Cache Me If You Can: Accuracy-Aware Inference Engine for DP Data Exploration

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SETUP

- Interactive DP.
- Data owner specifies privacy budget ε.
- Data analysts care about accuracy; they do not know DP techniques.
- Accuracy: A set of queries (workload) has two accuracy parameters: α, β.

MOTIVATION

Status quo: Ge et al.'s APEx [1] provides accurate responses to workloads, while minimizing ε.

Drawbacks:

- Does not consider prior noisy responses ⇒ more ε spent than required.
- Analysts may not know how to plan workloads for minimial ε usage.

CONTRIBUTIONS

An inference engine, **CacheDP** for interactive DP queries that offers the following:

- Accurate responses to all analysts' workloads using lower privacy budget.
- Removes the need for analysts to plan subsequent workloads by maintaining a cache.
- Uses the cache to save privacy budget for related workloads.

WORK-IN-PROGRESS

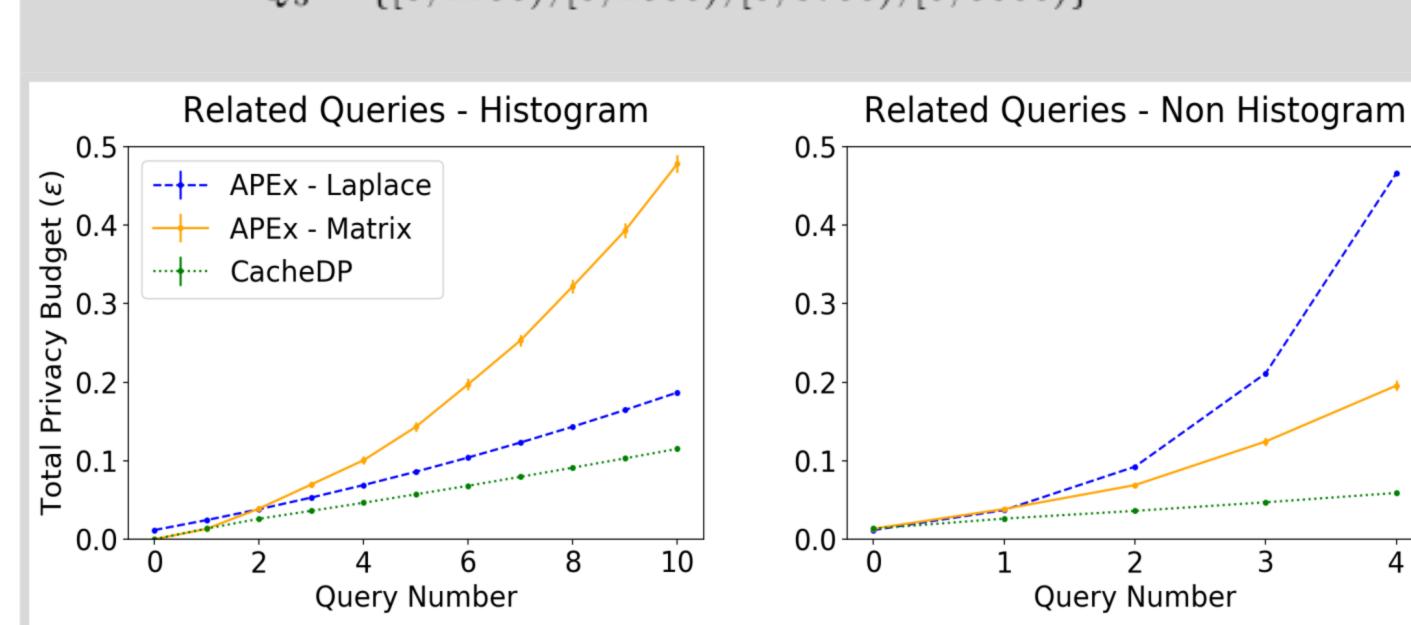
- Choosing a cache-aware optimal strategy matrix.
- Implementing our algorithm and cache structures.
- Evaluating our implementation for the following use-cases:
 - Entity resolution tasks.
 - Private spatial data exploration.
- Extending to multiple dimensions [4].

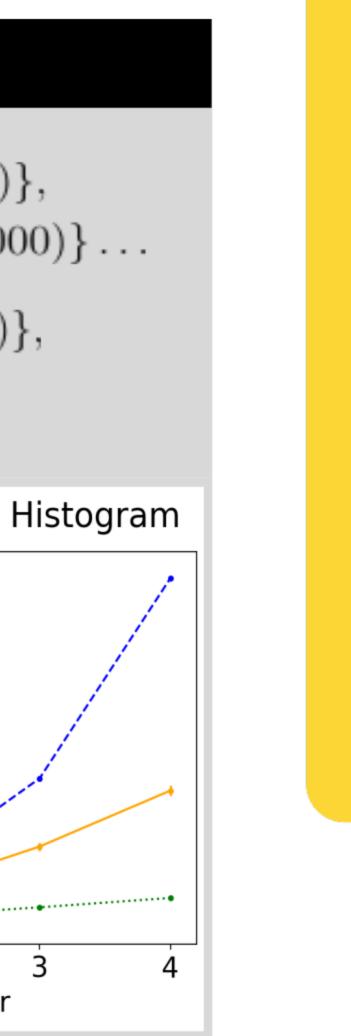
BACKGROUND **Accuracy Params:** error bound α, failure rate β. $Pr[\|Q - \tilde{Q}\|_{\infty} \ge \alpha] \le \beta$ Matrix Mechanism [3]: Example Strategy 1.25K < 2.5K 2.5K < 3.75K 3.75K < 5K

RELATED QUERIES EXAMPLES

Histogram: $Q_1 = [0, 5000), Q_2 = \{[0, 2500), [2500, 5000)\},\$ $Q_3 = \{[0, 1250), [1250, 2500), [2500, 3750), [3750, 5000)\}...$

Non-Histogram: $Q_1 = [0, 5000), Q_2 = \{[0, 2500), [0, 5000)\},\$ $Q_3 = \{[0, 1250), [0, 2500), [0, 3750), [0, 5000)\}...$





SYSTEM DIAGRAM Query to Workload Choose optimal Add Proactive Queries? (Modified Matrix Mechanism) Cache strategy queries in cache? Yes Paid: sample from Paid: sample from Free: get from cache Free: get < from cache Estimate ε using MC [1]. Relax privacy [2]: ϵ spent $\leq \epsilon$ stimate get from and Output: $\varepsilon = \text{only cost of}$ (Construct responses) as if **didn't** occur. to workload using Output: output

REFERENCES

- [2] Fragkiskos Koufogiannis, Shuo Han, and George J. Pappas. 2016. Gradual Release of Sensitive Data under Differential Privacy. Journal of Privacy & Confidentiality [3] Chao Li, Gerome Miklau, Michael Hay, Andrew Mcgregor, and Vibhor Rastogi. 2015. The Matrix Mechanism: Optimizing Linear Counting Queries under Differential Privacy. The VLDB Journal.
- [4] Ryan McKenna, Gerome Miklau, Michael Hay, and Ashwin Machanavajjhala. 2018. Optimizing Error of High-Dimensional Statistical Queries under Differential Privacy. VLDB '18.

PROACTIVE APPROACH

- Includes disjoint queries within the workload.
- Exploits the parallel composition theorem.
- Threshold to bound a minor increase in the privacy budget (for meeting the accuracy guarantee of overall workload)

DISJOINT AND REPEATED EXAMPLES

Disjoint: $Q_1 = [0, 5000), Q_2 = [0, 500), Q_3 = [500, 1000),$ $Q_4 = [500, 1000) \dots Q_{11} = [4500, 5000)$

Repeated: $Q_1 = [0, 5000), \alpha = 0.25|D|, \beta = 0.0005.$ $Q_2 = [0, 5000), \alpha = 0.125|D|, \beta = 0.0005.$

