [LZSW+20]

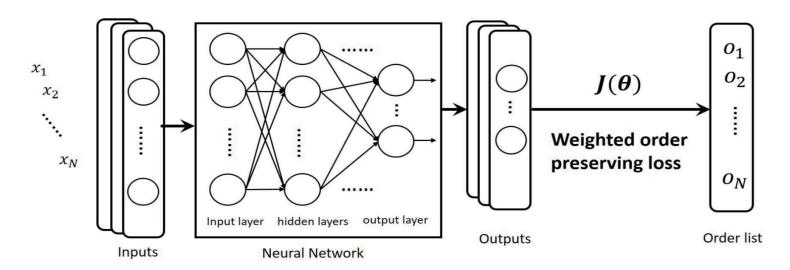
- Focus on external I/O
  - Use B-trees and maximize the use of sequential I/Os
- Scheme:
  - $\blacksquare$  H: R<sup>d</sup>  $\rightarrow$  R<sup>M</sup>
  - Index each dimension of H(X) in a clustered B-tree
- Query processing
  - Collect candidates
     on each of the M
     projected
     dimensions
  - When T candidates are seen on all M lists, rerank them and return top-k





#### Function family

- Consider
  - linear functions
    - $\blacksquare H(x)[m] = w_m^T x$
  - non-linear functions



19

Consider the linear functions:  $H(x)[m] = w_m^T x$ 

- □ Goal:
  - Encourage segment-order preserving mappings

Part of the Loss function

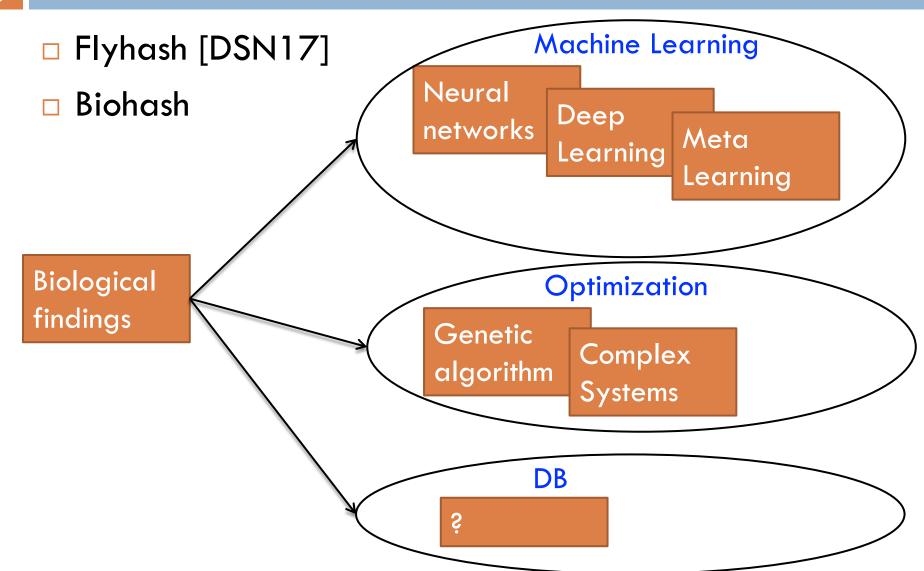
$$J^*(\mathbf{w}_m) = \sum_{i=1}^L \sum_{\tilde{x} \in l_i^o} \mathbf{1}_{r(\tilde{x})} \in [t \cdot (i-1), t \cdot i)$$

Continuous Relaxation

mapped x in the i-th segment

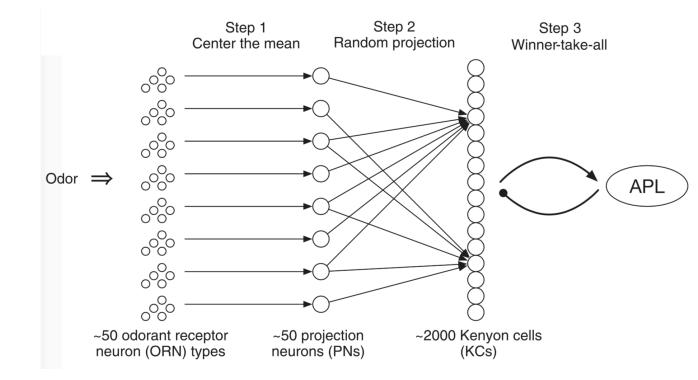
x in the i-th segment

# Biology Inspired Hashing



### FlyHash

 The fly olfactory circuit generates a lowoverlapping, sparse neuron activation pattern when an odor is presented

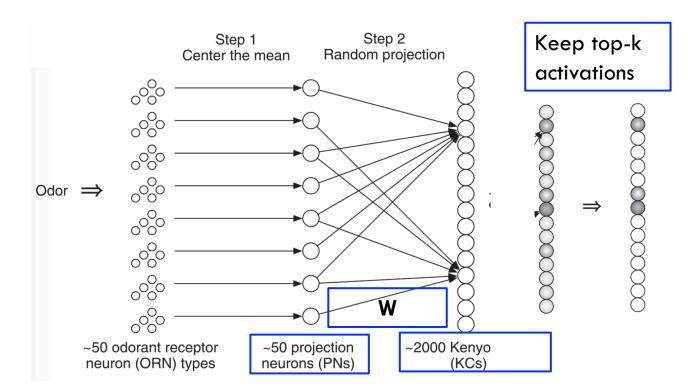


# FlyHash

Difference with LSH

 $\mathbf{z} = \sigma(\mathbf{W}\mathbf{x})$ 

- W is a sparse binary random matrix
- Dimensionality expansion !!
- Sparsification
- L2 distance approximately preserved in expectation

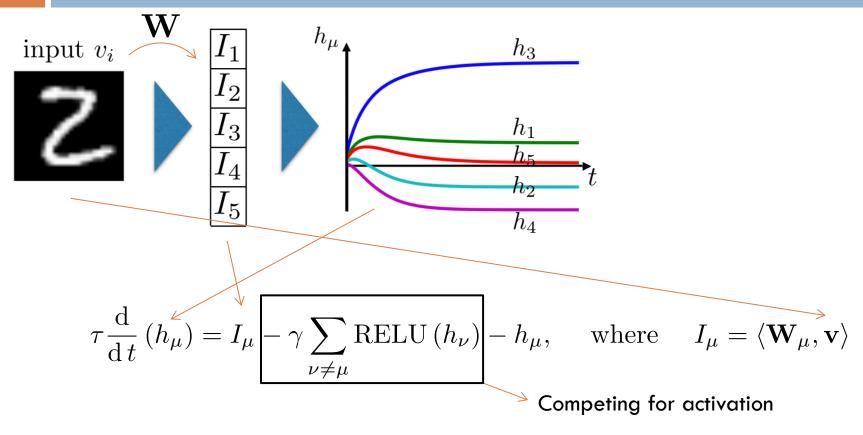


**Enable** 

#### [KH19]

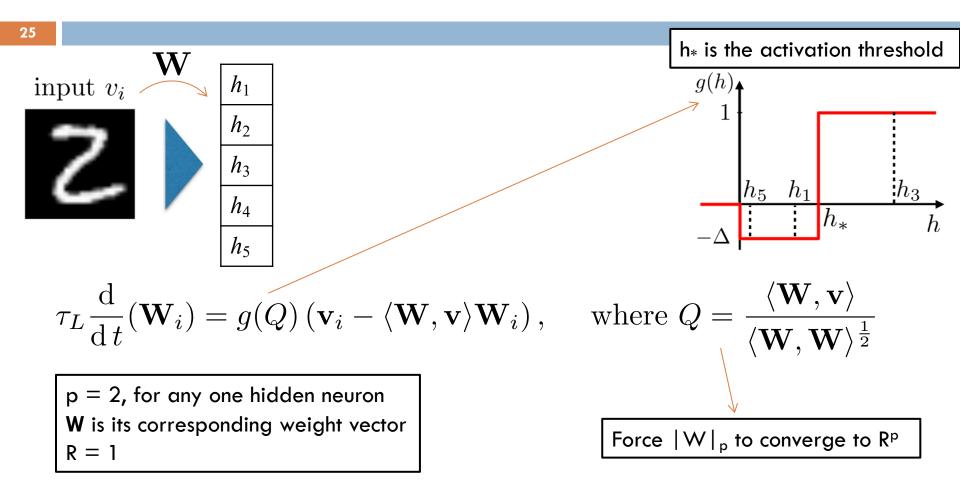
- Unsupervised learning inspired by biological synaptic plasticity rules
- Overview
  - □ (Given W) Stabilizing the hidden competing neurons
  - Learning the projection matrix W

#### Learning h



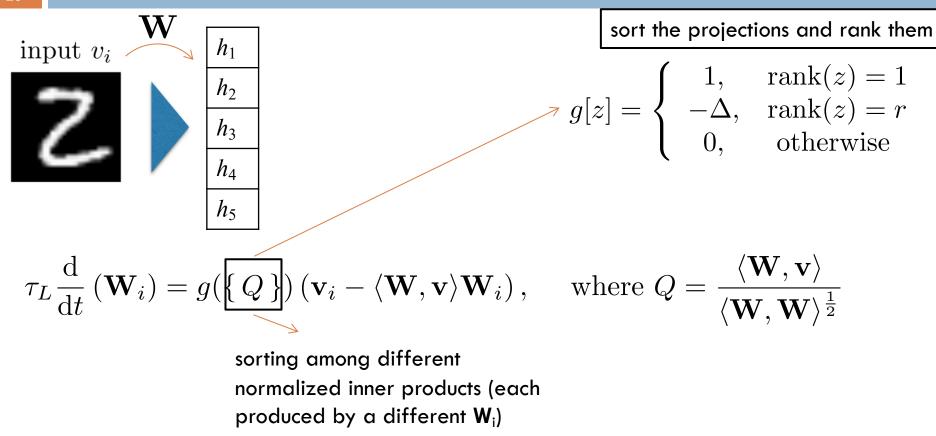
 Fixing the W, the dynamical equation will converge to a stable hidden vector h

### Learning W



 Fixing the h, the dynamical equation will converge to a final weight matrix W

### BioHash - Learning W



The rest is the same as FlyHash (i.e., WTA sparsification)