

## Computing Theory HW Set 3

### Chapter 7

4. Construct npda's that accept the following languages on  $\Sigma = \{a, b, c\}$ .

(a)  $L = \{a^n b^{2n} : n \geq 0\}$ .

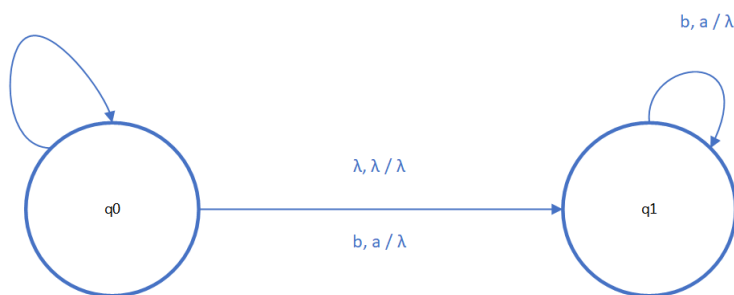
$((q_0, \lambda, \lambda) = (q_1, \lambda))$

$((q_0, a, \lambda) = (q_0, aa))$

$((q_0, b, a) = (q_1, \lambda))$

$((q_1, b, a) = (q_1, \lambda))$

$a, \lambda / aa$



(c)  $L = \{a^n b^m c^{n+m} : n \geq 0, m \geq 0\}$ .

$((q_0, \lambda, \lambda) = (q_2, \lambda))$

$((q_0, a, \lambda) = (q_0, a))$

$((q_0, c, a) = (q_2, \lambda))$

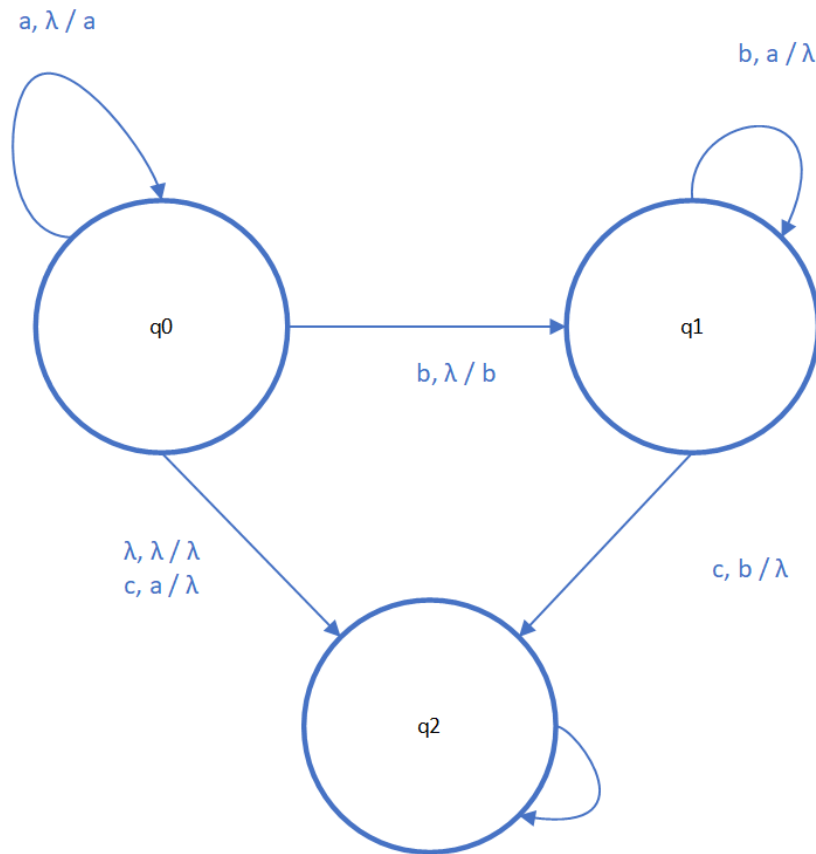
$((q_0, b, \lambda) = (q_1, b))$

$((q_1, b, \lambda) = (q_1, b))$

$((q_1, c, b) = (q_2, \lambda))$

$((q_2, c, b) = (q_2, \lambda))$

$((q_2, c, a) = (q_2, \lambda))$



(d)  $L = \{a^n b^{n+m} c^m : n \geq 0, m \geq 1\}$ .

$((q_0, a, z) = (q_1, 1z)$

$(q_0, b, z) = (q_3, 1z)$

$(q_1, a, 1) = (q_1, 11)$

$(q_1, b, 1) = (q_2, \lambda)$

$(q_2, b, 1) = (q_2, \lambda)$

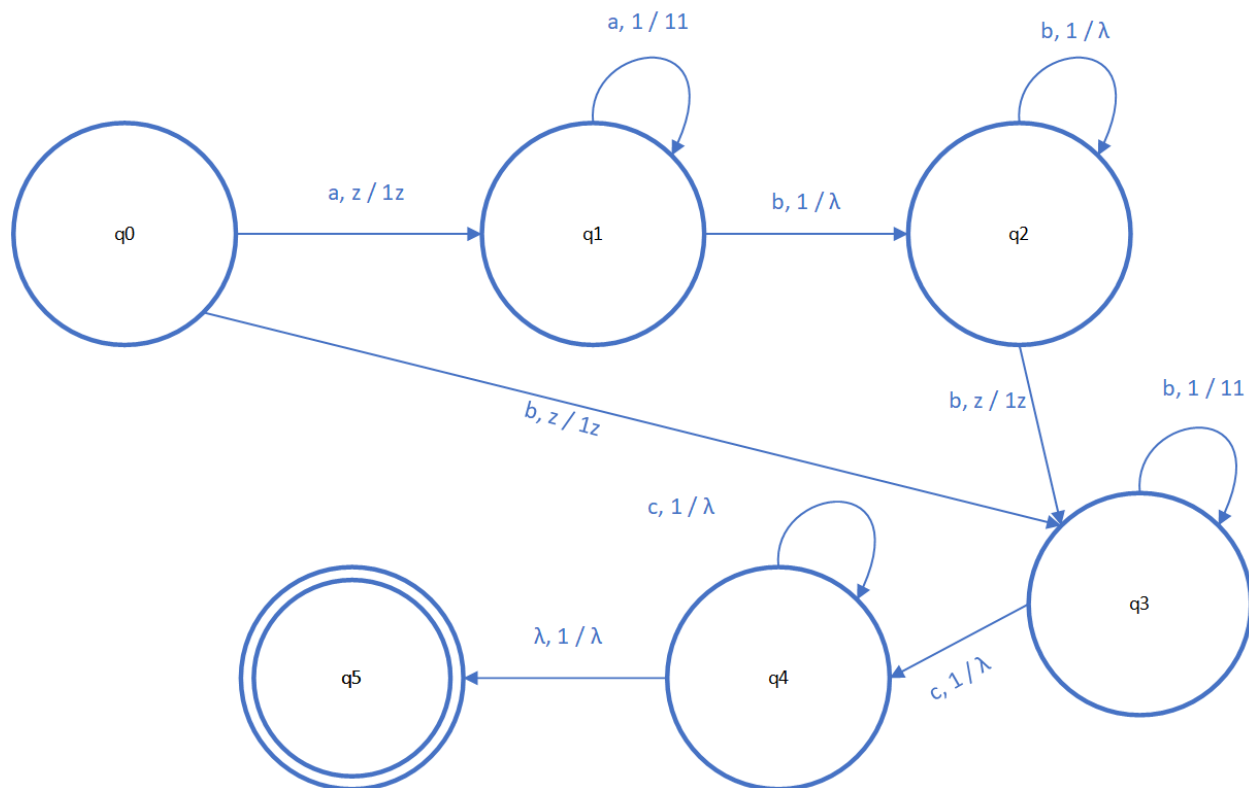
$(q_2, b, z) = (q_3, 1z)$

$(q_3, b, 1) = (q_3, 11)$

$(q_3, c, 1) = (q_4, \lambda)$

$(q_4, c, 1) = (q_4, \lambda)$

$(q_4, \lambda, 1) = (q_5, \lambda)$



(f)  $L = \{a^n b^m : n \leq m \leq 3n\}$ .

$(q_0, a, z) = (q_0, az)$

$(q_0, a, z) = (q_0, aaz)$

$(q_0, a, z) = (q_0, aaaz)$

$(q_0, a, a) = (q_0, aa)$

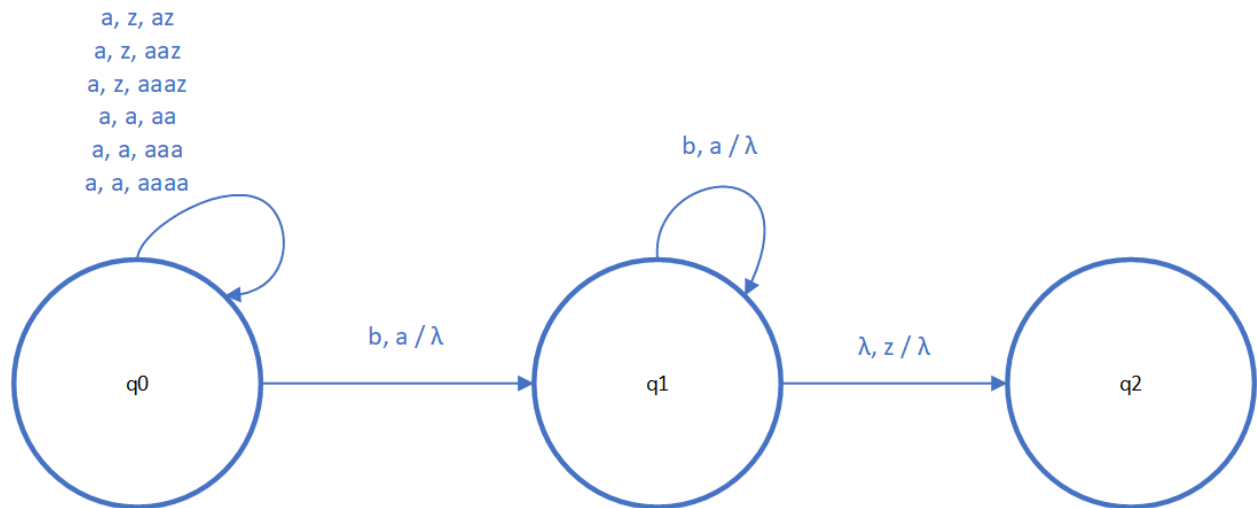
$(q_0, a, a) = (q_0, aaa)$

$(q_0, a, a) = (q_0, aaaa)$

$(q_0, b, a) = (q_1, \lambda)$

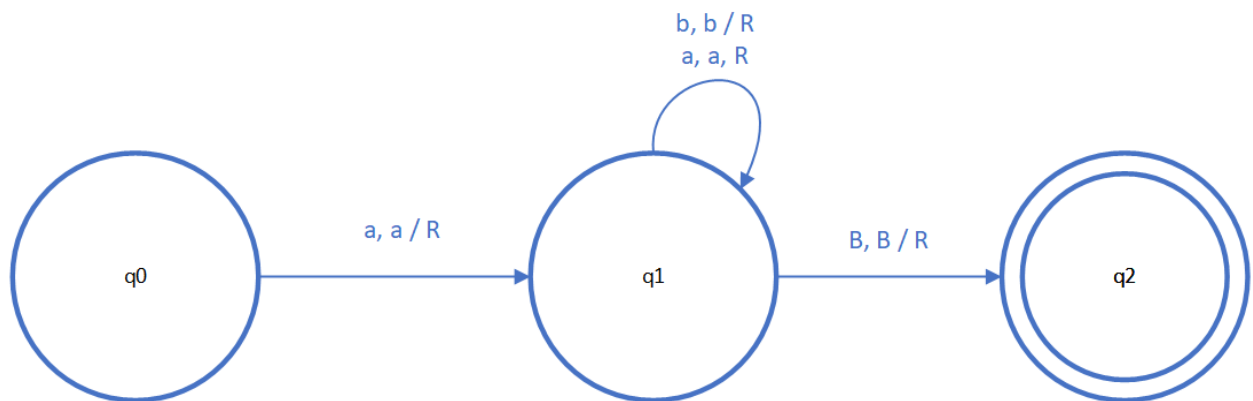
$(q_1, b, a) = (q_1, \lambda)$

$(q_1, \lambda, z) = (q_2, \lambda)$



## Chapter 9

2. Design a Turing machine with no more than three states that accepts the language  $L(a(a+b)^*)$ . Assume that  $\Sigma = \{a, b\}$ . Is it possible to do this with a two-state machine?



It is not possible with just two states, because two states are required to validate inputs 'a' and  $(a+b)^*$  and after that you need an accept state.

5. What language is accepted by the Turing machine whose transition graph is in the figure below?

$L = ab^* + bb^* a \dots$  strings either start or end with a and rest of the elements are all b

7. Construct Turing machines that will accept the following languages on  $\{a, b\}$ .

(a)  $L = L(aba^*b)$ .

