

## Test 2

Last Name Davis First Name Trevor Grade       /100

1. (10%) Give a formal definition of Context Free Grammar  $G = (V, T, S, P)$ . Define each of the 4 major components in a sentence, and form of  $P$  in set notation.

A grammar  $G = (V, T, S, P)$  is said to be context-free if all productions in  $P$  have the form  $A \rightarrow x$ , where  $A \in V$  and  $x \in (V \cup T)^*$

$V$  is the set of non-terminal symbols

$T$  is a set of terminals and  $V \cap T$  is null. They don't appear on left side of production

$S$  is the start symbol. It appears in the initial string

$P$  is the set of rules. They are rules for replacing non terminal symbols.

2. (40 %)

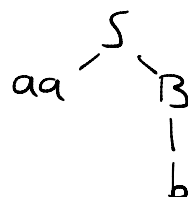
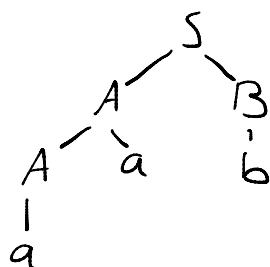
- (a) Show that the following grammar is ambiguous by using parsing trees.

$S \rightarrow AB \mid aaB$

$A \rightarrow a \mid Aa$

$B \rightarrow b$

String:  $aaab$



The grammar is ambiguous because we are able to draw 2 different parse trees for the same string

- (b) Find a CFG for the following language:  $L = \{a^n b^m : n \neq m\}$

$$n \neq m = n > m \cup n < m$$

$n > m$   
 $S_1 \rightarrow aS_1b \mid a$   
 $S_1 \rightarrow aS_1$

$n < m$   
 $S_2 \rightarrow aS_2b \mid b$   
 $S_2 \rightarrow S_2b$

$S \rightarrow S_1 \mid S_2$

$S_1 \rightarrow aS_1b \mid a \mid aS_1$

$S_2 \rightarrow aS_2b \mid b \mid S_2b$

3. (40%) Prove the following language is non-regular using Pumping Lemma step by step and lead to a contradiction at the end:  $L = \{a^n b^{n+3} : n \geq 0\}$

Assume that  $L$  is regular

String:  $\{abbbb, aabbbbb, aaa bbbbbb\}$

$$w = xyz$$

Case 1:  $y = a^i$

$$xyz = \underset{x}{a} \underset{y}{aa} \underset{z}{bbbbbb}$$

$$xy^2z = aaaaa bbbbbb \quad w \notin L$$

Case 2:  $y = a^r b^m$

$$xyz = \underset{x}{aaa} \underset{y}{abb} \underset{z}{bbb}$$

$$xy^2z = aaa abb abb bbb \quad w \notin L$$

Case 3:  $y = b^m$

$$xyz = \underset{x}{aaa} \underset{y}{bbb} \underset{z}{bbb}$$

$$xy^2z = aaa bbb bbb bbb \quad w \notin L$$

Hence because of the Pumping lemma,  $L$  is not a regular language.

4. (10 %) Prove  $L = \{a^n b^k : n \geq 2, k \leq 14\}$  is regular.

This language is regular because there exists a regular expression which defines it as well as an FA

$$aaa^*(\lambda + b + bb + \dots + b^{14})$$